

FIT1045: Algorithms and Programming Fundamentals in Python

Lecture 4

Loops and Euclid's Algorithm



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Recap

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Clayton: **AXXULH**
Malaysia: **LWERDE**

Boolean expressions:

- Can you translate the following sentence to a Boolean expression?

A good fruit salad contains one main fruit which can be either oranges or melons and a second fruit that can be either strawberries or pineapple, but it should never contain avocado.

Boolean operators have precedence:

- Which parentheses can be avoided?

`((fruit1=='orange') and (fruit2=='orange')) or ((fruit1=='apple') and (fruit2=='apple'))`

This lecture

Learn about loops to implement our first textbook algorithm in Python

Learning outcomes

- 2 (choose and implement appropriate problem solving strategies in Python)
- 5 (determine limitations of algorithms)

Concrete goal: An efficient algorithm for computing the greatest common divisor

Where am I?

1. Greatest Common Divisor
2. While loops
3. Euclid's Algorithm

Motivation: simplifying fractions

$$\frac{18}{24} = \frac{3}{4}$$

The diagram illustrates the simplification of the fraction $\frac{18}{24}$ to $\frac{3}{4}$. Two curved arrows, one above and one below the equals sign, point from the numerator and denominator of the first fraction to those of the second. Both arrows are labeled $/\text{gcd}(18,24)$, indicating that the greatest common divisor of 18 and 24 is used to simplify the fraction.

$$\frac{18480831109}{9231418071} = ?$$

Greatest Common Divisor Problem

Input: two positive integers m and n

Output: greatest common divisor, $\text{gcd}(m, n)$

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Let's find an algorithm

Greatest Common Divisor Problem

Input: two positive integers m and n

Output: greatest common divisor, $\text{gcd}(m, n)$

Observations

- the greatest possible common divisor is the smaller of the two numbers, e.g. $\text{gcd}(178, 89) = 89$
- the smallest possible divisor is 1, e.g. $\text{gcd}(97, 53) = 1$
- we are after the **greatest** divisor, e.g. $\text{gcd}(24, 18) = 6$, not 1, 2, or 3

“Brute force” Algorithm

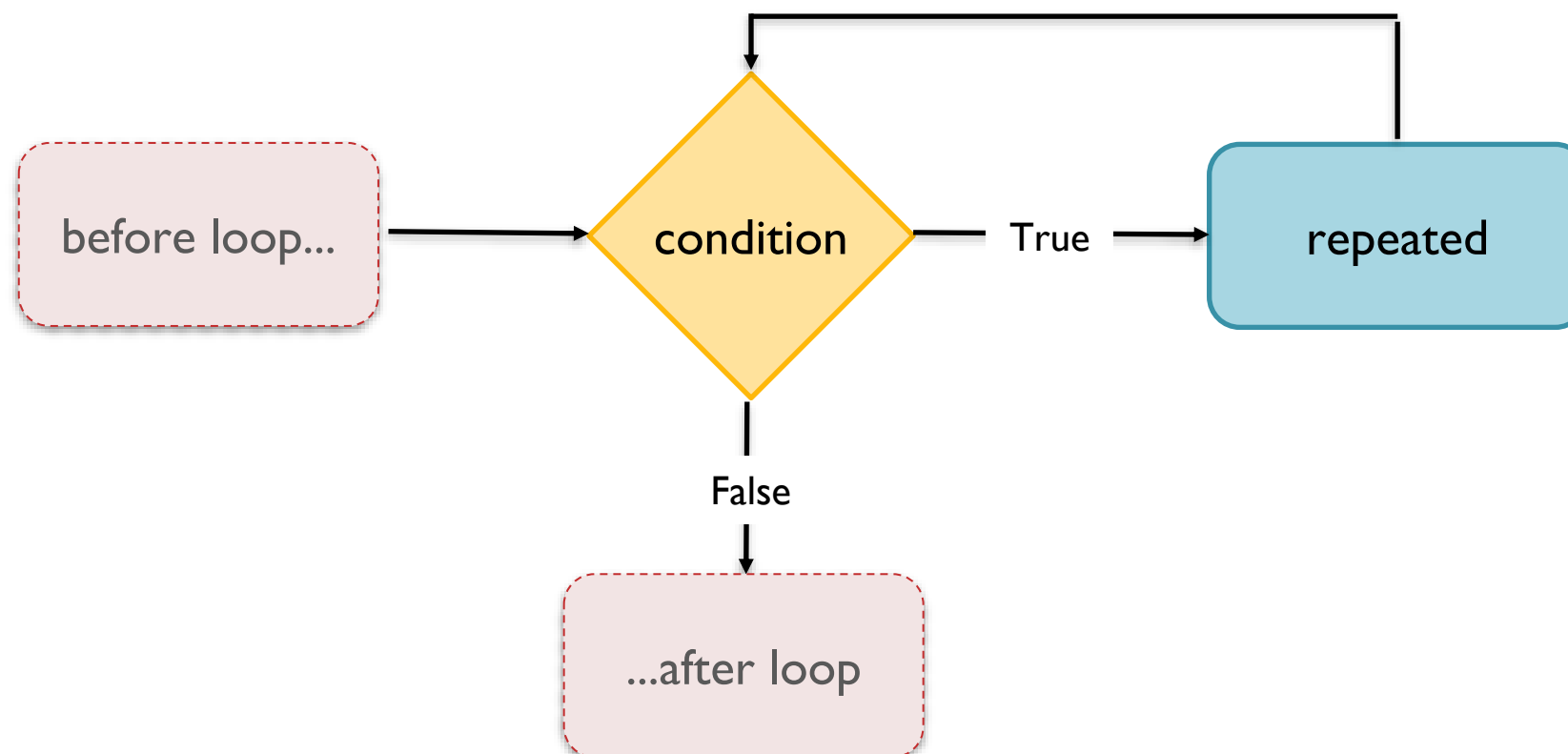
check all integers between $\min(m, n)$ and 1 (from big to small), output first common divisor encountered

How to “check every integer”?

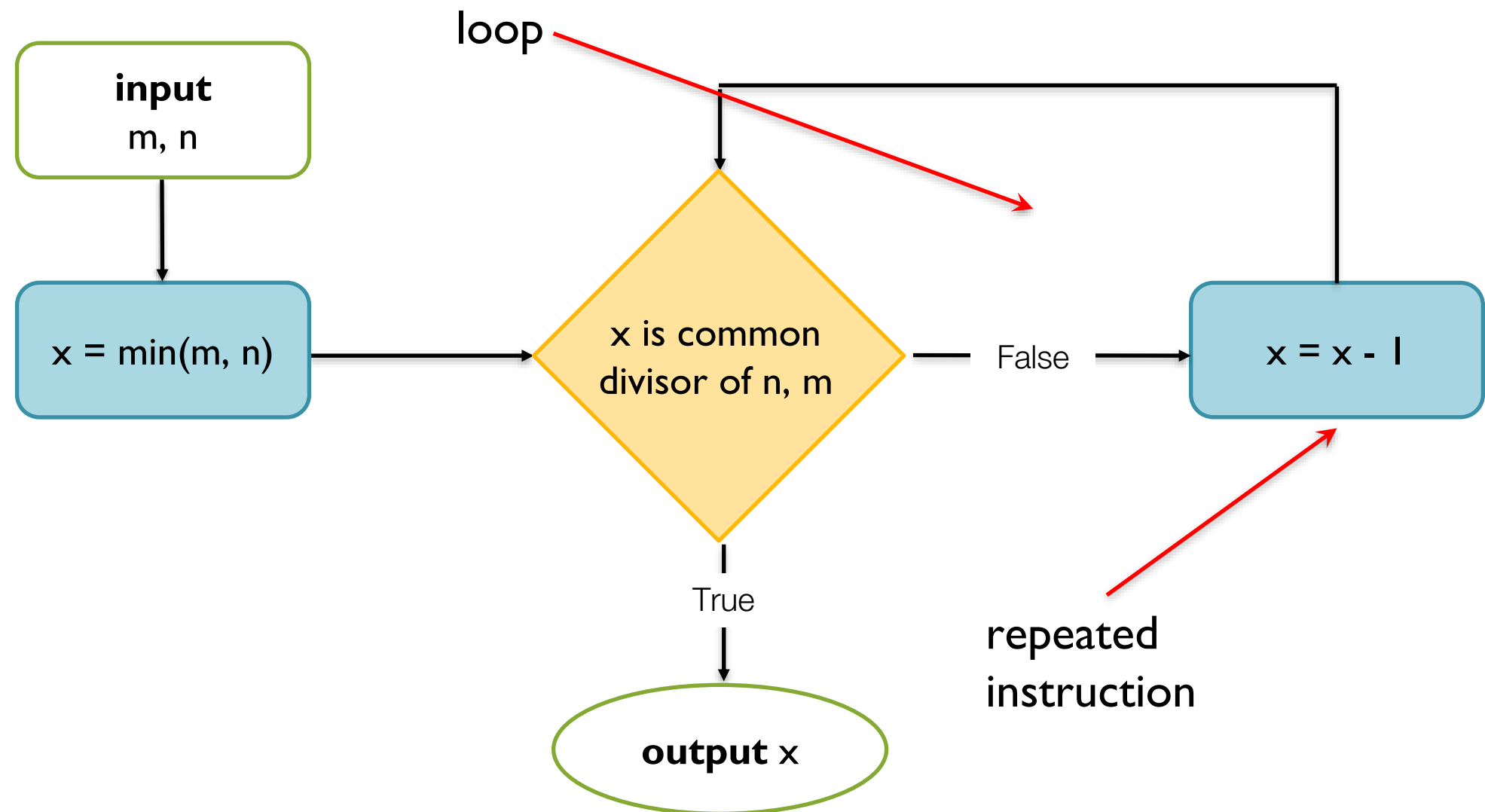
Observations

- Depending on input there can be an arbitrary number of integers to check
- Program will always have only a fixed number of instructions

Need to repeat some instructions many times **in a loop**.



How to “check every integer”?



“Brute force” Algorithm

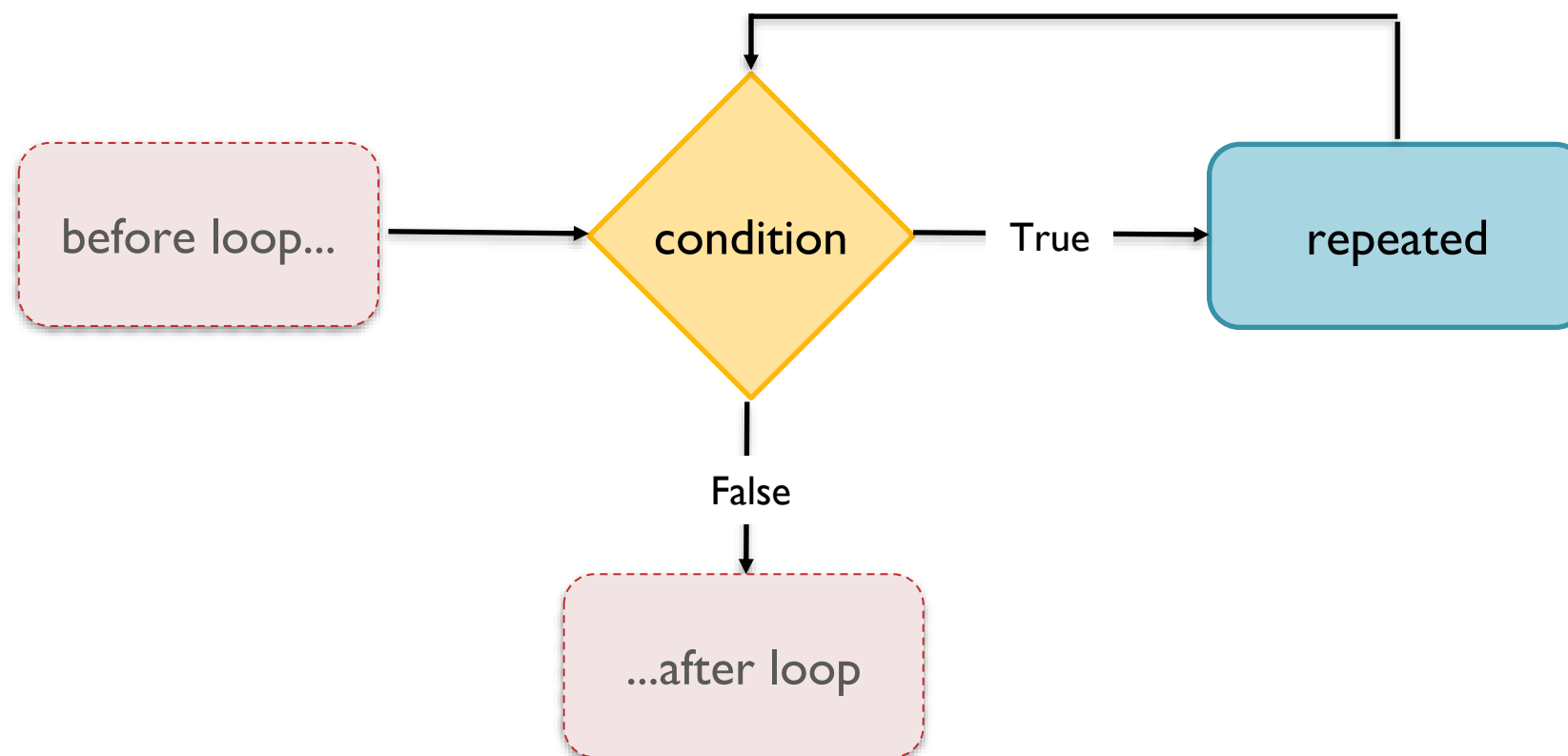
check all integers between $\min(m, n)$ and 1 (from big to small),
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Where am I?

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2. While loops
3. Euclid's Algorithm

While statement in Python for loopy control flows

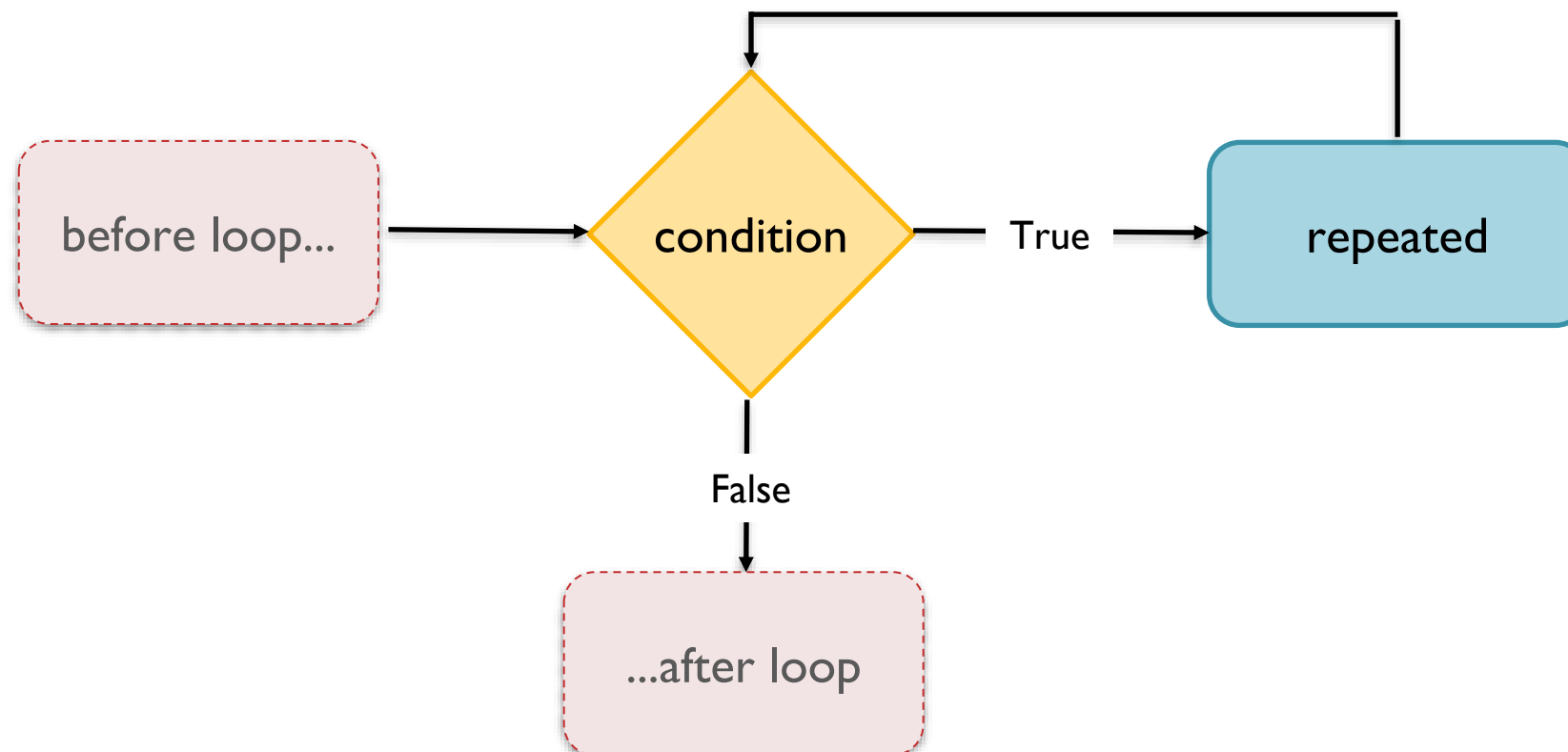
```
# before loop
# ...
#
while condition:
    # repeated
    # ...
    #
# after loop
# ...
#
```



Example: summing first n integers

```
def sum_of_first_n_ints(n):  
    """  
    Input : positive integer n  
    Output: sum of pos. integers up to n"""
```

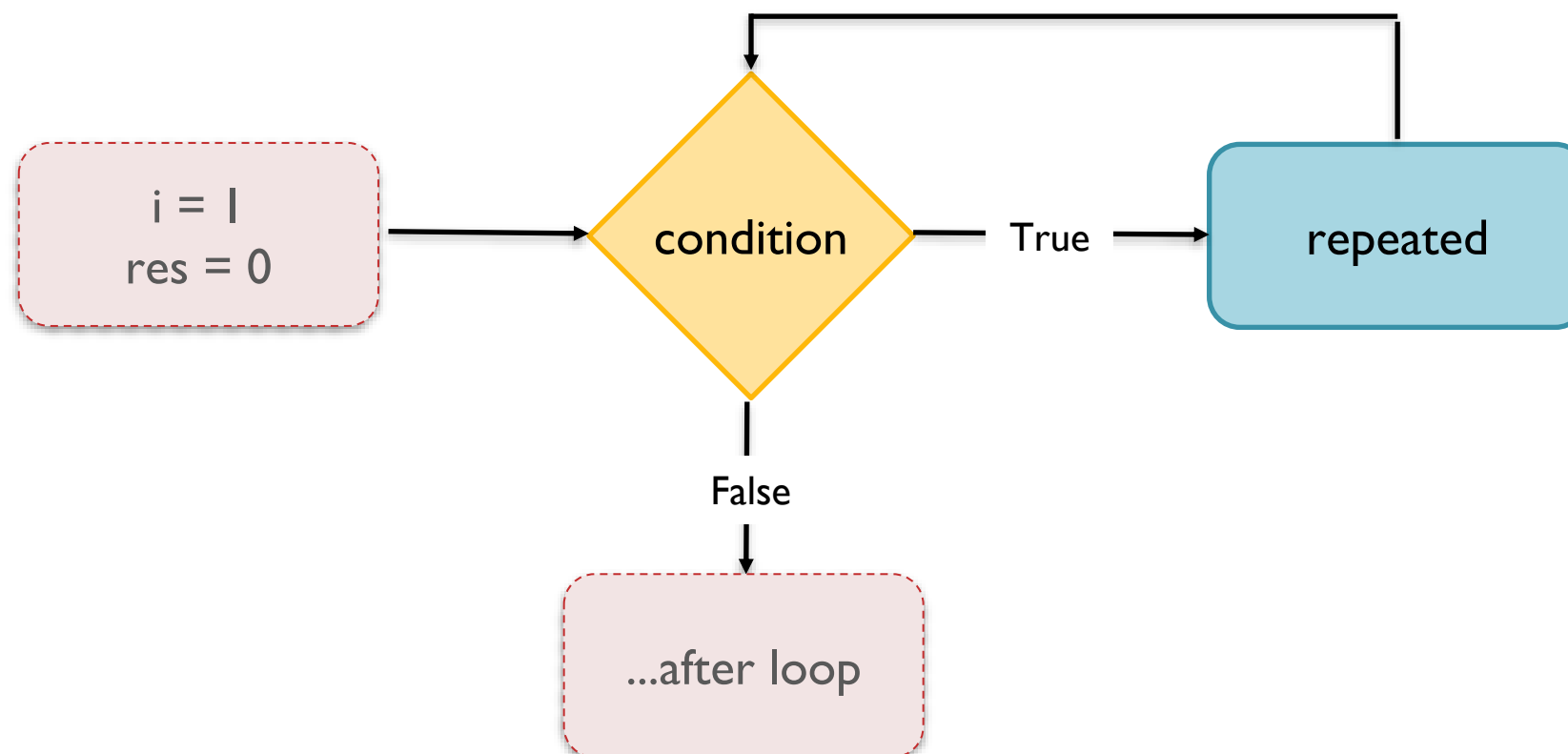
$$1 + 2 + \cdots + n = ?$$



Example: summing first n integers

```
def sum_of_first_n_ints(n):  
    """  
    Input : positive integer n  
    Output: sum of pos. integers up to n"""  
    i = 1 #iteration variable  
    res = 0 #accumulation variable
```

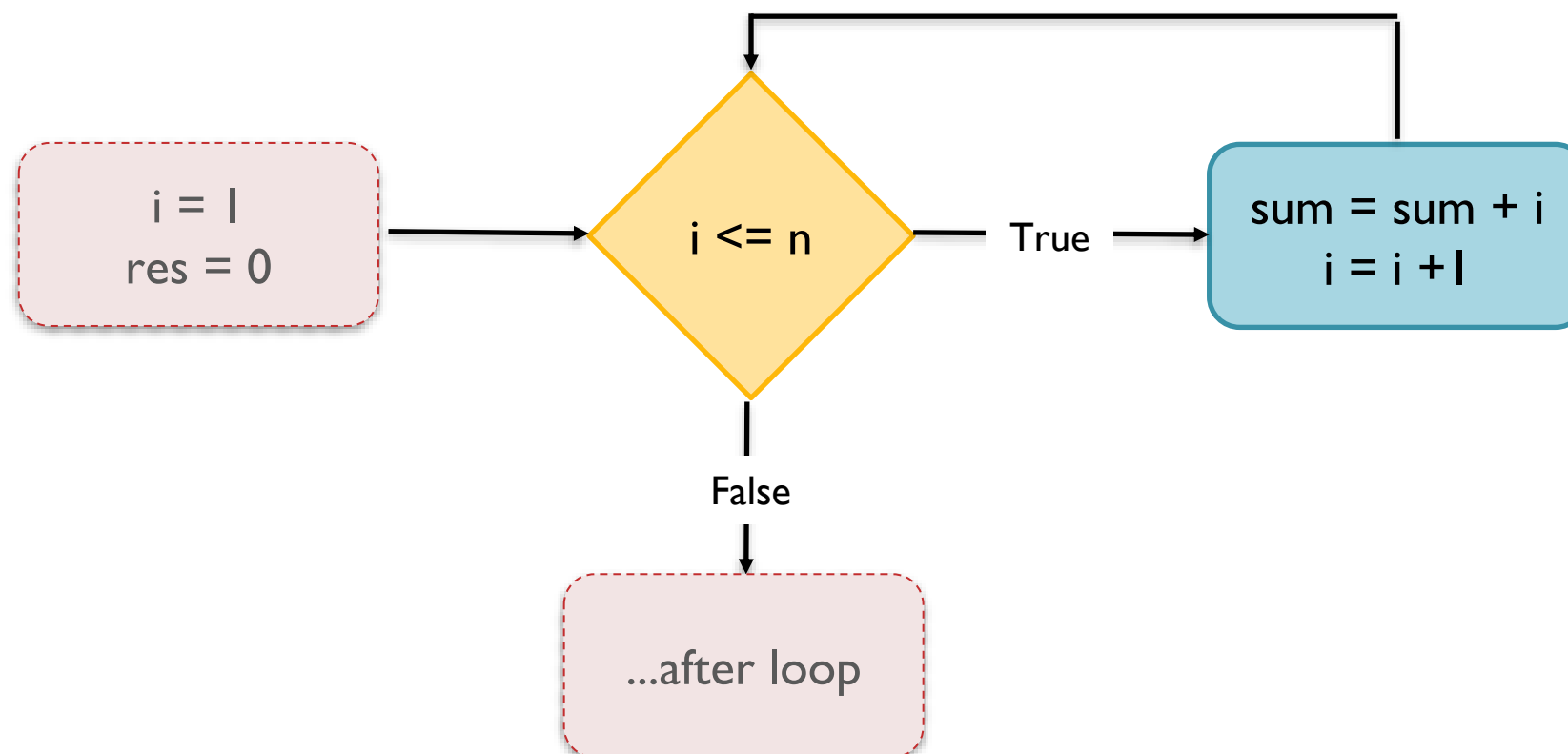
$$1 + 2 + \dots + n = ?$$



Example: summing first n integers

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def sum_of_first_n_ints(n):  
    """  
    Input : positive integer n  
    Output: sum of pos. integers up to n"""  
    i = 1 #iteration variable  
    res = 0 #accumulation variable  
    while i <= n:  
        res = res + i  
        i = i + 1
```

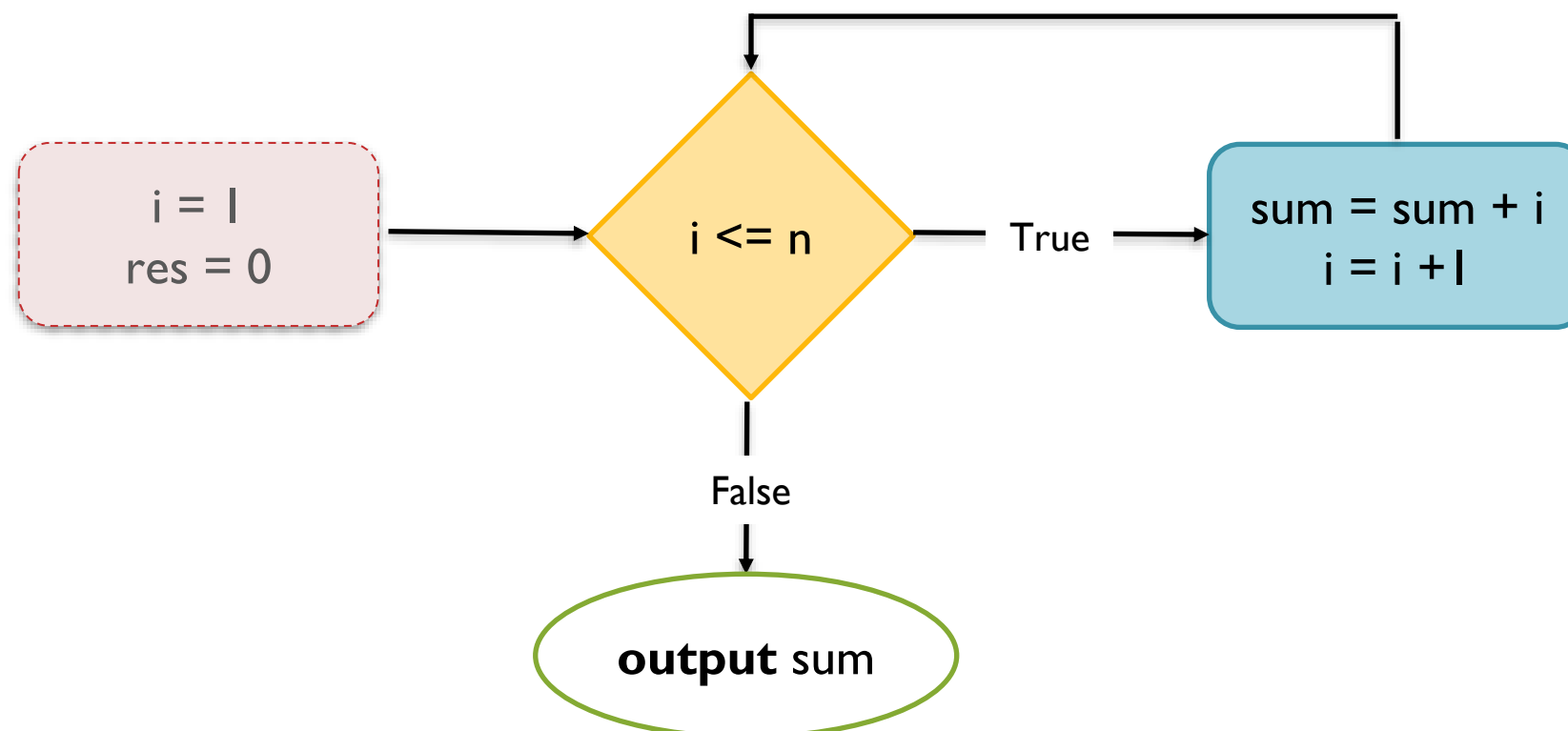
$$1 + 2 + \dots + n = ?$$



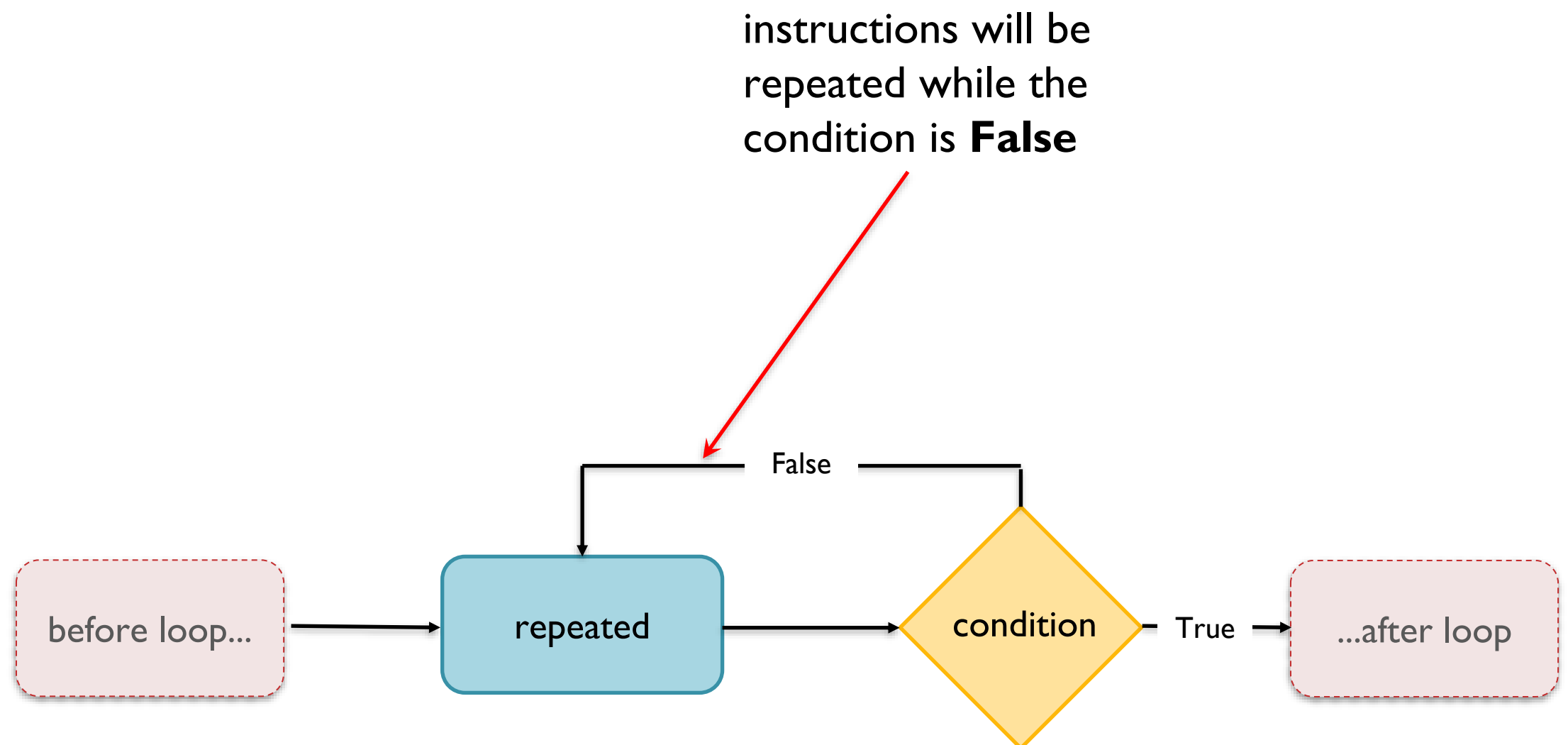
Example: summing first n integers

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    res = 0 #accumulation variable  
    while i <= n:  
        res = res + i  
        i = i + 1  
    return res
```

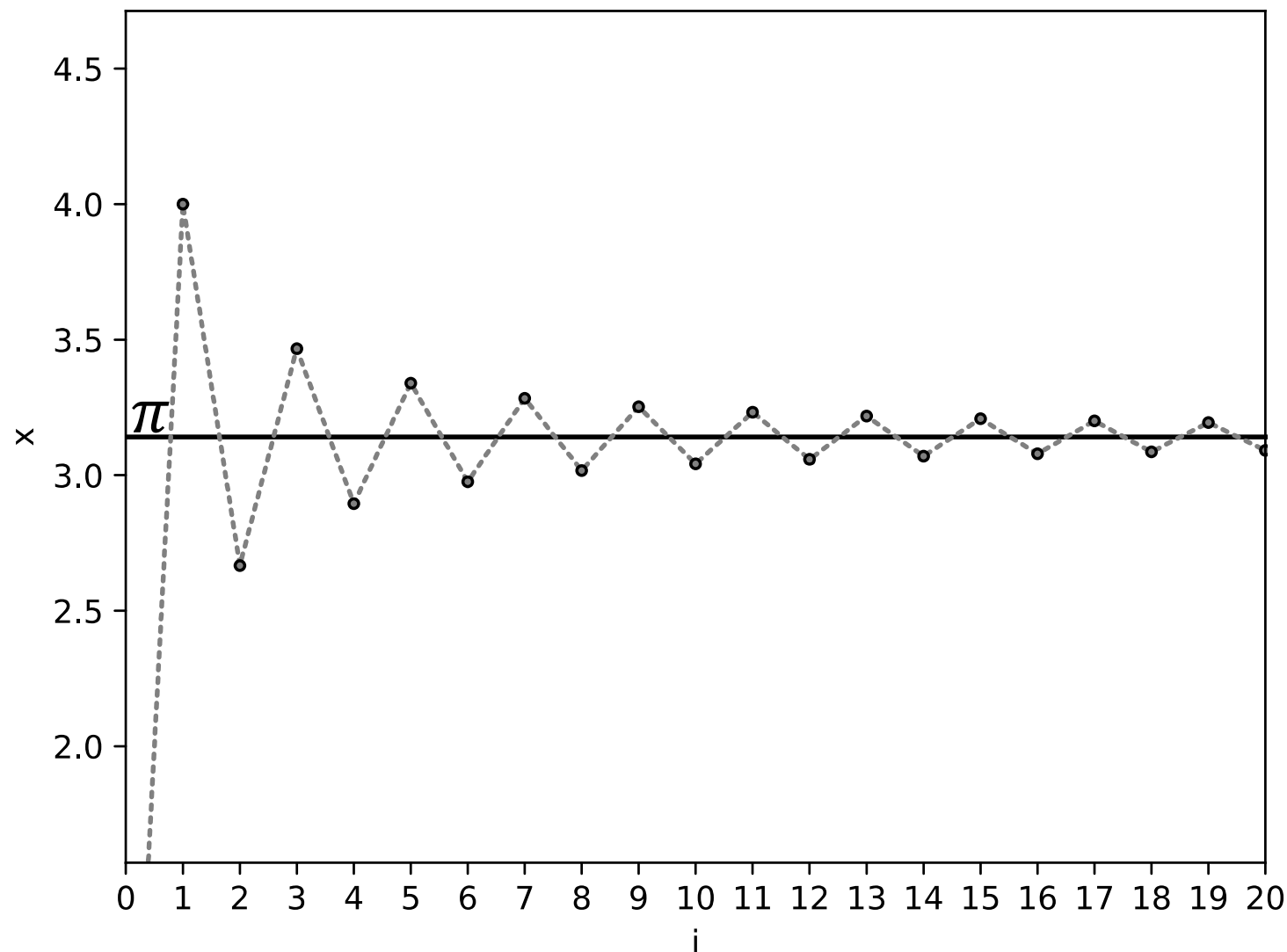
$$1 + 2 + \dots + n = ?$$



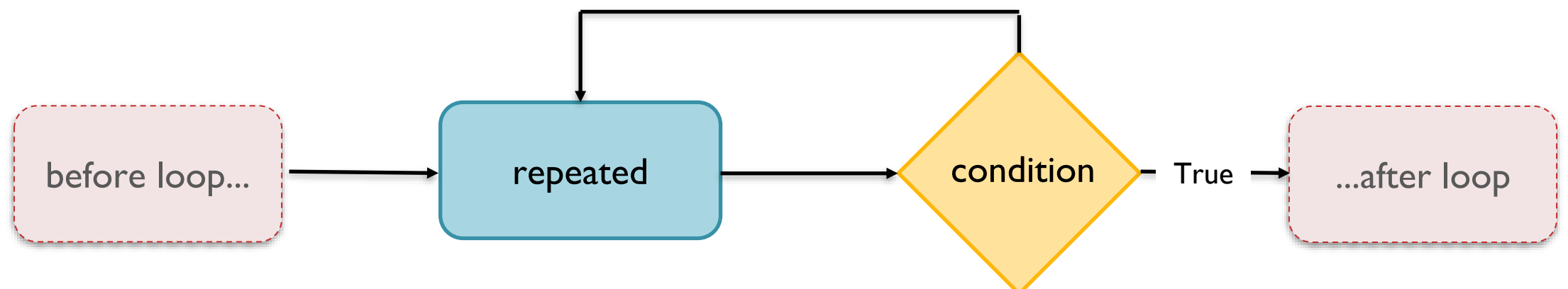
Sometimes we want condition last



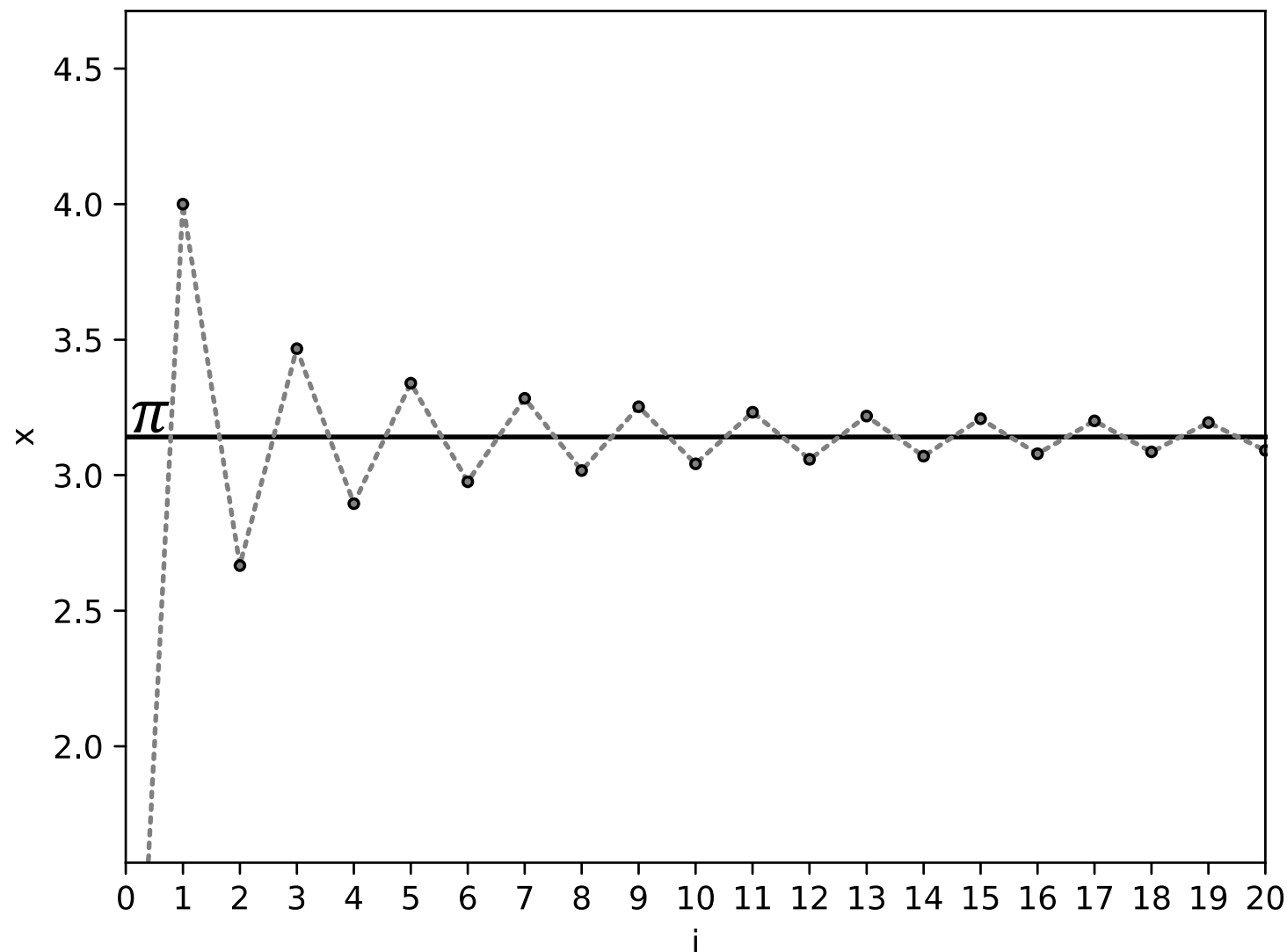
Example: approximating π



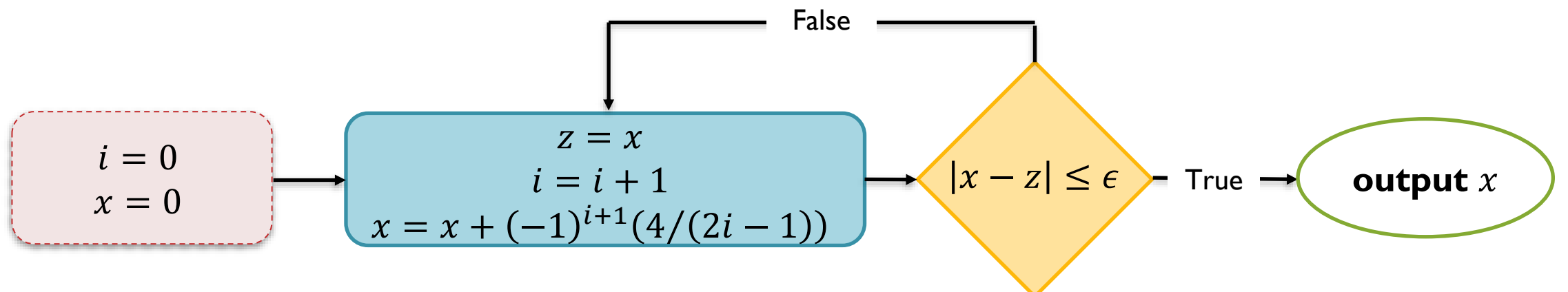
$$\pi = \frac{4}{1} - \frac{4}{3} + \frac{4}{5} - \frac{4}{7} \dots$$
$$= \sum_{i=1}^{\infty} (-1)^{i+1} \frac{4}{2i-1}$$



Example: approximating π



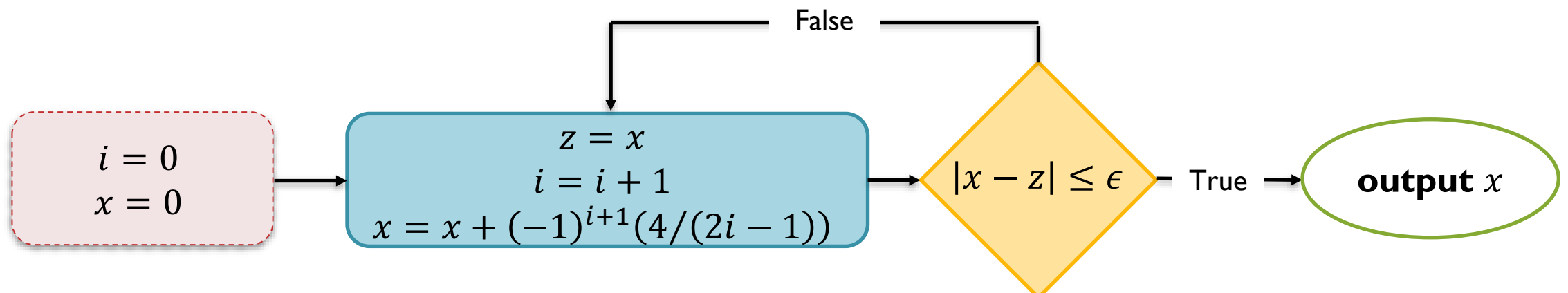
$$\pi = \frac{4}{1} - \frac{4}{3} + \frac{4}{5} - \frac{4}{7} \cdots$$
$$= \sum_{i=1}^{\infty} (-1)^{i+1} \frac{4}{2i-1}$$



To realise condition at end we can use conditional **break** statement

```
def pi_approximation(eps):  
    """  
    Input : positive float eps (accuracy)  
    Output: approximation x to pi with abs(x-pi)<=eps  
    """  
    i = 0    #iteration variable  
    x = 0    #candidate solution  
    while True:   
        i += 1  
        z = x  
        x = x + (-1)**(i+1)*4/(2*i-1)  
        # exit loop here if abs(x-z)<=eps  
  
    return x
```

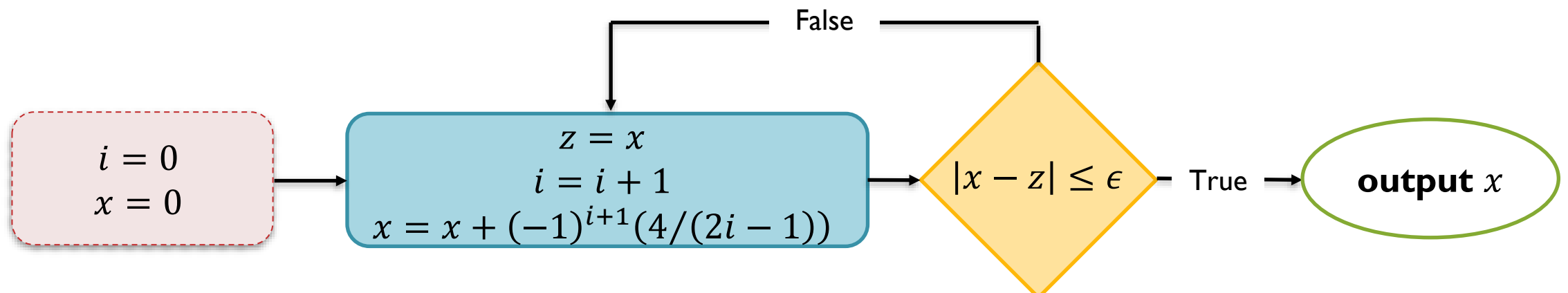
loop never
exits here



To realise condition at end we can use conditional **break** statement

```
def pi_approximation(eps):  
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    Input : positive float eps (accuracy)  
    Output: approximation x to pi with abs(x-pi)<=eps  
    """  
    i = 0 #iteration variable  
    x = 0 #candidate solution  
    while True:  
        i += 1  
        z = x  
        x = x + (-1)**(i+1)*4/(2*i-1)  
        if abs(x - z) <= eps:  
            break  
    return x
```

break statement
exits surrounding
loop



Loops pose a new kind of danger

```
def sum_of_first_n_ints(n):  
    """  
    Input : positive integer n  
    Output: sum of pos. integers up to n"""  
    i = 1 #iteration variable  
    res = 0  
    while i <= n:  
        res = res + i  
        i = i + 1  
    return sum
```

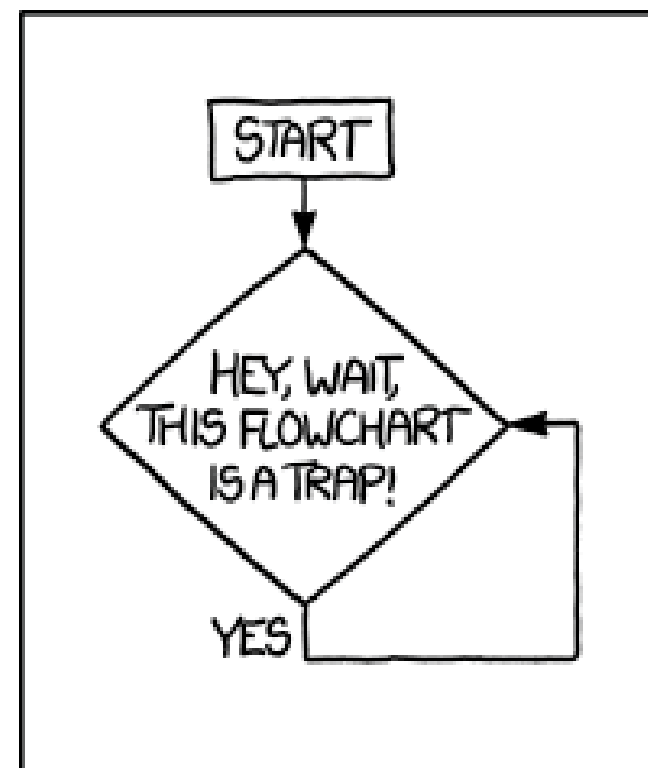
$$1 + 2 + \dots + n = ?$$

Forgetting just one line
results in this flowchart
☺

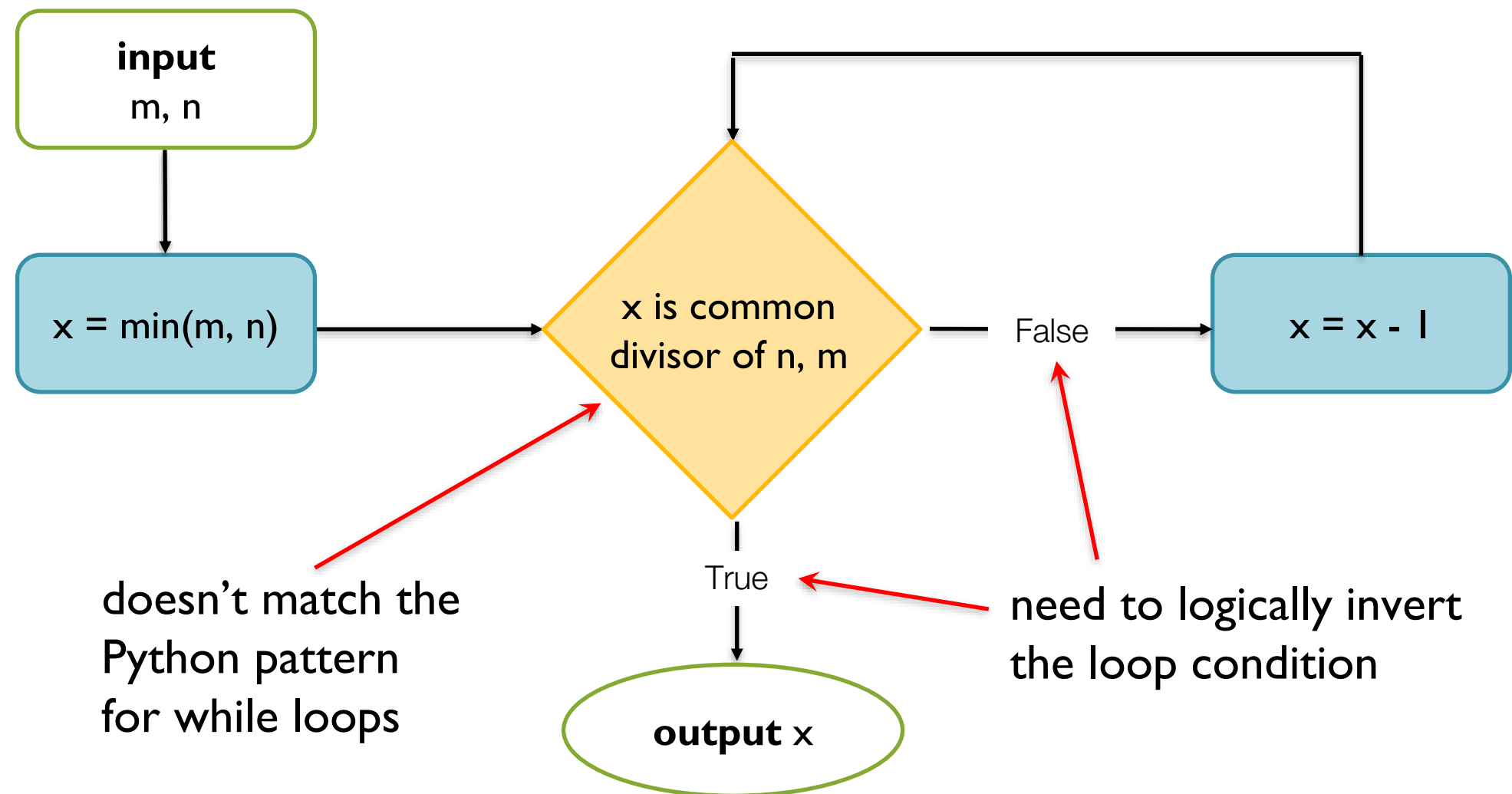
Everyone in this unit (including staff) will write an **infinite loop** by accident at some point

Common mistakes:

- Counter not incremented
- Tautological condition
- Forgot break



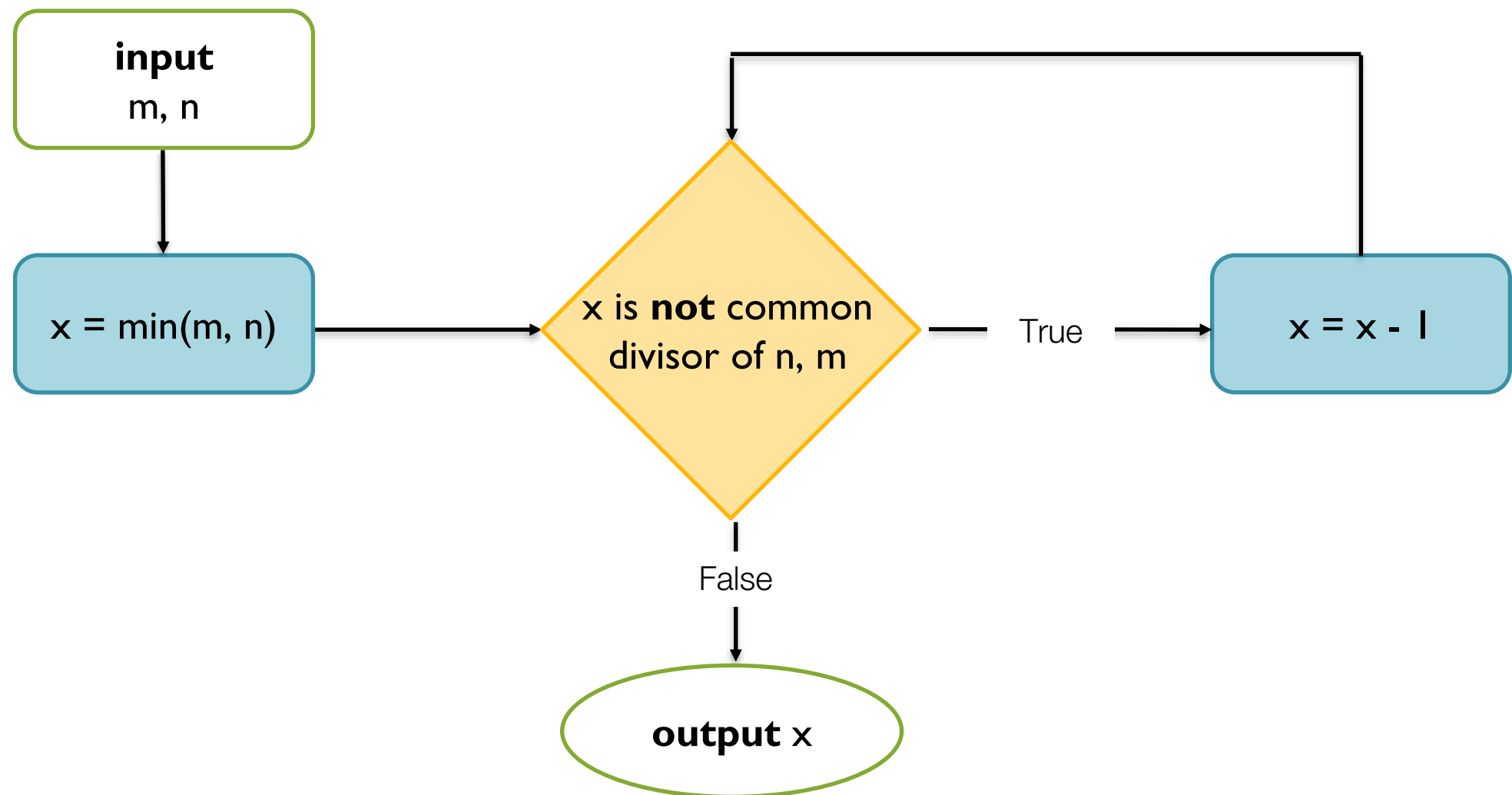
Using **while** to implement brute force GCD algorithm



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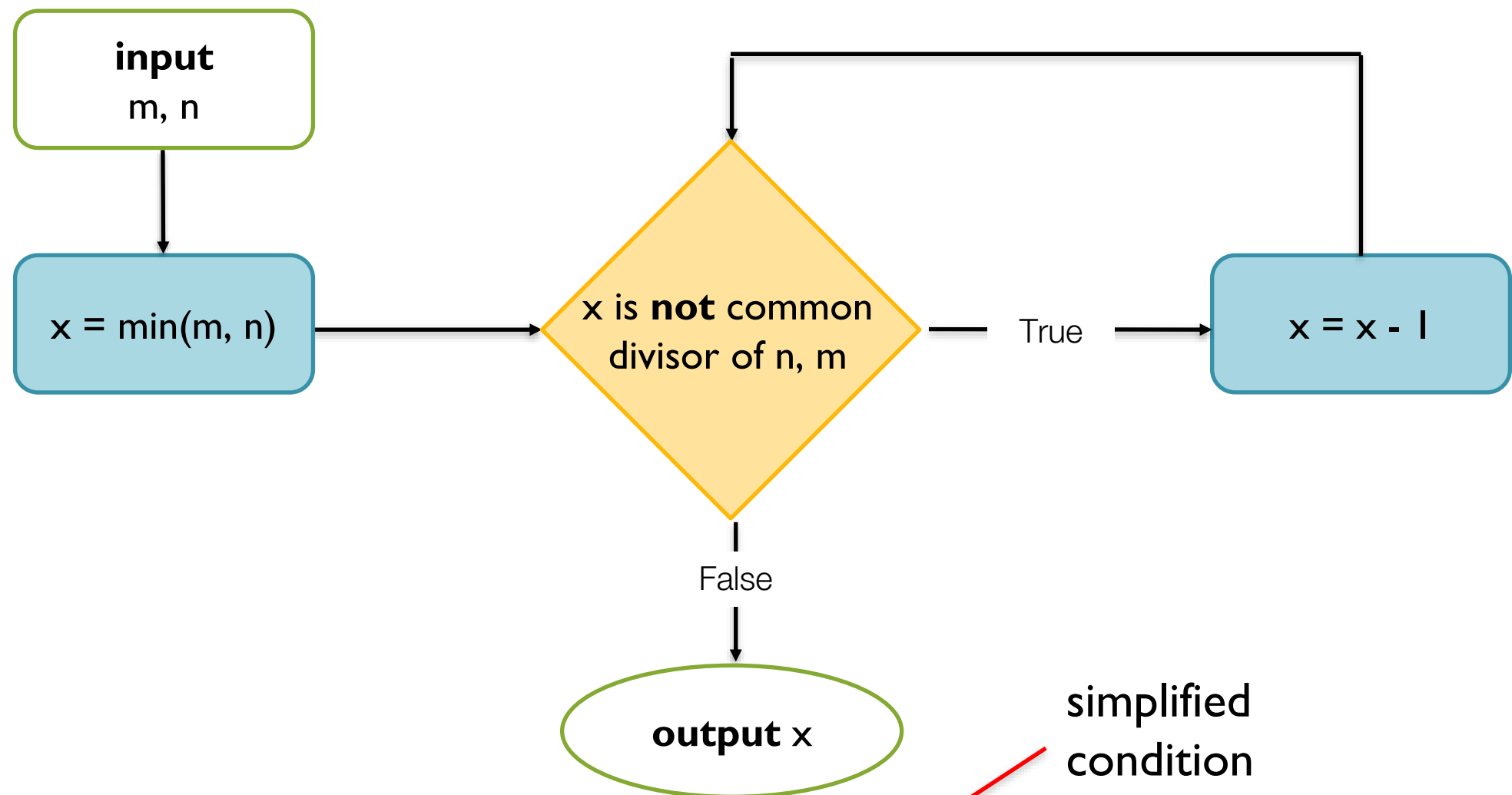
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Malaysia: **LWERDE**

Using **while** to implement brute force GCD algorithm



```
def gcd_brute_force(m, n):  
    x = min(m, n)  
    while not (m % x == 0 and n % x == 0):  
        x = x - 1  
    return x
```

Using **while** to implement brute force GCD algorithm



```
def gcd_brute_force(m, n):  
    x = min(m, n)  
    while (m % x != 0) or (n % x != 0):  
        x = x - 1  
    return x
```



Analysing our GCD algorithm

We have implemented an algorithm to find the GCD of two numbers.

But our algorithm is very inefficient!

Is there a better algorithm?

Where am I?

1. Greatest Common Divisor
2. While loops
3. Euclid's Algorithm

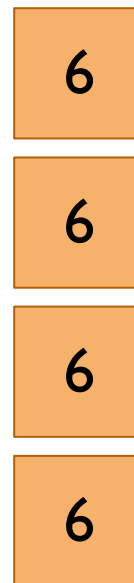
Let's *analyse the problem* to derive smarter algorithm

Orange: $24 = 4 \times 6$

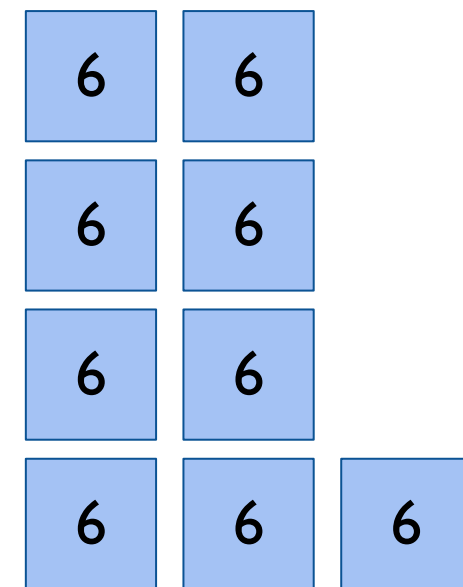
Blue: $54 = 9 \times 6$

Both stacks are
made of 6s.

24

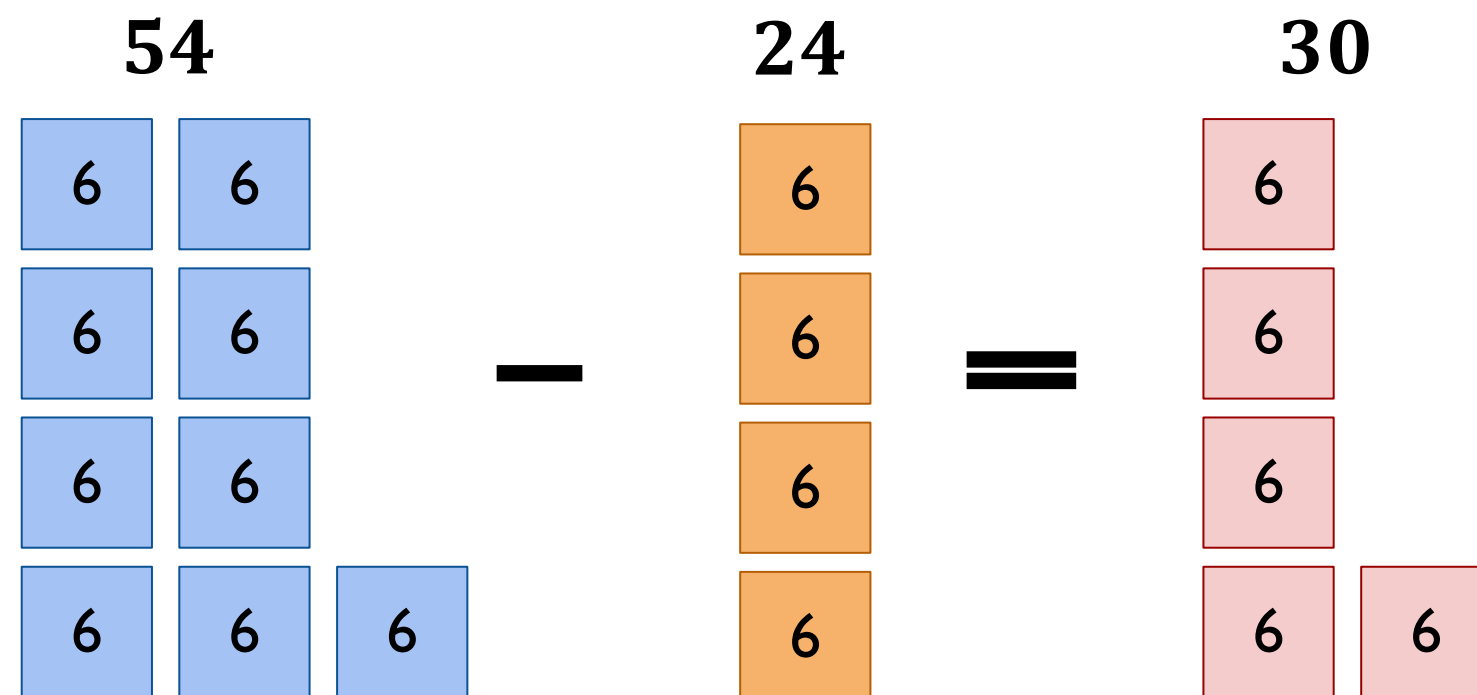


54



Input can be decreased to *smaller input with same output*

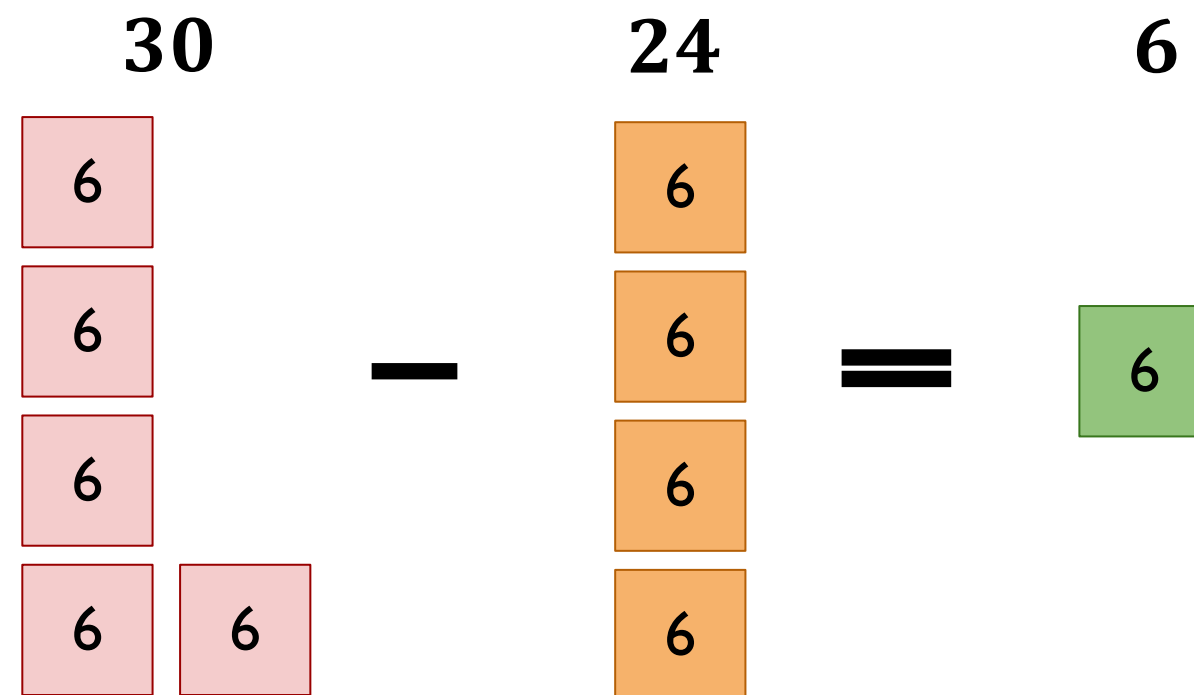
If we *subtract* the smaller stack from the bigger stack, the result will also be made of 6s:



We have (almost) shown that $\gcd(54, 24) = \gcd(24, 30)$

This reduction can be applied repeatedly

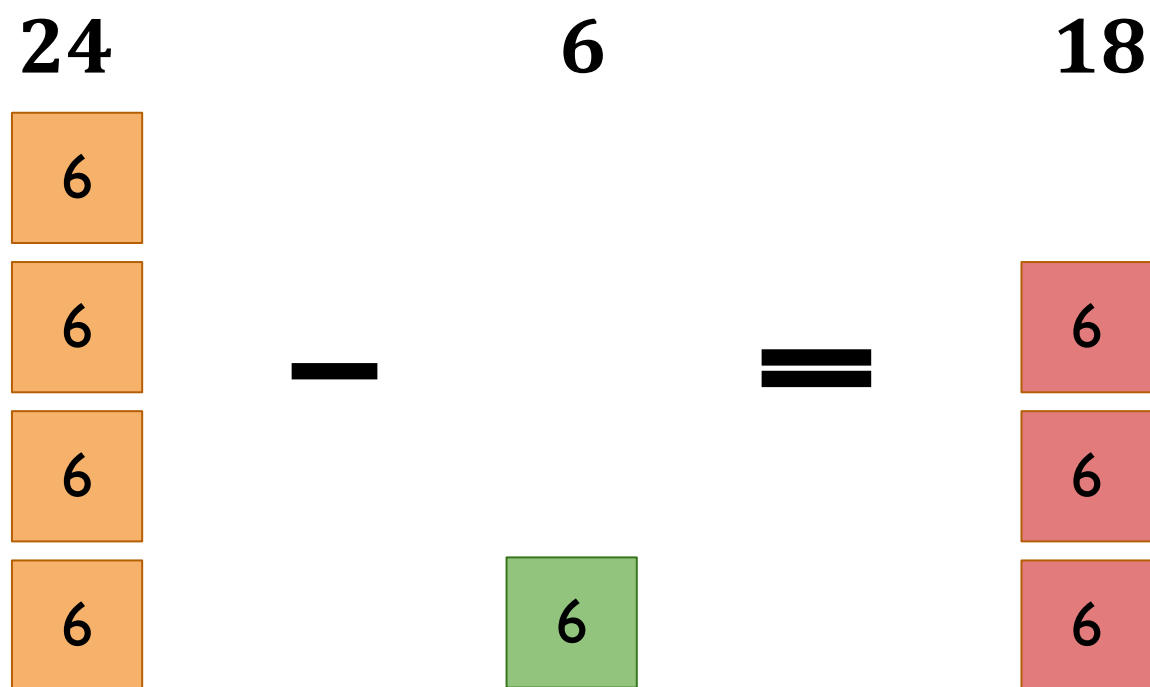
We repeat the process, subtracting the new smallest stack from the previous smallest stack:



We have shown that $\gcd(54, 24) = \gcd(24, 6)$

...but we don't have to stop there

We repeat the process, subtracting the new smallest stack from the previous smallest stack:



We have shown that $\gcd(54, 24) = \gcd(24, 6)$

Eventually we end up at a simple case

The of a non-zero m number and zero is simply m

6

0

6

6

—

=

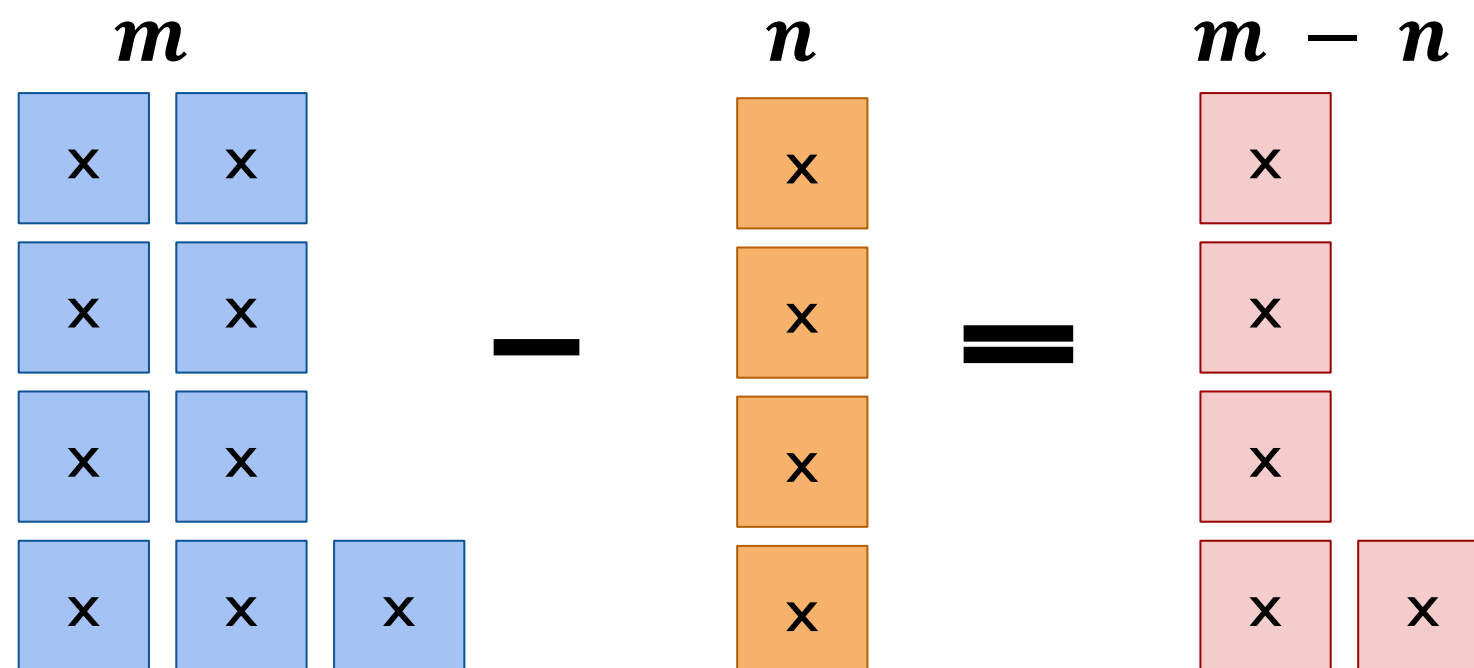
6

We have shown that $\gcd(6, 0) = 6$

Pattern also holds for *unknown number* in each box

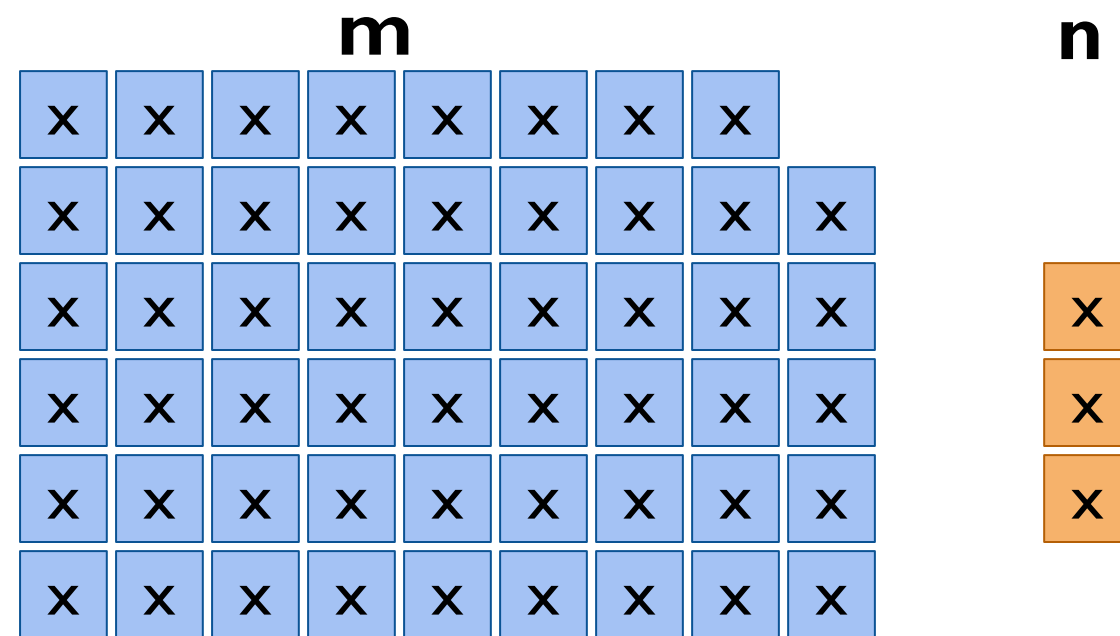
We have shown that: $\gcd(9x, 4x) = \gcd(4x, 9x - 4x)$

We can generalise: $\gcd(m, n) = \gcd(n, m - n)$



We can improve efficiency

What would happen if our m and n started like this:

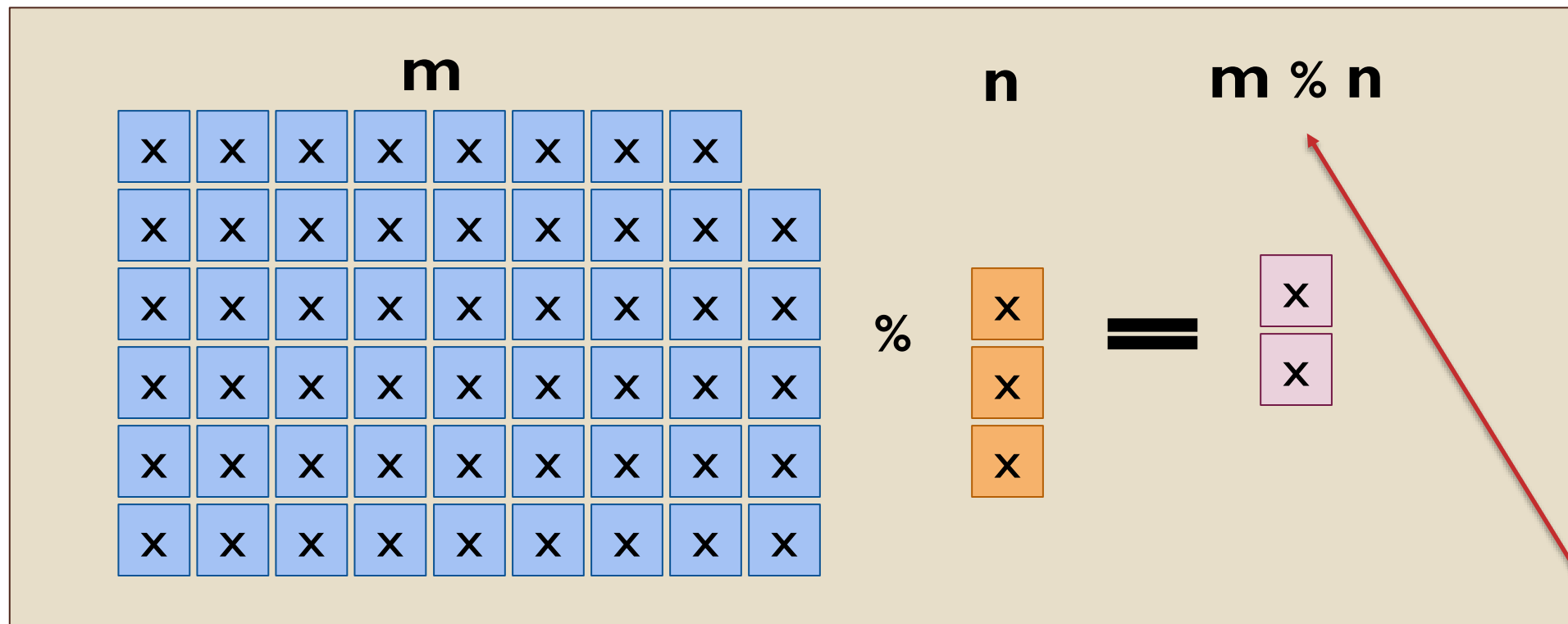


We would have to subtract n 17 times!

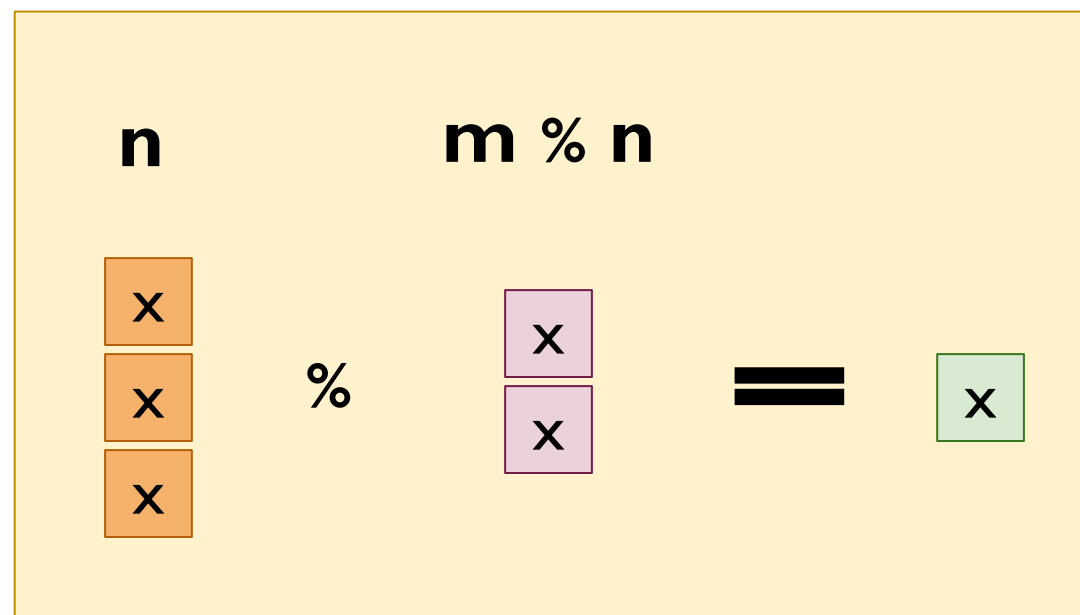
Instead of subtracting, we get the same result if we take the **integer remainder** of dividing m by n , i.e., $m \% n$ in Python.

(This operation is also called **modulo**).

Final problem reduction:

$$\gcd(m, n) = \gcd(n, m \% n)$$


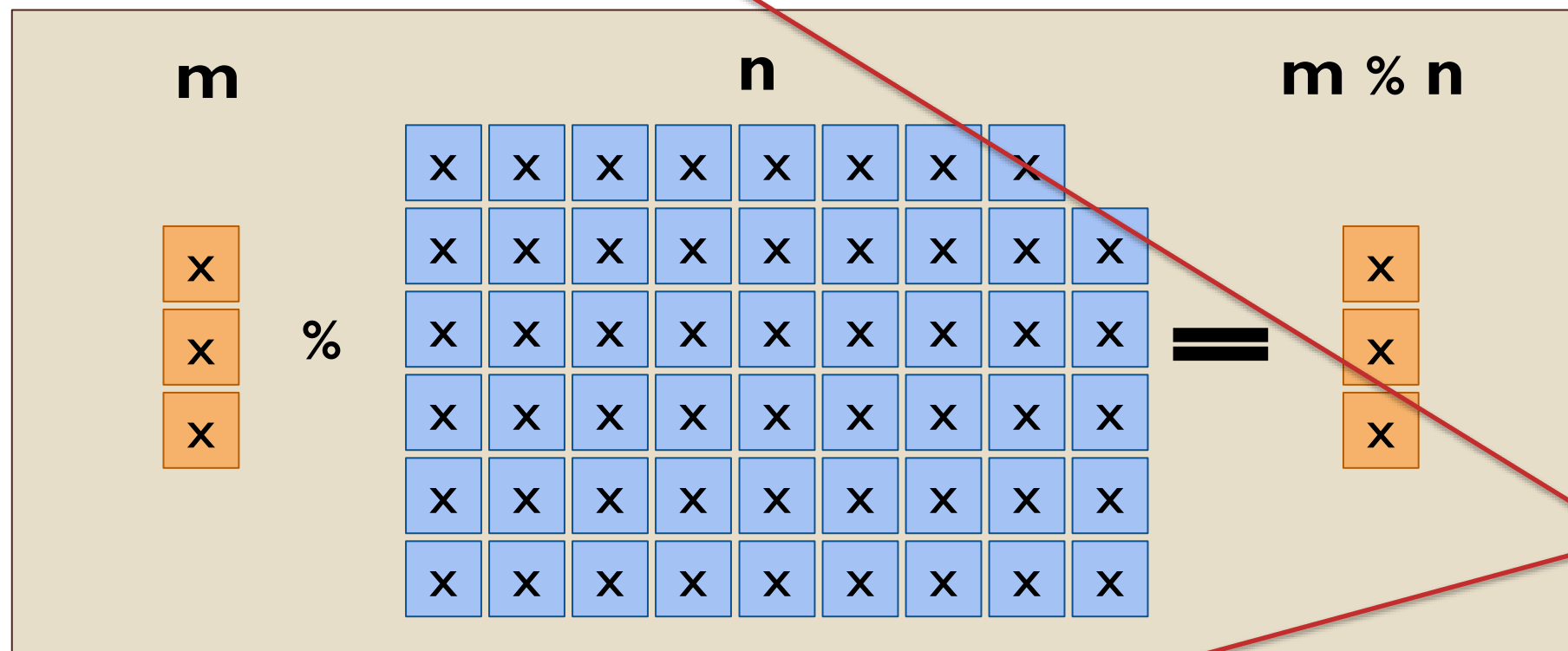
Seems we need
 $m \geq n$
for this to work.



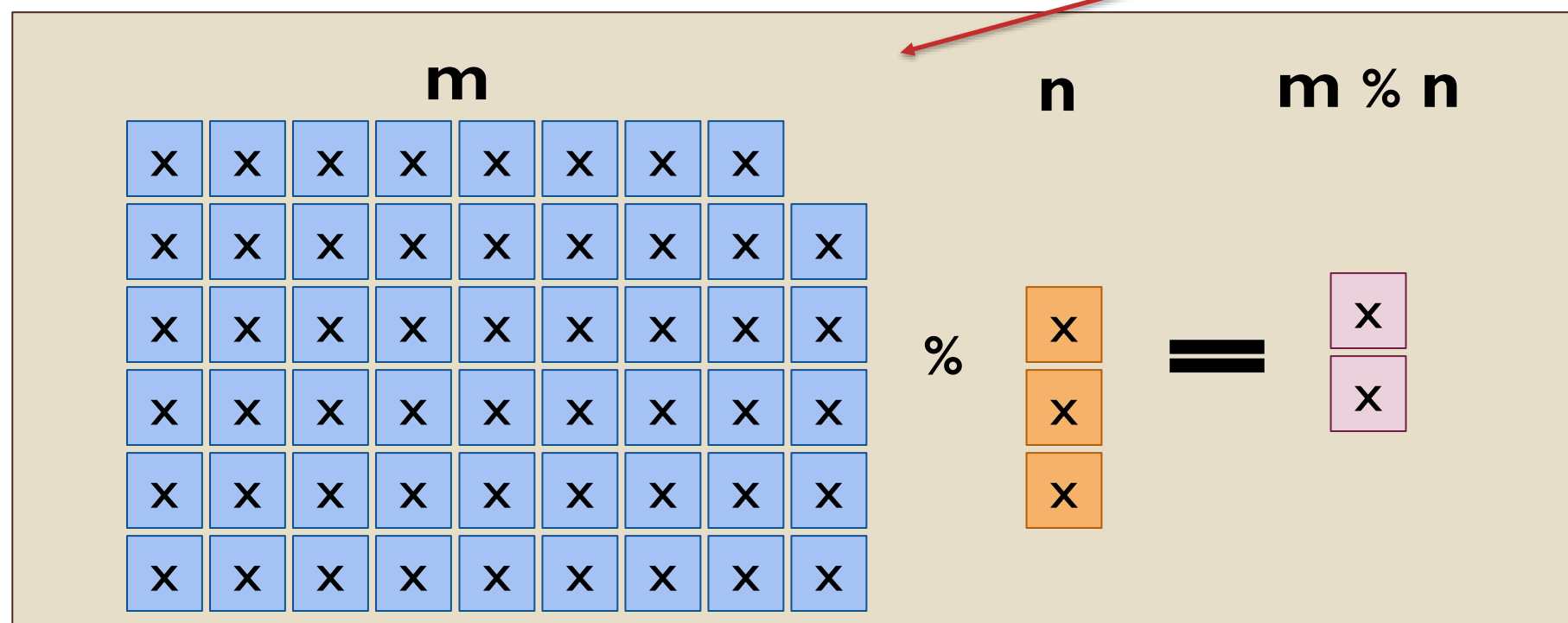
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Final problem reduction:

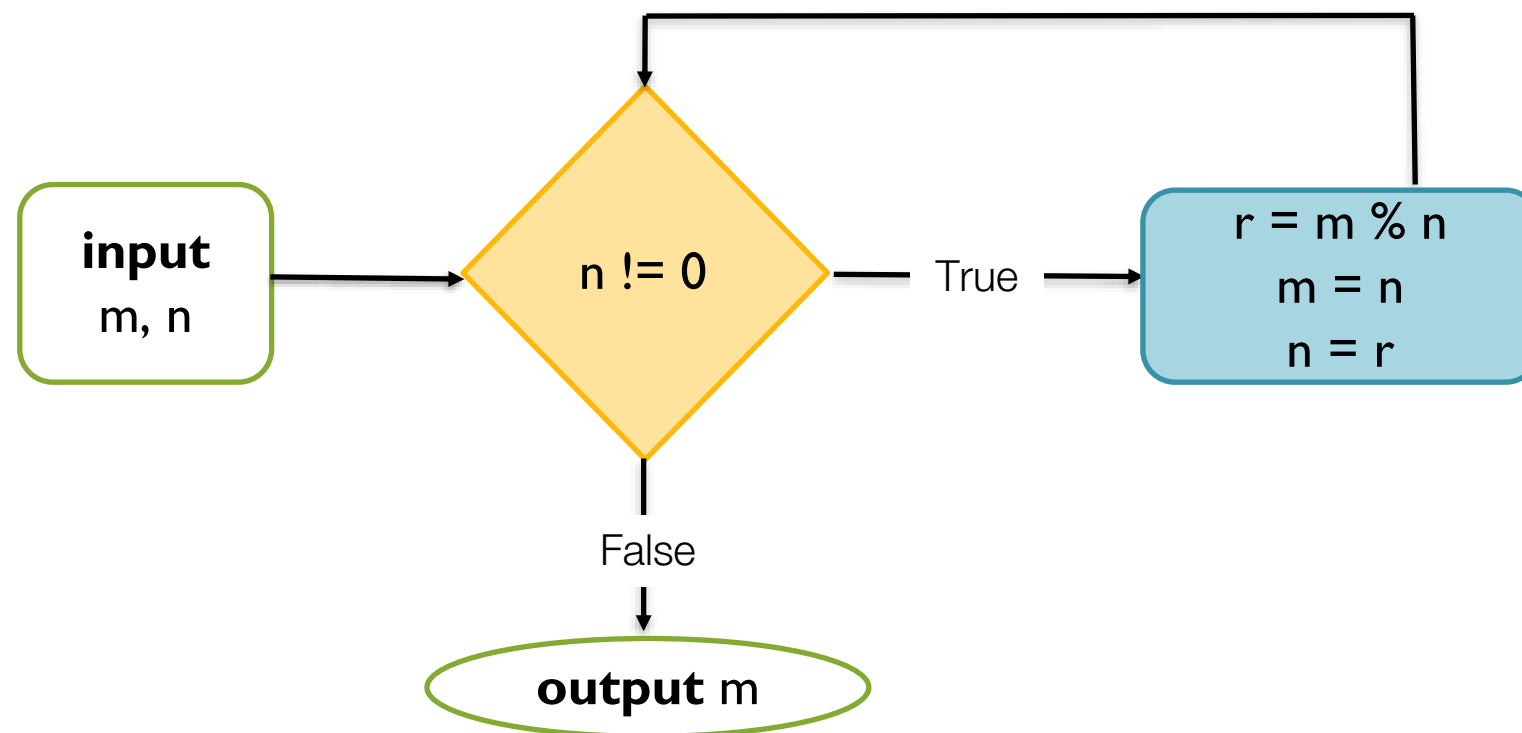
$$\gcd(m, n) = \gcd(n, m \% n)$$


But if $m \leq n$
reduction
simply flips
arguments



Thus
 $\gcd(m, n) = \gcd(n, m \% n)$
is generally
correct and
reduces
problem size
after at most
two
applications

Euclid's Algorithm



Eukleides of Alexandria
3xx BC – 2xx BC

```
def gcd(m, n):  
    """  
    Input : integers m and n such that not n==m==0  
    Output: the greatest common divisor of m and n  
    """  
    while n != 0:  
        r = m % n  
        m = n  
        n = r  
    return m
```

Recommended reading

“Introduction to Computing using Python: An Application Development Focus”, by L. Perkovic

- §2.3
- §5.3

FIT1045/53 Workbook

- Chapter 2, §2.2.1
- Chapter 3, §§3.1-3.3

Check point for this week

- By the end of this week you should be able to do the following:
- Implement Python programs to:
 - Calculate the average of a list
 - Find a given item in a list
 - Compute specific sums and products

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Coming Up

- More loops and sequence types
- Tables and matrices