

A proof of Doug Lea's memory manager

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Chapter 1

Glossary of macros, typedefs and minor routines

```
MALLOC_ALIGNMENT = 8
MAX_SIZE_T       =  $FFFF\ FFFF_h$ 
SIZE_T_SIZE      = 4
SIZE_T_BITSIZE   = 32
SIZE_T_ZERO      = 0
SIZE_T_ONE       = 1
SIZE_T_TWO       = 2
SIZE_T_FOUR      = 4
TWO_SIZE_T_SIZES = 8
FOUR_SIZE_T_SIZES = 16
SIX_SIZE_T_SIZES = 24
HALF_MAX_SIZE_T  =  $7FFF\ FFFF_h$ 
CHUNK_ALIGN_MASK =  $111_b$ 
mchunk           = struct malloc_chunk
mchunkptr        = mchunk*
sbinptr          = mchunk*
bindx_t          = unsigned int
binmap_t         = unsigned int
flag_t           = unsigned int
MCHUNK_SIZE      = 16
CHUNK_OVERHEAD   = 4
MIN_CHUNK_SIZE   = 16
chunk2mem(p)     =  $p + 8$ 
mem2chunk(mem)   =  $\text{mem} - 8$ 
MAX_REQUEST      =  $2^{32} - 63$ 
MIN_REQUEST      = 11
pad_request(req) =  $\lceil \text{req} + 4 \rceil_8$ 
request2size(req) =  $\max\{16, \lceil \text{req} + 4 \rceil_8\}$ 
```

```

PINUSE_BIT           =  $1_b$ 
CINUSE_BIT           =  $10_b$ 
FLAG4_BIT            =  $100_b$ 
INUSE_BITS            =  $11_b$ 
FLAG_BITS            =  $111_b$ 
cinuse(p)            =  $flags(p) = \nabla\_$ 
pinuse(p)            =  $flags(p) = \_ \blacktriangle$ 
is_inuse(p)          =  $flags(p) \in \{\nabla\_, \nabla\Delta\}$ 
is_mmapped(p)        =  $flags(p) = \nabla\Delta$ 
chunksize(p)         =  $size(p)$ 
 $\{flags(p) = C\_ \}$  clear_pinuse(p)  $\{flags(p) = C\Delta \}$ 
chunk_plus_offset(p,s) =  $p + s$ 
chunk_minus_offset(p,s) =  $p - s$ 
next_chunk(p)        =  $next(p)$ 
prev_chunk(p)        =  $prev(p)$ 
next_pinuse(p)       =  $flags(next(p)) = \_ \blacktriangle$ 
get_foot(p,s)        =  $prev\_foot(p + s)$ 
 $\{prev\_foot(p + s) = \_ \}$  set_foot(p,s)  $\{prev\_foot(p + s) = s \}$ 
 $\left\{ \begin{array}{l} size(p) = \_ \wedge flags(p) = -- \\ \wedge prev\_foot(p + s) = \_ \end{array} \right\}$  set_size_and_pinuse_of_free_chunk(p,s)  $\left\{ \begin{array}{l} size(p) = s \wedge flags(p) = \nabla\blacktriangle \\ \wedge prev\_foot(next(p)) = s \end{array} \right\}$ 
 $\left\{ \begin{array}{l} size(p) = \_ \wedge flags(p) = -- \\ \wedge prev\_foot(p + s) = \_ \\ \wedge flags(p + s) = -- \end{array} \right\}$  set_free_with_pinuse(p,s,n)  $\left\{ \begin{array}{l} size(p) = s \wedge flags(p) = \nabla\blacktriangle \\ \wedge prev\_foot(next(p)) = s \\ \wedge flags(next(p)) = \_ \Delta \end{array} \right\}$ 
tchunk              = malloc_tree_chunk
tchunkptr           = tchunk*
tbinptr             = tchunk*
leftmost_child(t)   =  $\begin{cases} child_0(*t) & \text{if } child_0(*t) \neq 0 \\ child_1(*t) & \text{otherwise} \end{cases}$ 
NSMALLBINS          = 32
NTREEBINS           = 32
SMALLBIN_SHIFT       = 3
SMALLBIN_WIDTH       = 8
TREEBIN_SHIFT        = 8
MIN_LARGE_SIZE       = 256
MAX_SMALL_SIZE       = 255
MAX_SMALL_REQUEST    = 244
mstate              = struct malloc_state
mparams              = struct malloc_params
is_small(s)          =  $s < 256$ 
small_index(s)       =  $\lfloor s/8 \rfloor$ 
small_index2size(i)  =  $8 \times i$ 
MIN_SMALL_INDEX      = 2

```

$$\begin{aligned}
& \{ \text{smallbins}[2i+2] \mapsto C_1 * \text{smallbins}[2i+3] \mapsto C_2 \} \quad x := \text{smallbin_at}(M, i) \quad \{ x.\text{fd} \mapsto C_1 * x.\text{bk} \mapsto C_2 \} \\
\text{treebin_at}(M, i) &= \text{treebins}[i] \\
\{ I = _ \} \text{compute_tree_index}(S, I) &= \left\{ I = \begin{cases} 0 & \text{if } S < 256 \\ 31 & \text{if } S > 2^{24} \\ 2(\log_2 \|S\| - 8) & \text{if } 0 \leq \{S\} < \frac{1}{2}\|S\| \\ 2(\log_2 \|S\| - 8) + 1 & \text{if } \frac{1}{2}\|S\| \leq \{S\} < \|S\| \end{cases} \right\} \\
\text{bin_for_tree_index}(i) &= \begin{cases} 31 & \text{if } i = 31 \\ \lfloor i/2 \rfloor + 6 & \text{otherwise} \end{cases} \\
\text{leftshift_for_tree_index}(i) &= \begin{cases} 0 & \text{if } i = 31 \\ 25 - \lfloor i/2 \rfloor & \text{otherwise} \end{cases} \\
\text{minsize_for_tree_index}(i) &= \begin{cases} 2 \ll (\lfloor i/2 \rfloor + 7) & \text{if } i \text{ even} \\ 3 \ll (\lfloor i/2 \rfloor + 7) & \text{if } i \text{ odd} \end{cases} \\
\text{idx2bit}(i) &= 1 \ll i \\
\{ \text{smallmap}[i] = _ \} \text{mark_smallmap}(M, i) &= \{ \text{smallmap}[i] = 1 \} \\
\{ \text{smallmap}[i] = _ \} \text{clear_smallmap}(M, i) &= \{ \text{smallmap}[i] = 0 \} \\
\text{smallmap_is_marked}(M, i) &= \text{smallmap}[i] = 1 \\
\{ \text{treemap}[i] = _ \} \text{mark_treemap}(M, i) &= \{ \text{treemap}[i] = 1 \} \\
\{ \text{treemap}[i] = _ \} \text{clear_treemap}(M, i) &= \{ \text{treemap}[i] = 0 \} \\
\text{treemap_is_marked}(M, i) &= \text{treemap}[i] = 1 \\
\text{least_bit}(x) &= \begin{cases} 0 \overset{i}{1} 0 & \text{if } x_i = 1 \wedge \forall j < i. x_j = 0 \\ 0 & \text{if } x = 0 \end{cases} \\
\text{left_bits}(x) &= \begin{cases} 1 \overset{i}{0} 0 & \text{if } x_i = 1 \wedge \forall j < i. x_j = 0 \\ 0 & \text{if } x = 0 \end{cases} \\
\text{same_or_left_bits}(x) &= \begin{cases} 1 \overset{i}{1} 0 & \text{if } x_i = 1 \wedge \forall j < i. x_j = 0 \\ 0 & \text{if } x = 0 \end{cases} \\
\{ I = _ \} \text{compute_bit2idx}(X, I) &= \{ X \neq 0 \Rightarrow I = \log_2 X \} \\
\{ p \} \text{mark_inuse_foot}(M, p, s) &= \{ p \} \\
\left\{ \begin{array}{l} \text{size}(p) = _ \wedge \text{flags}(p) = _P \\ \wedge \text{flags}(p+s) = C_ \end{array} \right\} \text{set_inuse}(M, p, s) &= \left\{ \begin{array}{l} \text{size}(p) = s \wedge \text{flags}(p) = \blacktriangledown P \\ \wedge \text{flags}(\text{next}(p)) = C \blacktriangle \end{array} \right\} \\
\left\{ \begin{array}{l} \text{size}(p) = _ \wedge \text{flags}(p) = _ _ \\ \wedge \text{flags}(p+s) = C_ \end{array} \right\} \text{set_inuse_and_pinuse}(M, p, s) &= \left\{ \begin{array}{l} \text{size}(p) = s \wedge \text{flags}(p) = \blacktriangledown \blacktriangle \\ \wedge \text{flags}(\text{next}(p)) = C \blacktriangle \end{array} \right\} \\
\left\{ \begin{array}{l} \text{size}(p) = _ \\ \wedge \text{flags}(p) = _ _ \end{array} \right\} \text{set_inuse_and_pinuse_of_inuse_chunk}(M, p, s) &= \left\{ \begin{array}{l} \text{size}(p) = s \\ \wedge \text{flags}(p) = \blacktriangledown \blacktriangle \end{array} \right\}
\end{aligned}$$

Chapter 2

Describing chunks

Define

$$\begin{aligned}
C_1 \curvearrowright C_2 &\Leftrightarrow C_1.\text{fd} \mapsto C_2 \\
C_1 \overset{S}{\curvearrowright} C_2 &\Leftrightarrow C_1.\text{fd} \mapsto C_2 * \exists S'. C_1.\text{size} \overset{5}{\mapsto} S' * S' \geq S \\
C_1 \curvearrowleft C_2 &\Leftrightarrow C_2.\text{bk} \mapsto C_1 \\
C_1 \curvearrowright C_2 &\Leftrightarrow C_1 \curvearrowright C_2 * C_1 \curvearrowleft C_2 \\
C_1 \overset{S}{\curvearrowright} C_2 &\Leftrightarrow C_1 \overset{S}{\curvearrowright} C_2 * C_1 \curvearrowleft C_2 \\
C_1 \overset{S}{\curvearrowright}^* C_2 &\Leftrightarrow C_1 \dot{=} C_2 \vee \exists C'. C_1 \overset{S}{\curvearrowright} C' * C' \overset{S}{\curvearrowright}^* C_2 \\
C_1 \overset{S}{\curvearrowleft}^* C_2 &\Leftrightarrow C_1 \dot{=} C_2 \vee \exists C'. C_1 \overset{S}{\curvearrowleft} C' * C' \overset{S}{\curvearrowleft}^* C_2
\end{aligned}$$

We begin by tackling a simplified scenario, in which we ignore treebins. Assume chunks have sizes that are multiples of 8 bytes up to 256 bytes (exclusive). All free chunks are thus contained in one of 32 smallbins. We shall also ignore the smallmap for now.

$$\begin{aligned}
node_S(C) &\Leftrightarrow busynode_S(C) \vee freenode_S(C) \\
busynode_S(C) &\Leftrightarrow \exists S', C_n. \\
&\quad * C.\text{cinuse} \mapsto 1 \\
&\quad * C.\text{size} \mapsto S' * S' \geq S \\
&\quad * C_n \dot{=} C + S' \\
&\quad * node_0(C_n) \\
&\quad * C_n.\text{pinuse} \mapsto 1 \\
freenode_S(C) &\Leftrightarrow \exists S', C_n. \\
&\quad * C.\text{cinuse} \mapsto 0 \\
&\quad * C.\text{size} \overset{5}{\mapsto} S' * S' \geq S \\
&\quad * snode_S(C) \\
&\quad * C_n \dot{=} C + S' \\
&\quad * busynode_0(C_n) \\
&\quad * C_n.\text{prevfoot} \mapsto S' \\
&\quad * C_n.\text{pinuse} \mapsto 0 \\
snode_S(C) &\Leftrightarrow \exists C' \neq C. C \overset{S}{\curvearrowright} C' * C' \overset{S}{\curvearrowright}^* C \\
sbinheader_S(B) &\Leftrightarrow B \curvearrowright B \vee \exists C \neq B. B \curvearrowright C * C \overset{S}{\curvearrowright}^* B
\end{aligned}$$

If **start** is the address of the first chunk in the arena, then the entire state can be described by:

$$node_0(\text{start}) * \forall^* i \in [2, 32). sbin_{8i}(\text{smallbins} + 2i)$$

TODO:

- make a special case for the first node (its PINUSE flag should be permanently set)

Chapter 3

Operations on smallbins and trees

Then let

$$\begin{aligned} \text{sbinempty}(B) &\Leftrightarrow B \dot{\subseteq} B \\ \text{sbinnonempty}(S, B) &\Leftrightarrow \exists C \neq B. B \dot{\subseteq} C * C \overset{S}{\dot{\subseteq}}^* B \\ \text{sbin}(S, B) &\Leftrightarrow \text{sbinempty}(B) \vee \text{sbinnonempty}(S, B) \end{aligned}$$

The following predicate describes a valid `smallbins` array and corresponding `smallmap` vector.

$$\begin{aligned} \text{sbinAndMap}(i) &\Leftrightarrow \exists S, B. \ i \in [2, 32) * S \dot{=} 8i * B \dot{=} \text{smallbins} + 2i * \\ &\quad (\&\text{smallmap}_i \mapsto 0 * \text{sbinempty}(B) \\ &\quad \vee \&\text{smallmap}_i \mapsto 1 * \text{sbinnonempty}(S, B)) \\ \text{sbinsOK} &\Leftrightarrow \forall i \in [2, 32). \text{sbinAndMap}(i) \end{aligned}$$

The `insert_small_chunk` macro can be specified like so. Note that it preserves the `sbinsOK` predicate:

$$\left\{ \exists i. i \in [2, 32) * S \dot{=} 8i * P \overset{S}{\dot{\subseteq}} _ * _ \dot{\subseteq} P * \text{sbinsOK} \right\} \text{insert_small_chunk}(M, P, S) \left\{ \text{sbinsOK} \right\}$$

The details are as follows.

$$\begin{aligned} &\left\{ \exists i. i \in [2, 32) * S \dot{=} 8i * P \overset{S}{\dot{\subseteq}} _ * _ \dot{\subseteq} P * \text{sbinsOK} \right\} \\ &\text{bindex_t } I = \text{small_index}(S); \\ &\left\{ P \overset{S}{\dot{\subseteq}} _ * _ \dot{\subseteq} P * \text{sbinsOK} * I \in [2, 32) * S \dot{=} 8I \right\} \\ &\left\{ P \overset{S}{\dot{\subseteq}} _ * _ \dot{\subseteq} P * (\text{sbinAndMap}(I) \dot{\rightarrow} \text{sbinsOK}) * \text{sbinAndMap}(I) * I \in [2, 32) * S \dot{=} 8I \right\} \\ &\text{mchunkptr } B = \text{smallbin_at}(M, I); \\ &\left\{ P \overset{S}{\dot{\subseteq}} _ * _ \dot{\subseteq} P * (\text{sbinAndMap}(I) \dot{\rightarrow} \text{sbinsOK}) * \text{sbinAndMap}(I) * I \in [2, 32) * S \dot{=} 8I * B \dot{=} \text{smallbins} + 2I \right\} \\ &\text{mchunkptr } F = B; \\ &\text{assert}(S \geq \text{MIN_CHUNK_SIZE}); \\ &\left\{ P \overset{S}{\dot{\subseteq}} _ * _ \dot{\subseteq} P * (\text{sbinAndMap}(I) \dot{\rightarrow} \text{sbinsOK}) * \text{sbinAndMap}(I) \right\} \\ &\left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} \text{smallbins} + 2I * F \dot{=} B \right\} \\ &\text{if } (!\text{smallmap_is_marked}(M, I)) \\ &\quad \left\{ P \overset{S}{\dot{\subseteq}} _ * _ \dot{\subseteq} P * (\text{sbinAndMap}(I) \dot{\rightarrow} \text{sbinsOK}) * \&\text{smallmap}_I \mapsto 0 * \text{sbinempty}(B) \right\} \\ &\quad \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} \text{smallbins} + 2I * F \dot{=} B \right\} \\ &\text{mark_smallmap}(M, I); \\ &\left\{ P \overset{S}{\dot{\subseteq}} _ * _ \dot{\subseteq} P * (\text{sbinAndMap}(I) \dot{\rightarrow} \text{sbinsOK}) * \&\text{smallmap}_I \mapsto 1 * \text{sbinempty}(B) \right\} \\ &\left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} \text{smallbins} + 2I * F \dot{=} B \right\} \end{aligned}$$

$$\begin{aligned}
& \left\{ P \overset{S}{\rightsquigarrow} _ * _ \rightsquigarrow P * (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * B \circlearrowleft F \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I * F \dot{=} B * B \dot{=} P \right\} \\
& \text{else if (RTCHECK(ok_address(M, B->fd)))} \\
& \left\{ P \overset{S}{\rightsquigarrow} _ * _ \rightsquigarrow P * (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * sbinnonempty(S, B) \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I \right\} \\
& \left\{ \exists C \neq B. P \overset{S}{\rightsquigarrow} _ * _ \rightsquigarrow P * (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * B \circlearrowleft C * C \overset{S}{\rightsquigarrow} * B \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I * B \dot{=} P \right\} \\
& F = B->fd; \\
& \left\{ P \overset{S}{\rightsquigarrow} _ * _ \rightsquigarrow P * (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * B \circlearrowleft F * F \overset{S}{\rightsquigarrow} * B \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I * B \dot{=} P * F \dot{=} B \right\} \\
& \text{else \{ } \\
& \quad CORRUPTION_ERROR_ACTION(M); \\
& \} \\
& \left\{ P \overset{S}{\rightsquigarrow} _ * _ \rightsquigarrow P * (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * B \circlearrowleft F * F \overset{S}{\rightsquigarrow} * B \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I * B \dot{=} P \right\} \\
& B->fd = P; \\
& \left\{ P \overset{S}{\rightsquigarrow} _ * _ \rightsquigarrow P * (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * B \rightsquigarrow P * B \rightsquigarrow F * F \overset{S}{\rightsquigarrow} * B \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I * B \dot{=} P \right\} \\
& F->bk = P; \\
& \left\{ P \overset{S}{\rightsquigarrow} _ * _ \rightsquigarrow P * (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * B \rightsquigarrow P * P \rightsquigarrow F * F \overset{S}{\rightsquigarrow} * B \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I * B \dot{=} P \right\} \\
& P->fd = F; \\
& \left\{ _ \rightsquigarrow P * (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * B \rightsquigarrow P * P \overset{S}{\rightsquigarrow} F * F \overset{S}{\rightsquigarrow} * B \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I * B \dot{=} P \right\} \\
& P->bk = B; \\
& \left\{ (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * B \circlearrowleft P * P \overset{S}{\rightsquigarrow} F * F \overset{S}{\rightsquigarrow} * B \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I * B \dot{=} P \right\} \\
& \left\{ (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * B \circlearrowleft P * P \overset{S}{\rightsquigarrow} * B \right\} \\
& \left\{ * I \in [2, 32) * S \dot{=} 8I * B \dot{=} smallbins + 2I * B \dot{=} P \right\} \\
& \left\{ (sbinAndMap(I) \multimap sbinsOK) * \&smallmap_I \mapsto 1 * sbinnonempty(S, smallbins + 2I) * I \in [2, 32) * S \dot{=} 8I \right\} \\
& \left\{ (sbinAndMap(I) \multimap sbinsOK) * sbinAndMap(I) \right\} \\
& \left\{ sbinsOK \right\}
\end{aligned}$$

Chapter 4

`tmalloc_small`

Chapter 5

`tmalloc_large`

Chapter 6

`sys_alloc`

Chapter 7

dlmalloc

We now give the entire source code of `dlmalloc`, interspersed with an explanatory commentary and annotations for the safety proof. We shall prove that a successful call to `dlmalloc(mem)` allocates a chunk with payload no less than `mem` bytes. I think it's probably sufficient to provide only a lower bound on the payload size – it certainly simplifies the proof!

```
 $\{p_1\}$ 
void* dlmalloc(size_t bytes) {
  #if USE_LOCKS
    ensure_initialization(); /* initialize in sys_alloc if not using locks */
  #endif
  if (!PREACTION(gm)) {
     $\{-\}$ 
    void* mem;
    size_t nb;
    if (bytes <= MAX_SMALL_REQUEST) {
       $\{P_{small}\}$  where  $P_{small} = \text{bytes} \leq 244$ 

```

Allocating small chunks

```
    bindex_t idx;
    binmap_t smallbits;
    nb = (bytes < MIN_REQUEST)? MIN_CHUNK_SIZE : pad_request(bytes);
     $\{P_{small} \wedge P_{nb}\}$  where  $P_{nb} = \max\{16, \lceil \text{bytes} + 4 \rceil_8\}$ 
    idx = small_index(nb);
     $\{P_{small} \wedge P_{nb} \wedge \text{idx} = \lfloor \text{nb}/8 \rfloor\}$ 
    smallbits = gm->smallmap >> idx;
     $\left\{ P_{small} \wedge P_{nb} \wedge \text{idx} = \lfloor \text{nb}/8 \rfloor \right.$ 
     $\left. \wedge \forall i \in [0, 32 - \text{idx}). \text{smallbits}[i] = 0 \Leftrightarrow \text{smallbin}(i + \text{idx}) = \emptyset \right\}$ 
    if ((smallbits & 0x3U) != 0) { /* Remainderless fit to a smallbin. */
       $\{P_{small} \wedge P_{nb} \wedge \text{idx} = \lfloor \text{nb}/8 \rfloor \wedge (\text{smallbin}(\text{idx}) \neq \emptyset \vee \text{smallbin}(\text{idx} + 1) \neq \emptyset)\}$ 

```

‘Remainderless’ fit to a smallbin

```
    mchunkptr b, p;
    idx += ~smallbits & 1; /* Uses next bin if idx empty */

```

```


$$\left\{ P_{small} \wedge P_{nb} \wedge (idx = \lfloor nb/8 \rfloor \vee idx = \lfloor nb/8 \rfloor + 1) \wedge smallbin(idx) \neq \emptyset \right\}$$

b = smallbin_at(gm, idx);

$$\left\{ P_{small} \wedge P_{nb} \wedge (idx = \lfloor nb/8 \rfloor \vee idx = \lfloor nb/8 \rfloor + 1) \wedge b = smallbin(idx) \wedge b \neq \emptyset \right\}$$

p = b->fd;

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge (idx = \lfloor nb/8 \rfloor \vee idx = \lfloor nb/8 \rfloor + 1) \\ \wedge b = smallbin(idx) \wedge b = p :: ps \wedge flags(p) = 01 \end{array} \right\}$$

assert(chunksize(p) == small_index2size(idx));
unlink_first_small_chunk(gm, b, p, idx);

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge (idx = \lfloor nb/8 \rfloor \vee idx = \lfloor nb/8 \rfloor + 1) \\ \wedge b = smallbin(idx) \wedge b = ps \\ \wedge flags(p) = \nabla \blacktriangle \end{array} \right\}$$

set_inuse_and_pinuse(gm, p, small_index2size(idx));

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge (idx = \lfloor nb/8 \rfloor \vee idx = \lfloor nb/8 \rfloor + 1) \\ \wedge b = smallbin(idx) \wedge b = ps \\ \wedge flags(p) = \nabla \blacktriangle \wedge flags(next(p)) = \blacktriangle \end{array} \right\}$$

mem = chunk2mem(p);
check_malloced_chunk(gm, mem, nb);

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge (idx = \lfloor nb/8 \rfloor \vee idx = \lfloor nb/8 \rfloor + 1) \\ \wedge b = smallbin(idx) \wedge b = ps \\ \wedge flags(p) = \nabla \blacktriangle \wedge flags(next(p)) = \blacktriangle \wedge mem = p - 2 \end{array} \right\}$$

goto postaction;
}
else if (nb > gm->dvsiz) {

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge idx = \lfloor nb/8 \rfloor \\ \wedge \forall i \in [0, 32 - idx). smallbits[i] = 0 \Leftrightarrow smallbin(i + idx) = \emptyset \\ \wedge smallbin(idx) = smallbin(idx + 1) = \emptyset \end{array} \right\}$$

if (smallbits != 0) { /* Use chunk in next nonempty smallbin */

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge idx = \lfloor nb/8 \rfloor \\ \wedge \forall i \in [0, 32 - idx). smallbits[i] = 0 \Leftrightarrow smallbin(i + idx) = \emptyset \\ \wedge \exists i \in [idx + 2, 32). smallbin(i) \neq \emptyset \end{array} \right\}$$

}
}

```

‘Remainderful’ fit to a smallbin

```

mchunkptr b, p, r;
size_t rsize;
bindex_t i;
binmap_t leftbits = (smallbits << idx) & left_bits(idx2bit(idx));

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge idx = \lfloor nb/8 \rfloor \\ \wedge \forall i \in [idx + 1, 32). leftbits[i] = 0 \Leftrightarrow smallbin(i) = \emptyset \\ \wedge \exists i \in [idx + 2, 32). smallbin(i) \neq \emptyset \end{array} \right\}$$

binmap_t leastbit = least_bit(leftbits);

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge idx = \lfloor nb/8 \rfloor \\ \wedge \forall i \in [idx + 1, 32). leftbits[i] = 0 \Leftrightarrow smallbin(i) = \emptyset \\ \wedge \lfloor \log_2(leastbit) \rfloor = \min\{i \in [idx + 2, 32) \mid smallbin(i) \neq \emptyset\} \end{array} \right\}$$

compute_bit2idx(leastbit, i);

```

```


$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge idx = \lfloor nb/8 \rfloor \\ \wedge \forall i \in [idx+1, 32). leftbits[i] = 0 \Leftrightarrow smallbin(i) = \emptyset \\ \wedge i = \min\{i \in [idx+2, 32) \mid smallbin(i) \neq \emptyset\} \end{array} \right\}$$


$$\left\{ P_{small} \wedge P_{nb} \wedge i \geq \lfloor nb/8 \rfloor + 2 \wedge smallbin(i) \neq \emptyset \right\}$$

b = smallbin_at(gm, i);

$$\left\{ P_{small} \wedge P_{nb} \wedge i \geq \lfloor nb/8 \rfloor + 2 \wedge b = smallbin(i) \wedge b \neq \emptyset \right\}$$

p = b->fd;

$$\left\{ P_{small} \wedge P_{nb} \wedge i \geq \lfloor nb/8 \rfloor + 2 \wedge b = smallbin(i) \wedge b = p :: ps \right\}$$

assert(chunksize(p) == small_index2size(i));
unlink_first_small_chunk(gm, b, p, i);

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge i \geq \lfloor nb/8 \rfloor + 2 \wedge b = smallbin(i) \wedge b = ps \\ \wedge flags(p) = \nabla \blacktriangle \end{array} \right\}$$

rsize = small_index2size(i) - nb;

$$\left\{ P_{small} \wedge P_{nb} \wedge flags(p) = \nabla \blacktriangle \wedge rsize = size(p) - nb \right\}$$

/* Fit here cannot be remainderless if 4byte sizes */
if (SIZE_T_SIZE != 4 && rsize < MIN_CHUNK_SIZE)

$$\left\{ false \right\}$$

    set_inuse_and_pinuse(gm, p, small_index2size(i));
else {
    set_size_and_pinuse_of_inuse_chunk(gm, p, nb);

$$\left\{ P_{small} \wedge P_{nb} \wedge flags(p) = \nabla \blacktriangle \wedge size(p) = nb \wedge rsize = size(p) - nb \right\}$$

    r = chunk_plus_offset(p, nb);

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge flags(p) = \nabla \blacktriangle \wedge size(p) = nb \wedge rsize = size(p) - nb \\ \wedge r = p + nb \end{array} \right\}$$

    set_size_and_pinuse_of_free_chunk(r, rsize);

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge flags(p) = \nabla \blacktriangle \wedge size(p) = nb \wedge rsize = size(p) - nb \\ \wedge r = p + nb \wedge flags(r) = \nabla \blacktriangle \wedge size(r) = rsize \end{array} \right\}$$

    replace_dv(gm, r, rsize);

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge flags(p) = \nabla \blacktriangle \wedge size(p) = nb \wedge rsize = size(p) - nb \\ \wedge r = p + nb \wedge flags(r) = \nabla \blacktriangle \wedge size(r) = rsize \\ \wedge dv = r \wedge dvsize = rsize \end{array} \right\}$$

}
mem = chunk2mem(p);
check_malloted_chunk(gm, mem, nb);

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge flags(p) = \nabla \blacktriangle \wedge size(p) = nb \wedge rsize = size(p) - nb \\ \wedge r = p + nb \wedge flags(r) = \nabla \blacktriangle \wedge size(r) = rsize \wedge mem = p + 2 \end{array} \right\}$$

goto postaction;
}

```

Using a treebin instead

```

else if (gm->treemap != 0 && (mem = tmalloc_small(gm, nb)) != 0) {

$$\left\{ \begin{array}{l} P_{small} \wedge P_{nb} \wedge idx = \lfloor nb/8 \rfloor \wedge \forall i \in [idx, 32). smallbin(i) = \emptyset \\ \wedge mem = p + 2 \wedge flags(p) = \nabla \blacktriangle \wedge size(p) \geq nb \end{array} \right\}$$

    check_malloted_chunk(gm, mem, nb);
    goto postaction;
}

```

```

    }
  }
}

```

Allocating large chunks

```

else if (bytes >= MAX_REQUEST)
  {bytes ≥ 232 - 63}
  nb = MAX_SIZE_T; /* Too big to allocate. Force failure (in sys alloc) */
  {nb = 232 - 1}
else {
  {Plarge} where Plarge = 244 < bytes < 232 - 63
  nb = pad_request(bytes);
  {Plarge ∧ Pnb}
  if (gm->treemap != 0 && (mem = tmalloc_large(gm, nb)) != 0) {
    {Plarge ∧ Pnb ∧ mem = p + 2 ∧ flags(p) = ▽▲ ∧ size(p) ≥ nb}
    check_malloced_chunk(gm, mem, nb);
    goto postaction;
  }
}

```

Using the designated victim

```

{Pnb}
if (nb <= gm->dvsizes) {
  {Pnb ∧ nb ≤ dvsizes}
  size_t rsize = gm->dvsizes - nb;
  {Pnb ∧ nb ≤ dvsizes ∧ rsize = dvsizes - nb}
  mchunkptr p = gm->dv;
  {Pnb ∧ nb ≤ size(p) ∧ rsize = size(p) - nb ∧ flags(p) = ▽▲}
  if (rsize >= MIN_CHUNK_SIZE) { /* split dv */
    {Pnb ∧ rsize = size(p) - nb ∧ rsize ≥ 16 ∧ flags(p) = ▽▲}
    mchunkptr r = gm->dv = chunk_plus_offset(p, nb);
    {Pnb ∧ rsize = size(p) - nb ∧ rsize ≥ 16 ∧ r = p + nb ∧ flags(p) = ▽▲}
    gm->dvsizes = rsize;
    set_size_and_pinuse_of_free_chunk(r, rsize);
    {Pnb ∧ rsize = size(p) - nb ∧ rsize ≥ 16 ∧ r = p + nb ∧ flags(p) = ▽▲}
    { ∧ flags(r) = ▽▲ ∧ size(r) = rsize }
    set_size_and_pinuse_of_inuse_chunk(gm, p, nb);
    {Pnb ∧ rsize ≥ 16 ∧ r = p + nb ∧ flags(p) = ▽▲ ∧ size(p) = nb}
    { ∧ flags(r) = ▽▲ ∧ size(r) = rsize }
  }
}
else { /* exhaust dv */
  {Pnb ∧ (size(p) = nb ∨ size(p) = nb + 8) ∧ flags(p) = ▽▲}
  size_t dvs = gm->dvsizes;
  gm->dvsizes = 0;
  gm->dv = 0;
}

```

```

    set_inuse_and_pinuse(gm, p, dvs);
    {  $P_{nb} \wedge (size(p) = nb \vee size(p) = nb + 8) \wedge flags(p) = \blacktriangledown \blacktriangle$  }
}
{  $P_{nb} \wedge (size(p) = nb \vee size(p) = nb + 8) \wedge flags(p) = \blacktriangledown \blacktriangle$  }
mem = chunk2mem(p);
check_malloted_chunk(gm, mem, nb);
{  $P_{nb} \wedge (size(p) = nb \vee size(p) = nb + 8) \wedge flags(p) = \blacktriangledown \blacktriangle \wedge mem = p + 2$  }
goto postaction;
}

```

Using the top chunk

```

else if (nb < gm->topsize) { /* Split top */
    {  $P_{nb} \wedge nb < size(top)$  }
    size_t rsize = gm->topsize - nb;
    {  $P_{nb} \wedge rsize = size(top) - nb \wedge rsize > 0$  }
    mchunkptr p = gm->top;
    {  $P_{nb} \wedge rsize = size(p) - nb \wedge rsize > 0$  }
    mchunkptr r = gm->top = chunk_plus_offset(p, nb);
    {  $P_{nb} \wedge rsize = size(p) - nb \wedge rsize > 0 \wedge r = p + nb$  }
    r->head = rsize | PINUSE_BIT;
    {  $P_{nb} \wedge size(r) = size(p) - nb \wedge size(r) > 0 \wedge flags(r) = \blacktriangledown \blacktriangle \wedge r = p + nb$  }
    set_size_and_pinuse_of_inuse_chunk(gm, p, nb);
    {  $P_{nb} \wedge size(p) = nb \wedge flags(p) = \blacktriangledown \blacktriangle \wedge size(r) > 0 \wedge flags(r) = \blacktriangledown \blacktriangle \wedge r = p + nb$  }
    mem = chunk2mem(p);
    {  $P_{nb} \wedge size(p) = nb \wedge flags(p) = \blacktriangledown \blacktriangle \wedge mem = p + 2$  }
    check_top_chunk(gm, gm->top);
    check_malloted_chunk(gm, mem, nb);
    goto postaction;
}

```

Obtaining memory from the system

```

    mem = sys_alloc(gm, nb);
postaction:
    POSTACTION(gm);
    return mem;
}
return 0;
}

```


Chapter 8

dlfree

```
void dlfree(void* mem) {
    if (mem != 0) {
        mchunkptr p = mem2chunk(mem);
    #if FOOTERS
        mstate fm = get_mstate_for(p);
        if (!ok_magic(fm)) {
            USAGE_ERROR_ACTION(fm, p);
            return;
        }
    #else /* FOOTERS */
    #define fm gm
    #endif /* FOOTERS */
        if (!PREACTION(fm)) {
            check_inuse_chunk(fm, p);
            if (RTCHECK(ok_address(fm, p) && ok_inuse(p))) {
                size_t psize = chunksize(p);
                mchunkptr next = chunk_plus_offset(p, psize);
                if (!pinuse(p)) {
                    size_t prevsize = p->prev_foot;
                    if (is_mmapped(p)) {
                        psize += prevsize + MMAP_FOOT_PAD;
                        if (CALL_MUNMAP((char*)p - prevsize, psize) == 0)
                            fm->footprint -= psize;
                        goto postaction;
                    }
                }
                else {
                    mchunkptr prev = chunk_minus_offset(p, prevsize);
                    psize += prevsize;
                    p = prev;
                    if (RTCHECK(ok_address(fm, prev))) { /* consolidate backward */
                        if (p != fm->dv) {
                            unlink_chunk(fm, p, prevsize);
                        }
                        else if ((next->head & INUSE_BITS) == INUSE_BITS) {
                            fm->dvsizes = psize;
                            set_free_with_pinuse(p, psize, next);
                            goto postaction;
                        }
                    }
                }
            }
        }
```

```

    }
    else
        goto erroraction;
}
}

if (RTCHECK(ok_next(p, next) && ok_pinuse(next))) {
    if (!cinuse(next)) { /* consolidate forward */
        if (next == fm->top) {
            size_t tsize = fm->topsize += psize;
            fm->top = p;
            p->head = tsize | PINUSE_BIT;
            if (p == fm->dv) {
                fm->dv = 0;
                fm->dvsiz = 0;
            }
            if (should_trim(fm, tsize))
                sys_trim(fm, 0);
            goto postaction;
        }
        else if (next == fm->dv) {
            size_t dsiz = fm->dvsiz += psize;
            fm->dv = p;
            set_size_and_pinuse_of_free_chunk(p, dsiz);
            goto postaction;
        }
        else {
            size_t nsiz = chunksize(next);
            psize += nsiz;
            unlink_chunk(fm, next, nsiz);
            set_size_and_pinuse_of_free_chunk(p, psize);
            if (p == fm->dv) {
                fm->dvsiz = psize;
                goto postaction;
            }
        }
    }
}
else
    set_free_with_pinuse(p, psize, next);

if (is_small(psize)) {
    insert_small_chunk(fm, p, psize);
    check_free_chunk(fm, p);
}
else {
    tchunkptr tp = (tchunkptr)p;
    insert_large_chunk(fm, tp, psize);
    check_free_chunk(fm, p);
    if (--fm->release_checks == 0)
        release_unused_segments(fm);
}
goto postaction;
}
}

```

```
    erroraction:
        USAGE_ERROR_ACTION(fm, p);
    postaction:
        POSTACTION(fm);
}
}
#if !FOOTERS
#undef fm
#endif /* FOOTERS */
}
```