Posteriori Analysis of Algorithms Through the Derivations of Growth Rate Based on Frequency Count

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Outline

- I. Overview of the Research
- II. Research Problem and Applications
- III. Overview of the Proposed System
- IV. Testing and Results

Research Description

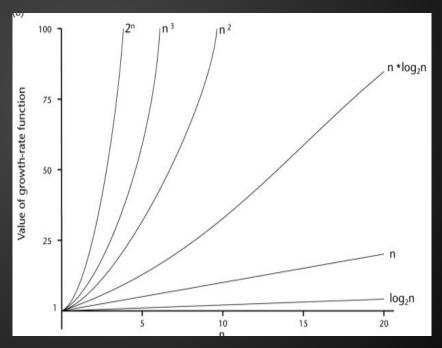
Developing an "across the board" method of

Algorithm Analysis

Output:

Time complexity

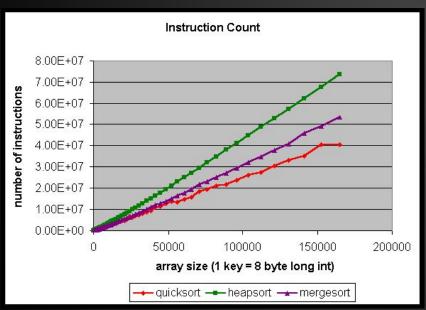
-Asymptotic Behavior



Research Description

Claim / Essence of the Study:

Performance measurements " → " Asymptotic Behavior



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Overview of the Current State of Technology

Two ways of analyzing algorithms:

- Posteriori Analysis
- Apriori Analysis

Posteriori Analysis

Idea: Analyze empirical performance of the algorithm

Advantages:

- Can be automated
- Somewhat trivial

Disadvantages:

- Varies with hardware and software
- Does not output asymptotic behavior
 - Necessary for generalizing complexity

Apriori Analysis

Idea: Analyze the logical structure of the algorithm

Advantages:

- Does not depend on external factors
- Does output asymptotic behavior

Disadvantages:

- Done by hand
- Limited by current mathematical methods

Statement of the Research Problem

There are algorithms that are too complex for current apriori and posteriori methods to accurately determine asymptotic behaviors which are necessary for measurement of the general efficiency of algorithms.

Difficult to Analyze Algorithms

- Algorithms with High-Order Linear Recurrence
 - Due to Abel's Impossibility Theorem

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + c_3 a_{n-3} + c_4 a_{n-4} + c_5 a_{n-5}$$

Difficult to Analyze Algorithms

- Algorithms where Master's Method fail
 - Difference between n/log(n) and n*log(2)/log(2)
 is not polynomial

$$T(n) = 2T\left(\frac{n}{2}\right) + \frac{n}{\log(n)}$$

Difficult to Analyze Algorithms

Lengthy algorithms (hard to keep track of)

```
){if(|a||a.node|ype||Wa.test(a.no.
n(a,b,d,e){if(!(a.nodeType===3||a.nodeType====o)
n=a; b=b.split(""); for(var x=0, r; l=b[x++];) {h=f?c.ex tend({i*, r*, r*, l=b[x++];})}
vent.global[]=true]a=null]}},global:{},remove:function(a,b,d,e){if
(), vo). join("(\.?: *\\.)?")+"(\\.\\\))]; if(\=[[[]])if(d){[=c.event.sp
; delete w. events; delete w. handle; if(typeof w=== "function")c. removeDat
B; a. target=d; b=c.makeArray(b); b. unshift(a)}a. currentTarget=d; (e=d.r
ecial[k]]][{};if((!x._default]]x._default.call(d,a)===false)&&!(e&
")+"(\\.\\\))}a.namespace=a.namespace||d.jain(".");f=c.data(this,thi
 c.expando])return a; var b=a; a=c. Event(b); for (var t=this.props.leng
 b&&b.clientTop||d&&d.clientTop||0)}if(a.which==null&&(a.char
ctor),a)}},beforeunload:{setup:function(a,b,d){if(c.isWindow(this))
do]=true]: . Event. prototype={preventDefault: function(){this.isDefault
r b=a.relatedlarget; try{for(; b&&b!==this;)b=b.parentNode; if(b!==thi
this, "click.specialSubmit", function(a) [var b=a.target, d=b.type;if(Cd
a.type, d=a.value; if(b=== "radio" | 1b=== "checkbox")d=a.checked; else if(
nner(a,b,d)}}};c.event.special.change={filters:{focusout:2,before
          ovent.add(this, a+".specialChange", V[a]); return ia.ta
                        "one"], function(a,b){c,fn[b]=functi
```

"Impossible" to Analyze Algorithms

Algorithms with General Nonlinear Recursion

$$F_n = (n)F_{n-1} + (n^2)c_2F_{n-2} + 3n$$

"Impossible" to Analyze Algorithms

Algorithms with Irreducible Double Recursion

$$Ack(m,n) = \begin{cases} n+1 & \text{if } m = 0\\ Ack(m-1,1) & \text{if } m > 0 \text{ and } n = 0\\ Ack(m-1,Ack(m,n-1)) & \text{if } m > 0 \text{ and } n > 0 \end{cases}$$

A way around the difficulties

- Generating automatic measurements
- Using a method that is numerical
 - Based on asymptotics
 - "Stumbles upon the answer"

Scope and Limitations

- Covers rates of growth from logarithmic to exponential
- Exact answer is not guaranteed
- Convergence test is not completely foolproof
- Results are for one parameter only

Theoretical Framework

The use of asymptotic equivalences

$$\lim_{n \to \infty} \frac{f(n)}{g(n)} < \infty \longleftrightarrow f(n) \in O(g(n))$$

$$0 < \lim_{n \to \infty} \frac{f(n)}{g(n)} < \infty \longleftrightarrow f(n) \in \Theta(g(n))$$

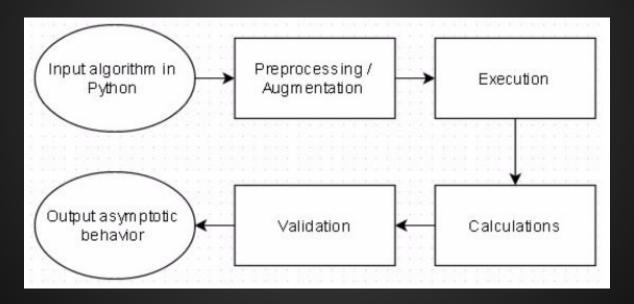
$$\lim_{n \to \infty} \frac{f(n)}{g(n)} > 0 \longleftrightarrow f(n) \in \Omega(g(n))$$

$$f(n) \sim j(n) \text{ as } n \to \infty$$

VS. or

$$\left(\lim_{n\to\infty}\frac{f(n)}{g(n)}=1\right)$$

Posteriori method that outputs asymp. behavior



- **Input :** Python file of the input algorithm The Python file must contain:
 - Function named "f" taking one parameter
 - "f" will be ran in experiments
 - Every function called by "f"

Sample input: File containing

```
def f(x):
    if(x<=1):
        return x
    else:
        return fun2(x, "hello world")

def fun2(x, str):
    print str
    return x/2.0</pre>
```

Preprocessing: Augmentation of the input file Inserting freqCount+=1 for lines with:

- Condition Check
- Function call
- Assignment statement
- Return statement

Sample augmentation: Resulting file

Execution: Broken down into five steps

- Running the augmented algorithm for a particular input size
- Storing freqCount measurements
- Setting freqCount variable back to 0
- Increment input size
- Repeat until input size = # of terms

Sample execution: Resulting data

Input Size	Measured Frequency Count
0	2
1	2
2	4
3	4
4	4

Calculation / Validation:

- Generating asymptotic behavior approx.
- Checking for discontinuities
- Determining the presence of log(n) growth
- Testing for asymptotic equivalence

Definition used:

$$\left(\lim_{n\to\infty}\frac{f(n)}{g(n)}=1\right)$$

Important property (independent):

$$\lim_{n \to \infty} (G_1(n)G_2(n)) = (\lim_{n \to \infty} G_1(n))(\lim_{n \to \infty} G_2(n))$$

Generalized scope:

$$A^n \cdot n^B \cdot Cn \cdot \sqrt[D]{n} \cdot log_E(n)$$

Generalized problem

$$a_n = \sum_{i=1}^m (e_i^n \cdot n^{p_i} \cdot c_i \ln(n, hasLog_i))$$

Simplification (Asymptotics)

$$a_n \sim e_1^n n^{p_1} c_1 \ln(n, hasLog_1)$$

Algebraic Manipulation

$$a_x = e^x x^p c \ln(x)$$

$$\frac{a_x}{e^x x^p \ln(x)} = c = \frac{a_y}{e^y y^p \ln(y)}$$

$$\begin{split} e &= exp(((log(y) - log(z))(log(a_x) - log(a_y) + log(ln(y)) - log(ln(x))) \\ &- (log(x) - log(y))(log(a_y) - log(a_z) + log(ln(z)) - log(ln(y)))) \\ &\div ((log(x) - log(y))(z - y) - (log(y) - log(z))(y - x))) \end{split}$$

$$p = ((y - z)(log(a_x) - log(a_y) + log(ln(y)) - log(ln(x)))$$
$$-(x - y)(log(a_y) - log(a_z) + log(ln(z)) - log(ln(y))))$$
$$\div((x - y)(log(z) - log(y)) - (y - z)(log(y) - log(x)))$$

$$\begin{split} c &= exp(((z\,\log(y) - ylog(z))(y(\log(a_x) - \log(\ln(x))) - x(\log(a_y) - \log(\ln(y)))) \\ &- (ylog(x) - xlog(y))(z(\log(a_y) - \log(\ln(y))) - y(\log(a_z) - \log(\ln(z))))) \\ &\div ((ylog(x) - xlog(y))(y - z) - (z\,\log(y) - ylog(z))(x - y))) \end{split}$$

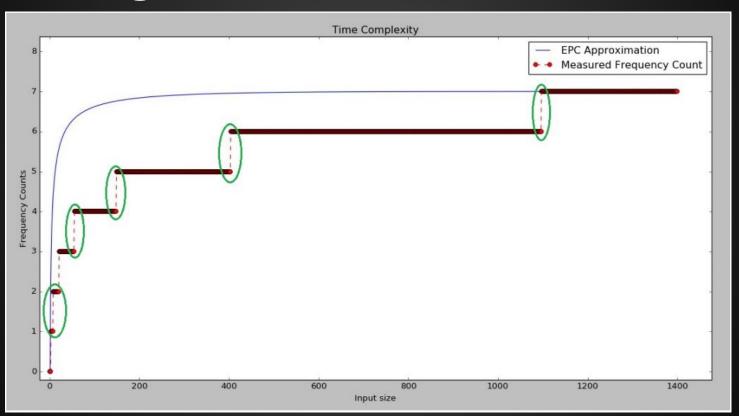
Fibo(n): Approximations

enlapprox. pnl approx. cnl approx. 1.617985612880 🗻 0.00087558705 T'480T30A0400; 1.786344173552 1.618002889798 0.000581974829 1.787167535114 1.618014025129 0.000385854481 1,618021191566 0.000255207181 1.787724293884 1.788099394175 1.618025795809 0.0001684242401.788351311609 1,618028749881 0.00011091875 1.788519984327 1.618030642550 0.00007290530i1.788632610764 1,618031853701 0.00004783128 1.788707618224 1.618032627815 0.00003132663 1.788757452260 1.618033122060 0.00002048350! 1.788790486382 1.618033437292 0.00001337285! 1.618033638157 0.000008717830 1.788812337826

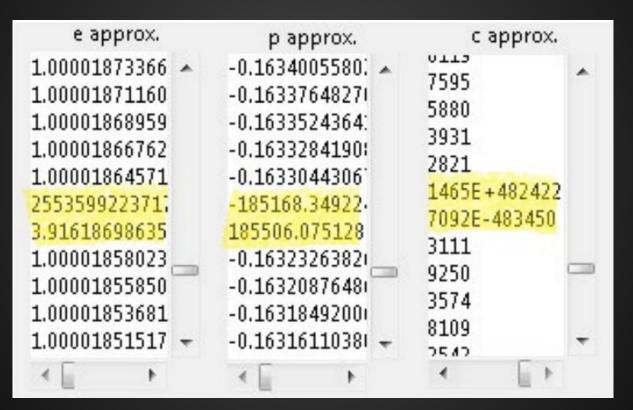
Actual Values

$$e_{Fibo} = \frac{1+\sqrt{5}}{2} \approx 1.618033988750$$
 $p_{Fibo} = 0$
 $c_{Fibo} = \frac{4}{\sqrt{5}} \approx 1.788854382000$
 $hasLog_{Fibo} = False$

Checking for Discontinuities



Checking for Discontinuities



Determining hasLog & Convergence

"Voting" heuristics

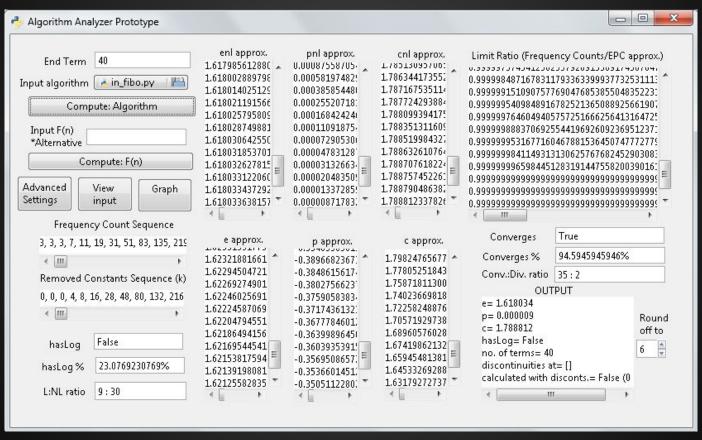
If $|a_n - e_{nl}^n * n^{p_{nl}} * c_{nl}| > |a_n - e^n * n^p * c \ln(n)|$ then "No Log" case gains 1 vote

If
$$|(n+1)*(F_{n+2}-F_{n+1})| \le |n*(F_{n+1}-F_n)|$$
.
then "Converges" case gains 1 vote

System UI

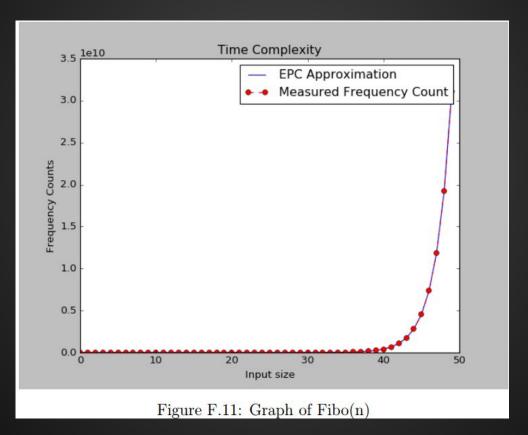
End Term 50	enl approx.	pnl approx.	cnl approx.	Limit Ratio (Frequency Counts/EPC approx.)
put algorithm (None)				
Compute: Algorithm				
Input F(n) *Alternative				
Compute: F(n)				
Advanced View Graph input				
Frequency Count Sequence				
	e approx.	p approx.	c approx.	Converges
				Converges % Conv.:Div. ratio
Removed Constants Sequence (k)				OUTPUT
				Roun off to
hasLog				6
hasLog %				
L:NL ratio				

System UI



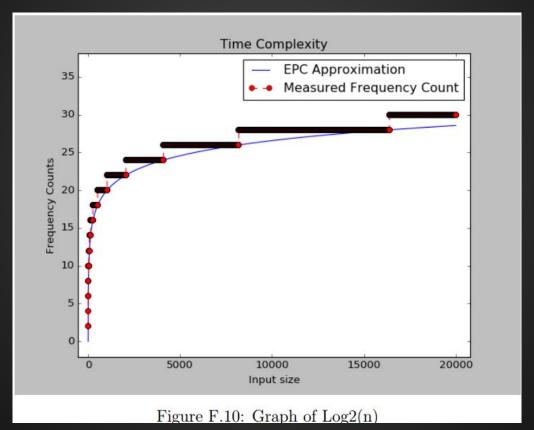
```
def f(n):
    if(n<=2):
        return 1
    else:
        return f(n-1) + f(n-2)</pre>
```

Algorithm	Expected Algorithm Output	Actual Algorithm Output
Fibo(n)	e = 1.618034 $p = 0$ $c = 1.788854$ $hasLog = False$	e = 1.618034 $p = 0.000000$ $c = 1.788854$ $hasLog = False$ (Using 50 terms)



```
def f(n):
    if(n<=2):
        return 1
    else:
        return 1+f(n/2.0)</pre>
```

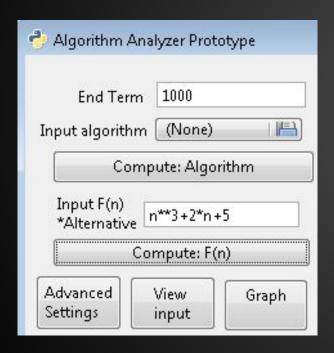
Algorithm	Expected Algorithm Output	Actual Algorithm Output
Log2(n)	e = 1 $p = 0$ $c = 2.885390$ $hasLog = True$	e = 1.000000 $p = 0.000035$ $c = 2.884487$ $hasLog = True$ (Using 20000 terms)

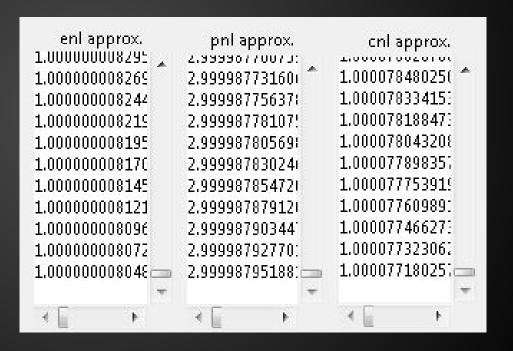


```
def f(n):
    if n==0:
        return 1
    else:
        for i in range(0, n):
        f(n-1)
```

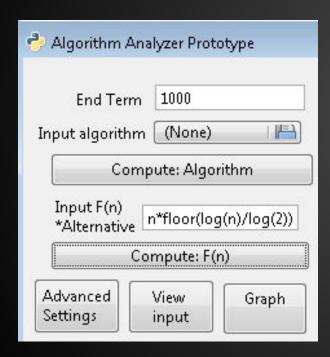
Algorithm	Expected Algorithm Output	Actual Algorithm Output
FactLike(n)	e = divergent p = divergent c = divergent hasLog = False	Converges = False (Using 30 terms)

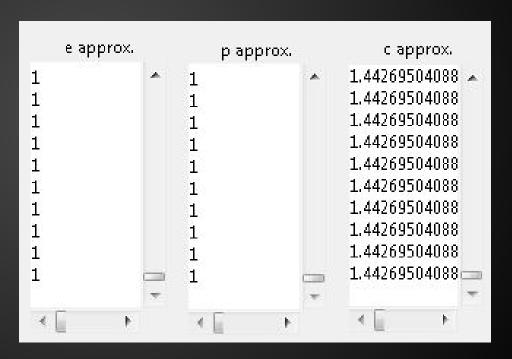
Convergent Continuous Approximation



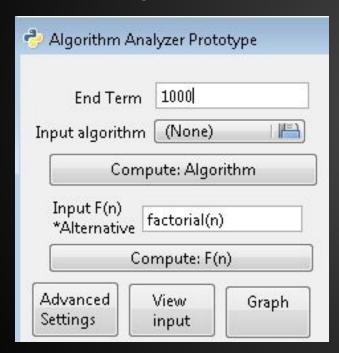


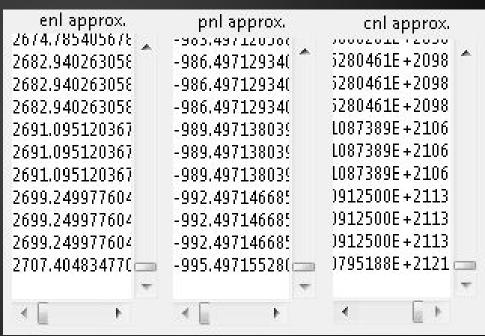
Convergent Discontinuous Approximation



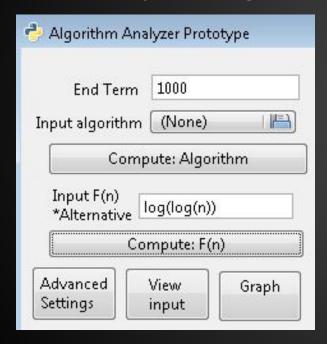


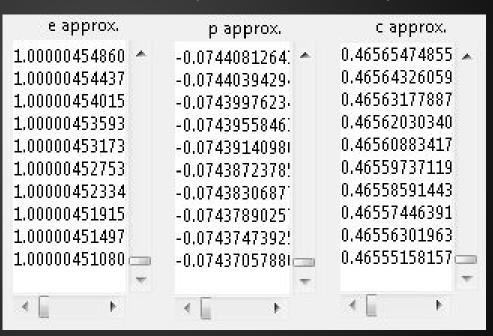
Divergent Approximation





Slowly Divergent Approximation (false positive)





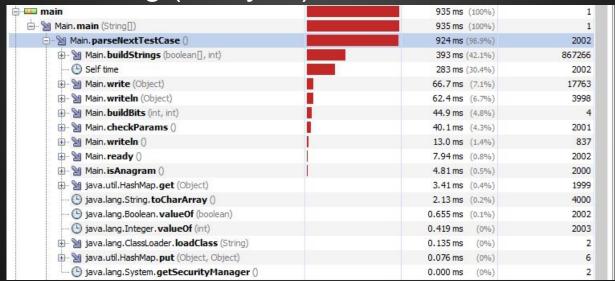
Questions & Answers

References

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Posteriori Sample - Profilers

- 1. Pre-Processing (Augmentation)
- 2. Processing (Running)
- 3. Post-Processing (Analysis)



Apriori Sample - Iterative Method

- 1. Expanding any iterations or recursions present
- 2. Finding a pattern or a series
- 3. Generalization of the pattern

$$T(n) = T(n-1) + n$$
, $T(0) = c$
 $T(n) = T(n-2) + n-1 + n$
 $T(n) = T(n-3) + n-2 + n-1 + n$
 $T(n) = T(n-i) + i*n - (0+1+2+...+i-1)$
 $\therefore T(n) = c + n^2 - (n-1)*n / 2$