Submit a short report containing the following. (40 points) Provide the table comparing the two versions by listing the running time for each data size.

| Input size | Version 1 (last element) | Version 2 (med of 3 random) |
|------------|--------------------------|-----------------------------|
| N = 10 | 0.0251 | 0.0313 |
| N = 100 | 1.5573 | 0.3088 |
| N = 200 | 5.4788 | 0.9123 |
| N = 500 | 35.6595 | 1.8977 |
| N = 1000 | 137.2155 | 3.8428 |
| N = 1500 | 315.1822 | 5.812 |
| N = 2000 | 545.6853 | 8.177 |

(Note: Time is in ms)

(30 point) Draw conclusion on why choice of pivot matters.

Quicksort has a worst case scenario of $O(n^2)$. This only happens when the pivot picked with each recursive call is the worst pivot possible. For our regular quicksort algorithm, a sorted list is the worst case scenario. When we implement quicksort where we pick each pivot based on 3 random elements, we are more likely to get a better pivot with each recursive call. Randomly picking indices from the list reduces the chance of getting a worst-case scenario to practically 0, especially for larger lists. It also guarantees that there is no predetermined list you can input to get a worst case scenario. Picking 3 random elements and using the median as the pivot guarantees that the runtime will be around $O(n\log n)$ no matter the input.

(30 points) Code Attach the source code of your implementation. If it is short enough, you may simply include your code as part of your report