

Next-Next-Gen Notes

Object-Oriented Maths

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Format: $characteristic((subjects), (dependencies)) \iff (conditions(dependencies)) \wedge (conditions(subjects))$

Note: All weaker objects automatically induces notions inherited from stronger objects.

TODO assign free variables as parameters

TODO define \parallel abs cross-product and other missing refs

TODO distinguish new condition vs implied proposition - separate propositions into new line thms

TODO silent link expressions! - e.g. *backslashsilentPLPL_X*

1 Logic and Set Theory

1.1 Logical Truths and Operators

$$truth[t] := t = \begin{cases} T \\ F \end{cases} \quad (1)$$

$$statement[s] := correctSyntaxSemantics[s] \quad (2)$$

$$proposition[s, t] := (statement[s], (truth[t])). \quad (3)$$

$$operatorOR[\vee][x, y] := (truth[x], (truth[y]), \left(truth[x \vee y] = \begin{cases} F & x=F, y=F \\ T & x=F, y=T \\ T & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (4)$$

$$operatorAND[\wedge][x, y] := (truth[x], (truth[y]), \left(truth[x \wedge y] = \begin{cases} F & x=F, y=F \\ F & x=F, y=T \\ F & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (5)$$

$$operatorNOT[\neg][x] := (truth[x], \left(truth[\neg x] = \begin{cases} T & x=F \\ F & x=T \end{cases} \right). \quad (6)$$

$$\begin{aligned} booleanAlgebra[\{T, F\}, \wedge, \vee, \neg] &:= POS-LCom((x \wedge y = y \wedge x), (x \vee y = y \vee x)) \# \text{Commutative,} \\ POS-LDis &\left((x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)), (x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)) \right) \# \text{Distributive,} \\ POS-LIdn &((x \wedge T = x), (x \vee F = x)) \# \text{Identity,} \\ POS-LCmp &((x \wedge \neg x = F), (x \vee \neg x = T)) \# \text{Complement.} \end{aligned} \quad (7)$$

$$\text{operator } XOR[\vee][x, y] := (\text{truth}[x][\square]), (\text{truth}[y][\square]), \left(\text{truth}[x \vee y][\square] = \begin{cases} F & x=F, y=F \\ T & x=F, y=T \\ T & x=T, y=F \\ F & x=T, y=T \end{cases} \right). \quad (8)$$

$$\text{operator } IF[\implies][x, y] := (\text{truth}[x][\square]), (\text{truth}[y][\square]), \left(\text{truth}[x \implies y][\square] = (\neg x) \vee y = \begin{cases} T & x=F, y=F \\ T & x=F, y=T \\ F & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (9)$$

$$\begin{aligned} & \text{THM-LExp-1} \text{ POS-LCmp } (F = x \wedge \neg x) \implies \\ & \quad \text{THM-LExp-2} (x), \\ & \quad \text{THM-LExp-3} (\neg x), \\ & \quad \text{THM-LExp-4} (x \vee y), \\ & \quad \text{THM-LExp-5} (y). \\ & \quad \text{THM-LExp-1} (F \implies y) \\ & \quad \text{THM-LExp-2} \\ & \quad \text{THM-LExp-3} \\ & \quad \text{THM-LExp-4} \\ & \quad \text{THM-LExp-5} \end{aligned}$$

$$\# \text{ The Principle of Explosion, anything follows from a false (F) premise} \quad (10)$$

$$\text{operator } OIF[\Leftarrow][x, y] := (\text{truth}[x][\square]), (\text{truth}[y][\square]), \left(\text{truth}[x \Leftarrow y][\square] = \begin{cases} T & x=F, y=F \\ F & x=F, y=T \\ T & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (11)$$

$$\text{operator } IIF[\Leftrightarrow][x, y] := (\text{truth}[x][\square]), (\text{truth}[y][\square]), \left(\text{truth}[x \Leftrightarrow y][\square] = \begin{cases} T & x=F, y=F \\ F & x=F, y=T \\ F & x=T, y=F \\ T & x=T, y=T \end{cases} \right). \quad (12)$$

1.2 Boolean Algebra Properties

$$\begin{aligned} & \text{THM-Dual-1} \text{ POS-LCom } \left(\text{booleanAlgebra}[\{T, F\}, \wedge, \vee, \neg][\square] \Leftrightarrow \right. \\ & \quad ((x \vee y = y \vee x), (x \wedge y = y \wedge x)) \# \text{ Reordered Commutative,} \\ & \quad ((x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)), (x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z))) \# \text{ Reordered Distributive,} \\ & \quad ((x \vee F = x), (x \wedge T = x)) \# \text{ Reordered Identity,} \\ & \quad ((x \vee \neg x = T), (x \wedge \neg x = F)) \# \text{ Reordered Complement. } \Leftrightarrow \\ & \quad \left. \text{booleanAlgebra}[\{F, T\}, \vee, \wedge, \neg][\square] \right) \\ & \text{THM-Dual} \\ & \text{THM-Dual-1} (\text{booleanAlgebra}[\{T, F\}, \wedge, \vee, \neg][\square] \Leftrightarrow \text{booleanAlgebra}[\{F, T\}, \vee, \wedge, \neg][\square]) \end{aligned}$$

$$\# \text{ Boolean Algebra Duality follows from the swap symmetry of } (\wedge, T) \text{ and } (\vee, F) \text{ within the axioms} \quad (13)$$

$$\begin{aligned}
& \text{THM-LUNt-1}((x \vee y = T = x \vee z) \wedge (x \wedge y = F = x \wedge z)) \implies \\
& \quad \text{THM-LUNt-2} \text{ POS-LIdn} (y = y \wedge T), \\
& \quad \text{THM-LUNt-3} (y \wedge T = y \wedge (x \vee z)), \\
& \quad \text{THM-LUNt-4} \text{ POS-LDis} (y \wedge (x \vee z) = (y \wedge x) \vee (y \wedge z)), \\
& \quad \text{THM-LUNt-5} \text{ POS-LCom} ((y \wedge x) \vee (y \wedge z) = (x \wedge z) \vee (y \wedge z)), \\
& \quad \text{THM-LUNt-6} \text{ POS-LCom} \text{ POS-LDis} ((x \wedge z) \vee (y \wedge z) = z \wedge (x \vee y)), \\
& \quad \text{THM-LUNt-7} \text{ THM-LUNt-1} (z \wedge (x \vee y) = z \wedge T), \\
& \quad \text{THM-LUNt-8} \text{ POS-LIdn} (z \wedge T = z). \\
& \text{THM-LUNt} \text{ THM-LUNt-1} \text{ THM-LUNt-2} \text{ THM-LUNt-3} \text{ THM-LUNt-4} \text{ THM-LUNt-5} \text{ THM-LUNt-6} \text{ THM-LUNt-7} \text{ THM-LUNt-8} \\
& \left(((x \vee y = T = x \vee z) \wedge (x \wedge y = F = x \wedge z)) \implies (y = z) \right)
\end{aligned}$$

Uniqueness of Complements (14)

$$\begin{aligned}
& \text{THM-LDom-1} \text{ POS-LIdn} (x \vee T = (x \vee T) \wedge T), \\
& \text{THM-LDom-2} \text{ POS-LCmp} ((x \vee T) \wedge T = (x \vee T) \wedge (x \vee \neg x)), \\
& \text{THM-LDom-3} \text{ POS-LDis} ((x \vee T) \wedge (x \vee \neg x) = x \vee (T \wedge \neg x)), \\
& \text{THM-LDom-4} \text{ POS-LIdn} (x \vee (T \wedge \neg x) = x \vee \neg x), \\
& \text{THM-LDom-5} \text{ POS-LCmp} (x \vee \neg x = T). \\
& \text{THM-LDom-6} \text{ THM-LDom-1} \text{ THM-LDom-2} \text{ THM-LDom-3} \text{ THM-LDom-4} \text{ THM-LDom-5} \\
& \text{THM-LDom} \text{ THM-LDom-6} \text{ THM-Dual} ((x \vee T = T), (x \wedge F = F)).
\end{aligned}$$

Domination (15)

$$\begin{aligned}
& \text{THM-LIdm-1} \text{ POS-LIdn} (x \vee x = (x \vee x) \wedge T), \\
& \text{THM-LIdm-2} \text{ POS-LCmp} ((x \vee x) \wedge T = (x \vee x) \wedge (x \vee \neg x)), \\
& \text{THM-LIdm-3} \text{ POS-LDis} ((x \vee x) \wedge (x \vee \neg x) = x \wedge (x \vee \neg x)), \\
& \text{THM-LIdm-4} \text{ POS-LCmp} (x \wedge (x \vee \neg x) = x \wedge T), \\
& \text{THM-LIdm-5} \text{ POS-LIdn} (x \wedge T = x), \\
& \text{THM-LIdm-6} \text{ THM-LIdm-1} \text{ THM-LIdm-2} \text{ THM-LIdm-3} \text{ THM-LIdm-4} \text{ THM-LIdm-5} \\
& \text{THM-LIdm} \text{ THM-LIdm-6} \text{ THM-Dual} ((x \vee x = x), (x \wedge x = x)).
\end{aligned}$$

Idempotent (16)

$$\begin{aligned}
& \text{THM-LInv-1} \text{ POS-LIdn} (\neg \neg x = \neg \neg x \vee F), \\
& \text{THM-LInv-2} \text{ POS-LCmp} (\neg \neg x \vee F = \neg \neg x \vee (x \wedge \neg x)), \\
& \text{THM-LInv-3} \text{ POS-LDis} (\neg \neg x \vee (x \wedge \neg x) = (\neg \neg x \vee x) \wedge (\neg \neg x \vee \neg x)), \\
& \text{THM-LInv-4} \text{ POS-LCmp} ((\neg \neg x \vee x) \wedge (\neg \neg x \vee \neg x) = (\neg \neg x \vee x) \wedge T), \\
& \text{THM-LInv-5} \text{ POS-LCmp} ((\neg \neg x \vee x) \wedge T = (\neg \neg x \vee x) \wedge (x \vee \neg x)), \\
& \text{THM-LInv-6} \text{ POS-LDis} ((\neg \neg x \vee x) \wedge (x \vee \neg x) = x \vee (\neg \neg x \wedge \neg x)),
\end{aligned}$$

$$\begin{array}{l} \textcolor{teal}{THM-LInv-7} \\ \textcolor{blue}{POS-LCmp} \end{array} (x \vee (\neg\neg x \wedge \neg x) = x \vee F),$$

$$\begin{array}{l} \textcolor{teal}{THM-LInv-8} \\ \textcolor{blue}{POS-LIdn} \end{array} (x \vee F = x),$$

$$\begin{array}{l} \textcolor{teal}{THM-LInv} \\ \textcolor{blue}{THM-LInv-1} \end{array} (\neg\neg x = x).$$

$$\textcolor{blue}{THM-LInv-2}$$

$$\textcolor{blue}{THM-LInv-3}$$

$$\textcolor{blue}{THM-LInv-4}$$

$$\textcolor{blue}{THM-LInv-5}$$

$$\textcolor{blue}{THM-LInv-6}$$

$$\textcolor{blue}{THM-LInv-7}$$

$$\textcolor{blue}{THM-LInv-8}$$

Involution (17)

$$\begin{array}{l} \textcolor{teal}{THM-LAbs-1} \\ \textcolor{blue}{POS-LIdn} \end{array} (x \vee (x \wedge y) = (x \wedge T) \vee (x \wedge y)),$$

$$\begin{array}{l} \textcolor{teal}{THM-LAbs-2} \\ \textcolor{blue}{POS-LDis} \end{array} ((x \wedge T) \vee (x \wedge y) = x \wedge (T \vee y)),$$

$$\begin{array}{l} \textcolor{teal}{THM-LAbs-3} \\ \textcolor{blue}{THM-LDom} \end{array} (x \wedge (T \vee y) = x \wedge T),$$

$$\begin{array}{l} \textcolor{teal}{THM-LAbs-4} \\ \textcolor{blue}{THM-LIdn} \end{array} (x \wedge T = x),$$

$$\begin{array}{l} \textcolor{teal}{THM-LAbs-5} \\ \textcolor{blue}{THM-LAbs-1} \end{array} (x \vee (x \wedge y) = x),$$

$$\textcolor{blue}{THM-LAbs-2}$$

$$\textcolor{blue}{THM-LAbs-3}$$

$$\textcolor{blue}{THM-LAbs-4}$$

$$\begin{array}{l} \textcolor{teal}{THM-LAbs} \\ \textcolor{blue}{THM-LAbs-5} \\ \textcolor{blue}{THM-Dual} \end{array} ((x \vee (x \wedge y) = x), (x \wedge (x \vee y) = x)).$$

Absorption (18)

$$\begin{array}{l} \textcolor{teal}{000} \\ \textcolor{blue}{000} \end{array} ()$$

Associative (19)

$$\begin{array}{l} \textcolor{teal}{000} \\ \textcolor{blue}{000} \end{array} ()$$

Boolean De Morgan's Laws (20)

000TODOIFPROPERTIES (21)