



Coq Proof Assistant: Propositions and Proofs

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Overview

Minimal Propositional Logic

Basics

Definition

Example

Definition

Demo

Structure

Tactics



Truth table

$(P \rightarrow Q)$

- Classical logic
- Assign to every variable a denotation true or false
- Formula is valid iff true in all cases
- Question "is the proposition P true?"



Coq system

$(P \rightarrow Q)$

- Intuitionistic logic
- Obtain a proof of Q from a proof of P
- Arbitrary proof of P constructs a proof of Q
- Question "what are the proof of P (if any)?"



Hypothesis

Hypothesis $h:P$

- Local declaration
- h is the name of the hypothesis
- P is its statement
- Synonymous to
- Use

Variable $h:P$

Hypotheses

or

Variables

to declare several at a time



Section

The section contains all Hypotheses / Variables from the Context

Start section sec1 with

`Section sec1`

End section sec1 with

`End sec1`



Axiom

Axiom $x:P$

- Global declaration
- Synonymous to

Parameter $x:P$

Environment contains axioms

Context contains hypotheses

$E, \Gamma \vdash \pi : P$



Goals and Tactics

Goals: what needs to be proven

Goal: $E, \Gamma \vdash P$

Construct a proof of P . Should be a well-formed term t in the environment E and context Γ

Term t is called a *solution*

Tactics: commands to decompose this goal into simpler goals

g is input goal and g_1, g_2, \dots, g_k are output goals

Possible to construct a solution of g from the solutions of goals g_i



intros

Goal: $(P \rightarrow Q) \rightarrow (Q \rightarrow R) \rightarrow (P \rightarrow R)$

`intros H H' p`

- Transform the task of construction a proof into proving R with those hypotheses added
- $H : P \rightarrow Q$
- $H' : Q \rightarrow R$ and
- $p : P$
- New subgoal: R

Simplifies the statement to prove and increases the resources available



apply

Subgoal: R

Hypothesis $H' : Q \rightarrow R$ and $H : P \rightarrow Q$

apply H'

- Use the hypothesis H' to advance our proof
- Argument has to be a premise and a conclusion
- Creates new goal for the premise
- New subgoal: Q

Applying hypothesis H gives the new subgoal P



assumption

Subgoal: P

Hypothesis $p : P$

assumption

- Statement to proof is exactly statement of hypothesis p
- Succeeds without generating any new goal

No more subgoals

Proof complete



Finish

Qed

- Saves the theorem's name, statement and proof term
- Displays the sequence of tactics.

```
intros H H' p.
apply H'.
apply H.
assumption.
```

Print theorem-name

- Shows the proof like any *Gallina* definition

```
theorem-name = fun (H:P -> Q)(H':Q -> R)(p:P) => H' (H p)
: (P -> Q) -> (Q -> R) -> P -> R
```



Transitivity

$$(P \rightarrow Q) \rightarrow (Q \rightarrow R) \rightarrow (P \rightarrow R)$$



One Shot

Not all the details for the proof is needed

A few automatic tactics are able to solve this goal

Theorem transitivity : $(P \rightarrow Q) \rightarrow (Q \rightarrow R) \rightarrow P \rightarrow R$.

Proof.

auto.

Qed.

Good usage requires the command Proof after Theorem or Lemma

Makes the proof documents more readable



exact

The exact tactic takes any term (with right type) as argument

The whole proof could be given as an argument

```
Theorem delta : (P→P→Q) → P → Q.
```

```
Proof.
```

```
exact (fun (H:P→P→Q)(p:P) => H p p).
```

```
Qed.
```

Or even shorter:

```
Theorem delta : (P→P→Q) → P → Q.
```

```
Proof (fun (H:P→P→Q)(p:P) => H p p).
```



Modus Ponens

apply tactic uses the Modus Ponens

$$\mathbf{App} \frac{E, \Gamma \vdash t : P \rightarrow Q \quad E, \Gamma \vdash t' : P}{E, \Gamma \vdash tt' : Q}$$

Term $t : P_1 \rightarrow P_2 \rightarrow \dots \rightarrow P_n \rightarrow Q$

Goal : $P_k \rightarrow P_{k+1} \rightarrow \dots \rightarrow P_n \rightarrow Q$

apply t

generates $k-1$ goals with statements P_1, \dots, P_{k-1}

if new goals have solution t_1, t_2, \dots, t_{k-1}

solution is $t \ t_1 \ t_2 \ \dots \ t_{k-1}$ for initial goal



intros

`intros v_1, v_2, ..., v_n`

is the same as

`intro v_1, intro v_2, ..., intro v_n`

`intro v_1` takes the first implication as a hypothesis called v_1

Proof Theorem K : $P \rightarrow Q \rightarrow P$

Goal is $P \rightarrow Q \rightarrow P$

`intro p.`

Hypothesis $p : P$ and new Goal $Q \rightarrow P$



Handling

Show i

Display goal i with complete context

Coq displays the goals after each proof step

Undo n

Go back n steps and try an alternative if goal can not be solved

Restart

Go back to the beginning of the proof

Abort

Abandon the proof

