

```

1)

public int longestCommonSubsequence(String firstText, String secondText){ 1 usage • jonathanp2004

    int a = firstText.length(); O(1)
    int b = secondText.length(); O(1)

    int [][] table = new int[a + 1][b + 1]; O(1)
        ↪ depends on the a = firstText
    for(int i = 1; i <= a; i++){ O(n) ↪ depends on the a = firstText
        for(int j = 1; j <= b; j++) ↪ now b times × a
        {
            if(firstText.charAt(i-1) == secondText.charAt(j-1)) O(1)
            {
                table [i][j] = table[i-1][j-1]+1; O(1)
            } else{
                table [i][j] = Math.max(table[i-1][j], table[i][j-1]); O(1)
            }
        }
    }

    return table[a][b]; O(1)
}

```

The time complexity of this algorithm would be $O(n^2)$
with n being the length of the input array.
There both have a time complexity of $O(n)$ and
since the inner loop is iterated b times $\times a$ times you
multiply both getting you $O(n^2)$

The space complexity is $O(n^2)$ because the int j
is initialized a timer, so the amount of initialization
depends on the $n = a.length$

```

public String longestCommonSubString(String fText, String sText){ 1 usage ▾ jonathanp2004

    int a = fText.length(); O(1)
    int b = sText.length(); O(1)

    int[][] ta = new int[a+1][b+1]; O(1)

    int maxLength = 0; O(1)
    int endIndex= 0; O(1)

    for(int i = 1; i <= a; i++){ O(n)
        for(int j = 1; j <= b; j++){ O(n)
            if(fText.charAt(i - 1) == sText.charAt(j - 1)) {
                ta[i][j] = ta[i-1][j-1]+1;
                if(ta[i][j] > maxLength) {
                    maxLength = ta[i][j];
                    endIndex = i;
                } else {
                    ta[i][j] = 0;
                }
            }
        }
        if(maxLength == 0) { O(1)
            return "";
        }
    }
    return fText.substring(endIndex - maxLength, endIndex); O(1)
}

```

$O(n^2)$

$O(n)$

The time complexity of this algorithm is $O(n^2)$

Because the inner loop is iterated through $b \times n$ again similar to the subsequence algorithm.

The space complexity is $O(n^2)$ because the int j is initialized a times.

3)

```

Scanner scan = new Scanner(System.in);

System.out.print("Sequence Amount: ");
int n = scan.nextInt();

if(n<= 0)
{
    System.out.println("Enter Positive");
    return;
}

long[] sq = new long[n];
//Initial value of 0,1
sq[0] = 0;
if(n>1)
{
    sq[1] = 2;
}

for(int i = 2; i < n; i++)  $\theta(n-2) = \Theta(n)$ 
{
    //same calc as show in the assignment
    double calc = (1.5* sq[i-1]) + (2* sq[i - 2]);  $\Theta(1)$ 

    sq[i] = (long) Math.floor(calc);  $\Theta(1)$ 
}

System.out.println("Input: " + n);

System.out.print("Output: ");
for(int i = 0; i < n; i++)
{
    System.out.print(sq[i]);
    if(i < n - 1)
    {
        System.out.print(", ");
    }
}

System.out.println();

```

The time complexity is $\Theta(n)$ where the n is the length of the sequence, because the loop runs n times.

The space complexity is constant because the total memory created stays the same - regardless of the n .

```

Scanner scan = new Scanner(System.in);

System.out.print("Enter Number: ");

long x = scan.nextLong();

if (x < 0){
    System.out.println("Output: -1");
    return;
}

if(x == 0){
    System.out.println("Output: 0");
}

if(x == 2){
    System.out.println("Output: 1");
}

```

```

long nm1 = 2;
long nm2 = 0;
int index = 1;

while(true){ ← goes on depending on x

    double calc = (1.5 * nm1) + (2*nm2); } O(1) -
    long c = (long)Math.floor(calc);
    index++;

    if(c == x){
        System.out.println("Input: " + x);
        System.out.println("Output: " + index); } O(1)
        return;
    }

    if(c > x){
        System.out.println("Input: " + x);
        System.out.println("Output: " + (index - 1)); } O(1)
        return;
    }

    nm2 = nm1; O(1)
    nm1 = c; O(1)
}

```

4

The time complexity for this algorithm is $O(n)$ where n depends on the value user inputs

The space complexity is $\Omega(1)$ because again the created variables don't change.

(5)

```

public class RemoveElements { ↳ jonathanp2004

    public static int removeElements(double[] nums, double val){ 1 usage ↳ jonathanp2004
        int k = 0;

        for(int i = 0; i < nums.length; i++){ O(n)
            if(nums[i] >= val){ ↳ O(1)
                nums[k] = nums[i]; O(1)
                k++;           O(1)
            }
        }
        return k;      O(n
    }

    public static void main(String[] args){ ↳ jonathanp2004

        double[] nums = {3.5, 1.2, 2.1, 6.7, 9.0, 4.4};
        double val = 3.0;

        System.out.print("Input: nums = {");
        for(int i = 0; i < nums.length; i++){
            System.out.print(nums[i]);
            System.out.print(",");
        }
        System.out.println("}, val = " + val);

        int s = removeElements(nums, val);

        System.out.println("Output = " + s);
        System.out.print("Modified array: {");
        for(int i = 0; i < s; i++){
            System.out.print(nums[i]);
            System.out.print(",");
        }
        System.out.println("}");
    }
}

```

The time complexity of this would be $O(n)$ because the algorithm runs a total of n time where n is the length of num .

The space complexity is $\Omega(1)$ because there are no extra variable being created.