# Mapping Integers to the Smallest Integer Possible

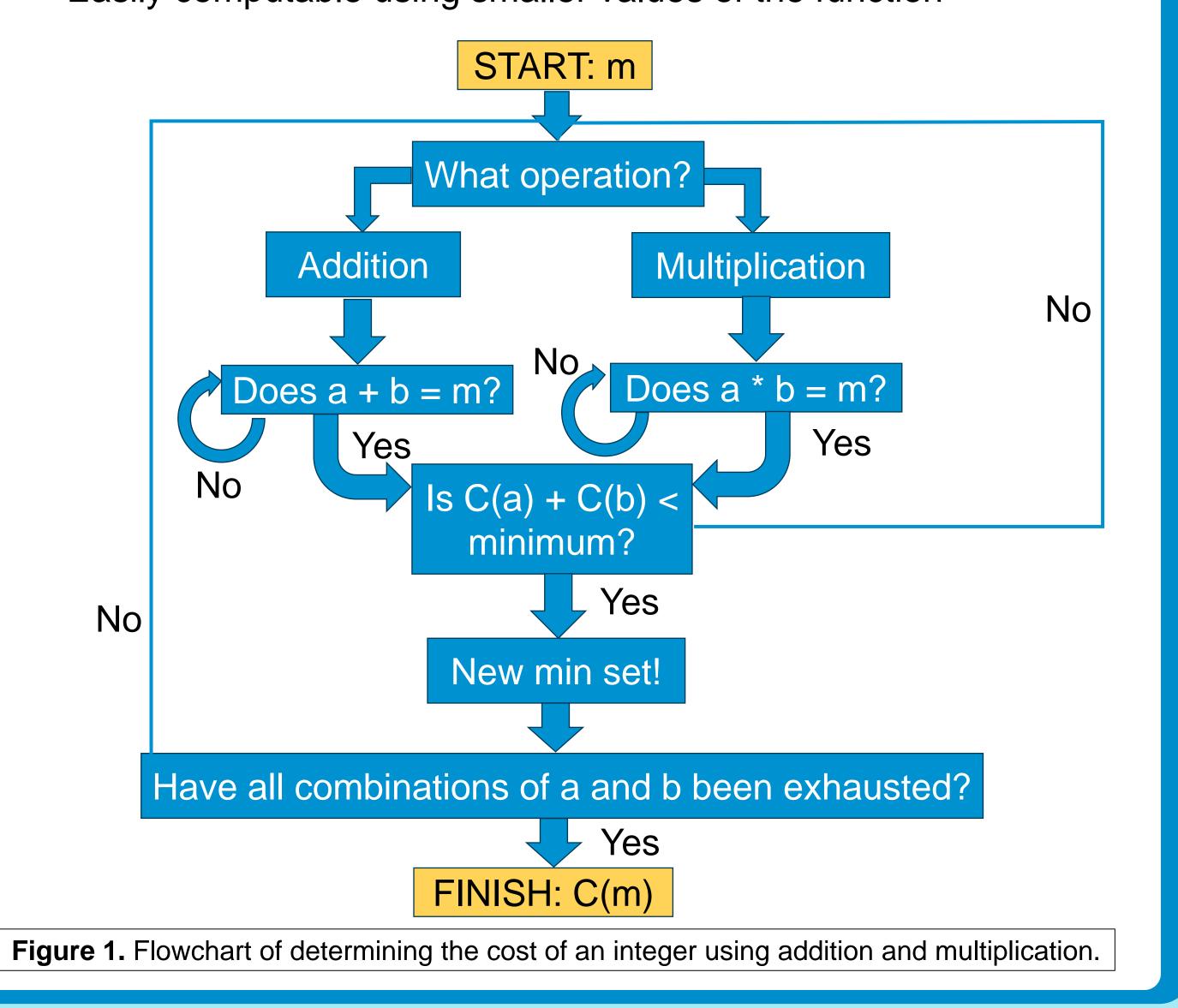
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## **Background and Motivation**

- If we are given an arbitrary computer program and an input, will a program always give an output?
- Cannot say "yes" to above due to real-world limitations:
  - Limited size of memory (e.g., disk drive)
  - Replicability on all inputs
  - Program execution time
  - Programs that take itself as input
- Project Goal:
  - Construct a mathematical function in Python mapping integers to integers based on a set of algorithms

# **Original Cost Function**

- Attempt at computability: let m be an input and C(m) its output
- Consider the original definition of the cost of an integer, m:
- Given two operands, a and b, such that if a + b = m, or a \* b = m,
   the goal is to minimize C(a) + C(b)
- The current minimum starts with the input and follows the chart
- Easily computable using smaller values of the function



### **Cost Function Implementation**

#### **Cost Function Framework**

- Norfolk's work proposes loose principles for the cost function:
  - Only a finite number of functions can be considered
  - One operation can be evaluated at a time
  - The function gives a charge to the operation itself
  - Any number of operands can be in an operation
  - The minimum can decrease according to the driver algorithm
- Any iteration (function-wide) has a fixed upper bound!
   Driver Algorithm
- The cost function contains several sub-algorithms
- If the cost, C(m), using just one algorithm is smaller than the minimum assigned, the cost becomes the minimum

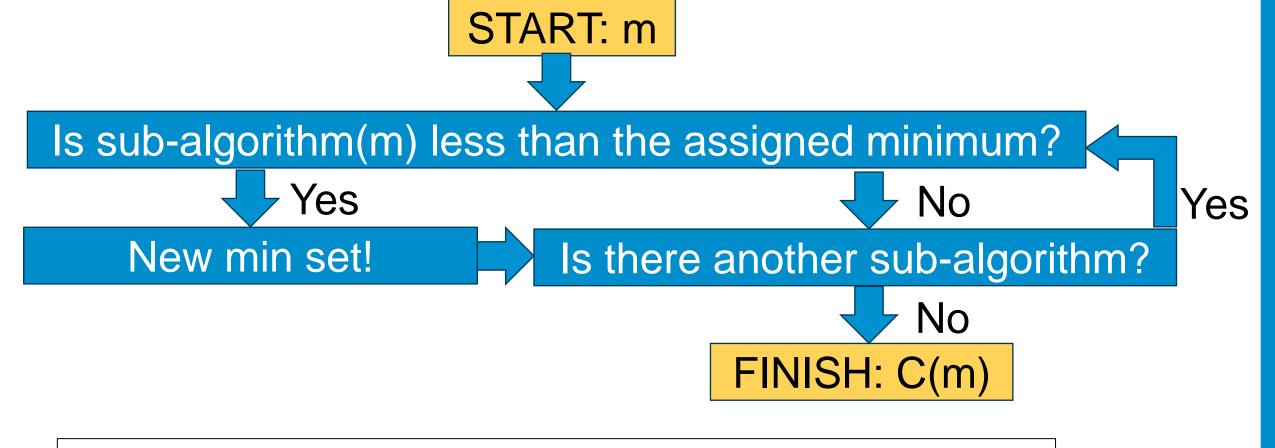


Figure 2. Flow chart explaining the driver algorithm for cost function.

Sub-algorithms can be used in any order leading to the same result

#### MULTIPLICATION:

Find the two factors of m, that minimizes C(a) + C(b) and add 3

#### ADDITION:

Find two addends summing to m that minimizes C(a1) + C(a2) and add 2

#### FIBONACCI SEQUENCE:

Start with 1, 1 and if m equals the previous two values, find the value in the sequence and add 4

#### BINARY REPRESENTATION:

Find the length of the binary representation and add 7

#### SUCCESSOR FUNCTION:

Find the cost of the previous number [C(m-1)] and add 1

#### POWER:

Find a base and exponent that equal m such that C(b) + C(e) is minimized and add 4

Figure 3. Diagram outlining the cost procedures for each sub-algorithm

#### Results

- We consider the lowest cost attained across all sub-algorithms
- Both models start increasing quickly and slow down for higher m
- Both show positive correlation with higher R<sup>2</sup> present in newer model

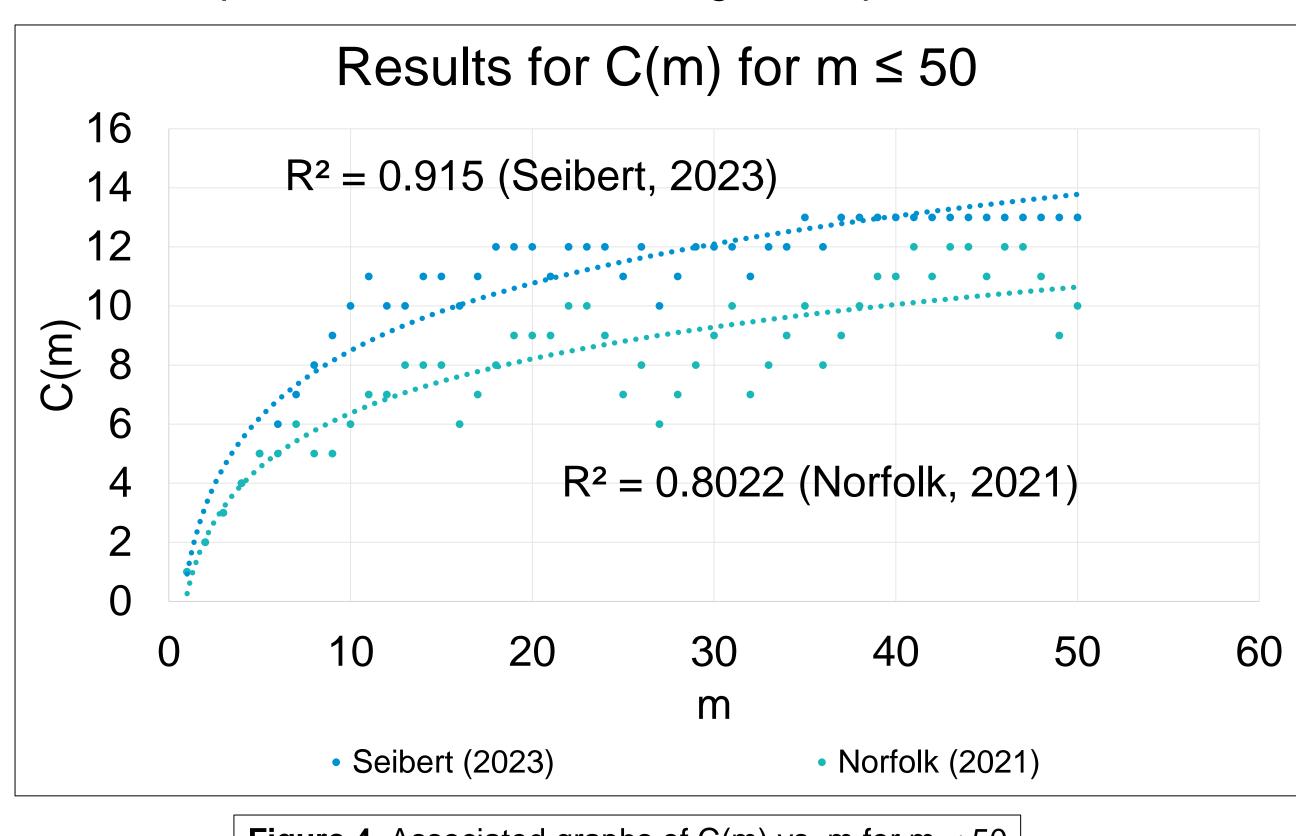


Figure 4. Associated graphs of C(m) vs. m for m < 50

#### Conclusions

- Outputs to C(m) mostly level out (except for powers of 2) in the newer model due to binary function being added
- Subtraction relies on cost values higher than numbers passed into the sub-algorithms, rendering it impossible for use
- More functions can be used when considering ones that use iteration on the output, but it requires changing the driver algorithm

#### Reference

 Norfolk, Maxwell (2021) "The Cost of a Positive Integer," Rose-Hulman Undergraduate Mathematics Journal: Vol. 22: Iss. 1, Article 9.

#### **Source Code**

