# **Today**

### Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

#### Structures

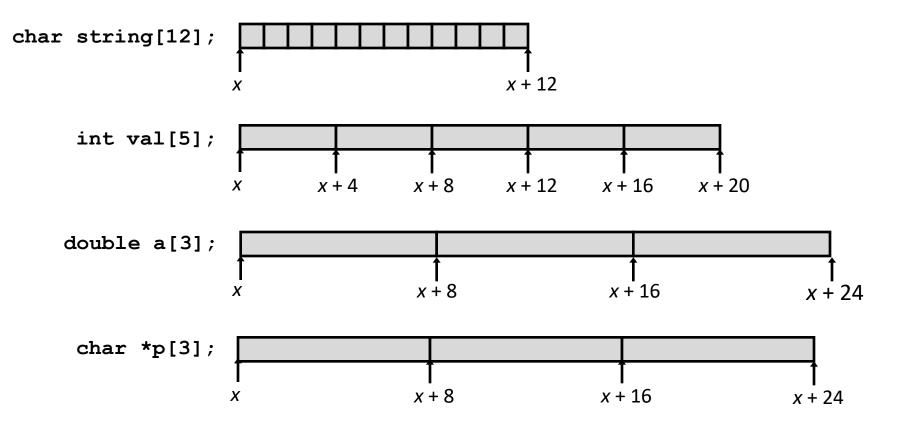
- Allocation
- Access
- Alignment
- Floating Point

# **Array Allocation**

### Basic Principle

```
T A[L];
```

- Array of data type T and length L
- Contiguously allocated region of L \* sizeof (T) bytes in memory

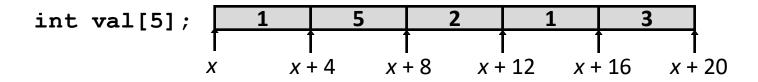


### **Array Access**

### Basic Principle

```
T A[L];
```

- Array of data type T and length L
- Identifier A can be used as a pointer to array element 0: Type T\*

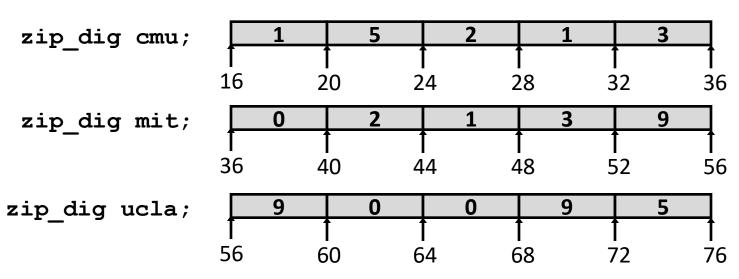


Reference	Type	Value
val[4]	int	
val	int *	
val+1	int *	
&val[2]	int *	
<b>val</b> [5]	int	
*(val+1)	int	
val + <i>i</i>	int *	

# **Array Example**

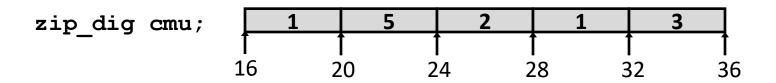
```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucla= { 9, 0, 0, 9, 5 };
```



- Declaration "zip\_dig cmu" equivalent to "int cmu[5]"
- Example arrays were allocated in successive 20 byte blocks
  - Not guaranteed to happen in general

# **Array Accessing Example**



```
int get_digit
  (zip_dig z, int digit)
{
  return z[digit];
}
```

```
# %rdi = z
# %rsi = digit
movl (%rdi,%rsi,4), %eax # z[digit]
```

- Register %rdi contains starting address of array
- Register %rsi contains array index
- Desired digit at %rdi + 4\*%rsi
- Use memory reference
  (%rdi,%rsi,4)

# **Array Loop Example**

```
void zincr(zip_dig z) {
   size_t i;
   for (i = 0; i < ZLEN; i++)
      z[i]++;
}</pre>
```

```
# %rdi = z
                         \# i = 0
 movl $0, %eax
                         # goto middle
 jmp .L3
                         # loop:
.L4:
 addl $1, (%rdi,%rax,4) # z[i]++
 addq $1, %rax
                         # 1++
.L3:
                         # middle
 cmpq $4, %rax
                         # i:4
                         # if <=, goto loop</pre>
 jbe .L4
 rep; ret
```

# Multidimensional (Nested) Arrays

#### Declaration

```
T \mathbf{A}[R][C];
```

- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

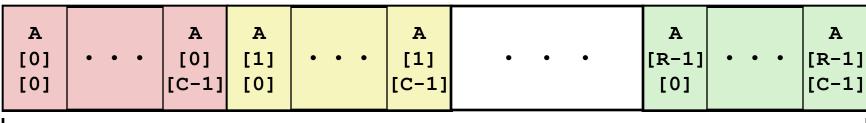
### Array Size

R \* C \* K bytes

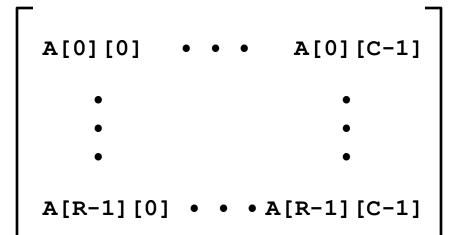
### Arrangement

Row-Major Ordering

### int A[R][C];

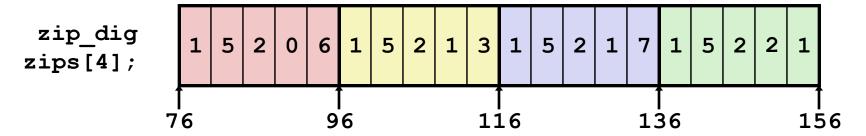


4\*R\*C Bytes



# **Nested Array Example**

```
#define PCOUNT 4
zip_dig zips[PCOUNT] =
   {{1, 5, 2, 0, 6},
    {1, 5, 2, 1, 3},
    {1, 5, 2, 1, 7},
    {1, 5, 2, 2, 1 }};
```



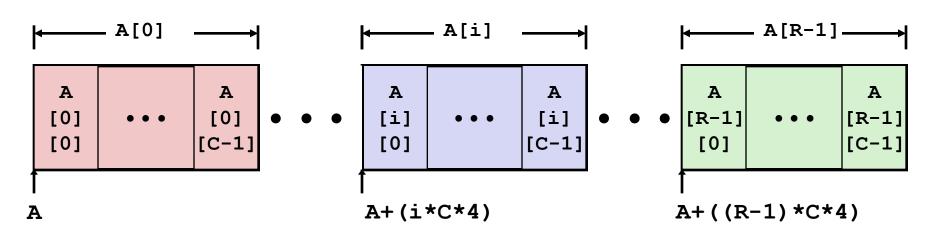
- "zip\_dig zips[4]" equivalent to "int zips[4][5]"
  - Variable zips: array of 4 elements, allocated contiguously
  - Each element is an array of 5 int's, allocated contiguously
- "Row-Major" ordering of all elements in memory

# **Nested Array Row Access**

#### Row Vectors

- A[i] is array of C elements
- Each element of type T requires K bytes
- Starting address A + i\* (C\* K)

#### int A[R][C];



# **Nested Array Row Access Code**

```
1 5 2 0 6 1 5 2 1 3 1 5 2 1 7 1 5 2 2 1

zips

int *get_zips_zip(int index)
{
    return zips[index];
}
```

```
# %rdi = index
leaq (%rdi,%rdi,4),%rax # 5 * index
leaq zips(,%rax,4),%rax # zips + (20 * index)
```

#### Row Vector

- zips[index] is array of 5 int's
- Starting address zips+20\*index

#### Machine Code

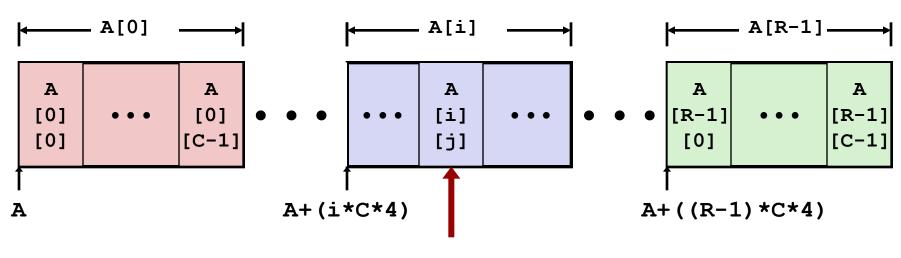
- Computes and returns address
- Compute as zips + 4\*(index+4\*index)

# **Nested Array Element Access**

### Array Elements

- A[i][j] is element of type T, which requires K bytes.
- Address **A** + i \* (C \* K) + j \* K = A + (i \* C + j) \* K

int A[R][C];



$$A+(i*C*4)+(j*4)$$

# **Nested Array Element Access Code**

```
leaq (%rdi,%rdi,4), %rax  # 5*index
addl %rax, %rsi  # 5*index+dig
movl zips(,%rsi,4), %eax  # M[zips + 4*(5*index+dig)]
```

### Array Elements

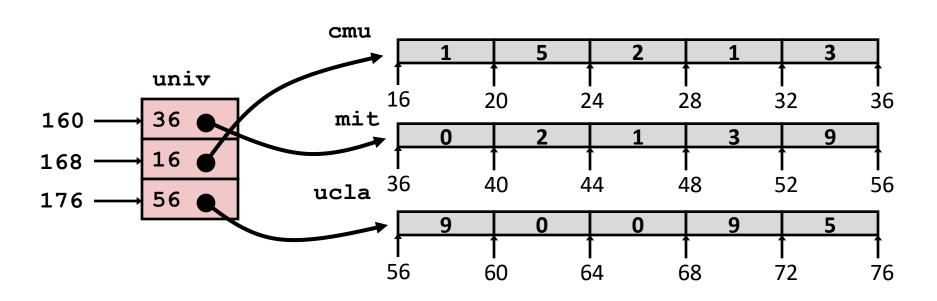
- zips[index][dig] is int
- Address: zips + 20\*index + 4\*dig
  - = zips + 4\*(5\*index + dig)

# **Multi-Level Array Example**

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucla = { 9, 0, 0, 9, 5 };
```

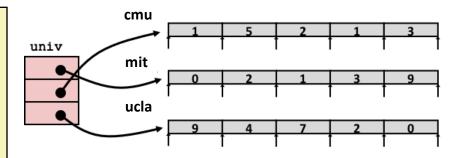
```
#define UCOUNT 3
int *univ[UCOUNT] = {mit, cmu, ucla};
```

- Variable univ denotes array of 3 elements
- Each element is a pointer
  - 8 bytes
- Each pointer points to array of int's



### **Element Access in Multi-Level Array**

```
int get_univ_digit
  (size_t index, size_t digit)
{
  return univ[index][digit];
}
```



```
salq $2, %rsi # 4*digit
addq univ(,%rdi,8), %rsi # p = univ[index] + 4*digit
movl (%rsi), %eax # return *p
ret
```

### Computation

- Element access Mem [Mem [univ+8\*index]+4\*digit]
- Must do two memory reads
  - First get pointer to row array
  - Then access element within array

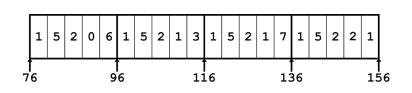
# **Array Element Accesses**

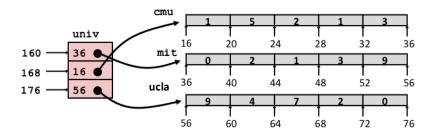
#### **Nested array**

```
int get_zips_digit
   (size_t index, size_t digit)
{
   return zips[index][digit];
}
```

#### Multi-level array

```
int get_univ_digit
  (size_t index, size_t digit)
{
  return univ[index][digit];
}
```





Accesses looks similar in C, but address computations very different:

Mem[zips+20\*index+4\*digit] Mem[Mem[univ+8\*index]+4\*digit]

# N X N Matrix Code

- Fixed dimensions
  - Know value of N at compile time
- Variable dimensions, explicit indexing
  - Traditional way to implement dynamic arrays
- Variable dimensions, implicit indexing
  - Now supported by gcc

### 16 X 16 Matrix Access

### Array Elements

- Address **A** + i \* (C \* K) + j \* K
- C = 16, K = 4

```
/* Get element a[i][j] */
int fix_ele(fix_matrix a, size_t i, size_t j) {
  return a[i][j];
}
```

```
# a in %rdi, i in %rsi, j in %rdx
salq $6, %rsi  # 64*i
addq %rsi, %rdi  # a + 64*i
movl (%rdi,%rdx,4), %eax # M[a + 64*i + 4*j]
ret
```

### n X n Matrix Access

### Array Elements

- Address **A** + i \* (C \* K) + j \* K
- C = n, K = 4
- Must perform integer multiplication

```
/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n], size_t i, size_t j)
{
  return a[i][j];
}
```

```
# n in %rdi, a in %rsi, i in %rdx, j in %rcx
imulq %rdx, %rdi  # n*i
leaq (%rsi,%rdi,4), %rax # a + 4*n*i
movl (%rax,%rcx,4), %eax # a + 4*n*i + 4*j
ret
```

# **Today**

### Arrays

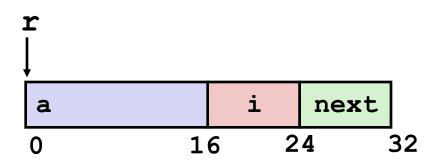
- One-dimensional
- Multi-dimensional (nested)
- Multi-level

#### Structures

- Allocation
- Access
- Alignment
- Floating Point

### **Structure Representation**

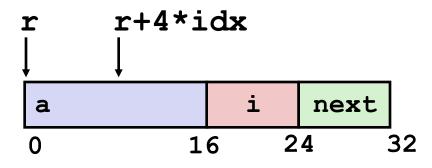
```
struct rec {
   int a[4];
   size_t i;
   struct rec *next;
};
```



- Structure represented as block of memory
  - Big enough to hold all of the fields
- Fields ordered according to declaration
  - Even if another ordering could yield a more compact representation
- Compiler determines overall size + positions of fields
  - Machine-level program has no understanding of the structures in the source code

# **Generating Pointer to Structure Member**

```
struct rec {
   int a[4];
   size_t i;
   struct rec *next;
};
```



### Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Compute as r + 4\*idx

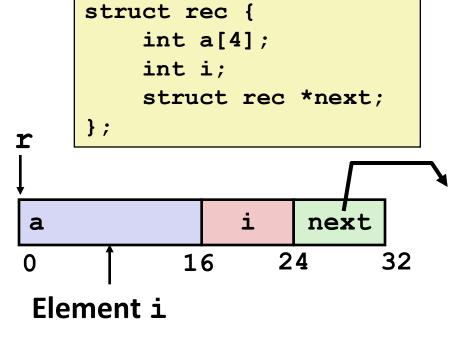
```
int *get_ap
  (struct rec *r, size_t idx)
{
   return &r->a[idx];
}
```

```
# r in %rdi, idx in %rsi
leaq (%rdi,%rsi,4), %rax
ret
```

# **Following Linked List**

C Code

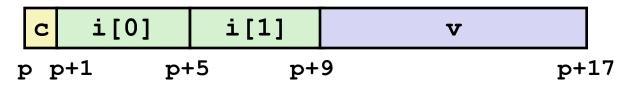
```
void set_val
  (struct rec *r, int val)
{
  while (r) {
    int i = r->i;
    r->a[i] = val;
    r = r->next;
  }
}
```



Register	Value
%rdi	r
%rsi	val

# **Structures & Alignment**

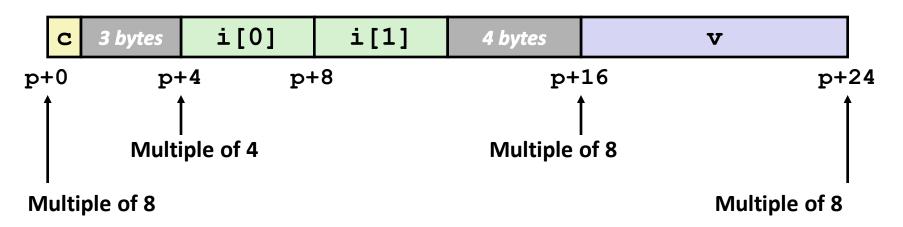
### Unaligned Data



```
struct S1 {
  char c;
  int i[2];
  double v;
} *p;
```

### Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K



# **Alignment Principles**

### Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K
- Required on some machines; advised on x86-64

### Motivation for Aligning Data

- Memory accessed by (aligned) chunks of 4 or 8 bytes (system dependent)
  - Inefficient to load or store datum that spans quad word boundaries
  - Virtual memory trickier when datum spans 2 pages

### Compiler

Inserts gaps in structure to ensure correct alignment of fields

# **Specific Cases of Alignment (x86-64)**

- 1 byte: char, ...
  - no restrictions on address
- 2 bytes: short, ...
  - lowest 1 bit of address must be 02
- 4 bytes: int, float, ...
  - lowest 2 bits of address must be 002
- 8 bytes: double, long, char \*, ...
  - lowest 3 bits of address must be 0002
- 16 bytes: long double (GCC on Linux)
  - lowest 4 bits of address must be 00002

# Satisfying Alignment with Structures

#### Within structure:

Must satisfy each element's alignment requirement

### Overall structure placement

- Each structure has alignment requirement K
  - **K** = Largest alignment of any element
- Initial address & structure length must be multiples of K

### Example:

K = 8, due to double element

```
        c
        3 bytes
        i [0]
        i [1]
        4 bytes
        v

        p+0
        p+4
        p+8
        p+16
        p+24

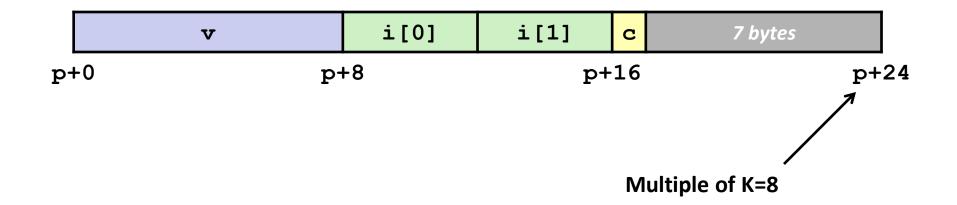
        Multiple of 4
        Multiple of 8
        Multiple of 8

Multiple of 8
```

# **Meeting Overall Alignment Requirement**

- For largest alignment requirement K
- Overall structure must be multiple of K

```
struct S2 {
  double v;
  int i[2];
  char c;
} *p;
```



# **Arrays of Structures**

- Overall structure length multiple of K
- Satisfy alignment requirement for every element

```
struct S2 {
  double v;
  int i[2];
  char c;
} a[10];
```



