給定
$$g^{\alpha}$$
, g^{β} , 可算 $g^{\alpha\beta}$ ______(1)
給定 g^{α} , 可算 g^{α^2} _____(2)

給定
$$g^{\alpha}$$
 和 $\alpha \neq 0$,可算 $g^{\frac{1}{\alpha}}$ _____(3)

給定
$$g^{\alpha}$$
, g^{β} , 且 $\beta \neq 0$, 可算 $g^{\frac{\alpha}{\beta}}$ ____(4)

當 $g^{\alpha}=g^{\beta}$, $g^{\alpha\beta}=g^{\alpha^2}$, 所以(1)等價於(2) -> 此作法需作一次,因此為 poly-time

設存在 γ 使得 $g^{\gamma\alpha} = g^{\frac{1}{\alpha}}$,

則
$$g^{\gamma\alpha^2} = g$$
, $g^{\gamma\alpha^4} = g^{\gamma\alpha^4}$
由(2)知,給定 g^{α^2} , 可算 g^{α^4}

因此當 $(g^{\alpha^4})^{\gamma}=g^{\alpha^2}$,則 $g^{\gamma\alpha}=g^{\frac{1}{\alpha}}$ —> 此作法需作 γ 次,且 γ < q,因此為 poly-time

設 $(g^{\frac{1}{\alpha}})^k$ 可解得 g,則 $(g^{\frac{1}{\alpha}})k^2$ 可解得 g^{α} ,可得 k -> 此作法需作 k 次,且 k < q,因此為 poly-time

由(4)可得 $g^{\frac{\alpha}{\beta}}$

設存在 k 使得 $(g^{\frac{\alpha}{\beta}})^k=g^{\alpha}$,則 $(g^{\frac{\alpha}{\beta}})k^2=g^{\alpha^2}$ —> 此作法需作 k 次,且 k < q ,因此為 poly-time

1. 2-out-of-2

選擇一數值作為 sk1

並設 $sk2 = s \oplus sk1 \rightarrow s = sk2 \oplus sk1$

此時即可將 s 拆成(sk1, sk2),並分別放在兩個 key server 中在解密時,需先將 c 分別交給兩個 server

得到

D(c, sk1)
$$\rightarrow$$
 c'₁
D(c, sk2) \rightarrow c'₂

 $C(c, c'_1, c'_2)$:

在取得 C'_1 和 C'_2 後,根據 ElGamal,透過 (C, C'_1) 與 (C, C'_2) 可反推 Sk1 與 Sk2 的結果

再由一開始的設計 $s = sk2 \oplus sk1$, 即可找到 s,最終透過 D(c, s)即可找出明文 m

因此
$$C(c, c'_1, c'_2) \rightarrow m$$

此處拆解 secret key 的作法使用 One time pad 的概念,因此根據 one time pad,在沒有重複使用 sk1 的情況下,不會透露任何部份 的 key。

2. 2-out-of-t

設計一方程式 y = ax + b, 其中 b = s, 隨機產生一個數字 a 在此方程式中,我們可以產生 t 組座標 (x_i, y_i) 設 $sk_1 = (x_1, y_1)$, $sk_2 = (x_2, y_2)$,, $sk_{t-1} = (x_{t-1}, y_{t-1})$

設
$$Sk_1 = (x_1, y_1)$$
, $Sk_2 = (x_2, y_2)$,, $Sk_{t-1} = (x_{t-1}, y_{t-1})$, $Sk_t = (x_t, y_t)$

 $D(c, sk_i) \rightarrow c'_i$

$$D(c, sk_j) \rightarrow c'_j$$

從任 2 個 Sk 中取得雨點座標 $C'_i = (x_i, y_i), C'_j = (x_j, y_j)$

由兩點座標帶入 y = ax + b 可求出此方程式的係數,即可找到 s,有 S 後即可解開 c,取得明文 m

參考:

https://en.wikipedia.org/wiki/Shamir%27s Secret Sharing

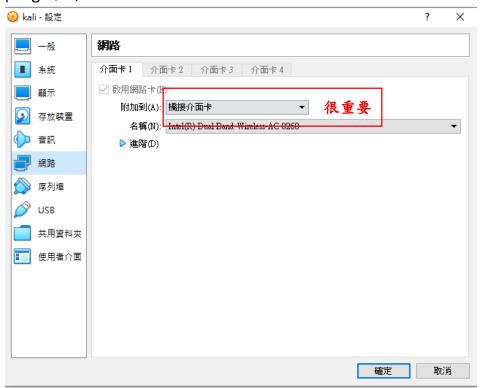
http://www.cs.columbia.edu/~tal/4261/F16/ss slides luke.pdf

作法說明:

這次的中間人攻擊實作主要是使用 ARP 攻擊,ARP 協定透過廣播功能,獲取 IP 位址和實體位置的對應關係(建表),而 ARP 攻擊主要是讓發送端誤以為 hacker 端是目標,讓發送端取得錯誤的 MAC 位址,從而監聽發送端的封包資訊

前處理:

由於我使用 virtualbox 建立 kali linux 環境來對本機進行 ARP 攻擊,因此要記得先將網路設定為橋接介面卡,否則兩邊根本連 ping 都ping 不到



確認 hacker 端 ip(千萬不要亂改,一開始想改個好辨認的,害得自己ping 的到卻找不到...)

```
(kali@kali)-[~]
$ ip a

1: lo: <LOOPBACK,UP,LOWER_UP> mtu 65536 qdisc noqueue state UNKNOWN group default qlen 1000
    link/loopback 00:00:00:00:00:00 brd 00:00:00:00:00
    inet 127.0.0.1/8 scope host lo
        valid_lft forever preferred_lft forever
    inet6 ::1/128 scope host
    valid_lft forever preferred_lft forever

2: eth0: <BROADCAST_MULTICAST_UP_LOWER_UP> mtu 1500 qdisc pfifo_fast state UP group default qlen 1000
    link/ether 08:00:27:77:66:1b brd ff:ff:ff:ff:ff
    inet 192.168.7.39/24 brd 192
    valid_lft 37:88ec preferred_lft 3578sec
    inet6 fe80::a00:27ff:fe77:661b/64 scope link noprefixroute
    valid_lft forever preferred_lft forever
```

確認本地端能否 ping 到 hacker 端,並且查看 arp 表格,以確保在同一個區網中(否則在 ettercap 中根本抓不到)

```
EX 命令提示字元
C:\Users\user>ping 192.168.7.39

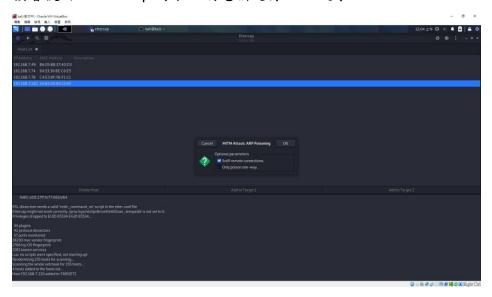
Ping 192.168.7.39 (使用 32 位元組的資料):
回覆自 192.168.7.39 (位元組=32 時間=lms TTL=64
回覆自 192.168.7.39 (位元組=32 時間<lms TTL=64
回覆自 192.168.7.39 (位元組=32 時間<lms TTL=64
回覆自 192.168.7.39 (位元組=32 時間<lms TTL=64
192.168.7.39 的 Ping 統計資料:
封包:已傳送 = 4,已收到 = 4,已遺失 = 0 (0% 遺失),
大約的來回時間 (毫秒):
最小值 = 0ms,最大值 = 1ms,平均 = 0ms

C:\Users\user>arp -a

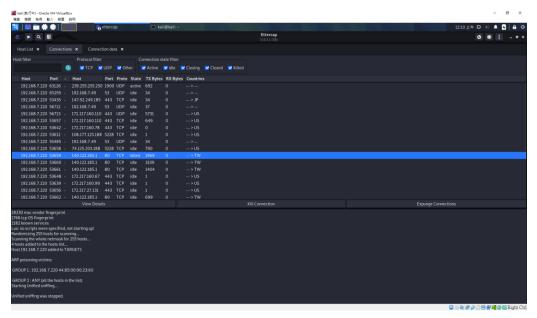
介面: 192.168.7.220 --- 0x7
網際網路網址 實體位址 類型
192.168.7.7 00-aa-aa-aa-aa 動態
192.168.7.39 08-00-27-77-66-1b 動態
192.168.7.49 b6-05-88-37-40-d3 動態
192.168.7.255 ff-ff-ff-ff-ff 靜態
224.0.0.22 01-00-5e-00-00-16 靜態
224.0.0.251 01-00-5e-00-00-fb 靜態
224.0.0.252 01-00-5e-00-00-fc 靜態
239.255.255.255.255 ff-ff-ff-ff-ff-ff

#態
```

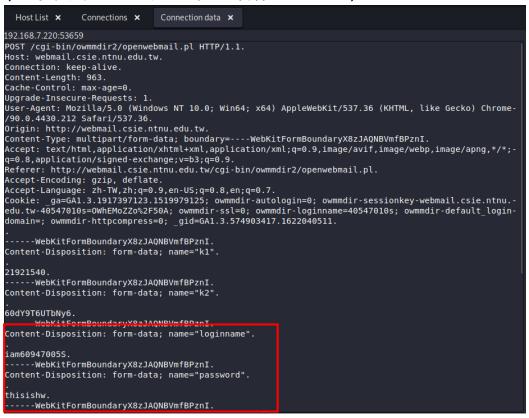
接著使用 ettercap 對本地端電腦進行 ARP 攻擊



透過 ARP 毒化,即可監聽此電腦的所有封包,找到師大類的 IP(140.122),開始看封包



由於師大資工 webmail 網站是 http,傳輸的過程都是以明文顯示,因此打開封包後,便可成功獲取師大資工 webmail 的輸入資訊 (此處為作業使用,所以沒有使用真實的帳號密碼)



參考:

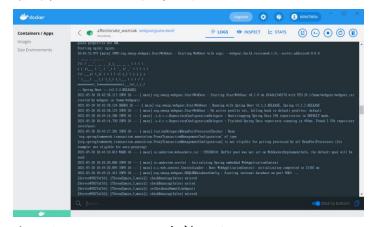
https://www.youtube.com/watch?v=5AhB6rLgPuA&ab channel=%E8%
95%AD-%E5%BF%97%E6%98%8E

https://www.youtube.com/watch?v=ogtWS6MfiWM&ab channel=Prof
essorAndrew

這個作者的攻擊手法,是透過 github,找到一些開源的 source code,並從裡面的一些 pattern(package.json),來推測有哪些人會使這類 API,藉此抓取目標族群,並針對這目標族群,仿造他們所需的 API,並在內部加入惡意程式碼,讓使用者以為這是該 API 的最新版本,當有人使用此 API 時,就會遭受 DNS 滲透攻擊,讓使用者在封包傳出去的時候,傳到作者的 server 中,藉此從這些封包分析,可能可以獲得該企業所使用的更多 API 以及版本資訊,就可以用同樣的方式偽造更多的 API,來竊取更多的訊息,甚至是在內部加入其他想做的事情。

我的防禦方式是規範工程師在每次使用新版本的 API 時,必須先查看要使用的 API 更新了甚麼內容,像是透過 git 版本控制功能,針對會用到的函式查看內文有沒有出現一些不該出現的內容,另外一點就是不要把一些pattern 以明文表示,可透過企業內部建立表格用對應的代號表示。

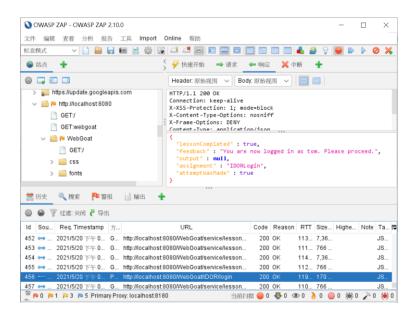
這題光是前置作業就花費我超久時間! 最後裝 docker,才成功連上



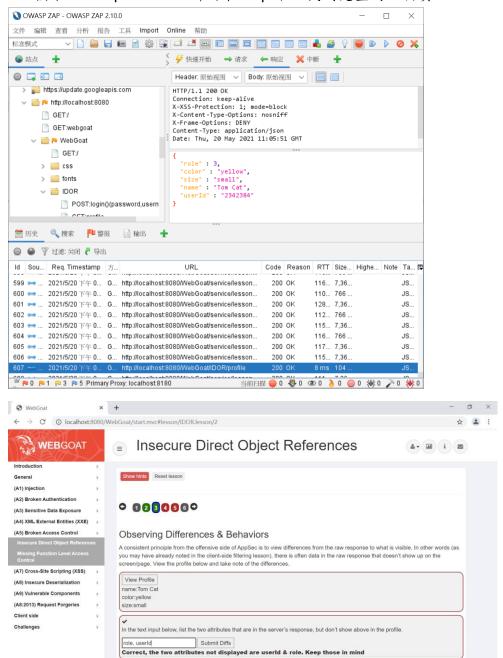
本題使用 OWASP ZAP 來幫忙監控

2. 輸入 tom cat, submit 後,會抓取到一個 post 封包

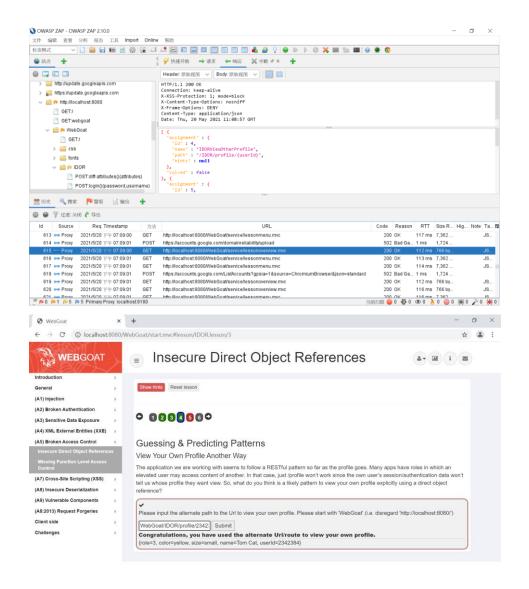




3. 按下 view profile,即可在 zap 中,找到完整的人物資訊



4. 根據上一題的 url,再加上 zap 中發現的一些格式規則(每 5 秒會發一組),即可找出此題的 pattern

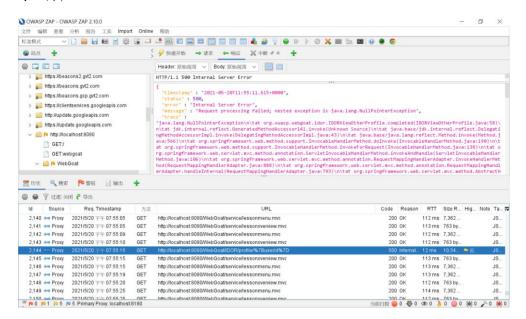


5_1. 老實說我本來完全不知道要幹嘛,按下 view profile 後,在 zap 也抓不到東西,後來另外開一個分頁,並將剛剛的 userId 不斷亂改,最終在 userid=2342388 時,找到另一位使用者

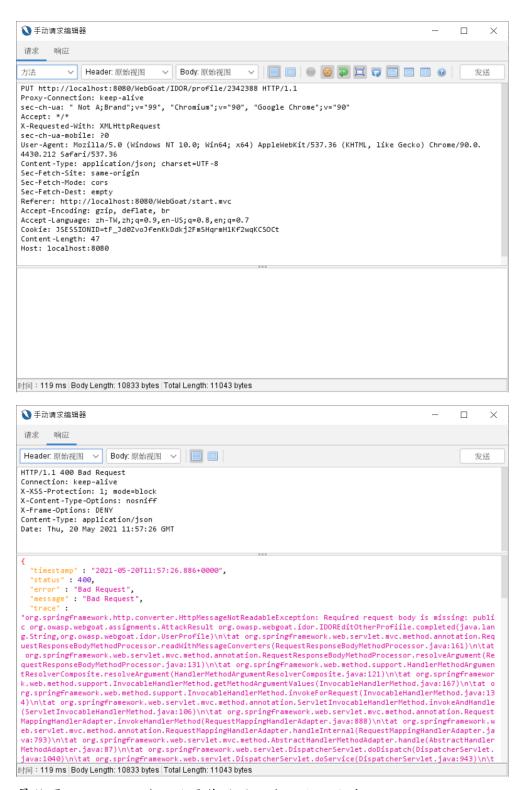


5_2. 按下第二個 view profile 後,在 zap 找到對應的封包,進行手動請

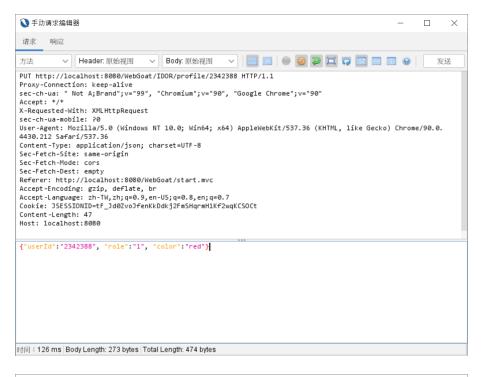
求編輯

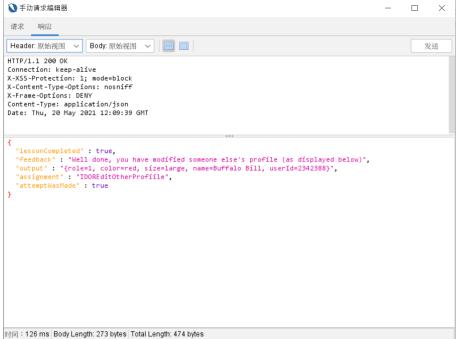


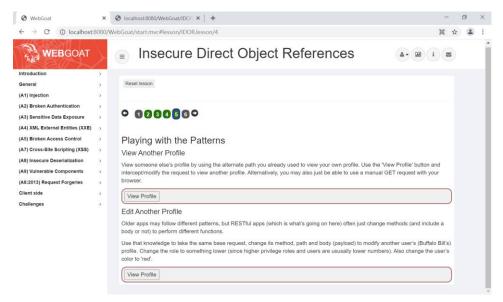
此處先修改 userid,將 content-type 改為 json,嘗試發送



最終再加上 userId 以及要修改的內容,即可完成







Task1

依照題目要求設定 $p \cdot q \cdot e$,並計算 phi(N) = (p-1)(q-1),最後計算 $ed \equiv 1 \pmod{phi(N)}$ 即可找出 d程式碼如下:

```
1#include<stdio.h>
 2 #include < openssl/bn.h>
 4 void printBN(char *msg, BIGNUM *a)
     char *number_str = BN_bn2hex(a);
 7
     printf("%s %s\n", msg, number str);
 8
     OPENSSL free(number str);
 9 }
10
11 int main()
12
13
     BIGNUM *p = BN new();
14
     BIGNUM *q = BN new();
     BIGNUM *e = BN_new();
15
    BN_hex2bn(&p, "F7E75FDC469067FFDC4E847C51F452DF");
BN_hex2bn(&q, "E85CED54AF57E53E092113E62F436F4F");
BN_hex2bn(&e, "0D88C3");
16
17
18
19
    BIGNUM *n = BN new();
    BN_CTX *ctx = BN_CTX_new();
20
21
     BN_mul(n, p, q, ctx);
22
     BIGNUM *p 1 = BN new();
23
     BIGNUM *q 1 = BN new();
24
     BN_sub(p_1, p, BN_value_one());
25
     BN_sub(q_1, q, BN_value_one());
26
     BIGNUM *phi = BN new();
     BN_mul(phi, p_1, q_1, ctx);
27
28
     BIGNUM *d = BN new();
29
     BN_mod_inverse(d, e, phi, ctx);
30
     printBN("Private key d =", d);
31
```

輸出結果:

```
[05/21/21]seed@VM:~/.../hw4_seedlab$ gcc task1.c -lcrypto
[05/21/21]seed@VM:~/.../hw4_seedlab$ ./a.out
Private key d = 3587A24598E5F2A21DB007D89D18CC50ABA5075BA19A33890FE7C28A9B496AEB
[05/21/21]seed@VM:~/.../hw4_seedlab$
```

Task2

使用 python 將明文轉成 hex string

```
[05/21/21]seed@VM:~/.../hw4_seedlab$ python3 -c 'print("A top secret!".encode("utf-8").hex())' 4120746f702073656372657421
```

程式碼:

```
1 #include<stdio.h>
 2 #include < openssl/bn.h>
 4 void printBN(char *msg, BIGNUM *a)
      char *number_str = BN_bn2hex(a);
      printf("%s %s\n", msg, number_str);
      OPENSSL_free(number_str);
 8
10
11 int main()
12
      BIGNUM *n = BN_new();
     BIGNUM *n = BN_new();

BIGNUM *e = BN_new();

BIGNUM *m = BN_new();

BN_hex2bn(&n, "DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5");

BN_hex2bn(&e, "010001");
14
15
16
18
19
      //input the plaintext
     BN_hex2bn(&m, "4120746f702073656372657421");
BN_CTX *ctx = BN_CTX_new();
BIGNUM *c = BN_new();
21
22
     BN_mod_exp(c, m, e, n, ctx);
printBN("Encrypt:", c);
24
25
     //check
BIGNUM *d = BN_new();
BN_hex2bn(&d, "74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D");
BN_mod_exp(m, c, d, n, ctx);
printBN("Decrypt:", m);
26
27
輸出結果:
[05/21/21]seed@VM:~/.../hw4_seedlab$ gcc task2.c -lcrypto
[05/21/21]seed@VM:-/.../hw4_seedlab$ ./a.out
Encrypt: 6FB078DA550B2650832661E14F4F8D2CFAEF475A0DF3A75CACDC5DE5CFC5FADC
Decrypt: 4120746F702073656372657421
```

Task3

程式碼:

```
1#include<stdio.h>
 2 #include < openssl/bn.h>
  4 void printBN(char *msg, BIGNUM *a)
 5 {
      char *number_str = BN_bn2hex(a);
      printf("%s %s\n", msg, number_str);
OPENSSL_free(number_str);
 8
  9}
10
11 int main()
12 {
      BIGNUM *n = BN_new();
BIGNUM *d = BN_new();
BIGNUM *m = BN_new();
BIGNUM *c = BN_new();
BN_hex2bn(&n, "DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5");
BN_hex2bn(&d, "74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D");
15
16
18
19
20
        //input the ciphertext
       BN_hex2bn(&c, "8C0F971DF2F3672B28811407E2DABBE1DA0FEBBBDFC7DCB67396567EA1E2493F");
21
22
23
      BN_CTX *ctx = BN_CTX_new();
BN_mod_exp(m, c, d, n, ctx);
printf("%s\n", BN_bn2hex(m));
25
26 }
```

```
此處我另外寫了一個 16 進位轉字串的程式(tostring.py):
                                                                                                                                                                                       tostring.py
                    Open ▼ 升
                  1s = input()
                   2 out =
                  3 for i in range(0, len(s), 2):
                           out += chr(int(s[i:i+2], 16))
                  5 print(out)
              輸出結果:
              [05/21/21]seed@VM:~/.../hw4_seedlab$ gcc task3.c -lcrypto
               [05/21/21]seed@VM:~/.../hw4 seedlab$ ./a.out
              50617373776F72642069732064656573
              [05/21/21]seed@VM:~/.../hw4 seedlab$ python3 tostring.py
              50617373776F72642069732064656573
              Password is dees
              [05/21/21]seed@VM:~/.../hw4_seedlab$
Task4
              分別將兩串明文轉成 hex string
              [05/21/21]seed@VM:~/.../hw4_seedlab$ python3 -c 'print("I owe you $2000.".encode("utf-8").hex())'
              [05/21/21]seed@VM:~/.../hw4_seedlab$ python3 -c 'print("I owe you $3000.".encode("utf-8").hex())'
              49206f776520796f752024333030302e
              簽名程式碼:
                1#include<stdio.h>
                 2 #include < openssl/bn.h>
                 4 void printBN(char *msq, BIGNUM *a)
                5 {
                        char *number_str = BN_bn2hex(a);
                         printf("%s %s\n", msg, number_str);
                        OPENSSL_free(number_str);
              10
              11 int main()
              12 {
                        BIGNUM *n = BN_new();
                        BIGNUM *H = BN_new();

BIGNUM *d = BN_new();

BN_hex2bn(&n, "DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5");

BN_hex2bn(&e, "010001");

BN_hex2bn(&d, "74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D");
              15
             17
18
19
20
21
22
23
24
25
26
27
28 }
                       //input the plaintext
BIGNUM *m = BN_new();
BN_hex2bn(&m, "49206f776520796f752024333030302e");
printBN("M =", m);
BN_CTX *ctx = BN_CTX_new();
BIGNUM *s = BN_new();
BN_mod_exp(s, m, d, n, ctx);
printBN("Signature:", s);
              輸出結果:
              [05/21/21]seed@VM:-/.../hw4_seedlab$ python3 -c 'print("I owe you $2000.".encode("utf-8").hex())'
              49206f776520796f752024323030302e
[05/21/21]seed@VM:~/.../hw4 seedlab$ python3 -c 'print("I owe you $3000.".encode("utf-8").hex())'
              | \(\frac{1}{21}\) \(\frac{1}\) \(\frac{1}{21}\) \(\frac{1}{21}\) \(\frac{1}{21}\) \(\frac{
             Signature: 55A4E7F17F04CCFE2766E1EB32ADDBA890BBE92A6FBE2D785ED6E73CCB35E4CB
[05/21/21]seed@VM:-/.../hw4_seedlab$ gcc task4.c -lcrypto
[05/21/21]seed@VM:-/.../hw4_seedlab$ ./a.out
M = 49206F776520796F752024333030302E
              Signature: BCC20FB7568E5D48E434C387C06A6025E90D29D848AF9C3EBAC0135D99305822
```

Task5

先將明文轉成 hex string,用於對照簽名

```
[05/21/21]seed@VM:-/.../hw4_seedlab$ python3 -c 'print("Launch a missile.".encode("utf-8").hex().upper()
4C61756E63682061206D697373696C652E
```

對照簽名程式碼:

```
1#include<stdio.h>
  2 #include < openssl/bn.h>
  4 void printBN(char *msg, BIGNUM *a)
        char *number_str = BN_bn2hex(a);
printf("%s %s\n", msg, number_str);
OPENSSL_free(number_str);
 9}
10
11 int main()
12 {
        BIGNUM *n = BN_new();
BIGNUM *e = BN_new();
BIGNUM *s = BN_new();
BN_hex2bn(&n, "AE1CD4DC432798D933779FBD46C6E1247F0CF1233595113AA51B450F18116115");
BN_hex2bn(&e, "010001");
BN_hex2bn(&s, "643D6F34902D9C7EC90CB0B2BCA36C47FA37165C0005CAB026C0542CBDB6802F");
13
14
15
17
18
19
        BN_CTX *ctx = BN_CTX_new();
BIGNUM *m = BN_new();
20
21
22
23
24
        BN_mod_exp(m, s, e, n, ctx);
printBN("M =", m);
char *target = "4C61756E63682061206D697373696C652E";
if (*BN_bn2hex(m) == *target)
25
26
27
        printf("The signature is indeed Alice's\n");
else
            printf("Not!!!\n");
```

輸出結果:

```
[05/21/21]seed@VM:~/.../hw4_seedlab$ gcc task5.c -lcrypto [05/21/21]seed@VM:~/.../hw4_seedlab$ ./a.out M = 4C61756E63682061206D697373696C652E The signature is indeed Alice's 將 2F 改成 3F 的結果:
```

```
[05/21/21]seed@VM:~/.../hw4_seedlab$ gcc task5.c -lcrypto [05/21/21]seed@VM:~/.../hw4_seedlab$ ./a.out M_2F = 4C61756E63682061206D697373696C652E M_3F = 91471927C80DF1E42C154FB4638CE8BC726D3D66C83A4EB6B7BE0203B41AC294 [05/21/21]seed@VM:~/.../hw4_seedlab$ python3 tostring.py 91471927C80DF1E42C154FB4638CE8BC726D3D66C83A4EB6B7BE0203B41AC294 ñä,0´cè¼rm=fÈ:N¶·¾´®Â
```

Task6

我嘗試的網站是 5278 網站,並將內容存入 task6.txt 中,將兩段憑證分別存入 c0.pem 與 c1.pem

透過以下指令,從中找出(e,n)以及簽名

```
[05/25/21]seed@VM:-/.../hw4_seedlab$ nano c0.pem
[05/25/21]seed@VM:-/.../hw4_seedlab$ nano c1.pem
[05/25/21]seed@WY:-/.../hw4_seedlab$ onessl x509 -in c1.pem -noout -modulus
[05/25/21]seed@VM:-/.../hw4_seedlab$ openssl x509 -in c1.pem -noout -modulus
Modulus=B8021528CCF6A094D30F12EC8D5592C3F882F199A67A4288A75D26AAB528B95254C81AF8E6BF975C8A3D70F4794145535578C9EA8A23919F582
3C42A94-GE6F538C32ED8B0C08D5C753938E7EDCF69F05A0B1B8EC094242587FA377IB313E71CACE19BEFDBE43B45524596A9C153CE34C852EEBSAEED8F
DE6070E2A554ABB66D0E97A540346B2B03BC66EB66347CFA6B8B8F572999F830175DBA726FFB81C5ADD2865B3D17C7E7098BF12BF786DCC1DA715DD446
E3CACAD25C18B8CG60F7566B3F118F7A25CE653FF3A88B647A5FF1318EA9809773F9D53F9CF01E5F5A6701714AF63A4FF99B3939DDC53A706FE48851DA1
69AE2575BB13CC5203F5ED51A18BDB15
[05/25/21]seed@VM:~/.../hw4_seedlab$ openssl x509 -in c1.pem -text -noout
Certificate:
      Data:
             Version: 3 (0x2)
             Serial Number:
                   40:01:75:04:83:14:a4:c8:21:8c:84:a9:0c:16:cd:df
             Signature Algorithm: sha256WithRSAEncryption
             Issuer: 0 = Digital Signature Trust Co., CN = DST Root CA X3
             Validity
                   Not Before: Oct 7 19:21:40 2020 GMT
             Not After : Sep 29 19:21:40 2021 GMT
Subject: C = US, O = Let's Encrypt, CN = R3
             Subject Public Key Info:
                   Public Key Algorithm: rsaEncryption
                          RSA Public-Key: (2048 bit)
                          Modulus:
                                00:bb:02:15:28:cc:f6:a0:94:d3:0f:12:ec:8d:55:
                                92:c3:f8:82:f1:99:a6:7a:42:88:a7:5d:26:aa:b5:
                                2b:b9:c5:4c:b1:af:8e:6b:f9:75:c8:a3:d7:0f:47:
                                94:14:55:35:57:8c:9e:a8:a2:39:19:f5:82:3c:42:
                                a9:4e:6e:f5:3b:c3:2e:db:8d:c0:b0:5c:f3:59:38:
                                e7:ed:cf:69:f0:5a:0b:1b:be:c0:94:24:25:87:fa:
                                37:71:b3:13:e7:1c:ac:e1:9b:ef:db:e4:3b:45:52:
                                45:96:a9:c1:53:ce:34:c8:52:ee:b5:ae:ed:8f:de:
                                60:70:e2:a5:54:ab:b6:6d:0e:97:a5:40:34:6b:2b:
d3:bc:66:eb:66:34:7c:fa:6b:8b:8f:57:29:99:f8:
                                30:17:5d:ba:72:6f:fb:81:c5:ad:d2:86:58:3d:17:
                                c7:e7:09:bb:f1:2b:f7:86:dc:c1:da:71:5d:d4:46:
                                e3:cc:ad:25:c1:88:bc:60:67:75:66:b3:f1:18:f7:
                                a2:5c:e6:53:ff:3a:88:b6:47:a5:ff:13:18:ea:98:
                                09:77:3f:9d:53:f9:cf:01:e5:f5:a6:70:17:14:af:
                                63:a4:ff:99:b3:93:9d:dc:53:a7:06:fe:48:85:1d:
                                a1:69:ae:25:75:bb:13:cc:52:03:f5:ed:51:a1:8b:
                                db . 15
                          Exponent: 65537 (0x10001)
```

將 co.pem 簽名內容另存,並去掉所有的空格與冒號

```
[05/25/21]seed@VM:-/.../hw4_seedlab$ nano signature [05/25/21]seed@VM:-/.../hw4_seedlab$ cat signature | tr -d '[:space:]:' 430ffaf8aldc514a547d75alb13a30ffcb625a9bd1c7d55a6a9b7a9dc5161885d7dab938339162621c0ed9f0a2a8091a4fdf107fc50ca3c87f4a8093fef d25a5e2112436056734e488090ac4a385db6bdd472de5859f21dbfdb17456c81cd09151ec85c8a02a1405f98656817515c9fd030680accab91d626858 3ddce632cc4fbeb9f6ffeff91c670ca427b3fa642de8c293dad909510d760cla2c1f8c0edef2c22fc302d7ee5b703131a816e2b86b6f3b12e4fbde4821 e9c1674620a621c108f53064d754496b0b8097a547851fd85e665ad91c195ae9f3271f6817c533027cfc4593d5af9f11fe11869745ac11le23c78a534e d6f7aa5f1eb794d4ba59c301[05/25/21]seed@WM:-/.../hw4_seedlab$
```

輸出憑證主體並計算 hash 值,憑證是從 4 開始,所以從 4 開始看

```
[05/25/21]seed@VM:-/.../hw4_seedlab$ openssl asnlparse -i -in c0.pem -strparse 4 -out c0_body.bin -noout [05/25/21]seed@VM:-/.../hw4_seedlab$ sha256sum c0_body.bin lc17f0624f5c572e7falfleaf234c2f44e6104f3e76bfd210a05e8c07d2932ec c0_body.bin
```

利用獲得的所有資訊進行驗證

程式碼:

```
1 #include<stdio.h>
2 #include<spensity.h.h>
3
4 void printBN(char *msg, BIGNUM *a)
5 {
6    char *number_str = BN_bn2hex(a);
    printf("%s %s\n", msg, number_str);
8    OPENSSL_Tree(number_str);
9 }
10
11 int main()
12 {
13    BIGNUM *n = BN_new();
14    BIGNUM *e = BN_new();
15    BIGNUM *s = BN_new();
16   BIGNUM *s = BN_new();
17
18    BN_hex2bn(&n,
**BB01258Cc(F6A094D30F12EC8D5592C3F882F199A67A4288A75D26AAB52BB9C54CB1AF8E6BF975C8A3D70F4794145535578C9EA8A23919F5823C42A9
19    BN_hex2bn(&e, "10001");
10    BN_hex2bn(&e, "10001");
11    BN_nex2bn(&e, "10001");
12    BN_hex2bn(&e, "10001");
13    BN_hex2bn(&e, "10001");
14    BN_hex2bn(&hash, "1c17f0624f5c572e7fa1fleaf234c2f44e6104f3e76bfd210a05e8c07d2932ec");
15    BN_hex2bn(&hash, "1c17f0624f5c572e7fa1fleaf234c2f44e6104f3e76bfd210a05e8c07d2932ec");
16    BN_hex2bn(&hash, "1c17f0624f5c572e7fa1fleaf234c2f44e6104f3e76bfd210a05e8c07d2932ec");
17    BN_hex2bn(&hash, "1c17f0624f5c572e7fa1fleaf234c2f44e6104f3e76bfd210a05e8c07d2932ec");
18    BN_crtx *ctx = BN_crtx_new();
18    BIGNUM *c = BN_new();
18    BIGNUM *c = BN_new();
19    BN_hex2bn(&hash, "1c17f0624f5c572e7fa1fleaf234c2f44e6104f3e76bfd210a05e8c07d2932ec");
19    BN_hex2bn(&hash, "1c17f0624f5c572e7fa1fleaf234c2f44e6104f3e76bfd210a05e8c07d2932ec");
19    BN_new();
10    BN_new();
11    BN_new();
12    BN_new();
12    BN_new();
13    BN_new();
14    BIGNUM *c = BN_new();
15    BN_new();
16    BN_new();
17    BN_new();
18    BN_new();
18    BN_new();
19    BN_new();
19    BN_new();
10    BN_new();
10    BN_new();
11    BN_new();
11    BN_new();
12    BN_new();
13    BN_new();
14    BN_new();
15    BN_new();
16    BN_new();
17    BN_new();
18    BN_new();
18    BN_new();
19    BN_new();
10    BN_new();
11    BN_new();
11    BN_new();
12    BN_new();
13    BN_new();
14    BN_new();
15    BN_new();
16    BN_new();
17    BN_new();
18    BN_new();
18
```

最終輸出結果包含 server's certificate hash 即表示驗證成功。



參考:

https://blog.csdn.net/weixin 48392428/article/details/107433009