Dynamic programming

July 10, 2022

```
\label{eq:https://www.youtube.com/watch?v=oBt53YbR9Kk} $$\# $$ memoization
```

write a function that takes the position n and returns the nth fib no basic wrong/inneficient implementation

the below program has a time complexity of O(2^n)

the space complexity is O(n)

because he function calls happen only one at a time in LNR

```
def fib(n):
    if n <= 2:
        return 1
        return fib(n - 1) + fib(n - 2)

print(fib(6))
print(fib(7))
print(fib(8))
print(fib(9))
# fib(50) takes forever to compute dont even try
# this takes 2~50 steps</pre>
```

```
8
```

13

21

34

better implementaion using memoization

```
time complexity O(2n) \rightarrow O(n)
```

space complexity \rightarrow O(n)

```
[]: def fib(n, memo={}):
    if n in memo:
        return memo[n]
    if n <= 2:
        return 1
        memo[n] = fib(n - 1, memo) + fib(n - 2, memo)
        return memo[n]

print(fib(2009))

"""

dict = {}
print(fib(50, dict))

print(dict)
"""</pre>
```

 $32113249982681582845238472032097086207406078264369213828207269738873528372164658\\13601607624368278410255618945484808988900205426234968094236646697230669809316249\\08775724739402394881483009611433054415348267177640331222672835631629564907858584\\30785050847615157489167011704773920617812432061755619769908742336741418110524930\\07153851287942959207659531735086437334610990020876787656487609216126565715584580\\25510084617009128909$

```
[]: '\ndict = {}\nprint(fib(50,dict))\n\nprint(dict)\n'
```

find the number of ways you can travel corner to corner in m*n grid provided you can only move right and down

finction takes in m and n rows and columns time complexity = $O(2^n+m)$ space complexity = O(n+m)

```
[]: def gridtravel(m, n, memo={}):
    if m == 1 and n == 1:
        return 1
    elif m == 0 or n == 0:
        return 0

    return gridtravel(m - 1, n) + gridtravel(m, n - 1)
```

[]: 184756

time complexity m*n space complexity n+m

```
[]: dict = {}
# joels better code
def gridtravel(m, n, memo={}):
    l = tuple(sorted((m, n)))
    if l in memo:
        return memo[l]
    elif 1 in l:
        return 1
    else:
        memo[l] = gridtravel(m - 1, n, memo) + gridtravel(m, n - 1, memo)
        return memo[l]

gridtravel(1001, 1001)
```

[]: 20481516269894897143351625029808250443964248879813970338203826376717481862020837 55828932994182610206201464766319998023692415481798004524792018047549769261578563 01289663432064714851152395251651227768588611539546256147907378668464154444533617 61377007385567381458963007130651045595951447988874620636871851455182855117316627 62536637730846829322553890497438594814317550307837964443708100851637248274627914 17016619883764840843541430817785947037746565188475514680749694674923803033101818 72329800966856745856025254991011811352535346588879419666536749045113061100963119 06270342502293155911108976733963991149120

memoization recipe

- 1. make it work (with recursion look for correctness)(this is harder than step 2) visualize as trees implement the tree using recursion test it
- 2. make it efficient add a memo object (hash object/dict) add a base case to return memo values store the return values into the memo

0.0.1 cansum()

```
can
sum(targetsum, numbers) eg, can
sum(7,[5,3,4,7]) -> true 3 + 4  7  can
sum(7,[2,4]) -> false
```

```
[]: # joels noobie implementation

def cansum(m, n, d={}):
    j = False
```

```
if m in d:
    return d[m]
if m == 0:
    return True
if m < 0:
    return False
for i in n:
    j = j or cansum(m - i, n, d)
    d[m] = j
    return d[m]</pre>
```

[]: True

 $\label{eq:mass_mass} \begin{aligned} &\text{time} = n \hat{\ } m \text{ space} = m \\ &\text{where } n \text{ is lenghth of array and } m \text{ is target sum} \\ &\text{with memo time} = n \hat{\ } m \end{aligned}$

```
[]: def cansum(m, n, d={}):
    if m in d:
        return d[m]
    if m == 0:
        return True
    if m < 0:
        return False
    for i in n:
        if cansum(m - i, n, d) == True:
            d[m] = True
            return True
    d[m] = False
    return False

cansum(1071, [33, 4])</pre>
```

[]: True

```
[]: def howsum(m, n, d={}):
    if m == 0:
        return []
    if m < 0:</pre>
```

```
return None

for i in n:
    if howsum(m - i, n) != None:
        return howsum(m - i, n) + [i]
    return None

print(howsum(300, [2, 4]))
```

with a memo time = $n*m^2$ space = m

```
[]: def howsum(m, n, d={}):
    if m in d:
        return d[m]
    if m == 0:
        return []
    if m < 0:
        return None

for i in n:
        if howsum(m - i, n, d) != None:
            d[m] = howsum(m - i, n, d) + [i]
        return d[m]
    d[m] = None
    return None

print(howsum(97, [2, 3, 30, 80, 99]))</pre>
```

[2, 3, 3, 3, 3, 80]

bestsum() return array with shortest combination of numbers that adds up to target sum without memo

```
timw = n^m * m \text{ space} = m * m = m^2 with memo time = m^* m * n \text{ space} = m^2
```

```
[]: def bestsum(m, n, d={}):
    if m in d:
        return d[m]
    if m == 0:
        return []
    if m < 0:
        return None</pre>
```

```
shortest = None
for i in n:
    if bestsum(m - i, n, d) != None:
        c = bestsum(m - i, n, d) + [i]
        if (shortest == None) or (len(c) < len(shortest)):
            shortest = c

d[m] = shortest
    return shortest

print(bestsum(100, [2, 33, 70, 1]))</pre>
```

[1, 33, 33, 33]

canconstruct(target,wordbank) accepts a target string and array of strings

should return a boolean indicating whether or not the target can be constructed by concatenating the elements of the word bank

you may reuse elements from the word bank

without memoization time complexity = $n^m m$ (slicinf gives an extra m) space complexity = m^2 where m = target length n = wordbakn length

```
[]: def canconstruct(m, n):
       if m == "":
           return True
       for i in n:
           if i in m:
              if m.index(i) == 0:
                  if canconstruct(m[len(i) :], n) == True:
                     return True
       return False
    print(canconstruct("abcdef", ["ab", "def", "cd", "cdef"]))
    print(
       canconstruct(
           eeeeeeeeeef",
           ["e", "ee", "eeee", "eeeee", "eeeeee", "eeeeeee"],
       )
    )
```

False False

)

print(

)

with memo

0.0.2 countconstruct()

canconstruct(

eeeeeeeeef",

find the total nuber of ways the target string can be constructed

["e", "ee", "eeee", "eeeee", "eeeeee", "eeeeeee"],

1 256

0.0.3 allconstruct()

return a 2d array containing all the possible ways the string cant be constructed blah bla random test code snippets below

```
[]: y = 3
x = 5
print(x + y)
```

8

```
[]: s1 = "AJYFAJYF"
s2 = "JY"
import re

if s1.startswith(s2):
        s3 = re.sub("^" + s2, "", s1)
s3
```

[]: 'YFAJYF'

```
[]: s1 = "AJYFAJYF"
s2 = "JY"
s1.index(s2)
s2 in s1
```

[]: True

```
[]: # for i in range(9):
     # print(i)
     item2 = "hello"
     print(type(item2))
     class item:
        h = 0
        def gun(self):
            pass
     print("koooooool")
     item1 = item()
     item1.name = "joel"
     print(type(item1))
    print(type(item1.gun))
    <class 'str'>
    kooooool
    <class '__main__.item'>
    <class 'method'>
[]: h = "hello"
     v = "varghese"
    print(h, "joel " + v)
    hello joel varghese
[]: h = "hello"
     v = "varghese"
    print(h, "joel " + v)
    hello joel varghese
```