CS 352 Spring 2015 Programming Project Part 1

1. Overview:

For part 1 of the project, your team will implement a simple go-back-N protocol similar to TCP. This protocol is called the 352 Reliable Data Protocol (RDP) version 1 (352 RDP v1). You will realize the 352 RDP v1 as a library of functions written in the C programming language. CS352 RDP uses UDP as the underlying transport protocol. Later versions will add port spaces, concurrency, and security functions.

As part of the project part 1, you will be given 4 files. You can also find them in the sakai site under "Resources" -> "Project resources" -> "Part 1".

- 1. **Makefile**: you must use a makefile that creates the client and server programs. You can extend this makefile but "make all", "make client", "make server" and "make clean" must operate as defined in this file. Your library should be called: sock352lib.o. If you want to extend the makefile, you can create a .a or .so file, but make sure the make commands end up building a client and server.
- 2. **sock352.h**: You may not alter this file. The file defines the interface for the cs 352 socket library that you will implement.
- 3. client.c: You may not alter the source code for this file. It must compile and link against your code.
- 4. **server.c**: You may not alter the source code for this file. It must compile and link against your code.

Your library must implement the following functions as defined in the sock352.h file:

```
int sock352_init(int udp_port);
int sock352_socket(int domain, int type, int protocol);
int sock352_bind(int fd, sockaddr_sock352_t *addr, socklen_t len);
int sock352_connect(int fd, sockaddr_sock352_t *addr, socklen_t len);
int sock352_listen(int fd, int n);
int sock352_accept(int _fd, sockaddr_sock352_t *addr, int *len);
int sock352_close(int fd);
int sock352_read(int fd, void *buf, int count);
int sock352_write(int fd, void *buf, int count);
```

These function map to the existing C-library functions for sockets. See chapter 4, pages 95-120 of Stevens et. al. for the definitions of these functions. The one exception is the <code>sock352_init()</code> call. This call takes a single parameter, which is the UDP port that the rest of the CS 352 RDP library will use for communication between hosts. Setting the udp_port to zero should use the default port of 27182.

For part 1 of the project, you will only need to make a *single connection* work over a *single port* for a *single thread*. The goal is to correctly implement a go-back-N protocol for one connection, for example, when sending a single file between a client and server. Later versions of the project will build

on part to add port-spaces and handle multiple simultaneous connections.

2. The 352 RDP v1 protocol:

Recall as in TCP, 352 RDP v1 maps the abstraction of a logical byte stream onto a model of an unreliable packet network. 352 RDP v1 thus closely follows TCP for the underlying packet protocol. A connection has 3 phases: Set-up, data transfer, and termination. 352 RDP v1 uses a much simpler timeout strategy than TCP for handling lost packets.

Packet structure:

The CS 352 RDP v1 packet as defined as a C-structure, as can be found in the sock352.h file:

typedef struct sock352 pkt hdr sock352 pkt hdr t; /* typedef shortcut */

Note that uintX_t is an X-bit unsigned integer., as defined in <sys/types.h>. At the packet level, all these fields are defined to be in network byte-order (big-endian, most significant byte first).

For part 1, in the packet, the version field should be set of 0x1. The protocol, opt_ptr, source_port and dest_port fields should all be set to zero. Future versions of the protocol will add port spaces and options. The header_len field will always be set to the size of the header, i.e., sizeof(sock352_pkt_hdr_t).

An address for the CS 352 RDP is slightly different from the normal socket address, as found in the sock352.h file. The main difference is the addition of a port layer on top of the UDP port space, as seen in the cs352_port field. This will be used in later versions of the protocol.

Connection Set-up:

352 RDP follows the same connection management protocol as TCP. See Chapter 3.5.6, pages 252-258 of Kurose and Ross for a more detailed description. The bit flags needed are set in the flags field of the packet header. The exact bit definitions of the flags are defined in the sock352.h file.

The client initiates a connection by sending a packet with the SYN bit set in the flags field, picking a random sequence number, and setting the sequence_no field to this number. If no connection is currently open, the server responds with both the SYN and ACK bits set, picks a random number for it's sequence_no field and sets the ack_no field to the client's incoming sequence_no+1. If there is an existing connection, the server responds with the sequence_no+1, but the RST flag set.

Data exchange:

352 RDP follows a simplified Go-Back-N protocol for data exchange, as described in section Kurose and Ross., Chapter 3.4.3, pages 218-223 and extended to TCP style byte streams as described in Chapter 3.5.2, pages 233-238.

When the client sends data, if it is larger than the maximum UDP packet size (64K bytes), it is first broken up into segments, that is, parts of the application byte-stream, of up to 64K. If the client makes a call smaller than 64K, then the data is sent in a single UDP packet of that size, with the payload_len field set appropriately. Segments are acknowledged as the last segment received inorder (that is, go-back-N). Data is delivered to the higher level application in-order based on the read() calls made. If insufficient data exists for a read() call, partial data can be returned and the number of bytes set in the call's return value.

For CS 352 RDP version 1, the server should always set the advertised window to 1024 KB. The client should thus not allow more than 1024 KB to be outstanding (not acknowledged).

Timeouts and retransmissions:

352 RDP v1 uses a single timer model of timeouts and re-transmission, similar to TCP in that there should be a *single timer per connection*, although each segment has a logical *timeout*. The timeout for a segment is 0.2 seconds. That is, if a packet has not been acknowledged after 0.2 seconds it should be re-transmitted, and the logical timeout would be set again set to 0.2 seconds in the future for that segment. The timeout used for a connection should be the timeout of the oldest segment.

There are two strategies for implementing timeouts. One approaches uses Unix signals and other uses a separate thread. These will be covered in class and recitation.

Connection termination:

Connection termination will follow a similar algorithm as TCP, although simpified. In this model, each side closes it's send side separtately, see pages 255-256 of Kurose and Ross and pages 39-40 of Stevens. In version 1, it is OK for the client to end the connection with a FIN bit set when it both gets the last ACK and close has been called. That is, close cannot terminate until the last ACK is received from the server. The sever can terminate the connection under the same confitions.

3. Grading:

Functionality: 80%

Style: 20%

Functionality:

We will run the client.c program linked to our library (called the 'course client)' against the server.c program linked against your library (the 'student server'), and the client.c linked to your library (the 'student client') against the server.c linked to our library ('course server'). We will send a file and see if the checksum on the client and server match the correct checksums. The client.c program opens a single file, sends to the server, and then both exit. See the source code for more details. The size of the file may range from a few bytes to many megabytes. There will be a total of 4 tests, as below, and each test is worth 20% of the total grade:

- (1) student client, course server, in-order packets.
- (2) course client, student server, in-order packets,
- (3) student client, course server, random 20% packets dropped by the course library.
- (4) course client, student server, random 20% packets dropped by the course library...

Style:

Style points are given by the instructor and TA after reading the code. Style is subjective, but will be graded on a scale from 1-5 where 1 is incomprehensible code and 5 means it is perfectly clear what the programmer intended.

4. What to hand in

You must hand in a single archived file, either zip, tar, gzipped tar, bzipped tar or WinRAR (.zip, .tar, .tgz, .rar) that contains: (1) a Makefile, as in part 1, (2) the client.c source code, (3) the server.c source code, (4) the sock352.h header file, and (4) all the files for your library source code.

5. Extra resources

For this project, you will need to keep lists and potentially a hash table. The uthash and utlist libraries are simple, easy to use C-ibraries for these purposes. The documentation for uthash is at https://troydhanson.github.io/uthash/ and for utlist is at https://troydhanson.github.io/uthash/ utlist is at https://troydhanson.github.io/uthash/

If you wish to use other 3^{rd} party source code in your project, you must clear them with the instructor first.