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-> Upper 4 Layous of TCP models are-

1. Application Layor.

PDV -> Dola

Protocols > SMTP, FTP, HTTP, DNS, SNMP, TELNET.

Functionality: It directly interacts withe the user. It

translates, encryptes and compresses data it necessary.

it enables data exchange over networks. Handles

tasks like, tile transfer, authentication, amail delivory

etc.

2. Transpord Layor:

PDU -> Segment, Dodagram.

Azotocols -> SCTP, TCP, VDP.

Functionality: Transpord layer ensures. rediable data fransfer between hosts. It can be both connectionless or connection orciented. It also ensures reliable delivery or bast delivery. Other functionalities at this layer are, flow condition, connection condition and france condition.

3. Interned Layer PDU > packet. (trolocal -> IP

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Fundionality: The	internel Ja	yer ensures	the address	exim
and satisfact	data packe	fg. By add	ma +1	,
il engures thre	connect	restinoutour.	1 Ticapara	V . C D
the transport lo	yer, for ea	asy and sec	wee tran	STRUMION
it divides the	data streeam	intosegmen	ds NISO	0/085
the reassembly.		,		

4. Host to Network Layer (Hayrical & Data Link):

PDU-> Arame

Protocols > poes not detine any specitic protocol. Supports all standwed prestocols.

Fundionality: It deals with frammission of brames over the Local network. Other tasks it handles are physical addressing, How, errore and access condition. It encupsulates IP packets into trames. Handles errors and reframmission.

-> Hore a computer is sending a trametotrom another computer using a bus topology LAN. Frame is the PDU of Data Link Layer. In this step a header is added to the kreame to define the sender/reciver which is physical addressing. Now, herce comes what happens when their physical destination addressing is corrected in Bub to pology LAN. Data Link Layer has a functionality freezon condition which is basically adding recliability to the physical layer. It is done by adding a trailer to the end of the treamer in order to delect and restreament damaged treames. Now, in but topology is the physical address is damaged the irdendend reciver will not get the fram. but if will still be broadcasted to all devices of the buts. No device will receive the treame as it don't matches with their address because of corraption and . Bus topology has no direct or builtin mechanism to understand uttomolity the senden that the mess frame was not delive

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In oreder to understand the ererore bus topology has to depend on ererore control mechanism which data Link Layer provides as discussed hore earlier. His Like in a speitic time it the sender doesn't receive the acknowledgement from the receiver it will understand that the attentions and unill rectransmit again.

Pros: 1. Top ensures reclability that all data sent trom
the sendor is received by the Hs receiver.

2. By including checksums and sequence numbers
it handles errore detection and handling.

3. It ensures that the order sender sent the message receiver receives exactly same order.

-> cons;

1. TOP is a slower protocol due to all its mechanisms like redramming data delivery acknowledgment it usages higher bandwith.

Now, which one will be best for Tom and Jordy depends on their requirement. It they want to choose recliability they need to use TCP. But it they want to that delivery they need to choose UDP.

-> According to the Question in a crucial level It game me and my troined have to work togethor and decides to save the game at regular intervals so that none of us lose the g progress in the game. Even it one of us lose can condinue krown the last saved point. This proceedure Tresembles the "session" Layer of OSI model. To be morre specific, session Layer have specific responsibilities Like-dialog conditrol and Syncroni-- zation. In the dialog section two system enloss indo a dialog. Ihrough a connection of either half-duplex on tull-duplex mode exactly like me and my trains need to work together. Again in the synctronization step the system adds check points to the stream of data to secure the transmission process. It works like this assection of data will be send to reciver but it a section greats lost it can simply wheck the checkpoint forom which point the data

is lost. After identifying the point sender can again send that speckics section of data. So, no need to send the entire data again. This is similar in case of me and my trained about saving at certain point so that it one loses can start from the previous saved point no need to start over again. So, the explained procedure resembles the Session Layer of OSI Model.



(a) -> Application Layer.

(B) -> Transport Layer and Data Link Layer.

Ø→ Physical Layor.