

CHAPTER 1

PROGRAM EXECUTION TIME

$$\text{EXECUTION TIME} = \text{ISA} \times \text{CPI} \times \frac{\text{CLOCK PERIOD}}{\text{PERIOD}}$$

$$\text{CLOCK PERIOD} = \frac{1}{\text{FREQUENCY}}$$

• THROUGHPUT: WORK DONE PER UNIT TIME

CHAPTER 2

R-FORMAT

OPCODE 11-bits	Rm 5	SHAMT 6	Rn 5	Rd 5
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- OPCODE - OPERATION CODE
- Rm - SECOND REGISTER SOURCE
- Shamt - SHIFT
- Rn - FIRST REGISTER SOURCE
- Rd - REGISTER DESIGNATION

B-TYPE

• B 1000 // GOTO LOCATION 10000_{TEN}

5	10000 _{TEN}
6	26

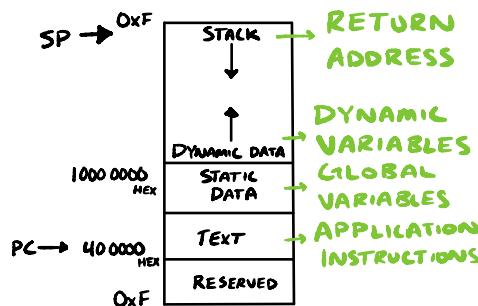
CB-TYPE

• CBNZ X19, EXIT // EXIT IF X19 != 0

181	EXIT	19
8	19	5

• Address = PC + OFFSET (2 Bits)

MEMORY ALLOCATION



CHAPTER 3

IEEE 754 FLOATING-POINT STANDARD

SINGLE PRECISION = 127

DOUBLE = 1023 (EXPONENT - BIAS)

$$(-1)^S \times (1 + \text{FRACTION}) \times 2^E$$

Floating-Point Addition

Now consider a 4-digit binary example

$$1.000_2 \times 2^{-1} + 1.110_2 \times 2^{-2} (0.5 + 0.4375)$$

1. Align binary points
Shift number with smaller exponent
 $1.000_2 \times 2^{-1} + 0.111_2 \times 2^{-1}$

2. Add significands

$$1.000_2 \times 2^{-1} + 0.111_2 \times 2^{-1} = 0.001_2 \times 2^{-1}$$

3. Normalize result & check for over/underflow
 $1.000_2 \times 2^{-4}$, with no over/underflow

4. Round and renormalize if necessary

$$1.000_2 \times 2^{-4} (\text{no change}) = 0.0625$$

DENORMAL NUMBERS

$$\text{EXponent} = 000...0$$

$$X = (-1)^S \times (0 + \text{FRACTION}) \times 2^{-E}$$

$$\text{DENORMAL FRACTION}$$

$$X = -1^S \times (0 + 0) \times 2^{-E}$$

SPEEDUP

$$\text{A ON B: } \frac{\text{TIME B}}{\text{TIME A}} = \frac{\text{PERFORMANCE A}}{\text{PERFORMANCE B}}$$

$$\text{PERFORMANCE} = \frac{1}{\text{CPU TIME}}$$

$$\text{CONVERSION} \\ 1 \text{ GHz} = \frac{1}{1 \times 10^9} \text{ SEC}$$

I-FORMAT

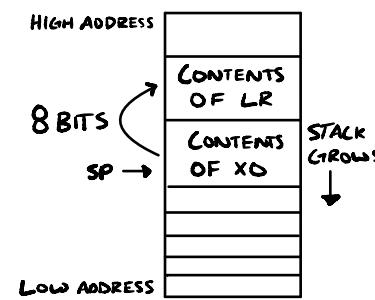
OPCODE 10	IMMEDIATE 12	Rn 5	Rd 5
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- OPCODE - OPERATION CODE
- Rn - SOURCE REGISTER
- Rd - REGISTER DESIGNATION
- IMMEDIATE - VALUE

LEGB8 REGISTERS

- X0-X7: PROCEDURE ARGUMENTS / RESULTS
- X8: INDIRECT RESULT LOCATION REGISTER.
- X9-X15: TEMPORARIES
- X16-X17: MAY BE USED BY LINKER AS A SCRATCH REGISTER, OTHER TIMES AS TEMPORARY REGISTER.
- X18: PLATFORM REG FOR PLATFORM IND. CODE. OTHERWISE A TEMP REG.
- X19-X27: SAVED
- X28(SP): STACK POINTER
- X29(FP): FRAME POINTER
- X30(LR): LINK REG (RETURN ADDRESS)
- X2R(REG 31): CONSTANT VALUE 0

DATA ON STACK



IEEE SINGLE PRECISION AND DOUBLE PRECISION FORMATS

S	EXPOENT	FRACTION
1	8	23

S	EXPOENT	FRACTION
1	11	52

Floating-Point Multiplication

Consider a 4-digit decimal example

$$1.110 \times 10^0 \times 9.200 \times 10^5$$

1. Add exponents

• For biased exponents, subtract bias from sum

$$\text{New exponent} = 10 + -5 = 5$$

2. Multiply significands

$$1.110 \times 9.200 = 10.212 \Rightarrow 10.212 \times 10^5$$

3. Normalize result & check for over/underflow

$$1.0212 \times 10^6$$

4. Round and renormalize if necessary

$$1.021 \times 10^6$$

5. Determine sign of result from signs of operands

$$+1.021 \times 10^6$$

INFINITY AND NaNs

- EXP = 111...1, FRACTION = 000...0
= ∞
- EXP = 111...1, FRACTION ≠ 000...0
= NOT A NUMBER

AMDAHL'S LAW

$$\text{SPEEDUP} = \frac{1}{(1-f) + \frac{f}{P}}$$

f = CAN IMPROVE
 $(1-f)$ = CAN NOT IMPROVE

D-FORMAT

OPCODE 11-bits	ADDRESS 9	Op2 2	Rn 5	Rd 5
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LOAD / STORE INSTRUCTIONS

- Rn: BASE REGISTER
- ADDRESS: CONSTANT OFFSET FROM CONTENTS OF BASE REGISTER
- Rt: DESIGNATION (LOAD) OR SOURCE (STORE) REGISTER #

CONDITIONAL OPERATIONS

- B.EQ
- B.NE
- B.LT (LESS THAN)
- B.LE (LESS THAN OR EQUAL)
- B.GT (GREATER THAN)
- B.GE (GREATER THAN, EQUAL TO)
- B.HS (UNSIGNED)

OTHER OPERATIONS

- LDUR X1, [X2, 40] • STUR X1, [X2, 40]
X1 = MEMORY[X2, 40] MEMORY[X2+40] = X1
- LSL BY i BITS MULTIPLIES BY 2^i
- LSR BY i BITS DIVIDES BY 2^i

NON LEAF PROCEDURE

- C CODE:


```
int fact (int n){  
    if (n < 1) return 1;  
    else return n * fact(n - 1);  
}
```
- ARGUMENT IN X0
- RESULT IN X1
- LEGV8 CODE:


```
fact:  
SUBI SP,SP,#16  
STUR LR,[SP,#8]  
STUR X0,[SP,#0]  
SUBIS XZR,X0,#1  
B.GE L1  
ADDI X1,XZR,#1  
ADDI SP,SP,#16  
BR LR  
L1: SUBI X0,X0,#1  
BL fact  
LDUR X0,[SP,#0]  
LDUR LR,[SP,#8]  
ADDI SP,SP,#16  
MUL X1,X0,X1  
BR LR
```

Save return address and n on stack
compare n and 1
if n >= 1, go to L1
Else, set return value to 1
Pop stack, don't bother restoring values
Return
n = n - 1
call fact(n-1)
Restore caller's n
Restore caller's return address
Pop stack
return n * fact(n-1)
return

MULTIPLICATION HARDWARE

