

Next-Next-Gen Notes

Object-Oriented Maths

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1 Mathematical Logic

1.1 NaiveMaster

(1)

set, element, \in

(2)

element[x] \in *set*[y] := x belongs to y

(3)

$x \subseteq y$:= x is included in y

(4)

$x = y$:= x is the same thing as y
:= $x \subseteq y, y \subseteq x$

(5)

$x \subset y, x \not\subseteq y$:= proper subset
:= $x \neq y, x \subseteq y$

(6)

$x \cup y$:= all elements in x or y

(7)

$x \cap y$:= all elements in x and y

(8)

disjoint(x, y) := disjoint sets
:= $x \cap y = \emptyset$

(9)

$\{e_1, e_2, e_3, \dots, e_n\}$:= unordered set containing $e_1, e_2, e_3, \dots, e_n$
 $\{e_1, e_2, e_3\} = \{e_3, e_1, e_2\}$

(10)

$\langle e_1, e_2, e_3, \dots, e_n \rangle$:= ordered tuple containing $e_1, e_2, e_3, \dots, e_n$
 $\langle e_1, e_2, e_3 \rangle \neq \langle e_2, e_3, e_1 \rangle$

(11)

$X^k = \{e_1, e_2, e_3, \dots, e_n\}^k$:= set of all ordered k -tuples from the elements of $e_1, e_2, e_3, \dots, e_n$
 $X^1 = \{e_1, e_2, e_3, \dots, e_n\}^1 = \{\langle e_1 \rangle, \langle e_2 \rangle, \langle e_3 \rangle, \dots, \langle e_n \rangle\} = \{e_1, e_2, e_3, \dots, e_n\} = X$

(12)

$Y \times Z = \{y_1, y_2, y_3, \dots, y_i\} \times \{z_1, z_2, z_3, \dots, z_j\}$:= Cartesian product
:= $\bigcup_{a \leq i, b \leq j} (\{\langle y_a, z_b \rangle\})$

(13)

$$R_Y^k \subseteq Y^k := \text{k-tuple relation } R \text{ on the set } Y \text{ takes only tuples that satisfy some relation}$$

$$P_Y \subseteq Y := \text{property } P \text{ of the set } Y \quad (14)$$

$$\begin{aligned} \langle y, z \rangle \in \text{binaryRelation}(R_X^2) &= yR_X^2 z \\ \text{domain}(Y), \text{range}(Z) \\ \text{field}(R) &= Y \cup Z \\ \langle a, b \rangle \in \text{inverse}(R^{-1}) &: \langle b, a \rangle \in R \\ \text{reflexive}(R_X^2) &: xR_X^2 x \\ \text{symmetric}(R_X^2) &: xR_X^2 y = yR_X^2 x \\ \text{transitive}(R_X^2) &: xR_X^2 y, yR_X^2 z : xR_X^2 z \\ \text{equivalenceRelation}(R_X^2) &:= \text{reflexive}(R_X^2), \text{symmetric}(R_X^2), \text{transitive}(R_X^2) \end{aligned} \quad (15)$$

$$\text{take this into shitt more srsly} \quad (16)$$

2 Logic and Set Theory

2.1 Logical Truths and Operators

$$(17)$$

$$\text{truth}[t] := t = \begin{cases} T \\ F \end{cases} \quad (18)$$

$$\text{statement}[s] := \text{correctSyntaxSemantics}^1[s] \quad (19)$$

$$\text{proposition}[s, t] := (\text{statement}[s]), (\text{truth}[t]). \quad (20)$$

$$\text{operatorOR}[\vee][x, y] := (\text{truth}[x]), (\text{truth}[y]), \left(\text{truth}[x \vee y] = \begin{cases} F & x=F, y=F \\ T & x=F, y=T \\ T & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (21)$$

$$\text{operatorAND}[\wedge][x, y] := (\text{truth}[x]), (\text{truth}[y]), \left(\text{truth}[x \wedge y] = \begin{cases} F & x=F, y=F \\ F & x=F, y=T \\ F & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (22)$$

$$\text{operatorNOT}[\neg][x] := (\text{truth}[x]), \left(\text{truth}[\neg x] = \begin{cases} T & x=F \\ F & x=T \end{cases} \right). \quad (23)$$

$$\begin{aligned} \text{booleanAlgebra}[\{T, F\}, \wedge, \vee, \neg] &:= \text{POS-LCom}((x \wedge y = y \wedge x), (x \vee y = y \vee x)) \# \text{Commutative,} \\ \text{POS-LDis}((x \wedge (y \vee z) &= (x \wedge y) \vee (x \wedge z)), (x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z))) \# \text{Distributive,} \\ \text{POS-LIdn}((x \wedge T &= x), (x \vee F = x)) \# \text{Identity,} \end{aligned}$$

$$POS-LCmp((x \wedge \neg x = F), (x \vee \neg x = T)) \# \text{ Complement.} \quad (24)$$

$$operatorXOR[\vee][x, y] := (\text{truth}[x][\square]), (\text{truth}[y][\square]), \left(\text{truth}[x \vee y][\square] = \begin{cases} F & x=F, y=F \\ T & x=F, y=T \\ T & x=T, y=F \\ F & x=T, y=T \end{cases} \right). \quad (25)$$

$$operatorIF[\implies][x, y] := (\text{truth}[x][\square]), (\text{truth}[y][\square]), \left(\text{truth}[x \implies y][\square] = (\neg x) \vee y = \begin{cases} T & x=F, y=F \\ T & x=F, y=T \\ F & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (26)$$

$$\begin{aligned} & THM-LExp-1(F = x \wedge \neg x) \implies \\ & POS-LCmp \\ & THM-LExp-2(x), \\ & THM-LExp-1(x), \\ & THM-LExp-3(\neg x), \\ & THM-LExp-1(\neg x), \\ & THM-LExp-4(x \vee y), \\ & THM-LExp-2(x \vee y), \\ & THM-LExp-5(y). \\ & THM-LExp-4(y). \\ & THM-LExp-3 \\ & THM-LExp-1(F \implies y) \\ & THM-LExp-1 \\ & THM-LExp-2 \\ & THM-LExp-3 \\ & THM-LExp-4 \\ & THM-LExp-5 \end{aligned}$$

$$\# \text{ The Principle of Explosion, anything follows from a false (F) premise} \quad (27)$$

$$operatorOIF[\Leftarrow][x, y] := (\text{truth}[x][\square]), (\text{truth}[y][\square]), \left(\text{truth}[x \Leftarrow y][\square] = (\neg y) \vee x = \begin{cases} T & x=F, y=F \\ F & x=F, y=T \\ T & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (28)$$

$$\begin{aligned} & operatorIIF[\Leftrightarrow][x, y] := (\text{truth}[x][\square]), (\text{truth}[y][\square]), \\ & \left(\text{truth}[x \Leftrightarrow y][\square] = (x \implies y) \wedge (y \implies x) = \begin{cases} T & x=F, y=F \\ F & x=F, y=T \\ F & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (29) \end{aligned}$$