

COMPULSORY EXERCISE 1

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1. PROBLEM 1

Let $A \in \mathbb{C}^{n \times n}$ be a nonsingular triangular matrix with inverse $B = [\mathbf{b}_1 \mathbf{b}_2 \cdots \mathbf{b}_n]$. Then we have

$$(1.1) \quad AB = [A\mathbf{b}_1 \ A\mathbf{b}_2 \ \cdots \ A\mathbf{b}_n] = I,$$

that is,

$$(1.2) \quad A\mathbf{b}_k = \mathbf{e}_k,$$

$\mathbf{e}_k \in \mathbb{R}^n$ denoting the k 'th unit vector.

As usual, denote entry (i, j) of a matrix A by $a_{i,j}$.

Case 1: A lower triangular. For $k = 1$, there is nothing to prove. Let $1 < k \leq n$. Define $A_{-[n-k+1]} \in \mathbb{C}^{n-k+1 \times n-k+1}$ by

$$(1.3) \quad A_{-[n-k+1]} := \begin{pmatrix} a_{k,k} & \cdots & a_{k,n} \\ \vdots & \vdots & \vdots \\ a_{n,k} & \cdots & a_{n,n} \end{pmatrix}.$$

Then, for some matrix $L \in \mathbb{C}^{k-1 \times n-k+1}$, we have

$$(1.4) \quad A\mathbf{b}_k = \begin{pmatrix} A_{[k-1]} & 0 \\ L & A_{-[n-k+1]} \end{pmatrix} \begin{pmatrix} \mathbf{b}_k^u \\ \mathbf{b}_k^l \end{pmatrix} = \begin{pmatrix} A_{[k-1]}\mathbf{b}_k^u \\ L\mathbf{b}_k^u + A_{-[n-k+1]}\mathbf{b}_k^l \end{pmatrix} = \begin{pmatrix} 0 \\ \mathbf{e}_1 \end{pmatrix},$$

where $\mathbf{b}_k^u \in \mathbb{R}^{k-1}$ denote the $k-1$ upper elements of \mathbf{b}_k and $\mathbf{b}_k^l \in \mathbb{R}^{n-k+1}$ denote the remaining $n-k+1$ lower elements.

Since A is nonsingular and triangular, so is $A_{[k]}$ (because necessarily the diagonal elements of $A_{[k]}$ are then all nonzero), so we conclude that $\mathbf{b}_k^u \equiv 0$. Thus, we are left only with the system of equations

$$(1.5) \quad A_{-[n-k+1]}\mathbf{b}_k^l = \mathbf{e}_1,$$

or in MATLAB notation, $A(k:n, k:n)b(k:n, k:n) = I(k:n, k)$, as was to be shown. \square

Case 2: A upper triangular. Similarly, for some matrix U , we write

$$(1.6) \quad A\mathbf{b}_k = \begin{pmatrix} A_{[k]} & U \\ 0 & A_{-[n-k]} \end{pmatrix} \begin{pmatrix} \mathbf{b}_k^u \\ \mathbf{b}_k^l \end{pmatrix} = \begin{pmatrix} A_{[k]}\mathbf{b}_k^u + U\mathbf{b}_k^l \\ A_{-[n-k]}\mathbf{b}_k^l \end{pmatrix},$$

leading us to conclude that $\mathbf{b}_k^l \equiv \mathbf{0}$, so we need only solve

$$(1.7) \quad A_{[k]}\mathbf{b}_k^u = \mathbf{e}_k,$$

or in MATLAB notation, $A(1:k, 1:k)\mathbf{b}_k(1:k) = I(1:k, k)$. \square

As a final remark, we note that the nondiagonal nonzero elements of B may be stored in place of the zero elements of A . That is, we may simply store $b_{i,j}$ as $a_{j,i}$, for all elements of B not on the diagonal or on the zero part of B .

The entries on the diagonal are not worth storing, as $b_{i,i} = 1/(a_{i,i})$.

PROBLEM 2

We have, for $j \geq k = 1, 2, \dots, n$, if A is lower triangular,

$$(1.8) \quad A_{-[n-k+1]}\mathbf{b}_k^l = \mathbf{e}_1$$

which is solved by

$$(\mathbf{b}_k^l)_1 = b_{k,k} = \frac{-1}{a_{k,k}}, \text{ 1 division,}$$

$$(\mathbf{b}_k^l)_2 = b_{k,k+1} = -\frac{a_{k+1,k}b_{k,k}}{a_{k+1,k+1}}, \text{ 1 multiplication, 1 division;}$$

...

$$(\mathbf{b}_k^l)_{n-k+1} = b_{k,n} = -\frac{1}{a_{n,n}} \sum_{j=n-k+1}^n a_{j,n}b_{k,j}, \text{ 2k-1 arithmetic operations.}$$

This gives j multiplications/divisions and $j - 1$ additions/subtractions for the j 'th element of the k 'th column of B , so upon adding all operations for all k rows, we have that the total number of arithmetic operations is

$$\begin{aligned} \sum_{k=1}^n \sum_{j=1}^{n-k+1} (2j-1) &= \sum_{k=1}^n (n-k+1)(n-k+2) - (n-k+1) = \sum_{k=1}^n (n-k+1)^2 \\ &\approx \int_1^n (n-k+1)^2 dk = -\frac{1}{3}[(n-k+1)^3]_{k=1}^n = -\frac{1}{9} + \frac{n^3}{3} \approx \frac{n^3}{3}. \end{aligned}$$

Hence, the algorithm is of order $\mathcal{O}(n^3)$.

PROBLEM 3

The code snippets compute the inverse of an upper diagonal matrix; it does this via a row-oriented back-solve very similar to Algorithm 2.7. Finally, the program prints the product AU , where U is the computed inverse matrix of A , hopefully giving the identity matrix as a result. The variables r and k denote the rows and columns of U respectively.