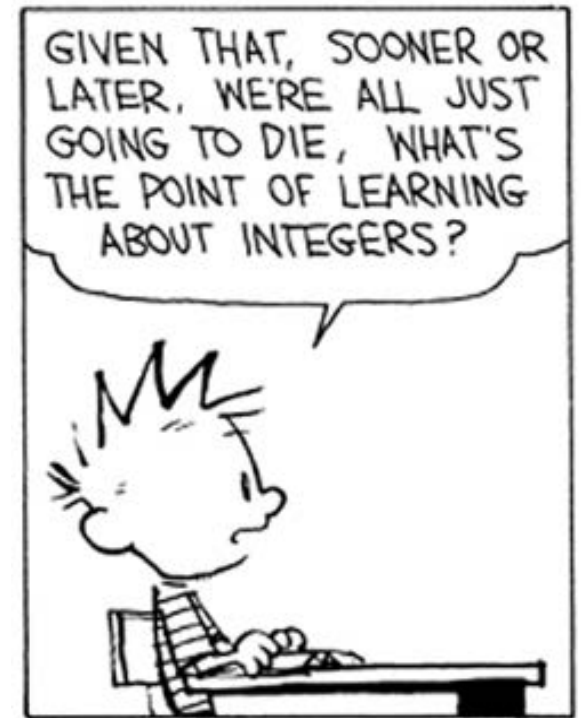


Advanced Mixed Integer Programming Formulation Techniques

MIP Representability

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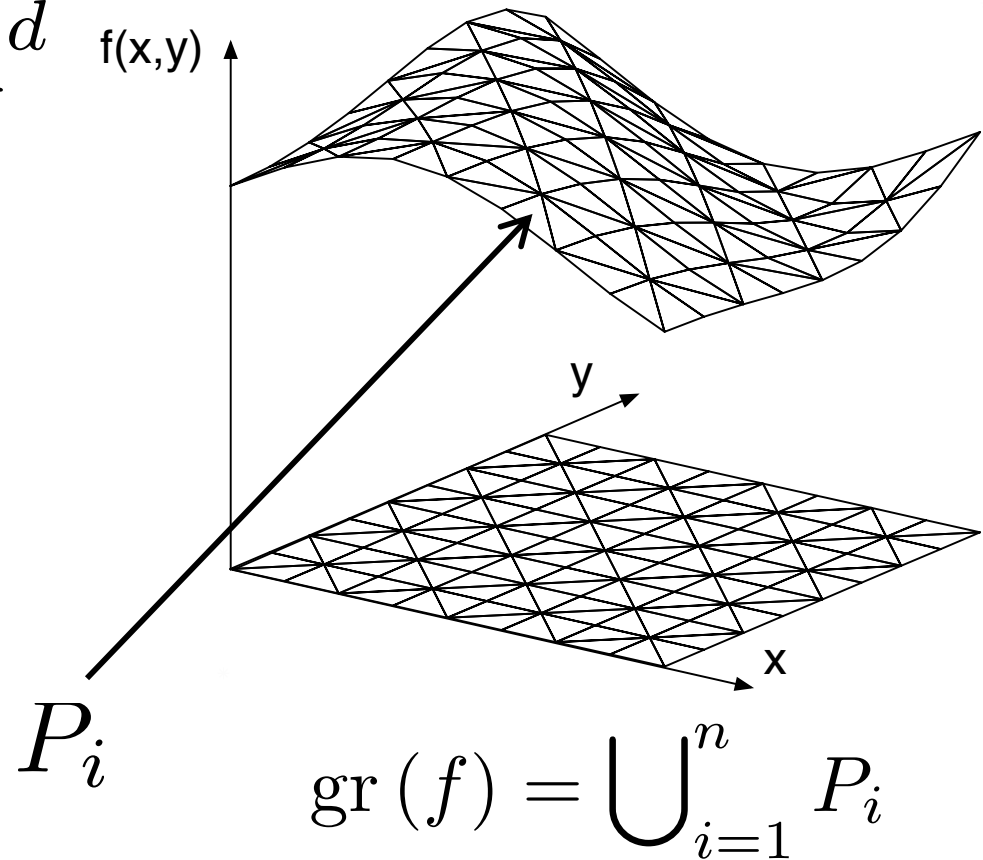
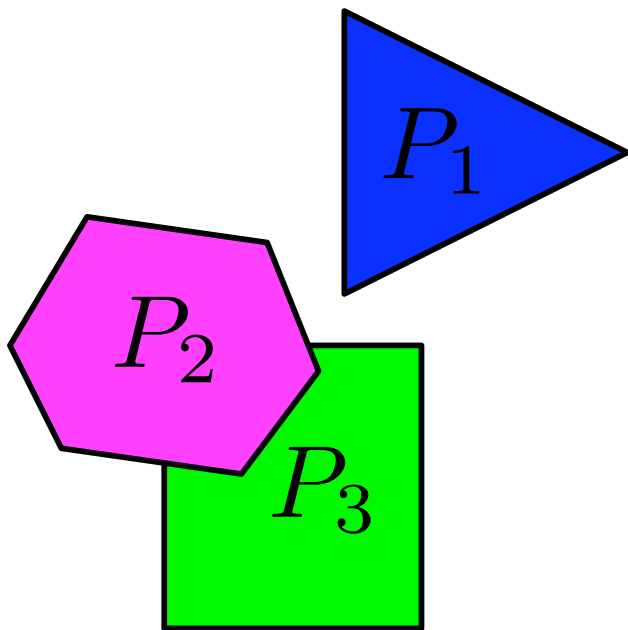
Spring School of the International Symposium on
Combinatorial Optimization (ISCO 2018)

Marrakesh, Morocco, April, 2018

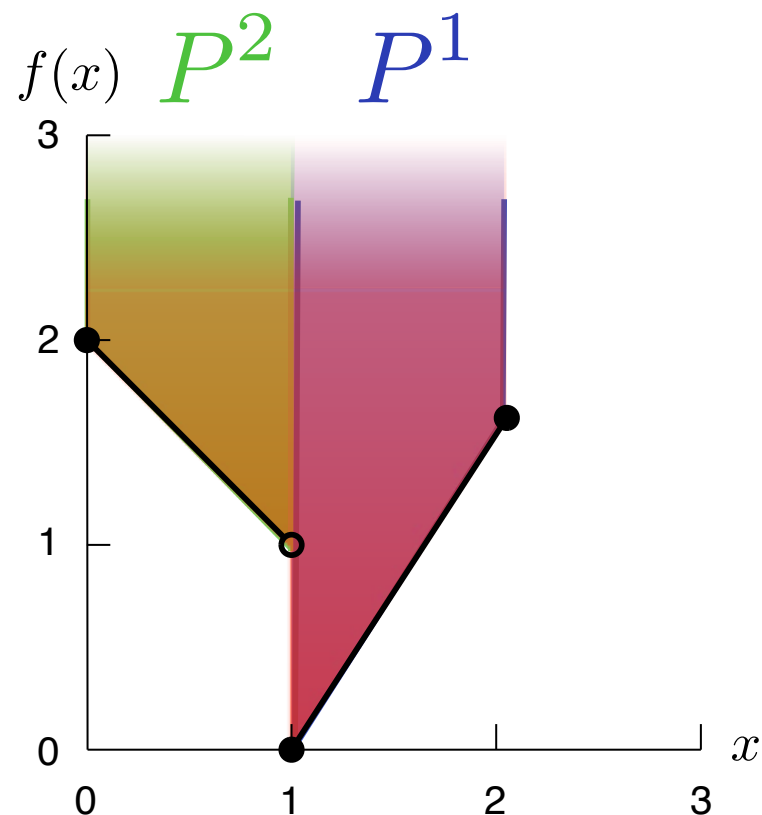
Mixed 0-1 Formulations

- Modeling Finite Alternatives = Unions of Polyhedra
 - Bounded or unbounded polyhedra, but bounded for now

$$x \in \bigcup_{i=1}^n P_i \subseteq \mathbb{R}^d$$



What About Discontinuous Piecewise Linear



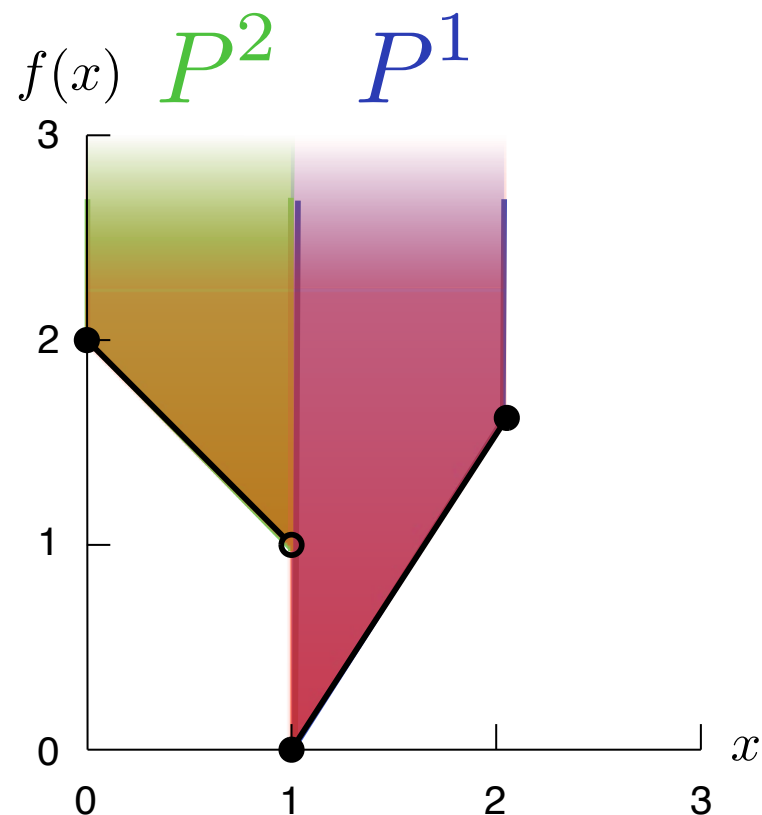
$$f(x) = \begin{cases} -x + 2 & x \in [0, 1) \\ x - 1 & x \in [1, 2] \end{cases}$$

$$\text{epi}(f) = \{(x, z) : f(x) \leq z\}$$

$$\text{epi}(f) = P^1 \cup P^2$$

- Works for piecewise linear functions:
 - Bounded domain
 - Lower semicontinuous = closed epigraph
 - Epigraph is union of **unbounded polyhedra**

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$$P_{\infty}^1 = P_{\infty}^2 = \{(x, z) \in \mathbb{R}^2 : x = 0, \quad z \geq 0\}$$

- Works for piecewise linear functions:
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 - Lower semicontinuous = closed epigraph
 - Epigraph is union of unbounded polyhedra

Extended Formulation for Unions of Polyhedra

- $P^i := \{x \in \mathbb{R}^n : A^i x \leq b^i\}$ rational polyhedra and

$$S = \bigcup_{i=1}^k P^i$$

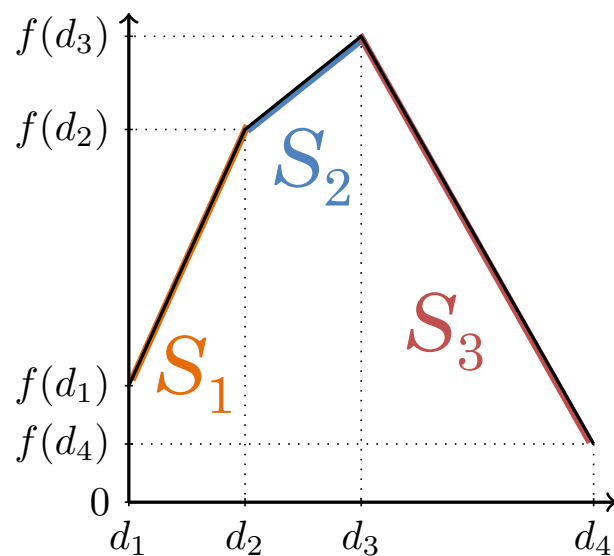
- If $P_\infty^i = P_\infty^j \quad \forall i, j \in [k]$
- then an ideal formulation for S is

$$\begin{aligned} A^i x^i &\leq b^i y_i \\ \sum_{i=1}^k x^i &= x \\ x^i &\in \mathbb{R}^n \quad \forall i \in [n] \\ \sum_{i=1}^k y_i &= 1 \\ y &\in \mathbb{Z}^n \end{aligned}$$

Extended Formulation for PWL Functions

$$S = \text{gr}(f) = \bigcup_{i=1}^k \left\{ (x, z) \in \mathbb{R}^2 : \begin{array}{l} d_i \leq x \leq d_{i+1} \\ m_i x + c_i = z \end{array} \right\}$$

MC Formulation:



$$d_i y_i \leq x^i \leq d_{i+1} y_i \quad \forall i \in [k]$$

$$m_i x^i + c_i y_i = z^i \quad \forall i \in [k]$$

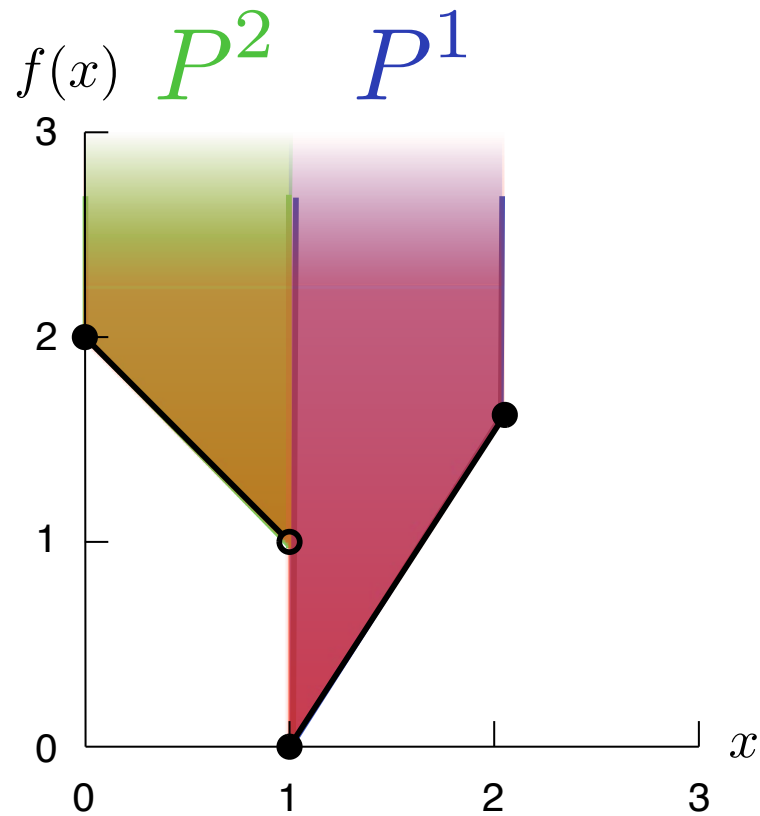
$$\sum_{i=1}^k x^i = x$$

$$\sum_{i=1}^k z^i = z$$

$$\sum_{i=1}^k y_i = 1$$

$$y \in \{0, 1\}^k$$

Example: Discontinuous Piecewise Linear



$$\text{epi}(f) = \{(x, z) : f(x) \leq z\}$$

$$\text{epi}(f) = P^1 \cup P^2$$

$$P_\infty^1 = P_\infty^2 = \{(x, z) \in \mathbb{R}^2 : x = 0, z \geq 0\}$$

$$f(x) = \begin{cases} -x + 2 & x \in [0, 1) \\ x - 1 & x \in [1, 2] \end{cases}$$

$$d_i y_i \leq x^i \leq d_{i+1} y_i \quad \forall i \in [k]$$

$$m_i x^i + c_i y_i \leq z^i \quad \forall i \in [k]$$

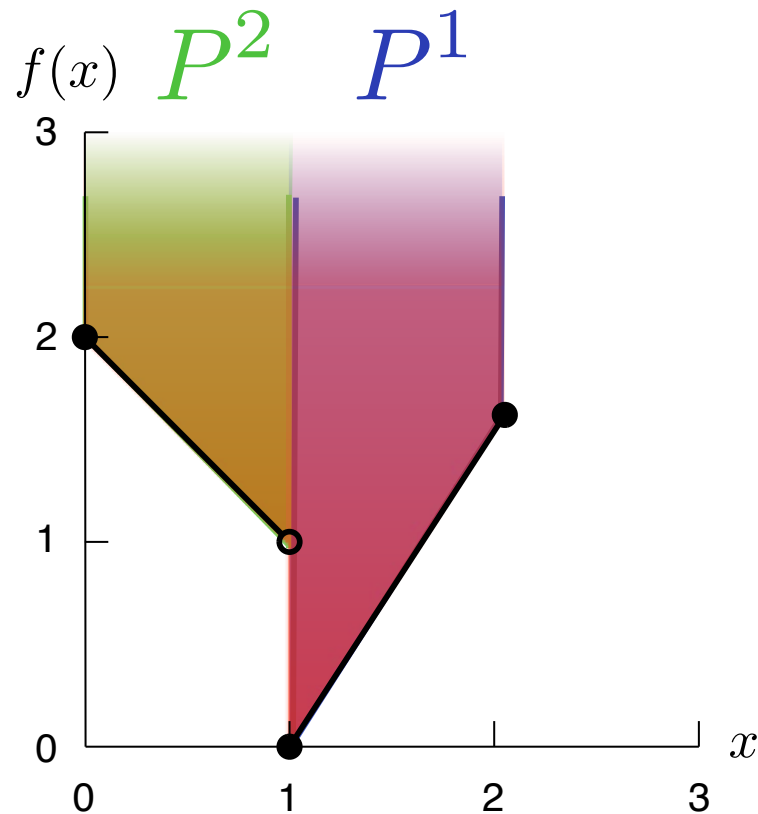
$$\sum_{i=1}^k x^i = x$$

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Example: Discontinuous Piecewise Linear



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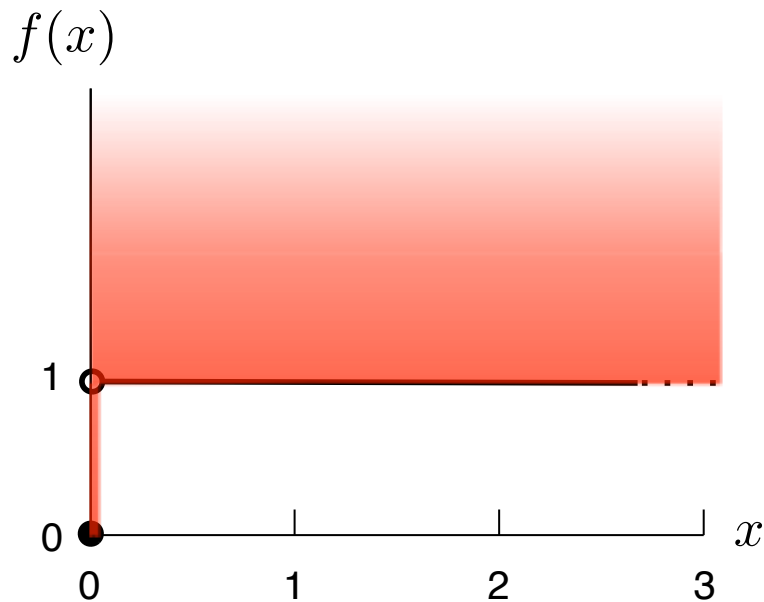
$$\text{epi}(f) = \{(x, z) : f(x) \leq z\}$$

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$$P_{\infty}^1 = P_{\infty}^2 = \{(x, z) \in \mathbb{R}^2 : x = 0, z \geq 0\}$$

- Works for piecewise linear functions:
 - Bounded domain
 - Lower semicontinuous = closed epigraph

What if Recession Cones are Different?



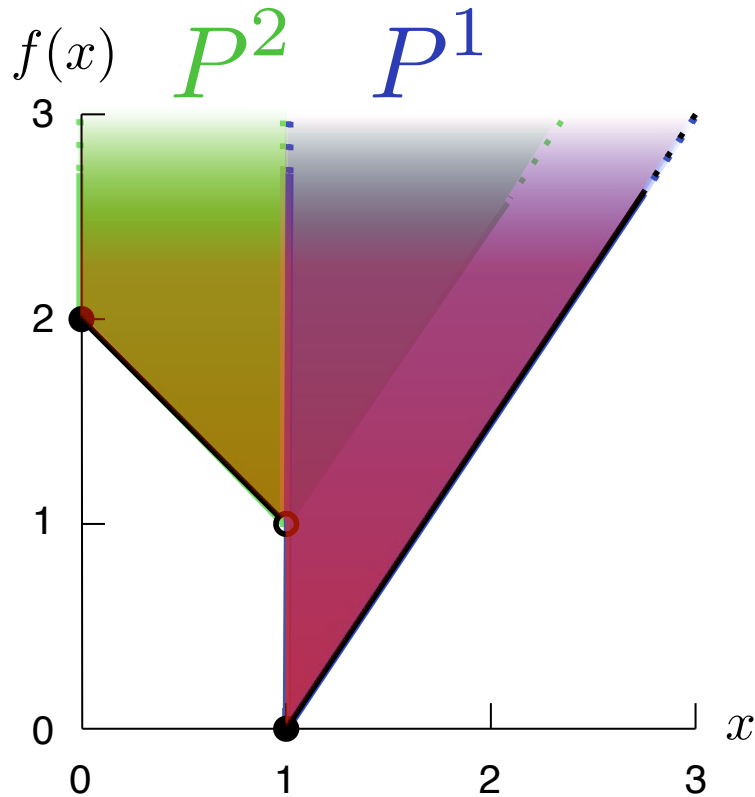
$$f : [0, \infty) \rightarrow \mathbb{R}$$

$$f(x) = \begin{cases} 0 & x = 0 \\ 1 & x > 0 \end{cases}$$

$$\text{epi}(f) = \{(x, z) : f(x) \leq z\}$$

- No linear MIP formulation:
 - 0-1 or general integer
 - More later

Bounded MIP Representability May be Hidden



$$f(x) = \begin{cases} -x + 2 & x \in [0, 1) \\ x - 1 & x \in [1, \infty) \end{cases}$$

$$\text{epi}(f) = \{(x, z) : f(x) \leq z\}$$

$$\text{epi}(f) = P^1 \cup P^2$$

$$P^2_\infty = \{(x, z) \in \mathbb{R}^2 : x = 0, z \geq 0\}$$

$$\nparallel \quad \{(x, z) \in \mathbb{R}^2 : x \geq 0, z \geq x\}$$

$$P^1_\infty = \{(x, z) \in \mathbb{R}^2 : x \geq 0, z \geq x\}$$

$$P^2 \rightarrow P^2 \cup P^1_\infty$$

- Obtaining a MIP formulation may require redundancy

Finite, Binary or 0-1 Linear MIP Representability

- Let
 - $S \subseteq \mathbb{R}^n$,
 - $n_1 + n_2 = n$, $p_1 + p_2 = p$, $A \in \mathbb{Q}^{m \times n}$, $D \in \mathbb{Q}^{m \times p}$, $b \in \mathbb{Q}^m$
 - $P := \{(x, w) \in \mathbb{R}^n \times \mathbb{R}^p : Ax + Dw \leq b\}$
 - $P_I := P \cap (\mathbb{R}^{n_1} \times \mathbb{Z}^{n_2} \times \mathbb{R}^{p_1} \times \mathbb{Z}^{p_2})$
 - P_I is a MIP formulation of S iff $S = \text{Proj}_x(P_I)$
- If $\exists M > 0$ s.t. $P \subseteq \mathbb{R}^{n_1} \times [-M, M]^{n_2} \times \mathbb{R}^{p_1} \times [-M, M]^{p_2}$
 - Then S is a finite union of rational polyhedra with the same recession cone
- S is finite/binary linear MIP representable if it has a linear MIP formulation
- Sometimes denoted “Bounded” MIP representable

Mixed-Integer Formulations: General Integer

Beyond Unions of Polyhedra

- $P^i := \{x \in \mathbb{R}^n : A^i x \leq b^i\}$ rational polyhedra and
 $S = \bigcup_{i=1}^k P^i \cap (\mathbb{R}^{n_1} \times \mathbb{Z}^{n_2})$
- If $P_\infty^i = P_\infty^j \quad \forall i, j \in [k]$
- then a formulation for S is

$$\begin{array}{ll}
 A^i x^i \leq b^i y_i & \\
 \sum_{i=1}^k x^i = x & \sum_{i=1}^k y_i = 1 \\
 x^i \in \mathbb{R}^n \quad \forall i \in [n] & y \in \mathbb{Z}^n \\
 & x \in \mathbb{R}^{n_1} \times \mathbb{Z}^{n_2}
 \end{array}$$

- Formulation is ideal if

$$\text{ext}(P^i) \subseteq \mathbb{R}^{n_1} \times \mathbb{Z}^{n_2}$$

General Linear MIP Representability

- $P^i := \{x \in \mathbb{R}^n : A^i x \leq b^i\}$ rational polyhedra and

$$S = \bigcup_{i=1}^k P^i \cap (\mathbb{R}^{n_1} \times \mathbb{Z}^{n_2})$$

- If $P_\infty^i = P_\infty^j \quad \forall i, j \in [k]$
- What about?
 - $S = \{0\} \cup ([2, \infty) \cap \mathbb{Z})$
 $= \{0, 2, 3, 4, \dots\}$
 - MIP formulation:

$$x = 2y_1 + 3y_2, \quad y_1, y_2 \geq 0, \quad y \in \mathbb{Z}^2$$

(Linear) Mixed Integer Programming Formulation

- Let
 - $S \subseteq \mathbb{R}^n$,
 - $n_1 + n_2 = n$, $p_1 + p_2 = p$, $A \in \mathbb{Q}^{m \times n}$, $D \in \mathbb{Q}^{m \times p}$, $b \in \mathbb{Q}^m$
 - $P := \{(x, w) \in \mathbb{R}^n \times \mathbb{R}^p : Ax + Dw \leq b\}$
 - $P_I := P \cap (\mathbb{R}^{n_1} \times \mathbb{Z}^{n_2} \times \mathbb{R}^{p_1} \times \mathbb{Z}^{p_2})$
- P_I is a MIP formulation of S iff

$$S = \text{Proj}_x (P_I)$$

- S is linear MIP representable if it has a linear MIP formulation

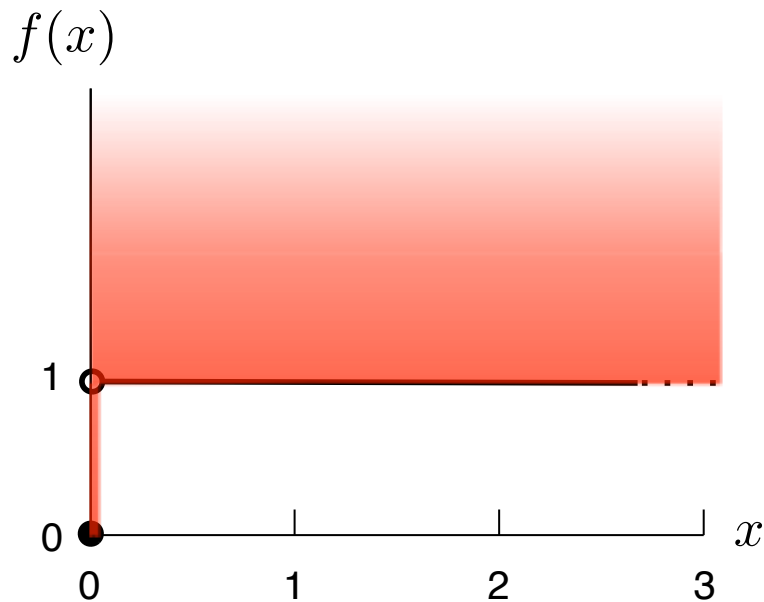
General Linear MIP Representability

- $S \subseteq \mathbb{R}^n$ is MIP representable if and only if there exist rational polyhedra $P_i \subseteq \mathbb{R}^n$ with the same recession cones and $\{r^j\}_{j=1}^L \subseteq \mathbb{Z}^n$ such that

$$- S = \bigcup_{i=1}^K P_i + \left\{ \sum_{j=1}^L \lambda_j r^j : \lambda \in \mathbb{Z}_+^L \right\}$$

- Can always assume the P_i are polytopes

What if Recession Cones are Different?



$$f : [0, \infty) \rightarrow \mathbb{R}$$

$$f(x) = \begin{cases} 0 & x = 0 \\ 1 & x > 0 \end{cases}$$

$$\text{epi}(f) = \{(x, z) : f(x) \leq z\}$$

$$\text{conv}(\text{epi}(f)) = \{(x, z) \in \mathbb{R}^2 : x \geq 0, z > 0\} \cup \{(0, 0)\}$$

Not MIP representable: bounded or general

- Useful lemma:
 - If S has a MIP formulation then S is closed and $\text{conv}(S)$ is a rational polyhedron

Rational Polyhedral Convex Hull is Not Sufficient

- $S = \left\{ x \in \mathbb{Q}^n : x_n = \max_{i \in [n-1]} x_i, \quad x_j \geq 0 \quad \forall j \in [n] \right\}$
 $= \bigcup_{i=1}^{n-1} \{x \in \mathbb{Q}^n : x_n = x_i, \quad x_n \geq x_j \geq 0 \quad \forall j \in [n]\}$
- $\text{conv}(S) = \left\{ x \in \mathbb{Q}^n : x_n \leq \sum_{i=1}^{n-1} x_i, \quad x_n \geq x_j \geq 0 \quad \forall j \in [n] \right\}$
- S does not have a MIP formulation

Rationality is Crucial

- $\text{conv}(\{x \in \mathbb{Z}^2 : x_2 - \sqrt{2} x_1 \leq 0, x_2 \geq 0\})$ is not even closed

(Convex) MICP Formulations and Representability

- A set $S \subseteq \mathbb{R}^n$ is MICP representable (MICPR) if it has an MICP formulation:
 - A closed convex set $M \subseteq \mathbb{R}^{n+p+d}$
 - auxiliary continuous variables $y \in \mathbb{R}^p$
 - auxiliary integer variables $z \in \mathbb{Z}^d$

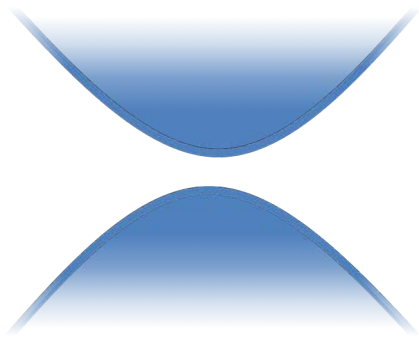
$$x \in S \quad \Leftrightarrow \quad \begin{array}{l} \exists (y, z) \in \mathbb{R}^p \times \mathbb{Z}^d \text{ s.t.} \\ (x, y, z) \in M \end{array}$$

or equivalently

$$S = \text{proj}_x \left(M \cap \left(\mathbb{R}^{n+p} \times \mathbb{Z}^d \right) \right)$$

What Sets are MICP Representable (MICPR) ?

Two sheet hyperbola?



$$\{x \in \mathbb{R}^2 : 1 + x_1^2 \leq x_2^2\}$$

Spherical shell?



$$\{x \in \mathbb{R}^2 : 1 \leq \|x\| \leq 2\}$$

- Discrete subsets of the real line or natural numbers:

– Dense discrete set? $\{\sqrt{2}x - \lfloor \sqrt{2}x \rfloor : x \in \mathbb{N}\} \subseteq [0, 1]$

– Set of prime numbers?



"God made the integers,
all else is the work of man"

- Leopold Kronecker

A Simple Lemma for non-MICP Representability

- Obstruction for MICP representability of S :

$$\text{infinite } R \subseteq S \quad \text{s.t.} \quad \frac{u+v}{2} \notin S \quad \forall u, v \in R, u \neq v$$

Proof: Assume for contradiction there exists M such that:

$$S = \text{proj}_x \left(M \cap \left(\mathbb{R}^{n+p} \times \mathbb{Z}^d \right) \right)$$

$$\begin{array}{l} (u, y_u, z_u) \in M \\ (v, y_v, z_v) \in M \end{array} \Rightarrow \frac{z_u + z_v}{2} \notin \mathbb{Z}^d$$

$$z_u \equiv z_v \pmod{2} \text{ component-wise} \Rightarrow \frac{z_u + z_v}{2} \in \mathbb{Z}^d$$

$$\text{component-wise parity classes} = 2^d < |R| = \infty \quad \Rightarrow \text{contradiction}$$

A Simple Lemma for non-MICP Representability

- Obstruction for MICP representability of S :

$$\text{infinite } R \subseteq S \quad \text{s.t.} \quad \frac{u+v}{2} \notin S \quad \forall u, v \in R, u \neq v$$

✗ **Spherical shell** $\{x \in \mathbb{R}^2 : 1 \leq \|x\| \leq 2\}$

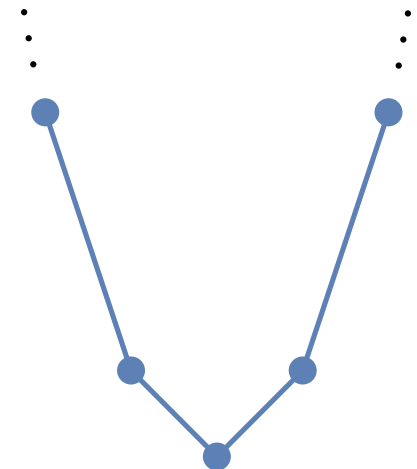


✗ **Set of prime numbers**

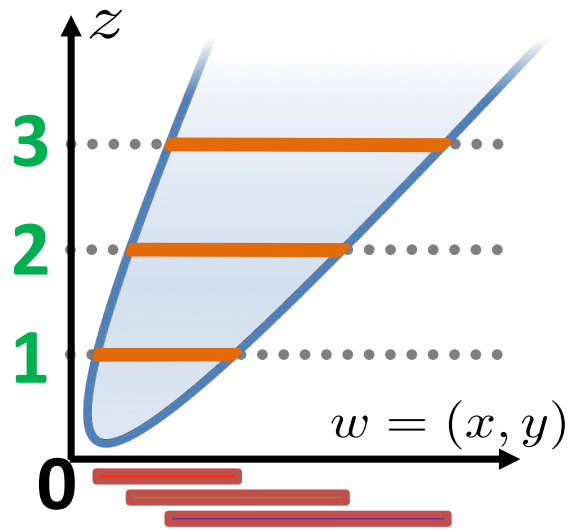
- Does have **non-convex** polynomial MIP

✗ Set of Matrices of rank at most k

✗ Piecewise linear interpolation of x^2 at all integers



MICPR = Convex Sets Indexed by Integers in Convex



$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x (B_z)$$

$$I = \text{proj}_z (M) \text{ convex}$$

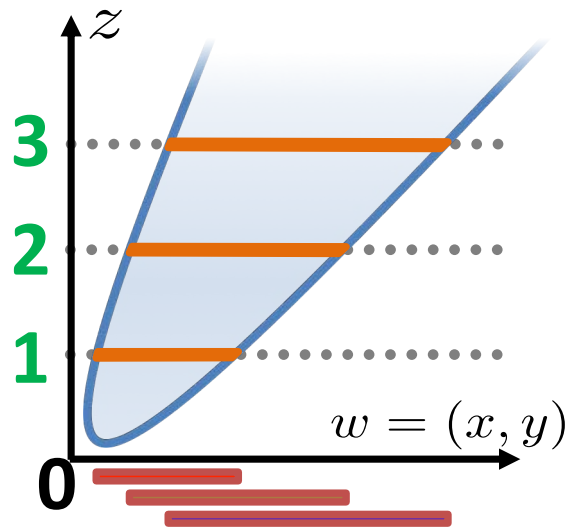
- For rational polyhedral M (Jeroslow and Lowe '84):

$$S = \bigcup_{i=1}^k P_i + \left\{ \sum_{i=1}^t \lambda_i r^i : \lambda \in \mathbb{Z}_+^t \right\}$$

1: Rational polyhedra with
the same recession cone

2: Finite # of shapes
+ periodic translations

MICPR = Convex Sets Indexed by Integers in Convex

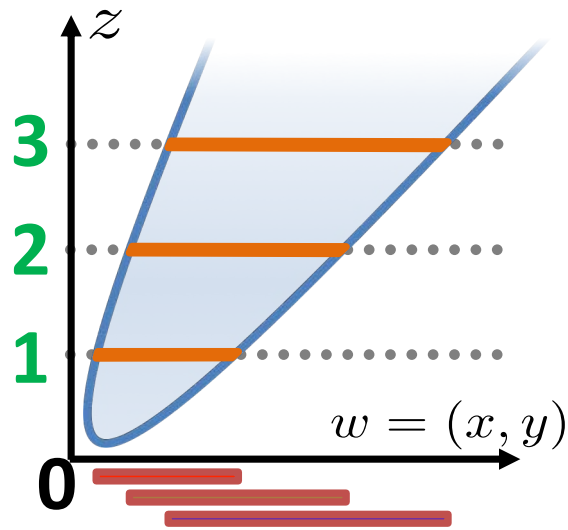


$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x (B_z)$$

$$I = \text{proj}_z (M) \text{ convex}$$

- Extensions $S = \bigcup_{i=1}^k P_i + \left\{ \sum_{i=1}^t \lambda_i r^i : \lambda \in \mathbb{Z}_+^t \right\}$
 - $M = \{x \in \mathbb{Z}^2 : x_1 \cdot x_2 \geq \alpha\} \Rightarrow P_i = \text{points}$ (Dey & Moran '13)
 - $M = \text{Rational Polyhedron} \cap \text{"Rational" Ellipsoidal Cylinder} \Rightarrow P_i = \text{Rational Ellipsoid} \cap \text{Polytope}$ (Del Pia & Poskin '16)
 - $M = \text{Compact Convex} + \text{Rational Polyhedron Cone} \Rightarrow P_i = \text{Compact Convex}$ (Lubin, Zadik & V. 17')

MICPR = Convex Sets Indexed by Integers in Convex



$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x (B_z)$$

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- For rational polyhedral M (Jeroslow and Lowe '84):

$$S = \bigcup_{i=1}^k P_i + \left\{ \sum_{i=1}^t \lambda_i r^i : \lambda \in \mathbb{Z}_+^t \right\}$$

1: Rational polyhedra with the same recession cone

2: Finite # of shapes + periodic translations

Extra from MICP 1: Non-Polyhedral Unions

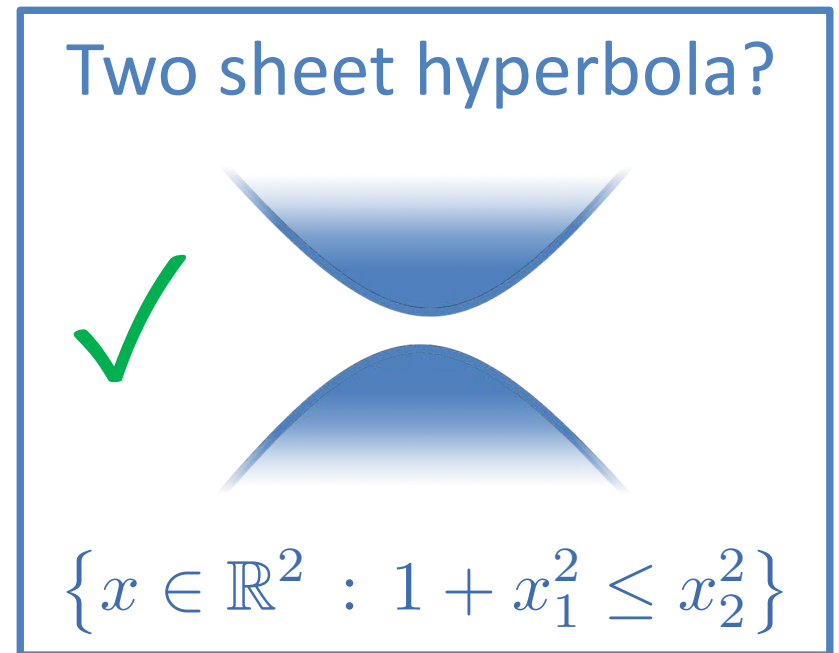
$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x (B_z)$$

1. Unions of **Non-Polyhedral sets**

Plus Projection:

2. Unions of **non-closed sets**
3. Unions of **convex sets with different recession cones**

$\left\{ \left\{ (x, t) : x \in S_i, \quad \|x\|_2^2 \leq t \right\} \right\}_{i=1}^k$ have the same recession cone



Extra from MICP 2: Non-Polyhedral I

- An **infinite set** S is **periodic** if and only if:

$$\exists r \in \mathbb{R}^n \quad \forall \lambda \in \mathbb{Z}_+, \quad x \in S \quad x + \lambda r \in S$$

- Non-periodic MICPR sets

- Dense discrete set $\left\{ \sqrt{2}x - \lfloor \sqrt{2}x \rfloor : x \in \mathbb{N} \right\} \subseteq [0, 1]$

$$\|(z_1, z_1)\|_2 \leq z_2 + 1, \quad \|(z_2, z_2)\|_2 \leq 2z_1, \quad x_1 = y_1 - z_2,$$

$$\|(z_1, z_1)\|_2 \leq y_1, \quad \|(y_1, y_1)\|_2 \leq 2z_1, \quad z \in \mathbb{Z}^2$$

- Set of naturals $\left\{ x \in \mathbb{N} : \sqrt{2}x - \lfloor \sqrt{2}x \rfloor \notin (\varepsilon, 1 - \sqrt{2}\varepsilon) \right\}$

$$\|(x_1, x_1)\|_2 \leq x_2 + \varepsilon,$$

$$\|(x_2, x_2)\|_2 \leq 2x_1 + 2\varepsilon, \quad x \in \mathbb{Z}_+^2$$

A Definition Rational MICPR (R-MICPR)

$$S = \bigcup_{z \in I \cap \mathbb{Z}^d} \text{proj}_x (B_z) \quad S = \text{proj}_x (M \cap (\mathbb{R}^{n+p} \times \mathbb{Z}^d))$$

$$I = \text{proj}_z (M)$$

- Any rational affine mapping of index set I :
 - Is bounded, or
 - Has an integer (rational) recession direction
- Irrational directions can hide!
 - R-MICPR $\not\Rightarrow$ $\text{span}(\text{rec}(I))$ and/or $\text{aff}(I)$ = rational space

$$(z_1 + \sqrt{2}z_2)^2 \leq z_3$$

$$\text{span}(\text{rec}(I)) = \text{span}(\{\mathbf{e}_3\})$$

$$(z_2 - \sqrt{2}z_1)^2 \leq 1$$

$$\text{rec}(\text{proj}_{z_1, z_2}(I)) = \text{span}(\{(1, \sqrt{2})\})$$

Properties of Rational MICPR (R-MICPR)

- For **compact** S :
 - **Finite unions of compact** convex sets
- For S **infinite unions of “uniformly bounded” closed convex sets** :
 - **Finite union of periodic**
 - **Dense discrete** and **non-periodic naturals** NOT R-MICPR
- Rational MICP Representability:
 - **Closed** under: **Finite Union, Cartesian Product and Minkowski sum**
 - **NOT Closed** under **intersection**.

R-MICPR: Periodicity for Natural #s

- An **infinite set** of naturals S is **periodic** if and only if:
 - $S = \bigcup_{i=1}^k \{s_i\} + \text{intcone}(\{r\})$
 - It is **rational MILP representable**
- A subset S of the naturals is **R-MICPR** if and only if:
 - It is the **union** of a **finite** and an **infinite periodic** set

R-MICPR does Not Imply Finite Shapes



- There exists increasing functions h such that:
 - $P_z \subseteq \mathbb{R}^2$ regular $h(z)$ -gon centered at $(z, 0)$
 - $P_z \cap P_{z'} = \emptyset, \quad z \neq z'$
 - $S = \bigcup_{z=1}^{\infty} P_z$ is R-MICPR and periodic
- Equal volume \Rightarrow Finite # of Shapes