

Image Coding

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Useful to code images for

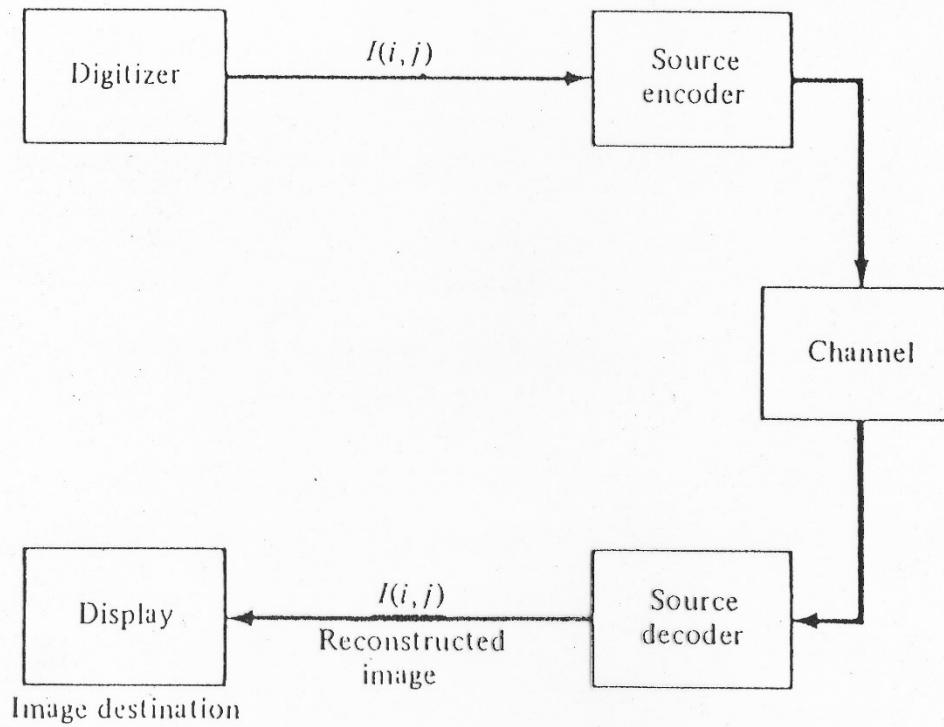
- i) digital storage ii) digital transmission

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For image sequences, we can take every nth frame. (lossy)

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Run-Length Encoding works well if most of the strings are long.

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Lossless

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11 pixels

1	2	18	1	3	1	4	1	2	18	1
A	B	R	A	C	A	D	A	B	R	A

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standard binary code 00001|00010|10010|00001|00011| ...
55 bits

1	2	18	1	3	1	4	1	2	18	1
A	B	R	A	C	A	D	A	B	R	A

Pixel Value	#Appearances
1	5
2	2
18	2
3	1
4	1

1	2	18	1	3	1	4	1	2	18	1
A	B	R	A	C	A	D	A	B	R	A

Pixel Value	#Appearances
1	5
2	2
18	2
3	1
4	1

Approach - assign shortest codes to most common pixel values

Pixel Value	Code
1	0
2	1
18	01
3	10
4	11

1	2	18	1	3	1	4	1	2	18	1				Pixel Value	Code
A	B	R	A	C	A	D	A	B	R	A				1	0
														2	1
														18	01
														3	10
														4	11

Pixels would get encoded as

010101001101010 (15 bits)

1	2	18	1	3	1	4	1	2	18	1				Pixel Value	Code
A	B	R	A	C	A	D	A	B	R	A				1	0
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But this doesn't work. The string could be decoded as

18 18 18 1 18 2 18 18 1

													Pixel Value	Code	
1	2	18	1	3	1	4	1	2	18	1			1	0	
A	B	R	A	C	A	D	A	B	R	A			2	1	
													18	01	
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We need to design the code so that no pixel code is the prefix of another.

													<u>Pixel Value</u>	<u>Code</u>
1	2	18	1	3	1	4	1	2	18	1		1	0	
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(Ex)	<u>Pixel Value</u>	<u>Code</u>
	1	11
	2	00
	18	011
	3	010
	4	10

String would be encoded as

11 00 011 11 010 11 10 11 00 011 11 (25 bits)

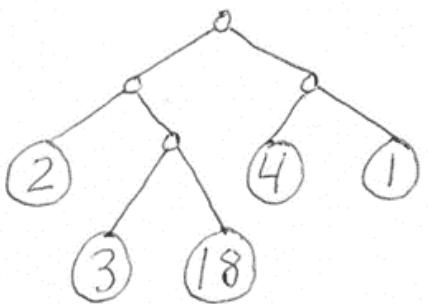
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Can represent this code using a trie



0 = go left

1 = go right

2 = 00

3 = 010

18 = 011

4 = 10

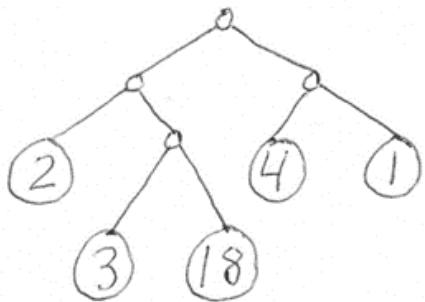
1 = 11

String would be encoded as

11 00 011 11 010 11 10 11 00 011 11 (25 bits)

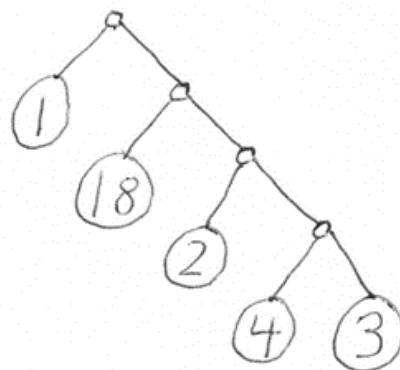
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0 = go left 2 = 00
1 = go right 3 = 010
 18 = 011
 4 = 10
 1 = 11

Another possibility



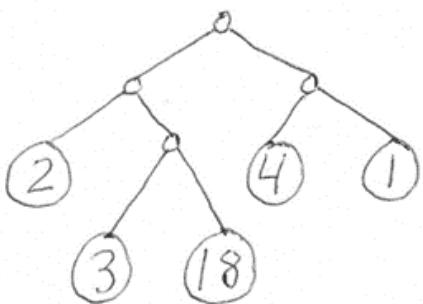
1 = 0
18 = 10
2 = 110
4 = 1110
3 = 1111

String would be encoded as

11 00 011 11 010 11 10 11 00 011 11 (25 bits)

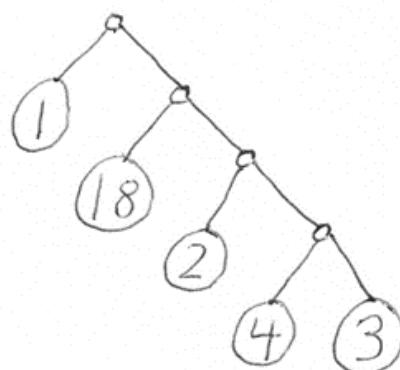
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0 = go left 2 = 00
1 = go right 3 = 010
 18 = 011
 4 = 10
 1 = 11

Another possibility



1 = 0
18 = 10
2 = 110
4 = 1110
3 = 1111

Pixels would get encoded as

0 110 10 0 1111 0 1110 0 110 10 0 (23 bits)

How do we determine the best code (trie)?

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Frequently used values should be near the top, rarely used values should be near the bottom.

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Huffman Coding

Ex "A SIMPLE STRING TO BE ENCODED USING A MINIMAL NUMBER OF BITS"

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Ex "A SIMPLE STRING TO BE ENCODED USING A MINIMAL NUMBER OF BITS"

First build a histogram.

	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	
k	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26
count[k]	11	3	3	1	2	5	1	2	0	6	0	0	2	4	5	3	1	0	2	4	3	2	0	0	0	0	

	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	
k	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26
count[k]	11	3	3	1	2	5	1	2	0	6	0	0	2	4	5	3	1	0	2	4	3	2	0	0	0	0	

Create a node for each nonzero count

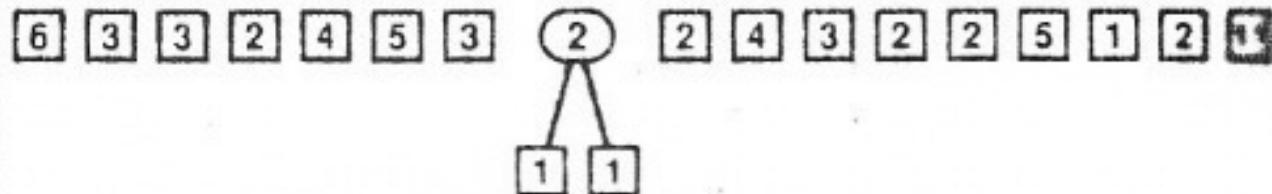
6 **3** **3** **2** **4** **5** **3** **1** **1** **2** **4** **3** **2** **2** **5** **1** **2** **11**

	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	
k	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26
count[k]	11	3	3	1	2	5	1	2	0	6	0	0	2	4	5	3	1	0	2	4	3	2	0	0	0	0	

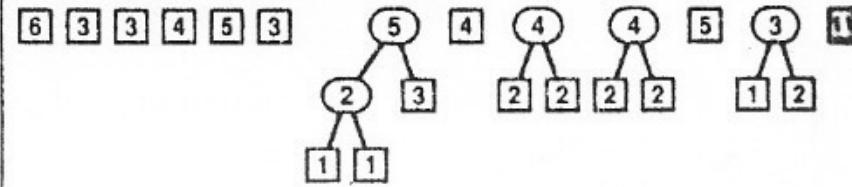
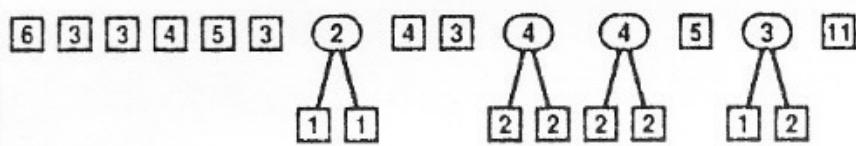
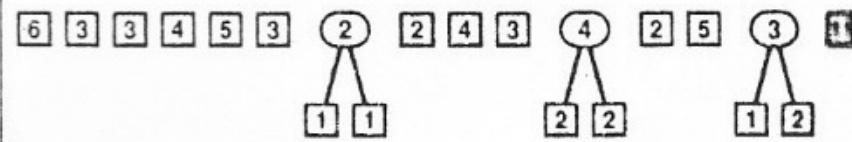
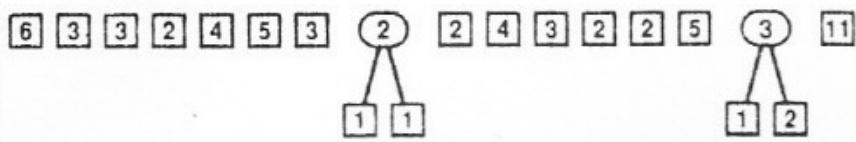
Create a node for each nonzero count

6 3 3 2 4 5 3 1 1 2 4 3 2 2 5 1 2 11

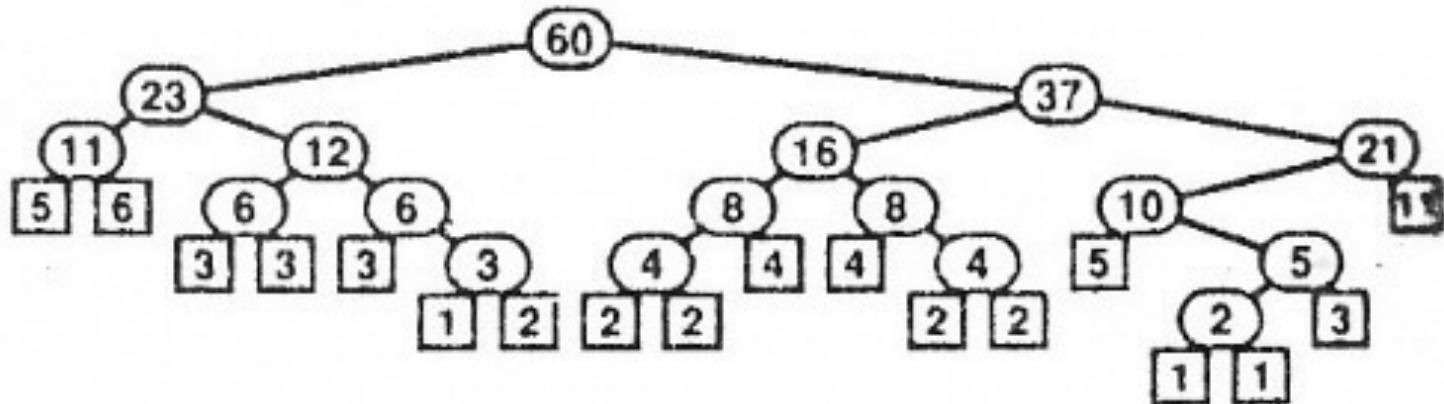
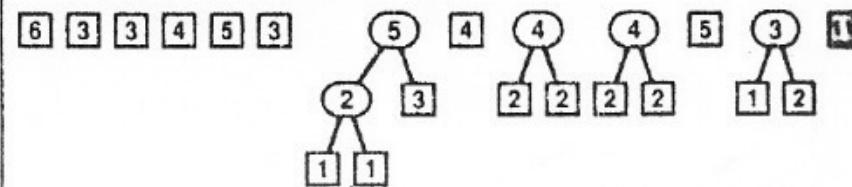
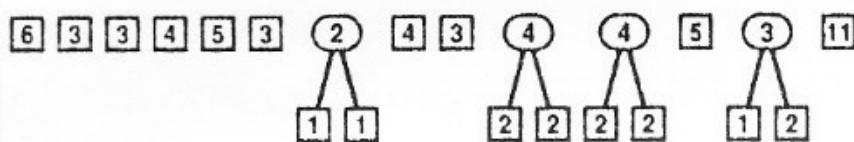
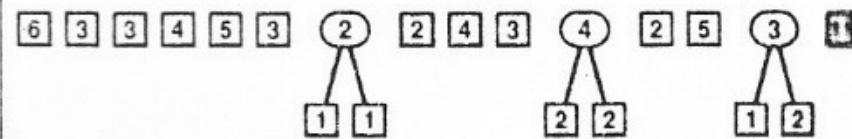
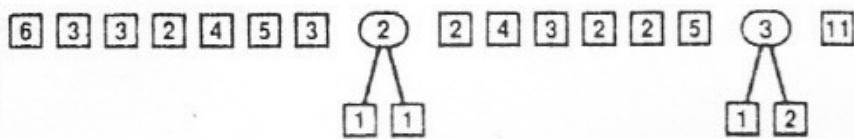
Find the 2 top level nodes with the smallest counts and
create a new node with these 2 nodes as children and
with a count the sum of the counts of the children.



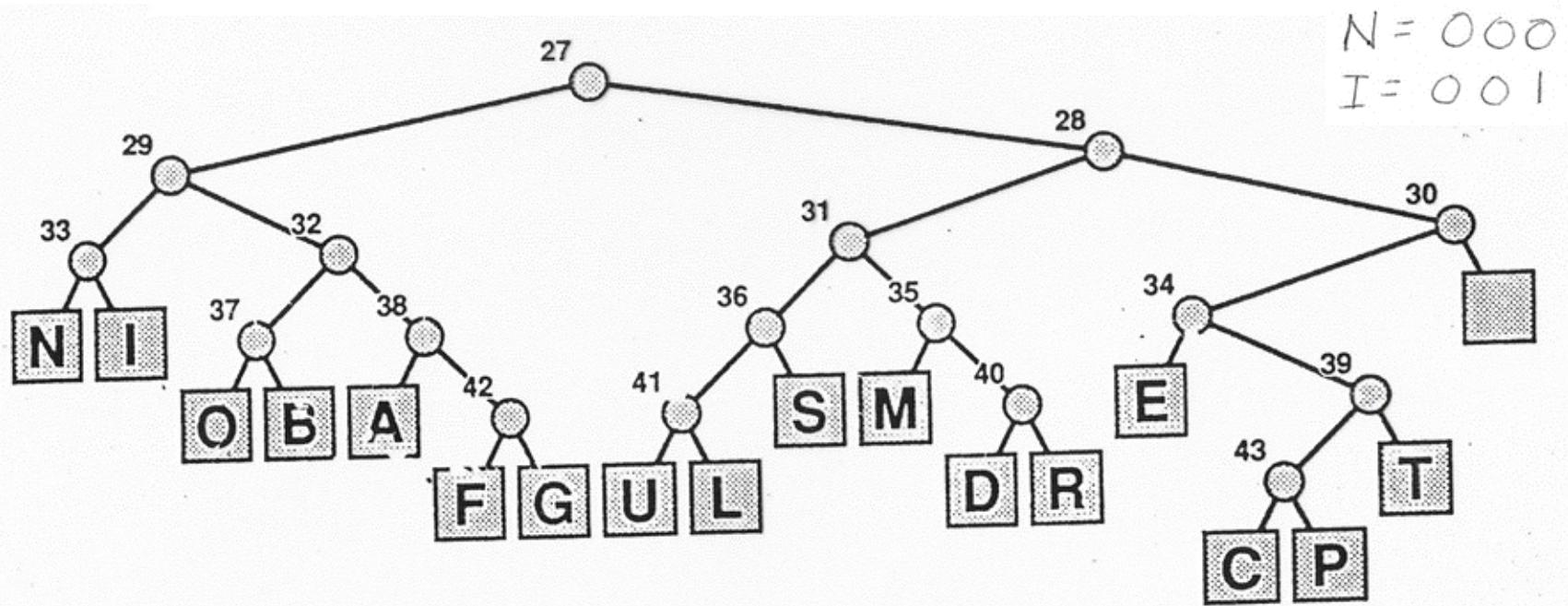
Repeat until we get a single trie.



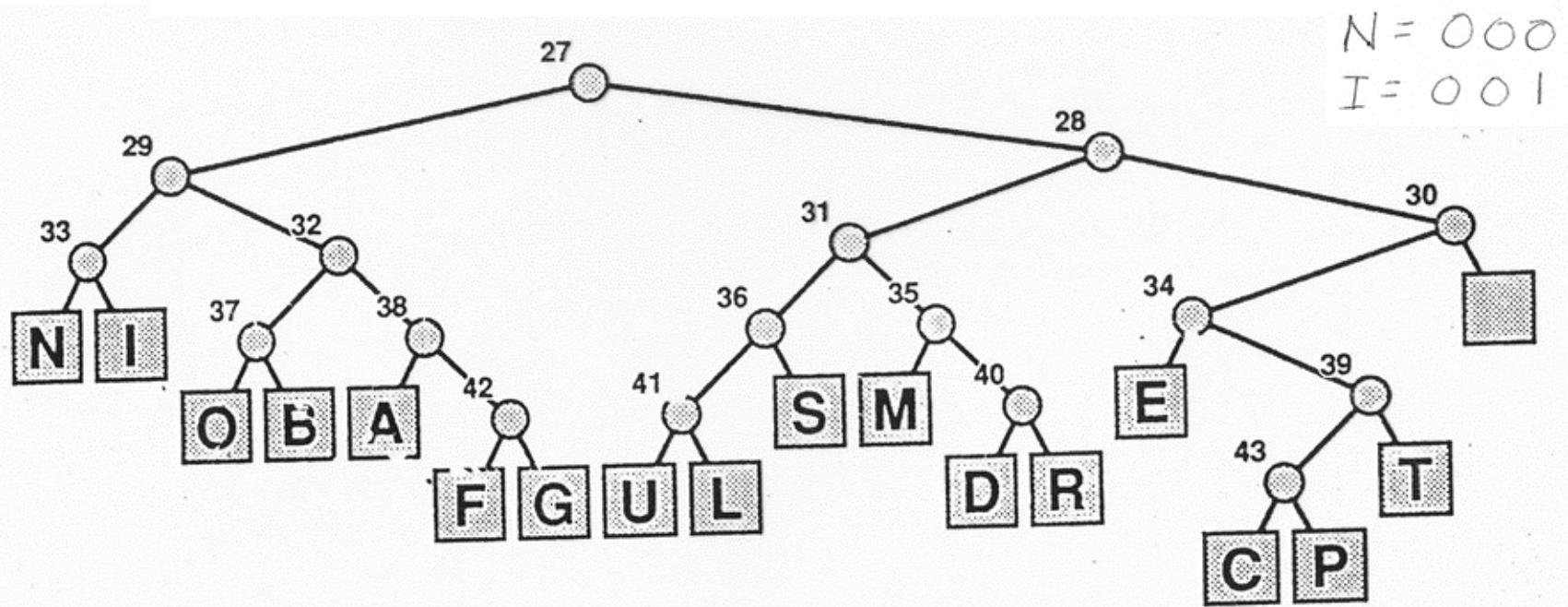
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Replace the frequencies at the bottom nodes with the associated letters and view as an encoding trie.



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The resulting trie gives the Huffman code which uses the minimum number of bits.