

# ECE551 - Homework 7

Khôi-Nguyễn Mac

November 5, 2016

## 1 Truncation as Filter Approximation

(a) Let  $\psi = \{\varphi_k\}$  be the basis of  $\mathbb{C}^{\mathbb{Z}}$

$$h_d \in \mathbb{C}^{\mathbb{Z}} \Rightarrow h_d = \sum_{\varphi_k \in \psi} \alpha_k \varphi_k$$

Since  $I \subset \mathbb{Z}$ ,  $\mathbb{C}^I \subset \mathbb{C}^{\mathbb{Z}}$ , where  $\mathbb{C}^I = \text{span}\{\phi^I\}$ ,  $\phi^I \subset \phi$ .

$$\begin{aligned} T_I h_d &= \sum_{\varphi_k \in \psi} w[k] \alpha_k \varphi_k \\ &= \sum_{\varphi_k \in \psi^I} 1 \cdot \alpha_k \varphi_k + \sum_{\varphi_k \in \psi / \psi^I} 0 \cdot \alpha_k \varphi_k \\ &= \sum_{\varphi_k \in \psi^I} \alpha_k \varphi_k \in \text{span}\{\psi^I\} = \mathbb{C}^I \\ \Rightarrow T_I h_d - h_d &= \sum_{\varphi_k \in \psi / \psi^I} \alpha_k \varphi_k \end{aligned}$$

$$\Rightarrow \langle T_I h_d - h_d, T_I h_d \rangle = 0 \Rightarrow T_I h_d - h_d \perp T_I h_d$$

By orthogonality principal,  $T_I h_d$  is the least square approximation of  $h_d$  on  $\ell_2(I)$ .

(b)  $\forall z \in \mathbb{C}^{\mathbb{Z}}$ , we have

$$\begin{aligned} T_I z &= \sum_{\varphi_k \in \psi} \beta_k \varphi_k \perp T_I h_d - h_d = \sum_{\varphi_k \in \psi / \psi^I} \alpha_k \varphi_k \\ \Rightarrow \langle T_I z, T_I h_d - h_d \rangle &= 0, \forall z \in \mathbb{C}^{\mathbb{Z}} \end{aligned}$$

Hence,  $T_I$  is an orthogonal projection.

(c) For  $I = \{0, \dots, 4\}$ ,

$$T_I h_d = [\dots \quad 0 \quad \text{sinc}0 \quad \text{sinc}\frac{\pi}{3} \quad \text{sinc}\frac{2\pi}{3} \quad \text{sinc}1 \quad \text{sinc}\frac{4\pi}{3} \quad 0 \quad \dots]^\top$$

(d) We can choose  $I$  as  $\{-2, -1, 0, 1, 2\}$ , so  $T_I h_d$  is

$$T_I h_d = [\dots \quad 0 \quad -\text{sinc}\frac{2\pi}{3} \quad -\text{sinc}\frac{\pi}{3} \quad \text{sinc}0 \quad \text{sinc}\frac{\pi}{3} \quad \text{sinc}\frac{2\pi}{3} \quad 0 \quad \dots]^\top$$

## 2 Lagrange Interpolation

(a)

(b)

## 3 Polynomial Spaces with Orthogonality

(a) Let  $v \in V_n$ , then

$$v = \sum_{j=0}^n \alpha_j v_j$$

$$\deg(v) = \max\{\deg(v_j)\}_{j=0}^n \leq n$$

Therefore  $v$  can be written as  $\sum_{j=0}^n \beta_j t^j$

$$\Rightarrow v \in W_n \Rightarrow V_n \subset W_n$$

We have

$$\dim(V_n) = n \quad \because \langle v_k, v_j \rangle = \delta[k - j]$$

$$\dim(W_n) = n \quad \because \{1, t^1, t^2, \dots, t^n\} \text{ are independent}$$

So  $\dim(V_n) = \dim(W_n)$ . Hence,  $V_n = W_n$ .

(b)  $p$  is a polynomial of degree  $m$ , so  $p \in V_n = W_n$ .

$$p = \sum_{j=0}^m \langle p, v_j \rangle v_j$$

For  $k > m$ ,

$$\begin{aligned} \langle p, v_k \rangle &= \left\langle \sum_{j=0}^m \langle p, v_j \rangle v_j, v_k \right\rangle \\ &= \sum_{j=0}^m \langle p, v_j \rangle \langle v_j, v_k \rangle \\ &= 0 \quad \because \langle v_k, v_k \rangle = 0 \end{aligned}$$

(c)  $v \in V_n = W_n \Rightarrow v(t) = \sum_{j=0}^n \alpha_j t^j$

$$\begin{aligned} \sum_{j=0}^n \alpha_j (t - t_0)^j &= \sum_{j=0}^n \alpha_j \left( \binom{j}{i} t^{j-i} (-t_0)^i \right) \\ &= \sum_{j=0}^n \alpha_j \binom{j}{i} t^j \frac{(-t_0)^i}{t^i} \\ &= \sum_{j=0}^n \left( \alpha_j \binom{j}{i} \frac{(-t_0)^i}{t^i} \right) t^j \end{aligned}$$

Since  $i \leq j$ ,  $\sum_{j=0}^n \left( \alpha_j \binom{j}{i} \frac{(-t_0)^i}{t^i} \right) t^j$  is a polynomial of degree up to  $n$ . So we can write it as

$$\sum_{j=0}^n \alpha_j (t - t_0)^j = \sum_{j=0}^n \beta_j t^j$$

Hence, it is shift-invariant.

## 4 Polynomial Spaces vs. Spline Spaces

(a) Figure 1 shows the graph of  $s_0, s_1 \in U$

(b)

(c)

## 5 Interpolation with Shifted Symmetric Functions

(a)

(b)

(c)

## 6 Python: Interpolation Games

$N = 5$

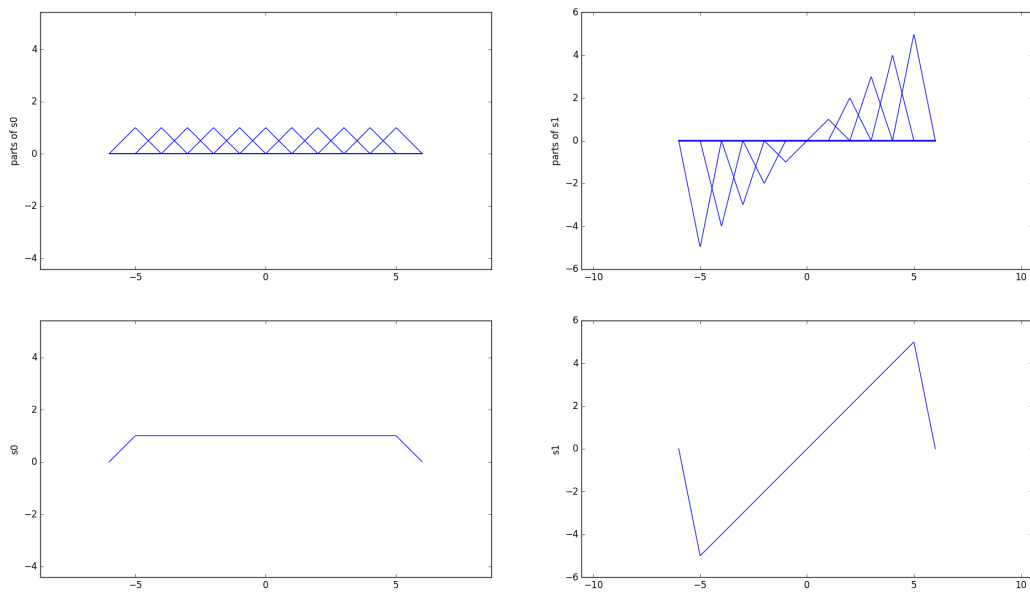


Figure 1:  $s_0$  and  $s_1$  with  $N = 5$