

Π -Ware: An Embedded Hardware Description Language using Dependent Types

Author: João Paulo Pizani Flor
<joaopizani@uu.nl>

Supervisor: Wouter Swierstra
<w.s.swierstra@uu.nl>

Department of Information and Computing Sciences
Utrecht University

Sunday 24th August, 2014

Background

Hardware Design
Functional Hardware
DTP

Research
Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Universiteit Utrecht

Table of Contents

Background

Hardware Design

Functional Hardware

DTP

Research Question

Question

Method

DTP / Agda

Big picture

Agda

Π -Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work

Background

Hardware Design

Functional Hardware

DTP

Research Question

Question

Method

DTP / Agda

Big picture

Agda

Π -Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Hardware design is hard(er)

- ▶ Strict(er) correctness requirements
 - You can't simply *update* a full-custom chip after production
 - Intel FDIV
 - Expensive verification / validation (up to 50% of development costs)
- ▶ Low-level details (more) important
 - Layout / area
 - Power consumption / fault tolerance

Hardware Design

Functional Hardware
DTP

Research Question

Question

Method

DTP / Agda

Big picture
Agda

Π-Ware

- Syntax
- Semantics
- Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Hardware design is growing

- ▶ Moore's law will still apply for some time
 - We can keep packing more transistors into same silicon area
- ▶ **But** optimizations in CPUs display diminishing returns
 - Thus, more algorithms *directly* in hardware

Background

Hardware Design

Functional Hardware

DTP

Research Question

Question

Method

DTP / Agda

Big picture

Agda

Π -Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Hardware Description Languages

- ▶ All started in the 1980s
- ▶ *De facto* industry standards: VHDL and Verilog
- ▶ Were intended for *simulation*, not modelling or synthesis
 - *Unsynthesizable* constructs
 - Widely variable tool support

Background

Hardware Design

Functional Hardware

DTP

Research Question

Question

Method

DTP / Agda

Big picture

Agda

Π-Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Functional Programming

- ▶ Easier to *reason* about program properties
- ▶ Inherently *parallel* and *stateless* semantics
 - In contrast to imperative programming

Background

Hardware Design

Functional Hardware

DTP

Research Question

Question

Method

DTP / Agda

Big picture

Agda

Π -Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Functional Hardware Description

- ▶ A functional program describes a circuit
- ▶ Several *functional* Hardware Description Languages (HDLs) during the 1980s
 - For example, μ FP [Sheeran, 1984]
- ▶ Later, *embedded* hardware Domain-Specific Languages (DSLs)
 - For example, Lava (Haskell) [Bjesse et al., 1998]

Background

Hardware Design

Functional Hardware

DTP

Research Question

Question

Method

DTP / Agda

Big picture

Agda

Π-Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Embedded DSLs for Hardware

- ▶ Lava
- ▶ Limitations
 - Low level types
 - Not guaranteeing size match

Background

Hardware Design

Functional Hardware

DTP

Research

Question

Question

Method

DTP / Agda

Big picture

Agda

Π-Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Dependently-Typed Programming

Dependently-Typed Programming (DTP) är en
programmationstechnik...

Background

- Hardware Design
- Functional Hardware
- DTP**

Research Question

- Question
- Method

DTP / Agda

- Big picture
- Agda

Π -Ware

- Syntax
- Semantics
- Proofs

Conclusions

- Limitations
- Future work



Research Question

“What are the improvements that DTP can bring to hardware design?”

Background

- Hardware Design
- Functional Hardware
- DTP

Research Question

- Question
- Method

DTP / Agda

- Big picture
- Agda

Π -Ware

- Syntax
- Semantics
- Proofs

Conclusions

- Limitations
- Future work



Methodology

- ▶ Develop a hardware DSL, *embedded* in a dependently-typed language (Agda)
 - Called **Π -Ware**
 - allowing simulation, synthesis and verification

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Dependently-Typed Programming

► Types can depend on values

- Example: `data Vec (α : Set) : $\mathbb{N} \rightarrow$ Set where...`
- Compare with Haskell (GADT style):
`data List :: * -> * where...`

► Types of arguments can depend on *values of previous arguments*

- Ensure a “safe” domain
- `take : (m : \mathbb{N}) \rightarrow Vec α ($m + n$) \rightarrow Vec α m`

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Dependently-Typed Programming

- ▶ Type checking requires *evaluation* of functions
 - We want `Vec Bool (2 + 2)` to unify with `Vec Bool 4`
- ▶ Consequence: all functions must be *total*
- ▶ Termination checker ensures (heuristics)
 - Structurally-decreasing recursion
 - This passes the check:
`add : $\mathbb{N} \rightarrow \mathbb{N} \rightarrow \mathbb{N}$`
`add zero y = y`
`add (suc x') y = suc (add x' y)`
 - This does not:
`silly : $\mathbb{N} \rightarrow \mathbb{N}$`
`silly zero = zero`
`silly (suc n') = silly [n' /2]`

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Universiteit Utrecht

Dependently-Typed Programming

- ▶ Dependent pattern matching can *rule out* impossible cases

Background

- Hardware Design
- Functional Hardware
- DTP

Research Question

- Question
- Method

DTP / Agda

- Big picture
- Agda

Π -Ware

- Syntax
- Semantics
- Proofs

Conclusions

- Limitations
- Future work



Dependently-Typed Programming

► Dependent pattern matching can *rule out* impossible cases

- Classic example: *safe head* function

$\text{head} : \text{Vec } \alpha \ (\text{suc } n) \rightarrow \alpha$

$\text{head } (x :: xs) = x$

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Dependently-Typed Programming

- ▶ Dependent pattern matching can *rule out* impossible cases

- Classic example: *safe head* function

$$\text{head} : \text{Vec } \alpha \ (\text{suc } n) \rightarrow \alpha$$
$$\text{head } (x :: xs) = x$$

- The **only** constructor returning $\text{Vec } \alpha \ (\text{suc } n)$ is $_::_$

Background

Hardware Design

Functional Hardware

DTP

Research Question

Question

Method

DTP / Agda

Big picture

Agda

Π-Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Dependent types as logic

- ▶ Programming language / Theorem prover
 - Types as propositions, terms as proofs [Wadler, 2014]

- ▶ Example:

- Given the relation (drawn triangle):

```
data __≤__ : ℕ → ℕ → Set where
  z≤n : ∀ {n}                → zero ≤ n
  s≤s  : ∀ {m n} → m ≤ n → suc m ≤ suc n
```

- Proposition:

```
twoLEQFour : 2 ≤ 4
```

- Proof:

```
twoLEQFour = s≤s (s≤s z≤n)
s≤s (s≤s (z≤n : 0 ≤ 4) : 1 ≤ 4) : 2 ≤ 4
```

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π-Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Universiteit Utrecht

Agda syntax for Haskell programmers

- ▶ Liberal identifier lexing (Unicode **everywhere**)
 - $a \equiv b + c$ is a valid identifier, $a \equiv b + c$ an expression
 - Actually used in Agda's standard library
 - And in Π -Ware: \mathbb{C} , $\llbracket c \rrbracket$, \Downarrow , \Uparrow
- ▶ *Mixfix* notation
 - $_[_]_ := _$ is the vector update function: $v \ [\ # \ 3 \] := \text{true}$.
 - $_[_]_ \ v \ (\# \ 3) \ \text{true} \iff v \ [\ # \ 3 \] := \text{true}$
- ▶ Almost nothing built-in
 - $_+_ _ : \mathbb{N} \rightarrow \mathbb{N} \rightarrow \mathbb{N}$ defined in **Data.Nat**
 - $\text{if_then_else_} : \text{Bool} \rightarrow \alpha \rightarrow \alpha \rightarrow \alpha$ defined in **Data.Bool**

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Universiteit Utrecht

Agda syntax for Haskell programmers

- ▶ Implicit arguments

- Don't have to be passed if Agda can **guess** it
- Syntax: $\epsilon : \{ \alpha : \text{Set} \} \rightarrow \text{Vec } \alpha \text{ zero}$

► “For all” syntax: $\forall n \iff (n : _)$

- Where `_` means: guess this type (based on other args)
- Example:
 - $\forall n \rightarrow \text{zero} \leq n$
 - `data < : ℕ → ℕ → Set`

- ▶ It's common to combine both:

- $\forall \{ \alpha \ n \} \rightarrow \text{Vec } \alpha \ (\text{suc } n) \rightarrow \alpha \iff$
 $\{ \alpha : _ \} \{ n : _ \} \rightarrow \text{Vec } \alpha \ n \rightarrow \alpha$

Background

- Hardware Design
- Functional Hardware
- DTP

Research Question

Question

Method

DTP / Agda

Big picture
Agda

Π-Ware

- Syntax
- Semantics
- Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Low-level circuits

- Structural representation
- Untyped but *sized*

data $\mathbb{C}' : \mathbb{N} \rightarrow \mathbb{N} \rightarrow \text{Set}$

data \mathbb{C}' where

Nil : $\mathbb{C}' \text{ zero zero}$

Gate : $(g\# : \text{Gates\#}) \rightarrow \mathbb{C}' (|\text{in}| g\#) (|\text{out}| g\#)$

Plug : $\forall \{i\ o\} \rightarrow (f : \text{Fin } o \rightarrow \text{Fin } i) \rightarrow \mathbb{C}' i\ o$

DelayLoop : $(c : \mathbb{C}' (i + l) (o + l)) \{\text{comb}'\ c\} \rightarrow \mathbb{C}' i\ o$

$_ \gg' _ : \mathbb{C}' i\ m \rightarrow \mathbb{C}' m\ o \rightarrow \mathbb{C}' i\ o$

$_ |' _ : \mathbb{C}' i_1\ o_1 \rightarrow \mathbb{C}' i_2\ o_2 \rightarrow \mathbb{C}' (i_1 + i_2) (o_1 + o_2)$

$_ |+' _ : \mathbb{C}' i_1\ o \rightarrow \mathbb{C}' i_2\ o \rightarrow \mathbb{C}' (\text{suc } i_1 \sqcup i_2) o$

Background

Hardware Design

Functional Hardware

DTP

Research

Question

Question

Method

DTP / Agda

Big picture

Agda

Π -Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Atoms

- ▶ How to carry values of an Agda type in *one* wire
- ▶ Defined by the **Atomic** type class in **PiWare.Atom**

record Atomic : **Set**₁ **where**

field

Atom : **Set**

|Atom|−1 : \mathbb{N}

n→atom : **Fin** (**suc** **|Atom|**−1) → *Atom*

atom→n : *Atom* → **Fin** (**suc** **|Atom|**−1)

inv-left : $\forall i \rightarrow \text{atom} \rightarrow n \ (n \rightarrow \text{atom} \ i) \equiv i$

inv-right : $\forall a \rightarrow n \rightarrow \text{atom} \ (\text{atom} \rightarrow n \ a) \equiv a$

|Atom| = **suc** **|Atom|**−1

Atom# = **Fin** **|Atom|**

Background

Hardware Design

Functional Hardware

DTP

Research

Question

Question

Method

DTP / Agda

Big picture

Agda

Π-Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Atomic instances

- ▶ Examples of types that can be **Atomic**
 - Bool, std_logic, other multi-valued logics
 - Predefined in the library: **PiWare.Atom.Bool**
- ▶ First, define how many atoms we are interested in

|B| - 1 = 1

|B| = suc **|B|** - 1

- ▶ Friendlier names for the indices (elements of **Fin 2**)

pattern **False#** = Fz

pattern **True#** = Fs Fz

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π-Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Atomic instance (Bool)

- Bijection between $\{n \in \mathbb{N} \mid n < 2\}$ (Fin 2) and Bool

$$n \rightarrow B = \lambda \{ \text{False\#} \rightarrow \text{false}; \text{True\#} \rightarrow \text{true} \}$$

$$B \rightarrow n = \lambda \{ \text{false} \rightarrow \text{False\#}; \text{true} \rightarrow \text{True\#} \}$$

- Proof that $n \rightarrow B$ and $B \rightarrow n$ are inverses

$$\text{inv-left-B} = \lambda \{ \text{False\#} \rightarrow \text{refl}; \text{True\#} \rightarrow \text{refl}; \}$$

$$\text{inv-right-B} = \lambda \{ \text{false} \rightarrow \text{refl}; \text{true} \rightarrow \text{refl} \}$$

- With all pieces at hand, we construct the instance

$$\begin{aligned} \text{Atomic-B} = \text{record} \{ & \text{Atom} = B \\ & ; |\text{Atom}|-1 = |B|-1 \\ & ; n \rightarrow \text{atom} = n \rightarrow B \\ & ; \text{atom} \rightarrow n = B \rightarrow n \\ & ; \text{inv-left} = \text{inv-left-B} \\ & ; \text{inv-right} = \text{inv-right-B} \} \end{aligned}$$

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Gates

- ▶ Circuits parameterized by collection of *fundamental gates*
- ▶ Examples:
 - {NOT, AND, OR} ([BoolTrio](#))
 - {NAND}
 - Arithmetic, Crypto, etc.
- ▶ The definition of what means to be such a collection is in [PiWare.Gates.Gates](#)

Background

- Hardware Design
- Functional Hardware
- DTP

Research Question

Question

Method

DTP / Agda

Big picture
Agda

П-Ware

- Syntax
- Semantics
- Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

The Gates type class

$W : \mathbb{N} \rightarrow \text{Set}$

$W = \text{Vec Atom}$

record Gates : Set where

field

|Gates| : \mathbb{N}

|in| |out| : Fin |Gates| $\rightarrow \mathbb{N}$

spec : (g : Fin |Gates|)
 $\rightarrow (W (|in| g) \rightarrow W (|out| g))$

Gates# = Fin |Gates|

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Gates instances

- ▶ Example: `PiWare.Gates.BoolTrio`
- ▶ First, how many gates are there in the library

`|BoolTrio| = 5`

- ▶ Then the friendlier names for the indices

```
pattern FalseConst# = Fz
pattern TrueConst#  = Fs Fz
pattern Not#        = Fs (Fs Fz)
pattern And#        = Fs (Fs (Fs Fz))
pattern Or#         = Fs (Fs (Fs (Fs Fz)))
```

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Gates instance (BoolTrio)

- ▶ Defining the *interfaces* of the gates

```
[in] FalseConst# = 0
```

```
|in| TrueConst# = 0
```

```
|in| Not# = 1
```

$$|in|_{And\#} = 2$$

|in| Or# = 2

$$|out| = 1$$

- And the specification function for each gate

```
spec=false == [ false ]
```

```
spec-true      _      = [ true  ]
```

$$\text{spec-not} \quad (x :: \varepsilon) \quad = \text{not } x$$

spec-and $(x :: y :: \varepsilon) = [x \wedge y]$

spec-or $(x :: y :: \varepsilon) = [x \vee y]$

Background

Hardware Design

Functional Hardware

DTP

Research

Question

Question

Method

DTP / Agda

Big picture

Agda

Π-Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Gates instance (BoolTrio)

- ▶ Mapping each gate index to its respective specification

specs-BoolTrio FalseConst# = spec-false

specs-BoolTrio TrueConst# = spec-true

specs-BoolTrio Not# = spec-not

specs-BoolTrio And# = spec-and

specs-BoolTrio Or# = spec-or

- ▶ With all pieces at hand, we construct the instance

BoolTrio : Gates

```
BoolTrio = record { |Gates| = |BoolTrio|  
                  ; |in|    = |in|  
                  ; |out|   = |out|  
                  ; spec    = specs-BoolTrio }
```

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



High-level circuits

- ▶ User is not supposed to describe circuits at low level (\mathbb{C}')
- ▶ The high level circuit type (\mathbb{C}) allows for *typed* circuit interfaces
 - The input and output indices are Agda types

```
data C (α β : Set) {i j : ℕ} : Set where
  MkC : { [ sα : ↓W↑ α {i} ] [ sβ : ↓W↑ β {j} ] }
        → C' i j → C α β {i} {j}
```

- ▶ MkC takes:
 - Low level description (\mathbb{C}')
 - Information on how to *synthesize* elements of α and β
 - Passed as *instance arguments*

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π-Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Synthesizable

- ▶ $\Downarrow W \Uparrow$ type class (pronounced Synthesizable)
 - Describes how to *synthesize* a given Agda type (α)
 - Two fields: from element of α to a *word* and back

```
record  $\Downarrow W \Uparrow$  ( $\alpha$  : Set) { $i$  :  $\mathbb{N}$ } : Set where
  constructor  $\Downarrow W \Uparrow$  [ $\_$ ,  $\_$ ]
  field
```

$$\Downarrow : \alpha \rightarrow W\ i$$
$$\Uparrow : W\ i \rightarrow \alpha$$

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Universiteit Utrecht

$\Downarrow W \Uparrow$ instances

- ▶ Any *finite* type can have such an instance
- ▶ Predefined in the library: `Bool`; `_×_`; `_⊔_`; `Vec`
- ▶ Example: instance for products (`_×_`)

$$\Downarrow W \Uparrow - \times : \{ \mid s\alpha : \Downarrow W \Uparrow \alpha \{i\} \} \{ \mid s\beta : \Downarrow W \Uparrow \beta \{j\} \} \rightarrow \Downarrow W \Uparrow (\alpha \times \beta)$$

$$\Downarrow W \Uparrow - \times \{ \mid s\alpha \} \{ \mid s\beta \} = \Downarrow W \Uparrow [\text{down} , \text{up}]$$

$$\text{where } \text{down} : (\alpha \times \beta) \rightarrow W (i + j) \\ \text{down } (a , b) = (\Downarrow a) ++ (\Downarrow b)$$

$$\text{up} : W (i + j) \rightarrow (\alpha \times \beta)$$

$$\text{up } w \text{ with splitAt } i \text{ } w$$

$$\text{up } .(\Downarrow a ++ \Downarrow b) \mid \Downarrow a , \Downarrow b , \text{refl} = \Uparrow \Downarrow a , \Uparrow \Downarrow b$$

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Synthesizable

- ▶ Both fields \Downarrow and \Uparrow should be inverses of each other

Background

- Hardware Design
- Functional Hardware
- DTP

Research Question

- Question
- Method

DTP / Agda

- Big picture
- Agda

Π -Ware

- Syntax**
- Semantics
- Proofs

Conclusions

- Limitations
- Future work



Synthesis

- ▶ Work-in-progress
- ▶ **Atom** and **Gates** with VHDL *abstract syntax*

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Simulation

- ▶ Combinational
- ▶ Sequential

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Examples

► AndN

Background

Hardware Design

Functional Hardware

DTP

Research

Question

Question

Method

DTP / Agda

Big picture

Agda

Π-Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Problems

- Definition of `[[_]]` blocks reduction

Background

- Hardware Design
- Functional Hardware
- DTP

Research Question

- Question
- Method

DTP / Agda

- Big picture
- Agda

Π -Ware

- Syntax
- Semantics

Proofs

Conclusions

- Limitations
- Future work



Summary

► Π-Ware is...

Background

- Hardware Design
- Functional Hardware
- DTP

Research Question

Question
Method

DTP / Agda

- Big picture
- Agda

Π-Ware

- Syntax
- Semantics
- Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht

Current limitations

- ▶ Problem with proofs (definition of $\llbracket_ \rrbracket$)
- ▶ Proofs on (infinite) **Streams**
- ▶ Bla

Background

Hardware Design
Functional Hardware
DTP

Research Question

Question
Method

DTP / Agda

Big picture
Agda

Π -Ware

Syntax
Semantics
Proofs

Conclusions

Limitations
Future work



Future work

- ▶ Proof by reflection for finite cases

Background

- Hardware Design
- Functional Hardware
- DTP

Research Question

- Question
- Method

DTP / Agda

- Big picture
- Agda

Π -Ware

- Syntax
- Semantics
- Proofs

Conclusions

- Limitations
- Future work**



Thank you!

Questions?



References I

 Bjesse, P., Claessen, K., Sheeran, M., and Singh, S. (1998).

Lava: hardware design in Haskell.

SIGPLAN Not., 34(1):174–184.

 Sheeran, M. (1984).

MuFP, a language for VLSI design.

In Proceedings of the 1984 ACM Symposium on LISP and Functional Programming, LFP '84, pages 104–112, New York, NY, USA. ACM.

 Wadler, P. (2014).

Propositions as types.

Unpublished note, <http://homepages.inf.ed.ac.uk/wadler/papers/propositions-as-types/propositions-as-types.pdf>.

Background

Hardware Design

Functional Hardware

DTP

Research

Question

Question

Method

DTP / Agda

Big picture

Agda

Π -Ware

Syntax

Semantics

Proofs

Conclusions

Limitations

Future work



Universiteit Utrecht