CAD for VLSI

Multi-level Logic Optimization

Outline

- What is multi-level logic synthesis
- What are the specific goals
- Stepwise transformations
- Algebraic model
- Algebraic division
- Kernel theory and applications

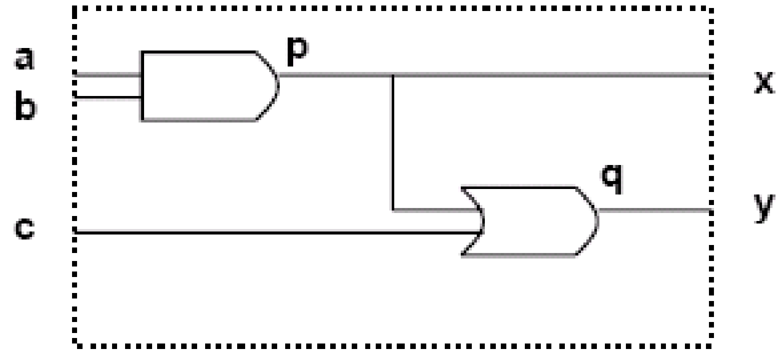
Motivation

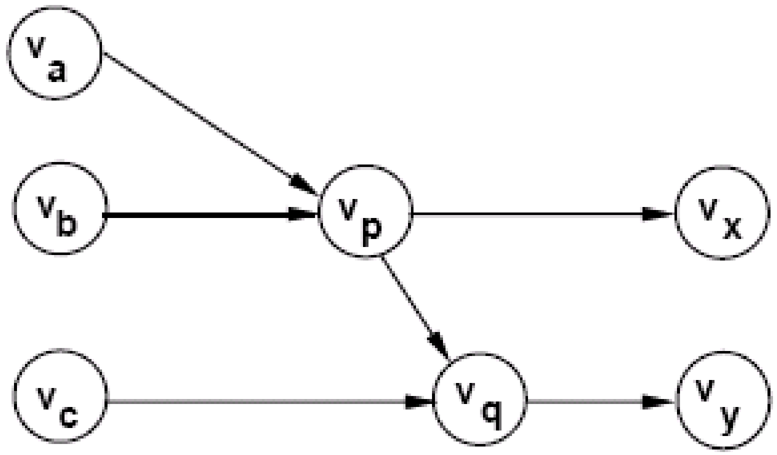
- Multiple-level logic networks
  - Semi-custom libraries
  - Logic gates versus macro-cells
    - More flexibility
    - Privilege specific paths on others
    - Better performance
- Applicable to a large variety of designs
- The importance of logic synthesis grew in parallel with the growth of foundries for the semi custom market

Circuit Model

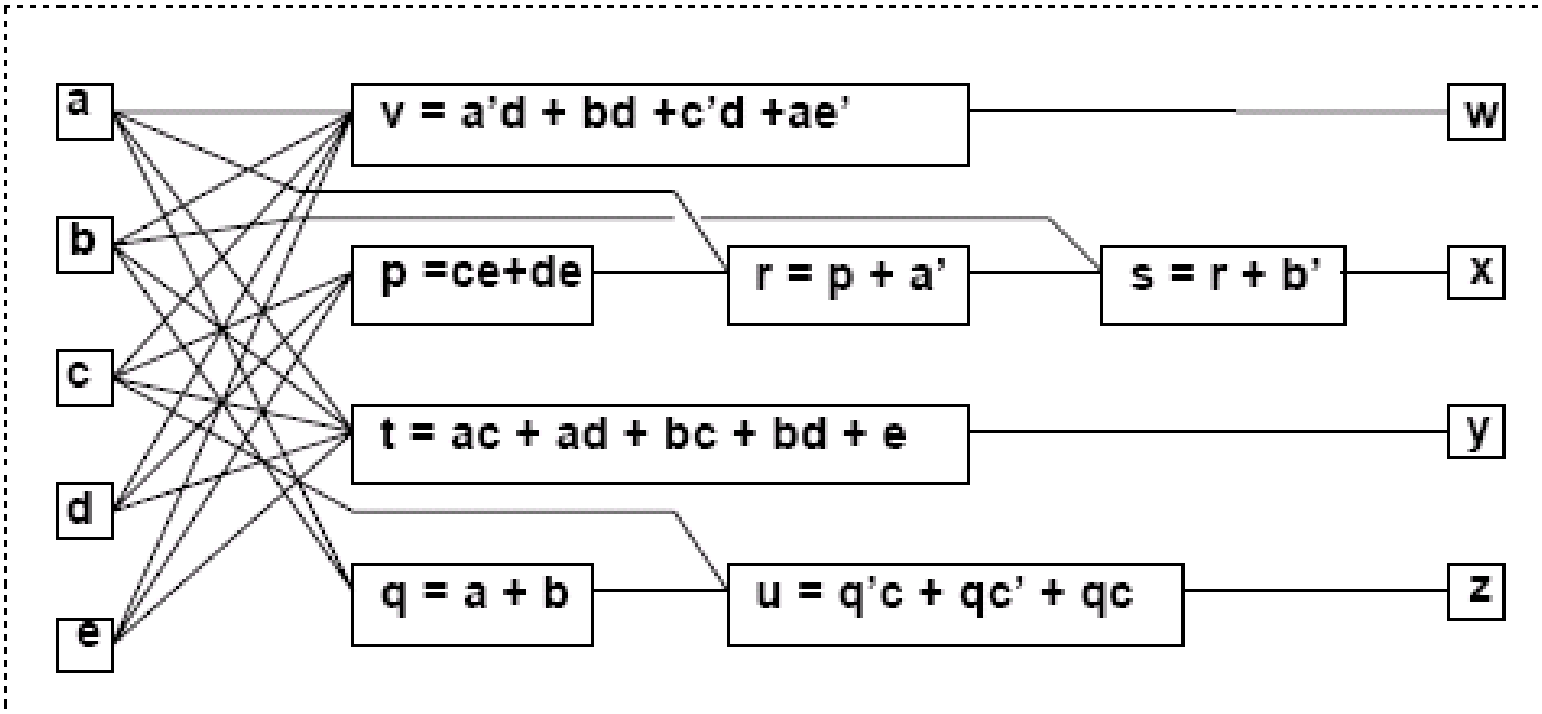
- Logic network
  - An interconnection of blocks
    - Each block modeled by a Boolean function
  - Usual restrictions:
    - Acyclic and memoryless
    - Single-output functions
- The model has a structural/behavioral semantics
  - The structure is induced by the interconnection
- Mapped network
  - Special case when the blocks correspond to library elements

Example of Mapped Network

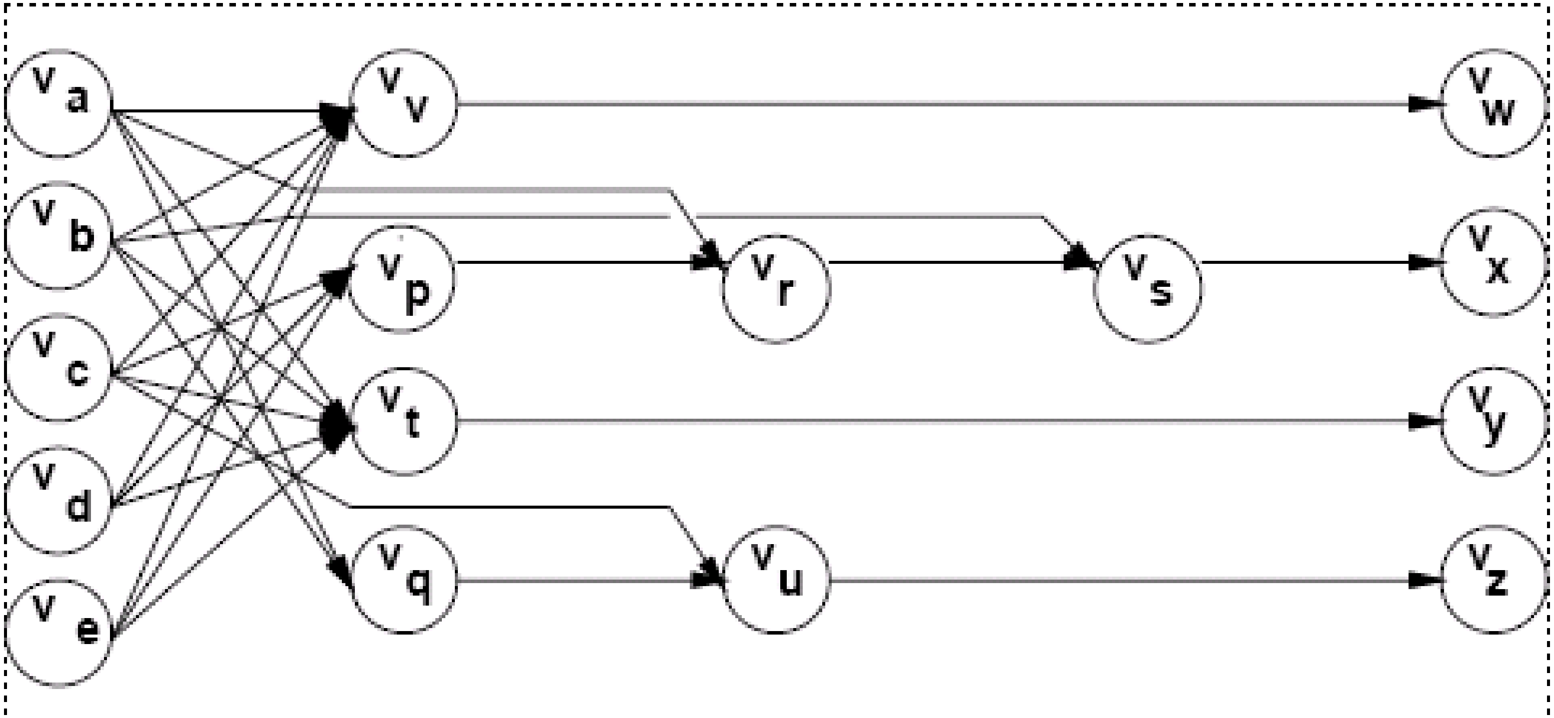




Example of General Network



Example of General Network Graph



Network Represented by Assignments

$$
\begin{aligned}
p &= ce + de \\
q &= a + b \\
r &= p + a' \\
s &= r + b' \\
t &= ac + ad + bc + bd + e \\
u &= q'c + qc' + qc \\
v &= a'd + bd + c'd + ae' \\
w &= v \\
x &= s \\
y &= t \\
z &= u
\end{aligned}
$$

Example of Terminal Behavior

- I/O functional behavior
  - Vector with as many entries as primary outputs
  - Each entry is a logic function

$$
f = \begin{bmatrix} a'd + bd + c'd + ae' \\ a' + b' + ce + de \\ ac + ad + bc + bd + e \\ a + b + c \end{bmatrix}
$$

Network Optimization

- Minimize maximum delay
  – Subject to area or power constraints
- Minimize area
  – Subject to delay constraints
- Minimize power consumption
  – Subject to timing constraints

Estimation

- Area:
  - Number of literals
    - Easy, widely accepted, good estimator
- Delay:
  - Number of stages
  - Gate delay models with wireloads
  - Sensitizable paths
- Power
  - Switching activity at each node
  - Capacitive loads

Problem Analysis

- Even the simplest problems are computationally hard
  - Ex: multi-input single-output network
- Few exact methods proposed
  - High complexity
  - Impractical
- Approximate optimization methods
  - Heuristic algorithms
  - Rule-based methods

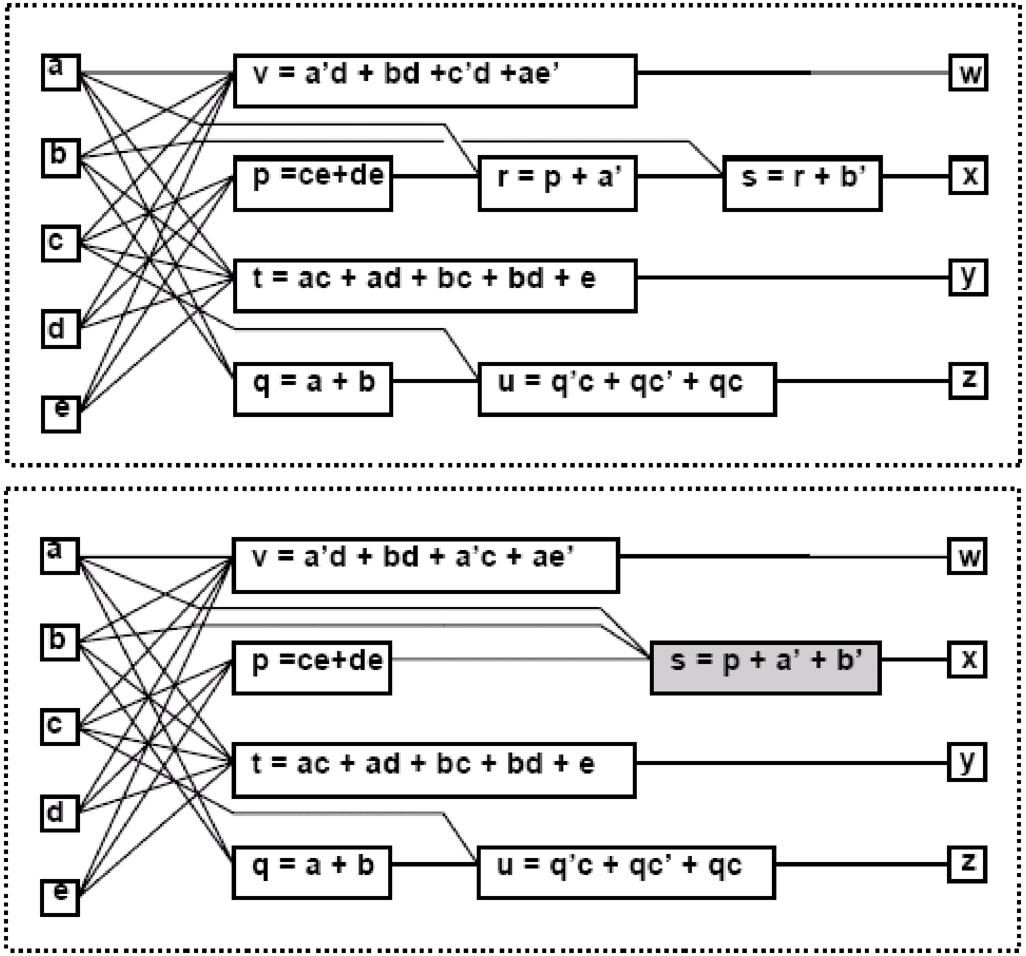
Strategies for Optimization

- Improve network step by step
  - Circuit transformations
- Preserve network I/O behavior
  - Exploit environment don't cares if desired
- Methods differ in:
  - Types of transformations applied
  - Selection and order of the transformations

Elimination

- Eliminate one function from the network
  - Similar to Gaussian elimination
- Perform variable substitution
- Example:
  - $s = r + b'$; $r = p + a'$;
  - $s = p + a' + b'$;

Example


a
b
c
d
e
v = a'd + bd +c'd +ae'
p =ce+de
r = p + a'
s = r + b'
t = ac + ad + bc + bd + e
q = a + b
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t = ac + ad + bc + bd + e
q = a + b
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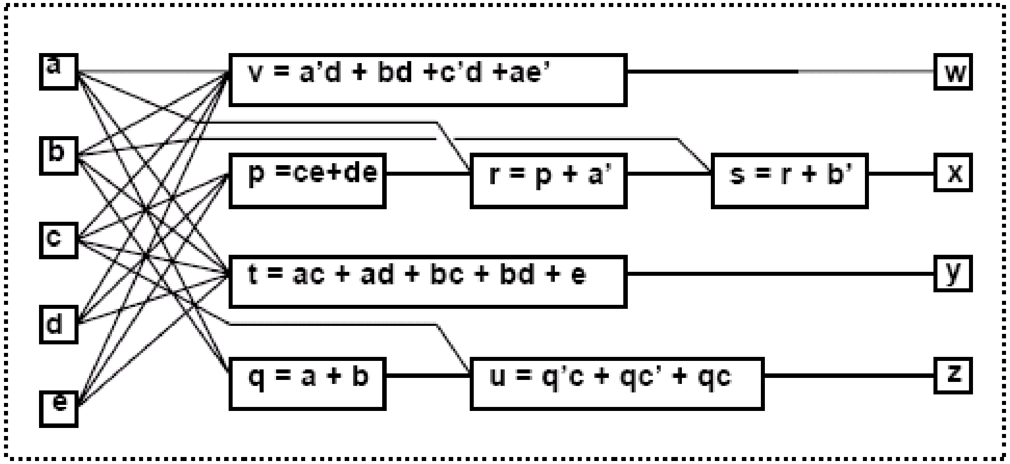
w
x
y
z

Decomposition

- Break a function into smaller ones
  - Opposite to elimination
- Introduce new variables/blocks into the network
- Example:
  - v = a'd + bd +c'd +ae'
  - j = a' + b + c';   v = jd + ae';

Example


a
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v = a'd + bd +c'd +ae'
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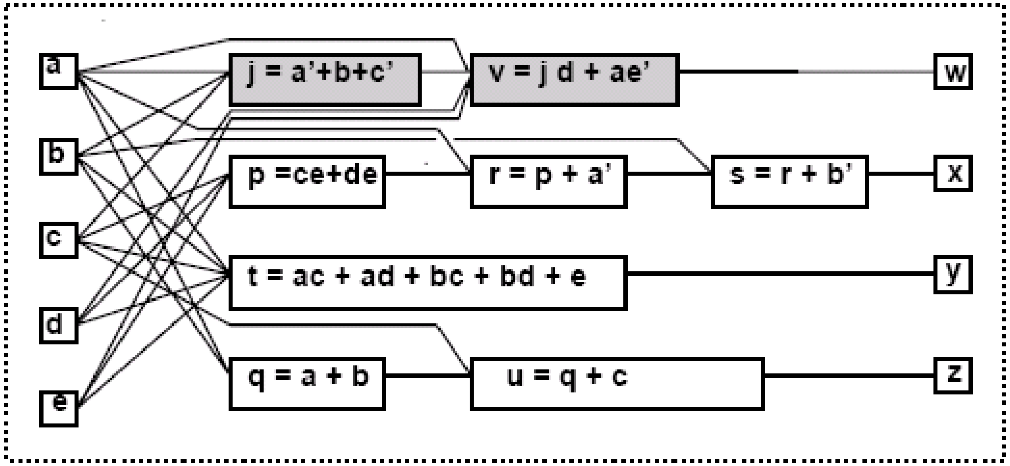
a
b
c
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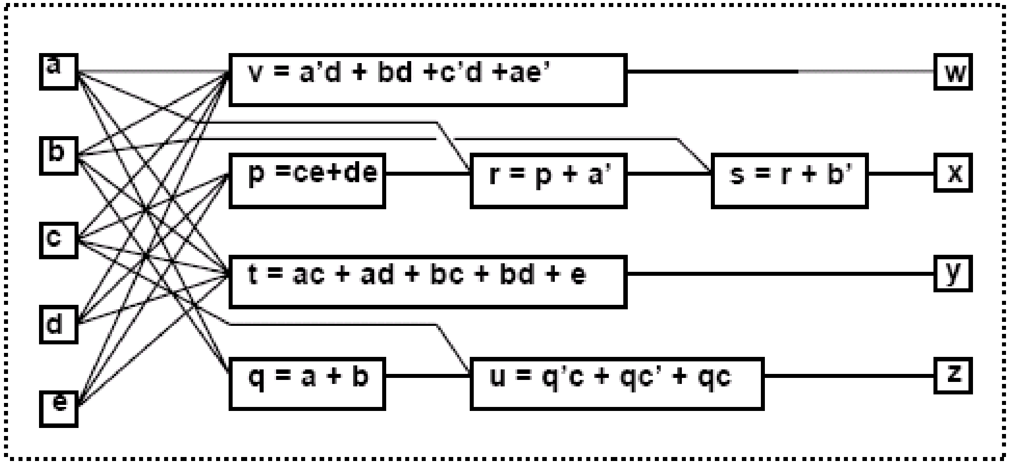
e
j = a'+b+c'
v = j d + ae'
p =ce+de
r = p + a'
s = r + b'
t = ac + ad + bc + bd + e
q = a + b
u = q + c
w
x
y
z

Extraction

- Find a common sub-expression of two (or more) expressions
  - Extract new sub-expression as new function
  - Introduce new block into the circuit
- Example
  - p = ce + de;   t = ac + ad + bc + bd + e;
  - p = (c + d) e;   t = (c + d) (a + b) + e;
  - k = c + d;   p = ke; t = ka + kb + e;

Example



a
b
c
d
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v = a'd + bd +c'd +ae'
p =ce+de
r = p + a'
s = r + b'
t = ac + ad + bc + bd + e
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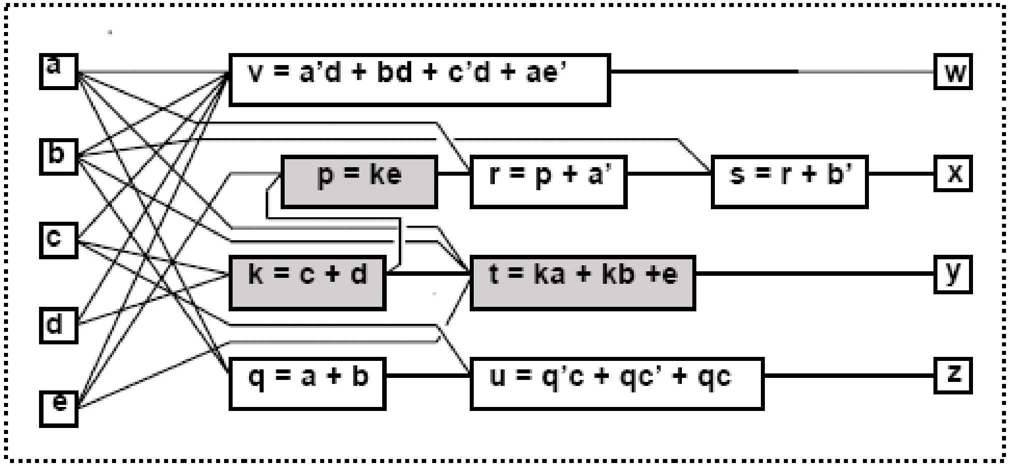
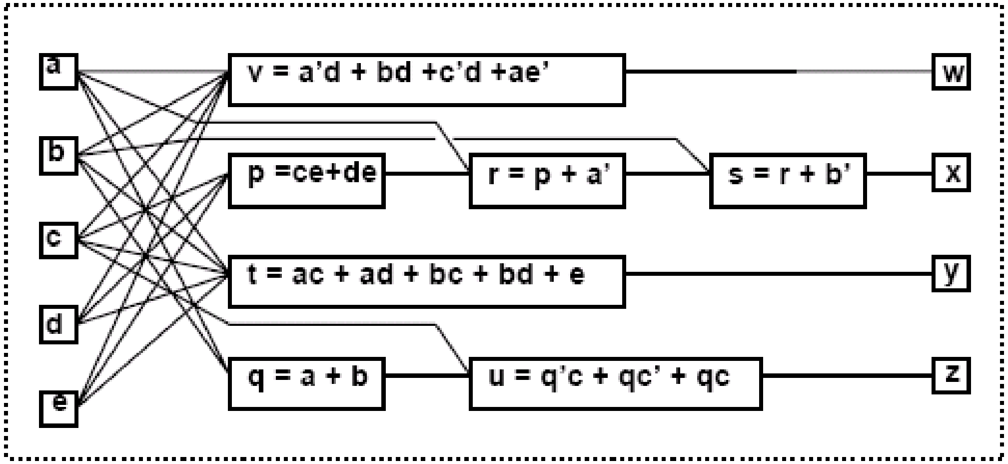


a
b
c
d
e
v = a'd + bd + c'd + ae'
p = ke
r = p + a'
s = r + b'
k = c + d
t = ka + kb +e
q = a + b
u = q'c + qc' + qc
w
x
y
z

Simplification

- Simplify local function
  - Use heuristic minimizer like Espresso
  - Modify fanin of target node
- Example:
  - u = q'c + qc' + qc;
  - u = q + c;

Example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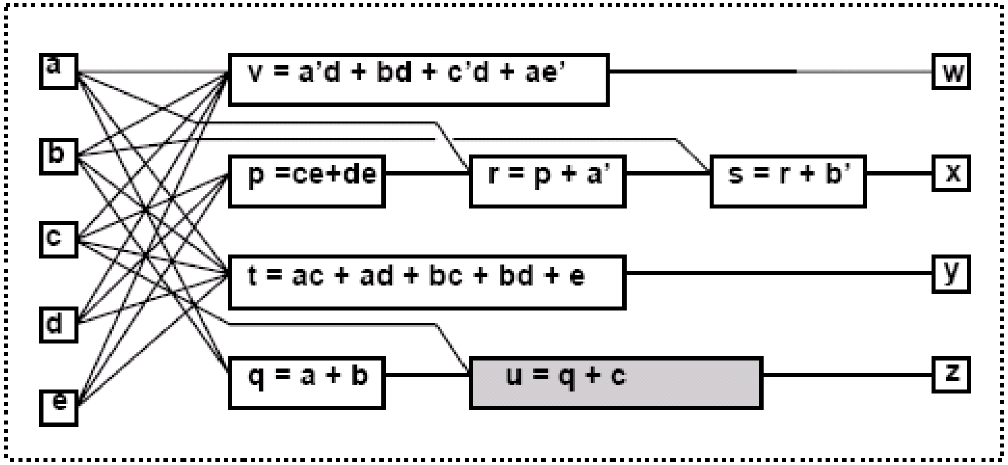
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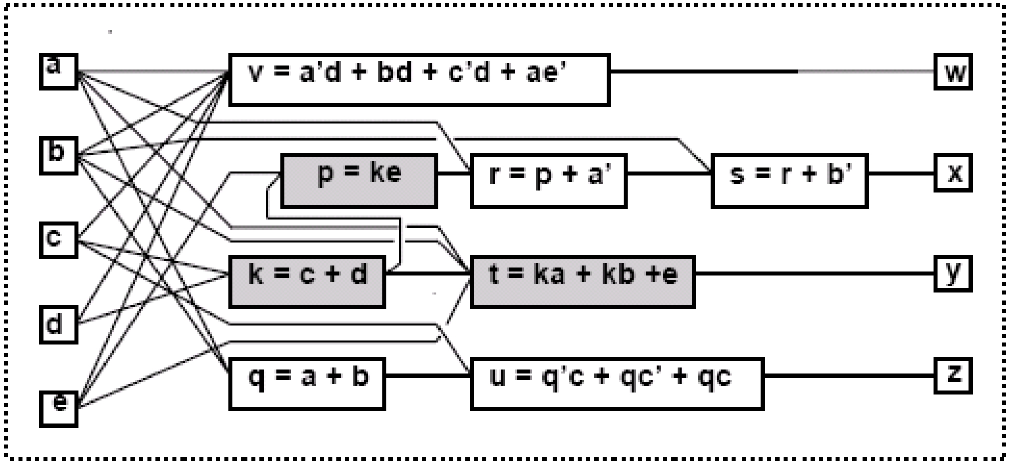
t = ac + ad + bc + bd + e
q = a + b
u = q + c
w
x
y
z

Substitution

- Simplify a local function by using an additional input that was not previously in its support set
- Example:
  - $t = ka + kb + e;$
  - $t = kq + e;$
  - Because $q = a + b$ is already part of the network

Example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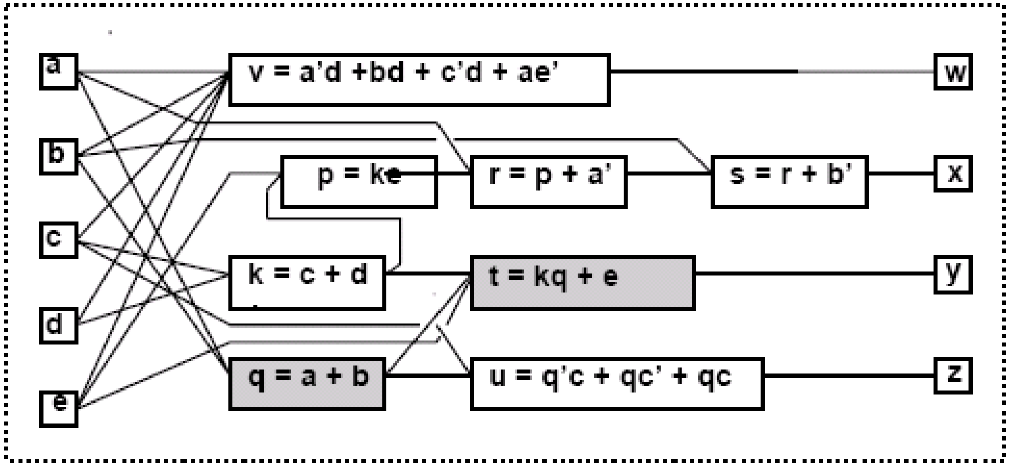
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p = ke
r = p + a'
s = r + b'
k = c + d
t = kq + e
q = a + b
u = q'c + qc' + qc
w
x
y
z

Example – Sequence of Transformations

j = a' + b + c

k = c + d

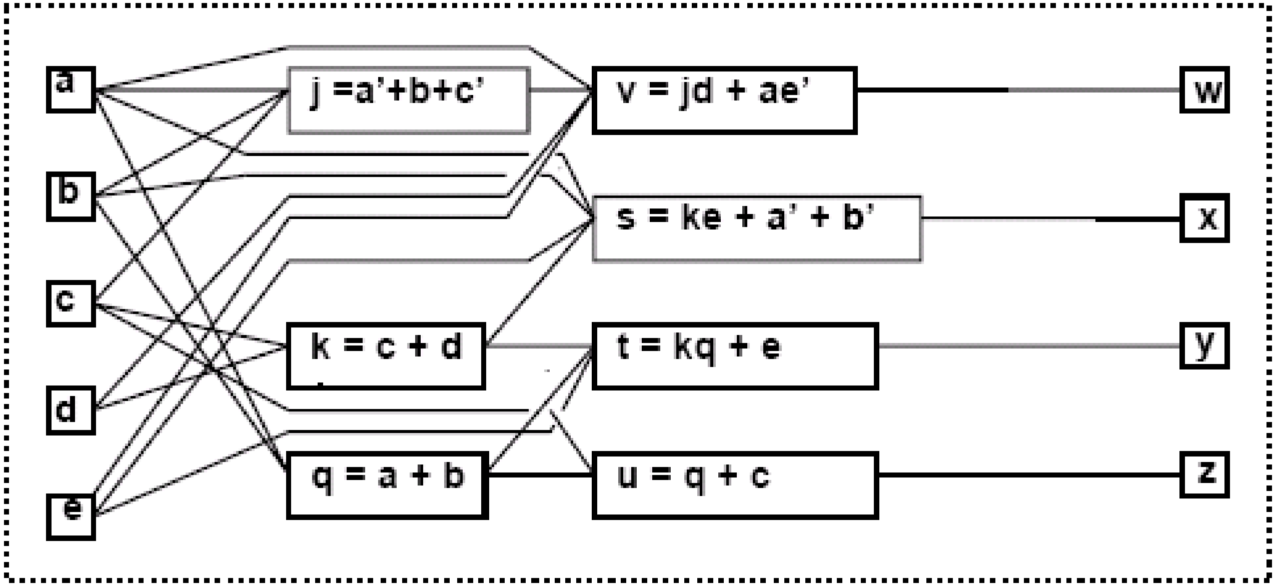
q = a + b

s = ke + a' + b'

t = kq + e

u = q + c

v = jd + ae'



Optimization Approaches

- Algorithmic approach
  - Define an algorithm for each transformation type
  - Algorithm is an operator on the network
  - Algorithms are sequenced by scripts
- Rule-based approach
  - Rule data base
    - Set of pattern pairs
  - Pattern replacement is driven by rules
- Most modern tools use the algorithmic approach to synthesis, even though rules are used to address specific issues

Boolean and Algebraic Methods

- Boolean methods for multilevel synthesis
  - Exploit properties of Boolean functions
  - Use don't care conditions
  - Computationally intensive
- Algebraic methods
  - Use polynomial abstraction of logic function
  - Simpler, faster, weaker
  - Widely used

Example

- Boolean substitution:
  - $h = a + bcd + c; \quad q = a + cd;$
  - $h = a + bq + c;$
  - Because $a + bq + c = a + b(a+cd) + c$
    $= a + bcd + c;$
- Algebraic substitution:
  - $t = ka + kb + e;$
  - $t = kq + e;$
  - Because $q = a + b;$

Outline

- What is multi-level logic synthesis
- What are the specific goals
- Stepwise transformations
- Algebraic model
- Algebraic division
- Kernel theory and applications

Algebraic Model

- Boolean algebra
  - Complement
  - Symmetric distribution laws
  - Don't care sets
- Algebraic models
  - Look at Boolean expressions as polynomials
  - Use sum of product forms
    - Minimal w.r.t. 1-cube containment
  - Use polynomial algebra

Algebraic Division

- Given two algebraic expressions
  - An expression divides algebraically the other
  - $f_{quotient} = f_{dividend} / f_{divisor}$ when:
  - $f_{dividend} = f_{divisor} f_{quotient} + f_{remainder}$
  - $f_{divisor} f_{quotient} \neq 0$
  - The support of $f_{divisor}$ and $f_{quotient}$ is disjoint
- Note that the $f_{quotient}$ and $f_{divisor}$ are interchangeable

Example

- Algebraic division
  - $f_{dividend} = ac + ad + bc + bd + e$
  - $f_{divisor} = a + b$
  - Then $f_{quotient} = c + d$ and $f_{remainder} = e$
    because $(a+b)\ (c+d) + e = f_{dividend}$
    and $\{a,b\} \cap \{c,d\} = \emptyset$
- Non-algebraic division:
  - $f_i = a + bc$ and $f_j = a+b$
  - Then $(a+b)\ (a+c) = f_i$
    but $\{a,b\} \cap \{a,c\} \neq \emptyset$

An Algorithm for Division

- Division can be performed in different way
  - Straightforward algorithm by literal sorting
    - Simple, quadratic complexity
  - Advanced algorithm using sorting
    - N-logN complexity
  - Typically algebraic division runs fast – small-sized problems
- Definitions
  - A = set of cubes C^A_j of the dividend, j = 1 ~ l
  - B = set of cubes C^B_i of the divisor, i = 1 ~ n
  - Q = quotient; R = remainder

An Algorithm for Division

```
ALGEBRAIC_DIVISION(A,B)
{
      for (i = 1 to n)
      {
            D = {C^A_j such that C^A_j ⊇ C^B_i };
            if (D == Ø) return(Ø,A);
            D_i = D with variables in sup(C^B_i) dropped;
            if i = 1
                        Q = D_i ;
            else
                        Q = Q ∩ D_i ;
      }
      R = A – Q x B;
      return(Q,R);
}
```

Example
$f_{dividend} = ac+ad+bc+bd+e; \quad f_{divisor} = a+b$

- A = {ac, ad, bc, bd, e} and B = {a, b}
  - i = 1:
  - C^B_1 = a, D = {ac, ad} and D_1 = {c, d}
    - Then Q = {c, d}
- i = 2 = n:
  - C^B_2 = b, D = {bc, bd} and D_2 = {c, d}
  - Then Q = {c, d} ∩ {c, d} = {c, d}
- Result:
  - Q = {c, d} and R = {e}
  - $f_{quotient}$ = c + d and $f_{remainder}$ = e

Theorem for Filtering

- Given algebraic expression f_i and f_j then f_i / f_j is empty when either:
  - f_j contains a variable not in f_i
  - f_j contains a cube whose support is not contained in that of any cube of f_i
  - f_j contains more terms than f_i
  - The count of any variable in f_j is higher than in f_i

Algebraic Substitution

- Consider expression pairs
- Apply division (in any order)
- If quotient is not void:
  - Evaluate area and delay gain
  - Substitute $f_{dividend}$ by $j \cdot f_{quotient} + f_{remainder}$ where j is the variable corresponding to $f_{divisor}$
    - i.e., use a single variable j to represent $f_{divisor}$
- Use filters based on previous theorem to reduce computation

Substitution Algorithm

```
SUBSTITUTE(Gn(V,E)){
    for (i = 1, 2, ..., |V|){
        for (j = 1, 2 ,..., |V|; j ≠ i){
            A = set of cubes of f_i;
            B = set of cubes of f_j;
            if (A, B pass the filter test){
                (Q, R) = ALGEBRAIC_DIVISION(A,B);
                if (Q ≠ Ø){
                    f_quotient = sum of cubes of Q;
                    f_remainder = sum of cubes of R;
                    if (substitution is favorable)
                        f_i = j  f_quotient + f_remainder;
                }
            }
        }
    }
}
```

Extraction

- Search for common sub-expressions
  - Single-cube extraction
  - Multiple-cube extraction (kernel extraction)
- Search for appropriate divisors
- Extraction is still done using the original kernel theory of Brayton and others [IBM]

Definitions

- Cube-free expression
  - Expression that cannot be factored by a cube
  - Example:
    - a + bc is cube free
    - abc and ab + ac are not
- Kernel of an expression
  - Cube-free quotient of the expression divided by a cube (The cube is called co-kernel)
  - Note that since divisors and quotients are interchangeable, kernels are just a subset of divisors
- Kernel set of an expression f is denoted by K(f)

Example

- f = ace + bce + de + g
- Trivial kernel search:
  - Divide f by a. Get ce. Not cube free
  - Divide f by b. Get ce. Not cube free
  - Divide f by c. Get ae + be. Not cube free
  - Divide f by ce. Get a + b. Cube free. KERNEL!
  - Divide f by d. Get e. Not cube free
  - Divide f by e. Get ac + bc + d. Cube free. KERNEL!
  - Divide f by g. Get 1. Not cube free
  - Divide f by 1. Get f. Cube free. KERNEL!
- K(f) ={(a+b), (ac+bc+d), (ace+bce+de+g)}
- CoK(f) = {ce, e, 1}

Theorem Brayton and McMullen

- Two expressions f_a and f_b have a common multiple-cube divisor f_d if and only if
  - There exist kernels k_a in $K(f_a)$ and k_b in $K(f_b)$ such that
    f_d is the sum of two (or more) cubes in $k_a \cap k_b$
- Consequences
  - If kernel intersection is void, then the search for common sub-expression can be dropped
  - If an expression has no kernels, it can be dropped from consideration
  - The kernel intersection is the basis for constructing the expression to extract

Example

- f_x = ace + bce + de + g
- f_y = ad + bd + cde + ge
- f_z = abc
- $K(f_x)$ = { (a+b); (ac+bc+d); (ace+bce+de+g) }
- $K(f_y)$ = { (a+b+ce); (cd+g); (ad+bd+cde+ge) }
- The kernel set of f_z is empty
- Select intersection (a+b)
  - f_w = a + b
  - f_x = wce + de + g
  - f_y = wd + cde + ge
  - f_z = abc

Kernel Set Computation

- Naïve method
  - Divide function by the elements of the power set of its support set
  - Weed out non cube-free quotients
- Smart way
  - Use recursion
    - Kernels of kernels are kernels
  - Exploit commutativity of multiplication

Recursive Algorithm

- The recursive algorithm is the first one proposed for kernel computation and still outperforms others
- It will be explained in two steps
  – R_KERNELS (with no pointer) to understand the concept
  – KERNELS (Complete algorithm)
- The algorithms use a subroutine for filtering
  – CUBES (f, C) which returns the cubes of f whose literals include those of cube C
  – Example: f = ace + bce + de + g
    CUBES(f, ce) = ace + bce

Simple Recursive Algorithm

```
Find maximal cube C such that CUBES(f, 1) = CUBES(f, C);
R_KERNELS(f / C);

R_KERNELS(f) {
    K = Ø;
    foreach variable x ∈ sup(f) {
        if (|CUBES(f, x)| ≥ 2) {
            C = maximal cube containing x,  s.t. CUBES(f, C) = CUBES(f, x);
            K = K ∪ R_KERNELS(f / C);
        }
    }
    K = K ∪ f;
    return(K);
}
```

Analysis

- The recursive algorithm does some redundant computation in the recursion
  - Example
    - Divide by a and then by b
    - Divide by b and then by a
  - Obtain duplicate kernels
- Improvement
  - Exploit commutativity of multiplication
  - Keep a pointer to the literals used so far

Recursive Kernel Computation

Find maximal cube C *such that* CUBES(f, 1) = CUBES(f, C);
KERNELS(f / C, 1);

```
KERNELS(f, j) {
    K = Ø;
    for i = j to n  {
        if (|CUBES(f, x_i)| ≥ 2) {
                C = maximal cube containing x_i,
                s.t. CUBES(f, C) = CUBES(f, x_i);
                if (C has no variable x_k, k < i )
                        K = K ∪ KERNELS(f / C, i+1);
        }
    }
    K = K ∪ f;
    return(K);
}
```

Example

- f = ace + bce + de + g
  - Literals a and b. No action required
  - Literal c. Select cube ce
    - Recursive call with argument f/ce= a + b. Pointer j = 3+1
    - Call considers variables {d, e, g}. No kernel.
    - Adds a + b to the kernel set at the last step.
  - Literal d. No action required.
  - Literal e. Select cube e
    - Recursive call with argument f/e = ac + bc + d. Pointer j = 5+1
      - Redundant computation of variable {c} is ignored since j = 6
    - Call considers variables {g}. No Kernel
    - Adds ac + bc + d to the kernel set at the last step of recursion
  - Literal g. No action required
  - Add f = ace + bce + de + g to kernel set
  - K(f) = {(ace + bce + de + g), (ac + bc + d), (a + b)}

Matrix Representation of Kernels

- f = ace + bce + de + g
- Incidence matrix
  - Cubes vs. variables
- Rectangle
  - Subset of rows/columns with all entries equal to 1
- Prime rectangle
  - Rectangle not included in another rectangle
- A co-kernel is a prime rectangle with at least two rows
- Example:
  - Prime rectangle ({1, 2}, {3, 5}), ({1, 2, 3}, {5})
  - Co-kernel ce, e

	var	a	b	c	d	e	g
cube	$R\backslash C$	1	2	3	4	5	6
ace	1	1	0	1	0	1	0
bce	2	0	1	1	0	1	0
de	3	0	0	0	1	1	0
g	4	0	0	0	0	0	1

Application of Kernel Methods

- Single cube extraction
  - – Extract one cube from two (or more) sub-expressions [Brayton]
- Multiple-cube extraction
  - – Extract a multiple-cube expression [Brayton]
- Double-cube extraction
  - – Newer, fast and efficient routine [Rajski]
- Kernel-based decomposition

Single-cube Extraction

- Form an auxiliary expression, which is the union (sum) of all local expression
- Find the largest co-kernel
  - Corresponding kernel must belong to two (or more) different expressions
  - Use additional variables to tag the expressions
- Extract chosen co-kernel
- The problem can be well visualized by a matrix representation and the extraction of a prime rectangle

Example

- Expressions:
  - $f_x = ace + bce + de + g$
  - $f_s = cde + b$
- Auxiliary function:
  - $f_{aux} = ace + bce + de + g + cde + b$
- Tagging:
  - $f_{aux} = xace + xbce + xde + xg + scde + sb$
- Co-kernel: ce
- After cube extraction
  - $f_z = ce$
  - $f_x = z\,(a+b) + de + g$
  - $f_s = zd + b$

		var	a	b	c	d	e	g
cube	ID	$R\backslash C$	1	2	3	4	5	6
ace	x	1	1	0	1	0	1	0
bce	x	2	0	1	1	0	1	0
de	x	3	0	0	0	1	1	0
g	x	4	0	0	0	0	0	1
cde	s	5	0	0	1	1	1	0
b	s	6	0	1	0	0	0	0

Multiple-cube Extraction

- We need a cube/kernel matrix
  - Relabel cubes by new variables
  - Kernels are now cubes in these new variables
- Find a prime rectangle
- Equivalently, find a co-kernel of the auxiliary expression that is the sum of the relabeled expressions

Example

- $f = ace + bce$
  - $K(f) = \{(a+b)\}$
- $g = ae + be + d$
  - $K(g) = \{(a+b); (ae + be + d)\}$
- Relabeling: $x_a=a$; $x_b=b$; $x_{ae}=ae$; $x_{be}=be$; $x_d=d$
  - Then $K(f) = \{\{x_a, x_b\}\}$ and $K(g) = \{\{x_a, x_b\}, \{x_{ae}, x_{be}, x_d\}\}$
  - $f_{aux} = y_f\, x_a\, x_b \;+\; y_g\, x_a\, x_b \;+\; y_g\, x_{ae}\, x_{be}\, x_d$
  - $CoK(f_{aux}) = x_a\, x_b$
- Go back to original variables
  - Extract $(a + b)$ from f and g

Double-cube Extraction

- Restrict extraction to:
  - Double-cube kernels
  - Single-cube kernels with two literals
  - Consider concurrently their complements
- Properties
  - These small kernels can be computed efficiently
  - Circuit testability is preserved
  - Very efficient in reducing network

Kernel-based Decomposition

- There are many different ways of performing decomposition
  - Several classic approaches (ex: Ashenhurst & Curtis)
- Algebraic decomposition
  - Find good algebraic divisors
  - Use kernels and decomposition recursively

Example

- Decompose $f = ace + bce + de + g$
- Select kernel $ac + bc + d$
- Decompose as: $f = te + g$; $t = ac + bc + d$
- Recur on quotient t
- Select kernel $a + b$
- Decompose $t = sc + d$; $s = a + b$; $f = te + g$;

Summary Algebraic Methods

- Algebraic methods abstract functions as polynomials
  - Polynomial division
- Methods are fast and widely applicable
- Algebraic methods miss opportunities for optimization
  - As compared to Boolean methods
- Algebraic transformations are reversible
  - Ease transformations back and forward to trade off area and speed