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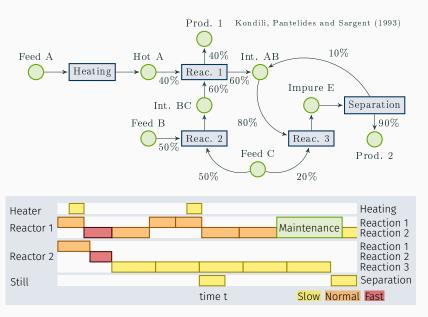
Data-driven optimization of processes with degrading equipment

Johannes Wiebe¹, Inês Cecílio², Ruth Misener¹ Wednesday 1st August, 2018

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Motivation: Why degradation matters



```
\begin{array}{ccc} \min & \cos(\boldsymbol{x}, \boldsymbol{m} &) \\ \text{s.t.} & \operatorname{process} \operatorname{model}(\boldsymbol{x}, \boldsymbol{m} &) & \text{(eg. balance equations)} \\ & & \operatorname{maintenance} \operatorname{model}(\boldsymbol{x}, \boldsymbol{m} &) & \text{(eg. types of maint.)} \end{array}
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where \boldsymbol{x} are process variables, \boldsymbol{m} are maintenance variables

```
\begin{array}{ll} \min \limits_{\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}} & \operatorname{cost}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}) \\ \text{s.t.} & \operatorname{process\ model}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}) & (\text{eg. balance\ equations}) \\ & \operatorname{maintenance\ model}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}) & (\text{eg. types\ of\ maint.}) \\ & \operatorname{health\ model}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}), & (\text{eq. prognosis\ model}) \end{array}
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where x are process variables, m are maintenance variables, and h are health related variables.

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Related Work

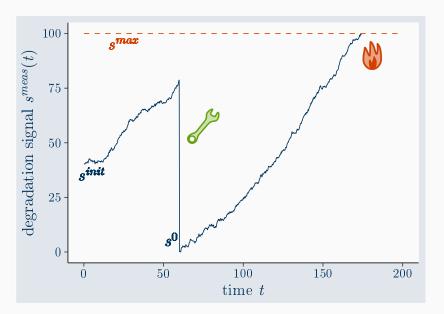
Vassiliads and Pistikopoulos (2001); Liu, Yahia and Papageorgiou (2014); Xenos, et int., Thornhill (2016); Aguirre and Papageorgiou (2018); Biondi, Sand and Harjunkoski (2017); Yildirim, Gebraeel and Sun (2017); Başçiftci, Ahmed, Gebraeel and Yildirim (2018)

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\begin{array}{ll} \min \limits_{\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}} & \operatorname{cost}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}) \\ \text{s.t.} & \operatorname{process\ model}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}) & (\text{eg. balance\ equations}) \\ & \operatorname{maintenance\ model}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}) & (\text{eg. types\ of\ maint.}) \\ & \operatorname{health\ model}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h}), & (\text{eq. prognosis\ model}) \end{array}
```

where x are process variables, m are maintenance variables, and h are health related variables.

Idea

Combine process level MI(N)LP scheduling & planning with more sophisticated (stochastic) degradation modelling and robust optimization.



The degradation signal $s^{meas}(t)$ can be modelled by a stochastic process :

$$S(t) = \{S_t : t \in T\},\$$

where S_t is a random variable (Alaswad and Xiang, 2017).

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Often used: Lévy type processes (Applebaum, 2004)

- · Independent increments: $S_{t_2} S_{t_1}, ..., S_{t_n} S_{t_{n-1}}$ are independent for any $0 < t_1 < t_2 < ... < t_n < \infty$
- Stationary increments: $S_t S_s$ and S_{t-s} have the same distribution for any s < t

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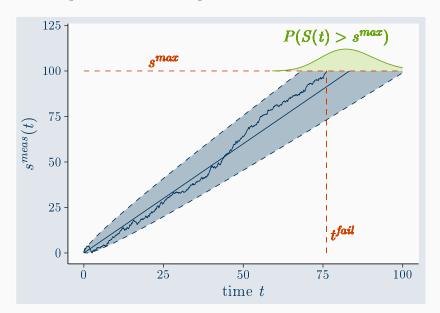
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- Stationary increments: $S_t S_s$ and S_{t-s} have the same distribution for any s < t

Therefore $S_t - S_{t-\Delta t} = D \sim \mathcal{D}(\Theta, \Delta t)$, where Θ are parameters of distribution \mathcal{D} .



A health model based on Lévy processes

Assumption

The health of each unit j can be described by a Lévy process $S_j(t)$ with increments $S_{j,t} - S_{j,t-\Delta t} = D_j \sim \mathcal{D}_j(\Theta, \Delta t)$.

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$$\min_{x,m,h} \quad \cos(x,m,h)$$

s.t. process model(x, m, h)

 $\text{maintenance model}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h})$

$$S_{j,t} \le s_j^{max} \qquad \forall t, j \in J$$

$$S_{j,t} = \begin{cases} S_{j,t-1} + D_j & , & \text{if } m_{j,t} = 0 \\ s_j^0, & \text{otherwise} \end{cases} \qquad \forall t, j \in J$$

where $m_{j,t} = 1$ if maintenance is performed on unit j at time t.

Accounting for effects of process variables

Assumption (Liao and Tian, 2013)

All relevant operating variables are piecewise constant – i.e. the process has a set of discrete operating modes $k \in K$.

where $x_{j,k,t} = 1$ if unit j operates in mode k at time t.

Deriving a robust counterpart (Lappas and Gounaris, 2016)

Idea

Replace random variables $D_{j,k}$ and $S_{j,t}$ by uncertain parameter $\tilde{d}_{j,k} \in \mathcal{U}$ and second stage variable $s_{j,t}(\tilde{d}_{j,k})$.

$$\min_{\boldsymbol{x}.\boldsymbol{m}.\boldsymbol{h}} \quad \operatorname{cost}(\boldsymbol{x}, \boldsymbol{m}, \boldsymbol{h})$$

s.t. process $\operatorname{model}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h})$ maintenance $\operatorname{model}(\boldsymbol{x},\boldsymbol{m},\boldsymbol{h})$ $s_{j,t}(\tilde{\boldsymbol{d}}_{j,k}) \leq s_j^{max} \qquad \forall t,j \in J$

$$s_{j,t} = \begin{cases} s_{j,t-1} + \sum_{k \in \mathcal{K}} x_{j,k,t} \cdot \tilde{\boldsymbol{d}}_{j,k}, & \text{if } m_{j,t} = 0 \\ s_j^0, & \text{otherwise} \end{cases} \quad \forall t, j \in J$$

$$\forall \tilde{d}_{i,k} \in \mathcal{U}.$$

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$$\begin{aligned} & \underset{\boldsymbol{x}, \boldsymbol{m}, \boldsymbol{h}}{\min} & & \operatorname{cost}(\boldsymbol{x}, \boldsymbol{m}, \boldsymbol{h}) \\ & & \text{s.t.} & & \operatorname{process model}(\boldsymbol{x}, \boldsymbol{m}, \boldsymbol{h}) \\ & & & \text{maintenance model}(\boldsymbol{x}, \boldsymbol{m}, \boldsymbol{h}) \\ & & & s_{j,t}(\tilde{\boldsymbol{d}}_{j,k}) \leq s_{j}^{max} & & \forall t, j \in J \\ & & s_{j,t} = \begin{cases} s_{j,t-1} + \sum\limits_{k \in \mathcal{K}} x_{j,k,t} \cdot \tilde{\boldsymbol{d}}_{j,k} , & \text{if } m_{j,t} = 0 \\ s_{j}^{0}, & & \text{otherwise} \end{cases} & \forall t, j \in J \end{aligned}$$

 $\forall \tilde{d}_{j,k} \in \mathcal{U}$. Approximate $s_{j,t}(\tilde{d}_{j,k})$ by linear decision rule. Utilize Robust Optimization reformulation techniques.

How do we choose \mathcal{U} ?

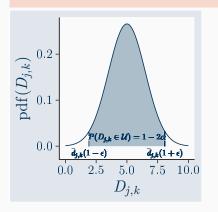
Assumption: \mathcal{U} is a box uncertainty set

$$\mathcal{U} = \{\tilde{d}_{j,k} | \bar{d}_{j,k} (1 - \epsilon_{j,k}) \le \tilde{d}_{j,k} \le \bar{d}_{j,k} (1 + \epsilon_{j,k})\}$$

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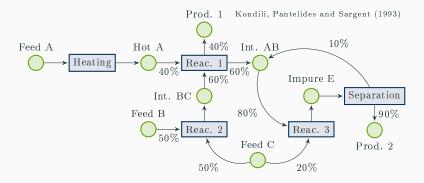


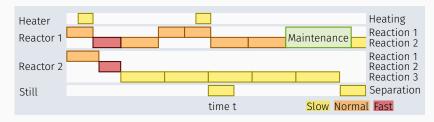
Choose $\epsilon_{j,k}$ from distribution $\mathcal{D}_{j,k}$:

$$\epsilon_{j,k} = 1 - F^{-1}(\alpha)/\bar{d}_{j,k}$$

Size of \mathcal{U} depends on a single parameter α !

Case study: State-Task-Network (Kondili et al., 1993)

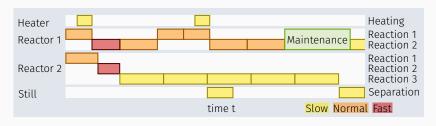




Case study: State-Task-Network (Kondili et al., 1993)

Biondi, Sand and Harjunkoski (2017) extend the STN to include...

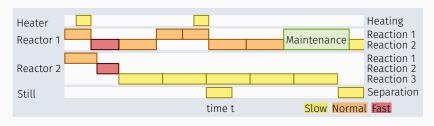
- · ...unit health and maintenance scheduling
- ...integrated scheduling and planning
- \cdot ...multiple operating modes per task



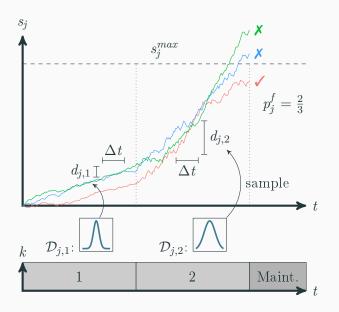
Case study: State-Task-Network (Kondili et al., 1993)

This work...

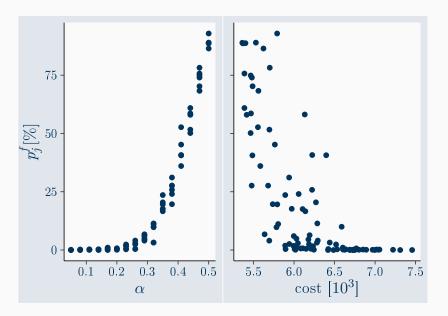
- ...replaces their deterministic health model by the proposed approach based on degradation modelling.
- ...utilizes robust optimization to obtain a solution that is likely to remain feasible.



Evaluating solution robustness



The price of robustness



Choosing α is its own optimization problem

We optimize α by solving

$$\min_{\alpha} c^*(\alpha) + \sum_{j} p_j^f(\alpha) \cdot c_j^f$$

- $c^*(\alpha)$ is the objective value of a MILP solution given α .
- $p_j^f(\alpha)$ is the corresponding probability of failure (of unit j).
- c_i^f is the cost of an unexpected failure.

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Idea: Use Bayesian Optimization (BO)

Both c^* and p_j^f can be viewed as expensive black box functions. BO is very suitable for this setting (Jones et al., 1998).

Saving time: a deterministic approximation

Assumption

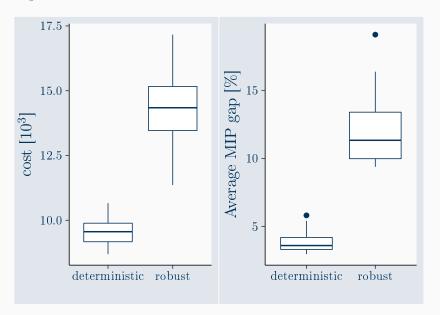
Only the health model depends on $\tilde{d}_{j,k}$ and $\tilde{d}_{j,k} \geq 0$.

Then we can prove that a solution to

$$\begin{aligned} & \underset{\boldsymbol{x}, \boldsymbol{m}, \boldsymbol{h}}{\min} & & \operatorname{cost}(\boldsymbol{x}, \boldsymbol{m}, \boldsymbol{h}) \\ & \text{s.t.} & & \operatorname{process model, maint. model}(\boldsymbol{x}, \boldsymbol{m}, \boldsymbol{h}) \\ & & s_{j,t} \leq s_{j}^{max} & & \forall t, j \in J \\ & & s_{j,t} = \begin{cases} s_{j,t-1} + \sum\limits_{k \in \mathcal{K}} x_{j,k,t} \cdot \boldsymbol{d}_{j,k}^{max}, & \text{if } m_{j,t} = 0 \\ s_{j}^{0}, & & \text{otherwise} \end{cases} & \forall t, j \in J \end{aligned}$$

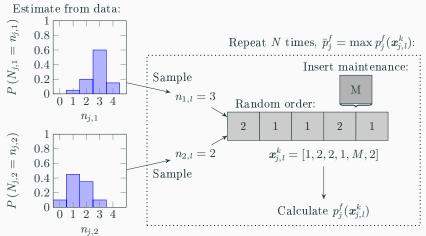
with $d_{j,k}^{max} = \max_{\mathcal{U}} \tilde{d}_{j,k}$ is also feasible in the robust problem.

Saving time: a deterministic approximation

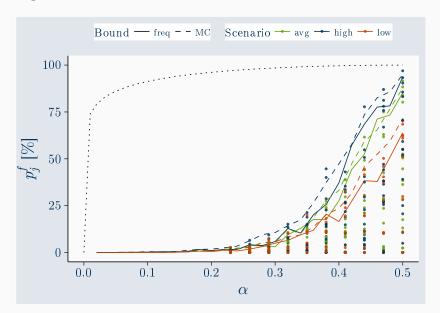


Saving time: data-driven approximations

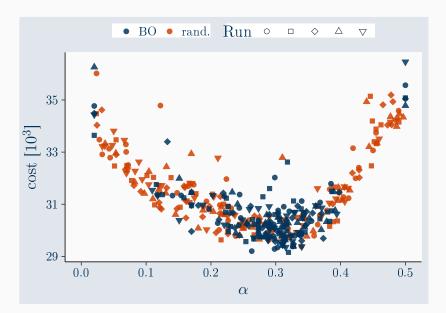
An upper bound on the probability of failure p_j^f can be estimated from data (using logistic regression).



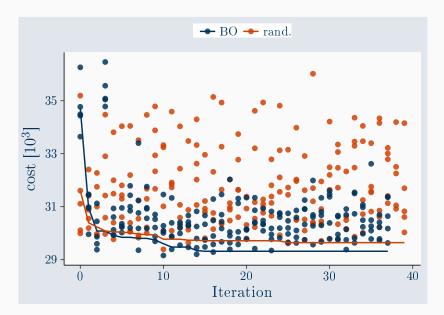
Saving time: data-driven approximations



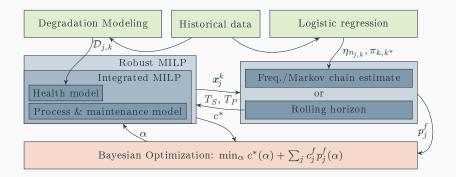
Bayesian Optimization



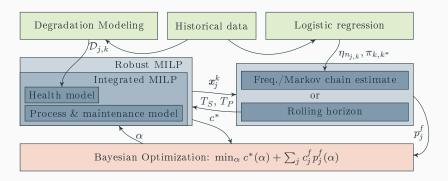
Bayesian Optimization



Conclusion



Conclusion



Thank You!

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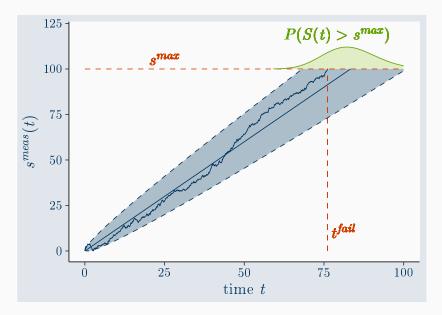
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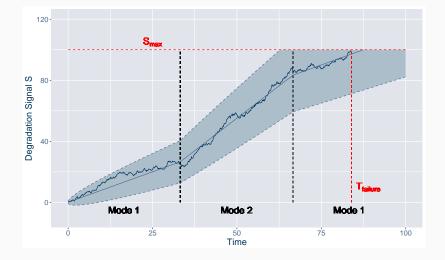
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Optimization for Opportunistic Maintenance and Operations in Wind Farms. IEEE Transactions on Power Systems, 32(6):4319-4328.

Degradation modelling



Degradation modelling with multiple operating modes



How does robust optimization work?

General idea

- · Make constraints hold for all values in \mathcal{U} : $\sum_{i} \tilde{a}_{ij} x_j \leq b_i, \forall \tilde{a}_{ij} \in \mathcal{U}$
- Reformulate semi-infinite constraint: $\sum_{i} a_{ij} x_j + \text{protection } (\mathcal{U}) \leq b_i$
- · How do we choose the right protection level?

Example: Soyster's method (worst case) [1973]

$$\max_{x_1, x_2} x_1 + x_2 \qquad \max_{x_1, x_2} x_1 + x_2$$
s.t. $\tilde{a}_{11}x_1 + \tilde{a}_{12}x_2 \le b_1$, $\forall \tilde{a}_{ij} \in \mathcal{U}$ s.t. $a_{11}x_1 + a_{12}x_2 + \sum_j \hat{a}_{ij} |x_j| \le b_1$
Given: $[a_{11}, a_{12}] = [1, 2], [\hat{a}_{11}, \hat{a}_{12}] = [0.1, 0.2], [b_1] = [2]$

Formulation

Scheduling

$$M_{j,t}S_{j,0} \leq S_{j,t} \leq S_{j,max} + M_{j,t} \cdot (S_{j,0} - S_{j,max}) \qquad \forall t, j \in J, D \in \mathcal{D}$$

$$S_{j,t} \geq S_{j,t-\Delta t} + \sum_{k} Z_{j,k,t}D_{j,k,t} + M_{j,t} \cdot (S_{j,0} - S_{j,max}) \qquad \forall t, j \in J, D \in \mathcal{D}$$

$$S_{j,t} \leq S_{j,t-\Delta t} + \sum_{k} Z_{j,k,t}D_{j,k,t} \qquad \forall t, j \in J, D \in \mathcal{D}$$

Planning

$$S_{j,t} \leq S_{j,max} \qquad \forall t, j \in J$$

$$S_{j,t} \geq S_{j,t-\Delta t} + \sum_{k} N_{j,k,t} D_{j,k,t} + M_{j,t} \cdot (S_{j,0} - S_{j,max}) \quad \forall t, j \in J$$

$$S_{j,t} \leq S_{j,t-\Delta t} + \sum_{k} N_{j,k,t} D_{j,k,t} \qquad \forall t, j \in J$$

Adjustable robust optimization

Affine decision rule

$$S_{j,t} = [S_{j,t}]_0 + \sum_k \sum_{t'=0}^t [S_{j,t}]_{k,t'} D_{j,k,t'}.$$

Size of toy problem

	deterministic	robust $D \neq f(t)$	robust $D = f(t)$	
# vars	913	3011	27719	
# binaries	338	338	338	
# constraints	1198	2356	13300	
time to solve [s]	2	0.3-10	0.3-10	
gap [%]	0	0	0	
scheduling periods	30	30	30	
planning periods	8	8	8	
task-unit-op. mode	6	6	6	
combinations				

Size of realistic problem

	deterministic	robust $D \neq f(t)$	robust $D = f(t)$
# vars	5389		397361
# binaries	2492		2492
# constraints	6798	6798	
time to solve [s]	7883		16756
gap [%]	3.62		31.02
scheduling periods	56	56	56
planning periods	24	24	24
task-unit-op. mode combinations	24	24	24

Deriving a robust counterpart

Replace $D_{j,k}$ by an uncertain parameter $\tilde{d}_{j,k}$ bounded by a set \mathcal{U} :

$$\begin{aligned} s_{j,t} &\leq s_j^{max} & \forall t,j \in J \\ s_{j,t} &= \begin{cases} s_{j,t-1} + \sum_{k \in \mathcal{K}} x_{j,k,t} \cdot \tilde{d}_{j,k}, & \text{if } m_{j,t} = 0 \\ s_j^0, & \text{otherwise} \end{cases} & \forall \tilde{d}_{j,k} \in \mathcal{U}, t,j \in J \end{aligned}$$

Reformulate:

$$m_{j,t}s_{j}^{0} \leq s_{j,t} \leq s_{j}^{max} + m_{j,t} \cdot (s_{j}^{0} - s_{j,max}) \qquad \forall t, j \in J, \tilde{d}_{j,k} \in \mathcal{U}$$

$$s_{j,t} \geq s_{j,t-\Delta t} + \sum_{k} x_{j,k,t} \tilde{d}_{j,k} + m_{j,t} \cdot (s_{j}^{0} - s_{j}^{max}) \qquad \forall t, j \in J, \tilde{d}_{j,k} \in \mathcal{U}$$

$$s_{j,t} \leq s_{j,t-\Delta t} + \sum_{k} x_{j,k,t} \tilde{d}_{j,k} \qquad \forall t, j \in J, \tilde{d}_{j,k} \in \mathcal{U},$$

Replace $s_{j,t}$ by linear decision rule $s_{j,t} = [s_{j,t}]_0 + \sum_k [s_{j,t}]_k \tilde{d}_{j,k}$.

Results: metrics data-driven approximation

instance	bound	${ m rms_all}$	${ m rms_max}$	p_out
toy	freq	8.00	1.53	29.40
toy	mc	10.41	3.08	21.27
P1	freq	12.61	3.52	17.54
P1	mc	17.25	4.39	9.62
P2	freq	7.40	2.31	18.08
P2	mc	13.68	4.98	10.13
P4	freq	9.17	3.27	47.78
P4	mc	11.43	2.84	32.50
P6	freq	18.75	8.94	12.17
P6	mc	20.84	10.09	10.98
all	freq	11.19	3.91	24.99
all	mc	14.72	5.08	16.90