## MAT 108 HW 6

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§2.5 1 (b) Proof. Let  $S=\{n\in\mathbb{N}|n>33 \land n=4s+5t \text{ for some } s,t\in\mathbb{N}, \text{ with } s\geq 3,t\geq 2\}$ 

Then we have:

• for 
$$s = 6, t = 2$$

$$4(6) + 5(2) = 24 + 10 = 34$$

• for 
$$s = 5, t = 3$$

$$4(5) + 5(3) = 20 + 15 = 35$$

• for 
$$s = 4, t = 4$$

$$4(4) + 5(4) = 16 + 20 = 36$$

• for 
$$s = 3, t = 5$$

$$4(3) + 5(5) = 12 + 25 = 37$$

So,  $S = \{34, 35, 36, 37, \dots, m-1\}$  for some  $m \in \mathbb{N}$ If m > 37, then  $m - 4 \in S$ .

Then we have for some  $s, t \in \mathbb{N}$ , with  $s \geq 3, t \geq 2$ 

$$m-4 = 4s + 5t$$
  
 $m = 4s + 5t + 4$   
 $m = 4s + 4 + 5t$   
 $m = 4(s + 1) + 5t$ 

And since  $s \geq 3, s+1 \geq 3$ . So  $m \in S$ .

Then by the Principle of Complete Induction, every natural number greater than 33 can be written as 4s + 5t, for some  $s, t \in \mathbb{N}$ , with  $s \geq 3, t \geq 2$ .

5 (d) *Proof.* We need to show the base cases and the inductive case.

• Base

$$- n = 1$$

$$f_1 = 1 = 2 - 1 = f_3 - 1 = f_{1+2} - 1$$

$$- n = 2$$

$$f_1 + f_2 = 1 + 1 = 2 = 3 - 1 = f_4 - 1 = f_{2+2} - 1$$

• Inductive Assume  $f_1 + f_2 + \cdots + f_n = f_{n+2} - 1$ Then we have:

$$f_1 + f_2 + \dots + f_n + f_{n+1} = (f_{n+2} - 1) + f_{n+1}$$

$$= f_{n+2} + f_{n+1} - 1$$

$$= f_{n+3} - 1$$

$$= f_{(n+1)+2} - 1$$

Thus for any  $n+1, f_1+f_2+\cdots+f_{n+1}=f_{(n+1)+2}-1$ Since we have shown both the base case and the inductive case, we have shown that for all natural numbers,  $f_1+f_2+\cdots+f_n=f_{n+2}-1$ 

- 6 (b) Proof. Let  $S = \{n \in \mathbb{N} | f_{n+6} = 4f_{n+3} + f_n\}$ . Then we have:
  - for n=1

$$f_{1+6} = f_7 = 13 = 12 + 1 = 4(3) + 1 = 4f_4 + f_1 = 4f_{1+3} + f_1$$

• for n=2

$$f_{2+6} = f_8 = 21 = 20 + 1 = 4(5) + 1 = 4f_5 + f_2 = 4f_{2+3} + f_2$$

• for n=3

$$f_{3+6} = f_9 = 34 = 32 + 2 = 4(8) + 2 = 4f_6 + f_3 = 4f_{3+3} + f_3$$

So  $S = \{1, 2, 3, \dots, m-1\}$  for some  $m \in \mathbb{N}$ . Then we have:

$$f_{(m-1)+6} = 4f_{(m-1)+3} + f_{m-1}$$

$$f_{m+5} = 4f_{m+2} + f_{m-1}$$

$$f_{m+5} + f_{m-2} = 4f_{m+2} + f_{m-1} + f_{m-2}$$

$$= 4f_{m+2} + f_m$$

$$f_{m+5} + f_{m-2} + 4f_{m+1} = 4f_{m+2} + f_m + 4f_{m+1}$$

$$= 4f_{m+2} + 4f_{m+1} + f_m$$

$$= 4(f_{m+2} + f_{m+1}) + f_m$$

$$= 4f_{m+3} + f_m$$

$$f_{m+5} + f_{m-2} + (f_{m+1} + 3f_{m+1}) =$$

$$f_{m+5} + f_{m-2} + (f_{m-1} + f_m) + 3f_{m+1} =$$

$$f_{m+5} + (f_{m-2} + f_{m-1}) + f_m + 3f_{m+1} =$$

$$f_{m+5} + f_m + f_m + (f_{m+1} + 2f_{m+1}) =$$

$$f_{m+5} + f_m + (f_m + f_{m+1}) + 2f_{m+1} =$$

$$f_{m+5} + f_m + f_{m+2} + (f_{m+1} + f_{m+1}) =$$

$$f_{m+5} + f_m + f_{m+2} + (f_{m+1} + f_{m+1}) =$$

$$f_{m+5} + f_m + (f_{m+2} + f_{m+1}) + f_{m+1} =$$

$$f_{m+5} + f_m + f_{m+3} + f_{m+1} =$$

$$f_{m+5} + f_{m+3} + (f_m + f_{m+1}) =$$

$$f_{m+5} + f_{m+3} + (f_m + f_{m+1}) =$$

$$f_{m+5} + f_{m+3} + (f_m + f_{m+1}) =$$

$$f_{m+5} + f_{m+3} + (f_{m+4} + f_{m+2}) =$$

$$f_{m+5} + f_{m+4} =$$

$$f_{m+6} = 4f_{m+3} + f_m$$

So  $m \in S$ .

Then by the Principle of Complete Induction,  $S = \mathbb{N}$ .

(c) Proof. Let  $S = \{n \in \mathbb{N} | \text{ for any } a \in \mathbb{N}, f_a f_n + f_{a+1} f_{n+1} = f_{a+n+1} \}$ Then we have, for any  $a \in \mathbb{N}$ :

• for n=1

$$f_a f_1 + f_{a+1} f_{1+1} = f_a(1) + f_{a+1}(1) = f_a + f_{a+1} = f_{a+2} = f_{a+1+1}$$

• for n=2

$$f_a f_2 + f_{a+1} f_{2+1} = f_a(1) + f_{a+1}(2) = f_a + f_{a+1} + f_{a+1} = f_{a+2} + f_{a+1} = f_{a+3}$$
  
=  $f_{a+2+1}$ 

So  $S = \{1, 2, ..., m - 1\}$  for some  $m \in \mathbb{N}$ . Then we have, for any  $a \in \mathbb{N}$ :

$$f_a (f_{m-1} + f_{m-2}) + f_{a+1} (f_m + f_{m-1}) = (f_a f_{m-1} + f_a f_{m-2}) + (f_{a+1} f_m + f_{a+1} f_{m-1})$$

$$= (f_a f_{m-1} + f_{a+1} f_m) + (f_a f_{m-2} + f_{a+1} f_{m-1})$$

$$= (f_a f_{m-1} + f_{a+1} f_m) + (f_a f_{m-2} + f_{a+1} f_{m-1})$$

$$= (f_{a+m}) + (f_{a+m-1})$$

$$= f_{a+m+1}$$

So  $m \in \mathbb{N}$ .

Thus, by the Principle of Complete Induction,  $S = \mathbb{N}$ .

- (d) Proof. Let  $S = \left\{ n \in \mathbb{N} | f_n = \frac{\alpha^n \beta^n}{\alpha \beta} \right\}$ . Then we have:
  - for n=1

$$f_1 = 1 = \frac{\alpha - \beta}{\alpha - \beta} = \frac{\alpha^1 - \beta^1}{\alpha - \beta}$$

• for n=2

$$f_2 = 1 = \frac{1}{2} + \frac{1}{2} = \frac{1 + \sqrt{5}}{2} + \frac{1 - \sqrt{5}}{2} = \alpha + \beta = \frac{\alpha - \beta}{\alpha - \beta} (\alpha + \beta) = \frac{\alpha^2 - \beta^2}{\alpha - \beta}$$

So  $S = \{1, 2, \dots, m-1\}$ , for some  $m \in \mathbb{N}$ . We should note that

$$\alpha^{2} = \alpha + 1$$

$$\alpha^{m-2}\alpha^{2} = \alpha^{m-2}(\alpha + 1)$$

$$\alpha^m = \alpha^{m-1} + \alpha^{m-2}$$

$$\alpha^{m-1} = \alpha^m - \alpha^{m-2}$$

$$\beta^{2} = \beta + 1$$

$$\beta^{m-2}\beta^{2} = \beta^{m-2} (\beta + 1)$$

$$\beta^{m} = \beta^{m-1} + \beta^{m-2}$$

$$\beta^{m-1} = \beta^{m} - \beta^{m-2}$$

Then we have:

$$f_{m-1} = \frac{\alpha^{m-1} - \beta^{m-1}}{\alpha - \beta}$$

$$= \frac{(\alpha^m - \alpha^{m-2}) - (\beta^m - \beta^{m-2})}{\alpha - \beta}$$

$$= \frac{(\alpha^m - \beta^m) - (\alpha^{m-2} - \beta^{m-2})}{\alpha - \beta}$$

$$= \frac{\alpha^m - \beta^m}{\alpha - \beta} - \frac{\alpha^{m-2} - \beta^{m-2}}{\alpha - \beta}$$

$$= \frac{\alpha^m - \beta^m}{\alpha - \beta} - f_{m-2}$$

$$f_{m-1} + f_{m-2} = \frac{\alpha^m - \beta^m}{\alpha - \beta}$$

$$f_m = \frac{\alpha^m - \beta^m}{\alpha - \beta}$$

So  $m \in \mathbb{N}$ .

Thus by the Principle of Complete Induction,  $S = \mathbb{N}$ .

§3.1 2 (f) 
$$Dom(W) = (-2, 2), Rng(W) = \{3\}.$$

(g) 
$$Dom(W) = \mathbb{R}, Rng(W) = \mathbb{R}.$$

(h) 
$$Dom(W) = \mathbb{R}, Rng(W) = \mathbb{R}.$$

$$x = -5y + 2 \implies 5y = -x + 2 \implies y = \frac{-x + 2}{5}$$
$$R_2^{-1} = \left\{ (x, y) \in \mathbb{R} \times \mathbb{R} | y = \frac{-x + 2}{5} \right\}$$

(d) 
$$x = y^2 + 2 \implies y^2 = x - 2 \implies y = \sqrt{x - 2}$$
 
$$R_4^{-1} = \left\{ (x, y) \in \mathbb{R} \times \mathbb{R} | y = \sqrt{x - 2} \right\}$$

(h) 
$$x = \frac{2y}{y-2} \implies xy - 2x = 2y \implies y(x-2) = 2x \implies y = \frac{2x}{x-2}$$
 
$$R_4^{-1} = \left\{ (x,y) \in \mathbb{R} \times \mathbb{R} | y = \frac{2x}{x-2} \right\}$$

5 (a) 
$$R \circ S = \{(3,5), (5,2)\}$$

(c) 
$$T \circ S = \{(2,1), (3,1), (3,4)\}$$

(e) 
$$S \circ R = \{(1,5), (2,4), (5,4)\}$$

9 For 
$$R \subseteq (A \times B), S \subseteq (B \times C)$$

(a) Proof.

$$\begin{aligned} \operatorname{Dom}(S \circ R) &= \{x \in A | \exists z \in A \times C : (x, z) \in S \circ R \} \\ &= \{x \in A | \exists y \in B, z \in C : (x, y) \in R \wedge (y, z) \in S \} \\ &\subseteq \{x \in A | \exists y \in B : (x, y) \in R \} \\ &= \operatorname{Dom}(R) \end{aligned}$$

So  $Dom(S \circ R) \subseteq Dom(R)$ .

- (b) Let  $R = \{(a, b)\}, S = \emptyset$ . Then  $S \circ R = \emptyset$ , Dom $(R) = \{a\}$ , Dom $(S \circ R) = \emptyset$ . So  $\emptyset \subseteq \{a\}$ , but  $\emptyset \not\supseteq \{a\}$ . So Dom $(S \circ R) \neq \text{Dom}(R)$ .
- (c) Proof.

$$\operatorname{Rng}(S \circ R) = \{ z \in C | \exists x \in A \times C : (x, z) \in S \circ R \}$$

$$= \{ z \in C | \exists x \in A, y \in B : (x, y) \in R \land (y, z) \in S \}$$

$$\subseteq \{ z \in C | \exists y \in B : (y, z) \in S \}$$

$$= \operatorname{Rng}(S)$$

So  $\operatorname{Rng}(S \circ R) \subseteq \operatorname{Rng}(S)$ .  $\square$ A counter example is: Let  $R = \emptyset, S = \{(a, b)\}$ . Then  $S \circ R = \emptyset$ ,  $\operatorname{Rng}(S) = \{b\}$ ,  $\operatorname{Rng}(S \circ R) = \emptyset$ . So  $\emptyset \subseteq \{b\}$ , but  $\emptyset \not\supseteq \{b\}$ . So  $\operatorname{Rng}(S) \not\subseteq \operatorname{Rng}(S \circ R)$ .

12 Proof. Given a set A with m elements and a set B with n elements.

We can form  $A \times B$  as a new set.

We can count the number of elements in  $A \times B$  as:

$$A \times B = \left\{ \underbrace{\underbrace{(a_1, b_1), \dots, (a_1, b_n)}_{n \text{ pairs}}, \underbrace{(a_2, b_1), \dots, (a_2, b_n)}_{n \text{ pairs}}, \dots, \underbrace{(a_m, b_1), \dots, (a_m, b_n)}_{n \text{ pairs}} \right\}$$

So there are  $m \cdot n$  elements in  $A \times B$ .

We can enumerate all subsets by taking the power set of  $A \times B$ .

And we have proved that for any set S of size k there are  $2^k$  subsets in  $\mathcal{P}(S)$ .

So we have  $2^{mn}$  subsets in  $\mathcal{P}(A \times B)$ .

Since each subset in  $\mathcal{P}(A \times B)$  is a subset of  $A \times B$ , every one is a relation from A to B.

So there are  $2^{mn}$  relations from A to B.