EXP NO: 9 DATE:

DEVELOP THE BACK-END OF A COMPILER THAT TAKES THREE-ADDRESS CODE (TAC) AS INPUT AND GENERATES CORRESPONDING 8086 ASSEMBLY LANGUAGE CODE AS OUTPUT.

AIM:

To design and implement the back-end of a compiler that takes three-address code (TAC) as input and produces 8086 assembly language code as output. The three-address code is an intermediate representation used by compilers to break down expressions and operations, while the 8086 assembly code is a machine-level representation of the program that can be executed by a processor.

ALGORITHM:

1. Parse the Three-Address Code (TAC):

Input: Three-Address Code, which is an intermediate representation. For example:

$$t0 = b + c$$

$$t1 = t0 * d$$

$$a = t1$$

Output: 8086 assembly language code. For example:

MOV AX, [b] ; Load b into AX ADD AX, [c] ; Add c to AX MOV [t0], AX ; Store result in t0

2. Process Each TAC Instruction:

1. Initialize Registers:

- o Set up the registers in 8086 assembly (e.g., AX, BX, CX, etc.) for storing intermediate results and final outputs.
- Maintain a temporary register counter for naming temporary variables in TAC (e.g., t0, t1).
- 2. For each TAC instruction, based on its operation:
 - o Identify the components: operands and operator.
 - Choose an appropriate register (AX, BX, etc.) for storing intermediate results.
 - o If the operation involves multiple operands or temporary variables, map them to registers.
- 3. Translating TAC to 8086 Assembly:
 - Addition/Subtraction (e.g., t0 = b + c):
 - o Load operands into registers and perform the operation:

MOV AX, [b] ; Load b into AX ADD AX, [c] ; Add c to AX MOV [t0], AX ; Store result in t0

- Multiplication (e.g., t1 = t0 * d):
 - o Load operands into registers and perform the operation:

MOV AX, [t0] ; Load t0 into AX MOV BX, [d] ; Load d into BX

MUL BX; Multiply AX by BX (result in AX)

MOV [t1], AX; Store result in t1

- Assignment (e.g., a = t1):
 - o Move the value from a temporary variable to the target variable:

MOV [a], [t1]; Move value of t1 into a

- Division (e.g., t2 = b / c):
 - Division is a bit more complex due to the 8086's limitations with the DIV instruction. For example, the result might need to be stored in AX or DX:AX (if it's a 32-bit result):

MOV AX, [b] ; Load b into AX

MOV DX, 0; Clear DX (important for division)

MOV BX, [c]; Load c into BX

DIV BX; AX = AX / BX (quotient in AX, remainder in DX)

MOV [t2], AX; Store quotient in t2

- 4. Manage Memory and Registers:
 - **Variables**: Variables like a, b, c are stored in memory, so you will use memory addressing modes such as [variable_name] to access them.
 - **Temporary Variables**: Temporary variables like t0, t1, t2, etc., are stored in registers (AX, BX, etc.) or memory if there are more variables than registers available.

5. Handle Control Flow (Optional):

If the TAC contains control structures (such as loops, if-else statements, or function calls), you will need to generate labels and jump instructions in 8086 assembly.

• If Statements: For example, if $(x > 0) \{ y = 1; \}$ could generate:

MOV AX, [x] CMP AX, 0

JG positive_case; Jump if greater

JMP end if

PROGRAM:

```
#include <stdio.h>
#include <string.h>
void generateAssembly(const char* tac) {
  char result[10], op1[10], op[2], op2[10];
  // Parse the TAC instruction
  sscanf(tac, "%s = %s %s %s", result, op1, op, op2);
  // Generate assembly code based on the operator
  if (strcmp(op, "+") == 0) {
     printf("MOV AX, [%s]\n", op1);
     printf("ADD AX, [%s]\n", op2);
     printf("MOV [%s], AX\n", result);
  } else if (strcmp(op, "-") == 0) {
     printf("MOV AX, [%s]\n", op1);
     printf("SUB AX, [%s]\n", op2);
     printf("MOV [%s], AX\n", result);
  else if (strcmp(op, "*") == 0) {
     printf("MOV AX, [%s]\n", op1);
     printf("MOV BX, [%s]\n", op2);
     printf("MUL BX\n");
     printf("MOV [%s], AX\n", result);
  } else if (\text{strcmp}(\text{op, "/"}) == 0)  {
     printf("MOV AX, [%s]\n", op1);
     printf("MOV BX, [\%s]\n", op2);
     printf("DIV BX\n");
     printf("MOV [%s], AX\n", result);
  } else {
     printf("Unsupported operation: %s\n", op);
int main() {
  const char* tacInstructions[] = {
```

```
"t0 = b + c", \\ "t1 = t0 * d", \\ "a = t1" \\ \}; \\ int numInstructions = size of(tacInstructions) / size of(tacInstructions[0]); \\ for (int i = 0; i < numInstructions; i++) \{ \\ generateAssembly(tacInstructions[i]); \\ printf("\n"); \\ \} \\ return 0; \\ \}
```

OUTPUT:

```
MOV AX, [b]
ADD AX, [c]
MOV [t0], AX

MOV AX, [t0]
MOV BX, [d]
MUL BX
MOV [t1], AX

MOV AX, [t1]
MOV AX, [t1]
```

| Implementation | |
|------------------|--|
| Output/Signature | |

RESULT:

Thus the above example provides a foundational approach to converting TAC to 8086 assembly using C. For a complete compiler back-end, you would need to handle additional aspects such as register allocation, memory management, and more complex control flow constructs.