

# A Note on Unforgeability of MuSig2 with Key Tweaking

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**Abstract.** *Key Tweaking* refers to the process of producing a new pair of secret and public key from a given pair. This is used, for example, to derive fresh keys from a master keypair or to create a commitment to a value such that the commitment is also a public key. In this note, we show that a variant of MuSig2 with naive tweaking is insecure and propose a variant that is not vulnerable against the attack.

## 1 The Vulnerable Scheme

Figure 1 shows MuSig2NaiveTweak, a multi-signature scheme that is identical to MuSig2 [NRS21] except that it has the additional capability of signing for a tweaked public key. There are multiple variants of tweaking, which differ mainly in how the tweak  $t$  is derived. We chose a variant of tweaking for MuSig2NaiveTweak that gives the adversary the minimal power necessary to make the attack work. In particular, the honest signer generates uniformly random tweaks without input from the adversary and outputs possible tweaks. MuSig2NaiveTweak uses *additive* tweaking, but the attack similarly applies to *multiplicative* tweaking.

## 2 Generalized Birthday Problem

The attack against MuSig2NaiveTweak makes use of Wagner’s algorithm [Wag02] for solving the Generalized Birthday Problem. It can be defined as follows for the purpose of this paper: Given a constant value  $t \in \mathbb{Z}_p$ , an integer  $k_{\max}$ , and access to random oracle  $H$  mapping onto  $\mathbb{Z}_p$ , find a set  $\{q_1, \dots, q_{k_{\max}}\}$  of  $k_{\max}$  queries such that  $\sum_{k=1}^{k_{\max}} H(q_k) = t$ . For  $k_{\max} = 2^{\sqrt{\log_2(p)}-1}$  the complexity of this algorithm is  $O(2^{2\sqrt{\log_2(p)}})$ .

## 3 Description of the Attack against MuSig2NaiveTweak

The adversary calls KeyTweak  $\ell_{\max} \in O(2^{2\sqrt{\log_2(p)}})$  times to obtain values  $t^{(1)}, \dots, t^{(\ell_{\max})}$  and computes the multiset of public keys  $L$  and aggregate key  $\tilde{X}$  for the (untweaked) public key of the honest signer  $X'_1 = g^{x_1}$  as

$$L = \{X'_1 g^{t^{(1)}}, \dots, X'_1 g^{t^{(\ell_{\max})}}\}$$
$$\tilde{X} = \text{KeyAgg}(L).$$

Then, the adversary opens  $k_{\max} = 2^{\sqrt{\log_2(p)}-1}$  concurrent signing sessions by requesting  $k_{\max}$  nonce tuples  $R_{1,1}^{(1)}, \dots, R_{1,\nu}^{(1)}, \dots, R_{1,1}^{(k_{\max})}, \dots, R_{1,\nu}^{(k_{\max})}$  from the honest signer and computes

$$R_j = \sum_{k=1}^{k_{\max}} R_{1,j}^{(k)}, \quad j \in [1, \nu]$$
$$b = H_{\text{non}}(\tilde{X}, (R_1, \dots, R_\nu), m)$$
$$R^* = \prod_{j=1}^{\nu} R_j^{b^{j-1}}.$$

Now it is possible to use Wagner’s algorithm to find a function  $f : [1, k_{\max}] \rightarrow [1, \ell_{\max}]$  that associates a value  $t^{(\ell)}$  to each session  $k$  such that

<b>Setup</b> ( $1^\lambda$ ) <hr/> $(\mathbb{G}, p, g) \leftarrow \text{GrGen}(1^\lambda)$ Select three hash functions $H_{\text{agg}}, H_{\text{non}}, H_{\text{sig}} : \{0, 1\}^* \rightarrow \mathbb{Z}_p$ $\text{par} := ((\mathbb{G}, p, g), H_{\text{agg}}, H_{\text{non}}, H_{\text{sig}})$ <b>return</b> $\text{par}$	<b>SignAgg</b> ( $\text{out}_1, \dots, \text{out}_n$ ) <hr/> <b>for</b> $i := 1 \dots n$ <b>do</b> $(R_{i,1}, \dots, R_{i,\nu}) := \text{out}_i$ <b>for</b> $j := 1 \dots \nu$ <b>do</b> $R_j := \prod_{i=1}^n R_{i,j}$ <b>return</b> $\text{out} := (R_1, \dots, R_\nu)$
<b>KeyGen</b> () <hr/> $x \leftarrow \$ \mathbb{Z}_p$ ; $X := g^x$ $sk := x$ ; $pk := X$ <b>return</b> $(sk, pk)$	<b>Sign'</b> ( $\text{state}_1, \text{out}, sk_1, m, (pk_2, \dots, pk_n), \mathbf{t}$ ) <hr/> // Sign' must be called at most once per $\text{state}_1$ . <b>if</b> $\mathbf{t} \neq \mathbf{0}$ and has not been output by <b>KeyTweak</b> <b>then return false</b> $(r_{1,1}, \dots, r_{1,\nu}) := \text{state}_1$ $x_1 := sk_1 + \mathbf{t} \bmod p$ ; $X_1 := g^{x_1}$ $(X_2, \dots, X_n) := (pk_2, \dots, pk_n)$ $L := \{X_1, \dots, X_n\}$ $a_1 := \text{KeyAggCoef}(L, X_1)$ $\tilde{X} := \text{KeyAgg}(L)$ $(R_1, \dots, R_\nu) := \text{out}$ $b := H_{\text{non}}(\tilde{X}, (R_1, \dots, R_\nu), m)$ $R := \prod_{j=1}^\nu R_j^{b^{j-1}}$ $c := H_{\text{sig}}(\tilde{X}, R, m)$ $s_1 := ca_1 x_1 + \sum_{i=1}^\nu r_{1,i} b^{i-1} \bmod p$ $\text{state}'_1 := R$ ; $\text{out}'_1 := s_1$ <b>return</b> $(\text{state}'_1, \text{out}'_1)$
<b>KeyTweak</b> () <hr/> $\mathbf{t} \leftarrow \$ \mathbb{Z}_p$ <b>return</b> $\mathbf{t}$	
<b>KeyAggCoef</b> ( $L, X_i$ ) <hr/> <b>return</b> $H_{\text{agg}}(L, X_i)$	
<b>KeyAgg</b> ( $L$ ) <hr/> $\{X_1, \dots, X_n\} := L$ <b>for</b> $i := 1 \dots n$ <b>do</b> $a_i := \text{KeyAggCoef}(L, X_i)$ <b>return</b> $\tilde{X} := \prod_{i=1}^n X_i^{a_i}$	
<b>Ver</b> ( $\tilde{pk}, m, \sigma$ ) <hr/> $\tilde{X} := \tilde{pk}$ ; $(R, s) := \sigma$ $c := H_{\text{sig}}(\tilde{X}, R, m)$ <b>return</b> $(g^s = R\tilde{X}^c)$	<b>SignAgg'</b> ( $\text{out}'_1, \dots, \text{out}'_n$ ) <hr/> $(s_1, \dots, s_n) := (\text{out}'_1, \dots, \text{out}'_n)$ $s := \sum_{i=1}^n s_i \bmod p$ <b>return</b> $\text{out}' := s$
<b>Sign</b> () <hr/> // Local signer has index 1. <b>for</b> $j := 1 \dots \nu$ <b>do</b> $r_{1,j} \leftarrow \$ \mathbb{Z}_p$ ; $R_{1,j} := g^{r_{1,j}}$ $\text{out}_1 := (R_{1,1}, \dots, R_{1,\nu})$ $\text{state}_1 := (r_{1,1}, \dots, r_{1,\nu})$ <b>return</b> $(\text{out}_1, \text{state}_1)$	<b>Sign''</b> ( $\text{state}'_1, \text{out}'$ ) <hr/> $R := \text{state}'_1$ ; $s := \text{out}'$ <b>return</b> $\sigma := (R, s)$

**Fig. 1.** The multi-signature scheme  $\text{MuSig2NaiveTweak}[\text{GrGen}, \nu]$ . The differences to  $\text{MuSig2}[\text{GrGen}, \nu]$  are displayed in **red**.

$$\sum_{k=1}^{k_{\max}} \underbrace{\text{H}_{\text{agg}}(L, X'_1 g^{t^{(f(k))}})}_{=: a_1^{(k)}} \underbrace{\text{H}_{\text{sig}}(\tilde{X}, R^*, m)}_{=: c^{(k)}} = \underbrace{\text{H}_{\text{sig}}(X'_1, R^*, m^*)}_{=: c^*}. \quad (1)$$

for a forgery target message  $m^*$ . For all  $k \in [1, k_{\max}]$  the honest signer is asked for a partial signature using value  $C^{f(k)}$  which is answered with  $s_1^{(k)} = r_{1,1}^{(k)} + br_{1,2}^{(k)} + c^{(k)} \cdot a_1^{(k)}(x_1 + t^{(f(k))})$ . This allows the adversary to compute

$$s_1^{*'} = \sum_{k=1}^{k_{\max}} s_1^{(k)} \quad (2)$$

$$= \sum_{k=1}^{k_{\max}} r_{1,1}^{(k)} br_{1,2}^{(k)} + \left( \sum_{k=1}^{k_{\max}} c^{(k)} a_1^{(k)} \right) \cdot x_1 + \sum_{k=1}^{k_{\max}} c^{(k)} a_1^{(k)} t^{(f(k))} \quad (3)$$

$$= \log_g(R^*) + c^* x_1 + \sum_{k=1}^{k_{\max}} c^{(k)} a_1^{(k)} t^{(f(k))} \quad (4)$$

where the last equality follows from Equation (1). The last summand can be subtracted as

$$s_1^* = s_1^{*'} - \sum_{k=1}^{k_{\max}} c^{(k)} a_1^{(k)} t^{(f(k))}$$

to obtain  $(R^*, s^*)$ , a valid forgery on message  $m^*$  for public key  $X'_1$ .

## 4 BLLOR attack

Benhamouda, Lepoint, Loss, Orrù, and Raykova [BLL+21] describe an algorithm that solves the ROS problem and can be applied to attack to break unforgeability of **MuSig2NaiveTweak**. If the adversary can open at least  $\log_2 p$  sessions, then the algorithm has complexity  $O(\log_2^2 p)$  and a negligible running time in practice (otherwise a variant of the algorithm can be applied that has a higher complexity). In contrast to the attack based on Wagner's algorithm, this attack allows using multisets of public keys that only have two elements.

## 5 Where the security proof of **MuSig2** fails against **MuSig2NaiveTweak**

**Jonas' note:** TODO:  $a \neq a$  but  $b = b$

## 6 A Fixed **MuSig2** Variant with Tweaking

The attack in [section 3](#) relies on selecting a tweak  $t^{(k)}$  for a signing session  $k$  *after* obtaining nonces  $R_{1,1}^{(k)}, \dots, R_{1,\nu}^{(k)}$ . If  $t^{(k)}$  was determined before the nonces of the session, then the function  $f$  in Equation (1) is fixed and, in particular, can not be influenced by the attacker. This is accomplished by the scheme shown in [Figure 1](#), which associates a fixed tweak when generating the nonces and is, therefore, not vulnerable to the attack.

Instead of generating the tweak in **Sign** as shown in [Figure 1](#), one can also prevent the attack by modifying **Sign** to take the tweaked secret key  $sk_1 + t^{(k)} \bmod p$  or tweaked public key  $X_1 g^{t^{(k)}}$  as input. This modified scheme would then write the tweaked secret or public key into  $state_1$  and refuse to sign for anything but the tweaked secret or public key with this  $state_1$ .

It is also possible to stop the attack without having to determine tweak  $t^{(k)}$  before outputting the session's nonces. This can be achieved by changing the scheme so that each signer  $i$  gets a different nonce coefficient  $b_i$  instead of using a single nonce coefficient  $b$  for all signers of a signing session. However, using a separate nonce coefficient  $b_i$  increases the communication complexity

of the scheme because it precludes nonce aggregation via **SignAgg**. All other signers' nonces  $R_{2,1}, \dots, R_{2,\nu}, \dots, R_{n,1}, \dots, R_{n,\nu}$  are required to be input to **Sign** and  $b_i$ ,  $R$  and  $s_1$  are computed as

$$\begin{aligned} b_i &:= H_{\text{non}}(\tilde{X}, (\mathbf{R}_{1,1}, \dots, \mathbf{R}_{1,\nu}, \dots, \mathbf{R}_{n,1}, \dots, \mathbf{R}_{n,\nu}), m, \mathbf{X}_i) \quad i \in [1, n] \\ R &:= \prod_{i=1}^n \prod_{j=1}^{\nu} R_{i,j}^{b_i^{j-1}} \\ s_1 &:= ca_1 x_1 + \sum_{i=1}^{\nu} r_{1,i} b_1^{i-1} \bmod p. \end{aligned}$$

## 7 Conclusion

Other multisignature schemes that use a key agg coefficient are similarly vulnerable. Other multisigs and fix?

- MuSig1 with naive tweaking: if tweak comes in
- multisig with PoK and naive tweaking: there's no “naive” tweaking with PoK

## Acknowledgments

We thank Yannick Seurin for identifying a vulnerability in a related scheme that ultimately led to the discovery of the attack discussed in this note.

## References

- [BLL+21] F. Benhamouda, T. Lepoint, J. Loss, M. Orrù, and M. Raykova. “On the (in)security of ROS”. In: *EUROCRYPT 2021*. 2021.
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<b>KeyTweak</b> () <hr/> $\mathbf{t} \leftarrow \$\mathbb{Z}_p$ <b>return</b> $\mathbf{t}$	<b>SignAgg'</b> ( $\text{out}'_1, \dots, \text{out}'_n$ ) <hr/> $(s_1, \dots, s_n) := (\text{out}'_1, \dots, \text{out}'_n)$ $s := \sum_{i=1}^n s_i \bmod p$ <b>return</b> $\text{out}' := s$
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**Fig. 2.** The multi-signature scheme  $\text{MuSig2Tweak}[\text{GrGen}, \nu]$ . The differences to  $\text{MuSig2}[\text{GrGen}, \nu]$  are displayed in **red**.