

MATH681  
Mathematical Logic

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February 15, 2023

## Abstract

This is a graduate-level mathematical logic course taught by [Matthew Harrison-Trainor](#), aiming to obtain insights into all other branches of mathematics, such as algebraic geometry, analysis, etc. Specifically, we will cover model theory beyond the basic foundational ideas of logic.

While there are no required textbooks, some books do cover part of the material in the class. For example, Marker's *Model Theory: An Introduction* [[Mar02](#)], Hodges's *A Shorter Model Theory* [[HH97](#)], and Hinman's *Fundamentals of Mathematical Logic* [[Hin05](#)].



# Contents

<b>1</b>	<b>Language, Logic, and Structures</b>	<b>2</b>
1.1	Syntax and Semantics . . . . .	3
1.2	Theories . . . . .	11
1.3	Completeness and Compactness . . . . .	14
1.4	A Detour to Algebraically Closed Fields . . . . .	27

# Chapter 1

## Language, Logic, and Structures

### Lecture 1: Introduction to Mathematical Logic

The goal of mathematical logic is to obtain insights into other areas of mathematics – algebra, analysis, combinatorics, and so on, by formalizing the **process** of mathematics. 5 Jan. 14:30

**Remark.** More concretely, there are different branches:

- (a) Model Theory: Study subsets of an object defined by a **formula** (i.e., first-order logic).
- (b) Computability Theory / Recursion Theory: Formalizing what it means to have an algorithm and studying relative computability.
- (c) Set Theory: Study the structure of the mathematical universe.
- (d) Proof Theory: Study the syntactic nature of **proofs**.

In this class, we study model theory in nature; specifically, we will cover

- basic definitions of logic:
  - What is a **formula**?
  - What does it mean for a **formula** to be **true**?
  - What is a **proof**?
- **Soundness** & completeness theorems:
  - Anything **provable** is **true**.
  - Anything **true** is **provable**.
- Compactness theorem:
  - Non-standard objects exist.
- Using compactness theorem for applications:
  - **Chevalley's theorem**.

The main theme of this course will be *syntax* v.s. *semantics*:

Syntax	v.s.	Semantics
<b>proofs</b>		<b>truth</b>
form of a <b>formula</b>		mathematical <b>structures</b>
number and type of quantifiers		<b>isomorphisms, embeddings</b>

## 1.1 Syntax and Semantics

### 1.1.1 Languages and Structures

Let's start with the fundamental object, [language](#).

**Definition 1.1.1 (Language).** A *language*  $\mathcal{L}$  consists of:

- a set  $\mathcal{F}$  of function symbols  $f$  with arities  $n_f$ ;
- a set  $\mathcal{R}$  of relation symbols  $R$  with arities  $n_R$ ;
- a set  $\mathcal{C}$  of constant symbols  $c$ .

A [language](#) is also sometimes called a *signature*, in which case we use  $\sigma$  rather than  $\mathcal{L}$ .

**Note.** A constant is the same as a 0-ary function.

**Remark.** Any or all sets in [Definition 1.1.1](#) might be empty.

**Example (Graph).** The [language](#) of graphs,  $\mathcal{L}_{\text{graph}} = \{E\}$  where  $E$  is a binary (2-ary) relation symbol.

**Example (Ring).** The [language](#) of rings,  $\mathcal{L}_{\text{ring}} = \{0, 1, +, \cdot, -\}$ , where  $0, 1$  are constants,  $+, \cdot$  are binary functions, and  $-$  is a unary function.

**Example (Ordered ring).** The [language](#) of ordered rings,  $\mathcal{L}_{\text{ord}} = \mathcal{L}_{\text{ring}} \cup \{\leq\}$  where  $\leq$  is the binary relation for an ordered ring.

Then, given a [language](#), we can now interpret it in the following way.

**Definition 1.1.2 (Structure).** Given a [language](#)  $\mathcal{L}$ , an  $\mathcal{L}$ -*structure*  $\mathcal{M}$  consists of:

- a non-empty set  $M$  called the *universe*, *domain*, or *underlying set* of  $\mathcal{M}$ ;
- for each function symbol  $f \in \mathcal{F}$ , a function  $f^{\mathcal{M}}: M^{n_f} \rightarrow M$ ;
- for each relation symbol  $R \in \mathcal{R}$ , a relation  $R^{\mathcal{M}} \subseteq M^{n_R}$ ;
- for each constant symbol  $c \in \mathcal{C}$ , an element  $c^{\mathcal{M}} \in M$ .

**Notation (Interpretation).** The *interpretation* of symbols  $f, R, c$  in  $\mathcal{M}$  is  $f^{\mathcal{M}}, R^{\mathcal{M}}, c^{\mathcal{M}}$ , respectively.

Basically, a [structure](#) gives meaning to the symbols from the [language](#), and we often write

$$\mathcal{M} = (M, f^{\mathcal{M}}, \dots, R^{\mathcal{M}}, \dots, c^{\mathcal{M}}, \dots) = (M, f^{\mathcal{M}}, R^{\mathcal{M}}, c^{\mathcal{M}}: f \in \mathcal{F}, R \in \mathcal{R}, c \in \mathcal{C}).$$

**Notation.** We usually use  $\mathcal{M}, \mathcal{N}, \dots, \mathcal{A}, \mathcal{B}, \dots$  to refer to [structures](#), and  $M, N, \dots, A, B, \dots$  for the domains.<sup>a</sup>

<sup>a</sup>Some people use  $|\mathcal{M}|$  for the domain of  $\mathcal{M}$ .

It's time to look at some examples.

**Example.** The rationals  $\mathbb{Q}$  and integers  $\mathbb{Z}$  are both  $\mathcal{L}_{\text{ring}}$ -structures.

**Proof.** Clearly, the domain is the set of rationals, and naively, we let  $+^{\mathbb{Q}} = +$  in  $\mathbb{Q}$ ,  $0^{\mathbb{Q}} = 0$  in

$\mathbb{Q}$ ,  $1^{\mathbb{Q}} = 1$  in  $\mathbb{Q}$ , etc. In this way,  $\mathbb{Q} = (\mathbb{Q}, 0, 1, +, \cdot, -)$  is an  $\mathcal{L}_{\text{ring}}$ -structure. Similarly,  $\mathbb{Z} = (\mathbb{Z}, 0, 1, +, \cdot, -)$  is as well.  $\circledast$

While the language we have seen are all intuitively correct with their name, e.g.,  $\mathcal{L}_{\text{ring}}$ ,  $\mathcal{L}_{\text{ord}}$ , and  $\mathcal{L}_{\text{graph}}$ , they are really just the high-level abstraction of the objects in the subscript.

**Example.** Nothing forces an  $\mathcal{L}_{\text{ring}}$ -structure to be a ring.

**Proof.** Since an  $\mathcal{L}_{\text{ring}}$ -structure is just any structure with two binary functions, a unary function, and two constants interpreting the symbols of the language; hence we can define an  $\mathcal{L}_{\text{ring}}$ -structure  $\mathcal{M}$  as

- $\mathcal{M} = \{0, 5, 11\}$ ;
- $0^{\mathcal{M}} = 5$ ;
- $1^{\mathcal{M}} = 11$ ;
- $+^{\mathcal{M}}$  is the constant function 0;
- $\cdot^{\mathcal{M}}$  is the function 5;
- $-^{\mathcal{M}}$  is the identity.

This is clearly not a ring since it fails nearly every axiom of a ring.  $\circledast$

**Note.** Later, we will talk about theories that let us restrict to structures we want.

### 1.1.2 Embeddings and Isomorphisms

We can now consider the relation between structures.

**Definition 1.1.3 (Embedding).** Let  $\mathcal{M}$  and  $\mathcal{N}$  be  $\mathcal{L}$ -structures. A map  $\eta: \mathcal{M} \rightarrow \mathcal{N}$  is an  $\mathcal{L}$ -embedding if it is one-to-one and preserves the interpretation of all symbols of  $\mathcal{L}$ :

- (a) for each function symbol  $f \in \mathcal{F}$  of arity  $n_f$ , and  $a_1, \dots, a_{n_f} \in M$ ,

$$\eta(f^{\mathcal{M}}(a_1, \dots, a_{n_f})) = f^{\mathcal{N}}(\eta(a_1), \dots, \eta(a_{n_f}));$$

- (b) for each relation symbol  $R \in \mathcal{R}$  of arity  $n_R$ , and  $a_1, \dots, a_{n_R} \in M$ ,

$$(a_1, \dots, a_{n_R}) \in R^{\mathcal{M}} \Leftrightarrow (\eta(a_1), \dots, \eta(a_{n_R})) \in R^{\mathcal{N}};$$

- (c) for each constant symbol  $c \in \mathcal{C}$ ,  $c^{\mathcal{M}} = c^{\mathcal{N}}$ .

From the definition, an  $\mathcal{L}$ -embedding is an injection, and naturally, we have the following.

**Definition 1.1.4 (Isomorphism).** An  $\mathcal{L}$ -isomorphism is a bijective  $\mathcal{L}$ -embedding.

**Definition 1.1.5 (Automorphism).** An  $\mathcal{L}$ -automorphism of  $\mathcal{M}$  is an  $\mathcal{L}$ -isomorphism from  $\mathcal{M}$  to  $\mathcal{M}$ .

**Definition.** Let  $\mathcal{M}$  and  $\mathcal{N}$  be  $\mathcal{L}$ -structures. Suppose  $M \subseteq N$  and the inclusion map  $\iota: \mathcal{M} \hookrightarrow \mathcal{N}$  is an  $\mathcal{L}$ -embedding.

**Definition 1.1.6 (Substructure).**  $\mathcal{M}$  is a substructure of  $\mathcal{N}$ .

**Definition 1.1.7 (Extension).**  $\mathcal{N}$  is an extension of  $\mathcal{M}$ .

**Example.** Ring embeddings are  $\mathcal{L}_{\text{ring}}$ -embeddings.

This generalizes the notions of embedding and isomorphism for many mathematical structures.

**Remark.** Asking that  $\eta$  be injective is the same as (b) in Definition 1.1.3 for the relation  $=$  since

$$a = b \in \mathcal{M} \Leftrightarrow \eta(a) = \eta(b) \in \mathcal{N}.$$

The notion of substructure is language sensitive. For groups, there are two possible languages:

- (a)  $\mathcal{L}_1 = \{e, \cdot\}$ ;
- (b)  $\mathcal{L}_2 = \{e, \cdot, {}^{-1}\}$ , i.e., with the unary inverse operation.

While both seem valid at the first glance, we should use the second one.

To see why, if we use  $\mathcal{L}_2$ , the substructure of a group is the same thing as a subgroup. But if we use  $\mathcal{L}_1$ , then  $(\mathbb{N}, +, 0)$  is a substructure of  $(\mathbb{Z}, +, 0)$ , while  $\mathbb{N}$  is not a group for sure.<sup>1</sup>

Similarly, we include  $-$  in  $\mathcal{L}_{\text{ring}}$  for a similar reason as in the previous example.

**Example.** An  $\mathcal{L}_{\text{ring}}$ -substructure of a field will be a subring, not a subfield. If we want subfields, use  $\mathcal{L}_{\text{ring}} \cup \{{}^{-1}\}$ .<sup>a</sup>

<sup>a</sup>We can set  $0^{-1} = 0$ , but never use this.

## Lecture 2: Formulas and First-Order Logic

We start by asking that given a function symbol  $f$  of arity  $n$ , could we replace  $f$  with an  $(n+1)$ -ary  $R$  relation to represent its graph? 10 Jan. 14:30

**Example.** Let  $\mathcal{L}$  be a language with only relation symbols. Let  $\mathcal{A}$  be an  $\mathcal{L}$ -structure. For any  $B \subseteq A$ , there is a substructure  $\mathcal{B}$  of  $\mathcal{A}$  with domain  $B$ .

**Proof.** For each relation symbol  $R$ , letting  $R^{\mathcal{B}} = R^{\mathcal{A}} \cap B^{n_R}$  will make  $\mathcal{B}$  a substructure of  $\mathcal{A}$ .  $\circledast$

The above is not true for function symbols though.

**Example.** If  $G = (\mathbb{Z}, 0, +)$ , then  $\mathbb{N}$  is not the domain of a subgroup. So if we took  $\mathcal{L} = \{0, +, {}^{-1}\}$ , where  $0$  is the unary relation,  $+$  is the ternary relation, and  ${}^{-1}$  is the binary relation, an  $\mathcal{L}$ -substructure of a group might not be a subgroup.

### 1.1.3 Terms

Intuitive, an  $\mathcal{L}$ -formula is an expression built using the symbols in a language  $\mathcal{L}$ ,  $=$ , the logical connectives  $\wedge, \vee, \neg$ , and variable symbols  $v_1, v_2, \dots, x, y, z$ , and also quantifiers  $\exists$  and  $\forall$ .

**Definition 1.1.8 (Term).** Given a language  $\mathcal{L}$ , the set of  $\mathcal{L}$ -terms are defined inductively by:

- (a) each constant symbol is a term;
- (b) each variable symbol  $v_1, \dots$  is a term;
- (c) if  $f$  is a function symbol, and  $t_1, \dots, t_{n_f}$  are terms, then  $f(t_1, \dots, t_{n_f})$  is a term.

If  $\mathcal{M}$  is an  $\mathcal{L}$ -structure, and  $t$  is a term involving only variables among  $v_1, \dots, v_n$ , then  $t$  has an interpretation  $t^{\mathcal{M}}: M^n \rightarrow M$  as a function as follows. On input  $a_1, \dots, a_n \in M$ ,

- (a) if  $t$  is a constant  $c$ ,  $t^{\mathcal{M}}(a_1, \dots, a_n) = c^{\mathcal{M}}$ .
- (b) if  $t$  is a variable  $v_i$ ,  $t^{\mathcal{M}}(a_1, \dots, a_n) = v_i$ ;

<sup>1</sup>Simply observe that both  $(\mathbb{N}, 0, +)$ ,  $(\mathbb{Z}, 0, +)$  are  $\mathcal{L}_1$ -structures.

(c) if  $t$  is  $f(s_1, \dots, s_k)$ , then  $t^{\mathcal{M}}(a_1, \dots, a_n) = f^{\mathcal{M}}(s_1^{\mathcal{M}}(a_1, \dots, a_n), \dots, s_k^{\mathcal{M}}(a_1, \dots, a_n))$ .

**Intuition.** We are basically substituting for variables and evaluating the expression.

**Example.** In  $(\mathbb{R}, 0, 1, +, \cdot, -)$ , a **term** is essentially just a polynomial with integer coefficients, assuming we interpret them in a ring. Technically, a **term** looks like

$$\cdot(+ (1, 1), + (x, y)),$$

but we will write **terms** the natural way, i.e.,

$$(1 + 1)(x + y).$$

Also, we will use  $\underline{n}$  or  $n$  to represent the **term**  $\underline{n} = \underbrace{1 + 1 + \dots + 1}_{n \text{ times}}$ . So we could write the above **term** as  $2 \cdot (x + y)$ .

#### 1.1.4 Formulas

**Definition 1.1.9 (Formula).** The set of  $\mathcal{L}$ -formulas is defined inductively:

- (a) If  $s, t$  are **terms**, then  $s = t$  is a *formula*.
- (b) If  $R$  is a relation symbol of arity  $n_R$  and  $s_1, \dots, s_{n_R}$  are **terms**, then  $R(s_1, \dots, s_{n_R})$  is a *formula*.
- (c) If  $f$  is a **formula**, then  $\neg f$  is a *formula*.
- (d) If  $\varphi$  and  $\psi$  are **formulas**, then  $\varphi \wedge \psi$  and  $\varphi \vee \psi$  are *formulas*.
- (e) If  $\varphi$  is a **formula** and  $v_i$  are variables, then  $\exists v_i \varphi$  and  $\forall v_i \varphi$  are *formulas*.

**Notation** (Atomic formula). **Definition 1.1.9 (a)** and **(b)** are called *atomic formulas*.

**Notation** (Quantifier-free formula). **Definition 1.1.9 (a), (b), (c), and (d)** are called *quantifier-free formulas*.

This logic is called *first-order logic* (FO logic), since the quantifiers range over elements of the **structures**, but not over, e.g., subsets.

**Example.** We can say that an element  $x$  of a ring has a square root by  $\exists y \ y^2 = x$ .

**Example.** A group is torsion of order 2 can be said by  $\forall x \ x \cdot x = e$ .

**Example.** We can write down all the field/group/... axioms as **formulas**.

Notice that for the first example, the **formula**  $\exists y \ y^2 = x$  only has meaning if we assign what  $x$  is. In this case, we say that  $y$  is *bound* by  $\exists y$ . But this is local:

**Example.** Consider

$$y = 1 \wedge \exists y \ y^2 = x,$$

while the first appearance of  $y$  is free, the second appearance of  $y$  is bound by (in the scope of)  $\exists y$ .

While our definitions work perfectly fine with the above example, but sometimes we don't want this to happen. In such a case, we simply replace the bound instances of  $y$  with a new variable  $z$ . This idea of variables being free or bound is defined formally as follows.



**Definition 1.1.10 (Free variable).** The *free variables*  $\text{FV}(\varphi)$  of a **formula**  $\varphi$  are defined inductively:

- (a)  $\text{FV}(s = t)$  is the set of variables showing up in  $s$  or  $t$ .
- (b)  $\text{FV}(R(s_1, \dots, s_{n_R}))$  is the set of variables showing up in  $s_1, \dots, s_{n_R}$ .
- (c)  $\text{FV}(\neg\varphi) = \text{FV}(\varphi)$ .
- (d)  $\text{FV}(\varphi \wedge \psi) = \text{FV}(\varphi \vee \psi) = \text{FV}(\varphi) \cup \text{FV}(\psi)$ .
- (e)  $\text{FV}(\exists x \varphi) = \text{FV}(\forall x \varphi) = \text{FV}(\varphi) \setminus \{x\}$ .

**Example.**  $\text{FV}(\exists y y^2 = x) = \{x\}$ .

**Example.**  $\text{FV}(\forall x x \cdot x = e) = \emptyset$ .

**Definition 1.1.11 (Sentence).** A **formula**  $\varphi$  is called a *sentence* if it has no **free variables**.

**Notation.** If  $\varphi$  is a **formula** with **free variables** among  $x_1, \dots, x_n$  we often write  $\varphi(x_1, \dots, x_n)$ .

**Remark.** So given  $\varphi(x_1, \dots, x_n)$ , we know that  $\varphi$  has no other **free variables** than  $x_1, \dots, x_n$ .

**Example.** It's valid to write  $\varphi(x, y, z) := x = y$ .

### 1.1.5 Truths

Finally, we define the notion of **truth**.

**Definition 1.1.12 (Truth).** Given an  $\mathcal{L}$ -**structure**  $\mathcal{M}$ , let  $\varphi(x_1, \dots, x_n)$  be an  $\mathcal{L}$ -**formula** and let  $a_1, \dots, a_n \in M$ . Then we say  $\varphi$  is *true* of  $\bar{a}$  in  $\mathcal{M}$ ,<sup>a</sup> denoted as  $\mathcal{M} \models \varphi(\bar{a})$ , as follows:

- (a) If  $\varphi$  is  $s = t$ , then  $\mathcal{M} \models \varphi(\bar{a})$  if  $s^{\mathcal{M}}(\bar{a}) = t^{\mathcal{M}}(\bar{a})$ .
- (b) If  $\varphi$  is  $R(t_1, \dots, t_{n_R})$ , then  $\mathcal{M} \models \varphi(\bar{a})$  if  $(t_1^{\mathcal{M}}(\bar{a}), \dots, t_{n_R}^{\mathcal{M}}(\bar{a})) \in R^{\mathcal{M}}$ .
- (c) If  $\varphi$  is  $\neg\psi$ , then  $\mathcal{M} \models \varphi(\bar{a})$  if  $\mathcal{M} \not\models \psi(\bar{a})$ .
- (d) If  $\varphi$  is  $\psi_1 \wedge \psi_2$ , then  $\mathcal{M} \models \varphi(\bar{a})$  if  $\mathcal{M} \models \psi_1(\bar{a})$  and  $\mathcal{M} \models \psi_2(\bar{a})$ .
- (e) If  $\varphi$  is  $\psi_1 \vee \psi_2$ , then  $\mathcal{M} \models \varphi(\bar{a})$  if  $\mathcal{M} \models \psi_1(\bar{a})$  or  $\mathcal{M} \models \psi_2(\bar{a})$ .
- (f) If  $\varphi$  is  $\exists y \psi(\bar{x}, y)$ , then  $\mathcal{M} \models \varphi(\bar{a})$  if there's  $b \in M$  such that  $\mathcal{M} \models \psi(\bar{a}, b)$ .
- (g) If  $\varphi$  is  $\forall y \psi(\bar{x}, y)$ , then  $\mathcal{M} \models \varphi(\bar{a})$  if for all  $b \in M$  such that  $\mathcal{M} \models \psi(\bar{a}, b)$ .

<sup>a</sup>Or  $\mathcal{M}$  satisfies  $\varphi(\bar{a})$ .

**Remark.** Every **formula** is **true**, or its negation is.

## Lecture 3: Logical Consequence and Equivalence

**Notation** (Material implication). The *material implication*  $\varphi \rightarrow \psi$  between two **formulas**  $\varphi, \psi$  is an abbreviation of  $\neg\varphi \vee \psi$ .

12 Jan. 14:30

**Notation.** We use  $\varphi \leftrightarrow \psi$  as an abbreviation of  $((\varphi \rightarrow \psi) \wedge (\psi \rightarrow \varphi))$ .

Essentially,  $\rightarrow$  and  $\leftrightarrow$  is different from  $\Rightarrow$  and  $\Leftrightarrow$ , where the former are only shown in [formula](#). Now, consider the [language of graphs](#)  $\mathcal{L}_{\text{graph}} = \{E\}$ , let's see some examples.

**Example.** An undirected graph can be written as

$$\forall x \forall y (xEy \rightarrow yEx).$$

**Example.** A vertex has at least three neighbors can be written as

$$\varphi(x) := \exists u \exists v \exists w (xEu \wedge xEv \wedge xEw \wedge u \neq v \wedge v \neq w \wedge u \neq w)$$

in non-reflexive graphs.

**Example.** For a vertex has exactly three neighbors,

$$\psi(x) := \exists u \exists v \exists w \forall y (xEu \wedge xEv \wedge xEw \wedge u \neq v \wedge v \neq w \wedge u \neq w \wedge (y = u \vee y = v \vee y = w \vee \neg yEx)).$$

**Problem.** Can we say that  $x$  has an even number of neighbors?

**Answer.** We can't. Some things are not expressible in FO logic. ⊛

**Example.** For a vertex  $x$  has a path of length 4 to  $y$ ,

$$\Theta(x, y) := \exists u \exists v \exists w (xEu \wedge uEv \wedge vEw \wedge wEy).$$

We can also express that there is a path of length at most 4.

**Problem.** Can we say that there is a path from  $x$  to  $y$ ?

**Answer.** We still can't! Not in FO logic (using [compactness theorem](#)). ⊛

**Remark.** When we prove results by induction on [formulas](#), we only need to prove for  $\neg, \wedge, \exists$ , instead of for both  $\wedge, \vee$ , and both  $\exists$  and  $\forall$ .

**Proof.** Since we can view  $\varphi \vee \psi$  as an abbreviation for  $\neg(\neg\varphi \wedge \neg\psi)$  and  $\forall x \varphi$  as an abbreviation for  $\neg(\exists x \neg\varphi)$ . ⊛

**Remark (Sheffer stroke).** In fact, we can get  $\wedge, \vee, \neg$  from one logical connective, e.g., the *sheffer stroke*  $\uparrow$ , which is defined as

$$\varphi \uparrow \psi := \neg(\varphi \wedge \psi),$$

and we can use  $\uparrow$  to define  $\neg, \vee, \wedge$ .

**Notation.** Let  $\Phi$  be a (possibly infinite) set of [sentences](#), we write  $\mathcal{M} \models \Phi$  if  $\mathcal{M} \models \varphi$  for all  $\varphi \in \Phi$ .

**Definition 1.1.13 (Logical consequence).** Let  $\Phi$  be a set of [sentences](#), and  $\varphi$  be a [sentence](#). We say that  $\varphi$  is a *logical consequence* of  $\Phi$ , written  $\Phi \models \varphi$ , if  $\mathcal{M} \models \varphi$  whenever  $\mathcal{M} \models \Phi$ .

If  $\Phi = \emptyset$  is the empty set, [Definition 1.1.13](#) is written as  $\models \varphi$ , i.e.,  $\varphi$  is [true](#) in all  $\mathcal{L}$ -structures.<sup>2</sup>

<sup>2</sup>Recall that we always have a [language](#)  $\mathcal{L}$  implicitly.

**Definition 1.1.14 (Equivalent).** Given two formulas  $\varphi, \psi$ ,  $\varphi(\bar{x})$  and  $\psi(\bar{x})$  are *equivalent* if

$$\models \forall \bar{x} (\varphi(\bar{x}) \leftrightarrow \psi(\bar{x})).$$

**Problem.** Two sentences  $\varphi$  and  $\psi$  are *equivalent* if and only if  $\varphi \models \psi$  and  $\psi \models \varphi$ .

DIY

**As previously seen.**  $\mathcal{A}$  is a *substructure* of  $\mathcal{B}$ , or  $\mathcal{A} \subseteq \mathcal{B}$ , means that  $A \subseteq B$  and  $\text{id}: A \hookrightarrow B$  is an  $\mathcal{L}$ -embedding.

**Proposition 1.1.1.** Suppose that  $\mathcal{A}$  is a *substructure* of  $\mathcal{B}$ , and  $\varphi(\bar{x})$  is a *quantifier-free formula*. Let  $\bar{a} \in \mathcal{A}$ ,<sup>a</sup> then  $\mathcal{A} \models \varphi(\bar{a})$  if and only if  $\mathcal{B} \models \varphi(\bar{a})$ .

<sup>a</sup>Formally, we need to write  $\mathcal{A}$  to be the Cartesian product with a fixed length.

**Proof.** We start with *terms* by proving that if  $t$  is a *term* and  $\bar{b} \in \mathcal{A}$ , then  $t^{\mathcal{A}}(\bar{b}) = t^{\mathcal{B}}(\bar{b})$ . The proof is induction on *terms*.

- (a) If  $t$  is a constant symbol  $c$ , then  $t^{\mathcal{A}}(\bar{b}) = c^{\mathcal{A}} = c^{\mathcal{B}} = t^{\mathcal{B}}(\bar{b})$ .
- (b) If  $t$  is a variable  $x_i$ , then  $t^{\mathcal{A}}(\bar{b}) = b_i = t^{\mathcal{B}}(\bar{b})$ .
- (c) If  $t$  is a function symbol  $f(s_1, \dots, s_n)$  where  $s_i$  are *terms*, then  $t^{\mathcal{A}}(\bar{b}) = f^{\mathcal{A}}(s_1^{\mathcal{A}}(\bar{b}), \dots, s_n^{\mathcal{A}}(\bar{b}))$ .  
By the induction hypothesis,  $s_i^{\mathcal{A}}(\bar{b}) = s_i^{\mathcal{B}}(\bar{b}) \in \mathcal{A}$ , and hence

$$t^{\mathcal{B}}(\bar{b}) = f^{\mathcal{B}}(s_1^{\mathcal{B}}(\bar{b}), \dots, s_n^{\mathcal{B}}(\bar{b})) = f^{\mathcal{A}}(s_1^{\mathcal{A}}(\bar{b}), \dots, s_n^{\mathcal{A}}(\bar{b})) = t^{\mathcal{A}}(\bar{b}),$$

i.e.,  $f^{\mathcal{B}} \upharpoonright_{\mathcal{A}} = f^{\mathcal{A}}$ , so  $t^{\mathcal{A}}(\bar{b}) = t^{\mathcal{B}}(\bar{b})$ .

Now we turn to *formulas*, and prove that for  $\varphi$  *quantifier-free*, then  $\mathcal{A} \models \varphi(\bar{a}) \Leftrightarrow \mathcal{B} \models \varphi(\bar{a})$  for  $\bar{a} \in \mathcal{A}$ . The proof is, again, induction on *formulas*.<sup>a</sup>

- (a) If  $\varphi$  is  $s = t$ , then  $s^{\mathcal{A}}(\bar{a}) = s^{\mathcal{B}}(\bar{a})$  and  $t^{\mathcal{A}}(\bar{a}) = t^{\mathcal{B}}(\bar{a})$ , so

$$\mathcal{A} \models \varphi(\bar{a}) \Leftrightarrow s^{\mathcal{A}}(\bar{a}) = t^{\mathcal{A}}(\bar{a}) \Leftrightarrow s^{\mathcal{B}}(\bar{a}) = t^{\mathcal{B}}(\bar{a}) \Leftrightarrow \mathcal{B} \models \varphi(\bar{a}).$$

- (b) If  $\varphi$  is  $R(s_1, \dots, s_n)$ , then

$$\mathcal{A} \models \varphi(\bar{a}) \Leftrightarrow (s_1^{\mathcal{A}}(\bar{a}), \dots, s_n^{\mathcal{A}}(\bar{a})) \in R^{\mathcal{A}} \Leftrightarrow (s_1^{\mathcal{B}}(\bar{a}), \dots, s_n^{\mathcal{B}}(\bar{a})) \in R^{\mathcal{B}} \Leftrightarrow \mathcal{B} \models \varphi(\bar{a}).$$

- (c) If  $\varphi$  is  $\neg\psi$ ,

$$\mathcal{A} \models \varphi(\bar{a}) \Leftrightarrow \mathcal{A} \not\models \psi(\bar{a}) \Leftrightarrow \mathcal{B} \not\models \psi(\bar{a}) \Leftrightarrow \mathcal{B} \models \varphi(\bar{a}),$$

where we use the induction hypothesis in the second  $\Leftrightarrow$ .

- (d) If  $\varphi$  is  $\psi_1 \vee \psi_2$ ,

$$\mathcal{A} \models \varphi(\bar{a}) \Leftrightarrow \mathcal{A} \models \psi_1(\bar{a}) \text{ or } \mathcal{A} \models \psi_2(\bar{a}) \Leftrightarrow \mathcal{B} \models \psi_1(\bar{a}) \text{ or } \mathcal{B} \models \psi_2(\bar{a}) \Leftrightarrow \mathcal{B} \models \varphi(\bar{a}),$$

where we use the induction hypothesis in the second  $\Leftrightarrow$ .

■

<sup>a</sup>Recall that we only need to show one of  $\vee$  or  $\wedge$ , and here we pick  $\vee$  and treat  $\wedge$  as an abbreviation.

**As previously seen (Characteristic).** Given a field  $K$ , the *characteristic*  $p$  of  $K$  is the number of 1 you need to add 1 in order to get 0, i.e.,  $\underbrace{1 + 1 + \dots + 1}_p = 0$ .

**Example.** Let  $L$  be a subfield of  $K$ , for each  $p > 0$ ,  $\varphi_p := \underbrace{1 + 1 + \dots + 1}_p = 0$ , which says the characteristic  $p$ .  $\varphi_p$  is **quantifier-free**, so

$$L \models \varphi_p \Leftrightarrow K \models \varphi_p.$$

**Example.** Consider  $\mathbb{Z} = (\mathbb{Z}, 0, 1, +, -, \cdot)$ , and let  $\varphi(x) := \neg \exists y \ y + y = x$ . We see that  $\mathbb{Z} \models \varphi(1)$  but  $\mathbb{Q} \models \neg \varphi(1)$ .

**Proposition 1.1.2.** Suppose that  $\mathcal{A}$  is a **substructure** of  $\mathcal{B}$ , and  $\varphi(\bar{x}, y_1, \dots, y_n)$  is a **quantifier-free formula**. Let  $\bar{a} \in \mathcal{A}$ , then

- (a) if  $\mathcal{A} \models \exists y_1 \dots \exists y_n \ \varphi(\bar{a}, y_1, \dots, y_n)$ , then  $\mathcal{B} \models \exists y_1 \dots \exists y_n \ \varphi(\bar{a}, y_1, \dots, y_n)$ ;
- (b) if  $\mathcal{B} \models \forall y_1 \dots \forall y_n \ \varphi(\bar{a}, y_1, \dots, y_n)$ , then  $\mathcal{A} \models \forall y_1 \dots \forall y_n \ \varphi(\bar{a}, y_1, \dots, y_n)$ .

**Proof.** Suppose that  $\mathcal{A} \models \exists y_1 \dots \exists y_n \ \varphi(\bar{a}, y_1, \dots, y_n)$ , so there are  $b_1, \dots, b_n \in \mathcal{A}$  such that  $\mathcal{A} \models \varphi(\bar{a}, b_1, \dots, b_n)$ . Since  $\varphi$  is **quantifier-free**, so  $\mathcal{B} \models \varphi(\bar{a}, b_1, \dots, b_n)$  from **Proposition 1.1.1**, and hence  $\mathcal{B} \models \exists y_1 \dots \exists y_n \ \varphi(\bar{a}, y_1, \dots, y_n)$ .

On the other hand, it's easy to see that (b) is implied by (a). ■

**Notation.** In **Proposition 1.1.2**, formulas as in (a) are called *existential* ( $\exists_1$  or  $\exists$ ) *formulas*; and formulas as in (b) are called *universal* ( $\forall_1$  or  $\forall$ ) *formulas*.

**Example.** Recall  $\mathcal{L}_1 = \{e, \cdot\}$ ,  $\mathcal{L}_2 = \{e, \cdot, {}^{-1}\}$ .

- Associativity:  $\forall x \forall y \forall z \ (xy)z = x(yz)$ .
- Identity:  $\forall x \ ex = xe$ .

These are  $\forall$ -formulas in either language.

- Inverses in  $\mathcal{L}_1$ :  $\forall x \exists y \ xy = yx = e$ , which is **not** an  $\forall$ -formula.
- Inverses in  $\mathcal{L}_2$ :  $\forall x \ xx^{-1} = x^{-1}x = e$ , which is an  $\forall$ -formula.

Hence, group axioms in  $\mathcal{L}_1$  are not universal, but in  $\mathcal{L}_2$  they are.

The above discrepancy is the reason why  $\mathcal{L}_2$  is better than  $\mathcal{L}_1$ , i.e.,  $\mathcal{L}_1$ -substructure might not be a group.

**Problem.** Show that  $\forall x \exists y \ xy = yx = e$  in the above example is not **equivalent** to an  $\forall$ -formula.

## Lecture 4: Theories and Axioms

**Example.** Let  $\mathcal{L}_1 = \{E\}$ , where  $E$  is a binary relation representing edge relation; and  $\mathcal{L}_2 = \{V, E, I\}$ , where  $V, E$  are unary relations and  $I$  is a binary relation representing incidence such that  $I(v, e)$  for  $v \in V, e \in E$  means that  $v$  is a vertex on edge  $e$ . Then,

- Let  $G$  be a graph, viewed as an  $\mathcal{L}_1$ -structure. A **substructure** of  $G$  is an induced subgraph  $H \subseteq G$  such that any edge in  $G$  between two vertices of  $H$  is in  $H$ .
- If we view  $G$  as an  $\mathcal{L}_2$ -substructure, a **substructure** is a subgraph  $H$  such that  $H$  has some vertices and edges from  $G$ .<sup>a</sup>

<sup>a</sup>But there might be edges in  $H$  with no vertices, which can be fixed by having two functions  $I_1(e) = v, I_2(e) = w$  when  $e: v \rightarrow w$ .

The difference is that for  $\mathcal{L}_1$ , having an edge is **quantifier-free**, while in  $\mathcal{L}_2$  is existential. To elaborate a bit further, for  $\mathcal{L}_2$ ,  $vEw$  is **quantifier-free**, while in  $\mathcal{L}_2$ ,

$$\exists (v \in V \wedge w \in V \wedge e \in E \wedge I(v, e) \wedge I(w, e))$$

is not **quantifier-free**.

## 1.2 Theories

Let's start by the notion of **theory**.

**Definition 1.2.1 (Theory).** An  $\mathcal{L}$ -theory is a set of  $\mathcal{L}$ -sentences.

**Definition 1.2.2 (Model).**  $\mathcal{M}$  is a *model* of a **theory**  $T$ , written as  $\mathcal{M} \models T$ , if  $\mathcal{M} \models \varphi$  for all  $\varphi \in T$ .

**Note.** Not every **theory** has a **model**, e.g.,  $\{\exists x x \neq x\}$ .

The above note motivates the following.

**Definition 1.2.3 (Satisfiable).** A **theory** is *satisfiable* if it has a **model**.

**Definition 1.2.4 (Elementary class).** A class  $\mathcal{K}$  of  $\mathcal{L}$ -structures  $\mathcal{M}$  is called an *elementary class* if there is an  $\mathcal{L}$ -theory  $T$  such that

$$\mathcal{K} = \{\mathcal{M} \mid \mathcal{M} \models T\}.$$

One way to get an **elementary class** is to take an  $\mathcal{L}$ -structure  $\mathcal{M}$  and take the **full theory**.

**Definition 1.2.5 (Full theory).** The *full theory*  $\text{Th}(\mathcal{M})$  of an  $\mathcal{L}$ -structure  $\mathcal{M}$  is defined as  $\text{Th}(\mathcal{M}) = \{\varphi \mid \mathcal{M} \models \varphi\}$ .

From the definition,  $\mathcal{M} \models \text{Th}(\mathcal{M})$ , and  $\text{Th}(\mathcal{M})$  characterizes the **structures** satisfying the same **sentences** as  $\mathcal{M}$ .

**Definition 1.2.6 (Complete).** A **theory**  $T$  is *complete* if for any **sentence**  $\varphi$ , either  $\varphi \in T$  or  $\neg\varphi \in T$ .

**Remark.**  $\text{Th}(\mathcal{M})$  is **complete**.

**Definition 1.2.7 (Elementarily equivalent).**  $\mathcal{M}$  and  $\mathcal{N}$  are *elementarily equivalent*  $\mathcal{M} \equiv \mathcal{N}$  if for all **sentences**  $\varphi$ ,

$$\mathcal{M} \models \varphi \Leftrightarrow \mathcal{N} \models \varphi.$$

**Remark.** There are  $\mathcal{N} \models \text{Th}(\mathbb{N})$ , but  $\mathcal{N}$  is not isomorphic to  $\mathbb{N}$ .  $\mathcal{N}$  is called a *non-standard model of arithmetic*, and  $\mathcal{N}$  might have *infinite element* larger than all of  $\mathbb{N}$ . Here,  $\mathbb{N} = (\mathbb{N}, 0, 1, +, \cdot, -)$

**Example.**  $\mathbb{Z} \oplus \mathbb{Z} \not\cong \mathbb{Z}$  as groups.

The other way to define a **theory** is to write down axioms.

**Example (Infinite set).** Let  $\mathcal{L} = \emptyset$ , and let  $T$  consist of

$$\varphi_n := \exists x_1 \dots \exists x_n \bigwedge_{i \neq j} x_i \neq x_j.$$

**Example (Linear order).** Let  $\mathcal{L} = \{\leq\}$ , and let  $T$  consist of the axioms of linear orders, e.g.,

$$\forall x \forall y (x \leq y \wedge y \leq x \rightarrow x = y).$$

There are other interesting theories of linear orders, e.g., dense ones.

**Example (Dense linear order).** Consider

$$\forall x \forall y (x < y \rightarrow \exists z x < z < y),$$

where we use  $a < b$  as shorthand of saying  $a \leq b \wedge a \neq b$ .

**Example (Group).** In  $\mathcal{L}_{\text{group}} = \{e, \cdot, {}^{-1}\}$ , let  $T$  be the group axioms.

Other theories of groups include Abelson group, divisible, etc.

**Definition 1.2.8 (Finitely axiomatizable).** A theory is *finitely axiomatizable* if it has a finite set of axioms.

Given a theory, consider  $T^{\models} = \{\varphi \mid T \models \varphi\}$ ,<sup>3</sup> so  $\mathcal{M} \models T$  if and only if  $\mathcal{M} \models T^{\models}$ . Often we think of  $T$  and  $T^{\models}$  as the same. A theory  $T$  is *finitely axiomatizable* if there is a finite  $\Phi$  such that  $T^{\models} = \Phi^{\models}$ .

### 1.2.1 Elementary Embeddings

Let's now consider the following notion.

**Definition 1.2.9 (Elementary embedding).** Let  $\mathcal{M}$  and  $\mathcal{N}$  be  $\mathcal{L}$ -structures, and  $f: \mathcal{M} \rightarrow \mathcal{N}$  an  $\mathcal{L}$ -embedding. Then  $f$  is an *elementary embedding* if for any formula  $\varphi(\bar{x})$  and  $\bar{a} \in M$ ,

$$\mathcal{M} \models \varphi(\bar{a}) \Leftrightarrow \mathcal{N} \models \varphi(f(\bar{a})).$$

**Definition 1.2.10 (Elementary substructure).** If  $f: \mathcal{M} \hookrightarrow \mathcal{N}$  is a *elementary embedding* where  $\mathcal{M}$  is a *substructure* of  $\mathcal{N}$ , then  $\mathcal{M}$  is an *elementary substructure* of  $\mathcal{N}$ .

**Example.** As groups,  $\mathbb{Z} \hookrightarrow \mathbb{Q}$  is not *elementary*. In fact,  $\mathbb{Z} \not\equiv \mathbb{Q}$ . Whereas, if  $f: \mathcal{M} \hookrightarrow \mathcal{N}$  is an *elementary embedding*,  $\mathcal{M} \equiv \mathcal{N}$ .<sup>a</sup>

<sup>a</sup>And also much more is true.

**Proposition 1.2.1.** Every *isomorphism* is an *elementary embedding*.

**Proof.** Let  $f: \mathcal{M} \rightarrow \mathcal{N}$  be an *isomorphism*. We will argue by induction on formulas  $\varphi$ , that for all  $\bar{a} \in M$ ,

$$\mathcal{M} \models \varphi(\bar{a}) \Leftrightarrow \mathcal{N} \models \varphi(f(\bar{a})).$$

Firstly, observe that all cases except quantifiers are the same as [Proposition 1.1.1](#). For quantifiers, suppose that  $\varphi(\bar{x})$  is  $\exists y \psi(\bar{x}, y)$  and  $\mathcal{M} \models \varphi(\bar{a})$ . This means that there is  $b \in M$  such that  $\mathcal{M} \models \psi(\bar{a}, b)$ . By the induction hypothesis,  $\mathcal{N} \models \psi(f(\bar{a}), f(b))$ , so  $\mathcal{N} \models \varphi(f(\bar{a}))$ .

Now suppose  $\mathcal{N} \models \varphi(f(\bar{a}))$ , then there is  $c \in N$  such that  $\mathcal{N} \models \psi(f(\bar{a}), c)$ . Since  $f$  is an *isomorphism*, so there is a  $b \in M$  such that  $f(b) = c$ . By the induction hypothesis,  $\mathcal{M} \models \psi(\bar{a}, b)$ , so  $\mathcal{M} \models \varphi(\bar{a})$ . ■

**Corollary 1.2.1.** If  $\mathcal{M} \cong \mathcal{N}$ , then  $\mathcal{M} \equiv \mathcal{N}$ .

<sup>3</sup>Recall [Definition 1.1.13](#).

### 1.2.2 Definable Sets

Consider the following.

**Definition 1.2.11 (Definable).** Let  $\mathcal{M}$  be an  $\mathcal{L}$ -structure, then  $X \subseteq M^n$  is *definable* if there is a formula  $\varphi(x_1, \dots, x_n, \bar{y})$  and  $\bar{b} \in M$  such that

$$X = \{\bar{a} \in M^n \mid \mathcal{M} \models \varphi(\bar{a}, \bar{b})\}.$$

**Notation (Define).** We say that  $\varphi(\bar{x}, \bar{b})$  *defines*  $X$  over  $\bar{b}$ , written as  $X = \varphi(\mathcal{M}, \bar{b})$ .

**Notation (Parameter).** The tuple  $\bar{b}$  is called the *parameters* when  $X$  is *definable* over  $\bar{b}$ .

**Remark.** Sometimes  $X$  is *definable* without *parameters*, or *definable* over  $\emptyset$ .

**Example.** Take  $\mathbb{R} = (\mathbb{R}, 0, 1, +, \cdot, -)$  in  $\mathcal{L}_{\text{ring}}$ , then

$$\leq = \{(a, b) : a \leq b\}$$

is *definable*.

**Example.** Let  $\mathbb{Z} = (\mathbb{Z}, +, -, \cdot, 0, 1)$ , then  $\mathbb{N}$  is  $\emptyset$ -*definable* in  $\mathbb{Z}$  by<sup>a</sup>

$$\mathbb{N} = \{z \in \mathbb{Z} : \exists u, v, x, y \ u^2 + v^2 + x^2 + y^2 = z\}.$$

<sup>a</sup>From the **Langrange's four-square theorem**, which says that every natural number is the sum of four squares.

**Example.**  $\mathbb{Z}$  is  $\emptyset$ -*definable* in  $\mathbb{Q} = (\mathbb{Q}, +, -, \cdot, 0, 1)$ . This is a result of Julia Robinson [Rob49], and the formulation is very complicated.

**Problem.** How does one show that a set is not *definable*? For example,  $\mathbb{R}$  is not *definable* in  $\mathbb{C} = (\mathbb{C}, 0, 1, +, \cdot, -)$ .

## Lecture 5: Hilbert-Style Deductive System

We start by asking whether  $\mathbb{R}$  is *definable* in  $\mathbb{C} = (\mathbb{C}, 0, 1, +, \cdot, -)$ ?

19 Jan. 14:30

**Proposition 1.2.2.** Let  $\mathcal{M}$  be an  $\mathcal{L}$ -structure, and let  $X \subseteq M^n$  be a set which is *definable* over  $\bar{a}$ . Then any *automorphism* of  $\mathcal{M}$  that fixes  $\bar{a}$  pointwise<sup>a</sup> fixes  $X$  setwise.<sup>b</sup>

<sup>a</sup>If  $\bar{a} = (a_1, \dots, a_m)$ , then  $f(a_i) = a_i$ .

<sup>b</sup>If  $b \in X$ , then  $f(b) \in X$ .

**Proof.** Let  $f$  be an *automorphism* of  $\mathcal{M}$  fixing  $\bar{a}$  pointwise, and  $X = \{\bar{b} \in M^n : \mathcal{M} \models \varphi(\bar{b}, \bar{a})\}$ . Fix  $\bar{b}$ , and suppose  $\bar{b} \in X$ , so  $\mathcal{M} \models \varphi(\bar{b}, \bar{a})$ . Because  $f$  is an *elementary embedding* from **Proposition 1.2.1**,

$$\mathcal{M} \models \varphi(f(\bar{b}), f(\bar{a})) \Rightarrow \mathcal{M} \models \varphi(f(\bar{b}), \bar{a}),$$

hence  $f(\bar{b}) \in X$ . Similarly, if  $\bar{b} \notin X$ ,  $\mathcal{M} \models \neg\varphi(\bar{b}, \bar{a}) \Rightarrow \mathcal{M} \models \neg\varphi(f(\bar{b}), \bar{a})$ , so  $f(\bar{b}) \notin X$ . ■

**Remark.** If  $X$  is  $\emptyset$ -*definable*, it is fixed setwise by any *automorphism*.

**Example.**  $\mathbb{N}$  is fixed setwise by any **automorphism** of the ring  $\mathbb{Z}$ . In fact, the only **automorphism** of  $\mathbb{Z}$  is the identity.

**Example.**  $\mathbb{N}$  is not  **$\emptyset$ -definable** in  $\mathbb{Z} = (\mathbb{Z}, 0, +)$ .

**Proof.** Consider an **automorphism**  $f(x) = -x$  of the group  $\mathbb{Z}$ , which does not fix  $\mathbb{N}$  setwise. ⊗

**Problem.** Is  $\mathbb{N}$  **definable** in  $\mathbb{Z} = (\mathbb{Z}, 0, +)$  over some parameters  $\bar{a}$ ?

**Answer.** For example, if  $\bar{a} = (1)$ , then  $f$  does not fix 1. In fact, any **automorphism** fixing 1 also fixes all of  $\mathbb{Z}$ , but  $\mathbb{N}$  is not **definable** in  $(\mathbb{Z}, 0, +)$ . To prove this we need **compactness**. ⊗

**As previously seen.** Given a field  $F$ , then  $F(a) \cong F(b)$  if  $a$  and  $b$  have the same minimal polynomial over  $F$  or if both do not satisfy any polynomial over  $F$ .

**Example.**  $\mathbb{Q}(\pi) \cong \mathbb{Q}(e)$  because  $\pi$  and  $e$  are both transcendental.

We now return to the big question: is  $\mathbb{R}$  **definable** in  $\mathbb{C} = (\mathbb{C}, 0, 1, +, \cdot, -)$ ? If  $f: \mathbb{Q}(a) \rightarrow \mathbb{Q}(b)$  such that  $a \mapsto b$ , then there is an **automorphism**  $\hat{f}: \mathbb{C} \rightarrow \mathbb{C}$  such that  $a \mapsto b$ , i.e.,  $\hat{f}$  extends  $f$ . In other words, we need to find such an  $f$  with  $a \in \mathbb{R}$  and  $b \notin \mathbb{R}$ .

**Example.**  $a = \pi$ ,  $b = i\pi$  are both transcendental.

**Example.**  $a$  is a real  $\sqrt[4]{2}$ ,  $b$  is a complex  $\sqrt[4]{2}$ .

The above two examples show that  $\mathbb{R}$  is not  **$\emptyset$ -definable** in  $\mathbb{C}$ . In fact,  $\mathbb{R}$  is not **definable** over any  $\bar{a}$  because there are elements of  $\mathbb{R}$  and  $\mathbb{C} \setminus \mathbb{R}$  transcendental over any  $\bar{a}$ .

**Intuition.** There are so many  $a, b$  such that given any  $\bar{a}$ , we can still find a pair that works.

## 1.3 Completeness and Compactness

In this section, we're going to formalize **proofs**.

### 1.3.1 Proofs

There are all sorts of different proof systems, and the one we use is the so-called Hilbert-style deductive system. Before that, we first see some common notions.

**Notation (Schema).** A *schema* is written in symbols for **formulas**, variables, etc.

**Example.**  $\varphi \rightarrow (\psi \rightarrow \varphi)$  is a **schema**, i.e., an infinite set with all possible choices of  $\varphi$  and  $\psi$ .

Specifically, every **logical axiom** is written in **schema**, meaning that any instance of a symbol for a **formula**, e.g.,  $\varphi$ , can be replaced by any **formula**.

**Definition 1.3.1 (Generalization).** A **formula**  $\varphi$  is a *generalization* of a **formula**  $\psi$  if  $\varphi$  is  $\forall x_1 \dots \forall x_n \psi$  where  $x_1, \dots, x_n$  are variables.

**Notation (Hypothesis).** *Hypotheses* are **formulas** that we may assume in a **proof**.



**Definition 1.3.2 (Proof).** A *proof* is a sequence of **formulas**  $\{\varphi_i\}_{i=1}^n$  such that  $\varphi_n$  is the conclusion, and each **formula** is either an **axiom** or is obtained from the previous **formulas** by a **rule of inference**.

Moreover, for a **proof** based on a set of **hypotheses**  $\Gamma$ , then in addition to a **logical axiom**, we can assert a **formula**  $\varphi \in \Gamma$ . If we prove  $\psi$  using  $\Gamma$  as **hypotheses**, we write  $\Gamma \vdash \psi$ .

**Definition 1.3.3 (Valid).** If we **prove**  $\psi$  without **hypotheses**, we write  $\vdash \psi$  and say  $\psi$  is *valid*.

**Definition 1.3.4 (Logical axioms).** The *logical axioms* are the following **formulas** written in **schema**, as well as all of their **generalizations**:

**Definition 1.3.5 (Propositional axioms).** The *propositional axioms* are

- (A1)  $\varphi \rightarrow (\psi \rightarrow \varphi)$ .
- (A2)  $(\varphi \rightarrow (\psi \rightarrow \theta)) \rightarrow ((\varphi \rightarrow \psi) \rightarrow (\varphi \rightarrow \theta))$ .
- (A3)  $(\neg\varphi \rightarrow \neg\psi) \rightarrow ((\neg\varphi \rightarrow \psi) \rightarrow \varphi)$ .
- (A4)  $\forall x \varphi(x, \dots) \rightarrow \varphi(t, \dots)$  where  $t$  is any **term**.
- (A5)  $[\forall x (\varphi \rightarrow \psi)] \rightarrow [(\forall x \varphi) \rightarrow (\forall x \psi)]$ .
- (A6)  $\varphi \rightarrow \forall x \varphi$ , where  $x$  is not **free** in  $\varphi$ .

**Definition 1.3.6 (Axioms for equality).** The *axioms for equality* is

- (A7) for any **terms**  $t, u, v, \dots$ , function symbols  $f$ , and relation symbols  $R$ ,
  - (a)  $t = t$ .
  - (b)  $t = u \rightarrow u = t$ .
  - (c)  $(t = u \wedge u = v) \rightarrow (t = v)$ .
  - (d)  $(u_1 = t_1 \wedge \dots \wedge u_{n_f} = t_{n_f}) \rightarrow f(u_1, \dots, u_{n_f}) = f(t_1, \dots, t_{n_f})$ .
  - (e)  $(u_1 = t_1 \wedge \dots \wedge u_{n_R} = t_{n_R}) \rightarrow (R(u_1, \dots, u_{n_R}) \leftrightarrow R(t_1, \dots, t_{n_R}))$ .

**Definition 1.3.7 (Rule of inference).** From  $\varphi$  and  $\varphi \rightarrow \psi$ , deduces  $\psi$ .<sup>a</sup>

<sup>a</sup>This is called **modus ponens**.

These **formulas** might have **free variables**.

**Example.** A **proof** from calculus of a limit, e.g.,  $\forall \epsilon \exists \delta \dots$ . And we start by stating

1. let  $\epsilon > 0$ ,
2. choose  $\delta = \epsilon$ ,
- $\vdots$
- $n$ .  $|f(x) - f(y)| < \epsilon$ .

We should interpret **free variables** as anything.

**As previously seen** (Propositional logic).  $(p \wedge q) \vee (r \wedge \neg q)$ .

**Remark.** We can check whether the **propositional axioms** are **true** with a truth table.

**Definition 1.3.8** (Propositional tautology). A *propositional tautology* is a boolean combination  $\vee, \wedge, \neg$  of **formulas**  $\varphi_1, \dots, \varphi_n$  which is **true** via a truth table assigning true or false to each of  $\varphi_1, \dots, \varphi_n$ .

So instead of using **propositional axioms**, we could instead allow as **logical axioms** any **propositional tautology**. To prove **completeness**, we will need 5 **propositional tautologies**. We will **prove** some of these, but take others on faith.

**Remark.** **Propositional axioms** are enough to **prove** all **propositional tautologies**.

**Notation.** We write  $\Gamma \vdash_{\mathcal{L}} \varphi$  if there is a **proof** of  $\varphi$  from  $\Gamma$  in the **language**  $\mathcal{L}$ .

**Note.** Passing to a larger **language** will not let you **prove** more, so we can just write  $\vdash$ .

## Lecture 6: Soundness Theorem

To see why **propositional axioms** are enough to **prove** all **propositional tautologies**, we see one example.

24 Jan. 14:30

**Problem.** **Prove**  $\varphi \rightarrow \varphi$  using **propositional axioms**.

**Answer.** We see that

1.  $\varphi \rightarrow ((\psi \rightarrow \varphi) \rightarrow \varphi)$  from (A1), where  $\psi$  is any **formula** (possibly  $\psi = \varphi$ ).
2.  $[\varphi \rightarrow ((\psi \rightarrow \varphi) \rightarrow \varphi)] \rightarrow [(\varphi \rightarrow (\psi \rightarrow \varphi)) \rightarrow (\varphi \rightarrow \varphi)]$  from (A2).
3.  $(\varphi \rightarrow (\psi \rightarrow \varphi)) \rightarrow (\varphi \rightarrow \varphi)$  from (MP) and the two above.
4.  $\varphi \rightarrow (\psi \rightarrow \varphi)$  from (A1).
5.  $\varphi \rightarrow \varphi$  from (MP) and the two above.

⊛

In general, we can **prove**

- |   |  |
|---|--|
| (a) $\varphi \rightarrow \varphi$ ;         | (d) $(\varphi \rightarrow \psi) \rightarrow ((\neg\varphi \rightarrow \psi) \rightarrow \psi)$ ; |
| (b) $\varphi \rightarrow \neg\neg\varphi$ ; |  |
| (c) $\neg\neg\varphi \rightarrow \varphi$ ; | (e) $\varphi \rightarrow (\psi \rightarrow (\varphi \rightarrow \psi))$ ,                        |

and so on.

**Note.** As we said, we may replace **propositional axioms** by every **propositional tautologies**.

Some **proof** system also have a second rule about universal quantifiers, but in our system, we have built this into the axioms. We can prove, as a theorem, what the other proof systems take as a rule.

**Theorem 1.3.1.** If  $\Gamma \vdash \varphi$ , and  $x$  does not occur **freely** in  $\Gamma$ , then  $\Gamma \vdash \forall x \varphi$ .

**Proof.** Fix  $\Gamma$  and  $x$ , we use *induction on proofs*. Consider the set  $\{\varphi \mid \Gamma \vdash \forall x \varphi\}$ , we will show that this set contains all the **logical axioms**, **formulas** from  $\Gamma$ , and is closed under **modus ponens**.<sup>a</sup>

- (a) If  $\varphi$  is a **logical axiom**, so is its **generalization**  $\forall x \varphi$ , so  $\Gamma \vdash \forall x \varphi$ .
- (b) If  $\varphi \in \Gamma$ , then  $x$  is not **free** in  $\varphi$ , so from (A6),  $\varphi \rightarrow \forall x \varphi$ , and from (MP),  $\forall x \varphi$ . The above

are based on  $\Gamma$ , hence  $\Gamma \vdash \forall x \varphi$ .

(c) Suppose  $\Gamma \vdash \forall x \varphi$  and  $\Gamma \vdash \forall x (\varphi \rightarrow \psi)$ , we want to show that  $\Gamma \vdash \forall x \psi$ .

1. By (A5),  $\forall x (\varphi \rightarrow \psi) \rightarrow (\forall x \varphi \rightarrow \forall x \psi)$ ,  $\Gamma$  proves this.
2. By (MP),  $\Gamma \vdash \forall x \varphi \rightarrow \forall x \psi$ .
3. By (MP) again,  $\Gamma \vdash \forall x \psi$ .

■

<sup>a</sup>Thus, if  $\Gamma \vdash \theta$ , then  $\theta \in \{\varphi \mid \Gamma \vdash \forall x \varphi\}$ .

**Corollary 1.3.1.** If  $\vdash \varphi$ , then  $\vdash \forall x \varphi$ . So the generalization of anything valid is also valid.

We now ask a critical question: is our proof system a good one?

### 1.3.2 Soundness Theorem

The first thing we should check is whether our proofs are sound.

**Definition 1.3.9 (Sound).** A proof system is *sound* if any provable sentence  $\varphi$  is true.

The idea is that if an  $\mathcal{L}$ -sentence  $\varphi$  is provable, then it is true in all  $\mathcal{L}$ -structures, i.e., every thing we prove should be true, in other words, we can't prove wrong things.

**Lemma 1.3.1 (Soundness).** If  $\Gamma$  is a set of  $\mathcal{L}$ -sentences and  $\varphi$  is a sentence, and  $\Gamma \vdash_{\mathcal{L}} \varphi$ , then  $\Gamma \models \varphi$ .

**Proof.** Suppose that  $\Gamma \vdash \varphi$ , let  $\psi_1, \psi_2, \dots, \psi_n = \varphi$  be such a proof.<sup>a</sup> Let  $\bar{x} = (x_1, \dots, x_m)$  be the free variable that appears in the  $\psi_i$ . Let  $\mathcal{M}$  be an  $\mathcal{L}$ -structure,  $\mathcal{M} \models \Gamma$ . To show  $\mathcal{M} \models \varphi$ , we show that by induction on  $i$ , for all  $\bar{a} \in M^m$ ,  $\mathcal{M} \models \psi_i(\bar{a})$ . For  $\psi_i$ , we have three cases.

- (a) If  $\psi_i \in \Gamma$ , then  $\mathcal{M} \models \Gamma$  so  $\mathcal{M} \models \psi_i$ .
- (b) If  $\psi_i$  is a (generalization of) a logical axiom, then we can check that  $\mathcal{M} \models \psi_i(\bar{a})$ . For example, if  $\psi_i$  is (A1),  $\theta \rightarrow (\gamma \rightarrow \theta)$ , it's easy to check that

$$\mathcal{M} \models \theta(\bar{a}) \rightarrow (\gamma(\bar{a}) \rightarrow \theta(\bar{a})).$$

- (c) If there are  $j, k < i$  such that  $\psi_k$  is  $\psi_j \rightarrow \psi_i$ , from inductive hypothesis, for all  $\bar{a}$ ,  $\mathcal{M} \models \psi_j(\bar{a}), \mathcal{M} \models \psi_k(\bar{a})$ , then  $\mathcal{M} \models \psi_j(\bar{a}) \rightarrow \psi_i(\bar{a})$ . Checking our definition of truth, we get  $\mathcal{M} \models \psi_i(\bar{a})$ .

■

<sup>a</sup>Some  $\psi_i$  might be formulas, but  $\varphi$  should be a sentence.

There are remarks to make about some obvious properties of  $\vdash_{\mathcal{L}}$ .

**Remark.** If  $\varphi \in \Gamma$ , then  $\Gamma \vdash \varphi$ .

**Remark.** If  $\Delta \subseteq \Gamma$ , and  $\Delta \vdash \varphi$ , then  $\Gamma \vdash \varphi$ .

**Remark.** If  $\Gamma \vdash_{\mathcal{L}} \varphi$ , and  $\mathcal{L}^+ \supseteq \mathcal{L}$ , then  $\Gamma \vdash_{\mathcal{L}^+} \varphi$ .

**Remark.** If  $\Gamma \vdash \varphi$ , then there is a finite  $\Delta \subseteq \Gamma$  such that  $\Delta \vdash \varphi$ .

We can prove the following.

**Theorem 1.3.2 (Deduction theorem).** For any set of formulas  $\Gamma$ , formulas  $\theta$  and  $\psi$ ,

$$\Gamma \cup \{\theta\} \vdash \psi \Leftrightarrow \Gamma \vdash \theta \rightarrow \psi.$$

**Proof.** The backward direction is easier. Suppose  $\Gamma \vdash \theta \rightarrow \psi$ , then  $\Gamma \cup \{\theta\} \vdash \psi$  since we can have a proof like:

1.  $\theta$
- $\vdots$  (the proof of  $\Gamma \vdash \theta \rightarrow \psi$ )
- $n$ .  $\theta \rightarrow \psi$
- $n + 1$ .  $\psi$ .

Now, suppose that  $\Gamma \cup \{\theta\} \vdash \psi$ , then there is a proof  $\psi_1, \dots, \psi_n = \psi$  from  $\Gamma \cup \{\theta\}$ . We argue inductively that  $\Gamma \vdash \theta \rightarrow \psi_i$ . For  $i$ , we have three cases.

- (a) If  $\psi_i \in \Gamma$  or it is a logical axiom. By (A1),  $\psi_i \rightarrow (\theta \rightarrow \psi_i)$ , so  $\Gamma \vdash \theta \rightarrow \psi_i$ .
- (b) If  $\psi_i = \theta$ . Then  $\Gamma \vdash \theta \rightarrow \theta$  by (A1) and (A2) from here, hence  $\Gamma \vdash \theta \rightarrow \psi_i$ .
- (c) If  $\psi_i$  follows from  $\psi_j, \psi_k = \psi_j \rightarrow \psi_i$ , using (MP) with  $j, k < i$ .
  1. From the induction hypothesis,  $\Gamma \vdash \theta \rightarrow \psi_j$  and  $\Gamma \vdash \theta \rightarrow (\psi_j \rightarrow \psi_i)$ .
  2. By (A2),  $\Gamma \vdash [\theta \rightarrow (\psi_j \rightarrow \psi_i)] \rightarrow [(\theta \rightarrow \psi_j) \rightarrow (\theta \rightarrow \psi_i)]$ .
  3. By (MP),  $\Gamma \vdash (\theta \rightarrow \psi_j) \rightarrow (\theta \rightarrow \psi_i)$ .
  4. By (MP),  $\Gamma \vdash \theta \rightarrow \psi_i$ .

■

## Lecture 7: Soundness, Completeness, and Compactness

**Proposition 1.3.1 (Contraposition).** If  $\Gamma \cup \{\varphi\} \vdash \neg\psi$ , then  $\Gamma \cup \{\psi\} \vdash \neg\varphi$ .

26 Jan. 14:30

**Proof.** Suppose  $\Gamma \cup \{\varphi\} \vdash \neg\psi$ , by the deduction theorem says that  $\Gamma \vdash \varphi \rightarrow \neg\psi$ . From (A1), (A2), and (A3), we can prove  $(\varphi \rightarrow \neg\psi) \rightarrow (\psi \rightarrow \neg\varphi)$ . By (MP),  $\Gamma \vdash \psi \rightarrow \neg\varphi$ , then from the deduction theorem,  $\Gamma \cup \{\psi\} \vdash \neg\varphi$ . ■

Now we introduce an important notion.

**Definition 1.3.10 (Consistent).** A theory  $T$  is *consistent* if for all  $\varphi$ , it is not the case that  $T \vdash \varphi$  and  $T \vdash \neg\varphi$ .

**Definition 1.3.11 (Inconsistent).** If a theory  $T$  is not consistent, then it's *inconsistent*.

We could make the same definition for a set of formulas.

**Proposition 1.3.2 (Proof by contradiction).** If  $\Gamma \cup \{\varphi\}$  is inconsistent, then  $\Gamma \vdash \neg\varphi$ .

**Proof.** There is  $\psi$  such that  $\Gamma \cup \{\varphi\} \vdash \psi$  and  $\Gamma \cup \{\varphi\} \vdash \neg\psi$ , so  $\Gamma \vdash \varphi \rightarrow \psi$  and  $\Gamma \vdash \varphi \rightarrow \neg\psi$  by the deduction theorem. Using (A1), (A2), and (A3), we can prove that

$$(\varphi \rightarrow \psi) \rightarrow ((\varphi \rightarrow \neg\psi) \rightarrow \neg\varphi).$$

By (MP),  $\Gamma \vdash (\varphi \rightarrow \neg\psi) \rightarrow \neg\varphi$ , and by (MP) again, we have  $\Gamma \vdash \neg\varphi$ . ■

**Proposition 1.3.3.** If a theory  $T$  is consistent, and  $\varphi$  is a sentence, then either  $T \cup \{\varphi\}$  or  $T \cup \{\neg\varphi\}$  is consistent.

**Proof.** If they were both inconsistent,  $T \vdash \neg\varphi$  and  $T \vdash \neg\neg\varphi$ , so  $T$  would be inconsistent  $\nmid$  ■

**Note.** The above is also true for formulas.

**Remark.** If  $T$  is inconsistent, then  $T \vdash \varphi$  for any  $\varphi$ .

**Proof.** If  $T$  is inconsistent, then  $T \cup \{\neg\varphi\}$  is inconsistent for all  $\varphi$ . Hence, from proof by contradiction,  $T \vdash \neg\neg\varphi$  for all  $\varphi$ , which is just  $T \vdash \varphi$ .  $\circledast$

**Definition 1.3.12 (Maximal).** A theory  $T$  is maximal if it is consistent and for all sentences  $\varphi$ , either  $\varphi \in T$  or  $\neg\varphi \in T$ .

In particular, if  $T \vdash \varphi$ , then  $\varphi \in T$ .

**Intuition.** Basically, a maximal consistent theory has opinion on everything.

Now, we want to see that given a consistent theory, whether we can extend it to a maximal one. To do this, we need the following.

**Definition.** Let  $(P, \leq)$  be a partially ordered set.

**Definition 1.3.13 (Chain).** A chain is a set  $C \subseteq P$  such that for every  $p, q \in C$ , either  $p \leq q$  or  $q \leq p$ .

**Definition 1.3.14 (Upper bound).** If  $X \subseteq P$  is a set, an upper bound for  $X$  is an element  $p \in P$  such that  $p \geq q$  for all  $q \in X$ .

**Definition 1.3.15 (Maximal).** An element  $p \in P$  is maximal if there is no  $q \in P$  with  $q > p$ .

**Note.** Note that a maximal element might not be greater than everything, there is just nothing greater than it.

**Theorem 1.3.3 (Zorn's lemma).** Let  $(P, \leq)$  be a partially ordered set. If every non-empty chain in  $P$  has an upper bound, then  $P$  has a maximal element.

**Theorem 1.3.4.** Any consistent theory  $T$  can be extended to a maximal consistent theory  $T' \supseteq T$ .

**Proof.** We first consider the case that  $T$  is countable by considering  $\mathcal{L}$  is countable since if  $\mathcal{L}$  is countable, then there are only countable many formulas since there are only countable many formulas of each length.

**Claim.** The result holds for  $\mathcal{L}$  being countable.

**Proof.** Firstly, list out all sentences  $\varphi_1, \varphi_2, \dots$ , start with  $T_0 = T$ . Given  $T_i$  consistent, one of  $T_i \cup \{\varphi_i\}$  or  $T_i \cup \{\neg\varphi_i\}$  is consistent from Proposition 1.3.3. Let  $T_{i+1}$  be one of these that is consistent. Let  $T^* = \bigcup_i T_i$ , which is maximal, and we now show that  $T^*$  is consistent.

Suppose not, then  $T^* \vdash \theta$  and  $T^* \vdash \neg\theta$  for some  $\theta$ . In this case, there is some  $T_i$  such that  $T_i \vdash \theta$  and  $T_i \vdash \neg\theta$  because proofs are finite, with  $T_i$  being consistent, a contradiction  $\nmid$   $\circledast$

**Claim.** The result holds for arbitrary  $\mathcal{L}$ .

**Proof.** For arbitrary  $\mathcal{L}$ , let  $(P, \leq)$  be the set of **consistent theories** extending  $T_i$  ordered by inclusion. Let  $C$  be a non-empty **chain**, and let  $T^* = \bigcup_{T' \in C} T' \supseteq T$ .

We see that  $T^*$  is **consistent** because if  $T^* \vdash \theta$  and  $T^* \vdash \neg\theta$ , there are finitely many **formulas** used in those **proofs**, from, say,  $T_1, \dots, T_n \in C$ . Because  $C$  is a **chain**, by reordering, we may assume that  $T_1 \subseteq \dots \subseteq T_n$ . So  $T_n \vdash \theta$  and  $T_n \vdash \neg\theta$ , contradicting the **consistency** of  $T_n$ , so  $T^*$  is **consistent**, i.e.,  $T^* \in P$ . Furthermore,  $T^*$  is an **upper bound** on  $C$ ,<sup>a</sup> so  $(P, \leq)$  has a **maximal consistent theory**  $T^* \supseteq T$  from **Zorn's lemma**.

If  $T^*$  is not **maximal**, then there is  $\varphi$  where  $\varphi \notin T^*$ ,  $\neg\varphi \notin T^*$ . From **Proposition 1.3.3**, one of  $T^* \cup \{\varphi\}$  or  $T^* \cup \{\neg\varphi\}$  is **consistent**, hence in  $P$ , contradicting to  $T^*$  being **maximal**  $\nmid \otimes$

<sup>a</sup>Note that  $C$  is arbitrary.

**Remark.** We can do that same proof for any  $\mathcal{L}$  using **transfinite recursion** for the uncountable case.

Motivated by **Lemma 1.3.1** and **Theorem 1.3.4**, we close this section with the following.

**Theorem 1.3.5 (Soundness).** Let  $T$  be a **theory** and  $\varphi$  be a **sentence**.

- (a) If  $T \vdash \varphi$ , then  $T \models \varphi$ .
- (b) If  $T$  is **satisfiable**, then it is **consistent**.

**Proof.** (a) is exactly **Theorem 1.3.5**. For (b), let  $\mathcal{M} \models T$ , suppose that  $T$  was **inconsistent**, then  $T \vdash \varphi$  and  $T \vdash \neg\varphi$  for some  $\varphi$ . By (a),  $T \models \varphi$  and  $T \models \neg\varphi$ , so  $\mathcal{M} \models \varphi$  and  $\mathcal{M} \models \neg\varphi$ . But  $\mathcal{M} \models \neg\varphi$  means  $\mathcal{M} \not\models \varphi$ , so this is a contradiction, hence  $T$  is **consistent**. ■

### 1.3.3 Completeness and Compactness Theorems

After knowing our **proof** system is **sound**, we now ask the converse: is our **proof** system **complete**?

**Definition 1.3.16 (Complete).** A **proof** system is *complete* if any **true sentence**  $\varphi$  is **provable**.

And indeed, this is the case.

**Theorem 1.3.6 (Completeness).** Let  $T$  be a **theory** and  $\varphi$  be a **sentence**.

- (a) If  $T \models \varphi$ , then  $T \vdash \varphi$ .
- (b) If  $T$  is **consistent**, then it is **satisfiable**.

(b) implies (a) is easy to see. Suppose that  $T \models \varphi$ , so  $T \cup \{\neg\varphi\}$  is not **satisfiable**. By (b),  $T \cup \{\neg\varphi\}$  is **inconsistent**. By **proof by contradiction**,  $T \vdash \varphi$ . One important consequence of the **completeness theorem** is the **compactness theorem**.

**Theorem 1.3.7 (Compactness).** Let  $T$  be a **theory** and  $\varphi$  be a **sentence**.

- (a) If  $T \models \varphi$ , then there is a finite  $T_0 \subseteq T$  such that  $T_0 \models \varphi$ .
- (b)  $T$  is **satisfiable** if and only if every finite subset of  $T$  is **satisfiable**.

**Proof.** Consider the following.

- (a\*) If  $T \vdash \varphi$ , then there is a finite  $T_0 \subseteq T$  such that  $T_0 \vdash \varphi$ .
- (b\*) If  $T$  is **consistent** if and only if every finite subset of  $T$  is **consistent**.

We see that (a\*) and (b\*) are true because proofs are finite, and soundness and completeness translate directly between (a) and (a\*) (and (b) and (b\*)). ■

**Remark.** The compactness theorem does have something to do with topological compactness; consider the topological space of complete satisfiable theories, with the basic open sets being the sets

$$U_\varphi := \{T : T \models \varphi\},$$

then this topological space is compact.

Let's see one cool example using compactness.

**Example.** Let  $\mathcal{L} = \{0, 1, +, \cdot, -, <\}$ , and  $\mathcal{L}^* = \mathcal{L} \cup \{c\}$ , where  $c$  is a new constant symbol. Let

$$T = \text{Th}_{\mathcal{L}}(\mathbb{N}) \cup \{c > \underline{n} \mid n \in \mathbb{N}\},$$

then  $T$  is finitely satisfiable.

**Proof.** Given  $T_0 \subseteq T$  finite,  $T_0 \subseteq \text{Th}_{\mathcal{L}}(\mathbb{N}) \cup \{c > \underline{n}, \dots, c > \underline{n}_\ell\}$ , and may assume they are equal and show that  $T_0$  is satisfiable. Let  $\mathcal{N}$  be the  $\mathcal{L} \cup \{c\}$ -structure which is the expansion of the  $\mathcal{L}$ -structure  $\mathbb{N}$ , with

$$c^{\mathcal{N}} = 1 + \max(n_1, \dots, n_\ell),$$

then  $\mathcal{N} \models T_0$ , and  $T_0$  is satisfiable. By compactness,  $T$  is satisfiable, say  $\mathcal{A} \models T$ . Then  $\mathcal{A} \equiv \mathbb{N}$  and  $\mathcal{A}$  contains an element  $c^{\mathcal{A}}$  bigger than  $1, 1+1, 1+1+1, \dots$ , but  $\mathcal{A} \not\equiv \mathbb{N}$ , so  $\mathcal{A}$  is a non-standard model of arithmetic. ⊛

We now start a long journey toward proving completeness theorem, specifically (b).

## Lecture 8: Henkin Constants

### 1.3.4 Henkin Construction

31 Jan. 14:30

To prove Theorem 1.3.6 (b), we need an additional definition and a technical lemma due to Henkin.

**Definition 1.3.17** (Henkin constant). An  $\mathcal{L}^*$ -theory  $T^*$  has *Henkin constants* if for each formula  $\varphi(x)$  with one free variable, there is a constant symbol  $c \in \mathcal{L}^*$  such that

$$(\exists x \varphi(x)) \rightarrow \varphi(c) \text{ is in } T^*.$$

We see that the above is equivalent to

$$(\neg \forall x \varphi(x)) \rightarrow \neg \varphi(c) \text{ is in } T^*,$$

and we will use this version ( $\forall$ ) and view  $\exists$  being a shorthand for  $\neg \forall \neg$ ; also, we will use  $\rightarrow$  and  $\neg$  as primitive, and  $\wedge, \vee$  are shorthand.

**Lemma 1.3.2.** If  $\Gamma \vdash \varphi(c)$ , and  $c$  does not occur in  $\Gamma$  or in  $\varphi(x)$ , then there is a variable  $y$  not appearing in  $\varphi(x)$ , such that  $\Gamma \vdash \forall y \varphi(y)$ . Moreover, there is a proof of  $\forall y \varphi(y)$  in which  $c$  does not appear.

**Proof.** Let  $\alpha_1(c), \dots, \alpha_n(c) = \varphi(c)$  be a proof of  $\varphi(c)$  from  $\Gamma$ , and let  $y$  be a variable not appearing in this proof. We claim that  $\alpha_1(y), \dots, \alpha_n(y) = \varphi(y)$  is still a valid proof of  $\varphi(y)$ . There are three cases to consider (for each  $i = 1, \dots, n$ ):

- (a) If  $\alpha_i(c)$  is in  $\Gamma$ , then  $c$  does not actually occur in  $\alpha_i(c)$  because it does not appear in  $\Gamma$ . So  $\alpha_i(y)$  is the same as  $\alpha_i(c)$ , hence in  $\Gamma$ .
- (b) If  $\alpha_i(c)$  is a logical axiom, then  $\alpha_i(y)$  is a logical axiom as well. For most of these it is easy to check, but for (A6), i.e.,  $\varphi \rightarrow \forall x \varphi$  if  $x$  is not free in  $\varphi$ , there is a little more. But  $y$  did

not appear in  $\alpha_i(c)$ , so  $y \neq x$ , and substituting  $y$  for  $c$  will not stop  $x$  from being not **free**.

- (c) If  $\alpha_i(c)$  follows by **(MP)** from  $\alpha_j(c)$  and  $\alpha_k(c) = \alpha_j(c) \rightarrow \alpha_i(c)$  for  $j, k < i$ , then  $\alpha_i(y)$  follows by **(MP)** from  $\alpha_j(y)$  and  $\alpha_k(y) = \alpha_j(y) \rightarrow \alpha_i(y)$ .

So  $\Gamma \vdash \varphi(y)$  and the **proof** does not involve  $c$ . Let  $\Phi \subseteq \Gamma$  be the subset of  $\Gamma$  that was used in the **proof**, so  $y$  does not appear in  $\Phi$ , hence  $\Phi \vdash \varphi(y)$  and  $\Phi \vdash \forall y \varphi(y)$ , so  $\Gamma \vdash \forall y \varphi(y)$ . ■

So **Lemma 1.3.2** implies that we have  $\Gamma \vdash \varphi(y)$  and the **proof** does not involve  $c$ . And sometimes, we want to be able to choose the variable  $y$  from above.

**Corollary 1.3.2.** If  $\Gamma \vdash \varphi(c)$ , and  $c$  does not occur in  $\Gamma$  or in  $\varphi(x)$ , then  $\Gamma \vdash \forall x \varphi(x)$ . Moreover, there is a **proof** of  $\forall x \varphi(x)$  not involving  $c$ .<sup>a</sup>

<sup>a</sup>Here,  $x$  is any variable that does not appear in  $\varphi(c)$ .

**Proof.** We know that for some  $y$ ,  $\Gamma \vdash \forall y \varphi(y)$ , **(A4)** says  $\forall y \varphi(y) \rightarrow \varphi(x)$ . So  $\forall y \varphi(y) \vdash \varphi(x)$  since  $x$  does not appear in  $\forall y \varphi(y)$ ,  $\forall y \varphi(y) \vdash \forall x \varphi(x)$ . ■

**Note.**  $x$  might appear in  $\Gamma$ .

**Theorem 1.3.8.** Let  $T$  be a **consistent  $\mathcal{L}$ -theory**. There is a **language  $\mathcal{L}^* \supseteq \mathcal{L}$**  and  $T^* \supseteq T$  a **consistent  $\mathcal{L}^*$ -theory** such that  $T^*$  has **Henkin constants**. We can choose  $\mathcal{L}^*$  such that  $|\mathcal{L}^*| = |\mathcal{L}| + \aleph_0$ , and all new symbols in  $\mathcal{L}^*$  are constants.

**Proof.** Let  $\mathcal{L}_0 = \mathcal{L}$  and  $T_0 = T$ . Let  $\mathcal{L}_1$  be the expansion of  $\mathcal{L}_0$  by adding a new constant symbol  $c_\varphi$  for each  **$\mathcal{L}_0$ -formula  $\varphi$**  w.r.t. the **Henkin** construction. First, we show that after this procedure,  $T_0$  is still a **consistent  $\mathcal{L}_1$ -theory**.

**Remark.** Technically,  $\vdash$  is really  $\vdash_{\mathcal{L}}$ , so this is a key step for seeing that it does not matter.

**Claim.**  $T_0$  is still a **consistent  $\mathcal{L}_1$ -theory** after the expansion of  $\mathcal{L}_0$ .

**Proof.** If not, there is a **proof** of a **contradiction** from  $T_0$ , and which uses only finitely many of the new constants symbols. By **Corollary 1.3.2**, we can replace these constants one-by-one by variables, e.g., if the original **contradiction** was  $\varphi(c_1, \dots, c_n)$  and  $\neg \varphi(c_1, \dots, c_n)$ , then  $T_0$  proves  $\forall x_1, \dots, \forall x_n \varphi(x_1, \dots, x_n)$  and  $\forall x_1, \dots, \forall x_n \neg \varphi(x_1, \dots, x_n)$ . Moreover, these **proofs** take place in  $\mathcal{L}_0$ , so by **(A4)**,  $T_0 \vdash_{\mathcal{L}_0} \varphi(x_1, \dots, x_n)$ , and  $T_0 \vdash_{\mathcal{L}_0} \neg \varphi(x_1, \dots, x_n) \not\vdash$  ⊗

To construct  $T_1$  w.r.t. the **Henkin** construction, it's natural to consider the following: if  $\varphi$  is of the form  $\neg \forall x \psi(x)$ , then let

$$\theta_\varphi := (\neg \forall x \psi(x)) \rightarrow \neg \psi(c_\varphi), \text{ i.e., } (\exists x \neg \psi(x)) \rightarrow \neg \psi(c_\varphi),$$

otherwise, let  $\theta_\varphi := \forall x (x = x)$  (trivially **true**). Let  $\Theta = \{\theta_\varphi \mid \varphi \text{ an } \mathcal{L}_0\text{-formula}\}$ , and we let that  $T_1 = T_0 \cup \Theta$ . We claim that  $T_1$  is still **consistent**.

**Claim.**  $T_1 = T_0 \cup \Theta$  is a **consistent  $\mathcal{L}_1$ -language** after the expansion of  $\mathcal{L}_0$ .

**Proof.** If not, then there are  $\varphi_1, \dots, \varphi_{m+1}$  such that  $T_0 \cup \{\theta_{\varphi_1}, \dots, \theta_{\varphi_m}, \theta_{\varphi_{m+1}}\}$  is **inconsistent**. Taking  $m$  to be as small as possible,  $T_0 \cup \{\theta_{\varphi_i}\}_{i=1}^m$  is **consistent**, so  $T_0 \cup \{\theta_{\varphi_i}\}_{i=1}^m \vdash \neg \theta_{\varphi_{m+1}}$  with  $\varphi_{m+1}$  being of the form  $\neg \forall x \psi(x)$ ,  $\theta_{\varphi_{m+1}}$  is  $\neg \forall x \psi(x) \rightarrow \neg \psi(c_{\varphi_{m+1}})$ . By **(A1)**, **(A2)**, **(A3)**,

$$T_0 \cup \{\theta_{\varphi_1}, \dots, \theta_{\varphi_m}\} \vdash \neg \forall x \psi(x) \text{ and } T_0 \cup \{\theta_{\varphi_1}, \dots, \theta_{\varphi_m}\} \vdash \psi(c_{\varphi_{m+1}}).$$

Since  $c_{\varphi_{m+1}}$  does not appear in  $T_0 \cup \{\theta_{\varphi_i}\}_{i=1}^m$ , so  $T_0 \cup \{\theta_{\varphi_i}\}_{i=1}^m \vdash \forall x \psi(x)$ , i.e.,  $T_0 \cup \{\theta_{\varphi_i}\}_{i=1}^m$  is **inconsistent**, contradicting to the fact that  $m$  is the smallest choice  $\not\vdash$  <sup>a</sup> ⊗

<sup>a</sup>If  $m = 0$ , then we violate the **consistency** of  $T_0$ .



It might be that  $T_1$  does not have Henkin constants since there are new  $\mathcal{L}_1$ -formulas which are not  $\mathcal{L}_0$ -formulas. But we know that  $T_1$  does have Henkin constants for  $\mathcal{L}_0$ -formulas, hence we can repeat that process and keep fixing things. In general, given  $T_i$  and  $\mathcal{L}_i$ , define a  $T_{i+1}$  and  $\mathcal{L}_{i+1}$  in the above way. Since each  $T_i$  is consistent, so  $T^* = \bigcup T_i$  is an  $\mathcal{L}^* = \bigcup \mathcal{L}_i$ -theory. Note that  $T^*$  is consistent as a nested union of consistent theories, and  $T^*$  has Henkin constants because every  $\mathcal{L}^*$ -formula  $\varphi$  is an  $\mathcal{L}_i$ -formula for some  $i$ , and  $\theta_\varphi \in T_{i+1} \subseteq T^*$ .

**Intuition.** This is like “chasing its own tail,” which basically fixes new errors introduced every time and then takes the union in the end.

Finally, we want to show that  $|\mathcal{L}^*| = |\mathcal{L}| + \aleph_0$ . Given  $\mathcal{L}_i$ , we define  $\mathcal{L}_{i+1}$  to be  $\mathcal{L}_i$  plus new constants  $c_\varphi$  for  $\varphi$  on  $\mathcal{L}_i$ -formula. Then, we have

$$|\mathcal{L}_{i+1}| \leq |\mathcal{L}_i| + \underbrace{|\mathcal{L}_i|}_{\# \text{ of } \mathcal{L}_i\text{-formulas}} + \aleph_0 = |\mathcal{L}_i| + \aleph_0.$$

So for all  $i$ ,  $|\mathcal{L}_i| \leq |\mathcal{L}| + \aleph_0$ , and  $\mathcal{L}^* = \bigcup_i \mathcal{L}_i$  is a countable union, so  $|\mathcal{L}^*| \leq |\mathcal{L}| + \aleph_0$ , and in fact,  $|\mathcal{L}^*| = |\mathcal{L}| + \aleph_0$ . ■

After proving Theorem 1.3.8, we see that to prove Theorem 1.3.6 (b), we can proceed by:

1. extend  $T^*$  to a maximal theory  $T^{**}$ ;<sup>4</sup>
2. turn  $T^{**}$  into a model. The elements of the model are constant symbols from  $\mathcal{L}^*$ , modulo the equivalence relation  $c \sim d$  if  $c = d$  is in  $T^{**}$ , i.e.,  $T^{**} \vdash c = d$ .

Thankfully, the first step is easy from Theorem 1.3.4, so we just need to show the second step, and we're done.

## Lecture 9: Proving the Completeness Theorem

To finish the proof of Theorem 1.3.6 (b), we follow the plan mentioned last lecture, and prove the following. 2 Feb. 14:30

**Theorem 1.3.9.** If  $T$  is a maximal consistent  $\mathcal{L}$ -theory with Henkin constants, then  $T$  has a model.

**Proof.** The model we build is called a “canonical model.” Let  $\mathcal{C}$  be the set of constants in  $\mathcal{L}$ , and define an equivalence relation  $\sim$  on  $\mathcal{C}$  by  $c \sim d$  if and only if  $c = d$  is in  $T$ .

**Claim.** The relation  $\sim$  on  $\mathcal{C}$  defined by  $c \sim d \Leftrightarrow c = d \in T$  is an equivalence relation.

**Proof.** We check the axioms for being an equivalence relation.

- (a)  $c \sim c$  because  $c = c$  is in  $T$  by (A7) (a).<sup>a</sup>
- (b) If  $c \sim d$ , then  $c = d$  is in  $T$  so  $d = c$  is in  $T$  by (A7) (b), i.e.,  $d \sim c$ .
- (c) If  $c \sim d$  and  $d \sim e$ , then  $c = d$  and  $d = e \in T \Rightarrow c = e \in T$  by (A7) (c), so  $c \sim e$ .

⊗

<sup>a</sup>Otherwise,  $c \neq c$  is in  $T$  from the maximality, so  $T \vdash c \neq c$  with  $T \vdash c = c$ , so  $T$  would be inconsistent.

Let  $[c]$  be the equivalence class of  $c$ . Define an  $\mathcal{L}$ -structure  $\mathcal{M}$  with domain  $M = \mathcal{C} / \sim = \{[c] \mid c \in \mathcal{C}\}$ , with functions, relations, and constants defined as follows:

- (a)  $c^{\mathcal{M}} = [c]$ .
- (b)  $R^{\mathcal{M}}([c_1], \dots, [c_n])$  true if  $R(c_1, \dots, c_n)$  is in  $T$ . This is well-defined by (A7) (e).

<sup>4</sup>Which still has Henkin constants.

- (c)  $f^{\mathcal{M}}([c_1], \dots, [c_n]) = [d]$  if  $f(c_1, \dots, c_n) = d$  is in  $T$ . Such a  $d$  exists because  $\exists x f(c_1, \dots, c_n) = x$ , i.e.,  $\neg \forall x f(c_1, \dots, c_n) \neq x$ , is in  $T$ .<sup>b</sup> If this is in  $T$ , then there is a **Henkin constant**  $d$  with  $f(c_1, \dots, c_n) = d$  in  $T$ . To show that this is well-defined, from (A7) (d), i.e.,

$$(t_1 = u_1 \wedge \dots \wedge t_n = u_n) \rightarrow f(t_1, \dots, t_n) = f(u_1, \dots, u_n).$$

So if  $[c_1] = [d_1], \dots, [c_n] = [d_n]$ , then  $c_1 = d_1, \dots, c_n = d_n$  are in  $T$ . So  $f(c_1, \dots, c_n) = f(d_1, \dots, d_n)$  is in  $T$  by (A7) (d). If  $a$  and  $b$  are constants such that  $f(c_1, \dots, c_n) = a$  and  $f(d_1, \dots, d_n) = b$  are in  $T$ , so  $a = b$  is in  $T$  by (A7) (c), i.e., the transitivity of  $=$ .

Now we need to show that  $\mathcal{M} \models T$ , i.e., we claim that

$$\mathcal{M} \models \varphi([c_1], \dots, [c_n]) \Leftrightarrow \varphi(c_1, \dots, c_n) \text{ is in } T.$$

We prove this by induction on **terms** and then **formulas**.

1. **Terms:** Show that  $t^{\mathcal{M}}([c_1], \dots, [c_n]) = [d]$  if and only if  $t(c_1, \dots, c_n) = d$  is in  $T$ .

- (a) If  $t$  is a constant  $e$ ,  $t^{\mathcal{M}}([c_1], \dots, [c_n]) = e^{\mathcal{M}} = [e]$ , and

$$[e] = t^{\mathcal{M}}([c_1], \dots, [c_n]) = [d] \Leftrightarrow [e] = [d] \Leftrightarrow e = d \text{ is in } T.$$

- (b) If  $t$  is  $x_i$ ,  $t^{\mathcal{M}}([c_1], \dots, [c_n]) = [c_i]$ . This is equal to  $[d]$  if and only if  $c_i = d$  is in  $T$ .

- (c) Suppose that  $t(x_1, \dots, x_n) = f(s_1(x_1, \dots, x_n), \dots, s_m(x_1, \dots, x_n))$ . Let

$$[d_i] = s_i^{\mathcal{M}}([c_1], \dots, [c_n]),$$

by the inductive hypothesis,  $d_i = s_i(c_1, \dots, c_n)$  is in  $T$ . Let  $[e] = f^{\mathcal{M}}([d_1], \dots, [d_m]) = t^{\mathcal{M}}([c_1], \dots, [c_n])$ . By the definition of  $f$ ,  $e = f(d_1, \dots, d_m)$  is in  $T$ . By (A7) (d),

$$e = f(s_1(c_1, \dots, c_n), \dots, s_m(c_1, \dots, c_n))$$

is in  $T$ . This is the direction ( $\Rightarrow$ ).

Now suppose that  $t(c_1, \dots, c_n) = e'$  is in  $T$ . We want to show that  $[e] = [e']$ , i.e.,  $e = e'$  is in  $T$ . Since  $e = t(c_1, \dots, c_n)$  is in  $T$ , and  $e' = t(c_1, \dots, c_n)$  is in  $T$ . By (A7) (c),  $e = e'$  is in  $T$ , so  $[e'] = [e] = t^{\mathcal{M}}([c_1], \dots, [c_n])$ . This is ( $\Leftarrow$ ).

2. **Formulas:** Show that  $\mathcal{M} \models \varphi([c_1], \dots, [c_n])$  if and only if  $\varphi(c_1, \dots, c_n)$  is in  $T$ .<sup>c</sup>

- (a) If  $\varphi$  is  $s(x_1, \dots, x_n) = t(x_1, \dots, x_n)$ :

$$(\Rightarrow) \text{ If } \mathcal{M} \models s([c_1], \dots, [c_n]) = t([c_1], \dots, [c_n]),$$

$$s^{\mathcal{M}}([c_1], \dots, [c_n]) = t^{\mathcal{M}}([c_1], \dots, [c_n]).$$

Let  $[d]$  be this element equal to the above, so  $d = s(c_1, \dots, c_n)$  and  $d = t(c_1, \dots, c_n)$  are in  $T$  so  $\underbrace{s(c_1, \dots, c_n) = t(c_1, \dots, c_n)}_{\varphi(c_1, \dots, c_n)}$  is in  $T$  by (A7) (c).

$$(\Leftarrow) \text{ If } s(c_1, \dots, c_n) = t(c_1, \dots, c_n) \text{ is in } T, \text{ let}$$

$$[d] = s^{\mathcal{M}}([c_1], \dots, [c_n]) \text{ and } [e] = t^{\mathcal{M}}([c_1], \dots, [c_n]),$$

so  $d = s(c_1, \dots, c_n)$  and  $e = t(c_1, \dots, c_n)$  are in  $t$ , so  $d = e$  is in  $t$ , and  $[e] = [d]$ .

(b) If  $\varphi$  is  $R(t_1(\bar{x}), \dots, t_m(\bar{x}))$ : Let  $[d_i] = t_i^{\mathcal{M}}([c_1], \dots, [c_n])$ ,

$$\begin{array}{ccc}
 R^{\mathcal{M}}([d_1], \dots, [d_m]) \text{ is true} & \iff & R(d_1, \dots, d_m) \text{ is in } T \\
 \updownarrow & & \updownarrow \\
 R^{\mathcal{M}}(t_1^{\mathcal{M}}[\bar{c}], \dots, t_m^{\mathcal{M}}[\bar{c}]) \text{ is true} & & R(t_1(\bar{c}), \dots, t_m(\bar{c})) \text{ is in } T \\
 \updownarrow & & \\
 \mathcal{M} \models \varphi([c_1], \dots, [c_n]) & & 
 \end{array}$$

(c) If  $\varphi$  is  $\neg\psi$ : Then

$$\mathcal{M} \models \varphi(\bar{c}) \Leftrightarrow \mathcal{M} \not\models \psi(\bar{c}) \Leftrightarrow \psi(\bar{c}) \text{ is not in } T \Leftrightarrow \varphi(\bar{c}) \text{ is in } T$$

where the last  $\Leftrightarrow$  follows from the fact that  $T$  is **maximal** and **consistent**.

(d) If  $\varphi$  is  $\psi \rightarrow \theta$ :

- If  $\psi(\bar{c}) \rightarrow \theta(\bar{c})$  is in  $T$ : then if  $\psi(\bar{c})$  is in  $T$ , then  $\theta(\bar{c})$  is in  $T$  by **(MP)**. then by the induction hypotheses, if  $\mathcal{M} \models \psi(\bar{c})$ , then  $\mathcal{M} \models \theta(\bar{c})$ .
- If  $\mathcal{M} \models \psi(\bar{c}) \rightarrow \theta(\bar{c})$ : then either  $\mathcal{M} \models \theta(\bar{c})$  or  $\mathcal{M} \models \neg\psi(\bar{c})$ . So either
  - i.  $\theta(\bar{c})$  is in  $T$ : by **(A1)**,  $\theta(\bar{c}) \rightarrow (\psi(\bar{c}) \rightarrow \theta(\bar{c}))$ , so  $\psi(\bar{c}) \rightarrow \theta(\bar{c})$  is in  $T$ .
  - ii.  $\neg\psi(\bar{c})$  is in  $T$ :  $T \cup \{\psi(\bar{c})\}$  is now **inconsistent**, so  $T \cup \{\psi(\bar{c})\} \vdash \theta(\bar{c})$ . From the **deductive theorem**,  $T \vdash \psi(\bar{c}) \rightarrow \theta(\bar{c})$ . Because  $T$  is **maximal** and **consistent**,  $\psi(\bar{c}) \rightarrow \theta(\bar{c})$  is in  $T$ .

<sup>b</sup>Otherwise,  $\forall x f(c_1, \dots, c_n) \neq x$  is in  $T$ . By **(A4)**,  $f(c_1, \dots, c_n) \neq f(c_1, \dots, c_n)$  is in  $T$ , contradicts to **(A7) (a)**.

<sup>c</sup>In particular, for a **sentence**  $\varphi$ ,  $\mathcal{M} \models \varphi \Leftrightarrow \varphi$  is in  $T$ , and so  $\mathcal{M} \models T$ .

## Lecture 10: Introduction to Model Theory

Let's start by finishing the proof of **Theorem 1.3.9**.

7 Feb. 14:30

**Proof of Theorem 1.3.9 (Continued).** There's one final case left:

- (e) If  $\varphi$  is  $\forall x \psi(x, \bar{y})$ : Because  $T$  has **Henkin constants**, there is  $d$  such that  $\neg\forall x \psi(x, \bar{c}) \rightarrow \neg\psi(d, \bar{c})$  is in  $T$ .
- If  $\varphi(c_1, \dots, c_n)$  is not in  $T$ , i.e.,  $\forall x \psi(x, \bar{c})$  is in  $T$ , then since  $T$  is **maximal**,  $\neg\forall x \psi(x, \bar{c})$  is in  $T$ . So by **(MP)**,  $\neg\psi(d, \bar{c})$  is in  $T$ . So,  $\mathcal{M} \models \neg\psi([d], [\bar{c}])$  by induction hypotheses, hence  $\mathcal{M} \models \neg\forall x \psi(x, [\bar{c}])$ , i.e.,  $\mathcal{M} \not\models \varphi([\bar{c}])$ .
  - If  $\mathcal{M} \models \varphi([\bar{c}])$ , then  $\mathcal{M} \models \forall x \varphi(x, [\bar{c}])$ , so there is  $[e]$  such that  $\mathcal{M} \models \neg\psi([e], [\bar{c}])$ . Hence,  $\neg\psi(e, \bar{c})$  is in  $T$ . Suppose for a contradiction that  $\varphi(\bar{c})$ , i.e.,  $\forall x \psi(x, \bar{c})$  is in  $T$ , by **(A4)**,  $\forall x \psi(x, \bar{c}) \rightarrow \psi(e, \bar{c})$ , so  $\psi(e, \bar{c})$  is in  $T$  by **maximality** and by **consistency**. But then  $T$  is **inconsistent**, a contradiction  $\nmid$  Hence  $\varphi(\bar{c})$  is not in  $T$ .

Thus,  $\mathcal{M} \models T$ , so  $T$  is **satisfiable**, proving the theorem. ■

**Remark.** We see that when proving the above, when we talk about  $\mathcal{M}$ , the witness comes for free, while for  $T$ , we need **Henkin constants** for getting a witness.

Now, we can complete the proof of **completeness theorem** by putting everything together.

**Claim.** The **completeness theorem (b)** holds.

**Proof.** We see that

1. **Theorem 1.3.8:** There is a **consistent**  $T^* \supseteq T$  and  $\mathcal{L}^*$ -theory (with  $\mathcal{L}^* \supseteq \mathcal{L}$ ) and  $T^*$  has **Henkin constants**.

2. **Theorem 1.3.4:** There is a maximal consistent  $\mathcal{L}^*$ -theory  $T^{**} \supseteq T^*$ , where  $T^{**}$  still has Henkin constants.
3. **Theorem 1.3.9:**  $T^{**}$  has a model  $\mathcal{M}^*$  an  $\mathcal{L}^*$ -structure. Let  $\mathcal{M}$  be the reduct of  $\mathcal{M}^*$  to an  $\mathcal{L}$ -structure.

Hence,  $\mathcal{M} \models T$ . ⊛

**As previously seen** (Problem set 1). Let  $\mathcal{L}^* \supseteq \mathcal{L}$ . If  $\mathcal{M}^*$  is an  $\mathcal{L}^*$ -structure, then by ignoring the interpretation of the symbols in  $\mathcal{L}^* - \mathcal{L}$ , we get an  $\mathcal{L}$ -structure  $\mathcal{M}$ .

**Notation** (Reduct).  $\mathcal{M}$  is a *reduct* of  $\mathcal{M}^*$ .

**Notation** (Expansion).  $\mathcal{M}^*$  is an *expansion* of  $\mathcal{M}$ .

**Remark.** We see that  $\vdash$  and  $\models$  are the same.

### 1.3.5 Consequences of Completeness Theorem

Now, let's step back and look at the proof of the completeness theorem, and ask the following.

**Problem.** When we did the Henkin construction of  $\mathcal{M}^* \models T^{**}$ , how big was  $M$ ?

This can be answered by the following.

**Theorem 1.3.10.** If  $T$  is a satisfiable  $\mathcal{L}$ -theory, then it has a model of size at most  $|\mathcal{L}| + \aleph_0$ .

**Proof.** Since  $|M| \leq |\mathcal{L}^*|$  since  $\mathcal{M} = \mathcal{C} / \sim$ , and in step one,  $|\mathcal{L}^*| \leq |\mathcal{L}| + \aleph_0$ , so  $|M| \leq |\mathcal{L}| + \aleph_0$ . ■

**Example.** Let  $\mathcal{L} = \{f\}$ ,  $T$  says that  $f$  is injective but not surjective.

**Example.** Let  $\mathcal{L} = \{\leq\}$ ,  $T$  says that  $\leq$  is a linear order with no greatest element.

**Example.** Let  $\mathcal{L} = \emptyset$ ,  $T$  says that there are at least  $n$  elements for each  $n$ .

**As previously seen.**  $\vdash$  and  $\models$  are actually  $\vdash_{\mathcal{L}^a}$  and  $\models_{\mathcal{L}^b}$

<sup>a</sup>Proofs can only use  $\mathcal{L}$ -formulas.

<sup>b</sup>Only looking at  $\mathcal{L}$ .

**Remark.** Suppose  $\mathcal{L} \supseteq \mathcal{L}_0$ , and  $\Gamma$  a set of  $\mathcal{L}_0$ -sentences,  $\varphi$  on  $\mathcal{L}_0$ -sentence.

(a)  $\Gamma \models_{\mathcal{L}_0} \varphi \Leftrightarrow \Gamma \models_{\mathcal{L}_1} \varphi$ .

(b)  $\Gamma \vdash_{\mathcal{L}_0} \varphi \Leftrightarrow \Gamma \vdash_{\mathcal{L}_1} \varphi$ .

**Proof.** (a) and (b) are equivalent by the completeness theorem, and we prove (a).

Suppose  $\Gamma \models_{\mathcal{L}_0} \varphi$ . Suppose  $\mathcal{M}_1$  is an  $\mathcal{L}_1$ -structure such that  $\mathcal{M}_1 \models \Gamma$ . Let  $\mathcal{M}_0$  be the reduct of  $\mathcal{M}_1$  to  $\mathcal{L}_0$  and  $\mathcal{M}_0 \models \Gamma$ , so  $\mathcal{M}_0 \models \varphi$ , then  $\mathcal{M}_1 \models \varphi$ , thus  $\Gamma \models_{\mathcal{L}_1} \varphi$ .

Now, suppose  $\Gamma \models_{\mathcal{L}_1} \varphi$ . Suppose  $\mathcal{M}_0$  is an  $\mathcal{L}_0$ -structure with  $\mathcal{M}_0 \models \Gamma$ . Expand  $\mathcal{M}_0$  to an  $\mathcal{L}_1$ -structure  $\mathcal{M}_1$  in any way.  $\mathcal{M}_1 \models \Gamma$ , so  $\mathcal{M}_1 \models \varphi$ . Thus,  $\mathcal{M}_0 \models \varphi$ , so  $\Gamma \models_{\mathcal{L}_0} \varphi$ . ⊛

What is important about the proof system?

**Definition 1.3.18** (Computably enumerable). A set is *computably enumerable* (ce) or *computable listable* if there is a program that lists out its elements.

- (1) Soundness and completeness,  $\vdash \Leftrightarrow \models$ .
- (2) Proofs are finite, and use only finitely many hypotheses  $\Rightarrow$  compactness.
- (3) Computational properties. If  $\mathcal{L}$  is finite, or computable (complete list of symbols and their arities).
  - (a) We can compute with formulas.
  - (b) Given a formula, it's computable to check whether it's a logical axiom.
  - (c) It's computable to check whether a proof is valid.
  - (d) If  $\Gamma$  is a ce set of sentences,  $\{\varphi: \Gamma \vdash \varphi\}$  is also ce.<sup>5</sup>
  - (e) There is no program that given  $\varphi$  can decide whether  $\vdash \varphi$  at least for  $\mathcal{L} = \{E\}$ ,  $E$  binary.

### 1.3.6 Complete Theories

**Proposition 1.3.4.** Let  $T$  be an  $\mathcal{L}$ -theory with an infinite model, and let  $\kappa$  be an infinite cardinal with  $\kappa \geq |\mathcal{L}|$ . Then  $T$  has a model of cardinality  $\kappa$ .

**Proof.** Let  $\mathcal{C}$  be a set of  $\kappa$ -many new constants, and let  $\mathcal{L}^* = \mathcal{L} \cup \mathcal{C}$ . Let

$$T^* = T \cup \{c \neq d \mid c, d \in \mathcal{C} \text{ distinct}\}.$$

If  $\mathcal{M} = T^*$ , then  $|\mathcal{M}| \geq \kappa$ . Also, if  $T^*$  is satisfiable, it has a model of size at most  $|\mathcal{L}^*| = \kappa$  since

$$\kappa = |\mathcal{C}| \leq |\mathcal{L}^*| \leq |\mathcal{C}| + |\mathcal{L}| \leq \kappa + \kappa = \kappa,$$

so if  $T^*$  is satisfiable, it has a model  $\mathcal{M}$  with  $|\mathcal{M}| = \kappa$ .

**Claim.**  $T^*$  is satisfiable.

**Proof.** By the compactness theorem, to show that  $T^*$  is satisfiable, it's equivalent to show that every finite  $\Gamma \subseteq T^*$  is satisfiable. Let  $\mathcal{M}$  be infinite, and  $\Gamma \subseteq T^*$  finite, then

$$\Gamma \subseteq T \cup \{c_i \neq c_j \mid i, j = 1, \dots, n, i \neq j\}$$

for  $c_1, \dots, c_n \in \mathcal{C}$  since only finitely many  $c_i$  are involved, and without loss of generality,  $\Gamma = T \cup \{c_i \neq c_j \mid i, j = 1, \dots, n, i \neq j\}$ . Pick  $a_1, \dots, a_n \in \mathcal{M}$ , distinct, we then turn  $\mathcal{M}$  into an  $\mathcal{L}^*$ -structure  $\mathcal{M}^*$  with  $c_i^{\mathcal{M}^*} = a_i$ .<sup>a</sup> So  $\mathcal{M}^* \models \Gamma$ , thus  $T^*$  is satisfiable.  $\oplus$

<sup>a</sup>And each other  $d \in \mathcal{C}$  with  $d^{\mathcal{M}^*} = a_1$ .

■

## Lecture 11: Algebraically Closed Fields

### 1.4 A Detour to Algebraically Closed Fields

9 Feb. 14:30

Algebraically closed fields are a great example of a *tame* theory (as opposed to e.g.,  $\mathbb{N}$ , which are not *tame*). We detour to discuss some important and related definitions for the future discussion.

<sup>5</sup>We can list out all the valid proofs from  $\Gamma$  of any  $\varphi$ .

### 1.4.1 Rings

All rings  $R$  we refer to will be commutative.

**Definition 1.4.1 (Ideal).** Let  $R$  be a ring. An *ideal*  $I$  of  $R$  is a set  $I \subseteq R$  such that

- (a)  $0 \in I$ ;
- (b) if  $a, b \in I$ , then  $a + b \in I$ ;
- (c) if  $a \in I$  and  $r \in R$ ,  $ra \in I$ .

**Definition 1.4.2 (Proper).** An *ideal* is *proper* if  $1 \notin I$ , equivalently,  $I \neq R$ .

**Definition.** Let  $I$  be a *proper ideal*.

**Definition 1.4.3 (Radical).**  $I$  is *radical* if  $a^n \in I$ , then  $a \in I$ .

**Definition 1.4.4 (Prime).**  $I$  is *prime* if  $ab \in I$ , then  $a \in I$  or  $b \in I$ .

**Definition 1.4.5 (Maximal).**  $I$  is *maximal* if there is no *proper ideal*  $J \supsetneq I$ .

**Remark.** *Maximal*  $\Rightarrow$  *Prime*  $\Rightarrow$  *Radical*.

**Definition 1.4.6 (Polynomial ring).** Let  $R$  be a ring. Then  $R[x_1, \dots, x_n]$  is the *polynomial ring* with coefficients in  $R$  on indeterminates  $x_1, \dots, x_n$ .

**Example.** Let  $K$  be a field,  $S \subseteq K^n$ , and  $I \subseteq K[x_1, \dots, x_n]$  defined as

$$I = \{f(\bar{x}) \mid f(\bar{s}) = 0 \text{ for all } \bar{s} \in S\}.$$

Then  $I$  is a *radical ideal*.

**Theorem 1.4.1.**  $K[x]$  is a principal ideal domain (PID): every *ideal* is generated by one element,  $I = (f(x)) = \{g(x)f(x) \mid g(x) \in K[x]\}$ .

**Proof.** We can let  $g$  be the polynomial of least degree in  $I$ . Then for any other  $h \in I$ , by long division,  $h = gs + r$ , with  $\deg(r) < \deg(g)$ . But then  $r = h - gs \in I$ , so if  $r$  has lower degree than  $g$ ,  $r = 0$ , hence  $h = gs \in (g)$ . ■

**Definition 1.4.7 (Noetherian).** A ring  $R$  is *Noetherian* if every *ideal*  $I$  of  $R$  is finitely generated.

**Remark.** Equivalently, there is no infinite proper ascending chain of *ideals*.

**Theorem 1.4.2 (Hilbert basis theorem).** If  $R$  is a *Noetherian* ring, then  $R[x]$  is also *Noetherian*. In particular,  $K[x_1, \dots, x_n]$  is *Noetherian* and so every *ideal* in  $K[x_1, \dots, x_n]$  is finitely generated.

**Theorem 1.4.3.** If  $\alpha: R \rightarrow S$  is a ring homomorphism, then  $\ker \alpha$  is an *ideal* of  $R$ , and the induced map  $\bar{\alpha}: R / \ker \alpha \rightarrow S$  is injective.

**Theorem 1.4.4.** Let  $R$  be a ring, and  $I$  an **ideal** of  $R$ .<sup>a</sup>

- (a)  $R/I$  is an integral domain<sup>b</sup> if and only if  $I$  is a **prime**.
- (b)  $R/I$  is a field if and only if  $I$  is **maximal**.

<sup>a</sup>Then  $\pi: R \rightarrow R/I$  is a ring homomorphism with kernel  $I$ .

<sup>b</sup>If  $ab = 0$ , then  $a = 0$  or  $b = 0$ .

### 1.4.2 Field Extensions

**Definition 1.4.8** (Field extension). If  $K \subseteq L$  is a subfield of  $L$ , we call  $L/K$  a *field extension*.

Given a **field extension**  $L/K$ , then we have that  $L$  is a  $K$ -vector space, which suggests the following natural notion.

**Definition 1.4.9** (Degree). The *degree* of  $L/K$  is the dimension of the  $K$ -vector space  $L$ .

**Notation** (Finite extension). If  $[L:K]$  is finite, then we say  $L/K$  is a *finite extension*.

**Theorem 1.4.5.** If  $M/L$  and  $L/K$  are **field extensions**, then

$$[M:K] = [M:L][L:K].$$

**Definition.** Let  $L/K$  be a **field extension**, and  $a \in L$ .

**Definition 1.4.10** (Algebraic). If there is a non-zero  $f(x) \in K[x]$  such that  $f(a) = 0$ , then  $a$  is *algebraic* over  $K$ .

**Definition 1.4.11** (Transcendental). If  $a$  is not **algebraic**, then it is *transcendental* over  $K$ .

**Definition 1.4.12** (Minimal polynomial). If  $a$  is **algebraic** over  $K$ , there is a non-zero, monic<sup>a</sup>  $f(x) \in K[x]$  of least degree such that  $f(a) = 0$  which we call the *minimal polynomial* of  $a$  over  $K$ .

<sup>a</sup>This is a common practice.

**Remark.** A **minimal polynomial** is irreducible.

**Remark.** If  $f(x)$  is a **minimal polynomial**, then

$$(f(x)) = \{g(x) \in K[x] \mid g(a) = 0\}.$$

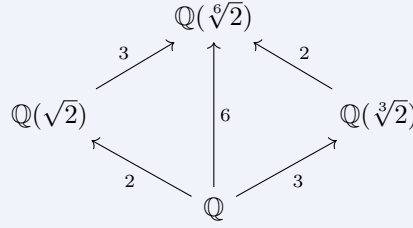
**Example.** Consider a **field extension**  $\mathbb{R}/\mathbb{Q}$  with  $a = \sqrt{2} \in \mathbb{R}$ . Then the **minimal polynomial** is  $f(x) = x^2 - 2$ .

**Theorem 1.4.6.** Let  $L/K$  be a **field extension** and  $a \in L$ , then  $a$  is **algebraic** over  $K$  if and only if  $n = [K(a):K] < \infty$ . Furthermore, if  $a$  is **algebraic** over  $K$ , then  $n$  is the degree of the **minimal polynomial** of  $a$ , and  $1, a, \dots, a^{n-1}$  is a basis for  $K(a)$  as a  $K$ -vector space.

**Intuition.** Think about  $f(a) = a^n + r_{n-1}a^{n-1} + \dots + r_1a + r_0 = 0$ .

The following example illustrates how can we combine **Theorem 1.4.5** and **Theorem 1.4.6**,

**Example.** Let  $f(x) = x^2 - 2$ ,  $\mathbb{Q}(\sqrt{2}) = \{a + b\sqrt{2} \mid a, b \in \mathbb{Q}\}$ .



**Theorem 1.4.7.** Let  $L/K$  be a field extension,  $a \in L$ , and  $f(x) \in K[x]$  be the minimal polynomial of  $a$  over  $K$ .

(a)  $K[x]/(f(x)) \cong K(a)$ .<sup>a</sup>

(b) If  $b \in L$  has the same minimal polynomial as  $a$ , then  $K(a) \cong K[x]/(f(x)) \cong K(b)$ .

<sup>a</sup>Let  $x \in K[x]$ , then  $\bar{x} = x + (f(x)) \in K[x]/(f(x))$ , i.e.,  $\bar{x}$  is a root of  $f$ , hence the isomorphism is given by  $\bar{x} \mapsto a$ .

**Example.** Let  $a = \sqrt{2}$ ,  $b = -\sqrt{2}$ , and  $f(x) = x^2 - 2$  with  $K = \mathbb{Q}$ . Then

$$\begin{aligned} \mathbb{Q}(\sqrt{2}) &\cong \mathbb{Q}[x]/(x^2 - 2) \cong \mathbb{Q}(-\sqrt{2}); \\ a + b\sqrt{2} &\mapsto [a + bx] \mapsto a - b\sqrt{2}. \end{aligned}$$

**Definition 1.4.13 (Algebraic extension).** Let  $L/K$  be a field extension. Then  $L$  is an algebraic extension of  $K$  if all  $a \in L$  are algebraic over  $K$ .

If  $a$  is algebraic over  $K$ , then  $K(a)/K$  is algebraic: If  $b \in K(a)$ , then  $K(b) \subseteq K(a)$ , so  $[K(b):K] \leq [K(a):K] < \infty$ , so  $b$  is algebraic over  $K$ .

**Theorem 1.4.8.** If  $M/L$  and  $L/K$  are algebraic extensions, then  $M/K$  is an algebraic extension.

**Proof.** Let  $a \in M$ , and let  $b_1, \dots, b_n \in L$  be the coefficients of the minimal polynomial of  $a$  over  $L$ . Then  $b_1, \dots, b_n$  are algebraic over  $K$ . Since

$$\begin{aligned} [K(a):K] &\leq [K(a, b_1, \dots, b_n):K] \\ &= [K(a, b_1, \dots, b_n):K(b_1, \dots, b_n)] \cdot [K(b_1, \dots, b_n):K(b_2, \dots, b_n)] \cdots [K(b_n):K]. \end{aligned}$$

Since each of these is a finite extension, so  $[K(a):K] < \infty$ . ■

**Definition 1.4.14 (Algebraically closed).** A field  $L$  is algebraically closed if any non-constant  $f(x) \in L[x]$  has a root in  $L$ .

**Definition 1.4.15 (Algebraic closure).** If  $L/K$ , then  $L$  is an algebraic closure of  $K$  if  $L$  is algebraically closed and an algebraic extension of  $K$ .

**Example.**  $\mathbb{C}$  is algebraically closed, while  $\mathbb{R}$  is not.

Given  $L/K$ ,  $L$  is algebraically closed of  $K$  if  $L$  is algebraic over  $K$ .

**Example.**  $\mathbb{C}$  is the algebraic closure of  $\mathbb{R}$ , and  $[\mathbb{C}:\mathbb{R}] = 2$ .



**Example.**  $\mathbb{Q}^{\text{alg}} = \{a \in \mathbb{C} \mid a \text{ is algebraic over } \mathbb{Q}\}$  is the algebraic closure of  $\mathbb{Q}$ .

If  $L$  is algebraically closed, any  $f(x) \in L[x]$  factors completely as  $f(x) = (x - a_1) \cdots (x - a_n)$  and  $a_1, \dots, a_n$  are the only roots of  $f$ .

**Theorem 1.4.9.** Every field  $K$  has an algebraic closure. If  $L/K$  and  $M/K$  are algebraic closures over  $K$ , then  $L \cong_K M$ .<sup>a</sup>

<sup>a</sup>There exists  $\alpha: L \rightarrow M$  such that  $\alpha(a) = a$  for  $a \in K$ .

**Proof.** First, we show the existence. Assume  $K$ . Let  $f_1, f_2, \dots$  be the polynomials over  $K$ . Start with  $K = K_0$ , let  $g_1(x)$  be an irreducible factor of  $f_1(x)$ . Let

$$K_1 = K_0[x] / (g_1(x)).$$

Since  $g_1$  is irreducible,  $(g_1(x))$  is maximal, so  $K_1$  is a field with a root of  $f_1$ . Now, we build

$$K_1 \subseteq K_1 \subseteq K_2 \subseteq \dots \subseteq K^* = \bigcup_i K_i.$$

Since any  $f(x) \in K$  has a root in  $K^*$ , so  $K^*/K$  is algebraic. Now, take

$$K \subseteq K^* \subseteq K^{**} \subseteq K^{***} \subseteq \dots \subseteq L = \bigcup K^{*...},$$

then  $L$  is the algebraic closure of  $K$ .

Now we prove the uniqueness.

**Lemma 1.4.1.** An algebraically closed field has no algebraic extension.

**Proof.** If  $a$  is algebraic over an algebraically closed  $L$ , the minimal polynomial of  $a$ ,  $f(x)$ , factors completely and is irreducible, so  $f(x) = x - r$ ,  $r \in L$ . Then  $f(a) = 0$  implies  $a = r \in L$ . ■

**Lemma 1.4.2.** Let  $L/K$  algebraic,  $M/K$  algebraically closed. Then there is an embedding  $\alpha: L \rightarrow M$  fixing  $K$ .

**Proof idea.** Consider the case that  $L = K(a)$ , then let  $f(x)$  be the minimal polynomial of  $a/K$ . Let  $b$  be a root of  $f$  in  $M$ .  $K(a) \cong K[x]/(f) \cong K(b) \subseteq M$ , this is our  $\alpha$ .

Then, the full proof is to keep doing this and perhaps by using Zorn's lemma or transfinite induction. ■

Now, if  $L/K$  and  $M/K$  are algebraic closures over  $K$ , then by Lemma 1.4.2, there is  $\alpha: L \rightarrow M$ .  $M/\alpha(L)$  is an algebraic extension, and  $\alpha(L) \cong L$  is algebraically closed, by Lemma 1.4.1,  $M = \alpha(L)$ , so  $\alpha$  is an isomorphism  $M \rightarrow L$  over  $K$ . ■

## Lecture 12

**Definition 1.4.16 (Characteristic).** A field  $F$  has finite characteristic  $p > 0$  if  $\underbrace{1 + \dots + 1}_{p \text{ times}} = 0$ .

14 Feb. 14:30

**Remark.**  $p$  is always prime, otherwise,  $F$  has characteristic  $p = 0$ , i.e.,  $1 + \dots + 1 \neq 0$ , always.

**Definition 1.4.17 (Prime field).** The prime field  $\mathbb{F}_p$  in characteristic  $p$  such that  $\mathbb{F}_p = \mathbb{Q}$  if  $p = 0$ ,  $\mathbb{F}_p = \mathbb{Z}/p\mathbb{Z}$  if  $p > 0$ .

**Definition 1.4.18** (Transcendence basis). Let  $L/K$  be a field extension. A set  $S \subseteq L$  is called a *transcendence basis* of  $L/K$  if:

- (a)  $S$  is algebraically independent: no  $a_1, \dots, a_n \in S$  have non-zero polynomial  $f(x_1, \dots, x_n) \in K[\bar{x}]$  with  $f(a_1, \dots, a_n) = 0$ ;
- (b)  $L$  is an algebraic extension of  $K(S)$ , i.e.,  $S$  is maximal.

**Remark.** Every field extension has<sup>a</sup> a transcendence basis, and any two transcendence basis have the same size.

<sup>a</sup>Same as vector spaces.

**Example.** Let  $K(t_1, \dots, t_n)$  be the fraction field of  $K[x_1, \dots, x_n]$ .  $\{t_1, \dots, t_n\}$  is a transcendence basis for  $K(t_1, \dots, t_n)$  over  $K$ .

**Definition 1.4.19** (Transcendence degree). The *transcendence degree* of  $L$  over  $K$  is the cardinality of any transcendence basis.

If we do not specify  $K$ , then  $K$  is the prime field  $K = \mathbb{F}_p$ .

**Theorem 1.4.10.** Any two algebraically closed fields of the same characteristic  $p$  and transcendence degree are isomorphic.

**Proof.** Let  $L, K$  be those fields, with transcendence basis  $S, T$  over  $\mathbb{F}_p$ .  $L$  is the algebraic closure of  $\mathbb{F}_p(S)$  and  $K$  is the algebraic closure of  $\mathbb{F}_p(T)$ . There is a bijection  $f: S \rightarrow T$ , and then  $f$  extends to  $\bar{f}: \mathbb{F}_p(S) \rightarrow \mathbb{F}_p(T)$  such that

$$\bar{f}\left(\frac{\sum_{\alpha} r_{\alpha} \bar{x}^{\alpha}}{\sum_{\alpha} s_{\alpha} \bar{x}^{\alpha}}\right) = \frac{\sum_{\alpha} r_{\alpha} f(\bar{x})^{\alpha}}{\sum_{\alpha} s_{\alpha} f(\bar{x})^{\alpha}},$$

where  $r_{\alpha}, s_{\alpha} \in \mathbb{F}_p$  and  $\bar{x}^{\alpha}$  is some monomial from  $S$ , e.g.,  $x_1^2 x_2$  for  $x_1, x_2 \in S$ .<sup>a</sup>

$\mathbb{F}_p(S)$  and  $\mathbb{F}_p(T)$  are the same (up to isomorphism), but the algebraic closures are unique, so  $K \cong L$  via an isomorphism extending  $\bar{f}$ . ■

<sup>a</sup> $\alpha$  can be thought as a tuple, in the case of  $x_1^2 x_2$ ,  $\alpha = (2, 1)$ .

**Definition 1.4.20** (ACF). ACF is the theory of algebraically closed fields consists of the following.

- (a) Field axioms.
- (b) For every  $n \geq 1$ ,  $\forall a_0 \dots \forall a_n (a_n \neq 0 \rightarrow \exists b a_n b^n + a_{n-1} b^{n-1} + \dots + a_0 = 0)$ .

**Remark.** The models of ACF are exactly the algebraically closed field, and the language  $\mathcal{L} = \mathcal{L}_{\text{ring}} = \{0, 1, +, -, \cdot\}$ .

For  $p > 0$  prime, let

$$\text{ACF}_p = \text{ACF} \cup \{ \underbrace{1 + \dots + 1}_p = 0 \},$$

and

$$\text{ACF}_p = \text{ACF} \cup \{ \underbrace{1 + \dots + 1}_n \neq 0 \mid n \in \mathbb{N} \}.$$

**Definition 1.4.21** (Categorical). Let  $\kappa$  be an infinite cardinal and  $T$  be an  $\mathcal{L}$ -theory.  $T$  is  $\kappa$ -categorical if any  $\mathcal{M}, \mathcal{N} \models T$  of size  $\kappa$  have  $\mathcal{M} \cong \mathcal{N}$ .

**Definition 1.4.22** (Countably categorical). If  $\kappa$  is countable, then  $T$  is *countably categorical*.

**Definition 1.4.23** (Uncountably categorical). If  $\kappa$  is uncountable, then  $T$  is *uncountably categorical*.

We see that for being *uncountably categorical*, we only need one uncountable  $\kappa$ .

**Example.**  $(\mathbb{Q}, \leq)$  is *countably categorical*.

**Lemma 1.4.3.** If  $K$  has *transcendence degree*  $\lambda$ , then  $|K| = \lambda + \aleph_0$ .

**Proof.**  $K$  is *algebraic* over  $\mathbb{F}_p(S)$ , where  $S$  is a *transcendence basis* of size  $\lambda$ . By counting,  $|\mathbb{F}_p(S)| = \lambda + \aleph_0$ , so  $|\mathbb{F}_p(S)[x]| = \lambda + \aleph_0$ . But since each element of  $K$  satisfies some polynomials, and each polynomial has finitely many roots in  $K$ , so  $|K| = \lambda + \aleph_0$ . ■

**Theorem 1.4.11.** Fix  $p$ .  $\text{ACF}_p$  is  *$\kappa$ -categorical* for every uncountable  $\kappa$ .

**Proof.** Let  $L, K$  be  $\text{ACF}_p$  for size  $\kappa$ , then  $L, K$  have *transcendence degree*  $\kappa$ , and hence are isomorphic from **Theorem 1.4.10**. With the application of **Lemma 1.4.3**, we're done. ■

**Example.**  $\mathbb{Q}^{\text{alg}}$ , the *algebraically closure* of  $\mathbb{Q}$ , has size  $\aleph_0$ , and has *transcendence degree* is 0.

**Example.**  $\mathbb{Q}(t)^{\text{alg}}$ , the *algebraically closure* of  $\mathbb{Q}(t) \cong \mathbb{Q}(\pi)$ , has size  $\aleph_0$ , and has *transcendence degree* is 1.

**Proof.** We see that

$$\mathbb{Q}(t)^{\text{alg}} = \{z \in \mathbb{C} \mid z \text{ is algebraic over } \mathbb{Q}(\pi)\}.$$

These are countable, but not isomorphic.  $\text{ACF}_0$  is not *countably categorical*. The same with  $\text{ACF}_0$  for  $p > 0$ . \*

**Note.**  $\text{ACF}$  is not *uncountably categorical*.

**Theorem 1.4.12** (Vaught's test). Let  $T$  be a *satisfiable  $\mathcal{L}$ -theory* with no finite *models*. If  $T$  is  *$\kappa$ -categorical* for some infinite  $\kappa \geq |\mathcal{L}|$ , then  $T$  is *complete*.

**Proof.** Suppose  $T$  was not *complete*, so pick  $\varphi$  with  $T \not\models \varphi$  and  $T \not\models \neg\varphi$ , and hence  $T \cup \{\varphi\}$  and  $T \cup \{\neg\varphi\}$  are *satisfiable*. By a consequence of the proof of *completeness theorem* (with a *compactness* argument),

- $T \cup \{\varphi\}$  has a *model*  $\mathcal{M}$  of size  $\kappa$ , and
- $T \cup \{\neg\varphi\}$  has a *model*  $\mathcal{N}$  of size  $\kappa$ .

But  $T$  is  *$\kappa$ -categorical*, so  $\mathcal{M} \cong \mathcal{N}$ , which is a contradiction ✗ ■

**Corollary 1.4.1.**  $\text{ACF}_p$  is *complete* for each  $p$ .

The axioms for  $\text{ACF}_p$  completely determines all of the first-order facts about *algebraically closed* fields of *characteristic*  $p$ .

**Remark (Fact).** The axioms for  $\text{ACF}$  or  $\text{ACF}_p$  can be *listed computably*. So  $\{\varphi \mid \text{ACF} \models \varphi\}$  and  $\{\varphi \mid \text{ACF}_p \models \varphi\}$  can be *listed computably*.

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**Definition 1.4.24 (Decidable).** A theory  $T$  is *decidable* if there is a program that given  $\varphi$ , it determines whether  $T \models \varphi$ .

**Remark.**  $\text{ACF}_p$  is decidable.

**Proof.** Given  $\varphi$ , either  $\text{ACF}_p \models \varphi$  or  $\text{ACF}_p \models \neg\varphi$  since  $\text{ACF}_p$  is complete. By looking for a proof of  $\varphi$  and a proof of  $\neg\varphi$ , eventually we will find one, telling us whether  $\text{ACF}_p \models \varphi$ .  $\circledast$

**Theorem 1.4.13.**  $\text{ACF}$  is decidable.

**Proof.** Given  $\varphi$ , simultaneously

- (a) Look for a proof of  $\text{ACF} \vdash \varphi$ , and
- (b) Look for  $p$  such that  $\text{ACF}_p \vdash \neg\varphi$  (so  $\text{ACF} \not\models \varphi$ ).<sup>a</sup>

The first case is fine. Suppose  $\text{ACF} \not\models \varphi$ , so there is  $\mathcal{M} \models \text{ACF}$ ,  $\mathcal{M} \models \neg\varphi$ . There is  $p$  such that  $\mathcal{M} \models \text{ACF}_p$ . Since  $\text{ACF}_p$  is complete,  $\text{ACF}_p \models \neg\varphi$ , so the search of the second case will halt, so the whole search will eventually halt.  $\blacksquare$

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<sup>a</sup>We don't know  $\text{ACF} \models \neg\varphi$ .

# Appendix

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