

Programming Project Checkpoint#4

姚林飛 109006243

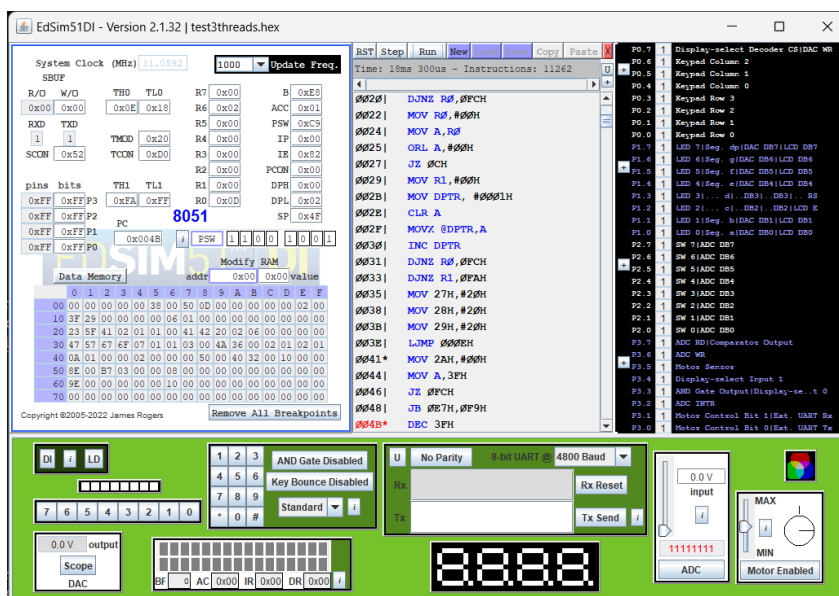
[1] Typescript for compilation

```
Ubuntu
(base) [2022-12-13 16:31:55] jojoelt@ZephyG14:~/OS-edsim/chkpnt4 (master *)
$ conda activate sdcc
(sdcc) [2022-12-13 16:31:59] jojoelt@ZephyG14:~/OS-edsim/chkpnt4 (master *)
$ make clean
rm *.hex *.ihx *.lnk *.lst *.map *.mem *.rel *.rst *.sym
rm: cannot remove '*.hex': No such file or directory
rm: cannot remove '*.ihx': No such file or directory
rm: cannot remove '*.lnk': No such file or directory
rm: cannot remove '*.lst': No such file or directory
rm: cannot remove '*.map': No such file or directory
rm: cannot remove '*.mem': No such file or directory
rm: cannot remove '*.rel': No such file or directory
rm: cannot remove '*.rst': No such file or directory
rm: cannot remove '*.sym': No such file or directory
make: *** [Makefile:25: clean] Error 1
(sdcc) [2022-12-13 16:32:01] jojoelt@ZephyG14:~/OS-edsim/chkpnt4 (master *)
$ make
sdcc -c test3threads.c
sdcc -c preemptive.c
preemptive.c:144: warning 85: in function ThreadCreate unreferenced function argument : 'fp'
sdcc -o test3threads.hex test3threads.rel preemptive.rel
(sdcc) [2022-12-13 16:32:02] jojoelt@ZephyG14:~/OS-edsim/chkpnt4 (master *)
$ |
```

[2] Screenshots and explanation

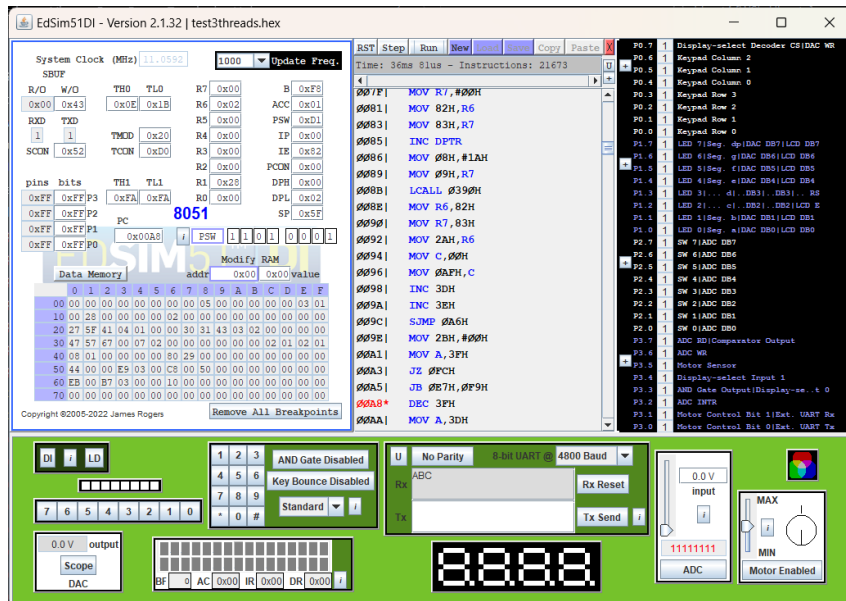
Mutex - 0x3D | Full - 0x3E | Empty - 0x3F | buffer = 0x27-0x29

Producer1 running



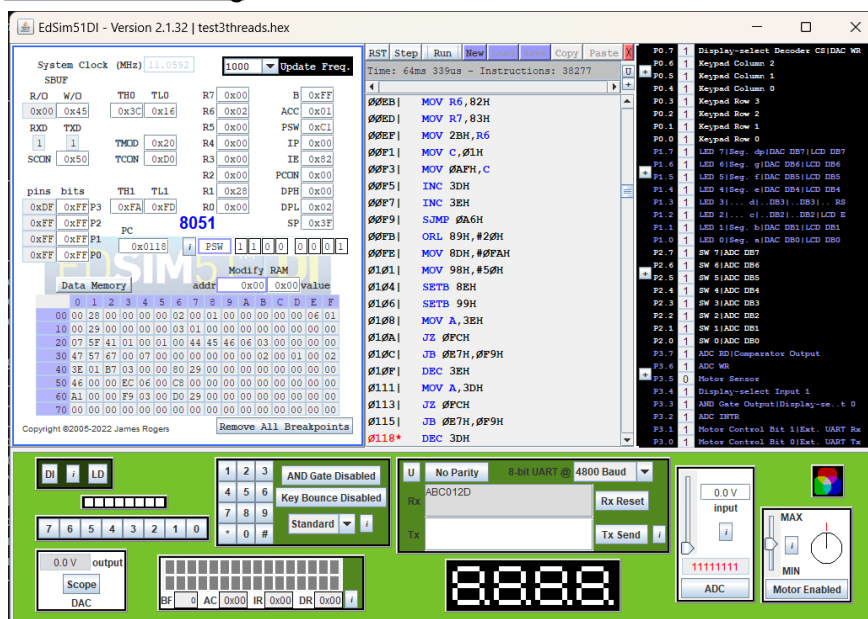
0x27, 0x28 now have values 'A' and 'B', and 0x29 is still empty (' '). 2 Slots occupied, 1 slot empty, at 3rd iteration, mutex not lock yet, Full now equals 2 and Empty now equals 1.

Producer2 running



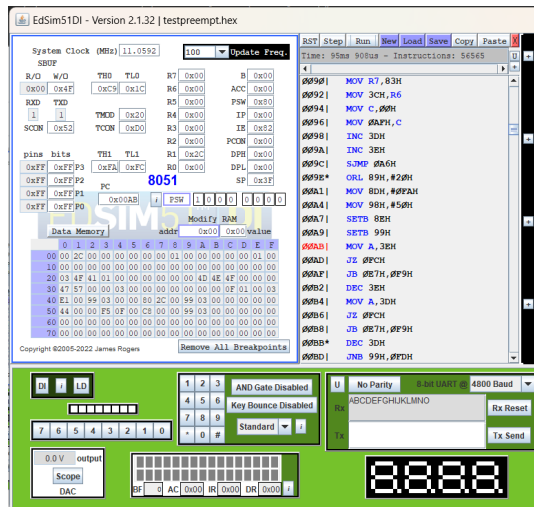
0x27, 0x28 now have values '0' and '1', and 0x29 is still empty (' '). 2 Slots occupied, 1 slot empty, at 3rd iteration, mutex not lock yet, Full now equals 2 and Empty now equals 1.

Consumer running

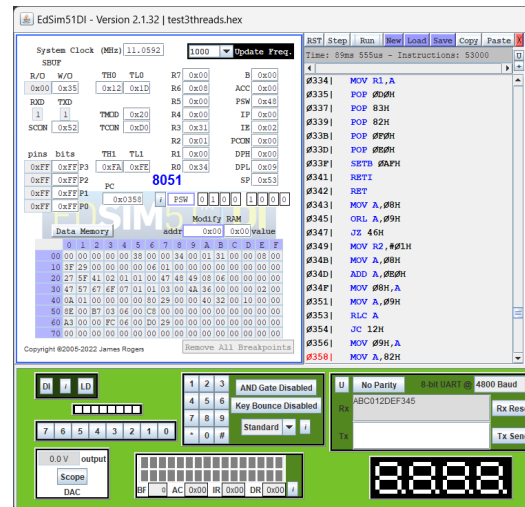


There is a delay when using breakpoint (D) is already out, when semaphore empty reaches 3, E and F will be out, in consumer, each time semaphore full decreases and semaphore empty increases, it mean the buffer has been consumed, so it will output on the screen.

Fairness



UART Output of unfair version



UART Output of fair version

It can be showed that in the unfair version, because of directly using the round-robin scheduling policy, (Letter Producer before Number Producer), the Letter producer will fill in the buffer, then preempts to the Number Producer but it get blocks by the empty Semaphore, By using the fair version (P1 C P2 C) where P stands for producer, and C stands for consumer, it solves the problem of the unfair version