

1.

$\text{map} :: (a \rightarrow b) \rightarrow [a] \rightarrow [b]$

$\text{map } f [] = []$  (map.1)

$\text{map } f (x: xs) = f x: \text{map } f xs$  (map.2)

$(++) :: [a] \rightarrow [a] \rightarrow [a]$

$[] ++ ys = ys$  (++.1)

$(x: xs) ++ ys = x: (xs ++ ys)$  (++.2)

$\text{map } f (xs ++ ys) = (\text{map } f xs) ++ (\text{map } f ys)$

Induction on xs

Base case  $xs = []$

$\text{map } f (xs ++ ys)$

$= \text{map } f ys$  (++.1)

$(\text{map } f xs) ++ (\text{map } f ys)$

$= [] ++ (\text{map } f ys)$  (map.1)

$= \text{map } f ys$  (++.1)

Induction step ( $xs = x: xs'$ ):

$\text{map } f (xs ++ ys)$

$= \text{map } f ((x: xs') ++ ys)$

$= \text{map } f (x: xs' ++ ys)$  (++.2)

$= f x: \text{map } f (xs' ++ ys)$  (map.2)

$(\text{map } f xs) ++ (\text{map } f ys)$

$= (\text{map } f (x: xs')) ++ (\text{map } f ys)$

$= (f x: (\text{map } f xs')) ++ (\text{map } f ys)$  (map.2)

$= f x: ((\text{map } f xs') ++ (\text{map } f ys))$  (++.2)

$= f x: \text{map } f (xs' ++ ys)$  (by IH)

2.

$\text{zip} :: [a] \rightarrow [b] \rightarrow [(a, b)]$

$\text{zip } [] \_ = []$  (zip.1)

$\text{zip } \_ [] = []$  (zip.2)

$\text{zip } (x: xs) (y: ys) = (x, y): \text{zip } xs ys$  (zip.3)

(1)

Induction on ps

Base case ( $ps = []$ ):

$\text{zip } (\text{fst } (\text{unzip } ps)) (\text{snd } (\text{unzip } ps))$

$= \text{zip } (\text{fst } ([], [])) (\text{snd } ([], []))$  (unzip.1)

$= \text{zip} ([ ] ) ([ ])$  (by definition of fst, snd)  
 $= [ ]$  (zip.1)

Induction step  $(ps = (p1, p2): ps')$ :

$\text{zip} (\text{fst} (\text{unzip} ((p1, p2): ps'))) (\text{snd} (\text{unzip} ((p1, p2): ps')))$   
 $= \text{zip} (\text{fst} ((p1: xs, p2: ys) \text{ where } (xs, ys) = \text{unzip } ps')) (\text{snd} (\text{unzip} ((p1, p2): ps'))) \quad (\text{unzip.2})$   
 $= \text{zip} (p1:xs \text{ where } (xs, ys) = \text{unzip } ps') (\text{snd} (\text{unzip} ((p1, p2): ps'))) \quad (\text{by definition of fst})$   
 $= \text{zip} (p1: \text{fst} (\text{unzip } ps')) ((\text{snd} (\text{unzip} ((p1, p2): ps')))) \quad (\text{by definition of fst})$   
 $= \text{zip} (p1: \text{fst} (\text{unzip } ps')) (\text{snd} ((p1: xs, p2: ys) \text{ where } (xs, ys) = \text{unzip } ps')) \quad (\text{unzip.2})$   
 $= \text{zip} (p1: \text{fst} (\text{unzip } ps')) (\text{snd} (p2: ys \text{ where } (xs, ys) = \text{unzip } ps')) \quad (\text{by definition of snd})$   
 $= \text{zip} (p1: \text{fst} (\text{unzip } ps')) (p2: \text{snd} (\text{unzip } ps')) \quad (\text{by definition of snd})$   
 $= (p1, p2): \text{zip} (\text{fst} (\text{unzip } ps')) (\text{snd} (\text{unzip } ps')) \quad (\text{zip.3})$   
 $= (p1, p2): ps' \quad (\text{by IH})$   
 $= ps$

(2)

$xs, ys$  同时为有限列表，并且二者长度相同

$xs, ys$  同时为有限列表，并且二者长度相同

Induction on  $xs, ys$

Base case  $xs = [ ], ys = [ ]$

$\text{unzip} (\text{zip } xs \text{ } ys)$   
 $= \text{unzip} [ ] \quad (\text{zip.1})$   
 $= ([ ], [ ]) \quad (\text{unzip.1})$   
 $= (xs, ys)$

Induction step  $(xs = x: xs', ys = y: ys', \text{length } xs' = \text{length } ys')$ :

$\text{unzip} (\text{zip} (x: xs') (y: ys'))$   
 $= \text{unzip} ((x, y): \text{zip } xs' \text{ } ys') \quad (\text{zip.3})$   
 $= (x: xs'', y: ys'') \text{ where } (xs'', ys'') = \text{unzip} (\text{zip} (xs', ys')) \quad (\text{unzip.2})$   
 $= (x: xs', y: ys') \quad (\text{by IH})$   
 $= (xs, ys)$