



MECH 6v29.002 – Model Predictive Control

L5 – LQ MPC

- Linear Quadratic MPC
- Three potential approaches
 - Formulations and comparison
- Infinite-horizon case
- HW #1

- Linear refers to **linear model**

$$x_{k+1} = Ax_k + Bu_k \quad x_0 = x(0) \quad x \in \mathbb{R}^n, u \in \mathbb{R}^m$$

- Quadratic refers to **quadratic cost function**

$$J_0(x_0, U_0) = \sum_{k=0}^{N-1} x_k^T Q x_k + u_k^T R u_k + x_N^T P x_N \quad U_0 = \begin{bmatrix} u_0 \\ u_1 \\ \vdots \\ u_{N-1} \end{bmatrix} \in \mathbb{R}^{m \cdot N}$$

- Assume $Q = Q^T \geq 0$ $P = P^T \geq 0$ $R = R^T > 0$

Keep an eye out for when this is needed

- Finite-time optimal control problem**

$$J_0^*(x_0) = \min_{U_0} \sum_{k=0}^{N-1} x_k^T Q x_k + u_k^T R u_k + x_N^T P x_N$$

s.t.

$$x_{k+1} = Ax_k + Bu_k, \quad k \in \{0, 1, \dots, N-1\}$$

$$x_0 = x(0)$$

- Three potential solution approaches
 - **Batch approach**
 - Based on the lifted form we derived last lecture
 - Produce an optimal input trajectory based on initial condition
 - **Recursive approach** (dynamic programming)
 - Produce an optimal feedback control policy
 - **Online optimization**
 - Not necessary in this most basic formulation
 - But the approach would not change when you add additional elements to the control formulation (e.g. constraints)

Batch Approach

- Last lecture we used $x_{k+1} = Ax_k + Bu_k$ to derive

$$X = S_x x_0 + S_u U_0 \quad \text{Previous notation} \quad X = Px_0 + HU$$

$$\begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} A \\ A^2 \\ \vdots \\ A^N \end{bmatrix} x_0 + \begin{bmatrix} B & 0 & \cdots & 0 \\ AB & B & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ A^{N-1}B & A^{N-2}B & \cdots & B \end{bmatrix} \begin{bmatrix} u_0 \\ u_1 \\ \vdots \\ u_{N-1} \end{bmatrix}$$

$\uparrow S_x$
 $\uparrow S_u$

- Using this notation, we can rewrite the cost function

$$J_0(x_0, U_0) = \sum_{k=0}^{N-1} x_k^T Q x_k + u_k^T R u_k + x_N^T P x_N = X^T \bar{Q} X + U_0^T \bar{R} U_0$$

$$\bar{Q} = \text{blkdiag}\{Q, \dots, Q, P\}$$

$$\bar{Q} \geq 0$$

$$\bar{R} = \text{blkdiag}\{R, \dots, R\}$$

$$\bar{R} > 0$$

N times

Batch Approach (cont.)

- Substitute lifted state trajectory into cost function

$$J_0(x_0, U_0) = X^T \bar{Q} X + U_0^T \bar{R} U_0$$

$$X = S_x x_0 + S_u U_0$$

$$J_0(x_0, U_0) = (S_x x_0 + S_u U_0)^T \bar{Q} (S_x x_0 + S_u U_0) + U_0^T \bar{R} U_0$$

transpose

$$= \left(x_0^T S_x^T + U_0^T S_u^T \right) \bar{Q} (S_x x_0 + S_u U_0) + U_0^T \bar{R} U_0$$

multiply and rearrange

$$= U_0^T \underbrace{\left(S_u^T \bar{Q} S_u + \bar{R} \right)}_H U_0 + 2x_0^T \underbrace{\left(S_x^T \bar{Q} S_u \right)}_F U_0 + x_0^T \underbrace{\left(S_x^T \bar{Q} S_x \right)}_Y x_0$$

$$= U_0^T H U_0 + 2x_0^T F U_0 + x_0^T Y x_0$$

- What can we say about H ?

$$\bar{Q} \geq 0, \quad \bar{R} > 0 \Rightarrow H > 0 \Rightarrow J_0(x_0, U_0) \text{ is pos. def. function of } U_0$$

Batch Approach (cont.)

$\Rightarrow J_0(x_0, U_0)$ is pos. def. function of U_0

- This means we can find the minimum of the cost function with respect to the input trajectory by **taking the derivative and setting it equal to zero**

$$J_0(x_0, U_0) = U_0^T H U_0 + 2x_0^T F U_0 + x_0^T Y x_0$$

$$\frac{dJ_0}{dU_0} = 2H U_0 + 2F^T x_0 = 0$$

- **Optimal input:** $U_0^* = -H^{-1} F^T x_0$
$$= -\left(S_u^T \bar{Q} S_u + \bar{R}\right)^{-1} \left(S_x^T \bar{Q} S_u\right)^T x_0$$
$$= -\left(S_u^T \bar{Q} S_u + \bar{R}\right)^{-1} S_u^T \bar{Q} S_x x_0$$
- **Optimal cost:** $J_0^*(x_0) = x_0^T \left(Y - F H^{-1} F^T\right) x_0$

Batch Approach (cont.)

- Batch approach produces an **optimal input trajectory**

$$\begin{aligned} U_0^* &= -H^{-1}F^T x_0 \\ &= -\left(S_u^T \bar{Q} S_u + \bar{R}\right)^{-1} S_u^T \bar{Q} S_x x_0 \end{aligned}$$

$$U_0^* = \begin{bmatrix} u_0^* \\ u_1^* \\ \vdots \\ u_{N-1}^* \end{bmatrix} \in \mathbb{R}^{m \cdot N}$$

- If we want just the **first input**, we can always multiply by a matrix

$$\begin{aligned} u_0^* &= [I \quad 0 \quad \cdots \quad 0] U_0^* \\ &= -[I \quad 0 \quad \cdots \quad 0] H^{-1} F^T x_0 \end{aligned}$$

- **Dynamic programming**

- In general, take a complicated problem and break it into simpler problems to be solved in a recursive manner

$$J_0(x_0, U_0) = \sum_{k=0}^{N-1} x_k^T Q x_k + u_k^T R u_k + x_N^T P x_N$$

Subject to
model
constraints

$$J_0(x_0, U_0) = x_0^T Q x_0 + u_0^T R u_0 + J_1^*(x_1)$$

$$J_1^*(x_1) = \min_{U_1} \sum_{k=1}^{N-1} x_k^T Q x_k + u_k^T R u_k + x_N^T P x_N$$

- Key idea: work your way **backwards**
 - Find the optimal solution at the last time step
 - Then find the optimal solution at the second-to-last time step based on the solution at the last time step
 - Repeat recursively until you reach time 0

Recursive Approach (cont.)

- Start at the very last time step

$$J_{N-1}^*(x_{N-1}) = \min_{u_{N-1}} x_{N-1}^T Q x_{N-1} + u_{N-1}^T R u_{N-1} + x_N^T \overset{P_N = P}{P_N} x_N$$

s.t.

$$x_N = A x_{N-1} + B u_{N-1}$$

- Plug in state constraint

$$\begin{aligned} J_{N-1}^*(x_{N-1}) &= \min_{u_{N-1}} x_{N-1}^T Q x_{N-1} + u_{N-1}^T R u_{N-1} + (A x_{N-1} + B u_{N-1})^T P_N (A x_{N-1} + B u_{N-1}) \\ &= \min_{u_{N-1}} x_{N-1}^T (A^T P_N A + Q) x_{N-1} + 2 x_{N-1}^T (A^T P_N B) u_{N-1} + u_{N-1}^T (B^T P_N B + R) u_{N-1} \end{aligned}$$

- Quadratic function of u_{N-1}

- Find optimal input $u_{N-1}^* = - \underbrace{(B^T P_N B + R)^{-1} B^T P_N A}_{F_{N-1}} x_{N-1}$

$$u_{N-1}^* = F_{N-1} x_{N-1}$$

Recursive Approach (cont.)

- Optimal one-step cost function

$$J_{N-1}^*(x_{N-1}) = x_{N-1}^T Q x_{N-1} + u_{N-1}^{*T} R u_{N-1}^* + x_N^T P_N x_N$$

$$x_N = A x_{N-1} + B u_{N-1}^* \quad u_{N-1}^* = F_{N-1} x_{N-1}$$



$$J_{N-1}^*(x_{N-1}) = x_{N-1}^T P_{N-1} x_{N-1}$$

$$P_{N-1} = A^T P_N A + Q - A^T P_N B (B^T P_N B + R)^{-1} B^T P_N A$$

Recursive Approach (cont.)

- Now we can do this for any time step k

$$u_{N-1}^* = -\left(B^T P_N B + R\right)^{-1} B^T P_N A x_{N-1} \quad \rightarrow \quad u_k^* = -\left(B^T P_{k+1} B + R\right)^{-1} B^T P_{k+1} A x_k \\ = F_k x_k, \quad k \in \{0, \dots, N-1\}$$

$$J_k^*(x_k) = x_k^T P_k x_k$$

$$P_k = A^T P_{k+1} A + Q - A^T P_{k+1} B \left(B^T P_{k+1} B + R\right)^{-1} B^T P_{k+1} A$$

- This is the **Discrete-time Riccati Equation**
 - Initialized at $P_N = P$
 - Solved backwards $P_{k+1} \rightarrow P_k$
- **Optimal feedback control policy** $u_k^* = F_k x_k$

- Or, if you don't care about the specific relationships between the current state and the optimal inputs, then you can just formulate and solve the optimization problem

$$\begin{aligned} J_0^*(x_0) = \min_{U_0} & \sum_{k=0}^{N-1} x_k^T Q x_k + u_k^T R u_k + x_N^T P x_N \\ \text{s.t.} & \\ & x_{k+1} = A x_k + B u_k, \quad k \in \{0, 1, \dots, N-1\} \\ & x_0 = x(0) \end{aligned}$$

- This is a **quadratic program** – relatively easy to solve
- Main advantage:
 - Approach is unchanged if now we have

$$\begin{aligned} J_0^*(x_0) = \min_{U_0} & \sum_{k=0}^{N-1} x_k^T Q x_k + u_k^T R u_k + x_N^T P x_N \\ \text{s.t.} & \\ & x_{k+1} = A x_k + B u_k, \quad k \in \{0, 1, \dots, N-1\} \\ & \mathbf{x}_{k+1} \in \mathcal{X}, \quad \mathbf{u}_k \in \mathcal{U}, \quad k \in \{0, 1, \dots, N-1\} \\ & x_0 = x(0) \end{aligned}$$

- Other approaches become significantly more difficult

Comparison of Approaches

- Remember that we are currently solving a finite-time optimization problem (not thinking about receding horizon MPC yet)

$$J_0^*(x_0) = \min_{U_0} \sum_{k=0}^{N-1} x_k^T Q x_k + u_k^T R u_k + x_N^T P x_N$$

s.t.

$$x_{k+1} = A x_k + B u_k, \quad k \in \{0, 1, \dots, N-1\}$$
$$x_0 = x(0)$$

- If model is perfect (no model error, no disturbances), then the three approaches will result in the same state and input trajectories
- But the batch approach is an open-loop optimization, while the recursive, dynamic programming approach is an optimal control policy based on the current state
- If there is model uncertainty (or disturbances), we would expect the recursive approach to be more robust

Infinite Horizon Problem

- What if our system operates forever? $N \rightarrow \infty$

$$J_0^*(x_0) = \min_{U_0} \sum_{k=0}^{\infty} x_k^T Q x_k + u_k^T R u_k \quad \text{No more terminal cost}$$

s.t.
 $x_{k+1} = Ax_k + Bu_k, \quad k \in \{0, 1, \dots, \infty\}$
 $x_0 = x(0)$

- Batch approach is no longer applicable
- Recursive approach can still work
- Using the Discrete-time Riccati Equation initialized as $P_0 = Q$

$$P_k = A^T P_{k+1} A + Q - A^T P_{k+1} B (B^T P_{k+1} B + R)^{-1} B^T P_{k+1} A \quad k \rightarrow -\infty$$

- If we assume that these iterations converge then $P_k \rightarrow P_\infty$

$$P_\infty = A^T P_\infty A + Q - A^T P_\infty B (B^T P_\infty B + R)^{-1} B^T P_\infty A$$

- **Discrete-time Algebraic Riccati Equation (DARE)**

Infinite Horizon Problem (cont.)

$$P_{\infty} = A^T P_{\infty} A + Q - A^T P_{\infty} B \left(B^T P_{\infty} B + R \right)^{-1} B^T P_{\infty} A$$

- **Optimal control input** $u_k^* = - \left(B^T P_{\infty} B + R \right)^{-1} B^T P_{\infty} A x_k$
 $= F_{\infty} x_k, \quad k \in \{0, \dots, \infty\}$
- **Optimal infinite horizon cost** $J_0^*(x_0) = x_0^T P_{\infty} x_0$
- This solution corresponds exactly to the discrete-time, infinite-horizon, **LQR solution**
 - In Matlab: `[K,S,e] = dlqr(A,B,Q,R,N)`
 - <https://www.mathworks.com/help/control/ref/dlqr.html>

$$u[n] = -Kx[n]$$

minimizes the quadratic cost function

$$J(u) = \sum_{n=1}^{\infty} (x[n]^T Q x[n] + u[n]^T R u[n] + 2x[n]^T N u[n])$$

for the discrete-time state-space mode

$$x[n+1] = Ax[n] + Bu[n]$$

- We assumed that the discrete-time Riccati equation integrations converged
- But, we can prove this
- **Theorem:** If (A, B) is a stabilizable pair and $(Q^{1/2}, A)$ is an observable pair, then the discrete-time Riccati equation with $P_0 \geq 0$ converges to the unique pos. def. solution P_∞ of the DARE and all the eigenvalues of $(A + BF_\infty)$ lie strictly inside the unit circle.
- **Stabilizable**
 - So that cost is finite – system can be driven to the origin
- **Observable**
 - Think about rewriting the cost function as $x_k^T Q x_k = (x_k^T Q^{1/2}) (Q^{1/2} x_k)$
 - Now the “output” is penalized $y_k = Q^{1/2} x_k$
 - Driving outputs to zero means that states converge to zero
- **Closed-loop system** $x_{k+1} = Ax_k + Bu_k \quad u_k = F_\infty x_k$
$$x_{k+1} = (A + BF_\infty) x_k$$

One Potential Approach

- What if we want to be able to **prove that MPC is closed-loop stable**?
 - In a previous lecture, we saw that adding the constraint $x_N = 0$ can help (and we can prove this)
 - But we can also add a **terminal costs** based on the DARE (discrete-time, infinite-horizon, LQR solution)

$$J_0^*(x_0) = \min_{U_0} \sum_{k=0}^{N-1} x_k^T Q x_k + u_k^T R u_k + x_N^T P_{\infty} x_N$$

s.t.

$$x_{k+1} = A x_k + B u_k, \quad k \in \{0, 1, \dots, N-1\}$$
$$x_0 = x(0)$$

- We will be able to use this to prove stability as well and avoids some of the issues associated with the terminal constraint $x_N = 0$

Homework #1 (Due: Sept. 15)



- Two problems
- Problem 1:
 - Implement and compare the batch and recursive approaches in Matlab.
 - Determine the stabilizing effects of a long prediction horizon.
- Problem 2:
 - Implement the online optimization approach in Matlab using the YALMIP toolbox
 - Become familiar with YALMIP example before attempting (let me know if you have any questions)
- In a single PDF, type your responses to the various questions, provide well formatted Matlab plots, and your Matlab code
 - All of this helps me provide you with more feedback