# STAT576 Empirical Process Theory

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#### Abstract

This is a graduate-level theoretical statistics course taught by Sabyasachi Chatterjee at University of Illinois Urbana-Champaign, aiming to provide an introduction to empirical process theory with applications to statistical M-estimation, non-parametric regression, classification and high dimensional statistics.

While there are no required textbooks, some books do cover (almost all) part of the material in the class, e.g., Van Der Vaart and Wellner's Weak Convergence and Empirical Processes [VW96].



This course is taken in Fall 2023, and the date on the covering page is the last updated time.

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# Chapter 1

# Introduction

#### Lecture 1: Introduction to Mathematical Statistics

# 1.1 What is Empirical Process Theory?

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This subject started in the 1930s with the study of the empirical CDF.

**Definition 1.1.1** (Empirical CDF). Given inputs i.i.d. data points  $X_1, \ldots, X_n \sim \mathbb{P}$ , the *empirical CDF* is

$$F_n(t) = \frac{1}{n} \sum_{i=1}^n \mathbb{1}_{X_i \le t}.$$

The classical result is that, fixing  $t, F_n(t) \to F(t)$  almost surely.

**Note.** At the same time,  $\sqrt{n}(F_n(t) - F(t)) \to \mathcal{N}(0, F(t)(1 - F(t)))$  in distribution.

On the other hand, we can also ask does this convergence happen if we jointly consider all possible  $t \in \mathbb{R}$ . By the Glivenko-Cantelli theorem,  $\sup_{t \leq \mathbb{R}} |F_n(t) - F(t)| \stackrel{n \to \infty}{\to} 0$  almost surely, so the answer is again ves.

Now, we're ready to see a "canonical" example of an empirical process.

**Example** (Canonical empirical process). The *canonical empirical process* is the family of random variables  $\{F_n(t)\}_{t\in\mathbb{R}}$ , i.e., a stochastic process.

By considering a general class of functions, we have the following.

**Definition 1.1.2** (Empirical process). Let  $\chi$  be the domain,  $\mathbb{P}$  be a distribution on  $\chi$ , and  $\mathscr{F}$  be the class of function such that  $\chi \to \mathbb{R}$ . The *empirical process* is the stochastic process indexed by functions in  $\mathscr{F}$ ,  $\{G_n(f): f \in \mathscr{F}\}$  where

$$G_n(f) = \frac{1}{n} \sum_{i=1}^n f(X_i) - \mathbb{E}\left[f(X)\right]$$

and  $X_1, \ldots, X_n \stackrel{\text{i.i.d.}}{\sim} \mathbb{P}$ .

**Remark.** The empirical process is a family of mutually dependent random variables, all of them being functions of the same inherent randomness in the i.i.d. data  $X_1, \ldots, X_n$ .

Now, two questions arises.

#### 1.1.1 Uniform Law of Large Numbers

As  $n \to 0$ , whether

$$S_n(\mathscr{F}) := \sup_{f \in \mathscr{F}} |G_n(f)| \to 0,$$

and if, at what rate?

**Remark.** The rate of convergence of law of large numbers uniformly over a class of functions  $\mathscr{F}$  determines the performance of many types of statistical estimators as we will see.

We will spend most of this course just on this topic with applications. We will show that  $S(\mathscr{F})$  concentrates around its expectation and will bound  $\mathbb{E}[S(\mathscr{F})]$ .

#### 1.1.2 Uniform Central Limit Theorem

The most general probabilistic question one can ask is the following:

**Problem.** What is the joint distribution of the empirical process?

Answer. For a given sample size, it's most often intractable to be able to calculate the joint distribution exactly. One can then use asymptotics when the sample size n is very large to derive limiting distributions. By the regular central limit theorem,  $\sqrt{n}G_n(f) \stackrel{d}{\to} \mathcal{N}(0, \text{Var}[f(X)])$  for any f. We want to understand if this holds uniformly (jointly) over  $f \in \mathscr{F}$  in some sense.

We first motivate this through an example.

**Example** (Uniform empirical process). Consider

- $X_1, \ldots, X_n$  i.i.d. from  $\mathcal{U}(0,1)$ .
- $\mathscr{F} = \{\mathbb{1}_{[-\infty,t]} : t \in \mathbb{R}\}$
- $U_n(t) = \sqrt{n}(F_n(t) t)$  where  $F_n$  is the empirical CDF.

We can view  $U_n(t)$  as collection of random variables one for each  $t \in (0,1)$ , or just as a random function. Then this stochastic process  $\{U_n(t): t \in (0,1)\}$  is called the *uniform empirical process*.

Then, the CLT states that for each  $t \in [0,1]$ ,  $U_n(t) \to \mathcal{N}(0,t-t^2)$  as  $n \to \infty$ . Moreover, for fixed  $t_1, \ldots, t_k$ , the multivariate CLT implies that  $(U_n(t_1), \ldots, U_n(t_k)) \stackrel{d}{\to} \mathcal{N}(0, \Sigma)$  where  $\Sigma_{ij} = \min(t_i, t_j) - t_i t_j$ .

 $^{a}\mathcal{U}$  denotes the uniform distribution.

From this example, one can ask question like the following.

**Problem.** Does the entire process  $\{U_n(t): t \in [0,1]\}$  converge in some sense? If so, what is the limiting process?

**Answer**. The limiting process is an object called the *Brownian Bridge*. This was conjectured by Doob and proved by Donsker.

Other than that, how do we characterize convergence of stochastic processes in distribution to another stochastic process? How do we generalize this result for a general function class  $\mathscr{F}$  defined on a probability space  $\chi$ ? What are some statistical applications of such process convergence results? This is a classical topic and in the last few weeks of this course, we will touch upon some of these questions.

# 1.2 Applications of Uniform Law of Large Numbers

Next, we see one major example where uniform law of large numbers can be applied.

#### 1.2.1 M-Estimators

Consider the class of estimators called "M-estimator", which is of the form

$$\hat{\theta} = \underset{\theta \in \Theta}{\operatorname{arg\,min}} \frac{1}{n} \sum_{i=1}^{n} M_{\theta}(X_i),$$

where  $X_1, \ldots, X_n$  taking values in  $\chi$ ,  $\Theta$  is the parameter space, and  $M_{\theta} \colon \chi \to \mathbb{R}$  for each  $\theta \in \Theta$ . Let's see some examples.

**Example** (Maximum log-likelihood).  $M_{\theta}(X) = -\log p_{\theta}(X)$  for a class of densities  $\{p_{\theta} : \theta \in \Theta\}$ , then  $\hat{\theta}$  is the Maximum log-likelihood of  $\theta$ .

There are lots of examples on "local estimators" as well.

**Example** (Mean).  $M_{\theta}(x) = (x - \theta)^2$ .

**Example** (Median).  $M_{\theta}(x) = |x - \theta|$ .

**Example** ( $\tau$  quantile).  $M_{\theta}(x) = Q_{\tau}(x - \theta)$  where  $Q_{\tau}(x) = (1 - \tau)x\mathbb{1}_{x < 0} + \tau x\mathbb{1}_{x \ge 0}$ .

**Example** (Mode).  $M_{\theta}(x) = -\mathbb{1}_{|X-\theta| \leq 1}$ .

Now, the target quantity for the estimator  $\hat{\theta}$  is

$$\theta_0 = \operatorname*{arg\,max}_{\theta \in \Theta} \mathbb{E}\left[M_{\theta}(X_1)\right]$$

where  $X_1, \ldots, X_n \overset{\text{i.i.d.}}{\sim} \mathbb{P}$ . In the asymptotic framework, the two key questions are the following.

**Problem.** Is  $\hat{\theta}$  consistent for  $\theta_0$ ? Does  $\hat{\theta}$  converge to  $\theta_0$  almost surely or in probability as  $n \to \infty$ ? I.e., is  $d(\hat{\theta}, \theta_0) \to 0$  for some metric d?

**Problem.** What is the rate of convergence of  $d(\hat{\theta}, \theta_0)$ ? For example is it  $O(n^{-1/2})$  or  $O(n^{-1/3})$ ?

To answer these questions, one is led to investigate the closeness of the empirical objective function to the population objective function in some uniform sense. Consider  $M_n(\theta) = \frac{1}{n} \sum_{i=1}^n M_{\theta}(X_i)$  and  $M(\theta) = \mathbb{E}[M_{\theta}(X_1)]$ , then

$$\mathbb{P}(d(\hat{\theta}, \theta_0) > \epsilon) \leq \mathbb{P}\left(\sup_{\theta \colon d(\theta, \theta_0) > \epsilon} M_n(\theta_0) - M_n(\theta) \geq 0\right)$$

$$= \mathbb{P}\left(\sup_{\theta \colon d(\theta, \theta_0) > \epsilon} (M_n(\theta_0) - M(\theta_0) - [M_n(\theta) - M(\theta)]) \geq \inf_{\theta \colon d(\theta, \theta_0) > \epsilon} (M(\theta) - M(\theta_0))\right)$$

$$\leq \mathbb{P}\left(2\sup_{\theta \in \Theta} |M_n(\theta) - M(\theta)| \geq \inf_{\theta \colon d(\theta, \theta_0) > \epsilon} (M(\theta) - M(\theta_0))\right).$$

We see that the left-hand side  $2\sup_{\theta\in\Theta}|M_n(\theta)-M(\theta)|$  is just  $S(\mathscr{F})$  for  $\mathscr{F}=\{f_\theta\colon\theta\in\Theta,f_\theta=M_\theta(\cdot)\}$ , while the right-hand side  $\inf_{\theta\colon d(\theta,\theta_0)>\epsilon}M(\theta)-M(\theta_0)$  is larger than 0.

Remark. The last step could be too loose in some problems.

### Lecture 2: Sub-Gaussian Random Variables and the MGF Trick

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# 1.3 Bounding Supremum of Empirical Process

Most of this course will focus on bounding suprema of the empirical process. Let's define it rigorously.

**Problem 1.3.1** (Bounding supremum of empirical process). Given a domain  $\chi$ , a probability measure  $\mathbb{P}$  on  $\chi$ , data  $X_1,\ldots,X_n \overset{\text{i.i.d.}}{\sim} \mathbb{P}$ , and a function class  $\mathscr{F} \ni f \colon \chi \to \mathbb{R}$ . We want to find an (non-asymptotically) bound on

$$S_n(\mathscr{F}) = \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^n f(X_i) - \mathbb{E}\left[f(X)\right] \right|.$$

Answer. To do this, broadly speaking, we will go through a route with three basic steps:

- (a)  $S_n(\mathscr{F})$  "concentrates" around its expectation  $\mathbb{E}[S_n(\mathscr{F})]$ .
- (b)  $\mathbb{E}[S_n(\mathscr{F})] \leq \text{the Rademacher complexity of } \mathscr{F} \text{ via "symmetrization"}.$
- (c) Bounding the Rademacher complexity's expected supremum of a "sub-Gaussian process" by a technique called *chaining*.

\*

Toward this end, we need some basic and fundamental concentration inequalities which are of wide interest and use.

# Chapter 2

# Concentration Inequalities

As we just saw, to solve Problem 1.3.1, we need some basic tools on concentration inequalities. The most celebrate concentration inequality might be the Gaussian tail, which achieve a quadratic exponential decay. Combine this with the classical central limit theorem, we can expect that as  $n \to \infty$ , approximately the Gaussian tail bound kicks in.

However, to get a concrete, non-asymptotic bound for  $S_n(\mathscr{F})$ , we would need more sophisticated tools. Let's start with the basics, i.e., the Gaussian distribution.

#### 2.1 Gaussian Distribution

For us, the gold standard for concentration would be the Gaussian distribution. The property of the Gaussian distribution we are interested in is its rapid tail decay as we mentioned:

**Lemma 2.1.1.** For  $Z \sim \mathcal{N}(0, 1)$ ,

$$\left(\frac{1}{t} - \frac{1}{t^3}\right) \frac{1}{\sqrt{2\pi}} e^{-t^2/2} \le \mathbb{P}(Z \ge t) \le \frac{1}{t} \cdot \frac{1}{\sqrt{2\pi}} e^{-t^2/2}.$$

**Proof.** We want to show

$$\left(\frac{1}{t} - \frac{1}{t^3}\right) \frac{1}{\sqrt{2\pi}} e^{-t^2/2} \le \int_t^\infty \frac{1}{\sqrt{2\pi}} e^{-x^2/2} \, \mathrm{d}x \le \frac{1}{t} \cdot \frac{1}{\sqrt{2\pi}} e^{-t^2/2}$$

$$\Leftrightarrow \left(\frac{1}{t} - \frac{1}{t^3}\right) e^{-t^2/2} \le \int_t^\infty e^{-x^2/2} \, \mathrm{d}x \le \frac{1}{t} \cdot e^{-t^2/2}.$$

Observe that from integration by part (with x/x introduced),

$$\int_{t}^{\infty} \frac{x}{x} \cdot e^{-x^{2}/2} \, \mathrm{d}x = -\frac{e^{-x^{2}/2}}{x} \bigg|_{t}^{\infty} - \int_{t}^{\infty} \frac{e^{-x^{2}/2}}{x^{2}} \, \mathrm{d}x = \frac{e^{-t^{2}/2}}{t} - \int_{t}^{\infty} \frac{e^{-x^{2}/2}}{x^{2}} \, \mathrm{d}x \le \frac{1}{t} \cdot e^{-t^{2}/2}$$

since the integrand  $e^{-x^2/2}/x^2$  is non-negative, which is the desired upper-bound. For the lower bound, if we again apply integration by part (with x/x introduced again), then

$$\int_{t}^{\infty} e^{-x^{2}/2} dx = \frac{e^{-t^{2}/2}}{t} - \int_{t}^{\infty} \frac{x}{x} \cdot \frac{e^{-x^{2}/2}}{x^{2}} dx$$

$$= \frac{e^{-t^{2}/2}}{t} - \left( -\frac{e^{-x^{2}/2}}{x^{3}} \Big|_{t}^{\infty} - \int_{t}^{\infty} 3 \frac{e^{-x^{2}/2}}{x^{4}} dx \right)$$

$$= \frac{e^{-t^{2}/2}}{t} - \frac{e^{-t^{2}/2}}{t^{3}} + \int_{t}^{\infty} 3 \frac{e^{-x^{2}/2}}{x^{4}} dx$$

$$\geq \left( \frac{1}{t} - \frac{1}{t^{3}} \right) e^{-t^{2}/2},$$

since, again, the integrand  $3e^{-x^2/2}/x^4$  is non-negative, so the last term is positive, hence we get the desired lower-bound.

**Corollary 2.1.1.** For all  $t \geq 1$ , we have

$$\mathbb{P}(\mathcal{N}(0,\sigma^2) > t) < e^{-t^2/2\sigma^2}$$

Now, as is suggested by CLT, the following question arises.

**Problem.** Does Corollary 2.1.1 hold for sums of independent random variables? That is, given i.i.d.  $X_1, \ldots, X_n$  with mean  $\mu$  and variance  $\sigma^2$ , whether for all  $t \ge 0$ ,

$$\mathbb{P}(\sqrt{n}(\overline{X} - \mu) \ge t) \le e^{-t^2/2\sigma^2}?$$

**Answer.** Just invoking CLT is not enough, we need to handle the error term in the normal approximation. We can show this directly for a class of distributions with fast tail decay.

To go beyond Gaussian tail bound, let start with the moment generating function (MGF) trick.

#### 2.2 MGF Trick

The MGF trick is easy to develop, but it gives a foundation of all the concentration inequalities we're going to develop. Hence, although it's short, it's worth to make it a separate section.

#### 2.2.1 Markov's Inequality

To start with, the most basic tool to bound tail probabilities is the Markov's inequality.

**Lemma 2.2.1** (Markov's inequality). For a non-negative random variable  $X \geq 0$ ,

$$\mathbb{P}(X \ge t) \le \frac{\mathbb{E}[X]}{t}.$$

**Note.** Markov's inequality is valid as soon as  $\mathbb{E}[X] < \infty$ . That is, it holds even when the second moment does not exist.

**Remark.** The rate of tail decay is slow (O(1/t)). For the Gaussian, by Lemma 2.1.1, it's  $O(e^{-t^2/2})$ .

By the above remark, one might ask the following.

**Problem.** Can we derive faster tail decay bounds in general?

Answer. Yes, if we assume more moments exist. If all moments exist and in particular the MGF exists, like for the Gaussian, then we can expect faster tail decay.

#### 2.2.2 Chebyshev Inequality

Continuing the discussion on the previous problem, for example, if we assume second moment exists, then we can get an  $O(1/t^2)$  tail decay by Chebyshev inequality.

**Lemma 2.2.2** (Generalized Chebyshev inequality). Given a random variable X,

$$\mathbb{P}(|X-\mu| \geq t) = \mathbb{P}(|X-\mu|^p \geq t^p) \leq \min_{p \geq 1} \frac{\mathbb{E}\left[|X-\mu|^p\right]}{t^p}.$$

**Proof.** This is directly implied by the Markov's inequality.

**Remark** (Chebyshev Inequality). For p = 2, we have the usual form

$$\mathbb{P}(|X - \mu| \ge t) \le \frac{\operatorname{Var}[X]}{t^2}$$

**Remark.** All tail bounds are derived using Markov's inequality; the clever part is to apply it to the right random variable. In this sense, every tail bound is just Markov's inequality.

#### 2.2.3 Crarmer-Chernoff Method

In the same vein, developed by Cramer and Chernoff, if we now assume the MGF exists and apply Markov's inequality, we get the MGF trick.

**Lemma 2.2.3** (MGF trick (Crarmer-Chernoff method)). Given a random variable X,

$$\mathbb{P}(X - \mu \ge t) = \mathbb{P}(e^{\lambda(X - \mu)} \ge e^{\lambda t}) \le \inf_{\lambda > 0} \frac{\mathbb{E}\left[e^{\lambda(X - \mu)}\right]}{e^{\lambda t}}.$$

We will use the MGF trick rather than the generalized Chebyshev's inequality to derive tail bounds because MGF of a sum of independent random variables decomposes as the product of the MGF's. It is messier to work with the  $p^{\text{th}}$  moment of a sum of independent random variables.

#### 2.3 Hoeffding's Inequality

#### 2.3.1 Sub-Gaussian Random Variables

We will now consider a class of distributions whose MGF is dominated by the MGF of a Gaussian. Then, in a very clean way, the MGF trick will give us Gaussian tail bounds for these distributions.

**Definition 2.3.1** (Sub-Gaussian). Given a random variable X with  $\mathbb{E}[X] = 0$ , we say X is sub-Gaussian with variance factor  $\sigma^a$  if for all  $\lambda \in \mathbb{R}$ ,

$$\mathbb{E}\left[e^{\lambda X}\right] \le e^{\frac{\sigma^2 \lambda^2}{2}}.$$

**Notation.** We write  $\operatorname{Subg}(\sigma^2)$  for a compact representation of the class of sub-Gaussian random variables with variance factor  $\sigma^2$ .

**Remark.** Observe that if  $X \in \text{Subg}(\sigma^2)$ :

- $-X \in \text{Subg}(\sigma^2);$
- $X \in \text{Subg}(t^2)$  if  $t^2 > \sigma^2$ ;
- $cX \in \text{Subg}(c\sigma^2)$ .

**Lemma 2.3.1** (Equivalent conditions). Given a random variable X with  $\mathbb{E}[X] = 0$ , the following are equivalent for absolute constants  $c_1, \ldots, c_5 > 0$ .

Add proof

- (a)  $\mathbb{E}\left[e^{\lambda X}\right] \leq e^{c_1^2 \lambda^2}$  for all  $\lambda \in \mathbb{R}$ .
- (b)  $\mathbb{P}(|X| \ge t) \le 2e^{-t^2/c_2^2}$ .
- (c)  $(\mathbb{E}[|X|^p])^{1/p} \le c_3\sqrt{p}$ .

<sup>&</sup>lt;sup>a</sup>Also called proxy, sub-Gaussian norm, etc.

- (d) For all  $\lambda$  such that  $|\lambda| \leq 1/c_4$ ,  $\mathbb{E}\left[e^{\lambda^2 X^2}\right] \leq e^{c_4^2 \lambda^2}$ .
- (e) For some  $c_5 < \infty$ ,  $\mathbb{E}\left[e^{X^2/c_5^2}\right] \le 2$ .

**Proof.** Let's just see the first implication from (a) to (b). Given  $X \in \text{Subg}(\sigma)$ ,

$$\mathbb{P}(X \ge t) \le \inf_{\lambda > 0} e^{\lambda^2 \sigma^2 / 2 - \lambda t} \le e^{-\frac{t^2}{2\sigma^2}}$$

where the last inequality follows from minimizing the quadratic function  $\lambda^2 \sigma^2 / 2 - \lambda t$  whose minimizer is  $\lambda^* = t/\sigma^2$ . The same bound holds for the left tail and a union bound gives the two-sided version.

Let's see some examples of the sub-Gaussian random variables.

**Example** (Rademacher random variable).  $\epsilon = \pm 1$  with probability 1/2 is a Subg(1) random variable.

**Proof.** We see that

$$\mathbb{E}\left[e^{\lambda\epsilon}\right] = \frac{1}{2}e^{\lambda} + \frac{1}{2}e^{-\lambda} = \frac{1}{2}\sum_{k=1}^{\infty} \left(\frac{\lambda^k}{k!} + \frac{(-\lambda)^k}{k!}\right) = \sum_{k=1}^{\infty} \frac{\lambda^{2k}}{(2k)!} \le 1 + \sum_{k=1}^{\infty} \frac{(\lambda^2)^k}{2^k k!} = e^{\lambda^2/2}$$

since  $(2k)! \ge 2^k \cdot k!$ .

In fact, the above can be generalized for any bounded random variable.

#### **Lemma 2.3.2.** Given $X \in [a, b]$ such that $\mathbb{E}[X] = 0$ . Then

Add proof

$$\mathbb{E}\left[e^{\lambda X}\right] \le \exp\left(\lambda^2 \frac{(b-a)^2}{8}\right)$$

for all  $\lambda \in \mathbb{R}$ , i.e.,  $X \in \text{Subg}((b-a)^2/4)$ .

**Proof.** We will prove this with a worse constant. Let  $X' \stackrel{\text{i.i.d.}}{\sim} X$  be an i.i.d. copy, then

$$\mathbb{E}\left[e^{\lambda X}\right] = \mathbb{E}\left[e^{\lambda (X - \mathbb{E}\left[X'\right])}\right] = \mathbb{E}\left[e^{\lambda X} \cdot e^{-\lambda (\mathbb{E}\left[X'\right])}\right] \leq \mathbb{E}\left[e^{\lambda X}\right] \cdot \mathbb{E}\left[e^{-\lambda X'}\right] = \mathbb{E}\left[e^{\lambda (X - X')}\right],$$

where we have used the Jensen's inequality for  $e^{-\lambda \mathbb{E}[X']} \leq \mathbb{E}\left[e^{-\lambda X'}\right]$ . Now we introduce a Rademacher random variable  $\epsilon = \pm 1$ , to further write

$$\mathbb{E}\left[e^{\lambda X}\right] \leq \mathbb{E}_{X,X'}\left[e^{\lambda(X-X')}\right] = \mathbb{E}_{X,X',\epsilon}\left[e^{\lambda \cdot \epsilon(X-X')}\right] = \mathbb{E}_{X,X'}\left[\mathbb{E}_{\epsilon}\left[e^{\lambda \epsilon(X-X')}\right]\right],$$

and  $\mathbb{E}_{\epsilon}\left[e^{\lambda\epsilon(X-X')}\right] \leq \mathbb{E}\left[e^{\frac{\lambda^2(X-X')}{2}}\right] \leq e^{\frac{\lambda^2(b-a)^2}{2}}$ , where we used the known bound on MGF of a Rademacher random variable, hence overall, we get

$$\mathbb{E}\left[e^{\lambda X}\right] \leq \mathbb{E}_{X,X'}\left[e^{\frac{\lambda^2(b-a)^2}{2}}\right] = e^{\frac{\lambda^2(b-a)^2}{2}}.$$

<sup>a</sup>This is a trick called symmetrization. A basic example is  $\operatorname{Var}[X] = \frac{1}{2}\mathbb{E}\left[(X - X')^2\right]$ .

**Note.** If a = -1 and b = 1, we get back to the earlier example.

Just like independent Gaussians, sums of independent sub-Gaussians remain sub-Gaussian.

**Lemma 2.3.3** (Closed under convolution). Let  $X_i$  be independent random variables with  $\mathbb{E}[X_i] = \mu_i$ ,

and  $X_i - \mu_i \in \text{Subg}(\sigma_i^2)$ . Then

$$\sum_{i=1}^{n} X_i - \sum_{i=1}^{n} \mu_i \in \text{Subg}\left(\sum_{i=1}^{n} \sigma_i^2\right).$$

**Proof.** We simply observe that

$$\mathbb{E}\left[e^{\lambda \sum_{i}(X_{i}-\mu_{i})}\right] = \prod_{i=1}^{n} \mathbb{E}\left[e^{\lambda(X_{i}-\mu_{i})}\right] \leq e^{\frac{\lambda^{2}(\sum_{i}\sigma_{i}^{2})}{2}}.$$

#### 2.3.2 Hoeffding's Inequality

We can now immediately prove the famous Hoeffding's inequality, which is the main tool in our interest.

**Theorem 2.3.1** (Hoeffding's inequality for sub-Gaussian random variables). Let  $X_i$  be independent random variables with  $\mathbb{E}[X_i] = \mu_i$ , and  $X_i - \mu_i \in \operatorname{Subg}(\sigma_i^2)$ . Then for all  $t \geq 0$ ,

$$\mathbb{P}\left(\left|\sum_{i=1}^{n} (X_i - \mu_i)\right| \ge t\right) \le 2\exp\left(\frac{-t^2}{2\sum_{i} \sigma_i^2}\right).$$

**Proof.** It's immediate from Lemma 2.3.3 and the equivalent condition (b) in Lemma 2.3.1.

## Lecture 3: Sub-Exponential Random Variables

For bounded random variables, we can apply Hoeffding's inequality to obtain the following.

25 Aug. 9:00

Corollary 2.3.1. Let  $X_i \in [a, b]$  be random variables with mean  $\mu_i$ ,

$$\mathbb{P}\left(\sum_{i}(X_{i}-\mu_{i}) \ge t\right) \le \exp\left(-\frac{2t^{2}}{n(b-a)^{2}}\right).$$

As a consequence, if  $X_i$  are i.i.d., then

$$\mathbb{P}(\sqrt{n}(\overline{X} - \mu) \ge t) \le \exp\left(-\frac{-2t^2}{(b-a)^2}\right)$$

Compare this with Gaussian approximation, we then have

$$\mathbb{P}(\sqrt{n}(\overline{X} - \mu) \ge t) \approx \mathbb{P}(\mathcal{N}(0, \sigma^2) \ge t) \le \exp\left(-\frac{t^2}{2\sigma^2}\right),$$

i.e.,  $\sigma^2 \sim (b-a)^2/4$ .

Remark (Comparison between Hoeffding's bound and Gaussian tail bound). We see that

- (a) Hoeffding's inequality can be used for any sample size, but Gaussian approximation can only be used when n is large.
- (b) As  $\sigma^2 \leq (b-a)^2/4$ , we see that Gaussian approximation gives a tighter tail bound.
- (c) Another way to state this is that from CLT we get the asymptotically valid confidence interval

 $<sup>^</sup>a$ One-sided version holds without the factor 2.

<sup>&</sup>lt;sup>1</sup>Actually,  $\sigma^2 \leq (b-a)^2/4$  always holds.

for  $\mu$  as

$$\left[ \overline{X} \pm \frac{\sigma}{\sqrt{n}} Z_{\alpha/2} \right],$$

while from the Hoffding's inequality, we have (finite sample valid) confidence interval

$$\left[ \overline{X} \pm \frac{b-a}{2\sqrt{n}} \sqrt{\log \frac{2}{\alpha}} \right],$$

which is much larger.

The above discussion suggests that if the range is very large compared to the variance, then Hoeffding's inequality may not perform very well. Clearly, such random variables exist. Here are some examples.

#### **Example.** Suppose

$$\mathbb{P}(X = 0) = 1 - 1/k^2$$
  
 $\mathbb{P}(X = \pm K) = 1/2k^2$ 

with  $\mathbb{E}[X] = 0$  and  $\text{Var}[X] \leq 1$ . The range is 2K, which is very large compared to the variance. This is a case where Hoeffding's inequality would not perform very well, in the sense that the confidence interval based on it would be too wide.

Another example is the following.

**Example.** Let  $X_1, \ldots, X_n$  be i.i.d. Bernoulli $(\lambda/n)$ , where each one of them has range 1, but its variance is at most  $\frac{\lambda}{n} \ll 1$ . Then a direct application of Hoeffding's inequality gives

$$\mathbb{P}\left(\sum_{i} X_{i} - \lambda \ge t\right) \le \exp\left(\frac{-2t^{2}}{n}\right).$$

This suggests that  $\sum_i X_i = O(\sqrt{n})$  whereas we know that in this case that the distribution of  $\sum_i X_i$  is close to the Poisson( $\lambda$ ) and thus should be O(1).

On the other hand, the CLT inspired bound would give the right order. This points out that we would like to be able to replace the range term by the variance in Hoeffding's inequality. This is what is done in Bernstein's inequality which we will discuss next.

Let's see some non-examples.

**Example** (Not sub-Gaussian). Some examples of random variables which are not sub-Gaussians random variables are Cauchy, exponential, and Possion random variables.

What about mixture?

**Problem.** Suppose  $Z_1, Z_2 \in \text{Subg}(\sigma^2)$  with mean 0, and consider

$$X = \begin{cases} Z_1, & \text{w.p. } p; \\ Z_2, & \text{w.p. } 1 - p. \end{cases}$$

Is this a sub-Gaussian random variable?

# 2.4 Bernstein's Inequality

#### 2.4.1 Sub-Exponential Random Variables

The main reason for considering the class of sub-Gaussian random variables is that the MGF is finite and thus the MGF trick works. So if we want to extend the MGF trick, we would like to ask the following:

**Problem.** How fat could the tails of a distribution be so that the MGF is finite?

Answer. It turns out that we can allow fatter tails than sub-Gaussian, essentially the PDF can decay no slower than an exponential with a proper exponent.

Consider the following example.

**Example.** Let  $Z^2 \sim \chi^2$ , then for all  $t \geq 1$ ,  $\mathbb{P}(Z^2 > t) = 2\mathbb{P}(Z \geq \sqrt{t}) \leq 2e^{-t/2}$ . It is seen that the rate of decrease of the  $\chi^2$  tail probability is slower than that of normal. In fact, the MGF of  $\chi^2$  is

$$\mathbb{E}\left[e^{\lambda(Z^2-1)}\right] = \begin{cases} \frac{e^{-\lambda}}{\sqrt{1-2\lambda}}, & \text{if } 0 \le \lambda < 1/2; \\ \infty, & \text{if } \lambda \ge 1/2, \end{cases}$$

where we see that the MGF exists in a neighborhood around 0, but not everywhere.

This motivates the following definition.

**Definition 2.4.1** (Sub-exponential). A random variable X is sub-exponential with parameters  $(\sigma^2, \alpha)$  with mean  $\lambda$  if for all  $|\lambda| < 1/\alpha$ 

$$\mathbb{E}\left[e^{\lambda(X-\mu)}\right] \le e^{\frac{\lambda^2\sigma^2}{2}}.$$

It's then immediate to see that  $\operatorname{SubExp}(\sigma^2, \alpha)$  random variables have the same bound on their MGF as a  $\operatorname{Subg}(\sigma^2)$  but only for  $\lambda$  in the interval  $(-\frac{1}{\alpha}, \frac{1}{\alpha})$ .

**Example.** For the  $\chi^2$  random variable  $Z^2$ , we have  $Z^2 \in \text{SubExp}(2,4)$ .

**Proof.** This is immediate from Definition 2.4.1 since For all  $|\lambda| < 1/4$ , we have

$$\frac{e^{-\lambda}}{\sqrt{1-2\lambda}} \le e^{2\lambda^2}.$$

(

With Definition 2.4.1, we can extend the MGF trick naturally.

**Lemma 2.4.1** (Tail decay for sub-exponential random variable). Let  $X \in \operatorname{SubExp}(\sigma^2, \alpha)$  with mean  $\mu$ . Then

$$\mathbb{P}(X - \mu \ge t) \le \begin{cases} e^{-\frac{t^2}{2\sigma^2}}, & \text{if } 0 \le t \le \frac{\sigma^2}{\alpha}; \\ e^{-\frac{t}{2\alpha}}, & \text{if } t > \frac{\sigma^2}{\alpha}. \end{cases}$$

**Proof.** We see that

$$\mathbb{P}(X-\mu \geq t) \leq \inf_{0 \leq \lambda < 1/\alpha} \frac{\mathbb{E}\left[e^{\lambda(X-\mu)}\right]}{e^{\lambda t}} \leq \inf_{0 \leq \lambda < 1/\alpha} e^{\frac{\lambda^2 \sigma^2}{2} - \lambda t}.$$

Now, we just need to minimize the exponent, which is a convex quadratic function, in the range  $(0, \frac{1}{\alpha})$ . The infimum depends on the value of  $\alpha$ :

- $\frac{t}{\sigma^2} < \frac{1}{\alpha}$ : we get the Gaussian bound.
- $\frac{t}{\sigma^2} \ge \frac{1}{\alpha}$ : the minimizer is  $1/\alpha$ , and we get the exponential bound.

**Corollary 2.4.1.** Let  $X \in \text{SubExp}(\sigma^2, \alpha)$  with mean  $\mu$ . Then

$$\mathbb{P}(|X - \mu| \ge t) \le 2 \exp\left(-\frac{t^2}{2(\sigma^2 + t\alpha)}\right)$$

for all  $t \geq 0$ .

**Proof.** We see that

$$\mathbb{P}(|X - \mu| \ge t) \le 2 \exp\left(-\min\left\{\frac{t^2}{2\sigma^2}, \frac{t}{2\alpha}\right\}\right) \le 2 \exp\left(-\frac{t^2}{2(\sigma^2 + t\alpha)}\right)$$

by observing  $\min(1/u, 1/v) \ge 1/(u+v)$ .

Just like Lemma 2.3.3 for sub-Gaussian random variables, sub-exponential random variables are also closed under convolution.

**Lemma 2.4.2** (Closed under convolution). Let  $X_i \in \text{SubExp}(\sigma_i^2, \alpha_i)$  be all independent with mean  $u_i$ , then

$$\sum_{i} (X_i - \mu_i) \in \text{SubExp}\left(\sum_{i} \sigma_i^2, \|\alpha\|_{\infty}\right).$$

**Proof.** Since

$$\mathbb{E}\left[e^{\lambda \sum_i (X_i - \mu_i)}\right] = \prod_{i=1}^n \mathbb{E}\left[e^{\lambda (X_i - \mu_i)}\right] \le \prod_{i=1}^n e^{\lambda^2 \sigma_i^2/2} = e^{\lambda^2 \sum_i \sigma_i^2/2}$$

where the inequality holds if  $|\lambda| < 1/\alpha_i$  for all i, i.e.,  $|\lambda| < 1/\|\alpha\|_{\infty}$ .

#### 2.4.2 Bernstein's Inequality

We are now ready to state the generalization of Hoeffding's inequality to sums of independent sub-exponential random variables.

**Theorem 2.4.1** (Bernstein's inequality for sub-exponential random variables). Let  $X_i \sim \operatorname{SubExp}(\sigma_i^2, \alpha_i)$  be all independent with mean  $\mu_i$ , then

$$\mathbb{P}\left(\left|\sum_{i=1}^{n} (X_i - \mu_i)\right| \ge t\right) \le 2 \exp\left(-\min\left\{\frac{t^2}{2\sum_i \sigma_i^2}, \frac{t}{2\|\alpha\|_{\infty}}\right\}\right).$$

**Proof.** This is immediate from Lemma 2.4.1 and Lemma 2.4.2.

We can restate Bernstein's inequality in a convenient way.

**Corollary 2.4.2.** Let  $X_i \sim \operatorname{SubExp}(\sigma_i^2, \alpha_i)$  be all independent with mean  $\mu_i$ , and let  $k \geq \sigma_i, \alpha_i$  for all i. Then for all  $a_i \in \mathbb{R}$ , we have

$$\left\| \mathbb{P}\left( \left| \sum_{i=1}^n a_i(X_i - \mu_i) \right| \ge t \right) \le 2 \exp\left( -\min\left\{ \frac{t^2}{k^2 \|a\|^2}, \frac{t}{k \|a\|_{\infty}} \right\} \right).$$

**Note.** If we let  $a_i = 1/\sqrt{n}$ , we obtain an absolute constant c (depending on k only)

$$\mathbb{P}\left(\left|\frac{1}{\sqrt{n}}\sum_{i=1}^{n}(X_i - \mu_i)\right| \ge t\right) \le \begin{cases} 2e^{-ct^2}, & \text{if } 0 < t < c\sqrt{n}; \\ 2e^{-t\sqrt{n}}, & \text{if } t > c\sqrt{n}. \end{cases}$$

**Remark.** Bernstein's inequality gives the sub-Gaussian tail decay expected from CLT for most t. Only in the very rare event regime, does the slower exponential tail decay come in.

## Lecture 4: McDiarmid's Inequality

#### 2.5 Bounded Difference Concentration Inequality

28 Aug. 9:00

#### 2.5.1 Applications of Berstein's Inequality to Bounded Random Variables

Now we see some applications of Bernstein's inequality, addressing weaknesses of Hoeffding's inequality.

**Lemma 2.5.1.** Let  $|X - \mu| \le b$  and  $X - \mu$  is  $\operatorname{Subg}(b^2)$ . It's also true that  $X - \mu \in \operatorname{SubExp}(2\sigma^2, 2b)$  where  $\operatorname{Var}[X] = \sigma^2$ .

**Proof.** From  $(X - \mu)^k \le (X - \mu)^2 |X - \mu|^{k-2} \le (X - \mu)^2 b^{k-2}$ , we have

$$\mathbb{E}\left[e^{\lambda(X-\mu)}\right] = 1 + \frac{\lambda^2}{2}\sigma^2 + \sum_{k=3}^{\infty} \lambda^k \frac{\mathbb{E}\left[X-\mu\right]^k}{k!} \le 1 + \frac{\lambda^2\sigma^2}{2} + \frac{\lambda\sigma^2}{2} \sum_{k=3}^{\infty} (|\lambda|b)^{k-2}.$$

The last sum is a geometric series, which converges if  $|\lambda| < 1/b$  to

$$1 + \frac{\lambda^2 \sigma^2}{2} \left( \frac{1}{1 - b|\lambda|} \right).$$

Then from  $1 + x \le e^x$ , we see that for  $|\lambda| < 1/2b$ .

$$\mathbb{E}\left[e^{\lambda(X-\mu)}\right] \le e^{\frac{\lambda^2\sigma^2}{2(1-b|\lambda|)}} \le e^{\lambda^2\sigma^2}.$$

From this, by directly apply Bernstein's inequality, we have the following.

**Corollary 2.5.1.** Let X be a random variable such that  $|X - \mu| \le b$ . For any t > 0,

$$\mathbb{P}(|X - \mu| \ge t) \le 2 \exp\left(\frac{-t^2}{2(2\sigma^2 + t \cdot 2b)}\right).$$

Furthermore, let  $X_1, \ldots, X_n$  be independent random variables with  $\mathbb{E}[X_i] = \mu_i$  and  $\text{Var}[X_i] = \sigma_i^2$  such that  $|X_i - \mu_i| \le b$  for all i. Then for any t > 0,

$$\left\| \mathbb{P}\left( \left| \sum_{i=1}^{n} (X_i - \mu_i) \right| \ge t \right) \le 2 \exp\left( \frac{-t^2}{4 \left( \sum_{i} \sigma_i^2 + tb \right)} \right).$$

In particular, if  $\mu_i = \mu$  for all i, then

$$\Pr\left(\left|\frac{1}{n}\sum_{i=1}^{n}X_{i}-\mu\right|\geq t\right)\leq 2\exp\left(-\frac{nt^{2}}{4(\sigma^{2}+tb)}\right).$$

**Remark.** Observe that in the last line of the proof of Lemma 2.5.1, the inequality is quite loose. This means that we can explicitly maximize the quantity in the exponent over  $|\lambda| \in (0, 1/2b)$  to get a higher bound and hence, a better variance factor. This leads to a tighter version of Corollary 2.5.1.

**Corollary 2.5.2.** Let  $X_1, \ldots, X_n$  be independent random variables with  $\mathbb{E}[X_i] = \mu$  and  $\text{Var}[X_i] = \sigma^2$  such that  $|X_i - \mu| \le b$  for all i. Then for any t > 0,

$$\mathbb{P}\left(\left|\sum_{i=1}^{n} X_i - \mu\right| \ge t\right) \le 2\exp\left(\frac{-t^2/2}{n\sigma^2 + bt/3}\right).$$

In particular,

$$\mathbb{P}\left(\left|\frac{1}{n}\sum_{i=1}^{n}X_{i}-\mu\right|\geq t\right)\leq 2\exp\left(\frac{-nt^{2}/2}{\sigma^{2}+bt/3}\right).$$

From Corollary 2.5.2:

- if  $t \leq 3\sigma^2/b$ , the tail of the sample mean behaves like a sub-Gaussian tail;
- if  $t > 3\sigma^2/b$ , the tail of the sample mean behaves like a sub-exponential tail.

**Remark.** In practice, since we know that sample mean is  $\sqrt{n}$ -consistent, we generally look at a sequence of quantiles of the sample mean that is of  $O(n^{-1/2})$ . Therefore, the tail behavior when t gets large, is practically irrelevant.

By choosing the appropriate t in the above tail bound, we can get the following confidence interval for  $\mu$ .

Corollary 2.5.3. Under the assumption of Corollary 2.5.2,

$$\mathbb{P}\left(\left|\frac{1}{n}\sum_{i=1}^{n}X_{i} - \mu\right| \le \frac{\sigma}{\sqrt{n}}\sqrt{2\log\frac{2}{\alpha}} + \frac{3b}{3n}\log\frac{2}{\alpha}\right) \ge 1 - \alpha$$

**Proof.** Let

$$\alpha = 2\exp\left(\frac{-t^2}{2(V+bt/3)}\right),\,$$

then

$$t^2 - \frac{2tb}{3}\log\frac{2}{\alpha} - 2V\log\frac{2}{\alpha} = 0.$$

In Corollary 2.5.3, we have an  $O(1/\sqrt{n})$  term, which is similar to the one derived from Hoeffding's inequality for bounded random variables. In contrary to the Hoeffding's bound, we have an additional lower order term here.

**Remark.** Observe that the higher order term in Corollary 2.5.3 involves the variance, whereas in the case of Hoeffding, it involves the range. Therefore, for random variables with large range but highly concentrated around its mean, the Hoeffding confidence interval would be much wider.

The above remark is demonstrated by the following example.

**Example.** Let  $X_1, \ldots, X_n \overset{\text{i.i.d.}}{\sim} \operatorname{Ber}(p)$ . Suppose we observe  $X_i = 0$  for all i, then  $\hat{p} = \overline{X} = 0$  and the estimate of  $\operatorname{Var}[X_1]$  would be  $\hat{p}(1-\hat{p}) = 0$ .

Hence, if we plug this estimate of variance into the confidence bound from Bernstein, the length of which would be O(1/n). However, in the case of Hoeffding (which works with the range, in this case, 1), the length would be  $O(1/\sqrt{n})$ .

#### 2.5.2 McDiarmid's Inequality

Now we go back to the discussion about empirical process. We do the first step, i.e., we want to show

$$S_n = \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^n f(X_i) - \mathbb{E}\left[f(X)\right] \right|$$

"concentrates" when  $\mathcal{F}$  is bounded provided that

$$\sup_{x \in \chi, f \in \mathscr{F}} |f(x)| \le B.$$

One simple example of bounded function class arises in the task of classification.

**Example** (Classification). Consider f(x) corresponds to the class label of an observation with feature value x, then the class is bounded.

However, since  $S_n$  falls neither into the category of Hoeffding nor Bernstein, we would need a more general concentration inequalities: the McDiarmid's inequality.<sup>2</sup>

**Theorem 2.5.1** (McDiarmid's inequality). Let  $X_1, \ldots, X_n$  be i.i.d. random variables on  $\chi$ , and let  $f: \chi^n \to \mathbb{R}$  satisfying the bounded difference property, i.e.,

$$\sup_{x_1, \dots, x_n, x_i'} |f(x_1, \dots, x_n) - f(x_1, \dots, x_i', \dots, x_n)| \le c_i$$

for all i. Then for any t > 0,

$$\mathbb{P}(f(X_1,\ldots,X_n) - \mathbb{E}\left[f(X_1,\ldots,X_n)\right] \ge t) \le \exp\left(\frac{-2t^2}{\sum_i c_i^2}\right).$$

The same bound holds for the left tail.

**Remark.** The qualitative statement for McDiarmid's inequality is that "a random variable that depends on the influence of many independent random variables but not too many on any one of them concentrates".

**Proof.** Typically,  $\sum_i c_i = O(1)$  concentration will happen if  $\sum_i c_i^2 = o(1)$ . For example, if each  $c_i = O(1/n)$ , then concentration happens but not when all  $c_i = 0$  except one of them is 1.

Remark. McDiarmid's inequality is a generalization of Hoffding's inequality.

**Proof.** Let

$$f(x_1, \dots, x_n) = \frac{1}{n}(x_1 + \dots + x_n).$$

When  $X_i$ 's are independent and  $X_i \in [a_i, b_i]$  for all i, it's easy to observe that when we change the i<sup>th</sup> argument of f, the value of f can change at most by  $(b_i - a_i)/n$ , i.e., McDiarmid's inequality is satisfied with  $c_i := (b_i - a_i)/n$ , plugging in, we get back Hoffding's inequality.

With McDiarmid's inequality, we can check that the following holds for bounded function classes F:

$$|S_n(x_1,\ldots,x_n)-S_n(x_1,\ldots,x_i',\ldots,x_n)|\leq \frac{2B}{n}=:c_i.$$

Then from McDiarmid's inequality, for any t > 0,

$$\mathbb{P}(S_n \ge \mathbb{E}\left[S_n\right] + t) \le \exp\left(\frac{-nt^2}{2B^2}\right) =: \delta,$$

or equivalently,  $S_n \leq \mathbb{E}\left[S_n\right] + B\sqrt{\frac{2}{n}\log\frac{1}{\delta}}$  with probability at least  $1 - \delta$ .

**Note.**  $B\sqrt{\frac{2}{n}\log\frac{1}{\delta}}$  is a lower order term, i.e.,  $\mathbb{E}\left[S_n\right]$  dominates it.

Proof. Since<sup>a</sup>

$$O(B) \ge \mathbb{E}\left[S_n\right] \ge \mathbb{E}\left[\left|\frac{1}{n}\sum_{i=1}^n f(x_i) - \mathbb{E}\left[f(X)\right]\right|\right] = O\left(\sqrt{\frac{\operatorname{Var}\left[f(X_1)\right]}{n}}\right) \approx O\left(\frac{1}{\sqrt{n}}\right).$$

<sup>a</sup>This upper bound is pretty weak, and we will eventually work on getting better bounds.

All these imply that it's enough to bound  $\mathbb{E}[S_n]$ .

\*

<sup>&</sup>lt;sup>2</sup>It's also known as the bounded difference inequality.

## Lecture 5: Proof of McDiarmid's Inequality

We should note that the usual proof of McDiarmid inequality involves martingale decomposition and 1 Sep. 9:00 Azuma-Hoeffding inequality, a generalization of Hoffding's inequality for martingale difference sequence.

**Definition 2.5.1** (Martingale difference sequence). A martingale difference sequence is a sequence of random variables  $\Delta_1, \ldots$  such that  $\mathbb{E} [\Delta_i \mid \Delta_{i-1}] = 0$  for all i.

However, we will not go with this route; instead, we prove something weaker but tricker.<sup>3</sup>

**Note.** The condition  $\sup_{x_1,\ldots,x_n,x_i'} |f(x_1,\ldots,x_n)-f(x_1,\ldots,x_i',\ldots,x_n)| \le c_i$  is equivalent to

$$|f(x_1,\ldots,x_n)-f(z_1,\ldots,z_n)| \le \sum_{i=1}^n c_i \mathbb{1}_{x_i \ne z_i}.$$

Now, we need one last lemma to prove McDiarmid inequality.

**Lemma 2.5.2.** For all  $x \neq y \in \mathbb{R}$ ,

$$\frac{e^x-e^y}{x-y} \leq \frac{e^x+e^y}{2} \Rightarrow |e^x-e^y| \leq |x-y| \left(\frac{e^x+e^y}{2}\right).$$

**Proof.** Since

$$\frac{e^x - e^y}{x - y} = \int_0^1 e^{sx + (1 - s)y} \, ds = \frac{1}{x - y} \int_x^y e^t \, dt$$

where we let t = sx + (1 - s)y. On the other hand, due to convexity, we also have

$$\frac{e^x - e^y}{x - y} = \int_0^1 e^{sx + (1 - s)y} \, ds \le \int_0^1 s \cdot e^x + (1 - s)e^y \, ds = \frac{e^x + e^y}{2}.$$

We're now ready.

**Proof of Theorem 2.5.1.** Firstly, we note that it's equivalent to show that  $f(X_1, \ldots, X_n) - \mathbb{E}[f] \in$  $\operatorname{Subg}(\sum_i c_i^2/4)$ . Without loss of generality, let  $\mathbb{E}[f] = 0$ , and we want to show that

$$\mathbb{E}\left[e^{\lambda(f(X) - \mathbb{E}[f])}\right] \le e^{\frac{\lambda^2 \sum_i c_i}{8}} \Leftrightarrow M(\lambda) = \mathbb{E}\left[e^{\lambda f(X)}\right] \le \exp\left(\frac{\lambda^2 \left(\sum_i c_i^2\right)}{8}\right) \Leftrightarrow \log M(\lambda) \le \lambda^2 \frac{\sum_i c_i^2}{8}.$$

Observe that since both sides of the inequality is 0 at  $\lambda = 0$ , it's enough to show

$$\frac{\mathrm{d}\log M(\lambda)}{\mathrm{d}\lambda} = \frac{M'(\lambda)}{M(\lambda)} \le \lambda \cdot \frac{\sum_{i} c_{i}^{2}}{4}$$

Let 
$$\mathbb{X}=(X_1,\ldots,X_n)$$
, and  $\mathbb{X}'\stackrel{\mathrm{i.i.d.}}{\sim}\mathbb{X}$  be the i.i.d. copy of  $\mathbb{X}$ . Then define the following.   
Notation.  $\mathbb{X}^{(i)}\coloneqq (X_1',\ldots,X_i',X_{i+1},\ldots,X_n)$  and  $\mathbb{X}^{[i]}\coloneqq (X_1,\ldots,X_{i-1},X_i',X_{i+1},\ldots,X_n)$ .

<sup>&</sup>lt;sup>3</sup>In fact, what we're going to prove is not even a weaker version: we prove something weaker while we really need the original (stronger) statement to hold.

Note that this implies  $\mathbb{X}^{(0)} = \mathbb{X}$  and  $\mathbb{X}^{(n)} = \mathbb{X}'$ . Then, we can show that

$$M'(\lambda) = \mathbb{E}\left[f(\mathbb{X})e^{\lambda f(\mathbb{X})}\right]$$
 As  $\mathbb{E}\left[f\right] = 0$  and  $\mathbb{X}$ ,  $\mathbb{X}'$  are independent 
$$= \mathbb{E}\left[\left(f(\mathbb{X}) - f(\mathbb{X}')\right)e^{\lambda f(\mathbb{X})}\right]$$
 
$$= \mathbb{E}\left[\sum_{i=1}^{n} (f(\mathbb{X}^{(i-1)}) - f(\mathbb{X}^{(i)})) \cdot e^{\lambda f(\mathbb{X})}\right]$$

if  $i^{\text{th}}$  position of  $\mathbb{X}$  and  $\mathbb{X}'$  are swapped, then for the new data  $\mathbb{X}^{(i-1)}$  and  $\mathbb{X}^{(i)}$  will also be swapped,

We note the following.

Note. The above proof doesn't even show a weaker version of McDiarmid's inequality.

**Proof.** While in the proof, we need to show

$$\frac{\mathrm{d}\log M(\lambda)}{\mathrm{d}\lambda} = \frac{M'(\lambda)}{M(\lambda)} \le \lambda \cdot \frac{\sum_{i} c_{i}^{2}}{4},$$

we only show

$$\frac{\mathrm{d} \log M(\lambda)}{\mathrm{d} \lambda} = \frac{M'(\lambda)}{M(\lambda)} \leq \lambda \cdot \frac{\sum_i c_i^2}{2}.$$

#### 2.5.3 Applications of McDiarmid's Inequality

#### U-Statistics

Let  $g: \mathbb{R}^2 \to \mathbb{R}$  be a symmetric function, and let  $X_1, \ldots, X_n \overset{\text{i.i.d.}}{\sim} \mathbb{P}$ . Consider

$$U(X) = \frac{1}{\binom{n}{2}} \sum_{j < k} g(X_j, X_k).$$

Here're some examples of g.

**Example.** 
$$g(x, y) = (x - y)^2$$
.

Example. 
$$g(x,y) = |x - y|$$
.

**Example** (Wilcoxm's ranksom test).  $g(x,y) = \mathbb{1}_{x_1+x_2>0}$ .

We're interested to know about  $\mathbb{E}[g(X_1, X_2)]$ . Assume g is bounded by B, then

$$U(X) - U(X^{[k]}) \le \frac{1}{\binom{n}{2}}(n-1)2B \le \frac{4B}{n},$$

\*

implying

$$\mathbb{P}(U - \mathbb{E}\left[U\right] \ge t) \le e^{-\frac{nt^2}{8b^2}}$$

from McDiarmid's inequality with  $c_i := 2B$ .

#### Beyond McDiarmid's Inequality

Let's see some more advanced inequalities. In many cases, we want variance to be small. While

$$\operatorname{Var}\left[X_{1}+\cdots+X_{n}\right] \leq \sum_{i=1}^{n} \operatorname{Var}\left[X_{i}\right],$$

to have an inequality for a non-linear function, we have the following.

**Theorem 2.5.2** (Efron-Stein inequality). Let  $X_1, \ldots, X_n$  be independent random variables, and  $X'_1, \ldots, X'_n$  be i.i.d. copies of  $X_i$ 's. Then

$$\operatorname{Var}\left[f(\mathbb{X})\right] \leq \frac{1}{2} \sum_{i=1}^{n} \mathbb{E}\left[\left(f(\mathbb{X}) - f(\mathbb{X}^{[i]})\right)^{2}\right].$$

**Note.** We see that since  $\operatorname{Var}[X] = \frac{1}{2}\mathbb{E}\left[(X - X')^2\right]$ , by letting  $f(X_1, \dots, X_n) = \sum_i X_i$ , if f satisfies bounded condition, then  $\operatorname{Var}[f] \leq \frac{1}{2}\sum_i c_i^2$ .

Now, recall that by using McDiarmid's inequality, we can show that for  $\mathscr{F} \ni f$  being B-bounded,

$$S_n \le \mathbb{E}\left[S_n\right] + B\sqrt{\frac{2}{n}\log\frac{1}{\delta}}$$

with probability at least  $1 - \delta$ . However, what if the variance Var[f(X)] is small, but the maximum spread (B) is very large? In this case, we would want to replace B in the inequality by Var[f(X)].

**Notation** (Empirical process notation). Let  $\mathbb{P}f = \mathbb{E}[f]$  and  $\mathbb{P}_n f = \sum_i f(X_i)/n$ .

This is achieved by the following, although it's much harder to prove [BLM13, §12].

**Theorem 2.5.3** (Talagrand's concentration inequality). Let  $\mathscr{F}$  is B-bounded, and  $S_n = \sup_{f \in \mathscr{F}} |\mathbb{P}_n f - \mathbb{P}_n f|$ . Then

$$S_n \le c \cdot \mathbb{E}\left[S_n\right] + c\sqrt{\frac{\sup_{f \in \mathscr{F}} \operatorname{Var}\left[f(X_1)\right]}{n}} \log \frac{1}{\alpha} + c \cdot \frac{B}{n} \log \frac{1}{\alpha}$$

with probability at least  $1 - \alpha$ .

**Remark.** We might encounter an explicit situation where Talagrand's concentration is more profitable to use than bounded differences inequality later in the course.

# Chapter 3

# Expected Supremum of Empirical Process

#### Lecture 6: A Glance at Statistical Learning Theory

#### 3.1 Goodness of Fit Testing

6 Sep. 9:00

Let's first see another motivation on studying uniform law of large numbers, i.e., the *goodness of fit* testing. Given  $X_1, \ldots, X_n \overset{\text{i.i.d.}}{\sim} \mathbb{P}$ , we want to distinguish between  $H_0 \colon \mathbb{P} = \mathbb{P}_0$  and  $H_1 \colon \mathbb{P} \neq \mathbb{P}_0$ .

Many tests are possible. One approach could be the Kolmogorov-Smirnov test: assume F is the CDF of  $\mathbb{P}_0$ , then consider the Kolmogorov-Smirnov statistics:

**Definition 3.1.1** (Kolmogorov-Smirnov statistics). The Kolmogorov-Smirnov statistics for a distribution  $\mathbb{P}$  is defined as

$$D_n = \sup_{t \in \mathbb{R}} |F_n(t) - F(t)|$$

where  $F_n(t)$  and F is the empirical CDF and the CDF of  $\mathbb{P}$ , respectively.

From Glivenko-Cantelli theorem,  $D_n \to 0$  under  $H_0$ , and  $D_n$  should not converge to 0, under some alternative. Assuming continuity of F, Kolmogorov showed that

- (a) the distribution  $D_n$  does not depend on F;
- (b)  $D_n = O_p(1/\sqrt{n});$
- (c)  $\sqrt{n}D_n \to \sup_{t \in [0,1]} |B(t)|$  where B(t) is the Broweian bridge on [0,1].
- (d)  $\mathbb{P}(\sqrt{n}D_n \le 2.4) \approx 0.999973$ .

We'll take a non-asymptotic approach to this problem, i.e., we may not get such sharp constants.

## 3.2 Statistical Learning

#### 3.2.1 Empirical Risk Minimization

Consider the following problem.

**Problem 3.2.1** (Empirical risk minimization). Let  $S = \{(x_1, y_1), \dots, (x_n, y_n)\}$  be n i.i.d. copies of  $(X, Y) \in \chi \times \mathcal{Y} \subseteq \mathbb{R}^d \times \mathbb{R}$  with distribution  $\mathbb{P} = \mathbb{P}_X \times \mathbb{P}_{Y|X}$ . Given a loss function  $\ell \colon \mathcal{Y} \times \mathcal{Y} \to \mathbb{R}$  and a function class  $\mathscr{F} = \{f \colon \chi \to \mathcal{Y}\}$ , the *empirical risk minimization* is

$$\hat{f} \in \underset{f \in \mathscr{F}}{\operatorname{arg\,min}} \frac{1}{n} \sum_{i=1}^{n} \ell(f(x_i), y_i).$$

**Example.**  $\mathscr{F}$  can be the set of neural networks, decision trees, linear functions.

**Example** (Linear regression). Consider  $\chi = \mathbb{R}^d$  and  $\mathcal{Y} = \mathbb{R}$ , with  $\mathscr{F} = \{x \to w^\top x \colon w \in \mathbb{R}^d\}$  and  $\ell(a,b) = (a-b)^2$ .

**Example** (Linear classification). Consider  $\chi = \mathbb{R}^d$  and  $\mathcal{Y} = \{0,1\}$ , with  $\mathscr{F} = \{x \to (\operatorname{sgn}(w^\top x) + 1)/2 \colon w \in B_2^d\}$  where  $B_2^d$  is the unit ball in d-dimension, and  $\ell(a,b) = \mathbb{1}_{a\neq b}$ .

We also define the following.

**Definition.** Consider the set-up of empirical risk minimization.

**Definition 3.2.1** (Expected loss). The expected loss  $^a$  of  $f \in \mathscr{F}$  is defined as

$$L(f) = \mathbb{E}_{(X|Y) \sim \mathbb{P}} \left[ \ell(f(X), Y) \right].$$

**Definition 3.2.2** (Empirical loss). The *empirical loss* is defined as

$$\hat{L}(f) = \frac{1}{n} \sum_{i=1}^{n} \ell(f(x_i), y_i).$$

The main question in statistical learning is that, what is an upper-bound on the expected loss of ERM? If we plug in  $\hat{f}$  instead of f, this is asking the test error of  $\hat{f}$ .

To be specific,  $\hat{f}$  is basically a function of training data S, but when we look at

$$L(\hat{f}) = \mathbb{E}_{(X,Y)} \left[ \ell(\hat{f}(x), Y) \right],$$

it is the expectation of future data points, i.e., it becomes a random variable, which is a function of S.

**Lemma 3.2.1.** For any  $\mathscr{F}$ , the ERM  $\hat{f}$  satisfies

$$\mathbb{E}[L(\hat{f})] - \inf_{f \in \mathscr{F}} L(f) \le \mathbb{E}\left[\sup_{f \in \mathscr{F}} \left(L(f) - \hat{L}(f)\right)\right].$$

**Proof.** Let  $f^* = \inf_{f \in \mathscr{F}} L(f)$ . Then

$$L(\hat{f}) - L(f^*) = [L(\hat{f}) - \hat{L}(\hat{f})] + [\hat{L}(\hat{f}) - \hat{L}(f^*)] + [\hat{L}(f^*) - L(f^*)].$$

We see that

- $\hat{L}(\hat{f}) \hat{L}(f^*) \le 0$  by definition;
- $\hat{L}(f^*) L(f^*) = 0$  in expectation since  $f^*$  is fixed,
- We can't say  $\mathbb{E}[L(\hat{f}) \hat{L}(\hat{f})] = 0$  since  $\hat{f}$  is also random.

Combine all these, we have

$$\mathbb{E}[L(\hat{f})] - \inf_{f \in \mathscr{F}} L(f) = \mathbb{E}[L(\hat{f}) - L(f^*)] \le \mathbb{E}[L(\hat{f}) - \hat{L}(\hat{f})] \le \mathbb{E}\left[\sup_{f \in \mathscr{F}} \left(L(f) - \hat{L}(f)\right)\right].$$

 $<sup>^</sup>a$ Also called *population loss* and *test error*.

**Note.** Let us decode what Lemma 3.2.1 is claiming.

- Since L(f) is the population error of f and  $\hat{L}(f)$  is the empirical loss of f,  $\sup_{f \in \mathscr{F}} \left( L(f) \hat{L}(f) \right)$  is the supremum of an empirical process.
- For the left-hand side, it represents the expected loss of  $\hat{f}$  and the best possible out-of-sample error.<sup>a</sup> This is often called the excess risk.

**Notation** (Excess risk).  $\mathbb{E}[L(\hat{f})] - \inf_{f \in \mathscr{F}} L(f)$  is often called the excess risk of an ERM.

Remark. For "curved" loss function like square loss, supremum can be further "localized".

Remark. The bound in Lemma 3.2.1 can be vacuumed for now, e.g., for linear regression.

**Example** (1-D classification with thresholds). Let  $\ell(a,b) = \mathbb{1}_{a\neq b} = a + (1-2a)b$  for  $a,b \in \{0,1\}$ . Then consider a=y and b=f(x),

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\left(L(f)-\hat{L}(f)\right)\right] = \mathbb{E}\left[\sup_{f\in\mathscr{F}}\left(\mathbb{E}\left[Y+(1-2Y)f(X)\right] - \frac{1}{n}\sum_{i=1}^{n}\left(y_i+(1-2y_i)f(x_i)\right)\right)\right],$$

which can be viewed essentially as a the empirical process on the function f instead of  $\ell$ ,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\left(\mathbb{E}\left[f(X)\right]-\frac{1}{n}\sum_{i=1}^{n}f(x_{i})\right)\right].$$

For 1-D case, assume that  $\mathscr{F} = \{x \mapsto \mathbb{1}_{x \leq \theta} : \theta \in \mathbb{R}\}$ , then

$$\mathbb{E}\left[\sup_{\theta\in\mathbb{R}}\left(\mathbb{P}(X\leq\theta)-\frac{1}{n}\sum_{i=1}^{n}\mathbb{1}_{x_i\leq\theta}\right)\right]=\mathbb{E}\left[\sup_{\theta\in\mathbb{R}}(F(\theta)-F_n(\theta))\right],$$

i.e.,  $P(X \leq \theta)$  is the CDF of the marginal distribution of X,  $F(\theta)$ , and  $\frac{1}{n} \sum_{i=1}^{n} \mathbb{1}_{x_i \leq \theta}$  is the empirical CDF  $F_n(\theta)$ . Therefore, we go back to the same problem we introduced in the beginning of the chapter, i.e., the Kolmogorov-Smirnov statistics.

Let the term  $\mathbb{P}(X \leq \theta) - \frac{1}{n} \sum_{i=1}^{n} \mathbb{1}_{x_i \leq \theta}$  to be a random variable  $U_{\theta}$ . One problem here is, we have infinitely many random variables, and they are also correlated with each other quite a lot. So how does this supremum behave?

Since each  $U_{\theta}$  is at most 1, for any  $\theta$ , i.e.,  $\sup U_{\theta} \leq 1$ . So the worst case here is 1, and probably the best case is  $O(1/\sqrt{n})$ .

# Lecture 7: Bracketing and Symmetrization

Our main empirical process is so far  $\mathbb{E}\left[\sup_{f\in\mathscr{F}}\mathbb{P}_nf-\mathbb{P}f\right]$ . Let's first focus on the 1-D thresholds 8 Sep. 9:00 classification, i.e., we want to bound the supremum

$$\mathbb{E}\left[\sup_{\theta\in\mathbb{R}}\left|\frac{1}{n}\sum_{i=1}^{n}\mathbb{1}_{x_i\leq\theta}-\mathbb{P}(X\leq\theta)\right|\right].$$

There are 2 approaches to bound this supremum: bracketing and symmetrization.

#### 3.2.2 Bracketing

The main idea of bracketing is the following.

<sup>&</sup>lt;sup>a</sup>Or the best possible prediction error of  $\mathscr{F}$ .

<sup>&</sup>lt;sup>a</sup>Since  $Y - \sum_{i} y_i/n$  is independent of f, so let's drop it; and 1 - 2Y is the sign, so can be dropped essentially.

Intuition. Reduce an infinite number of random variables to finite, which will be more manageable.

Assume that  $\mathbb{P}$  is continuous, and consider a finite set  $\{\theta_i\}_{i=0}^{N+1}$  with  $\theta_0 = -\infty$ ,  $\theta_{N+1} = \infty$ , such that they correspond to quantile of  $\mathbb{P}$ , i.e.,

$$\mathbb{P}(\theta_i \le X \le \theta_{i+1}) = \frac{1}{N+1}.$$

Given a  $\theta$ , X will lie in between two adjacent  $\theta_i$ 's in the sequence. Denote the upper-bound as  $u(\theta)$  and the lower-bound as  $\ell(\theta)$  for this  $\theta$ , then

$$\begin{split} \mathbb{P}(X \leq \theta) - \frac{1}{n} \sum_{i=1}^{n} \mathbb{1}_{x_i \leq \theta} \leq \mathbb{P}(X \leq u(\theta)) - \frac{1}{n} \sum_{i=1}^{n} \mathbb{1}_{x_i \leq \ell(\theta)} \\ \leq \mathbb{E} \left[ \mathbb{1}_{X \leq u(\theta)} \right] - \frac{1}{n} \sum_{i=1}^{n} \mathbb{1}_{x_i \leq \ell(\theta)} \\ \leq \mathbb{E} \left[ \mathbb{1}_{X \leq \ell(\theta)} \right] - \frac{1}{n} \sum_{i=1}^{n} \mathbb{1}_{x_i \leq \ell(\theta)} + \mathbb{P}(\ell(\theta) \leq X \leq u(\theta)) \\ \leq \mathbb{E} \left[ \mathbb{1}_{X \leq \ell(\theta)} \right] - \frac{1}{n} \sum_{i=1}^{n} \mathbb{1}_{x_i \leq \ell(\theta)} + \frac{1}{N+1} \end{split}$$

if we take the supremum over  $\ell(\theta) \in \mathbb{R}$  instead of  $\theta$ ,

$$\leq \frac{1}{N+1} + \mathbb{E}\left[\max_{0 \leq j \leq N} \mathbb{E}\left[\mathbb{1}_{X \leq \theta_j}\right] - \frac{1}{n} \sum_{i=1}^n \mathbb{1}_{x_i \leq \theta_j}\right]. \tag{3.1}$$

To further bound Equation 3.1, recall the following.

As previously seen. If  $X_i \sim \operatorname{Subg}(\sigma^2)$  independent,  $\sum_i a_i X_i \sim \operatorname{Subg}((\sum_i a_i^2)\sigma^2)$  from Lemma 2.3.3.

**Remark.** Let  $a_i = 1/n$ , we see that  $\mathbb{E}\left[\mathbb{1}_{X \leq \theta_j}\right] - \frac{1}{n} \sum_{i=1}^n \mathbb{1}_{x_i \leq \theta_j} \in \text{Subg}(1/n)$ .

<sup>a</sup>Since it's bounded between 0 and 1.

Finally, recall what we have proved in the homework.

**Lemma 3.2.2.** Let  $X_1, \ldots, X_n \sim \operatorname{Subg}(\sigma^2)$ , a then  $\mathbb{E}[\max_i X_i] \leq \sqrt{2\sigma^2 \log n}$ .

<sup>a</sup>Not necessary independent.

Then, we can show the final bound.

**Proposition 3.2.1** (Bracketing). Let  $x_1, \ldots, x_n \overset{\text{i.i.d.}}{\sim} \mathbb{P}$ , and  $\mathscr{F} = \{\mathbb{1}_{X \leq \theta} : \theta \in \mathbb{R}\}$ . Then

$$\mathbb{E}_X \left[ \sup_{f \in \mathscr{F}} \left( \mathbb{P}(X \leq \theta) - \frac{1}{n} \sum_{i=1}^n \mathbbm{1}_{x_i \leq \theta} \right) \right] = O\left(\sqrt{\frac{\log n}{n}}\right).$$

**Proof.** From Lemma 3.2.2, since we have (N+1) random variables with variance factor 1/n, by choosing N+1 := n, a Equation 3.1 can be further bounded by

$$\sqrt{\frac{2\log(N+1)}{n}} + \frac{1}{N+1} = O\left(\sqrt{\frac{\log n}{n}}\right).$$

<sup>a</sup>Recall that n is the sample size, so we can choose the corresponding n to meet the requirement.

#### 3.2.3 Symmetrization

Another technique called symmetrization, which is essentially stated in the following lemma.

**Lemma 3.2.3** (Symmetrization). Given a function class  $\mathscr{F} = \{f : \chi \to \mathcal{Y}\}$  and  $X_1, \ldots, X_n \overset{\text{i.i.d.}}{\sim} \mathbb{P}$ , and  $\epsilon_1, \ldots, \epsilon_n$  be i.i.d. Rademacher random variables. Then

$$\max\left(\mathbb{E}\left[\sup_{f\in\mathscr{F}}\mathbb{P}_nf-\mathbb{P}f\right],\mathbb{E}\left[\sup_{f\in\mathscr{F}}\mathbb{P}f-\mathbb{P}_nf\right]\right)\leq 2\mathbb{E}_{\epsilon_i,X_i}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^n\epsilon_if(X_i)\right].$$

In particular,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\left|\mathbb{P}_nf-\mathbb{P}f\right|\right] \leq 2\mathbb{E}_{\epsilon_i,X_i}\left[\sup_{f\in\mathscr{F}}\left|\frac{1}{n}\sum_{i=1}^n\epsilon_if(X_i)\right|\right].$$

**Proof.** Let  $X_i'$ 's be i.i.d. copies of  $X_i$ 's for all i. Since adding a sign  $\epsilon_i$  won't change the expectation,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\mathbb{E}\left[f(X)\right] - \frac{1}{n}\sum_{i=1}^{n}f(X_{i})\right] = \mathbb{E}\left[\sup_{f\in\mathscr{F}}\mathbb{E}_{X_{i}'}\left[\frac{1}{n}\sum_{i=1}^{n}f(X_{i}') - \frac{1}{n}\sum_{i=1}^{n}f(X_{i})\right]\right]$$

$$\leq \mathbb{E}_{X_{i}}\left[\mathbb{E}_{X_{i}'}\left[\sup_{f\in\mathbb{F}}\frac{1}{n}\sum_{i=1}^{n}(f(X_{i}') - f(X_{i}))\right]\right]$$

$$= \mathbb{E}_{X_{i},X_{i}',\epsilon_{i}}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^{n}(f(X_{i}') - f(X_{i}))\epsilon_{i}\right]$$

$$\leq \mathbb{E}_{X_{i}',\epsilon_{i}}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^{n}f(X_{i}')\epsilon_{i}\right] + \mathbb{E}_{X_{i},\epsilon_{i}}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^{n}f(X_{i})\epsilon_{i}\right]$$

$$= 2\mathbb{E}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^{n}\epsilon_{i}f(X_{i})\right].$$

<sup>a</sup>Since the distributions of  $f(X_i') - \sum_i f(X_i)$  and  $f(X_i) - \sum_i f(X_i')$  are the same.

**Intuition.** If we condition on  $X_i$ 's, the bound can be seen as linear combination of Rademacher random variables. Thus, we can refer to properties of sub-Gaussian random variables.

The upper-bound deserves a special name.

**Definition 3.2.3** (Rademacher complexity). Let  $X_i \overset{\text{i.i.d.}}{\sim} \mathbb{P}$  be independent and  $\epsilon_i$  be i.i.d. Rademacher random variables. The *Rademacher complexity* of a function class  $\mathscr{F}$  w.r.t.  $\mathbb{P}$  is

$$R_n(\mathscr{F}) := \mathbb{E}_{\epsilon_i, X_i} \left[ \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^n \epsilon_i f(X_i) \right| \right].$$

On the other hand, the opposite direction of symmetrization lemma also holds.

**Lemma 3.2.4.** Given a function class  $\mathscr{F} = \{f \colon \chi \to \mathcal{Y}\}$  and  $X_1, \ldots, X_n \overset{\text{i.i.d.}}{\sim} \mathbb{P}$ , and  $\epsilon_1, \ldots, \epsilon_n$  be i.i.d. Rademacher random variables. Then

$$\mathbb{E}_{X_i,\epsilon_i} \left[ \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^n \epsilon_i f(X_i) \right| \right] \le 2\mathbb{E} \left[ \sup_{f \in \mathscr{F}} |\mathbb{P}_n f - \mathbb{P} f| \right] + \frac{1}{\sqrt{n}} \sup_{f \in \mathscr{F}} |\mathbb{P} f|.$$

**Proof.** This technique is so-called desymmetrization: Consider

$$\mathbb{E}_{\epsilon_{i},X_{i}} \left[ \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^{n} \epsilon_{i} f(X_{i}) \right| \right] \\
\leq \mathbb{E}_{\epsilon_{i},X_{i}} \left[ \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^{n} \epsilon_{i} (f(X_{i}) - \mathbb{E}\left[f(X)\right]) \right| \right] + \mathbb{E}_{\epsilon_{i}} \left[ \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^{n} \epsilon_{i} \mathbb{E}\left[f(X)\right] \right| \right] \\
= \mathbb{E}_{\epsilon_{i},X_{i},X'_{i}} \left[ \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^{n} \epsilon_{i} (f(X_{i}) - \mathbb{E}\left[f(X'_{i})\right]) \right| \right] + \mathbb{E} \left[ \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^{n} \epsilon_{i} \mathbb{E}_{\epsilon_{i}} \left[f(X_{i})\right] \right| \right].$$

The first term can be further bounded by

$$\mathbb{E}_{\epsilon_{i},X_{i},X_{i}'}\left[\sup_{f\in\mathscr{F}}\left|\frac{1}{n}\sum_{i=1}^{n}\epsilon_{i}(f(X_{i})-\mathbb{E}\left[f(X_{i}')\right])\right|\right] \leq \mathbb{E}_{\epsilon_{i},X_{i},X_{i}'}\left[\sup_{f\in\mathscr{F}}\left|\frac{1}{n}\sum_{i=1}^{n}\epsilon_{i}(f(X_{i})-f(X_{i}'))\right|\right]$$

$$=\mathbb{E}\left[\sup_{f\in\mathscr{F}}\left|\frac{1}{n}\sum_{i=1}^{\infty}(f(X_{i})-f(X_{i}'))\right|\right]$$

$$=\mathbb{E}\left[\sup_{f\in\mathscr{F}}\left|\frac{1}{n}\sum_{i=1}^{n}\left(f(X_{i})-f(X_{i}')+(\mathbb{E}\left[f\right]-\mathbb{E}\left[f\right])\right)\right|\right]$$

$$=2\mathbb{E}\left[\sup_{f\in\mathscr{F}}\left|\mathbb{P}_{n}f-\mathbb{P}f\right|\right],$$

and the second term ca be bounded by

$$\mathbb{E}_{\epsilon_i} \left[ \sup_{f \in \mathscr{F}} \left| \frac{1}{n} \sum_{i=1}^n \epsilon_i \mathbb{E}\left[ f(X) \right] \right| \right] \leq \sup_{f \in \mathscr{F}} |\mathbb{E}\left[ f(X) \right]| \cdot \mathbb{E}\left[ \left| \frac{1}{n} \sum_{i=1}^n \epsilon_i \right| \right] \leq \frac{1}{\sqrt{n}} \sup_{f \in \mathscr{F}} |\mathbb{P}f|$$

where  $\mathbb{E}\left[\left|\frac{1}{n}\sum_{i}\epsilon_{i}\right|\right] \leq \frac{c}{\sqrt{n}}$  with c=1. Combine them together, we have the final result.

# Lecture 8: Symmetrization on 1-D Threshold Classification

Analogous to the Rademacher complexity defined for a function class w.r.t. P, we can define it on a set. 11 Sep. 9:00

**Definition 3.2.4** (Rademacher width). Let  $\epsilon_i$  be i.i.d. Rademacher random variables. Then the Rademacher width<sup>a</sup> of a set  $A \subseteq \mathbb{R}^n$  is defined as

$$R_n(A) = \mathbb{E}_{\epsilon_i} \left[ \sup_{a \in A} \frac{1}{n} \sum_{i=1}^n \epsilon_i a_i \right].$$

<sup>a</sup>Also called Rademacher average.

Notation. People sometimes just say "Rademacher complexity" for Rademacher width.

Now, applying the symmetrization lemma to  $\mathscr{F} = \{\mathbb{1}_{X \leq \theta} : \theta \in \mathbb{R}\}$ , we have the following result that is comparable to Proposition 3.2.1.

**Proposition 3.2.2.** Let 
$$x_1, \ldots, x_n \overset{\text{i.i.d.}}{\sim} \mathbb{P}$$
, and  $\mathscr{F} = \{\mathbb{1}_{x \leq \theta} : \theta \in \mathbb{R}\}$ . Then

$$\mathbb{E}_X \left[ \sup_{f \in \mathscr{F}} \left( \mathbb{P}(X \leq \theta) - \frac{1}{n} \sum_{i=1}^n \mathbb{1}_{x_i \leq \theta} \right) \right] = O\left(\sqrt{\frac{\log n}{n}}\right).$$

**Proof.** From the symmetrization lemma,

$$\mathbb{E}_{X,x_i} \left[ \sup_{\theta \in \mathbb{R}} \left( \mathbb{P}(X \le \theta) - \frac{1}{n} \sum_{i=1}^n \mathbb{1}_{x_i \le \theta} \right) \right] \le 2\mathbb{E}_{\epsilon_i,x_i} \left[ \sup_{\theta \in \mathbb{R}} \frac{1}{n} \sum_{i=1}^n \epsilon_i \mathbb{1}_{x_i \le \theta} \right] \quad \text{condition on } x_1, \dots, x_n$$

$$= 2\mathbb{E}_{x_i} \left[ \mathbb{E}_{\epsilon_i \mid x_i} \left[ \sup_{f \in \mathscr{F}} \frac{1}{n} \sum_{i=1}^n \epsilon_i \mathbb{1}_{x_i \le \theta} \middle| x_1, \dots, x_n \right] \right].$$

Let  $V_{\theta} := \frac{1}{n} \sum_{i=1}^{n} \epsilon_{i} \mathbb{1}_{x_{i} \leq \theta}$ , we see that there are only n+1 distinct  $V_{\theta}$ 's, and it's constant in the intervals  $\theta \in [X_{(k)}, X_{(k+1)})$  for  $k = 0, \ldots, n-1$  where  $X_{(k)}$  are the order statistics with  $X_{(0)} := -\infty$ . Now, define  $\theta_{k} := X_{(k)}$ , we can then write

$$\sup_{\theta \in \mathbb{R}} \frac{1}{n} \sum_{i=1}^{n} \epsilon_i \mathbb{1}_{x_i \le \theta} = \max_{k=0,\dots,n} \frac{1}{n} \sum_{i=1}^{n} \epsilon_i \mathbb{1}_{x_i \le \theta_k},$$

hence,

$$2\mathbb{E}_{x_i}\left[\mathbb{E}_{\epsilon_i|x_i}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^n\epsilon_i\,\mathbb{1}_{x_i\leq\theta}\bigg|x_1,\ldots,x_n\right]\right]=2\mathbb{E}_{x_i}\left[\mathbb{E}_{\epsilon_i|x_i}\left[\max_{k=0,\ldots,n}V_{\theta_k}\bigg|x_1,\ldots,x_n\right]\right]$$

with  $V_{\theta_k} \sim \text{Subg}(1/n)$  and Lemma 3.2.2,

$$\leq 2\mathbb{E}_{x_i}\left[\sqrt{\frac{2}{n}\log(n+1)}\right] = O\left(\sqrt{\frac{\log n}{n}}\right).$$

**Remark.** Looking back to the example of 1-D thresholds classification, we see that the excess risk of ERM is  $O(\sqrt{\log n/n})$ .

# 3.3 Vapnik-Chervonenkis Dimension

#### 3.3.1 Glivenko-Cantelli Class

From bracketing and symmetrization, we see that there are classes of functions such that

$$\sup_{f \in \mathscr{F}} |\mathbb{P}f - \mathbb{P}_n f| \to 0$$

as  $n \to \infty$ . They deserve their own name.

**Definition 3.3.1** (Glivenko-Cantelli). A function class  $\mathscr{F} = \{f : \chi \to \mathbb{R}\}$  is *Glivenko-Cantelli* w.r.t.  $\mathbb{P}$  if as  $n \to \infty$ ,

$$\sup_{f \in \mathscr{F}} |\mathbb{P}f - \mathbb{P}_n f| \to 0.$$

From bracketing and symmetrization, we know that  $\mathscr{F} = \{\mathbb{1}_{X \leq \theta} : \theta \in \mathbb{R}\}$  is Glivenko-Cantelli. Let's see some counterexamples.

**Example.** Let  $\chi = \mathbb{R}$ ,  $\mathscr{F} = \{\mathbb{1}_A : A \subseteq \chi, |A| < \infty\}$ , and  $\mathbb{P}$  be any continuous measure on  $\chi$ . Then  $\mathscr{F}$  is not Glivenko-Cantelli w.r.t.  $\mathbb{P}$ .

**Proof.** For  $f = \mathbbm{1}_A$ ,  $\mathbb{P} f = \mathbb{P}(X \in A) = 0$  since  $|A| < \infty$ . On the other hand, let  $A_0 = \{X_1, \dots, X_n\}$  be the observed empirical data,  $\mathbb{P}_n f = 1$ , i.e.,  $\sup_{f \in \mathscr{F}} |\mathbb{P} f - \mathbb{P}_n f| = 1$  for all  $n \in \mathbb{N}$ .

**Example.** Let  $\chi = \mathbb{R}$ ,  $\mathscr{F} = \{f : \chi \to \mathbb{R} \text{ bounded and continuous}\}$ , and  $\mathbb{P} = \mathcal{U}[0,1]$ . Then  $\mathscr{F}$  is not Glivenko-Cantelli.

**Proof.** Consider  $f(X_i) = 1$  for i = 1, ..., n and f = 0 elsewhere (continuously), a then we can make  $\int_0^1 f(t) dt < \delta$  for some  $\delta \in (0, 1)$ . This implies  $\sup_{f \in \mathscr{F}} |\mathbb{P}f - \mathbb{P}_n f| \ge 1 - \delta$  for all  $n \in \mathbb{N}$ .

#### 3.3.2 Vapnik-Chervonenkis Dimension

**Notation.** Let  $\mathscr{F}(x_1,\ldots,x_n) := \{(f(x_1),\ldots,f(x_n))\}_{f\in\mathscr{F}} \subseteq \mathbb{R}^n$ .

We can relate the Rademacher width of  $\mathscr{F}(X_1,\ldots,X_n)$  to the Rademacher complexity of  $\mathscr{F}$  since

$$\mathbb{E}_{X_i}\left[R_n(\mathscr{F}(X_1,\ldots,X_n))\right] = \mathbb{E}_{X_i,\epsilon_i}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^n\epsilon_i f(X_i)\right] = R_n(\mathscr{F}).$$

Moreover, we see that if  $\mathscr{F}(X_1,\ldots,X_n)$  is finite, by the same proof as in Proposition 3.2.2,

$$\mathbb{E}_{X_i}\left[R_n(\mathscr{F}(X_1,\ldots,X_n))\right] \le 2\sqrt{\frac{2\log|\mathscr{F}(X_1,\ldots,X_n)|}{n}}.$$

The up-shot is the following.

**Remark.** If  $|\mathscr{F}(X_1,\ldots,X_n)| \leq n^d$  for some  $d \in \mathbb{N}^+$ , then we again get an  $O(\sqrt{\log n/n})$  bound.

This is captured by the polynomial discrimination, where we're going to focus on boolean functions.

**Definition 3.3.2** (Polynomial discrimination). We say that a boolean function class  $\mathscr{F}$  on  $\chi$  has a polynomial discrimination if for all  $x_1, \ldots, x_n \in \chi$ ,  $|\mathscr{F}(x_1, \ldots, x_n)| \leq \mathsf{poly}(n)$ .

To characterize  $|\mathscr{F}(x_1,\ldots,x_n)|$ , we will look at the VC dimension of  $\mathscr{F}$ , which is related to the size of the discrimination of  $\mathscr{F}$  in a non-trivial way.

**Definition.** Let  $\mathscr{F}$  be a boolean function class on  $\chi$ .

**Definition 3.3.3** (Shatter). A finite set  $\{x_1, \ldots, x_D\} \subseteq \chi$  is *shattered* by  $\mathscr{F}$  if  $\mathscr{F}(x_1, \ldots, x_D) = \{0, 1\}^D$ .

 ${}^{a}$ We take the convention that  $\varnothing$  is always shattered.

**Definition 3.3.4** (Vapnik-Chervonenkis dimension). The VC dimension of  $\mathscr F$  on  $\chi$  is the maximum integer D such that there exists a size D finite set  $A\subseteq \chi$  shattered by  $\mathscr F$ .

Let's consider some examples on  $\chi = \mathbb{R}$ .

**Example.** The VC dimension of  $\mathscr{F} = \{\mathbb{1}_{X \leq \theta} : \theta \in \mathbb{R}\}$  is 1.

**Example.** The VC dimension of  $\mathscr{F} = \{\mathbb{1}_{[a,b]} : a, b \in \mathbb{R}\}$  is 2.

Let's look at one example with  $\chi = \mathbb{R}^2$ .

**Example.** The VC dimension of  $\mathscr{F} = \{\mathbb{1}_{[a,b] \times [c,d]} : a,b,c,d \in \mathbb{R}\}$  is 4.

#### Lecture 9: VC Dimension

Firstly, given VC dimension, we can upper-bound the size of the discrimination.

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<sup>&</sup>lt;sup>1</sup>This is why people overload  $R_n$  for both Rademacher width and Rademacher complexity.

**Lemma 3.3.1** (Sauer-Shelah lemma). Let  $\mathscr{F}$  be a boolean function class such that  $VC(\mathscr{F}) = d$ , then for every  $\{x_1, \ldots, x_n\} \subseteq \chi$  such that  $n \ge d$ ,

$$|\mathscr{F}(x_1,\ldots,x_n)| \le \binom{n}{0} + \binom{n}{1} + \cdots + \binom{n}{d} \le \left(\frac{en}{d}\right)^d.$$

To prove Sauer-Shelah lemma, we need Pajor's lemma.

**Lemma 3.3.2** (Pajor's lemma). Given a boolean function class  $\mathscr{F}$  on a finite set  $\Omega$ , then

$$|\mathscr{F}| < |\{S \subseteq \Omega \colon S \text{ shattered by } \mathscr{F}\}|.$$

**Proof.** We prove this by induction on n. For n=1 (base case), it holds trivially since

$$|\mathscr{F}| = 2 \le |\{S \subseteq \Omega \colon S \text{ shattered by } \mathscr{F}\}|.$$

Assume the statement holds for all  $\Omega$  such that  $|\Omega| = n$ . For  $|\Omega| = n + 1$ , we write  $\Omega = (\Omega \setminus \{x_0\}) \cup \{x_0\} =: \Omega_0 \cup \{x_0\}$  and let  $\mathscr{F}_0$  and  $\mathscr{F}_1$  be two boolean function classes defined on  $\Omega_0$  as

$$\mathscr{F}_0 = \{ f \in \mathscr{F} : f(x_0) = 0 \}, \quad \mathscr{F}_1 = \{ f \in \mathscr{F} : f(x_0) = 1 \}.$$

We further define  $S_{\mathscr{F}'}$  as  $S_{\mathscr{F}'} = \{S \subseteq \Omega' : S \text{ shattered by } \mathscr{F}'\}$  for any function class  $\mathscr{F}'$  defined on  $\Omega'$ . Then, by induction hypothesis,  $|\mathscr{F}_i| \leq |S_{\mathscr{F}_i}|$ , hence  $|\mathscr{F}| = |\mathscr{F}_0| + |\mathscr{F}_1| \leq |S_{\mathscr{F}_0}| + |S_{\mathscr{F}_1}|$ . Finally, we claim the following.

**Claim.**  $|S_{\mathscr{F}_0}| + |S_{\mathscr{F}_1}| \le |S_{\mathscr{F}}|$ .

**Proof.** Let  $S \subseteq \Omega_0$  shattered by both  $\mathscr{F}_0$  and  $\mathscr{F}_1$ , then S is shattered by  $\mathscr{F}$  too. Moreover, Observe that  $S \cup \{x_0\}$  is shattered by  $\mathscr{F}$  but not  $\mathscr{F}_i$   $(f(x_0))$  is fixed for  $f \in \mathscr{F}_i$ . Now, when

- S is shattered by only one of the  $\mathscr{F}_i$ 's: S contributes one unit both to  $|S_{\mathscr{F}}|$  and  $|S_{\mathscr{F}_i}|$ ;
- S is shattered by both  $\mathscr{F}_i$ 's, S and  $S \cup \{x_0\}$  are shattered by  $\mathscr{F}$ : S contributes two units to  $|S_{\mathscr{F}}|$  and one unit to both  $|S_{\mathscr{F}_i}|$ 's.

By counting, we're done (it's possible that S is shattered by  $\mathscr{F}$  but not  $\mathscr{F}_i$ 's, so  $\leq$ ).

This implies  $|\mathscr{F}| \leq |S_{\mathscr{F}}|$  for  $|\Omega| = n + 1$ , i.e., the induction is done.

We can then prove the Sauer-Shelah lemma.

**Proof of Lemma 3.3.1.** Let  $\Omega$  be a set of size n, then the number of subsets with size  $\leq d$  is  $\binom{n}{0} + \binom{n}{1} + \cdots + \binom{n}{d}$ , hence by the definition of VC dimension,

$$|\{S \subseteq \Omega \colon S \text{ shattered by } \mathscr{F}\}| \le \binom{n}{0} + \binom{n}{1} + \dots + \binom{n}{d}.$$

Then, as our motivation suggests, the same proof of Proposition 3.2.2 applies, giving the following.

**Proposition 3.3.1.** For any function class  $\mathscr{F}$ , if  $n \geq VC(\mathscr{F})$ , for some constant c,

$$R_n(\mathscr{F}) \le c\sqrt{\frac{\mathrm{VC}(\mathscr{F})}{n}\log\left(\frac{en}{\mathrm{VC}(\mathscr{F})}\right)}.$$

**Remark.** We see that Proposition 3.3.1 is independent of  $\mathbb{P}$ , i.e., the bounds still holds after taking  $\sup_{\mathbb{P}}$  on the left-hand side. However, if  $VC(\mathscr{F}) = \infty$ , then this "distribution-free" uniform convergence fails.

However, if we don't care about distribution-free property, we do have examples that the uniform convergence holds for a particular  $\mathbb{P}$  when  $VC(\mathscr{F}) = \infty$ .

**Example.** For  $\mathscr{F} = \{\mathbb{1}_A : \text{compact convex } A \subseteq [0,1]^d\}$ ,  $VC(\mathscr{F}) = \infty$ . If  $\mathbb{P}$  is continuous w.r.t. Lebesgue's measure, then the uniform law of large number still holds.

**Remark.** The  $\sqrt{\log n}$  factors in Proposition 3.3.1 is superfluous (Corollary 3.4.5).

**Example.** Let V be a vector space of real function on  $\chi$  with  $\dim(V) = D$ , and  $\mathscr{F} = \{\mathbb{1}_{f \geq 0} : f \in V\}$ . Then  $VC(\mathscr{F}) \leq D$ .

**Proof.** We want to show that for any  $\{x_1,\ldots,x_{D+1}\}$  can't be shattered. Let

$$T = \{ (f(x_1), \dots, f(x_{D+1})) : f \in V \},$$

which is a linear subspace of  $\mathbb{R}^{D+1}$  such that  $\dim(T) \leq D$ . Hence, there exists a non-zero  $y \in \mathbb{R}^{D+1}$  such that  $\sum_{i=1}^{D+1} y_i f(x_i) = 0$  for all  $f \in V$ . Now, without loss of generality, there exists an index k such that  $y_k > 0$ . If  $\mathscr{F}$  shatters  $\{x_1, \dots, x_{D+1}\}$ , then there exists  $f \in V$  such that

$$\begin{cases} f(x_i) < 0, & \forall i : y_i > 0; \\ f(x_i) \ge 0, & \forall i : y_i \le 0. \end{cases}$$

But then  $\sum_{i} y_i f(x_i) < 0$ , which is a contradiction.

**Example** (Half-space). For  $\mathscr{F} = \{1_H : \text{half space } H \subseteq \mathbb{R}^d\}, \text{ VC}(\mathscr{F}) = d+1.$ 

Although it seems like  $VC(\mathscr{F}) \approx \#$ parameters of  $\mathscr{F}$ ; however, it's not true in general.

**Example.** Consider  $\mathscr{F} = \{x \mapsto \mathbb{1}_{\sin tx \geq 0} : t \in \mathbb{R}^+\}$ , then  $VC(\mathscr{F}) = \infty$ .

## Lecture 10: Discretization of a Space

# 3.4 Metric Entropy Methods

15 Sep. 9:00

We have been focusing on boolean function class with finite VC dimension, and our goal now is to generalize beyond the boolean case. This can be done by discretizing of a space.

**Intuition** (Informal principle). We want to bound  $\mathbb{E}[\sup_{t \in T} X_t]$ . If  $\{X_t\}_{t \in T}$  is "sufficiently continuous", then  $\mathbb{E}[\sup_{t \in T} X_t]$  is governed by metric properties of T (metric entropy!).

**Definition 3.4.1** (Pseudo-metric). Given a space T, a function  $d: T \times T \to \mathbb{R}^+$  is a pseudo-metric if

- (a) d(x,x) = 0 for all  $x \in T$ ;
- (b) d(x,y) = d(y,x) for all  $x, y \in T$ ;
- (c)  $d(x,y) \le d(x,z) + d(y,z)$  for all  $x, y, z \in T$ .

**Note.** The motivation of looking at pseudo-metric instead of the usual metric is because, consider observed data  $x_1, \ldots, x_n$  at hands, the most natural distance might be

$$(f,g) \mapsto \sqrt{\frac{1}{n} \sum_{i=1}^{n} (f(x_i) - g(x_i))^2},$$

<sup>&</sup>lt;sup>a</sup>If d further satisfies that d(x,y) > 0 for all  $x \neq y$ , then it becomes a metric.

which is a pseudo-metric since f and g can agree only on  $x_i$ 's and vary elsewhere.

#### 3.4.1 Covering Number and Packing Number

Now, let (T,d) denote a pseudo-metric space in the remaining of this section, unless specified.

**Definition 3.4.2** ( $\epsilon$ -net). A set N is an  $\epsilon$ -net of (T,d) if for all  $t \in T$ , there exists  $\pi(t) \in N$  such that  $d(t,\pi(t)) \leq \epsilon$ .

**Definition 3.4.3** (Covering number). The  $\epsilon$ -covering number  $N(T,d,\epsilon)$  of (T,d) is defined as  $N(T,d,\epsilon) := \inf\{|N| : N \text{ is an } \epsilon\text{-net for } (T,d)\}.$ 

**Remark.** N is not necessary a subset of T for convenience. Furthermore, if  $N \nsubseteq T$ , one can construct another net N' such that  $N' \subseteq T$  and N' is a  $2\epsilon$ -net.

**Definition 3.4.4** (Totally bounded). (T,d) is totally bounded if for all  $\epsilon > 0$ ,  $N(T,d,\epsilon) < \infty$ .

**Definition 3.4.5** ( $\epsilon$ -packing). A set  $N \subseteq T$  is an  $\epsilon$ -packing of (T, d) if for all  $t \neq t'$  in N,  $d(t, t') > \epsilon$ .

**Definition 3.4.6** (Packing number). The  $\epsilon$ -pacing number  $M(T,d,\epsilon)$  of (T,d) is defined as  $M(T,d,\epsilon) = \sup\{|N| \colon N \text{ is an } \epsilon\text{-packing of } (T,d)\}.$ 

As the title suggests, we define the following metric properties, which is an essential notion helps us bound the expected empirical process supremum.

**Definition 3.4.7** (Metric entropy). The metric entropy of (T,d) is defined as  $\log M(T,d,\epsilon)$ .

The fact that we're using packing number  $M(T, d, \epsilon)$  when defining metric entropy is not relevant here due to the following.

**Lemma 3.4.1.** For any  $\epsilon > 0$ ,

$$M(T, d, 2\epsilon) \le N(T, d, \epsilon) \le M(T, d, \epsilon).$$

**Proof.** We show them one by one.

Claim.  $M(T, d, 2\epsilon) \leq N(T, d, \epsilon)$ .

**Proof.** Take  $\mathcal{M}$  to be a  $2\epsilon$ -packing and  $\mathcal{N}$  to be an  $\epsilon$ -net. Then for any  $t \in \mathcal{N}$ , consider  $B(t, \epsilon)$ . We see that there is at most one  $x \in \mathcal{M}$  such that  $d(t, x) \leq \epsilon$  since otherwise, if  $x, x' \in \mathcal{M}$  such that  $x \neq x'$  and  $d(t, x), d(t, x') \leq \epsilon$ , then  $d(x, x') \leq 2\epsilon$ , a contradiction to  $\mathcal{M}$ .

Claim.  $N(T, d, \epsilon) \leq M(T, d, \epsilon)$ .

**Proof.** Take  $\mathcal{M}$  to be a maximum  $\epsilon$ -packing, it suffices to show that  $\mathcal{M}$  is also an  $\epsilon$ -net, i.e., for all  $t \in T$ , there exists  $x \in \mathcal{M}$  such that  $d(x,t) \leq \epsilon$ . Suppose not, then  $d(t,x) > \epsilon$  for all  $x \in \mathcal{M}$ , i.e., we can add x to  $\mathcal{M}$ , contradiction.

For simplicity, we will use the following notations.

**Notation.** If (T,d) and  $\epsilon$  are clear from the context, we write  $N:=N(T,d,\epsilon)$  and  $M:=M(T,d,\epsilon)$ .

Turns out that there's a characterization of the packing number of the unit ball in euclidean space.

**Proposition 3.4.1.** Consider  $(\mathbb{R}^d, \|\cdot\|)$  where  $\|\cdot\|$  is any norm. Denote  $B = \{x : \|x\| \le 1\}$ , then for all  $\epsilon > 0$ ,

$$(1/\epsilon)^d \le M(B, \|\cdot\|, \epsilon) \le (1 + 2/\epsilon)^d.$$

**Proof.** For the lower-bound, we see that

$$N \operatorname{Vol}(\epsilon B) \ge \operatorname{Vol}(B) \Rightarrow N \epsilon^d \ge 1.$$

With  $N \leq M$  from Lemma 3.4.1, we get the lower-bound.

For the upper-bound, since  $\epsilon/2$  balls around points in M are disjoint, union of these  $\epsilon/2$  balls will lie in  $(1 + \epsilon/2)B$ . This implies

$$M \times \left(\frac{\epsilon}{2}\right)^d \times \operatorname{Vol}(B) \le \left(1 + \frac{\epsilon}{2}\right)^d \times \operatorname{Vol}(B) \Rightarrow M \le \left(1 + \frac{2}{\epsilon}\right)^d.$$

**Note.** From Proposition 3.4.1,  $\log M(\mathbb{R}^d, \|\cdot\|, \epsilon) \approx d \log 1/\epsilon$ .

#### 3.4.2 Hölder Smooth Functions

We are interested in looking at function spaces, and the following are the canonical smooth function classes studied in *nonparametric regression*.

**Definition 3.4.8** (Hölder smooth function class). Fix  $\alpha > 0$ , and  $\beta$  is the greatest integer  $< \alpha$ . Then the *Hölder smooth function class*  $S_{\alpha}$  is defined to be the class of functions on [0,1] such that

- (a) f continuous on [0, 1];
- (b) f is  $\beta$ -times differentiable;
- (c)  $|f^{(k)}| \le 1$  for all  $k = 0, \dots, \beta$ ;
- (d)  $|f^{(\beta)}(x) f^{(\beta)}(y)| \le |x y|^{\alpha \beta}$  for all  $x, y \in [0, 1]$ .

**Note.** When  $\alpha = 1$ ,  $S_{\alpha}$  is a class of 1-Lipschitz functions.

**Remark.** The Hölder smooth function classes are nested, so it's not surprising that the metric entropies decrease as  $\alpha$  increases.

Now, let  $d(f,g) = \sup_{x \in [0,1]} |f(x) - g(x)|$ , then  $(S_{\alpha}, d)$  is a pseudo-metric space.

**Theorem 3.4.1.** There exists  $c_1, c_2$  such that for all  $\epsilon > 0$ ,

$$\exp\left(c_2\epsilon^{-1/\alpha}\right) \le M(S_\alpha, d, \epsilon) \le \exp\left(c_1\epsilon^{-1/\alpha}\right).$$

**Proof sketch.** Here we illustrate the basic idea when  $\alpha = 1$ , i.e., the set of [0,1] valued 1-Lipschitz functions on [0,1]. We only sketch the proof of the upper-bound, since the lower-bound is similar.

Firstly, we partition both the domain and the range of f with small intervals with width  $\epsilon$ , resulting in  $1/\epsilon$  small intervals on both the x-axis and the y-axis.

Take any function  $f \in \mathcal{F}$ . We construct a piece-wise constant function  $\widetilde{f}$  which approximates f. On each small interval in the x-axis, we can define  $\widetilde{f}$  to be constant, taking value equal to the

\*

midpoint of the interval in the y-axis where the value of f at the left endpoint of this interval (in the x-axis) lies. Then, we have the following.

Claim.  $\sup_{x \in [0,1]} |f(x) - \widetilde{f}(x)| \le C\epsilon$ .

**Proof.** Since f cannot vary by more than  $\epsilon$  in any interval of length  $\epsilon$ .

Now, as we vary  $f \in \mathcal{F}$ , consider the following.

**Problem.** What is the number of distinct  $\widetilde{f}$  we can get?

A trivial bound is that, in each small interval on the x-axis, it takes one of the midpoints of the intervals on the y-axis and hence, the number of such functions is bounded by  $(\frac{1}{\epsilon})^{\frac{1}{\epsilon}}$ .

We can do slightly better. Note that, for the first interval, the number of possible values of  $\tilde{f}$  is  $\frac{1}{\epsilon}$ . However, after that, in the next interval, the value of  $\tilde{f}$  can only go up one interval, down one interval, or stay the same (due to 1-Lipschitzness of f), i.e., there are only 3 choices afterward for every interval, going from left to right, resulting an upper bound on the number of distinct  $\tilde{f}$  as

$$\frac{1}{\epsilon} 3^{\frac{1}{\epsilon} - 1} \le \exp\left(\frac{C}{\epsilon}\right).$$

**Remark.** Comparing Proposition 3.4.1 and Theorem 3.4.1, we see that the metric entropy is logarithmic in  $1/\epsilon$  versus some exponent of  $1/\epsilon$ . This is typically the hallmark of a parametric versus a nonparametric function class.

#### Lecture 11: Gaussian and Sub-Gaussian Process

#### 3.4.3 Sub-Gaussian Process

18 Sep. 9:00

As previously seen. Given a stochastic process  $\{X_t\}_{t\in T}$  with (T,d), we want to bound  $\mathbb{E}[\sup_{t\in T} X_t]$ .

Recall our informal principle, i.e., if  $\{X_t\}_{t\in T}$  is "sufficiently continuous" w.r.t. d, then  $\mathbb{E}\left[\sup_{t\in T}X_t\right]$  is governed by metric properties (e.g., metric entropy) of T. We start by considering the Gaussian process.

**Definition 3.4.9** (Gaussian process). A stochastic process  $\{X_t\}_{t\in T}$  is a Gaussian process if for any finite set of indices  $t_1,\ldots,t_k,\,(X_{t_1},\ldots,X_{t_k})\sim\mathcal{N}(0,\Sigma)$ .

Clearly, this is a very strong notion due to the following.

**Note.** For  $d(t,t') = \sqrt{\mathbb{E}[(X_t - X_{t'})^2]}$ , we have

$$\mathbb{E}\left[e^{\lambda(X_t - X_{t'})}\right] = e^{\lambda^2/2\mathbb{E}[X_t - X_{t'}]} = \exp\left(\frac{\lambda^2}{2}d^2(t, t')\right).$$

The following generalized process characterizes the concept of "sufficiently continuous".

**Definition 3.4.10** (Sub-Gaussian process). A stochastic process  $\{X_t\}_{t\in T}$  is a *sub-Gaussian process* w.r.t. d if  $X_t - X_s \sim \operatorname{Subg}(d^2(t,s))$ . Assume  $\mathbb{E}[X_t] = 0$  for all  $t \in T$ , then equivalently, for all  $t \neq s \in T$  and  $\lambda \in \mathbb{R}$ ,

$$\mathbb{E}\left[e^{\lambda(X_t - X_s)}\right] \le \exp\left(\frac{\lambda^2}{2}d^2(t, s)\right).$$

It's clear that the sub-Gaussian condition encodes a strong notion of continuity (in probability) of the stochastic process  $\{X_t\}_{t\in T}$  w.r.t. d.

**Example** (Gaussian process). We see that  $d(t,t') = \sqrt{\mathbb{E}[(X_t - X_{t'})^2]}$  is the naturally induced pseudometric such that a Gaussian process is sub-Gaussian.

Another interesting example is the following.

**Example** (Rademacher process). Consider the unnormalized Rademacher width of a set  $T \subseteq \mathbb{R}^n$ ,

$$R_n(T) = \mathbb{E}\left[\sup_{t \in \mathbb{R}^n} \sum_{i=1}^n \epsilon_i t_i\right].$$

Let  $X_t = \langle \epsilon, t \rangle$ , then from Lemma 2.3.3,  $X_t - X_{t'} = \langle \epsilon, t - t' \rangle \sim \text{Subg}(\|t - t'\|_2^2)$ , i.e.,  $X_t \sim \text{Subg}(\|t - t'\|_2^2)$ 

Inspired by the above example, one can also define the Gaussian width.

**Definition 3.4.11** (Gaussian width). Let  $g_i \overset{\text{i.i.d.}}{\sim} \mathcal{N}(0,1)$ . Then the Gaussian width of a set  $A \subseteq \mathbb{R}^n$  is defined as

$$GW_n(A) = \mathbb{E}\left[\sup_{a \in A} \sum_{i=1}^n \frac{1}{n} g_i a_i\right].$$

This means that the Rademacher process can be slightly modified as follows.

**Example.** Consider  $X_t = \langle g, t \rangle$  where g is a random Gaussian vector. We then have  $X_t \sim \text{Subg w.r.t. } \|\cdot\|_2$ .

**Theorem 3.4.2** (Gaussian width and Rademacher width are similar). For any  $n \geq 1$  and any set  $T \subset \mathbb{R}^n$ ,

$$R_n(T) \le GW_n(T) \le \sqrt{\log n} R_n(T).$$

Let's look at some examples of (unnormalized) Rademacher width.

**Example.**  $R(B_{\infty}^n) = n$ ,  $R(B_2^n) = \sqrt{n}$ , and  $R(B_1^n) = 1$ .

**Proof.** We see that

- for  $\ell_{\infty}$ , the supremum is achieved by matching signs of  $\epsilon$ , which gives  $R_n(B_{\infty}^n) = n$ ;
- for  $\ell_2$  the supremum is achieved by choosing  $t = \epsilon/\|\epsilon\|_2$ , then we get  $R(B_2^n) = \mathbb{E}\left[\epsilon\right] = \sqrt{n}$ ;
- for  $\ell_1$ , from Hölder's inequality,  $\langle \epsilon, t \rangle \leq \|\epsilon\|_{\infty} \|t\|_1 = 1$ .

\*

**Example** (Supremum of empirical process). Let  $\mathscr{F}$  be a class of functions bounded by 1. Let  $X_f = \sqrt{n}(\mathbb{P}_n f - \mathbb{P}f)$ , and consider  $\{X_f\}_{f \in \mathscr{F}}$ . Then,

$$X_f - X_g = \sqrt{n} \cdot \frac{1}{n} \sum_{i=1}^n (\underbrace{f(x_i) - g(x_i) - \mathbb{P}f - \mathbb{P}g}) \sim \operatorname{Subg} \left( 4\|f - g\|_{\infty}^2 \right),$$

hence  $\{X_f\}_{f\in\mathscr{F}}\sim \text{Subg w.r.t. }d(f,g)=2\|f-g\|_{\infty}.$ 

These are all simple sets. For an arbitrary set however, we need more general tools in order to compute the Rademacher width. Firstly, recall the following.

**Definition 3.4.12** (Diameter). The diameter of (T,d) is defined as  $Diam(T) = \sup_{t,t' \in T} d(t,t')$ .

#### 3.4.4 Single Scale Bound for Expected Supremum of Sub-Gaussian Process

We're going to see the most sophisticated tools in this course. We first see a preliminary version of which and generalize it later.

**Lemma 3.4.2** (Single scale bound). Let  $\{X_t\}_{t\in T}$  be a centered sub-Gaussian process on (T,d) w.r.t. d. Then

$$\mathbb{E}\left[\sup_{t\in T} X_t\right] \leq \inf_{\epsilon>0} \left(\mathbb{E}\left[\sup_{\substack{t,t'\in T:\\d(t,t')\leq \epsilon}} X_t - X_{t'}\right] + \mathrm{Diam}(T)\sqrt{2\log N(T,d,\epsilon)}\right).$$

**Proof.** We first note that  $\mathbb{E}\left[\sup_{t\in T}X_t\right] = \mathbb{E}\left[\sup_{t\in T}X_t - X_{t_0}\right]$  for some fixed  $t_0\in T$ . Now, take an  $\epsilon$ -net N with  $\pi(t)\in N$  denotes the point such that  $d(t,\pi(t))\leq \epsilon$ , then

$$\mathbb{E}\left[\sup_{t\in T} X_t - X_{t_0}\right] \leq \mathbb{E}\left[\sup_{t\in T} X_t - X_{\pi(t)}\right] + \mathbb{E}\left[\sup_{t\in T} X_{\pi(t)} - X_{t_0}\right]$$

Observe that  $X_{\pi(t)} - X_{t_0} \sim \text{Subg}(\text{Diam}^2(T))$ , then the second term is a finite maximum such that

$$\mathbb{E}\left[\sup_{t\in T} X_{\pi(t)} - X_{t_0}\right] \leq \sqrt{2\operatorname{Diam}^2(T)\log N(T,d,\epsilon)} = \operatorname{Diam}(T)\sqrt{2\log N(T,d,\epsilon)}$$

from Lemma 3.2.2. By rewriting the first term, we have

$$\mathbb{E}\left[\sup_{t\in T} X_t\right] \leq \inf_{\epsilon>0} \left(\mathbb{E}\left[\sup_{\substack{t,t'\in T:\\d(t,t')\leq \epsilon}} X_t - X_{t'}\right] + \operatorname{Diam}(T)\sqrt{2\log N(T,d,\epsilon)}\right).$$

**Notation** (Approximation error). The first term in the single scale bound is the approximation error.

We see that the first term in the <u>single scale bound</u> is still an infinite maximum, so it is not clear how to bound it. Typically, we have to do something crude here. There are some exceptions, though.

**Example.** For Rademacher processes, we have 
$$\mathbb{E}\left[\sup_{t,t'\in T: \|t-t'\|\leq \delta}\langle \epsilon, t-t'\rangle\right] \leq \|\epsilon\|\delta \leq \sqrt{n}\delta$$
.

**Remark.** As  $\epsilon$  decreases, the approximation error should get smaller and the finite maximum increases. Therefore, when we use the single scale bound we can then choose an optimum  $\epsilon$  to minimize the sum of these two.

Let's see some applications of single scale bound which show that the single scale bound may not get the optimal rate.

**Example.** Consider a finite set  $T = \{(0,0,\ldots,0),(1,0,\ldots,0),\ldots,(1,1,\ldots,1)\} \subseteq \mathbb{R}^n$ , i.e., the footprint of the boolean function class on  $\mathbb{R}$  given by  $\{\mathbb{1}_{x<\theta}\}_{\theta\in\mathbb{R}}$ . By Lemma 3.2.2,  $R_n(T) \leq \sqrt{n\log n}$ .

As previously seen. We claimed that  $\log n$  is superfluous.

We still can't remove the  $\sqrt{\log n}$ : from the single scale bound, with Diam $(T) = \sqrt{n}$ ,

$$R_n(T) \le \sqrt{n\epsilon} + \sqrt{n\sqrt{\log N(T, \|\cdot\|_2, \epsilon)}}.$$

To remove  $\log n$ ,  $\epsilon$  needs to be O(1) for the first term. But then  $\log N(T, \|\cdot\|_2, \epsilon) \to \infty$ , and we fail.

Now, let's revisit the previous example, and recall the following.

As previously seen. For a class of functions bounded by  $1, X_f \sim \operatorname{Subg}(2^2 \|f - g\|_{\infty}^2)$ , i.e.,  $X_f - X_g \leq c\sqrt{n}\|f - g\|_{\infty}$  almost surely.

**Example** (Empriical process supremum of  $S_1$ ). Consider  $X_f = \sqrt{n}(\mathbb{P}_n f - \mathbb{P} f)$  on  $\mathscr{F} = S_1$ , i.e., functions bounded by 1 and are 1-Lipschitz on [0,1]. So in particular,  $X_f - X_g \leq c\sqrt{n}\|f - g\|_{\infty}$ . From the single scale bound and Theorem 3.4.1,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}X_f\right] \le c\left(\sqrt{n}\epsilon + \sqrt{1/\epsilon}\right) = c(\sqrt{n}\cdot n^{-1/3})$$

by letting  $\epsilon = n^{-1/3}$  (where this bound is minimized), giving us

$$\mathbb{E}\left[\sup_{f\in S_1} \mathbb{P}_n f - \mathbb{P}f\right] \le \frac{c}{n^{1/3}}.$$

This is the first non-trivial bound we have shown besides boolean function classes.

However, observe that  $X_f - X_g \leq C\sqrt{n}\|f - g\|_{\infty}$  implies  $X_f - X_g \leq \|f - g\|_{\infty}$  in probability. The fact that we are stuck with the above almost surely bound and don't know how to incorporate this additional information, suggests that this bound is still not optimal.

**Remark.** The optimal bound for  $S_1$  is  $c/\sqrt{n}$ , i.e., the CLT rate.

It's perhaps surprising that for the class of functions  $S_1$ , we get the  $O(n^{-1/2})$  rate for the supremum of the empirical process, because even for a single function  $f \in S_1$ , we would still have got the  $O(n^{-1/2})$  rate. This is not always the case, though. For Lipschitz function defined on  $[0,1]^d$ , the rates are slower. We state this result without proof for now.

**Lemma 3.4.3.** Let  $S_{1,d}$  to be the set of 1-bounded 1-Lipschitz functions w.r.t. the Euclidean norm defined on  $[0,1]^d$ . Then there exists a universal constant C>0 such that

$$\mathbb{E}\left[\sup_{f \in S_{1,d}} \mathbb{P}_n f - \mathbb{P}f\right] \le \begin{cases} Cn^{-1/2}, & \text{if } d = 1; \\ Cn^{-1/2} \log n, & \text{if } d = 2; \\ Cn^{-1/d} \log n, & \text{if } d > 2. \end{cases}$$

These rates are tight and corresponding lower bounds are also known [Han16, Problem 5.11 (d)].

## Lecture 12: Chaining Method and Dudley's Entropy Bound

#### 3.4.5 Dudley's Entropy Bound

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To overcome the limitation of the single scale bound, we can repeatedly taking  $\epsilon$ -net, which is considered as a *multi-scale bound*. The theorem requires one technical assumption of the stochastic process.

**Definition 3.4.13** (Separable). We say that  $\{X_t\}_{t\in T}$  is a *separable* process if there exists a countable  $T_0\subseteq T$  such that (outside a null set) for all  $t\in T$ , there exists  $\{t_n\in T_0\}_n$  such that  $d(t_n,t)\to 0$  satisfying  $\lim_{n\to\infty} X_{t_n}=X_t$ .

It's clear that  $\sup_{t \in T_0} X_t = \sup_{t \in T} X_t$ . Moreover, this notion is consistent with the separability of a topological space<sup>2</sup> we saw in real analysis.

**Example** (Separable metric space). If (T, d) is separable (as a topological space),  $\{X_t\}$  has countable sample path almost surely, then  $\{X_t\}$  is separable.

<sup>&</sup>lt;sup>2</sup>A topological space is *separable* if it contains a countable dense subset.

Now, we can state the bound we want.

**Theorem 3.4.3** (Dudley's entropy bound). Let  $\{X_t\}_{t\in T}$  be a centered and separable sub-Gaussian process on (T,d) w.r.t. d. Then

$$\mathbb{E}\left[\sup_{t\in T} X_t\right] \le 6\sum_{k\in\mathbb{Z}} 2^{-k} \sqrt{\log N(T, d, 2^{-k})}.$$

**Proof.** Consider the case that  $|T| < \infty$  and  $|T| = \infty$ .

**Claim.** The result holds for  $|T| < \infty$ .

**Proof.** Let  $K_0$  be the largest integer such that  $2^{-K_0} \ge \text{Diam}(T)$ , and let  $K_1$  be the smallest integer such that  $0 < 2^{-K_1} < \min_{s \ne t \in T} d(s,t)$ . Then we let  $N_k$  be a  $2^{-k}$ -net of T such that

- $k = K_0$ :  $N_{K_0} = \{t_0\}$  is a  $2^{-K_0}$ -net of T for a fixed  $t_0 \in T$ .
- $k = K_1$ :  $N_{K_1} = T$  is a  $2^{-K_1}$ -net of T.

Write  $\pi_k(t)$  for the closest element in  $N_k$  to t, in particular,  $d(t, \pi_k(t)) \leq 2^{-k}$ . By writing

$$X_t = X_{\pi_{K_1}(t)} - X_{\pi_{K_0}(t)} = X_{\pi_{K_1}(t)} - X_{\pi_{K_1-}(t)} + X_{\pi_{k_{(1)}}(t)} - \dots + X_{\pi_{K_0+1}(t)} - X_{\pi_{K_0}(t)}$$

we have

$$\mathbb{E}\left[\sup_{t\in T} X_t\right] = \mathbb{E}\left[\sup_{t\in T} X_t - X_{t_0}\right]$$

$$= \mathbb{E}\left[\sup_{t\in T} \sum_{k=K_0+1}^{K_1} \left(X_{\pi_k(t)} - X_{\pi_{k-1}(t)}\right)\right] \le \sum_{k=K_0+1}^{K_1} \mathbb{E}\left[\sup_{t\in T} \left(X_{\pi_k(t)} - X_{\pi_{k-1}(t)}\right)\right].$$

Since the cardinality of  $\{X_{\pi_k(t)} - X_{\pi_{k-1}(t)}\}_{t \in T}$  is  $|N_k||N_{k-1}| \le |N_k|^2$ , with

$$X_{\pi_k(t)} - X_{\pi_{k-1}(t)} \sim \text{Subg}(d(\pi_k(t), \pi_{k-1}(t)))$$

where  $d(\pi_k(t), t) + d(t, \pi_{k-1}(t)) \le 2^{-k} + 2^{-(k+1)} \le 3 \cdot 2^{-k}$ . From Lemma 3.2.2, for each k,

$$\mathbb{E}\left[\sup_{t \in T} \left( X_{\pi_k(t)} - X_{\pi_{k-1}(t)} \right) \right] \le 3 \times 2^{-k} \sqrt{2\log|N_k|^2} = 6 \times 2^{-k} \sqrt{\log|N_k|^2}$$

**Claim.** The result holds for  $|T| = \infty$ .

**Proof.** From separability, there exists a countable  $T_0$  such that  $\mathbb{E}\left[\sup_{t\in T_0}X_t\right]=\mathbb{E}\left[\sup_{t\in T}X_t\right]$ . Let  $T_k$  be a countable approximation of  $T_0$ , then  $\sup_{t\in T_k}X_t\to\sup_{t\in T_0}X_t$  as  $k\to\infty$ , so

$$\mathbb{E}\left[\sup_{t\in T_k}X_t\right]\to\mathbb{E}\left[\sup_{t\in T_0}X_t\right]=\mathbb{E}\left[\sup_{t\in T}X_t\right] \text{ as } k\to\infty$$

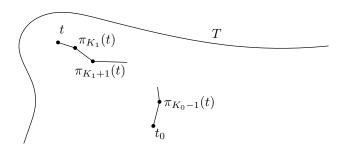
from monotone convergence theorem. Hence, it suffices to bound  $\mathbb{E}\left[\sup_{t\in T_k}X_t\right]$  instead of  $\mathbb{E}\left[\sup_{t\in T}X_t\right]$  for each k. As  $|T_k|<\infty$  and  $N(T_k,d,2^{-k})\leq N(T_0,d,2^{-k})$  for all k,

$$6\sum_{k\in\mathbb{Z}} 2^{-k} \sqrt{\log N(T_k, d, 2^{-k})} \le 6\sum_{k\in\mathbb{Z}} 2^{-k} \sqrt{\log N(T_0, d, 2^{-k})}.$$

\*

\*

**Note** (Chaining method). This method is called *chaining* since we're constructing a chain of  $X_{\pi_k(t)}$ , with smaller and smaller distances.



An alternative integral form of Dudley's entropy bound is given by the following.

Corollary 3.4.1 (Dudley integral entropy bound). Let  $\{X_t\}_{t\in T}$  be a centered and separable sub-Gaussian process on (T,d) w.r.t. d. Then

$$\mathbb{E}\left[\sup_{t \in T} X_t\right] \le 12 \int_0^{\mathrm{Diam}(T)} \sqrt{\log N(T, d, \epsilon)} \, \mathrm{d}\epsilon.$$

**Proof.** Observe that

$$\begin{split} \sum_{k \in \mathbb{Z}} 2^{-k} \sqrt{\log N(T, d, 2^{-k})} &= 2 \sum_{k \in \mathbb{Z}} \int_{2^{-k-1}}^{2^{-k}} \sqrt{\log N(T, d, 2^{-k})} \, \mathrm{d}\epsilon \\ &\leq 2 \sum_{k \in \mathbb{Z}} \int_{2^{-k-1}}^{2^{-k}} \sqrt{\log N(T, d, \epsilon)} \, \mathrm{d}\epsilon \qquad \qquad N(T, d, \epsilon) \nearrow \text{ as } \epsilon \searrow \\ &= 2 \int_0^\infty \sqrt{\log N(T, d, \epsilon)} \, \mathrm{d}\epsilon \\ &= 2 \int_0^{\mathrm{Diam}(T)} \sqrt{\log N(T, d, \epsilon)} \, \mathrm{d}\epsilon. \qquad \epsilon > \mathrm{Diam}(T), \, N(T, d, \epsilon) = 1 \end{split}$$

Now, we note that we have finally reached the optimal bound for  $S_1$ , solving the problems we saw in the previous example.

**Example** (Supremum of empirical process for  $S_1$ ). Consider the separable sub-Gaussian process  $X_f = \sqrt{n}(\mathbb{P}_n f - \mathbb{P}f)$  for  $\mathscr{F} = S_1$ . In particular,  $f, g \in \mathscr{F}$  are 1-Lipschitz on [0,1] satisfying  $|f|, |g| \leq 1$  and  $X_f - X_g \in \operatorname{Subg}(2^2 ||f - g||_{\infty}^2)$ . Since  $\operatorname{Diam}(\mathscr{F}) = 2$  and for all  $\epsilon < 1/2$ ,

$$N(S_1, \|\cdot\|_{\infty}, \epsilon) = \exp(c/\epsilon)$$

from Theorem 3.4.1. Then by the Dudley's integral entropy bound,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}X_f\right] \leq 12\int_0^2 \sqrt{\log N(S_1,\|\cdot\|_{\infty},\epsilon)}\,\mathrm{d}\epsilon = 12\int_0^2 \sqrt{\frac{c}{\epsilon}}\,\mathrm{d}\epsilon = 24\sqrt{2c} < O_n(1).$$

Dividing both sides by  $\sqrt{n}$ , we achieve the optimal rate  $\mathbb{E}\left[\sup_{f\in\mathscr{F}}(\mathbb{P}_nf-\mathbb{P}f)\right]=O(1/\sqrt{n})$ .

**Remark.** The Dudley's integral entropy bound for  $S_{\alpha}$  is also finite; while for function classes with covering number as  $\exp(c/\epsilon^2)$  is divergent.

## Lecture 13: More on Chaining

Let's see some alternate forms of Dudley's integral entropy bound. In the following, assume that  $\{X_t\}_{t\in T}$  22 Sep. 9:00 is a centered and separable sub-Gaussian process on (T,d) w.r.t. d.

**Corollary 3.4.2** (Difference form). The same bound as the Dudley's integral entropy bound holds for  $\mathbb{E}\left[\sup_{t\in T}|X_t-X_{t_0}|\right]$  and  $\mathbb{E}\left[\sup_{s,t\in T}|X_s-X_t|\right]$ .

**Proof.** This can be proved by the same chaining argument with triangle inequality.

Corollary 3.4.3 (High probability form). The high probability bound version holds:

$$\mathbb{P}\left(\sup_{s,t\in T}|X_s-X_t|\leq C\left(\int_0^\infty\sqrt{\log N(T,d,\epsilon)}\,\mathrm{d}\epsilon+u\,\mathrm{Diam}(T)\right)\right)\geq 1-2e^{-u^2}.$$

**Proof.** The proof is the same, where we first show the high probability bound for finite case.

Corollary 3.4.4 (Finite resolution form). The following generalizes the Dudley's integral entropy bound in the sense that  $\delta > 0$ :

$$\mathbb{E}\left[\sup_{t\in T} X_t\right] \le C\left(\left[\sup_{\substack{t,t'\in T\\d(t,t')\le \delta}} X_t - X_{t'}\right] + \int_{\delta}^{\infty} \sqrt{\log N(T,d,\epsilon)} \,\mathrm{d}\epsilon\right).$$

**Proof.** The proof is still the same, but instead we start with finite resolution.

The finite resolution version is useful since the entropy, integral can diverge, e.g., if  $\log N(t, d, \epsilon) = \Omega(1/\epsilon^2)$ . Moreover, this can be used to show Lemma 3.4.3.

Remark. We can moreover write

$$\begin{split} \mathbb{E}\left[\sup_{t \in T} X_t\right] &\leq C \int_0^{\mathrm{Diam}(T)} \sqrt{\log N(T,d,\epsilon)} \,\mathrm{d}\epsilon \\ &\leq C \left(\int_0^{\mathrm{Diam}(T)/2} \sqrt{\log N(T,d,\epsilon)} \,\mathrm{d}\epsilon + \int_{\mathrm{Diam}(T)/2}^{\mathrm{Diam}(T)} \sqrt{\log N(T,d,\epsilon)} \,\mathrm{d}\epsilon\right) \\ &\leq 2C \int_0^{\mathrm{Diam}(T)/2} \sqrt{\log N(T,d,\epsilon)} \,\mathrm{d}\epsilon. \end{split}$$

#### 3.4.6 Uniform Entropy Integral Bound

Let's discuss some limitation of the Dudley's integral entropy bound. First, recall the following.

As previously seen. In the example of the optimal rate for  $S_1$ , whenever

$$\int_0^\infty \sqrt{\log N(T, d, \epsilon)} \, \mathrm{d}\epsilon < \infty \Rightarrow \mathbb{E} \left[ \sup_f \mathbb{P}_n f - \mathbb{P} f \right] \le c / \sqrt{n}.$$

Note that we're doing chaining w.r.t.  $\|\cdot\|_{\infty}$  on  $\mathscr{F}$  so far. To see its limitation, consider again the boolean function classes  $\mathscr{F}$  and let  $X_f = \sqrt{n}(\mathbb{P}_n f - \mathbb{P} f) \sim \operatorname{Subg}(c^2 \|\cdot\|_{\infty}^2)$ . From Proposition 3.3.1,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\mathbb{P}_n f - \mathbb{P}f\right] \leq \sqrt{\frac{\mathrm{VC}(\mathscr{F})\log n}{n}}.$$

Now, observe that for any  $f \neq g$  in  $\mathscr{F}$ ,  $||f - g||_{\infty} = 1$ . This implies that by taking  $\epsilon \in (0,1)$ ,

$$N(\mathscr{F}, \|\cdot\|_{\infty}, \epsilon) = |\mathscr{F}| = \infty$$

for any interesting case, e.g.,  $\mathscr{F} = \{\mathbbm{1}_{x \leq \theta}\}_{\theta \in \mathbb{R}}$ , i.e., chaining w.r.t.  $\|\cdot\|_{\infty}$  only gives a vacuous bound.

**Intuition.** To fix this, we can use the idea of the symmetrization.

Firstly, given some observed i.i.d. data  $x_1, \ldots, x_n$ , recall the following.

As previously seen. By conditioning on the data  $x_1, \ldots, x_n$ , symmetrization shows that

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\sqrt{n}(\mathbb{P}_nf-\mathbb{P}f)\right] \leq \frac{2}{\sqrt{n}}R_n(\mathscr{F}) = 2\mathbb{E}_{x,\epsilon}\left[\sup_{f\in\mathscr{F}}\frac{1}{\sqrt{n}}\sum_{i=1}^n\epsilon_i f(x_i)\right].$$

Specifically, we want to look at  $\mathbb{E}_x \left[ \mathbb{E}_{\epsilon} \left[ \sup_{f \in \mathscr{F}} X_f \right] \right]$  and compute the Rademacher width. Let  $X_f = \frac{1}{\sqrt{n}} \sum_{i=1}^n \epsilon_i f(x_i)$ , we have

$$X_f - X_g = \frac{1}{\sqrt{n}} \sum_{i=1}^n \epsilon_i (f(x_i) - g(x_i)) \sim \text{Subg}(\|(f(x_i))_i - (g(x_i))_i\|_2^2) = \text{Subg}\left(\frac{1}{n} \sum_{i=1}^n (f(x_i) - g(x_i))^2\right),$$

where  $(f(x_i))_i = (f(x_1), \dots, f(x_n))$ . Hence,  $X_f \sim \text{Subg}\left(\frac{1}{n}\sum_i (f(x_i) - g(x_i))^2\right)$ .

**Note.** We're already doing better since  $\sqrt{\frac{1}{n}\sum_{i}(f(x_i)-g(x_i))^2} \leq \|f-g\|_{\infty}$ .

We see that  $\sqrt{\frac{1}{n}\sum_{i}(f(x_i)-g(x_i))^2}$  is similar to  $||f-g||_2$ , but just on the empirical measure (with i.i.d. data  $x_i$ 's). Hence, consider the following notation.

**Notation.** Let  $L_2(\mathbb{P}_n)$  denote the metric w.r.t.  $\mathbb{P}_n^a$  such that

$$L_2^2(\mathbb{P}_n)(f,g) := \frac{1}{n} \sum_{i=1}^n (f(x_i) - g(x_i))^2.$$

In our new notation,  $X_f \sim \operatorname{Subg}(L_2(\mathbb{P}_n))$ . Now, we can do the chaining argument on  $L_2(\mathbb{P}_n)$  and get

$$\mathbb{E}\left[\sup \sqrt{n}(\mathbb{P}_n f - \mathbb{P}f)\right] \le C \int_0^{\mathrm{Diam}(\mathscr{F})} \sqrt{\log N(\mathscr{F}, L_2(\mathbb{P}_n), \epsilon)} \, \mathrm{d}\epsilon,$$

where

$$\operatorname{Diam}(\mathscr{F}) = \sup_{f,g} L_2(\mathbb{P}_n)(f,g) = \frac{1}{n} \sum_{i=1}^n \left( f(x_i) - g(x_i) \right)^2 \le \sup_{f \in \mathscr{F}} \sqrt{\mathbb{P}_n f^2},$$

hence we have

$$\mathbb{E}\left[\sup \sqrt{n}(\mathbb{P}_n f - \mathbb{P}f)\right] \le C \cdot \mathbb{E}_x \left[ \int_0^{\sup \sqrt{\mathbb{P}_n f^2}} \sqrt{\log N(\mathscr{F}, L_2(\mathbb{P}_n), \epsilon)} \, \mathrm{d}\epsilon \right].$$

However, there's a problem.

**Problem.**  $L_2(\mathbb{P}_n)$  is a "random" metric, so  $N(\mathscr{F}, L_2(\mathbb{P}_n), \epsilon)$  is hard to compute.

**Answer.** To resolve this, we take the supremum over all measures  $\mu$  supported on  $\chi$ , i.e.,

$$C\mathbb{E}_x \left[ \int_0^{\sup_{f \in \mathscr{F}} \sqrt{\mathbb{P}_n f^2}} \sqrt{\log N(\mathscr{F}, L_2(\mathbb{P}_n), \epsilon)} \, \mathrm{d}\epsilon \right] \le C\mathbb{E}_x \left[ \int_0^{\sup_{f \in \mathscr{F}} \sqrt{\mathbb{P}_n f^2}} \sqrt{\sup_{\mu} \log N(\mathscr{F}, L_2(\mu), \epsilon)} \, \mathrm{d}\epsilon \right].$$

This might seem very bad, but actually it's not since  $L_2(\mu) < L_{\infty}$ . Specifically, to bound this supremum over all measures, consider the following.

<sup>&</sup>lt;sup>a</sup>Formally,  $\mathbb{P}_n$  is the empirical measure uniform on  $\{x_i\}_{i=1}^n$ .

**Definition 3.4.14** (Koltchinskii-Pollard entropy). The Koltchinskii-Pollard entropy of  $\mathscr{F}$  is defined as  $\sup_{\mu} \log N(\mathscr{F}, L_2(\mu), \epsilon).$ 

**Example.** For boolean function classes,  $\sup_{f} \sqrt{\mathbb{P}_n f^2} \leq 1$ .

We then have the following for the boolean function classes.

Intuition (Main bound). Let  $\mathscr{F}$  be a boolean function class, then since  $\sup_{f \in \mathscr{F}} \sqrt{\mathbb{P}_n f^2} \leq 1$ ,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\sqrt{n}|\mathbb{P}_nf-\mathbb{P}f|\right]\leq C\mathbb{E}_x\left[\int_0^1\sqrt{\sup_{\mu}\log N(\mathscr{F},L_2(\mu),\epsilon)}\,\mathrm{d}\epsilon\right].$$

More generally, if we have  $F \geq f$  (called envelope) for all  $f \in \mathscr{F}$  such that  $\mathbb{P}F^2 < \infty$ , this holds.

**Problem.** How can we compute the Koltchinskii-Pollard entropy?

Answer. We can use some notions of combinatorial dimension (e.g., VC dimension) upper-bounds the Koltchinskii-Pollard entropy such that

$$\sup_{\mu} N(\mathscr{F}, L_2(\mu), \epsilon) \le (c_1/\epsilon)^{c_2 \times \text{VC}(\mathscr{F})} \approx \epsilon^{-d}$$

for d being "dimension" (parametric).

**Remark.** This implies a  $\sqrt{\text{VC}(\mathcal{F})/n}$  rate (without a log term!) for  $\mathbb{E}[\sup(\mathbb{P}_n f - \mathbb{P}f)]$ .

## Lecture 14: Uniform Entropy Integral Bound

As previously seen. Motivated by the fact that boolean function classes are not totally bounded w.r.t.  $\ell_{\infty}$ ,  $||f - g||_{\infty} = 1$ , we're trying to establish the bound on

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\sqrt{n}(\mathbb{P}_nf-\mathbb{P}f)\right] \leq 2\mathbb{E}_x\left[\mathbb{E}_{\epsilon}\left[\sup_{f\in\mathscr{F}}\frac{1}{\sqrt{n}}\sum_{i=1}^n\epsilon_if(x_i)\right]\right] = \frac{1}{\sqrt{n}}R_n(\mathscr{F})$$

where the inner expectation is just the Rademacher width  $R_n(\{f(x_1),\ldots,f(x_n)\}_{f\in\mathscr{F}})$ . Let  $X_f=$  $\langle \epsilon, f \rangle / \sqrt{n}$ , then  $\{X_f\}_{f \in \mathscr{F}}$  is sub-Gaussian w.r.t.  $L_2(\mathbb{P}_n)$ .

The obstacle we're facing is the lack of control of  $\sup_{f \in \mathscr{F}} \sqrt{\mathbb{P}f^2}$ . Consider the following notion.

**Definition 3.4.15** (Envelope). A non-negative valued function  $F: \chi \to [0, \infty]$  is an *envelope* for  $\mathscr{F}$ if  $\sup_{f \in \mathscr{F}} |f(x)| \leq F(x)$  for all  $x \in \chi$ .

**Example.** For boolean function classes, F(x) = 1 is an envelope.

**Remark.** Let F be an envelope of  $\mathscr{F}$ , then  $\sup_{f\in\mathscr{F}}\sqrt{\mathbb{P}_nf^2}\leq\sqrt{\mathbb{P}_nF^2}$ , as we want.

With this new notion, we can state the main bound we want, i.e., the uniform entropy integral bound.

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<sup>&</sup>lt;sup>a</sup>We can also abuse the notation  $R_n(\mathscr{F})$  by  $R_n(\{f(x_1)/\sqrt{n},\ldots,f(x_n)/\sqrt{n}\}_{f\in\mathscr{F}})$ , but let's not do this. <sup>b</sup>Recall that  $L_2^2(\mathbb{P}_n)(f,g)=\frac{1}{n}\sum_{i=1}^n \left(f(X_i)-g(X_i)\right)^2$ , compared to  $L_2(\mathbb{P})(f,g)=\int \left(f(x)-g(x)\right)^2 d\mathbb{P}$ .

**Theorem 3.4.4** (Uniform entropy integral bound). Given a function class  $\mathscr{F}$  and an envelope F of  $\mathscr{F}$  such that  $\mathbb{P}F^2 < \infty$ , then for  $x_1, \ldots, x_n \overset{\text{i.i.d.}}{\sim} \mathbb{P}$ ,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\sqrt{n}|\mathbb{P}_nf-\mathbb{P}f|\right]\leq \frac{2}{\sqrt{n}}R_n(\mathscr{F})\leq C\|F\|_{L_2(\mathbb{P})}\int_0^1\sqrt{\log\sup_{\mu}N(\mathscr{F},L_2(\mu),\epsilon\sqrt{\mu F^2})}\,\mathrm{d}\epsilon.$$

**Proof.** Summarizing what we have established, we have

$$\begin{split} \mathbb{E}\left[\sup_{f\in\mathscr{F}}\sqrt{n}(\mathbb{P}_nf-\mathbb{P}f)\right] &\leq 2\mathbb{E}_x\left[\mathbb{E}_{\epsilon}\left[\sup_{f\in\mathscr{F}}\frac{1}{\sqrt{n}}\sum_{i=1}^n\epsilon_if(x_i)\right]\right] & \text{symmetrization} \\ &= \frac{2}{\sqrt{n}}R_n(\mathscr{F}) \\ &\leq C\mathbb{E}_x\left[\int_0^{\sup_{f,g\in\mathscr{F}}\frac{L_2(\mathbb{P}_n)(f,g)}{2}}\sqrt{\log N(\mathscr{F},L_2(\mathbb{P}_n),\epsilon)}\,\mathrm{d}\epsilon\right] & \mathrm{Dudley's\ bound} \\ &\leq C\mathbb{E}_x\left[\int_0^{\sup_{f\in\mathscr{F}}\sqrt{\mathbb{P}_nf^2}}\sqrt{\log N(\mathscr{F},L_2(\mathbb{P}_n),\epsilon)}\,\mathrm{d}\epsilon\right] \\ &\leq C\mathbb{E}_x\left[\int_0^{\sqrt{\mathbb{P}_nF^2}}\sqrt{\log N(\mathscr{F},L_2(\mathbb{P}_n),\epsilon)}\,\mathrm{d}\epsilon\right] \\ &= C\mathbb{E}_x\left[\sqrt{\mathbb{P}_nF^2}\int_0^1\sqrt{\log N(\mathscr{F},L_2(\mathbb{P}_n),\epsilon\sqrt{\mathbb{P}_nF^2}})\,\mathrm{d}\epsilon\right] & \epsilon\leftarrow\sqrt{\mathbb{P}_nF^2}\epsilon \\ &\leq C\mathbb{E}_x\left[\sqrt{\mathbb{P}_nF^2}\int_0^1\sqrt{\sup_{\mu}\log N(\mathscr{F},L_2(\mu),\epsilon\sqrt{\mu F^2})}\,\mathrm{d}\epsilon\right] \\ &\leq C\left[\int_0^1\sqrt{\sup_{\mu}\log N(\mathscr{F},L_2(\mu),\epsilon\sqrt{\mu F^2})}\,\mathrm{d}\epsilon\right] \mathbb{E}_x\left[\sqrt{\mathbb{P}_nF^2}\right] \\ &\leq C\|F\|_{L_2(\mathbb{P})}\left[\int_0^1\sqrt{\sup_{\mu}\log N(\mathscr{F},L_2(\mu),\epsilon\sqrt{\mu F^2})}\,\mathrm{d}\epsilon\right]. \end{split}$$

**Notation.** We sometimes denote  $\int_0^1 \sqrt{\log \sup_{\mu} N(\mathscr{F}, L_2(\mu), \epsilon \sqrt{\mu F^2})} \, d\epsilon$  by  $\boldsymbol{J}(F, \mathscr{F})$ .

 $\mathscr{F}$  needs not to be bounded, instead, what we really need is an envelope:

**Remark.** If we apply the above bound to a bounded function class then we do not really need the notion of an envelope. The assumption of an envelope is slightly more general than assuming boundedness.

**Remark.** The Koltchinskii-Pollard entropy integral is free from  $\mathbb{P}$ , while  $||F||_{L_2(\mathbb{P})}$  depends on  $\mathbb{P}$ . Thus, the overall bound is distribution free, as has been all of our bounds so far, if the function class is bounded.

**Example.** For boolean function classes,  $||F||_{L_2(\mathbb{P})} \leq 1$ , so the bound is uniform over  $\mathbb{P}$ .

#### 3.4.7 Uniform $L_2$ Entropy is Bounded by Combinatorial Dimension

We're now ready to revisit the problem we asked in the previous lecture:

**Problem.** How can be bound the uniform  $L_2$ -entropy, i.e., Koltchinskii-Pollard entropy?

**Answer.** For boolean function classes, Koltchinskii-Pollard entropy can be bounded in terms of VC dimension; for non-boolean function classes, an extended notion of VC dimension is needed.

**Theorem 3.4.5** (Dudley). Let  $\mathscr{F}$  be a boolean function class, then there exist absolute constants  $c_1, c_2$  such that for all  $0 < \epsilon < 1$ ,

$$\sup_{\mu} N(\mathscr{F}, L_2(\mu), \epsilon) \le \left(\frac{c_1}{\epsilon}\right)^{c_2 \operatorname{VC}(\mathscr{F})}$$

**Proof.** It suffices to bound the packing number instead of the covering number from Lemma 3.4.1. To upper-bound the packing number via  $d := VC(\mathscr{F})$ , first fix a probability measure  $\mu$  on  $\chi$ , consider a maximum  $\epsilon$ -packing  $M = \{f_1, \ldots, f_N\}$  of  $\mathscr{F}$  w.r.t.  $L_2(\mu)$ . Then for all  $i \neq j$ ,

$$\int (f_i - f_j)^2 d\mu = \mu(f_i \neq f_j) > \epsilon^2.$$

Then, sample K points  $W_1, \ldots, W_K$  i.i.d. from  $\mu$ , we want that all  $\{f_i\}_{i=1}^N$  to have different values on  $(w_1, \ldots, w_K)$ . Note that for  $i \neq j$ ,

$$\mu(f_i = f_i \text{ on } w_1, \dots, w_K) \le (1 - \epsilon^2)^K \le e^{-K\epsilon^2}$$

from  $(1-x)^k \le e^{-kx}$ . This implies

$$\mu(\exists \text{ at least one pair } i \neq j \text{ such that } f_i = f_j \text{ on } w_1, \dots, w_K) \leq \binom{N}{2} e^{-K\epsilon^2},$$

hence

$$\mu(\text{all the } f_i\text{'s are distinct on } w_1,\dots,w_k\ ) \geq 1-\binom{N}{2}e^{-K\epsilon^2} \geq \frac{1}{2}$$

by choosing  $K\epsilon^2 \approx 2 \log N$ . We conclude that there exists K points  $w_1, \ldots, w_K$  such that all the  $f_i$ 's are distinct on  $\{w_1, \ldots, w_K\}$ . From Sauer-Shelah lemma,

$$N = |\mathscr{F}(w_1, \dots, w_K)| \le \left(\frac{eK}{d}\right)^d = \left(\frac{2e \log N}{\epsilon^2 d}\right)^d.$$

We see that  $N \leq (\log N)^d$ . To further bound N, consider

$$N^{1/d} \le \frac{4e \log N}{2d\epsilon^2} = \frac{4e}{\epsilon^2} \log N^{1/2d} \le \frac{4e}{\epsilon^2} N^{1/2d}$$

from  $\log x \le x$ , hence  $N^{1/2d} \le 4e/\epsilon^2$ , or equivalently,

$$N \le (4e)^{2d} \left(\frac{1}{\epsilon}\right)^{4d} = \left(\frac{2\sqrt{e}}{\epsilon}\right)^{4d} =: \left(\frac{c_1}{\epsilon}\right)^{c_2d}.$$

**Remark.** For boolean function classes, while  $N(\mathcal{F}, L_{\infty}, \epsilon) = \infty$ , the above shows that

$$\sup_{\mu} N(\mathscr{F}, L_2(\mu), \epsilon) < \infty.$$

<sup>&</sup>lt;sup>a</sup>Note that we have only shown the case for  $K \geq \mathrm{VC}(\mathscr{F})$ . However, for  $K < \mathrm{VC}(\mathscr{F})$ , it's also easy to show.

<sup>&</sup>lt;sup>b</sup>We want to make the exponent a of  $\log N^a$  to be less than 1/d.

We can now finally generalize Proposition 3.3.1, where for boolean function classes, we have

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\sqrt{n}|\mathbb{P}_nf-\mathbb{P}f|\right]\leq c\sqrt{\mathrm{VC}(\mathscr{F})\log\frac{en}{\mathrm{VC}(\mathscr{F})}}.$$

**Corollary 3.4.5.** Let  $\mathscr{F}$  be a boolean function class, for some constant C, we have

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\sqrt{n}|\mathbb{P}_nf-\mathbb{P}f|\right]\leq C\sqrt{\mathrm{VC}(\mathscr{F})}.$$

**Proof.** Applying uniform entropy integral bound, with  $||F||_{L_2(\mathbb{P})} = 1$  and Theorem 3.4.5,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\sqrt{n}|\mathbb{P}_n f - \mathbb{P}f|\right] \leq C\int_0^1 \sqrt{c_2 \operatorname{VC}(\mathscr{F})\log\frac{c_1}{\epsilon}} \,\mathrm{d}\epsilon \leq C' \sqrt{\operatorname{VC}(\mathscr{F})}\int_0^1 \log\frac{1}{\epsilon} \,\mathrm{d}\epsilon \leq C' \sqrt{\operatorname{VC}(\mathscr{F})}.$$

**Remark.** Compare Proposition 3.3.1 to Corollary 3.4.5, the extra  $\sqrt{\log n}$  factor in the bound which we have now got rid of thanks to chaining. The bound really holds for Rademacher complexity and as a consequence for suprema of empirical process.

**Note.** Consider the classification problem in the statistical learning setting, where  $\hat{f}$  is the ERM over a given boolean function class. Then the excess risk is also bounded by  $C\sqrt{\text{VC}(\mathcal{F})/n}$ .

**Proof.** From symmetrization, the excess risk can be bounded as

$$\mathbb{E}\left[L(\hat{f})\right] - \inf_{f \in \mathscr{F}} \mathbb{E}\left[L(\hat{f})\right] \leq \mathbb{E}\left[\sup_{f \in \mathscr{F}} \left(\mathbb{E}\left[f(X)\right] - \frac{1}{n}\sum_{i=1}^n f(x_i)\right)\right] \leq 2R_n(\mathscr{F}).$$

From Corollary 3.4.5, we finally show that in this example, the  $\sqrt{\log n}$  factor is superfluous as well.

Remark. Corollary 3.4.5 is a distribution-free result, i.e.,

$$\sup_{\mathbb{P}} \mathbb{E} \left[ \sup_{f \in \mathscr{F}} |\mathbb{P}_n f - \mathbb{P} f| \right] \le C \sqrt{\frac{\text{VC}(\mathscr{F})}{n}}.$$

This means that boolean function classes with finite VC dimension are uniform Glivenko-Cantelli. We note that the "machinery" we have done is the following:

- 1. Bound uniform entropy w.r.t.  $L_2$  (uniform entropy integral bound).
- 2. Uniform  $L_2$  entropy  $\leq$  VC dimension (Theorem 3.4.5).

**Problem.** How to extend this "machinery" to non-boolean function classes?

Answer. Define VC dimension for non-boolean function classes.

#### Lecture 15: Parametric v.s. Non-Parametric

#### 3.4.8 Parametric versus Non-Parametric Function Classes

We start by asking the following question.

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**Problem.** What makes a function class "parametric" or "non-parametric"?

Answer. If the function class is a vector space, we can usually use the linear algebra notion of dimensio.

Consider the following (not very precise) definition in terms of the uniform  $L_2$  entropy.

**Definition 3.4.16** (Parametric). A function class  $\mathcal{F}$  is parametric if there exists a notion of dimension  $\dim \mathcal{F}$  and a constant C such that

$$\sup_{\mu} N(\mathscr{F}, L_2(\mu), \epsilon) \le \left(\frac{C}{\epsilon}\right)^{\dim(\mathscr{F})}.$$

**Example.** Boolean function class on  $\chi$  with finite VC dimension is parametric.

**Proof.** Dudley's result directly applies.

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**Definition 3.4.17** (Non-parametric). A function class  $\mathcal{F}$  is non-parametric if there is a p > 0 and a constant C such that

$$\sup_{\mu} \log N(\mathscr{F}, L_2(\mu), \epsilon) \le \left(\frac{C}{\epsilon}\right)^p.$$

Let's consider any parametric class  $\mathscr{F}$  uniformly bounded by 1 with dim  $\mathscr{F} = d$ . Then from Dudley's theorem, the Rademacher complexity can be bounded as

$$\mathbb{E}_x \left[ \mathbb{E}_{\epsilon} \left[ \sup_{f \in \mathscr{F}} \frac{1}{n} \sum_{i=1}^n \epsilon_i f(x_i) \right] \right] \leq \frac{12}{\sqrt{n}} \int_0^1 \sqrt{\sup_{\mu} \log N(\mathscr{F}, L_2(\mu), \epsilon)} \, \mathrm{d}\epsilon \leq \frac{12}{\sqrt{n}} \int_0^1 \sqrt{d \log \frac{C}{\epsilon}} \, \mathrm{d}\epsilon \leq C' \sqrt{\frac{d}{n}}.$$

Hence, we get the parametric rate  $O(\sqrt{d/n})$ , and this result is distribution-free.

Analogously, we want to know what we will get for a non-parametric function class (uniform bounded by 1)? Now, since for a non-parametric class, the uniform  $L_2$  entropy is  $\leq (C/\epsilon)^p$ ,

$$\mathbb{E}_{\epsilon} \left[ \sup_{f \in \mathscr{F}} \frac{1}{\sqrt{n}} \sum_{i=1}^{n} \epsilon_{i} f(x_{i}) \right] \qquad \frac{1}{\sqrt{n}} \sum_{i=1}^{n} \epsilon_{i} f(x_{i}) = X_{f} \sim \operatorname{Subg}(L_{2}(\mathbb{P}_{n}))$$

$$\leq \mathbb{E} \left[ \sup_{\substack{f,g \in \mathscr{F}: \\ L_{2}(\mathbb{P}_{n})(f,g) \leq \delta}} X_{f} - X_{g} \right] + \int_{\delta}^{1} \sqrt{\sup_{\mu} \log N(\mathscr{F}, L_{2}(\mu), \epsilon)} \, d\epsilon \qquad \text{modified Corollary 3.4.4}$$

$$\leq \sqrt{n} \cdot \delta + \int_{\delta}^{1} \left( \frac{C}{\epsilon} \right)^{p/2} \, d\epsilon$$

Since the choice of  $\delta$  is arbitrary, we can optimize in terms of  $\delta$ . There are three cases:

• p < 2: Take  $\delta = 0$  because the integral converges, and we get a parametric rate bound for  $R_n$  with some constant c:

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^{n}\epsilon_{i}f(x_{i})\right]\leq\frac{c}{\sqrt{n}}.$$

**Remark.**  $O(1/\sqrt{n})$  is a parametric rate.

Example. Consider Hölder smooth classes. Even though these function classes are nonparametric according to Definition 3.4.17, in terms of  $R_n$  or supremum of empirical process, the rate is still parametric.

• p > 2: We see that for all  $\delta \in (0, 1)$ ,

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^{n}\epsilon_{i}f(x_{i})\right] \leq \delta + \frac{1}{\sqrt{n}}\int_{\delta}^{1}\left(\frac{C}{\epsilon}\right)^{p/2}\mathrm{d}\epsilon$$

$$= \delta + \frac{C^{p/2}}{\sqrt{n}}\cdot\frac{\epsilon^{-p/2+1}}{1-p/2}\bigg|_{\delta}^{1}$$

$$\approx \delta + \frac{1}{\sqrt{n}}\left(-\epsilon^{-p/2+1}\right)\bigg|_{\delta}^{1} \qquad \text{dropping constant } (1-p/2<0)$$

$$\approx \delta + \frac{1}{\sqrt{n}}\delta^{1-p/2}.$$

Now by optimizing over  $\delta$ , setting  $\delta = \delta^{1-p/2}/\sqrt{n}$ , we get the bound  $O(n^{-1/p})$ .

**Remark.**  $O(n^{-1/p})$  is a non-parametric rate, strictly slower than the parametric rate  $O(n^{-1/2})$ .

This upper bound is also tight for certain function classes.

**Example.** For 1-bounded and 1-Lipschitz functions on  $[0,1]^d$ , the uniform  $L_2$  entropy (in fact the  $L_\infty$  entropy) grows like  $(1/\epsilon)^d$ . **Proof.** Since  $|f(x) - f(y)| \le ||x - y||_2$ ,  $O(n^{-1/d})$  rate here is tight for d > 2.

**Proof.** Since 
$$|f(x) - f(y)| \le ||x - y||_2$$
,  $O(n^{-1/d})$  rate here is tight for  $d > 2$ .

• p = 2: From the same calculation, we have

$$\mathbb{E}\left[\sup_{f\in\mathscr{F}}\frac{1}{n}\sum_{i=1}^{n}\epsilon_{i}f(x_{i})\right] \leq \delta + \frac{1}{\sqrt{n}}\int_{\delta}^{1}\frac{C}{\epsilon}\,\mathrm{d}\epsilon$$

$$= \delta + \frac{C}{\sqrt{n}}\ln\epsilon\Big|_{\delta}^{1}$$

$$= \delta + \frac{C}{\sqrt{n}}\ln\frac{1}{\delta} = O\left(\frac{1}{\sqrt{n}}\log n\right)$$

by setting  $\delta = O(1/\sqrt{n})$ . Hence, the Rademacher complexity is bounded by  $\log n/\sqrt{n}$ , which is "almost" the parametric rate up to the extra log factor. This might not be tight.

Remark. To summarize, we have the following bounds on the Rademacher complexity:

- Parametric class:  $C\sqrt{d/n}$
- Non-parametric class:

$$\begin{array}{ll} - \ p < 2 \colon \ C \sqrt{1/n}; \\ - \ p = 2 \colon \ C \log n / \sqrt{n}; \end{array}$$

$$- p > 2$$
:  $C \cdot n^{-1/p}$ .

**Example** (Linear function class). Let  $\chi = B_2^d$ , and  $\mathscr{F} = \{x \mapsto w^\top x \colon w \in B_2^d\}$ . For a given data  $x_1, \ldots, x_n \in \mathbb{R}^d$ ,

$$\mathscr{F}|_{x_1,...,x_n} = \left\{ Xw \colon w \in B_2^d, X_{n \times d} = \begin{bmatrix} x_1^\top \\ \vdots \\ x_n^\top \end{bmatrix} \right\}.$$

To determine whether  $\mathscr{F}$  is parametric or non-parametric, we need to bound  $N(\mathscr{F}, L_2(\mathbb{P}_n), \epsilon)$ :

$$\sqrt{\frac{1}{n} \sum_{i=1}^{n} (\langle w, x_i \rangle - \langle w', x_i \rangle)^2}$$

$$\leq \max_{i \in [n]} |\langle w - w', x_i \rangle|$$

$$\leq \max_{x \in B_2^d} |\langle w - w', x \rangle|$$

$$\leq M(\mathscr{F}, L_{\infty}(\mathbb{P}_n), \epsilon)$$

$$\leq M(\mathscr{F}, \|\cdot\|_{\infty}, \epsilon)$$

$$\leq \|w - w'\|_2$$

$$\leq N(\mathcal{F}, \|\cdot\|_{\infty}, \epsilon)$$

where  $\max_{x \in B_2^d} |\langle w - w', x \rangle| \le ||w - w'||_2$  since  $||x||_2 \le 1$ . Then, from Proposition 3.4.1,

$$N(B_2^d, \|\cdot\|_2, \epsilon) \le \left(1 + \frac{2}{\epsilon}\right)^d,$$

so we get a  $\sqrt{d/n}$  rate since this satisfies parametric condition.

It turns out that one can use the norm constraints to also show that

$$\sup_{\mu} \log N(\mathscr{F}, L_2(\mathbb{P}_n), \epsilon) \le \frac{C}{\epsilon^2},$$

i.e., a dimension-free bound, hence  $\mathscr{F}$  also behaves like a non-parametric class! This bound will be useful in high dimensional setting when d is not small compared to n. Thus, we make this an important remark.

Remark. A function class can be viewed as parametric and non-parametric at the same time.

There are other examples.

**Example.** Neural networks are like this: we can either measure its complexity by the *number of parameters*, or get dimension independent bounds on its Rademacher complexity by using *norm constraints*.

**Problem.** For what function classes can be bound in terms of uniform  $L_2$  entropy?

**Answer.** We have seen boolean function classes with finite VC dimension, and Hölder smooth function classes.

### Lecture 16: Beyond VC Dimension: Fat-Shattering Dimension

## 3.5 Fat-Shattering Dimension

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To generalize VC dimension, consider the following.

**Definition.** Let  $\mathscr{F}$  be a real-valued function class on  $\chi$ .

**Definition 3.5.1** ( $\epsilon$ -shattered). A set  $\{x_1, \ldots, x_n\}$  of  $\chi$  is  $\epsilon$ -shattered by  $\mathscr F$  if there exists  $t_1, \ldots, t_n$  such that for all  $S \subseteq [n]$ , there exists  $f \in \mathscr F$  such that

$$\begin{cases} f(x_s) \le t_s, & \text{if } s \in S; \\ f(x_s) \ge t_s + \epsilon, & \text{if } s \notin S. \end{cases}$$

<sup>&</sup>lt;sup>a</sup>In high dimension situation, this bound can be loose.

**Definition 3.5.2** (Fat-shattering dimension). The *fat-shattering dimension*  $VC(\mathscr{F}, \epsilon)$  of  $\mathscr{F}$  on  $\chi$  is the maximum integer D such that there exists a size D finite set  $A \subseteq \chi$   $\epsilon$ -shattered by  $\mathscr{F}$ .

**Remark.**  $VC(\mathscr{F}, \epsilon)$  is a non-increasing function of  $\epsilon$ .

**Proof.** IF a set is  $\epsilon$ -shattered then for any  $\delta \leq \epsilon$ , it's also  $\delta$ -shattered.

**Note.** If  $\mathscr{F}$  is boolean, then  $VC(\mathscr{F}, \epsilon) = VC(\mathscr{F})$  for all  $\epsilon \in (0, 1)$  by setting  $t_s = 0$  for all s.

**Example.** Consider  $\mathcal{M} = \{f : I \to [-1, 1] \text{ non-decreasing} \}$ . Then for all  $\epsilon > 0$ ,  $VC(\mathcal{M}, \epsilon) \le 1 + 2/\epsilon$ .

**Definition 3.5.3** (Total variation). The *total variation* of a function  $f: \mathbb{R} \supseteq I \to \mathbb{R}$  on an interval I is defined as

$$TV(f) := \sup_{n \ge 1} \sup_{x_1 < \dots < x_n \in I} \sum_{i=1}^n |f(x_i) - f(x_{i-1})|.$$

**Intuition.** The total variation is some measure of smoothness of functions. It's more general than differentiability since TV(f) can also be defined for discontinuous functions.

Remark. Total variation is actually a norm, i.e., triangle inequality holds.

Let's first see one example.

**Example.** Consider BV(2) :=  $\{f : \text{TV}(f) \leq 2\}$ . If f'(x) exists, then  $\text{TV}(f) = \int |f'(x)| \, dx$ . We see that BV(2)  $\supseteq \mathcal{M}$  since for any non-decreasing function f ranging in [a, b], TV(f) = b - a.

In general, we have the following.

**Lemma 3.5.1.** Let  $\mathscr{F} \ni f \colon I \to \mathbb{R}$  with  $\mathrm{TV}(f) \leq v$  for all  $f \in \mathscr{F}$ , i.e.,  $\mathscr{F} = \mathrm{BV}(v)$ . Then

$$VC(BV(v), \epsilon) = 1 + \left| \frac{v}{\epsilon} \right|.$$

**Proof.** We prove this by proving two directions.

Claim.  $VC(BV(v), \epsilon) \le 1 + \lfloor v/\epsilon \rfloor$ .

**Proof.** Let  $\{x_1,\ldots,x_n\}$  be  $\epsilon$ -shattered by  $\mathscr{F}$ , then there exists  $t_1,\ldots,t_n$  and  $f_1,f_2$  such that

$$\begin{cases} f_1(x_i) \le t_i, & \text{if } i \text{ is odd;} \\ f_1(x_i) \ge t_i + \epsilon, & \text{if } i \text{ is even;} \end{cases}, \begin{cases} f_2(x_i) \le t_i, & \text{if } i \text{ is even;} \\ f_2(x_i) \ge t_i + \epsilon, & \text{if } i \text{ is odd.} \end{cases}$$

Consider  $f = (f_1 - f_2)/2$ , then

$$\begin{cases} f(x_i) \le -\epsilon/2, & \text{if } i \text{ is odd;} \\ f(x_i) \ge \epsilon/2, & \text{if } i \text{ is even;} \end{cases} \Rightarrow \text{TV}(f) \ge (n-1)\epsilon$$

by considering this particular partition  $\{x_i\}$  of I. Furthermore, since TV is a norm,

$$TV(f) \le \frac{TV(f_1) + TV(f_2)}{2} \le v$$

from triangle inequality, hence  $(n-1)\epsilon \leq v$ , i.e.,  $n \leq 1 + \lfloor v/\epsilon \rfloor$ .

\*

Claim.  $VC(BV(v), \epsilon) \ge 1 + |v/\epsilon|$ .

**Proof.** Let  $d = \lfloor v/\epsilon \rfloor$ , and consider  $y_1 < y_2 < \cdots < y_d$ , which induces d+1 intervals

$$I_0 = (-\infty, y_1), \quad I_j = [y_j, y_{j+1}), \quad I_d = [y_d, \infty)$$

Let  $\mathcal{G}$  be the set of piece-wise continuous functions on  $I_0, \ldots, I_d$  taking values between  $\{0, \epsilon\}$ , so  $|\mathcal{G}| = 2^{d+1}$ . Then, the set

$$\{x_1,\ldots,x_{d+1}\colon x_j\in I_{j-1}\}$$

is  $\epsilon$ -shattered by  $\mathcal{G}$ . Finally, since  $\mathrm{TV}(g) \leq d\epsilon \leq v$  for all  $g \in \mathcal{G}$ , we're done.

Consider the following result (which we will not prove).

**Theorem 3.5.1** (Mendelson-Vershynin). Let  $\mathscr{F}$  be a class of functions that is uniformly bounded by 1. Then there exists c > 1 such that for every  $0 < \epsilon \le 1$ ,

$$\sup_{\mu} M(\mathscr{F}, \epsilon, L_2(\mu)) \leq \left(\frac{2}{\epsilon}\right)^{c \operatorname{VC}(\mathscr{F}, \frac{\epsilon}{c})}.$$

**Remark.** For BV(2), Mendelson-Vershynin theorem gives  $\sup_{\mu} M(BV(2), \epsilon, L_2(\mu)) \leq \exp(\frac{c}{\epsilon} \log 2\epsilon)$ .

Now, consider the classification such that  $\chi = B_2^d$ , and  $\mathscr{F} = \{x \mapsto w^\top x \colon w \in B_2^d\}$ . Then  $VC(\mathscr{F}, \epsilon) \leq d$ ; and if we consider  $sgn(w^\top x)$ , we get  $VC(\mathscr{F}, \epsilon) = d + 1$ .

<sup>&</sup>lt;sup>a</sup>Note that  $\log 2\epsilon$  is superfluous again.

# Appendix

## Bibliography

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