

Queens' College Cambridge

## Further Java



Alistair O'Brien

Department of Computer Science

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# 1 Networking

## 1.1 Sockets

- TCP connections are abstracted using UNIX *sockets*.
- Socket implemented `Closeable` interface  $\implies$  Socket is closed on exception w/ try-with-resources.
- Client sockets `Socket`:

— `Socket(String host, int port)`:

```
1  try (final Socket s = new Socket(host, port)) {
2      ...
3  } catch (UnknownHostException e) {
4      // The IP address of the host could not be determined
5      // using DNS.
6      ...
7  } catch (IllegalArgumentException e) {
8      // Argument port doesn't satisfy  $0 \leq port \leq 65535$ 
9      ...
10 } catch (IOException e) {
11     // An I/O error occurs when creating the socket
12     ...
13 } catch (SecurityException e) {
14     // Security manager exists and has insufficient
15     // permissions
16     ...
17 }
```

— `Socket(InetAddress address, int port)`:

```
1  try (final Socket s = new Socket(address, port)) {
2      ...
3  } catch (NullPointerException e) {
4      // The argument address is null
5      ...
6  } catch (IllegalArgumentException e) {
7      // Argument port doesn't satisfy  $0 \leq port \leq 65535$ 
8      ...
9  } catch (SecurityException e) {
10     // Security manager exists and has insufficient
11     // permissions
12 }
```

```

12     ...
13 } catch (IOException e) {
14     // An I/O error occurs when creating the socket
15     ...
16 }

```

- Server-side `ServerSocket`:

```

1  try (final ServerSocket serverSocket = new ServerSocket(port)) {
2      ...
3      while (true) {
4          // Blocking listening call: accept
5          // Exceptions:
6          // - IOException | SecurityException
7          Socket s = serverSocket.accept();
8          ...
9      }
10 } catch (IllegalArgumentException e) {
11     // Argument port doesn't satisfy  $0 \leq port \leq 65535$ 
12     ...
13 } catch (SecurityException e) {
14     // Security manager exists and has insufficient
15     // permissions
16     ...
17 } catch (IOException e) {
18     // An I/O error occurs when opening the socket
19     ...
20 }

```

- Input Streams:

```

1  try (InputStream in = s.getInputStream()) {
2      byte[] buffer = new byte[BUFFER_SIZE];
3      int length;
4
5      // Blocking read method
6      // Exceptions:
7      // - IOException: bytes cannot be read from socket e.g. ↘
8      //               closed
9      // - NullPointerException: buffer is null
10     while ((length != in.read(buffer)) != -1) {
11         // length bytes in array buffer
12         ...
13     }
14 } catch (IOException e) {
15     // Socket s is closed, or not connected
16 }

```

- Output Streams:

```

1  try (OutputStream out = s.getOutputStream()) {
2      // bytes is a byte[]
3
4      // Blocking write method e.g. full sliding window
5      // Exceptions:

```

```

6      // - IOException: bytes cannot be written to socket buffer e.\
      g. closed
7      out.write(bytes);
8
9      // Forces any buffered bytes in out to be written to socket. (\
      blocking)
10     out.flush();
11 } catch (IOException e) {
12     // Socket s is closed, or not connected
13 }

```

- Stream Decorators (and *Cheat Sheet*):

- `InputStreamReader`: Decorates a input stream, reading bytes and decoding them into characters (using a charset).

```
1  InputStreamReader in = new InputStreamReader(System.in);
```

- `BufferedReader`: Reads streams of characters, buffering them into strings / arrays. Improves performance for block-based IO (e.g. disk):

```

1  BufferedReader r = new BufferedReader(new InputStreamReader(\
    System.in));
2  String in;
3
4  while ((in = r.readLine()) != null) {
5      ...
6  }

```

- `FileReader`: Read contents of a text file:

```
1  FileReader fileReader = new FileReader(new File(path));
```

Reading all bytes of a file:

```

1  // or Paths.get(path)
2  byte[] bytes = Files.readAllBytes(new File(path).toPath());

```

## 1.2 Serialization

- Serialized classes implement the `Serializable` interface:

```

1  public class ChatMessage extends Message implements Serializable {\
    ... }

```

- Has no fields / methods. Simply tags class w/ serializable semantics.
- During serialization:

- Serialize primitives: e.g. `int`, `boolean`, etc.
- Serialize references to `o`: Recurse and serialize object `o`. If `o` doesn't implement `Serializable`, then `NotSerializableException` is thrown.
- Nonserializable state (e.g. a `Socket`) has the `transient` modifier:
 

```
1 private transient String dontSerializeMe;
```
- `static` fields are also not serialized.
- JVM uses versioning w/ `.class` files. Creating a UID for each class by hashing the `.class` file.
- JVM can determine whether two classes have the same version  $\implies$  detect whether object has a different version to on-disk `.class` version.
- Declare explicit UID using:
 

```
1 private static final long serialVersionUID = ...;
```

#### Advantages:

- Changing `.class` definition no-longer changes version implicitly  $\implies$  differing definitions:
  - \* removed fields are ignored
  - \* new fields are initialized w/ default value (0 / null / user-defined)
- UID only needs to be modified when:
  - `static` or `transient` fields are modified
  - Changing type of primitive field.
  - Changing inheritance hierarchy
- `ObjectInputStream` and `ObjectOutputStream` are used to serialize and deserialize objects, respectively.
 

```
1 try (ObjectInputStream in = new ObjectInputStream(s.getInputStream()
2     )) {
3     // Blocking readObject method
4     // Exceptions:
5     // - ClassNotFoundException: Class of serialized object isn't
6     //   loaded in JVM class loader
7     // - InvalidClassException: Deserialization error (see UID
8     //   conditions)
```

```

7      // - IOException: Error when reading stream
8      T o = (T) in.readObject();
9      ...
10 } catch (NullPointerException e) {
11     // s.getInputStream() is null
12     ...
13 } catch (IOException e) {
14     // Error when reading stream (s is closed or not connected)
15 }

1 try (ObjectOutputStream out = new ObjectOutputStream(s.
    getOutputStream())) {
2
3     // Blocking writeObject method
4     // Exceptions:
5     // - NotSerializableException: See above.
6     // - IOException: Error when writing to stream
7     out.writeObject(o);
8     ...
9 } catch (NullPointerException e) {
10     // s.getOutputStream() is null
11     ...
12 } catch (IOException e) {
13     // Error when writing to stream (s is closed or not connected)
14 }

```

### 1.2.1 Class Loaders

- A Java class loader is responsible for loading Java classes (compiled to bytecode, `.class` definitions) during the runtime of the JVM. Initially the default class loader is executed upon JVM initialization and loads classes from the local filesystem.
- **Problem:** Dynamically load classes during the runtime of the JVM from other locations e.g. socket
- **Solution:** Dynamic class loader. See `DynamicObjectInputStream` (Ticklet-2)
- **Problem:** Cannot guarantee non-byzantine behavior from server  $\implies$  send malicious Java bytecode containing static initializers that execute the malicious code (since static initializers are executed when the class is loaded).

### 1.2.2 Reflection

- Reflection permits inspection of classes at runtime:

```
1 Class<?> cls = o.getClass();
```

w/ generic wildcards (See IA OOP notes).

- List of *declared* fields (including private):

```
1 // Exceptions:
2 // - SecurityException: thrown if security manager doesn't have ↘
   accessDeclaredMembers
3 // permission
4 Field[] fs = cls.getDeclaredFields();
5 for (Field f : fs) {
6     ...
7 }
```

- Accessing a declared field:

```
1 // Suppresses all access modifier checks for f
2 // Exceptions:
3 // - SecurityException: Security manager doesn't have permission.
4 f.setAccessible(true);
5
6 ...
7
8 // Gets the value of the field in the reflected object o
9 // Exceptions:
10 // - IllegalAccessException: The underlying field f is ↘
   inaccessible (private, etc)
11 // - IllegalArgumentException: o doesn't have the field f or ↘
   class cls (to which f belongs)
12 // - NullPointerException: o is null and f is an instance field
13 f.get(o);
14
15 ...
```

- List of public methods:

```
1 Method[] ms = cls.getMethods();
2 for (Method m : ms) {
3     ...
4 }
```

- Get method from name:

```
1 try {
2     Method m = cls.getMethod(name);
3 } catch (NoSuchMethodException e) {
4     // Method w/ name name is not found
5     ...
6 } catch (NullPointerException e) {
7     // name is null
8     ...
9 }
```



- Invoke method:

```

1  // Exceptions:
2  // - IllegalAccessException: The underlying method m is ↘
   //   inaccessible
3  // - IllegalArgumentException: o is null, m is not in cls, number ↘
   //   of parameters and arguments differ
4  m.invoke(o, arg1, ..., argn);
```

### 1.2.3 Annotations

- Annotations add metadata to objects in a way which is accessible to the program at runtime.

- Prefixed w/ `@`:

```

1  @Override
2  public method foo() {
3      ...
4  }
```

- Annotations declared using `@interface` keyword:

```

1  public @interface SomeAnnotation {
2      ... // metadata fields
3  }
```

- Built-in annotations:

- `@Deprecated`: Annotated class, method, field is deprecated. Compiler warning if program uses deprecated code.
- `@Override`: Annotated method must override method in super-class.
- `@SuppressWarnings(warning)`: Compiler suppresses warnings it should generate. e.g. `@SuppressWarnings("deprecation")`

- Reflection support:

- Annotation: `@Retention(RetentionPolicy.RUNTIME)`
- `o.isAnnotationPresent(annotationClass)`: True if object *o* (class, field, method) has annotation w/ class `annotationClass`.
- `Annotation[] as = o.getAnnotations()`: Returns list of runtime annotations present on *o*.
- `Annotation a = o.getAnnotation(annotationClass)`: Returns annotation w/ class `annotationClass` on *o* if present. Otherwise `null`.

## 2 Concurrency

### 2.1 Multithreading

- Threads either extend `Thread` or implement `Runnable`:

– Extend `Thread` (`A extends Thread`):

1. Implement the `run()` method.
2. Instantiate `A` and execute `a.start()`

e.g. Idle thread: thread that schedules another thread and sleeps:

```
1 public class Idle extends Thread {
2
3     @Override
4     public void run() {
5         while (true) {
6             // Schedule new thread
7             Scheduler.schedule();
8
9             // Block idle until it's re-run by the scheduler
10            sleep(200);
11        }
12    }
13
14 }
```

– Implement `Runnable`:

1. Implement `Runnable` interface:

```
1 public interface Runnable {
2     void run();
3 }
```

`Runnable` is a *functional interface* (see IA OOP notes)  $\implies$  use a *lambda function*.

2. Create a new instance of the `Runnable` Class
3. Instantiate `Thread`, passing newly created `Runnable` instance as argument, then execute `start()`.

```

1 Thread idleThread = new Thread(() -> {
2     while (true) {
3         Scheduler.schedule();
4         Thread.sleep(200);
5     }
6 });
7 idleThread.start();

```

– **Advantages of Runnable:**

- \* No multiple inheritance  $\implies$  `extends Thread` prevents the class from extending any other classes.
  - \* `Runnable` yields a *task* abstraction  $\implies$  executed using `run()` or `Thread`. Hence reuse. Using `extends Thread` produces tightly coupled code which would benefit from being loosely coupled.
  - \* extending `Thread`  $\implies$  each extended thread will have a unique object w/ `run` implementation associated. implementing `Runnable`  $\implies$  many threads can share the same `Runnable` instance. Reduces overhead.
  - \* `Runnable` shared instance local resources for a group of threads. Easy state sharing w/ out global state.
- *Rule of Thumb:* extending a class  $\implies$  adding new behavior / state. If we're not modifying the `Thread` class then we should be using `Runnable`.
- Daemon threads: JVM will only exit when all threads running are daemons:

```
1 thread.setDaemon(true);
```

- Thread interrupts:

- Thread contains a `boolean` interrupted flag. `t.interrupt()` sets the flag to `true`.
- `Thread.interrupted()` returns `interrupted` of the current thread and sets it to `false`.
- Blocking thread methods, e.g `Thread.sleep`, check `interrupted` and throw `InterruptedException` if set:

```

1 public static void sleep(long millis) throws ↘
    InterruptedException {
2     while (/* still sleeping */) {
3         if (Thread.interrupted()) throw new ↘
            InterruptedException();
4     }
5 }

```

## 2.2 Monitors

- All objects implement an *implicit monitor*

**Definition 2.2.1. (Monitor)** A construct consisting of a lock  $\ell$  and a set of condition variables  $cv_1, \dots, cv_n$ .

**Definition 2.2.2. (Condition Variable)** A condition variable  $cv$ , or *condition queue*, is a queue  $Q$  associated w/ a condition  $c$  and lock  $\ell$  w/ operations:

- `wait( $cv, \ell$ )`: Thread  $t$  blocks until  $c$  is asserted. Atomically: releases  $\ell$ , moves  $t$  to  $Q$  and blocks. Once  $t$  is signaled,  $t$  is resumed and  $\ell$  is acquired.
- `signal( $cv$ )`: Called to indicate  $c$  is asserted. Dequeues thread  $t$  from  $Q$  and unblocks  $t$ .
- `broadcast( $cv$ )`: Called to indicate  $c$  is asserted. Dequeues *all* threads in  $Q$ .
- Java's monitors consist of a lock  $\ell$  and a single *non-blocking* condition variable  $cv$  (associated condition is defined by Programmer) w/ operations `wait()`, `notify()`, `notifyAll()`.
- Non-blocking condition variables: Signaling thread doesn't block, signaled thread is unblocked (and scheduled at the scheduler's whim).  
**Problem:** Condition  $c$  may not hold when signaled thread is scheduled.  
**Solution:** `while (!c) wait();`

### 2.2.1 Synchronized Statement

- Monitor locks are also acquired w/ **synchronized** statements:

```

1  synchronized (object) {
2      // Code w/ mutual exclusion on object
3      ...
4  }
```

Acquires *object*'s lock  $\ell$  and releases it after leaving the block.

- Syntactic Sugar:

<pre>public synchronized void lockMeUp() {     ... }</pre>	$\triangleq$	<pre>public void lockMeUp() {     synchronized (this) {         ...     } }</pre>
--	--------------	---

- Static methods are associated w/ a *class* lock  $cls.\ell$  not object lock  $o.\ell$ .

## 2.2.2 Deadlock

**Definition 2.2.3. (Deadlock)** A set of threads  $t_1, \dots, t_n$  are deadlocked if the following conditions hold:

- (i) **Mutual Exclusion:** Only a single thread  $t_i$  can use a resource  $R_j$  at a time.
- (ii) **Hold & Wait:** A thread  $t_i$  is holding the resources  $R_a^i, \dots, R_b^i$  is waiting to acquire resource  $R_c^j$  held by thread  $t_j$ .
- (iii) **No Preemption:** A resource  $R^i$  can only be released by the thread  $t_i$  holding it.
- (iv) **Circular Wait:** A cycle of thread requests exists:  $t_1 \rightarrow t_2 \rightarrow \dots \rightarrow t_1$

- **Problem:** Preventing Deadlock.
- **Idea:** Eliminate one (or more) conditions required for deadline  $\implies$  deadlock can never occur.
- **Solutions:**
  - **Mutual Exclusion:** Not required for shared resources (unsafe : ( ). Mutual exclusion must only hold for non-sharable resources.
  - **Hold & Wait:** Require that when threads request resource, they don't hold any other resources. Require threads to request and allocate all it's resource at once. Results in low resource utilization and difficult to know maximal resource set (in advance).
  - **No Preemption:** If thread  $t_i$  is holding resources ( $R_j^i$ ) and requests  $R_k$  and cannot acquire  $R_k$ , then ( $R_j^i$ ) are released. Pre-empted resources added to list of resources that  $t_i$  is waiting on.  $t_i$  is only scheduled again iff all old and new resources can be acquired.

- **Circular Wait:** Impose a partial ordering of all resources, requiring each thread  $t_i$  to acquire  $(R_j)$  in order.
- *Rule of Thumb:* **Circular Wait** solution is used. Define partial ordering  $\prec$  on objects.