

Expanders

CS292F

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Lecture 14

Complete graph K_n : $\binom{n}{2} = \frac{n(n-1)}{2}$ edges

$$\text{Laplacian} = \begin{pmatrix} n-1 & & -1 \\ & \ddots & \\ -1 & & n-1 \end{pmatrix} = nI - \mathbf{1}\mathbf{1}^T$$

Eigenvalues: $\lambda_1 = 0$

If $\mathbf{1}^T x = 0$ then

$$K_n x = nx - \mathbf{1}\mathbf{1}^T x = nx - \mathbf{1} \cdot 0 = nx$$

so $\lambda_2 = \lambda_3 = \dots = \lambda_n = n$.

Expanders

Intuition: Sparse approximation of K_n

Definitions:

TRADITIONAL: sets of vtxs have large neighborhoods (large conductance)

SPIELMAN: Eigenvalues of adjacency mtr. (small)

US: Laplacian eigenvalues.

DEF: $A_n(\varepsilon, d)$ -expander is a graph with n vertices with:

→ every vertex has degree d .

→ $(1-\varepsilon)d \leq \lambda_i \leq (1+\varepsilon)d$

for $i=2, 3, \dots, n$

$$\# \text{edges} = \frac{dn}{2} = \frac{d}{2}n.$$

Theorem: If G is an (ε, d) -expander
and $H = \frac{d}{n} K_n$, then

$$(1-\varepsilon) H \preceq G \preceq (1+\varepsilon) H$$

" G is an ε -approximation of H "

Proof: will compare $x^T G x$ and $x^T H x$
for $\mathbf{1}^T x = 0$.

If $\mathbf{1}^T x = 0$,

$$(1-\varepsilon)d \leq \lambda_2 \leq \frac{x^T G x}{x^T x} \leq \lambda_n \leq (1+\varepsilon)d$$

$$\Rightarrow (1-\varepsilon)d x^T x \leq x^T G x \leq (1+\varepsilon)d x^T x.$$

$$x^T H x = x^T \frac{d}{n} K_n x = x^T \frac{d}{n} n x = d x^T x.$$

$$\Rightarrow (1-\varepsilon)x^T H x \leq x^T G x \leq (1+\varepsilon)x^T H x.$$

$$\Rightarrow (1-\varepsilon)H \preceq G \preceq (1+\varepsilon)H.$$

G is an ε - "spectral approximation" of H .

Recall $\|A\| = \max_{x \neq 0} \frac{\|Ax\|}{\|x\|}$
 $= \max_i |\lambda_i|$ (if A symmetric)

$$(1-\varepsilon)H \leq G \leq (1+\varepsilon)H$$

$$\Rightarrow -\varepsilon H \leq G - H \leq \varepsilon H$$

since evals of H are all 0 or d ,

$$\text{evals}(G - H) \leq \varepsilon d$$

$$\Rightarrow \|G - H\| \leq \varepsilon d$$

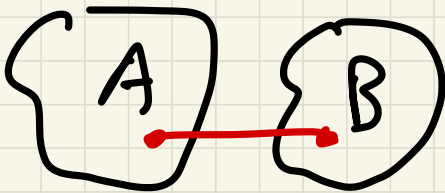
"Random" d -regular graphs are expanders. (WHP)

Explicit constructions of expanders

date to ~ 1980

Spielman does one in ch. 30

Expanders act a lot like random grfs.
Take 2 disjoint sets of vertices



$$|A| = \alpha n \quad |B| = \beta n$$

$$E(A, B) = \# \text{edges with one end in } A \text{ and one in } B \\ = -\mathbf{1}_A^T G \mathbf{1}_B \quad (\text{exercise})$$

$$\text{Complete graph: } E_K(A, B) = \alpha \beta n^2$$

$$H = \frac{d}{n} K_n : -\mathbf{1}_A^T H \mathbf{1}_B = \alpha \beta d n$$

Pick $A+B$ at random in any d -regular grf:

$$E_{\text{expect}}(A, B) = (\# \text{edges}) \cdot (\alpha \beta + \beta \alpha) \\ = \frac{dn}{2} \cdot 2\alpha\beta = \alpha\beta dn.$$

THM (Spielman 27.3.1):

In an expander, $E(A, B)$ is close to $\alpha\beta dn$ for all sets $A \subseteq B$.

Precisely,

$$\left| E(A, B) - \alpha\beta dn \right| \leq \varepsilon dn \sqrt{(\alpha - \alpha^2)(\beta - \beta^2)}.$$

Why?

$$E(A, B) = -\mathbf{1}_A^T G \mathbf{1}_B$$

and $G \approx H$ means

$$-\mathbf{1}_A^T G \mathbf{1}_B \approx -\mathbf{1}_A^T H \mathbf{1}_B = \alpha\beta dn$$

$$N(A) = \{ \text{vertices adj to vtxs in } A \}$$

Theorem (Tanner): Let G be a d -regular graph that ε -approximates $H = \frac{d}{n} K_n$.

Then for all $A \subset V$: ($|A| = \alpha n$)

$$|N(A)| \geq \frac{|A|}{\varepsilon^2(1-\alpha) + \alpha} \quad \text{"expander"}$$

Proof: Let $R = N(A)$, $B = V - R$.

$$E(A, B) = 0. \text{ Say } |B| = \beta n.$$

Then

$$\alpha \beta d n \leq \varepsilon d n \sqrt{(\alpha - \alpha^2)(\beta - \beta^2)}$$

Say $|R| = |N(A)| = \gamma n$.

$$\alpha^2 \beta^2 \leq \varepsilon^2 (\alpha - \alpha^2) (\beta - \beta^2)$$

$$\alpha \beta \leq \varepsilon^2 (1 - \alpha) (1 - \beta)$$

$$\alpha\beta \leq \varepsilon^2 (1-\alpha)(1-\beta)$$

$$\frac{\beta}{1-\beta} \leq \varepsilon^2 \frac{1-\alpha}{\alpha}$$

$$\begin{aligned} \beta &= V - N(A) \\ \beta_n &= V - N(A_n) \\ \text{so } \beta &= 1 - \gamma \end{aligned}$$

$$\frac{1-\gamma}{\gamma} \leq \varepsilon^2 \left(\frac{1-\alpha}{\alpha} \right)$$

$$\frac{1}{\gamma} - 1 \leq \varepsilon^2 \left(\frac{1-\alpha}{\alpha} \right)$$

$$\frac{1}{\gamma} \leq \frac{\varepsilon^2(1-\alpha) + \alpha}{\alpha}$$

$$\gamma \geq \frac{\alpha}{\varepsilon^2(1-\alpha) + \alpha}$$

(times n)

$$|N(A)| \geq \frac{|A|}{\varepsilon^2(1-\alpha) + \alpha}$$

QED