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Data Types

Consider data types in the C language. Here is an example related to SHA 512:

```
struct sha512_state_st {
   uint64_t h[8];
   uint64_t Nl, Nh;
   uint8_t p[128];
   unsigned num, md_len;
};
```

This structure consists of an 8 element 64 bit unsigned integer array named h, 64 bit unsigned integer numbers N1 and Nh, a 128 element byte array named p and 32 bit unsigned integers num and md_len. Creation and initialization of a variable of this type might look like this:

A simple function to manipulate a SHA512 state and return a fresh copy might look like this (please forgive the lack of protections – this is just for illustration):

```
SHA512_CTX *stateManip (SHA512_CTX state) {
   int i;
   SHA512_CTX *s = (SHA512_CTX*)malloc(sizeof(SHA512_CTX));
   for (i=0 ; i < 8 ; i++) s->h[i] = state.h[i];
   s->Nl = state.Nl;
   s->Nh = state.Nh;
   for (i=0 ; i < 128 ; i++) s->p[i] = state.p[i];
   s->num = state.num;
   s->md_len = 12;
   return s;
}
```

A function to print a SHA512 state might look like this:

```
void printState (SHA512_CTX state) {
   int i;
   printf("h=");
   for (i=0 ; i < 8 ; i++) printf("%ld,",state.h[i]);
   printf("\nNl=%ld, Nh=%ld\n", state.Nl, state.Nh);
   printf("p=");
   for (i=0 ; i < 128 ; i++) printf("%d,",state.p[i]);</pre>
```

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```
printf("\nnum=%d, md_len=%d\n", state.num, state.md_len);
}
```

Cryptol versions of these would be the following. For the data structure SHAState:

For an initialization of a variable of type SHAState:

A function that manipulates a state might look like this:

Observe that in later lessons there will be attempts to show that C function are functionally equivalent to Cryptol specifications. But Cryptol has no pointers to manipuate as C does. Hence something needs to be added to make that happen.

Cryptol allows parameterization of types. This is done by creating a section called parameter and then declaring type parameters with constraints. For example:

```
parameter
   type w : #
   type constraint (fin w, w <= 64, 32 <= w)
H : [2][w]
H = [100,101]</pre>
```

results in a variable H which is an array of size 2 with width between 32 and 64. Note that

```
v1 : [2][32]
```

and

```
type v1 = [2][32]
```

express the same type.

Example:

Consider the Chicken, Fox, Grain riddle. A farmer has to get a fox, a chicken, and a sack of corn across a river, all intact. To do this the farmer has only a rowboat that is big enough to carry just the farmer and one other thing. If the fox and the chicken are left together, the fox

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will eat the chicken. If the chicken and the corn are left together, the chicken will eat the corn. The fox does not like corn. What sequence of river crossings will allow the farmer to arrive at the opposite bank without anything eating anything?

Begin setting up a Cryptol solution by considering problem state. There is a left bank and a right bank. Use a field of 4 bits to identify who or what, called a player, is on each bank. Let the farmer's status be represented by the least significant bit, the chicken's status by the 2^{nd} bit, the corn's status by the 3^{rd} bit, and the fox's status but the most significant bit. If a player is on the left bank then it's bit is 1 on the 4 bit field representing the left bank and it's bit is 0 on the 4 bit field representing the right bank and vice versa. So, in Cryptol, one can declare a type such as the following to represent bank state:

```
type OneBank = [4]
For the state of the problem, two banks can be combined like this:
   type BankState = { left : OneBank, right : OneBank }
Say the farmer starts on the left bank. The the initial state is
   startState = { left = Ob1111, right = Ob0000 }
the end state is
   endState = { left = Ob0000, right = Ob1111 }
```

A OneBank state is not *safe* if it is 0b1010, or 0b0110. One of those states means something has eaten something on the corresponding bank. The following function <code>safeAndValidState</code> returns <code>True</code> given <code>BankState</code> s if and only if both left and right <code>OneBank</code> states of s are safe and the <code>BankState</code> is valid which means a player with bit of value 1 on one bank has value 0 on the opposite bank.

```
safeAndValidState : BankState -> Bit safeAndValidState s = \sim(s.left == 10 \/ s.left == 6 \/ s.right == 10 \/ s.right == 0xF
```

The next function takes as input a (safe and valid) state s plus a position (0,1,2) idx and outputs a possible next state. If s.left is even then s.right is odd and represents the move to opposite bank. Then the next state left will be s.left $|\cdot|$ s.right masked with 0b0011 if idx = 0, 0b0101 if idx = 1 and 0b1001 if idx = 2. The next state right will be s.right && 0b0001 && 0b1100 if idx = 0, 0b1010 if idx = 1, and 0b0110 if idx = 2. Similarly for the case that s.left is odd and s.right is even. Here is the function:

```
{left = 0b1000, right = 0b0111}
Main> nextState t 2
{left = 0b0000, right = 0b1111}
```

If states represents the sequence of BankStates that is the solution, there are a few checks that must be made on states. First, all successive BankStates in the sequence states must be a legitimate nextState of the previous BankState. That check is done by the following:

Second, all states must be safe and valid. That may be taken care of by the following

```
allStatesSafeAndValid : {n} (fin n) => [n]BankState -> Bit allStatesSafeAndValid states = z \cdot 0 where z = [True]\#[safeAndValidState y / x | y <- states | x <- z ]
```

The first BankState should be the startState and the last BankState should be the endState. These may be encoded like this:

```
(states ! 0 == endState) /\ (states @ 0 == startState)
```

All BankStates in the solution sequence should be unique. The following can be used for that:

```
allStatesUnique states = \sim(z ! 0) where z = [False]#[ (states@i == states@j /\ \sim(i == j)) \/ k | k <- z | i <- [0..7] , j <- [0..7] ]
```

The 7 shows up in the above because a minimum solution contains 8 BankStates (a well known number). Putting all of the above together gives:

```
solution : [8]BankState -> Bit
property solution states =
   (neighborsConsistent states) /\
   (allStatesSafeAndValid states) /\
   (allStatesUnique states) /\
   (states ! 0 == endState) /\
   (states @ 0 == startState)
```

Notice the number 8 again indicating the expected solution size. Run it like this:

```
Main> :s base=2
Main> :s satNum=all  // print all solutions
Main> :sat solution
// Satisfiable
// solution
// [{left = 0b1111, right = 0b0000}, {left = 0b1100, right = 0b0011},
// {left = 0b1101, right = 0b0010}, {left = 0b1000, right = 0b0111},
// {left = 0b1011, right = 0b0100}, {left = 0b0010, right = 0b1101},
// {left = 0b0011, right = 0b1100}, {left = 0b0000, right = 0b1111}]
// = True
```

```
// solution
// [{left = 0b1111, right = 0b0000}, {left = 0b1100, right = 0b0011},
// {left = 0b1101, right = 0b0010}, {left = 0b0100, right = 0b1011},
// {left = 0b0111, right = 0b1000}, {left = 0b0010, right = 0b1101},
// {left = 0b0011, right = 0b1100}, {left = 0b0000, right = 0b1111}]
// = True
// Models found: 2
// (Total Elapsed Time: 0.108s, using "Z3")
```

The use of list comprehensions as seen above in all States Not Unique, all States Safe And Valid, and neighbors Consistent will be discussed with plenty of examples in the next lesson.