

CMSC351: Prerequisite

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# 1 Symbols

1. For all ( $\forall$ )
2. There exists ( $\exists$ )

## 2 Proofs

### 2.1 Weak Induction

First, we need to prove  $\forall n \geq n_0 P(n)$  we first prove  $P(n_0)$  (which is the base case) and then we prove  $\forall k \geq n_0 P(k) \rightarrow P(k+1)$  (which is the inductive step) (you can also prove  $P(k-1)$ ). The assumption of  $P(k)$  in the inductive step is the inductive hypothesis.

Ultimately, for the inductive step we are trying to find that for any  $k \geq n_0$  if  $P(k)$  is true, then  $P(k+1)$  is also true.

**Example:** Suppose we have a set of nested Russian dolls (Matryoshka dolls). Each doll is contained with another doll, each doll is labeled 1, 2, 3 and so on.

In this hypothetical, there are two things that are true. Let's say  $M(n)$  is true iff (if and only if) doll  $n$  has another doll contained.

(a) The first doll has another doll contained. That is,  $M(n)$  is true.

(b) For every doll  $k$ , if doll  $k$  has another doll inside then doll  $k+1$  has another doll inside.

$$\forall k \geq 1, M(k) \rightarrow M(k+1)$$

We can now conclude that  $\forall n \geq 1, M(n)$ .

### 2.2 Strong Induction

The goal of strong induction is we need to find some property, we can say  $P(n)$  and we need to find some  $n$  greater than equal to a ( $n \geq a$ ).

How can we accomplish this?

**Step 1 (Basis step):** We are going to prove for  $P(a)$  (we would just prove  $P(a)$  for weak induction),  $P(a+1)$ , ..., for some finite number say  $P(b)$ .

( $P(a)$ ,  $P(a+1)$ , ...,  $P(b)$ )  $\leftarrow$  we prove each of these.

**Step 2 (Induction):** We can now assume  $P(i)$  where  $a \leq i \leq k$ . (Assume  $P(i)$ ). Then, we prove  $P(k+1)$ .

**Example: Fibonacci Sequence**

*Proof.* Claim: The Fibonacci Sequence is defined as the following:  $F(0) = 0$ ,  $F(1) = 1$  and for  $n \geq 2$ ,  $F(n) = F(n-1) + F(n-2)$ . We want to prove that for all  $n \geq 0$ ,  $F(n) \leq 2^n$

Base cases:

For  $n = 0$ :  $F(0) = 0 \leq 2^0 = 1$

For  $n = 1$ :  $F(1) = 1 \leq 2^1 = 2$

For both these base case, they hold true.

Inductive Step: Assume that for all  $i$ ,  $0 \leq i \leq k$ , we have  $F(k) \leq 2^k$ .

(This could also be known as the inductive hypothesis)

We now need to prove  $F(k+1) \leq 2^{k+1}$ .

By the inductive hypothesis, we have  $F(k) \leq 2^k$  and  $F(k-1) \leq 2^{k-1}$ .

We can prove by which:

$$\begin{aligned} F(k+1) &= F(k) + F(k-1) \\ &\leq 2^{k-1} + 2^{k-2} && \text{(IH)} \\ &\leq 2^{k-1} + 2^{k-1} && (2^{k-2} \text{ is less than } 2^{k-1}) \\ &= 2^k && \text{(Simplify)} \end{aligned}$$

□