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Course: CS590-A Algorithms

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Description: Homework 1 Algorithms

## Problem 1

**Definition 1.**  $a \mid b$  ("a divides b") if and only if there exists some integer  $k$  such that  $b = ak$ . Equivalently,  $a \mid b$  if and only if  $b$  has a remainder of 0 when divided by  $a$  (see question 2).

1. Using the formal definition of divisibility above, prove that there exist positive integers  $a$ ,  $b$ , and  $c$  such that  $a \mid bc$ , but  $a \nmid b$  and  $a \nmid c$ .

**Proof.** Let  $a = 6$ ,  $b = 2$ , and  $c = 3$ . Then  $bc = 6$  and  $a \mid bc$  since  $6 = 6 \cdot 1$ . However,  $a \nmid b$  and  $a \nmid c$  since  $2 = 6 \cdot 0 + 2$  and  $3 = 6 \cdot 0 + 3$ . ■

**Theorem 1.** *The Division Algorithm. For any integers  $a$  and  $b$  where  $b \neq 0$ , there exist a unique pair of integers  $q$  and  $r$  such that  $a = qb + r$  and  $0 \leq r < b$ . The integers  $q$  and  $r$  are known as the quotient and remainder of  $a \div b$ , respectively.*

2. Using the formal definition of the remainder above, prove that if  $n$  and  $m$  are positive integers such that  $n$  has a remainder of  $r$  when divided by  $m$  and  $r < \sqrt{m}$ ,  $n^2$  has a remainder of  $r^2$  when divided by  $m$ .

**Proof.** Let  $n$ , and  $m$  be positive integers such that  $n \div m$  has a remainder  $r$ .

By Theorem 1,  $n = qm + r$  and  $m = qn^2 + r^2$ .

$$\sqrt{m} = \sqrt{qn^2 + r^2}$$

$$\sqrt{m} = \sqrt{qn^2} + \sqrt{r^2}$$

$$\sqrt{m} = n\sqrt{q} + r.$$

3. Use the formal definition of Big-Oh to prove that if  $f(n) = n^x + an^y$ , where  $a$ ,  $x$ , and  $y$  are positive integers such that  $x > y$ ,  $f(n) = O(n^x)$ .

*Big-Oh is defined as  $f(n) = O(g(n))$  if there exists a positive constant  $c$  and a positive integer  $n_0$  such that  $f(n) \leq cg(n)$  for all  $n \geq n_0$ .*

**Solution.** In this case,  $f(n) = n^x + an^y$  and  $g(n) = n^x$ . Let us assume  $y = x$ . Then  $f(n) = n^x + an^x$  which can be reduced to  $f(n) = n^x(1 + a)$ . Let  $1 + a = c$  and  $n_0 = 1$  to establish the upper bound. Then  $f(n) = n^x(1 + a) \leq cn^x$  for all  $n \geq n_0$ . Therefore,  $f(n) = O(n^x)$ . ■

4. Use the formal definition of Big-Omega to prove that if  $f_1(n)$ ,  $f_2(n)$ ,  $g_1(n)$ , and  $g_2(n)$  are functions such that  $f_1(n) = \Omega(g_1(n))$  and  $f_2(n) = \Omega(g_2(n))$ ,  $f_1(n) + f_2(n) = \Omega(\max(g_1(n), g_2(n)))$ .

**Solution.** Big Omega is defined as  $f(n) = \Omega(g(n))$  if there exists a positive constant  $c$  and a positive integer  $n_0$  such that  $f(n) \geq cg(n)$  for all  $n \geq n_0$ .

For  $f_1(n)$ ,  $n \geq c_1n$  for say  $n_0 = 0$  and  $c_1 = 0.5$ . The same is true for  $f_2(n)$  and  $g_2(n)$ . All of the functions are separated by a constant factor. Given that  $f_1(n)$  and  $f_2(n)$  were both larger than  $g_1(n)$  and  $g_2(n)$  the sum of the two must be larger than the maximum of either  $g_1(n)$  or  $g_2(n)$ . ■