

Local Instances and Type Synonyms

We've seen that declaring instances has deep consequences: typically, there must be a unique instance of a certain class on a given type, and once we declare such an instance things are stuck forever.

This might be too rigid. Two ways out are *local instances* and *type synonyms*.

Local Instances

Inside a `section` we can use the `attribute [instance] myStructure` or `attribute [-instance] myStructure` to upgrade `myStructure` to an instance or remove it from the list of declared instances. Finding who is `myStructure` in a specific use-case can be non-trivial (won't be too hard neither).

An equivalent syntax is `local instance` instead of `instance` (but this does not work to *remove* instances).

+++ An example from Mathlib

In Mathlib, `N` is endowed with the discrete uniformity, coming from the discrete metric:

1. The metric, induced from that on `R`, satisfies $\forall n : N, \text{Metric.ball } n \ (1/2 : R) = \{n\}$.
2. The uniformity (i. e. a filter on `N × N`) is the principal filter containing the diagonal: `Uniformity N = P (idRel)` where
 - `idRel` is the identity relation, so the subset $\{p : N \times N \mid p.1 = p.2\}$;
 - `P (idRel)` is the collection of all subsets in `N × N` that contain `idRel`, seen as a filter;
 - It can be proven that the uniformity induced by the discrete metric is indeed `P (idRel)`;
 - Filters and uniformities are ordered, and one can prove that `P (idRel) = ⊥`, the bottom element.

Since the discrete metric induces the discrete topology, `UniformSpace.toTopologicalSpace N = ⊥` where now `⊥` is the discrete topology.

GOAL : Provide another non-discrete uniform structure on `N` that still induces the discrete topology.

Reference : This is actually [a counterexample](#) in Mathlib.

Idea : Set

$$\text{dist } n \ m := |2^{(-n : Z)} - 2^{(-m : Z)}| : R$$

We're identifying `N` with the subset $2^{\{-N\}} \subseteq R$, inheriting the distance from this embedding and looking at the induced topology.

Consequence This new uniformity will be so crazy that the identity sequence $\text{id} : \mathbb{N} \rightarrow \mathbb{N}$ is actually Cauchy (Cauchy sequences in discrete uniform spaces are only the eventually constant ones).

+++ The problems

- How can we "replace" the discrete uniformity on \mathbb{N} with another one?
- How can we check that our results (for instance about id being Cauchy) re-become *false* in the usual setting where $\text{UniformSpace } \mathbb{N} := \langle \perp \rangle$?
- How can we check that the topology remained the same, namely the discrete one?

Solution Use local instances.

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Type Synonyms

Another strategy that works more globally is to use *type synonyms*. The idea is to create a copy of a type, in a way that this copy inherits some instances of the original type, but not all of them.

+++ Difference between `abbrev` and `def`

This is probably beyond what is meant for this course, and certainly beyond my paygrade.

You can think that "`abbrev` is a reducible `def`", whatever this means.

Concretely: Lean "looks deeper" inside the definition of an `abbrev` than a `def`.

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Suppose `X` is a type, and that

```
instance : ClassOne X := ...
```

up to

```
instance : Class_n X := ...`.
```

We want a new type `newX` that has some of the above instances (and to perform this **fast**).

+++ The **wrong** way: `abbrev newX_bad := X`.

For Lean, `newX_bad` and `X` are **equal**: so, every declaration with variable `newX_bad` will accept a variable of type `X`. In particular, an `instance : MyClass newX_bad := ...` will result in an `instance : MyClass X :=`

We are also changing the *old* type `X`. This is **not** what we wanted.

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+++ The **good** way: `def newX_good := X`.

We're creating a completely new type `newX_good`. The problem is that it has no property at all, whereas we might want to inherit some properties from `X` (although probably not all of them).

We can use the syntax

```
instance : myClass newX_good := inferInstanceAs (myClass X)
```

that instructs Lean to *copy* the instance term of `myClass` from `X` to `newX_good`.

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Structures

- Main reference: [The Lean Language Reference](#), in particular § 3.4.2.

The usual way to define a **structure** is to write its name, then **where** (or `:=`) and then the list of fields that we want a term of the structure to be made of

```
structure MyStructure where
  firstfield : firstType
  secondfield : secondType
  ...
  lastfield : lastType
```

or equivalently

```
structure MyStructure :=
  firstfield : firstType
  secondfield : secondType
  ...
  lastfield : lastType
```

where each field is a term in some known type. Every field can depend upon the previous ones.

- The `nth` field of a structure can be *any* term (of the right type...) but if we write

```
nth_field : nth_Type := myterm
```

we are declaring that, if left unspecified, the `nth` term will be `myterm`. This is typically what we do when `nthType = Prop`: we do not want that *some property* is satisfied, but that *our sought-for property* is satisfied.

Declaring a structure as above automatically creates several terms:

1. A term `MyStructure.mk : firstType → secondType → ... → lastType → MyStructure` to *construct* terms; the name `.mk` can be overridden with the syntax `constructor_name ::` on the second line (so starting the list of fields on the third line).

2. A term `MyStructure.nthfield : MyStructure → nthType`: this *projects* a term of type `MyStructure` onto its `nth` field.

3. If the attribute `@[ext]` is prepended on the line before the declaration, a theorem `MyStructure.ext` is created, of type

$$\forall \{x\ y : \text{MyStructure}\}, x.\text{firstfield} = y.\text{firstfield} \rightarrow \dots \rightarrow x.\text{lastfield} = y.\text{lastfield} \rightarrow x = y$$

saying that if all fields of two terms coincide, the terms themselves coincide.

- If `nthType = Prop`, the arrow `x.(n-1)stfield = y.(n-1)stfield → x.nthfield = y.nthfield` *→ is skipped thanks to proof irrelevance*. Another theorem `MyStructure.ext_iff` is also added, that adds the reverse implication.

4. If the `@[class]` attribute is added (possibly with syntax `@[ext, class]`), a new class is created as well so that `instance : MyStructure := someterm` becomes accessible.

The call `whatsnew in` on the line preceding the structure makes Lean show all newly created declarations.

+++ Use of parameters

It is also possible to define structures that depend on parameters. The syntax is the usual as for `def` or `theorem`.

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The call `#print MyStructure` has Lean print all fields, parameters and constructors.

Examples

We will define a structure `OneNat`, that "packs" a single natural number; the structures `TwoNat` and `Couple` that pack to numbers; or the structure of order pairs that pack two numbers where the second is larger or equal than the first, so it is a `Prop`: this is called a *mixin*.


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Constructing terms


To look at the details, let's try to build some terms of the above structures.

When doing so, `VSCode` comes at rescue: once we declare that we are looking for a term in a structure `MyStructure` (i. e. in an inductive type with one constructor, itself a function with several arguments), we can type

```
def MyTerm : MyStructure :=  
—
```

(beware that the underscore `_` **must not be indented**), and a (blue) bulb  appears. Click on it to generate a *skeleton* of the structure at hand, so you do not need to remember all fields by heart.

Either using  or not, there are three ways to define a term of a structure:

1. `myTerm : MyStructure :=`, followed either by
 - `by constructor` and then you're in tactic mode; or
 - `{firstfield := firstterm, secondfield := secondterm, ..., lastfield := lastterm}`.
 2. `myTerm : MyStructure where` and then the list `nthfield := nthterm`, each one a new (indented) line (observe that the  -action replaces `:=` with `where` automatically).
 3. Using the so-called *anonymous constructor* provided by `{` and `}`: just insert the list of terms `(firstterm, secondterm, ..., lastterm)` after `myTerm : MyStructure :=` and Lean will understand.
- Remember that `classes` are a special case of `structures`: so, defining an `instance` as we did in the last lecture really boils down to constructing a term of a certain `structure`. Points
1. – 3. above are crucial for this.

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Now, constructing terms of a structure with many fields is particularly

1. boring;
2. error-prone; and
3. far from mathematical usage: to construct a term of a complicated structure we might want to use a term of a simpler one and "only add what is left to update the simpler one to the richer".

There are two ways, somewhat parallel to the `MyStructure := ...` vs `Mystructure where ...` syntaxes.

- The syntax `with` instructs Lean to take all possible labels from that term and to only ask for the remaining ones: it works when using the `:=` construction. Calling `with` triggers both
 - collecting all useful fields from a term; and
 - discharging all useless ones.

Both can be used independently.

- The syntax `___` has the same behaviour, and works when using the `where` construction.

In both cases, the "extra-fields" are forgotten, and thrown away.

+++ Labels Matter

The big difference between `TwoNat`, and `Couple` are the names of the fields:

```
structure TwoNat where
  fst : ℕ
  snd : ℕ

structure Couple where
  left : ℕ
  right : ℕ
```

These names **are relevant!** You might think of a term of type `TwoNat` (or `Couple`) as a pair of *labelled* naturals, and that a structure is a collection of *labelled* terms. So, the terms `t := {fst := 2, snd := 1} : TwoNat` and the term `t' := {left := 2, right := 1} : Couple` have **nothing to do with each other**.

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+++ More about `with`

Technically, `with` updates a value: so `{fst := 1, snd := 2} with fst := 3` is `{fst := 3, snd := 2}`.

Using `with` without specifying a new value simply instructs Lean to consider all fields on their own without changing them (but possibly picking some of them if needed).

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Exercises

1. Define the class of metric spaces (but call them `SpaceWithMetric` to avoid conflict with the existing library) as defined in https://en.wikipedia.org/wiki/Metric_space#Definition, and deduce an instance of `TopologicalSpace` on every `SpaceWithMetric`.

Explain why this is the *wrong* choice, on an explicit example, and fix it.

The other exercises have been moved to the file `Structures4`.