

Next-Next-Gen Notes

Object-Oriented Maths

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undefined terms: $(:=), (=), (,), (arg_(_)), (.)$

1 Logic and Set Theory

1.1 Logical Truths and Operators

$$truth[t] := \left(t = \begin{cases} T \\ F \end{cases} \right) \quad (1)$$

$$operatorOR[\vee][x, y] := {}_1(truth[x]), {}_1(truth[y]), {}_1 \left(truth[x \vee y] = \begin{cases} F & x=F, y=F \\ T & x=F, y=T \\ T & x=T, y=F \\ T & x=T, y=T \end{cases} \right) {}_1 \quad (2)$$

$$operatorAND[\wedge][x, y] := {}_1(truth[x]), {}_1(truth[y]), {}_1 \left(truth[x \wedge y] = \begin{cases} F & x=F, y=F \\ F & x=F, y=T \\ F & x=T, y=F \\ T & x=T, y=T \end{cases} \right) {}_1 \quad (3)$$

$$operatorNOT[\neg][x] := {}_1(truth[x]), {}_1 \left(truth[\neg x] = \begin{cases} T & x=F \\ F & x=T \end{cases} \right) {}_1 \quad (4)$$

$$\begin{aligned} booleanAlgebra[\{T, F\}, \wedge, \vee, \neg] &:= {}_1^{POS-LCom}((x \wedge y = y \wedge x), {}_1(x \vee y = y \vee x)) \# \text{Commutative}, {}_1 \\ &{}^{POS-LDis}((x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)), {}_1(x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z))) \# \text{Distributive}, {}_1 \\ &{}^{POS-LIdn}((x \wedge T = x), {}_1(x \vee F = x)) \# \text{Identity}, {}_1 \\ &{}^{POS-LComp}((x \wedge \neg x = F), {}_1(x \vee \neg x = T)) \# \text{Complement}, {}_1 \end{aligned}$$

Note: I sometimes get too lazy to refer to *POS-LCom*.



(5)

$$operatorXOR[\underline{\vee}][x, y] := {}_1(truth[x]), {}_1(truth[y]), {}_1 \left(truth[x \underline{\vee} y] = \begin{cases} F & x=F, y=F \\ T & x=F, y=T \\ T & x=T, y=F \\ F & x=T, y=T \end{cases} \right) {}_1 \quad (6)$$

$$\text{operator } IF[\Rightarrow][x, y] :=_1 (\text{truth}[x]),_1 (\text{truth}[y]),_1 \left(\text{truth}[x \Rightarrow y] = (\neg x) \vee y = \begin{cases} T & x=F, y=F \\ T & x=F, y=T \\ F & x=T, y=F \\ T & x=T, y=T \end{cases} \right) \cdot_1 \quad (7)$$

$$\begin{aligned} & \text{THM-LExp-1} (F = x \wedge \neg x) \Rightarrow_1 \\ & \text{POS-LCmp} \\ & \text{THM-LExp-2} (x),_1 \\ & \text{THM-LExp-3} (\neg x),_1 \\ & \text{THM-LExp-4} (x \vee y),_1 \\ & \text{THM-LExp-5} (y),_1 \\ & \text{THM-LExp-3} \\ & \text{THM-LExp-1} (F \Rightarrow y) \\ & \text{THM-LExp-2} \\ & \text{THM-LExp-3} \\ & \text{THM-LExp-4} \\ & \text{THM-LExp-5} \end{aligned}$$

The Principle of Explosion, anything follows from a false (F) premise (8)

$$\text{operator } OIF[\Leftarrow][x, y] :=_1 (\text{truth}[x]),_1 (\text{truth}[y]),_1 \left(\text{truth}[x \Leftarrow y] = (\neg y) \vee x = \begin{cases} T & x=F, y=F \\ F & x=F, y=T \\ T & x=T, y=F \\ T & x=T, y=T \end{cases} \right) \cdot_1 \quad (9)$$

$$\text{operator } IIF[\Leftrightarrow][x, y] :=_1 (\text{truth}[x]),_1 (\text{truth}[y]),_1 \left(\text{truth}[x \Leftrightarrow y] = (x \Rightarrow y) \wedge (y \Rightarrow x) = \begin{cases} T & x=F, y=F \\ F & x=F, y=T \\ F & x=T, y=F \\ T & x=T, y=T \end{cases} \right) \cdot_1 \quad (10)$$

1.2 Boolean Algebra Properties

$$\begin{aligned} & \text{THM-Dual-1} \left(\text{booleanAlgebra}[\{T, F\}, \wedge, \vee, \neg] \Leftrightarrow_1 \right. \\ & ((x \wedge y = y \wedge x),_1 (x \vee y = y \vee x)) \# \text{Commutative},_1 \\ & ((x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)),_1 (x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z))) \# \text{Distributive},_1 \\ & ((x \wedge T = x),_1 (x \vee F = x)) \# \text{Identity},_1 \\ & ((x \wedge \neg x = F),_1 (x \vee \neg x = T)) \# \text{Complement},_1 \Leftrightarrow_2 \\ & ((x \vee y = y \vee x),_2 (x \wedge y = y \wedge x)) \# \text{Reordered Commutative},_2 \\ & ((x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)),_2 (x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z))) \# \text{Reordered Distributive},_2 \\ & ((x \vee F = x),_2 (x \wedge T = x)) \# \text{Reordered Identity},_2 \\ & ((x \vee \neg x = T),_2 (x \wedge \neg x = F)) \# \text{Reordered Complement},_2 \Leftrightarrow \\ & \left. \text{booleanAlgebra}[\{F, T\}, \vee, \wedge, \neg] \right) \\ & \text{THM-Dual} \\ & \text{THM-Dual-1} (\text{booleanAlgebra}[\{T, F\}, \wedge, \vee, \neg] \Leftrightarrow \text{booleanAlgebra}[\{F, T\}, \vee, \wedge, \neg]) \\ & \# \text{Boolean Algebra Duality follows from the swap symmetry of } (\wedge, T) \text{ and } (\vee, F) \text{ within the axioms} \quad (11) \end{aligned}$$

$$\begin{aligned}
& \text{THM-LUNt-1}((x \vee y = T = x \vee z) \wedge (x \wedge y = F = x \wedge z)) \implies_1 \\
& \quad \text{THM-LUNt-2} \text{ POS-LIdn} (y = y \wedge T),_1 \\
& \quad \text{THM-LUNt-3} \text{ THM-LUNt-1} (y \wedge T = y \wedge (x \vee z)),_1 \\
& \quad \text{THM-LUNt-4} \text{ POS-LDis} (y \wedge (x \vee z) = (y \wedge x) \vee (y \wedge z)),_1 \\
& \quad \text{THM-LUNt-5} \text{ POS-LCom} ((y \wedge x) \vee (y \wedge z) = (x \wedge z) \vee (y \wedge z)),_1 \\
& \quad \text{THM-LUNt-4} \text{ THM-LUNt-5} \\
& \quad \text{THM-LUNt-6} \text{ POS-LCom} \text{ POS-LDis} ((x \wedge z) \vee (y \wedge z) = z \wedge (x \vee y)),_1 \\
& \quad \text{THM-LUNt-7} \text{ THM-LUNt-1} (z \wedge (x \vee y) = z \wedge T),_1 \\
& \quad \text{THM-LUNt-8} \text{ POS-LIdn} (z \wedge T = z),_1 \\
& \text{THM-LUNt} \text{ THM-LUNt-1} \text{ THM-LUNt-2} \text{ THM-LUNt-3} \text{ THM-LUNt-4} \text{ THM-LUNt-5} \text{ THM-LUNt-6} \text{ THM-LUNt-7} \text{ THM-LUNt-8} \\
& \left(((x \vee y = T = x \vee z) \wedge (x \wedge y = F = x \wedge z)) \implies (y = z) \right)
\end{aligned}$$

Uniqueness of Complements (12)

$$\begin{aligned}
& \text{THM-LDom-1} \text{ POS-LIdn} (x \vee T = (x \vee T) \wedge T) \\
& \text{THM-LDom-2} \text{ POS-LCmp} ((x \vee T) \wedge T = (x \vee T) \wedge (x \vee \neg x)) \\
& \text{THM-LDom-3} \text{ POS-LDis} ((x \vee T) \wedge (x \vee \neg x) = x \vee (T \wedge \neg x)) \\
& \text{THM-LDom-4} \text{ POS-LIdn} (x \vee (T \wedge \neg x) = x \vee \neg x) \\
& \text{THM-LDom-5} \text{ POS-LCmp} (x \vee \neg x = T) \\
& \text{THM-LDom-6} \text{ THM-LDom-1} \text{ THM-LDom-2} \text{ THM-LDom-3} \text{ THM-LDom-4} \text{ THM-LDom-5} \\
& \text{THM-LDom} \text{ THM-LDom-6} \text{ THM-Dual} ((x \vee T = T), (x \wedge F = F))
\end{aligned}$$

Domination (13)

$$\begin{aligned}
& \text{THM-LIdm-1} \text{ POS-LIdn} (x \vee x = (x \vee x) \wedge T) \\
& \text{THM-LIdm-2} \text{ POS-LCmp} ((x \vee x) \wedge T = (x \vee x) \wedge (x \vee \neg x)) \\
& \text{THM-LIdm-3} \text{ POS-LDis} ((x \vee x) \wedge (x \vee \neg x) = x \wedge (x \vee \neg x)) \\
& \text{THM-LIdm-4} \text{ POS-LCmp} (x \wedge (x \vee \neg x) = x \wedge T) \\
& \text{THM-LIdm-5} \text{ POS-LIdn} (x \wedge T = x) \\
& \text{THM-LIdm-6} \text{ THM-LIdm-1} \text{ THM-LIdm-2} \text{ THM-LIdm-3} \text{ THM-LIdm-4} \text{ THM-LIdm-5} \\
& \text{THM-LIdm} \text{ THM-LIdm-6} \text{ THM-Dual} ((x \vee x = x), (x \wedge x = x))
\end{aligned}$$

Idempotent (14)

$$\begin{aligned}
& \text{THM-LInv-1} \text{ POS-LIdn} (\neg x = \neg x \vee F) \\
& \text{THM-LInv-2} \text{ POS-LCmp} (\neg x \vee F = \neg x \vee (x \wedge \neg x)) \\
& \text{THM-LInv-3} \text{ POS-LDis} (\neg x \vee (x \wedge \neg x) = (\neg x \vee x) \wedge (\neg x \vee \neg x)) \\
& \text{THM-LInv-4} \text{ POS-LCmp} ((\neg x \vee x) \wedge (\neg x \vee \neg x) = (\neg x \vee x) \wedge T) \\
& \text{THM-LInv-5} \text{ POS-LCmp} ((\neg x \vee x) \wedge T = (\neg x \vee x) \wedge (x \vee \neg x)) \\
& \text{THM-LInv-6} \text{ POS-LDis} ((\neg x \vee x) \wedge (x \vee \neg x) = x \vee (\neg x \wedge \neg x))
\end{aligned}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LInv-7} \\
\textcolor{blue}{POS-LCmp} \left(x \vee (\neg \neg x \wedge \neg x) = x \vee F \right) \\
\textcolor{teal}{THM-LInv-8} \\
\textcolor{blue}{POS-LIdn} \left(x \vee F = x \right) \\
\textcolor{teal}{THM-LInv} \\
\textcolor{teal}{THM-LInv-1} \left(\neg \neg x = x \right) \\
\textcolor{teal}{THM-LInv-2} \\
\textcolor{teal}{THM-LInv-3} \\
\textcolor{teal}{THM-LInv-4} \\
\textcolor{teal}{THM-LInv-5} \\
\textcolor{teal}{THM-LInv-6} \\
\textcolor{teal}{THM-LInv-7} \\
\textcolor{teal}{THM-LInv-8}
\end{array}$$

Involution (15)

$$\begin{array}{l}
\textcolor{teal}{THM-LAbs-1} \\
\textcolor{blue}{POS-LIdn} \left(x \vee (x \wedge y) = (x \wedge T) \vee (x \wedge y) \right) \\
\textcolor{teal}{THM-LAbs-2} \\
\textcolor{blue}{POS-LDis} \left((x \wedge T) \vee (x \wedge y) = x \wedge (T \vee y) \right) \\
\textcolor{teal}{THM-LAbs-3} \\
\textcolor{blue}{THM-LDom} \left(x \wedge (T \vee y) = x \wedge T \right) \\
\textcolor{teal}{THM-LAbs-4} \\
\textcolor{blue}{POS-LIdn} \left(x \wedge T = x \right) \\
\textcolor{teal}{THM-LAbs-5} \\
\textcolor{teal}{THM-LAbs-1} \left(x \vee (x \wedge y) = x \right) \\
\textcolor{teal}{THM-LAbs-2} \\
\textcolor{teal}{THM-LAbs-3} \\
\textcolor{teal}{THM-LAbs-4}
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAbs} \\
\textcolor{teal}{THM-LAbs-5} \\
\textcolor{blue}{THM-Dual}
\end{array}
\left((x \vee (x \wedge y) = x), (x \wedge (x \vee y) = x) \right)$$

Absorption (16)

$$\textcolor{teal}{THM-LAsc-1} \left((A = x \vee (y \vee z)), (B = (x \vee y) \vee z) \right) \Rightarrow_1$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-2} \\
\textcolor{blue}{THM-LAsc-1} \left(x \wedge A = x \wedge (x \vee (y \vee z)) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-3} \\
\textcolor{blue}{THM-LAbs} \left(x \wedge (x \vee (y \vee z)) = x \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-4} \\
\textcolor{blue}{THM-LAsc-1} \left(x \wedge B = x \wedge ((x \vee y) \vee z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-5} \\
\textcolor{blue}{POS-LDis} \left(x \wedge ((x \vee y) \vee z) = (x \wedge (x \vee y)) \vee (x \wedge z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-6} \\
\textcolor{blue}{THM-LAbs} \left((x \wedge (x \vee y)) \vee (x \wedge z) = x \vee (x \wedge z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-7} \\
\textcolor{blue}{THM-LAbs} \left(x \vee (x \wedge z) = x \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-8} \\
\textcolor{blue}{THM-LAbs} \left((x \wedge (x \vee y)) \vee (x \wedge z) = x \vee (x \wedge z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-9} \\
\textcolor{teal}{THM-LAsc-2} \left(x \wedge A = x = x \wedge B \right),_1 \\
\textcolor{teal}{THM-LAsc-3} \\
\textcolor{teal}{THM-LAsc-4} \\
\textcolor{teal}{THM-LAsc-5} \\
\textcolor{teal}{THM-LAsc-6} \\
\textcolor{teal}{THM-LAsc-7} \\
\textcolor{teal}{THM-LAsc-8}
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-10} \\
\textcolor{blue}{THM-LAsc-1} \left(\neg x \wedge A = \neg x \wedge (x \vee (y \vee z)) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-11} \\
\textcolor{blue}{POS-LDis} \left(\neg x \wedge (x \vee (y \vee z)) = (\neg x \wedge x) \vee (\neg x \wedge (y \vee z)) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-12} \\
\textcolor{blue}{POS-LCmp} \left((\neg x \wedge x) \vee (\neg x \wedge (y \vee z)) = F \vee (\neg x \wedge (y \vee z)) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-13} \\
\textcolor{blue}{POS-LIdn} \left(F \vee (\neg x \wedge (y \vee z)) = \neg x \wedge (y \vee z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-14} \\
\textcolor{blue}{THM-LAsc-1} \left(\neg x \wedge B = \neg x \wedge ((x \vee y) \vee z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-15} \\
\textcolor{blue}{POS-LDis} \left(\neg x \wedge ((x \vee y) \vee z) = (\neg x \wedge (x \vee y)) \vee (\neg x \wedge z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-16} \\
\textcolor{blue}{POS-LDis} \left((\neg x \wedge (x \vee y)) \vee (\neg x \wedge z) = ((\neg x \wedge x) \vee (\neg x \wedge y)) \vee (\neg x \wedge z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-17} \\
\textcolor{blue}{POS-LCmp} \left(((\neg x \wedge x) \vee (\neg x \wedge y)) \vee (\neg x \wedge z) = (F \vee (\neg x \wedge y)) \vee (\neg x \wedge z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-18} \\
\textcolor{blue}{POS-LIdn} \left((F \vee (\neg x \wedge y)) \vee (\neg x \wedge z) = (\neg x \wedge y) \vee (\neg x \wedge z) \right),_1
\end{array}$$

$$\begin{array}{l}
\textcolor{teal}{THM-LAsc-19} \\
\textcolor{blue}{POS-LDis} \left((\neg x \wedge y) \vee (\neg x \wedge z) = \neg x \wedge (y \vee z) \right),_1
\end{array}$$

$$\begin{aligned}
& \begin{array}{l} \textcolor{teal}{THM-LAsc-20} \\ \textcolor{blue}{THM-LAsc-10} \\ \textcolor{blue}{THM-LAsc-11} \\ \textcolor{blue}{THM-LAsc-12} \\ \textcolor{blue}{THM-LAsc-13} \\ \textcolor{blue}{THM-LAsc-14} \\ \textcolor{blue}{THM-LAsc-15} \\ \textcolor{blue}{THM-LAsc-16} \\ \textcolor{blue}{THM-LAsc-17} \\ \textcolor{blue}{THM-LAsc-18} \\ \textcolor{blue}{THM-LAsc-19} \end{array} (\neg x \wedge A = \neg x \wedge (y \vee z) = \neg x \wedge B),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-21} \\ \textcolor{blue}{POS-LDis} \end{array} (A = A \wedge T),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-22} \\ \textcolor{blue}{POS-LCmp} \end{array} (A \wedge T = A \wedge (x \vee \neg x)),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-23} \\ \textcolor{blue}{POS-LDis} \end{array} (A \wedge (x \vee \neg x) = (x \wedge A) \vee (\neg x \wedge A)),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-24} \\ \textcolor{blue}{THM-LAsc-9} \end{array} ((x \wedge A) \vee (\neg x \wedge A) = (x \wedge B) \vee (\neg x \wedge B)),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-25} \\ \textcolor{blue}{THM-LAsc-20} \end{array} ((x \wedge B) \vee (\neg x \wedge A) = (x \wedge B) \vee (\neg x \wedge B)),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-26} \\ \textcolor{blue}{POS-LDis} \end{array} ((x \wedge B) \vee (\neg x \wedge B) = B \wedge (x \vee \neg x)),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-27} \\ \textcolor{blue}{POS-LCmp} \end{array} (B \wedge (x \vee \neg x) = B \wedge T),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-27} \\ \textcolor{blue}{POS-LIdn} \end{array} (B \wedge T = B),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-28} \\ \textcolor{blue}{THM-LAsc-21} \\ \textcolor{blue}{THM-LAsc-22} \\ \textcolor{blue}{THM-LAsc-23} \\ \textcolor{blue}{THM-LAsc-24} \\ \textcolor{blue}{THM-LAsc-25} \\ \textcolor{blue}{THM-LAsc-26} \\ \textcolor{blue}{THM-LAsc-27} \end{array} (A = B),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc-29} \\ \textcolor{blue}{THM-LAsc-28} \\ \textcolor{blue}{THM-LAsc-1} \end{array} (x \vee (y \vee z) = (x \vee y) \vee z),_1 \\
& \begin{array}{l} \textcolor{teal}{THM-LAsc} \\ \textcolor{blue}{THM-LAsc-29} \\ \textcolor{blue}{THM-Dual} \end{array} \left((x \vee (y \vee z) = (x \vee y) \vee z), (x \wedge (y \wedge z) = (x \wedge y) \wedge z) \right)
\end{aligned}$$

Associative (17)

$$\begin{aligned}
& \begin{array}{l} \textcolor{teal}{THM-LDMr-1} \\ \textcolor{blue}{POS-LDis} \end{array} \left((x \vee y) \vee (\neg x \wedge \neg y) = ((x \vee y) \vee \neg x) \wedge ((x \vee y) \vee \neg y) \right) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-2} \\ \textcolor{blue}{POS-LCom} \\ \textcolor{blue}{THM-LAsc} \end{array} \left(((x \vee y) \vee \neg x) \wedge ((x \vee y) \vee \neg y) = ((x \vee \neg x) \vee y) \wedge ((\neg y \vee y) \vee x) \right) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-3} \\ \textcolor{blue}{POS-LCmp} \end{array} \left(((x \vee \neg x) \vee y) \wedge ((\neg y \vee y) \vee x) = (T \vee y) \wedge (T \vee x) \right) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-4} \\ \textcolor{blue}{THM-LDom} \end{array} ((T \vee y) \wedge (T \vee x) = T \wedge T) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-5} \\ \textcolor{blue}{THM-LIdm} \end{array} (T \wedge T = T) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-6} \\ \textcolor{blue}{THM-LDMr-1} \\ \textcolor{blue}{THM-LDMr-2} \\ \textcolor{blue}{THM-LDMr-3} \\ \textcolor{blue}{THM-LDMr-4} \\ \textcolor{blue}{THM-LDMr-5} \end{array} ((x \vee y) \vee (\neg x \wedge \neg y) = T) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-7} \\ \textcolor{blue}{POS-LDis} \end{array} ((x \vee y) \wedge (\neg x \wedge \neg y) = (x \wedge \neg x \wedge \neg y) \vee (y \wedge \neg x \wedge \neg y)) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-8} \\ \textcolor{blue}{POS-LCom} \\ \textcolor{blue}{THM-LAsc} \end{array} \left((x \wedge \neg x \wedge \neg y) \vee (y \wedge \neg x \wedge \neg y) = ((x \wedge \neg x) \wedge \neg y) \vee ((y \wedge \neg y) \wedge \neg x) \right) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-9} \\ \textcolor{blue}{POS-LCmp} \end{array} \left(((x \wedge \neg x) \wedge \neg y) \vee ((y \wedge \neg y) \wedge \neg x) = (F \wedge \neg y) \vee (F \wedge \neg x) \right) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-10} \\ \textcolor{blue}{THM-LDom} \end{array} ((F \wedge \neg y) \vee (F \wedge \neg x) = F \vee F) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-11} \\ \textcolor{blue}{THM-LIdm} \end{array} (F \vee F = F) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-12} \\ \textcolor{blue}{THM-LDMr-7} \\ \textcolor{blue}{THM-LDMr-8} \\ \textcolor{blue}{THM-LDMr-9} \\ \textcolor{blue}{THM-LDMr-10} \\ \textcolor{blue}{THM-LDMr-11} \end{array} ((x \vee y) \wedge (\neg x \wedge \neg y) = F) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-13} \\ \textcolor{blue}{THM-LDMr-6} \\ \textcolor{blue}{THM-LDMr-12} \\ \textcolor{blue}{POS-LCmp} \end{array} \left(((x \vee y) \vee (\neg x \wedge \neg y) = T = (x \vee y) \vee \neg(x \vee y)), ((x \vee y) \wedge (\neg x \wedge \neg y) = F = (x \vee y) \wedge \neg(x \vee y)) \right) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr-14} \\ \textcolor{blue}{THM-LDMr-13} \\ \textcolor{blue}{THM-LUNt} \end{array} (\neg x \wedge \neg y = \neg(x \vee y)) \\
& \begin{array}{l} \textcolor{teal}{THM-LDMr} \\ \textcolor{blue}{THM-LDMr-14} \\ \textcolor{blue}{THM-Dual} \end{array} \left((\neg x \wedge \neg y = \neg(x \vee y)), (\neg x \vee \neg y = \neg(x \wedge y)) \right)
\end{aligned}$$

Boolean De Morgan's Laws (18)

$$\begin{aligned}
& \text{THM-ContrP-1} \text{ operatorIF } (x \implies y = (\neg x) \vee y) \\
& \text{THM-ContrP-2} \text{ POS-LCom THM-LInv } \left((\neg x) \vee y = ((\neg y) \vee (\neg x)) \right) \\
& \text{THM-ContrP-3} \text{ operatorIF } \left((\neg y) \vee (\neg x) = (\neg y) \implies (\neg x) \right) \\
& \text{THM-ContrP} \text{ THM-ContrP-1 THM-ContrP-2 THM-ContrP-3 } (x \implies y = (\neg y) \implies (\neg x)) \\
& \# \text{ Contrapositive Law} \quad (19)
\end{aligned}$$

MISC IMPLICATION LAWS:

$$\begin{aligned}
& (T \implies x = x) \\
& (F \implies x = T) \\
& (x \implies T = T) \\
& (x \implies F = \neg x) \\
& ((x \vee y) \implies z) = (x \implies z) \wedge (y \implies z) \\
& (x \implies (y \wedge z)) = (x \implies y) \wedge (x \implies z) \quad (20)
\end{aligned}$$

1.3 Predicate Logic

$$\text{predicate}[P] := \text{truth}[P(v_{free})] \quad (21)$$

$$\begin{aligned}
\text{universalQuantifier}[\forall][Q, P] &:= (\text{predicate}[Q]),_1 (\text{predicate}[P]),_1 \\
&(\forall_{arg_x(Q(x))} (P(x)) = Q(y_{free}) \implies P(y_{free})) \cdot_1 \quad (22)
\end{aligned}$$

$$\text{existentialQuantifier}[\exists][Q, P] := (\exists_{arg_x(Q(x))} (P(x)) = \neg \forall_{arg_x(Q(x))} (\neg P(x))) \quad (23)$$

$$\text{uniquenessQuantifier}[\exists!][Q, P] := (\exists!_{arg_x(Q(x))} (P(x)) = \exists_{arg_x(Q(x))} (P(x)) \wedge \forall_{arg_y(P(y))} \forall_{arg_z(P(z))} (y = z)) \quad (24)$$

$$0 \quad (25)$$

2 Dry Run

2.1 NaiveMaster

undefined terms: $:=, set, tuple, element, nnumber, \in, \subseteq, =, \not\subseteq, \cup, \cap, \emptyset, \{, \}, \langle, \rangle, |, \wedge, \times, relation, property, binaryRelation, domain, range, field, \forall, \exists, \wedge, \vee, \implies, \iff, \iff,$

$$(26)$$

$$\begin{aligned}
& \text{element}[x] \in \text{set}[y] \\
& \# \text{ x belongs to y} \quad (27)
\end{aligned}$$

$$\begin{aligned}
& \text{set}[x] \subseteq \text{set}[y] \\
& \# \text{ x is included in y} \quad (28)
\end{aligned}$$

$$(\text{set}[x] = \text{set}[y]) := (\text{set}[x] \subseteq \text{set}[y] \wedge \text{set}[y] \subseteq \text{set}[x])$$

x is the same set as y (29)

$$(\text{set}[x] \subset \text{set}[y]) := (\text{set}[x] \subseteq \text{set}[y] \wedge \text{set}[x] \neq \text{set}[y])$$

x is a proper subset of y (30)

$$\text{set}[x] \cup \text{set}[y]$$

all elements in x or y (31)

$$\text{set}[x] \cap \text{set}[y]$$

all elements in x and y (32)

$$\text{disjoint}[x, y] := \text{set}[x] \cap \text{set}[y] = \emptyset$$

disjoint sets do not intersect (33)

$$\text{set}[E] = \{e_1, e_2, e_3, \dots, e_n\}$$

unordered set containing $e_1, e_2, e_3, \dots, e_n$ (34)

$$\{e_1, e_2, e_3\} = \{e_3, e_1, e_2\}$$

$$\text{tuple}[E] = \langle e_1, e_2, e_3, \dots, e_n \rangle$$

ordered tuple containing $e_1, e_2, e_3, \dots, e_n$ (35)

$$\langle e_1, e_2, e_3 \rangle \neq \langle e_2, e_3, e_1 \rangle$$

$$\text{set}[X]^{\text{number}[k]}$$

set of all ordered k-tuples from the elements in X (36)

$$X^1 = \{e_1, e_2, e_3, \dots, e_n\}^1 = \{\langle e_1 \rangle, \langle e_2 \rangle, \langle e_3 \rangle, \dots, \langle e_n \rangle\} = \{e_1, e_2, e_3, \dots, e_n\} = X$$

$$\text{set}[Y] \times \text{set}[Z]$$

Cartesian product (37)

$$\text{relation}[R][S, k] := R \subseteq \text{set}[S]^{\text{number}[k]}$$

k-tuple relation R on the set S takes only tuples that satisfy some relation (38)

$$\text{property}[P][S] := \text{relation}[P][S, 1] \subseteq \text{set}[S]^1 = S$$

property P of the set S (39)

use defeq to keep the object typing instead of
ref-eq which reduces it to a propositional truth value.

But, how do I chain ref-eq tho? (40)

$$\text{binaryRelation}[B][S] = \text{relation}[B][S, 2] \subseteq \text{set}[S]^2$$

$xBy = \langle x, y \rangle \in B$ (41)

$$\text{domain}[X][B, S] = \{x \mid \langle x, y \rangle \in \text{binaryRelation}[B][S]\}$$

(42)

$$\text{range}[Y][B, S] = \{y \mid \langle x, y \rangle \in \text{binaryRelation}[B][S]\}$$

(43)

$$\textit{field}[F][B, S] = \textit{domain}[X][B, S] \cup \textit{range}[Y][B, S] \quad (44)$$

$$\textit{inverseRelation}[B^{-1}][B, S] := \{\langle y, x \rangle \mid \langle x, y \rangle \in \textit{binaryRelation}[B][S]\} \quad (45)$$

$$\textit{reflexive}[B][S] := \forall_{x \in \textit{field}[F][B, S]} (xBx) \quad (46)$$

$$\textit{symmetric}[B][S] := \forall_{x, y \in S} (xB y \implies yB x) \quad (47)$$

$$\textit{transitive}[B][S] := \forall_{x, y, z \in S} ((xB y \wedge yB z) \implies xB z) \quad (48)$$

$$\textit{equivalenceRelation}[B][S] := (\textit{reflexive}[B][S] \wedge \textit{symmetric}[B][S] \wedge \textit{transitive}[B][S]) \quad (49)$$

$$\begin{aligned} & \textit{equivalenceClass}[[y]][y, B, S] := \{z \in \textit{field}[F][B, S] \mid yB z\} \\ & (\textit{equivalenceClass}[[u]][u, B, S] = \textit{equivalenceClass}[[v]][v, B, S]) \iff (uB v) \\ & (\textit{equivalenceClass}[[u]][u, B, S] \neq \textit{equivalenceClass}[[v]][v, B, S]) \implies ([u] \cap [v] = \emptyset) \\ & \# \text{ a set can be partitioned by equivalence classes} \end{aligned} \quad (50)$$

$$\textit{function}[f][S] := (\forall_{x, y, z \in S} (\langle x, y \rangle \in f \wedge \langle x, z \rangle \in f) \implies (y = z)) \quad (51)$$