

# Data Structures

Built-in python data structures and relevant notes:

Structure	Python	Relevant Notes
Vector	list()	<code>append</code> , <code>pop</code> , <code>insert</code> , <code>remove</code> , <code>extend</code> , <code>index</code> , <code>clear</code>
HashMap	<code>dict()</code> , <code>collections.defaultdict(lambda: 0)</code>	<code>d[k]=v</code> , <code>d.pop(k)</code> . CPython uses open addressing and random probing to solve collisions.
HashSet	set()	<code>add</code> , <code>update</code> , <code>remove</code> , <code>clear</code> , <code>union</code> , <code>intersection</code>
Stack	list()	
Deque	<code>collections.deque</code>	<code>rotate</code> , <code>append</code> , <code>appendleft</code> , <code>pop</code> , <code>popleft</code> , <code>extend</code> , <code>extendleft</code>
Priority Queue	heapq	<code>heapify</code> , <code>heappush</code> , <code>heappop</code> , <code>nlargest</code> , <code>nsmallest</code>

Function (Python)	Relevant Notes
<code>sorted(iterable, key=key, reverse=reverse)</code>	Ascending sort of an iterable collection
<code>reversed(sequence)</code>	Reverses a sequence (lists, strings, tuples, ...)
<code>bin(number)</code>	Binary string representation of a number

## Tree

- acyclic graph (root + children)
- given the height of tree as H:
  - $O(H)$  lookup
  - $O(H)$  insert
  - $O(H)$  delete

```
class Node:
    def __init__(self, val, children = []):
        self.val = val
        self.children = children
```

## Binary Tree

- a tree with at most 2 children
- no certainty regarding tree height, hence:
  - $O(H)$  lookup
  - $O(H)$  insert

- $O(H)$  delete

```
class Node:
    def __init__(self, val, left = None, right = None):
        self.val = val
        self.left = left
        self.right = right
```

## Binary Search Tree

- a binary tree where  $\text{left} < \text{root} < \text{right}$
- no certainty regarding tree height, hence:
  - $O(H)$  lookup
  - $O(H)$  insert
  - $O(H)$  delete

```
class Node:
    def __init__(self, val, left = None, right = None):
        if left: assert(left.val < val)
        if right: assert(val <= right.val)

        self.val = val
        self.left = left
        self.right = right
```

## Balanced Binary Search Tree

- a binary search tree where the height difference between subtrees is at most 1
- the height  $H$  is balanced, hence with  $N$  nodes height is  $\log N$ , thus:
  - $O(\log N)$  lookup
  - $O(\log N)$  insert
  - $O(\log N)$  delete
- insertions and deletions possibly make the tree unbalanced, self-balancing trees correct this through rotations (e.g. AVL)

```
class Node:
    def __init__(self, val, left = None, right = None):
        if left: assert(left.val <= val)
        if right: assert(val <= right.val)
        if left and right: assert(abs(left.height() - right.height()) <= 1)

        self.val = val
        self.left = left
        self.right = right
```

```
def height(self):
    left_h = self.left.height() if self.left else 0
    right_h = self.right.height() if self.right else 0
    return max(left_h, right_h) + 1
```

## Trie

- trees of characters
- terminal nodes (leaves) represent words
- allows caching of current prefix and current node for efficient search
- given the prefix length of K:
  - $O(K)$  lookup

```
class Trie:
    def __init__(self):
        self.children = {}
        self.terminal = False

    def insert(self, word):
        cur = self
        for c in word:
            if c not in cur.children:
                cur.children[c] = Trie()
            cur = cur.children[c]
        cur.terminal = True

    def remove(self, word):
        cur = self
        for c in word:
            if c not in cur.children:
                break
            cur = cur.children[c]
        cur.terminal = False

    def search(self, word):
        cur = self
        for c in word:
            if c not in cur.children:
                return False
            cur = cur.children[c]
        return cur.terminal
```

## Heap (Max)

- balanced binary tree
- root is bigger than children (recursive definition meaning maximum is at the top)
- insertion is done by inserting new element in the last spot and bubbling it up, swapping with parent if needed

- deletion is done by removing element and replacing by the last element added, swapping it down with the max child
- great for sorting or priority queues
- balanced binary tree:
  - $O(\log N)$  insertions/deletions
  - $O(1)$  query max

## Disjoint Set

- keeps track of multiple sets of elements, disjoint at first
- allows fast check of disjoint sets of elements
- implemented as a simple array that keeps track of set parents
- `union(x, y)` should set `x` and `y` to the same set ( $O(1)$ )
- `find(x)` should return the set `x` belongs to ( $O(N)$ )
  - can be made  $O(\log N)$  if we track the size and chain to the smallest, guaranteeing at max  $\log N$  length of each chain

```
class DisjointSet:
    def __init__(self):
        self.groups = dict()

    def find(self, x):
        if x not in self.groups:
            self.groups[x] = x

        while x != self.groups[x]:
            x = self.groups[x]
        return x

    def union(self, x, y):
        root_x = self.find(x)
        root_y = self.find(y)
        self.groups[root_x] = root_y
```

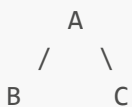
## Graph

- collection of vertices ( $V$ ) and edges ( $E$ )
- adjacency matrix representation (good for dense graphs):  $V * V$  matrix with distances (0, inf, x)
- adjacency list representation (good for sparse graphs): list of lists of neighbors

## Algorithms

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Example tree:





## Tree Traversal

Method	Order	Example
Pre (dfs)	<b>root</b> left right	A B D E C F G
In	left <b>root</b> right	D B E A G F C
Post	left right <b>root</b>	D E B G F C A

```
def preorder(root):
    if root == None: return
    print(root)
    preorder(root.left)
    preorder(root.right)

def inorder(root):
    if root == None: return
    inorder(root.left)
    print(root)
    inorder(root.right)

def postorder(root):
    if root == None: return
    postorder(root.left)
    postorder(root.right)
    print(root)
```

## Minimum Spanning Tree (MST)

- a tree that contains all nodes of the original one with a minimal sum of edge weights

### Kruskal's Algorithm

- select minimum cost edges that do not form a cycle
- pop them one by one, using those that do not connect two already used vertices (disjoint set)
- stop when all vertices are connected

```
def kruskal(edges):
    edges.sort()

    mst = []
```

```

disjoint_set = DisjointSet()
while len(edges) > 0:
    cost, src, dst = edges.pop(0)

    # disjoint set keeps track of connectivity
    if disjoint_set.find(src) != disjoint_set.find(dst):
        disjoint_set.union(src, dst)
        mst.append((cost, src, dst))

return mst

```

## Binary Search

- cut the search space in half each iteration (logarithmic complexity)
- requires a sorted collection and monotonicity
- $O(\log N)$

```

def bin_search(nums, target):
    lb, ub = 0, len(nums) - 1
    while lb <= ub:
        mid = lb + (ub - lb) // 2
        if nums[mid] < target:
            lb = mid + 1
        elif nums[mid] > target:
            ub = mid - 1
        else:
            return mid
    return -1

```

## Depth-First Search (DFS)

- LIFO approach
- search leftmost first, backtracking when needed
- useful to detect graph cycles too
- $O(V)$
- example: A B D E C F G (tree pre-order)

```

# recursive
def dfs(root):
    if root == None: return
    print(root)
    for child in root.children:
        dfs(child)

# stack based
def dfs(root):
    stack = [root]

```

```

while len(stack) > 0:
    top = stack.pop()
    print(top)
    for child in reversed(top.children):
        if child == None: continue
        stack.append(child)

```

## Breadth-First Search (BFS)

- FIFO approach
- explore all nodes in a "level" before going deeper
- $O(V)$
- example: A B C D E F G

```

from collections import deque

def bfs(root):
    queue = deque([root])
    while len(queue) > 0:
        front = queue.popleft()
        print(front)
        for child in front.children:
            if child == None: continue
            queue.append(child)

```

## Dijkstra

- greedy algorithm to find the shortest path from one node to all others
- no negative weight edges allowed
- $O((V + E) \log V)$  with min-heap:
  - heappush once per edge  $\rightarrow E \log V$
  - heappop once per node  $\rightarrow V \log V$

```

from heapq import heappush, heappop

def dijkstra(graph, src):
    dists = [float("inf")] * len(graph)
    dists[src] = 0

    visited = set()
    pq = [(0, src)]
    while len(pq) > 0:
        (_, cur) = heappop(pq)

        if cur in visited:
            continue
        visited.add(cur)

```

```

    # for each neighbor check if the cost of going
    # from current to neighbor is lower than neighbor distance
    for (neighbor, cost) in enumerate(graph[cur]):
        alt = dists[cur] + cost
        if alt < dists[neighbor]:
            dists[neighbor] = alt
            heappush(pq, (dists[neighbor], neighbor))

    return dists

```

## Bellman-Ford

- finds the shortest path from one node to all others
- relaxes edges  $V-1$  times, quitting early if no distance improves
- works for negative edges
- does not work with negative cycles but detects them

```

def bellman_ford(graph, src):
    n_vertices = len(graph)

    dists = [float("inf")] * n_vertices
    dists[src] = 0

    for _ in range(n_vertices - 1):
        # for each neighbor check if the cost of going
        # from current to neighbor is lower than neighbor distance
        for i in range(n_vertices):
            for j in range(n_vertices):
                alt = dists[i] + graph[i][j]
                if alt < dists[j]:
                    dists[j] = alt

    return dists

```

## Floyd-Warshall

- shortest path between all nodes
- $O(V^3)$

```

def floyd_warshall(graph):
    n_vertices = len(graph)

    dists = [[float("inf") for _ in range(n_vertices)] for _ in range(n_vertices)]
    for i in range(n_vertices):
        for j in range(n_vertices):
            dists[i][j] = graph[i][j]

    for intermediate in range(n_vertices):

```



```

        for i in range(n_vertices):
            for j in range(n_vertices):
                alt = dists[i][intermediate] + dists[intermediate][j]
                if dists[i][j] > alt:
                    dists[i][j] = alt

    return dists

```

## Cycle Detection

Detecting cycles in a graph can be done in several ways:

- DFS: check if a node has been visited twice

```

def has_cycle(root):
    visited = set()
    stack = [root]
    while len(stack) > 0:
        top = stack.pop()

        if top in visited:
            return True

        for child in reversed(top.children):
            if child == None:
                continue
            stack.append(child)

    return False

```

- Disjoint Set: union nodes for each edge and quit if same set is found

```

def has_cycle(graph):
    n_vertices = len(graph)
    disjoint_set = DisjointSet(n_vertices)

    for i in range(n_vertices):
        for j in range(n_vertices):
            if i == j or graph[i][j] == float("inf"):
                continue

            if disjoint_set.find(i) == disjoint_set.find(j):
                return True

            disjoint_set.union(i, j)

    return False

```

- Bellman-Ford: run an extra cycle and if it improves there is a cycle

```
def has_cycle(graph, src):
    n_vertices = len(graph)
    dists = bellman_ford(graph, src)

    # run an extra cycle to see if anything improves
    for i in range(n_vertices):
        for j in range(n_vertices):
            alt = dists[i] + graph[i][j]
            if alt < dists[j]:
                return True

    return False
```

- Tortoise & Hare: if both pointers meet, there is a cycle

```
def has_cycle(root):
    slow, fast = root, root
    while fast and fast.next:
        slow = slow.next
        fast = fast.next.next
        if slow == fast:
            return True
    return False
```

## Dynamic Programming

- applicable when optimal solution depends on the optimal solution for subproblems
- bottom-up: solve base cases and compound results
- top-down: memoization, cache results and avoid recomputation, easily applied to recursive solutions

## Quick Sort

- recursively sort halves, partitioned by a pivot
- swap left and right elements of the pivot and call quick sort on both halves
- $O(N * \log N)$

```
def quicksort(collection):
    return _quicksort(collection, 0, len(collection) - 1)

def _quicksort(collection, left, right):
    if left >= right:
        return

    pivot = collection[(left + right) // 2]
```

```
split = partition(collection, left, right, pivot)
_quicksort(collection, left, split - 1)
_quicksort(collection, split, right)
return collection

def partition(collection, left, right, pivot):
    while left <= right:
        while collection[left] < pivot:
            left += 1

        while collection[right] > pivot:
            right -= 1

        if left <= right:
            tmp = collection[left]
            collection[left] = collection[right]
            collection[right] = tmp
            left += 1
            right -= 1

    return left
```

## Merge Sort

- recursively sort halves, call merge sort on each
- copy elements in order to a new array
- $O(N * \log N)$

```
def mergesort(collection):
    if len(collection) <= 1:
        return collection

    middle = len(collection) // 2
    left = mergesort(collection[:middle])
    right = mergesort(collection[middle:])
    merged = merge(left, right)
    return merged

def merge(left, right):
    merged = []

    l, r = 0, 0
    while l < len(left) and r < len(right):
        if left[l] < right[r]:
            merged.append(left[l])
            l += 1
        else:
            merged.append(right[r])
            r += 1
```

```
while l < len(left):
    merged.append(left[l])
    l += 1

while r < len(right):
    merged.append(right[r])
    r += 1

return merged
```

## Heap Sort

- build an heap (heapify  $O(N)$ )
- keep popping the min element into a new array
- $O(N * \log N)$ :
  - the popped top element will be replaced by a leaf and bubbled down

```
from heapq import heapify, heappush, heappop

def heapsort(collection):
    heapify(collection)
    return [heappop(collection) for _ in range(len(collection))]
```

## Object Oriented Programming (OOP)

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### SOLID

**Single Responsibility** - classes should do one thing and do it well, having one reason to change

**Open-Closed** - classes should be open for extension and closed for modifications

**Liskov Substitution** - classes should be substituted for parent classes or interfaces they implement

**Interface Segregation** - keep interfaces thin, split big ones into smaller contracts, each client implements what is needed

**Dependency Inversion** - entities depend on abstractions and not on concretions

### Design Patterns

Typical solutions for common software OOP design problems.

**Creational** - objects' creation

- **Factory** - interface for creating objects, simplifying and centralizing logic

```

class Burger:
    def __init__(self, ingredients):
        self.ingredients = ingredients

class BurgerFactory:
    @classmethod
    def create_cheese_burger(cls):
        return Burger(["bun", "cheese", "beef-patty"])

    @classmethod
    def create_deluxe_burger(cls):
        return Burger(["bun", "cheese", "beef-patty", "tomatoe", "lettuce"])

```

- **Builder** - construct complex objects step by step

```

class Burger:
    def __init__(self):
        self.buns = None
        self.patty = None

    def set_buns(self, buns):
        self.buns = buns

    def set_patty(self, patty):
        self.patty = patty

class BurgerBuilder:
    def __init__(self):
        self.burger = Burger()

    def build(self):
        return self.burger

    def add_buns(self, buns):
        self.burger.set_buns(buns)
        return self

    def add_patty(self, patty):
        self.burger.set_patty(patty)
        return self

```

- **Singleton** - ensure a single instance of a class

```

class Singleton:
    _instance = None

    @classmethod
    def instance(cls):
        if cls._instance == None:

```

```
cls._instance = cls()
return cls._instance
```

**Behavioral** - objects' communication (events / state changes)

- **Iterator** - defines how the values in a collection are iterated through

```
class LinkedList:
    def __init__(self):
        self.head = None
        self.cur = None

    def __iter__(self):
        self.cur = self.head
        return self

    def __next__(self):
        if self.cur == None:
            raise StopIteration

        val = self.cur.val
        self.cur = self.cur.next
        return val
```

- **Command** - turns actions into objects (e.g. useful for queues, delays, undo/redo, event sourcing, ...)

```
class Command:
    def execute(self):
        pass

class KillCommand(Command):
    def __init__(self, program):
        self.program = program

    def execute(self):
        self.program.kill()

class RestartCommand(Command):
    def __init__(self, program):
        self.program = program

    def execute(self):
        self.program.restart()
```

- **Observer** - subscription/notification of objects to events

```
class Publisher:
    def __init__(self):
        self.subscribers = []

    def subscribe(self, sub):
        self.subscribers.append(sub)

    def notify(self, event):
        for sub in self.subscribers:
            sub.notify(event)
```

- **Strategy** - define a family of interchangeable algorithms

```
class FilterStrategy:
    def filter(self, val):
        pass

class FilterPositives(FilterStrategy):
    def filter(self, val):
        return val > 0

class FilterNegatives(FilterStrategy):
    def filter(self, val):
        return val < 0

def filter_fn(values, strategy):
    return [x for x in values if strategy.filter(x)]
```

## Structural - objects' assembly

- **Facade** - a wrapper used to abstract lower-level details

```
class Vector:
    ...
```

- **Adapter** - allow objects with incompatible interfaces to communicate

```
class UsbCable:
    def __init__(self):
        self.port = None

    def plugUsb(self, port):
        self.port = port

class UsbPort:
    def plug(self, cable):
```

```
        cable.plugUsb(self)

class MicroUsbCable:
    def __init__(self):
        self.port = None

    def plugMicroUsb(self, port):
        self.port = port

class MicroToUsbAdapter(UsbCable):
    def __init__(self, microUsbCable):
        self.microUsbCable = microUsbCable

    def plugUsb(self, port):
        super().plugUsb(port)
        self.microUsbCable.plugMicroUsb(port)
```

- **Decorator** - wrap objects with additional functionality

```
class Text:
    def __init__(self, text):
        self.text = text

    def render(self):
        return self.text

class UnderlineText(Text):
    def __init__(self, wrapped):
        self.wrapped = wrapped

    def render(self):
        return "<u>" + self.wrapped.render() + "<u>"

class BoldText(Text):
    def __init__(self, wrapped):
        self.wrapped = wrapped

    def render(self):
        return "<b>" + self.wrapped.render() + "<b>"
```