

FIT9136: Algorithms and programming foundations in Python

Week 7: Binary Trees and BSTs



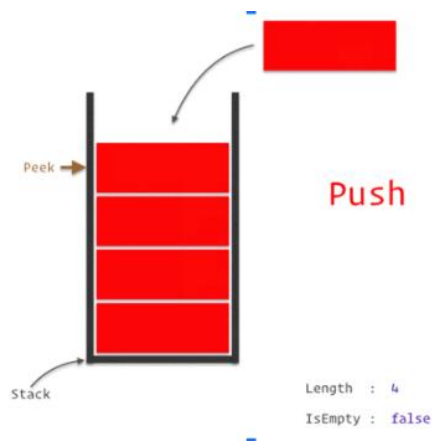
Agenda

- Review of Week 6
- Synopsis
- Learning Objectives
- Fundamental Class Update
 - Nodes and Links
- Advanced Data Structure:
 - Binary Trees
- Abstract Data Type:
 - Binary Search Trees

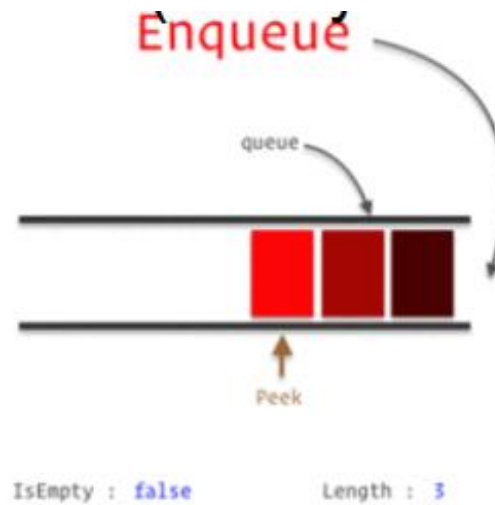
Review of Week 6

Review of Week 6

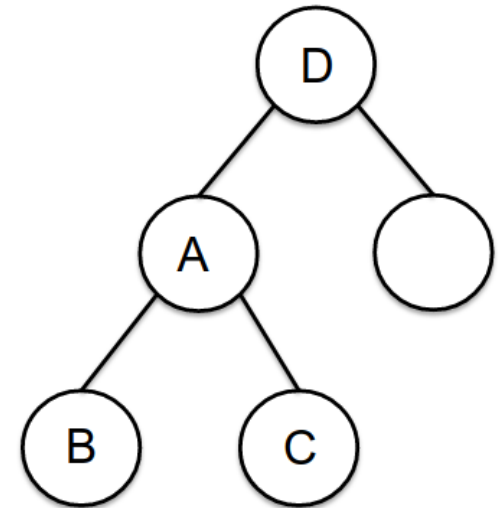
Stack



Queue



Heap



- Week 7 is aimed to provide you with:
 - Advanced Data Structure and ADT:
 - Binary Trees and Binary Search Trees

- Understand the concept of Binary Trees and Binary Search Tree

Fundamental Class Update: Nodes and Links

- Node:
 - Nodes for the data structure for today require two links rather than one.
 - There is a need for a left, and a right element, thus, the definition has changed:

```
class Node:
```

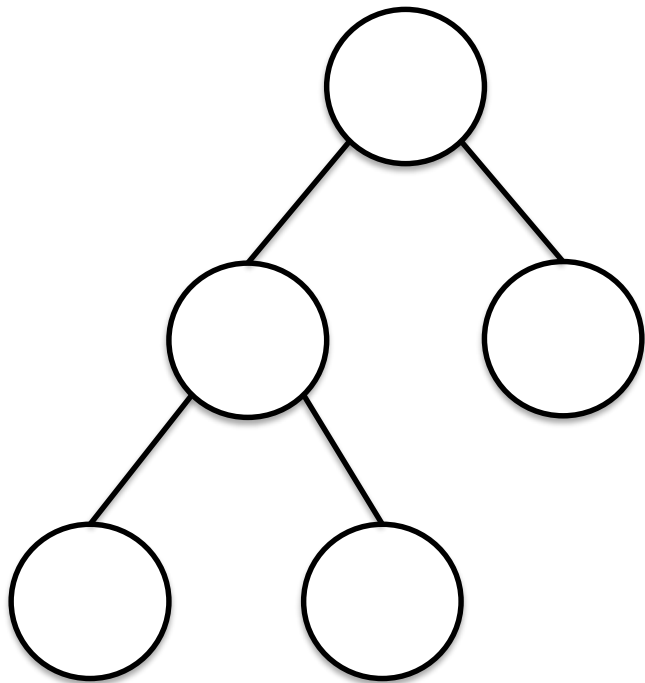
```
    def __init__(self, item, left = None, right = None):  
        self.item = item  
        self.left = left  
        self.right = right
```

- This now allows for Binary Trees to be constructed.
 - If we wanted to make Trinary trees there would need to be three links. (Tri meaning 3)

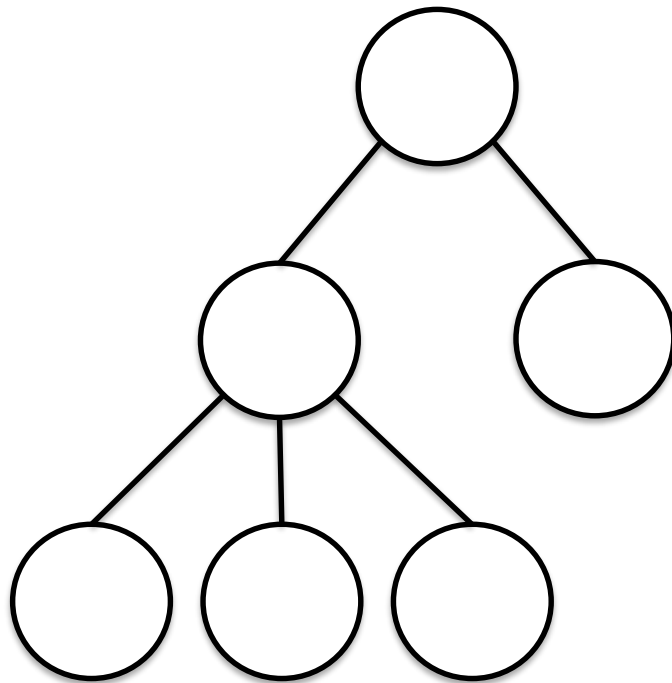
Advanced Data Structure: Binary Trees

Binary Trees

- A binary tree is a tree whose nodes have at most 2 children.



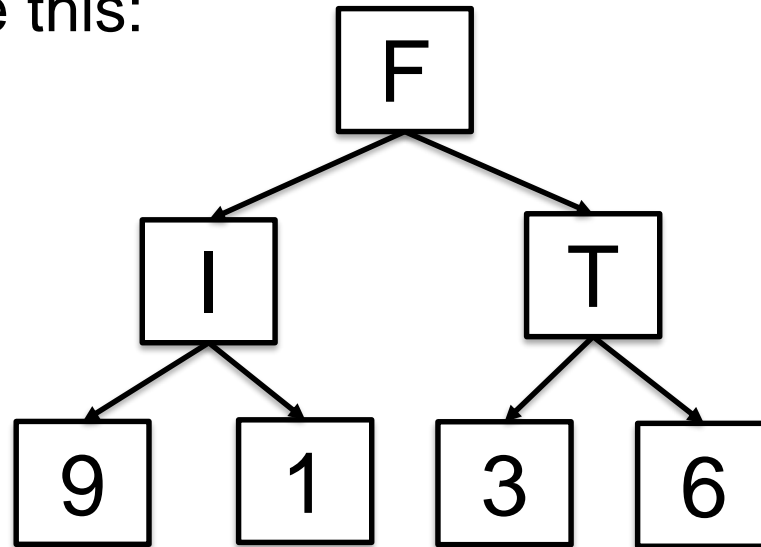
Binary Tree



Non-Binary Tree

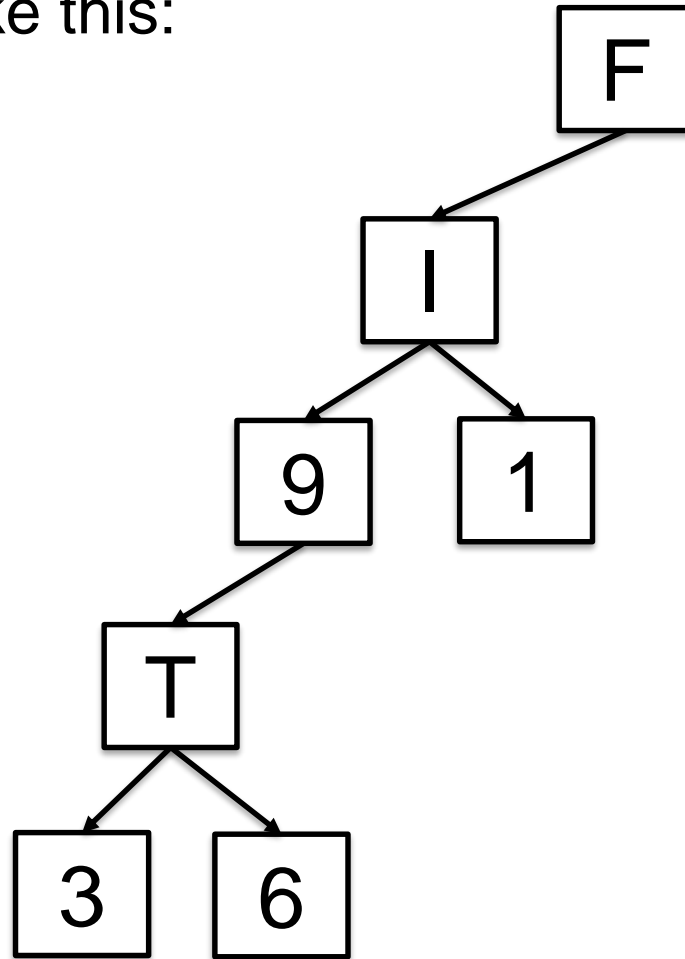
Binary Trees

- Using a Binary Tree structure: our unit's name could look like this:



Binary Trees

- Or like this:

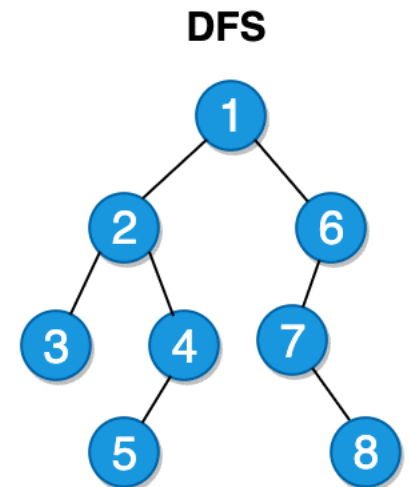
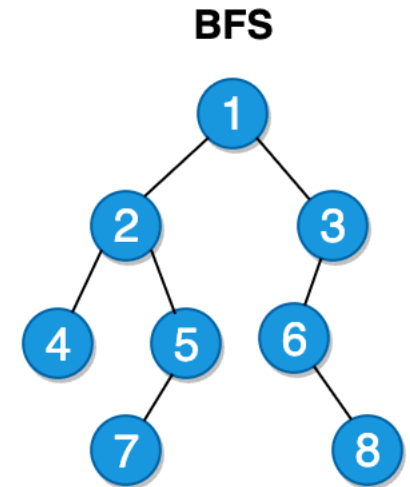


- Ultimately the structure isn't useful without rules in place to mandate where items are placed.
- Trees (not just binary trees) are useful to show chains of actions from a preceding state.
 - Not just actions chosen, but all possible choices.
 - The rules imparted on the tree gives the tree its usefulness.

Abstract Data Type: Binary Search Trees

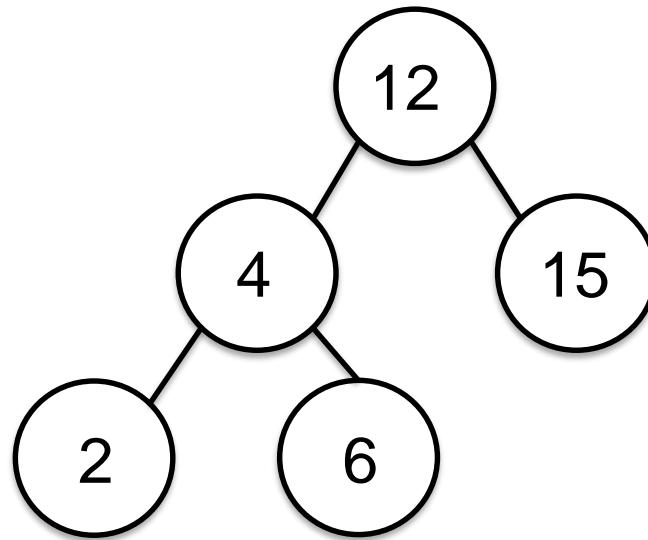
Tree Search

- Tree search refers to the process of finding one node in a tree data structure.
- Two tree search steps:
 - Breadth First Search (BFS)
 - Depth First Search (DFS)



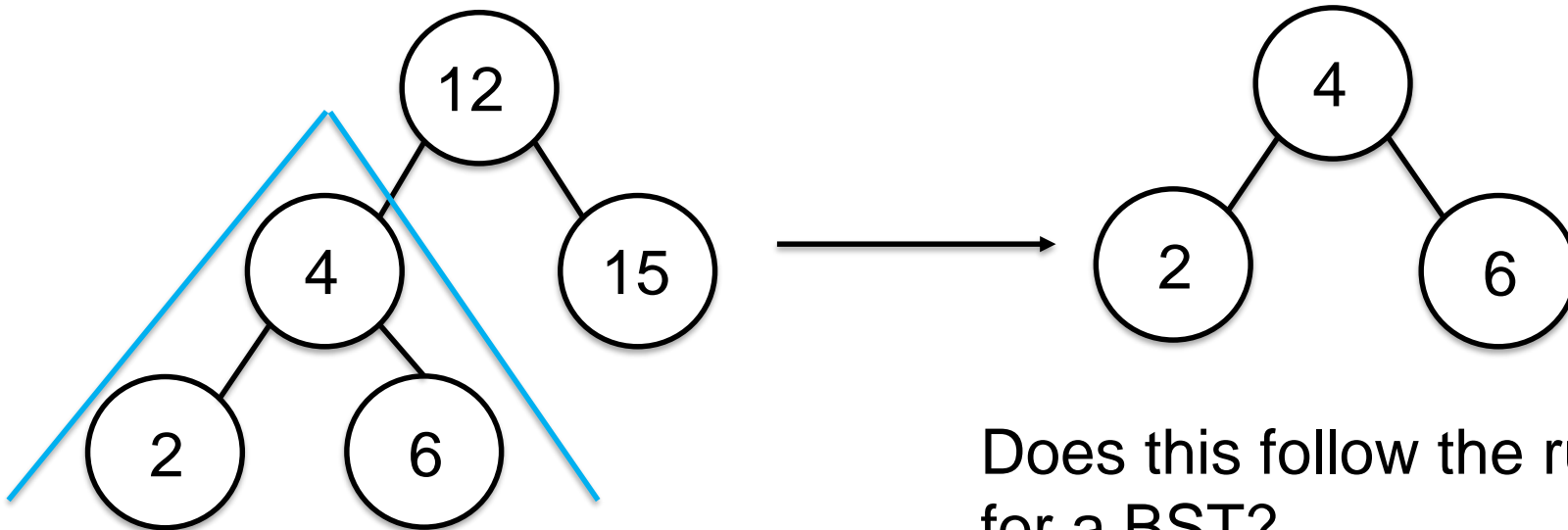
- **Binary Search Tree (BST)** is a node-based binary tree data structure which has the following properties:
 - The left subtree of a node contains only nodes with values smaller than the root.
 - The right subtree of a node contains only nodes with values larger than the root.
 - The left and right subtree each must also be a binary search tree
- Binary Search Trees exists to allow for fast searching for a dataset.

An Example BST



Interesting Fact

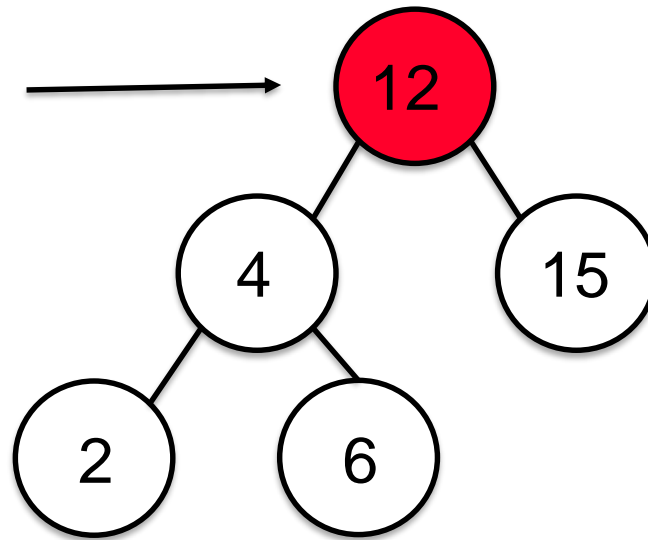
- The sub-trees of a BST are also a BST.



Does this follow the rules for a BST?

How do we find an element?

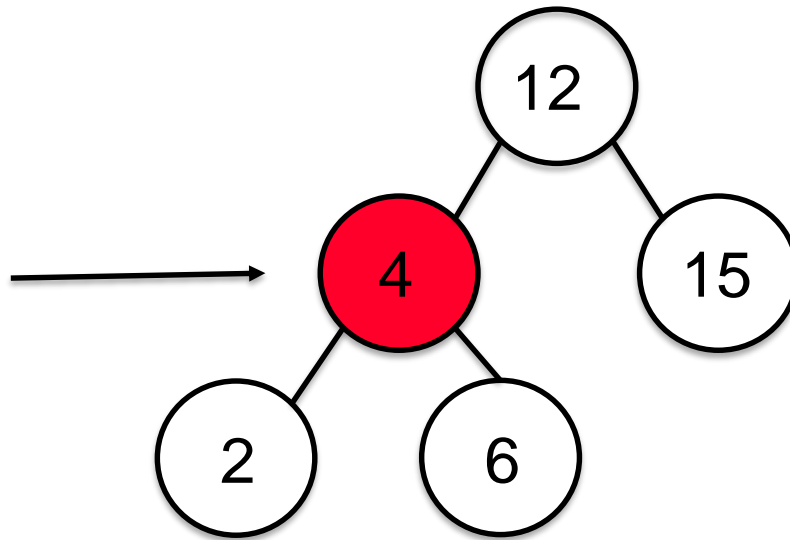
- Let's look for 6. To start searching, we start at the root:



- Is the root what we want? No.
- Is $6 > 12$? No.
 - 6 must exist to the left of 12.
- Go left.

How do we find an element?

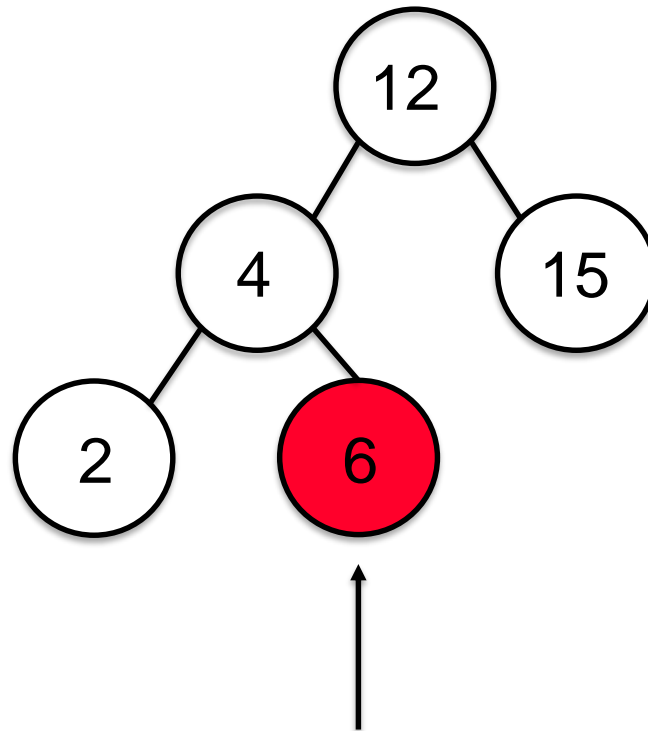
- Continue the quest for 6:



- Is the root what we want? No.
- Is $6 > 4$? Yes.
 - 6 must exist to the right of 4.
- Go right.

How do we find an element?

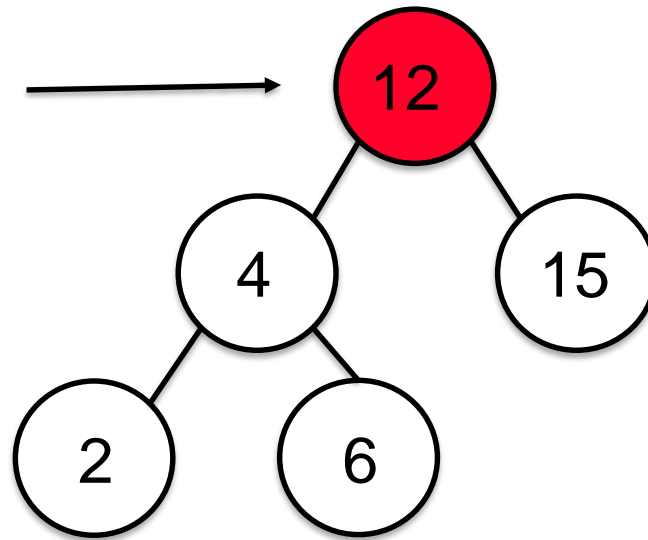
- Continue the quest for 6:



- Is the root what we want? Yes.
- Return True.

How do we find an element?

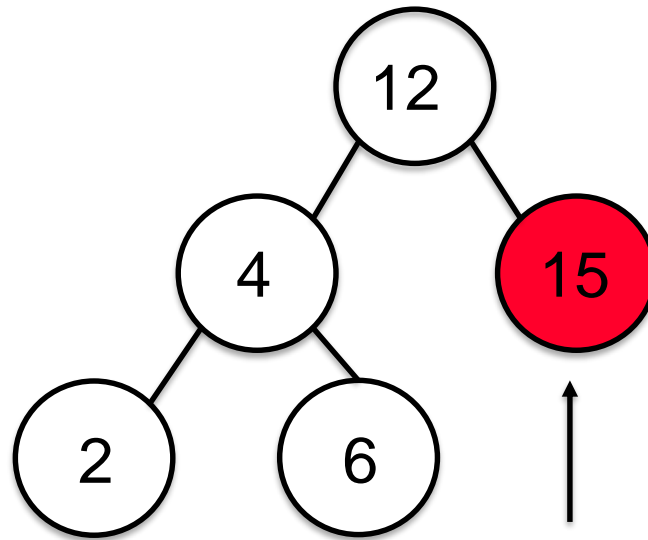
- Let's look for 14. To start searching, we start at the root:



- Is the root what we want? No.
- Is $14 > 12$? Yes.
 - 14 must exist to the right of 12.
- Go right.

How do we find an element?

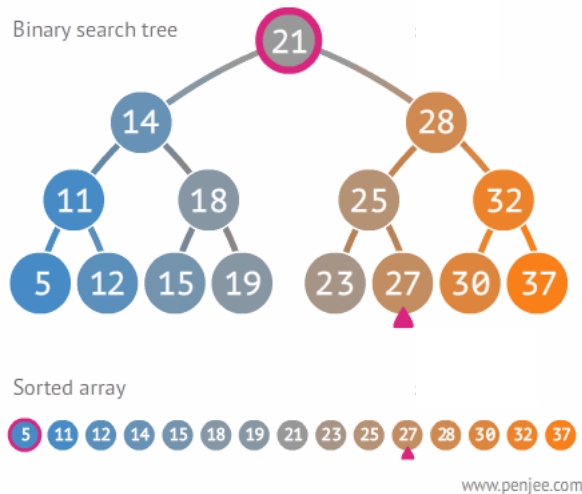
- Continue the quest for 14:



- Is the root what we want? No.
- Is $14 > 15$? No.
 - 14 must exist to the left of 15.
- Cannot keep searching.
- Return False.

Why do we use BST?

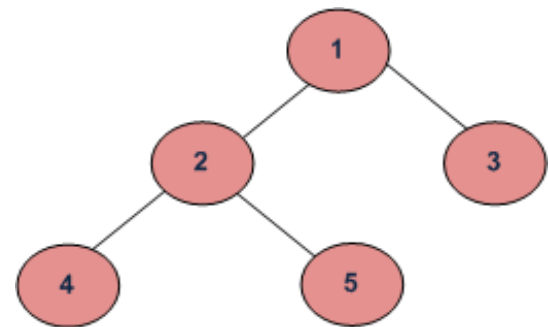
- Binary Search Trees is much more efficient for searching one element in a list



<https://blog.penjee.com/5-gifs-to-understand-binary-search-tree/>

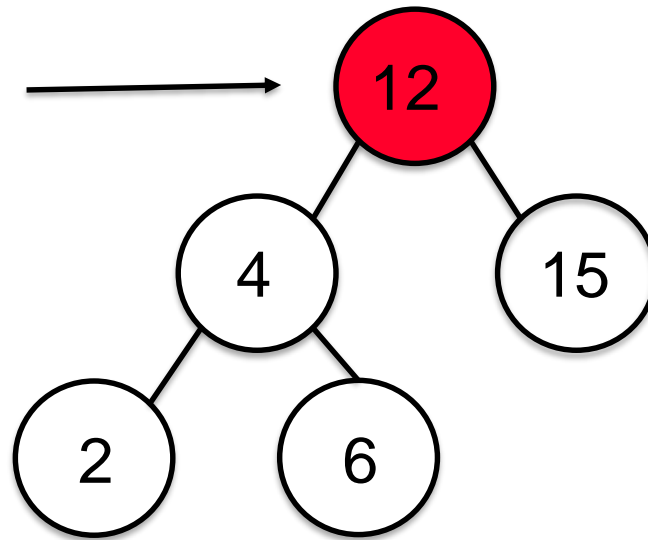
Binary Search Tree

- Find the **minimum** element:
 - All way to the left.
- Find the **maximum** element:
 - All way to the right.
- Tree traversal:
 - Refers to the process of visiting each node in a tree data structure, exactly once
 - Depth First Traversal (DFT),
 - Breadth First Traversal (BFT)



How do you add an element?

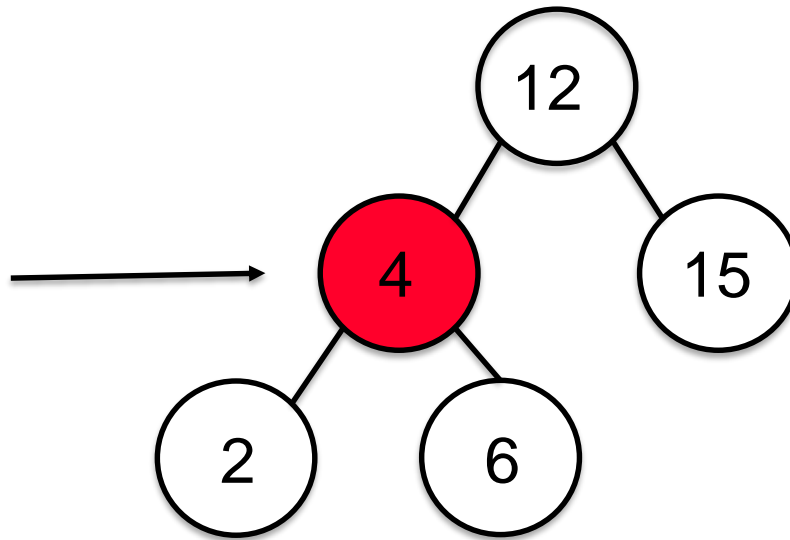
- First find the location. This is what we have already done. Lets try to add 5.



- Is the root what we want? No.
- Is $5 > 12$? No.
 - 5 must be inserted to the left of 12.
- Go left.

How do you add an element?

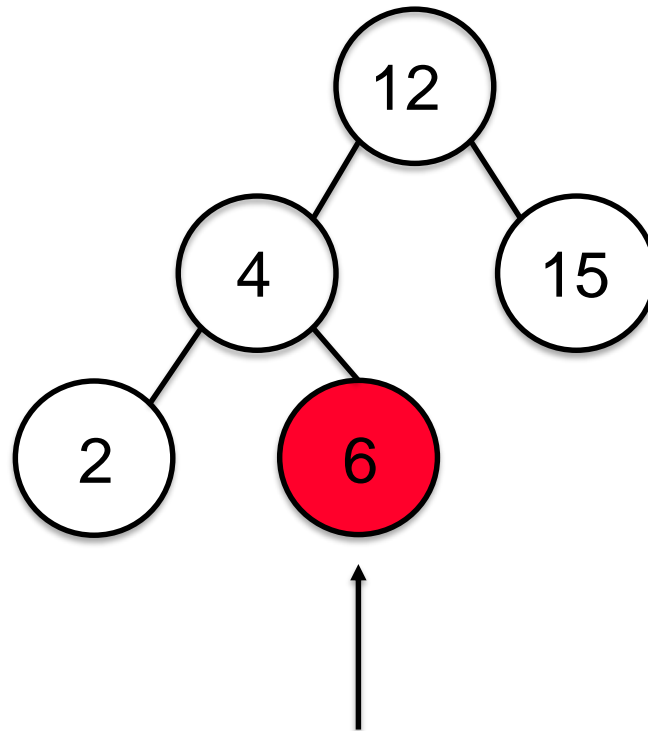
- Continue the search:



- Is the root what we want? No.
- Is $5 > 4$? Yes.
 - 5 must be inserted to the right of 4.
- Go right.

How do you add an element?

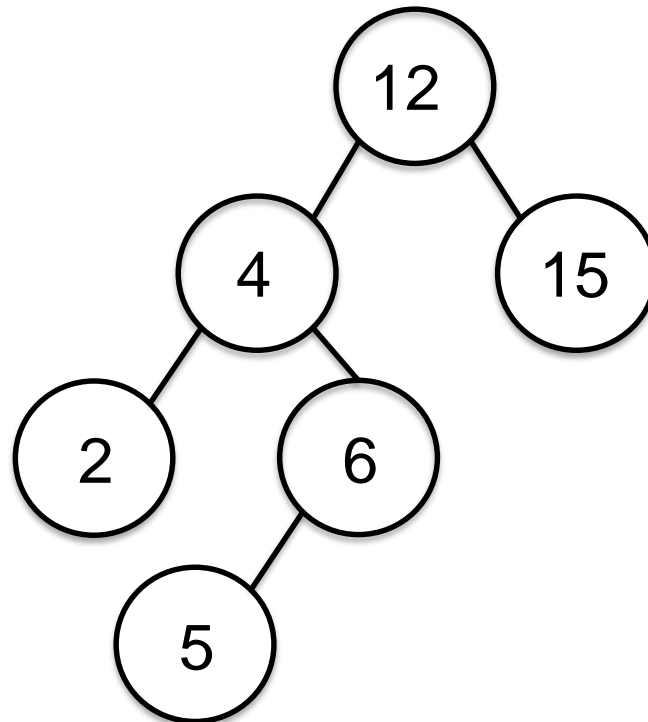
- Continue the search:



- Is the root what we want? No.
- Is $5 > 6$? No.
 - 5 must be inserted to the left of 6.

How do you add an element?

- We have the location!



- The left node does not exist.
- We have found the location for 5.
- Insert 5 here.

How do you construct from base?

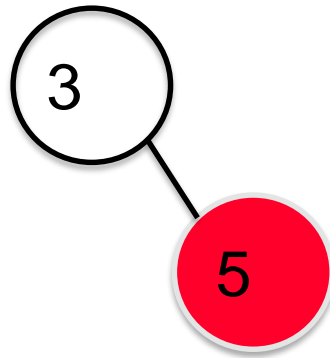
- We have the list: [3,5,4,7,6,1]



- 3 is the root

How do you construct from base?

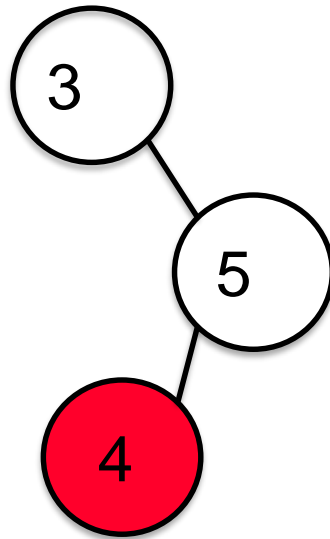
- We have the list: [3,5,4,7,6,1]



- Is $5 > 3$? Yes.
 - 5 must be inserted to the right of 3.
- Go to the next element.

How do you construct from base?

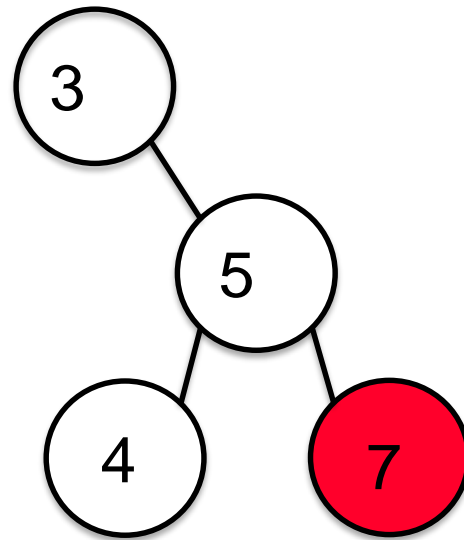
- We have the list: [3,5,4,7,6,1]



- Is $4 > 3$? Yes.
 - 4 must be inserted to the right of 3.
- Is $4 > 5$? No.
 - 4 must be inserted to the left of 5.

How do you construct from base?

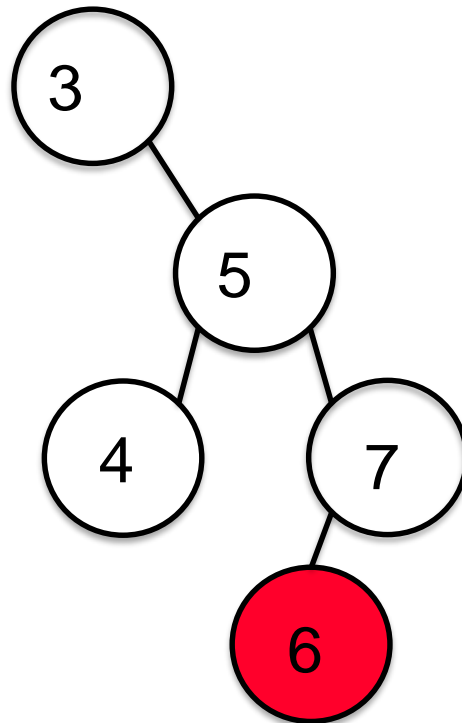
- We have the list: [3,5,4,7,6,1]



- Is $7 > 3$? Yes.
 - 7 must be inserted to the right of 3.
- Is $7 > 5$? Yes.
 - 7 must be inserted to the right of 5.

How do you construct from base?

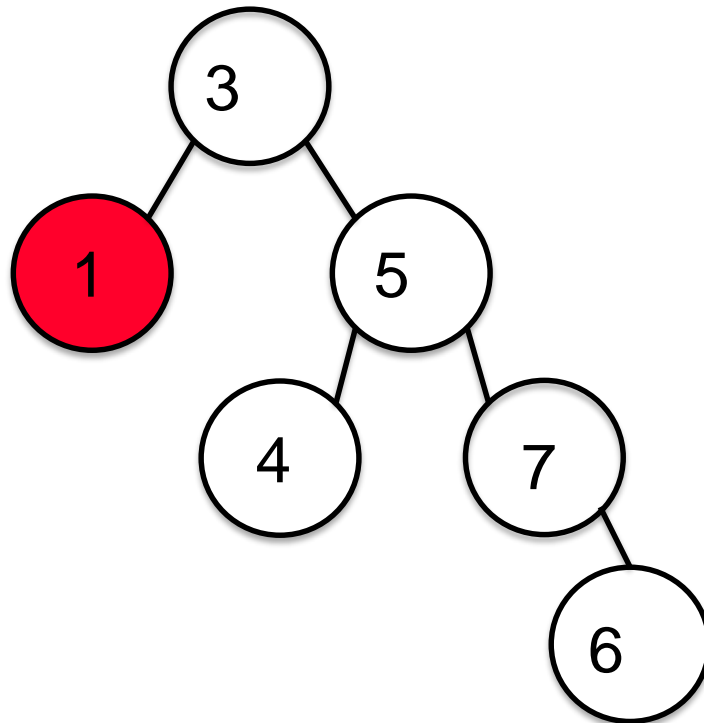
- We have the list: [3,5,4,7,6,1]



- Is $6 > 3$? Yes.
 - 6 must be inserted to the right of 3.
- Is $6 > 5$? Yes.
 - 6 must be inserted to the right of 5.
- Is $6 > 7$? No.
 - 6 must be inserted to the left of 7.

How do you construct from base?

- We have the list: [3,5,4,7,6,1]

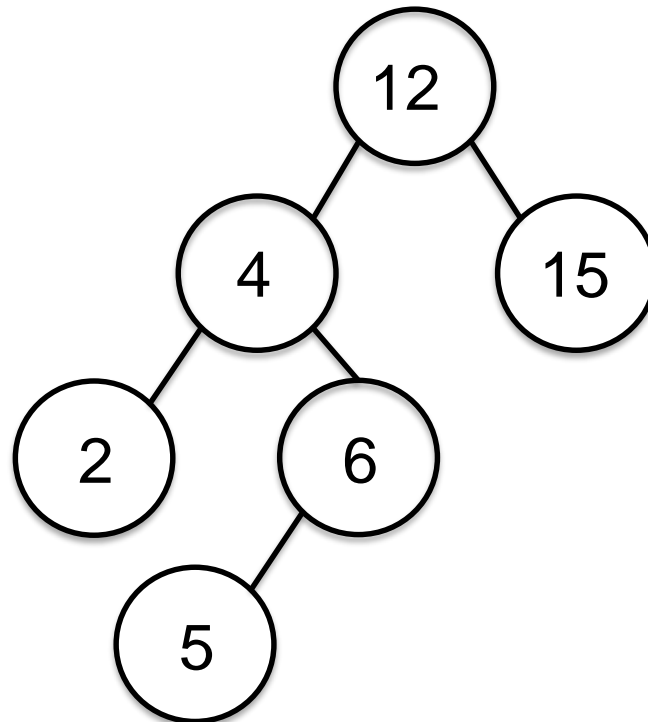


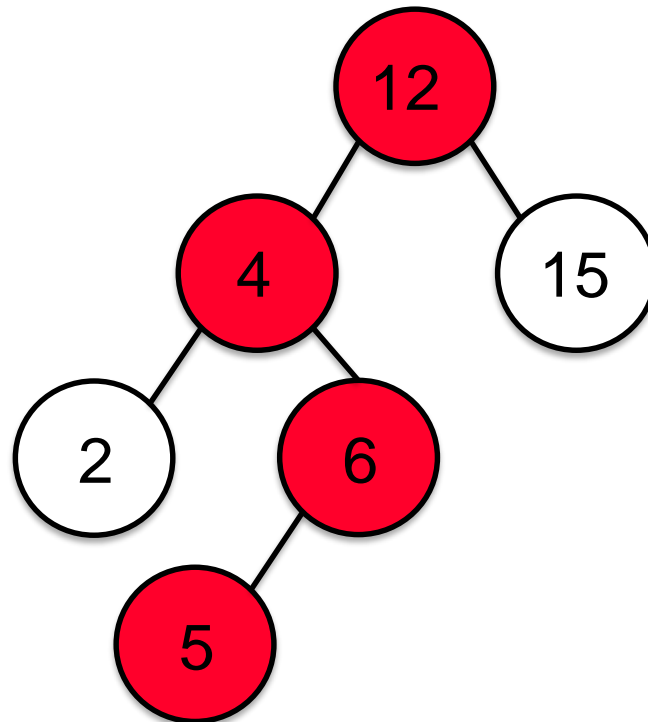
- Is $1 > 3$? No.
 - 1 must be inserted to the left of 3.

- What happens if we try to insert an element already there?
 - When we try to search for a location, if we come across the item in a node, stop searching.
 - Do not insert two elements of the same value into a search tree.

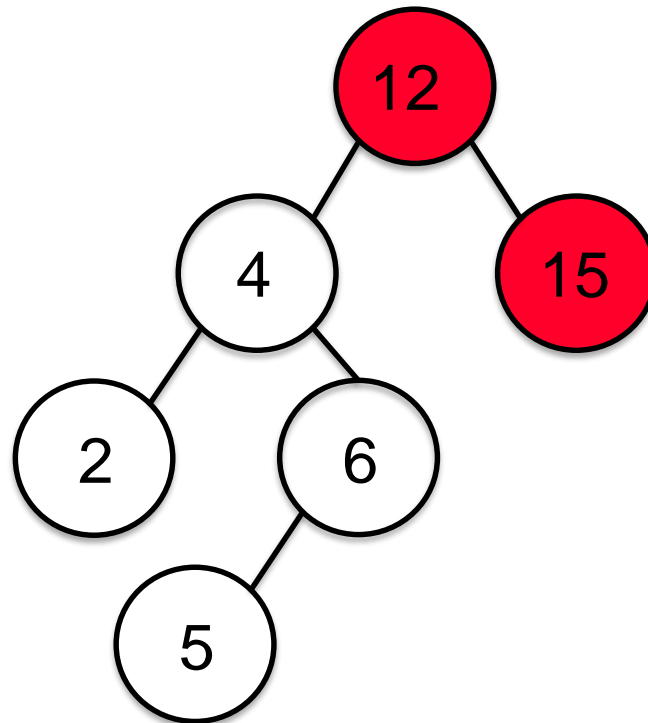
- BST allows for fast search times to discover if elements are present.
- How much time does it take? It depends on “balance”.
- A tree is considered balanced if the length of a branch is “close” to the length of all other branches.
 - For the sake of this unit, “close” will refer to ± 1 .

- The tree we just made is considered NOT balanced.





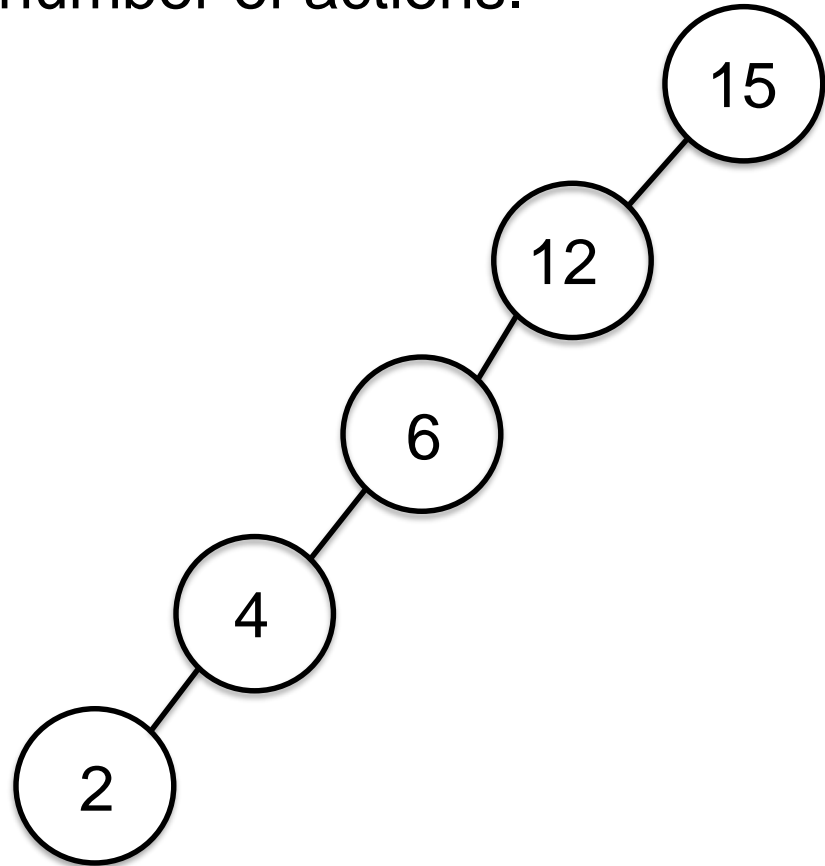
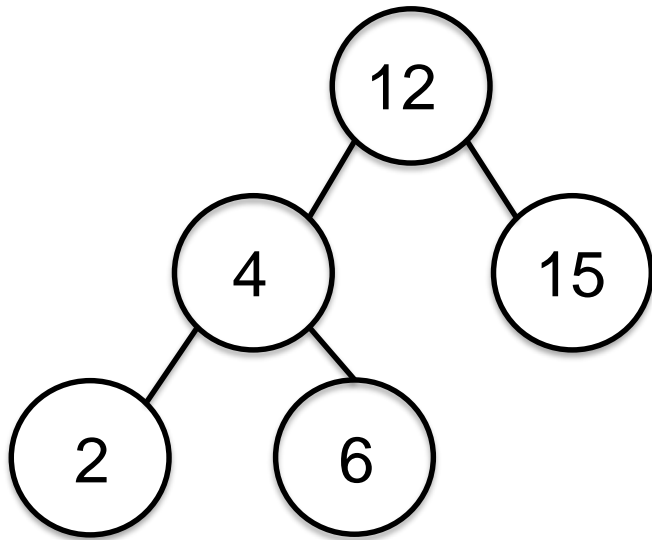
The depth of this branch is 3.



The depth of this branch is 1.

- $3-1 = 2$, thus, because of the difference in depth, the tree is not balanced.
- Why should we care? Consider the following situations.

- Search for 2, count the number of actions.



- For the balance tree, there are 3 probes to find 2.
- For the imbalanced tree there are 5 probes.
- This suggests that balance has an impact on complexity for BSTs (and it does).
- Right now, the numbers are quite small, but for very large trees there is a significant increase in time taken to find elements.

- For a balanced tree, the complexity is $\log_2 n$ because with every step down the tree, the search space possibilities are halved.
- For an imbalanced tree, the complexity is n because probing for an element that is not there, in the worst case, will require each element to be examined.

- Is it worth rebalancing trees? Think about it for your workshop.
- Rebalancing trees is outside the scope of the unit but is a very interesting topic to go into.