Detailed explanation of WCH TMOS usage

Original post: https://www.cnblogs.com/debugdabiaoge/p/15775521.html

In order to connect with multiple devices and achieve multi-function and multi-tasking, Bluetooth has a scheduling problem. Although the software and protocol stack can be expanded, after all, there is only one bottom-level execution department. In order to realize multi-event and multi-task switching, it is necessary to correspond events and tasks, and a TMOS name operating system abstraction layer is set up for this application.

TMOS is the scheduling core, and the BLE protocol stack, profile definition, and all applications are implemented around it. TMOS is not the traditional operating system that everyone uses, but a cycle that allows software to create and execute events.

In the multi-task management mode, only one task is actually running, but multiple tasks can be scheduled using the task scheduling strategy, and each task takes a certain amount of time (exclusive, exit after executing the current task, and continue to query other executable tasks), all tasks are processed by time slicing.

The **TMOS** system clock unit is **625us**, and all the required system time is obtained based on the RTC.

Software functions are realized by task events. Creating a task event requires the following work:

1. Create task identifier task ID

For example:

```
| The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation. | The id definition is a global variable, initialisation is a global variable,
```

2. Write a task initialization routine process and add it to the TMOS initialization process

Which means that the function cannot be dynamically added after the system starts (new Task ID).

```
PRINT("%s\n", VER_LIB);

CHSTY_BIRINI(');

HAL_Init(');

GAPRole_PeripheralInit(');

Peripheral_Init(');
```

3. Write a task handler

```
Peripheral ProcessEvent
     * @fn
     * @brief Peripheral Application Task event processor. This function
300
301
                  is called to process all events for the task.
                 include timers, messages and any other user defined events.
303
      @param task_id - The TMOS assigned task ID.
@param events - events to process. This is
304
                                                 This is a bit map and can
306
                           contain more than one event.
307
308
     * @return events not processed
                                                                          Parameters
309 */
310@uint16 Peripheral_ProcessEvent( uint8 task_id, uint16 events ) 参数
312
313 // VOID task_id; // TMOS required parameter that isn't used in this function
3150 if ( events & SYS_EVENT_MSG ) {
316    uint8 *pMsg;
```

4. Define task events and write user function codes

The event name is defined by bit, each **taskID** contains at most 1 message event and 15 task events (total 16 bits). For example:

Define EVT task events bit by bit, as shown in the figure below:

There are two ways to start the task event (the task is only executed once after the task is started, if it is executed repeatedly, the task needs to be restarted):

1) The task is started immediately, and the event time is executed immediately after the call

```
20569/***
     * @fn
2057
                   tmos_set_event
2058
     * @brief
                   start a event immediately
2059
2060 *
2061
     * input parameters
2062
     * @param
2063
                    taskID - task ID of event
2064
     * @param
                   event - event value
2065
2066
     * output parameters
2067
      * @param
2068
                    None.
2069
     * @return
                    0 - success.
2072 extern bStatus_t tmos_set_event( tmosTaskID taskID, tmosEvents event );
```

For example:

```
// Setup a delayed profile startup
tmos_set_event(Peripheral_TaskID, SBP_START_DEVICE_EVT);
278 }
```

2) Set a delay to start a task, and start timing after the setting is completed

For example: the custom **SBP_PERIODIC_EVT** task under the **Peripheral_TaskID** function runs after a delay of (**SBP_READ_RSSI_EVT_PERIOD***625)us.

User code function. When generating **TaskID**, you need to register the EVT processing function pointer with TMOS. After the EVT execution conditions are met, TMOS will automatically call this function, as shown in the following figure:

```
310@uint16 Peripheral_ProcessEvent( uint8 task_id, uint16 events )
313 // VOID task_id; // TMOS required parameter that isn't used in this function
3150 if ( events & SYS EVENT MSG ) {
        uint8 *pMsg;
        if ( (pMsg = tmos_msg_receive( Peripheral_TaskID )) != NULL ){[]
318⊕
         // return unprocessed events
return (events ^ SYS_EVENT_MSG);
324
 325
 3270 if ( events & SBP START DEVICE EVT | {...
3330 if ( events & SBP PERIODIC EVT )
344⊕ if ( events & SBP_PARAM_UPDATE_EVT ).
357⊕ if ( events & SBP_READ_RSSI_EVT ).
364
       // Discard unknown events
      return 0;
```

extern bStatus t tmos stop task(tmosTaskIDtaskID, tmosEvents event);

This function will stop a task named event that will take effect at the **taskID** layer. After calling this function, the **event** task will not take effect.

5. The main loop calls TMOS_SystemProcess continuously to query executable events

If **HAL_SLEEP** starts, the chip turns on low-power sleep mode, TMOS will turn on the RTC wakeup function, and it will automatically wake up before the event is executed, and run the event code.

Precautions for using the task scheduling function:

- 1. It is forbidden to call in interrupt
- 2. It is recommended not to execute tasks that exceed half the connection interval in a single task, otherwise it will affect Bluetooth communication
- 3. In the same way, it is recommended not to perform tasks that exceed half the connection interval during the interruption, otherwise it will affect the Bluetooth communication
- 4. When the delayed execution function is called in the code executed by the event, the delay time is offset from the current event effective time point, so there is no requirement for the position of the delayed execution function called in the executed code.
- 5. Tasks have priority, which is determined according to the sequence of judgments in the **xxx_ProcessEvent** function. Tasks that are effective at the same time are executed first and judged first, and then judged after execution. Note that after executing the first-judgment event task, the last-judgment event task will not be executed until the task scheduling system takes turns.
- 6. The event name is defined by bit. Each **taskID** contains at most 1 message event and 15 task events (total 16 bits)

I talked about the application of a Task ID earlier. In order to reduce the coupling between C files or functions, it is generally better to put the same function or similar events under the same Task ID, which creates a problem; different Task IDs may have data that needs to be interacted with, and TMOS provides functions for data interaction between different Task IDs.

For example: Take two of the Task IDs in the peripheral as examples, **halTaskID** and **Peripheral TaskID**, assuming that data interaction between these two Tasks is required.

It was said before that each Task ID has 1 message event

```
186 */
1870 tmosEvents HAL ProcessEvent( tmosTaskID task_id, tmosEvents events )
189
     uint8 * msgPtr;
                                       Task ID-2
190
1910
    if ( events & SYS_EVENT_MSG )
            处埋HAL层消息,调用tmos_ms
192
                                  receive读歌消息,处理完成后静脉消息。
       msgPtr = tmos_msg_receive( task_id );
193
1948
       if ( msgPtr )
195
      {
        /* De-allocate */
196
        tmos_msg_deallocate( msgPtr );
197
198
199
       return events ^ SYS_EVENT_MSG;
     }
200
```

The above figure demonstrates receiving messages, which mainly uses 2 functions:

```
2183 * @fn
                   tmos_msg_receive
2184
      * @brief
                    receive a msg
 2187
      * input parameters
                    taskID - task ID of task need to receive msg
 2189
      * @param
                                              Receive a message
 2190
 2191
      * output parameters
                                              收消息
 2192
      * @param
 2194
      * @return *u8 - message information or NULL if no message
 2196 */
 2197 extern u8 *tmos msg receive( tmosTaskID taskID );
 * @fn
                tmos_msg_allocate
 * @brief
                 allocate buffer for msg when need to send msg
 * input parameters
                                               Deallocate message buffer
 * @param
                 len - length of msg
 * output parameters
 * @param
 * @return
                 pointer to allocated buffer or NULL if allocation failed.
extern u8 *tmos msg allocate( u16 len );
57 void HAL KEY RegisterForKeys ( tmosTaskID id )
                                                   应用层调用的注册函数, 保存传递进来
                                                   的taskID值
     registeredKeysTaskID = id;
 59
                                                   Registration function called by the application
 60
                                                   layer, saves the passed taskID value.
 void HalKeyConfig (uint8 interruptEnable, halKeyCBack t cback)
102 uint8 OnBoard SendKeys( uint8 keys, uint8 state )
                                                   分配要发送的消息内存, 如果申请内存
103 ⊟ {
                                                  成功,再进行下面的赋值发送。
104
      keyChange_t *msgPtr;
                                                   Allocate memory for the message to be sent,
105
                                                   and if the allocation is successful, send the
106
     if ( registeredKeysTaskID != TASK_NO_TASK ) {
107
      msgPtr = (keyChange_t *)tmos_msg_allocate( sizeof(keyChange_t) );
108
109
       if ( msgPtr ) {
         msgPtr->hdr.event = KEY_CHANGE;
110
       msgPtr->state = state;
msgPtr->keys = keys;
111
112
         tmos_msg_send( registeredKeysTaskID, (uint8 *)msgPtr );
113
114
                               调用发送消息函数,参数为此前注册函数保存的应用层taskID,以
       return ( SUCCESS );
115
                               及消息指针
116
                               Call the send message function. The parameters are the application layer taskID
117
                               saved by the registration function, and the message point
118
       return ( FAILURE );
119
```

```
2194 * @fn
2195 *
               tmos_msg_allocate
2196 * @brief
               allocate buffer for msg when need to send msg
2198 * input parameters
2199 *
                                        Allocate message buffer
2199 *
2200 * & param len - length of msg
2201 *
2202 * output parameters
2203 *
                                        申请内存
2204 * @param
                None.
2206 * @return pointer to allocated buffer or NULL if allocation failed. 2207 */
2208 extern u8 *tmos msg allocate ( u16 len );
2142 * @fn
2143 *
                tmos_msg_send
2144 * @brief
                send msg to a task, callback events&SYS_EVENT_MSG
2145 *
2146 * input parameters
2147
2148 * @param
2149 * @param
                 taskID - task ID of task need to send msg
                *msg_ptr - point of msg
Send a message
2150 *
2151 * output parameters
2152 *
2153 * @param
                None.
2154 * 2155 * @return 0 - success. 2156 */
2157 extern bStatus_t tmos_msg_send( tmosTaskID taskID, u8 *msg_ptr );
```