

Practical Concurrent and Parallel Programming IX

Performance and Scalability

Jørgen Staunstrup

Agenda



- Executor and Future
- Scalability, speed-up and loss (of scalablity) classification
 Example: QuickSort
- Lock striping
 - A case study with Hash maps

Prime counter task (skeleton)

```
5
```

```
public class countPrimesTask implements Runnable {
 private final int low;
                                            [low..high]
  private final int high;
  private final ExecutorService pool;
  @Override public void run() {
    int mid= low+(high-low)/2;
    pool.submit( new countPrimesTask(low, mid, pool) );
    pool.submit( new countPrimesTask(mid+1, high, pool) );
               Shortcomings:
                1. How to stop?
               2. Will create too many "small" tasks
               3. Returning result (# primes)
```

Splitting tasks



```
@Override
  public void run() {
    if ((high-low) < threshold) {
        for (int i=low; i<=high; i++) if (isPrime(i)) lc.increment();
    } else {
        int mid= low+(high-low)/2;
        pool.submit(new countPrimesTask(lc, low, mid, pool, threshold) );
        pool.submit(new countPrimesTask(lc, mid+1, high, pool, threshold) );
}
</pre>
```

1. How to stop?

Shortcomings:

- 2. Will create too many "small" tasks
- 3. Returning result (# primes)

Future



Futures are executors that deliver a result



```
T2:
  public Future<Integer> calculate(Integer input) {
    return executor.submit(() -> {
      ... // compute result
      return result;
    });
T1:
  Future<Integer> future = T2.calculate( ); // start
  future.get();
                                 get() is a blocking call !!!!
```

Splitting tasks

```
**
```

```
public void run() {
  if ((high-low) < threshold) {</pre>
    } else {
      int mid= low+(high-low)/2;
      Future<?> f1= pool.submit( new countPrimesTask(lc, low, mid,
                                  pool, threshold) );
      Future<?> f2= pool.submit( new countPrimesTask(lc, mid+1, high,
                                  pool, threshold) );
```

But how do we get the results from a future?

```
T_1
t_1
t_2
t_3
t_4
t_4
t_5
t_7
t_8
```

Combining tasks

```
9
```

```
public void run() {
  if ((high-low) < threshold) { ...
  } else {
    int mid= low+(high-low)/2;
    Future<?> f1= pool.submit( new countPrimesTask(lc, low, mid, pool, threshold) );
    Future<?> f2= pool.submit( new countPrimesTask(lc, mid+1, high, pool, threshold) );
    try { f1.get();f2.get(); }
    catch (InterruptedException | ExecutionException e) { }
}
```

1: Does the order of f1.get and f2.get matter?

Splitting tasks

```
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```

```
public void run() {
  if ((high-low) < threshold) {</pre> .
    } else {
      int mid= low+(high-low)/2;
      Future <?> f1= pool.submit( new countPrimesTask(lc, low, mid,
                                 pool, threshold) );
      Future<?> f2= pool.submit( new countPrimesTask(lc, mid+1, high,
                                 pool, threshold) );
      try { f1.get(); f2.get(); }
      catch (InterruptedException | ExecutionException e) {
       Shortcomings:
```

1. How to stop?

- 2. Will create too many "small" tasks
- 3. Returning result (# primes)

Combining tasks

```
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```

```
public void run() {
  if ((high-low) < threshold) {</pre>
    } else {
      int mid= low+(high-low)/2;
      Future<?> f1= pool.submit( new countPrimesTask(lc, low, mid,
                                pool, threshold) );
      Future<?> f2= pool.submit( new countPrimesTask(lc, mid+1, high,
                                pool, threshold) );
      try { f1.get();f2.get(); }
      catch (InterruptedException | ExecutionException e) {     }
                                       How do we get the result
       Shortcomings:
                                               # primes?
       1. How to stop?
       2. Will create too many "small" tasks
       3. Returning result (# primes)
```

Counting the primes

```
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```

```
public void run() {
  if ((high-low) < threshold) { ... lc.increment();</pre>
    } else {
      int mid= low+(high-low)/2;
      Future<?> f1= pool.submit( new countPrimesTask(lc, low, mid,
                                pool, threshold) );
      Future<?> f2= pool.submit( new countPrimesTask(lc, mid+1, high,
                                pool, threshold) );
      try { f1.get();f2.get(); }
      catch (InterruptedException | ExecutionException e) {     }
       Shortcomings:
       1. How to stop?
       2. Will create too many "small" tasks
```

3. Returning result (# primes) Final value in Ic

countPrimesTask



```
public class countPrimesTask implements Runnable {
 private final int low;
 private static boolean isPrime(int n) {
 public countPrimesTask(PrimeCounter lc, int low, int high,
   ExecutorService pool, int threshold) {
    this.lc = lc;
  @Override
 public void run() {
```

Java Executors - Tasks



Ideally, tasks should be independent i.e. not update shared variables

```
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```

```
public class countPrimesTask implements Runnable
private final int low; private final int high
private final ExecutorService pool;
private final PrimeCounter lc;

private static boolean isPrime(int n) {...)
}

public countPrimesTask(PrimeCounter lc...) {this.lc= lc;...
@Override
```

2: Are countPrimesTasks independent?

```
public countPrimesTask(PrimeCounter lc...) {this.lc= lc;...}

@Override
public void run() {
   if ((high-low) < threshold) {
      for (int i=low; i<=high; i++) if (isPrime(i)) lc.increment();
} else {
   int mid= low+(high-low)/2;
   Future<?> f1= pool.submit( new countPrimesTask(lc, ...));
   Future<?> f2= pool.submit( new countPrimesTask(lc, ...));
   try { f1.get();f2.get();
   }
   catch (InterruptedException | ExecutionException e) { e.printStackTrace();
}
}
Code in countPrimesTask.java
```

```
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```

```
public class countPrimesTask implements Runnable {
  private final int low; private final int high; priva
  private final ExecutorService pool;
                                                             Not really, they all
  private final PrimeCounter lc;
                                                                   update lc
  private static boolean isPrime(int n) {...)
  public countPrimesTask(PrimeCounter lc...) {this.lc= lc;... }
  @Override
  public void run() {
    if ((high-low) < threshold) {</pre>
      for (int i=low; i<=high; i++) if (isPrime(i)) lc.increment();</pre>
   } else {
      int mid= low+(high-low)/2;
      Future<?> f1= pool.submit( new countPrimesTask(lc, ... ) );
      Future<?> f2= pool.submit( new countPrimesTask(lc, ... ) );
      try { f1.get();f2.get();
      catch (InterruptedException | ExecutionException e) {    e.printStackTrace();
```

PrimeCounter



```
class PrimeCounter {
 private int count= 0;
  public synchronized void increment() {
    count= count + 1;
  public synchronized int get() {
    return count;
  public synchronized void setZero() {
    count= 0;
```

- •19
- Kick-off class for the program
- It initializes the Executor service

```
class PrimeCountExecutor {
   private ExecutorService pool;
    ...
   public PrimeCountExecutor () {
      pool= new ForkJoinPool();
      Future<?> done= pool.submit(new countPrimesTask( ... ));

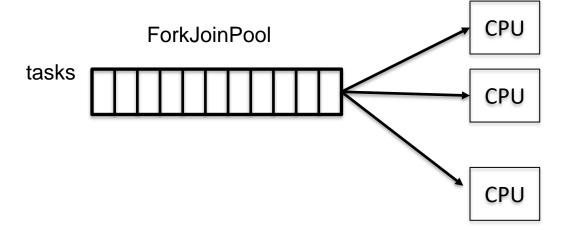
      try { done.get(); }
   }
}
```

https://docs.oracle.com/javase/tutorial/essential/concurrency/forkjoin.html

Thread pools

```
• 20
```

```
class PrimeCountExecutor {
   private ExecutorService pool;
   ...
   public PrimeCountExecutor () {
      pool= new ForkJoinPool();
      ...
}
```



ForkJoinPool



https://docs.oracle.com/javase/8/docs/api/java/util/concurrent/ForkJoinPool.html

Count prime results



Executor	Threads			
1 120.6 s	1 126.7 s			
2 68.0 s	2 82.4 s			
4 37.7 s	4 47.7 s			
8 32.2 s	8 38.2 s			
16 32.4 s	16 37.2 s			

1/8: 3.9

range 2.. 1_000_000

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1/8: 3.7

Agenda



- Executors and Future
- Scalability, speed-up and loss (of scalablity) classification
 Example: QuickSort
- Lock striping
 - A case study with Hash maps

Quicksort



1	2	43	78	19	54	33	21	64	52	17	53
1	2	43	78	19	54	33	21	64	52	17	53
1	2	43 †	78	19	54	33	21	64	52	17	53
1	2	17 †	78	19	54	33	21	64	52	43	53
1	2	17	78 †	19	54	33	21 †	64	52	43	53
1	2	17	21	19	33	54	78	64	52	43	53

17 21 19 33 54 78 64 52 43 53

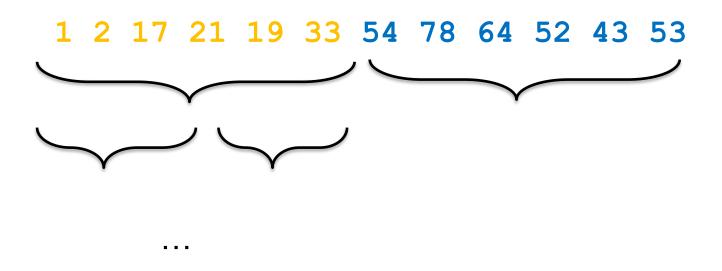
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Two parts can be

sorted independently

Distributing work as tasks





No further splitting when the sorting problem is smaller than a threshold (similarly to what we did for prime counting)

These tasks may differ in size!!

Quicksort executor (pseudo code)

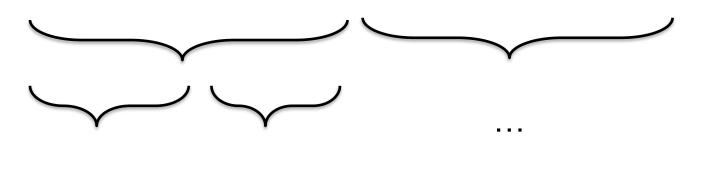
```
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```

```
class QuicksortTask implements Runnable {
  Task p; // low and high boundaries
 ExecutorService pool;
  @Override public void run() { gsort(p, pool, ...); }
 public static void qsort(Task p, ExecutorService pool, ...) {
     //split task in two: Low and High as shown on previous slides
    if (Low.size>= threshold) pool.submit( new QuicksortTask( pLow, pool, ... ))
     else Quicksort(pLow); //sequential sort
    if (High.size>= threshold) pool.submit( new QuicksortTask( pHigh, pool, ... ))
     else Quicksort(pHigh);
```

Code in QuickSortTask.java

Termination (Quicksort)

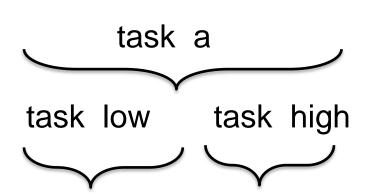






How do we know when all task are done?

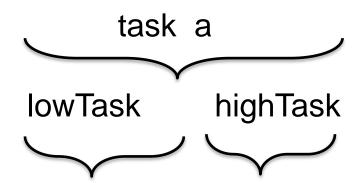
Termination



3: When can task a finish?

Termination

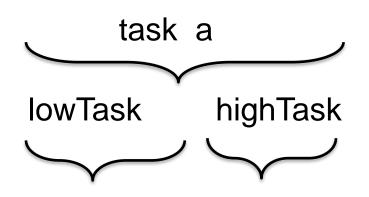




task a can finish when
both task low and task high
have finished

Termination





```
lowTask.get() returns when lowTask
has finished
highTask.get() returns when highTask
has finished
```



```
@Override
public void run() { // modified Quicksort using Executor tasks
  int a= low; int b= high;
  Future<?> lowTask= null; Future<?> highTask= null;
  if (a < b) { ... // split array in two independent part
    if ((j-a)>= threshold)
       lowTask= pool.submit(new QuickSortTask(arr, a, j, pool, threshold));
    else // all remaining work done without starting more tasks
       SearchAndSort.qsort(arr, a, j);
    if ((b-i)>= threshold)
       highTask= pool.submit(new QuickSortTask(arr, i, b, pool, threshold));
    else // all remaining work done without starting more tasks
       SearchAndSort.gsort(arr, i, b);
  //Waiting for longest running subtask to finish
  try {
    if (lowTask != null ) lowTaskF.get();
    if (highTask != null) highTaskF.get();
   } catch (InterruptedException | ExecutionException e) { e.printStackTrace(); }
```

Goetz p. 224



We use Mark8Setup to measure runtime

```
private static void runSize(ExecutorService pool, int pSize, int threshold, int n) {
  final int[] intArray= fillIntArray(pSize);
  Benchmark.Mark8Setup ("Quicksort Executor", String.format("%2d", n),
    new Benchmarkable() {
      public void setup() {
        shuffle(intArray);
      public double applyAsDouble(int i) {
        Future<?> done= pool.submit(new QuickSortTask(intArray, 0, pSize-1, pool, threshold));
        PoolFinish (done);
        //testSorted(intArray); //only needed while testing
        return 0.0;
```

Code in PoolSortingBenchmarkable.java

Sorting results

Does not scale perfectly



E	xecutor
4	~ -

- $8.5\,\mathrm{S}$ 2 4.8 s
- 2.6 s
- 8 2.2 s
- 16 2.2 s

Speed-up

- - 1/8: 3.9

- - 4 3.8 s
 - 8 3.2 s 16 3.5 s

Threads

2 6.4 s

11.2 s

- Speed-up
- 1/8: 3.9

Sorting 100_000 numbers

What limits scalability?



Example: growing a crop

- 4 months growth + 1 month harvest if done by 1 person
- Growth (sequential) cannot be speeded up
- Using 30 people to harvest, takes 1/30 month = 1 day
- Speed-up using many harvesters: 5/(4+1/30) = 1.24 times faster

Amdahl's law (Goetz 11.2)

F =sequential fraction of problem e.g. 4/5 = 0.8

N = number of threads (people) e.g. 30

Speed-up $\leq 1/(F+(1-F)/N) = 1/(0.8+0.2/30) = 1.24$

Maximum speed-up Quicksort



Best case

$$R(n) = n + n/2 + n/4 + (n/8)\log_2(n/8)$$

$$1,75n + (n/8)\log_2(n/8)$$

Amdahl: Starvation loss

Max speed-up=
$$\frac{1,75n + (n/8)\log_2(n/8)}{n \log_2(n)}$$

n= 100.000 Max speed-up ~ 4.8

Loss of scalability (much more than Amdahl)

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- Starvation loss (Amdahl)
 - Minimize the time that the task pool is empty
- QuickSort

- Separation loss (best threshold)
 - Find a good threshold to distribute workload evenly

Prime count

- Saturation loss (locking common data structure)
 - Minimize high thread contention in the problem
- Braking loss
 - Stop all tasks as soon as the problem is solved

String search

Møller-Nielsen, P and Staunstrup, J, Problem-heap. A paradigm for multiprocessor algorithms. *Parallel Computing*, 4:63-74, 1987

Shut down



The ExecutorService can be shut down.

```
// Executor body
...
...
pool.shutdown();
```

The challenge is often when to shut down

After shutdown the pool cannot be reused, but you may assign it a new value (of type **ExecutorService**)

Thread pools

```
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```

```
class PrimeCountExecutor {
   private ExecutorService pool;
   public PrimeCountExecutor () {
     pool= newFixedThreadPool(3);
tasks
                                                 CPU
                                                 CPU
```

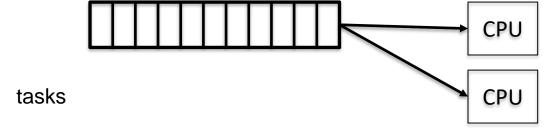
https://docs.oracle.com/javase/tutorial/essential/concurrency/pools.html

Thread pools

```
-40
```

```
class PrimeCountExecutor {
   private ExecutorService pool;
    ...
   public PrimeCountExecutor () {
      pool= newWorkStealingPool(x);
      ...
}
```

Quite a few more types of ExecutorService





Both are used to specify the code of a task.

- Runnable cannot return a result
 - Overrides run()
- <u>Callable</u> returns a result (via a Future)
 - Overrides call()

As illustrated by the Quicksort and countPrimes examples, Runnables may use shared data (e.g., to deliver a result)



Both are used to spawn a task

- pool.execute does not return a result may complicate determining when to finish
- pool.submit returns a result (via a Future)

https://docs.oracle.com/javase/8/docs/api/java/util/concurrent/Executor.html

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Java Collections



A *collection* is simply an object that groups multiple elements into a single unit

Package: java.util

Examples: ArrayList, HashMap, TreeSet, ...

https://docs.oracle.com/javase/tutorial/collections/intro/index.html

Methods: add, remove, size, contains, ...

Many of the classes have thread-safe/concurrent implementations

https://www.baeldung.com/java-synchronized-collections

Example: synchronizedCollection

```
-48
```

```
import java.util.*;
public class syncCollectionExample {
  public static void main(String[] args) {    new syncCollectionExample(); }
  public String getLast(ArrayList<String> 1) {
    int last= 1.size()-1;
    return l.get(last);
                                                        Thread-safe
  public static void delete(ArrayList<String> 1) {
                                                        (but no locking !!!!)
    int last= 1.size()-1;
    1.remove(last);
  public syncCollectionExample() {
    ArrayList<String> a= new ArrayList<String>(); // Collection
    a.add("A"); ...
    Collection<String> synColl = Collections.synchronizedCollection(a);
```

Example: synchronizedCollection

```
.49
```

```
import java.util.*;
public class syncCollectionExample {
  public static void main(String[] args) {    new syncCollectionExample(); }
  public String getLast(ArrayList<String> 1) {
    int last= 1.size()-1;
    return l.get(last);
                                                           4: Is this version thread-safe?
  public static void delete (ArrayList<String> 1) {
    int last= 1.size()-1;
    1.remove(last);
  public syncCollectionExample() {
    ArrayList<String> a= new ArrayList<String>();
    a.add("A"); ...
    Collection<String> synColl = Collections.synchronizedCollection(a);
```

Example: synchronizedCollection

```
•50
```

```
import java.util.*;
public class syncCollectionExample {
  public static void main(String[] args) {    new syncCollectionExample(); }
  public String getLast(ArrayList<String> 1) {
    int last= 1.size()-1;
    return l.get(last);
  public static void delete(ArrayList<String> 1) {
    int last= 1.size()-1;
    1.remove(last);
  public syncCollectionExample() {
    ArrayList<String> a= new ArrayList<String>();
    a.add("A"); ...
    Collection<String> synColl = Collections.synchronizedCollection(a);
```

No

(because there is not a happens-before relation between the writes and reads in the two methods)

Thread-safety (from week 3)



It is very important to note that for a program p:

Making the synchronized ArrayList thread safe



```
import java.util.*;
public class syncCollectionExample {
 public static void main(String[] args) {
                                          new syncCollectionExample(); }
 public String getLast(ArrayList<String> 1) {
    synchronized(1) {
     int last= 1.size()-1;
                                        But then the advantage of the
     return l.get(last);
                                        synchronized collections is lost !!
 public static void delete(ArrayList<String> 1) {
    synchronized(1) {
      int last= 1.size()-1;
     1.remove(last);
 public syncCollectionExample() {
```

What if the data structure is huge?



and used by many threads?

for example:

a bank

Facebook updates

•••

Would not work if everything is "synchronized"

What can we do?

Reduce locking !!

Example: A huge HashMap

Key value pairs: <k1, v1>, <k2, v2>, ...

```
class HashMap<K,V> {
    ... // data structure
    public V get(K k) { ... }
    public V put(K k, V v) { ... }
    public boolean containsKey(K k) { ... }
    public int size() { return cachedSize; }
    public V remove(K k) { ... }
    ...
}
```

How to make it thread-safe? (without making all the methods synchronized)

У	Value
ter	20487612
na	51251218
na	34458318
lger	89545010
a	94959500

Pet

An

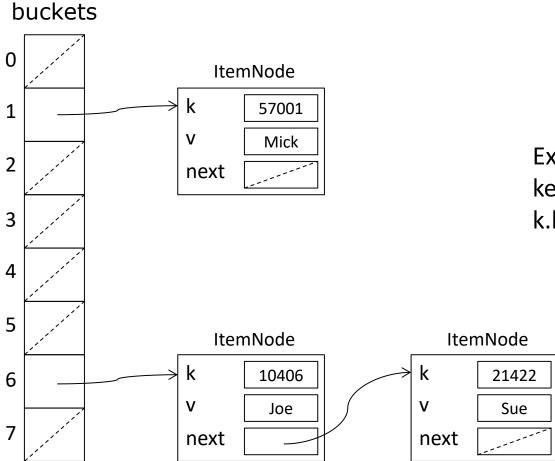
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HashMap implementation

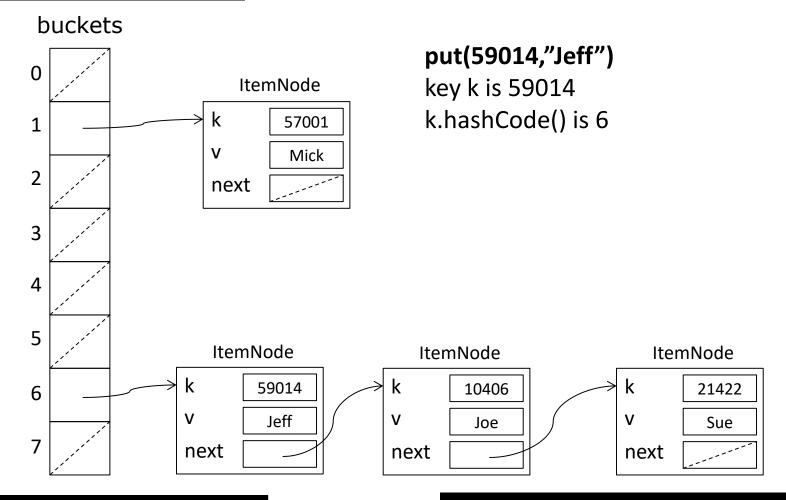




Example get(10406) key k is 10406 k.hashCode() is 6

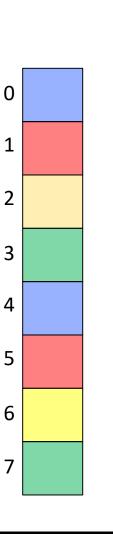
HaspMap put





Improving scalability – Lock striping

- Guarding the table with a single lock works
- -... but does not scale well (actually **very** badly)
- Idea: Each bucket could have its own lock
- In practice
- -use fewer, to illustrate we use 4, locks
- -guard every 4th bucket with the same lock
- -locks[0] guards bucket 0, 4, 8, ... (stripe 0)
- -locks[1] guards bucket 1, 5, 9, ... (stripe 1) et
- -With high probability
- -two operations will work on different stripes
- -hence will take different locks
- Less lock contention, better scalability

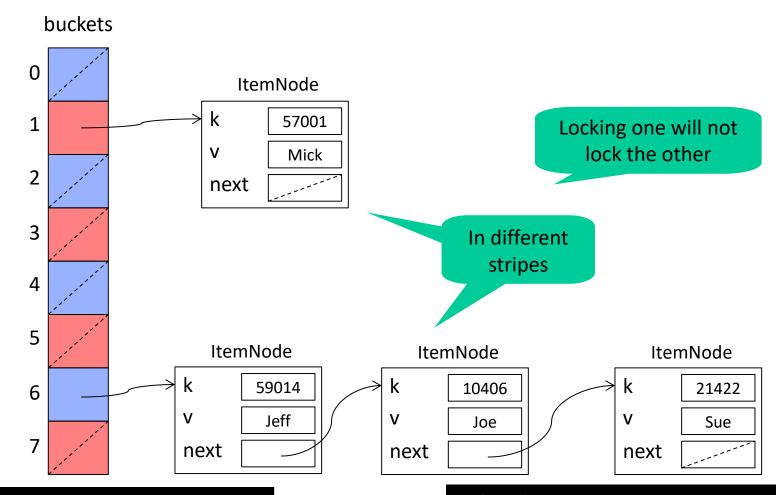


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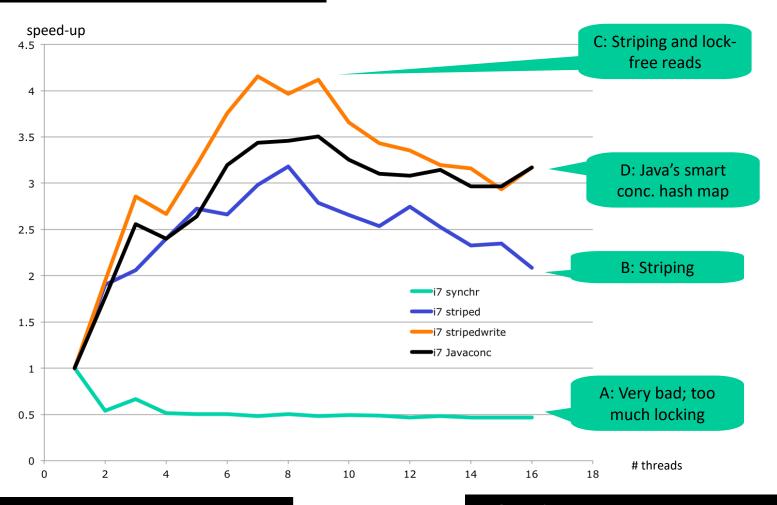
Bucket idea





Reducing locking





Ultimate scalability



A web-shop, Facebook, ...

We must give up thread safety,

but still maintain some sort of consistency