ECS 122A Homework 2

Hardy Jones
999397426
Professor Bai
Spring 2014

1. (a) We can expand some of the polynomial

$$(n+2)^{10} = n^{10} + 20n^9 + \cdots + 1024$$

Looking at everything except the first term, we see that the greatest power is n^9. We can set a *little-oh* bound of n^{10}, because $\forall c > 0, \exists n_0 > 0$ such that $0 \leq 20n^9 + \cdots + 1024 < cn^{10}, \forall n \geq n_0$. That is:

$$(n+2)^{10} = n^{10} + o(n^{10})$$

Now, this is $\Theta(n^{10})$ because
$\exists c_1 > 0, c_2 > 0, n_0 > 0$ such that $0 \leq c_1 n^{10} \leq n^{10} + o(n^{10}) \leq c_2 n^{10},\ \forall n \geq n_0$
This simplifies to $0 \leq c_1 \leq 1 + \frac{o(n^{10})}{n^{10}} \leq c_2$.
Now, since $o(n^{10})$ is an upper bound we know the largest the term $\frac{o(n^{10})}{n^{10}}$ can be is 1. So, this further siplifies to $0 \leq c_1 \leq 2 \leq c_2$.
So this is $\Theta(n^{10})$ if we choose appropriate constants. For instance we can choose $c_1 = c_2 = 2$.
Thus, $(n+2)^{10} = \Theta(n^{10})$

(b) We can apply reasoning similar to the previous problem.
We can expand some of the polynomial

$$(n+a)^b = n^b + \binom{b}{1} n^{b-1} a + \cdots + a^b$$

Looking at everything except the first term, we see that the greatest power is n^{b-1}. We can set a *little-oh* bound of n^b, because $\forall c > 0, \exists n_0 > 0$ such that $0 \leq \binom{b}{1} n^{b-1} a + \cdots + a^b < cn^b, \forall n \geq n_0$. That is:

$$(n+a)^b = n^b + o(n^b)$$

Now, this is $\Theta(n^b)$ because
$\exists c_1 > 0, c_2 > 0, n_0 > 0$ such that $0 \leq c_1 n^b \leq n^b + o(n^b) \leq c_2 n^b,\ \forall n \geq n_0$
This simplifies to $0 \leq c_1 \leq 1 + \frac{o(n^b)}{n^b} \leq c_2$.
Now, since $o(n^b)$ is an upper bound we know the largest the term $\frac{o(n^b)}{n^b}$ can be is 1. So, this further siplifies to $0 \leq c_1 \leq 2 \leq c_2$.
So this is $\Theta(n^b)$ if we choose appropriate constants. For instance we can choose $c_1 = c_2 = 2$.
Thus, $(n+a)^b = \Theta(n^b)$

2. "The running time of algorithm A is at least $O(n^2)$," states that the running time of "A" can be greater than $O(n^2)$. This has no meaning because *big-oh* notation states an asymptotic upper bound on a function. Asserting that the running time can be greater than this value negates the reason for using a bound in the first place.

3. (a) Yes. We need to find some positive constants c, n_0 such that $0 \leq 2^{n+1} \leq c2^n \forall n \geq n_0$.

$$0 \leq 2^{n+1} = 2 \cdot 2^n$$

. So we have our $c = 2, n = 1$.
Thus $2^{n+1} = O(2^n)$.

(b) No. We need to show that no constants exist that satisfy the definition.

$$2^{2n} = 2^n \cdot 2^n \leq c2^n$$

.
It's easy to see that no constant will make this equation true for all values of n.
Thus $2^{2n} \neq O(2^n)$.

4. (a) From smallest to largest:

Function	Notes
$\lg n$	
$n \lg n$	
n	
$n^2, n^2 + \lg n$	These are the same asymptotically because the $\lg n$ term will "drop-out"
n^3	
$n - n^3 + 7n^5$	
2^n	

(b) From smallest to largest:

Function	Notes
1	
$\lg \lg n$	
$\lg n$, $\ln n$	The change of base merely changes the value by a constant factor
$(\lg n)^2$	
$n \lg n$	
$\sqrt{n}$, $\sqrt{2}^{\lg n}$, $n^{1+\epsilon}$	$\sqrt{n} = \sqrt{2}^{\lg n}$, $n^{1+\epsilon}$ differs by constant factor
n	
$n^2, n^2 + \lg n$	These are the same asymptotically because the $\lg n$ term will "drop-out"
n^3	
$n - n^3 + 7n^5$	
2^n, 2^{n-1}	See Problem 3.a
e^n	
$n!$	