|          | Floyd (WA)rshall                     |                  |   | 4.11 DFS                             | 43<br>43<br>44  |
|----------|--------------------------------------|------------------|---|--------------------------------------|-----------------|
|          | Yorkorka                             |                  |   | 4.14 dominator tree                  | $\frac{45}{46}$ |
| C        | Contents                             |                  |   | 4.16 dsu rollback                    | 46              |
| 1        | Binary Search and Ternary Search     | 5                |   | 4.17 erdos gallai                    | 48<br>48        |
| 1        | 1.1 Aplications                      | 5                |   | 4.19 flow with minimum capacities    | 49              |
|          | 1.2 BS                               | 5                |   | 4.20 Floyd Warshall                  | 50              |
|          | 1.3 LowerBound                       | 6                |   | 4.21 Ford Fulkerson                  | 51<br>51        |
|          | 1.4 parallel binary search           | 6                |   | 4.23 hall theorem                    | 52              |
|          | 1.5 recursive parallel binary search | 8                |   | 4.24 hld                             | 53              |
|          | 1.7 TS                               | 8                |   | 4.25 hld edge                        | 54              |
|          | 1.8 UpperBound                       | 8                |   | 4.26 hopcroft karp                   | 56<br>57        |
|          |                                      |                  |   | 4.27 hungarian                       | 57<br>58        |
| <b>2</b> |                                      | 8                |   | 4.29 kuhn                            | 58              |
|          | 2.1 aliens trick                     | 8                |   | 4.30 LCA                             | 59              |
|          | 2.2 bitwise digit dp                 | 9                |   | 4.31 link cut tree edge              | 60              |
|          | 2.3 broken profile                   | 10               |   | 4.32 link cut tree vertex            | 62              |
|          |                                      | 11               |   | 4.33 MatrixDijkstra                  | 63<br>64        |
|          |                                      | 11               |   | 4.34 max matching without one vertex | 66              |
|          |                                      | 12               |   | 4.36 mo dsu                          | 66              |
|          |                                      | 12               |   | 4.37 mo trees                        | 67              |
|          |                                      | 13               |   | 4.38 mo trees edges                  | 69              |
|          |                                      | 14<br>15         |   | 4.39 Prim                            | 70              |
|          |                                      | 15               |   | 4.40 push relabel                    | 71              |
|          |                                      | 15               |   | 4.42 rmg tree                        | 73              |
|          |                                      | 16               |   | 4.43 sack                            | 73              |
|          |                                      | 16               |   | 4.44 scc                             | 74              |
|          |                                      | 17<br>17         |   | 4.45 segtree graph                   | 75              |
|          |                                      | 17<br>18         |   | 4.46 stable matching                 | 76              |
|          |                                      | 19               |   | 4.47 strong orientation              | 77              |
|          |                                      | 19               |   | 4.49 TreeDiameter                    | 78              |
|          |                                      | 20               |   | 4.50 tree isomorfism                 | 78              |
|          |                                      | 21               |   | 4.51 two sat                         | 79              |
|          | 2.23 subset sum                      | $\frac{22}{23}$  |   | 4.52 virtual tree                    | 80              |
|          |                                      | 93               |   | 1 10 1 1                             | 0.4             |
|          | •                                    | Ð                | ) | kactl forked                         | 81              |
| 3        |                                      | $\frac{23}{2}$ 6 | 3 | kactl forked/combinatorial           | 81              |
|          |                                      | 23<br>24         |   | 6.1 IntPerm                          | 81              |
|          |                                      | $\frac{24}{25}$  |   | 6.2 multinomial                      | 82              |
|          |                                      | 26               |   |                                      |                 |
|          | 3.5 kd free                          | 27 <b>7</b>      |   | kactl_forked/data-structures         | 82              |
|          |                                      | 28               |   | 7.1 FenwickTree                      | 82              |
|          |                                      | 29               |   | 7.2 FenwickTree2d                    | 82              |
|          |                                      | 29<br>31         |   | 7.3 HashMap                          | 83              |
|          |                                      | 31               |   | 7.5 LineContainer                    | 83              |
|          | 3.11 polygon area                    | 33               |   | 7.6 Matrix                           | 83              |
|          | 3.12 polygon isomorfism              | 33               |   | 7.7 MoQueries                        | 84              |
|          | 3.13 smallest enclosing circle       | 34               |   | 7.8 OrderStatisticTree               | 84              |
|          |                                      | ~=               |   | 7.9 RMQ                              | 84              |
| 4        | 1                                    | 35               |   | 7.10 SegmentTree                     | 85<br>85        |
|          |                                      | $\frac{35}{36}$  |   | 7.12 Treap                           | 85              |
|          |                                      | 36               |   | 7.13 UnionFind                       | 85              |
|          |                                      | 37               |   | 7.14 UnionFindRollback               | 86              |
|          |                                      | 39               |   |                                      |                 |
|          | 4.6 caminhoeuleriano                 | 39 <b>8</b>      | 3 | kactl forked/geometry                | 86              |
|          |                                      | 40               |   | 8.1 3dHull                           | 86              |
|          |                                      | 41               |   | 8.2 Angle                            | 86              |
|          | 4.9 centroid decomposition2          | 42               |   | 8.3 CircleIntersection               | 01              |

|    | 8.5  | CirclePolygonIntersection  |  | 10.9 ModInverse  |   |
|----|--|--|--|--|---|
|    | 8.6  | CircleTangents   | 87   | 10.10 ModLog   | 107   |
|    | 8.7  | circumcircle   | 88   | 10.11 ModMulLL   |   |
|    | 8.8  | ClosestPair  | 88   | 10.12 ModPow   | 108   |
|    | 8.9  | ConvexHull   | 88   | 10.13 ModSqrt  | 108   |
|    | 8.10   | DelaunayTriangulation  | 88   | 10.14 ModSum   |   |
|    | 8.11   | FastDelaunay   | 89   | 10.15 Modular Arithmetic   |   |
|    |  | HullDiameter   | 89   | 10.16 phiFunction  |   |
|    |  | InsidePolygon  |  | Total part and other transfer and transfer a | 100   |
|    | 8.14   | kdTree   | 90   | 11 141 6-1-1/  | 100   |
|    | 8.15   | linearTransformation   | 90   |  | 109   |
|    |  | lineDistance   | 90   | 11.1 BerlekampMassey   | 109   |
|    |  |  |  | 11.2 Determinant   |   |
|    | 8.17   | LineHullIntersection   | 91   | 11.3 FastFourierTransform  | 110   |
|    | 8.18   | lineIntersection   | 91   | 11.4 FastFourierTransformMod   | 110   |
|    | 8.19   | LineProjectionReflection   | 91   | 11.5 FastSubsetTransform   | 111   |
|    | 8.20   | ManhattanMST   | 92   | 11.6 GoldenSectionSearch   | 111   |
|    | 8.21   | Minimum Enclosing Circle   | 92   | 11.7 HillClimbing  |   |
|    | 8.22   | OnSegment  | 92   | 11.8 IntDeterminant  | 111   |
|    | 8.23   | Point  | 92   | 11.9 Integrate   |   |
|    | 8.24   | Point3D  | 93   | 11.10 IntegrateAdaptive  |   |
|    | 8.25   | PointInsideHull  | 93   | 11.11 LinearRecurrence   |   |
|    | 8.26   | PolygonArea  | 93   | 11.12 MatrixInverse-mod  |   |
|    | 8.27   | PolygonCenter  | 93   |  |   |
|    | 8.28   | PolygonCut   | 93   | 11.13 MatrixInverse  | 113   |
|    | 8.29   | PolygonUnion   | 94   | 11.14 NumberTheoreticTransform   |   |
|    | 8.30   | PolyhedronVolume   | 94   | 11.15 PolyInterpolate  |   |
|    | 8.31   | SegmentDistance  | 94   | 11.16 Polynomial   |   |
|    | 8.32   | SegmentIntersection  | 94   | 11.17 PolyRoots  |   |
|    | 8.33   | sideOf   | 95   | 11.18 Simplex  |   |
|    |  | sphericalDistance  | 95   | 11.19 SolveLinear  | 115   |
|    | 0.54   | sphericalDistance  | 30   | 11.20 SolveLinear2   | 115   |
| _  | 1 .  |  | 0.5  | 11.21 SolveLinearBinary  | 115   |
| 9  | kact   | $1~{ m forked/graph}$  | 95   | 11.22 Tridiagonal  | 115   |
|    | 9.1  | 2sat   | 95   |  |   |
|    | 0.0  | BellmanFord  | 96   | 101 110 1 1/11   | 116   |
|    | 9.2  |  |  | 12 kactl torked/strings  |   |
|    | $9.2 \\ 9.3$   | BiconnectedComponents  | 96   |  |   |
|    |  |  |  | 12.1 AhoCorasick   | 116   |
|    | 9.3  | BiconnectedComponents  | 96   | 12.1 AhoCorasick   | $\frac{116}{117}$   |
|    | $9.3 \\ 9.4$   | BiconnectedComponents  | 96<br>96   | 12.1 AhoCorasick   | $\frac{116}{117}$   |
|    | 9.3<br>9.4<br>9.5  | BiconnectedComponents  | 96<br>96<br>97   | 12.1 AhoCorasick   | 116<br>117<br>117<br>117  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic   | 96<br>96<br>97<br>97<br>97   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher   | 116<br>117<br>117<br>117<br>118   |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8   | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST   | 96<br>96<br>97<br>97   | 12.1 AhoCorasick   | 116<br>117<br>117<br>117<br>118<br>118  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring  | 96<br>96<br>97<br>97<br>97<br>98<br>98   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray   | 116<br>117<br>117<br>117<br>118<br>118<br>118   |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp  | 96<br>96<br>97<br>97<br>97<br>98<br>98   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>118  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>118  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>118  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>118  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>99   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>118<br>119   |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>118<br>119   |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocator 13.2 BumpAllocatorSTL   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119   |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>100<br>101   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>100<br>101   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocator STL 13.3 ConstantIntervals 13.4 DivideAndConquerDP  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>119   |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.19  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>100<br>101<br>101  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>119<br>119<br>120<br>120  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.19  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>101   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocator 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>119<br>120<br>120   |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.19<br>9.20<br>9.21  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumCliques MaximumCliques  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>120<br>120<br>120<br>120  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.19<br>9.20<br>9.21<br>9.21  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximulCliques MaximumIndependentSet  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>120<br>120<br>120<br>120  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximumCliques MaximumCliques MaximumIndependentSet MinCostMaxFlow  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102<br>103   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalCover  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>120<br>120<br>120<br>121  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximumClique MaximumClique MaximumIndependentSet MinCostMaxFlow MinCut   | 96<br>96<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102<br>103<br>103  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121                                    |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25   | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumIndependentSet MinCott MinCut MinCut MinimumVertexCover   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>103<br>103<br>104  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS  | 1116<br>1117<br>1117<br>1117<br>1118<br>1118<br>1118<br>1119<br>1119  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26   | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximumClique MaximumIndependentSet MinCut MinCut MinCut Mincut MinimumVertexCover PushRelabel  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>103<br>103<br>104<br>104   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD  | 1116<br>1117<br>1117<br>1117<br>1118<br>1118<br>1118<br>1119<br>1119  |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.19<br>9.20<br>9.21<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp Euler Walk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumClique MaximumIndependentSet MinCut MinimumVertexCover PushRelabel SCC   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>103<br>103<br>104<br>104<br>104  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>121               |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximumCliques MaximumCliques MaximumClique MaximumIndependentSet MinCostMaxFlow MinCut MinimumVertexCover PushRelabel SCC TopoSort   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102<br>103<br>103<br>104<br>104<br>104<br>104  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalContainer 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.13 SmallPtr 13.14 TernarySearch  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>121        |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp Euler Walk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumClique MaximumIndependentSet MinCut MinimumVertexCover PushRelabel SCC   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102<br>103<br>103<br>104<br>104<br>104<br>104  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>121        |
|    | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximumCliques MaximumCliques MaximumClique MaximumIndependentSet MinCostMaxFlow MinCut MinimumVertexCover PushRelabel SCC TopoSort   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102<br>103<br>103<br>104<br>104<br>104<br>104  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalContainer 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.13 SmallPtr 13.14 TernarySearch  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>121        |
| 10 | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.27<br>9.29  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumIndependentSet MinCott MinCut MinCut MinimumVertexCover PushRelabel SCC TopoSort WeightedMatching   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102<br>103<br>103<br>104<br>104<br>104<br>104  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocator 13.2 BumpAllocatorsTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.14 TernarySearch 13.15 Unrolling  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>121        |
| 10 | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumIndependentSet MinCut MinimumVertexCover PushRelabel SCC TopoSort WeightedMatching  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>103<br>103<br>104<br>104<br>104<br>105   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.14 TernarySearch 13.15 Unrolling   | 1116<br>1117<br>1117<br>1118<br>1118<br>1118<br>1119<br>1119<br>1119  |
| 10 | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.20<br>9.21<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.20<br>9.20<br>9.21<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.20<br>9.21<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.29<br>9.20<br>9.21<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.20<br>9.21<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.20<br>9.21<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9 | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumClique MaximumIndependentSet MinCostMaxFlow MinCut MinimumVertexCover PushRelabel SCC TopoSort Weighted/number-theory ContinuedFractions  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>103<br>103<br>104<br>104<br>104<br>105<br>105  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalCover 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.14 TernarySearch 13.15 Unrolling  14 Math 14.1 baby step gigant step   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>122<br>122        |
| 10 | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.29<br>9.29<br>9.20<br>9.21<br>9.25<br>9.20<br>9.21<br>9.25<br>9.20<br>9.21<br>9.25<br>9.20<br>9.21<br>9.25<br>9.20<br>9.21<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.29<br>9.29<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.21<br>9.29<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9.29<br>9 | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumClique MaximumIndependentSet MinCostMaxFlow MinCut MinimumVertexCover PushRelabel SCC TopoSort Weighted/number-theory ContinuedFractions CRT  | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102<br>103<br>104<br>104<br>104<br>105<br>105  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocator 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.14 TernarySearch 13.15 Unrolling  14 Math 14.1 baby step gigant step 14.2 berlekamp massey  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>120<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>122<br>122 |
| 10 | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.20<br>9.21<br>9.21<br>9.21<br>9.22<br>9.23  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumIndependentSet MinCostMaxFlow MinCut MinimumVertexCover PushRelabel SCC TopoSort WeightedMatching  I forked/number-theory ContinuedFractions CRT Eratosthenes   | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>103<br>103<br>104<br>104<br>104<br>105<br>105<br>105   | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalCover 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.14 TernarySearch 13.15 Unrolling  14 Math 14.1 baby step gigant step 14.2 berlekamp massey 14.3 binomial theorem   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>122<br>122        |
| 10 | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.20<br>9.20<br>9.21<br>9.21<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.20<br>9.21<br>9.21<br>9.21<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.20<br>9.21<br>9.20<br>9.21<br>9.21<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29<br>9.20<br>9.21<br>9.20<br>9.21<br>9.21<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.26<br>9.27<br>9.28<br>9.29  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumIndependentSet MinCostMaxFlow MinCut MinimumVertexCover PushRelabel SCC TopoSort WeightedMatching I forked/number-theory ContinuedFractions CRT Eratosthenes euclid                                       | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>103<br>103<br>104<br>104<br>104<br>105<br>105  | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.14 TernarySearch 13.15 Unrolling  14 Math 11 baby step gigant step 14.2 berlekamp massey 14.3 binomial theorem 14.4 catalan  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>122<br>122               |
| 10 | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.20<br>9.21<br>9.23<br>9.24<br>9.25<br>9.27<br>9.28<br>9.29<br>9.20<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10<br>10.10  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp Euler Walk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximulCliques MaximumClique MaximumIndependentSet MinCostMaxFlow MinCut MinimumVertexCover PushRelabel SCC TopoSort WeightedMatching I forked/number-theory ContinuedFractions CRT Eratosthenes euclid Factor                 | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>103<br>103<br>104<br>104<br>104<br>105<br>105<br>105                                     | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.14 TernarySearch 13.15 Unrolling  14 Math 14.1 baby step gigant step 14.2 berlekamp massey 14.3 binomial theorem 14.4 catalan 14.5 crivo   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>122<br>122<br>122 |
| 10 | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.17<br>9.18<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.29<br>9.29<br>9.20<br>9.21<br>10.2<br>10.3<br>10.4<br>10.5<br>10.5<br>10.6  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp EulerWalk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximalCliques MaximumClique MaximumIndependentSet MinCostMaxFlow MinCut MinimumVertexCover PushRelabel SCC TopoSort WeightedMatching I forked/number-theory ContinuedFractions CRT Eratosthenes euclid Factor FastEratosthenes | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>102<br>103<br>103<br>104<br>104<br>105<br>105<br>105<br>106<br>106<br>106<br>106         | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.14 TernarySearch 13.15 Unrolling  14 Math 14.1 baby step gigant step 14.2 berlekamp massey 14.3 binomial theorem 14.4 catalan 14.5 crivo 14.6 crt  | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>122<br>122<br>122        |
| 10 | 9.3<br>9.4<br>9.5<br>9.6<br>9.7<br>9.8<br>9.9<br>9.10<br>9.11<br>9.12<br>9.13<br>9.14<br>9.15<br>9.16<br>9.19<br>9.20<br>9.21<br>9.22<br>9.23<br>9.24<br>9.25<br>9.27<br>9.28<br>9.29<br>9.20<br>10.1<br>10.2<br>10.1<br>10.3<br>10.4<br>10.5<br>10.6<br>10.7  | BiconnectedComponents BinaryLifting CompressTree DFSMatching Dinic DirectedMST EdgeColoring EdmondsKarp Euler Walk FloydWarshall GeneralMatching GlobalMinCut GomoryHu HLD hopcroftKarp LCA LinkCutTree MaximulCliques MaximumClique MaximumIndependentSet MinCostMaxFlow MinCut MinimumVertexCover PushRelabel SCC TopoSort WeightedMatching I forked/number-theory ContinuedFractions CRT Eratosthenes euclid Factor                 | 96<br>96<br>97<br>97<br>97<br>98<br>98<br>99<br>99<br>100<br>100<br>101<br>101<br>101<br>102<br>103<br>103<br>104<br>104<br>104<br>105<br>105<br>105<br>106<br>106<br>106<br>106<br>107<br>107 | 12.1 AhoCorasick 12.2 Hashing-codeforces 12.3 Hashing 12.4 KMP 12.5 Manacher 12.6 MinRotation 12.7 SuffixArray 12.8 SuffixTree 12.9 Zfunc  13 kactl forked/various 13.1 BumpAllocator 13.2 BumpAllocatorSTL 13.3 ConstantIntervals 13.4 DivideAndConquerDP 13.5 FastInput 13.6 FastKnapsack 13.7 FastMod 13.8 IntervalContainer 13.9 IntervalCover 13.10 KnuthDP 13.11 LIS 13.12 SIMD 13.13 SmallPtr 13.14 TernarySearch 13.15 Unrolling  14 Math 14.1 baby step gigant step 14.2 berlekamp massey 14.3 binomial theorem 14.4 catalan 14.5 crivo   | 116<br>117<br>117<br>117<br>118<br>118<br>118<br>119<br>119<br>119<br>120<br>120<br>120<br>121<br>121<br>121<br>121<br>121<br>122<br>122        |

| 14.9 division trick   |   |
|---|---|
| 14.10 divisors  |   |
| 14.11 extended euclidean  | 129 18.10 mergesorttree   |
| 14.12 fft   | 129 18.11 min queue   |
| 14.13 fraction  | 130 18.12 mo  |
| 14.14 fwht  | 130 18.13 mo update   |
| 14.15 gaussian elimination  | 131 18.14 persistent seg  |
| 14.16 gaussian elimination2   |   |
| 14.17 lagrange  | 133 18.16 rmq   |
| 14.18 lucas theorem   | 134 18.17 SegTree   |
| 14.19 markov  | 134 18.18 Segtree2  |
|   |   |
| 14.20 matrix exponentiation   | 136 18.20 segtree lazy  |
| 14.21 matrix exponentiation2  | 100 10.20 Segiree lazy  |
| 14.22 matrix inverse and determinant  | 137 18.21 segtree max seg sum   |
| 14.23 max xor subsequence   |   |
| 14.24 mobius  | 138 18.23 sparsetable   |
| 14.25 mobius2   |   |
| 14.26 modular arithmetic  | $140 	 18.25 	ext{ treap2} \dots \dots 182$   |
| 14.27 ntt   | 140   |
| 14.28 operadores binarios   | 141 <b>19 Theorems and Formulas</b> 183   |
| 14.29 pollard rho   |   |
| 14.30 primefactors  |   |
| 14.31 primefactors2   | 149   |
| 14.32 segmentedsieve  |   |
| 14.33 simplex   |   |
|   | 1.10  |
| 14.34 stars and bars  |   |
| 14.35 totient   | 140   |
| 14.36 xor trie  | 21 ufmg forked/DP   |
|   | 01.1 -1-D /   |
| 15 Miscellaneous  | 47 21.1 dcDp  |
| 15.1 bitmasks   |   |
| 15.2 coordinate compression   | 144 21.3 mochila  |
| 15.3 inversion count  | $\begin{array}{cccccccccccccccccccccccccccccccccccc$  |
| 15.4 max plus convolution   |   |
| 15.5 meetinthemiddle  | 150   |
| 15.6 prefix sum 2d  | 187 22 utmg forked/Estruturas   |
|   |   |
|   |   |
| 15.7 rectangle union  | $\begin{array}{cccccccccccccccccccccccccccccccccccc$  |
| 15.7 rectangle union  | $\begin{array}{cccccccccccccccccccccccccccccccccccc$  |
| 15.7       rectangle union         15.8       segment covering         15.9       sprague grundy  | 151       22.1 bit  |
| 15.7       rectangle union         15.8       segment covering         15.9       sprague grundy         15.10       stack trick  | 151       22.1 bit       187         152       22.2 bit2d       187         153       22.3 bitRange       187         153       22.4 bitSortTree       188  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash   | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.5     cht     188  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi  | 151       22.1       bit       187         152       22.2       bit2d       187         153       22.3       bitRange       187         153       22.4       bitSortTree       188         154       22.5       cht       188         154       22.6       chtDinamico       188         154       22.6       chtDinamico       188   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash   | 151       22.1 bit       187         152       22.2 bit2d       187         153       22.3 bitRange       187         153       22.4 bitSortTree       188         154       22.5 cht       188         154       22.6 chtDinamico       188         155       22.7 dsu       189   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi  | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers   | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189       55     22.9     lichaoLazy     190  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers   | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       22.10     mergeSortTree     190  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set  | $\begin{array}{cccccccccccccccccccccccccccccccccccc$  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers   | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       155     22.7     dsu     188       155     22.7     dsu     189       22.8     lichao     189       55     22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL   | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichao Lazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       22.12     minqueueStack     191       22.13     orderStatisticSet     192   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17 Strings  | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       56     22.14     priorityOueueDs     192   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17 Strings 17.1 aho corasick  | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       156     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       156     22.14     priorityQueueDs     192       156     22.15     rangeColor     192  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17 Strings 17.1 aho corasick 17.2 de bruijin  | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       155     22.7     dsu     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       156     22.14     priorityQueueDs     192       156     22.15     rangeColor     192       157     22.16     rmq     193  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17 Strings 17.1 aho corasick  | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       156     22.14     priorityQueueDs     192       157     22.15     rangeColor     192       157     22.16     rmq     193       158     22.17     slopeTrick     193       158     22.17     slopeTrick     193  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17 Strings 17.1 aho corasick 17.2 de bruijin  | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       156     22.14     priorityQueueDs     192       157     22.15     rangeColor     192       157     22.16     rmq     193       158     22.17     slopeTrick     193       158     22.17     slopeTrick     193  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp   | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       56     22.14     priorityQueueDs     192       157     22.16     rmq     193       158     22.17     slopeTrick     193       158     22.17     slopeTrick     193       158     22.18     sparseTable     194   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix   | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitSortTree     188       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       156     22.14     priorityQueueDs     192       157     22.16     rmq     193       158     22.17     slopeTrick     193       158     22.18     sparseTable     194       159     22.19     sparseTableDisjunta     194   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp  | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       156     22.14     priorityQueueDs     192       156     22.15     rangeColor     192       157     22.16     rmq     193       158     22.17     slopeTrick     193       158     22.18     sparseTable     194       159     22.20     splaytree     194       159     22.20     splaytree     194   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing   | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       156     22.14     priorityQueueDs     192       157     22.16     rmq     193       158     22.17     slopeTrick     193       158     22.18     sparseTable     194       159     22.19     sparseTableDisjunta     194       159     22.20     splaytree     194       160     22.21     splaytreeImplicita     195   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.8 stringhashing 17.8 stringhashing 17.8 stringhashing2   | 151       22.1       bit       187         152       22.2       bit2d       187         153       22.3       bitRange       188         153       22.4       bitSortTree       188         154       22.5       cht       188         154       22.6       chtDinamico       188         155       22.7       dsu       189         22.8       lichao       189         22.9       lichao Lazy       190         155       22.10       mergeSortTree       190         155       22.11       minqueueDeque       191         156       22.12       minqueueStack       191         22.13       orderStatisticSet       192         156       22.14       priorityQueueDs       192         156       22.15       rangeColor       192         157       22.16       rmq       193         158       22.17       slopeTrick       193         158       22.18       sparseTable       194         159       22.20       splaytree       194         160       22.21       splaytreeImplicita       195         161   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fift  | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitSortTree     188       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.11     minqueueStack     191       22.12     22.12     minqueueStack     191       22.13     orderStatisticSet     192       156     22.14     priorityQueueDs     192       157     22.16     rmq     192       158     22.17     slopeTrick     193       158     22.17     slopeTrick     193       159     22.19     sparseTableDisjunta     194       159     22.20     splaytree     194       160     22.21     splaytreeImplicita     195       161     22.22     splitMergeSet     196       161     22.23     splitMergeSet   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fit 17.10 suffix array   | 151     22.1     bit     187       152     22.2     bit2d     187       153     22.3     bitRange     187       153     22.4     bitSortTree     188       154     22.5     cht     188       154     22.6     chtDinamico     188       155     22.7     dsu     189       22.8     lichao     189       22.9     lichaoLazy     190       155     22.10     mergeSortTree     190       155     22.11     minqueueDeque     191       156     22.12     minqueueStack     191       22.13     orderStatisticSet     192       156     22.14     priorityQueueDs     192       157     22.16     rmq     192       157     22.16     rmq     193       158     22.17     slopeTrick     193       158     22.18     sparseTable     194       159     22.20     splaytree     194       160     22.21     splaytreeImplicita     195       161     22.22     splitMergeSet     196       162     22.24     splitMergeSetLazy     197       162     22.24     spltTree     199   <  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fit 17.10 suffix array 17.11 suffix automaton   | 151       22.1       bit       187         152       22.2       bit2d       187         153       22.3       bitRange       188         153       22.4       bitSortTree       188         154       22.5       cht       188         154       22.6       chtDinamico       188         155       22.7       dsu       189         22.8       lichao       189         22.9       lichao 189       189         22.10       mergeSortTree       190         155       22.10       mergeSortTree       190         155       22.11       minqueueDeque       191         156       22.12       minqueueStack       191         22.13       orderStatisticSet       192         156       22.14       priorityQueueDs       192         157       22.16       rmq       193         158       22.15       rangeColor       193         158       22.17       slopeTrick       193         158       22.18       sparseTable Disjunta       194         159       22.20       splaytree       194         160       22.21  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fit 17.10 suffix array   | 151       22.1       bit       187         152       22.2       bit2d       187         153       22.4       bitSortTree       188         153       22.4       bitSortTree       188         154       22.5       cht       188         154       22.6       chtDinamico       188         155       22.7       dsu       189         22.8       lichao       189         22.9       lichaoLazy       190         155       22.10       mergeSortTree       190         155       22.11       minqueueDeque       191         156       22.12       minqueueStack       191         22.13       orderStatisticSet       192         256       22.12       priorityQueueDs       192         157       22.16       rmq       193         158       22.17       slopeTrick       193         158       22.17       slopeTrick       193         159       22.19       sparseTable Disjunta       194         159       22.20       splaytree Implicita       195         161       22.22       splitMergeSet Lazy       196 <t< th=""></t<>                    |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fft 17.10 suffix array 17.11 suffix automaton 17.12 z-function  | 151       22.1       bit       187         152       22.2       bit2d       187         153       22.3       bitRange       187         153       22.4       bitSortTree       188         154       22.5       cht       188         154       22.6       chtDinamico       188         155       22.7       dsu       189         22.8       lichao       189         22.9       lichaoLazy       190         155       22.10       mergeSortTree       190         155       22.11       minqueueDeque       191         156       22.12       minqueueStack       191         22.13       orderStatisticSet       192         56       22.14       priorityQueueDs       192         156       22.15       rangeColor       192         157       22.16       rmq       193         158       22.17       slopeTrick       193         158       22.17       slopeTrick       193         159       22.19       sparseTable       194         160       22.21       splaytreeImplicita       195         161  |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fit 17.10 suffix array 17.11 suffix automaton 17.12 z-function   | 151       22.1       bit       187         152       22.2       bit2d       187         153       22.3       bitRange       187         153       22.4       bitSortTree       188         154       22.5       cht       188         154       22.6       chtDinamico       188         155       22.7       dsu       189         22.8       lichao       189         22.9       lichaoLazy       190         155       22.10       mergeSortTree       190         155       22.11       minqueueDeque       191         156       22.11       minqueueStack       191         22.12       minqueueStack       191         22.13       orderStatisticSet       192         256       22.14       priorityQueueDs       192         157       22.15       rangeColor       192         157       22.16       rmq       193         158       22.17       slopeTrick       193         158       22.18       sparseTable       194         159       22.20       splaytree       194         160       22.21   |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fft 17.10 suffix array 17.11 suffix automaton 17.12 z-function  | 151       22.1 bit       187         152       22.2 bit2d       187         153       22.3 bitRange       187         153       22.4 bitSortTree       188         154       22.5 cht       188         155       22.6 chtDinamico       188         155       22.7 dsu       189         22.8 lichao       189         55       22.9 lichaoLazy       190         155       22.10 mergeSortTree       190         155       22.11 minqueueDeque       191         22.12 minqueueStack       191         22.13 orderStatisticSet       192         156       22.14 priorityQueueDs       192         156       22.15 rangeColor       192         157       22.16 rmq       193         158       22.17 slopeTrick       193         158       22.18 sparseTable       194         159       22.19 sparseTableDisjunta       194         159       22.20 splaytree       194         160       22.21 splaytreeImplicita       195         161       22.23 splitMergeSet       196         162       22.24 sqrtTree       199         163       22.25 treap    |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fit 17.10 suffix array 17.11 suffix automaton 17.12 z-function   | 151       22.1       bit       187         152       22.2       bitRange       187         153       22.4       bitSortTree       188         154       22.5       cht       188         154       22.6       chtDinamico       188         155       22.7       dsu       189         22.8       lichao       189         22.9       lichaoLazy       190         155       22.10       mergeSortTree       190         155       22.11       minqueueDeque       191         156       22.12       minqueueStack       191         22.13       orderStatisticSet       192         156       22.14       priorityQueueDs       192         157       22.15       rangeColor       192         157       22.16       rmq       193         158       22.17       slopeTrick       193         158       22.17       slopeTrick       193         158       22.19       sparseTableDisjunta       194         159       22.19       sparseTableDisjunta       194         160       22.21       splitMergeSet       196                                       |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL 17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fit 17.10 suffix array 17.11 suffix automaton 17.12 z-function  18 Structures 18.1 binary lifting                               | 151       22.1       bit       187         152       22.2       bitRange       187         153       22.4       bitSortTree       188         154       22.5       cht       188         154       22.6       chtDinamico       188         155       22.7       dsu       189         22.8       lichao       189         22.9       lichaoLazy       190         155       22.10       mergeSortTree       190         155       22.11       minqueueDeque       191         156       22.12       minqueueStack       191         22.13       orderStatisticSet       192         156       22.14       priorityQueueDs       192         157       22.15       rangeColor       192         157       22.16       rmq       193         158       22.17       slopeTrick       193         158       22.17       slopeTrick       193         158       22.19       sparseTableDisjunta       194         159       22.19       sparseTableDisjunta       194         160       22.21       splitMergeSet       196                                       |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.9 substring fft 17.10 suffix array 17.11 suffix automaton 17.12 z-function  18 Structures 18.1 binary lifting 18.2 bit2d                                      | 151       22.1       bit       187         152       22.2       bitRange       187         153       22.3       bitRange       188         153       22.4       bitSortTree       188         154       22.5       cht       188         154       22.6       chtDinamico       188         155       22.7       dsu       189         22.8       lichao       189         22.8       lichao       189         22.9       lichaoLazy       190         155       22.1       minqueueSortTree       190         155       22.1       minqueueStack       191         22.12       minqueueStack       191         22.13       orderStatisticSet       192         156       22.14       priorityQueueDs       192         157       22.16       rmq       192         157       22.16       rmq       193         158       22.17       slopeTrick       193         158       22.17       slopeTrick       193         158       22.17       slopeTrick       194         159       22.20       splaytree <t< th=""></t<>                                      |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers   16 STL  16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fit 17.10 suffix automaton 17.12 z-function  18 Structures 18.1 binary lifting 18.2 bit2d 18.3 color update 18.4 fenwick     | 151       22.1 bit       187         152       22.2 bit2d       187         153       22.3 bitRange       188         153       22.4 bitSortTree       188         154       22.5 cht       188         155       22.6 chtDinamico       188         155       22.7 dsu       189         22.8 lichao       189         22.9 lichaoLazy       190         155       22.10 mergeSortTree       190         155       22.11 minqueueDeque       191         22.12 minqueueStack       191         22.13 orderStatisticSet       192         156       22.14 priorityQueueDs       192         157       22.16 rmq       192         158       22.17 slopeTrick       193         158       22.17 slopeTrick       193         158       22.18 sparseTable       194         159       22.20 splaytree       194         160       22.21 splaytreeImplicita       195         161       22.22 splitMergeSet       196         161       22.23 splitMergeSetLazy       197         162       22.24 sqrtTree       199         163       22.25 treap       190     |
| 15.7 rectangle union 15.8 segment covering 15.9 sprague grundy 15.10 stack trick 15.11 sum hash 15.12 tower of hanoi 15.13 two pointers  16 STL 16.1 ordered set 16.2 STL  17 Strings 17.1 aho corasick 17.2 de bruijin 17.3 kmp 17.4 manacher 17.5 min suffix 17.6 rabin-karp 17.7 stringhashing 17.8 stringhashing 17.8 stringhashing 17.9 substring fft 17.10 suffix array 17.11 suffix automaton 17.12 z-function  18 Structures 18.1 binary lifting 18.2 bit2d 18.3 color update | 151       22.1 bit       187         152       22.2 bit2d       187         153       22.3 bitRange       188         153       22.4 bitSortTree       188         154       22.5 cht       188         155       22.6 chtDinamico       188         155       22.7 dsu       189         22.8 lichao       189         22.9 lichaoLazy       190         155       22.10 mergeSortTree       190         155       22.11 minqueueDeque       191         22.12 minqueueStack       191         22.13 orderStatisticSet       192         256       22.14 priorityQueueDs       192         157       22.16 rmq       193         158       22.15 rangeColor       192         157       22.16 rmq       193         158       22.17 slopeTrick       193         158       22.18 sparseTable       194         159       22.19 sparseTableDisjunta       194         159       22.20 splaytree       194         160       22.21 splaytreeImplicita       195         161       22.22 splitMergeSet       196         161       22.23 splitMergeSet       20 |

| 23.0   | segTreeEsparsa   |   |   |
|--|--|---|---|
| 23.7   |  |   |   |
| 23.8   |  | 206   | 26. 28.7 cycleDetection   |
| 23.9   | segTreeIterativaComLazy  | 206   | 06 28.8 diofantina  |
| 23.1   | .0 segTreePa   | 206   | 06 28.9 divisionTrick   |
| 23.1   | 1 segTreePersistent  | 207   | 07 28.10 evalInterpol   |
| 23.1   | 2 segTreePersistentComLazy   | 208   |   |
|  |  |   | 28.12 fwht  |
| 24 ufn   | ng forked/Extra  | 208   | 98 28.13 gauss  |
| 24.1   | debug  |   | $28.14 \text{ gaussZ} 2 \dots 232$  |
|  |  |   |   |
| 24.2   |  |   | 00 10 T O 1 1   |
| 24.4   |  |   | 90 17 internal  |
| $24.4 \\ 24.5$   |  |   | 000   |
| 24.6   |  |   | ' 90 10 lomDiagneta 924   |
|  |  |   | 00.00 11 0.04   |
| 24.7   |  |   | 00.01 11  |
| 24.8   | 3 timer  | 209   | 28.22 mulmod  |
|  |  |   | 28.23 multipoint Evaluation 234   |
| 25 ufn   |  | 209   | 99 98 94 ntt 935  |
| 25.1   | articulationPoints   | 209   | 29 28.25 pollardrho   |
| 25.2   | bellmanFord  | 209   | 09 28.26 powerSeries  |
| 25.3   | B blockCutTree   | 210   | 20.20 powerberies 236 28.27 probabilityBinomial   |
| 25.4   |  |   | 28.28 simplex   |
| 25.5   |  |   | 11 26.26 shiplex  |
| 25.6   |  |   |   |
| 25.7   |  |   | 11  |
| 25.8   |  |   | 12 29 umg forked/Primitivas 237   |
| 25.9   |  |   | 29.1 bigint   |
|  | 0 dijkstra   |   |   |
|  | 1 dinitz   |   |   |
|  | 2 directedMst  |   |   |
|  |  |   | 20 5 20 20 20 20 20 20 20 20 20 20 20 20 20   |
|  | 3 dominatorTree  |   | 00.6  |
|  | 4 eulerPath  |   | 00.7  |
|  | 5 eulerTourTree  |   | 00.0  |
| 25.1   | 6 floydWarshall  | 216   | 20.0 madulan Anithmatia   |
|  | 7 functionalGraph  |   | 10  |
|  |  |   |   |
| 25.1   | 8 hopcroftKarp   | 217   | 17 00 f f 1 1/D 11  |
| $25.1 \\ 25.1$   | 8 hopcroftKarp   | 217   | 17 30 ufmg forked/Problemas 247   |
| $25.1 \\ 25.1 \\ 25.2$                                       | 8 hopcroftKarp   | $\frac{217}{218}$   | 17 <b>30 ufmg forked/Problemas</b> 247<br>18 30.1 additionChain   |
| 25.1<br>25.1<br>25.2<br>25.2                                 | 8 hopcroftKarp   | 217 $218$ $218$   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247  |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2                         | 8 hopcroftKarp   | 217<br>218<br>218<br>218  | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2                 | 8 hopcroftKarp 9 johnson 9 johnson 9 loksaraju 11 kruskal 12 kuhn 13 linetree 13 linetree 14 kruskal 15 kruskal 16 kruskal 17 kruskal 17 kruskal 18 linetree 18 kruskal 18  | 217<br>218<br>218<br>218<br>218   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247  |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2                 | 8 hopcroftKarp   | 217<br>218<br>218<br>218<br>218   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp   | 217<br>218<br>218<br>218<br>218<br>219<br>219   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 10 kosaraju 12 kuskal 13 linetree 14 lowerBoundMaxFlow  | 217<br>218<br>218<br>218<br>218<br>219<br>219   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249   |
| 25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp   | 217<br>218<br>218<br>218<br>218<br>219<br>219<br>220  | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp   | 217<br>218<br>218<br>218<br>218<br>219<br>219<br>220<br>220   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage   | 217<br>218<br>218<br>218<br>218<br>219<br>219<br>220<br>220<br>221  | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250   |
| 25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan   | 217<br>218<br>218<br>218<br>218<br>219<br>219<br>220<br>220<br>221<br>221   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort   | 217<br>218<br>218<br>218<br>218<br>219<br>219<br>220<br>220<br>221<br>221   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf  | 217<br>218<br>218<br>218<br>218<br>219<br>219<br>220<br>220<br>221<br>221<br>221<br>221   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         30.12 distinctUpdate       251  |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort   | 217<br>218<br>218<br>218<br>218<br>219<br>219<br>220<br>220<br>221<br>221<br>221<br>221   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         30.12 distinctUpdate       251  |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree   | 217<br>218<br>218<br>218<br>218<br>219<br>220<br>220<br>220<br>221<br>221<br>221<br>221<br>222                                    | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         21       30.12 distinctUpdate       251         22       30.13 dominacao3d       252         30.14 dominatorPoints       253  |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  ng forked/Grafos/LCA-HLD   | 217<br>218<br>218<br>218<br>218<br>219<br>219<br>220<br>220<br>221<br>221<br>221<br>221<br>222<br>222                             | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.8 conectividadeDinamica       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         21       30.12 distinctUpdate       251         22       30.13 dominacao3d       252         30.14 dominatorPoints       253         30.15 dynamicHull       253         30.14 driver       253  |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  18 forked/Grafos/LCA-HLD 18 hldAresta  | 217<br>218<br>218<br>218<br>218<br>219<br>219<br>220<br>220<br>221<br>221<br>221<br>221<br>222<br>222                             | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.8 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.10 delaunay       250         21       30.11 distinct       251         22       30.13 dominacao3d       252         30.14 dominatorPoints       253         30.15 dynamicHull       253         30.16 graph/Triangles       253         30.17 correctales       253         30.17 correctales       253   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta 1 hldSemUpdate  | 217<br>218<br>218<br>218<br>219<br>219<br>220<br>220<br>221<br>221<br>221<br>221<br>222<br>222<br>222<br>222                      | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         21       30.12 distinctUpdate       251         30.13 dominacao3d       252         30.14 dominatorPoints       253         30.15 dynamicHull       253         30.16 graphTriangles       253         30.18 calfClustureraction       254  |
| 25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  11 mg forked/Grafos/LCA-HLD 11 hldAresta 12 hldSemUpdate 13 hldVertice   | 217<br>218<br>218<br>218<br>219<br>220<br>220<br>221<br>221<br>221<br>221<br>222<br>222<br>222<br>223<br>223                      | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         21       30.12 distinctUpdate       251         22       30.13 dominacao3d       252         30.14 dominatorPoints       253         30.15 dynamicHull       253         30.16 graphTriangles       253         30.18 halfPlaneIntersection       254         30.18 halfPlaneIntersection       254  |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  18 forked/Grafos/LCA-HLD 18 hldSemUpdate 18 hldVertice 19 lca  | 217<br>218<br>218<br>218<br>219<br>220<br>220<br>221<br>221<br>221<br>221<br>222<br>222<br>222<br>223<br>223                      | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         22       30.12 distinctUpdate       251         23       30.13 dominacao3d       252         30       30.15 dynamicHull       253         32       30.16 graphTriangles       253         33       30.18 halfPlaneIntersection       254         30       19 heapSort       254   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 33 linetree 14 lowerBoundMaxFlow 55 minCostMaxFlow 66 prufer 77 sack 88 stableMarriage 99 tarjan 60 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta hldVertice lca lca lcaComBinaryLifting  | 217 218 218 218 218 219 220 220 221 221 221 222 223 223 223 223 224 224   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         22       30.13 dominacao3d       252         30.14 dominatorPoints       253         30.15 dynamicHull       253         30.3 30.16 graphTriangles       253         30.3 30.18 halfPlaneIntersection       254         30.3 30.19 heapSort       254         30.20 hungarian       254         30.20 hungarian       254         30.20 hungarian       254   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  18 forked/Grafos/LCA-HLD 18 hldSemUpdate 18 hldVertice 19 lca  | 217 218 218 218 218 219 220 220 221 221 221 222 223 223 223 223 224 224   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         22       30.13 dominacao3d       252         30.14 dominatorPoints       253         30.2       30.15 dynamicHull       253         30.3 30.16 graphTriangles       253         30.3 30.18 halfPlaneIntersection       254         30.4 30.20 hungarian       254         30.21 intervalGraphColoring       254         30.21 intervalGraphColoring       254   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta hldVertice lca lca lcaComBinaryLifting lcaComHld  | 217 218 218 218 218 219 220 221 221 221 222 222 222 222 222 223 223   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         22       30.13 dominacao3d       252         30.14 dominatorPoints       253         30.15 dynamicHull       253         30.3 30.16 graphTriangles       253         30.17 grayCode       253         30.3 30.18 halfPlaneIntersection       254         30.19 heapSort       254         30.21 intervalGraphColoring       254         30.22 intervalGraphIndSet       255         30.22 intervalGraphCount       255  |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta hldVertice lca lca lcaComBinaryLifting lcaComHld  | 217 218 218 218 218 219 220 220 221 221 221 222 223 223 223 223 224 224   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         22       30.13 dominacao3d       252         23       30.14 dominatorPoints       253         25       30.15 dynamicHull       253         22       30.15 dynamicHull       253         23       30.16 graphTriangles       253         23       30.17 grayCode       253         24       30.20 hungarian       254         25       30.21 intervalGraphColoring       254         30.21 intervalGraphIndSet       255         30.23 inversionCount       255  |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta 1 hldSemUpdate 1 hldVertice 1 lca 1 lcaComBinaryLifting 1 lcaComHld  mg forked/Grafos/LCT   | 217<br>218<br>218<br>218<br>219<br>220<br>220<br>221<br>221<br>221<br>221<br>222<br>222<br>222<br>223<br>223<br>224<br>225<br>225 | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         21       30.12 distinctUpdate       251         22       30.13 dominacao3d       252         30.13 dominacao3d       252         30.14 dominatorPoints       253         32       30.15 dynamicHull       253         23       30.18 halfPlaneIntersection       254         24       30.20 hungarian       254         30.21 intervalGraphColoring       254         30.22 intervalGraphIndSet       255         30.23 inversionCount       255         30.24 lis       255  |
| 25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree 18 hldSemUpdate 18 hldSemUpdate 19 lcaComBinaryLifting 19 lcaComHld 19 lcaComHld 19 lcaComHld 19 lcaComHld 19 lcaComHld 10 | 217 218 218 218 218 219 220 220 221 221 221 221 222 223 223 223 223 224 224 225 225   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.11 distinct       251         21       30.12 distinctUpdate       251         22       30.13 dominacao3d       252         30.14 dominatorPoints       253         30.15 dynamicHull       253         30       30.16 graphTriangles       253         33       30.16 graphTriangles       253         23       30.18 halfPlaneIntersection       254         24       30.20 hungarian       254         25       30.21 intervalGraphColoring       254         25       30.23 inversionCount       255         25       30.23 inversion          |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 33 linetree 14 lowerBoundMaxFlow 55 minCostMaxFlow 66 prufer 77 sack 88 stableMarriage 99 tarjan 60 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta hldVertice lca lcaComBinaryLifting lcaComHld  mg forked/Grafos/LCT lct lct lctAresta  | 217 218 218 218 218 219 220 220 221 221 221 222 223 223 223 223 224 224 225 225 225 225   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.5 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.9 deBrujin       250         21       30.10 delaunay       250         21       30.12 distinct Update       251         22       30.13 dominacao3d       252         23       30.14 dominatorPoints       253         25       30.15 dynamicHull       253         25       30.16 graphTriangles       253         23       30.18 halfPlaneIntersection       254         24       30.20 hungarian       254         25       30.21 intervalGraphColoring       254         25       30.23 inversionCount       255         25       30.23 inversionCount       255         25       30.25 lis2       255                            |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree 18 hldSemUpdate 18 hldSemUpdate 19 lcaComBinaryLifting 19 lcaComHld 19 lcaComHld 19 lcaComHld 19 lcaComHld 19 lcaComHld 10 | 217 218 218 218 218 219 220 220 221 221 221 222 223 223 223 223 224 224 225 225 225 225   | 17       30 ufmg forked/Problemas       247         18       30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         20       30.5 conectividadeDinamica       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.10 delaunay       250         21       30.11 distinct       251         21       30.12 distinctUpdate       251         22       30.13 dominacao3d       252         30.14 dominatorPoints       253         23       30.15 dynamicHull       253         23       30.16 graphTriangles       253         23       30.17 grayCode       253         23       30.18 halfPlaneIntersection       254         24       30.21 intervalGraphColoring       254         25       30.21 intervalGraphColoring       254         30.22 intervalGraphIndSet       255         30.23 intervalGraphIndSet       255         30.24 lis       2 |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta 1 hldSemUpdate 1 hldVertice 1 lca 1 caComBinaryLifting 1 caComHld  mg forked/Grafos/LCT   | 217 218 218 218 218 219 220 221 221 221 222 222 223 223 224 224 225 225 225 225 226   | 17       30 ufmg forked/Problemas       247         30.1 additionChain       247         18       30.2 angleRange       247         18       30.3 areaHistograma       247         18       30.4 areaUniaoRetangulo       247         19       30.5 binomial       248         19       30.6 closestPairOfPoints       248         20       30.7 conectividadeDinamica       249         20       30.8 conectividadeDinamica2       249         21       30.10 delaunay       250         21       30.11 distinct       251         21       30.12 distinctUpdate       251         22       30.13 dominacao3d       252         30.13 dominacroPoints       253         23       30.15 dynamicHull       253         23       30.15 dynamicHull       253         23       30.17 grayCode       253         23       30.18 halfPlaneIntersection       254         24       30.20 hungarian       254         25       30.23 inversionCount       255         25       30.24 lis       255         30.25 lis2       255         30.26 maxDist       256         30.27 minCirc<                           |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treelsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta hldSemUpdate hldVertice lca lcaComBinaryLifting lcaComHd  mg forked/Grafos/LCT lct lct lctAresta lctVertice mg forked/Matematica  | 217 218 218 218 218 219 220 220 221 221 221 222 222 222 223 224 224 225 225 225 225 225 225 227                                   | 17     30 ufmg forked/Problemas     247       30.1 additionChain     247       18     30.2 angleRange     247       18     30.3 areaHistograma     247       18     30.4 areaUniaoRetangulo     247       19     30.5 binomial     248       20     30.7 conectividadeDinamica     249       20     30.8 conectividadeDinamica2     249       21     30.9 deBrujin     250       21     30.10 delaunay     250       21     30.11 distinct     251       22     30.13 dominacao3d     252       30.14 dominatorPoints     253       23     30.15 dynamicHull     253       23     30.16 graphTriangles     253       23     30.17 grayCode     253       23     30.18 halfPlaneIntersection     254       24     30.20 hungarian     254       25     30.21 intervalGraphColoring     254       25     30.23 inversionCount     255       30     30.24 lis     255       30     30.28 minkowski     256       30.28 minkowski     256       30.29 mo     256       30.29 mo     256   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 90 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta hldSemUpdate hldVertice lca lcaComBinaryLifting lcaComHld  mg forked/Grafos/LCT lct lctAresta lctVertice  mg forked/Matematica 2sat   | 217 218 218 218 218 219 220 220 221 221 221 221 222 222 223 223 223 223   | 17     30 ufmg forked/Problemas     247       30.1 additionChain     247       18     30.2 angleRange     247       18     30.3 areaHistograma     247       18     30.4 areaUniaoRetangulo     248       19     30.5 binomial     248       19     30.6 closestPairOfPoints     248       20     30.7 conectividadeDinamica     249       20     30.8 conectividadeDinamica2     249       21     30.9 deBrujin     250       21     30.10 delaunay     250       21     30.12 distinct Update     251       22     30.13 dominacao3d     252       30.14 dominatorPoints     253       30.15 dynamicHull     253       32     30.16 graphTriangles     253       33     30.18 halfPlaneIntersection     254       44     30.20 hungarian     254       25     30.21 intervalGraphColoring     254       30     30     256       30.23 inversionCount     255       30     30     256       30     30     256       30     25     30       30     25     30       30     25     30       30     25     30       30   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta hldSemUpdate hldVertice lca lcaComBinaryLifting lcaComHld  mg forked/Grafos/LCT lct lct lctAresta lctVertice  mg forked/Matematica 2sat berlekampMassey   | 217 218 218 218 218 219 220 221 221 221 221 222 223 223 223 223 224 225 225 225 226 227 227 228                                   | 17     30 ufmg forked/Problemas     247       30.1 additionChain     247       18     30.2 angleRange     247       18     30.3 areaHistograma     247       18     30.4 areaUniaoRetangulo     247       19     30.5 binomial     248       19     30.6 closestPairOfPoints     248       20     30.7 conectividadeDinamica     249       20     30.8 conectividadeDinamica     249       21     30.9 deBrujin     250       21     30.10 delaunay     250       21     30.11 distinct     251       21     30.12 distinctUpdate     251       22     30.13 dominacao3d     252       23     30.15 dynamicHull     253       25     30.14 dominatorPoints     253       22     30.15 grayCode     253       23     30.17 grayCode     253       23     30.18 halfPlaneIntersection     254       24     30.20 hungarian     254       25     30.21 intervalGraphColoring     254       25     30.23 inversionCount     255       30     30.24 lis     255       30.25 lis2     255       30     30.26 maxDist     256       30     29 mo     256   |
| 25.1<br>25.1<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2<br>25.2 | 8 hopcroftKarp 9 johnson 10 kosaraju 11 kruskal 12 kuhn 13 linetree 14 lowerBoundMaxFlow 15 minCostMaxFlow 16 prufer 17 sack 18 stableMarriage 19 tarjan 10 topoSort 11 treeIsomorf 12 virtualTree  mg forked/Grafos/LCA-HLD hldAresta hldVertice lca lca lcaComBinaryLifting lcaComHld  mg forked/Grafos/LCT lct lctAresta lctVertice  mg forked/Matematica 2sat berlekampMassey lchinese   | 217 218 218 218 218 219 220 221 221 221 222 222 223 223 223 224 225 225 225 225 226 227 227 228 228                               | 17     30 ufing forked/Problemas     247       30.1 additionChain     247       18     30.2 angleRange     247       18     30.3 areaHistograma     247       19     30.4 areaUniaoRetangulo     247       19     30.5 binomial     248       20     30.6 closestPairOfPoints     248       20     30.7 conectividadeDinamica     249       20     30.8 conectividadeDinamica2     249       21     30.9 deBrujin     250       21     30.10 delaunay     250       21     30.11 distinct     251       22     30.13 dominacao3d     252       23     30.14 dominatorPoints     253       25     30.15 dynamicHull     253       25     30.16 graphTriangles     253       23     30.17 grayCode     253       23     30.18 halfPlaneIntersection     254       24     30.20 hungarian     254       25     30.21 intervalGraphColoring     254       25     30.22 intervalGraphIndSet     255       25     30.25 lis2     255       30.28 minkowski     256       27     30.29 mo     256       30.28 minkowski     256       27     30.30 moDsu     257   |

|    | 30.34     | rmqOffline          |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 259               |
|----|-----------|---------------------|---|---|---|-------|---|---|---|---|-------|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|-------------------|
|    |           | segmentIntersection |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 259               |
|    | 30.36     | simplePolygon       |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 259               |
|    |           | steinerTree         |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |                   |
|    |           | sweepDirection      |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 260               |
|    | 50.50     | sweepDirection ::   | • | • | • | <br>• | • | • | • | • | <br>• | • | • | • | • | • | • | • | • | • | • | • | • | • | • | • | 200               |
| 31 | ufm       | g forked/Strings    |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 261               |
|    | 31.1      | ahocorasick         |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 261               |
|    | 31.2      | dynamicSuffixArray  |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 261               |
|    | 31.3      | eertree             |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 262               |
|    | 31.4      | hashing             |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 263               |
|    | 31.5      | hashingLargeMod .   |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 263               |
|    | 31.6      | kmp                 |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 263               |
|    | 31.7      | manacher            |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | $\frac{264}{264}$ |
|    | 31.8      | minMaxSuffixCyclic  |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 264               |
|    | 31.9      | suffixArray         |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 264               |
|    |           | suffixArray2        |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 265               |
|    |           | suffixAutomaton     |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |                   |
|    |           | trie                |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 266               |
|    |           |                     |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 267               |
|    | 31.13     | z                   | • | • | • | <br>• | • | • | • | • | <br>• | • | • | • | • | • | • | • | • | • | • | • | • | • | • | • | 207               |
|    | T T 4 • 1 |                     |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 007               |
| 32 | Util      |                     |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 267               |
|    | 32.1      | execution time      |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 267               |
|    | 32.2      | rand                |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 267               |
|    | 32.3      | runner              |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 267               |
|    | 32.4      | runner2             |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 268               |
|    | 32.5      | int128              |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   | 268               |
|    |           |                     |   |   |   |       |   |   |   |   |       |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |                   |

# 1 Binary Search and Ternary Search

## 1.1 Aplications

```
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define in insert
#define pi pair<int, int>
#define pii pair<int, pi>
#define mp make_pair
#define fir first
#define sec second
#define MAXN 1001
// 1 - ts para double
long double ts()
  long double 1 = 0, r = DBL_MAX;
  for (int i = 0; i < 2000; i++)
    long double 11 = (1 * 2 + r) / 3.0;
    long double 12 = (1 + 2 * r) / 3.0;
    if (possible(11))
     r = 12;
    else
     1 = 11;
  return 1;
// 2- bb para double
long double bb()
  long double i = 0, f = DBL_MAX, m;
  while (f - i > 0.000000001)
    m = (i + f) / 2.0;
    if (possible(m))
     f = m;
    else
     i = m;
```

```
return i;
// 3 - bb pra int
lli bb()
  lli i = 0, f = INT_MAX, m;
  while (i < f)
    \mathbf{m} = (\mathbf{i} + \mathbf{f}) / 2;
   if (possible(m))
     f = m;
    else
      i = m + 1;
  return i;
// 4 - ts pra int (valor minimo da funcao f(x)), sendo x um inteiro
int 1 = 1, r = INT_MAX;
while (r - 1 > 15)
  int 11 = (1 * 2 + r) / 3;
  int 12 = (1 + 2 * r) / 3;
  (calc(11) < calc(12)) ? r = 12 : 1 = 11;
for (int i = 1; i <= r; i++)
// vejo qual a melhor opcao de l ate r em o(n)
// busca ternaria para int, usando busca binaria:
int 1 = 0, r = 1e9;
while (1 < r)
  int mid = (1 + r) >> 1;
  (calc(mid) < calc(mid + 1)) ? r = mid : 1 = mid + 1;
return calc(1);
```

#### 1.2 BS

```
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
vector <int> v;
int binarysearch (int n , int x)
   int i = 0;
   int f = n - 1;
   int m ;
   while(i <= f)</pre>
       m = (i + f) / 2;
       if(v[m] == x) return m + 1;
       if(v[m] < x) i = m + 1;
       if(v[m] > x) f = m - 1;
   return 0 ;
int main ()
   int n , aux , m ;
   cin >> n ;
    for (int i = 0; i < n; i++)
        cin >> aux ;
        v.pb(aux);
```

```
sort(v.begin() , v.end());
cin >> m;
cout << binarysearch(n , m) << endl;
return 0;</pre>
```

#### 1.3 LowerBound

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 1000001
#define mod 1000000007
// first element >= x
vector<int> k (MAXN);
int lower(int 1, int r, int x) // first element >= x
  while (1 < r)
    int mid = (1 + r) >> 1;
    (x \le k[mid]) ? r = mid : 1 = mid + 1;
  return k[1];
```

## 1.4 parallel binary search

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 300006
#define mod 1000000007
int n;
int b[MAXN];
void reset()
  for (int i = 0; i < MAXN; i++)</pre>
   b[i] = 0;
int sum(int r)
```

```
int ret = 0;
  for (; r >= 0; r = (r & (r + 1)) - 1)
   ret += b[r];
  return ret;
void add(int idx, int delta)
  for (; idx < MAXN; idx = idx | (idx + 1))
   b[idx] += delta;
void update(int 1, int r, int x)
  add(1, x);
  add(r + 1, -x);
void upd(int 1, int r, int x)
  if (1 \le r)
    update(1, r, x);
   return;
  update(1, MAXN - 2, x);
  update(0, r, x);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, m;
  cin >> n >> m;
  vector<vector<int>> adj(n);
  for (int i = 0; i < m; i++)
    int x;
   cin >> x;
   adj[x].pb(i);
  vector<int> need(n);
  for (int i = 0; i < n; i++)
   cin >> need[i];
  int q;
  cin >> q;
  vector<pii> qry(q);
  for (int i = 0; i < q; i++)
    cin >> qry[i].sec.fir >> qry[i].sec.sec >> qry[i].fir;
   qry[i].sec.fir--, qry[i].sec.sec--;
  vector<int> 1(n);
  vector<int> r(n);
  vector<vector<int>> on(q);
  for (int i = 0; i < n; i++)
   1[i] = 0;
   r[i] = q;
  while (1)
   bool ok = 1;
for (int i = 0; i < n; i++)</pre>
     if (l[i] < r[i])
        ok = 0;
        int mid = (l[i] + r[i]) >> 1;
        on[mid].pb(i);
    if (ok)
     break;
    reset();
    for (int mid = 0; mid < q; mid++)</pre>
```

```
upd(qry[mid].sec.fir, qry[mid].sec.sec, qry[mid].fir);
      for (auto const &j : on[mid])
        int val = 0;
        for (auto const &k : adj[j])
          val += sum(k);
          if (val >= need[j])
           break:
        (val >= need[j]) ? r[j] = mid : l[j] = mid + 1;
      on[mid].clear();
 for (int i = 0; i < n; i++)
   if (l[i] >= q)
      cout << "NIE\n";</pre>
     cout << 1[i] + 1 << endl;
 return 0;
// busca binaria paralela
// https://www.spoj.com/problems/METEORS/
// tem n member states e m sectors
// cada sector ta associado a uma member state
// cada query incrementa o range [l[i], r[i]] de sectors por a[i]
// seja q[i] a soma de todos os v[j], sendo j um sector associado ao member
    state i
// qual o primeiro momento no qual q[i] >= min_qt[i]
// para todos os i
// a sagacidade vai ser fzr uma busca binaria pra cada resposta
// primeiro vc faz todas a primeira iteracao de cada busca binaria
// depois cada segunda iteracao de cada bb
// e assim vai
// ai a bb eh so tipo
// a soma apos a query mid ja deu bom pra aquele member state?
```

## 1.5 recursive parallel binary search

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define pb push_back
#define fir first
#define sec second
#define MAXN 300006
#define mod 1000000007
int n, ptr;
int b[MAXN];
int ll[MAXN];
int rr[MAXN];
int qryl[MAXN];
int qryr[MAXN];
int gryg[MAXN];
int need[MAXN];
vector<int> adj[MAXN];
int sum(int r)
  int ret = 0;
  for (; r >= 0; r = (r & (r + 1)) - 1)
   ret += b[r];
  return ret;
void add(int idx, int delta)
  for (; idx < MAXN; idx = idx | (idx + 1))
```

```
b[idx] += delta;
void update(int 1, int r, int x)
  add(1, x);
  add (r + 1, -x);
void upd(int 1, int r, int x)
  if (1 \le r)
    update(1, r, x);
   return;
  update(1, MAXN - 2, x);
  update(0, r, x);
void solve(vector<int> &v, int 1, int r)
  if (!(l < r) || !v.size())</pre>
    return:
  int mid = (1 + r) / 2;
  while (ptr <= mid)</pre>
    upd(qryl[ptr], qryr[ptr], qryg[ptr]);
   ptr++;
  while ((ptr - 1) > mid)
    ptr--;
   upd(qryl[ptr], qryr[ptr], -qryg[ptr]);
  vector<int> to_left, to_right;
  for (auto const & j : v)
    int val = 0;
    for (auto const &k : adj[j])
      val += sum(k);
      if (val >= need[j])
       break;
    if (val >= need[i])
      to_right.pb(j);
      rr[j] = mid;
    else
      to_left.pb(j);
      11[j] = mid + 1;
  v.clear();
  solve(to_left, mid + 1, r);
  solve(to_right, 1, mid);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cout.tie(NULL);
  int n, m;
  cin >> n >> m;
  for (int i = 0; i < m; i++)
    int x:
    cin >> x;
    adj[x].pb(i);
  for (int i = 0; i < n; i++)
   cin >> need[i];
  int q;
```

```
cin >> q;
 for (int i = 0; i < q; i++)
   cin >> qryl[i] >> qryr[i] >> qryg[i];
   qryl[i]--, qryr[i]--;
 vector<int> vec;
 for (int i = 0; i < n; i++)
   11[i] = 0;
   rr[i] = q;
   vec.pb(i);
 solve(vec, 0, q);
 for (int i = 0; i < n; i++)
   if (ll[i] >= q)
     cout << "NIE\n";
    else
     cout << 11[i] + 1 << endl;</pre>
 return 0;
// busca binaria paralela
// https://www.spoj.com/problems/METEORS/
// so que a recursiva
```

#### 1.6 STL

```
// lower - primeiro maior ou iqual a x
// upper - ultimo menor ou igual a x
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push back
vector <int> v ;
int main()
    int n , aux ;
    cin >> n ;
    for (int i = 0; i < n; i++)
        cin >> aux ;
        v.pb(aux);
    sort(v.begin() , v.end());
    cin >> q;
    while (q--)
        vector <int> :: iterator low = lower_bound (v.begin() , v.end() , aux) ;
        vector <int> :: iterator up = upper_bound (v.begin() , v.end() , aux) ;
        cout << (low - v.begin()) << " " << (up - v.begin()) - 1 << endl ;</pre>
    return 0 ;
```

#### $1.7 ext{ TS}$

```
// achar valor maximo da funcao
for (int i = 0; i < 400; i++)
{</pre>
```

```
double m1 = 1 + (r - 1) / 3;
double m2 = r - (r - 1) / 3;
double f1 = f(m1);
double f2 = f(m2);
if (f1 < f2)
1 = m1;
else
r = m2;
```

## 1.8 UpperBound

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 1000001
#define mod 1000000007
// last element <= x
vector<int> k (MAXN);
int upper(int 1, int r, int x)
  while (1 < r)
    int mid = (1 + r + 1) >> 1;
    (k[mid] \le x) ? 1 = mid : r = mid - 1;
  return k[1];
```

# 2 Dynamic programming and common problems

#### 2.1 aliens trick

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
int n, k, 1;
string s;
```

```
pi solve(vector<int> &v, int lambda)
  // associar um custo lambda para ser subtraido quando realizamos uma operacao
  // dp[i] - melhor profit que tivemos considerando as i primeiras posicoes
  // cnt[i] - quantas operacoes utilizamos para chegarno valor de dp[i]
 vector < int > dp(n + 1);
 vector<int> cnt(n + 1);
 dp[0] = 0;
cnt[0] = 0;
  for (int i = 1; i \le n; i++)
    dp[i] = dp[i - 1];
   cnt[i] = cnt[i - 1];
    int id = i - 1;
    dp[i] += v[id];
    int lo = max(011, id - 1 + 1);
    int s = dp[lo] + (id - lo + 1) - lambda;
    if (s > dp[i])
      dp[i] = s;
      cnt[i] = cnt[lo] + 1;
  return {dp[n], cnt[n]};
int aliens_trick(vector<int> &v)
  int 1 = 0, r = n;
 while (1 < r)
   int mid = (1 + r) >> 1;
   pi ans = solve(v, mid);
(ans.sec > k) ? l = mid + 1 : r = mid;
 pi ans = solve(v, 1);
 return ans.fir + (1 * k);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 cin >> n >> k >> 1 >> s;
 vector<int> a(n);
  vector<int> b(n);
 for (int i = 0; i < n; i++)
    a[i] = 1, b[i] = 0;
   if (s[i] >= 'A' \&\& s[i] <= 'Z')
      a[i] ^= 1;
     b[i] ^= 1;
 cout << n - max(aliens_trick(a), aliens_trick(b)) << endl;</pre>
 return 0;
// https://codeforces.com/contest/1279/problem/F
```

## 2.2 bitwise digit dp

```
#define sec second
#define MAXN 65
#define mod 1000000007
// current bit / i is bigger than 1 / i is lower than r / j is bigger than 1 / j
      is lower than r
int dp[MAXN][2][2][2][2];
int solve(int i, int j, int k, int l, int m)
  if (i < 0)
   return (j && k && l && m) ? 1 : 0;
  if (dp[i][j][k][l][m] != -1)
   return dp[i][j][k][l][m];
  int ret = 0:
  int 11 = a & (1LL << i);</pre>
  int rr = b & (1LL << i);</pre>
  if ((j || !11) && (1 || !11))
   ret += solve(i - 1, j, (rr) ? 1 : k, l, (rr) ? 1 : m);
  if ((k || rr) && (1 || !11))
    ret += solve(i - 1, (!11) ? 1 : j, k, l, (rr) ? 1 : m);
  if ((m || rr) && (j || !11))
  ret += solve(i - 1, j, (rr) ? 1 : k, (!11) ? 1 : 1, m);
return dp[i][j][k][1][m] = ret;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int q;
  cin >> q;
  while (q--)
   cin >> a >> b;
    a--, b++;
    memset(dp, -1, sizeof(dp));
    if (a == -1)
      cout << solve(60, 1, 0, 1, 0) << endl;
    else
      cout << solve(60, 0, 0, 0, 0) << endl;
  return 0;
// https://codeforces.com/contest/1245/problem/F
// https://codeforces.com/blog/entry/88064
// count the number of pairs (i, j) which (i \& j) == 0
```

## 2.3 broken profile

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define pi pair<int, int>
#define pii pair<int, pair<int, pi>>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 1001
#define mod 1000000007
int n;
vector<int> validmasks;
int dp[MAXN][1 << 4];</pre>
```

```
void init() // preprocess valid masks
 for (int mask = 0; mask < (1 << 7); mask++)
    int nxt_mask = 0, prev_mask = 0, valid = true;
    for (int k = 0; k < 7; k++)
      if (mask & (1 << k))
        if (k \le 3)
          int idx = k, idx2 = k;
          if (nxt_mask & (1 << idx) || prev_mask & (1 << idx2))</pre>
            valid = false;
          prev mask = prev mask | (1 << idx);</pre>
          nxt mask = nxt mask | (1 << idx2);
        else
          int idx = k - 4, idx2 = idx + 1;
          if (nxt_mask & (1 << idx) || nxt_mask & (1 << idx2))</pre>
            valid = false;
          nxt_mask = nxt_mask | (1 << idx);</pre>
          nxt_mask = nxt_mask | (1 << idx2);</pre>
    if (valid)
      validmasks.pb(mask);
int solve(int i, int j)
 if (i == n)
 return (j == ((1 << 4) - 1)) ? 1 : 0;
if (dp[i][j] != -1)
   return dp[i][j];
  int ret = 0;
 for (auto const &mask : validmasks)
    int nxt_mask = 0, prev_mask = j, valid = true;
    for (int k = 0; k < 7; k++)
      if (mask & (1 << k))
        if (k \le 3)
          int idx = k, idx2 = idx;
          if (prev_mask & (1 << idx) || nxt_mask & (1 << idx2))</pre>
            valid = false;
          prev_mask = prev_mask | (1 << idx);</pre>
          nxt mask = nxt mask | (1 << idx2);
        else
          int idx = k - 4, idx2 = idx + 1;
          if (nxt mask & (1 << idx) || nxt mask & (1 << idx2))</pre>
            valid = false;
          nxt_mask = nxt_mask | (1 << idx);</pre>
          nxt_mask = nxt_mask | (1 << idx2);</pre>
    if (valid && prev_mask == ((1 << 4) - 1))</pre>
      ret += solve(i + 1, nxt mask);
 return dp[i][j] = ret;
signed main()
 ios base::svnc with stdio(false);
 cin.tie(NULL);
 int q;
 cin >> q;
  init();
 for (int i = 1; i \le q; i++)
```

```
{
    cin >> n;
    memset(dp, -1, sizeof(dp));
    cout << i << " " << solve(0, (1 << 4) - 1) << endl;
}
return 0;
}
// broken profile dp
// if you can fully fill an area with some figures
// finding number of ways to fully fill an area with some figures
// finding a way to fill an area with minimum number of figures
// ...
// https://www.spoj.com/problems/GNY07H/
// We wish to tile a 4xN grid with rectangles 2xl (in either orientation)
// dp[i][mask]
// i denotes the current column
// mask denotes the situation of the previous column
// our mission is to fill all of the units of
// the previous column in a state [i][mask]</pre>
```

#### 2.4 cht

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1000005
#define mod 1000000007
struct line
  int m, b, p;
  line(int m, int b) : m(m), b(b) {}
  bool operator<(const line &o) const
   if (m != o.m)
     return m < o.m;</pre>
    return b < o.b;</pre>
  bool operator<(const int x) const { return p < x; }</pre>
  int eval(int x) const { return m * x + b; }
  int inter(const line &o) const
    int x = b - o.b, y = o.m - m;
   return (x / y) - ((x ^ y) < 0 && x % y);
};
struct cht
  int ptr:
  vector<line> a;
  cht() { ptr = 0;
  void add(line 1)
    while (1)
      if (a.size() >= 1 && a.back().m == 1.m && 1.b > a.back().b)
        a.pop_back();
      else if (a.size() >= 1 && a.back().m == 1.m && 1.b <= a.back().b)</pre>
        break;
```

```
else if (a.size() \ge 2 \&\& a.back().inter(l) \ge a[a.size() - 2].inter(a.
          back()))
        a.pop_back();
      else
        a.pb(1);
        break;
  int get(int x)
   if (!a.size())
      return -inf;
    while (ptr + 1 < a.size() && a[ptr].eval(x) <= a[ptr + 1].eval(x))
    return a[ptr].eval(x);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 return 0;
// queries ordenadas em ordem decrescente
// linhas ordenadas em ordem decrescente
```

## 2.5 Digitdp

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define fir first
#define sec second
#define MAXN 2001
#define mod 1000000007
int dp[20][20 * 9][2]; // a,b <= 10^18</pre>
vector<int> dig;
int solve(int i, int j, int k)
  if (i == dig.size())
   return (k) ? dp[i][j][k] = j : dp[i][j][k] = 0;
  if (dp[i][j][k] != -1)
   return dp[i][j][k];
  int sum = 0;
  if (k)
    for (int f = 0; f <= 9; f++)</pre>
      sum += solve(i + 1, j + f, k);
  if (!k)
    for (int f = 0; f <= dig[i]; f++)</pre>
      sum += solve(i + 1, j + f, (dig[i] != f) ? 1 : 0);
  return dp[i][j][k] = sum;
void get_digits(int n)
  dig.clear();
  while (n)
   dig.pb(n % 10);
   n = n / 10;
  reverse(dig.begin(), dig.end());
signed main()
```

```
ios_base::sync_with_stdio(false);
cin.tie(NULL);
int a, b;
cin >> a >> b;
get_digits(a);
memset(dp, -1, sizeof(dp));
int aa = solve(0, 0, 0);
get_digits(b + 1);
memset(dp, -1, sizeof(dp));
int bb = solve(0, 0, 0);
cout << bb - aa << endl;
return 0;</pre>
```

### 2.6 divideandconquer

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 1000000007
int s[8005];
int dp[3005][8005];
int cost(int 1, int r)
  return (s[r + 1] - s[1]) * (r - 1 + 1);
void compute(int 1, int r, int opt1, int optr, int i)
  if (1 > r)
   return;
  int mid = (1 + r) >> 1;
  pair<int, int> ans = \{1e18, -1\}; // dp, k
  for (int q = optl; q <= min(mid, optr); q++)</pre>
      ans = min(ans, {dp[i - 1][q - 1] + cost(q, mid), q});
    else
      ans = min(ans, {cost(q, mid), q});
  dp[i][mid] = ans.fir;
  compute(1, mid - 1, optl, ans.sec, i);
  compute(mid + 1, r, ans.sec, optr, i);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, g;
  cin >> n >> q;
  for (int i = 1; i <= n; i++)</pre>
    cin >> s[i];
    s[i] += s[i - 1];
  for (int i = 0; i \le g; i++)
    for (int j = 0; j <= n; j++)</pre>
      dp[i][j] = 1e18;
```

```
for (int i = 1; i <= g; i++)
    compute(0, n - 1, 0, n - 1, i);
cout << dp[g][n - 1] << endl;
return 0;
}
// https://codeforces.com/gym/103536/problem/A
// https://codeforces.com/contest/321/problem/E

// otimizacao de dp usando divide and conquer
// para dps do tipo:
// dp[i][j] = min(dp[i - 1][k] + c(k, j)), para algum k <= j
// considerando opt(i, j) o menor valor de k que minimiza dp[i][j]
// podemos calcular opt(i, j) usando divide and conquer
// isso diminuiria a complexidade para O(k * n * log(n))</pre>
```

## 2.7 dynamic cht

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pf push_front
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1000005
#define mod 1000000007
struct line
 mutable int m, b, p;
  bool operator<(const line &o) const
    if (m != o.m)
     return m < o.m;</pre>
    return b < o.b;</pre>
  bool operator<(const int x) const { return p < x; }</pre>
  int eval(int x) const { return m * x + b; }
  int inter(const line &o) const
    int x = b - o.b, y = o.m - m;
    return (x / y) - ((x ^ y) < 0 && x % y);
struct cht
  int INF = 1e18;
  multiset<line, less<>> 1;
  void add(int m, int b)
    auto y = 1.insert({m, b, INF});
    auto \bar{z} = next(y);
    if (z != 1.end() && y->m == z->m)
      1.erase(y);
      return;
    if (y != l.begin())
      auto x = prev(y);
      if (x->m == y->m)
        x = 1.erase(x);
```

```
while (1)
      if (z == 1.end())
        y->p = INF;
       break;
      y->p = y->inter(*z);
      if (y->p < z->p)
       break;
      else
        z = 1.erase(z);
    if (y == l.begin())
      return;
    z = y;
    auto x = --y;
    while (1)
      int ninter = x->inter(*z);
      if (ninter <= x->p)
       x->p = ninter;
      else
        1.erase(z);
       break:
      if (x == 1.begin())
       break;
      y = x;
      if (x->p < y->p)
       break;
      else
        1.erase(y);
  int get(int x)
   if (1.empty())
      return 0;
   return 1.lower_bound(x)->eval(x);
};
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 return 0;
// https://github.com/pauloamed/Training/blob/master/PD/cht.cpp
// https://github.com/brunomaletta/Biblioteca/blob/master/Codigo/DP/CHT-Dinamico
    .cpp
// cht dinamico
// dado uma coordenada x
// e um conjunto com varias equacoes lineares da forma: y = mx + c
// retorna o maior valor de y entre as equações do conjunto
// para o menor valor, multiplicar m e c de cada equacao por -1
// e multiplicar o resultado da query por -1
// problemas iniciais:
// https://atcoder.jp/contests/dp/tasks/dp_z
// https://codeforces.com/contest/1083/problem/E
```

# 2.8 exchange arguments

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
```

```
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1001
#define mod 1000000009
const int inf = 1e18;
int n;
vector<pi> v:
int dp[MAXN][MAXN];
bool vis[MAXN][MAXN];
int solve(int i, int j)
  if (j == 0)
   return inf;
  if (i == n)
   return -inf;
  if (vis[i][j])
   return dp[i][j];
  int ans = -inf;
  ans = max(ans, solve(i + 1, j));
  int ot = min(v[i].sec, solve(i + 1, j - 1) - v[i].fir);
  ans = max(ans, ot);
  vis[i][j] = 1;
  return dp[i][j] = ans;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n;
  v.resize(n);
  for (int i = 0; i < n; i++)</pre>
    cin >> v[i].fir >> v[i].sec;
    v[i].sec -= v[i].fir;
  auto cmp = [&] (pi a, pi b)
    return (a.sec - b.fir) < (b.sec - a.fir);</pre>
  sort(v.begin(), v.end(), cmp);
  memset(dp, -1, sizeof(dp));
  int ans = 0:
  for (int i = n; i >= 0; i--)
    if (solve(0, i) >= 0)
      ans = i:
      break;
  cout << ans << endl;
  return 0;
// problema:
// existem n caixas, cada uma tem um peso w[i] e uma resistencia r[i]
// voce deve escolher um subset de caixas e empilhar na ordem que vc quiser
// tal que: a soma dos pesos de todas as caixas acima de uma caixa seja menor ou
      iqual a resistencia dessa caixa
// dp[i][j] - estou na caixa i e quero escolher mais j caixas para botar na
    pilha
// qual a maior resistencia restante que eu posso obter escolhendo essas j
    caixas
// a grande sacada pra achar a ordenacao otima antes da dp:
// para duas caixas a e b
// quando vai ser stonks botar a antes de b?
```

```
// r[a] - w[b] > r[b] - w[a]
// pois a resistencia reestante vai ser maior

// pra demais problemas de exchange argument, essa ideia pode se aplicar
// do tipo, ver o jeito otimo de resolver pro n = 2
// e fazer a ordenacao baseada nisso
```

### 2.9 expected value

```
//https://atcoder.jp/contests/dp/tasks/dp_j
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define mp make_pair
#define pi pair<int, int>
#define pii pair<pi, int>
#define pci pair<char, int>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 301
#define mod 1000000007
int n;
vector<int> v;
vector<int> cnt(3);
double dp[MAXN][MAXN][MAXN];
double solve(int i, int j, int k)
  if (!i && !i && !k)
    return dp[i][j][k] = 0;
  if (dp[i][j][k] != −1)
   return dp[i][j][k];
  It is well-known from statistics that for the geometric distribution
  (counting number of trials before a success, where each independent trial is
      probability p)
  the expected value is i / p
  double p = ((double)(i + j + k) / n);
  double ret = 1 / p; // expected number of trials before a success
  if (i)
    double prob = (double)i / (i + j + k); // probabilidade de ser um prato com
        um sushi
    ret += (solve(i - 1, j, k) * prob);
  if (j)
    double prob = (double) j / (i + j + k); // probabilidade de ser um prato com
        dois sushis
    ret += (solve(i + 1, j - 1, k) * prob);
  if (k)
    double prob = (double)k / (i + j + k); // probabilidade de ser um prato com
        tres sushis
    ret += (solve(i, j + 1, k - 1) * prob);
  return dp[i][j][k] = ret;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
```

```
cin >> n;
v.resize(n);
for (int i = 0; i < n; i++)
    cin >> v[i], cnt[v[i] - 1]++;
for (int i = 0; i < MAXN; i++)
    for (int j = 0; j < MAXN; j++)
        for (int k = 0; k < MAXN; k++)
            dp[i][j][k] = -1;
cout << setprecision(15) << solve(cnt[0], cnt[1], cnt[2]) << endl;
return 0;
}
```

#### 2.10 inclusion exclusion

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first.
#define sec second
#define MAXN 300005
#define mod 1000000007
struct modint
  int val;
 modint(int v = 0) { val = v % mod; }
  int pow(int y)
    modint x = val;
    modint z = 1:
    while (y)
      if (y & 1)
       z *= x;
      x *= x;
     y >>= 1;
    return z.val;
  int inv() { return pow(mod - 2); }
  void operator=(int o) { val = o % mod; }
  void operator=(modint o) { val = o.val % mod; }
  void operator+=(modint o) { *this = *this + o;
  void operator = (modint o) { *this = *this - o;
  void operator*=(modint o) {
                              *this = *this * 0;
                              *this = *this / o;
  void operator/=(modint o) {
  bool operator==(modint o) { return val == o.val;
  bool operator!=(modint o) { return val != o.val; }
  int operator*(modint o) { return ((val * o.val) % mod); }
  int operator/(modint o) { return (val * o.inv()) % mod;
  int operator+(modint o) { return (val + o.val) % mod; }
  int operator-(modint o) { return (val - o.val + mod) % mod; }
modint f[MAXN];
modint invfat[MAXN];
void calc()
  f[0] = 1;
  invfat[0] = 1;
  for (int i = 1; i < MAXN; i++)</pre>
    f[i] = f[i - 1] * i;
    invfat[i] = f[i].inv();
```

```
modint ncr(int n, int k) // combinacao
  modint ans = f[n] * invfat[k];
  ans \star = invfat[n - k];
  return ans;
modint arr(int n, int k) // arranjo
  modint ans = f[n] * invfat[n - k];
  return ans;
int h, w, n;
pi v[3005];
modint m1 = 1000000006; // -1
modint dp[3005][3005][2];
bool vis[3005][3005][2];
modint solve(int i, int j, int par)
  if (i > n)
    modint ans = ncr((h - v[j].fir) + (w - v[j].sec), h - v[j].fir);
    if (!par)
      return ans;
    return ans * m1;
  if (vis[i][j][par])
    return dp[i][j][par];
  modint ans = solve(i + 1, j, par);
  if (v[i].sec >= v[j].sec)
    modint nxt = solve(i + 1, i, (par + 1) % 2);
    modint curr = ncr((v[i].fir - v[j].fir) + (v[i].sec - v[j].sec), v[i].fir -
         v[j].fir);
    curr *= nxt;
   ans = ans + curr;
  vis[i][j][par] = 1;
  dp[i][j][par] = ans;
  return ans;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  calc();
  cin >> h >> w >> n;
  // tenho que fazer essa inclusao exclusao numa dp
  // for (int i = 0; i < n; i++)
  // cin >> v[i].fir >> v[i].sec;
  // sort (v, v + n);
  // modint ans = 0;
  // for (int mask = 0; mask < (1 << n); mask++)
       pi prv = \{1, 1\};
       for (int j = 0; j < n; j++)
         if (mask & (1 << j))
           if (v[j].sec < prv.sec)
             at = 0;
             break;
           at \star = ncr((v[j].fir - prv.fir) + (v[j].sec - prv.sec), v[j].fir - prv
       .fir);
           prv = v[j];
      at *= ncr((h - prv.fir) + (w - prv.sec), h - prv.fir);
cout << mask << " " << at.val << endl;
       ans = (__builtin_popcount(mask) % 2) ? ans - at : ans + at;
```

```
// cout << ans.val << endl;
 v[0] = \{1, 1\};
  for (int i = 1; i <= n; i++)</pre>
   cin >> v[i] fir >> v[i] sec;
  sort(v, v + n + 1);
 cout << solve(1, 0, 0).val << endl;
// https://atcoder.jp/contests/dp/tasks/dp_y
// dado um grid com h linhas e w colunas
// h, w <= 10^5
// existem n <= 3000 posicoes com wall
// quantos caminhos tem do square (1, 1) ate o (h, w)
// com isso, podemos fzr inclusao exclusao
// calculando quantos caminhos passam por um conjunto de walls
// mas nao podemos fazer 2^30
// maaas, podemos fazer uma dp pra computar essa inclusao exclusao
// e dale
```

### 2.11 Knapsack

```
//O problema mais classico de Programacao Dinamica talvez seja o Knapsack.
//De maneira geral, um ladrao ira roubar uma casa com uma mochila
//que suporta um peso s. Ele ve n objetos na casa e sabe estimar o peso pi
    o valor vi
//de cada objeto i. Com essas informacoes, qual o maior valor que o ladrao pode
    roubar sem rasgar sua mochila?
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define in insert
#define pi pair<int, int>
#define pii pair<int, pi>
#define mp make_pair
#define fir first
#define sec second
#define MAXN 1001
#define INF 1000000000
int n, 1;
int value[MAXN];
int peso[MAXN];
int dp[MAXN][MAXN];
int knapsack(int i, int limit)
  if (dp[i][limit] >= 0) // se ja foi calculado
    return dp[i][limit];
  if (i == n or !limit) // se chegou no fim do array ou chegou no limite
    return dp[i][limit] = 0;
  int nao_coloca = knapsack(i + 1, limit); // recursivamente pra caso eu nao
       coloque o objeto i
  if (peso[i] <= limit) // se eu consigo botar o objeto i</pre>
    int coloca = value[i] + knapsack(i + 1, limit - peso[i]);
   return dp[i][limit] = max(coloca, nao_coloca);
  return dp[i][limit] = nao_coloca;
signed main()
  cin >> 1 >> n;
  for (int i = 0; i < n; i++)
```

```
cin >> peso[i] >> value[i];
}
memset(dp, -1, sizeof(dp));
cout << knapsack(0, 1) << endl;
return 0;</pre>
```

### 2.12 largest-sum-contiguous-subarray

```
// dada uma sequencia s qual a maior soma que podemos obter escolhendo um
     subconjunto de termos adjacentes de s
// nesse caso o temos apenas duas opcoes
// nao usar o elemento v[i]
// usamos, adicionando a maior soma possivel que antes dele
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200001
#define mod 1000000007
int kadane(vector<int> v)
  int n = v.size(), ans = 0, max_here = 0;
  for (int i = 0; i < n; i++)
    max_here += v[i];
    if (ans < max_here)</pre>
      ans = max_here;
    if (max_here < 0)</pre>
     max_here = 0;
  return ans:
int kadane_circular(vector<int> v)
  int n = v.size(), max_kadane = kadane(v);
  int max_wrap = 0, i;
  for (i = 0; i < n; i++)
    max_wrap += v[i];
   v[i] = -v[i];
  max_wrap += kadane(v);
  return max(max_wrap, max_kadane);
signed main()
  int n;
  cin >> n;
  vector<int> v(n);
  for (int i = 0; i < n; i++)
  cout << kadane_circular(v) << endl;</pre>
  return 0;
```

#### 2.13 largest square

```
#include <bits/stdc++.h>
using namespace std;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define double long double
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 1001
#define mod 1000000007
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  int v[n][n];
  int dp[n][n];
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++)
      cin >> v[i][j];
  int ans = 0;
  for (int i = 0; i < n; i++)
     for (int j = 0; j < n; j++)
       dp[i][j] = v[i][j];
       \begin{array}{lll} \textbf{if} & (i & \& \& j & \& \& & dp[i][j]) \\ & dp[i][j] &= min(\{dp[i][j-1], \ dp[i-1][j], \ dp[i-1][j-1]\}) \ + \ 1; \\ \end{array} 
       ans = max(ans, dp[i][j]);
  cout << ans * ans << endl;</pre>
  return 0;
```

#### 2.14 lcs

dp[i][j] = 0;

```
//Dadas duas sequencias s1 e s2, uma de tamanho n e outra de tamanho m, qual a
    maior subsequencia comum as duas?
// uma subsequencia de s e um subconjunto dos elementos de s na mesma ordem em
    que apareciam antes.
// isto significa que {1, 3, 5} e uma subsequencia de {1, 2, 3, 4, 5}, mesmo 1
    nao estando do lado do 3.
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define in insert
#define pi pair<int, int>
#define pii pair<int, pi>
#define mp make_pair
#define fir first
#define sec second
#define MAXN 1001
#define INF 1000000000
int v1[MAXN];
int v2[MAXN];
int dp[MAXN][MAXN];
void lcs(int m, int n)
  for (int i = 0; i \le m; i++)
    for (int j = 0; j \le n; j++)
      if (i == 0 || j == 0) //se uma das sequencias for vazia
```

```
else if (v1[i-1] == v2[j-1]) // se eh igual, adiciono a lcs e subtraio
           dos dois
        dp[i][j] = dp[i - 1][j - 1] + 1;
      else
        dp[i][j] = max(dp[i-1][j], dp[i][j-1]); // se nao retorno o maximo
            entre tirar um dos dois caras
 cout << dp[m][n] << endl;</pre>
signed main()
 int n, m;
 cin >> n >> m;
  for (int i = 0; i < n; i++)
   cin >> v1[i];
  for (int i = 0; i < m; i++)
   cin >> v2[i];
  lcs(n, m);
 return 0;
```

#### 2.15 lichao

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define pci pair<char, int>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
const int inf = 1e18;
struct line
  int a, b, id;
 int ch[2];
  line()
   a = 0, b = \inf, id = -1;
   ch[0] = -1, ch[1] = -1;
  line(int aa, int bb, int i)
    a = aa, b = bb, id = i;
   ch[0] = -1, ch[1] = -1;
  int f(int x) { return a * x + b; }
struct save
  int a, b, id, new_id, p;
struct lichao
  int lo, hi, curr;
  vector<line> t;
  stack<save> st; // se nao precisar de rollback, pode tirar a stack (pra nao
       usar tanta memoria)
  lichao(int ll, int rr)
```

```
lo = ll, hi = rr;
    t.emplace_back();
  int child(int p, int d)
    if (t[p].ch[d] == -1)
      t[p].ch[d] = t.size();
      t.emplace_back();
    return t[p].ch[d];
 bool cmp(line a, line b, int x)
    if (a.f(x) != b.f(x)) // menor valor em x
     return a.f(x) < b.f(x);
    return a.id > b.id; // desempata pelo maior id
 void add(int 1, int r, line s, int p)
    int mid = (1 + r) >> 1;
   bool fl = cmp(s, t[p], 1);
    bool fm = cmp(s, t[p], mid);
    bool fr = cmp(s, t[p], r);
      st.push({t[p].a, t[p].b, t[p].id, curr, p});
      swap(t[p], s);
      swap(t[p].ch, s.ch);
    if (s.b == inf)
     return;
    if (fl != fm)
     add(1, mid - 1, s, child(p, 0));
    else if (fr != fm)
      add(mid + 1, r, s, child(p, 1));
  pi query(int 1, int r, int x, int p)
    int mid = (1 + r) >> 1;
   pi ans = {t[p].f(x), -t[p].id}; // como eu quero o maior id, basta negar e
        pegar o menor
    if (ans.fir == inf)
     return ans;
    if (x < mid)
     return min(ans, query(1, mid - 1, x, child(p, 0)));
    return min(ans, query(mid + 1, r, x, child(p, 1)));
 void add(line s)
    curr = s.id:
    add(lo, hi, s, 0);
 pi qry(int x)
    return query(lo, hi, x, 0);
  void rollback(int id)
    while (!st.empty() && st.top().new_id == id)
      int p = st.top().p;
     t[p].a = st.top().a;
      t[p].b = st.top().b;
      t[p].id = st.top().id;
     st.pop();
signed main()
  lichao 1t(0, 1e9 + 2);
  lt.add(line(3, 2, 0));
 lt.add(line(5, -6, 1));
 cout << lt.qry(10).fir << " " << -lt.qry(10).sec << endl;</pre>
// li-chao tree
// dado uma coordenada x
// e um conjunto com varias equacoes lineares da forma: y = ax + b
```

```
// retorna o menor valor de y entre as equacoes do conjunto
// O(log(hi - lo))
// no qual:
// lo -> menor valor possivel de um x que vai ser passado pra uma query
// hi -> maior valor possivel de um x que vai ser passado pra uma query
```

#### 2.16 lis

```
// dada uma sequencia s qualquer, descobrir o tamanho da maior subsequencia
     crescente de s
   uma subsequencia de s e qualquer subconjunto de elementos de s.
// Para cada novo numero, voce tem duas operacoes possiveis:
// 1 - Colocar o novo numero no topo de uma pilha se ele nao superar o que ja
     esta em seu topo;
// ou
// 2 - Criar uma nova pilha a direita de todas as outras e colocar o novo numero
      la.
// ao final do processo a nossa pilha tera os elementos da lis.
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push back
#define in insert
#define pi pair<int, int>
#define pd pair<double, int>
#define pib pair<pi, bool>
#define mp make pair
#define fir first
#define sec second
#define MAXN 200001
#define MAXL 1000001
#define mod 1000000007
vector<int> v;
int lis()
  vector<int> q;
  for (int i = 0; i < v.size(); i++)</pre>
    vector<int>::iterator it = lower_bound(g.begin(), g.end(), v[i]);
    if (it == q.end())
      q.pb(v[i]);
    else
      *it = v[i];
  for (int i = 0; i < q.size(); i++)
  cout << q[i] << " ";</pre>
  cout << endl;
  return q.size();
signed main()
  int n:
  cin >> n;
  v.resize(n);
  for (int i = 0; i < n; i++)
   cin >> v[i];
  cout << lis() << endl;</pre>
  return 0;
```

## 2.17 max matrix path

```
#include <bits/stdc++.h>
using namespace std;

#define PI acos(-1)
#define pb push_back
#define int long long int
#define mp make_pair
```

```
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 301
#define MAXL 20
#define mod 1000000007
#define INF 100000001
int grid[MAXN][MAXN];
int dp[MAXN][MAXN];
int solve(int i, int j)
 if (i == n - 1 \&\& j == n - 1)
   return grid[i][j];
 if (dp[i][j] != -1)
   return dp[i][j];
  if (i + 1 < n &  i + 1 < n)
   return dp[i][j] = grid[i][j] + max(solve(i + 1, j), solve(i, j + 1));
 if (i + 1 < n)
   return dp[i][j] = grid[i][j] + solve(i + 1, j);
  if (j + 1 < n)
   return dp[i][j] = grid[i][j] + solve(i, j + 1);
signed main()
  cin >> n;
 for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n; j++)
     cin >> grid[i][j];
 memset(dp, -1, sizeof(dp));
 cout << solve(0, 0) << endl;</pre>
 return 0;
```

## 2.18 optimization with fft

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
// #define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 998244353
struct modint
  modint(int v = 0) { val = ((v % mod) + mod) % mod; }
  int pow(int y)
    modint x = val;
    modint z = 1:
    while (y)
     if (y & 1)
       z *= x;
      x *= x;
      y >>= 1;
    return z.val;
```

```
int inv() { return pow(mod - 2); }
  void operator=(int o) { val = o % mod; }
  void operator=(modint o) { val = o.val % mod; }
  void operator+= (modint o) { *this = *this + o; ]
  void operator = (modint o) { *this = *this - o;
  void operator*=(modint o) { *this = *this * o;
  void operator/=(modint o) { *this = *this / o; ]
  bool operator== (modint o) { return val == o.val; }
  bool operator!=(modint o) { return val != o.val; }
  int operator*(modint o) { return ((val * 111 * o.val) % mod); }
  int operator/(modint o) { return (val * 111 * o.inv()) % mod; }
  int operator+(modint o) { return (val + o.val) % mod; }
  int operator-(modint o) { return (val - o.val + mod) % mod; }
namespace fft
  int n;
  int root = -31;
  int root_1 = 128805723;
  int pw = builtin ctz(mod - 1);
  int root_pw = (1 << pw);</pre>
  void ntt(vector<modint> &a, bool invert)
    int n = a.size();
    for (int i = 1, j = 0; i < n; i++)
      int bit = n >> 1;
      for (; j & bit; bit >>= 1)
        j ^= bit;
^= bit;
      if (i < j)
        swap(a[i], a[j]);
    for (int len = 2; len <= n; len <<= 1)</pre>
      modint wlen = (invert) ? root_1 : root;
      for (int i = len; i < root_pw; i <<= 1)</pre>
        wlen *= wlen;
      for (int i = 0; i < n; i += len)</pre>
        modint w = 1;
        for (int j = 0; j < len / 2; j++)
          modint u = a[i + j];
          modint v = a[i + j + len / 2] * w;
          a[i + j] = u + v;
          a[i + j + len / 2] = u - v;
          w \star = wlen;
    if (invert)
      modint n_1 = modint(n).inv();
      for (int i = 0; i < a.size(); i++)</pre>
        a[i] *= n_1;
  vector<modint> mul(vector<modint> &a, vector<modint> &b)
    n = 1:
    while (n < 2 * max(a.size(), b.size()))
     n <<= 1;
    a.resize(n);
    b.resize(n);
    ntt(a, false);
    ntt(b, false);
    for (int i = 0; i < n; i++)
     a[i] *= b[i];
    ntt(a, true);
    return a;
// int dp[3005][3005 + 3005];
signed main()
```

```
ios_base::sync_with_stdio(false);
 cin.tie(NULL);
  // dp naive:
  // auto get = [&] (int i, int j)
     return dp[i][j + 3005];
  // dp[1][3005] = 1;
  // for (int i = 2; i < 3005; i++)
       for (int j = -i; j \le i; j++)
         dp[i][j + 3005] = get(i - 1, j - 1);
         dp[i][j + 3005] += get(i - 1, j + 1);
        dp[i][j + 3005] += (i - 2) * get(i - 1, j);
dp[i][j + 3005] %= mod;
  // for (int i = 0; i \le 5; i++)
      for (int j = -i; j \le i; j++)
        cout << dp[i][j + 3005] << ", ";
      cout << endl;
  int n, k;
 cin >> n >> k;
  queue<vector<modint>> vecs;
  vecs.push({0, 1, 0}); // dp[1][-1], dp[1][0], dp[1][1]
 for (int i = 2; i <= n; i++)
    // dp[i] = fft::conv(dp[i-1], \{1, i-2, 1\})
   vecs.push(\{1, i - 2, 1\});
  // faz as convolucoes pegando smp os dois de menor tamanho
  while (vecs.size() > 1)
    vector<modint> v = vecs.front();
   vecs.pop();
    vector<modint> v2 = vecs.front();
   vecs.pop();
   vector<modint> v3 = fft::mul(v, v2);
   vecs.push(v3);
  // quero o valor de dp[n][k], o k dado na entrada pode ser negativo
 cout << vecs.back()[k + n].val << endl;</pre>
// https://atcoder.jp/contests/abc385/tasks/abc385_g
// a dp eh tipo:
// dp[i][j] = dp[i-1][j] * (i-2) + dp[i-1][j-1] + dp[i-1][j+1]
// considerando um certo i, j ta no intervalo [-i,\ i] // mas n eh na casa de 10^5
```

## 2.19 permutations

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 300005
#define mod 1000000007
int n;
string s;
```

```
int dp[3005][3005];
int sum[3005][3005];
int solve(int i, int j);
void solve2(int i)
  if (sum[i][0] != -1)
    return;
  sum[i][0] = 0;
  for (int j = 0; j <= i; j++)
  sum[i][j + 1] = (sum[i][j] + solve(i, j)) % mod;</pre>
int get(int i, int l, int r)
  return (sum[i][r + 1] - sum[i][l] + mod) % mod;
int solve(int i, int j)
  if (i == s.size())
    return 1;
  if (dp[i][j] != -1)
    return dp[i][j];
  int nums = i + 1, ans = 0;
  solve2(i + 1);
  if (s[i] == '>')
    ans = get(i + 1, 0, min(nums, j));
    ans = get(i + 1, j + 1, nums);
  return dp[i][j] = ans;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> s;
  memset(dp, -1, sizeof(dp));
  memset(sum, -1, sizeof(sum));
  cout << solve(0, 0) << endl;</pre>
// dado uma string s de tamanho n - 1 e um inteiro n
// conte quantas permutacoes de tamanho n satisfazem as seguintes restricoes // se s[i] == '>', logo p[i] > p[i+1] // ou se s[i] == '<', logo p[i] < p[i+1]
// se eu calculei dp[i][j]
// significa que eu preenchei o prefixo de tamanho i com uma permutacao
// e o ultimo elemento que eu coloquei foi o j (note que j <= i)
// se o proximo elemento precisa ser > j
// entao dp[i][j] = dp[i + 1][j + 1] + ... dp[i + 1][i + 1]
// caso contrario
// note que essa dp eh correta pq
// se eu tou botando um elemento x < n na transicao
// pro prefixo que eu tinha resolvido, posso incrementar todo mundo que for >= x
      que ta tudo certo
// tipo se eu tinha construido a permutacao 2 1 4 3
// tou calculando dp[4][3] entao pq o 3 foi o ultimo elemento
// mas na transicao tou considerando dp[4][3] = ... dp[5][2] ... // a permutacao do dp[5][2] seria a 3 1 5 4 2 que continua sendo valida
// soh codar com soma de prefixo pra otimizar a dp e dale nessa questao
```

## $2.20 \quad sos dp$

```
#define int long long int
#define endl '\n'
#define pb push back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100001
#define mod 1000000007
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 // exemplos de sos dp para calcular f[x] para cada mask x
  // a[x] eh o valor de uma funcao a para uma mask x
  // complexidade: O(M * 2^M), M = numero de bits
  // Exemplo 1:
  // nesse caso, f[x] eh a funcao que soma:
  // todos os a[i], tal que, (x \& i) == i)
  // isso eh, i eh uma "mask filha" de x
  // pois todos os bits de i estao setados em x
  for (int mask = 0; mask < (1 << m); mask++)</pre>
    f[mask] = a[mask];
  for (int i = 0; i < m; ++i)
    for (int mask = 0; mask < (1 << m); mask++)
     if (mask & (1 << i))
        f[mask] += f[mask ^ (1 << i)];
  // Exemplo 2:
  // nesse caso, f[x] eh a funcao que soma:
  // todos os a[i], tal que, (x \& i) == x)
  // isso eh, i eh uma "mask pai" de x
  // pois todos os bits de x estao setados em i
 for (int mask = 0; mask < (1 << m); mask++)</pre>
    f[mask] = a[mask];
  for (int i = 0; i < m; ++i)
    for (int mask = 0; mask < (1 << m); mask++)</pre>
     if (!(mask & (1 << i)))</pre>
        f[mask] += f[mask ^ (1 << i)];
 return 0;
// https://codeforces.com/blog/entry/45223
```

#### 2.21 steiner tree

```
#include <bits/stdc++.h>
using namespace std;

#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 4007
#define mod 998244353

const int inf = le18;
int n, m, k, sz;
```

```
int dx[] = \{-1, 1, 0, 0\};
int dy[] = \{0, 0, 1, -1\};
int a[101][101];
vector<int> q[201];
int w[201][201];
int dist[201][201];
int par_floyd[201][201];
int dp[201][1 << 8];
char anss[201][201];
vector<pi> par[201][1 << 8];</pre>
int pos[10];
int steiner()
  // flovd warshall
  for (int i = 0; i < sz; i++)
    for (int j = 0; j < sz; j++)
      if (i == i)
       dist[i][j] = 0;
      else
        dist[i][j] = w[i][j];
        par_floyd[i][j] = j;
  for (int k = 0; k < sz; k++)
    for (int i = 0; i < sz; i++)
      for (int j = 0; j < sz; j++)
        if (dist[i][j] > dist[i][k] + dist[k][j])
          dist[i][j] = dist[i][k] + dist[k][j];
          par_floyd[i][j] = par_floyd[i][k];
  for (int i = 0; i < sz; i++)
    for (int j = 0; j < (1 << k); j++)
      dp[i][j] = inf;
  for (int i = 0; i < k; i++)
    for (int j = 0; j < sz; j++)
      dp[j][1 << i] = dist[pos[i]][j];
      par[j][1 << i] = {\{pos[i], 0\}\};}
  for (int mask = 2; mask < (1 << k); mask++)
    for (int i = 0; i < sz; i++)
      for (int mask2 = mask; mask2 > 0; mask2 = (mask2 - 1) & mask)
        int mask3 = mask ^ mask2;
        if (dp[i][mask] > dp[i][mask2] + dp[i][mask3])
          dp[i][mask] = dp[i][mask2] + dp[i][mask3];
          par[i][mask] = {\{i, mask2\}, \{i, mask3\}\};}
      for (int j = 0; j < sz; j++)
        if (dp[j][mask] > dp[i][mask] + dist[i][j])
          dp[j][mask] = dp[i][mask] + dist[i][j];
          par[j][mask] = {{i, mask}};
```

```
// preciso somar a[i / m][i % m] pq tou usando a[i][j]
  // como peso de uma aresta (i, j) \rightarrow (x, y)
  // mas eh especifico para esse problema do garden
 int ans = inf, best = -1;
 for (int i = 0; i < sz; i++)</pre>
   int curr = a[i / m][i % m] + dp[i][(1 << k) - 1];</pre>
    if (curr < ans)</pre>
      ans = curr;
     best = i;
  // recuperar resposta
  queue<pi> q;
  q.push({best, (1 << k) - 1});
 while (!q.empty())
    auto [i, mask] = q.front();
    q.pop();
    anss[i / m][i % m] = 'X';
    for (auto [j, mask2] : par[i][mask])
      // marcar o caminho de j para i feito pelo floyd warshall
      int st = j, en = i;
      while (st != en)
        st = par_floyd[st][en];
        anss[st / m][st % m] = 'X';
      q.push({j, mask2});
 return ans;
signed main()
 ios::sync_with_stdio(false);
 cin.tie(0);
 cin >> n >> m >> k;
 for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < m; j++)
     cin >> a[i][j];
      anss[i][j] = '.';
  sz = (n * m);
 for (int i = 0; i < sz; i++)</pre>
    for (int j = 0; j < sz; j++)
      w[i][j] = inf;
  // montando o grafo
  // (i, j) \rightarrow (x, y) com peso a[i][j]
 for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < m; j++)
      for (int d = 0; d < 4; d++)
        int x = i + dx[d];
        int y = j + dy[d];
        if (x >= 0 \&\& x < n \&\& y >= 0 \&\& y < m)
          w[(i * m) + j][(x * m) + y] = a[i][j];
  // posicoes importantes
 for (int i = 0; i < k; i++)
    int x, y;
    cin >> \bar{x} >> y;
   pos[i] = (x * m) + y;
```

```
cout << steiner() << endl;</pre>
 for (int i = 0; i < n; i++)
   for (int j = 0; j < m; j++)
     cout << anss[i][j];</pre>
   cout << endl;
 return 0;
// https://codeforces.com/problemset/problem/152/E
// dada uma matriz representando um jardim
// a posicao (i, j) tem a[i][j] flores
// com isso, quero cobrir algumas posicoes com concreto
// quando cobrimos uma posicao (i, j) com concreto, "matamos" as a[i][j] flores
    daquela posicao
// existem k <= 7 posicoes importantes, que devem ser cobertas com concreto
// alem disso, posso cobrir qualquer outra posicao com concreto
// alem disso, para duas posicoes a e b que sao importantes, deve existir um
    caminho
// de a ate b passando somente por posicoes cobertas por concreto.
// quero minimizar o numero de flores mortas, satisfazendo essas condicoes
// o que queremos nesse caso, eh uma steiner tree
// dado um grafo, com peso nas arestas
// e um subconjunto de vertices
// queremos achar uma arvore de menor peso que contenha todos os vertices do
// mas essa arvore pode conter tambem outros vertices que nao estao no
    subconjunto
// minimizando o peso total das arestas da arvore
// um outro problema de steiner tree: https://codeforces.com/gym/101908/problem/
```

## 2.22 subsequences string

```
#include <bits/stdc++.h>
using namespace std;
#define PI acos(-1)
#define int long long int
#define pb push_back
#define mp make_pair
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100
#define MAXL 20
#define mod 998244353
void count(string a, string b)
  int m = a.size();
  int n = b.size();
  int dp[m + 1][n + 1] = \{\{0\}\};
  for (int i = 0; i <= n; ++i)</pre>
    dp[0][i] = 0;
  for (int i = 0; i <= m; ++i)</pre>
    dp[i][0] = 1;
  for (int i = 1; i <= m; i++)</pre>
    for (int j = 1; j \le n; j++)
      if (a[i - 1] == b[j - 1])
        dp[i][j] = dp[i - 1][j - 1] + dp[i - 1][j];
        dp[i][j] = dp[i - 1][j];
    }
  cout << dp[m][n] << endl;</pre>
signed main()
  string a, b;
```

```
cin >> a >> b;
count(a, b);
return 0;
```

#### 2.23 subset sum

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
// subset sum com bitset de tamanho variado
// usada no https://codeforces.com/contest/1856/problem/E2
// with n <= 10^{^{\circ}}6
template <int len = 1>
int subset_sum(int n, int h)
  if (n >= len)
    return subset_sum<std::min(len * 2, (int)MAXN)>(n, h);
 bitset<len> dp;
  dp[0] = 1;
  for (auto const &x : w)
   dp = dp \mid (dp << x);
  return dp._Find_next(max(011, h - 1)); // retorna o proximo bit setado apos a
       posicao passada como parametro
int solve(vector<int> &w, int tot, int h)
  // tot -> soma de todos os elementos de w
  // h -> valor desejado
  // quero retornar o menor valor x \ge h, tal que existe um subset com soma x em
  if (!w.size())
   return 0;
  sort(w.rbegin(), w.rend());
  if (w[0] * 2 >= tot)
   return w[0];
  int n = w.size();
  \mathbf{w.pb}(0);
  vector<int> aux;
  int p = 0;
  for (int i = 1; i <= n; i++)
    if (w[i] != w[i - 1])
      int cnt = i - p;
      int x = w[i - 1];
      int j = 1;
      while (j < cnt)
        aux.pb(x * j);
        cnt -= j;
        j *= 2;
      aux.pb(x * cnt);
      p = i;
```

```
swap(aux, w);
  return subset_sum(tot, h);
int f[MAXN];  // f[i] -> quantos "itens" com valor i tem
bitset<MAXN> dp; // dp[i] = 1, se existe um subset com soma i
// garantir que a soma de todo mundo seja < MAXN
void subset_sum(vector<int> &v)
  for (auto const &i : v)
    f[i]++;
  dp[0] = 1;
  for (int i = 1; i < MAXN; i++)
    while (f[i] > 2)
      f[i * 2]++;
      f[i] -= 2;
    while (f[i]--)
      dp \mid = (dp \ll i);
// https://github.com/gabrielpessoal/ICPC-Library/blob/master/code/Miscellaneous
     /SubsetSum.cpp
 Given N non-negative integer weights w and a non-negative target t,
 computes the maximum S \le t such that S is the sum of some subset of the
     weights.
 O(N * max(w[i]))
int knapsack(vector<int> w, int t)
  int a = 0, b = 0;
  while (b < w.size() \&\& a + w[b] <= t)
    a += w[b++];
  if (b == w.size())
    return a;
  int m = *max_element(w.begin(), w.end());
  vector<int> u, v(2 * m, -1);
  v[a + m - t] = b;
  for (int i = b; i < w.size(); i++)</pre>
    for (int x = 0; x < m; x++)
      v[x + w[i]] = max(v[x + w[i]], u[x]);
    for (int x = 2 * m; --x > m;)
      for (int j = max(011, u[x]); j < v[x]; j++)
        v[x - w[j]] = max(v[x - w[j]], j);
  a = t;
  while (v[a + m - t] < 0)
   a--;
  return a:
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
```

#### 2.24 suffix sum

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
int n, k;
int a[101];
int dp[101][100005];
int sum[101][100005];
int solve(int i, int j);
void calc(int i)
  sum[i][0] = solve(i, 0);
  for (int j = 1; j <= k; j++)</pre>
    sum[i][j] = (sum[i][j-1] + solve(i, j)) % mod;
int solve(int i, int j)
  if (i == n)
   return (j == 0) ? 1 : 0;
  if (dp[i][j] != -1)
  return dp[i][j];
int ans = 0, limit = min(j, a[i]);
  if (sum[i + 1][j] == -1)
   calc(i + 1);
  ans = (ans + sum[i + 1][j]) % mod;
  if (j - limit - 1 >= 0)
   ans = (ans - sum[i + 1][j - limit - 1] + mod) % mod;
  return dp[i][j] = ans;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> k;
  for (int i = 0; i < n; i++)
   cin >> a[i];
  memset(dp, -1, sizeof(dp));
  memset(sum, -1, sizeof(sum));
  cout << solve(0, k) << endl;</pre>
  return 0;
// uma dp que tem uma recorrencia do tipo:
// dp[i][i] = dp[i + 1][1] + dp[i + 1][1 + 1] + ... + dp[i + 1][r]
// o jeito de fazer isso com soma de sufixo sem ter que codar a dp iterativa :)
```

### 2.25 tip

```
#define lli long long int
#define pb push_back
#define in insert
#define pi pair<int, int>
#define pd pair<double, int>
#define pib pair<pi, bool>
#define mp make_pair
#define fir first
#define sec second
#define MAXN 200001
#define MAXL 10001
#define mod 1000000007
int dp[MAXN];
vector<int> v;
int solve(int rem)
  if (rem == 0)
   return 1;
  if (rem < 0)
   return 0;
  if (dp[rem] >= 0)
    return dp[rem];
  for (int i = 0; i < v.size(); i++)</pre>
   if (solve(rem - v[i]))
      return dp[rem - v[i]] = 1;
  return dp[rem] = 0;
signed main()
  int n, m;
  cin >> n >> m;
  v.resize(n);
  for (int i = 0; i < n; i++)</pre>
   cin >> v[i];
  memset(dp, -1, sizeof(dp));
  (solve(m)) ? cout << "Yes\n" : cout << "No\n";
  return 0:
```

# 3 Geometry

#### 3.1 ConvexHull

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define MAXN 100005
struct point
  int x, y, id;
  point (int x, int y, int id) : x(x), y(y), id(id) {}
  point operator-(point const &o) const { return {x - o.x, y - o.y, -1}; }
  bool operator<(point const &o) const
    if (x == 0.x)
      return y < o.y;</pre>
    return x < o.x;
  int operator (point const &o) const { return x * o.v - y * o.x; }
int ccw (point const &a, point const &b, point const &x)
  auto p = (b - a) \cdot (x - a);
  return (p > 0) - (p < 0);
vector<point> convex_hull(vector<point> P) // sem colineares
```

```
sort(P.begin(), P.end());
 vector<point> L, U;
 for (auto p : P)
    while (L.size() >= 2 && ccw(L.end()[-2], L.end()[-1], p) == -1)
     L.pop_back();
    L.push_back(p);
  reverse(P.begin(), P.end());
 for (auto p : P)
    while (U.size() >= 2 && ccw(U.end()[-2], U.end()[-1], p) == -1)
     U.pop_back();
   U.push_back(p);
 L.insert(L.end(), U.begin(), U.end() - 1);
 return L:
vector<point> convex_hull_no_collinears(vector<point> P) // com colineares
 sort(P.begin(), P.end());
 vector<point> L, U;
 for (auto p : P)
   while (L.size() >= 2 && ccw(L.end()[-2], L.end()[-1], p) <= 0)</pre>
     L.pop_back();
   L.push_back(p);
 reverse(P.begin(), P.end());
 for (auto p : P)
   while (U.size() \ge 2 \&\& ccw(U.end()[-2], U.end()[-1], p) \le 0)
     U.pop_back();
   U.push_back(p);
 L.insert(L.end(), U.begin(), U.end() - 1);
 return L;
signed main()
  // int n:
  // cin >> n;
 // vector<point> v(n);
  // for (int i = 0; i < n; i++)
     cin >> v[i].x >> v[i].y;
     v[i].id = i;
  // vector<point> ans = convex_hull(v);
  // vector<int> ids;
  // for (auto const &i : ans)
      ids.pb(i.id);
  // sort(ids.begin(), ids.end());
  // ids.erase(unique(ids.begin(), ids.end()), ids.end());
  // for (auto const &i : ids)
  // cout << i + 1 << " ";
  // cout << endl;
 int n;
 while (cin >> n)
   if (n == 0)
    vector<point> v(n);
    for (int i = 0; i < n; i++)
     cin >> v[i].x >> v[i].y;
     v[i].id = i;
    vector<point> ans = convex_hull_no_collinears(v);
    vector<pi> resp;
    for (auto const &i : ans)
     if (!resp.size() || (pi(i.x, i.y) != resp.back()))
        resp.pb({i.x, i.y});
```

```
cout << resp.size() << endl;
for (auto [x, y] : resp)
     cout << x << " " << y << endl;
}
}
// https://codeforces.com/gym/104555/problem/G
// https://open.kattis.com/problems/convexhull (sem colinear)</pre>
```

### 3.2 convex hull point location

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 5005
#define mod 998244353
struct pt
  int x, y;
  pt operator+(pt p) { return {x + p.x, y + p.y};
  pt operator-(pt p) { return {x - p.x, y - p.y}; }
  bool operator == (pt p) { return (x == p.x && y == p.y); }
  int cross(pt p) { return x * p.y - y * p.x; }
  int cross(pt a, pt b) { return (a - *this).cross(b - *this); }
  int dot(pt p) { return x * p.x + y * p.y; }
bool cmp_x(pt a, pt b)
  if (a.x != b.x)
   return a.x < b.x;
  return a.y < b.y;</pre>
// acha o convex hull
vector<pt> convex_hull(vector<pt> pts)
  if (pts.size() <= 1)
   return pts;
  sort(pts.begin(), pts.end(), cmp_x);
  vector<pt> h(pts.size() + 1);
  int s = 0, t = 0;
  for (int it = 2; it--; s = --t, reverse(pts.begin(), pts.end()))
    for (auto const &p : pts)
      while (t >= s + 2 \&\& h[t - 2].cross(h[t - 1], p) <= 0)
      h[t++] = p;
  return \{h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])\};
int sgn(int x)
  return (x > 0) - (x < 0);
int side_of(pt s, pt e, pt p)
  return sgn(s.cross(e, p));
bool on_segment(pt s, pt e, pt p)
  return p.cross(s, e) == 0 && (s - p).dot(e - p) <= 0;
// retorna se o ponto p esta dentro ou nao do convex hull 1
// caso strict = true, entao considera true se tiver na borda
```

```
// caso strict = false, entao considera false se tiver na borda
bool is_hull(vector<pt> &l, pt p, bool strict = true)
  int a = 1, b = 1.size() - 1, r = !strict;
 if (1.size() < 3)
   return r && on_segment(1[0], 1.back(), p);
  if (side_of(1[0], 1[a], 1[b]) > 0)
   swap(a, b);
  if (side_of(1[0], 1[a], p) >= r || side_of(1[0], 1[b], p) <= -r)
    return false;
  while (abs(a - b) > 1)
    int c = (a + b) / 2;
    (side\_of(1[0], 1[c], p) > 0 ? b : a) = c;
  return sqn(l[a].cross(l[b], p)) < r;</pre>
signed main()
 int n;
  cin >> n;
  vector<pt> v(n);
  for (int i = 0; i < n; i++)
    cin >> v[i].x >> v[i].y;
  vector<pt> ans = convex_hull(v);
  for (int i = 0; i < n; i + +)
    if (!is_hull(ans, v[i]))
      cout << i + 1 << " ";
  cout << endl;
// h da subregional - https://codeforces.com/gym/104555/problem/G
```

### 3.3 dynamic ch

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define double long double
#define mod 1000000007
const double eps = 1e-9;
struct pt
  double x, y;
  pt operator-(pt p) { return {x - p.x, y - p.y}; }
 bool eq(double a, double b) const
    return abs(a - b) <= eps;</pre>
  double operator^(const pt p) const { return x * p.y - y * p.x; }
  bool operator<(const pt p) const</pre>
    if (!eq(x, p.x))
      return x < p.x;
    if (!eq(y, p.y))
```

```
return y < p.y;</pre>
    return 0;
  bool operator == (const pt p) const
    return eq(x, p.x) and eq(y, p.y);
double sarea(pt p, pt q, pt r)
  return ((q - p) ^ (r - q)) / 2;
bool ccw(pt p, pt q, pt r)
  return sarea(p, q, r) > eps;
// https://github.com/brunomaletta/Biblioteca/blob/master/Codigo/Problemas/
    dvnamicHull.cpp
struct upper
  set<pt> se;
  set<pt>::iterator it;
  // 0 - fora
  // 1 - dentro
  // 2 - na borda
  int is_under(pt p)
    it = se.lower bound(p);
    if (it == se.end())
     return 0:
    if (it == se.begin())
     return p == *it ? 2 : 0;
    if (ccw(p, *it, *prev(it)))
     return 1;
   return ccw(p, *prev(it), *it) ? 0 : 2;
  void insert(pt p)
    if (is_under(p))
      return;
    if (it != se.end())
      while (next(it) != se.end() and !ccw(*next(it), *it, p))
        it = se.erase(it);
    if (it != se.begin())
      while (--it != se.begin() and !ccw(p, *it, *prev(it)))
        it = se.erase(it);
    se.insert(p);
};
struct dyn_hull
  upper U, L;
  int is_inside(pt p)
    int u = U.is_under(p), l = L.is_under({-p.x, -p.y});
    if (!u || !1)
     return 0;
    return max(u, 1);
  void insert(pt p)
    U.insert(p);
    L.insert(\{-p.x, -p.y\});
  int size()
    int ans = U.se.size() + L.se.size();
    return ans <= 2 ? ans / 2 : ans - 2;
};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
// convex hull dinamico
// problema para usar: https://open.kattis.com/problems/hiringhelp
```

## 3.4 halfplane intersection

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
const long double eps = 1e-9;
const long double inf = 1e9;
struct pt
 long double x, y; pt(long double x = 0, long double y = 0) : x(x), y(y) {}
  friend pt operator+(pt p, pt q)
    return pt(p.x + q.x, p.y + q.y);
  friend pt operator-(pt p, pt q)
    return pt(p.x - q.x, p.y - q.y);
  friend pt operator*(pt p, long double k)
    return pt(p.x * k, p.y * k);
  friend long double dot(pt p, pt q)
    return p.x * q.x + p.y * q.y;
  friend long double cross(pt p, pt q)
    return p.x * q.y - p.y * q.x;
struct halfplane
  pt p, pq;
  long double angle;
  halfplane() {}
  halfplane(pt a, pt b) : p(a), pq(b - a)
    angle = atan21(pq.y, pq.x);
  bool out (const pt &r)
    return cross(pq, r - p) < -eps;</pre>
  bool operator<(halfplane e) const
    return angle < e.angle;
  friend pt inter(halfplane s, halfplane t)
    long double alpha = cross((t.p - s.p), t.pq) / cross(s.pq, t.pq);
    return s.p + (s.pq * alpha);
};
vector<pt> hp_intersect (vector<halfplane> &h)
  pt box[4] = {pt(inf, inf), pt(-inf, inf), pt(-inf, -inf), pt(inf, -inf)}; //
```

```
Bounding box in CCW order
  for (int i = 0; i < 4; i++)
   halfplane aux(box[i], box[(i + 1) % 4]);
   h.pb(aux);
  sort(h.begin(), h.end());
  deque<halfplane> dq;
  int len = 0:
  for (int i = 0; i < h.size(); i++)</pre>
    while (len > 1 && h[i].out(inter(dq[len - 1], dq[len - 2])))
      dq.pop_back();
      --len;
    while (len > 1 && h[i].out(inter(dq[0], dq[1])))
      dq.pop_front();
      --len;
    if (len > 0 && fabsl(cross(h[i].pq, dq[len - 1].pq)) < eps)</pre>
      if (dot(h[i].pq, dq[len - 1].pq) < 0.0)</pre>
        return vector<pt>();
      if (h[i].out(dq[len - 1].p))
        dq.pop_back();
        --len;
      else
        continue:
    dq.push_back(h[i]);
  while (len > 2 && dq[0].out(inter(dq[len - 1], dq[len - 2])))
    dq.pop_back();
    --len;
  while (len > 2 && dq[len - 1].out(inter(dq[0], dq[1])))
    dq.pop_front();
    --len;
  if (len < 3)
   return vector<pt>();
  vector<pt> ret(len);
  for (int i = 0; i + 1 < len; i++)
   ret[i] = inter(dq[i], dq[i + 1]);
  ret.back() = inter(dq[len - 1], dq[0]);
  return ret;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int q;
 cin >> q; // quantidade de poligonos
  vector<halfplane> h;
  while (q--)
   int n;
   cin >> n;
    vector<pt> v(n);
    for (int i = 0; i < n; i++)
      cin >> v[i].x >> v[i].y;
```

for (int i = 0; i < n; i++)

```
int j = (i + 1) % n;
     h.pb(halfplane(v[i], v[j]));
  vector<pt> ans = hp_intersect(h);
 if (ans.size() == 0)
   cout << "0.0\n";
    return 0;
 long double res = 0;
 for (int i = 0; i < ans.size(); i++) // area da interseccao</pre>
   pt p = (i) ? ans[i - 1] : ans.back();
   pt q = ans[i];
   res += (p.x - q.x) * (p.y + q.y);
 double resp = abs(res) / 2;
 cout << fixed << setprecision(15) << resp << endl;</pre>
 return 0;
// half-plane intersection
// definicoes:
// half-plane - regiao planar que consiste de todos os pontos que estao de um
    lado de uma reta
// geralmente podem ser descritos da seguninte forma
// conjuntos dos pontos (x, y) que satisfazem algo do tipo:
// ax + by + c \le 0 ou ax + by + c \ge 0
// da pra representar as retas e os half-planes atraves de um ponto (que ta na
    reta) e o vetor de direcao
// e dai pros half-planes, considerando que e a regiao da esquerda em relacao ao
     vetor de direcao
// alem disso, considerar uma bounding box sendo um retangulo, pra caso a
    inserseccao dos halfplanes nao seja "fechada"
// https://open.kattis.com/problems/bigbrother
// qual a area que voce pode botar uma camera dentro do poligono
// tal que de um ponto escolhido, e possivel ver todos o poligono
// dai considerar todos os halfplanes de arestas do poligono
// e achar a interseccao de todos esses halfplanes
// https://www.codechef.com/problems/CHN02
// achar a area da interseccao de varios poligonos convexos
// considerar todos os halfplanes de arestas do poligono
// e achar a interseccao de todos esses halfplanes
```

#### 3.5 kd tree

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 300005
#define mod 998244353
#define inf LLONG_MAX
struct pt
  int x, y, id;
  pt() {}
```

```
pt(int xx, int yy) { x = xx, y = yy; }
  pt operator-(pt p) const { return pt(x - p.x, y - p.y); }
  bool operator<(pt p) const { return x < p.x; }</pre>
  int dist() const { return x * x + y * y; }
bool on_x(const pt &a, const pt &b) { return a.x < b.x; }
bool on_y(const pt &a, const pt &b) { return a.y < b.y; }</pre>
struct node
  int id;
  int x0 = inf, x1 = -inf, y0 = inf, y1 = -inf;
  node *first = 0, *second = 0;
  int distance(const pt &p)
    int x = (p.x < x0 ? x0 : p.x > x1 ? x1
                                       : p.x);
    int y = (p.y < y0 ? y0 : p.y > y1 ? y1
                                       : p.y);
    return (pt(x, y) - p).dist();
  node(vector<pt> &&vp) : pp(vp[0])
    for (pt p : vp)
      x0 = min(x0, p.x);
      x1 = max(x1, p.x);
      y0 = min(y0, p.y);
      y1 = max(y1, p.y);
    if (vp.size() > 1)
      sort(vp.begin(), vp.end(), x1 - x0 >= y1 - y0 ? on_x : on_y);
      int half = vp.size() / 2;
      first = new node({vp.begin(), vp.begin() + half});
      second = new node({vp.begin() + half, vp.end()});
};
struct kd_tree
  node *root:
  kd_tree(const vector<pt> &vp) : root(new node({vp.begin(), vp.end()})) {}
  pi search(node *n, const pt &p)
    if (!n->first)
      if (n->pp.x == p.x && n->pp.y == p.y)
       return make_pair(inf, n->pp.id); // distancia infinita pra pontos iquais
      return make_pair((p - n->pp).dist(), n->pp.id);
    node *f = n \rightarrow first, *s = n \rightarrow second;
    int bfirst = f->distance(p), bsec = s->distance(p);
    if (bfirst > bsec)
      swap(bsec, bfirst), swap(f, s);
    auto best = search(f, p);
    if (bsec < best.first || (!f->first))
     best = min(best, search(s, p));
    return best;
  pi nearest (const pt &p)
    return search (root, p);
};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  vector<pt> v(n);
  for (int i = 0; i < n; i++)
    cin >> v[i].x >> v[i].y;
   v[i].id = i;
  kd_tree t(v);
```

```
pii ans = {inf, {inf, inf}};
 for (int i = 0; i < n; i++)
   pi curr = t.nearest(v[i]);
   ans = min(ans, {curr.fir, {i, curr.sec}});
 cout << fixed << setprecision(6) << ans.sec.fir << " " << ans.sec.sec << " "</pre>
      << sqrt(ans.fir) << endl;
 return 0;
// closest pair of points com kdtree
// da pra ser adaptado pro 3d tbm
// quando um ponto (x, y) pode aparecer em mais de um indice, tratar antes
// fonte: https://github.com/kth-competitive-programming/kactl/blob/main/kactl.
// testei em:
// https://codeforces.com/contest/429/problem/D
// https://www.spoj.com/problems/CLOPPAIR/
// https://vjudge.net/problem/UVA-10245
// https://codeforces.com/gym/104020/problem/L (3D)
```

## 3.6 LineSweep

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 1000000007
#define PI acos(-1)
const double EPS = 1e-9;
struct pt
  double x, y;
struct seq
  pt p, q;
  double get_y (double x) const
    if (abs(p.x - q.x) < EPS)
     return p.y;
    return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x);
bool intersect1d(double 11, double r1, double 12, double r2)
  if (11 > r1)
   swap(11, r1);
  if (12 > r2)
    swap(12, r2);
  return max(11, 12) <= min(r1, r2) + EPS;</pre>
int vec(const pt &a, const pt &b, const pt &c)
  double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y) * (c.x - a.x);
  return abs(s) < EPS ? 0 : s > 0 ? +1
```

```
: -1;
bool intersect (const seg &a, const seg &b)
  return intersect1d(a.p.x, a.q.x, b.p.x, b.q.x) &&
         intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) &&
         vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <= 0 &&
         vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <= 0;
bool operator<(const seg &a, const seg &b)</pre>
  double x = max(min(a.p.x, a.q.x), min(b.p.x, b.q.x));
  return a.get_y(x) < b.get_y(x) - EPS;</pre>
struct event
  double x;
  int tp, id;
  event() {}
  event (double x, int tp, int id) : x(x), tp(tp), id(id) {}
  bool operator<(const event &e) const</pre>
   if (abs(x - e.x) > EPS)
      return x < e.x;</pre>
    return tp > e.tp;
};
set<seg> s;
set<seg>::iterator prev(set<seg>::iterator it)
  return it == s.begin() ? s.end() : --it;
set<seq>::iterator next(set<seq>::iterator it)
  return ++it:
pi line_sweep(vector<seg> v)
  vector<event> e;
  for (int i = 0; i < v.size(); i++)</pre>
    e.push_back({min(v[i].p.x, v[i].q.x), 1, i});
   e.push_back({max(v[i].p.x, v[i].q.x), 0, i});
  sort(e.begin(), e.end());
  for (int i = 0; i < e.size(); i++)
    int id = e[i].id;
    if (e[i].tp == 1)
      auto nxt = s.lower_bound(v[id]), prv = prev(nxt);
      if (nxt != s.end() && intersect(*nxt, v[id]))
        return {(*nxt).id, id};
      if (prv != s.end() && intersect(*prv, v[id]))
        return {(*prv).id, id);
      s.insert(nxt, v[id]);
    else
      auto where = s.lower_bound(v[id]);
      auto nxt = next(where), prv = prev(where);
      if (nxt != s.end() && prv != s.end() && intersect(*nxt, *prv))
        return {(*prv).id, (*nxt).id};
      s.erase(where);
  return {-1, -1};
signed main()
  int n:
  cin >> n;
  vector<seg> v(n);
  for (int i = 0; i < n; i++)
   cin >> v[i].p.x >> v[i].p.y >> v[i].q.x >> v[i].q.y;
   v[i].id = i;
```

```
}
pi ans = line_sweep(v);
if (ans.fir == -1)
{
    cout << "NO\n";
}
else
{
    cout << "YES\n";
    cout << ans.fir + 1 << " " << ans.sec + 1 << endl;
}
return 0;
}
// https://cp-algorithms.com/geometry/intersecting_segments.html
// https://acm.timus.ru/problem.aspx?space=1&num=1469</pre>
```

#### 3.7 line trick

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500005
#define mod 998244353
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
pi get_line(pi x, pi y)
  int xx = x.fir - y.fir;
  int yy = x.sec - y.sec;
  int g = \underline{gcd(abs(xx), abs(yy))};
  if (q != 0)
    xx /= g, yy /= g;
  if (xx < 0)
    xx += -1:
    yy \star = -1;
  return {xx, yy};
void solve()
 int n;
  cin >> n;
  vector<pi> v(n);
  for (int i = 0; i < n; i++)
    cin >> v[i].fir >> v[i].sec;
 map<pi, int> mp;
  map<pi, int> repr;
  int cc = 0;
  // pega 2 indices i e j, e acha a reta que passa por esses 2 pontos
  while (cc < (2 * n))
    int x = rng() % n;
    int y = rng() % n;
    if (x == y)
      continue;
    pi l = get_line(v[x], v[y]);
```

```
mp[1]++;
   repr[1] = x;
  vector<pii> vec;
  for (auto const &i : mp)
   vec.pb({i.sec, i.fir});
 sort(vec.rbegin(), vec.rend());
  // agora considerando as duas retas de maior frequencia
  // veja quantos pontos do conjunto estao nela
  int ans = n;
  for (int i = 0; i < min(211, (int)vec.size()); i++)</pre>
    // quardo um "representante" dessa reta, que eu sei que ta nessa reta
    // pq dai eh so iterar por cada ponto e ver se o get_line bate com a reta
        que queremos
    int guy = repr[vec[i].sec];
    int cnt = 0;
    for (int j = 0; j < n; j++)
     if (j == guy || get_line(v[guy], v[j]) == vec[i].sec)
    ans = min(ans, n - cnt);
  cout << ans << endl;
signed main()
 ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int q;
 cin >> q;
  for (int i = 1; i \le q; i++)
   cout << "Case #" << i << ": ";
   solve();
 return 0;
// problema dahora
// https://www.facebook.com/codingcompetitions/hacker-cup/2024/practice-round/
    problems/C
// tem n pontos (x, y)
// ache a quantidade maxima de pontos, tal que, existe uma reta que passa por
    todos eles
// e imprima n - essa quantidade
// mas nao quero printar a resposta exata
// se a resposta otima for igual a m
// posso printar ate m * 2
// o que significa que da pra fazer umas heuristicas
```

#### 3.8 minkowski

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 15
#define mod 1000000007
```

```
struct pt
  int x, y;
  bool operator<(pt ot)</pre>
    if (x != ot.x)
     return x < ot.x;</pre>
    return y < ot.y;</pre>
  void operator=(pt p) { x = p.x, y = p.y; }
  bool operator==(pt p) { return (x == p.x && y == p.y); }
 bool operator!=(pt p) { return (x != p.x || y != p.y); }
  pt operator+(const pt &p) { return {x + p.x, y + p.y}; }
  pt operator-(const pt &p) { return {x - p.x, y - p.y}; }
  pt operator*(int d) { return {x * d, y * d};
  pt operator/(int d) { return {x / d, y / d};
  int cross(pt ot) const { return x * ot.y - y * ot.x; }
  int cross(pt a, pt b) const { return (a - *this) cross(b - *this); }
enum type
  outside,
  inside.
  boundary
int cross(pt v, pt w)
  return v.x * w.y - v.y * w.x;
bool ccw(pt a, pt b, pt c)
  return cross(b - a, c - b) > 0;
void radial_sort(vector<pt> &a)
  pt pivot = *min_element(a.begin(), a.end());
  auto cmp = [&] (pt p, pt q)
    if (p == pivot || q == pivot)
     return q != pivot;
    return ccw(pivot, p, q) > 0;
  sort(a.begin(), a.end(), cmp);
vector<pt> trata(vector<pt> p)
  vector<pt> ans;
  for (int i = 0; i < p.size(); i++)</pre>
    while (ans.size() \ge 2 \&\& ans.back().cross(p[i], ans.end()[-2]) == 0)
     ans.pop_back();
    ans.pb(p[i]);
  if (ans.size() > 2 \&\& ans.back().cross(p[0], ans.end()[-2]) == 0)
   ans.pop back();
  return ans;
void prepare(vector<pt> &p)
  radial_sort(p); // sort points in counter-clockwise order
  p = trata(p); // and the polygon dont have 3 colinear points
int sqn(int val)
  if (val > 0)
   return 1:
  else if (val < 0)</pre>
   return -1;
  return 0;
bool in_seg(pt p, pt a, pt b)
  // check if point p is in the line segment formed by a and b
  if (a.cross(b, p) == 0)
    return (p.x \ge min(a.x, b.x) && p.x \le max(a.x, b.x) && p.y \ge min(a.y, b.y)
          && p.y <= max(a.y, b.y));
  return 0;
```

```
bool in_tri(pt p, pt a, pt b, pt c)
    // check if point p is in the triangle formed by a, b and c
    int a1 = abs(a.cross(b, c));
    int a2 = abs(p.cross(a, b)) + abs(p.cross(a, c)) + abs(p.cross(b, c));
    return a1 == a2;
int in_polygon(vector<pt> &poly, pt p)
    int n = poly.size();
    if (n == 1)
        return (p == poly[0]) ? type::boundary : type::outside;
    if (n == 2)
    \label{eq:continuous_continuous_continuous} \textbf{return} \text{ (in\_seg(p, poly[0], poly[1])) ? type::boundary : type::outside;} \\ \textbf{if (poly[0].cross(poly[1], p) != 0 && sgn(poly[0].cross(poly[1], p)) != sgn(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(poly[0].cross(
               poly[0].cross(poly[1], poly[n - 1])))
        return type::outside;
    if (poly[0].cross(p, poly[n-1]) != 0 && sgn(poly[0].cross(p, poly[n-1]))
               != sgn(poly[0].cross(poly[1], poly[n - 1])))
         return type::outside;
     int 1 = 2, r = n - 1;
    if (poly[0].cross(poly[1], p) > 0)
        while (1 < r)
              int mid = (1 + r) >> 1;
              (poly[0].cross(poly[mid], p) \le 0) ? r = mid : 1 = mid + 1;
    if (!in_tri(p, poly[0], poly[1 - 1], poly[1]))
        return type::outside;
    if (in_seg(p, poly[1 - 1], poly[1]))
        return type::boundary;
    if (in_seg(p, poly[0], poly[1]))
        return type::boundary;
    if (in_seg(p, poly[0], poly[n - 1]))
        return type::boundary;
     return type::inside;
vector<pt> minkowski(vector<pt> a, vector<pt> b)
    prepare(a);
    prepare (b);
    a.push_back(a[0]);
    a.push_back(a[1]);
    b.push back(b[0]);
    b.push_back(b[1]);
    vector<pt> ans;
    int i = 0, j = 0;
    while (i < a.size() - 2 || j < b.size() - 2)
        ans.pb(a[i] + b[j]);
         auto c = cross(a[i + 1] - a[i], b[j + 1] - b[j]);
         if (c >= 0)
              i++;
         if (c <= 0)
              j++;
    return ans:
signed main()
    ios_base::sync_with_stdio(false);
    cin.tie(NULL);
    vector<pt> v;
    for (int _ = 0; _ < 3; _++)
        int n:
         cin >> n;
        vector<pt> p(n);
for (int i = 0; i < n; i++)</pre>
             cin >> p[i].x >> p[i].y;
         if ( == 0)
             v = p;
         else
              v = minkowski(v, p);
```

```
prepare(v);
int q;
cin >> q;
while (q--)
{
  pt p;
  cin >> p.x >> p.y;
  p.x *= 3, p.y *= 3;
    // ve se o ponto (3x, 3y) esta na bora, dentro ou fora do poligono v
    (in_polygon(v, p) != type::outside) ? cout << "YES\n" : cout << "NO\n";
}
return 0;
}
// problema exemplo:
// https://codeforces.com/contest/87/problem/E</pre>
```

#### 3.9 points and vectors

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
#define PI acos (-1)
namespace p
  struct pt
    double x, y;
    pt operator+(pt p) { return {x + p.x, y + p.y}; } // soma de pontos
    pt operator-(pt p) { return {x - p.x, y - p.y}; } // subtracao de pontos
    pt operator*(double d) { return {x * d, y * d}; } // multiplicacao por um
    pt operator/(double d) { return {x / d, y / d}; } // divisao por um double
  double dot(pt v, pt w) // produto escalar (dot product)
    return v.x * w.x + v.y * w.y;
  bool is_perp(pt v, pt w) // retorna se dois vetores sao perpendiculares (
       angulo 90 graus)
    return dot(v, w) == 0;
  double cross(pt v, pt w) // produto vetorial (cross product)
    return v.x * w.y - v.y * w.x;
  double dist(pt a, pt b) // distancia entre 2 pontos
    pt c = a - b;
    return sqrt(c.x * c.x + c.y * c.y);
  double dist2(pt a, pt b) // retorna o quadrado da distancia entre dois pontos
    pt c = a - b;
    return c.x * c.x + c.y * c.y;
  bool is_colinear(pt a, pt b, pt c) // retorna se os pontos a, b e c sao
       colineares
```

```
return cross(b - a, c - a) == 0;
  bool ccw(pt a, pt b, pt c) // retorna se os pontos a,b e c estao no sentido
       anti horario
    return cross(b - a, c - b) > 0;
  bool cw(pt a, pt b, pt c) // retorna se os pontos a,b e c estao no sentido
    return cross(b - a, c - b) < 0;
  double modulo (pt v) //|v| = sqrt(x2 + y2)
    return sqrt(v.x * v.x + v.y * v.y);
  double angle (pt a, pt b, pt c) // angulo entre os vetores ab e ac
    // dot(ab , ac) / |ab| * |ac|
    pt ab = b - a; // vetor ab
    pt ac = c - a; // vetor ac
    double m1 = modulo(ab);
    double m2 = modulo(ac);
    double m3 = m1 * m2;
    return (dot(ab, ac) / m3); // retorna o cos do angulo em graus
  pt rotate(pt p, double a) // rotacionar o ponto p em relacao a origem, em a
       graus, no sentido anti-horario
    a = (a * PI) / 180;
    double xx = (\cos(a) * p.x) + ((\sin(a) * -1) * p.y);
    double yy = (\sin(a) * p.x) + (\cos(a) * p.y);
    pt ans = \{xx, yy\};
    return ans:
  double polar(pt p) // polar angle
    return atan21(p.y, p.x);
  bool cmp(pt a, pt b) // ordenar pontos pelo polar angle
    return polar(a) < polar(b);</pre>
  bool cmp_x(pt a, pt b) // ordenar os pontos pela coordenada x
    if (a.x != b.x)
      return a.x < b.x;
    return a.y < b.y;</pre>
  pt polar_to_cartesian(double r, double theta) // r - distancia do centro,
       theta - polar angle
    pt ans:
    ans.x = r * cos(double(theta) / 180 * PI); // assumindo que theta ta em
        graus, transforma pra radiano
    ans.y = r * sin(double(theta) / 180 * PI);
    return ans;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
```

### 3.10 polygons distance

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
#define int long long int
```

```
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 5005
#define mod 998244353
struct pt
  double x, y;
  pt operator+(const pt p) const { return pt(x + p.x, y + p.y);
  pt operator-(const pt p) const { return pt(x - p.x, y - p.y);
  pt operator*(const double c) const { return pt(x * c, y * c);
  pt operator/(const double c) const { return pt(x / c, y / c);
  double operator*(const pt p) const { return x * p.x + y * p.y;
  double operator^(const pt p) const { return x * p.y - y * p.x;
struct line
  pt p, q;
  line() {}
  line(pt p_, pt q_) : p(p_), q(q_) {}
struct building
  // 0 - circulo, 1 - quadrado, 2 - triangulo
  int type, r;
  vector<pt> v;
 building()
   v.clear();
   \mathbf{r} = 0;
    type = 0;
  void find_vertices()
    pt a, b;
    double vx = v[1].x - v[0].x;
    double vy = v[1].y - v[0].y;
   a.x = (v[0].x + v[1].x) / 2 + (-vy) / 2;
a.y = (v[0].y + v[1].y) / 2 + (vx) / 2;
    b.x = (v[0].x + v[1].x) / 2 - (-vy) / 2;
    b.y = (v[0].y + v[1].y) / 2 - (vx) / 2;
    v.pb(a);
    swap(v[1], v[2]);
    v.pb(b);
};
int c, q, t, n;
vector<br/>building> v;
vector<int> vc, vq, vt;
double sarea(pt p, pt q, pt r)
  return ((q - p) ^ (r - q)) / 2;
double dist(pt p, pt q)
  return hypot(p.y - q.y, p.x - q.x);
double disttoline (pt p, line r)
 return 2 * abs(sarea(p, r.p, r.q)) / dist(r.p, r.q);
double disttoseg(pt p, line r)
  if ((r.q - r.p) * (p - r.p) < 0)
   return dist(r.p, p);
  if ((r.p - r.q) * (p - r.q) < 0)
   return dist(r.q, p);
  return disttoline(p, r);
double dist_circ_seg(pt p1, pt p2, pt p, int r)
  double dist = disttoseg(p, line(p2, p1));
  dist -= r;
```

```
return dist;
double dist_seg_seg(line a, line b)
  double ret = DBL_MAX;
  ret = min(ret, disttoseg(a.p, b));
  ret = min(ret, disttoseg(a.q, b));
  ret = min(ret, disttoseg(b.p, a));
  ret = min(ret, disttoseg(b.g, a));
  return ret;
double dist_square_tri(int i, int j)
  double ans = DBL_MAX;
  for (int x = 0; x < 4; x++)
    int v = (x + 1) % 4;
    for (int x2 = 0; x2 < 3; x2++)
      int y2 = (x2 + 1) % 3;
      ans = min(ans, dist_seg_seg(line(v[i].v[x], v[i].v[y]), line(v[j].v[x2], v[x2])
           [i].v[v2]));
  return ans:
double dist_square_circ(int i, int j)
  double ans = DBL_MAX;
  for (int x = 0; x < 4; x++)
    int y = (x + 1) \% 4;
   ans = min(ans, dist_circ_seg(v[i].v[x], v[i].v[y], v[j].v[0], v[j].r));
  return ans:
double dist_tri_circ(int i, int j)
  double ans = DBL MAX:
  for (int x = 0; x < 3; x++)
    int v = (x + 1) % 3;
   ans = min(ans, dist_circ_seg(v[i].v[x], v[i].v[y], v[j].v[0], v[j].r));
  return ans:
double dist_circ_circ(int i, int j)
  double dist = (v[i].v[0].x - v[j].v[0].x) * (v[i].v[0].x - v[j].v[0].x);
  dist += (v[i].v[0].y - v[j].v[0].y) * (v[i].v[0].y - v[j].v[0].y);
  dist = sqrtl(dist);
  dist = (v[i].r + v[j].r);
  return (dist < 0) ? 0 : dist;</pre>
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> c >> q >> t;
  int n = c + q + t;
  int id = 0;
  for (int i = 0; i < c; i++)
   int x, y, r;
    cin >> x >> y >> r;
    building b;
   b.v.pb(\{x, y\});
   b.r = r;
   b.type = 0;
    v.pb(b);
    vc.pb(id);
    id++;
  for (int i = 0; i < q; i++)
    building b;
   b.type = 1;
    for (int j = 0; j < 2; j++)
```

```
int x, y;
cin >> x >> y;
    b.v.pb({x, y});
  b.find_vertices();
  v.pb(b);
  vq.pb(id);
  id++;
for (int i = 0; i < t; i++)
  building b;
  b.type = 2;
  for (int j = 0; j < 3; j++)
   int x, y;
    cin >> x >> y;
    b.v.pb(\{x, y\});
  v.pb(b);
  vt.pb(id);
  id++;
vector<vector<pair<double, int>>> adj(n + 2);
double ans = DBL_MAX;
for (auto const &i : vq)
  for (auto const & i : vt)
    ans = min(ans, dist_square_tri(i, j));
for (auto const &i : vq)
  for (auto const & j : vc)
    double curr = dist_square_circ(i, j);
    adj[i].pb({curr, j});
    adj[j].pb({curr, i});
for (auto const &i : vt)
  for (auto const &j : vc)
    double curr = dist_tri_circ(i, j);
    adj[i].pb({curr, j});
    adj[j].pb({curr, i});
for (auto const &i : vc)
  for (auto const &j : vc)
    double curr = dist_circ_circ(i, j);
    adj[i].pb({curr, j});
    adj[j].pb({curr, i});
int src = n, sink = n + 1;
for (auto const &i : vt)
  adj[src].pb({0, i});
for (auto const &i : vq)
  adj[i].pb({0, sink});
vector<double> dist(n + 2, 1e18);
vector<bool> vis(n + 2, 0);
dist[src] = 0;
priority_queue<pair<double, int>, vector<pair<double, int>>, greater<pair<</pre>
     double, int>>> pq;
pq.push({dist[src], src});
while (!pq.empty())
  int x = pq.top().sec;
  pq.pop();
  if (vis[x])
```

```
continue;
vis[x] = 1;
for (auto [d, y] : adj[x])
{
    if (dist[y] > dist[x] + d)
        {
             dist[y] = dist[x] + d;
             pq.push({dist[y], y});
        }
    }
    ans = min(ans, dist[sink]);
    cout << fixed << setprecision(15) << ans << endl;
    return 0;
}
// solution for: https://codeforces.com/gym/104603/problem/I</pre>
```

## 3.11 polygon area

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 998244353
double area(vector<pi> fig)
  double res = 0;
  for (unsigned i = 0; i < fig.size(); i++)</pre>
   pi p = i ? fig[i - 1] : fig.back();
   pi q = fiq[i];
   res += (p.fir - q.fir) * (p.sec + q.sec);
  return fabs(res) / 2;
int cross (pi a, pi b)
  return a.fir * b.sec - a.sec * b.fir;
double area2(vector<pi> fig)
  double res = 0;
  for (unsigned i = 0; i < fig.size(); i++)</pre>
   pi p = i ? fiq[i - 1] : fiq.back();
   pi q = fig[i];
    res += cross(p, q);
  return fabs(res) / 2;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
// achar area de um poligono
// tomar cuiddado com a ordem
// percorrer os vertices em sentido horario ou anti-horario
```

## 3.12 polygon isomorfism

```
#include <bits/stdc++.h>
using namespace std;
```

```
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 4007
#define mod 998244353
struct pt
  int x, y;
  pt operator+(pt p) { return {x + p.x, y + p.y}; }
  pt operator-(pt p) { return {x - p.x, y - p.y}; }
  bool operator==(pt p) { return (x == p.x && y == p.y); }
  int cross(pt p) { return x * p.y - y * p.x; }
  int cross(pt a, pt b) { return (a - *this).cross(b - *this); }
  int dot(pt p) { return x * p.x + y * p.y; }
bool cmp_x(pt a, pt b)
  if (a.x != b.x)
    return a.x < b.x;</pre>
  return a.y < b.y;</pre>
vector<pt> convex_hull(vector<pt> pts)
  if (pts.size() <= 1)
   return pts;
  sort(pts.begin(), pts.end(), cmp_x);
  vector<pt> h(pts.size() + 1);
  int s = 0, t = 0;
  for (int it = 2; it--; s = --t, reverse(pts.begin(), pts.end()))
    for (auto const &p : pts)
      while (t >= s + 2 \&\& h[t - 2].cross(h[t - 1], p) <= 0)
      h[t++] = p;
  return {h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])};
int max_suffix(vector<pair<long long, long double>> s, bool mi = false)
  s.push_back(*min_element(s.begin(), s.end()));
  s.back().first -= 1;
  s.back().second -= 1;
  int ans = 0;
  for (int i = 1; i < s.size(); i++)</pre>
    while (ans + j < i \text{ and } s[i + j] == s[ans + j])
    if(s[i+j] > s[ans+j])
      if (!mi or i != s.size() - 2)
        ans = i:
    else if (i)
      i += j - 1;
  return ans;
vector<pair<long long, long double>> max_cyclic_shift(vector<pair<long long,</pre>
     long double>> s)
  int n = s.size();
  for (int i = 0; i < n; i++)
   s.pb(s[i]);
  int id = max_suffix(s);
  vector<pair<long long, long double>> ans;
  for (int i = 0; i < n; i++)
    ans.pb(s[id]);
    id = (id + 1) % n;
  return ans;
```

```
int sqr(int x)
  return x * x:
int dd(pt a, pt b)
  return sqr(a.x - b.x) + sqr(a.y - b.y);
long long dot(pt a, pt b)
  return a.x * b.x + a.y * b.y;
vector<pair<long long, long double>> get_sides_and_dots(vector<pt> v)
  int n = (int) v.size();
  vector<pair<long long, long double>> ans;
  for (int i = 0; i < n; i++)
    pt prv = v[i ? i - 1 : n - 1];
    pt nxt = v[i + 1 < n ? i + 1 : 0];
    long long dist = dd(v[i], v[(i + 1) % n]);
    long double angle = dot(prv - v[i], nxt - v[i]);
    ans.emplace_back(dist, angle);
  return ans:
signed main()
  ios::sync_with_stdio(false);
  cin.tie(0);
  int n, m;
  cin >> n >> m;
  vector<pt> a(n);
  for (int i = 0; i < n; i++)
   cin >> a[i].x >> a[i].y;
  auto cha = convex_hull(a);
  auto distsA = get_sides_and_dots(cha);
  vector<pt> b(m);
  for (int i = 0; i < m; i++)
   cin >> b[i].x >> b[i].y;
  auto chb = convex_hull(b);
  auto distsB = get_sides_and_dots(chb);
  vector<pair<long long, long double>> aa = max_cyclic_shift(distsA);
  vector<pair<long long, long double>> bb = max_cyclic_shift(distsB);
(aa == bb) ? cout << "YES\n" : cout << "NO\n";</pre>
// https://codeforces.com/problemset/problem/1017/E
// dados dois conjuntos de pontos
// achar o convex hull de cada conjunto
// e em seguida ver se os poligonos sao isomorfos
// podemos checar olhando para o comprimento de cada aresta e o dot product
```

## 3.13 smallest enclosing circle

```
#define double long double
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
struct pt
  double x, y;
  pt operator+(pt p) { return {x + p.x, y + p.y}; } // soma de pontos
pt operator-(pt p) { return {x - p.x, y - p.y}; } // subtracao de pontos
  pt operator*(double d) { return {x * d, y * d}; } // multiplicacao por um
  pt operator/(double d) { return {x / d, y / d}; } // divisao por um double
struct circle
  double r;
bool inside(circle c, pt p)
  double dist = (c.c.x - p.x) * (c.c.x - p.x) + (c.c.y - p.y) * (c.c.y - p.y);
  return dist <= c.r;</pre>
circle get_circle(pt a, pt b)
  pt c = \{(a.x + b.x) / 2.0, (a.y + b.y) / 2.0\};
  double dist = sqrt((a.x - b.x) * (a.x - b.x) + (a.y - b.y) * (a.y - b.y));
  dist /= 2.0;
 dist *= dist;
  return {c, dist};
pt get_center(pt b, pt c)
  double bb = b.x * b.x + b.y * b.y;
  double cc = c.x * c.x + c.y * c.y;
  double dd = b.x * c.y - b.y * c.x;
  return \{(c.y * bb - \hat{b}.y * \hat{c}c) / (2 * dd), (b.x * cc - c.x * bb) / (2 * dd)\};
circle get_circle(pt a, pt b, pt c)
 b = b - a;
  c = c - a;
  pt p = get_center(b, c);
  double dist = (a.x - p.x) * (a.x - p.x) + (a.y - p.y) * (a.y - p.y);
  return {p, dist};
circle solve2 (vector<pt> &v)
  if (v.empty())
    return {{0, 0}, 0};
  if (v.size() == 1)
    return {v[0], 0};
  if (v.size() == 2)
    return get_circle(v[0], v[1]);
  for (int i = 0; i < 3; i++)
    for (int j = i + 1; j < 3; j++)
      circle c = get_circle(v[i], v[j]);
      bool ok = 1;
      for (auto const &k : v)
        ok &= inside(c, k);
      if (ok)
        return c:
  return get_circle(v[0], v[1], v[2]);
circle solve(vector<pt> &v, vector<pt> r, int n)
  if (n == 0 || r.size() == 3)
    return solve2(r);
  int idx = rand() % n;
```

```
pt p = v[idx];
  swap(v[idx], v[n-1]);
  circle c = solve(v, r, n - 1);
  if (inside(c, p))
   return c;
  r.pb(p);
  return solve(v, r, n - 1);
circle welz1(vector<pt> v)
  random_shuffle(v.begin(), v.end());
  return solve(v, {}, v.size());
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  srand(time(NULL));
  int n;
  cin >> n;
  vector<pt> v(n);
  for (int i = 0; i < n; i++)
   cin >> v[i].x >> v[i].y;
  circle ans = welzl(v);
  cout << fixed << setprecision(3) << ans.c.x << " " << ans.c.y << endl;
  cout << fixed << setprecision(3) << sqrt(ans.r) << endl;</pre>
  return 0;
// acmicpc.net/problem/2626
// achar uma circuferencia
// minimizando o raio
// que cobre todos os pontos dela
// ai oq tem q printar eh o centro dessa circuferencia e o raio
// Minimum enclosing circle
// Welzl's algorithm
// complexidade O(n)
```

# 4 Graph

## 4.1 articulation points

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 400005
#define mod 1000000007
int n, m, timer;
vector<int> adj[MAXN];
bool is_cutpoint[MAXN];
int tin[MAXN];
int low[MAXN];
bool vis[MAXN];
void dfs(int v, int p)
  vis[v] = true;
  tin[v] = timer, low[v] = timer++;
  int childs = 0;
  for (auto const &u : adj[v])
```

```
if (u == p)
      continue;
    if (vis[u])
      low[v] = min(low[v], tin[u]);
    else
     dfs(u, v);
low[v] = min(low[v], low[u]);
if (low[u] >= tin[v] && p != -1)
        is_cutpoint[v] = true;
      childs++;
 if (p == -1 \&\& childs > 1)
    is_cutpoint[v] = true;
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 cin >> n >> m;
 for (int i = 0; i < m; i++)
    int a, b;
   cin >> a >> b;
    a--, b--;
   adj[a].pb(b);
   adj[b].pb(a);
 memset(tin, -1, sizeof(tin));
 memset(low, -1, sizeof(low));
for (int i = 0; i < n; i++)</pre>
   if (!vis[i])
      dfs(i, -1);
 return 0;
```

#### 4.2 BFS

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define DEBUG 1
#define MAXN 1001
#define mod 1000000007
int n, m;
vector<int> adj[MAXN];
bool visited[MAXN];
void bfs(int s)
  queue<int> q;
  q.push(s);
  while (!q.empty())
    int v = q.front();
    q.pop();
```

```
if (visited[v])
      continue;
    visited[v] = true;
    for (auto const &u : adj[v])
     if (!visited[u])
        q.push(u);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> m;
  for (int i = 0; i < m; i++)</pre>
   int a, b, c;
    cin >> a >> b >> c;
   a--, b--;
   adj[a].pb(b);
   adj[b].pb(a);
  bfs(0);
```

## 4.3 bipartite

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
int n, m;
vector<int> adj[MAXN];
bool is()
  vector<int> c(n, -1);
  bool is = 1;
  queue<int> q;
  for (int st = 0; st < n; st++)
    if (c[st] == -1)
      q.push(st);
      c[st] = 0;
      while (!q.empty())
        int v = q.front();
        q.pop();
        for (int u : adj[v])
          if (c[u] == -1)
            c[u] = c[v] ^1;
            q.push(u);
          else
            is &= (c[u] != c[v]);
```

### 4.4 block-cut-tree

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
struct dsu
  vector<pi> parent;
  vector<int> rank;
  vector<int> bipartite;
  void reset(int v)
    parent[v] = \{v, 0\};
    rank[v] = 0;
    bipartite[v] = 1;
  dsu(int n)
    parent.resize(n);
    rank.resize(n);
    bipartite.resize(n);
    for (int v = 0; v < n; v++)
      reset(v);
  dsu() {}
  pi find_set(int v)
    if (v != parent[v].fir)
      int parity = parent[v].sec;
      parent[v] = find_set(parent[v].fir);
      parent[v].sec ^= parity;
    return parent[v];
```

```
void add_edge(int a, int b)
   pi pa = find_set(a);
    a = pa.fir;
   int x = pa.sec;
   pi pb = find_set(b);
    b = pb.fir;
    int y = pb.sec;
   if (a == b)
      if (x == y)
        bipartite[a] = 0;
    else
      if (rank[a] < rank[b])</pre>
        swap(a, b);
      parent[b] = {a, x ^ y ^ 1};
bipartite[a] &= bipartite[b];
      if (rank[a] == rank[b])
        rank[a]++;
 bool is_bipartite(int v)
   return bipartite[find_set(v).fir];
};
struct block_cut_tree
  // Source: https://github.com/brunomaletta/Biblioteca/blob/master/Codigo/
      Grafos/blockCutTree.cpp
  // Cria a block-cut tree, uma arvore com os blocos
  // e os pontos de articulação
  // Blocos sao componentes 2-vertice-conexos maximais
  // Uma 2-coloracao da arvore eh tal que uma cor sao
  // os blocos, e a outra cor sao os pontos de art.
  // Funciona para grafo nao conexo
  // art[i] responde o numero de novas componentes conexas
  // criadas apos a remocao de i do grafo g
  // Se art[i] >= 1, i eh ponto de articulação
  // Para todo i <= blocks.size()</pre>
  // blocks[i] eh uma componente 2-vertce-conexa maximal
  // edgblocks[i] sao as arestas do bloco i
  // tree[i] eh um vertice da arvore que corresponde ao bloco i
  // tree - eh a propia block-cut tree
  // pos[i] responde a qual vertice da arvore vertice i pertence
  // Arvore tem no maximo 2n vertices
  // O(n + m)
 vector<vector<int>> g, blocks, tree;
 vector<vector<pi>>> edgblocks;
  stack<int> s;
 stack<pi> s2;
 vector<int> id, art, pos;
 block_cut_tree(vector<vector<int>> q_) : q(q_)
   int n = q.size():
   id.resize(n, -1), art.resize(n), pos.resize(n);
   build();
  int dfs (int i, int &t, int p = -1)
   int lo = id[i] = t++;
   s.push(i);
    if (p != −1)
      s2.emplace(i, p);
    for (int j : q[i])
      if (j != p \text{ and } id[j] != -1)
        s2.emplace(i, j);
    for (int j : g[i])
```

```
if (j != p)
        if (id[j] == -1)
          int val = dfs(j, t, i);
          lo = min(lo, val);
          if (val >= id[i])
            art[i]++;
            blocks.emplace_back(1, i);
            while (blocks.back().back() != j)
              blocks.back().pb(s.top());
              s.pop();
            edgblocks.emplace_back(1, s2.top());
            s2.pop();
            pi aux = {j, i};
            while (edgblocks.back().back() != aux)
              edgblocks.back().pb(s2.top());
              s2.pop();
          // if (val > id[i]) aresta i-j eh ponte
        else
          lo = min(lo, id[j]);
    if (p == -1 and art[i])
     art[i]--;
    return lo;
  void build()
    int t = 0;
    for (int i = 0; i < q.size(); i++)</pre>
     if (id[i] == -1)
        dfs(i, t, -1);
    tree.resize(blocks.size());
    for (int i = 0; i < g.size(); i++)</pre>
     if (art[i])
        pos[i] = tree.size(), tree.emplace_back();
    for (int i = 0; i < blocks.size(); i++)</pre>
     for (int j : blocks[i])
        if (!art[j])
         pos[j] = i;
          tree[i].pb(pos[j]), tree[pos[j]].pb(i);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n, m;
 cin >> n >> m;
 vector<vector<int>> adj(n);
 for (int i = 0; i < m; i++)</pre>
   int a, b;
   cin >> a >> b;
   a--, b--;
   adj[a].pb(b);
   adj[b].pb(a);
```

```
block_cut_tree bt(adj);
 vector<vector<int>> g(n);
  dsu d(n);
  for (auto const &v : bt.edgblocks)
    vector<int> guys;
    for (auto const & j : v)
      guys.pb(j.fir);
      guys.pb(j.sec);
      d.add_edge(j.fir, j.sec);
   bool bip = 1;
    for (auto const &j : guys)
     bip &= d.is_bipartite(j);
    if (bip)
      for (auto const & i : v)
        g[j.fir].pb(j.sec);
        g[j.sec].pb(j.fir);
    for (auto const &j : guys)
      d.reset(j);
  vector<bool> vis(n, 0);
  vector<bool> c(n, 0);
  int a = 0, b = 0;
  for (int i = 0; i < n; i++)
   if (vis[i])
     continue;
    int x = 1, y = 0;
    queue<int> q;
    q.push(i);
    vis[i] = 1;
    while (!q.empty())
      int k = q.front();
      q.pop();
      for (auto const &i : q[k])
        if (!vis[i])
          vis[i] = 1;
          c[i] = c[k] ^1;
          (c[i] == 1) ? y++ : x++;
          q.push(i);
      }
   a += (x * (x - 1)) / 2;
   a += (y * (y - 1)) / 2;
   b += (x * y);
 cout << a << " " << b << endl;
 return 0;
// https://codeforces.com/gym/103934/problem/M
// pares (a, b) com a < b
// contar pares (a, b) tal que todos os caminhos de a para b possuem distancia
// contar pares (a, b) tal que todos os caminhos de a para b possuem distancia
    par
// grafo biconexo (ou 2-vertice-conexo) - nao tem ponto de articulação
// blocos - sao subgrafos biconexos maximais (sem ponto de articulação)
// block graph
// grafo que tem um vertice para cada bloco do grafo G
// e uma aresta entre dois vertices tal que os blocos correspondentes tem um
    vertice em comum
// block-cut tree
```

```
// um ponto de articulação eh um vertice que esta em dois ou mais blocos
// a estrutura dos blocos e dos pontos de articulação de um grafo conectado pode
     ser descrita por uma arvore chamada de arvore de block-cut tree
// essa arvore tem um vertice para cada bloco e para cada ponto de articulacao
    do grafo dado.
// tem uma aresta na block-cut tree para cada par (bloco, ponto de articulacao),
     tal que esse ponto de articulação ta no bloco
// para o problema:
// para um grafo nao bipartido que e biconexo, tem caminhos de tamanho impar e
    par entre qualquer par de vertices
// um caminho em um grafo G, tem meio que um caminho equivalente na sua block-
// da pra pensar em resolver para cada bloco
// resolvendo pra cada bloco:
// o bloco tem que ser bipartido
// quando o bloco nao en bipartido, eu nao considero as arestas dele
// considerando o grafo restante sendo bipartido
// da pra resolver pra cada componente conexa
// caminhos entre vertices de mesma cor tem paridade impar
// caminhos entre vertices de cor diferente tem paridade par
// https://codeforces.com/gym/102512/problem/A
// ter queries do tipo
// quantos pontos de articulação desconectam u e v
// dai monta a block cut tree
// para cada ponto de articulacao, seta a pos[i] dele como 1 na arvore
// e o valor dos demais vertices como 0
// dai responde uma query com hld (ou com lca tbm sai)
```

## 4.5 bridges

```
#include <bits/stdc++.h>
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 400005
#define mod 1000000007
int n, m, timer;
vector<pi> edges;
vector<bool> is_bridge;
vector<pi> adj[MAXN];
int tin[MAXN];
int low[MAXN];
bool vis[MAXN];
void dfs(int v, int p)
  vis[v] = true;
  tin[v] = timer, low[v] = timer++;
  for (auto const &u : adj[v])
   if (u.fir == p)
      continue;
    if (vis[u.fir])
      low[v] = min(low[v], tin[u.fir]);
      continue;
```

```
dfs(u.fir, v);
   low[v] = min(low[v], low[u.fir]);
    if (low[u.fir] > tin[v])
     is_bridge[u.sec] = 1;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> m;
  is_bridge.resize(m);
  for (int i = 0; i < m; i++)
   int a, b;
   cin >> a >> b;
    a--, b--;
    edges.pb({a, b});
   adj[a].pb({b, i});
   adj[b].pb({a, i});
 memset(tin, -1, sizeof(tin));
 memset(low, -1, sizeof(low));
  for (int i = 0; i < n; i++)
   if (!vis[i])
     dfs(i, -1);
  return 0;
```

### 4.6 caminhoeuleriano

```
// caminho euleriano em um grafo
// passa por todas as arestas apenas uma unica vez e percorre todas elas
// condicao de existencia:
// todos os vertices possuem grau par (ciclo euleriano) comeca e acaba no mesmo
    vertice
// ou
// apenas 2 vertices possuem grau impar, todos os outros possuem grau par ou ==
    0.
// comeca num vertice de grau impar e termina num vertice de grau impar nesse
    caso.
// solucao:
// rodar um dfs com map de visited para as arestas
// no final por o source no vector path
// ao final teremos o caminho inverso no vector path
// note que o caminho inverso tambem e um caminho valido
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define in insert
#define pi pair<int, int>
#define pd pair<double, int>
#define pib pair<pi, bool>
#define mp make_pair
#define fir first
#define sec second
#define MAXN 10001
#define MAXL 1000001
#define mod 1000000007
int n, m, start;
vector<int> path;
vector<int> adj[MAXN];
map<pi, bool> visited;
void dfs(int s)
  for (int i = 0; i < adj[s].size(); i++)</pre>
    int v = adj[s][i];
    if (!visited[mp(s, v)])
```

```
visited[mp(s, v)] = true;
      visited[mp(v, s)] = true;
      dfs(v);
  path.pb(s);
bool check()
  int odd = 0;
 for (int i = 0; i < n; i++)</pre>
   if (adj[i].size() & 1)
      odd++, start = i;
  return (odd == 0 || odd == 2);
signed main()
  cin >> n >> m;
  for (int i = 0; i < m; i++)
    int a, b;
    cin >> a >> b;
   adj[a].pb(b);
    adj[b].pb(a);
  start = 0;
  bool ok = check();
  (ok) ? cout << "Yes\n" : cout << "No\n";</pre>
  if (ok)
    dfs(start);
    for (int i = 0; i < path.size(); i++)</pre>
     cout << path[i] << " ";
    cout << "\n";
  return 0;
```

### 4.7 caminhoeuleriano2

if (!vis[v.sec])

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500005
#define mod 998244353
int deg[MAXN];
bool vis[MAXN];
vector<int> path;
vector<pi> edges, edges2;
vector<pi> ans;
vector<pi> adi[MAXN]:
vector<pi> a[MAXN];
void dfs2(int s)
  while (a[s].size() > 0)
    auto v = a[s].back();
    a[s].pop_back();
```

```
vis[v.sec] = 1;
      dfs2(v.fir);
  path.pb(s);
void dfs(int i)
  vis[i] = 1;
  for (auto const &j : adj[i])
    if (!vis[j.fir])
      dfs(j.fir);
      if (deg[j.fir])
       ans.pb(edges[j.sec]);
deg[j.fir] ^= 1;
        deg[i] ^= 1;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int q;
  cin >> q;
  while (q--)
    int n, m;
    cin >> n >> m;
    edges.clear();
    edges2.clear();
    ans.clear();
    for (int i = 0; i < n; i++)
      adj[i].clear();
      vis[i] = 0;
      deg[i] = 0;
    for (int i = 0; i < m; i++)
      int a, b, c;
      cin >> a >> b >> c;
      a--, b--;
      if (c == 1) // edges obrigatorios
        edges2.pb({a, b});
        deg[a] ^= 1;
deg[b] ^= 1;
      else // edges nao obrigatorios
        edges.pb({a, b});
        adj[a].pb({b, edges.size() - 1});
        adj[b].pb({a, edges.size() - 1});
    // ajeita pra ver se consegue fazer todo mundo ficar com grau par
    // considerando so os edges nao obrigatorios
    for (int i = 0; i < n; i++)
      if (!vis[i])
        dfs(i);
   bool ok = 1;
    for (int i = 0; i < n; i++)
      if (deg[i])
        ok = 0;
    if (!ok)
      cout << "NO\n";
      continue;
```

```
for (int i = 0; i < n; i++)
     a[i].clear();
    // se ajeitei, agora dale
    // monta o grafo final e acha o ciclo euleriano
    // funciona para grafos com self loops e multiple edges
    int id = 0;
    for (auto [u, v] : ans)
     a[u].pb({v, id});
     a[v].pb({u, id});
     id++;
    for (auto [u, v] : edges2)
     a[u].pb({v, id});
     a[v].pb({u, id});
     id++;
    for (int i = 0; i < m; i++)
     vis[i] = 0;
    path.clear();
    dfs2(0);
   cout << "YES\n";</pre>
    cout << path.size() - 1 << endl;</pre>
    for (int i = 0; i < path.size(); i++)</pre>
     cout << path[i] + 1 << " ";
    cout << endl;
 return 0;
// https://codeforces.com/contest/1994/problem/F
// dado um grafo, tem edges que sao obrigatorios de manter
// e outros q posso remover
// quero fazer com que um grafo tenha um ciclo euleriano
// no qual o grau de cada vertice eh par
// se tiver solucao, eu quero imprimir o ciclo euleriano
```

## 4.8 centroid decomposition

```
if (heaviest_child == -1 || subtree_size[v] > subtree_size[
                                                                                              heaviest child])
// centroid de uma arvore -> e um no que ao ser removido da arvore, separaria as
                                                                                            heaviest_child = v;
// arvores resultantes de modo com que a maior arvore desse conjunto teria no
// (n / 2) nos, sendo n o numero de nos da arvore. Para qualquer arvore com n
                                                                                       return (is_centroid && sz - subtree_size[s] <= sz / 2) ? s : getCentroid(
                                                                                           heaviest_child, s);
// o centroid sempre existe.
                                                                                     int decompose_tree(int s)
    dfs(s, s);
                                                                                       int cend_tree = getCentroid(s, s);
// centroid decomposition -> muito util para tentar diminuir a complexidade em
                                                                                       visited[cend_tree] = true;
                                                                                       for (auto const &v : adj[cend_tree])
// tipos de consultas a serem feitas, uma maneira melhor de organizar a arvore.
                                                                                        if (!visited[v])
// algoritimo:
// 1) o centroid e a raiz dessa nova arvore
                                                                                          int cend_subtree = decompose_tree(v);
// 2) achar o centroid das arvores menores que surgiram com a remocao do
                                                                                          adjl[cend_tree].pb(cend_subtree);
    centroid "pai"
                                                                                          adjl[cend_subtree].pb(cend_tree);
// 3) por uma aresta entre o centroid "filho" e o centroid "pai"
                                                                                          father[cend_subtree] = cend_tree;
// 4) repetir isso ate todos os nos serem removidos
// 5) ao final teremos a centroid tree
                                                                                      return cend_tree;
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
                                                                                     void init()
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
                                                                                      subtree_size.resize(n);
using namespace __gnu_pbds;
                                                                                      visited.resize(n);
                                                                                      father.assign(n, -1);
template <class T>
                                                                                      decompose_tree(0);
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
                                                                                   signed main()
#define PI acos (-1)
```

#define pb push\_back

#define fir first

#define DEBUG 0

namespace cd

int sz;

int n;

#define sec second

#define MAXN 100001

#define mod 1000000007

vector<int> adj[MAXN];

vector<int> adjl[MAXN];

vector<bool> visited;

void dfs(int s, int f)

dfs(v, s);

subtree\_size[s] = 1;

vector<int> father, subtree\_size;

for (auto const &v : adj[s])

int getCentroid(int s, int f)

bool is\_centroid = true;

int heaviest child = -1;

for (auto const &v : adj[s])
{
 if (v != f && !visited[v])
 {
 if (subtree\_size[v] > sz / 2)

is\_centroid = false;

if (v != f && !visited[v])

subtree\_size[s] += subtree\_size[v];

#define int long long int

#define pi pair<int, int>
#define pii pair<int, pi>

```
{
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n;
  for (int i = 0; i < n - 1; i++)
  {
    int a, b;
    cin >> a >> b;
    a--, b--;
    adj[a].pb(b);
    adj[b].pb(a);
  }
  cd::init();
  return 0;
}
```

# 4.9 centroid decomposition2

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 998244353
int n, k, resp;
vector<int> adj[MAXN];
vector<int> cnt;
namespace cd
 vector<int> subtree_size;
 vector<bool> visited;
 void dfs(int s, int f)
   sz++:
    subtree_size[s] = 1;
    for (auto const &v : adj[s])
      if (v != f && !visited[v])
        subtree_size[s] += subtree_size[v];
  int get_centroid(int s, int f)
    bool is centroid = true;
    int heaviest child = -1;
    for (auto const &v : adj[s])
      if (v != f && !visited[v])
        if (subtree_size[v] > sz / 2)
         is_centroid = false;
        if (heaviest_child == -1 || subtree_size[v] > subtree_size[
            heaviest childl)
          heaviest_child = v;
    return (is_centroid && sz - subtree_size[s] <= sz / 2) ? s : get_centroid(
        heaviest child, s);
  void dfs2(int s, int f, int d)
    while (d >= cnt.size())
```

```
cnt.pb(0);
    cnt[d]++;
    for (auto const &v : adj[s])
      if (v != f && !visited[v])
        dfs2(v, s, d + 1);
  void solve(int s)
    vector<int> tot;
    for (auto const &v : adj[s])
      if (visited[v])
        continue;
      cnt.clear();
      dfs2(v, s, 1);
      for (int i = 1; i < cnt.size(); i++)</pre>
        if (k - i < tot.size() && k - i >= 1)
          resp += (cnt[i] * tot[k - i]);
      for (int i = 1; i < cnt.size(); i++)</pre>
        while (i >= tot.size())
          tot.pb(0);
        tot[i] += cnt[i];
    if (k < tot.size())</pre>
      resp += tot[k];
  int decompose_tree(int s)
    sz = 0:
    dfs(s, s);
    int cend_tree = get_centroid(s, s);
    visited[cend_tree] = true;
    solve(cend_tree);
    for (auto const &v : adj[cend_tree])
      if (!visited[v])
        decompose_tree(v);
    return cend tree:
  void init()
    subtree_size.resize(n);
    visited.resize(n);
    decompose_tree(0);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> k;
  for (int i = 1; i < n; i++)
    int a, b;
    cin >> a >> b;
    a--, b--;
   adj[a].pb(b);
   adj[b].pb(a);
  cd::init();
  cout << resp << endl;</pre>
  return 0;
// https://codeforces.com/contest/161/problem/D
// problema: contar quantos pares de vertices (u, v) existem tal que dist(u, v)
// durante a decomposicao
// pega o centroid atual e resolve o problema pra ele
// isso eh:
// para cada centroid que eu achei, devo contar quantos caminhos
// de tamanho k passam por esse centroid
```

// somando todas essas respostas, a gente tem a resposta final

## 4.10 cycle detection

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 205
#define MAXP 100001
#define mod 1000000007
int n, m, idx;
vector<int> cycles[MAXN];
vector<int> adj[MAXN];
int color[MAXN];
int parent[MAXN];
int ans[MAXN];
void dfs(int u, int p)
  if (color[u] == 2)
   return:
  if (color[u] == 1)
    idx++;
    int curr = p;
    ans[curr] = idx;
    cycles[idx].pb(curr);
    while (curr != u)
      curr = parent[curr];
      cycles[idx].pb(curr);
      ans[curr] = idx;
    return;
  parent[u] = p;
  color[u] = 1;
  for (auto const &v : adj[u])
    if (v != parent[u])
     dfs(v, u);
  color[u] = 2;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> m;
  for (int i = 0; i < m; i++)</pre>
   int a, b;
   cin >> a >> b;
    a--, b--;
    adj[a].pb(b);
    adj[b].pb(a);
  for (int i = 0; i < n; i++)
   if (!color[i])
     dfs(i, -1);
  cout << idx << endl;</pre>
  for (int i = 1; i <= idx; i++)</pre>
    cout << cycles[i].size() << endl;</pre>
    for (auto const &j : cycles[i])
```

```
cout << j + 1 << " ";
cout << endl;
}
return 0;</pre>
```

#### 4.11 DFS

```
#include <bits/stdc++.h>
using namespace std;
#define MAXN 500000
int n , m ;
int visited [MAXN] ;
vector <int> adj_list [MAXN] ;
void dfs (int x)
    for (int i = 0 ; i < adj_list[x].size() ; i++)</pre>
        int v = adj_list[x][i] ;
        if(visited[v] == -1)
            visited[v] = visited[x];
            dfs(v);
void initialize ()
    for (int i = 1; i \le n; i++)
        visited[i] = -1;
int main ()
    int a , b ;
    cin >> n >> m ;
    initialize();
    for (int i = 1; i \le m; i++)
        cin >> a >> b;
        adj_list[a].push_back(b);
        adj_list[b].push_back(a);
    dfs(1);
    return 0;
```

# 4.12 Dijkstra

```
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define DEBUG 1
#define MAXN 1001
#define mod 1000000007
int n, m;
vector<pi> adj[MAXN];
bool visited[MAXN];
int dist[MAXN];
void dijkstra(int s)
 for (int i = 0; i < n; i++)</pre>
    dist[i] = INT_MAX;
visited[i] = false;
  priority_queue<pi, vector<pi>, greater<pi>> q;
  dist[s] = 0;
  q.push({dist[s], s});
  while (!q.empty())
    int v = q.top().second;
    q.pop();
if (visited[v])
      continue;
    visited[v] = true;
    for (auto const &u : adj[v])
      if (dist[u.sec] > dist[v] + u.fir)
        dist[u.sec] = dist[v] + u.fir;
        q.push({dist[u.sec], u.sec});
    }
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> m;
  for (int i = 0; i < m; i++)
    int a, b, c;
    cin >> a >> b >> c;
    a--, b--;
    adj[a].pb({c, b});
    adj[b].pb({c, a});
  dijkstra(0);
```

## 4.13 dinic

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 705
```

```
#define mod 1000000007
#define INF 1e9
struct edge
  int to, from, flow, capacity, id;
struct dinic
  int n, src, sink;
  vector<vector<edge>> adj;
  vector<int> level;
  vector<int> ptr;
  dinic(int sz)
   n = sz;
    adj.resize(n);
    level.resize(n);
   ptr.resize(n);
  void add_edge(int a, int b, int c, int id)
    adj[a].pb({b, (int)adj[b].size(), c, c, id});
   adj[b].pb({a, (int)adj[a].size() - 1, 0, 0, id});
  bool bfs()
    level.assign(n, -1);
    level[src] = 0;
    queue<int> q;
    q.push(src);
    while (!q.empty())
      int u = q.front();
      q.pop();
      for (auto at : adj[u])
        if (at.flow && level[at.to] == -1)
          q.push(at.to);
          level[at.to] = level[u] + 1;
    return level[sink] != -1;
  int dfs(int u, int flow)
    if (u == sink || flow == 0)
      return flow;
    for (int &p = ptr[u]; p < adj[u].size(); p++)</pre>
      edge &at = adj[u][p];
      if (at.flow && level[u] == level[at.to] - 1)
        int kappa = dfs(at.to, min(flow, at.flow));
        at.flow -= kappa;
        adj[at.to][at.from].flow += kappa;
        if (kappa != 0)
          return kappa;
    return 0;
  int run()
    int max_flow = 0;
    while (bfs())
      ptr.assign(n, 0);
      while (1)
        int flow = dfs(src, INF);
       if (flow == 0)
         break;
        max_flow += flow:
```

```
return max_flow;
  vector<pii> cut edges() // arestas do corte minimo
   bfs();
    vector<pii> ans;
    for (int i = 0; i < n; i++)
      for (auto const &j : adj[i])
        if (level[i] != -1 && level[j.to] == -1 && j.capacity > 0)
          ans.pb({j.capacity, {i, j.to}});
   return ans:
  vector<int> flow_edges(int n, int m) // fluxo em cada aresta, na ordem da
       entrada
    vector<int> ans(m);
    for (int i = 0; i < n; i++)
      for (auto const &j : adj[i])
        if (!j.capacity)
          ans[j.id] = j.flow;
   return ans:
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n, m;
 cin >> n >> m;
 dinic d(n);
for (int i = 0; i < m; i++)</pre>
   int a, b, c;
   cin >> a >> b >> c;
   a--, b--;
   d.add_edge(a, b, c, i);
 d.src = 0, d.sink = n - 1;
 cout << d.run() << endl;</pre>
 vector<int> ans = d.flow_edges(n, m);
 for (auto const &i : ans)
   cout << i << endl;</pre>
 return 0;
```

### 4.14 dominator tree

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 998244353
struct dominator_tree
  int n, t;
```

```
vector<vector<int>> g, tree, rg, bucket;
vector<int> dfs_1, dfs_r, idom, sdom, prv, pre, ancestor, label, preorder;
dominator tree(vector<vector<int>> &adj, int source)
  int n = adj.size();
  g = adj;
  tree.resize(n);
  rq.resize(n);
  bucket.resize(n);
  dfs_l.resize(n);
  dfs_r.resize(n);
  idom.resize(n);
  sdom.assign(n, -1);
  prv.resize(n);
  pre.assign(n, -1);
  ancestor.assign(n, -1);
  label.resize(n);
  build(source);
void dfs(int v)
 pre[v] = ++t;
  sdom[v] = label[v] = v;
  preorder.pb(v);
  for (auto const &nxt : g[v])
    if (sdom[nxt] == -1)
      prv[nxt] = v;
      dfs(nxt);
    rg[nxt].pb(v);
int eval(int v)
 if (ancestor[v] == -1)
   return v;
  if (ancestor[ancestor[v]] == -1)
   return label[v];
  int u = eval(ancestor[v]);
  if (pre[sdom[u]] < pre[sdom[label[v]]])</pre>
   label[v] = u;
  ancestor[v] = ancestor[u];
  return label[v];
void dfs2(int v)
  dfs_l[v] = t++;
  for (auto const &nxt : tree[v])
    dfs2(nxt);
  dfs_r[v] = t++;
void build(int s)
  t = 0:
  dfs(s);
  if (preorder.size() == 1)
    return;
  int sz = preorder.size();
  for (int i = sz - 1; i >= 1; i --)
    int w = preorder[i];
    for (auto const &v : rg[w])
      int u = eval(v);
      if (pre[sdom[u]] < pre[sdom[w]])</pre>
        sdom[w] = sdom[u];
    bucket[sdom[w]].push_back(w);
    ancestor[w] = prv[w];
    for (auto const &v : bucket[prv[w]])
      int u = eval(v);
```

```
idom[v] = (u == v) ? sdom[v] : u;
      bucket[prv[w]].clear();
    for (int i = 1; i < preorder.size(); i++)</pre>
      int w = preorder[i];
     if (idom[w] != sdom[w])
  idom[w] = idom[idom[w]];
      tree[idom[w]].push_back(w);
    idom[s] = sdom[s] = -1;
    t = 0;
   dfs2(s);
 bool dominates(int u, int v)
    if (pre[v] == -1)
     return 1:
    return dfs_1[u] <= dfs_1[v] && dfs_r[v] <= dfs_r[u];</pre>
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
  int n, m;
 cin >> n >> m;
 vector<vector<int>> adj(n);
 for (int i = 0; i < m; i++)
    int a, b;
   cin >> a >> b;
    a--, b--;
   adj[a].pb(b);
 dominator_tree d(adj, 0);
  vector<int> ans;
 for (int i = 0; i < n; i++)
    if (d.dominates(i, n - 1))
     ans.pb(i);
 cout << ans.size() << endl;</pre>
 for (int i = 0; i < ans.size(); i++)
cout << ans[i] + 1 << " \n"[i == ans.size() - 1];</pre>
// https://tanujkhattar.wordpress.com/2016/01/11/dominator-tree-of-a-directed-
// https://cses.fi/problemset/task/1703/ (problema desse codigo)
// https://codeforces.com/gym/100513/problem/L
// https://codeforces.com/contest/757/problem/F
// dado um vertice source s
// dizemos que u domina w, se todos os caminhos de
// s ate w passam pelo vertice u
// dizemos que u e um dominador imediato de w se u domina w
// e todos os demais dominadores de w, dominam u
// 1 - todo vertice (tirando o source) tem um dominador
// pois o source domina todos os demais vertices
// 2 - todo vertice (tirando o source) tem exatamente um
// unico dominador imediato
// se eu crio um grafo com todas as arestas do tipo
// (dominador imediato de w) - w
// para todos os vertices w que nao sao a source
// esse grafo eh uma arvore
// e eh a dominator tree
```

#### 4.15 dsu

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
```

```
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 2001
#define mod 1000000007
struct dsu
  int tot:
  vector<int> parent;
  vector<int> sz;
  dsu(int n)
    parent.resize(n);
    sz.resize(n);
    tot = n;
    for (int i = 0; i < n; i++)
     parent[i] = i;
      sz[i] = 1;
  int find_set(int i)
    return parent[i] = (parent[i] == i) ? i : find_set(parent[i]);
  void make_set(int x, int y)
    x = find_set(x), y = find_set(y);
    if (x != y)
      if (sz[x] > sz[y])
       swap(x, y);
      parent[x] = y;
      sz[y] += sz[x];
      tot--:
  }
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  dsu d(n);
  int a, b;
  cin >> a >> b;
  d.make_set(a, b);
  d.find_set(a);
```

### 4.16 dsu rollback

```
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 600005
#define mod 1000000007
namespace dsu
  struct rollback
    int u, v, ranku, rankv;
  } ;
  int num_sets;
  int parent[MAXN];
  int rank[MAXN];
  stack<rollback> op;
  int Find(int i)
    return (parent[i] == i) ? i : Find(parent[i]);
  bool Union(int x, int y)
    int xx = Find(x);
    int yy = Find(y);
    if (xx != yy)
      num sets--:
      if (rank[xx] > rank[yy])
        swap(xx, yy);
      op.push({xx, yy, rank[xx], rank[yy]});
      parent[xx] = yy;
      if (rank[xx] == rank[yy])
        rank[yy]++;
      return true;
    return false;
  void do rollback()
    if (op.empty())
     return;
    rollback x = op.top();
    op.pop();
    num_sets++;
    parent[x.v] = x.v;
    rank[x.v] = x.rankv;
    parent[x.u] = x.u;
    rank[x.u] = x.ranku;
  void init(int n)
    for (int i = 0; i < n; i++)
      parent[i] = i;
      rank[i] = 0;
    num\_sets = n;
namespace seg
  struct query
    int v, u, is_bridge;
  };
  vector<vector<query>> t(4 * MAXN);
  int ans[MAXN];
  void add(int i, int l, int r, int ql, int qr, query q)
    if (1 > r || 1 > qr || r < q1)
```

```
return;
    if (1 >= q1 && r <= qr)
      t[i].push_back(q);
      return;
    int mid = (1 + r) >> 1;
   add((i << 1), 1, mid, ql, qr, q);
   add((i << 1) | 1, mid + 1, r, ql, qr, q);
  void dfs(int i, int 1, int r)
    for (query &q : t[i])
      if (dsu::Union(q.v, q.u))
        q.is_bridge = 1;
    if (1 == r)
     ans[1] = dsu::num_sets;
    else
      int mid = (1 + r) >> 1;
      dfs((i << 1), 1, mid);
      dfs((i << 1) | 1, mid + 1, r);
    for (query q : t[i])
      if (q.is_bridge)
        dsu::do_rollback();
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, q;
  cin >> n >> q;
  int time = 0;
  map<pi, int> tin;
  vector<int> queries;
  while (q--)
    char t;
    cin >> t;
    if (t == '?')
      queries.pb(++time);
    else if (t == '+')
      int a, b;
      cin >> a >> b;
      a--, b--;
      if (a > b)
        swap(a, b);
      tin[{a, b}] = ++time;
    else
      int a, b;
      cin >> a >> b;
      a--, b--;
      if (a > b)
       swap(a, b);
      seg::query kappa = {a, b, 0};
      seg::add(1, 0, MAXN - 1, tin[{a, b}], ++time, kappa);
      tin[{a, b}] = -1;
  for (auto const &i : tin)
   if (i.sec != -1)
      seg::query kappa = {i.fir.fir, i.fir.sec, 0};
      seg::add(1, 0, MAXN - 1, i.sec, ++time, kappa);
  dsu::init(n);
  seg::dfs(1, 0, MAXN - 1);
  for (auto const &i : queries)
    cout << seg::ans[i] << endl;</pre>
```

```
return 0;
// https://codeforces.com/edu/course/2/lesson/7/3/practice/contest/289392/
    problem/C
// conectividade dinamica
// para uma query (u, v)
// podemos descrever em um intervalo [1, r]
// 1 = quando a aresta (u, v) foi adicionada
// r = quando a aresta (u, v) foi removida
// dai agora que temos um intervalo, podemos adicionar
// a query (u, v) em uma segtree "adaptada"
// no final rodamos um dfs nessa segtree e vamos atualizando as repostas das
    queries
// quando estamos em uma posicao na seg, dou union em todos os caras daquela
// e em seguida chamo pros meus filhos, quando chego em uma folha, ela eh
    equivalente
// a uma unidade de "tempo", logo a resposta para aquele tempo eh a resposta
    atual no dsu
// e ao sair recursivamente, vou dando rollbacks no dsu
```

## 4.17 erdos gallai

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define pci pair<char, int>
#define fir first
#define sec second
#define MAXN 100005
#define mod 998244353
bool erdos_gallai(vector<int> &v)
  int sum = 0, n = v.size();
  for (auto const &i : v)
    sum += i;
  if (sum % 2)
    return false;
  sort(v.rbegin(), v.rend());
  vector<int> suf(n + 1, 0);
  int qt = 0, ptr = n;
  sum = 0;
  for (int i = n - 1; i >= 0; i--)
    if (v[i] >= i)
      qt++;
    else
      sum += v[i];
      ptr = i;
    while (ptr < n \&\& v[ptr] >= i)
      at++;
      sum -= v[ptr];
      ptr++;
    suf[i] = sum + (qt * i);
  sum = 0;
  bool ok = 1;
  for (int i = 0; i < n; i++)
```

```
sum += v[i];
   int curr = i * (i + 1) + suf[i + 1];
   ok &= (sum <= curr);
  return ok;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int q;
  cin >> q;
  while (q--)
   int n;
   cin >> n;
    vector<int> v(n);
    for (int i = 0; i < n; i++)
     cin >> v[i];
    (erdos_gallai(v)) ? cout << "Y\n" : cout << "N\n";</pre>
  return 0;
// https://codeforces.com/gym/101726/problem/A
// erdos gallai
// dado uma sequencia de n inteiros d[0], d[1], ..., d[n - 1]
// quero saber se existe um grafo simples e undirected com n vertices
// tal que o grau do i-esimo vertice e igual a d[i]
```

#### 4.18 eulertour

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG 1
#define MAXN 100001
#define mod 1000000009
#define d 31
int n, idx;
vector<int> adj[MAXN];
int euler[2 * MAXN];
int entrei[MAXN];
int sai[MAXN];
void euler_tour(int s, int f)
  euler[idx] = s;
  entrei[s] = idx;
  idx++;
  for (auto const &v : adj[s])
    if (v == f)
      continue;
    euler_tour(v, s);
  euler[idx] = s;
  sai[s] = idx;
  idx++;
signed main()
  ios_base::sync_with_stdio(false);
```

```
cin.tie(NULL);
 int n:
  cin >> n;
 for (int i = 0; i < n - 1; i++)
   int a, b;
   cin >> a >> b;
   a--, b--;
adj[a].pb(b);
   adj[b].pb(a);
  euler_tour(0, -1);
  for (int i = 0; i < 2 * n; i++)</pre>
   cout << euler[i] << " ";
  cout << endl;</pre>
 return 0:
// euler tour of a tree
// muito util para algumas coisas
// exemplos:
// 1- soma da subarvore de v(com update)
// usando segment trees, podemos fazer uma query(entrei[v], sai[v])
// lca(u, v) = query(entrei[u], entrei[v])
// usando uma query de minimo e considerando as profundidade dos vertices
// a resposta sera o vertice de profundidade minima que encontrarmos no
    intervalo
// 3- agilidade para remover arestas/vertices/subtrees da arvore
// basta apenas tratar o segmento equivalente do jeito que for necessario
// 4- reroot a tree
// basta apenas rotacionar o euler path
```

## 4.19 flow with minimum capacities

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
#define INF 1e9
struct edge
  int dest, back, f, c, id;
struct push_relabel
 int n;
  vector<vector<edge>> g;
  vector<int> ec;
  vector<edge *> cur;
  vector<vector<int>> hs:
  vector<int> H:
  push_relabel(int sz) : g(sz), ec(sz), cur(sz), hs(2 * sz), H(sz) { n = sz; }
  void add_edge(int s, int t, int cap, int rcap, int id)
    if (s == t)
      return;
    g[s].pb({t, (int)g[t].size(), 0, cap, id});
    g[t].pb({s, (int)g[s].size() - 1, 0, rcap, -1});
```

```
void add_flow(edge &e, int f)
    edge &back = g[e.dest][e.back];
    if (!ec[e.dest] && f)
     hs[H[e.dest]].push_back(e.dest);
    e.f += f;
    e.c -= f;
    ec[e.dest] += f;
    back.f -= f;
    back.c += f;
    ec[back.dest] -= f;
  int calc(int s, int t)
    int v = q.size();
    H[s] = v;
    ec[t] = 1;
    vector<int> co(2 * v);
co[0] = v - 1;
    for (int i = 0; i < v; i++)
     cur[i] = g[i].data();
    for (edge &e : q[s])
      add_flow(e, e.c);
    for (int hi = 0;;)
      while (hs[hi].empty())
        if (!hi--)
          return -ec[s];
      int u = hs[hi].back();
      hs[hi].pop_back();
      while (ec[u] > 0)
        if (cur[u] == g[u].data() + g[u].size())
          H[u] = INF;
           for (edge &e : g[u])
            if (e.c && H[u] > H[e.dest] + 1)
  H[u] = H[e.dest] + 1, cur[u] = &e;
           if (++co[H[u]], !--co[hi] && hi < v)</pre>
             for (int i = 0; i < v; i++)</pre>
               if (hi < H[i] && H[i] < v)
                 --co[H[i]], H[i] = v + 1;
          hi = H[u];
        else if (cur[u] \rightarrow c \&\& H[u] == H[cur[u] \rightarrow dest] + 1)
          add_flow(*cur[u], min(ec[u], cur[u]->c));
        else
           ++cur[u];
  vector<int> flow_edges(int m) // fluxo em cada aresta
    vector<int> ans(m):
    for (int i = 0; i < n; i++)
      for (auto const &j : g[i])
        if (j.id != -1)
          ans[j.id] = j.f;
    return ans;
struct flow with demands
  push relabel pr;
  vector<int> in, out;
  flow_with_demands(int sz) : n(sz), pr(sz + 2), in(sz), out(sz) {}
  void add_edge(int u, int v, int cap, int dem, int id)
    pr.add_edge(u, v, cap - dem, 0, id);
    out[u] += dem, in[v] += dem;
  int run(int s, int t)
```

```
pr.add_edge(t, s, INF, 0, -1);
    for (int i = 0; i < n; i++)
      pr.add_edge(n, i, in[i], 0, -1);
     pr.add_edge(i, n + 1, out[i], 0, -1);
    return pr.calc(n, n + 1);
 bool check() // todas as constraints foram satisfeitas?
    for (auto const &i : pr.g[n])
      if (i.c > 0)
        return 0;
   return 1:
};
int dx[] = \{1, -1, 0, 0\};
int dy[] = \{0, 0, 1, -1\};
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
  int h, w;
 cin >> h >> w;
 vector<string> v(h);
 for (int i = 0; i < h; i++)
   cin >> v[i];
 vector<pi> s[2];
  for (int i = 0; i < h; i++)
    for (int j = 0; j < w; j++)
     if (v[i][j] != '1')
        s[(i + j) % 2].pb({i, j});
  for (int i = 0; i < 2; i++)
    sort(s[i].begin(), s[i].end());
  flow_with_demands \ mf(s[0].size() + s[1].size() + 2);
  int src = s[0].size() + s[1].size();
  int sink = s[0].size() + s[1].size() + 1;
  for (int x = 0; x < s[0].size(); x++)
   int i = s[0][x].fir, j = s[0][x].sec;
if (v[i][j] == '2')
     mf.add_edge(src, x, 1, 1, -1);
    else
     mf.add_edge(src, x, 1, 0, -1);
  for (int x = 0; x < s[1].size(); x++)
    int i = s[1][x].fir, j = s[1][x].sec;
    if (v[i][j] == '2')
     mf.add\_edge(s[0].size() + x, sink, 1, 1, -1);
     mf.add\_edge(s[0].size() + x, sink, 1, 0, -1);
  for (int x = 0; x < s[0].size(); x++)
    for (int d = 0; d < 4; d++)
      pi \ curr = \{s[0][x].fir + dx[d], s[0][x].sec + dy[d]\};
      if (binary_search(s[1].begin(), s[1].end(), curr))
        int y = lower_bound(s[1].begin(), s[1].end(), curr) - s[1].begin();
        mf.add\_edge(x, s[0].size() + y, 1, 0, -1);
 mf.run(src, sink);
  // existe um jeito de passar fluxo que satisfaz todas as constraints?
```

```
(mf.check()) ? cout << "Yes\n" : cout << "No\n";</pre>
  return 0:
  problema exemplo
// https://atcoder.jp/contests/abc285/tasks/abc285_g
// as celulas com 1 eu posso ignorar
// agr pras celulas com 2, eu preciso achar um matching dela com alguem
// so considerando os 2 e as ?
// entao a missao vira achar um matching (nao necessariamente maximo)
// mas que englobe todos os 2
// pode ter 2 de um lado e pode ter 2 do outro
// e se eu pudesse adicionar a sequinte constraint para algumas arestas:
// a quantidade de fluxo passada naquela aresta tem que ser entre [1, r]
// Maximum flow problem with minimum capacities, tambem conhecido como flow with
      demands
// ai da pra dale em resolver
// https://cp-algorithms.com/graph/flow_with_demands.html
```

### 4.20 Floyd Warshall

```
#include <bits/stdc++.h>
using namespace std;
#define pb push_back
#define lli long long int
#define MAXN 10000
#define INF 999999
int n , m , a , b , c ;
int dist [MAXN] [MAXN] ;
void floyd_warshall ()
    for (int k = 0; k < n; k++)
        for (int i = 0; i < n; i++)
           for (int j = 0; j < n; j++)
                dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j]);
void initialize ()
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
           if (i == i)
                dist[i][j] = 0;
           else
               dist[i][j] = INF;
int main()
   cin >> n >> m ;
   initialize ();
    for (int i = 0; i < m; i++)
        cin >> a >> b >> c;
```

```
dist [a][b] = min (dist[a][b], c);
    dist [b][a] = min (dist[b][a], c);
}
floyd_warshall ();
return 0;
```

#### 4.21 Ford Fulkerson

```
// ford-fulkerson: obter qual o fluxo maximo de um vertice s ate um vertice d
// 1 - rodar um bfs para descobrir um novo caminho de s ate d
// 2 - apos isso pego a aresta de menor custo desse caminho e subtraio o valor
    dela nas outras arestas do caminho
// 3 - fluxo_maximo += custo da aresta de menor custo desse caminho
// 4 - rodar isso ate nao existirem mais caminhos disponiveis (com fluxo
    diferente de 0) entre s e d
// 5 - o fluxo maximo de s ate d sera a soma das arestas de menor custo de cada
    caminho feito
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define MAXN 10000
#define INF 999999
int n , m , a , b , c , s , d , max_flow , flow ;
vector <int> parent ;
vector <int> adj [MAXN] ;
int cost [MAXN] [MAXN] ;
bool visited [MAXN] ;
void get_menor_custo (int v , int mincost)
    if (v == s)
        flow = mincost;
        return :
    else if (parent[v] != -1)
        get_menor_custo(parent[v] , min(mincost , cost[parent[v]][v])) ;
        cost[parent[v]][v] -= flow;
        cost[v][parent[v]] += flow ;
void bfs ()
    visited[s] = true ;
    queue <int> q ;
    q.push(s);
    parent assign (MAXN , -1) ;
    while (!q.empty())
        int u = q.front();
        q.pop();
        if (u == d)
            break :
        for (int j = 0 ; j < adj[u].size() ; j++)</pre>
            int v = adj[u][j] ;
            if (cost[u][v] > 0 && !visited[v])
                visited[v] = true ;
                q.push(v);
                parent[v] = u;
```

```
int ford_fulkerson ()
   max_flow = 0;
    while (1)
        flow = 0;
        memset(visited , false , sizeof(visited));
        get_menor_custo(d , INF) ;
        if (flow == 0)
            break ;
        max_flow += flow ;
    return max_flow ;
int main ()
    ios_base::sync_with_stdio(false) ;
    cin.tie(NULL);
    cin >> n >> m ;
    for (int i = 0; i < m; i++)
        cin >> a >> b >> c;
        adj[a].pb(b);
        adj[b].pb(a);
        cost[a][b] = c;
    cin >> s >> d;
    cout << ford fulkerson() << endl ;</pre>
    return 0 ;
```

# 4.22 Grafo Bipartido

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200006
#define mod 1000000007
struct dsu
  vector<pi> parent;
  vector<int> rank;
  vector<int> bipartite;
```

```
dsu(int n)
   parent.resize(n);
    rank.resize(n);
    bipartite.resize(n);
    for (int v = 0; v < n; v++)
     parent[v] = \{v, 0\};
      rank[v] = 0;
     bipartite[v] = 1;
 dsu() {}
 pi find_set(int v)
    if (v != parent[v].fir)
      int parity = parent[v].sec;
     parent[v] = find set(parent[v].fir);
      parent[v].sec ^= parity;
    return parent[v];
 void add_edge(int a, int b)
    pi pa = find_set(a);
    a = pa.fir;
    int x = pa.sec;
    pi pb = find_set(b);
    b = pb.fir;
    int y = pb.sec;
    if (a == b)
     if (x == y)
        bipartite[a] = 0;
    else
     if (rank[a] < rank[b])</pre>
        swap(a, b);
     parent[b] = {a, x ^ y ^ 1};
bipartite[a] &= bipartite[b];
      if (rank[a] == rank[b])
        rank[a]++;
 bool is_bipartite(int v)
    return bipartite[find_set(v).fir];
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 return 0;
```

## 4.23 hall theorem

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;

#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pi pair<int, pi>
#define sec second
#define sec second
#define mod 1000000007
```

```
int cnt[6];
int cnt_mask[1 << 6];</pre>
int tot_mask[1 << 6];</pre>
bool halls()
  for (int mask = 1; mask < (1 << 6); mask++)</pre>
    int x = 0;
    for (int i = 0; i < 6; i++)
      if (mask & (1 << i))
        x += cnt[i];
    if (x < tot mask[mask])</pre>
      return 0;
  return 1;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  string s;
  cin >> s;
  for (auto const &i : s)
   cnt[i - 'a']++;
  int n = s.size();
  vector<int> mask(n, (1 << 6) - 1); // a mask que diz quais chars nao podem</pre>
      aparecer naquela posicao
  int m;
  cin >> m;
  while (m--)
   int i;
    string t;
    cin >> i >> t;
    i--;
    mask[i] = 0:
    for (auto const &j : t)
      mask[i] = (1 << (j - 'a'));
  for (auto const &i : mask)
    cnt_mask[i]++;
  for (int m = 0; m < (1 << 6); m++)
    for (int s = m; s; s = (s - 1) & m) // soma dos cnt_mask de todas as
         submasks
      tot_mask[m] += cnt_mask[s];
  if (!halls())
    cout << "Impossible\n";</pre>
    return 0;
  for (int i = 0; i < n; i++)
    for (int m = 0; m < (1 << 6); m++)
      if ((m & mask[i]) == mask[i])
        tot_mask[m]--;
    for (int j = 0; j < 6; j++)
      if ((mask[i] & (1 << j)) && cnt[j] > 0)
        if (halls()) // faz s[i] = j e ve se o matching continua a existir
          cout << (char) (j + 'a');
          break;
```

```
    cnt[j]++;
    cout << endl;
    return 0;
}
// https://codeforces.com/contest/1009/problem/G
// problema bem legal, que usa o teorema de hall
// dada uma string s, no qual cada char eh a, b, c, d, e ou f
// eu quero permutar essa string s de alguma forma, tal que a seguinte condicao eh satisfeita para todo indice i:
// cada indice tem um set de chars que sao proibidos de estar naquela posicao
// se existir multiplas solucoes, printe a menor string lexicograficamente que puder ser a resposta</pre>
```

### 4.24 hld

```
// https://codeforces.com/contest/343/problem/D
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 500001
#define mod 1000000007
struct segtree
  int n;
  vector<int> v;
  vector<int> seq;
  vector<int> lazy;
  segtree() {}
  segtree(int sz)
    seg.assign(4 * n, 0);
    lazy.assign(4 * n, -1);
  int single(int x)
    return x;
  int neutral()
    return 0;
  int merge(int a, int b)
    return a + b;
  void add(int i, int l, int r, int diff)
    seg[i] = (r - 1 + 1) * diff;
    if (1 != r)
      lazy[i << 1] = diff;</pre>
      lazy[(i << 1) | 1] = diff;</pre>
    lazy[i] = -1;
  void update(int i, int l, int r, int ql, int qr, int diff)
    if (lazy[i] != -1)
      add(i, l, r, lazy[i]);
    if (1 > r || 1 > qr || r < q1)
```

```
if (1 >= q1 && r <= qr)
      add(i, 1, r, diff);
      return;
    int mid = (1 + r) >> 1;
    update(i << 1, 1, mid, q1, qr, diff);
    update((i << 1) | 1, mid + 1, r, ql, qr, diff);
    seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
  int query(int 1, int r, int q1, int qr, int i)
    if (lazy[i] != -1)
      add(i, 1, r, lazy[i]);
    if (1 > r || 1 > qr || r < q1)</pre>
     return neutral();
    if (1 >= ql && r <= qr)</pre>
      return seg[i];
    int mid = (1 + r) >> 1;
    return merge(query(1, mid, q1, qr, i << 1), query(mid + 1, r, q1, qr, (i <<</pre>
         1) | 1));
struct hld
  int n, cur_pos;
  segtree seg;
  vector<vector<int>> adj;
  vector<int> parent, depth, heavy, head, pos, sz;
  int dfs(int s)
    int size = 1, max_c_size = 0;
    for (auto const &c : adj[s])
      if (c != parent[s])
        parent[c] = s;
        depth[c] = depth[s] + 1;
        int c size = dfs(c);
        size += c_size;
        if (c_size > max_c_size)
          max_c_size = c_size, heavy[s] = c;
    return sz[s] = size;
  void decompose(int s, int h)
   head[s] = h;
    pos[s] = cur_pos++;
    if (heavy[s] != -1)
     decompose(heavy[s], h);
    for (int c : adj[s])
      if (c != parent[s] && c != heavy[s])
        decompose(c, c);
  hld(vector<vector<int>> &g)
   n = g.size();
    adj = g;
    seg = segtree(n);
    parent.assign(n, -1);
    depth.assign(n, -1);
    heavy.assign(n, -1);
    head.assign(n, -1);
    pos.assign(n, -1);
    sz.assign(n, 1);
    cur_pos = 0;
    dfs(0);
    decompose(0, 0);
  int query_path(int a, int b)
    int res = 0;
    for (; head[a] != head[b]; b = parent[head[b]])
```

```
if (depth[head[a]] > depth[head[b]])
       swap(a, b);
     res += seg.query(0, n - 1, pos[head[b]], pos[b], 1);
    if (depth[a] > depth[b])
     swap(a, b);
    res += seg.query(0, n - 1, pos[a], pos[b], 1);
   return res;
 void update_path(int a, int b, int x)
    for (; head[a] != head[b]; b = parent[head[b]])
     if (depth[head[a]] > depth[head[b]])
       swap(a, b);
      seg.update(1, 0, n - 1, pos[head[b]], pos[b], x);
   if (depth[a] > depth[b])
     swap(a, b);
    seg.update(1, 0, n - 1, pos[a], pos[b], x);
 void update_subtree(int a, int x)
   seg.update(1, 0, n - 1, pos[a], pos[a] + sz[a] - 1, x);
 int query_subtree(int a)
    return seg.query(0, n - 1, pos[a], pos[a] + sz[a] - 1, 1);
 int lca(int a, int b)
   if (pos[a] < pos[b])
     swap(a, b);
   return (head[a] == head[b]) ? b : lca(parent[head[a]], b);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n;
 cin >> n;
 vector<vector<int>> adj(n);
 for (int i = 0; i < n - 1; i++)
   int a, b;
   cin >> a >> b;
   a--, b--;
   adj[a].pb(b);
   adj[b].pb(a);
 hld h(adj);
 int q;
 cin >> q;
 while (q--)
   int a, b;
   cin >> a >> b;
   if (a == 1)
     h.update_subtree(b, 1);
    if (a == 2)
     h.update_path(0, b, 0);
   if (a == 3)
     cout << h.query_path(b, b) << endl;</pre>
 return 0;
```

## 4.25 hld edge

```
//https://www.spoj.com/problems/QTREE/
//Don't use cin/cout in this problem (gives TLE)
#include <bits/stdc++.h>
using namespace std;
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 10001
#define mod 1000000007
int n;
vector<pi> adj[MAXN];
vector<pi> edges;
namespace seg
  int seg[4 * MAXN];
  int lazy[4 * MAXN];
  int v[MAXN];
  int single(int x)
    return x;
  int neutral()
    return -1:
  int merge(int a, int b)
    return max(a, b);
  void add(int i, int l, int r, int diff)
    seg[i] = (r - 1 + 1) * diff;
    if(1 != r)
      lazy[i << 1] = diff;</pre>
     lazy[(i << 1) | 1] = diff;
    lazv[i] = -1;
  void update(int i, int l, int r, int ql, int qr, int diff)
    if (lazy[i] != -1)
     add(i, 1, r, lazy[i]);
    if (1 > r || 1 > qr || r < q1)
      return;
    if (1 >= ql \&\& r <= qr)
      add(i, 1, r, diff);
      return:
    int mid = (1 + r) >> 1;
    update(i << 1, 1, mid, q1, qr, diff);
    update((i << 1) | 1, mid + 1, r, ql, qr, diff);
    seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
  int query(int 1, int r, int q1, int qr, int i)
    if (lazy[i] != -1)
      add(i, 1, r, lazy[i]);
    if (1 > r || 1 > qr || r < q1)</pre>
     return neutral();
    if (1 >= q1 && r <= qr)</pre>
     return seg[i];
    int mid = (1 + r) >> 1;
    return merge(query(1, mid, q1, qr, i << 1), query(mid + 1, r, q1, qr, (i <<</pre>
         1) | 1));
  void build(int 1, int r, int i)
```

```
if (1 == r)
      seg[i] = single(v[1]);
      lazy[i] = -1;
      return;
    int mid = (1 + r) >> 1;
    build(1, mid, i << 1);
    build(mid + 1, r, (i << 1) | 1);
    seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
    lazy[i] = -1;
} // namespace seg
namespace hld
  int cur_pos;
  vector<int> parent, depth, heavy, head, pos, sz, up;
  int dfs(int s)
    int size = 1, max_c_size = 0;
    for (auto const &c : adj[s])
      if (c.fir != parent[s])
        parent[c.fir] = s;
        depth[c.fir] = depth[s] + 1;
        int c_size = dfs(c.fir);
        size += c_size;
        if (c_size > max_c_size)
          max_c_size = c_size, heavy[s] = c.fir;
    return sz[s] = size;
  void decompose(int s, int h)
    head[s] = h;
    pos[s] = cur_pos++;
    seg::v[pos[s]] = up[s];
    for (auto const &c : adj[s])
      if (c.fir != parent[s] && c.fir == heavy[s])
        up[c.fir] = c.sec;
        decompose (heavy[s], h);
    for (auto const &c : adj[s])
      if (c.fir != parent[s] && c.fir != heavy[s])
        up[c.fir] = c.sec;
        decompose(c.fir, c.fir);
  void init()
    parent.assign(MAXN, -1);
    depth.assign(MAXN, -1);
    heavy.assign(MAXN, -1);
    head.assign(MAXN, -1);
    pos.assign(MAXN, -1);
    sz.assign (MAXN, 1);
    up.assign(MAXN, 0);
    cur_pos = 0;
    dfs(0);
    decompose(0, 0);
    seg::build(0, n - 1, 1);
  int query_path(int a, int b)
    int res = -1:
    for (; head[a] != head[b]; b = parent[head[b]])
      if (depth[head[a]] > depth[head[b]])
```

```
swap(a, b);
      res = max(res, seg::query(0, n - 1, pos[head[b]], pos[b], 1));
    if (depth[a] > depth[b])
     swap(a, b);
    res = \max(\text{res, seq::query}(0, n - 1, pos[a] + 1, pos[b], 1));
   return res;
 void update_path(int a, int b, int x)
    for (; head[a] != head[b]; b = parent[head[b]])
      if (depth[head[a]] > depth[head[b]])
        swap(a, b);
      seg::update(1, 0, n - 1, pos[head[b]], pos[b], x);
    if (depth[a] > depth[b])
     swap(a, b);
    seg::update(1, 0, n - 1, pos[a] + 1, pos[b], x);
  void update_subtree(int a, int x)
   seg::update(1, 0, n - 1, pos[a] + 1, pos[a] + sz[a] - 1, x);
  int query_subtree(int a, int x)
   return seg::query(0, n - 1, pos[a] + 1, pos[a] + sz[a] - 1, 1);
} // namespace hld
signed main()
 int q;
scanf("%d", &q);
  while (q--)
    scanf("%d", &n);
    for (int i = 0; i < n; i++)
     adj[i].clear();
    edges.clear();
    for (int i = 0; i < n - 1; i++)
     int a, b, c;
scanf("%d %d %d", &a, &b, &c);
      a--, b--;
      adj[a].pb({b, c});
      adj[b].pb({a, c});
      edges.pb({a, b});
    hld::init();
    while (true)
      char k[10];
scanf("%s", k);
      if (k[0] == 'Q')
        int a, b;
        scanf("%d %d", &a, &b);
        printf("%d\n", hld::query_path(a, b));
      else if (k[0] == 'C')
        int a, b;
        scanf("%d %d", &a, &b);
        hld::update_path(edges[a].fir, edges[a].sec, b);
      else
        break;
  return 0;
```

## 4.26 hopcroft karp

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
// #define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1000003
#define mod 998244353
#define INF 1e9
struct hopcroft_karp
  vector<int> match;
  vector<int> dist;
  vector<vector<int>> adj;
  int n, m, t;
  hopcroft_karp(int a, int b)
   n = a, m = b;
    t = n + m + 1;
    match.assign(t, n + m);
    dist.assign(t, 0);
    adj.assign(t, vector<int>{});
  void add_edge(int u, int v)
    adj[u].pb(v);
    adj[v].pb(u);
  bool bfs()
    queue<int> q;
    for (int u = 0; u < n; u++)
      if (match[u] == n + m)
        dist[u] = 0, q.push(u);
      else
        dist[u] = INF;
    dist[n + m] = INF;
    while (!q.empty())
      int u = q.front();
      q.pop();
      if (dist[u] < dist[n + m])</pre>
        for (auto const &v : adj[u])
          if (dist[match[v]] == INF)
            dist[match[v]] = dist[u] + 1;
            q.push(match[v]);
    return dist[n + m] < INF;</pre>
  bool dfs(int u)
    if (u < n + m)
      for (auto const &v : adj[u])
```

```
if (dist[match[v]] == dist[u] + 1 && dfs(match[v]))
        match[v] = u;
       match[u] = v;
        return true;
    dist[u] = INF;
   return false;
 return true;
vector<pi> run()
 int cnt = 0;
  while (bfs())
    for (int u = 0; u < n; u++)
     if (match[u] == n + m && dfs(u))
       cnt++;
  vector<pi> ans;
  for (int v = n; v < n + m; v++)
    if (match[v] < n + m)
     ans.pb({match[v], v});
  return ans;
vector<int> mvc() // minimum vertex cover
  vector<pi> ans = run();
  vector<bool> vis(n + m, 0);
  for (int i = 0; i < n; i++)
    if (match[i] == n + m)
      queue<int> q;
      q.push(i);
      while (!q.empty())
        int x = q.front();
        q.pop();
        vis[x] = 1;
        for (auto const &y : adj[x])
          if (!vis[y])
           vis[y] = 1;
           q.push(match[y]);
  vector<int> vc;
  for (int i = 0; i < n; i++)
   if (!vis[i])
      vc.pb(i);
  for (int i = n; i < n + m; i++)
    if (vis[i])
      vc.pb(i);
 return vc;
vector<pi> mec() // minimum edge cover
  vector<pi> ans = run();
  for (int i = 0; i < n + m; i++)
    if (match[i] == n + m \&\& adj[i].size() > 0)
      if (i < n)
       ans.pb({i, adj[i][0]});
        ans.pb({adj[i][0], i});
 return ans;
```

```
};
// minimum path cover on dag
// minimum set of paths such that each of the vertices belongs to exactly one
vector<vector<int>> mpc(int n, vector<pi> &e)
  hopcroft_karp h(n, n);
  for (auto const &i : e)
   h.add_edge(i.fir, n + i.sec);
  vector<pi> mat = h.run();
  vector<int> prv(n, -1);
  vector<int> nxt(n, -1);
  for (int i = 0; i < mat.size(); i++)</pre>
   nxt[mat[i].fir] = mat[i].sec - n;
    prv[mat[i].sec - n] = mat[i].fir;
  vector<vector<int>> ans;
  for (int i = 0; i < n; i++)
    if (prv[i] == -1 && nxt[i] == -1)
      ans.pb({i});
    else if (prv[i] == -1)
      vector<int> curr;
      int x = i;
      while (1)
        curr.pb(x);
        if (nxt[x] == -1)
         break;
        x = nxt[x];
      ans.pb(curr);
  return ans;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, m;
  cin >> n >> m;
  vector<pi> e;
  for (int i = 0; i < m; i++)
    int a, b;
   cin >> a >> b;
    a--, b--;
    e.pb({a, b});
  vector<vector<int>> ans = mpc(n, e);
  cout << ans.size() << endl;</pre>
  for (auto const &v : ans)
    for (auto const &i : v)
     cout << i + 1 << " ";
    cout << endl;</pre>
  return 0;
```

# 4.27 hungarian

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;

template <class T>
```

```
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 505
#define mod 998244353
struct hungarian
 int n, inf;
 vector<vector<int>> a;
 vector<int> u, v, p, way;
  hungarian(int n_{-}): n(n_{-}), u(n + 1), v(n + 1), p(n + 1), way(n + 1)
   a = vector<vector<int>>(n, vector<int>(n));
   inf = numeric_limits<int>::max();
  void add_edge(int x, int y, int c)
   a[x][y] = c;
  pair<int, vector<int>> run()
    for (int i = 1; i <= n; i++)</pre>
      p[0] = i;
      int j0 = 0;
      vector<int> minv(n + 1, inf);
      vector<int> used(n + 1, 0);
      do
        used[j0] = true;
        int i\bar{0} = p[j0], j1 = -1;
        int delta = inf;
        for (int j = 1; j \le n; j++)
          if (!used[j])
            int cur = a[i0 - 1][j - 1] - u[i0] - v[j];
            if (cur < minv[j])</pre>
              minv[j] = cur, way[j] = j0;
            if (minv[j] < delta)</pre>
              delta = minv[j], j1 = j;
        for (int j = 0; j \le n; j++)
          if (used[i])
            u[p[j]] += delta, v[j] -= delta;
          else
            minv[j] -= delta;
        i0 = j1;
        while (p[j0] != 0);
      do
        int j1 = way[j0];
        p[j0] = p[j1];
        j0 = j1;
      } while (j0);
    vector<int> ans(n);
    for (int j = 1; j \le n; j++)
      ans[p[j] - 1] = j - 1;
    return make_pair(-v[0], ans);
};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
 cin >> n;
```

### 4.28 Kruskal

```
// Algoritimo de kruskal - Achar a mst
// 1 - listar todas as arestas em ordem crescente.
// 2 - Cada aresta liga dois vertices x e y, checar se eles ja estao na mesma
    componente conexa
// (aqui, consideramos apenas as arestas ja colocadas na arvore).
// 3 - Se x e y estao na mesma componente, ignoramos a aresta e continuamos o
    procedimento
// (se a usassemos, formariamos um ciclo). Se estiverem em componentes distintas
    , colocamos a aresta
//na arvore e continuamos o procedimento.
// OBS: como a prioridade eh ordenar pelas menores distancias, basta botar o
    custo da aresta como
// first no vector das arestas para poder ordenar
// em suma: ordeno as arestas em ordem crescente com prioridade no custo, depois
     para cada aresta,
// se o find(x) != find(y) sendo x e y os vertices das arestas, eu adiciono eles
     a mst e dou um join
// nos dois, como as arestas tao ordenadas em ordem crescente, o primeiro que eu
// eh necessariamente a melhor opcao e assim a mst eh formada.
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define mp make_pair
#define fir first
#define sec second
#define MAXN 100001
int n, m, a, b, c;
vector<pii> ar;
vector<pii> mst;
int pai[MAXN];
int peso[MAXN];
int find(int x)
  if (pai[x] == x)
    return x;
  return pai[x] = find(pai[x]);
void join(int a, int b)
  a = find(a);
  b = find(b);
  if (peso[a] < peso[b])</pre>
   pai[a] = b;
```

```
else if (peso[b] < peso[a])</pre>
   pai[b] = a;
 else
   pai[a] = b;
   peso[b]++;
void initialize()
  for (int i = 1; i <= n; i++)</pre>
   pai[i] = i;
int main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 cin >> n >> m;
  for (int i = 0; i < m; i++)
   cin >> a >> b >> c;
   ar.pb(mp(c, mp(a, b)));
  sort(ar.begin(), ar.end());
  initialize();
  int size = 0;
  for (int i = 0; i < m; i++)
   if (find(ar[i].sec.fir) != find(ar[i].sec.sec))
      join(ar[i].sec.fir, ar[i].sec.sec);
      mst.pb(mp(ar[i].fir, mp(ar[i].sec.fir, ar[i].sec.sec)));
  for (int i = 0; i < mst.size(); i++)</pre>
    cout << mst[i].sec.fir << " " << mst[i].sec.sec << " " << mst[i].fir << endl
  return 0;
```

## 4.29 kuhn

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 2001
#define mod 998244353
```

```
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
struct kuhn
  int n, m, at;
  vector<vector<int>> q;
  vector<int> vis, ma, mb;
  kuhn(int n_, int m_) : n(n_), m(m_), at(0), g(n),
                          vis(n), ma(n, -1), mb(m, -1) {}
  void add_edge(int a, int b) { g[a].push_back(b); }
  bool dfs(int i)
    vis[i] = at;
    for (int j : g[i])
      if (mb[j] == -1)
        ma[i] = j, mb[j] = i;
        return true;
    for (int j : g[i])
      if (vis[mb[j]] != at && dfs(mb[j]))
        ma[i] = j, mb[j] = i;
        return true;
    return false;
  int augment (int i) // pode usar quando tou adicionando o vertice i pela
       primeira vez no grafo
    shuffle(g[i].begin(), g[i].end(), rng);
    if (dfs(i))
     at++;
    return at;
  int matching() // calcula o max matching
    int aum = 1;
    while (aum)
      for (auto &i : g)
        shuffle(i.begin(), i.end(), rng);
      for (int j = 0; j < n; j++)
       vis[j] = 0;
      aum = 0;
      for (int i = 0; i < n; i++)
  if (ma[i] == -1 and dfs(i))</pre>
          at++, aum = 1;
    return at;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, m;
  cin >> n >> m;
  vector<int> p(n);
  for (int i = 0; i < n; i++)
    cin >> p[i];
  vector<int> c(n);
  vector<vector<pi>> v(5001);
  for (int i = 0; i < n; i++)
    cin >> c[i];
    c[i]--;
    v[p[i]].pb({i, c[i]});
  int q;
  cin >> q;
```

```
vector<int> qry(q);
  vector<bool> rem(n, 0);
  for (int i = 0; i < q; i++)
    cin >> qry[i];
   qry[i]--;
   rem[qry[i]] = 1;
  int sz = m + 1;
  kuhn h(5001, sz);
  vector<int> ans(q);
  int mex = 0;
  for (auto const & j : v[0])
   if (!rem[j.fir])
      h.add_edge(0, j.sec);
  h.augment(0);
  for (int i = q - 1; i >= 0; i--)
    int x = qry[i];
   while (mex < m \&\& h.at == mex + 1)
      mex++:
      for (auto const & j : v[mex])
       if (!rem[j.fir])
          h.add_edge(mex, j.sec);
     h.augment (mex);
    ans[i] = mex;
    rem[x] = 0;
    if (p[x] \le mex)
     h.add\_edge(p[x], c[x]);
     h.matching();
  for (auto const &i : ans)
   cout << i << endl;</pre>
 return 0;
// problema de matching incremental, ir adicionando arestas e calculando o max
    matching
// https://codeforces.com/contest/1139/problem/E
// selecionar um student de cada club, caso o club possua pelo menos um caba
// pros selecionados, maximizar o mex
// fazer as queries ao contrario talvez seja o caminho
```

### 4.30 LCA

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 100001
#define mod 1000000007
int n;
```

```
vector<int> adj[MAXN];
namespace lca
 int 1, timer;
 vector<int> tin, tout, depth;
 vector<vector<int>> up;
 void dfs(int v, int p)
   tin[v] = ++timer;
    up[v][0] = p;
    for (int i = 1; i <= 1; i++)
     up[v][i] = up[up[v][i - 1]][i - 1];
    for (auto const &u : adj[v])
     if (p == u)
       continue;
      depth[u] = depth[v] + 1;
     dfs(u, v);
    tout[v] = ++timer;
 bool is_ancestor(int u, int v)
    return tin[u] <= tin[v] && tout[u] >= tout[v];
  int binary_lifting(int u, int v)
    if (is_ancestor(u, v))
      return u;
    if (is_ancestor(v, u))
     return v;
    for (int i = 1; i >= 0; --i)
     if (!is_ancestor(up[u][i], v))
        u = up[u][i];
    return up[u][0];
  void init()
    tin.resize(n);
    tout.resize(n):
    depth.resize(n);
   timer = 0;
   1 = ceil(log2(n));
    up.assign(n, vector<int>(1 + 1));
    dfs(0, 0);
  int dist(int s, int v)
    int at = binary_lifting(s, v);
   return (depth[s] + depth[v] - 2 * depth[at]);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 cin >> n;
  for (int i = 0; i < n - 1; i++)
   int a, b;
   cin >> a >> b;
    a--, b--;
   adj[a].pb(b);
   adj[b].pb(a);
  lca::init();
 return 0;
```

## 4.31 link cut tree edge

```
#include <bits/stdc++.h>
using namespace std;
```

```
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
struct node
  int p, ch[2];
  pii val, sub;
  bool rev;
  int sz, ar;
  int lazy;
  node() {}
  node(pii v, int ar_) : p(-1), val(v), sub(v), rev(0), sz(ar_), ar(ar_), lazy
       (0)
   ch[0] = ch[1] = -1;
};
struct link_cut_tree
  vector<node> t:
  map<pair<int, int>, int> aresta;
  int sz, n;
  link_cut_tree(int nn)
    t.clear();
   n = nn;
   sz = 0;
  static pii neutral()
    return {0, {0, 0}};
  pii merge(pii a, pii b)
    return max(a, b);
  void prop(int x)
    if (t[x].lazy)
      if (t[x].ar)
       t[x].val.fir = t[x].lazy;
      t[x].sub = merge(t[x].sub, t[x].val);
      if (t[x].ch[0] + 1)
        t[t[x].ch[0]].lazy = t[x].lazy;
      if (t[x].ch[1] + 1)
        t[t[x].ch[1]].lazy = t[x].lazy;
    if (t[x].rev)
      swap(t[x].ch[0], t[x].ch[1]);
      if (t[x].ch[0] + 1)
        t[t[x].ch[0]].rev ^= 1;
      if (t[x].ch[1] + 1)
        t[t[x].ch[1]].rev ^= 1;
    t[x].lazy = 0, t[x].rev = 0;
  void update(int x)
    t[x].sz = t[x].ar, t[x].sub = t[x].val;
    for (int i = 0; i < 2; i++)
      if (t[x].ch[i] + 1)
        prop(t[x].ch[i]);
        t[x].sz += t[t[x].ch[i]].sz;
        t[x].sub = merge(t[x].sub, t[t[x].ch[i]].sub);
```

```
bool is_root(int x)
  return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].ch[1] != x);
void rotate(int x)
  int p = t[x].p, pp = t[p].p;
  if (!is_root(p))
    t[pp].ch[t[pp].ch[1] == p] = x;
  bool d = t[p].ch[0] == x;
  t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
  if (t[p].ch[!d] + 1)
    t[t[p].ch[!d]].p = p;
  t[x].p = pp, t[p].p = x;
  update(p), update(x);
int splay(int x)
  while (!is root(x))
    int p = t[x].p, pp = t[p].p;
    if (!is_root(p))
      prop(pp);
    prop(p), prop(x);
    if (!is_root(p))
      rotate((t[pp].ch[0] == p) ^ (t[p].ch[0] == x) ? x : p);
    rotate(x);
  return prop(x), x;
int access(int v)
  int last = -1;
  for (int w = \dot{v}; w + 1; update(last = \dot{w}), splay(\dot{v}), w = t[\dot{v}].p) splay(\dot{w}), t[\dot{w}].ch[1] = (last == -1 ? -1 : \dot{v});
  return last;
void make_tree(int v, pii w = neutral(), int ar = 0)
  while (t.size() <= v)</pre>
    t.pb(node(neutral(), 0));
  t[v] = node(w, ar);
int find_root(int v)
  access(v), prop(v);
  while (t[v].ch[0] + 1)
    v = t[v].ch[0], prop(v);
  return splay(v);
bool conn(int v, int w)
  access(v), access(w);
  return v == w ? true : t[v].p != -1;
void rootify(int v)
  access(v);
  t[v].rev ^= 1;
pii query(int v, int w)
  rootify(w), access(v);
  return t[v].sub;
void update(int v, int w, int x)
  rootify(w), access(v);
  t[v].lazy += x;
void link_(int v, int w)
  rootify(w);
  t[w].p = v;
void link(int v, int w, pii x)
  int id = n + sz++;
```

```
aresta[make_pair(v, w)] = id;
    make_tree(id, x, 1);
    link_(v, id), link_(id, w);
  void cut_(int v, int w)
    rootify(w), access(v);
    t[v].ch[0] = t[t[v].ch[0]].p = -1;
  void cut(int v, int w)
    int id = aresta[make_pair(v, w)];
   cut_(v, id), cut_(id, w);
  int lca(int v, int w)
    access(v);
   return access(w);
};
struct dsu
  int tot;
  vector<int> parent;
  vector<int> sz;
  dsu(int n)
   parent.resize(n);
    sz.resize(n);
    tot = n:
    for (int i = 0; i < n; i++)
     parent[i] = i;
     sz[i] = 1;
  int find_set(int i)
    return parent[i] = (parent[i] == i) ? i : find_set(parent[i]);
  void make_set(int x, int y)
    x = find_set(x), y = find_set(y);
    if (x != y)
      if (sz[x] > sz[y])
       swap(x, y);
      parent[x] = y;
      sz[y] += sz[x];
      tot--;
};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
 int q;
  cin >> q;
  while (q--)
    int n, m, qq;
    cin >> n >> m >> qq;
    vector<pii> e;
    for (int i = 0; i < m; i++)
      int a, b, c;
      cin >> a >> b >> c;
      a--, b--;
     e.pb({c, {a, b}});
    sort(e.begin(), e.end());
    dsu d(n);
    link_cut_tree 1(n);
    for (int i = 0; i < n; i++)
      1.make_tree(i);
```

```
int cost = 0;
    for (auto const &i : e)
      if (d.find set(i.sec.fir) != d.find set(i.sec.sec))
        d.make_set(i.sec.fir, i.sec.sec);
        1.link(i.sec.fir, i.sec.sec, i);
        cost += i.fir;
    while (qq--)
      int a, b, c;
      cin >> a >> b >> c;
      pii mx = 1.query(a, b);
      if (c < mx.fir)</pre>
        cost -= mx.fir;
        cost += c;
        1.cut(mx.sec.fir, mx.sec.sec);
        // 1.link(a, b, {c, {a, b}}); poderia fazer assim, mas quero testar o
        1.link(a, b, {0, {a, b}});
        l.update(a, b, c);
      cout << cost << endl;
 return 0;
// link cut tree com peso nas arestas
// solucao para o: https://codeforces.com/gym/101047/problem/I
// problema onde e dado um grafo inicial e algumas queries
// cada query adiciona uma nova aresta nesse grafo
// e o objetivo e achar a mst apos cada adicao de aresta
// implementacao baseada na: https://github.com/brunomaletta/Biblioteca/blob/
    master/Codigo/Grafos/LCT/lctAresta.cpp
// make_tree(v) cria uma nova arvore com um um unico vertice
// rootify(v) torna v a raiz de sua arvore
// cut(u, v) apaga a aresta u, v
// link(u, v, c) adiciona a aresta de u ate v com peso c
// query(v, w) retorna a aresta de maior peso no caminho de v ate w
// update(v, w, x) faz com que as arestas do caminho de v ate w passem a ter
    peso x
// operacoes tem complexidade O(log(n)) amortizado
```

#### 4.32 link cut tree vertex

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 998244353
struct node
  int p, ch[2];
  int val, sub, lazy;
  bool rev;
  int sz;
```

```
node() {}
 node(int v) : p(-1), val(v), sub(v), rev(0), sz(1), lazy(0)
    ch[0] = ch[1] = -1;
};
struct link_cut_tree
 vector<node> t:
 link_cut_tree()
    t.clear();
  static int neutral()
    return 0;
  int merge(int a, int b)
    return a ^ b;
  void prop(int x)
    if (t[x].lazy)
      t[x].val = t[x].lazy;
      t[x].sub = t[x].val;
      if (t[x].ch[0] + 1)
        t[t[x].ch[0]].lazy = t[x].lazy;
      if (t[x].ch[1] + 1)
        t[t[x].ch[1]].lazy = t[x].lazy;
    if (t[x].rev)
      swap(t[x].ch[0], t[x].ch[1]);
      if (t[x].ch[0] + 1)
        t[t[x].ch[0]].rev ^= 1;
      if (t[x].ch[1] + 1)
        t[t[x].ch[1]].rev ^= 1;
    t[x].lazy = 0, t[x].rev = 0;
  void update(int x)
    t[x].sz = 1, t[x].sub = t[x].val;
    for (int i = 0; i < 2; i++)
      if (t[x].ch[i] + 1)
        prop(t[x].ch[i]);
        t[x].sz += t[t[x].ch[i]].sz;
        t[x].sub = merge(t[x].sub, t[t[x].ch[i]].sub);
 bool is_root(int x)
    return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].ch[1] != x);
  void rotate(int x)
    int p = t[x].p, pp = t[p].p;
    if (!is_root(p))
     t[pp].ch[t[pp].ch[1] == p] = x;
    bool d = t[p].ch[0] == x;
    t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
if (t[p].ch[!d] + 1)
     t[t[p].ch[!d]].p = p;
   t[x].p = pp, t[p].p = x;
update(p), update(x);
  int splay(int x)
    while (!is_root(x))
      int p = t[x].p, pp = t[p].p;
if (!is_root(p))
        prop(pp);
```

```
prop(p), prop(x);
      if (!is_root(p))
        rotate((t[pp].ch[0] == p) ^ (t[p].ch[0] == x) ? x : p);
      rotate(x);
    return prop(x), x;
  int access(int v)
    int last = -1;
    for (int w = v; w + 1; update(last = w), splay(v), w = t[v].p)
     splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
    return last:
  void make_tree(int v, int w = neutral()) // cria uma arvore com um unico
      vertice, sendo o vertice i
    while (t.size() <= v)</pre>
      t.pb(node(neutral()));
    t[v] = node(w);
  int find_root(int v) // acha a raiz da arvore do vertice v
    access(v), prop(v);
   while (t[v].ch[0] + 1)
     v = t[v].ch[0], prop(v);
    return splay(v);
 bool connected(int v, int w) // checa se v e w estao na mesma arvore (aka
      componente conexa)
   access(v), access(w);
   return v == w ? true : t[v].p != -1;
 void rootify(int v) // torna v a raiz de sua arvore
    access(v);
   t[v].rev ^= 1;
  int query(int v, int w) // query no caminho de v ate w
    rootify(w), access(v);
   return t[v].sub;
 void update(int v, int w, int x) // aplica o update em todos os vertices no
      caminho de v ate w
    rootify(w), access(v);
   t[v].lazy = x;
  void link(int v, int w) // adiciona a aresta v - w
    rootify(w);
   t[w].p = v;
 void cut(int v, int w) // remove a aresta v - w
    rootify(w), access(v);
   t[v].ch[0] = t[t[v].ch[0]].p = -1;
  int lca(int v, int w) // acha o lca(v, w)
    access(v);
    return access(w);
};
set < int > adi[MAXN];
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n, q;
 cin >> n >> q;
 int x = 0;
  link_cut_tree 1;
 for (int i = 0; i < n; i++)
```

```
// l.make_tree(i, i) // poderia fazer isso tbm, para criar o vertice i com
    l.make_tree(i);
   1.update(i, i, i); // testa o update
  while (q--)
   int a, b, c;
   cin >> a >> b >> c;
    a = ((a * (1 + x)) % mod) % 2;
   b = ((b * (1 + x)) % mod) % n;
    c = ((c * (1 + x)) % mod) % n;
    if (a == 0)
      adj[b].insert(c);
      adj[c].insert(b);
      1.link(b, c);
    else
      if (!l.connected(b, c))
        \mathbf{x} = 0:
        cout << x << endl;</pre>
        continue;
      int d = 1.query(b, c);
     d = b;
     d ^= c;
      x = 0;
      if (d >= 0 \&\& d < n)
       bool ok = 1;
        ok &= (adj[b].find(d) != adj[b].end());
        ok &= (adj[c].find(d) != adj[c].end());
        if (ok)
          x = d + 1;
      cout << x << endl;
  }
// problema exemplo: https://atcoder.jp/contests/abc350/tasks/abc350_g
// nesse caso, a link cut tree tem lazy para assign de valor
// e de queries de xor
```

## 4.33 MatrixDijkstra

```
#include <bits/stdc++.h>
using namespace std;
\#define \ lli \ long \ long \ int
#define pb push_back
#define MAXN 10000000
typedef pair <int , int> pii ;
int t :
int dist [MAXN] ;
bool visited [MAXN] ;
vector <pii> adj_list [MAXN] ;
void dijkstra (int s)
    dist[s] = 0;
    priority_queue <pii , vector<pii> , greater<pii>> q ;
    q.push(pii(dist[s], s));
    while(1)
        int davez = -1;
        int menor = INT_MAX ;
        while(!q.empty())
```

```
int atual = q.top().second ;
            q.pop();
            if(!visited[atual])
                davez = atual;
                break;
        if(davez == -1)
            break ;
        visited[davez] = true ;
        for(int i = 0 ; i < adj list[davez].size() ; i++)</pre>
            int distt = adj_list[davez][i].first;
            int atual = adj_list[davez][i].second;
            if(dist[atual] > dist[davez] + distt)
                dist[atual] = dist[davez] + distt;
                q.push(pii(dist[atual] , atual)) ;
void initialize ()
    for (int i = 0; i < t; i++)
        visited[i] = false ;
        dist[i] = INT_MAX ;
int main()
    ios base::sync with stdio(false);
    cin.tie(NULL);
    int n , m ;
    cin >> n >> m ;
    t = n * m ;
    char array [t] ;
    for (int i = 0; i < t; i++)
        cin >> array[i] ;
    for (int i = 0; i < t; i++)
        if (i >= m && array[i] != '#')
            adj_list[i].pb(pii(1 , (i - m)));
        if (i < (n * m) - m && array[i] != '#')</pre>
            adj_list[i].pb(pii(1 , (i + m)));
        if (i % m != 0 && array[i] != '#')
            adj_list[i].pb(pii(1 , (i - 1))) ;
        if ((i + 1) % m != 0 && array[i] != '#')
            adj_list[i].pb(pii(1 , (i + 1)));
    int q;
    cin >> q;
```

```
while (q--)
{
    int a , b , c , d , e ;
    cin >> a >> b >> c >> d >> e ;
    a-- , b-- , c-- , d-- ;

    int index1 = (m * a) + b ;
    int index2 = (m * c) + d ;

    adj_list[index1].pb(pii(e , index2)) ;
    adj_list[index2].pb(pii(e , index1)) ;
}

initialize () ;
dijkstra(0) ;
cout << dist[t - 1] << endl ;
return 0 ;
}</pre>
```

# 4.34 max matching without one vertex

```
using namespace std;
#define int long long int
#define mod 1000000007
#define MAXN 100005
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define pb push_back
#define INF 1e9
struct hopcroft_karp
 vector<int> match;
 vector<int> dist;
 vector<vector<int>> adj;
 int n, m, t;
 hopcroft_karp(int a, int b)
   n = a, m = b;
   t = n + m + 1;
   match.assign(t, n + m);
   dist.assign(t, 0);
   adj.assign(t, vector<int>{});
 void add_edge(int u, int v)
    adj[u].pb(v);
   adj[v].pb(u);
 bool bfs()
    queue<int> q;
    for (int u = 0; u < n; u++)
      if (match[u] == n + m)
       dist[u] = 0, q.push(u);
      else
        dist[u] = INF;
    dist[n + m] = INF;
    while (!q.empty())
      int u = q.front();
     q.pop();
      if (dist[u] < dist[n + m])</pre>
```

#include <bits/stdc++.h>

```
for (auto const &v : adj[u])
        if (dist[match[v]] == INF)
          dist[match[v]] = dist[u] + 1;
          q.push(match[v]);
  return dist[n + m] < INF;</pre>
bool dfs(int u)
  if (u < n + m)
    for (auto const &v : adj[u])
      if (dist[match[v]] == dist[u] + 1 && dfs(match[v]))
        match[v] = u;
        match[u] = v;
        return true;
    dist[u] = INF;
    return false;
  return true;
vector<pi> run()
  int cnt = 0;
  while (bfs())
    for (int u = 0; u < n; u++)
      if (match[u] == n + m && dfs(u))
        cnt++;
  vector<pi> ans;
  for (int v = n; v < n + m; v++)
    if (match[v] < n + m)
      ans.pb({match[v], v});
  return ans;
vector<int> solve()
  vector<pi> ans = run();
  vector<bool> vis(n + m, 0);
  vector<bool> can_remove(n + m, 0);
  for (int i = 0; \overline{i} < n; i++)
    if (match[i] == n + m)
      queue<int> q;
      q.push(i);
      while (!q.empty())
        int x = q.front();
        q.pop();
        vis[x] = 1;
        can\_remove[x] = 1;
        for (auto const &y : adj[x])
          if (!vis[y])
            vis[y] = 1;
            q.push(match[y]);
  vis = vector < bool > (n + m, 0);
  for (int i = n; i < n + m; i++)
    if (match[i] == n + m)
      queue<int> q;
      q.push(i);
      while (!q.empty())
```

```
int x = q.front();
          q.pop();
          vis[x] = 1;
          can\_remove[x] = 1;
          for (auto const &y : adj[x])
            if (!vis[y])
              vis[y] = 1;
              q.push(match[y]);
    vector<int> resp;
    for (int i = 0; i < n + m; i++)
      if (can_remove[i])
        resp.pb(i);
   return resp;
};
int dx[] = \{-1, 1, 0, 0\};
int dy[] = \{0, 0, 1, -1\};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  vector<pi> v(n);
  vector<int> sz(2, 0);
  map<pi, int> mp;
  map<int, pi> rev;
  for (int i = 0; i < n; i++)</pre>
    cin >> v[i].fir >> v[i].sec;
   int id = sz[abs(v[i].fir + v[i].sec) % 2]++;
   mp[v[i]] = id;
  for (int i = 0; i < n; i++)
    if (abs(v[i].fir + v[i].sec) % 2)
      mp[v[i]] += sz[0];
   rev[mp[v[i]]] = v[i];
  vector<pi> edges;
  for (auto [i, j] : v)
    if (!(abs(i + j) % 2))
      for (int dir = 0; dir < 4; dir++)
        int x = i + dx[dir];
        int y = j + dy[dir];
        if (mp.find({x, y}) != mp.end())
          edges.pb(\{mp[\{i, j\}], mp[\{x, y\}]\});
  hopcroft_karp h(sz[0], sz[1]);
  for (auto [x, y] : edges)
   h.add_edge(x, y);
  vector<int> ans = h.solve();
  cout << ans.size() << endl;</pre>
  vector<pi> sorted;
  for (auto const &i : ans)
```

```
sorted.pb(rev[i]);
}
sort(sorted.begin(), sorted.end());
for (auto [x, y] : sorted)
{
   cout << x << " " << y << endl;
}
// https://codeforces.com/group/TFrGcBYYxs/contest/583964/problem/E
// https://codeforces.com/gym/105053/problem/C
// https://codeforces.com/group/TFrGcBYYxs/contest/584239/problem/B
// dado um grafo bipartido
// basicamente responde o seguinte problema para cada vertice v:
// apos remover v do grafo, o tamanho do max matching continua o mesmo ou diminui em um?</pre>
```

### 4.35 mincostflow

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 301
#define mod 1000000007
#define INF 1e9
namespace mcf
  struct edge
    int to, capacity, cost, res;
  };
  int source, destiny;
  vector<edge> adj[MAXN];
  vector<int> dist;
  vector<int> parent;
  vector<int> edge_index;
  vector<bool> in_queue;
  void add_edge(int a, int b, int c, int d)
                                                  // aresta normal
    adj[a].pb({b, c, d, (int)adj[b].size()});
    adj[b].pb({a, 0, -d, (int)adj[a].size() - 1}); // aresta do grafo residual
  bool dijkstra(int s) // rodando o dijkstra, terei o caminho de custo minimo
                       // que eu consigo passando pelas arestas que possuem
       capacidade > 0
    dist.assign(MAXN, INF);
    parent.assign(MAXN, -1);
    edge_index.assign(MAXN, -1);
    in_queue.assign(MAXN, false);
    dist[s] = 0;
    queue<int> q;
    q.push(s);
    while (!q.empty())
      int u = q.front(), idx = 0;
      q.pop();
      in_queue[u] = false;
      for (auto const &v : adj[u])
        if (v.capacity && dist[v.to] > dist[u] + v.cost)
```

```
dist[v.to] = dist[u] + v.cost;
          parent[v.to] = u;
          edge_index[v.to] = idx;
          if (!in_queue[v.to])
            in_queue[v.to] = true;
            q.push(v.to);
        idx++;
    return dist[destiny] != INF; // se eu chequei em destiny por esse caminho,
        ainda posso passar fluxo
  int get_cost()
    int flow = 0, cost = 0;
    while (dijkstra(source)) // rodo um dijkstra para saber qual o caminho que
        irei agora
      int curr_flow = INF, curr = destiny;
      while (curr != source) // com isso, vou percorrendo o caminho encontrado
          para achar a aresta "gargalo"
        int p = parent[curr];
        curr_flow = min(curr_flow, adj[p][edge_index[curr]].capacity);
        curr = p;
      flow += curr_flow;
                                         // fluxo que eu posso passar por esse
          caminho = custo da aresta "gargalo"
      cost += curr_flow * dist[destiny]; // quanto eu gasto para passar esse
          fluxo no caminho encontrado
      curr = destiny;
      while (curr != source) // apos achar a aresta gargalo, passamos o fluxo
          pelo caminho encontrado
        int p = parent[curr];
        int res_idx = adj[p][edge_index[curr]].res;
        adj[p][edge_index[curr]].capacity -= curr_flow;
        adj[curr][res_idx].capacity += curr_flow;
        curr = p;
    return cost; // ao final temos a resposta :)
} // namespace mcf
signed main()
  int n;
  cin >> n;
  int v[n][n];
  mcf::source = 0, mcf::destiny = (2 * n) + 1;
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++)
      cin >> v[i][j];
      mcf::add\_edge(i + 1, j + n + 1, 1, v[i][j]);
  for (int i = 1; i <= n; i++)
   mcf::add_edge(mcf::source, i, 1, 0);
  for (int i = n + 1; i \le n + n; i++)
   mcf::add_edge(i, mcf::destiny, 1, 0);
  cout << mcf::get_cost << endl;</pre>
```

### 4.36 mo dsu

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
```

```
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 50001
#define mod 1000000007
const int block = 224;
struct query
  int 1, r, i;
};
vector<query> queries[block];
namespace dsu
  struct rollback
    int u, v, ranku, rankv;
  };
  int num_sets;
  int parent[MAXN];
  int rank[MAXN];
  stack<rollback> op;
  int Find(int i)
    return (parent[i] == i) ? i : Find(parent[i]);
  bool Union(int x, int y, bool can_rollback)
    int xx = Find(x);
    int yy = Find(y);
    if (xx != yy)
      num_sets--;
      if (rank[xx] > rank[yy])
        swap(xx, yy);
      if (can_rollback)
        op.push({xx, yy, rank[xx], rank[yy]});
      parent[xx] = yy;
      if (rank[xx] == rank[yy])
        rank[yy]++;
      return true;
    return false;
  void do_rollback()
    if (op.empty())
     return;
    rollback x = op.top();
    op.pop();
    num_sets++;
    parent[x.v] = x.v;
    rank[x.v] = x.rankv;
    parent[x.u] = x.u;
    rank[x.u] = x.ranku;
  void rollback all()
    while (!op.empty())
      do_rollback();
  void init(int n)
    for (int i = 0; i < n; i++)
```

parent[i] = i;

```
rank[i] = 0;
    num\_sets = n;
} // namespace dsu
bool cmp(query a, query b)
  return a.r < b.r;</pre>
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, m;
  cin >> n >> m;
  vector<pi> edges(m);
  for (int i = 0; i < m; i++)
    cin >> edges[i].fir >> edges[i].sec;
   edges[i].fir--;
   edges[i].sec--;
  int q;
  cin >> q;
  for (int i = 0; i < q; i++)
    int 1, r;
    cin >> 1 >> r;
    1--, r--;
    queries[l / block].pb({l, r, i});
  for (int i = 0; i < block; i++)</pre>
    if (queries[i].size())
      sort(queries[i].begin(), queries[i].end(), cmp);
  vector<int> ans(q);
  for (int i = 0; i < block; i++)</pre>
    if (!queries[i].size())
      continue;
    dsu::init(n);
    int limit = (i + 1) * block;
    for (auto const &j : queries[i])
      while (j.r >= limit)
        dsu::Union(edges[limit].fir, edges[limit].sec, false), limit++;
      for (int k = j.1; k \le \min(((i + 1) * block) - 1, j.r); k++)
       dsu::Union(edges[k].fir, edges[k].sec, true);
      ans[j.i] = dsu::num_sets;
      dsu::rollback_all();
  for (auto const &i : ans)
    cout << i << endl;</pre>
  return 0;
// https://codeforces.com/edu/course/2/lesson/7/3/practice/contest/289392/
    problem/B
   temos que fazer algo parecido com um mo algorithm
// sendo que a operacao eh um union/rollback do dsu
// podemos aplicar a seguinte ideia:
   - separamos os queries em blocos (pelo 1) de tamanho sgrt(n)
// - para cada bloco, ordenamos esse bloco em ordem crescente do r
// - com isso, em cada query eu posso fazer:
// - de l ate (limite dâquele bloco) - 1, adiciono na marra, podendo dar
    rollback
// - como o r eh crescente, para os valores de r que forem maior do que o limit
    daquele block
// - eu ja posso deixar adicionado para sempre sem precisar dar rollback
// fica algo que funciona em coisas que voce pode dar rollback
```

#### 4.37 mo trees

#include <bits/stdc++.h>
#include <ext/pb\_ds/assoc\_container.hpp>

```
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 998244353
struct gry
  int 1, r, lca, id;
};
int n, q;
vector<int> adj[MAXN];
int v[MAXN];
int cnt[MAXN];
int freq[MAXN];
int tin[MAXN];
int tout[MAXN];
int depth[MAXN];
int up[MAXN][25];
vector<int> t;
vector<qry> qq;
void dfs(int s, int p)
 tin[s] = t.size();
  up[s][0] = p;
  for (int i = 1; i < 25; i++)
   up[s][i] = up[up[s][i - 1]][i - 1];
  t.pb(s);
  for (auto const &i : adj[s])
   if (i == p)
     continue;
    depth[i] = depth[s] + 1;
    dfs(i, s);
  tout[s] = t.size();
  t.pb(s);
bool is(int u, int v)
  return tin[u] <= tin[v] && tout[u] >= tout[v];
int lca(int u, int v)
  if (is(u, v))
   return u;
  if (is(v, u))
   return v;
  for (int i = 24; i >= 0; i--)
   if (!is(up[u][i], v))
      u = up[u][i];
  return up[u][0];
void compress()
  vector<int> vals;
  for (int i = 0; i < n; i++)</pre>
   vals.pb(v[i]);
  sort(vals.begin(), vals.end());
  vals.erase(unique(vals.begin(), vals.end()), vals.end());
  for (int i = 0; i < n; i++)
    v[i] = lower_bound(vals.begin(), vals.end(), v[i]) - vals.begin();
signed main()
```

```
ios_base::sync_with_stdio(false);
cin.tie(NULL);
while (cin >> n >> q)
  t.clear();
  qq.clear();
  depth[0] = 0;
  memset(cnt, 0, sizeof(cnt));
  memset(freq, 0, sizeof(freq));
for (int i = 0; i < n; i++)</pre>
    adj[i].clear();
    cin >> v[i];
  compress();
  for (int i = 0; i < n - 1; i++)
    int a, b;
    cin >> a >> b;
    a--, b--;
    adj[a].pb(b);
    adj[b].pb(a);
  dfs(0, 0);
  for (int i = 0; i < q; i++)
    int x, y;
    cin >> x >> y;
    x--, y--;
int 1 = lca(x, y);
    if (tin[x] > tin[y])
      swap(x, y);
    if (1 == x)
      qq.pb({tin[x], tin[y], -1, i});
    else
      qq.pb({tout[x], tin[y], 1, i});
  int block = sqrt(n) + 1;
  auto cmp = [&](qry x, qry y)
    if (x.1 / block != y.1 / block)
      return x.1 / block < y.1 / block;
    return x.r < v.r;</pre>
  sort(qq.begin(), qq.end(), cmp);
  vector<int> ans(q);
  int c1 = 0, cr = 0, resp = 0;
  auto add2 = [&](int x)
    freq[v[x]]++;
    if (freq[v[x]] == 1)
      resp++;
  auto rem2 = [\&] (int x)
    freq[v[x]]--;
    if (freq[v[x]] == 0)
      resp--;
  auto add = [&](int x)
    cnt[x]++;
    if (cnt[x] == 2)
      rem2(x);
    else
      add2(x);
  auto rem = [\&] (int x)
    cnt[x]--;
    if (cnt[x] == 1)
      add2(x);
    else
      rem2(x);
  for (int i = 0; i < q; i++)
    int idx = qq[i].id;
```

```
int 1 = qq[i].1;
      int r = qq[i].r;
      int lc = qq[i].lca;
      while (cl < 1)
        rem(t[cl++]);
      while (cl > 1)
        add(t[--cl]);
      while (cr <= r)</pre>
       add(t[cr++]);
      while (cr > r + 1)
        rem(t[--cr]);
      if (1c != -1)
       add(lc);
      ans[idx] = resp;
      if (1c != -1)
        rem(lc);
    for (auto const &i : ans)
      cout << i << endl;</pre>
 return 0;
// https://www.spoj.com/problems/COT2/
// quantos caras distintos em um path entre u e v
// mo em arvores
// acha o euler tour da arvore com tin e tout
// desconsidera no mo os indices duplicados no range
// para queries em subtree eh mais simples:
// apenas saber o tamanho da subtree de i
// fazer o euler tour apenas com o tin
// e fzr a query pro range tin[i] ate tin[i] + sz[i] - 1
// pra queries de path com peso nos edges:
// https://codeforces.com/gym/100962/attachments (problema F)
// considera v[i] -> peso do edge que liga ao meu pai na arvore
// dai pra query com o lca == u, nao tenho que considerar v[u] ([tin[u], tin[v
    ]], dps removendo v[u])
// e pra query com o lca != u, so fazer ela normalmente ([tout[u], tin[v]])
```

## 4.38 mo trees edges

int depth[MAXN];

```
#include <bits/stdc++.h>
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200001
#define mod 998244353
struct qry
  int 1, r, ini, id;
int n, q;
vector<pi> adj[MAXN];
int v[MAXN];
int cnt[MAXN];
int freq[MAXN];
int in_block[MAXN];
int tin[MAXN];
int tout[MAXN];
```

```
int up[MAXN][25];
vector<int> t;
vector<qry> qq;
void dfs(int s, int p, int par_edge)
  v[s] = par_edge;
  tin[s] = t.size();
  up[s][0] = p;
  for (int i = 1; i < 25; i++)
   up[s][i] = up[up[s][i-1]][i-1];
  t.pb(s);
  for (auto const &i : adj[s])
   if (i.fir == p)
     continue;
    depth[i.fir] = depth[s] + 1;
   dfs(i.fir, s, i.sec);
  tout[s] = t.size();
  t.pb(s);
bool is(int u, int v)
  return tin[u] <= tin[v] && tout[u] >= tout[v];
int lca(int u, int v)
  if (is(u, v))
   return u;
  if (is(v, u))
   return v;
  for (int i = 24; i >= 0; i--)
    if (!is(up[u][i], v))
      u = up[u][i];
  return up[u][0];
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> q;
  for (int i = 0; i < n - 1; i++)</pre>
    int a, b, c;
    cin >> a >> b >> c;
    a--, b--;
    adj[a].pb({b, c});
    adj[b].pb({a, c});
  dfs(0, 0, 0);
  for (int i = 0; i < q; i++)
    int x, y;
    cin >> x >> y;
    x--, y--;
    int 1 = lca(x, y);
    if (tin[x] > tin[y])
     swap(x, y);
    if (1 == x)
      qq.pb({tin[x], tin[y], x, i});
    else
      qq.pb({tout[x], tin[y], -1, i});
  int block = sqrt(n) + 1;
  auto cmp = [&] (qry x, qry y)
    if (x.1 / block != y.1 / block)
      return x.1 / block < y.1 / block;
    return x.r < y.r;</pre>
  sort(qq.begin(), qq.end(), cmp);
  vector<int> ans(q);
  int cl = 0, cr = 0, resp = 0;
  auto add2 = [&](int x)
   if (v[x] >= MAXN)
```

```
return;
    freq[v[x]]++;
    if (freq[v[x]] == 1)
     in_block[v[x] / block]++;
 auto rem2 = [&](int x)
    if (v[x] >= MAXN)
      return;
    freq[v[x]]--;
    if (freq[v[x]] == 0)
      in_block[v[x] / block]--;
  auto add = [&](int x)
    cnt[x]++;
    if (cnt[x] == 2)
     rem2(x);
    else
     add2(x);
  auto rem = [&](int x)
    cnt[x]--;
    if (cnt[x] == 1)
     add2(x);
    else
      rem2(x);
  for (int i = 0; i < q; i++)
   int idx = qq[i].id;
int l = qq[i].1;
    int r = qq[i].r;
    int ini = qq[i].ini;
while (cl < 1)</pre>
      rem(t[cl++]);
    while (cl > 1)
      add(t[--c1]);
    while (cr \ll r)
      add(t[cr++]);
    while (cr > r + 1)
      rem(t[--cr]);
    if (ini != -1)
      rem(ini);
    for (int b = 0;; b++)
      if (in_block[b] != block)
        ans[idx] = b * block;
        while (freq[ans[idx]])
         ans[idx]++;
        break;
    if (ini != -1)
      add(ini);
 for (auto const &i : ans)
   cout << i << endl;
  return 0;
// https://codeforces.com/gym/100962/attachments (problema F)
// mo em arvore com peso nos edges
// nesse problema em especifico: dado uma arvore, responder queries de mex
// no caminho entre u e v, considerando os pesos de arestas no caminho de u pra
// mo em arvores
// acha o euler tour da arvore com tin e tout
// desconsidera no mo os indices duplicados no range
// e bem parecido com o de peso nos vertices
// considera v[i] -> peso do edge que liga ao meu pai na arvore
// dai pra query com o lca == u, nao tenho que considerar v[u] ([tin[u], tin[v
    ]], dps removendo v[u])
// e pra query com o lca != u, so fazer ela normalmente ([tout[u], tin[v]])
```

## 4.39 Prim

```
// algoritimo de prim
// 1 - definir a distancia de cada vertice como infinito (similar ao dijkstra). // 2 - definir a distancia de 0 para o source(0).
// 3 - Em cada passo, encontrar o vertice u, que ainda nao foi processado, que
     possua a menor das distancias.
// 4 - ao termino fazer a soma de todas as distancias e encontrar qual a soma
    das distancias na MST.
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define pii pair<int, int>
#define mp make_pair
#define MAXN 100001
#define INF 999999
#define sec second
#define fir first
int n, m, a, b, c;
vector<pii> adj[MAXN];
int dist[MAXN];
bool processed[MAXN];
void prim()
  for (int i = 0; i < n; i++)
    dist[i] = INF;
  dist[0] = 0;
  priority_queue<pii, vector<pii>, greater<pii>> q;
  q.push(pii(dist[0], 0));
  while (1)
    int davez = -1:
    while (!q.empty())
      int atual = q.top().sec;
      q.pop();
      if (!processed[atual])
        davez = atual;
        break;
    if (davez == -1)
      break;
    processed[davez] = true;
    for (int i = 0; i < adj[davez].size(); i++)</pre>
      int distt = adj[davez][i].fir;
      int atual = adj[davez][i].sec;
      if (dist[atual] > distt && !processed[atual])
        dist[atual] = distt;
        q.push(pii(dist[atual], atual));
```

```
int ans = 0;
 for (int i = 0; i < n; i++)
   ans += dist[i];
 cout << ans << endl;
int main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 cin >> n >> m;
 for (int i = 0; i < m; i++)
   cin >> a >> b >> c;
   a--;
   b--;
   adj[a].pb(mp(c, b));
   adj[b].pb(mp(c, a));
 prim();
 return 0:
```

# 4.40 push relabel

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 300005
#define mod 1000000007
#define INF 1e9
struct edge
 int dest, back, f, c, id;
struct push_relabel
 int n;
  vector<vector<edge>> q;
  vector<int> ec;
  vector<edge *> cur;
  vector<vector<int>> hs;
  vector<int> H:
  push_relabel(int sz) : g(sz), ec(sz), cur(sz), hs(2 * sz), H(sz) { n = sz; }
  void add_edge(int s, int t, int cap, int rcap, int id)
   if (s == t)
     return;
    g[s].pb({t, (int)g[t].size(), 0, cap, id});
    g[t].pb({s, (int)g[s].size() - 1, 0, rcap, -1});
  void add_flow(edge &e, int f)
    edge &back = g[e.dest][e.back];
```

```
if (!ec[e.dest] && f)
     hs[H[e.dest]].push_back(e.dest);
    e.f += f;
    e.c -= f;
    ec[e.dest] += f;
    back f -= f;
    back.c += f;
    ec[back.dest] -= f;
  int calc(int s, int t)
    int v = q.size();
    H[s] = v;
    ec[t] = 1;
    vector<int> co(2 * v);
    co[0] = v - 1;
    for (int i = 0; i < v; i++)
      cur[i] = g[i].data();
    for (edge &e : g[s])
      add_flow(e, e.c);
    for (int hi = 0;;)
      while (hs[hi].empty())
        if (!hi--)
          return -ec[s];
       int u = hs[hi].back();
      hs[hi].pop_back();
      while (ec[u] > 0)
         if (cur[u] == g[u].data() + g[u].size())
           H[u] = INF;
           for (edge &e : g[u])
             if (e.c && H[u] > H[e.dest] + 1)
           H[u] = H[e.dest] + 1, cur[u] = &e;
if (++co[H[u]], !--co[hi] && hi < v)
             for (int i = 0; i < v; i++)</pre>
               if (hi < H[i] && H[i] < v)
                  --co[H[i]], H[i] = v + 1;
           hi = H[u];
         else if (\operatorname{cur}[u] \rightarrow \operatorname{c \&\& H}[u] == \operatorname{H}[\operatorname{cur}[u] \rightarrow \operatorname{dest}] + 1)
           add_flow(*cur[u], min(ec[u], cur[u]->c));
           ++cur[u];
  vector<int> flow_edges(int m) // fluxo em cada aresta
    vector<int> ans(m);
    for (int i = 0; i < n; i++)</pre>
       for (auto const &j : g[i])
         if (j.id != -1)
           ans[j.id] = j.f;
    return ans;
struct flow_with_demands
  push_relabel pr;
  vector<int> in, out;
  int n;
  flow_with_demands(int sz) : n(sz), pr(sz + 2), in(sz), out(sz) {}
  void add_edge(int u, int v, int cap, int dem, int id)
    pr.add_edge(u, v, cap - dem, 0, id);
    out[u] += dem, in[v] += dem;
  int run(int s, int t)
    pr.add_edge(t, s, INF, 0, -1);
    for (int i = 0; i < n; i++)
```

```
pr.add_edge(n, i, in[i], 0, -1);
     pr.add_edge(i, n + 1, out[i], 0, -1);
   return pr.calc(n, n + 1);
 bool check() // todas as constraints foram satisfeitas?
    for (auto const &i : pr.q[n])
     if (i.c > 0)
        return 0;
   return 1;
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n;
 cin >> n;
 vector<pi> v(n);
 vector<int> x_vals, y_vals;
 for (int i = 0; i < n; i++)
   cin >> v[i].fir >> v[i].sec;
   x_vals.pb(v[i].fir);
   y_vals.pb(v[i].sec);
 sort(x_vals.begin(), x_vals.end());
 sort(y_vals.begin(), y_vals.end());
 x_vals.erase(unique(x_vals.begin(), x_vals.end()), x_vals.end());
 y_vals.erase(unique(y_vals.begin(), y_vals.end()), y_vals.end());
 int xx = x_vals.size();
 int yy = y_vals.size();
vector<int> cntx(xx, 0);
 vector<int> cnty(yy, 0);
 for (int i = 0; i < n; i++)
   v[i].fir = lower_bound(x_vals.begin(), x_vals.end(), v[i].fir) - x_vals.
   v[i].sec = lower_bound(y_vals.begin(), y_vals.end(), v[i].sec) - y_vals.
        begin();
   cntx[v[i].fir]++;
   cnty[v[i].sec]++;
  flow_with_demands mf(xx + yy + 2);
  int src = xx + yy;
 int sink = xx + yy + 1;
 int edge_id = 0;
 for (int i = 0; i < xx; i++)
   int half = cntx[i] / 2;
   int can_pass = cntx[i] - half;
   mf.add_edge(src, i, can_pass, half, edge_id++);
 for (int i = 0; i < yy; i++)
   int half = cnty[i] / 2;
   int can_pass = cnty[i] - half;
   mf.add_edge(xx + i, sink, can_pass, half, edge_id++);
  vector<int> middle_edges(n);
 for (int i = 0; i < n; i++)</pre>
   middle_edges[i] = edge_id;
   mf.add_edge(v[i].fir, xx + v[i].sec, 1, 0, edge_id++);
 mf.run(src, sink);
 assert (mf.check());
 vector<int> flow_edges = mf.pr.flow_edges(edge_id);
 for (int i = 0; i < n; i++)</pre>
   if (flow_edges[middle_edges[i]])
     cout << "L";
   else
     cout << "F";
 cout << endl;
```

### 4.41 reroot

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 200001
#define mod 1000000007
int n:
vector<int> adj[MAXN];
int sz[MAXN];
int dp[MAXN];
int dfs(int u, int v)
  sz[u] = 1;
  for (auto const &i : adj[u])
    if (i != v)
     sz[u] += dfs(i, u);
  return sz[u];
void reroot(int u, int v)
  for (auto const &i : adj[u])
    if (<u>i</u> != v)
      int a = sz[u], b = sz[i];
      dp[i] = dp[u];
      dp[i] = sz[u], dp[i] = sz[i];
      sz[u] = sz[i], sz[i] = n;
      dp[i] += sz[u], dp[i] += sz[i];
      reroot(i, u);
      sz[u] = a, sz[i] = b;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n;
  for (int i = 0; i < n - 1; i++)
   int a, b;
    cin >> a >> b;
    a--, b--;
   adj[a].pb(b);
   adj[b].pb(a);
  dfs(0, -1);
  for (int i = 0; i < n; i++)
```

```
dp[0] += sz[i]; // answer when tree is rooted on vertex 0
  reroot(0, -1);
  cout << *max_element(dp, dp + n) << endl;
  return 0;
}
// https://codeforces.com/contest/1187/problem/E
// f(v) = when tree is rooted at vertex v, the current
// answer is the sum of all subtrees sizes
// final answer = max(f(0), f(1), f(2), ..., f(n))
// easy approach: O(N^2)
// with reroot: O(N)
// 1 - run a dfs and calculate f(0)
// 2 - let be dp[i] = f(i)
// 3 - now, lets run a another dfs, and re-calculate the
// answer when tree is rooted at vertex i (dp[i])
// 4 - the final answer is the maximum value of dp[i]</pre>
```

# 4.42 rmq tree

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 5005
#define mod 998244353
struct lca
  int n:
  vector<vector<int>> adj;
  vector<int> vec:
  int 1, timer;
  vector<int> tin, tout, depth;
  vector<vector<pi>>> up;
  void dfs(int v, int p)
    tin[v] = ++timer;
    up[v][0] = {p, min(vec[v], vec[p])};
    for (int i = 1; i <= 1; i++)
      up[v][i].fir = up[up[v][i - 1].fir][i - 1].fir;
      up[v][i].sec = min(up[v][i - 1].sec, up[up[v][i - 1].fir][i - 1].sec);
    for (auto const &u : adj[v])
      if (p == u)
       continue;
      depth[u] = depth[v] + 1;
      dfs(u, v);
    tout[v] = ++timer;
  bool is_ancestor(int u, int v)
    return tin[u] <= tin[v] && tout[u] >= tout[v];
  int find_lca(int u, int v)
    if (is_ancestor(u, v))
      return u;
    if (is_ancestor(v, u))
      return v;
    for (int i = 1; i >= 0; i--)
      if (!is_ancestor(up[u][i].fir, v))
        u = up[u][i].fir;
```

```
return up[u][0].fir;
  int dist(int s, int v)
   int at = find_lca(s, v);
   return (depth[s] + depth[v] - 2 * depth[at]);
  int solve(int u, int d)
    int ans = vec[u];
    for (int i = 1; i >= 0; i--)
      if (d & (1 << i))
       ans = min(ans, up[u][i].sec);
       u = up[u][i].fir;
   return ans:
  int rmq(int u, int v)
   int 1 = find_lca(u, v);
   return min(solve(u, dist(u, 1)), solve(v, dist(v, 1)));
  lca(vector<vector<int>> &_adj, vector<int> &_vec)
   adj = _adj;
   vec = _vec;
   n = adj.size();
   tin.resize(n);
   tout.resize(n);
   depth.resize(n);
   timer = 0;
    1 = ceil(log2(n));
   up.assign(n, vector<pi>(1 + 1));
   dfs(0, 0);
};
signed main()
// valores nos vertices
// rmg considerando o caminho entre os vertices u e v
```

#### 4.43 sack

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
vector<int> adj[MAXN];
vector<int> v[MAXN];
int c[MAXN];
int cnt[MAXN];
int sz[MAXN];
void dfs_sz(int x, int p)
  sz[x] = 1;
```

```
for (auto const &i : adj[x])
   if (i != p)
     dfs_sz(i, x);
     sz[\bar{x}] += sz[i];
void modify(int c, int val)
 cnt[c] += val;
void dfs(int x, int p, bool keep)
 int best = -1, big_child = -1;
 for (auto const &i : adj[x])
   if (i != p && sz[i] > best)
     best = sz[i];
     big_child = i;
 for (auto const &i : adj[x])
   if (i != p && i != big_child)
     dfs(i, x, 0);
 if (big_child != -1)
   dfs(big_child, x, 1);
   swap(v[x], v[big\_child]); // O(1)
 v[x].pb(x);
 modify(c[x], 1); // adiciona
 for (auto const &i : adj[x])
   if (i != p && i != big_child)
      for (auto const &j : v[i])
        v[x].pb(j);
        modify(c[j], 1); // adiciona
  // a cor c aparece cnt[c] vezes na subtree de x
  // dai vc pode fazer algo tendo essa informacao
  // seja responser queries ou algo do tipo aqui
 if (!keep)
    for (auto const &i : v[x])
     modify(c[i], -1); // remove
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n;
 cin >> n:
 for (int i = 0; i < n; i++)
   cin >> c[i];
 for (int i = 0; i < n - 1; i++)
   int a, b;
   cin >> a >> b;
   a--, b--;
   adj[a].pb(b);
   adj[b].pb(a);
 dfs_sz(0, -1);
dfs(0, -1, 0);
 cout << endl;
// https://codeforces.com/blog/entry/44351
// https://codeforces.com/blog/entry/67696
```

#### 4.44 scc

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500005
#define mod 1000000007
int n, m;
bool vis[MAXN];
int root[MAXN];
vector<int> order;
vector<int> roots;
vector<int> comp;
vector<vector<int>> comps;
vector<int> adj[MAXN];
vector<int> adj_rev[MAXN];
vector<int> adj_scc[MAXN];
void dfs(int v)
  vis[v] = true;
  for (auto const &u : adj[v])
   if (!vis[u])
     dfs(u);
  order.pb(v);
void dfs2(int v)
  comp.pb(v);
  vis[v] = true;
  for (auto const &u : adj_rev[v])
   if (!vis[u])
      dfs2(u);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> m;
  for (int i = 0; i < m; i++)
   int a, b;
    cin >> a >> b;
    adj[a].pb(b);
    adj_rev[b].pb(a);
  for (int i = 0; i < n; i++)
   if (!vis[i])
      dfs(i);
```

```
reverse(order.begin(), order.end());
 memset(vis, false, sizeof(vis));
 for (auto const &v : order)
   if (!vis[v])
     comp.clear();
     dfs2(v);
     comps.pb(comp);
     // making condensation graph
     int r = comp.back();
     for (auto const &u : comp)
      root[u] = r;
     roots.push_back(r);
     */
  // making condensation graph
  for (int v = 0; v < n; v++)
    for (auto const &u : adj[v])
     int root_v = roots[v];
     int root_u = roots[u];
     if (root_u != root_v)
       adj_scc[root_v].pb(root_u);
 // printing scc
 cout << comps.size() << endl;</pre>
 for (auto const &comp : comps)
   cout << comp.size() << " ";
   for (auto const &u : comp)
     cout << u << " ";
   cout << endl;</pre>
 return 0:
// to test: https://judge.yosupo.jp/problem/scc
```

# 4.45 segtree graph

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500005
#define mod 998244353
struct segtree_graph
 int n;
  vector<vector<int>> adj;
  vector<int> id_rev, id, deg;
  segtree_graph(int sz)
   n = sz;
```

```
adj.resize(4 * n);
  id.resize(n);
  id_rev.assign(4 * n, -1);
  deg.assign(4 * n, 0);
 build(0, n - 1, 1);
void build(int 1, int r, int i)
 if (1 == r)
    id[1] = i;
    id_rev[i] = 1;
    return;
  int mid = (1 + r) >> 1;
 build(1, mid, i << 1);
 build(mid + 1, r, (i << 1) | 1);
  adj[i << 1].pb(i);
  adj[(i << 1) | 1].pb(i);
 deg[i] += 2;
void add(int 1, int r, int q1, int qr, int i, int x)
  if (1 > r || 1 > qr || r < q1)
    return;
  if (1 >= ql \&\& r <= qr)
    adi[i].pb(x);
    deg[x]++;
    return;
  int mid = (1 + r) >> 1;
  add(1, mid, q1, qr, i << 1, x);
  add(mid + 1, r, ql, qr, (i << 1) | 1, x);
void add2(int a, int b)
  adi[a].pb(b);
  deg[b]++;
void topological_sort()
  vector<bool> vis(4 * n, 0);
  queue<int> q;
  for (int i = 0; i < (4 * n); i++)
    if (!deg[i])
      q.push(i);
  int qt = 0, xx = 1;
  vector<int> ans(n);
  while (!q.empty())
    int x = q.front();
    q.pop();
    if (id_rev[x] != -1)
      ans[id_rev[x]] = xx++;
    for (auto const &i : adj[x])
      deg[i]--;
      if (!deg[i])
        q.push(i);
  if (qt != (4 * n))
    cout << "-1\n";
    return;
  for (auto const &i : ans)
    cout << i << " ";
```

```
cout << endl;
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int q;
 cin >> q;
 while (q--)
    int n;
   cin >> n;
    segtree_graph st(n);
    for (int i = 0; i < n; i++)
      cin >> x; // next[i]
      if (x != -1)
        if (x > (i + 1))
         st.add(0, n - 1, i + 1, x - 1, 1, st.id[i]);
        if (x < n)
         st.add2(st.id[i], st.id[x]);
   st.topological sort();
// https://codeforces.com/contest/1158/problem/C
// se next[i] = j
// entao p[j] > p[i]
// e para todo k tal que: i < k < j, p[k] < p[i]
// em outras palavras o primeiro caba a direita de i que eh maior do que p[i] eh
     no indice j
// montar o grafo de implicacoes e fazer topsort
// pq se next[i] = j
// entao implica que p[k] < p[i] para i < k < j
// e implica que p[i] < p[j]</pre>
// mas tem que ser o grafo de segtree ne
// faz algo similar ao problema legacy: https://codeforces.com/problemset/
    problem/786/B
```

### 4.46 stable matching

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
// #define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define pci pair<char, int>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
bool in[1005][1005];
int tt[1005][1005];
int b[1005];
int ini[1005];
// a.size() <= b.size()
vector<int>> stable_marriage(vector<vector<int>> &a, vector<vector<int>> &b)
```

```
int n = a.size(), m = b.size();
 assert(a[0].size() == m and b[0].size() == n and n <= m);
  vector<int> match(m, -1), it(n, 0);
  vector<vector<int>> inv_b(m, vector<int>(n));
  for (int i = 0; i < m; \overline{i}++)
   for (int j = 0; j < n; j++)</pre>
     inv_b[i][b[i][j]] = j;
 queue<int> q;
  for (int i = 0; i < n; i++)
   q.push(i);
  while (q.size())
   int i = q.front();
    q.pop();
    int j = a[i][it[i]];
    if (match[j] == -1)
     match[j] = i;
    else if (inv_b[j][i] < inv_b[j][match[j]])
      q.emplace(match[j]);
      it[match[j]]++;
     match[j] = i;
   else
     q.emplace(i), it[i]++;
  vector<int> ret(n);
  for (int i = 0; i < m; i++)
   if (match[i] != -1)
      ret[match[i]] = i;
  return ret;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
 int n, m, d, e;
  cin >> n >> m >> d >> e;
  if (n > m)
   cout << "impossible\n";</pre>
   return 0;
  vector<vector<int>> adj(n);
  vector<vector<int>> adj2(m);
 memset(b, -1, sizeof(b));
  memset(ini, -1, sizeof(ini));
  for (int i = 0; i < e; i++)
   int s, k, t;
   cin >> s >> k >> t;
    k--, t--;
    if (t == -1)
      tt[k][b[k]] += (s - ini[k]);
     b[k] = -1;
      ini[k] = -1;
    else
      if (b[k] != -1)
        tt[k][b[k]] += (s - ini[k]);
      if (!in[k][t])
        in[k][t] = 1;
       adj[k].pb(t);
     b[k] = t;
      ini[k] = s;
  for (int k = 0; k < n; k++)
   if (b[k] != -1)
      tt[k][b[k]] += (d - ini[k]);
```

```
for (int k = 0; k < n; k++)
    for (int t = 0; t < m; t++)
      if (!in[k][t])
        adj[k].pb(t);
  for (int t = 0; t < m; t++)</pre>
    vector<pi> curr;
    for (int k = 0; k < n; k++)
      curr.pb({tt[k][t], k});
    sort(curr.begin(), curr.end());
    for (auto const &i : curr)
     adj2[t].pb(i.sec);
  vector<int> ans = stable_marriage(adj, adj2);
 for (auto const &i : ans)
  cout << i + 1 << " ";</pre>
 cout << endl;</pre>
// solucao pro: https://open.kattis.com/problems/jealousyoungsters
// stable marriage
// voce quer achar um matching em um grafo bipartido
// que todos os caras de um lado dao matching com algum do outro lado
// cada vertice tem um vector
// que diz qual a ordem de preferencia dele
// o que ele mais quer dar matching eh com v[0]
// o segundo com que ele mais quer dar matching eh com v[1]
// quando a.size() <= b.size(), entao sempre existe um stable matching</pre>
```

### 4.47 strong orientation

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1000005
#define mod 1000000007
int n, m, timer, comps, bridges;
vector<pi> edges;
vector<pi> adj[MAXN];
int tin[MAXN];
int low[MAXN];
bool vis[MAXN];
char orient[MAXN];
void find_bridges(int v)
  low[v] = timer, tin[v] = timer++;
  for (auto const &p : adj[v])
    if (vis[p.sec])
      continue;
    vis[p.sec] = true;
```

```
orient[p.sec] = (v == edges[p.sec].first) ? '>' : '<';
   if (tin[p.fir] == -1)
      find_bridges(p.fir);
      low[v] = min(low[v], low[p.fir]);
      if (low[p.fir] > tin[v])
       bridges++;
   else
      low[v] = min(low[v], low[p.fir]);
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 cin >> n >> m;
  for (int i = 0; i < m; i++)
   int a, b;
   cin >> a >> b;
    a--, b--;
    edges.pb({a, b});
   adj[a].pb({b, i});
   adj[b] pb({a, i});
 memset(tin, -1, sizeof(tin));
  memset(low, -1, sizeof(low));
  for (int v = 0; v < n; v++)
   if (tin[v] == -1)
      comps++:
      find_bridges(v);
  // numero minimo de scc = numero de componentes + numero de pontes
  cout << comps + bridges << endl;</pre>
  // > - a aresta foi orientada da esquerda pra direita
  // < - a aresta foi orientada da direita pra esquerda
  for (int i = 0; i < m; i++)
   cout << orient[i];</pre>
  cout << endl;</pre>
  return 0;
// to_test: https://szkopul.edu.pl/problemset/problem/nldsb4EW1YuZykB1f4lcZL1Y/
    site/?key=statement
// strong orientation:
// encontrar uma orientacao para as arestas tal que o numero
// minimo de scc e o menor possivel
```

# 4.48 Topological Sort

```
#include <bits/stdc++.h>
using namespace std;

#define lli long long int
#define pb push_back
#define MAXN 10000

int n , m , a , b ;
vector <int> adj [MAXN];
int grau [MAXN];
vector <int> order;

bool topological_sort ()
{
   int ini = 0;
   while (ini < order.size())
   {
      int atual = order[ini];
      ini++;
   }
}</pre>
```

```
for (int i = 0 ; i < adj[atual].size() ; i++)</pre>
            int v = adj[atual][i] ;
            grau[v]-- ;
            if (grau[v] == 0)
                order.pb(v);
    return (order.size() == n) ? true : false ;
int main ()
    ios_base::sync_with_stdio(false) ;
    cin.tie(NULL);
    cin >> n >> m ;
    for (int i = 1; i \le m; i++)
        cin >> a >> b ;
        grau[a]++ ;
        adj[b].pb(a);
    for (int i = 1; i \le n; i++)
        if (grau[i] == 0)
            order.pb(i);
    if (topological_sort())
        for (int i = 0 ; i < order.size() ; i++)</pre>
            cout << order[i] << " ";</pre>
        cout << endl ;
    else
        cout << "Impossible\n";</pre>
    return 0 :
```

### 4.49 TreeDiameter

```
#include <bits/stdc++.h>
using namespace std;
#define PI acos(-1)
#define int long long int
#define pb push back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 100001
#define mod 1000000007
int diameter, best;
vector<int> adj[MAXN];
bool visited[MAXN];
void dfs(int s, int c)
  if (c > diameter)
```

```
diameter = c:
   best = s;
 visited[s] = true;
 for (auto const &i : adj[s])
   if (!visited[i])
     dfs2(i, c + 1);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int q;
 cin >> q;
 while (q--)
   int n;
   cin >> n;
   for (int i = 0; i < n; i++)
     adj[i].clear();
    for (int i = 0; i < n - 1; i++)
     int a, b;
     cin >> a >> b;
     a--, b--;
     adj[b].pb(a);
     adj[a].pb(b);
   diameter = 0, best = 0;
   memset(visited, false, sizeof(visited));
   dfs(1, 0);
                                     // achar o vertice mais distante a partir
        do vertice 0
    memset(visited, false, sizeof(visited));
   dfs(best, 0);
                                    // achar o mais distante a partir do
        primeiro vertice que achamos
   cout << diameter << endl;</pre>
 return 0;
```

#### 4.50 tree isomorfism

```
#include <bits/stdc++.h>
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 501
#define mod 1000000007
int curr_hash = 1;
map<vector<int>, int> mp;
struct hash tree
  pi h;
  int n;
  vector<int> c, sz, large_comp;
  vector<vector<int>> adj;
  hash_tree (vector<vector<int>> &a)
   n = a.size();
```

```
adj = a;
 void dfs(int s, int p)
   sz[s] = 1;
   large\_comp[s] = 0;
   for (auto const &v : adj[s])
     if (v != p)
       dfs(v, s);
        sz[s] += sz[v];
        large_comp[s] = max(large_comp[s], sz[v]);
    large_comp[s] = max(large_comp[s], n - sz[s]);
 int dfs2(int s, int p)
   if (s == -1)
     return -1;
    vector<int> child;
   for (auto const &v : adj[s])
     if (v != p)
       child.pb(dfs2(v, s));
   sort(child.begin(), child.end());
   if (!mp[child])
     mp[child] = curr_hash++;
   return mp[child];
 pi get_hash()
   sz.assign(n, 0);
   large_comp.assign(n, 0);
   dfs(0, -1);
   int best = 1e18;
   for (int i = 0; i < n; i++)
     if (large_comp[i] < best)</pre>
       best = large_comp[i];
       c.clear();
     if (large_comp[i] == best)
       c.pb(i);
   while (c.size() < 2)</pre>
     c.pb(-1);
   h.fir = dfs2(c[0], -1);
   h.sec = dfs2(c[1], -1);
   if (h.fir > h.sec)
     swap(h.fir, h.sec);
   return h;
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int q;
 cin >> q;
 while (q--)
   int n;
   cin >> n;
   vector<vector<int>> a(n);
   vector<vector<int>> b(n);
    for (int i = 0; i < n - 1; i++)
     int x, y;
     cin >> x >> y;
     x--, y--;
     a[x].pb(y);
     a[y].pb(x);
    for (int i = 0; i < n - 1; i++)
```

```
int x, y;
      cin >> \bar{x} >> y;
     x--, y--;
b[x].pb(y);
      b[y].pb(x);
    (hash_tree(a).get_hash() == hash_tree(b).get_hash()) ? cout << "YES\n" :</pre>
         cout << "NO\n";
  return 0;
// https://www.spoj.com/problems/TREEISO/
// https://www.beecrowd.com.br/judge/en/problems/view/1229
// hash de arvores
// para descobrir se duas arvores sao isomorfas
// 1 - achar todos os centroides da arvore (toda arvore tem no maximo 2
    centroides)
// 2 - achar o hashing com a arvore enraizada em cada centroid
// 3 - dai o hashing da arvore eh um pair ordenado, indicando o hashing de cada
    enraizamento no centroid
```

#### 4.51 two sat

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 1000000007
struct two sat
  vector<vector<int>> q, qr; // qr is the reversed graph
  vector<int> comp, ord, ans; // comp[v]: ID of the SCC containing node v
  vector<bool> vis:
  two_sat() {}
  two_sat(int sz)
    g.assign(2 * n, vector<int>());
    gr.assign(2 * n, vector<int>());
   comp.resize(2 * n);
   vis.resize(2 * n);
   ans.resize(2 * n);
  void add_edge(int u, int v)
    q[u].push_back(v);
    gr[v].push_back(u);
  // int x, bool val: if 'val' is true, we take the variable to be x. Otherwise
       we take it to be x's complement (not x).
  void implies (int i, bool f, int j, bool g) // a \rightarrow b
    add_edge(i + (f ? 0 : n), j + (g ? 0 : n));
    add_edge(j + (g ? n : 0), i + (f ? n : 0));
  void add_clause_or(int i, bool f, int j, bool g) // At least one of them is
       true
```

```
add_edge(i + (f ? n : 0), j + (g ? 0 : n));
   add_edge(j + (g ? n : 0), i + (f ? 0 : n));
 void add_clause_xor(int i, bool f, int j, bool g) // only one of them is true
   add_clause_or(i, f, j, g);
   add_clause_or(i, !f, j, !g);
 void add_clause_and(int i, bool f, int j, bool g) // both of them have the
    add_clause_xor(i, !f, j, g);
 void set(int i, bool f) // Set a variable
    add_clause_or(i, f, i, f);
 void top_sort(int u)
    vis[u] = 1;
    for (auto const &v : g[u])
     if (!vis[v])
        top_sort(v);
    ord.push_back(u);
 void scc(int u, int id)
   vis[u] = 1;
    comp[u] = id;
    for (auto const &v : gr[u])
     if (!vis[v])
       scc(v, id);
 bool solve()
    fill(vis.begin(), vis.end(), 0);
    for (int i = 0; i < 2 * n; i++)
     if (!vis[i])
        top sort(i);
    fill(vis.begin(), vis.end(), 0);
    reverse(ord.begin(), ord.end());
    int id = 0;
    for (const auto &v : ord)
     if (!vis[v])
        scc(v, id++);
    for (int i = 0; i < n; i++)</pre>
     if (comp[i] == comp[i + n])
       return 0;
     ans[i] = (comp[i] > comp[i + n]) ? 1 : 0;
    return 1;
signed main()
// https://codeforces.com/blog/entry/92977
   https://codeforces.com/blog/entry/16205
   https://cp-algorithms.com/graph/2SAT.html
```

#### 4.52 virtual tree

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
```

```
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define pci pair<char, int>
#define fir first
#define sec second
#define MAXN 200005
#define mod 998244353
int n, ans;
vector<int> adj[MAXN];
vector<int> with_color[MAXN];
vector<pi> virt[MAXN];
int a[MAXN];
int dp[MAXN];
int dp2[MAXN];
int pot[MAXN];
namespace lca
  int 1, timer;
  vector<int> tin, tout, depth;
  vector<vector<int>> up;
  void dfs(int v, int p)
   tin[v] = ++timer;
   up[v][0] = p;
for (int i = 1; i <= 1; i++)
  up[v][i] = up[up[v][i - 1]][i - 1];</pre>
    for (auto const &u : adj[v])
      if (p == u)
       continue;
      depth[u] = depth[v] + 1;
      dfs(u, v);
    tout[v] = ++timer;
  bool is_ancestor(int u, int v)
    return tin[u] <= tin[v] && tout[u] >= tout[v];
  int find_lca(int u, int v)
    if (is_ancestor(u, v))
      return u;
    if (is_ancestor(v, u))
      return v;
    for (int i = 1; i >= 0; i--)
      if (!is_ancestor(up[u][i], v))
        u = up[u][i];
    return up[u][0];
  void init()
    tin.resize(n);
    tout.resize(n);
    depth.resize(n);
    timer = 0;
    1 = ceil(log2(n));
    up.assign(n, vector<int>(1 + 1));
    dfs(0, 0);
  int dist(int s, int v)
    int at = find_lca(s, v);
    return (depth[s] + depth[v] - 2 * depth[at]);
// dp naive (fazer O(n) para cada cor)
// fixa uma cor c e faz uma dp on tree pra calcular
```

```
// quantas subtrees tem tal que todas as folhas sao da cor c
// e chamando o dfs saindo de qualquer vertice
// void dfs(int s, int p, int c)
    // dado que eu calculo o dp2[i] para cada filho
    // para cada possivel subset (i1, i2, ..., ik) nao vazio de filhos
    // acha o valor de dp2[i1] * dp2[i2] * ... * dp2[ik]
    // no final eu quero a soma de todos esses valores, isso vai tar nessa
    variavel prod
    int prod = 1;
    for (auto const &i : adj[s])
      if (i != p)
        dfs(i, s, c);
        prod = (prod * (dp2[i] + 1)) % mod;
    prod = (prod - 1 + mod) % mod;
    dp[s] = prod;
    dp2[s] = prod;
    if (a[s] == c)
       dp[s] = (dp[s] + 1) % mod;
      dp2[s] = (dp2[s] + 1) % mod;
    for (auto const &i : adj[s])
      if (i == p)
        continue;
      if (a[s] != c)
        dp[s] = (dp[s] - dp2[i] + mod) % mod;
// virtual tree
// dado um conjunto de vertices v
// montar uma arvore comprimida
// tal que escolhendo dois vertices do conjunto v[i] e v[j]
// lca(v[i], v[j]) tambem ta na arvore
// se o conjunto v tem k vertices
// entao a arvore comprimida tem menos do que 2k vertices
// O(k log(k)), sem considerar a complexidade de achar lca
int build virt(vector<int> v)
 auto cmp = [&](int i, int j)
    return lca::tin[i] < lca::tin[j];
  sort(v.begin(), v.end(), cmp);
  for (int i = v.size() - 1; i > 0; i--)
    v.pb(lca::find_lca(v[i], v[i - 1]));
  sort(v.begin(), v.end(), cmp);
  v.erase(unique(v.begin(), v.end()), v.end());
  for (int i = 0; i < v.size(); i++)</pre>
    virt[v[i]].clear();
  for (int i = 1; i < v.size(); i++)</pre>
    virt[lca::find_lca(v[i - 1], v[i])].clear();
  for (int i = 1; i < v.size(); i++)</pre>
    int parent = lca::find_lca(v[i - 1], v[i]);
    int d = lca::dist(parent, v[i]);
    virt[parent].pb({v[i], d});
 return v[0];
void dfs(int s, int c) // dp naive, so que fazer isso na virtual tree
  int prod = 1;
 for (auto const &i : virt[s])
   dfs(i.fir, c);
   prod = (prod * (dp2[i.fir] + 1)) % mod;
```

```
prod = (prod - 1 + mod) % mod;
  dp[s] = prod;
  dp2[s] = prod;
  if (a[s] == c)
    dp[s] = (dp[s] + 1) % mod;
   dp2[s] = (dp2[s] + 1) % mod;
  for (auto const &i : virt[s])
   if (a[s] != c)
      dp[s] = (dp[s] - dp2[i.fir] + mod) % mod;
  ans = (ans + dp[s]) % mod;
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 pot[0] = 1;
  for (int i = 1; i < MAXN; i++)</pre>
   pot[i] = (pot[i - 1] * 2) % mod;
  for (int i = 0; i < n; i++)</pre>
   cin >> a[i];
   with_color[a[i]].pb(i);
  for (int i = 1; i < n; i++)
    int a, b;
   cin >> a >> b;
    a--, b--;
   adj[a].pb(b);
   adj[b].pb(a);
  for (int i = 1; i \le n; i++)
   if (with_color[i].size() > 0)
      int root = build_virt(with_color[i]);
      dfs(root, i);
 cout << ans << endl;
// https://atcoder.jp/contests/abc340/tasks/abc340_g
// problema legal
// dado uma arvore com n vertices (n <= 2 * 10^5)
// cada vertice tem uma cor a[i]
// conte quantas subarvores existem tal que:
// todas as folhas nessa subarvore tem a mesma cor
// sei codar em O(n^2), rodando um dfs pra cada cor
// dai montar a virtual tree para cada cor
// e rodar a dp naive na virtual tree
// tambem resolve o https://codeforces.com/gym/103960/problem/L da sub de 2022
```

### 5 kactl forked

# 6 kactl forked/combinatorial

### 6.1 IntPerm

```
/**
  * Author: Simon Lindholm
  * Date: 2018-07-06
```

#### 6.2 multinomial

# 7 kactl forked/data-structures

#### 7.1 FenwickTree

```
* Author: Lukas Polacek
 * Date: 2009-10-30
 * License: CC0
 * Source: folklore/TopCoder
 * Description: Computes partial sums a[0] + a[1] + ... + a[pos - 1], and
     updates single elements a[i],
 \star taking the difference between the old and new value.
 * Time: Both operations are $0(\log N)$.
 * Status: Stress-tested
#pragma once
struct FT {
        vector<ll> s;
        void update(int pos, ll dif) { // a[pos] += dif
                for (; pos \langle sz(s); pos |= pos + 1) s[pos] += dif;
        11 query(int pos) { // sum of values in [0, pos)
                11 \text{ res} = 0;
                for (; pos > 0; pos \&= pos - 1) res += s[pos-1];
                return res:
        int lower_bound(ll sum) {// min pos st sum of [0, pos] >= sum
                // Returns n if no sum is >= sum, or -1 if empty sum is.
                if (sum \leq 0) return -1;
                int pos = 0;
                for (int pw = 1 << 25; pw; pw >>= 1) {
                        if (pos + pw \le sz(s) \&\& s[pos + pw-1] < sum)
                                 pos += pw, sum -= s[pos-1];
                return pos;
```

#### 7.2 FenwickTree2d

};

```
* Author: Simon Lindholm
 * Date: 2017-05-11
 * License: CC0
 * Source: folklore
 \star Description: Computes sums a[i,j] for all i<I, j<J, and increases single
     elements a[i,j].
   Requires that the elements to be updated are known in advance (call
     fakeUpdate() before init()).
 * Time: \$O(\log^2 N)$. (Use persistent segment trees for \$O(\log N)$.)
 * Status: stress-tested
 */
#pragma once
#include "FenwickTree.h"
struct FT2 {
        vector<vi> ys; vector<FT> ft;
        FT2(int limx) : ys(limx) {}
        void fakeUpdate(int x, int y) {
                for (; x < sz(ys); x |= x + 1) ys[x].push_back(y);
        void init()
                for (vi& v : ys) sort(all(v)), ft.emplace_back(sz(v));
        int ind(int x, int y) {
                return (int) (lower_bound(all(ys[x]), y) - ys[x].begin()); }
        void update(int x, int y, ll dif) {
                for (; x < sz(ys); x | = x + 1)
                        ft[x].update(ind(x, y), dif);
        11 query(int x, int y) {
                11 \text{ sum} = 0;
                for (; x; x &= x - 1)
                        sum += ft[x-1].query(ind(x-1, y));
                return sum;
};
```

# 7.3 HashMap

```
* Author: Simon Lindholm, chilli
 * Date: 2018-07-23
 * License: CCO
 * Source: http://codeforces.com/blog/entry/60737
 * Description: Hash map with mostly the same API as unordered\_map, but \tilde
 * 3x faster. Uses 1.5x memory.
 * Initial capacity must be a power of 2 (if provided).
 */
#pragma once
#include <bits/extc++.h> /** keep-include */
// To use most bits rather than just the lowest ones:
struct chash { // large odd number for C
        const uint64_t C = 11(4e18 * acos(0)) | 71;
        11 operator()(ll x) const { return __builtin_bswap64(x*C); }
__gnu_pbds::gp_hash_table<ll,int,chash> h({},{},{},{},{1<<16});
/** For CodeForces, or other places where hacking might be a problem:
const int RANDOM = chrono::high_resolution_clock::now().time_since_epoch().count
struct chash { // To use most bits rather than just the lowest ones:
        const uint64_t C = 11(4e18 * acos(0)) | 71; // large odd number
        11 operator()(11 x) const { return __builtin_bswap64((x^RANDOM)*C); }
```

```
}; 
gp_hash_table<11, int, chash> h({},{},{},{},{},{});
```

### 7.4 LazySegmentTree

```
* Author: Simon Lindholm
 * Date: 2016-10-08
 * License: CCO
 * Source: me
 * Description: Segment tree with ability to add or set values of large
     intervals, and compute max of intervals.
 * Can be changed to other things.
 * Use with a bump allocator for better performance, and SmallPtr or implicit
     indices to save memory.
 * Time: O(\log N).
 * Usage: Node* tr = new Node(v, 0, sz(v));
 * Status: stress-tested a bit
#pragma once
#include "../various/BumpAllocator.h"
const int inf = 1e9;
struct Node {
        Node *1 = 0, *r = 0;
        int lo, hi, mset = inf, madd = 0, val = -inf;
Node(int lo,int hi):lo(lo),hi(hi){} // Large interval of -inf
        Node (vi& v, int lo, int hi) : lo(lo), hi(hi) {
                if (lo + 1 < hi) {
                        int mid = lo + (hi - lo)/2;
                         l = new Node(v, lo, mid); r = new Node(v, mid, hi);
                        val = max(1->val, r->val);
                else val = v[lo];
        int query(int L, int R) {
                if (R <= lo || hi <= L) return -inf;
                if (L <= lo && hi <= R) return val;
                push();
                return max(1->query(L, R), r->query(L, R));
        void set(int L, int R, int x) {
                if (R <= lo || hi <= L) return;
                if (L <= lo && hi <= R) mset = val = x, madd = 0;
                else {
                         push(), l->set(L, R, x), r->set(L, R, x);
                         val = max(1->val, r->val);
        void add(int L, int R, int x) {
                if (R <= lo || hi <= L) return;
                if (L <= lo && hi <= R) {
                        if (mset != inf) mset += x;
                         else madd += x;
                         val += x;
                         push(), l->add(L, R, x), r->add(L, R, x);
                         val = max(1->val, r->val);
        void push() {
                if (!1) {
                         int mid = lo + (hi - lo)/2;
                         l = new Node(lo, mid); r = new Node(mid, hi);
                if (mset != inf)
                        l->set(lo,hi,mset), r->set(lo,hi,mset), mset = inf;
                         1- add (lo, hi, madd), r- add (lo, hi, madd), madd = 0;
};
```

#### 7.5 LineContainer

```
* Author: Simon Lindholm
 * Date: 2017-04-20
 * License: CC0
 * Source: own work
 \star Description: Container where you can add lines of the form kx+m, and query
     maximum values at points x.
 * Useful for dynamic programming (''convex hull trick'').
 * Time: O(\log N)
 * Status: stress-tested
#pragma once
struct Line {
        mutable 11 k, m, p;
        bool operator<(const Line& o) const { return k < o.k; }</pre>
        bool operator<(ll x) const { return p < x; }</pre>
};
struct LineContainer : multiset<Line, less<>>> {
        // (for doubles, use inf = 1/.0, div(a,b) = a/b)
        static const ll inf = LLONG_MAX;
        ll div(ll a, ll b) { // floored division
               return a / b - ((a ^ b) < 0 && a % b); }
        bool isect(iterator x, iterator y) {
                if (y == end()) return x \rightarrow p = inf, 0;
                if (x-k == y-k) x-p = x-m > y-m ? inf : -inf;
                else x->p = div(y->m - x->m, x->k - y->k);
                return x->p >= y->p;
        void add(ll k, ll m) {
                auto z = insert(\{k, m, 0\}), y = z++, x = y;
                while (isect(y, z)) z = erase(z);
                if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(y));
                while ((y = x) != begin() && (--x)->p >= y->p)
                        isect(x, erase(y));
        11 query(11 x) {
                assert(!empty());
                auto 1 = *lower_bound(x);
                return l.k * x + l.m;
};
```

#### 7.6 Matrix

```
* Author: Ulf Lundstrom
* Date: 2009-08-03
* License: CC0
* Source: My head
* Description: Basic operations on square matrices.
* Usage: Matrix<int, 3> A;
* A.d = \{\{\{1,2,3\}\}, \{\{4,5,6\}\}, \{\{7,8,9\}\}\}\};
 * array<int, 3> vec = {1,2,3};
* vec = (A^N) * vec;
 * Status: tested
#pragma once
template < class T, int N> struct Matrix {
        typedef Matrix M;
        array<array<T, N>, N> d{};
        M operator* (const M& m) const {
                M a:
                rep(i,0,N) rep(j,0,N)
                        rep(k,0,N) \ a.d[i][j] += d[i][k]*m.d[k][j];
        array<T, N> operator*(const array<T, N>& vec) const {
                array<T, N> ret{};
```

```
rep(i,0,N) rep(j,0,N) ret[i] += d[i][j] * vec[j];
return ret;
}
M operator^(ll p) const {
    assert(p >= 0);
    M a, b(*this);
    rep(i,0,N) a.d[i][i] = 1;
    while (p) {
        if (p&1) a = a*b;
        b = b*b;
        p >>= 1;
    }
    return a;
};
```

### 7.7 MoQueries

```
* Author: Simon Lindholm
 * Date: 2019-12-28
 * License: CC0
 * Source: https://github.com/hoke-t/tamu-kactl/blob/master/content/data-
      structures/MoQueries.h
 * Description: Answer interval or tree path queries by finding an approximate
      TSP through the queries,
 * and moving from one query to the next by adding/removing points at the ends.
 * If values are on tree edges, change \texttt{step} to add/remove the edge $(a,
       c)$ and remove the initial \texttt{add} call (but keep \texttt{in}).
 * Time: O(N \sqrt Q)
 * Status: stress-tested
 */
#pragma once
void add(int ind, int end) { ... } // add a[ind] (end = 0 or 1)
void del(int ind, int end) { ... } // remove a[ind]
int calc() { ... } // compute current answer
vi mo(vector<pii> Q) {
        int L = 0, R = 0, blk = 350; // N/sqrt(Q)
        vi s(sz(Q)), res = s;
#define K(x) pii(x.first/blk, x.second ^ -(x.first/blk & 1))
        iota(all(s), 0);
        sort(all(s), [\&](int s, int t) \{ return K(Q[s]) < K(Q[t]); \});
        for (int qi : s) {
                pii q = Q[qi];
while (L > q.first) add(--L, 0);
                 while (R < q.second) add(R++, 1);
                 while (L < q.first) del(L++, 0);
                 while (R > q.second) del(--R, 1);
                 res[qi] = calc();
        return res;
vi moTree(vector<array<int, 2>> Q, vector<vi>& ed, int root=0){
        int N = sz(ed), pos[2] = {}, blk = 350; // N/sqrt(Q)
        vi s(sz(Q)), res = s, I(N), L(N), R(N), in(N), par(N);
        add(0, 0), in[0] = 1;
auto dfs = [&](int x, int p, int dep, auto& f) -> void {
                par[x] = p;
                L[x] = N;
                if (dep) I[x] = N++;
                for (int y : ed[x]) if (y != p) f(y, x, !dep, f);
                 if (!dep)^{T}[x] = N++;
                R[x] = N;
iota(all(s), 0);
        sort(all(s), [\&](int s, int t) \{ return K(Q[s]) < K(Q[t]); \});
        for (int qi : s) rep(end, 0, 2) {
                int &a = pos[end], b = Q[qi][end], i = 0;
\#define step(c) { if (in[c]) { del(a, end); in[a] = 0; } \
                  else { add(c, end); in[c] = 1; } a = c; }
```

#### 7.8 OrderStatisticTree

```
* Author: Simon Lindholm
 * Date: 2016-03-22
 * License: CC0
 * Source: hackIT, NWERC 2015
 * Description: A set (not multiset!) with support for finding the n'th
 * element, and finding the index of an element.
 * To get a map, change \texttt{null\_type}.
 * Time: O(\log N)
 */
#pragma once
#include <bits/extc++.h> /** keep-include */
using namespace __gnu_pbds;
template < class T>
using Tree = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
void example() {
        Tree<int> t, t2; t.insert(8);
        auto it = t.insert(10).first;
        assert(it == t.lower_bound(9));
        assert(t.order_of_key(10) == 1);
        assert (t.order_of_key(11) == 2);
        assert(*t.find_by_order(0) == 8);
        t.join(t2); // assuming T < T2 or T > T2, merge t2 into t
```

### 7.9 RMQ

```
* Author: Johan Sannemo, pajenegod
 * Date: 2015-02-06
 * License: CC0
 * Source: Folklore
 * Description: Range Minimum Queries on an array. Returns
 * min(V[a], V[a + 1], ... V[b - 1]) in constant time.
 * RMQ rmq(values);
 * rmq.query(inclusive, exclusive);
 * Time: $O(|V| \log |V| + Q)$
 * Status: stress-tested
 */
#pragma once
template<class T>
struct RMQ {
        vector<vector<T>> jmp;
        RMQ(const vector<T>& V) : jmp(1, V) {
                for (int pw = 1, k = 1; pw * 2 <= sz(V); pw *= 2, ++k) {
                        jmp.emplace_back(sz(V) - pw * 2 + 1);
                        rep(j,0,sz(jmp[k]))
                                 jmp[k][j] = min(jmp[k - 1][j], jmp[k - 1][j + pw
                                     ]);
        T query(int a, int b) {
                assert(a < b); // or return inf if a == b
                int dep = 31 - \underline{\quad}builtin_clz(b - a);
                return min(jmp[dep][a], jmp[dep][b - (1 << dep)]);
```

};

### 7.10 SegmentTree

```
* Author: Lucian Bicsi
 * Date: 2017-10-31
 * License: CC0
 * Source: folklore
 * Description: Zero-indexed max-tree. Bounds are inclusive to the left and
     exclusive to the right.
 \star Can be changed by modifying T, f and unit.
 * Time: O(\log N)
 * Status: stress-tested
#pragma once
struct Tree {
        typedef int T;
        static constexpr T unit = INT_MIN;
        T f(T a, T b) { return max(a, b); } // (any associative fn)
        vector<T> s; int n;
        Tree(int n = 0, T def = unit) : s(2*n, def), n(n) {}
        void update(int pos, T val) {
                for (s[pos += n] = val; pos /= 2;)
                        s[pos] = f(s[pos * 2], s[pos * 2 + 1]);
        T query(int b, int e) { // query [b, e)
                T ra = unit, rb = unit;
                for (b += n, e += n; b < e; b /= 2, e /= 2) {
                        if (b % 2) ra = f(ra, s[b++]);
                        if (e \% 2) rb = f(s[--e], rb);
                return f(ra, rb);
};
```

#### 7.11 SubMatrix

```
* Author: Johan Sannemo
 * Date: 2014-11-28
 * License: CC0
 * Source: Folklore
 * Description: Calculate submatrix sums quickly, given upper-left and lower-
     right corners (half-open).
 * SubMatrix<int> m(matrix);
 * m.sum(0, 0, 2, 2); // top left 4 elements
 * Time: O(N^2 + Q)
 * Status: Tested on Kattis
#pragma once
template < class T>
struct SubMatrix {
        vector<vector<T>> p;
        SubMatrix(vector<vector<T>>& v) {
                int R = sz(v), C = sz(v[0]);
                p.assign(R+1, vector<T>(C+1));
                rep(r, 0, R) rep(c, 0, C)
                        p[r+1][c+1] = v[r][c] + p[r][c+1] + p[r+1][c] - p[r][c];
        T sum(int u, int l, int d, int r) {
                return p[d][r] - p[d][l] - p[u][r] + p[u][l];
};
```

#### 7.12 Treap

/\*\*

```
* Author: someone on Codeforces
 * Date: 2017-03-14
 * Source: folklore
 * Description: A short self-balancing tree. It acts as a
 * sequential container with log-time splits/joins, and
 * is easy to augment with additional data.
 * Time: $\tilde{O}(\log \tilde{N})$
 * Status: stress-tested
#pragma once
struct Node {
        Node *1 = 0, *r = 0;
        int val, y, c = 1;
        Node(int val) : val(val), y(rand()) {}
        void recalc();
};
int cnt(Node* n) { return n ? n->c : 0; }
void Node::recalc() { c = cnt(1) + cnt(r) + 1; }
template<class F> void each(Node* n, F f) {
        if (n) { each (n->1, f); f(n->val); each (n->r, f); }
pair<Node*, Node*> split(Node* n, int k) {
        if (!n) return {};
        if (cnt(n->1) >= k) { // "n->val >= k" for lower_bound(k)}
                auto [L,R] = split(n->1, k);
                n->1 = R;
                n->recalc();
                return {L, n};
        } else {
                auto [L,R] = split(n->r,k-cnt(n->l)-1); // and just "k"
                n->r = L:
                n->recalc();
                return {n, R};
Node* merge(Node* 1, Node* r) {
        if (!1) return r;
        if (!r) return 1;
        if (1->y > r->y) {
                1->r = merge(1->r, r);
                return 1->recalc(), 1;
        } else {
                r->1 = merge(1, r->1);
                return r->recalc(), r;
Node* ins(Node* t, Node* n, int pos) {
        auto [l,r] = split(t, pos);
        return merge (merge (1, n), r);
// Example application: move the range [1, r) to index k
void move(Node*& t, int 1, int r, int k) {
        Node *a, *b, *c;
        tie(a,b) = split(t, 1); tie(b,c) = split(b, r - 1);
        if (k \le 1) t = merge(ins(a, b, k), c);
        else t = merge(a, ins(c, b, k - r));
```

### 7.13 UnionFind

```
/**
    * Author: Lukas Polacek
    * Date: 2009-10-26
    * License: CC0
    * Source: folklore
```

#### 7.14 UnionFindRollback

```
* Author: Lukas Polacek, Simon Lindholm
 * Date: 2019-12-26
 * License: CC0
 * Source: folklore
 * Description: Disjoint-set data structure with undo.
 * If undo is not needed, skip st, time() and rollback().
 * Usage: int t = uf.time(); ...; uf.rollback(t);
 * Time: $0(\log(N))$
 * Status: tested as part of DirectedMST.h
#pragma once
struct RollbackUF {
        vi e; vector<pii> st;
        RollbackUF(int n) : e(n, -1) {}
        int size(int x) { return -e[find(x)]; }
        int find(int x) { return e[x] < 0 ? x : find(e[x]); }
        int time() { return sz(st); }
        void rollback(int t) {
                for (int i = time(); i \longrightarrow t;)
                        e[st[i].first] = st[i].second;
                st.resize(t);
        bool join(int a, int b) {
                a = find(a), b = find(b);
                if (a == b) return false;
                if (e[a] > e[b]) swap(a, b);
                st.push back({a, e[a]});
                st.push_back({b, e[b]});
                e[a] += e[b]; e[b] = a;
                return true;
};
```

# 8 kactl forked/geometry

#### 8.1 3dHull

```
/**
 * Author: Johan Sannemo
 * Date: 2017-04-18
 * Source: derived from https://gist.github.com/msg555/4963794 by Mark Gordon
 * Description: Computes all faces of the 3-dimension hull of a point set.
 * *No four points must be coplanar*, or else random results will be returned.
 * All faces will point outwards.
 * Time: O(n^2)
 * Status: tested on SPOJ CH3D
```

```
#pragma once
#include "Point3D.h"
typedef Point3D<double> P3;
struct PR {
        void ins(int x) { (a == -1 ? a : b) = x; }
        void rem(int x) { (a == x ? a : b) = -1; }
        int cnt() { return (a !=-1) + (b !=-1); }
};
struct F { P3 q; int a, b, c; };
vector<F> hull3d(const vector<P3>& A) {
        assert(sz(A) >= 4);
        vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
\#define E(x,y) E[f.x][f.y]
        vector<F> FS;
        auto mf = [&](int i, int j, int k, int l) {
     P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
                 if (q.dot(A[1]) > q.dot(A[i]))
                q = q * -1;
F f{q, i, j, k};
                 E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);
                 FS.push_back(f);
        rep(i,0,4) rep(j,i+1,4) rep(k,j+1,4)
                 mf(i, j, k, 6 - i - j - k);
        rep(i,4,sz(A)) {
                 rep(j,0,sz(FS))
                         if(f.q.dot(A[i]) > f.q.dot(A[f.a])) {
                                  E(a,b).rem(f.c);
                                  E(a,c).rem(f.b);
                                  E(b,c).rem(f.a);
                                  swap(FS[j--], FS.back());
                                  FS.pop_back();
                 int nw = sz(FS);
                 rep(j,0,nw) {
#define C(a, b, c) if (E(a,b).cnt() != 2) mf(f.a, f.b, i, f.c);
                         C(a, b, c); C(a, c, b); C(b, c, a);
        for (F& it: FS) if ((A[it.b] - A[it.a]).cross(
                A[it.c] - A[it.a]).dot(it.q) <= 0) swap(it.c, it.b);
        return FS;
};
```

# 8.2 Angle

```
int x, y;
        int t:
        Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
        Angle operator-(Angle b) const { return {x-b.x, y-b.y, t}; }
        int half() const {
                assert(x || y);
                return y < 0 \mid | (y == 0 && x < 0);
        Angle t90() const { return \{-y, x, t + (half() \&\& x >= 0)\}; \}
        Angle t180() const { return \{-x, -y, t + half()\}; }
        Angle t360() const { return \{x, y, t + 1\}; }
bool operator<(Angle a, Angle b) {</pre>
        // add a.dist2() and b.dist2() to also compare distances
        return make_tuple(a.t, a.half(), a.y * (ll)b.x) <
               make_tuple(b.t, b.half(), a.x * (ll)b.y);
// Given two points, this calculates the smallest angle between
// them, i.e., the angle that covers the defined line segment.
pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
        if (b < a) swap(a, b);
        return (b < a.t180() ?
                make_pair(a, b) : make_pair(b, a.t360()));
Angle operator+(Angle a, Angle b) { // point a + vector b
        Angle r(a.x + b.x, a.y + b.y, a.t);
        if (a.t180() < r) r.t--;
        return r.t180() < a ? r.t360() : r;
Angle angleDiff(Angle a, Angle b) { // angle b - angle a
        int tu = b.t - a.t; a.t = b.t;
        return \{a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu - (b < a)\};
```

#### 8.3 CircleIntersection

```
* Author: Simon Lindholm
 * Date: 2015-09-01
 * License: CC0
 * Description: Computes the pair of points at which two circles intersect.
 * Returns false in case of no intersection.
 * Status: stress-tested
#pragma once
#include "Point.h"
typedef Point < double > P;
bool circleInter(P a,P b,double r1,double r2,pair<P, P>* out) {
        if (a == b) { assert(r1 != r2); return false; }
        P \text{ vec} = b - a;
        double d2 = vec.dist2(), sum = r1+r2, dif = r1-r2,
               p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1 - p*p*d2;
        if (sum*sum < d2 || dif*dif > d2) return false;
        P \text{ mid} = a + \text{vec*p, per} = \text{vec.perp()} * \text{sqrt(fmax(0, h2) / d2);}
        *out = {mid + per, mid - per};
        return true:
```

#### 8.4 CircleLine

```
/**
    * Author: Victor Lecomte, chilli
    * Date: 2019-10-29
    * License: CC0
    * Source: https://vlecomte.github.io/cp-geo.pdf
    * Description: Finds the intersection between a circle and a line.
    * Returns a vector of either 0, 1, or 2 intersection points.
    * P is intended to be Point<double>.
    * Status: unit tested
```

```
#pragma once
#include "Point.h"

template<class P>
vector<P> circleLine(P c, double r, P a, P b) {
    P ab = b - a, p = a + ab * (c-a).dot(ab) / ab.dist2();
    double s = a.cross(b, c), h2 = r*r - s*s / ab.dist2();
    if (h2 < 0) return {};
    if (h2 == 0) return {p};
    P h = ab.unit() * sqrt(h2);
    return {p - h, p + h};
}</pre>
```

### 8.5 CirclePolygonIntersection

```
* Author: chilli, Takanori MAEHARA
* Date: 2019-10-31
 * License: CC0
 * Source: https://github.com/spaghetti-source/algorithm/blob/master/geometry/
     _geom.cc#L744
 * Description: Returns the area of the intersection of a circle with a
 * ccw polygon.
 * Time: O(n)
* Status: Tested on GNYR 2019 Gerrymandering, stress-tested
#pragma once
#include "../../content/geometry/Point.h"
typedef Point < double > P;
#define arg(p, g) atan2(p.cross(g), p.dot(g))
double circlePoly(P c, double r, vector<P> ps) {
        auto tri = [&] (P p, P q) {
                auto r2 = r * r / 2;
                P d = q - p;
                auto a = d.dot(p)/d.dist2(), b = (p.dist2()-r*r)/d.dist2();
                auto det = a * a - b;
                if (det \le 0) return arg(p, q) * r2;
                auto s = max(0., -a-sqrt(det)), t = min(1., -a+sqrt(det));
                if (t < 0 \mid | 1 \le s) return arg(p, q) * r2;
                Pu = p + d * s, v = q + d * (t-1);
                return arg(p,u) * r2 + u.cross(v)/2 + arg(v,q) * r2;
        auto sum = 0.0;
        rep(i,0,sz(ps))
                sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);
```

# 8.6 CircleTangents

```
/**
    * Author: Victor Lecomte, chilli
    * Date: 2019-10-31
    * License: CC0
    * Source: https://vlecomte.github.io/cp-geo.pdf
    * Description: Finds the external tangents of two circles, or internal if r2 is negated.
    * Can return 0, 1, or 2 tangents -- 0 if one circle contains the other (or overlaps it, in the internal case, or if the circles are the same);
    * 1 if the circles are tangent to each other (in which case .first = .second and the tangent line is perpendicular to the line between the centers).
    * .first and .second give the tangency points at circle 1 and 2 respectively.
    * To find the tangents of a circle with a point set r2 to 0.
    * Status: tested
    */
#pragma once
```

```
#include "Point.h"

template<class P>
vector<pair<P, P>> tangents(P c1, double r1, P c2, double r2) {
    P d = c2 - c1;
    double dr = r1 - r2, d2 = d.dist2(), h2 = d2 - dr * dr;
    if (d2 == 0 || h2 < 0) return {};
    vector<pair<P, P>> out;
    for (double sign : {-1, 1}) {
        P v = (d * dr + d.perp() * sqrt(h2) * sign) / d2;
        out.push_back({c1 + v * r1, c2 + v * r2});
    }
    if (h2 == 0) out.pop_back();
    return out;
}
```

#### 8.7 circumcircle

```
* Author: Ulf Lundstrom
 * Date: 2009-04-11
 * License: CC0
* Source: http://en.wikipedia.org/wiki/Circumcircle
 * Description: \\
\begin{minipage}{75mm}
The circumcirle of a triangle is the circle intersecting all three vertices.
    ccRadius returns the radius of the circle going through points A, B and C
    and ccCenter returns the center of the same circle.
\end{minipage}
\begin{minipage} {15mm}
\vspace{-2mm}
\includegraphics[width=\textwidth] {content/geometry/circumcircle}
\end{minipage}
* Status: tested
#pragma once
#include "Point.h"
typedef Point < double > P;
double ccRadius (const P& A, const P& B, const P& C) {
        return (B-A).dist()*(C-B).dist()*(A-C).dist()/
                        abs((B-A).cross(C-A))/2;
P ccCenter(const P& A, const P& B, const P& C) {
        P b = C-A, c = B-A;
        return A + (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
```

#### 8.8 ClosestPair

```
* Author: Simon Lindholm
 * Date: 2019-04-17
 * License: CC0
 * Source: https://codeforces.com/blog/entry/58747
 * Description: Finds the closest pair of points.
 * Time: O(n \log n)
 * Status: stress-tested
#pragma once
#include "Point.h"
typedef Point<ll> P;
pair<P, P> closest(vector<P> v) {
        assert(sz(v) > 1);
        set<P> S:
        sort(all(v), [](P a, P b) { return a.y < b.y; });
        pair<ll, pair<P, P>> ret{LLONG_MAX, {P(), P()}};
        int j = 0;
        for (P p : v) {
```

#### 8.9 ConvexHull

```
* Author: Stjepan Glavina, chilli
 * Date: 2019-05-05
 * License: Unlicense
 * Source: https://github.com/stjepang/snippets/blob/master/convex_hull.cpp
 * Description:
\\begin{minipage}{75mm}
Returns a vector of the points of the convex hull in counter-clockwise order.
Points on the edge of the hull between two other points are not considered part
    of the hull.
\end{minipage}
\begin{minipage} {15mm}
\vspace{-6mm}
\includegraphics[width=\textwidth] {content/geometry/ConvexHull}
\vspace{-6mm}
\end{minipage}
* Time: O(n \log n)
 * Status: stress-tested, tested with kattis:convexhull
#pragma once
#include "Point.h"
typedef Point<11> P;
vector<P> convexHull(vector<P> pts) {
        if (sz(pts) \le 1) return pts;
        sort(all(pts));
        vector<P> h(sz(pts)+1);
        int s = 0, t = 0;
        for (int it = 2; it--; s = --t, reverse(all(pts)))
                for (P p : pts) {
                        while (t \ge s + 2 \&\& h[t-2].cross(h[t-1], p) \le 0) t--;
                        h[t++] = p;
        return \{h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])\};
```

### 8.10 DelaunayTriangulation

```
* Author: Mattias de Zalenski
 * Date: Unknown
 * Source: Geometry in C
 * Description: Computes the Delaunay triangulation of a set of points.
 * Each circumcircle contains none of the input points.
 * If any three points are collinear or any four are on the same circle,
     behavior is undefined.
 * Time: O(n^2)
 * Status: stress-tested
#pragma once
#include "Point.h"
#include "3dHull.h"
template<class P, class F>
void delaunay(vector<P>& ps, F trifun) {
        if (sz(ps) == 3) { int d = (ps[0].cross(ps[1], ps[2]) < 0);
                trifun(0,1+d,2-d); }
        vector<P3> p3;
```

### 8.11 FastDelaunay

```
* Author: Philippe Legault
 * Date: 2016
 * License: MIT
 * Source: https://github.com/Bathlamos/delaunay-triangulation/
 * Description: Fast Delaunay triangulation.
 * Each circumcircle contains none of the input points.
 * There must be no duplicate points.
 * If all points are on a line, no triangles will be returned.
 * Should work for doubles as well, though there may be precision issues in '
 * Returns triangles in order \{t[0][0], t[0][1], t[0][2], t[1][0], \dots\}, all
       counter-clockwise.
 * Time: O(n \log n)
 * Status: stress-tested
#pragma once
#include "Point.h"
typedef Point<11> P;
typedef struct Quad* Q;
typedef __int128_t lll; // (can be ll if coords are < 2e4)
P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
struct Quad {
        Q rot, o; P p = arb; bool mark;
        P& F() { return r()->p; }
        O& r() { return rot->rot; }
        Q prev() { return rot->o->rot;
        Q next() { return r()->prev(); }
} *H;
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
        111 p2 = p.dist2(), A = a.dist2()-p2,
            B = b.dist2()-p2, C = c.dist2()-p2;
        return p.cross(a,b) \starC + p.cross(b,c) \starA + p.cross(c,a) \starB > 0;
Q makeEdge(P orig, P dest) {
        Q r = H ? H : new Quad{new Quad{new Quad{0}}}};
        H = r -> 0; r -> r() -> r() = r;
        rep(i,0,4) r = r->rot, r->p = arb, r->o = i & 1 ? r : r->r();
        r\rightarrow p = orig; r\rightarrow F() = dest;
        return r;
void splice(Q a, Q b) {
        swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
Q connect(Q a, Q b) {
        Q = makeEdge(a->F(), b->p);
        splice(q, a->next());
        splice(q->r(), b);
        return q;
pair<Q,Q> rec(const vector<P>& s) {
        if (sz(s) \le 3) {
                Q = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back());
                if (sz(s) == 2) return { a, a->r() };
                splice(a->r(), b);
                auto side = s[0].cross(s[1], s[2]);
                Q c = side ? connect(b, a) : 0;
                return {side < 0 ? c->r() : a, side < 0 ? c : b->r() };
\#define H(e) e->F(), e->p
#define valid(e) (e->F().cross(H(base)) > 0)
```

```
Q A, B, ra, rb;
        int half = sz(s) / 2;
        tie(ra, A) = rec({all(s) - half});
        tie(B, rb) = rec({sz(s) - half + all(s)});
        while ((B\rightarrow p.cross(H(A)) < 0 \&\& (A = A\rightarrow next())) | |
                (A->p.cross(H(B)) > 0 && (B = B->r()->o)));
        Q base = connect(B \rightarrow r(), A);
        if (A->p == ra->p) ra = base->r();
        if (B->p == rb->p) rb = base;
#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
                while (circ(e->dir->F(), H(base), e->F())) { \setminus
                         Q t = e->dir;
                         splice(e, e->prev()); \
                         splice(e->r(), e->r()->prev()); \
                         e->o = H; H = e; e = t; \
        for (;;) {
                DEL(LC, base->r(), o); DEL(RC, base, prev());
                if (!valid(LC) && !valid(RC)) break;
                if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
                         base = connect(RC, base->r());
                         base = connect(base->r(), LC->r());
        return { ra, rb };
vector<P> triangulate(vector<P> pts) {
        sort(all(pts)); assert(unique(all(pts)) == pts.end());
        if (sz(pts) < 2) return \{\};
        Q e = rec(pts).first;
        vector < Q > q = \{e\};
        int qi = 0;
        while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
\#define ADD { Q c = e; do { c->mark = 1; pts.push_back(c->p); \
        q.push_back(c->r()); c = c->next(); } while (c != e); }
        ADD; pts.clear();
        while (qi < sz(q)) if (!(e = q[qi++]) \rightarrow mark) ADD;
        return pts;
```

### 8.12 HullDiameter

```
* Author: Oleksandr Bacherikov, chilli
 * Date: 2019-05-05
 * License: Boost Software License
 * Source: https://codeforces.com/blog/entry/48868
* Description: Returns the two points with max distance on a convex hull (ccw,
 * no duplicate/collinear points).
 * Status: stress-tested, tested on kattis:roberthood
 * Time: O(n)
#pragma once
#include "Point.h"
typedef Point<ll> P;
array<P, 2> hullDiameter(vector<P> S) {
        int n = sz(S), j = n < 2 ? 0 : 1;
        pair<ll, array<P, 2>> res({0, {S[0], S[0]}});
        rep(i,0,j)
                 for (;; j = (j + 1) % n) {
                        res = \max(\text{res}, \{(S[i] - S[j]).dist2(), \{S[i], S[j]\}\});
                        if ((S[(j + 1) % n] - S[j]).cross(S[i + 1] - S[i]) >= 0)
        return res.second:
```

# 8.13 InsidePolygon

```
* Author: Victor Lecomte, chilli
 * Date: 2019-04-26
 * License: CC0
 * Source: https://vlecomte.github.io/cp-geo.pdf
 * Description: Returns true if p lies within the polygon. If strict is true,
 \star it returns false for points on the boundary. The algorithm uses
 * products in intermediate steps so watch out for overflow.
 * Time: O(n)
 * vector<P> v = \{P\{4,4\}, P\{1,2\}, P\{2,1\}\};
 * bool in = inPolygon(v, P{3, 3}, false);
 * Status: stress-tested and tested on kattis:pointinpolygon
#pragma once
#include "Point.h"
#include "OnSegment.h"
#include "SegmentDistance.h"
template<class P>
bool inPolygon(vector<P> &p, P a, bool strict = true) {
        int cnt = 0, n = sz(p);
        rep(i,0,n) {
                P q = p[(i + 1) % n];
                if (onSegment(p[i], q, a)) return !strict;
                //or: if (segDist(p[i], q, a) <= eps) return !strict;</pre>
                cnt \hat{} = ((a.y<p[i].y) - (a.y<q.y)) * a.cross(p[i], q) > 0;
        return cnt;
```

#### 8.14 kdTree

```
* Author: Stanford
 * Date: Unknown
 * Source: Stanford Notebook
 * Description: KD-tree (2d, can be extended to 3d)
 * Status: Tested on excellentengineers
#pragma once
#include "Point.h"
typedef long long T;
typedef Point<T> P;
const T INF = numeric_limits<T>::max();
bool on_x(const P& a, const P& b) { return a.x < b.x; }
bool on_y(const P& a, const P& b) { return a.y < b.y; }
struct Node {
        P pt; // if this is a leaf, the single point in it
        T x0 = INF, x1 = -INF, y0 = INF, y1 = -INF; // bounds
        Node *first = 0, *second = 0;
        T distance(const P& p) { // min squared distance to a point
                T x = (p.x < x0 ? x0 : p.x > x1 ? x1 : p.x);
                T y = (p.y < y0 ? y0 : p.y > y1 ? y1 : p.y);
                return (P(x,y) - p).dist2();
        Node(vector<P>&& vp) : pt(vp[0]) {
                for (P p : vp) {
                        x0 = min(x0, p.x); x1 = max(x1, p.x);
                        y0 = min(y0, p.y); y1 = max(y1, p.y);
                if (vp.size() > 1) {
                        // split on x if width >= height (not ideal...)
                        sort(all(vp), x1 - x0 >= y1 - y0 ? on_x : on_y);
                        // divide by taking half the array for each child (not
                        // best performance with many duplicates in the middle)
                        int half = sz(vp)/2;
                        first = new Node({vp.begin(), vp.begin() + half});
                        second = new Node({vp.begin() + half, vp.end()});
```

```
};
struct KDTree {
       Node* root:
        KDTree(const vector<P>& vp) : root(new Node({all(vp)})) {}
        pair<T, P> search(Node *node, const P& p) {
                if (!node->first) {
                        // uncomment if we should not find the point itself:
                        // if (p == node->pt) return {INF, P()};
                        return make_pair((p - node->pt).dist2(), node->pt);
                Node *f = node -> first, *s = node -> second;
                T bfirst = f->distance(p), bsec = s->distance(p);
                if (bfirst > bsec) swap(bsec, bfirst), swap(f, s);
                // search closest side first, other side if needed
                auto best = search(f, p);
                if (bsec < best.first)
                        best = min(best, search(s, p));
                return best;
        // find nearest point to a point, and its squared distance
        // (requires an arbitrary operator< for Point)
        pair<T, P> nearest(const P& p) {
                return search(root, p);
};
```

#### 8.15 linearTransformation

```
* Author: Per Austrin, Ulf Lundstrom
 * Date: 2009-04-09
 * License: CC0
 * Source:
 * Description: \\
\begin{minipage} {75mm}
Apply the linear transformation (translation, rotation and scaling) which takes
       line p0-p1 to line q0-q1 to point r.
\end{minipage}
\begin{minipage} {15mm}
\vspace{-8mm}
\includegraphics[width=\textwidth] {content/geometry/linearTransformation}
\vspace{-2mm}
\end{minipage}
* Status: not tested
#pragma once
#include "Point.h"
typedef Point < double > P;
P linearTransformation (const P& p0, const P& p1,
                const P& q0, const P& q1, const P& r) {
        P dp = p1-p0, dq = q1-q0, num(dp.cross(dq), dp.dot(dq));
        return q0 + P((r-p0).cross(num), (r-p0).dot(num))/dp.dist2();
```

#### 8.16 lineDistance

```
/**

* Author: Ulf Lundstrom

* Date: 2009-03-21

* License: CC0

* Source: Basic math

* Description:\\
begin{minipage} {75mm}
```

```
Returns the signed distance between point p and the line containing points a and
Positive value on left side and negative on right as seen from a towards b. a == b
     gives nan.
P is supposed to be Point<T> or Point3D<T> where T is e.g. double or long long.
It uses products in intermediate steps so watch out for overflow if using int or
      long long.
Using Point3D will always give a non-negative distance. For Point3D, call .dist
    on the result of the cross product.
\end{minipage}
\begin{minipage}{15mm}
\includegraphics[width=\textwidth] {content/geometry/lineDistance}
\end{minipage}
 * Status: tested
 */
#pragma once
#include "Point.h"
template<class P>
double lineDist(const P& a, const P& b, const P& p) {
        return (double) (b-a).cross(p-a)/(b-a).dist();
```

#### 8.17 LineHullIntersection

```
* Author: Oleksandr Bacherikov, chilli
 * Date: 2019-05-07
 * License: Boost Software License
 * Source: https://github.com/AlCash07/ACTL/blob/master/include/actl/geometry/
     algorithm/intersect/line_convex_polygon.hpp
 * Description: Line-convex polygon intersection. The polygon must be ccw and
     have no collinear points.
 * lineHull(line, poly) returns a pair describing the intersection of a line
     with the polygon:
    \begin{itemize*}
      \item (-1, -1) if no collision,
      \item (i, -1) if touching the corner i,
      \item (i, i) if along side (i, i+1),
     \item (i, j) if crossing sides (i, i+1) and (j, j+1).
    \end{itemize*}
   In the last case, if a corner $i$ is crossed, this is treated as happening
     on side (i, i+1).
   The points are returned in the same order as the line hits the polygon.
 * \texttt{extrVertex} returns the point of a hull with the max projection onto
 * Time: O(\log n)
 * Status: stress-tested
#pragma once
#include "Point.h"
#define cmp(i,j) sgn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
#define extr(i) cmp(i + 1, i) >= 0 && cmp(i, i - 1 + n) < 0
template <class P> int extrVertex(vector<P>& poly, P dir) {
        int n = sz(poly), lo = 0, hi = n;
        if (extr(0)) return 0;
        while (lo + 1 < hi) \{
                int m = (lo + hi) / 2;
                if (extr(m)) return m;
                int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
                (1s < ms \mid | (1s == ms \&\& 1s == cmp(1o, m)) ? hi : 1o) = m;
        return lo:
#define cmpL(i) sgn(a.cross(poly[i], b))
template <class P>
array<int, 2> lineHull(P a, P b, vector<P>& poly) {
        int endA = extrVertex(poly, (a - b).perp());
        int endB = extrVertex(poly, (b - a).perp());
        if (cmpL(endA) < 0 \mid \mid cmpL(endB) > 0)
                return {-1, -1};
```

#### 8.18 lineIntersection

```
* Author: Victor Lecomte, chilli
 * Date: 2019-05-05
 * License: CC0
 * Source: https://vlecomte.github.io/cp-geo.pdf
 * Description: \\
\begin{minipage}{75mm}
If a unique intersection point of the lines going through s1,e1 and s2,e2 exists
      \{1, point\} is returned.
If no intersection point exists \{0, (0,0)\} is returned and if infinitely many
    exists \{-1, (0,0)\} is returned.
The wrong position will be returned if P is Point<ll> and the intersection point
      does not have integer coordinates.
Products of three coordinates are used in intermediate steps so watch out for
    overflow if using int or 11.
\end{minipage}
\begin{minipage}{15mm}
\includegraphics[width=\textwidth] {content/geometry/lineIntersection}
\end{minipage}
        auto res = lineInter(s1,e1,s2,e2);
        if (res.first == 1)
                cout << "intersection point at " << res.second << endl;</pre>
 * Status: stress-tested, and tested through half-plane tests
#pragma once
#include "Point.h"
template<class P>
pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
        auto d = (e1 - s1).cross(e2 - s2);
        if (d == 0) // if parallel
                return \{-(s1.cross(e1, s2) == 0), P(0, 0)\};
        auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
        return \{1, (s1 * p + e1 * q) / d\};
```

### 8.19 LineProjectionReflection

```
/**
    * Author: Victor Lecomte, chilli
    * Date: 2019-10-29
    * License: CC0
    * Source: https://vlecomte.github.io/cp-geo.pdf
    * Description: Projects point p onto line ab. Set refl=true to get reflection
    * of point p across line ab instead. The wrong point will be returned if P is
    * an integer point and the desired point doesn't have integer coordinates.
    * Products of three coordinates are used in intermediate steps so watch out
    * for overflow.
    * Status: stress-tested
```

```
*/
#pragma once
#include "Point.h"

template<class P>
P lineProj(P a, P b, P p, bool refl=false) {
        P v = b - a;
        return p - v.perp()*(1+refl)*v.cross(p-a)/v.dist2();
}
```

#### 8.20 ManhattanMST

```
* Author: chilli, Takanori MAEHARA
 * Date: 2019-11-02
 * License: CC0
 * Source: https://github.com/spaghetti-source/algorithm/blob/master/geometry/
     rectilinear_mst.cc
 * Description: Given N points, returns up to 4*N edges, which are guaranteed
 * to contain a minimum spanning tree for the graph with edge weights w(p, q) =
 * |p.x - q.x| + |p.y - q.y|. Edges are in the form (distance, src, dst). Use a
 * standard MST algorithm on the result to find the final MST.
 * Time: O(N \log N)
 * Status: Stress-tested
#pragma once
#include "Point.h"
typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
        vi id(sz(ps));
        iota(all(id), 0);
        vector<array<int, 3>> edges;
        rep(k,0,4) {
                sort(all(id), [&](int i, int j) {
                     return (ps[i]-ps[j]).x < (ps[j]-ps[i]).y;});
                map<int, int> sweep;
                for (int i : id) {
                        for (auto it = sweep.lower_bound(-ps[i].y);
                                        it != sweep.end(); sweep.erase(it++)) {
                                int j = it->second;
                                P d = ps[i] - ps[j];
                                if (d.y > d.x) break;
                                edges.push_back({d.y + d.x, i, j});
                        sweep[-ps[i].y] = i;
                for (P& p : ps) if (k & 1) p.x = -p.x; else swap(p.x, p.y);
        return edges;
```

# 8.21 MinimumEnclosingCircle

### 8.22 OnSegment

```
/**
    * Author: Victor Lecomte, chilli
    * Date: 2019-04-26
    * License: CC0
    * Source: https://vlecomte.github.io/cp-geo.pdf
    * Description: Returns true iff p lies on the line segment from s to e.
    * Use \text1t{(segDist(s,e,p)<=epsilon)} instead when using Point<double>.
    * Status:
    */
#pragma once
#include "Point.h"

template<class P> bool onSegment(P s, P e, P p) {
        return p.cross(s, e) == 0 && (s - p).dot(e - p) <= 0;
}</pre>
```

#### 8.23 Point

```
* Author: Ulf Lundstrom
 * Date: 2009-02-26
* License: CC0
* Source: My head with inspiration from tinyKACTL
 * Description: Class to handle points in the plane.
        T can be e.g. double or long long. (Avoid int.)
* Status: Works fine, used a lot
#pragma once
template <class T> int sqn(T x) { return (x > 0) - (x < 0); }
template<class T>
struct Point {
        typedef Point P;
        explicit Point (T x=0, T y=0) : x(x), y(y) {}
        bool operator < (P p) const { return tie(x,y) < tie(p.x,p.y); }
        bool operator == (P p) const { return tie(x,y) == tie(p.x,p.y); }
        P operator+(P p) const { return P(x+p.x, y+p.y); }
        P operator-(P p) const { return P(x-p.x, y-p.y); }
        P operator* (T d) const { return P(x*d, y*d);
        P operator/(T d) const { return P(x/d, y/d); }
        T dot(P p) const { return x*p.x + y*p.y; }
        T cross(P p) const { return x*p.y - y*p.x; }
        T cross(P a, P b) const { return (a-*this).cross(b-*this); }
        T dist2() const { return x*x + y*y; }
        double dist() const { return sqrt((double)dist2()); }
        // angle to x-axis in interval [-pi, pi]
        double angle() const { return atan2(y, x); }
        P unit() const { return *this/dist(); } // makes dist()=1
P perp() const { return P(-y, x); } // rotates +90 degrees
        P normal() const { return perp().unit(); }
        // returns point rotated 'a' radians ccw around the origin
        P rotate(double a) const {
                 return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a));}
        friend ostream& operator<<(ostream& os, P p) {
```

```
return os << "(" << p.x << "," << p.y << ")"; }
```

#### 8.24 Point3D

};

```
* Author: Ulf Lundstrom with inspiration from tinyKACTL
 * Date: 2009-04-14
 * License: CC0
 * Source:
 \star Description: Class to handle points in 3D space.
      T can be e.g. double or long long.
 * Status: tested, except for phi and theta
#pragma once
template<class T> struct Point3D {
        typedef Point3D P;
        typedef const P& R;
        T x, y, z;
        explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y), z(z) {}
        bool operator<(R p) const
                return tie(x, y, z) < tie(p.x, p.y, p.z); }
        bool operator == (R p) const {
                return tie(x, y, z) == tie(p.x, p.y, p.z); }
        P operator+(R p) const { return P(x+p.x, y+p.y, z+p.z);
        P operator-(R p) const { return P(x-p.x, y-p.y, z-p.z); }
        P operator*(T d) const { return P(x*d, y*d, z*d);
        P operator/(T d) const { return P(x/d, y/d, z/d); }
        T dot(R p) const { return x*p.x + y*p.y + z*p.z; }
        P cross(R p) const {
                return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x);
        T dist2() const { return x*x + y*y + z*z; }
        double dist() const { return sqrt((double)dist2()); }
        //Azimuthal angle (longitude) to x-axis in interval [-pi, pi]
        double phi() const { return atan2(y, x); }
        //Zenith angle (latitude) to the z-axis in interval [0, pi]
        double theta() const { return atan2(sqrt(x*x+y*y),z); }
        P unit() const { return *this/(T)dist(); } //makes dist()=1
        //returns unit vector normal to *this and p
        P normal(P p) const { return cross(p).unit(); }
        //returns point rotated 'angle' radians ccw around axis
        P rotate (double angle, P axis) const {
                double s = sin(angle), c = cos(angle); P u = axis.unit();
                return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
};
```

### 8.25 PointInsideHull

```
/**
  * Author: chilli
  * Date: 2019-05-17
  * License: CC0
  * Source: https://github.com/ngthanhtrung23/ACM_Notebook_new
  * Description: Determine whether a point t lies inside a convex hull (CCW
  * order, with no collinear points). Returns true if point lies within
  * the hull. If strict is true, points on the boundary aren't included.
  * Usage:
  * Status: stress-tested
  * Time: O(\log N)
  */
  #pragma once

#include "Point.h"
#include "sideOf.h"
#include "onSegment.h"

typedef Point<11> P;
```

### 8.26 PolygonArea

```
/**
 * Author: Ulf Lundstrom
 * Date: 2009-03-21
 * License: CC0
 * Source: tinyKACTL
 * Description: Returns twice the signed area of a polygon.
 * Clockwise enumeration gives negative area. Watch out for overflow if using int as T!
 * Status: Stress-tested and tested on kattis:polygonarea
 */
#pragma once
#include "Point.h"

template<class T>
T polygonArea2(vector<Point<T>>& v) {
    T a = v.back().cross(v[0]);
        rep(i,0,sz(v)-1) a += v[i].cross(v[i+1]);
        return a;
}
```

# 8.27 PolygonCenter

### 8.28 PolygonCut

```
/**
    * Author: Ulf Lundstrom
    * Date: 2009-03-21
    * License: CC0
    * Source:
```

```
* Description:\\
\begin{minipage}{75mm}
 Returns a vector with the vertices of a polygon with everything to the left of
      the line going from s to e cut away.
\end{minipage}
\begin{minipage} {15mm}
\vspace{-6mm}
\includegraphics[width=\textwidth] {content/geometry/PolygonCut}
\vspace{-6mm}
\end{minipage}
 * Usage:
        vector < P > p = ...;
        p = polygonCut(p, P(0,0), P(1,0));
 * Status: tested but not extensively
 */
#pragma once
#include "Point.h"
typedef Point < double > P;
vector<P> polygonCut(const vector<P>& poly, P s, P e) {
        vector<P> res;
        rep(i,0,sz(poly))
                P cur = poly[i], prev = i ? poly[i-1] : poly.back();
                auto a = s.cross(e, cur), b = s.cross(e, prev);
                if ((a < 0) != (b < 0))
                         res.push_back(cur + (prev - cur) * (a / (a - b)));
                if (a < 0)
                         res.push_back(cur);
        return res;
```

### 8.29 PolygonUnion

```
* Author: black_horse2014, chilli
 * Date: 2019-10-29
 * License: Unknown
 * Source: https://codeforces.com/gym/101673/submission/50481926
 * Description: Calculates the area of the union of $n$ polygons (not
      necessarily
 * convex). The points within each polygon must be given in CCW order.
 * (Epsilon checks may optionally be added to sideOf/sgn, but shouldn't be
     needed.)
 * Time: $O(N^2)$, where $N$ is the total number of points
 * Status: stress-tested, Submitted on ECNA 2017 Problem A
#pragma once
#include "Point.h"
#include "sideOf.h"
typedef Point < double > P;
double rat(P a, P b) { return sgn(b.x) ? a.x/b.x : a.y/b.y; }
double polyUnion(vector<vector<P>>& poly) {
        double ret = 0;
        rep(i,0,sz(poly)) rep(v,0,sz(poly[i])) {
                PA = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
                vector<pair<double, int>> segs = {{0, 0}, {1, 0}};
                rep(j,0,sz(poly)) if (i != j) {
                        rep(u, 0, sz(poly[j])) {
                                P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[
                                     j])];
                                int sc = sideOf(A, B, C), sd = sideOf(A, B, D);
                                if (sc != sd) {
                                        double sa = C.cross(D, A), sb = C.cross(
                                             D, B);
                                        if (\min(sc, sd) < 0)
                                                segs.emplace_back(sa / (sa - sb)
                                                      , sgn(sc - sd));
                                } else if (!sc && !sd && j<i && sgn((B-A).dot(D-
                                     C))>0){
                                        segs.emplace_back(rat(C - A, B - A), 1);
                                        segs.emplace_back(rat(D - A, B - A), -1)
```

```
}

}
sort(all(segs));
for (auto& s : segs) s.first = min(max(s.first, 0.0), 1.0);
double sum = 0;
int cnt = segs[0].second;
rep(j,1,sz(segs)) {
    if (!cnt) sum += segs[j].first - segs[j - 1].first;
        cnt += segs[j].second;
}
ret += A.cross(B) * sum;
}
return ret / 2;
}
```

### 8.30 PolyhedronVolume

# 8.31 SegmentDistance

```
* Author: Ulf Lundstrom
 * Date: 2009-03-21
 * License: CC0
 * Source:
* Description: \\
\begin{minipage} {75mm}
Returns the shortest distance between point p and the line segment from point s
\end{minipage}
\begin{minipage}{15mm}
\vspace{-10mm}
\includegraphics[width=\textwidth] {content/geometry/SegmentDistance}
\end{minipage}
 * Usage:
        Point<double> a, b(2,2), p(1,1);
        bool onSegment = segDist(a,b,p) < 1e-10;
 * Status: tested
 */
#pragma once
#include "Point.h"
typedef Point < double > P;
double segDist(P& s, P& e, P& p) {
        if (s==e) return (p-s).dist();
        auto d = (e-s).dist2(), t = min(d, max(.0, (p-s).dot(e-s)));
        return ((p-s)*d-(e-s)*t).dist()/d;
```

# 8.32 SegmentIntersection

```
* Author: Victor Lecomte, chilli
 * Date: 2019-04-27
 * License: CC0
 * Source: https://vlecomte.github.io/cp-geo.pdf
 * Description:\\
\begin{minipage} {75mm}
If a unique intersection point between the line segments going from s1 to e1 and
      from s2 to e2 exists then it is returned.
If no intersection point exists an empty vector is returned. If infinitely many exist a vector with 2 elements is returned, containing the \frac{1}{2}
     endpoints of the common line segment.
The wrong position will be returned if P is Point<11> and the intersection point
      does not have integer coordinates.
Products of three coordinates are used in intermediate steps so watch out for
     overflow if using int or long long.
\begin{minipage}{15mm}
\includegraphics[width=\textwidth] {content/geometry/SegmentIntersection}
 * vector<P> inter = segInter(s1,e1,s2,e2);
 \star if (sz(inter)==1)
 * cout << "segments intersect at " << inter[0] << endl;
 * Status: stress-tested, tested on kattis:intersection
#pragma once
#include "Point.h"
#include "OnSegment.h"
template<class P> vector<P> segInter(P a, P b, P c, P d) {
        auto oa = c.cross(d, a), ob = c.cross(d, b),
             oc = a.cross(b, c), od = a.cross(b, d);
         // Checks if intersection is single non-endpoint point.
        if (sgn(oa) * sgn(ob) < 0 && sgn(oc) * sgn(od) < 0)
                return { (a * ob - b * oa) / (ob - oa) };
        set<P> s;
        if (onSegment(c, d, a)) s.insert(a);
        if (onSegment(c, d, b)) s.insert(b);
        if (onSegment(a, b, c)) s.insert(c);
        if (onSegment(a, b, d)) s.insert(d);
        return {all(s)};
```

#### 8.33 sideOf

```
* Author: Ulf Lundstrom
 * Date: 2009-03-21
 * License: CC0
 * Description: Returns where p is as seen from s towards e. 1/0/-1
     Leftrightarrow$ left/on line/right.
 * If the optional argument $eps$ is given 0 is returned if $p$ is within
     distance $eps$ from the line.
 * P is supposed to be Point<T> where T is e.g. double or long long.
 * It uses products in intermediate steps so watch out for overflow if using int
 * Usage:
       bool left = sideOf(p1, p2, q) ==1;
 * Status: tested
#pragma once
#include "Point.h"
template<class P>
int sideOf(P s, P e, P p) { return sgn(s.cross(e, p)); }
int sideOf(const P& s, const P& e, const P& p, double eps) {
        auto a = (e-s).cross(p-s);
        double l = (e-s).dist()*eps;
        return (a > 1) - (a < -1);
```

### 8.34 sphericalDistance

```
* Author: Ulf Lundstrom
 * Date: 2009-04-07
 * License: CC0
 * Source: My geometric reasoning
 * Description: Returns the shortest distance on the sphere with radius radius
     between the points
 * with azimuthal angles (longitude) f1 ($\phi_1$) and f2 ($\phi_2$) from x axis
       and zenith angles
 * (latitude) t1 (\hat{z}) and t2 (\hat{z}) from z axis (0 = north pole).
      All angles measured
 * in radians. The algorithm starts by converting the spherical coordinates to
     cartesian coordinates
 * so if that is what you have you can use only the two last rows. dx*radius is
     then the difference
 \star between the two points in the x direction and d*radius is the total distance
     between the points.
 * Status: tested on kattis:airlinehub
#pragma once
double sphericalDistance(double f1, double t1,
                double f2, double t2, double radius) {
        double dx = \sin(t2) \cdot \cos(f2) - \sin(t1) \cdot \cos(f1);
        double dy = sin(t2)*sin(f2) - sin(t1)*sin(f1);
        double dz = cos(t2) - cos(t1);
        double d = sqrt(dx*dx + dy*dy + dz*dz);
        return radius*2*asin(d/2);
```

# 9 kactl forked/graph

#### 9.1 2sat

```
* Author: Emil Lenngren, Simon Lindholm
 * Date: 2011-11-29
* License: CC0
 * Source: folklore
 * Description: Calculates a valid assignment to boolean variables a, b, c,...
     to a 2-SAT problem,
 * so that an expression of the type (a||b)\\(a||c)\\(a||c)\\(a||b)\\(a|c...
 * becomes true, or reports that it is unsatisfiable.
 * Negated variables are represented by bit-inversions (\texttt{\tilde{}x}).
 * TwoSat ts(number of boolean variables);
 * ts.either(0, \tilde3); // Var 0 is true or var 3 is false
 * ts.setValue(2); // Var 2 is true
 * ts.atMostOne({0, \tilde{2}}); // <= 1 of vars 0, \tilde{2} and 2 are true
 * ts.solve(); // Returns true iff it is solvable
 \star ts.values[0..N-1] holds the assigned values to the vars
 \star Time: O(N+E), where N is the number of boolean variables, and E is the number
      of clauses.
 * Status: stress-tested
#pragma once
struct TwoSat {
        vector<vi> ar:
        vi values; // 0 = false, 1 = true
        TwoSat(int n = 0) : N(n), gr(2*n) {}
        int addVar() { // (optional)
                gr.emplace_back();
                gr.emplace_back();
                return N++;
```

```
void either(int f, int j) {
        f = \max(2*f, -1-2*f);
        j = \max(2*j, -1-2*j);
        gr[f].push_back(j^1);
        gr[j].push_back(f^1);
void setValue(int x) { either(x, x); }
void atMostOne(const vi& li) { // (optional)
        if (sz(li) <= 1) return;
int cur = ~li[0];
        rep(i,2,sz(li)) {
                int next = addVar();
                either(cur, ~li[i]);
                either(cur, next);
                either(~li[i], next);
                cur = ~next;
        either(cur, ~li[1]);
vi val, comp, z; int time = 0;
int dfs(int i) {
        int low = val[i] = ++time, x; z.push_back(i);
        for(int e : gr[i]) if (!comp[e])
                low = min(low, val[e] ?: dfs(e));
        if (low == val[i]) do {
                x = z.back(); z.pop_back();
                comp[x] = low;
                if (values[x>>1] == -1)
                        values[x>>1] = x&1;
        } while (x != i);
        return val[i] = low;
bool solve() {
        values.assign(N, -1);
        val.assign(2*N, 0); comp = val;
        rep(i,0,2*N) if (!comp[i]) dfs(i);
        rep(i,0,N) if (comp[2*i] == comp[2*i+1]) return 0;
        return 1:
```

#### 9.2 BellmanFord

};

```
* Author: Simon Lindholm
 * Date: 2015-02-23
 * License: CC0
 * Source: http://en.wikipedia.org/wiki/Bellman-Ford_algorithm
 * Description: Calculates shortest paths from $s$ in a graph that might have
     negative edge weights.
 * Unreachable nodes get dist = inf; nodes reachable through negative-weight
     cycles get dist = -inf.
 * Assumes $V^2 \max |w_i| < \tilde{} 2^{63}$.
 * Time: O(VE)
 * Status: Tested on kattis:shortestpath3
#pragma once
const ll inf = LLONG_MAX;
struct Ed { int a, b, w, s() { return a < b ? a : -a; }};
struct Node { 11 \text{ dist} = \inf; \text{ int prev} = -1; };
void bellmanFord(vector<Node>& nodes, vector<Ed>& eds, int s) {
        nodes[s].dist = 0;
        sort(all(eds), [](Ed a, Ed b) { return a.s() < b.s(); });</pre>
        int lim = sz(nodes) / 2 + 2; // /3+100 with shuffled vertices
        rep(i,0,lim) for (Ed ed : eds) {
                Node cur = nodes[ed.a], &dest = nodes[ed.b];
                if (abs(cur.dist) == inf) continue;
                11 d = cur.dist + ed.w;
```

# 9.3 BiconnectedComponents

```
* Author: Simon Lindholm
 * Date: 2017-04-17
 * License: CC0
 * Source: folklore
* Description: Finds all biconnected components in an undirected graph, and
 * runs a callback for the edges in each. In a biconnected component there
 * are at least two distinct paths between any two nodes. Note that a node can
 * be in several components. An edge which is not in a component is a bridge,
 * i.e., not part of any cycle.
 * Usage:
 * int eid = 0; ed.resize(N);
 * for each edge (a,b)
     ed[a].emplace_back(b, eid);
      ed[b].emplace_back(a, eid++); }
 * bicomps([\&](const vi\& edgelist) {...});
 * Time: O(E + V)
 * Status: tested during MIPT ICPC Workshop 2017
#pragma once
vi num, st;
vector<vector<pii>> ed;
int Time;
template<class F>
int dfs(int at, int par, F& f) {
   int me = num[at] = ++Time, top = me;
        for (auto [y, e] : ed[at]) if (e != par) {
                if (num[y]) {
                        top = min(top, num[y]);
                        if (num[y] < me)</pre>
                                 st.push_back(e);
                } else {
                        int si = sz(st);
                        int up = dfs(y, e, f);
                        top = min(top, up);
                        if (up == me) {
                                 st.push_back(e);
                                 f(vi(st.begin() + si, st.end()));
                                 st.resize(si);
                        else if (up < me) st.push_back(e);
                        else { /* e is a bridge */ }
        return top;
template<class F>
void bicomps(F f) {
        num.assign(sz(ed), 0);
        rep(i,0,sz(ed)) if (!num[i]) dfs(i,-1,f);
```

### 9.4 BinaryLifting

```
/**
 * Author: Johan Sannemo
 * Date: 2015-02-06
```

```
* License: CC0
 * Source: Folklore
 * Description: Calculate power of two jumps in a tree,
 * to support fast upward jumps and LCAs.
 * Assumes the root node points to itself.
 * Time: construction $0(N \log N)$, queries $0(\log N)$
 * Status: Tested at Petrozavodsk, also stress-tested via LCA.cpp
 */
#pragma once
vector<vi> treeJump(vi& P){
        int on = 1, d = 1;
        while (on < sz(P)) on *= 2, d++;
        vector<vi> jmp(d, P);
        rep(i,1,d) rep(j,0,sz(P))
                jmp[i][j] = jmp[i-1][jmp[i-1][j]];
        return jmp;
int jmp(vector<vi>& tbl, int nod, int steps){
        rep(i, 0, sz(tbl))
                if(steps&(1<< i)) nod = tbl[i][nod];
        return nod;
int lca(vector<vi>& tbl, vi& depth, int a, int b) {
        if (depth[a] < depth[b]) swap(a, b);
        a = jmp(tbl, a, depth[a] - depth[b]);
        if (a == b) return a;
        for (int i = sz(tbl); i--;) {
                int c = tbl[i][a], d = tbl[i][b];
                if (c != d) a = c, b = d;
        return tbl[0][a];
```

### 9.5 CompressTree

```
* Author: Simon Lindholm
 * Date: 2016-01-14
 * License: CC0
 * Description: Given a rooted tree and a subset S of nodes, compute the minimal
 * subtree that contains all the nodes by adding all (at most |S|-1)
 * pairwise LCA's and compressing edges.
 * Returns a list of (par, orig\_index) representing a tree rooted at 0.
 * The root points to itself.
 * Time: $0(|S| \log |S|)$
 * Status: Tested at CodeForces
#pragma once
#include "LCA.h"
typedef vector<pair<int, int>> vpi;
vpi compressTree(LCA& lca, const vi& subset)
        static vi rev; rev.resize(sz(lca.time));
        vi li = subset, &T = lca.time;
        auto cmp = [\&] (int a, int b) { return T[a] < T[b]; };
        sort(all(li), cmp);
        int m = sz(li)-1;
        rep(i,0,m) {
                int a = li[i], b = li[i+1];
                li.push_back(lca.lca(a, b));
        sort(all(li), cmp);
        li.erase(unique(all(li)), li.end());
        rep(i, 0, sz(li)) rev[li[i]] = i;
        vpi ret = {pii(0, li[0])};
        rep(i, 0, sz(li) -1) {
                int a = li[i], b = li[i+1];
                ret.emplace_back(rev[lca.lca(a, b)], b);
        return ret;
```

### 9.6 DFSMatching

```
* Author: Lukas Polacek
 * Date: 2009-10-28
 * License: CC0
 * Source:
 * Description: Simple bipartite matching algorithm. Graph $g$ should be a list
 * of neighbors of the left partition, and $btoa$ should be a vector full of
 * -1's of the same size as the right partition. Returns the size of
 * the matching. \theta will be the match for vertex \theta on the right side,
 * or $-1$ if it's not matched.
 * Time: O(VE)
 * Usage: vi btoa(m, -1); dfsMatching(g, btoa);
 * Status: works
#pragma once
bool find(int j, vector<vi>& g, vi& btoa, vi& vis) {
        if (btoa[j] == -1) return 1;
        vis[j] = 1; int di = btoa[j];
        for (int e : g[di])
    if (!vis[e] && find(e, g, btoa, vis)) {
                         btoa[e] = di;
                         return 1;
        return 0;
int dfsMatching(vector<vi>& g, vi& btoa) {
        vi vis;
        rep(i,0,sz(g)) {
                vis.assign(sz(btoa), 0);
                 for (int j : g[i])
                         if (find(j, g, btoa, vis)) {
    btoa[j] = i;
                                 break;
        return sz(btoa) - (int)count(all(btoa), -1);
```

### 9.7 Dinic

```
* Author: chilli
 * Date: 2019-04-26
 * License: CC0
 * Source: https://cp-algorithms.com/graph/dinic.html
 * Description: Flow algorithm with complexity $O(VE\log U)$ where $U = \max |\
      text{cap}|$.
 * (\min(E^{1/2}, V^{2/3})E) if U = 1; (\sqrt{V}E) for bipartite
     matching.
 * Status: Tested on SPOJ FASTFLOW and SPOJ MATCHING, stress-tested
#pragma once
struct Dinic {
        struct Edge {
                int to, rev;
                11 flow() { return max(oc - c, OLL); } // if you need flows
       vi lvl, ptr, q;
vector<vector<Edge>> adj;
        Dinic(int n) : lvl(n), ptr(n), q(n), adj(n) {}
        void addEdge(int a, int b, ll c, ll rcap = 0) {
                adj[a].push_back({b, sz(adj[b]), c, c});
                adj[b].push_back({a, sz(adj[a]) - 1, rcap, rcap});
        11 dfs(int v, int t, 11 f) {
                if (v == t \mid \mid !f) return f;
                for (int& i = ptr[v]; i < sz(adj[v]); i++) {
                        Edge& e = adj[v][i];
```

```
if (lvl[e.to] == lvl[v] + 1)
                                 if (ll p = dfs(e.to, t, min(f, e.c))) {
                                         e.c -= p, adj[e.to][e.rev].c += p;
                                         return p;
                return 0;
        11 calc(int s, int t) {
                 11 \text{ flow} = 0; q[0] = s;
                 rep(L,0,31) do { // 'int L=30' maybe faster for random data
                         lvl = ptr = vi(sz(q));
                         int qi = 0, qe = lvl[s] = 1;
                         while (qi < qe && !lvl[t]) {
                                 int \hat{v} = q[qi++];
                                 for (Edge e : adj[v])
                                         if (!lvl[e.to] && e.c >> (30 - L))
                                                  q[qe++] = e.to, lvl[e.to] = lvl[
                                                       v1 + 1:
                         while (ll p = dfs(s, t, LLONG_MAX)) flow += p;
                 } while (lvl[t]);
                return flow;
        bool leftOfMinCut(int a) { return lvl[a] != 0; }
};
```

#### 9.8 DirectedMST

```
* Author: chilli, Takanori MAEHARA, Beng, Simon Lindholm
 * Date: 2019-05-10
 * License: CC0
 * Source: https://github.com/spaghetti-source/algorithm/blob/master/graph/
     arborescence.cc
 * and https://github.com/bgi343/USACO/blob/42
      d177dfb9d6ce350389583cfa71484eb8ae614c/Implementations/content/graphs
      %20(12)/Advanced/DirectedMST.h for the reconstruction
 * Description: Finds a minimum spanning
 * tree/arborescence of a directed graph, given a root node. If no MST exists,
     returns -1.
 * Time: O(E \log V)
 * Status: Stress-tested, also tested on NWERC 2018 fastestspeedrun
#pragma once
#include "../data-structures/UnionFindRollback.h"
struct Edge { int a, b; ll w; };
struct Node { /// lazy skew heap node
        Edge key;
        Node *1, *r;
        11 delta;
        void prop() {
                key.w += delta;
                if (1) 1->delta += delta;
                if (r) r->delta += delta;
                delta = 0;
        Edge top() { prop(); return key; }
Node *merge(Node *a, Node *b) {
        if (!a || !b) return a ?: b;
        a->prop(), b->prop();
        if (a->key.w > b->key.w) swap(a, b);
        swap(a->1, (a->r = merge(b, a->r)));
        return a:
void pop(Node*\& a) { a->prop(); a = merge(a->1, a->r); }
pair<11, vi> dmst(int n, int r, vector<Edge>& g) {
        RollbackUF uf(n);
        vector<Node*> heap(n);
        for (Edge e : g) heap[e.b] = merge(heap[e.b], new Node{e});
        11 \text{ res} = 0;
```

```
vi seen(n, -1), path(n), par(n);
seen[r] = r;
vector<Edge> Q(n), in(n, \{-1,-1\}), comp;
deque<tuple<int, int, vector<Edge>>> cycs;
rep(s,0,n) {
        int u = s, qi = 0, w;
        while (seen[u] < 0) {
                if (!heap[u]) return {-1,{}};
                Edge e = heap[u]->top();
                heap[u]->delta -= e.w, pop(heap[u]);
                Q[qi] = e, path[qi++] = u, seen[u] = s;
                res += e.w, u = uf.find(e.a);
                if (seen[u] == s) { /// found cycle, contract
                        Node* cyc = 0;
                        int end = qi, time = uf.time();
                        do cyc = merge(cyc, heap[w = path[--qi]]);
                        while (uf.join(u, w));
                        u = uf.find(u), heap[u] = cyc, seen[u] = -1;
                        cycs.push_front({u, time, {&Q[qi], &Q[end]}});
        rep(i, 0, qi) in[uf.find(Q[i].b)] = Q[i];
for (auto& [u,t,comp] : cycs) { // restore sol (optional)
        uf.rollback(t);
        Edge inEdge = in[u];
        for (auto& e : comp) in[uf.find(e.b)] = e;
        in[uf.find(inEdge.b)] = inEdge;
rep(i,0,n) par[i] = in[i].a;
return {res, par};
```

# 9.9 EdgeColoring

```
* Author: Simon Lindholm
 * Date: 2020-10-12
 * License: CC0
 * Source: https://en.wikipedia.org/wiki/Misra_%26_Gries_edge_coloring_algorithm
 * https://codeforces.com/blog/entry/75431 for the note about bipartite graphs.
 * Description: Given a simple, undirected graph with max degree $D$, computes a
 * $(D + 1)$-coloring of the edges such that no neighboring edges share a color.
 * ($D$-coloring is NP-hard, but can be done for bipartite graphs by repeated
      matchings of
 * max-degree nodes.)
 * Time: O(NM)
 * Status: stress-tested, tested on kattis:gamescheduling
#pragma once
vi edgeColoring(int N, vector<pii> eds) {
        vi cc(N + 1), ret(sz(eds)), fan(N), free(N), loc;
         for (pii e : eds) ++cc[e.first], ++cc[e.second];
         int u, v, ncols = *max_element(all(cc)) + 1;
        vector<vi> adj(N, vi(ncols, -1));
         for (pii e : eds) {
                 tie(u, v) = e;
fan[0] = v;
                 loc.assign(ncols, 0);
                 int at = u, end = u, d, c = free[u], ind = 0, i = 0; while (d = free[v], !loc[d] && (v = adj[u][d]) != -1)
                          loc[d] = ++ind, cc[ind] = d, fan[ind] = v;
                  cc[loc[d]] = c;
                 for (int cd = d; at != -1; cd ^= c ^ d, at = adj[at][cd])
    swap(adj[at][cd], adj[end = at][cd ^ c ^ d]);
                  while (adj[fan[i]][d] != -1) {
                          int left = fan[i], right = fan[++i], e = cc[i];
                          adj[u][e] = left;
                          adi[left][e] = u;
                          adj[right][e] = -1;
                          free[right] = e;
                  adj[u][d] = fan[i];
```

### 9.10 EdmondsKarp

```
* Author: Chen Xing
 * Date: 2009-10-13
 * License: CCO
 * Source: N/A
 * Description: Flow algorithm with quaranteed complexity $0(VE^2)$. To get edge
      flow values, compare
 * capacities before and after, and take the positive values only.
 * Status: stress-tested
 */
#pragma once
template<class T> T edmondsKarp(vector<unordered_map<int, T>>&
                graph, int source, int sink) {
        assert (source != sink);
        T flow = 0:
        vi par(sz(graph)), q = par;
        for (;;) {
                fill(all(par), -1);
                par[source] = 0;
                int ptr = 1;
                q[0] = source;
                rep(i,0,ptr) {
                        int x = q[i];
                        for (auto e : graph[x]) {
                                if (par[e.first] == -1 && e.second > 0) {
                                        par[e.first] = x;
                                        q[ptr++] = e.first;
                                        if (e.first == sink) goto out;
                return flow;
out:
                T inc = numeric limits<T>::max();
                for (int y = sink; y != source; y = par[y])
                        inc = min(inc, graph[par[y]][y]);
                flow += inc;
                for (int y = sink; y != source; y = par[y]) {
                        int p = par[y];
                        if ((graph[p][y] -= inc) \le 0) graph[p].erase(y);
                        graph[y][p] += inc;
```

### 9.11 EulerWalk

```
/**
  * Author: Simon Lindholm
  * Date: 2019-12-31
  * License: CC0
  * Source: folklore
  * Description: Eulerian undirected/directed path/cycle algorithm.
  * Input should be a vector of (dest, global edge index), where
  * for undirected graphs, forward/backward edges have the same index.
  * Returns a list of nodes in the Eulerian path/cycle with src at both start and end, or
```

```
* empty list if no cycle/path exists.
* To get edge indices back, add .second to s and ret.
 * Time: O(V + E)
* Status: stress-tested
#pragma once
vi eulerWalk(vector<vector<pii>>& gr, int nedges, int src=0) {
        int n = sz(qr);
        vi D(n), its(n), eu(nedges), ret, s = {src};
        D[src]++; // to allow Euler paths, not just cycles
        while (!s.empty()) {
                int x = s.back(), y, e, &it = its[x], end = sz(gr[x]);
                if (it == end) { ret.push_back(x); s.pop_back(); continue; }
                tie(y, e) = gr[x][it++];
                if (!eu[e]) {
                        D[x] --, D[y] ++;
                        eu[e] = 1; s.push_back(y);
        for (int x : D) if (x < 0 \mid \mid sz(ret) != nedges+1) return {};
        return {ret.rbegin(), ret.rend()};
```

### 9.12 FloydWarshall

```
* Author: Simon Lindholm
 * Date: 2016-12-15
 * License: CC0
 * Source: http://en.wikipedia.org/wiki/Floyd-Warshall_algorithm
 * Description: Calculates all-pairs shortest path in a directed graph that
     might have negative edge weights.
 * Input is an distance matrix $m$, where $m[i][j] = \textt{inf}$ if $i$ and
     $j$ are not adjacent.
 * As output, m[i][j] is set to the shortest distance between i and j, \
     texttt{inf} if no path,
 * or \texttt{-inf} if the path goes through a negative-weight cycle.
 * Time: O(N^3)
 * Status: slightly tested
#pragma once
const 11 inf = 1LL << 62;
void floydWarshall(vector<vector<ll>>& m) {
        int n = sz(m);
        rep(i,0,n) m[i][i] = min(m[i][i], OLL);
        rep(k,0,n) rep(i,0,n) rep(j,0,n)
               if (m[i][k] != inf && m[k][j] != inf) {
                        auto newDist = max(m[i][k] + m[k][j], -inf);
                        m[i][j] = min(m[i][j], newDist);
        rep(k,0,n) if (m[k][k] < 0) rep(i,0,n) rep(j,0,n)
                if (m[i][k] != inf && m[k][j] != inf) m[i][j] = -inf;
```

### 9.13 GeneralMatching

```
* Author: Simon Lindholm
* Date: 2016-12-09
* License: CC0
* Source: http://www.mimuw.edu.pl/~mucha/pub/mucha_sankowski_focs04.pdf
* Description: Matching for general graphs.
* Fails with probability $N / mod$.
* Time: O(N^3)
* Status: not very well tested
*/
#pragma once
#include "../numerical/MatrixInverse-mod.h"

vector<pii> generalMatching(int N, vector<pii>& ed) {
```

```
vector<vector<ll>> mat(N, vector<ll>(N)), A;
for (pii pa : ed) {
        int a = pa.first, b = pa.second, r = rand() % mod;
       mat[a][b] = r, mat[b][a] = (mod - r) % mod;
int r = matInv(A = mat), M = 2*N - r, fi, fj;
assert(r % 2 == 0);
if (M != N) do {
        mat.resize(M, vector<11>(M));
        rep(i,0,N) {
                mat[i].resize(M);
                rep(j,N,M) {
                        int r = rand() % mod;
                        mat[i][j] = r, mat[j][i] = (mod - r) % mod;
} while (matInv(A = mat) != M);
vi has(M, 1); vector<pii> ret;
rep(it, 0, M/2)
        rep(i, 0, M) if (has[i])
                rep(j,i+1,M) if (A[i][j] && mat[i][j]) {
                       fi = i; fj = j; goto done;
        } assert(0); done:
        if (fj < N) ret.emplace_back(fi, fj);</pre>
        has[fi] = has[fj] = 0;
        rep(sw,0,2) {
                11 a = modpow(A[fi][fj], mod-2);
                rep(i,0,M) if (has[i] && A[i][fj]) {
                        ll b = A[i][fj] * a % mod;
                        rep(j, 0, M) A[i][j] = (A[i][j] - A[fi][j] * b) %
                swap(fi,fj);
return ret:
```

#### 9.14 GlobalMinCut

```
* Author: Simon Lindholm
 * Date: 2021-01-09
 * License: CC0
 * Source: https://en.wikipedia.org/wiki/Stoer%E2%80%93Wagner_algorithm
 * Description: Find a global minimum cut in an undirected graph, as represented
      by an adjacency matrix.
 * Time: O(V^3)
 * Status: Stress-tested together with GomoryHu
#pragma once
pair<int, vi> globalMinCut(vector<vi> mat) {
        pair<int, vi> best = {INT_MAX, {}};
        int n = sz(mat);
        vector<vi> co(n);
        rep(i,0,n) co[i] = {i};
        rep(ph, 1, n) {
                vi w = mat[0];
                size_t s = 0, t = 0;
                rep(it,0,n-ph) { // O(V^2) -> O(E log V) with prio. queue
                        w[t] = INT_MIN;
                        s = t, t = max\_element(all(w)) - w.begin();
                        rep(i, 0, n) w[i] += mat[t][i];
                best = min(best, \{w[t] - mat[t][t], co[t]\});
                co[s].insert(co[s].end(), all(co[t]));
                rep(i,0,n) mat[s][i] += mat[t][i];
                rep(i, 0, n) mat[i][s] = mat[s][i];
                mat[0][t] = INT_MIN;
        return best;
```

### 9.15 GomoryHu

```
* Author: chilli, Takanori MAEHARA
 * Date: 2020-04-03
 * License: CC0
 * Source: https://github.com/spaghetti-source/algorithm/blob/master/graph/
     gomory_hu_tree.cc#L102
 * Description: Given a list of edges representing an undirected flow graph,
 * returns edges of the Gomory-Hu tree. The max flow between any pair of
 * vertices is given by minimum edge weight along the Gomory-Hu tree path.
 * Time: $0(V)$ Flow Computations
 * Status: Tested on CERC 2015 J, stress-tested
 \star Details: The implementation used here is not actually the original
 * Gomory-Hu, but Gusfield's simplified version: "Very simple methods for all
 * pairs network flow analysis". PushRelabel is used here, but any flow
 * implementation that supports 'leftOfMinCut' also works.
 */
#pragma once
#include "PushRelabel.h"
typedef array<11, 3> Edge;
vector<Edge> gomoryHu(int N, vector<Edge> ed) {
        vector<Edge> tree;
        vi par(N);
        rep(i,1,N) {
                PushRelabel D(N); // Dinic also works
                for (Edge t : ed) D.addEdge(t[0], t[1], t[2], t[2]);
                tree.push_back({i, par[i], D.calc(i, par[i])});
                rep(j,i+1,N)
                        if (par[j] == par[i] && D.leftOfMinCut(j)) par[j] = i;
        return tree;
```

#### 9.16 HLD

```
* Author: Benjamin Qi, Oleksandr Kulkov, chilli
 * Date: 2020-01-12
 * Source: https://codeforces.com/blog/entry/53170, https://github.com/bqi343/
     USACO/blob/master/Implementations/content/graphs%20(12)/Trees%20(10)/HLD
      %20(10.3).h
 * Description: Decomposes a tree into vertex disjoint heavy paths and light
 * edges such that the path from any leaf to the root contains at most log(n)
 * light edges. Code does additive modifications and max queries, but can
 * support commutative segtree modifications/queries on paths and subtrees.
 * Takes as input the full adjacency list. VALS\_EDGES being true means that
 * values are stored in the edges, as opposed to the nodes. All values
 * initialized to the segtree default. Root must be 0.
 * Time: O((\log N)^2)
 * Status: stress-tested against old HLD
#pragma once
#include "../data-structures/LazySegmentTree.h"
template <bool VALS_EDGES> struct HLD {
       int N, tim = 0;
        vector<vi> adj;
        vi par, siz, rt, pos;
       Node *tree;
        HLD(vector<vi> adj )
                : N(sz(adj_)), adj(adj_), par(N, -1), siz(N, 1),
                  rt(N),pos(N),tree(new Node(0, N)) { dfsSz(0); dfsHld(0); }
        void dfsSz(int v) {
                for (int& u : adj[v]) {
                        adj[u].erase(find(all(adj[u]), v));
                        par[u] = v;
                        dfsSz(u);
```

```
siz[v] += siz[u];
               if (siz[u] > siz[adj[v][0]]) swap(u, adj[v][0]);
void dfsHld(int v) {
       pos[v] = tim++;
        for (int u : adj[v]) {
               rt[u] = (u == adj[v][0] ? rt[v] : u);
               dfsHld(u);
template <class B> void process(int u, int v, B op) {
       for (;; v = par[rt[v]]) {
                if (pos[u] > pos[v]) swap(u, v);
               if (rt[u] == rt[v]) break;
               op(pos[rt[v]], pos[v] + 1);
       op(pos[u] + VALS_EDGES, pos[v] + 1);
void modifyPath(int u, int v, int val) {
       process(u, v, [&](int 1, int r) { tree->add(1, r, val); });
int queryPath(int u, int v) { // Modify depending on problem
       int res = -1e9;
       process(u, v, [&](int l, int r) {
                        res = max(res, tree->query(1, r));
       return res;
int querySubtree(int v) { // modifySubtree is similar
        return tree->query(pos[v] + VALS_EDGES, pos[v] + siz[v]);
```

# 9.17 hopcroftKarp

};

```
/**
* Author: Chen Xing
 * Date: 2009-10-13
 * License: CC0
 * Source: N/A
 * Description: Fast bipartite matching algorithm. Graph $g$ should be a list
* of neighbors of the left partition, and \theta should be a vector full of -1's of the same size as the right partition. Returns the size of
 * the matching. $btoa[i]$ will be the match for vertex $i$ on the right side,
 * or $-1$ if it's not matched.
 * Usage: vi btoa(m, -1); hopcroftKarp(g, btoa);
 * Time: O(\sqrt{V}E)
 * Status: stress-tested by MinimumVertexCover, and tested on oldkattis.
      adkbipmatch and SPOJ:MATCHING
#pragma once
bool dfs(int a, int L, vector<vi>& g, vi& btoa, vi& A, vi& B) {
        if (A[a] != L) return 0;
        A[a] = -1;
        for (int b : g[a]) if (B[b] == L + 1) {
                B[b] = 0;
                 if (btoa[b] == -1 \mid \mid dfs(btoa[b], L + 1, g, btoa, A, B))
                         return btoa[b] = a, 1;
        return 0:
int hopcroftKarp(vector<vi>& g, vi& btoa) {
        int res = 0:
        vi A(g.size()), B(btoa.size()), cur, next;
        for (;;) {
                 fill(all(A), 0);
                 fill(all(B), 0);
                 /// Find the starting nodes for BFS (i.e. layer 0).
                 cur.clear();
                 for (int a : btoa) if (a !=-1) A[a] = -1;
                 rep(a, 0, sz(g)) if(A[a] == 0) cur.push_back(a);
                 /// Find all layers using bfs.
```

```
for (int lay = 1;; lay++) {
        bool islast = 0;
        next.clear();
        for (int a : cur) for (int b : g[a]) {
                if (btoa[b] == -1) {
                         B[b] = lay;
                         islast = 1;
                 else if (btoa[b] != a && !B[b]) {
                         B[b] = lay;
                         next.push_back(btoa[b]);
        if (islast) break;
        if (next.empty()) return res;
        for (int a : next) A[a] = lay;
        cur.swap(next);
^{^{\prime}}/// Use DFS to scan for augmenting paths.
rep(a, 0, sz(q))
        res += dfs(a, 0, g, btoa, A, B);
```

#### 9.18 LCA

```
* Author: chilli, pajenegod
 * Date: 2020-02-20
 * License: CC0
 * Source: Folklore
 * Description: Data structure for computing lowest common ancestors in a tree
 \star (with \hat{0} as root). C should be an adjacency list of the tree, either directed
 * or undirected.
 * Time: $0(N \log N + 0)$
 * Status: stress-tested
 * /
#pragma once
#include "../data-structures/RMQ.h"
struct LCA {
        int T = 0;
        vi time, path, ret;
        RMQ<int> rmq;
        LCA(vector < vi > \& C) : time(sz(C)), rmq((dfs(C,0,-1), ret)) {}
        void dfs(vector<vi>& C, int v, int par) {
                time[v] = T++;
                for (int y : C[v]) if (y != par) {
                        path.push_back(v), ret.push_back(time[v]);
                        dfs(C, y, v);
        int lca(int a, int b) {
                if (a == b) return a;
                tie(a, b) = minmax(time[a], time[b]);
                return path[rmq.query(a, b)];
        //dist(a,b) {return depth[a] + depth[b] - 2*depth[lca(a,b)];}
};
```

### 9.19 LinkCutTree

```
/**
    * Author: Simon Lindholm
    Date: 2016-07-25
    * Source: https://github.com/ngthanhtrung23/ACM_Notebook_new/blob/master/
        DataStructure/LinkCutTree.h
    * Description: Represents a forest of unrooted trees. You can add and remove
    * edges (as long as the result is still a forest), and check whether
```

```
\star two nodes are in the same tree.
 * Time: All operations take amortized O(\log N).
 * Status: Stress-tested a bit for N <= 20
#pragma once
struct Node { // Splay tree. Root's pp contains tree's parent.
        Node *p = 0, *pp = 0, *c[2];
        bool flip = 0;
        Node() { c[0] = c[1] = 0; fix(); }
        void fix() {
                 if (c[0]) c[0] -> p = this;
                if (c[1]) c[1] -> p = this;
                // (+ update sum of subtree elements etc. if wanted)
        void pushFlip() {
                 if (!flip) return;
                flip = 0; swap(c[0], c[1]); if (c[0]) c[0]->flip ^= 1;
                if (c[1]) c[1]->flip ^= 1;
        int up() { return p ? p \rightarrow c[1] == this : -1; }
        void rot(int i, int b) {
                 int h = i \hat{b};
                 Node *x = c[i], *y = b == 2 ? x : x -> c[h], *z = b ? y : x;
                if ((y->p = p)) p->c[up()] = y;
c[i] = z->c[i ^ 1];
                 if (b < 2) {
                         x->c[h] = y->c[h ^ 1];
                         y - > c[h ^ 1] = x;
                 z\rightarrow c[i ^1] = this;
                 fix(); x->fix(); y->fix();
                 if (p) p->fix();
                swap(pp, y->pp);
        void splay() { /// Splay this up to the root. Always finishes without
             flip set.
                 for (pushFlip(); p; ) {
                         if (p->p) p->p->pushFlip();
                         p->pushFlip(); pushFlip();
                         int c1 = up(), c2 = p->up();
                         if (c2 == -1) p->rot (c1, 2);
                         else p->p->rot(c2, c1 != c2);
        Node* first() { /// Return the min element of the subtree rooted at this
             , splayed to the top.
                pushFlip();
                 return c[0] ? c[0]->first() : (splay(), this);
};
struct LinkCut {
        vector<Node> node;
        LinkCut(int N) : node(N) {}
        void link(int u, int v) { // add an edge (u, v)
                 assert(!connected(u, v));
                 makeRoot(&node[u]);
                node[u].pp = &node[v];
        void cut(int u, int v) { // remove an edge (u, v)
                 Node *x = &node[u], *top = &node[v];
                 makeRoot(top); x->splay();
                 assert(top == (x-pp ?: x-c[0]));
                 if (x->pp) x->pp = 0;
                 else {
                         x->c[0] = top->p = 0;
                         x->fix();
        bool connected(int u, int v) { // are u, v in the same tree?
                Node* nu = access(&node[u]) ->first();
                 return nu == access(&node[v])->first();
        void makeRoot(Node* u) { /// Move u to root of represented tree.
                 access (u):
                 u->splay();
```

```
if(u->c[0]) {
                          u - > c[0] - > p = 0;
                          u - c[0] - flip ^= 1;
                          u - c[0] - pp = u;
                          u \rightarrow c[0] = 0;
                          u->fix();
        Node* access(Node* u) { /// Move u to root aux tree. Return the root of
             the root aux tree.
                 u->splav();
                 while (Node* pp = u->pp) {
                          pp->splay(); u->pp = 0;
                          if (pp->c[1]) {
                                   pp - c[1] - p = 0; pp - c[1] - pp = pp; 
                          pp - c[1] = u; pp - fix(); u = pp;
                 return u:
};
```

### 9.20 MaximalCliques

```
* Author: Simon Lindholm
 * Date: 2018-07-18
 * License: CC0
 * Source: https://en.wikipedia.org/wiki/Bron%E2%80%93Kerbosch algorithm
 * Description: Runs a callback for all maximal cliques in a graph (given as a
 * symmetric bitset matrix; self-edges not allowed). Callback is given a bitset
 * representing the maximal clique.
 * Time: O(3^{n/3}), much faster for sparse graphs
 * Status: stress-tested
 */
#pragma once
/// Possible optimization: on the top-most
/// recursion level, ignore 'cands', and go through nodes in order of increasing
/// degree, where degrees go down as nodes are removed.
/// (mostly irrelevant given MaximumClique)
typedef bitset<128> B;
template<class F>
void cliques(vectorB eds, F f, B P = B(), B X=\{\}, B R=\{\}) {
        if (!P.any()) { if (!X.any()) f(R); return; }
        auto q = (P | X)._Find_first();
        auto cands = P & ~eds[q];
rep(i,0,sz(eds)) if (cands[i]) {
                R[i] = 1;
                cliques(eds, f, P & eds[i], X & eds[i], R);
                R[i] = P[i] = 0; X[i] = 1;
```

# 9.21 MaximumClique

```
struct Vertex { int i, d=0; };
typedef vector<Vertex> vv;
vv V;
vector<vi> C;
vi qmax, q, S, old;
void init(vv& r) {
        for (auto& v : r) v.d = 0;
        for (auto& v : r) for (auto j : r) v.d += e[v.i][j.i];
        sort(all(r), [](auto a, auto b) { return a.d > b.d; });
        int mxD = r[0].d;
        rep(i, 0, sz(r)) r[i].d = min(i, mxD) + 1;
void expand(vv& R, int lev = 1) {
        S[lev] += S[lev - 1] - old[lev];
        old[lev] = S[lev - 1];
        while (sz(R)) {
                 if (sz(q) + R.back().d <= sz(qmax)) return;
                 q.push_back(R.back().i);
                 for (auto v:R) if (e[R.back().i][v.i]) T.push_back(\{v.i\})
                         if (S[lev]++ / ++pk < limit) init(T);</pre>
                         int j = 0, mxk = 1, mnk = max(sz(qmax) - sz(q) + 1, 1);
                         C[1].clear(), C[2].clear();
                         for (auto v : T) {
                                 int k = 1;
                                 auto f = [\&] (int i) \{ return e[v.i][i];
                                  while (any_of(all(C[k]), f)) k++;
                                  if (k > mxk) mxk = k, C[mxk + 1].clear()
                                  if (k < mnk) T[j++].i = v.i;
                                 C[k].push_back(v.i);
                         if (j > 0) T[j - 1].d = 0;
rep(k,mnk,mxk + 1) for (int i : C[k])
                                 T[j].i = i, T[j++].d = k;
                         expand(T, lev + 1);
                 } else if (sz(q) > sz(qmax)) qmax = q;
                 q.pop_back(), R.pop_back();
vi maxClique() { init(V), expand(V); return qmax; }
Maxclique(vb conn) : e(conn), C(sz(e)+1), S(sz(C)), old(S) {
        rep(i,0,sz(e)) V.push_back({i});
```

### 9.22 MaximumIndependentSet

```
/**
 * Author: chilli
 * Date: 2019-05-17
 * Source: Wikipedia
 * Description: To obtain a maximum independent set of a graph, find a max
 * clique of the complement. If the graph is bipartite, see MinimumVertexCover.
 */
```

### 9.23 MinCostMaxFlow

};

```
/**
  * Author: Stanford
  * Date: Unknown
  * Source: Stanford Notebook
  * Description: Min-cost max-flow.
  * If costs can be negative, call setpi before maxflow, but note that negative cost cycles are not supported.
  * To obtain the actual flow, look at positive values only.
```

```
* Status: Tested on kattis:mincostmaxflow, stress-tested against another
      implementation
 * Time: \$O(F E \setminus \log(V))\$ where F is max flow. \$O(VE)\$ for setpi.
#pragma once
// #include <bits/extc++.h> /// include-line, keep-include
const ll INF = numeric_limits<ll>::max() / 4;
struct MCMF {
        struct edge {
                int from, to, rev;
                11 cap, cost, flow;
        int N:
        vector<vector<edge>> ed;
        vi seen;
vector<ll> dist, pi;
        vector<edge*> par;
        MCMF(int N) : N(N), ed(N), seen(N), dist(N), pi(N), par(N) {}
        void addEdge(int from, int to, ll cap, ll cost) {
                if (from == to) return;
                ed[from].push_back(edge{ from,to,sz(ed[to]),cap,cost,0 });
                ed[to].push_back(edge{ to,from,sz(ed[from])-1,0,-cost,0 });
        void path(int s) {
                fill(all(seen), 0);
                fill(all(dist), INF);
dist[s] = 0; ll di;
                 __gnu_pbds::priority_queue<pair<11, int>> q;
                vector<decltype(q)::point_iterator> its(N);
                q.push({ 0, s });
                while (!q.empty()) {
                        s = q.top().second; q.pop();
                         seen[s] = 1; di = dist[s] + pi[s];
                         for (edge& e : ed[s]) if (!seen[e.to]) {
                                 11 val = di - pi[e.to] + e.cost;
                                 if (e.cap - e.flow > 0 && val < dist[e.to]) {
                                         dist[e.to] = val;
                                         par[e.to] = &e;
                                         if (its[e.to] == q.end())
                                                 its[e.to] = q.push({ -dist[e.to
                                                      ], e.to });
                                         else
                                                 q.modify(its[e.to], { -dist[e.to
                                                      ], e.to });
                rep(i, 0, N) pi[i] = min(pi[i] + dist[i], INF);
        pair<11, 11> maxflow(int s, int t) {
                 11 totflow = 0, totcost = 0;
                while (path(s), seen[t]) {
                        11 fl = INF;
                         for (edge * x = par[t]; x; x = par[x->from])
                                 fl = min(fl, x->cap - x->flow);
                         totflow += fl;
                         for (edge* x = par[t]; x; x = par[x->from]) {
                                 ed[x->to][x->rev].flow -= fl;
                rep(i,0,N) for(edge& e : ed[i]) totcost += e.cost * e.flow;
                return {totflow, totcost/2};
        // If some costs can be negative, call this before maxflow:
        void setpi(int s) { // (otherwise, leave this out)
                fill(all(pi), INF); pi[s] = 0;
                int it = N, ch = 1; ll v;
```

#### 9.24 MinCut

```
/**
 * Author: Simon Lindholm
 * Date: 2015-05-13
 * Source: Wikipedia
 * Description: After running max-flow, the left side of a min-cut from $s$ to
    $t$ is given
 * by all vertices reachable from $s$, only traversing edges with positive
    residual capacity.
 * Status: works
 */
```

#### 9.25 MinimumVertexCover

```
* Author: Johan Sannemo, Simon Lindholm
 * Date: 2016-12-15
 * License: CC0
 * Description: Finds a minimum vertex cover in a bipartite graph.
 * The size is the same as the size of a maximum matching, and
 \star the complement is a maximum independent set.
 * Status: stress-tested
#pragma once
#include "DFSMatching.h"
vi cover(vector<vi>& g, int n, int m) {
        vi match(m, -1);
        int res = dfsMatching(g, match);
        vector<bool> lfound(n, true), seen(m);
        for (int it : match) if (it != -1) lfound[it] = false;
        vi q, cover;
        rep(i,0,n) if (lfound[i]) q.push_back(i);
        while (!q.empty()) {
                int i = q.back(); q.pop_back();
                lfound[i] = 1;
                for (int e : g[i]) if (!seen[e] && match[e] != -1) {
                        seen[e] = true;
                        q.push_back(match[e]);
        rep(i,0,n) if (!lfound[i]) cover.push_back(i);
        rep(i,0,m) if (seen[i]) cover.push_back(n+i);
        assert(sz(cover) == res);
        return cover;
```

### 9.26 PushRelabel

```
/**
  * Author: Simon Lindholm
  * Date: 2015-02-24
  * License: CC0
  * Source: Wikipedia, tinyKACTL
  * Description: Push-relabel using the highest label selection rule and the gap heuristic. Quite fast in practice.
  * To obtain the actual flow, look at positive values only.
  * Time: $0(V^2\sqrt E)$
```

```
* Status: Tested on Kattis and SPOJ, and stress-tested
 * /
#pragma once
struct PushRelabel {
        struct Edge {
                 int dest, back;
                 11 f, c;
         vector<vector<Edge>> q;
        vector<11> ec;
         vector<Edge*> cur;
         vector<vi> hs; vi H;
         PushRelabel(int n): g(n), ec(n), cur(n), hs(2*n), H(n) {}
         void addEdge(int s, int t, ll cap, ll rcap=0) {
                 if (s == t) return;
                 g[s].push_back({t, sz(g[t]), 0, cap});
                 g[t].push_back({s, sz(g[s])-1, 0, rcap});
         void addFlow(Edge& e, ll f) {
                 Edge &back = g[e.dest][e.back];
                 if (!ec[e.dest] && f) hs[H[e.dest]].push_back(e.dest);
                 e.f += f; e.c -= f; ec[e.dest] += f;
                 back.f -= f; back.c += f; ec[back.dest] -= f;
         ll calc(int s, int t) {
                 int v = sz(g); H[s] = v; ec[t] = 1;
                 vi co(2*v); co[0] = v-1;
                 rep(i,0,v) cur[i] = g[i].data();
                 for (Edge& e : g[s]) addFlow(e, e.c);
                 for (int hi = 0;;) {
                          while (hs[hi].empty()) if (!hi--) return -ec[s];
                          int u = hs[hi].back(); hs[hi].pop_back();
                          while (ec[u] > 0) // discharge u
                                   if (\operatorname{cur}[u] == \operatorname{g}[u].\operatorname{data}() + \operatorname{sz}(\operatorname{g}[u])) {
                                           H[u] = 1e9;
                                            for (Edge& e : g[u]) if (e.c && H[u] > H
                                                 [e.dest]+1)
                                                    H[u] = H[e.dest]+1, cur[u] = &e;
                                            if (++co[H[u]], !--co[hi] && hi < v)
                                                    rep(i, 0, v) if (hi < H[i] \&\& H[i]
                                                           < v)
                                                             --co[H[i]], H[i] = v +
                                                                  1;
                                            hi = H[u];
                                   } else if (cur[u] \rightarrow c \&\& H[u] == H[cur[u] \rightarrow dest
                                           addFlow(*cur[u], min(ec[u], cur[u]->c));
                                   else ++cur[u];
        bool leftOfMinCut(int a) { return H[a] >= sz(g); }
};
```

### 9.27 SCC

```
/**

* Author: Lukas Polacek

* Date: 2009-10-28

* License: CCO

* Source: Czech graph algorithms book, by Demel. (Tarjan's algorithm)

* Description: Finds strongly connected components in a

* directed graph. If vertices $u, v$ belong to the same component,

* we can reach $u$ from $v$ and vice versa.

* Usage: scc(graph, [\&](vi\& v) { ... }) visits all components

* in reverse topological order. comp[i] holds the component

* index of a node (a component only has edges to components with

* lower index). ncomps will contain the number of components.

* Time: O(E + V)

* Status: Bruteforce-tested for N <= 5

*/

#pragma once
```

```
vi val, comp, z, cont;
int Time, ncomps;
template < class G, class F> int dfs (int j, G& g, F& f) {
        int low = val[j] = ++Time, x; z.push_back(j);
        for (auto e : g[j]) if (comp[e] < 0)
                low = min(low, val[e] ?: dfs(e,g,f));
        if (low == val[j]) {
                        x = z.back(); z.pop_back();
                        comp[x] = ncomps;
                        cont.push_back(x);
                } while (x != j);
                f(cont); cont.clear();
                ncomps++;
        return val[j] = low;
template < class G, class F> void scc(G& g, F f) {
        int n = sz(g);
        val.assign(n, 0); comp.assign(n, -1);
        Time = ncomps = 0;
        rep(i, 0, n) if (comp[i] < 0) dfs(i, g, f);
```

### 9.28 TopoSort

```
* Author: Unknown
 * Date: 2002-09-13
 * Source: predates tinyKACTL
 * Description: Topological sorting. Given is an oriented graph.
 * Output is an ordering of vertices, such that there are edges only from left
      to right.
 \star If there are cycles, the returned list will have size smaller than n\ --
      nodes reachable
 * from cycles will not be returned.
 * Time: $0(|V|+|E|)$
 * Status: stress-tested
#pragma once
vi topoSort(const vector<vi>& gr) {
         vi indeg(sz(gr)), q;
        for (auto@ li : gr) for (int x : li) indeg[x]++; rep(i,0,sz(gr)) if (indeg[i] == 0) q.push_back(i);
         rep(j, 0, sz(q)) for (int x : gr[q[j]])
                 if (--indeg[x] == 0) q.push_back(x);
```

### 9.29 WeightedMatching

```
if (a.empty()) return {0, {}};
int n = sz(a) + 1, m = sz(a[0]) + 1;
vi u(n), v(m), p(m), ans(n-1);
rep(i,1,n) {
        int j0 = 0; // add "dummy" worker 0
        vi dist(m, INT_MAX), pre(m, -1);
        vector<bool> done(m + 1);
        do { // dijkstra
                done[j0] = true;
                int i\bar{0} = p[j0], j1, delta = INT_MAX;
                rep(j,1,m) if (!done[j]) {
                        auto cur = a[i0 - 1][j - 1] - u[i0] - v[j];
                        if (cur < dist[j]) dist[j] = cur, pre[j] = j0;
                        if (dist[j] < delta) delta = dist[j], j1 = j;</pre>
                rep(j,0,m)
                        if (done[j]) u[p[j]] += delta, v[j] -= delta;
                        else dist[j] -= delta;
                j0 = j1;
        } while (p[j0]);
        while (j0) { // update alternating path
                int j1 = pre[j0];
                p[j0] = p[j1], j0 = j1;
rep(j,1,m) if (p[j]) ans[p[j] - 1] = j - 1;
return {-v[0], ans}; // min cost
```

# 10 kactl forked/number-theory

#### 10.1 ContinuedFractions

```
* Author: Simon Lindholm
 * Date: 2018-07-15
 * License: CC0
 * Source: Wikipedia
 * Description: Given $N$ and a real number $x \ge 0$, finds the closest
      rational approximation p/q with p, q \le N.
 * It will obey p/q - x \le 1/qN.
 * For consecutive convergents, p_{k+1}q_k - q_{k+1}p_k = (-1)^k.
 * (p_k/q_k alternates between x>x and x<.)
 * If $x$ is rational, $y$ eventually becomes $\infty$;
 * if $x$ is the root of a degree $2$ polynomial the $a$'s eventually become
 * Time: O(\log N)
 * Status: stress-tested for n <= 300
typedef double d; // for N \tilde{\ } 1e7; long double for N \tilde{\ } 1e9
pair<11, 11> approximate(d x, 11 N) {
        11 LP = 0, LQ = 1, P = 1, Q = 0, inf = LLONG_MAX; d y = x;
                ll lim = min(P ? (N-LP) / P : inf, Q ? (N-LQ) / Q : inf),
                   a = (ll) floor(y), b = min(a, lim),
                   NP = b*P + LP, NQ = b*Q + LQ;
                if (a > b) {
                        // If b > a/2, we have a semi-convergent that gives us a
                        // better approximation; if b = a/2, we *may* have one.
                        // Return {P, Q} here for a more canonical approximation
                        return (abs(x - (d)NP / (d)NQ) < abs(x - (d)P / (d)Q)) ?
                                make_pair(NP, NQ) : make_pair(P, Q);
                if (abs(y = 1/(y - (d)a)) > 3*N) {
                        return {NP, NQ};
                LP = P; P = NP;
                LQ = Q; Q = NQ;
```

### 10.2 CRT

```
/**
 * Author: Simon Lindholm
 * Date: 2019-05-22
 * License: CC0
 * Description: Chinese Remainder Theorem.
 * \texttt{crt(a, m, b, n)} computes $x$ such that $x\equiv a \pmod m$, $x\equiv
      b \pmod n$.
 * If |a| < m and |b| < n, x will obey 0 \le x < \text{(m, n)}.
 * Assumes mn < 2^{62}.
 * Time: $\log(n)$
 * Status: Works
#pragma once
#include "euclid.h"
ll crt(ll a, ll m, ll b, ll n) {
        if (n > m) swap(a, b), swap(m, n);
        ll x, y, g = euclid(m, n, x, y);
        assert((a - b) % g == 0); // else no solution
        x = (b - a) % n * x % n / q * m + a;
       return x < 0 ? x + m*n/q : x;
```

#### 10.3 Eratosthenes

```
* Author: Hakan Terelius
 * Date: 2009-08-26
 * License: CC0
 * Source: http://en.wikipedia.org/wiki/Sieve_of_Eratosthenes
 \star Description: Prime sieve for generating all primes up to a certain limit.
     isprime$[i]$ is true iff $i$ is a prime.
 * Time: lim=100'000'000 $\approx$ 0.8 s. Runs 30\% faster if only odd indices
      are stored.
 * Status: Tested
#pragma once
const int MAX_PR = 5'000'000;
bitset<MAX_PR> isprime;
vi eratosthenesSieve(int lim)
        isprime.set(); isprime[0] = isprime[1] = 0;
        for (int i = 4; i < \lim; i += 2) isprime[i] = 0;
        for (int i = 3; i * i < 1im; i += 2) if (isprime[i])
                for (int j = i*i; j < lim; j += i*2) isprime[j] = 0;
        rep(i,2,lim) if (isprime[i]) pr.push_back(i);
        return pr;
```

### 10.4 euclid

```
/**
    * Author: Unknown
    * Date: 2002-09-15
    * Source: predates tinyKACTL
    * Description: Finds two integers $x$ and $y$, such that $ax+by=\gcd(a,b)$. If
    * you just need gcd, use the built in \texttt{\\_\gcd} instead.
    * If $a$ and $b$ are coprime, then $x$ is the inverse of $a \pmod{b}$.
    */
#pragma once
```

```
ll euclid(ll a, ll b, ll &x, ll &y) {
    if (!b) return x = 1, y = 0, a;
    ll d = euclid(b, a % b, y, x);
    return y -= a/b * x, d;
}
```

#### 10.5 Factor

```
* Author: chilli, SJTU, pajenegod
 * Date: 2020-03-04
 * License: CC0
 * Source: own
 * Description: Pollard-rho randomized factorization algorithm. Returns prime
 * factors of a number, in arbitrary order (e.g. 2299 \rightarrow \{11, 19, 11\}).
 * Time: 0(n^{1/4}), less for numbers with small factors.
 * Status: stress-tested
\star Details: This implementation uses the improvement described here
 * (https://en.wikipedia.org/wiki/Pollard%27s_rho_algorithm#Variants), where
 * one can accumulate gcd calls by some factor (40 chosen here through
 \star exhaustive testing). This improves performance by approximately 6-10x
 * depending on the inputs and speed of gcd. Benchmark found here:
 * (https://ideone.com/nGGD9T)
* GCD can be improved by a factor of 1.75x using Binary GCD
 * (https://lemire.me/blog/2013/12/26/fastest-way-to-compute-the-greatest-common
 \star However, with the gcd accumulation the bottleneck moves from the gcd calls
 * to the modmul. As GCD only constitutes ~12% of runtime, speeding it up
 * doesn't matter so much.
* This code can probably be sped up by using a faster mod mul - potentially
 * montgomery reduction on 128 bit integers.
 * Alternatively, one can use a quadratic sieve for an asymptotic improvement,
 * which starts being faster in practice around 1e13.
 * Brent's cycle finding algorithm was tested, but doesn't reduce modmul calls
 * significantly.
 * Subtle implementation notes:
 \star - prd starts off as 2 to handle the case n = 4; it's harmless for other n
    since we're guaranteed that n > 2. (Pollard rho has problems with prime
    powers in general, but all larger ones happen to work.)
  - t starts off as 30 to make the first gcd check come earlier, as an
    optimization for small numbers.
 \star - we vary f between restarts because the cycle finding algorithm does not
     find the first element in the cycle but rather one at distance k*|cycle|
     from the start, and that can result in continual failures if all cycles
     have the same size for all prime factors. E.g. fixing f(x) = x^2 + 1 would
     loop infinitely for n = 352523 * 352817, where all cycles have size 821.
\star - we operate on residues in [i, n + i) which modmul is not designed to
    handle, but specifically modmul(x, x) still turns out to work for small
     enough i. (With reference to the proof in modmul-proof.tex, the argument
    for "S is in [-c, 2c)" goes through unchanged, while S < 2^63 now follows from S < 2c and S = x^2 (mod c) together implying S < c + i^2.)
#pragma once
#include "ModMulLL.h"
#include "MillerRabin.h"
ull pollard(ull n) {
        ull x = 0, y = 0, t = 30, prd = 2, i = 1, q;
        auto f = [\&](ull x) \{ return modmul(x, x, n) + i; \};
        while (t++ % 40 || __gcd(prd, n) == 1) { if (x == y) | x = ++i, y = f(x);
                if ((q = modmul(prd, max(x,y) - min(x,y), n))) prd = q;
                x = f(x), y = f(f(y));
        return __gcd(prd, n);
vector<ull> factor(ull n) {
        if (n == 1) return {};
        if (isPrime(n)) return {n};
```

```
ull x = pollard(n);
auto l = factor(x), r = factor(n / x);
l.insert(l.end(), all(r));
return l;
```

#### 10.6 FastEratosthenes

```
* Author: Jakob Kogler, chilli, pajenegod
 * Date: 2020-04-12
 * License: CC0
 * Description: Prime sieve for generating all primes smaller than LIM.
 * Time: LIM=1e9 $\approx$ 1.5s
 * Status: Stress-tested
 * Details: Despite its n log log n complexity, segmented sieve is still faster
 * than other options, including bitset sieves and linear sieves. This is
 * primarily due to its low memory usage, which reduces cache misses. This
 * implementation skips even numbers.
 * Benchmark can be found here: https://ideone.com/e7TbX4
 * The line 'for (int i=idx; i<S+L; idx = (i += p))' is done on purpose for
     performance reasons.
 * Se https://github.com/kth-competitive-programming/kactl/pull/166#
     discussion_r408354338
#pragma once
const int LIM = 1e6;
bitset<LIM> isPrime;
vi eratosthenes() {
        const int S = (int)round(sqrt(LIM)), R = LIM / 2;
        vi pr = \{2\}, sieve(S+1); pr.reserve(int(LIM/log(LIM)*1.1));
        vector<pii> cp;
        for (int i = 3; i \le s; i += 2) if (!sieve[i]) {
                cp.push_back(\{i, i * i / 2\});
                for (int j = i * i; j \le S; j += 2 * i) sieve[j] = 1;
        for (int L = 1; L <= R; L += S) {
                array<bool, S> block{};
                for (auto &[p, idx] : cp)
                        for (int i=idx; i < S+L; idx = (i+=p)) block[i-L] = 1;
                rep(i, 0, min(S, R - L))
                        if (!block[i]) pr.push_back((L + i) * 2 + 1);
        for (int i : pr) isPrime[i] = 1;
        return pr;
```

### 10.7 FracBinarySearch

```
* Author: Lucian Bicsi, Simon Lindholm
 * Date: 2017-10-31
 * License: CC0
 * Description: Given f and f, finds the smallest fraction f in f, in f
 * such that f(p/q) is true, and p, q \le N.
 * You may want to throw an exception from $f$ if it finds an exact solution,
 * in which case $N$ can be removed.
 * Usage: fracBS([](Frac f) { return f.p>=3*f.q; }, 10); // {1,3}
 * Time: O(\log(N))
 * Status: stress-tested for n <= 300
struct Frac { ll p, q; };
template<class F>
Frac fracBS(F f, 11 N) {
        bool dir = 1, A = 1, B = 1;
        Frac 10\{0, 1\}, hi\{1, 1\}; // Set hi to 1/0 to search (0, N]
        if (f(lo)) return lo;
```

#### 10.8 MillerRabin

```
* Author: chilli, c1729, Simon Lindholm
 * Date: 2019-03-28
 * License: CC0
 * Source: Wikipedia, https://miller-rabin.appspot.com/
 * Description: Deterministic Miller-Rabin primality test.
 * Guaranteed to work for numbers up to $7 \cot 10^{-18}; for larger numbers,
     use Python and extend A randomly.
 * Time: 7 times the complexity of $a^b \mod c$.
 * Status: Stress-tested
 */
#pragma once
#include "ModMulLL.h"
bool isPrime(ull n) {
        if (n < 2 \mid | n % 6 % 4 != 1) return (n \mid 1) == 3;
        ull A[] = {2, 325, 9375, 28178, 450775, 9780504, 1795265022},
        s = \_builtin\_ctzll(n-1), d = n >> s; for (ull a : A) { // ^ count trailing zeroes
                 ull p = modpow(a%n, d, n), i = s;
                 while (p != 1 && p != n - 1 && a % n && i--)
                         p = modmul(p, p, n);
                 if (p != n-1 && i != s) return 0;
        return 1;
```

### 10.9 ModInverse

```
/**
  * Author: Simon Lindholm
  * Date: 2016-07-24
  * License: CC0
  * Source: Russian page
  * Description: Pre-computation of modular inverses. Assumes LIM $\le$ mod and that mod is a prime.
  * Status: Works
  */
#pragma once

// const ll mod = 1000000007, LIM = 200000; ///include-line
ll* inv = new ll[LIM] - 1; inv[1] = 1;
rep(i,2,LIM) inv[i] = mod - (mod / i) * inv[mod % i] % mod;
```

### 10.10 ModLog

```
* Author: Bjorn Martinsson
 * Date: 2020-06-03
 * License: CCO
 * Source: own work
 * Description: Returns the smallest x > 0 s.t. a^x = b \mod m, or
 * \$-1\$ if no such \$x\$ exists. modLog(a,1,m) can be used to
 * calculate the order of $a$.
 * Time: $0(\sqrt m)$
 * Status: tested for all 0 <= a,x < 500 and 0 < m < 500.
 * Details: This algorithm uses the baby-step giant-step method to
 * find (i,j) such that a^(n * i) = b * a^j \pmod{m}, where n > sqrt(m)
 * and 0 < i, j <= n. If a and m are coprime then a^{j} has a modular
 * inverse, which means that a^(i * n - j) = b \pmod{m}.
 * However this particular implementation of baby-step giant-step works even
 * without assuming a and m are coprime, using the following idea:
 * Assume p^x is a prime divisor of m. Then we have 3 cases
         1. b is divisible by p^x
         2. b is divisible only by some p^y, 0 < y < x
         3. b is not divisible by p
 * The important thing to note is that in case 2, modLog(a,b,m) (if
 * it exists) cannot be > sqrt(m), (technically it cannot be >= log2(m)).
 * So once all exponenents of a that are <= sqrt(m) has been checked, you
 * cannot have case 2. Case 2 is the only tricky case.
 * So the modification allowing for non-coprime input involves checking all
 * exponents of a that are <= n, and then handling the non-tricky cases by
 * a simple gcd(a^n,m) == gcd(b,m) check.
#pragma once
11 modLog(ll a, ll b, ll m) {
        11 n = (11) sqrt(m) + 1, e = 1, f = 1, j = 1;
        unordered_map<ll, 11> A;
        while (j \le n \&\& (e = f = e * a % m) != b % m)
               A[e * b % m] = j++;
        if (e == b % m) return j;
        if (\underline{gcd}(m, e) == \underline{gcd}(m, b))
                rep(i,2,n+2) if (A.count(e = e * f % m))
                        return n * i - A[e];
        return -1;
```

#### 10.11 ModMulLL

```
* Author: chilli, Ramchandra Apte, Noam527, Simon Lindholm
 * Date: 2019-04-24
 * License: CC0
 * Source: https://github.com/RamchandraApte/OmniTemplate/blob/master/src/
     number_theory/modulo.hpp
 * Description: Calculate $a\cdot b\bmod c$ (or $a^b \bmod c$) for $0 \le a, b \
     le c \le 7.2\cdot 10^{18}$.
 * Time: O(1) for \texttt{modmul}, O(\log b) for \texttt{modpow}
 * Status: stress-tested, proven correct
 * This runs ^{\sim}2x faster than the naive (__int128_t)a * b % M.
 * A proof of correctness is in doc/modmul-proof.tex. An earlier version of the
 * from when the code used a * b / (long double)M, is in doc/modmul-proof.md.
 * The proof assumes that long doubles are implemented as x87 80-bit floats; if
     thev
 * are 64-bit, as on e.g. MSVC, the implementation is only valid for
 * 0 \le a, b \le c < 2^{52} \ge 4.5 \le 10^{15}$.
 */
#pragma once
typedef unsigned long long ull;
ull modmul(ull a, ull b, ull M) {
        ll ret = a * b - M * ull(1.L / M * a * b);
       return ret + M * (ret < 0) - M * (ret >= (11)M);
ull modpow(ull b, ull e, ull mod) {
```

#### 10.12 ModPow

# 10.13 ModSqrt

```
* Author: Simon Lindholm
* Date: 2016-08-31
 * License: CC0
* Source: http://eli.thegreenplace.net/2009/03/07/computing-modular-square-
     roots-in-python/
 * Description: Tonelli-Shanks algorithm for modular square roots. Finds $x$ s.t
      . x^2 = a \neq 0 ($-x$ gives the other solution).
 * Time: O(\log^2 p) worst case, O(\log p) for most $p$
 * Status: Tested for all a,p <= 10000
#pragma once
#include "ModPow.h"
ll sqrt(ll a, ll p) {
        a \% = p; if (a < 0) a += p;
        if (a == 0) return 0;
        assert (modpow(a, (p-1)/2, p) == 1); // else no solution
        if (p % 4 == 3) return modpow(a, (p+1)/4, p);
        // a^(n+3)/8 or 2^(n+3)/8 * 2^(n-1)/4 works if p % 8 == 5
        11 s = p - 1, n = 2;
int r = 0, m;
        while (s % 2 == 0)
                ++r, s /= 2;
        /// find a non-square mod p
        while (modpow(n, (p-1) / 2, p) != p-1) ++n;
        11 x = modpow(a, (s + 1) / 2, p);
        11 b = modpow(a, s, p), g = modpow(n, s, p);
        for (;; r = m)
                11 t = b;
                for (m = 0; m < r \&\& t != 1; ++m)
                       t = t * t % p;
                if (m == 0) return x;
                11 \text{ gs} = \text{modpow}(g, 1LL \ll (r - m - 1), p);
                g = gs * gs % p;
                x = x * gs % p;
                b = b * q % p;
```

#### 10.14 ModSum

```
/**
   * Author: Simon Lindholm
    * Date: 2015-06-23
    * License: CC0
    * Source: own work
    * Description: Sums of mod'ed arithmetic progressions.
   * \texttt{modsum(to, c, k, m)} = \sum_{i=0}^{\min_{i=0}^{\min_{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}^{i=0}
    * \texttt{divsum} is similar but for floored division.
    * Time: $\log(m)$, with a large constant.
    * Status: Tested for all |k|, |c|, to, m \le 50, and on kattis:aladin
    */
 #pragma once
typedef unsigned long long ull;
ull sumsq(ull to) { return to /2 * ((to-1) | 1); }
/// ^ written in a weird way to deal with overflows correctly
ull divsum(ull to, ull c, ull k, ull m) {
                           ull res = k / m * sumsq(to) + c / m * to;
                          k %= m; c %= m;
                           if (!k) return res;
                           ull to2 = (to * k + c) / m;
                           return res + (to - 1) * to2 - divsum(to2, m-1 - c, m, k);
11 modsum(ull to, 11 c, 11 k, 11 m) {
                        C = ((C \% m) + m) \% m;
                          k = ((k % m) + m) % m;
                           return to * c + k * sumsq(to) - m * divsum(to, c, k, m);
```

### 10.15 Modular Arithmetic

```
* Author: Lukas Polacek
 * Date: 2009-09-28
 * License: CC0
 * Source: folklore
 * Description: Operators for modular arithmetic. You need to set {\tt mod} to
 * some number first and then you can use the structure.
#pragma once
#include "euclid.h"
const 11 mod = 17; // change to something else
struct Mod {
        11 x;
         Mod(ll xx) : x(xx) \{ \}
         Mod operator+(Mod b) { return Mod((x + b.x) % mod); }
         Mod operator-(Mod b) { return Mod((x - b.x + mod) % mod); }
        Mod operator*(Mod b) { return Mod((x * b.x) % mod); }
Mod operator/(Mod b) { return *this * invert(b); }
         Mod invert (Mod a) {
                 ll x, y, g = euclid(a.x, mod, x, y);
                 assert(g == 1); return Mod((x + mod) % mod);
         Mod operator (11 e) {
                 if (!e) return Mod(1);
                 Mod r = *this ^ (e / 2); r = r * r;
                 return e&1 ? *this * r : r;
};
```

## 10.16 phiFunction

```
/**
 * Author: Hakan Terelius
```

```
* Date: 2009-09-25
    * License: CCO
    * Source: http://en.wikipedia.org/wiki/Euler's_totient_function
    * Description: \emph{Euler's $\phi$} function is defined as $\phi(n):=\#$ of
                   positive integers $\leq n$ that are coprime with $n$.
    * \hat{p}_i(1)=1, \hat{
                   Rightarrow \phi(mn) = \phi(m) \phi(n) $.
    * If p_1^{k_1}p_2^{k_2}...p_r^{k_r} then \phi(n) = (p_1-1)p_1^{k_1}-1,...(p_r-1)p_r^{k_r}.
    * \phi(n) = n \cdot \phi(n) = n \cdot \phi(n) \cdot (1-1/p) .
    * \sum_{d|n} \phi(d) = n, \sum_{d|n} \phi(d) = n, \sum_{d|n} \phi(d) = n, \sum_{d|n} \phi(d) = n
                   /2, n>1$
    * \textbf{Euler's thm}: $a,n$ coprime $\Rightarrow a^{\phi(n)} \equiv 1 \pmod{n}
    * \textbf{Fermat's little thm}: $p$ prime $\Rightarrow a^{p-1} \equiv 1 \pmod{p}
                   }$ $\forall a$.
    * Status: Tested
 #pragma once
 const int LIM = 5000000;
int phi[LIM];
void calculatePhi() {
                           rep(i,0,LIM) phi[i] = i&1 ? i : i/2;
                           for (int i = 3; i < LIM; i += 2) if (phi[i] == i)
                                                     for (int j = i; j < LIM; j += i) phi[j] -= phi[j] / i;
```

# 11 kactl forked/numerical

## 11.1 BerlekampMassey

```
* Author: Lucian Bicsi
 * Date: 2017-10-31
 * License: CC0
 * Source: Wikipedia
 * Description: Recovers any $n$-order linear recurrence relation from the first
 \star $2n$ terms of the recurrence.
 * Useful for quessing linear recurrences after brute-forcing the first terms.
 * Should work on any field, but numerical stability for floats is not
     guaranteed.
 * Output will have size $\le n$.
 * Usage: berlekampMassey({0, 1, 1, 3, 5, 11}) // {1, 2}
 * Time: O(N^2)
 \star Status: bruteforce-tested mod 5 for n <= 5 and all s
 */
#pragma once
#include "../number-theory/ModPow.h"
vector<ll> berlekampMassey(vector<ll> s) {
        int n = sz(s), L = 0, m = 0;
        vector<ll> C(n), B(n), T;
        C[0] = B[0] = 1;
        11 b = 1;
        rep(i,0,n) \{ ++m;
                11 d = s[i] % mod;
                rep(j,1,L+1) d = (d + C[j] * s[i - j]) % mod;
                if (!d) continue;
                T = C; 11 coef = d * modpow(b, mod-2) % mod;
                rep(j,m,n) C[j] = (C[j] - coef * B[j - m]) % mod;
                if (2 * L > i) continue;
                L = i + 1 - L; B = T; b = d; m = 0;
        C.resize(L + 1); C.erase(C.begin());
        for (11& x : C) x = (mod - x)^{-8} mod;
```

return C;

#### 11.2 Determinant

```
* Author: Simon Lindholm
 * Date: 2016-09-06
 * License: CC0
 * Source: folklore
 * Description: Calculates determinant of a matrix. Destroys the matrix.
 * Time: $0(N^3)$
 * Status: somewhat tested
 */
#pragma once
double det(vector<vector<double>>& a) {
        int n = sz(a); double res = 1;
        rep(i,0,n) {
                int b = i;
                rep(j,i+1,n) if (fabs(a[j][i]) > fabs(a[b][i])) b = j;
                if (i != b) swap(a[i], a[b]), res *= -1;
                res *= a[i][i];
                if (res == 0) return 0;
                rep(j,i+1,n) {
                        double v = a[j][i] / a[i][i];
                        if (v != 0) rep(k,i+1,n) a[j][k] -= v * a[i][k];
        return res;
```

#### 11.3 FastFourierTransform

```
* Author: Ludo Pulles, chilli, Simon Lindholm
 * Date: 2019-01-09
 * License: CC0
 * Source: http://neerc.ifmo.ru/trains/toulouse/2017/fft2.pdf (do read, it's
      excellent)
   Accuracy bound from http://www.daemonology.net/papers/fft.pdf
 * Description: fft(a) computes $\hat f(k) = \sum_x a[x] \exp(2\pi i \cdot k x /
       N)$ for all $k$. N must be a power of 2.
   Useful for convolution:
   \text{texttt}\{\text{conv}(a, b) = c\}, \text{ where } c[x] = \sum_{i=1}^{n} b[x-i] 
   For convolution of complex numbers or more than two vectors: FFT, multiply
   pointwise, divide by n, reverse(start+1, end), FFT back.
   Rounding is safe if (\sum_{i=0}^{4} + \sum_{i=0}^{4} + \sum_{i=0}^{4} ) \log_2{N} < 9 \cdot (14)$
   (in practice $10^{16}$; higher for random inputs).
   Otherwise, use NTT/FFTMod.
 * Time: O(N \log N) with N = |A| + |B| ($\tilde 1s$ for N=2^{22}$)
 * Status: somewhat tested
 * Details: An in-depth examination of precision for both FFT and FFTMod can be
 * here (https://github.com/simonlindholm/fft-precision/blob/master/fft-
      precision.md)
#pragma once
typedef complex<double> C;
typedef vector<double> vd;
void fft(vector<C>& a) {
        int n = sz(a), L = 31 - _builtin_clz(n);
        static vector<complex<long double>> R(2, 1);
        static vector<C> rt(2, 1); // (^ 10% faster if double)
        for (static int k = 2; k < n; k *= 2) {
                R.resize(n); rt.resize(n);
                 auto x = polar(1.0L, acos(-1.0L) / k);
                rep(i,k,2*k) rt[i] = R[i] = i&1 ? R[i/2] * x : R[i/2];
        vi rev(n);
        rep(i,0,n) \ rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
```

```
rep(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);
        for (int k = 1; k < n; k *= 2)
                 for (int i = 0; i < n; i += 2 * k) rep(j,0,k) {
                         // C z = rt[j+k] * a[i+j+k]; // (25% faster if hand-rolled) /// include-line
                          auto x = (double *) &rt[j+k], y = (double *) &a[i+j+k];
                                       /// exclude-line
                          C z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]*y[0]);
                                         /// exclude-line
                          a[i + j + k] = a[i + j] - z;
                          a[i + j] += z;
vd conv(const vd& a, const vd& b) {
        if (a.empty() || b.empty()) return {};
        vd res(sz(a) + sz(b) - 1);
int L = 32 - __builtin_clz(sz(res)), n = 1 << L;
        vector<C> in(n), out(n);
        copy(all(a), begin(in));
        rep(i,0,sz(b)) in[i].imag(b[i]);
        fft(in);
        for (C& x : in) x *= x;
        rep(i,0,n) out[i] = in[-i & (n-1)] - conj(in[i]);
        fft (out);
        rep(i, 0, sz(res)) res[i] = imag(out[i]) / (4 * n);
        return res;
```

### 11.4 FastFourierTransformMod

```
/**
* Author: chilli
 * Date: 2019-04-25
 * License: CC0
 * Source: http://neerc.ifmo.ru/trains/toulouse/2017/fft2.pdf
 * Description: Higher precision FFT, can be used for convolutions modulo
      arbitrary integers
 * as long as N\log_2N\cdot dot \text{ } < 8.6 \cdot 10^{14} (in practice $10
      ^{16}$ or higher).
 * Inputs must be in $[0, \text{mod})$.
 * Time: O(N \log N), where N = |A| + |B| (twice as slow as NTT or FFT)
 * Status: stress-tested
 * Details: An in-depth examination of precision for both FFT and FFTMod can be
 * here (https://github.com/simonlindholm/fft-precision/blob/master/fft-
      precision.md)
#pragma once
#include "FastFourierTransform.h"
typedef vector<ll> vl;
template<int M> vl convMod(const vl &a, const vl &b) {
         if (a.empty() || b.empty()) return {};
        vl res(sz(a) + sz(b) - 1);
        int B=32-_builtin_clz(sz(res)), n=1<<B, cut=int(sqrt(M));
        vector<C> L(n), R(n), outs(n), outl(n);
rep(i,0,sz(a)) L[i] = C((int)a[i] / cut, (int)a[i] % cut);
         rep(i,0,sz(b)) R[i] = C((int)b[i] / cut, (int)b[i] % cut);
         fft(L), fft(R);
         rep(i,0,n) {
                 int j = -i \& (n - 1);
                 outl[j] = (L[i] + conj(L[j])) * R[i] / (2.0 * n);
                 outs[j] = (L[i] - conj(L[j])) * R[i] / (2.0 * n) / 1i;
         fft (outl), fft (outs);
         rep(i,0,sz(res)) {
                 11 \text{ av} = 11(\text{real}(\text{outl}[i]) + .5), \text{ cv} = 11(\text{imag}(\text{outs}[i]) + .5);
                 11 \text{ bv} = 11(\text{imag}(\text{outl}[i]) + .5) + 11(\text{real}(\text{outs}[i]) + .5);
                 res[i] = ((av % M * cut + bv) % M * cut + cv) % M;
         return res;
```

#### 11.5 FastSubsetTransform

```
/**
 * Author: Lucian Bicsi
 * Date: 2015-06-25
 * License: GNU Free Documentation License 1.2
 * Source: csacademy
 * Description: Transform to a basis with fast convolutions of the form
 * $\displaystyle c[z] = \sum\nolimits_{z = x \oplus y} a[x] \cdot b[y]$,
 * where $\oplus$ is one of AND, OR, XOR. The size of $a$ must be a power of two
 * Time: O(N \log N)
 * Status: stress-tested
#pragma once
void FST(vi& a, bool inv) {
        for (int n = sz(a), step = 1; step < n; step *= 2) {
                for (int i = 0; i < n; i += 2 * step) rep(j,i,i+step) {
                        int &u = a[j], &v = a[j + step]; tie(u, v) =
                                inv ? pii(v - u, u) : pii(v, u + v); // AND
                                // inv ? pii(v, u - v) : pii(u + v, u); // OR
                                     /// include-line
                                // pii(u + v, u - v);
                                     /// include-line
        // if (inv) for (int& x : a) x /= sz(a); // XOR only /// include-line
vi conv(vi a, vi b) {
        FST(a, 0); FST(b, 0);
        rep(i, 0, sz(a)) a[i] *= b[i];
        FST(a, 1); return a;
}
```

#### 11.6 GoldenSectionSearch

```
* Author: Ulf Lundstrom
 * Date: 2009-04-17
 * License: CC0
 * Source: Numeriska algoritmer med matlab, Gerd Eriksson, NADA, KTH
 \star Description: Finds the argument minimizing the function $f$ in the interval $
 * assuming $f$ is unimodal on the interval, i.e. has only one local minimum and
      no local
 * maximum. The maximum error in the result is $eps$. Works equally well for
     maximization
 * with a small change in the code. See TernarySearch.h in the Various chapter
     for a
 * discrete version.
 * Usage:
        double func(double x) { return 4+x+.3*x*x; }
        double xmin = gss(-1000, 1000, func);
 * Time: O(\log((b-a) / \epsilon))
 * Status: tested
#pragma once
/// It is important for r to be precise, otherwise we don't necessarily maintain
     the inequality a < x1 < x2 < b.
double gss(double a, double b, double (*f)(double)) {
        double r = (sgrt(5)-1)/2, eps = 1e-7;
        double x1 = b - r*(b-a), x2 = a + r*(b-a);
        double f1 = f(x1), f2 = f(x2);
        while (b-a > eps)
                if (f1 < f2) { //change to > to find maximum
                        b = x2; x2 = x1; f2 = f1;
                        x1 = b - r*(b-a); f1 = f(x1);
                } else {
                        a = x1; x1 = x2; f1 = f2;
                        x2 = a + r*(b-a); f2 = f(x2);
        return a;
```

## 11.7 HillClimbing

```
* Author: Simon Lindholm
* Date: 2015-02-04
* License: CC0
* Source: Johan Sannemo
* Description: Poor man's optimization for unimodal functions.
 * Status: used with great success
#pragma once
typedef array<double, 2> P;
template<class F> pair<double, P> hillClimb(P start, F f) {
        pair<double, P> cur(f(start), start);
        for (double jmp = 1e9; jmp > 1e-20; jmp /= 2) {
                rep(j,0,100) rep(dx,-1,2) rep(dy,-1,2) {
                        P p = cur.second;
                        p[0] += dx * jmp;
                        p[1] += dy * jmp;
                        cur = min(cur, make_pair(f(p), p));
        return cur;
```

### 11.8 IntDeterminant

```
* Author: Unknown
 * Date: 2014-11-27
 \star Source: somewhere on github
 * Description: Calculates determinant using modular arithmetics.
 * Modulos can also be removed to get a pure-integer version.
 * Time: $0(N^3)$
 * Status: bruteforce-tested for N <= 3, mod <= 7
#pragma once
const 11 mod = 12345;
11 det(vector<vector<11>>& a) {
        int n = sz(a); ll ans = 1;
        rep(i,0,n) {
                rep(j,i+1,n) {
                        while (a[j][i] != 0) { // gcd step}
                                ll t = a[i][i] / a[j][i];
                                 if (t) rep(k,i,n)
                                         a[i][k] = (a[i][k] - a[j][k] * t) % mod;
                                 swap(a[i], a[j]);
                                ans *= -1;
                ans = ans * a[i][i] % mod;
                if (!ans) return 0;
        return (ans + mod) % mod;
```

## 11.9 Integrate

```
/**

* Author: Simon Lindholm

* Date: 2015-02-11

* License: CC0

* Source: Wikipedia

* Description: Simple integration of a function over an interval using

* Simpson's rule. The error should be proportional to $h^4$, although in

* practice you will want to verify that the result is stable to desired

* precision when epsilon changes.
```

### 11.10 IntegrateAdaptive

```
* Author: Simon Lindholm
 * Date: 2015-02-11
 * License: CC0
 * Source: Wikipedia
 * Description: Fast integration using an adaptive Simpson's rule.
        double sphereVolume = quad(-1, 1, [](double x) {
        return quad(-1, 1, [\&](double y)
        return quad(-1, 1, [\&] (double z)
        return x*x + y*y + z*z < 1; }); }); }); }
 * Status: mostly untested
#pragma once
typedef double d:
#define S(a,b) (f(a) + 4*f((a+b) / 2) + f(b)) * (b-a) / 6
template <class F>
d rec(F& f, da, db, deps, dS) {
        dc = (a + b) / 2;
        d S1 = S(a, c), S2 = S(c, b), T = S1 + S2;
        if (abs(T - S) \le 15 * eps | | b - a < 1e-10)
                return T + (T - S) / 15;
        return rec(f, a, c, eps / 2, S1) + rec(f, c, b, eps / 2, S2);
template<class F>
d \text{ quad}(d \text{ a, } d \text{ b, } F \text{ f, } d \text{ eps} = 1e-8)  {
        return rec(f, a, b, eps, S(a, b));
```

### 11.11 LinearRecurrence

```
* Author: Lucian Bicsi
 * Date: 2018-02-14
 * License: CC0
 * Source: Chinese material
 * Description: Generates the $k$'th term of an $n$-order
 * linear recurrence $S[i] = \sum_j S[i-j-1]tr[j]$,
 * given S[0 \leq n-1] and tr[0 \leq n-1].
 * Faster than matrix multiplication.
 * Useful together with Berlekamp--Massey.
 * Usage: linearRec({0, 1}, {1, 1}, k) // k'th Fibonacci number
 * Time: O(n^2 \log k)
 * Status: bruteforce-tested mod 5 for n <= 5
#pragma once
const ll mod = 5; /** exclude-line */
typedef vector<ll> Poly;
ll linearRec(Poly S, Poly tr, ll k) {
       int n = sz(tr);
        auto combine = [&] (Poly a, Poly b) {
               Poly res(n \star 2 + 1);
                rep(i, 0, n+1) rep(j, 0, n+1)
```

#### 11.12 MatrixInverse-mod

```
* Author: Simon Lindholm
* Date: 2016-12-08
* Source: The regular matrix inverse code
 * Description: Invert matrix $A$ modulo a prime.
* Returns rank; result is stored in $A$ unless singular (rank < n).
 * For prime powers, repeatedly set A^{-1} = A^{-1} (2I - AA^{-1})\ (\text{mod}
      p^k where A^{-1} starts as
 * the inverse of A mod p, and k is doubled in each step.
* Time: O(n^3)
* Status: Slightly tested
#pragma once
#include "../number-theory/ModPow.h"
int matInv(vector<vector<ll>>& A) {
        int n = sz(A); vi col(n);
        vector<vector<ll>> tmp(n, vector<ll>(n));
        rep(i,0,n) tmp[i][i] = 1, col[i] = i;
        rep(i,0,n) {
                int r = i, c = i;
                rep(j,i,n) rep(k,i,n) if (A[j][k]) {
                       r = j; c = k; goto found;
                return i;
found:
                A[i].swap(A[r]); tmp[i].swap(tmp[r]);
                rep(j,0,n)
                        swap(A[j][i], A[j][c]), swap(tmp[j][i], tmp[j][c]);
                swap(col[i], col[c]);
                11 v = modpow(A[i][i], mod - 2);
                rep(j,i+1,n) {
                        11 f = A[j][i] * v % mod;
                        A[j][i] = 0;
                        rep(k, i+1, n) A[j][k] = (A[j][k] - f*A[i][k]) % mod;
                        rep(k, 0, n) tmp[j][k] = (tmp[j][k] - f*tmp[i][k]) % mod;
                rep(j,i+1,n) A[i][j] = A[i][j] * v % mod;
                rep(j, 0, n) tmp[i][j] = tmp[i][j] * v % mod;
                A[i][i] = 1;
        for (int i = n-1; i > 0; --i) rep(j,0,i) {
                11 v = A[j][i];
                rep(k, 0, n) tmp[j][k] = (tmp[j][k] - v*tmp[i][k]) % mod;
        rep(i,0,n) rep(j,0,n)
                A[col[i]][col[j]] = tmp[i][j] % mod + (tmp[i][j] < 0) *mod;
        return n;
```

#### 11.13 MatrixInverse

```
/**
 * Author: Max Bennedich
 * Date: 2004-02-08
 * Description: Invert matrix $A$. Returns rank; result is stored in $A$ unless
     singular (rank < n).
 \star Can easily be extended to prime moduli; for prime powers, repeatedly
 * set A^{-1} = A^{-1} (2I - AA^{-1}) (\text{text} \text{ mod } p^k) where A^{-1} starts
 * the inverse of A mod p, and k is doubled in each step.
 * Time: O(n^3)
 * Status: Slightly tested
#pragma once
int matInv(vector<vector<double>>& A) {
        int n = sz(A); vi col(n);
        vector<vector<double>> tmp(n, vector<double>(n));
        rep(i,0,n) tmp[i][i] = 1, col[i] = i;
        rep(i,0,n) {
                 int r = i, c = i;
                 rep(j,i,n) rep(k,i,n)
                         if (fabs(A[j][k]) > fabs(A[r][c]))
                r = j, c = k;
if (fabs(A[r][c]) < 1e-12) return i;
                 A[i].swap(A[r]); tmp[i].swap(tmp[r]);
                rep(j,0,n)
                         swap(A[j][i], A[j][c]), swap(tmp[j][i], tmp[j][c]);
                 swap(col[i], col[c]);
double v = A[i][i];
                 rep(j,i+1,n) {
                         double f = A[j][i] / v;
                         A[j][i] = 0;
                         rep(k, i+1, n) A[j][k] -= f*A[i][k];
                         rep(k,0,n) tmp[j][k] -= f*tmp[i][k];
                 rep(j,i+1,n) A[i][j] /= v;
                 rep(j,0,n) tmp[i][j] /= v;
                A[i][i] = 1;
        /// forget A at this point, just eliminate tmp backward
        for (int i = n-1; i > 0; --i) rep(j,0,i) {
                double v = A[j][i];
                 rep(k,0,n) tmp[j][k] -= v*tmp[i][k];
        rep(i,0,n) rep(j,0,n) A[col[i]][col[j]] = tmp[i][j];
        return n;
```

### 11.14 NumberTheoreticTransform

```
/**
* Author: chilli
 * Date: 2019-04-16
 * License: CC0
 * Source: based on KACTL's FFT
 * Description: ntt(a) computes \hat{s} = \sum_{x \in \mathbb{Z}} q^{xk} for all \hat{s},
      where g=\text{text{root}}^{\hat{}} { (mod-1)/N}$.
 * N must be a power of 2.
 * Useful for convolution modulo specific nice primes of the form $2^a b+1$,
 * where the convolution result has size at most 2^a. For arbitrary modulo,
   \text{texttt}\{\text{conv}(a, b) = c\}, \text{ where } c[x] = \sum_{i=1}^{n} b[x-i].
   For manual convolution: NTT the inputs, multiply
   pointwise, divide by n, reverse(start+1, end), NTT back.
 * Inputs must be in [0, mod).
 * Time: O(N \log N)
 * Status: stress-tested
```

```
#pragma once
#include "../number-theory/ModPow.h"
const 11 mod = (119 \ll 23) + 1, root = 62; // = 998244353
// For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 << 21
// and 483 << 21 (same root). The last two are > 10^9.
typedef vector<ll> vl;
void ntt(vl &a) {
        int n = sz(a), L = 31 - _builtin_clz(n);
        static vl rt(2, 1);
        for (static int k = 2, s = 2; k < n; k *= 2, s++) {
                rt.resize(n);
                ll z[] = \{1, modpow(root, mod >> s)\};
                rep(i,k,2*k) rt[i] = rt[i / 2] * z[i & 1] % mod;
        vi rev(n);
        rep(i,0,n) \ rev[i] = (rev[i / 2] | (i \& 1) << L) / 2;
        rep(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
        for (int k = 1; k < n; k *= 2)
                for (int i = 0; i < n; i += 2 * k) rep(j,0,k) {
                        ll z = rt[j + k] * a[i + j + k] % mod, &ai = a[i + j];
                        a[i + j + k] = ai - z + (z > ai ? mod : 0);
                        ai += (ai + z >= mod ? z - mod : z);
vl conv(const vl &a, const vl &b) {
        if (a.empty() || b.empty()) return {};
        int s = sz(a) + sz(b) - 1, B = 32 - _builtin_clz(s),
            n = 1 << B;
        int inv = modpow(n, mod - 2);
        v1 L(a), R(b), out(n);
        L.resize(n), R.resize(n);
        ntt(L), ntt(R);
        rep(i,0,n)
                out [-i \& (n - 1)] = (l1)L[i] * R[i] % mod * inv % mod;
        ntt(out);
        return {out.begin(), out.begin() + s};
```

## 11.15 PolyInterpolate

```
* Author: Simon Lindholm
 * Date: 2017-05-10
 * License: CC0
 * Source: Wikipedia
 * Description: Given $n$ points (x[i], y[i]), computes an n-1-degree polynomial
    passes through them: p(x) = a[0] * x^0 + ... + a[n-1] * x^{n-1}.
 * For numerical precision, pick x[k] = c \cdot \cos(k/(n-1) \cdot pi), k=0 \cdot n-1.
 * Time: O(n^2)
#pragma once
typedef vector<double> vd;
vd interpolate(vd x, vd y, int n) {
        vd res(n), temp(n);
        rep(k,0,n-1) rep(i,k+1,n)
        y[i] = (y[i] - y[k]) / (x[i] - x[k]);
double last = 0; temp[0] = 1;
        rep(k,0,n) rep(i,0,n) {
                 res[i] += y[k] * temp[i];
                 swap(last, temp[i]);
                 temp[i] -= last * x[k];
        return res;
```

## 11.16 Polynomial

```
* Author: David Rydh, Per Austrin
 * Date: 2003-03-16
 * Description:
#pragma once
struct Poly {
        vector<double> a;
        double operator()(double x) const {
                double val = 0;
                for (int i = sz(a); i--;) (val *= x) += a[i];
                return val;
        void diff() {
                rep(i,1,sz(a)) a[i-1] = i*a[i];
                a.pop_back();
        void divroot(double x0) {
                double b = a.back(), c; a.back() = 0;
                for(int i=sz(a)-1; i--;) c = a[i], a[i] = a[i+1]*x0+b, b=c;
                a.pop_back();
};
```

## 11.17 PolyRoots

```
* Author: Per Austrin
 * Date: 2004-02-08
 * License: CCO
 * Description: Finds the real roots to a polynomial.
 * Usage: polyRoots(\{\{2,-3,1\}\},-1e9,1e9) // solve x^2-3x+2=0
 * Time: O(n^2 \log(1/\epsilon))
#pragma once
#include "Polvnomial.h"
vector<double> polyRoots(Poly p, double xmin, double xmax) {
        if (sz(p.a) == 2) \{ return \{-p.a[0]/p.a[1]\}; \}
        vector<double> ret;
        Poly der = p;
        der.diff();
        auto dr = polyRoots(der, xmin, xmax);
        dr.push_back(xmin-1);
        dr.push_back(xmax+1);
        sort(all(dr));
        rep(i, 0, sz(dr) - 1) {
                double l = dr[i], h = dr[i+1];
                bool sign = p(1) > 0;
                if (sign ^ (p(h) > 0))
                        rep(it, 0, 60) { // while (h - 1 > 1e-8)
                                 double m = (1 + h) / 2, f = p(m);
                                 if ((f \le 0) \hat{sign}) l = m;
                                 else h = m;
                        ret.push_back((1 + h) / 2);
        return ret:
```

## 11.18 Simplex

```
* Returns -inf if there is no solution, inf if there are arbitrarily good solutions, or the maximum value of c^T x otherwise.
 \star The input vector is set to an optimal $x$ (or in the unbounded case, an
      arbitrary solution fulfilling the constraints).
 * Numerical stability is not guaranteed. For better performance, define
      variables such that x = 0 is viable.
 * Usage:
* vvd A = \{\{1,-1\}, \{-1,1\}, \{-1,-2\}\};
* vd b = \{1,1,-4\}, c = \{-1,-1\}, x;
 * T val = LPSolver(A, b, c).solve(x);
 \star Time: O(NM \star \#pivots), where a pivot may be e.g. an edge relaxation. O(2^n)
      in the general case.
 * Status: seems to work?
#pragma once
typedef double T; // long double, Rational, double + mod<P>...
typedef vector<T> vd;
typedef vector<vd> vvd;
const T eps = 1e-8, inf = 1/.0;
#define MP make_pair
#define ltj(X) if(s == -1 \mid \mid MP(X[j], N[j]) < MP(X[s], N[s])) s=j
struct LPSolver {
        int m. n:
        vi N, B;
        vvd D;
        LPSolver(const vvd& A, const vd& b, const vd& c) :
                 m(sz(b)), n(sz(c)), N(n+1), B(m), D(m+2), vd(n+2)) {
                          rep(i, 0, m) rep(j, 0, n) D[i][j] = A[i][j];
                          rep(i, 0, m) { B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[i]
                          rep(j, 0, n) \{ N[j] = j; D[m][j] = -c[j]; \}
                          N[n] = -1; D[m+1][n] = 1;
        void pivot(int r, int s)
                 T *a = D[r].data(), inv = 1 / a[s];
rep(i,0,m+2) if (i != r && abs(D[i][s]) > eps) {
                          T *b = D[i].data(), inv2 = b[s] * inv;
                          rep(j, 0, n+2) b[j] -= a[j] * inv2;
                          b[s] = a[s] * inv2;
                 rep(j, 0, n+2) if (j != s) D[r][j] *= inv;
                 rep(i, 0, m+2) if (i != r) D[i][s] *= -inv;
                 D[r][s] = inv;
                 swap(B[r], N[s]);
        bool simplex(int phase) {
                 int x = m + phase - 1;
                 for (;;) {
                          int s = -1;
                          rep(j,0,n+1) if (N[j] != -phase) ltj(D[x]);
                          if (D[x][s] >= -eps) return true;
                          int r = -1;
                          rep(i,0,m) {
                                   if (D[i][s] <= eps) continue;
                                   if (r == -1 \mid | MP(D[i][n+1] / D[i][s], B[i])
                                                 < MP(D[r][n+1] / D[r][s], B[r])) r
                                                       = i;
                          if (r == -1) return false;
                          pivot(r, s);
        T solve(vd &x) {
                 int r = 0;
                 rep(i,1,m) if (D[i][n+1] < D[r][n+1]) r = i;
                 if^{-}(D[r][n+1] < -eps) {
                          pivot(r, n);
                          if (!simplex(2) \mid \mid D[m+1][n+1] < -eps) return -inf;
                          rep(i, 0, m) if (B[i] == -1) {
                                   int s = 0;
                                   rep(j,1,n+1) ltj(D[i]);
```

```
pivot(i, s);
}
bool ok = simplex(1); x = vd(n);
rep(i,0,m) if (B[i] < n) x[B[i]] = D[i][n+1];
return ok ? D[m][n+1] : inf;
}
};</pre>
```

#### 11.19 SolveLinear

```
/**
* Author: Per Austrin, Simon Lindholm
 * Date: 2004-02-08
 * License: CC0
 * Description: Solves A \times x = b. If there are multiple solutions, an
      arbitrary one is returned.
 * Returns rank, or -1 if no solutions. Data in $A$ and $b$ is lost.
 * Time: O(n^2 m)
 * Status: tested on kattis:equationsolver, and bruteforce-tested mod 3 and 5
      for n,m \le 3
#pragma once
typedef vector<double> vd;
const double eps = 1e-12;
int solveLinear(vector<vd>& A, vd& b, vd& x) {
        int n = sz(A), m = sz(x), rank = 0, br, bc;
        if (n) assert(sz(A[0]) == m);
        vi col(m); iota(all(col), 0);
        rep(i,0,n) {
                double v, bv = 0;
                rep(r,i,n) rep(c,i,m)
                         if ((v = fabs(A[r][c])) > bv)
                                 br = r, bc = c, bv = v;
                if (bv <= eps) {
                         rep(j,i,n) if (fabs(b[j]) > eps) return -1;
                         break:
                swap(A[i], A[br]);
swap(b[i], b[br]);
                swap(col[i], col[bc]);
                rep(j,0,n) swap(A[j][i], A[j][bc]);
                bv = 1/A[i][i];
                rep(j,i+1,n) {
                         double fac = A[j][i] * bv;
                         b[j] -= fac * b[i];
                         rep(k,i+1,m) A[j][k] -= fac*A[i][k];
                rank++;
        x.assign(m, 0);
        for (int i = rank; i--;) {
                b[i] /= A[i][i];
                x[col[i]] = b[i];
                rep(j, 0, i) b[j] -= A[j][i] * b[i];
        return rank; // (multiple solutions if rank < m)
```

### 11.20 SolveLinear2

```
* Status: tested on kattis:equationsolverplus, stress-tested
*/
#pragma once
#include "SolveLinear.h"

rep(j,0,n) if (j != i) // instead of rep(j,i+1,n)
// ... then at the end:
x.assign(m, undefined);
rep(i,0,rank) {
    rep(j,rank,m) if (fabs(A[i][j]) > eps) goto fail;
    x[col[i]] = b[i] / A[i][i];
fail:; }
```

## 11.21 SolveLinearBinary

```
* Author: Simon Lindholm
 * Date: 2016-08-27
 * License: CC0
 * Source: own work
 * Description: Solves $Ax = b$ over $\mathbb F_2$. If there are multiple
      solutions, one is returned arbitrarily.
 * Returns rank, or -1 if no solutions. Destroys $A$ and $b$.
 * Time: O(n^2 m)
 \star Status: bruteforce-tested for n, m <= 4
#pragma once
typedef bitset<1000> bs;
int solveLinear(vector<bs>& A, vi& b, bs& x, int m) {
        int n = sz(A), rank = 0, br;
        assert(m \le sz(x));
        vi col(m); iota(all(col), 0);
        rep(i,0,n) {
                 for (br=i; br<n; ++br) if (A[br].any()) break;
                 if (br == n) {
                         rep(j,i,n) if(b[j]) return -1;
                         break:
                 int bc = (int)A[br]._Find_next(i-1);
                 swap(A[i], A[br]);
                 swap(b[i], b[br]);
                 swap(col[i], col[bc]);
                 rep(j,0,n) if (A[j][i] != A[j][bc]) {
                         A[j].flip(i); A[j].flip(bc);
                rep(j,i+1,n) if (A[j][i]) {
    b[j] ^= b[i];
    A[j] ^= A[i];
                 rank++;
        x = bs();
        for (int i = rank; i--;) {
                 if (!b[i]) continue;
                x[col[i]] = 1;
                 rep(j,0,i) b[j] ^= A[j][i];
        return rank; // (multiple solutions if rank < m)
```

## 11.22 Tridiagonal

```
/**
  * Author: Ulf Lundstrom, Simon Lindholm
  * Date: 2009-08-15
  * License: CCO
  * Source: https://en.wikipedia.org/wiki/Tridiagonal_matrix_algorithm
  * Description: $x=\textrm{tridiagonal}(d,p,q,b)$ solves the equation system
```

```
\left(\frac{3}{c}b_0\right)_1\left(\frac{3}{c}b_1\right)_2\left(\frac{3}{c}b_1\right)_1
\left(\begin{array}{cccccc}
d_0 & p_0 & 0 & 0 & \cdots & 0\\
q_0 & d_1 & p_1 & 0 & \cdots & 0\\
0 & q_1 & d_2 & p_2 & \cdots & 0\\
\vdots & \vdots & \ddots & \ddots & \vdots\\
0 & 0 & \cdots & q_{n-3} & d_{n-2} & p_{n-2}\
0 & 0 & \cdots & 0 & q_{n-2} & d_{n-1}\\
\end{array}\right)
\left(\left(\frac{n-1}{\alpha ray}\right)^{c}x_0\right)^{-1}\right).
This is useful for solving problems on the type
[a_i=b_ia_{i-1}+c_ia_{i+1}+d_i, ,1\leq i\leq n, ]
where a_0, a_{n+1}, b_i, c_i and d_i are known. a can then be
    obtained from
\begin{align*}
\{a_i^i\}=\text{textrm}\{\text{tridiagonal}\} (\&\{1,-1,-1,...,-1,1\}, \\{0,c_1,c_2,\dots,c_n\},\
\{b_1, b_2, dots, b_n, 0\}, \{a_0, d_1, d_2, dots, d_n, a_{n+1}\}\}
\end{align*}
Fails if the solution is not unique.
If |d_i| > |p_i| + |q_{i-1}| for all |s_i|, or |d_i| > |p_{i-1}| + |q_i|, or
    the matrix is positive definite,
the algorithm is numerically stable and neither \texttt{tr} nor the check for \
    texttt{diag[i] == 0} is needed.
 * Status: Brute-force tested mod 5 and 7 and stress-tested for real matrices
      obeying the criteria above.
#pragma once
typedef double T;
vector<T> tridiagonal(vector<T> diag, const vector<T>& super,
                const vector<T>& sub, vector<T> b) {
        int n = sz(b); vi tr(n);
        rep(i, 0, n-1) {
                if (abs(diag[i]) < 1e-9 * abs(super[i])) { // diag[i] == 0}
                        b[i+1] = b[i] * diag[i+1] / super[i];
                        if (i+2 < n) b[i+2] = b[i] * sub[i+1] / super[i];
                        diag[i+1] = sub[i]; tr[++i] = 1;
                } else {
                        diag[i+1] -= super[i]*sub[i]/diag[i];
                        b[i+1] = b[i] * sub[i] / diag[i];
        for (int i = n; i--;) {
                if (tr[i]) {
                        swap(b[i], b[i-1]);
diag[i-1] = diag[i];
                        b[i] /= super[i-1];
                } else {
                        b[i] /= diag[i];
                        if (i) b[i-1] -= b[i] *super[i-1];
        return b;
```

# 12 kactl forked/strings

there, or -1 if none.

### 12.1 AhoCorasick

```
/**

* Author: Simon Lindholm

* Date: 2015-02-18

* License: CCO

* Source: marian's (TC) code

* Description: Aho-Corasick automaton, used for multiple pattern matching.

* Initialize with AhoCorasick ac(patterns); the automaton start node will be at index 0.

* find(word) returns for each position the index of the longest word that ends
```

```
* findAll(\$-\$, word) finds all words (up to \$N \rightarrow \$ many if no duplicate
      patterns)
 * that start at each position (shortest first).
 * Duplicate patterns are allowed; empty patterns are not.
 * To find the longest words that start at each position, reverse all input.
 * For large alphabets, split each symbol into chunks, with sentinel bits for
      symbol boundaries.
 * Time: construction takes $O(26N)$, where $N = $ sum of length of patterns.
 * find(x) is O(N), where N = length of x. findAll is O(NM).
 * Status: stress-tested
#pragma once
struct AhoCorasick {
        enum {alpha = 26, first = 'A'}; // change this!
        struct Node {
                // (nmatches is optional)
                int back, next[alpha], start = -1, end = -1, nmatches = 0;
                Node(int v) { memset(next, v, sizeof(next)); }
        vector<Node> N;
        vi backp;
        void insert(string& s, int j) {
                assert(!s.empty());
                int n = 0;
                for (char c : s) {
                        int& m = N[n].next[c - first];
                        if (m == -1) { n = m = sz(N); N.emplace_back(-1); }
                        else n = m;
                if (N[n].end == -1) N[n].start = j;
                backp.push_back(N[n].end);
                N[n].end = j;
                N[n].nmatches++;
        AhoCorasick(vector<string>& pat) : N(1, -1) {
                rep(i, 0, sz(pat)) insert(pat[i], i);
                N[0].back = sz(N);
                N.emplace_back(0);
                queue<int> q;
                for (q.push(0); !q.empty(); q.pop()) {
                        int n = q.front(), prev = N[n].back;
                        rep(i,0,alpha) {
                                int &ed = N[n].next[i], y = N[prev].next[i];
                                if (ed == -1) ed = y;
                                else {
                                         N[ed].back = y;
                                         (N[ed].end == -1 ? N[ed].end : backp[N[
                                             ed].start])
                                                 = N[y].end;
                                         N[ed].nmatches += N[y].nmatches;
                                         q.push(ed);
        vi find(string word) {
                int n = 0;
                vi res; // 11 count = 0;
                for (char c : word) {
                        n = N[n].next[c - first];
                        res.push_back(N[n].end);
                        // count += N[n].nmatches;
                return res:
        vector<vi> findAll(vector<string>& pat, string word) {
                vi r = find(word);
                vector<vi> res(sz(word));
                rep(i,0,sz(word)) {
                        int ind = r[i];
                        while (ind !=-1)
                                res[i - sz(pat[ind]) + 1].push_back(ind);
                                ind = backp[ind];
                return res;
```

};

## 12.2 Hashing-codeforces

```
* Author: Simon Lindholm
 * Date: 2015-03-15
 * License: CCO
 * Source: own work
 * Description: Various self-explanatory methods for string hashing.
 * Use on Codeforces, which lacks 64-bit support and where solutions can be
 * Status: stress-tested
#pragma once
typedef uint64_t ull;
static int C; // initialized below
// Arithmetic mod two primes and 2^32 simultaneously.
// "typedef uint64_t H;" instead if Thue-Morse does not apply.
template<int M, class B>
struct A {
        int x; B b; A(int x=0) : x(x), b(x) {}
        A(int x, B b) : x(x), b(b) {}
        A operator+(A o) {int y = x+o.x; return{y - (y>=M)*M, b+o.b};}
        A operator-(A o) {int y = x-o.x; return{y + (y< 0)*M, b-o.b};}
        A operator*(A o) { return {(int)(1LL*x*o.x % M), b*o.b}; }
        explicit operator ull() { return x ^ (ull) b << 21; }
        bool operator==(A o) const { return (ull) *this == (ull) o; }
        bool operator<(A o) const { return (ull)*this < (ull)o; }</pre>
typedef A<1000000007, A<1000000009, unsigned>> H;
struct HashInterval {
        vector<H> ha, pw;
        HashInterval(string& str) : ha(sz(str)+1), pw(ha) {
                pw[0] = 1;
                rep(i,0,sz(str))
                        ha[i+1] = ha[i] * C + str[i],
                        pw[i+1] = pw[i] * C;
        H hashInterval(int a, int b) { // hash [a, b)
                return ha[b] - ha[a] * pw[b - a];
};
vector<H> getHashes(string& str, int length) {
        if (sz(str) < length) return {};
        H h = 0, pw = 1;
        rep(i,0,length)
                h = h * C + str[i], pw = pw * C;
        vector<H> ret = {h};
        rep(i,length,sz(str))
                ret.push_back(h = h * C + str[i] - pw * str[i-length]);
        return ret:
H hashString(string& s){H h{}; for(char c:s) h=h*C+c; return h;}
#include <sys/time.h>
int main() {
        timeval tp;
        gettimeofday(&tp, 0);
        C = (int)tp.tv_usec; // (less than modulo)
        assert((ull)(H(1)*2+1-3) == 0);
        // ...
```

## 12.3 Hashing

```
* Author: Simon Lindholm
 * Date: 2015-03-15
* License: CC0
 * Source: own work
 * Description: Self-explanatory methods for string hashing.
 * Status: stress-tested
 */
#pragma once
// Arithmetic mod 2^64-1. 2x slower than mod 2^64 and more
// code, but works on evil test data (e.g. Thue-Morse, where
// ABBA... and BAAB... of length 2^10 hash the same mod 2^64).
// "typedef ull H;" instead if you think test data is random,
// or work mod 10^9+7 if the Birthday paradox is not a problem.
typedef uint64_t ull;
struct H {
        ull x; H(ull x=0) : x(x) {}
        H operator+(H o) { return x + o.x + (x + o.x < x); }
        H operator-(H o) { return *this + ~o.x; }
        H operator*(H o) { auto m = (\underline{uint128\_t})x * o.x;
                return H((ull)m) + (ull)(m >> 64); }
        ull get() const { return x + !~x; }
        bool operator==(H o) const { return get() == o.get(); }
        bool operator<(H o) const { return get() < o.get(); }</pre>
};
static const H C = (11)1e11+3; // (order ~ 3e9; random also ok)
struct HashInterval
        vector<H> ha, pw;
        HashInterval(string& str) : ha(sz(str)+1), pw(ha) {
                pw[0] = 1;
                rep(i, 0, sz(str))
                        ha[i+1] = ha[i] * C + str[i],
                        pw[i+1] = pw[i] * C;
        H hashInterval(int a, int b) { // hash [a, b)
                return ha[b] - ha[a] * pw[b - a];
};
vector<H> getHashes(string& str, int length) {
        if (sz(str) < length) return {};
        H h = 0, pw = 1;
        rep(i,0,length)
                h = h * C + str[i], pw = pw * C;
        vector<H> ret = {h};
        rep(i,length,sz(str)) {
                ret.push_back(h = h * C + str[i] - pw * str[i-length]);
        return ret;
H hashString(string& s) {H h{}; for(char c:s) h=h*C+c; return h;}
```

### 12.4 KMP

```
* Author: Johan Sannemo
* Date: 2016-12-15
* License: CC0
* Description: pi[x] computes the length of the longest prefix of s that ends
 * other than s[0...x] itself (abacaba -> 0010123).
 * Can be used to find all occurrences of a string.
 * Time: O(n)
 * Status: Tested on kattis:stringmatching
*/
#pragma once
vi pi(const string& s) {
        vi p(sz(s));
        rep(i,1,sz(s)) {
                int q = p[i-1];
                while (g \& \& s[i] != s[g]) g = p[g-1];
                p[i] = g + (s[i] == s[g]);
```

### 12.5 Manacher

```
* Author: User adamant on CodeForces
 * Source: http://codeforces.com/blog/entry/12143
 * Description: For each position in a string, computes p[0][i] = half length of
 * longest even palindrome around pos i, p[1][i] = longest odd (half rounded
      down).
 * Time: O(N)
 * Status: Stress-tested
#pragma once
array<vi, 2> manacher(const string& s) {
        int n = sz(s);
        array < vi, 2 > p = {vi(n+1), vi(n)};
        rep(z, 0, 2) for (int i=0, l=0, r=0; i < n; i++) {
                 int t = r-i+!z;
                 if (i<r) p[z][i] = min(t, p[z][l+t]);
int L = i-p[z][i], R = i+p[z][i]-!z;</pre>
                 while (L>=1 \&\& R+1< n \&\& s[L-1] == s[R+1])
                          p[z][i]++, L--, R++;
                 if (R>r) l=L, r=R;
        return p;
```

### 12.6 MinRotation

```
/**
  * Author: Stjepan Glavina
  * License: Unlicense
  * Source: https://github.com/stjepang/snippets/blob/master/min_rotation.cpp
  * Description: Finds the lexicographically smallest rotation of a string.
  * Time: O(N)
  * Usage:
  * rotate(v.begin(), v.begin()+minRotation(v), v.end());
  * Status: Stress-tested
  */
#pragma once

int minRotation(string s) {
    int a=0, N=sz(s); s += s;
    rep(b,0,N) rep(k,0,N) {
        if (a+k == b || s[a+k] < s[b+k]) {b += max(0, k-1); break;}
        if (s[a+k] > s[b+k]) {a = b; break;}
    }
    return a;
}
```

## 12.7 SuffixArray

```
/**
 * Author: Luo Sui Qian , chilli
 * Date: 2019-04-11
 * License: Unknown
 * Source: Suffix array - a powerful tool for dealing with strings
 * (Chinese IOI National team training paper, 2009)
```

```
* Description: Builds suffix array for a string.
 * \texttt{sa[i]} is the starting index of the suffix which
 * is $i$'th in the sorted suffix array.
 * The returned vector is of size n+1, and \text{texttt}\{sa[0] = n\}.
 * The \texttt{lcp} array contains longest common prefixes for
 * neighbouring strings in the suffix array:
 * \text{texttt}\{lcp[i] = lcp(sa[i], sa[i-1])\}, \text{texttt}\{lcp[0] = 0\}.
 * The input string must not contain any nul chars.
 * Time: O(n \log n)
 * Status: stress-tested
#pragma once
struct SuffixArray {
        vi sa, lcp;
        SuffixArray(string s, int lim=256) { // or vector<int>
                 s.push\_back(0); int n = sz(s), k = 0, a, b;
                 vi x(all(s)), y(n), ws(max(n, lim));
                 sa = lcp = y, iota(all(sa), 0);
                 for (int j = 0, p = 0; p < n; j = max(1, j * 2), lim = p) {
                         p = j, iota(all(y), n - j);
                         rep(i,0,n) if (sa[i] >= j) y[p++] = sa[i] - j;
                         fill(all(ws), 0);
                         rep(i, 0, n) ws[x[i]] ++;
                         rep(i, 1, lim) ws[i] += ws[i - 1];
                         for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
                         swap(x, y), p = 1, x[sa[0]] = 0;

rep(i,1,n) = a = sa[i-1], b = sa[i], x[b] =
                                  (y[a] == y[b] && y[a + j] == y[b + j]) ? p - 1 :
                 for (int i = 0, j; i < n - 1; lcp[x[i++]] = k)
                         for (k \& \& k--, j = sa[x[i] - 1];
                                          s[i + k] == s[j + k]; k++);
};
```

### 12.8 SuffixTree

```
* Author: Unknown
* Date: 2017-05-15
 * Source: https://e-maxx.ru/algo/ukkonen
 * Description: Ukkonen's algorithm for online suffix tree construction.
 * Each node contains indices [1, r) into the string, and a list of child nodes
 \star Suffixes are given by traversals of this tree, joining [1, r) substrings.
   The root is 0 (has 1 = -1, r = 0), non-existent children are -1.
 * To get a complete tree, append a dummy symbol -- otherwise it may contain
 * an incomplete path (still useful for substring matching, though).
 * Time: $0(26N)$
 * Status: stress-tested a bit
#pragma once
struct SuffixTree {
        enum { N = 200010, ALPHA = 26 }; // N \sim 2*maxlen+10
        int toi(char c) { return c - 'a'; }
        string a; // v = cur node, q = cur position
        int t[N][ALPHA], 1[N], r[N], p[N], s[N], v=0, q=0, m=2;
        void ukkadd(int i, int c) { suff:
                if (r[v] \leq q)
                        if (t[v][c]==-1) { t[v][c]=m; l[m]=i;
                        p[m++]=v; v=s[v]; q=r[v]; goto suff; }
v=t[v][c]; q=l[v];
                if (q==-1 || c==toi(a[q])) q++; else {
                        l[m+1]=i; p[m+1]=m; l[m]=l[v]; r[m]=q;
                        p[m]=p[v]; t[m][c]=m+1; t[m][toi(a[q])]=v;
                        l[v]=q; p[v]=m; t[p[m]][toi(a[l[m]])]=m;
                        v=s[p[m]]; q=l[m];
                        while (q < r[m]) \{ v = t[v][toi(a[q])]; q + = r[v] - l[v]; \}
                        if (q==r[m]) s[m]=v; else s[m]=m+2;
                        q=r[v]-(q-r[m]); m+=2; goto suff;
```

```
SuffixTree(string a) : a(a) {
        fill(r,r+N,sz(a));
        memset(s, 0, sizeof s);
        memset(t, -1, size of t);
       fill(t[1],t[1]+ALPHA,0);
s[0] = 1; 1[0] = 1[1] = -1; r[0] = r[1] = p[0] = p[1] = 0;
        rep(i,0,sz(a)) ukkadd(i, toi(a[i]));
// example: find longest common substring (uses ALPHA = 28)
int lcs(int node, int i1, int i2, int olen) {
        if (1[node] <= i1 && i1 < r[node]) return 1;
        if (1[node] <= i2 && i2 < r[node]) return 2;
        int mask = 0, len = node ? olen + (r[node] - 1[node]) : 0;
        rep(c, 0, ALPHA) if (t[node][c] != -1)
                mask |= lcs(t[node][c], i1, i2, len);
        if (mask == 3)
                best = max(best, {len, r[node] - len});
        return mask;
static pii LCS(string s, string t) {
        SuffixTree st(s + (char)('z' + 1) + t + (char)('z' + 2));
        st.lcs(0, sz(s), sz(s) + 1 + sz(t), 0);
        return st.best;
```

#### 12.9 Zfunc

};

```
* Author: chilli
 * License: CC0
 \star Description: z[i] computes the length of the longest common prefix of s[i:]
 \star \text{ except } z[0] = 0. \text{ (abacaba -> 0010301)}
 * Time: O(n)
 * Status: stress-tested
 */
#pragma once
vi Z(const string& S) {
        vi z(sz(S));
        int 1 = -1, r = -1;
                 z[i] = i >= r ? 0 : min(r - i, z[i - 1]);
                 while (i + z[i] < sz(S) \&\& S[i + z[i]] == S[z[i]])
                        z[i]++;
                 if (i + z[i] > r)
                         1 = i, r = i + z[i];
        return z;
```

# 13 kactl forked/various

## 13.1 BumpAllocator

```
/**
    * Author: Simon Lindholm
    Date: 2015-09-12
    License: CC0
    * Source: me
    Description: When you need to dynamically allocate many objects and don't care about freeing them.
    "new X" otherwise has an overhead of something like 0.05us + 16 bytes per allocation.
```

### 13.2 BumpAllocatorSTL

```
* Author: Simon Lindholm
 * Date: 2016-07-23
 * License: CC0
 * Source: me
 * Description: BumpAllocator for STL containers.
 * Usage: vector<vector<int, small<int>>> ed(N);
 * Status: tested
#pragma once
char buf[450 << 20] alignas(16);
size_t buf_ind = sizeof buf;
template<class T> struct small {
        typedef T value_type;
        small() {}
        template<class U> small(const U&) {}
        T* allocate(size_t n) {
                buf_ind -= n * sizeof(T);
                buf_ind &= 0 - alignof(T);
                return (T*) (buf + buf_ind);
        void deallocate(T*, size_t) {}
};
```

#### 13.3 ConstantIntervals

```
* Author: Simon Lindholm
 * Date: 2015-03-20
 * License: CC0
 * Description: Split a monotone function on [from, to) into a minimal set of
     half-open intervals on which it has the same value.
 * Runs a callback g for each such interval.
 * Usage: constantIntervals(0, sz(v), [\ (int x){return v[x];}, [\ (int lo,
     int hi, T val) {...});
 * Time: O(k\log\frac{n}{k})
 * Status: tested
 */
#pragma once
template<class F, class G, class T>
void rec(int from, int to, F& f, G& q, int& i, T& p, T q) {
       if (p == q) return;
        if (from == to) {
               g(i, to, p);
                i = to; p = q;
                int mid = (from + to) >> 1;
                rec(from, mid, f, g, i, p, f(mid));
                rec(mid+1, to, f, g, i, p, q);
template < class F, class G>
void constantIntervals(int from, int to, F f, G g) {
```

```
if (to <= from) return;
int i = from; auto p = f(i), q = f(to-1);
rec(from, to-1, f, g, i, p, q);
g(i, to, q);
```

### 13.4 DivideAndConquerDP

```
* Author: Simon Lindholm
 * License: CC0
 * Source: Codeforces
 * Description: Given a[i] = \min_{0 \le k \le k} (i) \le k \le hi(i) (f(i, k))  where the (
     minimal)
 * optimal $k$ increases with $i$, computes $a[i]$ for $i = L..R-1$.
 * Time: O((N + (hi-lo)) \log N)
 * Status: tested on http://codeforces.com/contest/321/problem/E
 */
#pragma once
struct DP { // Modify at will:
        int lo(int ind) { return 0; }
        int hi(int ind) { return ind; }
        11 f(int ind, int k) { return dp[ind][k]; }
        void store(int ind, int k, ll v) { res[ind] = pii(k, v); }
        void rec(int L, int R, int LO, int HI) {
                if (L >= R) return;
                int mid = (L + R) \gg 1;
                pair<11, int> best(LLONG_MAX, LO);
                rep(k, max(LO,lo(mid)), min(HI,hi(mid)))
                        best = min(best, make_pair(f(mid, k), k));
                store (mid, best.second, best.first);
                rec(L, mid, LO, best.second+1);
                rec(mid+1, R, best.second, HI);
        void solve(int L, int R) { rec(L, R, INT_MIN, INT_MAX); }
};
```

## 13.5 FastInput

```
/**
 * Author: chilli
 * License: CC0
 * Source: Own work
 * Description: Read an integer from stdin. Usage requires your program to pipe
 * input from file.
 * Usage: ./a.out < input.txt
 * Time: About 5x as fast as cin/scanf.
 * Status: tested on SPOJ INTEST, unit tested
#pragma once
inline char gc() { // like getchar()
        static char buf[1 << 16];
        static size_t bc, be;
        if (bc >= be) {
                buf[0] = 0, bc = 0;
                be = fread(buf, 1, sizeof(buf), stdin);
        return buf[bc++1; // returns 0 on EOF
int readInt() {
        while ((a = gc()) < 40);
if (a == '-') return -readInt();
        while ((c = gc()) >= 48) a = a * 10 + c - 480;
        return a - 48;
```

### 13.6 FastKnapsack

```
* Author: Marten Wiman
* License: CC0
 * Source: Pisinger 1999, "Linear Time Algorithms for Knapsack Problems with
      Bounded Weights"
 * Description: Given N non-negative integer weights w and a non-negative target
 \star computes the maximum S <= t such that S is the sum of some subset of the
     weights.
 * Time: O(N \max(w i))
 * Status: Tested on kattis:eavesdropperevasion, stress-tested
#pragma once
int knapsack(vi w, int t) {
        int a = 0, b = 0, x;
        while (b < sz(w) && a + w[b] <= t) a += w[b++];
       if (b == sz(w)) return a;
        int m = *max_element(all(w));
        vi u, v(2*m, -1);
        v[a+m-t] = b;
        rep(i,b,sz(w))
                u = v:
                rep (x, 0, m) v[x+w[i]] = max(v[x+w[i]], u[x]);
                for (x = 2*m; --x > m;) rep(j, max(0,u[x]), v[x])
                        v[x-w[j]] = max(v[x-w[j]], j);
        for (a = t; v[a+m-t] < 0; a--);
        return a;
```

### 13.7 FastMod

```
* Author: Simon Lindholm
 * Date: 2020-05-30
 * License: CC0
 * Source: https://en.wikipedia.org/wiki/Barrett_reduction
 * Description: Compute $a \% b$ about 5 times faster than usual, where $b$ is
      constant but not known at compile time.
 * Returns a value congruent to $a \pmod b$ in the range $[0, 2b)$.
 * Status: proven correct, stress-tested
 * Measured as having 4 times lower latency, and 8 times higher throughput, see
     stress-test.
 * More precisely, it can be proven that the result equals 0 only if a = 0,
 * and otherwise lies in \{[1, (1 + a/2^64) * b)\}.
#pragma once
typedef unsigned long long ull;
struct FastMod {
        ull b, m;
        FastMod(ull b) : b(b), m(-1ULL / b) {}
        ull reduce(ull a) { // a % b + (0 or b)
                return a - (ull) ((__uint128_t(m) * a) >> 64) * b;
};
```

### 13.8 IntervalContainer

```
/**
  * Author: Simon Lindholm
  * License: CC0
  * Description: Add and remove intervals from a set of disjoint intervals.
  * Will merge the added interval with any overlapping intervals in the set when adding.
```

```
* Intervals are [inclusive, exclusive).
 * Time: O(\log N)
 * Status: stress-tested
#pragma once
set<pii>::iterator addInterval(set<pii>& is, int L, int R) {
        if (L == R) return is.end();
        auto it = is.lower_bound({L, R}), before = it;
        while (it != is.end() && it->first <= R) {
                R = max(R, it->second);
                before = it = is.erase(it);
        if (it != is.begin() && (--it)->second >= L) {
                L = min(L, it->first);
                R = max(R, it->second);
                is.erase(it);
        return is.insert(before, {L,R});
void removeInterval(set<pii>& is, int L, int R) {
        if (L == R) return;
        auto it = addInterval(is, L, R);
        auto r2 = it->second;
        if (it->first == L) is.erase(it);
        else (int&)it->second = L;
        if (R != r2) is.emplace(R, r2);
```

#### 13.9 IntervalCover

```
* Author: Johan Sannemo
 * License: CC0
 * Description: Compute indices of smallest set of intervals covering another
 * Intervals should be [inclusive, exclusive). To support [inclusive, inclusive
     ],
 * change (A) to add \texttt{|| R.empty()}. Returns empty set on failure (or if
     G is empty).
 * Time: O(N \log N)
 * Status: Tested on kattis:intervalcover
 */
#pragma once
template<class T>
vi cover(pair<T, T> G, vector<pair<T, T>> I) {
        vi S(sz(I)), R;
        iota(all(S), 0);
        sort(all(S), [&](int a, int b) { return I[a] < I[b]; });</pre>
        T cur = G.first;
        int at = 0;
        while (cur < G.second) { // (A)
                pair<T, int> mx = make pair(cur, -1);
                while (at < sz(I) \&\& I[S[at]].first <= cur) {
                        mx = max(mx, make_pair(I[S[at]].second, S[at]));
                        at++;
                if (mx.second == -1) return {};
                cur = mx.first;
                R.push_back (mx.second);
        return R;
}
```

### 13.10 KnuthDP

```
/**

* Author: Simon Lindholm

* License: CC0

* Source: http://codeforces.com/blog/entry/8219
```

#### 13.11 LIS

```
* Author: Johan Sannemo
 * License: CC0
 * Description: Compute indices for the longest increasing subsequence.
* Time: $0(N \log N)$
 * Status: Tested on kattis:longincsubseq, stress-tested
#pragma once
template<class I> vi lis(const vector<I>& S) {
        if (S.empty()) return {};
        vi prev(sz(S));
        typedef pair<I, int> p;
        vector res;
        rep(i,0,sz(S)) {
                // change 0 -> i for longest non-decreasing subsequence
                auto it = lower_bound(all(res), p{S[i], 0});
                if (it == res.end()) res.emplace_back(), it = res.end()-1;
                *it = {S[i], i};
                prev[i] = it == res.begin() ? 0 : (it-1) -> second;
        int L = sz(res), cur = res.back().second;
        vi ans(L);
        while (L--) ans[L] = cur, cur = prev[cur];
       return ans;
```

#### 13.12 SIMD

```
* Author: Simon Lindholm
 * Date: 2015-03-18
 * License: CC0
 * Source: https://software.intel.com/sites/landingpage/IntrinsicsGuide/
 * Description: Cheat sheet of SSE/AVX intrinsics, for doing arithmetic on
     several numbers at once.
 * Can provide a constant factor improvement of about 4, orthogonal to loop
     unrolling.
 * Operations follow the pattern \texttt{"\_mm(256)?\_name\_(si(128|256)|epi
      (8|16|32|64)|pd|ps)"}. Not all are described here;
 * grep for \text{texttt}_{\text{mm}} in \text{texttt}_{\text{usr/lib/gcc/}}^{4}/4.9/\text{include/} for more.
 \star If AVX is unsupported, try 128-bit operations, "emmintrin.h" and \#define \
     * For aligned memory use \texttt{\_mm\_malloc(size, 32)} or \texttt{int buf[N]
     alignas(32)}, but prefer loadu/storeu.
#pragma once
#pragma GCC target ("avx2") // or sse4.1
#include "immintrin.h" /** keep-include */
typedef m256i mi;
#define L(x) _mm256_loadu_si256((mi*)&(x))
// High-level/specific methods:
// load(u)?_si256, store(u)?_si256, setzero_si256, _mm_malloc
// blendv_(epi8|ps|pd) (z?y:x), movemask_epi8 (hibits of bytes)
```

```
// i32gather_epi32(addr, x, 4): map addr[] over 32-b parts of x
// sad_epu8: sum of absolute differences of u8, outputs 4xi64
// maddubs_epi16: dot product of unsigned i7's, outputs 16xi15
// madd_epi16: dot product of signed i16's, outputs 8xi32
// extractf128_si256(, i) (256->128), cvtsi128_si32 (128->lo32)
// permute2f128_si256(x,x,1) swaps 128-bit lanes
// shuffle_epi32(x, 3*64+2*16+1*4+0) == x for each lane
// shuffle_epi8(x, y) takes a vector instead of an imm
// Methods that work with most data types (append e.g. _epi32):
// set1, blend (i8?x:y), add, adds (sat.), mullo, sub, and/or,
// andnot, abs, min, max, sign(1,x), cmp(gt|eg), unpack(lo|hi)
int sumi32(mi m) { union {int v[8]; mi m;} u; u.m = m;
        int ret = 0; rep(i,0,8) ret += u.v[i]; return ret; }
mi zero() { return _mm256_setzero_si256(); }
mi one() { return _mm256_set1_epi32(-1); }
bool all_zero(mi m) { return _mm256_testz_si256(m, m); }
bool all_one(mi m) { return _mm256_testc_si256(m, one()); }
11 example_filteredDotProduct(int n, short* a, short* b) {
        int i = 0; 11 r = 0;
        mi zero = _mm256_setzero_si256(), acc = zero;
        while (i + 16 \le n)
                mi \ va = L(a[i]), \ vb = L(b[i]); \ i += 16;
                va = _mm256_and_si256(_mm256_cmpgt_epi16(vb, va), va);
                mi vp = _mm256_madd_epi16(va, vb);
                acc = _mm256_add_epi64(_mm256_unpacklo_epi32(vp, zero),
                        _mm256_add_epi64(acc, _mm256_unpackhi_epi32(vp, zero)));
        union {ll v[4]; mi m;} u; u.m = acc; rep(i,0,4) r += u.v[i];
        for (;i < n; ++i) if (a[i] < b[i]) r += a[i]*b[i]; // <- equiv
        return r;
```

#### 13.13 SmallPtr

```
* Author: Simon Lindholm
 * Date: 2016-08-23
 * License: CCO
 * Source: me
 * Description: A 32-bit pointer that points into BumpAllocator memory.
 * Status: tested
#pragma once
#include "BumpAllocator.h"
template<class T> struct ptr {
        unsigned ind;
        ptr(\bar{T}*p=0) : ind(p ? unsigned((char*)p - buf) : 0) {
                assert(ind < sizeof buf);
        T& operator*() const { return *(T*)(buf + ind); }
        T* operator->() const { return &**this; }
        T& operator[](int a) const { return (&**this)[a]; }
        explicit operator bool() const { return ind; }
};
```

## 13.14 TernarySearch

```
/**
    * Author: Simon Lindholm
    * Date: 2015-05-12
    * License: CC0
    * Source: own work
    * Description:
    * Find the smallest i in $[a,b]$ that maximizes $f(i)$, assuming that $f(a) < \
        dots < f(i) \ge \dots \ge f(b)$.

* To reverse which of the sides allows non-strict inequalities, change the <
        marked with (A) to <=, and reverse the loop at (B).</pre>
```

```
* To minimize $f$, change it to >, also at (B).
* Usage:
    int ind = ternSearch(0,n-1,[\&](int i){return a[i];});
* Time: O(\log(b-a))
* Status: tested
*/
#pragma once

template<class F>
int ternSearch(int a, int b, F f) {
    assert(a <= b);
    while (b - a >= 5) {
        int mid = (a + b) / 2;
            if (f(mid) < f(mid+1)) a = mid; // (A)
            else b = mid+1;
    }
    rep(i,a+1,b+1) if (f(a) < f(i)) a = i; // (B)
    return a;
}</pre>
```

## 13.15 Unrolling

### 14 Math

## 14.1 baby step gigant step

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class string>
using ordered_set = tree<string, null_type, less<string>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 2001
#define mod 1000000007
int bsgs(int a, int b, int m)
  if (a == 0 && b == 0)
   return 1;
  a \% = m, b \% = m;
  int k = 1, add = 0, g;
  while ((g = __gcd(a, m)) > 1) // fazer a e m serem coprimos
    if (b == k)
```

```
return add;
   if (b % g)
     return -1;
   b /= g, m /= g, ++add;
k = (k * 111 * a / g) % m;
 int n = sqrt(m) + 1;
 int an = 1;
 for (int i = 0; i < n; i++)
   an = (an * 111 * a) % m;
  unordered_map<int, int> vals;
 for (int q = 0, cur = b; q \le n; q++)
    vals[cur] = q;
   cur = (cur * 111 * a) % m;
 for (int p = 1, cur = k; p \le n; p++)
    cur = (cur * 111 * an) % m;
   if (vals.count(cur))
     int ans = n * p - vals[cur] + add;
     return ans;
 return -1;
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int q;
 cin \rightarrow q;
 while (q--)
   int a, b, m;
   cin >> a >> b >> m;
   cout << bsgs(a, b, m) << endl;</pre>
 return 0;
// menor x tal que: (a^x) % m = b % m
// a e m sao coprimos
// se nao forem coprimos tem como tratar
// complexidade: sqrt(m)
```

## 14.2 berlekamp massey

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
struct modint
  modint(int v = 0) { val = ((v % mod) + mod) % mod; }
  int pow(int y)
   modint x = val;
```

```
modint z = 1;
    while (y)
      if (y & 1)
       z *= x;
      x *= x;
     y >>= 1;
   return z.val:
  int inv() { return pow(mod - 2); }
  void operator=(int o) { val = o % mod; }
  void operator=(modint o) { val = o.val % mod; }
  void operator+= (modint o) { *this = *this + o; }
  void operator = (modint o) { *this = *this - o;
  void operator*=(modint o)
                             { *this = *this * 0;
  void operator/=(modint o)
                            { *this = *this / o; }
  bool operator== (modint o) { return val == o.val;
  bool operator!=(modint o) { return val != o.val; }
  int operator*(modint o) { return ((val * o.val) % mod);
  int operator/(modint o) { return (val * o.inv()) % mod;
  int operator+(modint o) {
                            return (val + o.val) % mod; }
  int operator-(modint o) { return (val - o.val + mod) % mod; }
// berlekamp massey (muito roubado)
// mas precisa que o mod seja primo (para poder achar inverso)
// dado os n primeiros termos de uma recorrencia linear
// a[0], a[1], a[2], ..., a[n - 1]
// ele acha a reccorrencia linear mais curta que da matching com os n primeiros
    valores
vector<modint> berlekamp_massey(vector<modint> x)
  vector<modint> ls, cur;
  int 1f, 1d;
  for (int i = 0; i < x.size(); i++)</pre>
    modint t = 0:
    for (int j = 0; j < cur.size(); j++)
      t += (x[i - j - 1] * cur[j]);
    if (modint(t - x[i]).val == 0)
      continue;
    if (cur.empty())
      cur.resize(i + 1);
      1d = (t - x[i]) % mod;
      continue;
    modint k = -(x[i] - t);
    k *= modint(ld).inv();
    vector<modint> c(i - 1f - 1);
    c.pb(k);
    for (auto const & j : 1s)
      modint curr = modint(j.val * -1) * k;
      c.pb(curr);
    if (c.size() < cur.size())</pre>
      c.resize(cur.size());
    for (int j = 0; j < cur.size(); j++)</pre>
      c[j] = c[j] + cur[j];
    if (i - lf + ls.size() >= cur.size())
      tie(ls, lf, ld) = make\_tuple(cur, i, t - x[i]);
   cur = c;
  return cur:
modint get_nth(vector<modint> rec, vector<modint> dp, int n)
  int m = rec.size();
  vector<modint> s(m), t(m);
  s[0] = 1;
```

```
if (m != 1)
   t[1] = 1;
  else
   t[0] = rec[0];
  auto mul = [&rec] (vector<modint> v, vector<modint> w)
    vector<modint> ans(2 * v.size());
    for (int j = 0; j < v.size(); j++)
      for (int k = 0; k < v.size(); k++)
        ans[j + k] += v[j] * w[k];
    for (int j = 2 * v.size() - 1; j >= v.size(); j--)
      for (int k = 1; k <= v.size(); k++)</pre>
        ans[j - k] += ans[j] * rec[k - 1];
    ans.resize(v.size());
    return ans;
  while (n)
    if (n & 1)
     s = mul(s, t);
    t = mul(t, t);
   n >>= 1;
  modint ret = 0;
  for (int i = 0; i < m; i++)
   ret += s[i] * dp[i];
  return ret;
modint guess_nth_term(vector<modint> x, int n)
  if (n < x.size())
   return x[n];
  vector<modint> coef = berlekamp_massey(x); // coeficientes da recorrencia
  /*for (auto const &i : coef)
   cout << i.val << " ";
  cout << endl; */
  if (coef.empty())
   return 0;
  return get_nth(coef, x, n);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  vector<modint> vals;
  vals.pb(0);
  vals.pb(1);
  for (int i = 2; i <= 200; i++)
   vals.pb(vals[vals.size() - 1] + vals[vals.size() - 2]);
  int n;
  cin >> n;
  cout << quess_nth_term(vals, n).val << endl;</pre>
  return 0;
// exemplo fibonacci
```

### 14.3 binomial theorem

```
#include <bits/stdc++.h>
using namespace std;

#pragma GCC optimize("03")
#pragma GCC optimize("unroll-loops")
#pragma gcc optimize("Ofast")

#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
```

```
#define MAXN 100005
#define mod 1000000007
struct modint
  modint(int v = 0) { val = (v < 0) ? v + mod : v; }
  int pow(int y)
   modint x = val;
    modint z = 1;
    while (y)
      if (y & 1)
       z \star = x;
      x *= x;
      y >>= 1;
    return z.val;
  int inv() { return pow(mod - 2); }
  void operator=(int o) { val = o % mod;
  void operator=(modint o) { val = o.val % mod; }
  void operator+= (modint o) { *this = *this + o; }
  void operator = (modint o) { *this = *this - o;
                             { *this = *this * 0;
  void operator*=(modint o)
  void operator/=(modint o) { *this = *this / o; }
  bool operator== (modint o) { return val == o.val;
  bool operator!=(modint o) { return val != o.val; }
  int operator*(modint o) { return ((val * o.val) % mod);
  int operator/(modint o) { return (val * o.inv()) % mod; ;
  int operator+(modint o) { return (val + o.val) % mod; }
  int operator-(modint o) { return (val - o.val + mod) % mod; }
modint f[MAXN];
modint invfat[MAXN];
void calc()
  f[0] = 1;
  for (int i = 1; i < MAXN; i++)
   f[i] = f[i - 1] * i;
  for (int i = 0; i < MAXN; i++)</pre>
    invfat[i] = f[i].inv();
modint ncr(int n, int k)
  modint ans = f[n] * invfat[k];
  ans \star = invfat[n - k];
  return ans;
struct strange_modint // escrever um numero na forma: a + (b * sqrt(5))
  modint a, b;
  strange_modint() { a = 0, b = 0; }
  strange_modint(modint v) { a = v, b = 0; }
  strange_modint(modint v, modint v2) { a = v, b = v2; }
  strange_modint operator*(strange_modint o)
    return strange_modint((a \star o.a) + (b \star o.b \star 5), (a \star o.b) + (b \star o.a));
  strange_modint operator+(strange_modint o)
    return strange_modint(a + o.a, b + o.b);
  strange modint operator-(strange modint o)
    return strange modint(a - o.a, b - o.b);
  strange_modint pow(int y)
    strange_modint x(a, b);
    strange_modint z(1, 0);
    while (y)
     if (y & 1)
        z = z * x;
```

```
x = x * x;
     y >>= 1;
    return z;
};
strange_modint matrix_ans[3][3];
void matrix_multiply(strange_modint a[3][3], strange_modint b[3][3])
  strange_modint res[3][3];
  for (int i = 0; i < 3; i++)
    for (int j = 0; j < 3; j++)
      for (int k = 0; k < 3; k++)
        res[i][j] = res[i][j] + (a[i][k] * b[k][j]);
  for (int i = 0; i < 3; i++)
    for (int j = 0; j < 3; j++)
      a[i][j] = res[i][j];
void matrix_pow(strange_modint mat[3][3], int m)
  for (int i = 0; i < 3; i++)
    for (int j = 0; j < 3; j++)
      matrix_ans[i][j] = strange_modint((i == j));
  while (m > 0)
    if (m & 1)
     matrix_multiply(matrix_ans, mat);
    \mathbf{m} = \mathbf{m} / 2;
    matrix_multiply(mat, mat);
signed main()
  ios base::sync with stdio(false);
  cin.tie(NULL);
  calc();
  int n, k;
  cin >> n >> k;
  strange_modint a(modint(1) / 2, modint(1) / 2);
  strange_modint b(modint(1) / 2, -(modint(1) / 2));
  strange_modint c(0, modint(1) / 5);
  strange_modint ans(0, 0);
  for (int j = 0; j \le k; j++)
    strange_modint curr(modint(1000000006).pow(j), 0); // (-1)^j
    curr = curr * strange_modint(ncr(k, j), 0);
    strange_modint prod = a.pow(k - j) * b.pow(j);
    strange_modint s0(0, 0);
    strange_modint s1(1, 0);
    strange_modint mat[3][3] = {{prod, s0, s1}, {s1, s0, s0}, {s0, s0, s1}};
    matrix_pow(mat, n);
    curr = curr * (matrix_ans[0][0] + matrix_ans[0][2]);
    ans = ans + curr;
  ans = ans \star c.pow(k);
  cout << ans.a.val << endl;</pre>
  return 0;
// https://codeforces.com/gym/104412/problem/F
```

### 14.4 catalan

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
```

```
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 400005
#define mod 998244353
struct modint
  int val;
  modint(int v = 0) { val = v % mod; }
  int pow(int y)
   modint x = val;
    modint z = 1;
    while (y)
      if (y & 1)
       z *= x;
      x \star = x;
     y >>= 1;
    return z.val;
  int inv() { return pow(mod - 2); }
  void operator=(int o) { val = o % mod; }
  void operator=(modint o) { val = o.val % mod; }
  void operator+=(modint o) { *this = *this + o;
  void operator = (modint o) { *this = *this - o;
  void operator*=(modint o) { *this = *this * o;
  void operator/=(modint o) { *this = *this / o; }
  bool operator== (modint o) { return val == o.val;
  bool operator!=(modint o) { return val != o.val; }
  int operator*(modint o) { return ((val * o.val) % mod);
  int operator/(modint o) { return (val * o.inv()) % mod;
  int operator+(modint o) { return (val + o.val) % mod; }
  int operator-(modint o) { return (val - o.val + mod) % mod; }
};
modint f[MAXN];
modint inv[MAXN];
modint invfat[MAXN];
void calc()
  f[0] = 1;
  for (int i = 1; i < MAXN; i++)</pre>
    f[i] = f[i - 1] * i;
  invfat[0] = 1;
  for (int i = MAXN - 1; i >= 1; i--)
    invfat[i] = modint(f[i]).inv();
modint ncr(int n, int k)
  modint ans = f[n] * invfat[k];
  ans *= invfat[n - k];
  return ans;
modint catalan(int x)
  modint ans = modint(1) / modint(x + 1);
  ans \star = ncr(2 \star x, x);
  return ans;
signed main()
```

```
ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  calc();
  cout << catalan(10).val << endl;
  return 0;
}

// catalan numbers:
// 1, 1, 2, 5, 14, 42, 132, 429, 1430, ...
// c[i] = (1 / (n + 1)) * ncr(2n, n)

// existem c[i] regular bracket sequences de tamanho 2 * n
// o numero de binary trees completas com n + 1 folhas
// o numero de possibilidades de conectar 2n numeros em um circulo em n arestas
// https://cp-algorithms.com/combinatorics/catalan-numbers.html</pre>
```

#### 14.5 crivo

```
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define in insert
#define pi pair<int, int>
#define pd pair<double, int>
#define pii pair<int, pi>
#define mp make_pair
#define fir first
#define sec second
#define MAXN 100001
#define mod 1000000007
bitset <MAXN> prime;
void crivo ()
  prime.set();
  prime[0] = false;
  prime[1] = false;
  for (int i = 2; i < MAXN; i++)
    if(prime[i])
      for (int j = 2; j * i < MAXN; j++)
        prime[j * i] = false;
signed main()
  crivo();
  int q;
  cin >> q;
  while (q--)
    int n;
    cin >> n;
    (prime[n]) ? cout << "YES\n" : cout << "NO\n" ;</pre>
  return 0;
```

### 14.6 crt

```
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 2000006
//#define mod 1000000007
namespace crt
  vector<pi> eq;
  int gcd(int a, int b, int &x, int &y)
   if (b == 0)
      x = 1, y = 0;
      return a;
    int x1, y1, d = gcd(b, a % b, x1, y1);
    x = y1, y = x1 - y1 * (a / b);
   return d;
  pi crt()
    int a1 = eq[0].fir, m1 = eq[0].sec;
    a1 %= m1:
    for (int i = 1; i < eq.size(); i++)</pre>
      int a2 = eq[i].fir, m2 = eq[i].sec;
      int g = __gcd(m1, m2);
if (a1 % g != a2 % g)
        return {-1, -1};
      int p, q;
gcd(m1 / g, m2 / g, p, q);
      int mod = m1 / g * m2;
      int x = (a1 * (m2 / g) * mod * q * mod + a2 * (m1 / g) * mod * p * mod) *
      a1 = x;
      if (a1 < 0)
       a1 += mod;
      m1 = mod;
   return {a1, m1};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  for (int i = 0; i < n; i++)</pre>
    int a, b;
   cin >> a >> b;
    crt::eq.pb({a, b});
  pi ans = crt::crt();
  if (ans.fir == -1)
    cout << "No solution\n";</pre>
   cout << ans.fir << " " << ans.sec << endl;</pre>
  return 0;
// https://forthright48.com/chinese-remainder-theorem-part-2-non-coprime-moduli/
// https://cp-algorithms.com/algebra/chinese-remainder-theorem.html
// https://www.geeksforgeeks.org/chinese-remainder-theorem-set-1-introduction/
// teorema chines do resto(crt)
// para resolver sistemas de congruencias modulares
// o menor inteiro a que satisfaz:
// a mod p1 = x1
// a mod p2 = x2
// ...
// a mod pn = xn
// a funcao crt retorna um pair {a, mod}
```

```
// dai a solucao pode ser descrita como
// x = a % mod
// entao os valores possiveis sao:
// a, (a + mod), a + (2 * mod), a + (3 * mod), ...
```

#### 14.7 crt trick

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 200005
vector<pi> eq;
map<int, int> by_mod;
bool check (pi curr)
  if (by_mod.find(curr.sec) != by_mod.end())
    return by_mod[curr.sec] == curr.fir;
  // no maximo O(sqrt(n)) mods distintos
  // ja que tem O(sqrt(n)) tamanhos de ciclos distintos na permutacao
  // ja que a soma dos tamanhos dos ciclos da permutacao eh iqual a n
  for (auto [x, mod] : eq)
    if ((curr.fir - x) % __gcd(curr.sec, mod))
      return 0:
  return 1:
void add(pi curr)
  eq.pb(curr);
  by_mod[curr.sec] = curr.fir;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  vector<int> p(n);
  for (int i = 0; i < n; i++)
    cin >> p[i];
   p[i]--;
  vector<int> a(n);
  for (int i = 0; i < n; i++)
    cin >> a[i];
    a[i]--;
  vector<bool> vis(n, 0);
  vector<int> comp(n);
  vector<int> pos(n);
  vector<vector<int>> vals(n);
  vector<vector<int>> vals2(n);
  for (int i = 0; i < n; i++)
    if (!vis[i])
```

```
int x = i;
      while (!vis[x])
        vis[x] = 1;
        comp[x] = i;
        pos[x] = vals[i].size();
        vals[i].pb(a[x]);
        vals2[i].pb(x);
        x = p[x];
   }
  set < int > st;
  vector<int> resp(n);
  for (int i = 0; i < n; i++)
    if (st.find(comp[i]) == st.end())
      st.insert(comp[i]);
      int x = comp[i];
      int sz = vals[x].size();
      int mn = 1e18, op = -1;
      for (int i = 0; i < sz; i++)
        // numeros x tal que x % sz = i
if (check({i, sz}) && vals[x][i] < mn)</pre>
          mn = vals[x][i];
          op = i;
      for (int i = 0; i < sz; i++)
        resp[vals2[x][i]] = vals[x][op];
        op = (op + 1) % sz;
      add({op, sz});
  for (auto const &i : resp)
    cout << i + 1 << " ";
  cout << endl;
// a[i] -> a[p[i]]
// i -> p[i]
// tenho varios ciclos
// e quero minimizar lexocograficamente a permutacao
// o que eu quero e:
// para cada ciclo da permutacao, ver se em um dado momento do ciclo
// eh possivel satisfazer o sistema de equacoes modulares
// dado um sistema com duas equacoes modulares
// x == n1 \pmod{m1}
// x == n2 \pmod{m2}
// existe solucao se (n1 - n2) % gcd(m1, m2) == 0
// para n equacoes podemos fazer esse check
// iterando por todos os pares de equacoes
// e checando se todo par satisfaz
// se eu ja tenho n equacoes no conjunto
// e eu sei que esse sistema de n equacoes tem solucao
// entao pra adicionar mais uma
// so preciso checar iterando pelas as que ja existem e dale
```

## 14.8 diophantine

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;

template <class T>
```

```
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200001
#define mod 998244353
namespace dio
 vector<pi> sols;
  int gcd(int a, int b, int &x, int &y)
    if (b == 0)
      x = 1, y = 0;
     return a;
    int x1, y1, d = gcd(b, a % b, x1, y1);
    x = y1, y = x1 - y1 * (a / b);
    return d:
  void one sol(int a, int b, int c)
    int x0, y0, g;
    g = gcd(abs(a), abs(b), x0, y0);
    if (c % q)
     return;
    x0 \star = (c / g);
    v0 \star = (c / g);
    if (a < 0)
     x0 *= -1;
    if (b < 0)
     y0 *= -1;
    sols.pb({x0, y0});
 void more_sols(int a, int b, int c)
    int g = \underline{gcd}(a, b);
    int x0 = sols[0].fir, y0 = sols[0].sec;
    for (int k = -200000; k \le 200000; k++)
      int x = x0 + k * (b / g);
     int y = y0 - k * (a / q);
      sols.pb({x, y});
 }
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
  int a, b, c;
 cin >> a >> b >> c;
  dio::one_sol(a, b, c);
 if (!dio::sols.size())
   cout << "No\n";</pre>
   return 0:
 dio::more_sols(a, b, c);
 bool can = false;
 for (auto const &i : dio::sols)
   can |= (i.fir >= 0 && i.sec >= 0);
  (can) ? cout << "Yes\n" : cout << "No\n";</pre>
 return 0;
// equacoes do tipo:
// ax + by = c
// o caso a = 0 e b = 0, nao eh tratado nesse codigo
// nesse caso quero checar se equacao diofantina tem uma solucao
// com x >= 0 e y >= 0
```

#### 14.9 division trick

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 300005
#define mod 1000000007
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  auto soma_pa = [&](int 1, int r)
   int x = (1 + r) % mod;
   x = (x * ((r - 1 + 1) % mod)) % mod;
   x = (x * 500000004) % mod; // divisao por 2
   return x;
  int ans = 0;
  for (int 1 = 1, r; 1 <= n; 1 = r + 1)
    // todos os numeros i no intervalo [1, r] possuem (n / i) == x
    r = n / (n / 1);
    int x = (n / 1);
   // for (int i = 1; i <= r; i++)
// assert((n / i) == x);
   ans += (x * soma_pa(1, r));
   ans %= mod;
  cout << ans << endl;</pre>
// https://cses.fi/problemset/task/1082
// dado um n, ache a soma dos divisores de todos os numeros de 1 ate n, (n <=
// se a gnt computar a contribuicao de cada divisor na resposta
// isso fica: achar a soma de i \star floor(n / i) para todos os numeros de 1 ate n
// a quatidade de valores distintos de floor(n / i) eh O(sqrt(n))
// entao o que queremos eh iterar por todos os valores possiveis, e ver o
     intervalo de numeros
// que tem o valor de floor(n / i) sendo um determinado x
```

#### 14.10 divisors

```
#include <bits/stdc++.h>
using namespace std;

#define PI acos(-1)
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define fir first
#define sec second
#define MAXN 5001
#define mod 1000000007

signed main()
```

```
ios_base::sync_with_stdio(false);
cin.tie(NULL);
int n;
cin >> n;
int ans = 0;
for (int i = 1; i <= sqrt(n); i++)
{
   if (!(n % i)) {
      ans++;
      if (n / i != i)
            ans++;
   }
}
cout << ans << endl;</pre>
```

#### 14.11 extended euclidean

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200001
#define mod 998244353
int gcd(int a, int b, int &x, int &y)
 if (b == 0)
   x = 1;
   \mathbf{y} = 0;
   return a;
 int x1, y1;
 int d = gcd(b, a % b, x1, y1);
 x = y1;
 y = x1 - y1 * (a / b);
 return d;
signed main()
 int n;
 cin >> n;
 int k = 2;
 while (1)
    int x, y;
   if (\gcd(k, n, x, y) == 1)
     x = ((x % n) + n) % n;
     cout << x << endl;</pre>
     return 0;
   k++;
 return 0;
// achar os numeros x e y tal que:
// a * x + b * y = gcd(a, b)
// problema exemplo:
// https://codeforces.com/group/btcK4I5D5f/contest/451372/problem/J
// dado um numero k
// quero achar um numero x, se possivel, tal que:
// (k * x) % n = 1
// k * x + n * y = 1
// se gcd(k, n) = 1, entao:
// k * x + n * y = gcd(k, n)
// note que, se gcd(k, n) > 1, logo nao existe solucao
```

#### 14.12 fft

```
#include <bits/stdc++.h>
using namespace std;
#define PI acos(-1)
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 125001
#define mod 1000000007
#define cd complex<double>
namespace fft
  int n;
  void fft(vector<cd> &a, bool invert)
    int n = a.size();
    for (int i = 1, j = 0; i < n; i++)
      int bit = n >> 1;
      for (; j & bit; bit >>= 1)
       j ^= bit;
j ^= bit;
      if (i < j)
        swap(a[i], a[j]);
    for (int len = 2; len <= n; len <<= 1)
      double ang = 2 * PI / len * (invert ? -1 : 1);
      cd wlen(cos(ang), sin(ang));
      for (int i = 0; i < n; i += len)
        for (int j = 0; j < len / 2; j++)
          cd u = a[i + j], v = a[i + j + len / 2] * w;
         a[i + j] = u + v;
          a[i + j + len / 2] = u - v;
          w \star = wlen;
    if (invert)
      for (cd &x : a)
        x /= n;
  vector<int> mul(vector<int> a, vector<int> b)
   vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());
    n = 1;
    while (n < a.size() + b.size())</pre>
     n <<= 1;
    fa.resize(n);
    fb.resize(n);
    fft(fa, false);
    fft(fb, false);
    for (int i = 0; i < n; i++)
  fa[i] *= fb[i];</pre>
    fft(fa, true);
    vector<int> ans(n);
    for (int i = 0; i < n; i++)</pre>
     ans[i] = round(fa[i].real());
   return ans;
signed main()
  ios::sync_with_stdio(false);
  cin.tie(0);
```

```
int n;
 cin >> n;
  int m = n + n;
 vector<int> a(m, 0);
 vector<int> aa(m, 0);
 for (int i = 0; i < n; i++)
   cin >> a[i];
   aa[i] = (a[i] == 47) ? 1 : 0;
   a[i + n] = a[i];
   aa[i + n] = aa[i];
 vector<int> b(m, 0);
 vector<int> bb(m, 0);
 for (int i = n - 1; i >= 0; i--)
   cin >> b[i + n];
   bb[i + n] = (b[i + n] == 47) ? 1 : 0;
 vector<int> ans1 = fft::mul(a, b);
 vector<int> ans2 = fft::mul(aa, bb);
 int ans = 0;
 for (int i = (m - 1); i < (m - 1) + n; i++)
   if (ans2[i] > 0)
     continue;
   ans = max(ans1[i], ans); // produto escalar de algum cyclic shift
 cout << ans << endl;
 return 0;
// https://algo.sk/br24/problem.php?problem=d3-badsquare
// exemplo do all possible scalar products
// dados dois arrays a e b de tamanho n
// quero computar o scalar product de todos os cyclics shifts de a com b
// duplicar o array a
// dar reverse no array b e adicionar n zeros no inicio
```

#### 14.13 fraction

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200006
#define mod 1000000007
struct fraction
  int x, y; // x / y
  fraction() {}
  fraction(int x, int y) : x(x), y(y) {}
 bool operator==(fraction o) { return (x * o.y == o.x * y); } bool operator!=(fraction o) { return (x * o.y != o.x * y); }
  bool operator>(fraction o) { return (x * o.y > o.x * y); }
  bool operator>=(fraction o) { return (x * o.y >= o.x * y); }
  bool operator<(fraction o) { return (x * o.y < o.x * y); }</pre>
  bool operator<=(fraction o) { return (x * o.y <= o.x * y); }</pre>
  fraction operator+(fraction o)
    fraction ans;
    ans.y = (y == o.y) ? y : y * o.y;
```

```
ans.x = (x) * (ans.y / y) + (o.x) * (ans.y / o.y);
    // ans.simplify();
    return ans;
  fraction operator*(fraction o)
    fraction ans;
    ans.x = x * o.x;
    ans.y = y * o.y;
    // ans.simplify();
    return ans;
  fraction inv()
    fraction ans = fraction(x, y);
    swap(ans.x, ans.y);
    return ans;
  fraction neg()
    fraction ans = fraction(x, y);
    ans.x *=-1;
   return ans;
  void simplify()
    if (abs(x) > 1e9 \mid \mid abs(y) > 1e9) // slow simplification
      int g = \underline{gcd}(y, x);
      x /= g;
      y /= g;
  // substraction and division can be easily done
};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
```

### 14.14 fwht

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1025
#define mod 998244353
struct modint
  modint(int v = 0) { val = ((v % mod) + mod) % mod; }
  int pow(int y)
    modint x = val;
    modint z = 1;
    while (y)
      if (y & 1)
```

```
z \star = x;
     x *= x;
     y >>= 1;
    return z.val;
 int inv() { return pow(mod - 2); }
 void operator=(int o) { val = o % mod; }
 void operator=(modint o) { val = o.val % mod; }
 void operator+=(modint o) { *this = *this + o;
 void operator = (modint o) { *this = *this - o;
 void operator*=(modint o) { *this = *this * o;
 void operator/=(modint o) { *this = *this / o; }
 bool operator== (modint o) { return val == o.val;
 bool operator!=(modint o) { return val != o.val; }
 int operator*(modint o) { return ((val * o.val) % mod); }
 int operator/(modint o) { return (val * o.inv()) % mod; }
 int operator+(modint o) { return (val + o.val) % mod; }
 int operator-(modint o) { return (val - o.val + mod) % mod; }
vector<modint> fwht(char op, vector<modint> f, bool inv = 0)
 int n = f.size();
 for (int k = 0; (n - 1) >> k; k++)
    for (int i = 0; i < n; i++)</pre>
     if (i >> k \& 1)
       int j = i ^ (1 << k);
if (op == '^')</pre>
         f[j] += f[i], f[i] = f[j] - modint(2) * f[i];
        if (op == ' | ')
         f[i] += modint(inv ? -1 : 1) * f[j];
        if (op == '&')
         f[j] += modint(inv ? -1 : 1) * f[i];
 if (op == '^' and inv)
   for (auto &i : f)
      i /= n;
 return f;
vector<modint> conv(char op, vector<modint> a, vector<modint> b)
 a = fwht(op, a, 0);
 b = fwht(op, b, 0);
 for (int i = 0; i < a.size(); i++)</pre>
   a[i] *= b[i];
 return fwht (op, a, 1);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n;
 cin >> n;
 n = 1 \ll n;
 vector<modint> a(n);
 for (int i = 0; i < n; i++)
   int x;
   cin >> x:
   a[i] = x;
 vector<modint> b(n);
 for (int i = 0; i < n; i++)
   int x;
   cin >> x;
   b[i] = x;
```

```
vector<modint> c = conv('^', a, b); // convolucao de xor
  for (auto const &i : c)
   cout << i.val << " ";
  cout << endl;</pre>
  vector<modint> d = conv('&', a, b); // convolucao de and
  for (auto const &i : d)
   cout << i.val << " ";
  cout << endl;
  return 0;
// o tipo ta como modint, mas tem como mudar para qualquer um
// usar preferencialmente tamanho como potencia de 2
// faz a convolucao de a com b
// c[k] = (a[i] * b[j]), com (i op j) = k
// op pode ser xor, and ou or
// para testar
// https://judge.yosupo.jp/problem/bitwise_xor_convolution
// https://judge.yosupo.jp/problem/bitwise_and_convolution
```

### 14.15 gaussian elimination

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define double long double
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define DEBUG 1
#define MAXN 2001
#define mod 1000000007
#define EPS 1e-9
vector<double> ans;
int gauss(vector<vector<double>> a)
  int n = a.size(), m = a[0].size() - 1, ret = 1;
  ans.assign(m, 0);
  vector<int> where (m, -1);
  for (int col = 0, row = 0; col < m && row < n; col++, row++)</pre>
    int sel = row;
    for (int i = row; i < n; i++)</pre>
      if (abs(a[i][col]) > abs(a[sel][col]))
        sel = i;
    if (abs(a[sel][col]) < EPS)</pre>
      continue;
    for (int i = col; i <= m; i++)</pre>
      swap(a[sel][i], a[row][i]);
    where[col] = row;
    for (int i = 0; i < n; i++)
      if (i != row)
        double c = a[i][col] / a[row][col];
        for (int j = col; j <= m; j++)</pre>
          a[i][j] = a[row][j] * c;
  for (int i = 0; i < m; i++)
```

```
if (where[i] != -1)
     ans[i] = (a[where[i]][m] / a[where[i]][i]);
    else
     ret = 2:
 for (int i = 0; i < n; i++)
    double sum = 0;
    for (int j = 0; j < m; j++)
      sum += (ans[j] * a[i][j]);
    if (abs(sum - a[i][m]) > EPS)
      ret = 0;
  return ret; // 0 = nao existe solucao, 1 = existe uma solucao, 2 = existem
      multiplas solucoes
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 vector<vector<double>> a = \{\{1.0, 1.0, 20.0\}, // 1x + 1y = 20\}
                              \{3.0, 4.0, 72.0\}\}; // 3x + 4y = 72
  cout << gauss(a) << endl;</pre>
 for (auto const &i : ans) // x = 8 e y = 12
   cout << i << " ";
  cout << endl;
// eliminacao gaussiana
// para resolver sistemas com n equacoes e m incognitas
// para isso iremos utilizar uma representacao usando
// matrizes, no qual uma coluna extra e adicionada,
// representando os resultados de cada equacao.
// algoritimo:
// ideia: qualquer equacao pode ser reescrita como uma combinacao linear dela
// 1- dividir a primeira linha(primeira equacao) por a[0][0]
// 2- adicionar a primeira linha as linhas restantes, de modo que, os
    coeficientes da primeira coluna se tornem todos zeros, para que
     isso aconteca, na i-esima linha devemos adicionar a primeira linha
    multiplicada por (a[i][0] * -1)
// 3- com isso, o elemento a[0][0] = 1 e os demais elementos da primeira coluna
    serao iguais a zero
// 4- continuamos o algoritimo a partir da etapa 1 novamente, dessa vez
     com a segunda coluna e a segunda linha, dividindo a linha por a[1][1]
     e assim sucessivamente
// 5- ao final, teremos a resposta
// complexidade O(min(n, m) * n * m);
// se n == m, logo a complexidade sera O(n^3)
```

## 14.16 gaussian elimination2

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 4007
#define mod 998244353
#define EPS 1e-9
bitset<MAXN> ans;
int gauss(vector<bitset<MAXN>> &a)
  ans.reset();
  int n = a.size(), m = a[0].size() - 1, ret = 1;
  vector<int> where(m, -1);
  for (int col = 0, row = 0; col < m && row < n; col++)</pre>
```

```
for (int i = row; i < n; i++)</pre>
      if (a[i][col])
        swap(a[i], a[row]);
       break;
    if (!a[row][col])
     continue;
    where[col] = row;
    for (int i = 0; i < n; i++)
      if (i != row && a[i][col])
       a[i] ^= a[row];
    ++row:
  for (int i = 0; i < m; i++)
   if (where[i] != -1)
     ans[i] = (a[where[i]][m] / a[where[i]][i]);
    else
     ret = 2;
  for (int i = 0; i < n; i++)
    double sum = 0;
    for (int j = 0; j < m; j++)</pre>
     sum += (ans[j] * a[i][j]);
    if (abs(sum - a[i][m]) > EPS)
     ret = 0;
 return ret:
signed main()
 ios::sync_with_stdio(false);
  cin.tie(0);
  int m, n;
 cin >> m >> n;
  string s;
  getline(cin, s);
 auto get_id = [&](string st)
   return n - stoi(st.substr(1));
  vector<bitset<MAXN>> v(n + m);
  for (int i = 0; i < m; i++)
   getline(cin, s);
    s = s.substr(1); // (
    if (i == m - 1)
      s = s.substr(0, s.size() - 1); //)
    else
     s = s.substr(0, s.size() - 5); // ) and
    istringstream input_stream(s);
    string t;
    while (input_stream >> t)
      if (t == "not") // not var
        input_stream >> t;
        v[i][n + qet_id(t)] = v[i][n + qet_id(t)]^1;
      else if (t != "or") // var
        v[i][get\_id(t)] = v[i][get\_id(t)] ^ 1;
    v[i][MAXN - 1] = v[i][MAXN - 1] ^ 1;
  for (int i = 0; i < n; i++)
   v[i + m][i] = v[i + m][i] ^ 1;
   v[i + m][n + i] = v[i + m][n + i] ^ 1;
   v[i + m][MAXN - 1] = v[i + m][MAXN - 1] ^ 1;
  if (gauss(v) == 0)
```

```
cout << "impossible\n";</pre>
   return 0:
 string resp(n, 'F');
 int id = n - 1;
 for (int i = 0; i < (n + n); i++)
   if (ans[i])
      (i < n) ? resp[id] = 'T' : resp[id] = 'F';
   id--;
   if (id < 0)
     id = n - 1;
 cout << resp << endl;</pre>
 return 0;
// exemplo de solucao para o https://codeforces.com/gym/101908/problem/M
// esse codigo ja acha a menor solucao lexicograficamente (caso exista)
// caso a gente queira a maior lexicograficamente (que e o caso desse problema
// basta considerar as variaveis na ordem contraria
```

## 14.17 lagrange

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 600005
#define mod 1000000007
struct modint
  int val;
 modint(int v = 0) { val = v % mod; }
  int pow(int y)
    modint x = val;
    modint z = 1;
    while (y)
      if (y & 1)
       z *= x;
      x *= x;
      y >>= 1;
    return z.val;
  int inv() { return pow(mod - 2); }
  void operator=(int o) { val = o % mod; }
  void operator=(modint o) { val = o.val % mod; }
  void operator+=(modint o) { *this = *this + o;
  void operator = (modint o) { *this = *this - o;
  void operator*=(modint o) { *this = *this * o;
  void operator/=(modint o) {
                              *this = *this / o;
  bool operator==(modint o) { return val == o.val;
  bool operator!=(modint o) { return val != o.val;
  int operator*(modint o) { return ((val * o.val) % mod); }
  int operator/(modint o) { return (val * o.inv()) % mod; }
  int operator+(modint o) { return (val + o.val) % mod; }
```

```
int operator-(modint o) { return (val - o.val + mod) % mod; }
struct lagrange
 int n;
  vector<modint> den;
  vector<modint> y;
  vector<modint> fat;
  vector<modint> inv_fat;
  lagrange(vector<modint> &v) // f(i) = v[i], gera um polinomio de grau n - 1
   n = v.size();
   calc(n);
   calc_den(n);
   y = v;
  void calc_den(int n)
    den.resize(n);
    for (int i = 0; i < n; i++)
      den[i] = inv_fat[n - i - 1] * inv_fat[i];
      if ((n - i - 1) % 2 == 1)
       int x = (mod - den[i].val) % mod;
        den[i] = x;
  void calc(int n)
    fat.resize(n + 1);
    inv_fat.resize(n + 1);
    fat[0] = 1;
    inv_fat[0] = 1;
    for (int i = 1; i <= n; i++)</pre>
      fat[i] = fat[i - 1] * i;
      inv_fat[i] = fat[i].inv();
  modint get_val(int x) // complexidade: O(n)
    x %= mod;
    vector<modint> 1(n);
    vector<modint> r(n);
    1[0] = 1, r[n - 1] = 1;
    for (int i = 1; i < n; i++)
      modint cof = (x - (i - 1) + mod);
      l[i] = l[i - 1] * cof;
    for (int i = n - 2; i >= 0; i--)
      modint cof = (x - (i + 1) + mod);
      r[i] = r[i + 1] * cof;
    modint ans = 0:
    for (int i = 0; i < n; i++)
      modint cof = l[i] * r[i];
      ans += modint(cof * y[i]) * den[i];
  vector<modint> find_coefs() // encontra os coeficientes do polinomio
    int nn = n;
    int d = nn - 1;
    vector<modint> c(nn, 0);
    for (int i = 0; i < y.size(); i++)</pre>
      c[d] += (y[i] * den[i]);
    for (int p = nn - 2; p >= 0; p--)
     nn--:
      calc_den(nn);
```

```
for (int i = 0; i \le p; i++)
        y[i] = (c[p + 1] * modint(i).pow(d));
        c[p] += (y[i] * den[i]);
     d--;
   return c;
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n, k;
 cin >> n >> k;
 vector<modint> v;
 v.pb(0);
 int lim = k + 1;
 for (int i = 1; i <= lim; i++)</pre>
   v.pb(v.back() + modint(i).pow(k));
 lagrange 1(v);
 cout << l.get_val(n).val << endl;</pre>
 return 0;
// https://codeforces.com/contest/622/problem/F
// https://codeforces.com/contest/1817/problem/C
// https://codeforces.com/gym/103388/problem/A
```

#### 14.18 lucas theorem

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1000006
#define mod 2
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int q;
  cin >> q;
  while (q--)
    int n, k;
    cin >> n >> k;
    for (int i = 0; i < n; i++)
     if (((n-1) \& i) == i)
       cout << k << " ";
      else
        cout << 0 << " ";
   cout << endl:
// https://codeforces.com/contest/2072/problem/F
// problema interessante de um div3, calcular ncr(n, k) mod 2
// ncr(n-1, i) mod 2 = (((n-1) & i) == i)
// sejam m e n numeros inteiros nao negativos e p um numero primo
// desenvolver n e m na base p
```

```
// ou seja:
// m = m[k]*p^(k) + m[k - 1]*p^(k - 1) + ... + m[0]*p^(0)
// n = n[k]*p^(k) + n[k - 1]*p^(k - 1) + ... + n[0]*p^(0)

// entao:
// ncr(m, n) mod p = produtorio de (ncr(m[i], n[i]) mod p)

// dai pra generalizar pro mod 2 eh deboas, pq se tiver um bit setado em n[i] que nao ta setado em m[i], entao miou, vai dar zero
```

#### 14.19 markov

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
struct modint
  int val:
  modint(int v = 0) { val = ((v % mod) + mod) % mod; }
  int pow(int y)
   modint x = val;
    modint z = 1;
    while (y)
      if (y & 1)
       z *= x;
      x *= x:
     y >>= 1;
    return z.val;
  int inv() { return pow(mod - 2); }
  void operator=(int o) { val = o % mod;
  void operator=(modint o) { val = o.val % mod; }
  void operator+= (modint o) { *this = *this + o; }
  void operator = (modint o) { *this = *this - o;
  void operator*=(modint o) { *this = *this * o;
  void operator/=(modint o) { *this = *this / o; }
  bool operator==(modint o) { return val == o.val;
  bool operator!=(modint o) { return val != o.val; }
  int operator*(modint o) { return ((val * o.val) % mod);
  int operator/(modint o) { return (val * o.inv()) % mod;
  int operator+(modint o) { return (val + o.val) % mod; }
  int operator-(modint o) { return (val - o.val + mod) % mod; }
};
vector<modint> ans;
int gauss(vector<vector<modint>> a)
  int n = a.size(), m = a[0].size() - 1, ret = 1;
  ans.assign(m, 0);
  vector<int> where (m, -1);
  for (int col = 0, row = 0; col < m && row < n; col++, row++)</pre>
    int sel = row;
    for (int i = row; i < n; i++)</pre>
      if (a[i][col].val > a[sel][col].val)
```

```
sel = i;
    if (!a[sel][col].val)
      continue;
    for (int i = col; i <= m; i++)</pre>
      swap(a[sel][i], a[row][i]);
    where[col] = row;
for (int i = 0; i < n; i++)</pre>
      if (i != row)
        modint c = a[i][col] / a[row][col];
         for (int j = col; j <= m; j++)</pre>
          a[i][j] -= modint(a[row][j] * c);
  for (int i = 0; i < m; i++)
    if (where[i] != -1)
      ans[i] = (a[where[i]][m] / a[where[i]][i]);
    else
      ret = 2;
  for (int i = 0; i < n; i++)
    modint sum = 0;
    for (int j = 0; j < m; j++)
  sum += (ans[j] * a[i][j]);</pre>
    if ((sum - a[i][m]) > 0)
      ret = 0;
  assert(ret == 1); // so uma solucao existe
  return ret;
int dx[] = \{-1, 0, 1, 0\};
int dy[] = \{0, -1, 0, 1\};
modint a[4];
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int r;
  cin >> r;
  modint sum = 0;
  for (int i = 0; i < 4; i++)
    cin >> a[i].val;
    sum += a[i];
  for (int i = 0; i < 4; i++)
    a[i] /= sum;
  vector<pi> in;
  queue<pi> q;
  set <pi> vis;
  q.push({0, 0});
  vis.insert({0, 0});
  // quais posicoes tao a distancia <= R da origem
  while (!q.empty())
    auto [x, y] = q.front();
    q.pop();
    if ((x * x) + (y * y) > (r * r))
      continue;
    in.pb({x, y});
    for (int dir = 0; dir < 4; dir++)</pre>
      int i = x + dx[dir];
      int j = y + dy[dir];
      if (vis.find({i, j}) == vis.end())
```

```
q.push({i, j});
       vis.insert({i, j});
   }
  sort(in.begin(), in.end());
  int sz = in.size() + 1;
 auto get_id = [&](int x, int y)
   if ((x \star x) + (y \star y) > (r \star r))
     return sz - 1;
    return (int) (lower_bound(in.begin(), in.end(), pi(x, y)) - in.begin());
 vector<vector<modint>> prob(sz, vector<modint>(sz, 0));
  // primeiro passo:
  // para cada nao terminal ii computamos prob[ii][j], probabilidade de ir de ii
       pra j em um move
  // todas as celulas com distancia > r, eu considero por um unico indice nessa
       representacao
  for (int ii = 0; ii < (sz - 1); ii++)
    auto [x, y] = in[ii];
   for (int dir = 0; dir < 4; dir++)</pre>
      int i = x + dx[dir];
     int j = y + dy[dir];
      int id = get_id(i, j);
     prob[ii][id] += a[dir];
  // agr quero calcular EV(i) para todo i
  // valor esperado de moves saindo da posicao i
  // monta o sisteminha de equações e dale
  vector<vector<modint>> eqs;
  for (int i = 0; i < sz; i++)
    vector<modint> curr(sz + 1, 0);
    if (i == sz - 1) // estado terminal, EV(i) = 0
      curr[sz - 1] = 1;
      eqs.pb(curr);
    else // nao terminal, EV(i) = 1 + soma dos (prob[i][j] * EV(j))
      for (int j = 0; j < sz; j++)
        if (i == i)
          curr[j] = modint(1) - prob[i][j];
          curr[j] = prob[i][j] * -1;
     curr[sz] = 1;
      eqs.pb(curr);
  gauss (eqs);
 for (int i = 0; i < in.size(); i++)</pre>
   if (in[i].fir == 0 && in[i].sec == 0) // ja que eu quero o valor esperado
        saindo da posicao (0,0)
      cout << ans[i].val << endl;</pre>
 return 0;
// solucao tle (porem correta) de um problema legal
// https://codeforces.com/contest/963/problem/E
// tem um chip inicialmente na posicao (0, 0)
// e ele vai comecar a se mover aleatoriamente
// tem a[i] / (a[1] + a[2] + a[3] + a[4]) para se mover para a direcao i
// ele pode ir pra (x + 1, y) (x, y + 1), (x - 1, y), (x, y - 1)
// quero saber o valor esperado ate ele checar em uma celula
// com distancia maior do que R do (0, 0)
```

### 14.20 matrix exponentiation

```
// https://codeforces.com/gym/102644/problem/C
// achar o n-esimo termo da sequencia de fibonacci mod (10^9 + 7) em O(log(n))
// n <= 10^18
// podemos escrever a recorrencia de fibonnaci como uma exponenciacao de matriz
    fib(n) ) (1 1) ^{\circ} (n - 1)
                                          (fib(1) = 1)
  (fib(n-1)) = (1 0)
                                       * (fib(0) = 1)
// e possivel fazer essa exponenciacao em O(log(n)) com um algoritimo muito
     similar ao de exponenciacao rapida
// dai calculamos o n-esimo termo da sequencia de fibonacci mod (10^9 + 7) em O(
    log(n))
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 201
#define mod 1000000007
namespace matrix
  vector<vector<int>> ans;
  int multi(int x, int y)
    return (x * y) % mod;
  int sum(int a, int b)
    return (a + b >= mod) ? a + b - mod : a + b;
  vector<vector<int>> multiply(vector<vector<int>> a, vector<vector<int>> b)
    vector<vector<int>> res(a[0].size(), vector<int>(b[0].size()));
    for (int i = 0; i < a.size(); i++)</pre>
      for (int j = 0; j < b[0].size(); j++)</pre>
        res[i][j] = 0;
        for (int k = 0; k < a[0].size(); k++)</pre>
          res[i][j] = sum(res[i][j], multi(a[i][k], b[k][j]));
    return res;
  vector<vector<int>> expo(vector<vector<int>> mat, int m)
    ans = vector<vector<int>>(mat.size(), vector<int>(mat[0].size()));
    for (int i = 0; i < mat.size(); i++)</pre>
      for (int j = 0; j < mat[0].size(); j++)</pre>
        ans[i][j] = (i == j);
    while (m > 0)
      if (m & 1)
       ans = multiply(ans, mat);
      m = m / 2;
     mat = multiply(mat, mat);
    return ans;
```

```
signed main()
{
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  vector<vector<int>> mat = {{1, 1}, {1, 0}};
  vector<vector<int>> ans = matrix::expo(mat, n);
  cout << ans[0][1] << endl;
  return 0;
}</pre>
```

### 14.21 matrix exponentiation2

```
// https://www.spoj.com/problems/ITRIX12E/
// count some \{f(0) + f(1) + \dots + f(n)\} with just one matrix exponentiation
// creates an extra dimension in the matrix and initializes that column with 1s
#include <bits/stdc++.h>
using namespace std;
#define PI acos (-1)
#define pb push_back
#define mp make_pair
#define int long long int
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 100001
#define MAXL 20
#define INF 200001
#define mod 1000000007
const int n = 11:
vector<vector<int>> ans(n, vector<int>(n));
vector<vector<int>> multiply(vector<vector<int>> a, vector<vector<int>> b)
  vector<vector<int>> res(n, vector<int>(n));
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++)
      res[i][j] = 0;
      for (int k = 0; k < n; k++)
        res[i][j] = (res[i][j] + (((a[i][k] % mod) * (b[k][j] % mod)) % mod)) %
  return res;
vector<vector<int>> expo(vector<vector<int>> mat, int m)
  for (int i = 0; i < n; i++)
   for (int j = 0; j < n; j++)
ans[i][j] = (i == j);</pre>
  while (m > 0)
   if (m & 1)
     ans = multiply(ans, mat);
   m = m / 2;
   mat = multiply(mat, mat);
  return ans;
bool is_prime(int n)
  for (int i = 2; i < n; i++)
   if (!(n % i))
      return false;
  return true;
signed main()
```

```
ios_base::sync_with_stdio(false);
cin.tie(NULL);
int q;
cin >> q;
while (q--)
  int k;
  cin >> k;
  int resp = 0;
  vector<vector<int>> mat(n, vector<int>(n, 0));
  for (int i = 1; i <= 9; i++)
  for (int j = 1; j <= 9; j++)</pre>
       if (is_prime(i + j))
          mat[i][j] = 1;
   for (int i = 0; i <= 10; i++)
    mat[i][10] = 1;
   vector<vector<int>> ans = expo(mat, k - 1);
  for (int i = 0; i < n; i++)
  for (int j = 0; j < n; j++)</pre>
  resp = (resp + ans[i][j]) % mod;
cout << resp - 7 << end1;</pre>
return 0;
```

#### 14.22 matrix inverse and determinant

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 998244353
struct modint
 int val:
 modint(int v = 0) { val = (v + mod) % mod; }
 int pow(int y)
    modint x = val;
    modint z = 1;
    while (y)
     if (y & 1)
       z *= x;
     x *= x;
     y >>= 1;
    return z.val;
 int inv() { return pow(mod - 2); }
  void operator=(int o) { val = o % mod; }
 void operator=(modint o) { val = o.val % mod; }
 void operator+=(modint o) { *this = *this + o; }
  void operator = (modint o) { *this = *this - o;
 void operator*=(modint o) { *this = *this * o;
  void operator/=(modint o) { *this = *this / o; }
 bool operator==(modint o) { return val == o.val;
 bool operator!=(modint o) { return val != o.val; }
  int operator*(modint o) { return ((val * o.val) % mod);
 int operator/(modint o) { return (val * o.inv()) % mod; }
  int operator+(modint o) { return (val + o.val) % mod; }
 int operator-(modint o) { return (val - o.val + mod) % mod; }
pair<int, modint> gauss(vector<vector<modint>> &a, int pivot_end)
 int n = a.size(), m = a[0].size();
 int rank = 0, curr = pivot_end;
 if (curr == -1)
```

```
curr = m;
  modint det = 1;
  for (int j = 0; j < curr; j++)</pre>
    int idx = -1;
    for (int i = rank; i < n; i++)</pre>
      if (a[i][j].val != 0)
        idx = i;
        break;
    if (idx == -1)
      det = 0:
      continue;
    if (rank != idx)
      det. *= -1:
      swap(a[rank], a[idx]);
    det *= a[rank][j];
    if (a[rank][j].val != 1)
      modint coef = a[rank][j].inv();
      for (int k = j; k < m; k++)</pre>
        a[rank][k] *= coef;
    int is = 0;
    for (int i = is; i < n; i++)</pre>
      if (i == rank)
        continue;
      if (a[i][j].val != 0)
        modint coef = a[i][j] / a[rank][j];
        for (int k = j; k < m; k++)
          a[i][k] -= a[rank][k] * coef;
   rank++;
  return {rank, det};
vector<vector<modint>> inverse_matrix(vector<vector<modint>> a)
  int n = a.size();
  vector<vector<modint>> m(n, vector<modint>(2 * n));
  for (int i = 0; i < n; i++)
    copy(begin(a[i]), end(a[i]), begin(m[i]));
   m[i][n + i] = 1;
  auto [rank, det] = gauss(m, n);
  if (rank != n)
   return {};
  vector<vector<modint>> b(n);
  for (int i = 0; i < n; i++)
   copy(begin(m[i]) + n, end(m[i]), back_inserter(b[i]));
  return b;
modint determinant(vector<vector<modint>> a)
  return gauss (a, -1).sec;
signed main()
  ios base::sync with stdio(false);
  cin.tie(NULL);
  cin >> n;
  vector<vector<modint>> v(n, vector<modint>(n));
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++)
      cin >> v[i][j].val;
```

```
if (determinant(v).val == 0)
{
   cout << "-1\n";
   return 0;
}
vector<vector<modint>> ans = inverse_matrix(v);
for (int i = 0; i < n; i++)
{
   for (int j = 0; j < n; j++)
      cout << ans[i][j].val << " ";
   cout << endl;
}
return 0;
}
// https://judge.yosupo.jp/problem/inverse_matrix
// https://judge.yosupo.jp/problem/matrix_det
// como precisa de divisao, entao o mod tem que ser primo</pre>
```

## 14.23 max xor subsequence

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 300005
#define mod 1000000007
int modpow(int a, int b)
  int res = 1;
  while (b > 0)
   if (b & 1)
     res = (res * a) % mod;
    a = (a * a) % mod;
   b >>= 1;
  return res;
int all, qt;
int dp[33];
void add(int x)
  all++;
  for (int i = 32; i >= 0; i--)
   if (x & (111 << i))</pre>
      if (dp[i] == 0)
        dp[i] = x;
        qt++;
        return;
        ^= dp[i];
int get(int x) // qual o x-esimo menor valor de xor de uma subsequencia
```

```
int tot = (111 << qt), ans = 0;</pre>
  for (int i = 32; i >= 0; i--)
    if (dp[i] > 0)
      int d = tot / 2;
      if (d < x && !(ans & (111 << i)))
       ans ^= dp[i];
      else if (d >= x \&\& (ans \& (111 << i)))
       ans ^= dp[i];
      if (d < x)
       x -= d;
      tot /= 2;
   }
  return ans:
bool check(int x) // se existe pelo menos uma subsequencia com xor x
  for (int i = 32; i >= 0; i--)
    if (x & (111 << i))</pre>
      if (!dp[i])
       return 0;
      x = dp[i];
  return 1;
int count(int x) // quantas subsequencias tem xor x
  if (!check(x))
   return 0;
  return modpow(2, all - qt);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  vector<int> v(n);
  for (int i = 0; i < n; i++)
    cin >> v[i];
   add(v[i]);
  int x = get(111 << gt); // maior xor possivel de uma subsequencia</pre>
  int y = get(1);
                           // maior xor possivel != 0 (o 0 sempre eh possivel -
      subsequencia vazia)
  return 0:
// referencia:
// https://codeforces.com/blog/entry/68953
// problemas:
// https://codeforces.com/gym/103708/problem/A
// https://codeforces.com/contest/959/problem/F
// https://codeforces.com/contest/1101/problem/G
// https://atcoder.jp/contests/abc283/tasks/abc283_g
```

## 14.24 mobius

```
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 5000005
#define mod 100000001
int lpf[MAXN];
int mobius[MAXN];
int g[MAXN];
void calc_lpf()
 for (int i = 2; i < MAXN; i++)</pre>
    if (!lpf[i])
      for (int j = i; j < MAXN; j += i)
        if (!lpf[j])
          lpf[j] = i;
   }
void calc_mobius()
 calc_lpf();
 mobius[1] = 1;
 for (int i = 2; i < MAXN; i++)</pre>
    if (lpf[i / lpf[i]] == lpf[i])
     mobius[i] = 0;
    else
     mobius[i] = -1 * mobius[i / lpf[i]];
int count_pairs(int n)
  // f(n) -> contar pares (i, j) com \_gcd(i, j) == 1 e 1 <= i, j <= n
 int ans = 0;
 for (int d = 1; d <= n; d++)</pre>
    // quadrado pq sao pares (2 caras)
    // mas se fossem x caras seria (n / d) x
   int sq = (n / d) * (n / d);
   int x = mobius[d] * sq;
   ans += x;
 return ans;
int gcd_sum(int n)
  // soma de todos os gcd(i, j) com 1 <= i, j <= n
 int ans = 0;
 for (int k = 1; k \le n; k++) // fixa o valor do gcd(i, j) e conta quantos
      pares com gcd(i, j) == k
    int lim = n / k;
   int curr = k * count_pairs(lim);
    ans += curr;
 return ans;
int lcm sum(int n)
  // soma de todos os lcm(i, j) com 1 <= i, j <= n
 for (int i = 1; i <= n; i++)
   q[i] = 0;
  for (int i = 1; i <= n; i++)
    for (int j = i; j \le n; j += i)
      g[j] += (mobius[i] * j * i);
 int ans = 0;
  for (int 1 = 1; 1 <= n; 1++)
   int cima = (1 + n / 1) * (n / 1);
```

```
int f = (cima / 2) * (cima / 2);
    f *= q[1];
    ans += f;
  return ans;
signed main()
 ios_base::sync_with_stdio(false);
  cin.tie(NULL);
 int q;
 cin >> q;
  calc mobius();
  for (int i = 1; i <= q; i++)
   int n;
    cin >> n;
    int ans = lcm_sum(n);
    for (int i = \overline{1}; i \le n; i++)
     ans -= i;
    ans /= 2;
    cout << "Case " << i << ": " << ans << endl;</pre>
  return 0;
// https://codeforces.com/blog/entry/53925
// mobius inversion
// sejam f(x) e g(x) funcoes
// e g(x) e definida da seguinte maneira
// g(x) = soma \ dos \ f(d), no qual d eh um divisor de x
// temos que:
// f(n) = soma dos (g(d) * u(n / d)), no qual d eh um divisor de x
// u(x) \rightarrow mobius function
// propiedade legal:
// seja l(x) -> soma de u(d), para cada divisor d de x
// 1(\bar{1}) = 1
// 1(x) = 0, x > 1
// problemas iniciais:
// https://vjudge.net/problem/AtCoder-abc162_e
// https://vjudge.net/problem/CodeChef-SMPLSUM
```

### 14.25 mobius2

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200001
#define mod 998244353
int lpf[MAXN];
int mobius[MAXN];
int mp[MAXN];
vector<int> d[MAXN];
void calc_lpf()
  for (int i = 2; i < MAXN; i++)
   if (!lpf[i])
      for (int j = i; j < MAXN; j += i)
        if (!lpf[j])
          lpf[j] = i;
```

```
void calc()
  for (int i = 2; i < MAXN; i++) // divisores</pre>
    for (int j = i; j < MAXN; j += i)
      d[j].pb(i);
  calc_lpf();
  mobius[1] = 1;
  for (int i = 2; i < MAXN; i++)
    if (lpf[i / lpf[i]] == lpf[i])
      mobius[i] = 0;
      mobius[i] = -1 * mobius[i / lpf[i]];
void add(int x, int dd) // adiciona dd em todos os val[i] que gcd(x, i) > 1
  for (auto const &i : d[x])
   mp[i] += dd;
int sum(int x) // valor de val[x]
  int ans = 0:
  for (auto const &i : d[x])
   ans += (mobius[i] * -1 * mp[i]);
  return ans;
signed main()
  return 0:
// mobius/inclusao-exclusao com os fatores primos
// a funcao de mobius eh definida como:
// mi(n) = 1, se n e um square-free com um numero par de fatores primos
// mi(n) = -1, se n e um square-free com um numero impar de fatores primos
// mi(n) = 0, caso nenhum dos dois
// square-free = nenhum fator primo aparece duas vezes ou mais
// dai pra problemas que da pra se fazer com inclusao-exclusao nos fatores
    primos
// tambem sai com mobius
```

### 14.26 modular arithmetic

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
struct modint
  modint(int v = 0) { val = v % mod; }
  int pow(int y)
    modint x = val;
    modint z = 1;
    while (y)
```

```
if (y & 1)
       z *= x;
      x *= x;
     y >>= 1;
    return z.val;
  int inv() { return pow(mod - 2); }
  void operator=(int o) { val = o % mod; }
  void operator=(modint o) { val = o.val % mod; }
  void operator+= (modint o) { *this = *this + o; }
  void operator-= (modint o) { *this = *this - o;
  void operator*=(modint o) { *this = *this * o;
  void operator/= (modint o)
                             { *this = *this / o;
  bool operator== (modint o) { return val == o.val;
  bool operator!=(modint o) { return val != o.val; }
  int operator*(modint o) { return ((val * o.val) % mod);
  int operator/(modint o) { return (val * o.inv()) % mod;
  int operator+(modint o) { return (val + o.val) % mod; }
  int operator-(modint o) { return (val - o.val + mod) % mod; }
};
modint f[MAXN];
modint inv[MAXN];
modint invfat[MAXN];
void calc()
  f[0] = 1;
  for (int i = 1; i < MAXN; i++)</pre>
    f[i] = f[i - 1] * i;
  for (int i = 2; i < MAXN; ++i)</pre>
    int val = mod / i;
    val = (inv[mod % i] * val) % mod;
    val = mod - val;
    inv[i] = val;
 invfat[0] = 1;
invfat[MAXN - 1] = modint(f[MAXN - 1]).inv();
  for (int i = MAXN - 2; i >= 1; i--)
    invfat[i] = invfat[i + 1] * (i + 1);
modint ncr(int n, int k) // combinacao
  modint ans = f[n] * invfat[k];
  ans *= invfat[n - k];
  return ans;
modint arr(int n, int k) // arranjo
  modint ans = f[n] * invfat[n - k];
  return ans;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
```

#### 14.27 ntt

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;

template <class T>
```

```
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 250005
#define mod 998244353
struct modint
 int val:
 modint(int v = 0) { val = ((v % mod) + mod) % mod; }
 int pow(int y)
   modint x = val;
   modint z = 1;
    while (y)
     if (y & 1)
      z *= x;
      x *= x;
     y >>= 1;
    return z.val;
  int inv() { return pow(mod - 2); }
 void operator=(int o) { val = o % mod;
 void operator=(modint o) { val = o.val % mod; }
  void operator+=(modint o) { *this = *this + o;
 void operator = (modint o) { *this = *this - o;
 void operator*=(modint o) { *this = *this * o;
  void operator/=(modint o) { *this = *this / o; }
 bool operator==(modint o) { return val == o.val;
 bool operator!=(modint o) { return val != o.val; }
 int operator*(modint o) { return ((val * o.val) % mod);
  int operator/(modint o) { return (val * o.inv()) % mod; }
 int operator+(modint o) { return (val + o.val) % mod; }
 int operator-(modint o) { return (val - o.val + mod) % mod; }
namespace fft
  // para o modulo ser valido
  // precisa ser primo
  // precisa possuir a forma c * 2^k + 1
  // 998244353 - possui a forma - c * 2^k + 1 e eh primo
 int root = -1;
 int root_1 = -1;
 int pw = __builtin_ctz(mod - 1);
 int root_pw = (1 << pw);</pre>
  void find_root()
   if (root != -1)
     return;
    while (! (modint(r).pow((1 << pw)) == 1 && modint(r).pow((1 << (pw - 1))) !=
        1))
     r++;
    root_1 = modint(root).inv();
  void ntt(vector<modint> &a, bool invert)
    find root();
    int \overline{n} = a.size();
    for (int i = 1, j = 0; i < n; i++)
      int bit = n >> 1;
      for (; j & bit; bit >>= 1)
       j ~= ~
^= bit;
          ^= bit;
      if (i < i)
        swap(a[i], a[j]);
```

```
for (int len = 2; len <= n; len <<= 1)</pre>
      modint wlen = (invert) ? root_1 : root;
      for (int i = len; i < root_pw; i <<= 1)</pre>
        wlen *= wlen;
      for (int i = 0; i < n; i += len)</pre>
        modint w = 1;
        for (int j = 0; j < len / 2; j++)
          modint u = a[i + j];
          modint v = a[i + j + len / 2] * w;
          a[i + j] = u + v;
          a[i + j + len / 2] = u - v;
          w \star = wlen:
    if (invert)
      modint n_1 = modint(n).inv();
      for (int i = 0; i < a.size(); i++)</pre>
        a[i] *= n_1;
  vector<modint> mul(vector<modint> a, vector<modint> b)
    while (n < 2 * max(a.size(), b.size()))
     n <<= 1;
    a.resize(n);
    b.resize(n):
    ntt(a, false);
    ntt(b, false);
    for (int i = 0; i < n; i++)
     a[i] \star = b[i];
    ntt(a, true);
    return a;
} // namespace fft
  // https://codeforces.com/contest/1613/problem/F
```

## 14.28 operadores binarios

```
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define in insert
#define pi pair<int, int>
#define pd pair<double, int>
#define pii pair<int, pi>
#define mp make_pair
#define fir first
#define sec second
#define MAXN 200001
#define mod 1000000007
void shifts ()
  bitset <4> bs;
  bs.reset();
  bs[2] = true;
  bs[3] = true;
  cout << bs << endl ; // 1100
  bs >>= 1; // 0110
  bs <<= 1; // 1100
  bs >>= 2; // 0011
  bs <<= 2; // 1100
  bs >>= 3; // 0001
  bs <<= 3; // 1000
  cout << bs << endl ;
```

```
void op_xor ()
  // 0 ^ 0 = 0
// 0 ^ 1 = 1
// 1 ^ 0 = 1
  // 1 ^ 1 = 0
  bitset <4> bs , bs2;
  bs.reset();
  bs2.reset();
  bs[2] = true;
  bs[3] = true;
  bs2[1] = true;
  bs2[3] = true;
 bs = bs2; //bs = bs \cdot bs2
  cout << bs.count() << endl ;</pre>
void op_and ()
  // 0 & 0 = 0
  // 0 & 1 = 0
  // 1 & 0 = 0
  // 1 & 1 = 1
  bitset <4> bs , bs2;
  bs.reset();
  bs2.reset();
  bs[2] = true;
  bs[3] = true;
  bs2[1] = true;
  bs2[3] = true;
  bs &= bs2; // bs = bs & bs2
  cout << bs.count() << endl ;</pre>
void op_or ()
  // 0 | 0 = 0
  // 0 | 1 = 1
  // 1 | 0 = 1
  // 1 | 1 = 1
 bitset <4> bs , bs2;
 bs.reset(); // poe tudo 0
 bs2.reset();
 bs[2] = true;
 bs[3] = true;
bs2[1] = true;
  bs2[3] = true;
 bs |= bs2; // bs = bs | bs2
  cout << bs.count() << endl ; // quantidade de 1</pre>
signed main()
 op_or();
  op_and();
 op_xor();
  shifts();
  return 0;
```

## 14.29 pollard rho

```
#define sec second
#define MAXN 1000001
#define mod 998244353
int read() // __int128 functions
  int x = 0, f = 1;
  char ch = getchar();
  while (ch < '0' || ch > '9')
   if (ch == '-')
     f = -1;
    ch = getchar();
  while (ch >= '0' && ch <= '9')
   x = x * 10 + ch - '0';
   ch = getchar();
  return x * f;
void print(__int128 x) // __int128 functions
  if (x < 0)
   cout << "-";
   x = -x;
  stack<char> s;
  while (x)
   s.push((x % 10) + '0');
   x = x / 10;
  while (!s.empty())
    cout << s.top();</pre>
   s.pop();
namespace pollard_rho
  int multiplicate(int x, int y, int m)
   return (x * y) % m;
  int modpow(int x, int y, int m)
    int z = 1;
   while (y)
     if (y & 1)
       z = (z * x) % m;
      x = (x * x) % m;
     y >>= 1;
   return z;
  bool is_composite(int n, int a, int d, int s)
    int x = modpow(a, d, n);
    if (x == 1 \text{ or } x == n - 1)
      return false;
    for (int r = 1; r < s; r++)
      x = multiplicate(x, x, n);
      if (x == n - 1LL)
        return false;
   return true;
  int miller_rabin(int n)
   if (n < 2)
     return false;
    int r = 0, d = n - 1LL;
    while ((d & 1LL) == 0)
```

```
d >>= 1;
     r++;
    for (int a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37})
     if (n == a)
       return true;
     if (is_composite(n, a, d, r))
       return false;
   return true;
 int f(int x, int m)
   return multiplicate(x, x, m) + 1;
 int rho(int n)
   int x0 = 1, t = 0, prd = 2;
   int x = 0, y = 0, q;
   while (t % 40 || __gcd(prd, n) == 1)
     if (x == y)
       x0++;
       x = x0:
       y = f(x, n);
     q = multiplicate(prd, max(x, y) - min(x, y), n);
     if (q != 0)
      prd = q;
     x = f(x, n);
     y = f(y, n);
     \dot{y} = f(y, n);
     t++;
   return __gcd(prd, n);
 vector<int> fact(int n)
   if (n == 1)
     return {};
   if (miller_rabin(n))
     return {n};
   int x = rho(n);
   auto l = fact(x), r = fact(n / x);
   l.insert(l.end(), r.begin(), r.end());
   return 1;
signed main()
 //ios_base::sync_with_stdio(false);
  //cin.tie(NULL);
 while (1)
   int n = read();
   if (n == 0)
     break;
   vector<int> factors = pollard_rho::fact(n);
    sort(factors.begin(), factors.end());
   int prev = -1, cnt = 0;
   for (auto const &i : factors)
     if (prev != i)
       if (prev != -1)
         print (prev);
         printf("^");
         print(cnt);
         printf(" ");
       prev = i;
       cnt = 0;
     cnt++;
```

```
if (prev != -1)
     print (prev);
     printf("^");
     print(cnt);
     printf(" ");
   printf("\n");
  return 0;
// https://github.com/PauloMiranda98/Competitive-Programming-Notebook/blob/
    master/code/math/prime.h
// https://github.com/brunomaletta/Biblioteca/blob/master/Codigo/Matematica/
    pollardrho.cpp
// fast integer factorization with pollard-rho
// https://www.spoj.com/problems/FACT0/ - ok
// https://www.spoj.com/problems/FACT1/ - ok
// https://www.spoj.com/problems/FACT2/ - sigkill
// since the limit is at most 29 digits(in FACT2), we need to use __int128
```

### 14.30 primefactors

```
#include <bits/stdc++.h>
using namespace std;
#define PI acos(-1)
#define int long long int
#define pb push_back
#define mp make_pair
#define pi pair<int, int>
#define fir first
#define sec second
#define MAXN 501
#define MAXL 20
#define mod 1000000007
vector<int> facts;
void primefactors(int n)
  while (n \% 2 == 0)
   facts.pb(2);
   n = n / 2;
  for (int i = 3; i \le sqrt(n); i += 2)
    while (n \% i == 0)
      facts.pb(i);
     n = n / i;
  if (n > 2)
   facts.pb(n);
signed main()
  int n;
  cin >> n;
  primefactors(n);
  sort(facts.begin(), facts.end());
  for (auto const &i : facts)
   cout << i << endl;</pre>
  return 0;
```

## 14.31 primefactors2

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
```

```
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 1000001
#define mod 1000000007
namespace primefactors
  bitset<MAXN> prime;
  vector<int> nxt (MAXN);
  vector<int> factors;
  void crivo()
    prime.set();
    prime[0] = false, prime[1] = false;
    for (int i = 2; i < MAXN; i++)</pre>
      if (prime[i])
        nxt[i] = i;
        for (int j = 2; j * i < MAXN; j++)
          prime[j * i] = false;
          nxt[j * i] = i;
  void fact(int n)
    factors.clear();
    while (n > 1)
      factors.pb(nxt[n]);
      n = n / nxt[n];
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
```

## 14.32 segmentedsieve

```
#define pii pair<int, pi>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 1000003
#define mod 1000000007
vector<int> prime;
void segmentedsieve(int 1, int r)
  int lim = sqrt(r);
  vector<bool> mark(lim + 1, false);
  vector<int> primes;
  for (int i = 2; i <= lim; ++i)</pre>
    if (!mark[i])
      primes.pb(i);
      for (int j = i * i; j <= lim; j += i)
        mark[j] = true;
  vector<bool> isprime(r - 1 + 1, true);
 for (int i : primes)
  for (int j = max(i * i, (1 + i - 1) / i * i); j <= r; j += i)</pre>
      isprime[j - 1] = false;
  if (1 == 1)
    isprime[0] = false;
  for (int i = 0; i < isprime.size(); i++)</pre>
    if (isprime[i])
      prime.pb(i + 1);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int 1, r;
  cin >> 1 >> r;
  segmentedsieve(1, r);
  for (auto const &i : prime)
   cout << i << " ";
  return 0;
```

## 14.33 simplex

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __qnu_pbds;
#define int long long int
#define pi pair<int, int>
#define fir first
#define sec second
#define mod 2147483647
#define pb push_back
#define double long double
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
const double eps = 1e-8;
const double inf = 1e18;
#define MP make_pair
#define ltj(X)
  if (s == -1 \mid | MP(X[j], N[j]) < MP(X[s], N[s])) \setminus
// resolve um problema de programacao linear para maximizar uma funcao c[0] * x
     [0] + c[1] * x[1] + ... <= b (c[i] eh o coeficiente do i-esimo cara na
     funcao objetiva)
```

```
// sujeito a restricoes que tem a seguinte forma:
// a[0] * x[0] + a[1] * x[1] + ... <= b, (a[i] eh o coeficiente)
// ai todas as restricoes sao passadas nos vectors a e b
// complexidade: 2^n, mas na pratica pode ser melhor do que isso
struct lp_solver
 int m, n;
 vector<int> N, B;
 vector<vector<double>> D;
 lp_solver(const vector<vector<double>> &A, const vector<double> &b, const
      vector < double > &c) : m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2,
      vector<double>(n + 2))
   for (int i = 0; i < m; i++)
     for (int j = 0; j < n; j++)
       D[i][j] = A[i][j];
    for (int i = 0; i < m; i++)
     B[i] = n + i;
     D[i][n] = -1;
     D[i][n + 1] = b[i];
   for (int j = 0; j < n; j++)
     N[j] = j;
     D[m][j] = -c[j];
   N[n] = -1;
   D[m + 1][n] = 1;
 void pivot(int r, int s)
   double *a = D[r].data(), inv = 1 / a[s];
   for (int i = 0; i < m + 2; i++)
     if (i != r && abs(D[i][s]) > eps)
       double *b = D[i].data(), inv2 = b[s] * inv;
       for (int j = 0; j < n + 2; j++)
         b[j] = a[j] * inv2;
       b[s] = a[s] * inv2;
   for (int j = 0; j < n + 2; j++)
     if (j != s)
       D[r][j] *= inv;
    for (int i = 0; i < m + 2; i++)
     if (i != r)
       D[i][s] *= -inv;
   D[r][s] = inv;
   swap(B[r], N[s]);
 bool simplex(int phase)
   int x = m + phase - 1;
   for (;;)
     int s = -1;
     for (int j = 0; j < n + 1; j++)
       if (N[j] != -phase)
         ltj(D[x]);
     if (D[x][s] \ge -eps)
      return true;
     int r = -1;
     for (int i = 0; i < m; i++)
       if (D[i][s] <= eps)
       ], B[r]))
         r = i;
     if (r == -1)
```

```
return false;
      pivot(r, s);
  double solve()
    int r = 0;
    for (int i = 1; i < m; i++)
      if (D[i][n + 1] < D[r][n + 1])
    if (D[r][n + 1] < -eps)
      pivot(r, n);
      if (!simplex(2) || D[m + 1][n + 1] < -eps)</pre>
        return -inf;
      for (int i = 0; i < m; i++)
        if (B[i] == -1)
          int s = 0;
          for (int j = 1; j < n + 1; j++)
           ltj(D[i]);
          pivot(i, s);
   bool ok = simplex(1);
    vector<double> x = vector<double>(n); // os valores escolhidos pra cada x[i]
          (se quiser eles tbm, so retornar)
    for (int i = 0; i < m; i++)
      if (B[i] < n)
        x[B[i]] = D[i][n + 1];
    return ok ? D[m][n + 1] : inf;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int k;
  cin >> k;
  vector < double > t(k), l(k), r(k);
  for (int i = 0; i < k; i++)
   cin >> t[i] >> 1[i] >> r[i];
  int q;
  cin >> q;
  while (q--)
   int aa, bb;
    cin >> aa >> bb;
    int p = aa * bb;
    vector<vector<double>> a;
    vector<double> b;
    vector<double> curr(k, 1);
    a.pb(curr);
    b.pb(bb);
    curr = vector<double>(k, -1);
    a.pb(curr);
    b.pb(-bb);
    a.pb(t);
    b.pb(p);
    curr = vector<double>(k, 0);
    for (int i = 0; i < k; i++)
      curr[i] = 1;
      a.pb(curr);
      b.pb(r[i]);
      curr[i] = 0;
    for (int i = 0; i < k; i++)
      curr[i] = -1;
      a.pb(curr);
```

b.pb(-1[i]);

```
curr[i] = 0;
}
int x = a.size();
lp_solver l(a, b, t);
int ans = round(l.solve());
if (ans == p)
    cout << "yes\n";
else
    cout << "no\n";
}
// solucao pro: https://open.kattis.com/problems/joiningflows
// source: https://github.com/kth-competitive-programming/kactl/blob/main/
    content/numerical/Simplex.h</pre>
// lembrete: quando eu quero adicionar algo com <= ao inves de >=, basta
    multiplicar os dois lados por -1 :)
// TODO: escrever melhor isso tudo depois
```

#### 14.34 stars and bars

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
struct modint
  int val:
  modint(int v = 0) { val = v % mod; }
  int pow(int y)
   modint x = val;
    modint z = 1;
    while (y)
      if (y & 1)
      z = z * x;
      x = x * x;
     y >>= 1;
    return z.val;
  int inv() { return pow(mod - 2); }
  int operator*(modint o) { return ((val * o.val) % mod); }
  int operator/(modint o) { return (val * o.inv()) % mod; }
  int operator+(modint o) { return (val + o.val) % mod; }
  int operator-(modint o) { return (val - o.val + mod) % mod; }
modint ncr(int n, int k)
  // calcular combinacao para n grande
  // nesse problema n <= 10^12
  // em O(k)
  modint num = 1;
  modint den = 1;
  for (int i = 0; i < k; i++)</pre>
   num = num * modint(n - i);
    den = den * modint(i + 1);
```

```
modint ans = num / den;
  return ans;
modint stars_and_bars(int n, int k)
  // para pares de inteiros n e k
  // enconte a quantidade de k-tuplas com soma == n
  // x1 + x2 + \dots + xk = n
  return ncr(n + k - 1, k - 1);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, s;
  cin >> n >> s;
  vector<int> v(n);
  for (int i = 0; i < n; i++)</pre>
   cin >> v[i];
  modint all = stars_and_bars(s, n);
  modint to_sub = 0;
  for (int mask = 1; mask < (1 << n); mask++)</pre>
    int sum = 0;
    for (int j = 0; j < n; j++)
      if (mask & (1 << j))</pre>
        sum += (v[j] + 1);
    if (sum <= s)
      modint curr = stars_and_bars(s - sum, n);
      to_sub = (__builtin_popcount(mask) % 2) ? to_sub + curr : to_sub - curr;
  all = all - to_sub;
  cout << all.val << endl;
  return 0:
// stars and bars
// dado dois inteiros positivos n e k
// conte o numero de k-tuplas (x1, x2, ..., xk) tal que
// x1 + x2 + ... + xk = n
// com x1, x2, ..., xk >= 0
// resposta = ncr(n + k - 1, k - 1)
// para k-tuplas com x1, x2, ..., xk > 0:
// resposta = ncr(n - 1, k - 1)
// problema exemplo:
// https://codeforces.com/contest/451/problem/E
// contar quantas k-tuplas com soma == n
// tal que: x[i] >= 0 e x[i] <= f[i]
// k \le 20
// solucao:
// conta tudo com stars and bars
// dai preciso subtrair todas as possibilidades invalidas (com pelo menos um i
    tal que x[i] > f[i])
// seja n(i) as possibilidades com x[i] > f[i]
// dai eu quero calcular a quantidade de elementos na uniao entre todos os n(i)
// dai da pra fzr usando a formulinha de uniao de conjuntos:
// n (A uniao B uniao C) = n (A) + n (B) + n (C) - n (A intersecao B) ... + n (A
     intersecao B intersecao C)
// itera por todos os 2^n subsets e calcula o que deve subtrair/somar com
    aqueles caras
```

#### 14.35 totient

```
#define MAXN 100000
int phi[MAXN];
void calc()
```

```
for (int i = 0; i < MAXN; i++)
   phi[i] = i;
  for (int i = 2; i < MAXN; i++)</pre>
   if (phi[i] == i)
      for (int j = i; j < MAXN; j += i)
        phi[j] -= phi[j] / i;
int calc_phi(int n)
 int ans = n;
 for (int i = 2; i * i <= n; i++)</pre>
   if (n % i == 0)
      while (n \% i == 0)
       n /= i;
     ans -= ans / i;
 if (n > 1)
   ans -= ans / n;
 return ans;
```

#### 14.36 xor trie

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
struct node
  int me, cnt, id;
  int down[2];
  node(int c = 0) : me(c)
    id = -1:
    fill(begin(down), end(down), -1);
struct trie xor
  vector<node> t;
  trie_xor()
    t.resize(1);
  void add(int n, int id)
    int v = 0;
    t[v].cnt++;
    for (int i = 30; i >= 0; i--)
```

int bit = (n & (1 << i)) ? 1 : 0;

```
if (t[v].down[bit] == -1)
        t[v].down[bit] = t.size();
        t.emplace_back(bit);
      v = t[v].down[bit];
      t[v].cnt++;
    t[v].id = id;
  void rem(int n, int id)
    int v = 0;
    t[v].cnt--;
    for (int i = 30; i >= 0; i--)
      int bit = (n & (1 << i)) ? 1 : 0;
      v = t[v].down[bit];
      t[v].cnt--;
  int gry(int n) // maximum xor with n
    if (t[0].cnt == 0) // no element
     return -1;
    int v = 0;
    for (int i = 30; i >= 0; i--)
      int bit = (n & (1 << i)) ? 0 : 1;
      int nxt = t[v].down[bit];
      if (nxt != -1 \&\& t[nxt].cnt > 0)
        v = nxt;
      else
        v = t[v].down[bit ^ 1];
    return t[v].id;
signed main()
// alguns problemas:
// https://codeforces.com/problemset/problem/706/D
// https://codeforces.com/contest/1625/problem/D
// https://codeforces.com/contest/888/problem/G
```

# 15 Miscellaneous

#### 15.1 bitmasks

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 998244353
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, mask;
```

```
vector<int> masks;
// quantidade de bits setados na mask
cout << __builtin_popcount(mask) << endl;</pre>
// para printar o valor do bit i
for (int i = 0; i < n; i++)
 cout << ((mask >> i) & 1) << " ";
cout << endl;
// quando eh necessario percorrer todas as submasks ate (1 << n)
// e fazer algo com todas as submasks dessa mask
// util em problemas de dp com mask por exemplo
for (int i = 0; i < n; i++)
  for (int j = 0; j < (1 << n); j++)
    if ((i) >> i & 1) == 0)
      //alguma coisa agui sabendo que a mask(j) eh uma submask de(j ^ 1 << i)
// para percorrer por todas as submasks de uma mask
for (int s = mask; s; s = (s - 1) & mask)
  // alguma coisa agui sabendo que s eh uma submask de mask
// quando eh necessario percorrer todas as submasks ate (1 << n)
// e fazer algo com todas as submasks dessa mask O(3^n)
// util em problemas de dp com mask por exemplo
for (int m = 0; m < (1 << n); m++)
  for (int s = m; s; s = (s - 1) \& m)
    // alguma coisa aqui sabendo que mask s eh uma submask de m
// comprimindo as masks de um vector baseada em uma mask qualquer
for (int i = 0; i < masks.size(); i++)</pre>
  int compressed = 0, curr_bit = 0;
  for (int j = 0; j < n; j++)
    if (!(mask & (1LL << j)))</pre>
      continue;
    if (masks[i] & (1LL << j))</pre>
      compressed |= (1LL << curr_bit);</pre>
    curr_bit++;
  // alguma coisa sabendo que a mask compressed eh a mask comprimida da mask
       atual
return 0;
```

## 15.2 coordinate compression

```
#define fir first
#define sec second
#define MAXN 500005
#define mod 1000000007

void compress(vector<int> &v)
{
   vector<int> val;
   for (auto const &i : v)
      val.pb(i);
   sort(val.begin(), val.end());
   val.erase(unique(val.begin(), val.end()), val.end());
   for (auto &i : v)
      i = lower_bound(val.begin(), val.end(), i) - val.begin();
}
```

#### 15.3 inversion count

```
// seja S = a1, a2, ..., an
// uma inversao S e um par (i,j) com i < j e ai > aj
// Solucao O(n2) nao ideal:
//for(int i=0;i<n;i++)
      for (int j=i+1; j<n; j++)
                if(v[i]>v[j]) ans++;
// Em vez de trabalharmos com o vetor inteiro(n2), vamos dividir o vetor ao meio
     e trabalhar com suas metades,
// que chamaremos de u1 e u2.
// Queremos saber o valor de inv, o numero de inversoes em v. Ha tres tipos de
    inversoes (i, j) (i, j) em v:
// aquelas em que i e j estao ambos em ul, aquelas em que i e j estao ambos em
    u2 e aquelas
// em que i esta em u1 e j esta em u2.
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define pb push_back
#define in insert
#define pi pair<int, int>
#define pii pair<int, pi>
#define mp make_pair
#define fir first
#define sec second
#define MAXN 100001
#define INF 1000000000
int merge_sort(vector<int> &v)
  int ans = 0;
  if (v.size() == 1)
    return 0:
  vector<int> u1, u2;
  for (int i = 0; i < v.size() / 2; i++)</pre>
    u1.pb(v[i]);
  for (int i = v.size() / 2; i < v.size(); i++)</pre>
   u2.pb(v[i]);
  ans += merge_sort(u1);
  ans += merge_sort(u2);
  u1.pb(INF);
  u2.pb(INF);
  int ini1 = 0, ini2 = 0;
```

```
for (int i = 0; i < v.size(); i++)</pre>
    if (u1[ini1] <= u2[ini2])</pre>
      v[i] = u1[ini1];
      ini1++;
    else
      v[i] = u2[ini2];
      ini2++;
      ans += u1.size() - ini1 - 1;
 return ans;
signed main()
 int n;
 cin >> n;
 vector<int> v(n);
 for (int i = 0; i < n; i++)
   cin >> v[i];
 cout << merge_sort(v) << endl;</pre>
 return 0;
```

## 15.4 max plus convolution

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 1000000007
bool needs[MAXN];
int dp0[MAXN];
multiset<int> dp[MAXN];
struct dsu
  vector<int> parent;
  vector<int> sz;
 vector<int> sum;
  dsu(int n)
    parent.resize(n);
    sz.resize(n);
    sum.assign(n, 0);
    tot = n;
    for (int i = 0; i < n; i++)</pre>
      parent[i] = i;
      sz[i] = 1;
      sum[i] = needs[i];
  int find_set(int i)
```

```
return parent[i] = (parent[i] == i) ? i : find_set(parent[i]);
  int make set(int x, int y)
   x = find_set(x), y = find_set(y);
   if (x != y)
     if (sz[x] > sz[y])
     swap(x, y);
parent[x] = y;
      sz[y] += sz[x];
     sum[y] += sum[x];
     tot--;
   return y;
};
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
  int q;
  cin >> q;
 while (q--)
   int n, m, p;
cin >> n >> m >> p;
    for (int i = 0; i < n; i++)
     needs[i] = 0;
     dp[i].clear();
    for (int i = 0; i < p; i++)
      int x;
     cin >> x;
     x--;
     needs[x] = 1;
    vector<pii> edges;
    for (int i = 0; i < m; i++)
      int u, v, w;
     cin >> u >> v >> w;
     u--, v--;
      edges.pb({w, {u, v}});
    for (int i = 0; i < n; i++)
      dp[i].insert(0);
     dp0[i] = 0;
    sort(edges.begin(), edges.end());
    for (auto const &ii : edges)
      int w = ii.fir;
      auto [u, v] = ii.sec;
      if (d.find_set(u) != d.find_set(v))
       u = d.find_set(u);
        v = d.find_set(v);
        if (d.sz[u] > d.sz[v])
          swap(u, v);
        // problema brabo e educational
        // o que funcionava ate a hard version era algo tipo isso
        // for (int i = 0; i <= szu; i++)
            for (int j = 0; j \le szv; j++)
               new_dp[i + j] = min(new_dp[i + j], dp[u][i] + dp[v][j]);
        // isso eh uma (min, +) convolution
        // e ambas as sequencias satisfazem:
        // dp[i + 1] - dp[i] \le dp[i] - dp[i - 1]
        // dai pra fazer a ideia desse blog, descrita em (max, +) convolution
        // https://codeforces.com/blog/entry/98663
```

```
// que eh basicamente ir "mergindo" as slopes de forma gulosa
        // entretanto a unica coisa que muda eh que em cada iteracao desse
        // temos que setar o valor de dp[0] para algo diferente antes de fzr a (
            min, +) convolution
        // ai tem que refletir isso no multiset das diferencas
         int ini = *dp[u].begin();
         dp[u].erase(dp[u].find(ini));
         int new_dp0 = d.sum[u] * w;
         // se antes diferenca era dp[1] - dp[0]
         // eu aumentei o dp[0] por x
         // agr vai ser (dp[1] - dp[0]) - x
         int x = new_dp0 - dp0[u];
         dp[u].insert(ini - x);
         dp0[u] = new_dp0;
         int ini = *dp[v].begin();
         dp[v].erase(dp[v].find(ini));
         int new_dp0 = d.sum[v] * w;
         // se antes diferenca era dp[1] - dp[0]
         // eu aumentei o dp[0] por x
         // agr vai ser (dp[1] - dp[0]) - x
         int x = \text{new\_dp0} - \text{dp0[v]};
         dp[v].insert(ini - x);
         dp0[v] = new_dp0;
       for (auto const &x : dp[u])
         dp[v].insert(x);
       dp0[v] += dp0[u];
       dp[u].clear();
       d.make_set(u, v);
   int par = d.find_set(0);
   int x = dp0[par];
    // como eh uma min+ convolution, vai da menor slope pra maior slope
   for (auto const &val : dp[par])
     x += val;
     cout << x << " ";
   cout << endl;
 return 0;
// https://codeforces.com/contest/2021/problem/E3
```

#### 15.5 meetinthemiddle

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 1000001
int n, t;
vector<int> v;
vector<int> a;
vector<int> b;
```

```
void solve2(int i, int j, int k)
  if (i == i)
   b.pb(k);
   return;
  solve2(i + 1, j, k);
  solve2(i + 1, j, k + v[i]);
void solve(int i, int j, int k)
  if (i == j)
   a.pb(k);
    return;
  solve(i + 1, j, k);
  solve(i + 1, j, k + v[i]);
int upper(int 1, int r, int x)
  while (1 < r)
   int mid = (1 + r + 1) >> 1;
    (b[mid] \le x) ? 1 = mid : r = mid - 1;
  return b[1];
int meetinthemiddle()
  solve(0, (n >> 1) + 1, 0);
  solve2((n >> 1) + 1, n, 0);
  sort(b.begin(), b.end());
  int ans = 0;
  for (auto const &i : a)
   if(i>t)
     continue;
    ans = max(ans, i);
    int kappa = i + upper(0, b.size() - 1, t - i);
    if (kappa <= t)</pre>
      ans = max(ans, kappa);
  return ans;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> t;
  v.resize(n);
  for (int i = 0; i < n; i++)</pre>
   cin >> v[i];
  cout << meetinthemiddle() << endl;</pre>
  return 0;
```

# 15.6 prefix sum 2d

```
// https://cses.fi/problemset/task/1652
#include <bits/stdc++.h>
using namespace std;

#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define sec second
#define MAXN 200005
#define mod 100000007
```

```
int v[1001][1001];
int p[1001][1001];
int qry(int x1, int y1, int x2, int y2)
  return p[x2 + 1][y2 + 1] - p[x2 + 1][y1] - p[x1][y2 + 1] + p[x1][y1];
signed main()
 int n, q;
cin >> n >> q;
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++)
      char c:
      cin >> c;
      v[i][j] = (c == '*');
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++)
      p[i + 1][j + 1] = p[i][j + 1] + p[i + 1][j] - p[i][j];

p[i + 1][j + 1] += v[i][j];
  while (q--)
    int a, b, c, d;
    cin >> a >> b >> c >> d;
    a--, b--, c--, d--;
    cout << qry(a, b, c, d) << endl;
  return 0;
// prefix sum 2d
// me enrolo pra codar toda vez, e bom deixar na lib
```

## 15.7 rectangle union

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200007
vector<int> x_vals;
struct segtree
  vector<int> seg, tag;
  segtree()
    seg.assign(8 * x_vals.size(), 0);
    tag.assign(8 * x_vals.size(), 0);
  void add(int ql, int qr, int x, int v, int l, int r)
    if (qr <= l || r <= ql)</pre>
      return;
```

```
if (ql <= l && r <= qr)
      tag[v] += x;
      if (tag[v] == 0)
        if (1 != r)
          seg[v] = seg[v << 1] + seg[(v << 1) | 1];
        else
          seg[v] = 0;
      else
        seg[v] = x_vals[r] - x_vals[1];
    else
      int mid = (1 + r) >> 1;
      add(q1, qr, x, (v << 1), 1, mid);
      add(q1, qr, x, ((v << 1) | 1), mid, r);
      if (tag[v] == 0 && 1 != r)
        seg[v] = seg[v << 1] + seg[(v << 1) | 1];
  int qry()
   return seq[1];
  void upd(int 1, int r, int x)
    add(1, r, x, 1, 0, x_vals.size());
struct rect
  int x1, y1, x2, y2;
};
struct event
  int time, 1, r, type;
  bool operator < (const event &b)
    if (time != b.time)
     return time < b.time;</pre>
    return type > b.type;
};
const int inf = 1e9;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  vector<rect> v(n);
  for (int i = 0; i < n; i++)</pre>
   cin >> v[i].x1 >> v[i].y1 >> v[i].x2 >> v[i].y2;
    x_vals.pb(v[i].x1);
   x_{vals.pb}(v[i].x2);
  // comprime o x
  sort(x_vals.begin(), x_vals.end());
  x_vals.erase(unique(x_vals.begin(), x_vals.end()), x_vals.end());
  vector<event> ev;
  for (int i = 0; i < n; i++)
    v[i].x1 = lower_bound(x_vals.begin(), x_vals.end(), v[i].x1) - x_vals.begin
        ();
    v[i].x2 = lower_bound(x_vals.begin(), x_vals.end(), v[i].x2) - x_vals.begin
    ev.pb({v[i].y1, v[i].x1, v[i].x2, 0}); // adicao
   ev.pb({v[i].y2, v[i].x1, v[i].x2, 1}); // remocao
  segtree s;
```

```
sort(ev.begin(), ev.end());
  int area = 0, 1 = -inf;
  for (auto const &i : ev)
    if (1 == -inf)
      1 = i.time;
      s.upd(i.1, i.r, 1);
    else if (i.type == 1)
      int curr = s.qry();
      s.upd(i.l, i.r, -1);
      if (s.qry() != curr)
        int new_t = (s.qry() == 0) ? -inf : i.time;
        int lo = 1, hi = i.time - 1;
        area += ((hi - lo + 1) * curr);
        1 = new_t;
    else
      int curr = s.qry();
      s.upd(i.1, i.r, 1);
if (s.qry() != curr)
        int lo = 1, hi = i.time - 1;
        area += ((hi - lo + 1) * curr);
        l = i.time;
 cout << area << endl;</pre>
 return 0;
// area da uniao de retangulos
// comprime coordenada no x pra montar a segtree dos valores de x
// faz o line sweep pelo y
// testado em dois judges:
// https://cses.fi/problemset/task/1741/
// n <= 10^5
// -10^6 \le x, y \le 10^6
// https://judge.yosupo.jp/problem/area_of_union_of_rectangles
// n <= 5 * 10^5
// 0 <= x, y <= 10^9
```

## 15.8 segment covering

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, pi>
#define fir first
#define sec second
#define MAXN 500005
#define mod 1000000007
int st[MAXN + 1][21];
void naive(vector<pi> &v) // guloso do point covering
```

```
// os segmentos precisam estar ordenados pelo .second
  int n = v.size(), last = -1;
  vector<int> ans;
  for (int i = 0; i < n; i++)
    if (v[i].fir > last)
      ans.pb(v[i].sec);
      last = v[i].sec;
bool can(int 1, int r, int x)
  for (int i = 20; i >= 0; i--)
    if (x & (1 << i))
      1 = st[1][i];
  return 1 > r;
void solve(vector<pi> &v, int a, int b) // segment covering com binary lifting (
    da pra fazer point covering de forma bem similar)
  for (int i = 0; i <= MAXN; i++)</pre>
    st[i][0] = i;
  for (auto const &i : v)
    st[i.fir][0] = max(st[i.fir][0], i.sec + 1);
  for (int i = 1; i <= MAXN; i++) // se um segmento com 1 menor tem um r maior
    st[i][0] = max(st[i][0], st[i - 1][0]);
  for (int i = 1; i < 21; i++)
    for (int v = 0; v \le MAXN; v++)
      st[v][i] = st[st[v][i - 1]][i - 1];
  int lo = 1, hi = v.size();
  while (lo < hi) // busca binaria na resposta</pre>
    int mid = (lo + hi) >> 1;
    (can(a, b, mid)) ? hi = mid : lo = mid + 1;
  if (can(a, b, lo))
   cout << lo << endl;
  else
    cout << "-1\n";
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
// um tipo de problema que eu achei bem legal
// https://codeforces.com/problemset/problem/1175/E
// https://codeforces.com/gym/101221 (Problem K)
// https://vjudge.net/contest/512192#problem/E
// problema exemplo:
// dado um conjunto de n segmentos [1[i], r[i]]
// qual o numero minimo de pontos que vc pode escolher, tal que:
// para cada segmento [1[i], r[i]], pelo menos um ponto escolhido ta nesse
    segmento
// tem a solucao gulosa em O(N)
// mas que da pra ser otimizada com binary lifting/sparse table (em caso de ter
    varias queries sobre o conjunto de segmentos)
// depois de adicionar um ponto a resposta, acho o nxt dele: o proximo ponto que
     irei colocar na resposta depois de adicionar ele
// outro problema exemplo:
// dado um conjunto de n segmentos [1[i], r[i]]
// voce quer selecionar o numero minimo de segmentos do conjunto
// para cobrir todo o segmento [a, b]
// bem parecido, tem uma solucao gulosa normal
```

### 15.9 sprague grundy

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500009
#define mod 1000000001
vector < int > v = \{2, 3, 4, 5, 6\};
unordered_map<int, bool> vis;
unordered_map<int, int> dp;
int q(int x) // achar o grundy number na marra
    if (x == 0)
        return 0;
    vector<bool> ok(4, 0);
    int mex = 0;
    for (auto const &i : v)
        int curr = g(x / i);
        if (curr < 4)
            ok[curr] = 1;
        while (ok[mex])
            mex++;
    vis[x] = 1;
    return dp[x] = mex;
int solve(int x) // padraozin
    vector<int> ini = {0, 1, 2, 2, 3, 3, 0, 0, 0, 0, 0, 0};
    while (x >= 12)
       x /= 12;
    return ini[x];
signed main()
    ios_base::sync_with_stdio(false);
    cin.tie(NULL);
    int q;
    cin >> q;
    while (q--)
        int n;
        cin >> n;
        int x = 0;
        for (int i = 0; i < n; i++)
            int k:
            cin >> k;
            x = solve(k);
        (x > 0) ? cout << "Henry\n" : cout << "Derek\n";
    return 0;
game theory (um exemplo simples de problema pra ficar no repo)
```

```
- pro nim classico
- existem n pilhas cada uma possui x[i] blocos
- em uma play posso escolher uma pilha e tirar uma quantidade qualquer de blocos
     dela
- quem ganha?
- o jogador que comeca ganha se o xor dos tamanhos das pilhas for != 0
- teorema sprague-grundy (transformar um jogo qualquer em nim)
- seja v um estado que eu tou do jogo, podemos calcular o grundy number desse
     estado
- seja o conjuntos de estados adjacentes a v {u1, u2, ..., un}
-g(v) = mex(g(u1), g(u2), ..., g(un))
- se v nao tem nenhum extado adjacente, entao g(v) = 0
- q(v) -> grundy number do estado v
- com isso se tivemos varios estados iniciais (varias pilhas)
- podemos simplesmente achar o grundy number de cada um deles e depois saber
     quem ganha
- pelo valor do xor dos grundy numbers
- exemplo: floor division game
- existem n numeros e em uma play posso escolher um deles e dividir por 2, 3, 4,
     5 ou 6
- quem ganha?
- achar o grundy number de cada um dos n numeros
- se o xor for != 0, ganha quem comeca jogando
- caso contrario, o outro jogador ganha
- as vzs e util tbm ver se existe um padrao (em caso de altas constantes)
- notando o padrao, da pra achar o grundy number de forma mais eficiente e
     resolver o problema
```

#### 15.10 stack trick

```
#include <bits/stdc++.h>
using namespace std;
#define PI acos(-1)
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define fir first
#define sec second
#define MAXN 300001
#define mod 1000000007
int n;
vector<int> v;
vector<int> ans;
void solve()
  stack<pi> s;
  for (int i = n - 1; i >= 0; i--)
    while (!s.empty() && s.top().fir <= v[i])</pre>
     s.pop();
    (!s.empty()) ? ans[i] = s.top().sec : ans[i] = -1;
   s.push({v[i], i});
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n:
  v.resize(n);
  ans.resize(n);
  for (int i = 0; i < n; i++)</pre>
   cin >> v[i];
  solve():
  for (auto const &i : ans)
   cout << i << " ";
  cout << endl;
```

```
// WITHOUT SEGMENT TREE // for each index (0 <= i < n), find another index (0 <= j < n) // which v(j) > v[i] and j > i and j is as close as possible to i. // if this index does not exist, print -1 /* 5  
1 3 3 4 5  
*/ /*  
1 3 3 4 -1  
*/
```

#### 15.11 sum hash

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 300005
#define mod 998244353
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  mt19937_64 rnq(chrono::steady_clock::now().time_since_epoch().count());
  int n, k;
  cin >> n >> k;
  vector<int> h(k + 1, 0);
  for (int i = 1; i < k; i++)
   h[i] = rng();
   h[k] = h[i];
  vector<int> v(n);
  int sum = 0, ans = 0;
  map<int, int> mp;
  mp[0] = 0;
  for (int i = 0; i < n; i++)</pre>
    cin >> v[i];
    sum += h[v[i]];
    if (mp.find(sum) != mp.end())
     ans = max(ans, i - mp[sum] + 1);
     mp[sum] = i + 1;
  cout << ans << endl;
// solucao pra C da final brasileira da maratona de 2023
// dado um array com n inteiros, cada a[i] ta entre 1 e k
// qual o maior tamanho de um subarray no qual todos os numeros de 1 ate k
// tem a mesma frequencia nesse subarray
```

## 15.12 tower of hanoi

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
```

```
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
// #define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 300005
#define mod 998244353
vector<pair<char, char>> ans;
void solve(int n, char a, char b, char c)
  if (n == 0)
   return;
  solve(n - 1, a, c, b);
  ans.pb({a, b});
solve(n - 1, c, b, a);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, k;
  cin >> n >> k;
  solve(n, 'A', 'C', 'B');
  if (ans.size() > k)
    cout << "N\n";
    return 0;
  cout << "Y\n";</pre>
  k -= ans.size();
  if (k \% 2 == 0)
    for (int i = 0; i < (k / 2); i++)
      cout << "A B\n";
      cout << "B A\n";
  else
    for (int i = 0; i < (k / 2) - 1; i++)
      cout << "A B\n";
     cout << "B A\n";
    cout << "A B\n";
    cout << "B C\n";
   cout << "C A\n";</pre>
  for (auto const &i : ans)
   cout << i.fir << " " << i.sec << endl;
  return 0;
// torre de hanoi
// 3 pilhas, sendo uma pilha com n discos e as outras duas pilhas vazias
// em cada movimento, vc tira o disco do topo de uma pilha e poe no topo de
     outra pilha
// desde que o raio do disco seja menor do que o raio do disco que ta no topo da
      outra pilha
// os n discos tem raios distintos aos pares
// fazer com que todos os discos vao parar em outra pilha
// https://codeforces.com/gym/101879/problem/I
// resolver a torre de hanoi com k movimentos
// se for possivel resolver, printar os movimentos feitos
// numero minimo pra resolver pros primeiros n
// 1, 3, 7, 15, 31, 63, 127, 255
```

#### 15.13 two pointers

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 3005
#define mod 1000000007
const int inf = LLONG_MAX;
stack<pii> s[2];
void add(int x, int i)
  int mn = inf, mx = -inf;
  if (!s[i].empty())
   mn = min(mn, s[i].top().sec.fir);
    mx = max(mx, s[i].top().sec.sec);
  mn = min(mn, x);
  mx = max(mx, x);
  s[i].push({x, {mn, mx}});
void change()
  while (!s[1].empty())
    int x = s[1].top().fir;
   s[1].pop();
    add(x, 0);
void rem()
  if (!s[0].size())
   change();
  s[0].pop();
int q()
  int mn = inf, mx = -inf;
  for (int i = 0; i < 2; i++)
    if (!s[i].empty())
      mn = min(mn, s[i].top().sec.fir);
      mx = max(mx, s[i].top().sec.sec);
  if (mn == inf)
   return 0:
  return mx - mn;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, k;
  cin >> n >> k;
```

```
vector<int> v(n);
  for (int i = 0; i < n; i++)
   cin >> v[i];
  int ans = 0, i = 0;
  for (int j = 0; j < n; j++)
    add(v[j], 1);
    while (q() > k)
      rem();
     <u>i</u>++;
   ans += (j - i + 1);
  cout << ans << endl;
 return 0;
// https://codeforces.com/edu/course/2/lesson/9/2/practice/contest/307093/
    problem/F
// Given an array of n integers, Let's say that a segment of this array is good
// if the difference between the maximum and minimum elements on this segment is
// Your task is to find the number of different good segments
// amazing trick using stack
```

## 16 STL

#### 16.1 ordered set

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __qnu_pbds;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG false
#define MAXN 200002
template <class T> // template do ordered set
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  ordered_set<int> s; // ordered_set
  s.insert(1);
  s.insert(1):
  s.insert(2);
  s.insert(4);
  s.insert(3);
  for (auto const &i : s) // nao adiciona elementos repetidos, que nem o set
      normal
    cout << i << " ";
  cout << endl;
  cout << *(s.find_by_order(0)) << endl; // iterator do elemento 0</pre>
  cout << *(s.find_by_order(1)) << endl; // iterator do elemento 1
cout << s.order_of_key(4) << endl; // quantidade de elementos que sao</pre>
      menores do que 4
  cout << s.order_of_key(6) << endl;</pre>
                                            // quantidade de elementos que sao
       menores do que 4
// find_by_order : O(log n), retorna (um iterator) qual o k-esimo elemento do
// order_of_key: O(log n), retorna qual a quantidade de elementos menores do que
      x no set
```

#### 16.2 STL

```
1) Vector
vector <int> v; //Criacao o vector
v.push_back(10); //Adiciono o elemento 10 no final do vector v
v.size() // retorna o tamanho do vector
v.resize(10); //Muda o tamanho do vector v para 10.
v.pop_back(); //Apaga o ultimo elemento do vector V.
v.clear(); // apaga todos os elementos do vector v .
sort(v.begin(), v.end()); //Ordena todo o vector v
2)Pair
pair <string, int> p; // criando a pair relacionando um first com um second
p.first = "Joao"; // adicionando elementos
p.second = 8 ; //adicionando elementos
// utilidade: vector de pair
vector< pair <int, string> > v; // criando o vector v de pair
v.push_back(make_pair(a,b)); // dando push back em uma pair no vector usando
sort(v.begin(), v.end()); // tambem e possivel ordenar o vector de pair
3) Queue / FIla
queue <int> f; // criando a queque
f.push(10); // adiciona alguem na fila
f.pop(); // remove o elemento que esta na frente da fila
f.front(); // olha qual o elemento esta na frete da fila
f.empty() // retorna true se a fila estiver vazia e false se nao estiver vazia
4) Stack / Pilha
stack <int> p ; // criando a stack
pilha.push(x); //Adiciona o elemento x no topo da pilha
pilha.pop(); //Remove elemento do topo da pilha
pilha.top(); // retorna o elemento do topo da pilha
pilha.empty(); // verifica se a pilha esta vazia ou nao
5) Set
set <int> s ; // criando a set
// obs: a set nao adiciona elementos repetidos
s.insert(10); //Adiciona o elemento 10 no set
s.find(10) // Para realizar uma busca no set utilizamos o comando find,
o find retorna um ponteiro que aponta para o elemento procurado caso o elemento
     esteja no set ou para o final do set, caso o elemento procurado nao esteja
     no set , em complexidade O(log n)
if(s.find(10) != s.end()) // procurando pelo 10, se ele estiver no set
s.erase(10); //Apaga o elemento 10 do set em O(log n)
s.clear(); // Apaga todos os elementos
s.size(); // Retorna a quantidade de elementos
s.begin(); // Retorna um ponteiro para o inicio do set
s.end(); // Retorna um ponteiro para o final do set
6)Map
map <string, int> m; //Cria uma variavel do tipo map que mapeia strings em int
// Em um map cada elemento esta diretamente ligado a um valor, ou seja, cada
     elemento armazenado no map possui um valor correspondente
// Se tivermos um map de strings em inteiros e inserimos os pair ("Joao", 1), ("
     Alana", 10), ("Rodrigo", 9)
// Caso facamos uma busca pela chave "Alana" receberemos o numero 10 como
     retorno.
m.insert(make_pair("Alana", 10)); //Inserimos uma variavel do tipo pair
     diretamente no map, O(log n)
```

```
M["Alana"] = 10; //Relacionando o valor 10 a chave "Alana"
if(m.find("Alana") != m.end()){ //Se a chave "Alana" foi inserida no map
cout << m["Alana"] << endl; //Imprime o valor correspondente a chave "Alana", no</pre>
     caso, o valor 10.
m.erase("Alana"); //Apaga o elemento que possui a chave "Alana" do map
m.clear(); // Apaga todos os elementos
m.size(); // Retorna a quantidade de elementos
m.begin(); // Retorna um ponteiro para o inicio do map
m.end(); // Retorna um ponteiro para o final do map
7) Priority Queue
priority_queue <int> q; // declarando a priority queue
// Para utilizar a priority_queue do C++ e importante apenas saber que o maior
     elemento sempre estara na primeiro posicao.
// Com execao disso, todos os outros metodos sao semelhantes ao uso de uma queue
     comum, porem para manter a estrutura organizada, a complexidade da
     operacao de insercao e O(logn).
p.push(i) // adiciono o elemento i na priority_queue
p.pop(); // apago o primeiro da fila
p.top(); // vejo quem esta no topo
```

# 17 Strings

#### 17.1 aho corasick

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 5001
#define mod 1000000007
namespace aho
  int go(int v, char ch);
  const int K = 26; // tamanho do alfabeto
  struct trie
    char me;
                        // char correspondente ao no atual
    int go[K];
                        // proximo vertice que eu devo ir estando em um estado (
        V. C)
    int down[K];
                        // proximo vertice da trie
    int is_leaf = 0;
                        // se o vertice atual da trie eh uma folha (fim de uma
        ou mais strings)
                        // no ancestral do no atual
    int parent = -1;
    int link = -1;
                        // link de sufixo do no atual (outro no com o maior
        matching de sufixo)
    int exit_link = -1; // folha mais proxima que pode ser alcancada a partir de
         v usando links de sufixo
    trie(int p = -1, char ch = '\$') : parent(p), me(ch)
      fill(begin(go), end(go), -1);
      fill(begin(down), end(down), -1);
  vector<trie> ac;
  void init() // criar a raiz da trie
    ac.resize(1);
```

```
void add_string(string s) // adicionar string na trie
   int v = 0;
   for (auto const &ch : s)
     int c = ch - 'a';
     if (ac[v].down[c] == -1)
       ac[v].down[c] = ac.size();
       ac.emplace_back(v, ch);
     v = ac[v].down[c];
   ac[v].is_leaf++;
 int get_link(int v) // pegar o suffix link saindo de v
   if (ac[v].link == -1)
     ac[v].link = (!v || !ac[v].parent) ? 0 : go(get_link(ac[v].parent), ac[v].
          me);
   return ac[v].link;
 int qo(int v, char ch) // proximo estado saindo do estado(v, ch)
   int c = ch - 'a';
   if (ac[v].go[c] == -1)
     if (ac[v].down[c] != -1)
       ac[v].go[c] = ac[v].down[c];
     else
       ac[v].go[c] = (!v) ? 0 : go(get\_link(v), ch);
   return ac[v].go[c];
 int get_exit_link(int v) // suffix link mais proximo de v que seja uma folha
   if (ac[v].exit_link == -1)
     int curr = get_link(v);
     if (!v || !curr)
       ac[v].exit_link = 0;
     else if (ac[curr].is_leaf)
       ac[v].exit_link = curr;
     else
       ac[v].exit_link = get_exit_link(curr);
   return ac[v].exit_link;
 int query(string s) // query O(n + ans)
   int ans = 0, curr = 0, at;
   for (auto const &i : s)
     curr = go(curr, i);
     ans += ac[curr].is leaf;
     at = get_exit_link(curr);
     while (at)
       ans += ac[at].is_leaf;
       at = get_exit_link(at);
   return ans;
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n, q;
 cin >> n >> q;
 aho::init();
 for (int i = 0; i < n; i++)
   string s;
   cin >> s;
   aho::add_string(s);
```

```
while (q--)
{
    string t;
    cin >> t;
    cout << aho::query(t) << endl;
}
    return 0;
}

automato de aho-corasick
// imagine o seguinte problema:
// temos um conjunto de n strings
// e q queries para processar
// em cada uma das q queries, voce recebe uma string s
// e quer saber, o numero de ocorrencias de
// alguma string do conjunto como
// substring de s e em tempo linear</pre>
```

### 17.2 de bruijin

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500005
#define mod 1000000009
int n, m, k, sz;
string ans, ss, path;
vector<int> d;
set<string> st;
void dfs(string s)
  if (ans.size() + path.size() == sz) // a sagacidade aqui
   ans += path;
   cout << ans << endl;</pre>
    exit(0);
  for (auto const &i : d)
    string t = s;
    t.pb('0' + i);
    if (!st.count(t))
      st.insert(t);
      string nxt = t.substr(1);
      path.pb('0' + i);
      dfs(nxt);
      path.pop_back();
      ans.pb('0' + i);
      if (ans.size() == sz)
        cout << ans << endl;</pre>
        exit(0);
signed main()
  ios_base::sync_with_stdio(false);
```

```
cin.tie(NULL);
 srand(time(NULL));
 cin >> n >> m >> k;
  d.resize(m);
 for (int i = 0; i < m; i++)
   cin >> d[i];
 sz = n + k - 1;
  if (n >= 40) // n grande -> a probabilidade de colisao eh muito baixa
    for (int i = 0; i < sz; i++)
     char c = '0' + d[rand() % m];
     s.pb(c);
    cout << s << endl; // vai uma string gerada no random e gg</pre>
    return 0:
  // n pequeno -> vamo achar um caminho euleriano
  for (int i = 1; i < n; i++)
    ss.pb('0' + d[0]);
  dfs(ss);
  ans += ss;
 while (ans.size() > sz)
   ans.pop_back();
  cout << ans << endl;
 return 0;
// vou escrever pg achei mto dahora esse problema
// https://codeforces.com/gym/102001/problem/C
// o problema basicamente eh:
// ache uma string s, minimizando o tamanho dessa string
// tal que ela tem k substrings distintas de tamanho n
// ai vai ser tipo:
// achar uma string na qual todas as substrings de tamanho n sao distintas
// alem disso, o alfabeto contem m letras
// essa string vai ter comprimento n + k - 1
// essa string eh chamada de de Bruijn sequence pro caso de k = m^n
// dai o que queremos eh basicamente achar um prefixo de uma de Bruijn sequence,
     pro k < m^n
// dai da pra transformar em um problema de achar um caminho euleriano num grafo
// montar um grafo no qual os vertices sao strings de tamanho n - 1
// e existe uma aresta direcionada u -> v se:
// v pode ser obtida adicionando um char no final de u, e tirando o primeiro
    char de u
```

## 17.3 kmp

```
#define MAXN 100005
#define mod 998244353
string s;
int n, m;
string a, b;
int c[MAXN] [26];
vector<int> kmp(string &s)
  int n = s.size();
  vector<int> p(n);
  for (int i = 1; i < n; i++)
    int j = p[i - 1];
    while (j > 0 && s[i] != s[j])
      j = p[j - 1];
    if(s[i] == s[j])
    p[i] = j;
  return p;
void compute(string s)
  s.pb('*');
  vector<int> p = kmp(s);
  for (int i = 0; i < s.size(); i++)</pre>
    for (int cc = 0; cc < 26; cc++)
      int j = i;
      while (j > 0 \&\& 'a' + cc != s[j])
       j = p[j - 1];
      if ('a' + cc == s[j])
      c[i][cc] = j;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  string s;
  cin >> s;
  compute(s);
  return 0;
// algoritmo eh online, vai coonstruindo da esquerda pra direita
// calcula pi[i], a seguinte funcao:
// seja a substring s.substr(0, i + 1)
// pi[i] = tamanho do maior prefixo que tbm eh um sufixo dessa substring
// dai por exemplo
// da pra contar a quantidade de matchings de s em t
// so concatenar as strings fazendo: t = s + "*" + t
// dai contar as posicoes com pi[i] = s.size()
// tambem eh possivel construir um automato do kmp
// do tipo
// se meu pi[i] == x, e leio a letra c
// dai devo ir pro estado p[i] == y
// as transicoes podem ser computadas e isso pode ser muito util
```

#### 17.4 manacher

```
#include <bits/stdc++.h>
using namespace std;
#define PI acos(-1)
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define fir first
```

```
#define sec second
#define MAXN 100001
#define mod 1000000007
vector<int> d1;
vector<int> d2;
void manacher(string s)
  d1.resize(s.size());
  d2.resize(s.size());
  int 1 = 0, r = -1;
  for (int i = 0; i < s.size(); i++)</pre>
    int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
    while (0 \le i - k \&\& i + k \le s.size() \&\& s[i - k] == s[i + k])
      k++;
    d1[i] = k;
    k = k - 1;
    if(i + k > r)
     1 = i - k, r = i + k;
  1 = 0, r = -1;
  for (int i = 0; i < s.size(); i++)</pre>
    int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1);
    while (0 \le i - k - 1 \&\& i + k \le s.size() \&\& s[i - k - 1] == s[i + k])
     k++;
    d2[i] = k;
    k = k - 1;
    if (i + k > r)
      1 = i - k - 1, r = i + k;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  string s;
  cin >> s;
  manacher(s);
 return 0;
// algoritimo de manacher
// motivacao: dada uma string s, encontre todos os pares (1, r) tal que, a
     substring s[l,r]
// e palindroma.
// para cada posicao (0 <= i < s.size()), vamos encontrar os valores de d1[i] e
    d2[i],
// sendo estes o numero de palindromos com comprimentos impares e com
    comprimentos pares
// e com i sendo a posicao central desses palindromos
// algoritimo mais facil:
// para cada posicao (0 <= i < s.size()), ele tenta aumentar a resposta em 1
// ate q nao seja mais possivel
// while(s[i - curr] == s[i + curr])
// complexidade O(N^2)
// algoritimo de manacher:
// para cada posicao (0 <= i < s.size()):</pre>
// seja o par (1, r) os extremos da substring palindroma que possui o maior r
    entre todas as encontradas ate entao
// se i > r, o fim do ultimo palindromo foi antes de i: iremos rodar o
    algoritimo mais facil mais facil e ir ate o limite.
// caso contrario, so precisamos rodar o algoritimo a partir de onde nao foi
    percorrido previamente.
// ao final se o r atual e maior do que o nosso antigo r, atualizamos o par (1,
// por incrivel que pareca, a complexidade e O(N)
// voltando para a motivacao:
// se temos os valores de d1[i] e d2[i]:
// a substring s[i - k, i + k] e palindroma, para todo (0 <= k < d1[i])
// a substring s[i-k-1, i+k] e palindroma, para todo (0 <= k < d2[i])
// dai temos todos os intervalos
```

```
// note que a complexidade do algoritimo de manacher e O(N), 
// mas como a quantidade maxima de palindromos em uma string e n^2, 
// imprimir todos os intervalos consequentemente teria complexidade O(N^2) no pior caso
```

#### 17.5 min suffix

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class string>
using ordered_set = tree<string, null_type, less<string>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 2001
#define mod 1000000007
int max suffix(string s, bool mi = false)
  s.push back(*min element(s.begin(), s.end()) - 1);
  int ans = 0;
  for (int i = 1; i < s.size(); i++)</pre>
    int j = 0;
    while (ans + j < i \text{ and } s[i + j] == s[ans + j])
      j++;
    if (s[i + j] > s[ans + j])
      if (!mi or i != s.size() - 2)
        ans = i;
    else if (j)
      i += j - 1;
  return ans:
int min_suffix(string s)
  for (auto &i : s)
   i *= -1;
  s.push_back(*max_element(s.begin(), s.end()) + 1);
  return max_suffix(s, true);
int max_cyclic_shift(string s)
  int n = s.size();
  for (int i = 0; i < n; i++)
    s.pb(s[i]);
  return max_suffix(s);
int min_cyclic_shift(string s)
  for (auto &i : s)
   i *= -1:
  return max_cyclic_shift(s);
// retorna a posicao de inicio menor/maior sufixo/shift de uma string
```

## 17.6 rabin-karp

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
```

```
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 100001
const int p = 31;
const int mod = 1e9 + 9;
int multiplicate(int x, int y)
  return (x * y) % mod;
int subtract(int a, int b)
  return (a - b < 0) ? a - b + mod : a - b;
int sum(int a, int b)
  return (a + b >= mod) ? a + b - mod : a + b;
vector<int> rabin_karp(string s, string t)
  int n = s.size(), m = t.size();
  vector<int> pot(n);
  pot[0] = 1;
  for (int i = 1; i < n; i++)
   pot[i] = multiplicate(pot[i - 1], p);
  vector<int> pref(n + 1, 0);
  for (int i = 0; i < n; i++)
    int val = multiplicate(pref[i], p);
    pref[i + 1] = sum(s[i], val);
  int hs = 0;
  for (int i = 0; i < m; i++)</pre>
    int val = multiplicate(hs, p);
   hs = sum(t[i], val);
  vector<int> ans;
  for (int i = 0; i + m - 1 < n; i++)
    int cur_h = subtract(pref[i + m], multiplicate(pref[i], pot[m]));
    if (cur_h == hs)
      ans.pb(i);
  return ans:
signed main()
  string s, t;
 cin >> s >> t;
  vector<int> ans = rabin_karp(s, t);
  for (auto const &i : ans)
   cout << i << " " << i + t.size() - 1 << endl;
  return 0;
// rabin-karp for pattern matching
// given two string s and t, determine all occurrences of t in s
// 1- calcule the hash of string t
// 2- calcule the prefix hash of string s
// 3- compare every substring of s with length |t|
// 4- store all occurrences in a vector and return this vector
// complexity: O(|t| + |s|)
```

### 17.7 stringhashing

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1000001
// https://github.com/brunomaletta/Biblioteca/blob/master/Codigo/Strings/
    hashingLargeMod.cpp
const int MOD = (111 << 61) - 1;
int P;
int mulmod(int a, int b)
  const static int LOWER = (111 << 30) - 1, GET31 = (111 << 31) - 1;</pre>
  int 11 = a & LOWER, h1 = a >> 30, 12 = b & LOWER, h2 = b >> 30;
  int m = 11 * h2 + 12 * h1, h = h1 * h2;
  int ans = 11 * 12 + (h >> 1) + ((h & 1) << 60) + (m >> 31) + ((m & GET31) <<
      30) + 1;
  ans = (ans \& MOD) + (ans >> 61), ans = (ans \& MOD) + (ans >> 61);
  return ans - 1;
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r)
  uniform int distribution < int > uid(1, r);
  return uid(rng);
struct string_hashing
  vector<int> h, p;
  string_hashing() {}
  string_hashing(string s) : h(s.size()), p(s.size())
    p[0] = 1, h[0] = s[0];
    for (int i = 1; i < s.size(); i++)</pre>
     p[i] = mulmod(p[i - 1], P), h[i] = (mulmod(h[i - 1], P) + s[i]) % MOD;
  int get(int 1, int r)
    int hash = h[r] - (1 ? mulmod(h[1 - 1], p[r - 1 + 1]) : 0);
    return hash < 0 ? hash + MOD : hash;</pre>
  int append(int h, int hb, int blen)
    return (hb + mulmod(h, p[blen])) % MOD;
};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n;
  cin >> n;
  P = uniform(256, MOD - 1);
  vector<string_hashing> v(n);
  vector<string_hashing> v_rev(n);
  vector<int> sz(n);
  int ans = 0;
  for (int i = 0; i < n; i++)
    string s;
    cin >> s;
    v[i] = string_hashing(s);
    sz[i] = s.size();
    ans += (s.size() * n);
```

```
ans += (s.size() * n);
   reverse(s.begin(), s.end());
    v_rev[i] = string_hashing(s);
 unordered_map<int, int> mp;
 for (int \overline{i} = 0; i < n; i++)
   for (int j = 1; j <= sz[i]; j++)</pre>
     mp[v[i].get(0, j-1)]++;
 for (int i = 0; i < n; i++)
    int acc = 0;
    for (int j = sz[i]; j >= 1; j--)
     int curr = mp[v_rev[i].get(0, j - 1)];
     ans -= ((curr - acc) * j * 2);
     acc = curr;
 cout << ans << endl;
// https://codeforces.com/contest/1902/problem/E
// solucao usando hash mod 2^61 - 1
```

### 17.8 stringhashing2

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 2001
#define mod 1000000009
struct modint
  modint(int v = 0) { val = v % mod; }
  int pow(int y)
   modint x = val;
    modint z = 1;
    while (y)
      if (v & 1)
      z \star = x;
      x \star = x;
      y >>= 1;
    return z.val;
  int inv() { return pow(mod - 2); }
  void operator=(int o) { val = 0 % mod;
  void operator=(modint o) { val = o.val % mod; }
  void operator+=(modint o) { *this = *this + o;
  void operator = (modint o) { *this = *this - o;
  void operator*=(modint o) { *this = *this * o;
  void operator/=(modint o) { *this = *this / o; }
  bool operator==(modint o) { return val == o.val;
  bool operator!=(modint o) { return val != o.val; }
  int operator*(modint o) { return ((val * o.val) % mod);
  int operator/(modint o) { return (val * o.inv()) % mod; }
```

```
int operator+(modint o) { return (val + o.val) % mod; }
 int operator-(modint o) { return (val - o.val + mod) % mod; }
struct string_hashing
 modint d;
 modint h;
  vector<modint> pref;
 vector<modint> pot;
 string_hashing() {}
  string_hashing(int base, string &s)
   d = base;
   pref.resize(s.size() + 1);
   pref[0] = 0;
    for (int i = 0; i < s.size(); i++)</pre>
      modint val = pref[i] * d;
     pref[i + 1] = val + s[i];
   h = pref[s.size()];
   pot.resize(s.size() + 1);
    pot[0] = 1;
   for (int i = 1; i <= s.size(); i++)
     pot[i] = pot[i - 1] * d;
 modint get(int 1, int r)
    return pref[r + 1] - (pref[1] * pot[r - 1 + 1]);
 modint append (modint hb, int blen)
   h = hb + (h * pot[blen]);
   return h;
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 string s;
 cin >> s;
 string hashing h(256, s); // (base, string)
  // string_hashing h(227, s); // (base, string)
```

## 17.9 substring fft

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define PI acos(-1)
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
#define cd complex<double>
const double eps = 1e-12;
const int alphabet_size = 26;
namespace fft
```

```
void dft(vector<cd> &a)
   int n = a.size();
   if (n == 1)
     return;
    vector<cd> a0 (n / 2), a1 (n / 2);
    for (int i = 0; 2 * i < n; i++)
     a0[i] = a[2 * i];
     a1[i] = a[2 * i + 1];
    dft(a0);
    dft(a1);
    double ang = 2 * PI / n;
    cd w(1), wn(cos(ang), sin(ang));
    for (int i = 0; 2 * i < n; i++)
     a[i] = a0[i] + w * a1[i];
     a[i + n / 2] = a0[i] - w * a1[i];
     w \star = wn;
 void inverse_dft(vector<cd> &a)
   int n = a.size();
   if (n == 1)
     return:
    vector<cd> a0(n / 2), a1(n / 2);
    for (int i = 0; 2 * i < n; i++)
     a0[i] = a[2 * i];
     a1[i] = a[2 * i + 1];
    inverse_dft(a0);
    inverse_dft(a1);
    double ang = 2 * PI / n * -1;
    cd w(1), wn(cos(ang), sin(ang));
    for (int i = 0; 2 * i < n; i++)
     a[i] = a0[i] + w * a1[i];
     a[i + n / 2] = a0[i] - w * a1[i];
     a[i] /= 2;
     a[i + n / 2] /= 2;
     w \star = wn:
  vector<double> mul(vector<cd> a, vector<cd> b)
    vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());
   while (n < a.size() + b.size())</pre>
     n <<= 1;
    fa.resize(n);
    fb.resize(n);
    dft(fa);
   dft(fb);
    for (int i = 0; i < n; i++)
     fa[i] *= fb[i];
    inverse_dft(fa);
    vector<double> ans(n);
    for (int i = 0; i < n; i++)</pre>
     ans[i] = fa[i].real();
    return ans;
} // namespace fft
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 string s, t;
 cin >> s >> t;
 int n = s.size(), m = t.size();
 reverse(t.begin(), t.end());
 vector<cd> a(n);
 vector<cd> b(m);
  for (int i = 0; i < n; i++)
   int ch = s[i] - 'a';
```

```
double ang = (2 * PI * ch) / alphabet_size;
    a[i] = cd(cos(ang), sin(ang));
}
for (int i = 0; i < m; i++)
{
    int ch = t[i] - 'a';
    double ang = (2 * PI * ch) / alphabet_size;
    b[i] = cd(cos(ang), -sin(ang));
}
vector<double> ans = fft::mul(a, b);
int matches = 0;
for (int i = m - 1; i < n; i++)
    matches += (abs(ans[i] - m) <= eps);
cout << matches << endl;
return 0;
}
// number of matches of a pattern in string
// using fft</pre>
```

### 17.10 suffix array

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push back
#define pi pair<int, int>
#define pii pair<pi, int>
#define pci pair<char, int>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
vector<int> suffix_array(string s)
  s += "$"; // menor do que todos os chars da string st
  int n = s.size(), N = max(n + 1, 26011);
  vector<int> sa(n), ra(n);
  for (int i = 0; i < n; i++)
    sa[i] = i, ra[i] = s[i];
  for (int k = 0; k < n; k ? k *= 2 : k++)
    vector<int> nsa(sa), nra(n), cnt(N);
    for (int i = 0; i < n; i++)
      nsa[i] = (nsa[i] - k + n) % n;
      cnt[ra[i]]++;
    for (int i = 1; i < N; i++)
      cnt[i] += cnt[i - 1];
    for (int i = n - 1; i + 1; i--)
      sa[--cnt[ra[nsa[i]]]] = nsa[i];
    for (int i = 1, r = 0; i < n; i++)
      nra[sa[i]] = r += (ra[sa[i]] != ra[sa[i - 1]] || ra[(sa[i] + k) % n] != ra
          [(sa[i-1] + k) % n]);
    ra = nra;
    if (ra[sa[n-1]] == n-1)
      break;
```

```
return vector<int>(sa.begin() + 1, sa.end());
vector<int> kasai(string s, vector<int> sa)
 int n = s.size(), k = 0;
 vector<int> ra(n), lcp(n);
for (int i = 0; i < n; i++)</pre>
   ra[sa[i]] = i;
 for (int i = 0; i < n; i++, k -= !!k)
    if (ra[i] == n - 1)
      k = 0;
      continue;
    int j = sa[ra[i] + 1];
    while (i + k < n \text{ and } j + k < n \text{ and } s[i + k] == s[j + k])
    lcp[ra[i]] = k;
 return lcp;
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 string s;
 cin >> s;
 int n = s.size();
 vector<int> p = suffix_array(s);
 vector<int> lcp = kasai(s, p);
 for (int i = 0; i < s.size(); i++) // sufix array</pre>
   cout << p[i] << " ";
  cout << endl;</pre>
 for (int i = 0; i + 1 < s.size(); i++) // lcp entre 2 suffixos adjacentes no</pre>
      suffix array
   cout << lcp[i] << " ";
 cout << endl;
 int q;
 cin >> q;
 while (q--)
    string t;
    cin >> t;
    int i = 0, f = n - 1, m, lb, ub;
    while (i < f)
      m = (i + f) / 2;
      (t \le s.substr(p[m], t.size())) ? f = m : i = m + 1;
    ub = i, i = 0, f = n - 1;
    while (i < f)
      m = (i + f) / 2;
      (t \ge s.substr(p[m], t.size())) ? i = m + 1 : f = m;
    if (s.substr(p[lb], t.size()) == t)
    cout << lb - ub << endl;
 return 0:
```

#### 17.11 suffix automaton

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
```

```
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100001
#define mod 998244353
namespace sa
  struct state
   int len, suf_link;
   map<char, int> nxt;
  vector<int> term;
  state st[2 * MAXN];
  int dp[2 * MAXN];
  int sz, last;
  void init()
   memset(dp, -1, sizeof(dp));
   st[0].len = 0;
    st[0].suf_link = -1;
    sz++;
   last = 0;
  void get_link(int curr, int p, char c)
    while (p != -1 && !st[p].nxt.count(c))
      st[p].nxt[c] = curr;
      p = st[p].suf_link;
    if (p == -1)
      st[curr].suf link = 0;
      return;
    int q = st[p].nxt[c];
    if (st[p].len + 1 == st[q].len)
      st[curr].suf_link = q;
      return;
    int clone = sz;
    sz++:
    st[clone].len = st[p].len + 1;
    st[clone].nxt = st[q].nxt;
    st[clone].suf_link = st[q].suf_link;
    while (p != -1 && st[p].nxt[c] == q)
      st[p].nxt[c] = clone;
      p = st[p].suf_link;
    st[q].suf_link = clone;
   st[curr].suf_link = clone;
  void build(string &s)
    for (auto const &c : s)
      int curr = sz;
      st[curr].len = st[last].len + 1;
      get_link(curr, last, c);
      last = curr;
    // achar os estados terminais
    // um estado terminal e aquele que representa um sufixo da string s
    int p = last;
    while (p != -1)
```

```
term.pb(p);
     p = st[p].suf_link;
 void dfs2(int v)
    if (dp[v] != -1)
     return:
    dp[v] = 1;
    for (auto const &u : st[v].nxt)
     if (!u.sec)
       continue:
     dfs2(u.sec);
     dp[v] += dp[u.sec];
 void dfs(int v, int k, int &at, string &curr)
    if (at == k)
     return;
    for (auto const &u : st[v].nxt)
     if (!u.sec)
       continue:
      if (at + dp[u.sec] < k)
        at += dp[u.sec];
       continue;
     curr.pb(u.fir);
     at++;
      dfs(u.sec, k, at, curr);
     if (at == k)
       return;
     curr.pop_back();
 void find_kth(int k)
    int at = 0:
   string curr = "";
   dfs(0, k, at, curr);
   cout << curr << endl;</pre>
1 // namespace sa
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 string s;
 cin >> s;
 sa::init();
 sa::build(s);
 sa::dfs2(0);
 int q;
 cin >> q;
 while (q--)
   cin >> k;
   sa::find_kth(k);
 return 0;
// https://cp-algorithms.com/string/suffix-automaton.html
// suffix automaton
// definicao: um suffix automaton de uma string s e um automato finito
    deterministico
// que aceita todos os suffixos da string s.
// um suffix automaton eh um grafo aciclico orientado
// tal que, um vertice representa um estado
// e uma aresta representa uma transicao (um caractere a mais em relacao ao
    estado(suffixo) atual)
// t0 -> estado inicial(string vazia), e todos os demais estados podem ser
    alcancados a partir de tÕ
// o suffix automaton minimiza o numero de vertices
```

```
// a propiedade mais importante de um suffix automaton eh a de que
// ele contem informacoes sobre todas as substrings de s
// pois, qualquer caminho comecando do estado t0 corresponde a uma substring de
// conceitos:
// 1 - endpos
// seja t uma substring de s, endpos(t) eh o conjunto de todas os indices(
    posicoes)
// na string s no qual todas as ocorrencias de t acabam
// por exemplo, se s = "abcbc" e t = "bc"
// logo endpos(t) = {2, 4}
// com isso se duas duas substrings t1 e t2 possuem os seus endpos iquais,
// chamamos de endpos-equivalent e dai podemos extrair algumas informacoes
// info 1: se duas substrings u e w u.size() \le w.size(), se u eh um sufixo de w
     , logo endpos(u) esta contido em endpos(w)
// info 2: se duas substrings u e w u.size() <= w.size(), se u nao eh um sufixo
    de w, logo nao existe interseccao entre endpos(u) e endpos(w)
// 2 - suffix link
// seja v algum estado != t0, sabemos que v corresponde a classe de strings que
      possui os mesmos endpos
// seja w a maior dessas strings, com isso, todas as demais sao suffixos de w
// com isso um suffix_link(v) corresponde ao maior suffix de w que esta em outra
     classe de equivalencia pelos endpos
// com isso podemos abstrair algumas informações:
// info 1: os suffix links foram uma arvore enraizada em t0
// info 2: se construirmos uma arvore usando os sets endpos, a estrutura sera a
    arvore com os suffix links
// com isso, vamos ao algoritimo
// 1 - vai ser online, e iremos adicionar os caracteres de 1 por 1, da esquerda
    para a direita
// 2 - com isso para adicionar um novo char, seja v o ultimo estado que
    adicionamos antes do atual, adicionamos uma aresta
// do proximo em relacao a ele e iremos procurar pelo suffix link para adicionar
// 3 - complexidade O(n) ou O(n \log k), se usarmos uma map para guardar as
     transicoes partindo de um estado
// exemplos de aplicacoes:
// 1 - checar se t aparece em s como substring:
// construa o suffix automaton de s, e vamos tentar fazer um caminho partindo de
// se em algum momento, nao existir transicao, logo nao existe
// se conseguir chegar no final, existe
// 2 - numero de substrings diferentes de s
// constura o suffix automaton de s, sabemos que, cada substring de s
    corresponde a um caminho no automato
   com isso, o numero de substrings distintas eh o numero de caminhos diferentes
     que comecam de t0
// e terminam em algum canto
// isso pode ser calculado facilmente com uma dpzinha
// 3 - tamanho total de todas as substrings distintas de s
// similar a solucao passada, podemos fazer isso com uma dpzinha :)
// 4 - a k-esima menor substring lexicografica
// a k-esima menor substring lexicograficamente corresponde ao k-esimo path no
// se considerarmos as transicoes sempre indo do menor char para o maior durante
     o percurso
// 5 - o menor cyclic shift
// construa o suffix automaton da string s + s (duplicada)
// com isso o suffix automaton vai conter todos os cyclic shifts da string s
// e agora o problema eh reduzido para: encontre o menor caminho
     lexicograficamente de tamanho s.size()
// 6 - numero de ocorrencias de uma substring t em s
// construa o suffix automaton da string s
// com isso, quando criamos um no que nao seja o t0 nem um clone
// inicializamos cnt[v] = 1
// depois vamos percorrer todo os estados em ordem decrescente de len
// e aplicando cnt[link(v)] += cnt[v]
```

// no final, para responder a query basta fazer o caminho ate o estado que

```
quisermos e printar o cnt dele

// e mais uma porrada de aplicacoes alem dessas :)

// example of a problem: https://www.spoj.com/problems/SUBLEX/
// ver qual a k-th string lexicografica sem repeticao
// note que o k pode ser gigante
// ideia: calcular dp[v] -> quantidade de caminhos que comecam em v
// dai para cada query roda um dfs, sendo que, so vou pro proximo estado se at +
dp[u] >= k
// caso contrario, posso ignorar
```

#### 17.12 z-function

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 2001
#define mod 1000000007
vector<int> z_function(string &s)
  int n = s.size();
  vector<int> z(n);
  z[0] = n;
  for (int i = 1, l = 0, r = 0; i < n; i++)
    if (i \le r)
     z[i] = min(r - i + 1, z[i - 1]);
    while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
     z[i]++;
    if (i + z[i] - 1 > r)
     1 = i, r = i + z[i] - 1;
  return z;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  string s;
  cin >> s;
  vector<int> z = z_function(s);
// z-function
// calcula para cada i:
// z[i] = o tamanho de lcp(s, s.substr(i, n - i))
// lcp -> longest comom prefix
```

## 18 Structures

## 18.1 binary lifting

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
```

```
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
struct item
  int nxt, sum;
};
int n, q;
int v[MAXN];
item st[MAXN][21];
signed main()
  cin >> n >> q;
  for (int i = 0; i < n; i++)
   cin >> v[i];
  for (int i = 0; i < n; i++)
    st[i][0].nxt = min(i + 1, n - 1);
   st[i][0].sum = v[st[i][0].nxt];
  for (int i = 1; i < 21; i++)
    for (int v = 0; v < n; v++)
      st[v][i].nxt = st[st[v][i - 1].nxt][i - 1].nxt;
     st[v][i].sum = st[v][i - 1].sum + st[st[v][i - 1].nxt][i - 1].sum;
  while (q--)
    int 1, r;
    cin >> 1 >> r;
    int ans = v[1], len = r - 1;
    for (int i = 20; i >= 0; i--)
      if (len & (1 << i))
        ans += st[l][i].sum;
        1 = st[1][i].nxt;
   cout << ans << endl;</pre>
  return 0;
// simple range sum query with binary lifting
// https://judge.yosupo.jp/problem/static_range_sum
```

## 18.2 bit2d

```
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1000001
#define mod 1000000007
// source: https://github.com/tfq50/Competitive-Programming/blob/master/
    Biblioteca/Data%20Structures/Bit2D.cpp
struct bit2d
 vector<int> ord;
 vector<vector<int>> t;
 vector<vector<int>> coord;
 bit2d(vector<pi> &pts) // recebe todos os pontos que vao ser inseridos pra
      construir, mas nao insere eles
    sort(pts.begin(), pts.end());
    for (auto const &a : pts)
      if (ord.empty() || a.fir != ord.back())
        ord.pb(a.fir);
    t.resize(ord.size() + 1);
    coord.resize(t.size());
    for (auto &a : pts)
      swap(a.fir, a.sec);
    sort(pts.begin(), pts.end());
    for (auto &a : pts)
      swap(a.fir, a.sec);
      for (int on = upper_bound(ord.begin(), ord.end(), a.fir) - ord.begin(); on
            < t.size(); on += on & -on)
        if (coord[on].empty() || coord[on].back() != a.sec)
          coord[on].push_back(a.sec);
    for (int i = 0; i < t.size(); i++)</pre>
      t[i].assign(coord[i].size() + 1, 0);
 void add(int x, int y, int v) // v[a][b] += v
    for (int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin(); xx < t.</pre>
        size(); xx += xx & -xx)
      for (int yy = upper_bound(coord[xx].begin(), coord[xx].end(), y) - coord[
           xx].begin(); yy < t[xx].size(); yy += yy & -yy)
        t [xx] [yy] += v;
  int qry(int x, int y) // soma de todos os v[a][b] com (a <= x && b <= y)
    int ans = 0;
    for (int xx = upper_bound(ord.begin(), ord.end(), x) - ord.begin(); xx > 0;
        xx = xx & -xx
      for (int yy = upper_bound(coord[xx].begin(), coord[xx].end(), y) - coord[
          xx].begin(); yy > 0; yy -= yy & -yy)
        ans += t[xx][yy];
    return ans;
 int gry2 (int x1, int y1, int x2, int y2)
    return qry(x2, y2) - qry(x2, y1 - 1) - qry(x1 - 1, y2) + qry(x1 - 1, y1 - 1)
  void add2(int x1, int y1, int x2, int y2, int v)
    add(x1, y1, v);
```

```
add(x1, y2 + 1, -v);
add(x2 + 1, y1, -v);
add(x2 + 1, y2 + 1, v);
};
signed main()
{
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
}
```

### 18.3 color update

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
const int inf = 1e15;
struct color upd
#define left fir
#define right sec.fir
#define color sec.sec
  set<pii> ranges:
  vector<pii> erased;
  color_upd(int n) // inicialmente, todo mundo pintado com a cor inf
    // nao usar cores negativas!!!!!!!!
    ranges.insert(\{0, \{n-1, inf\}\}\);
  int get(int i) // qual a cor do elemento na posicao i
    auto it = ranges.upper_bound({i, {1e18, 1e18}});
    if (it == ranges.begin())
      return -1;
    return (*(it)).color;
  void del(int 1, int r) // apaga o intervalo [1, r]
    erased.clear();
    auto it = ranges.upper_bound({1, {0, 0}});
    if (it != ranges.begin())
      it --:
    while (it != ranges.end())
      if ((*(it)).left > r)
       break:
      else if ((*(it)).right >= 1)
        erased.push_back(*it);
      it++;
    if (erased.size() > 0)
```

```
int sz = erased.size();
     auto it = ranges.lower_bound({erased[0].left, {0, 0}});
      auto it2 = ranges.lower_bound({erased[sz - 1].left, {0, 0}});
      pii ini = *it, fim = *it2;
      it2++;
     ranges erase(it, it2);
     pii upd1 = {ini.left, {1 - 1, ini.color}};
     pii upd2 = {r + 1, {fim.right, fim.color}};
     erased[0].left = max(erased[0].left, 1);
      erased[sz - 1].right = min(erased[sz - 1].right, r);
     if (upd1.left <= upd1.right)</pre>
        ranges.insert(upd1);
      if (upd2.left <= upd2.right)</pre>
        ranges.insert(upd2);
 void add(int a, int b, int c) // nao ter dois intervalos adjacentes com a
      mesma cor no set de ranges
   auto it = ranges.lower bound({a, {b, 0}});
   pii aa = \{-1, \{-1, -1\}\};
    pii bb = \{-1, \{-1, -1\}\};
   if (it != ranges.end())
     if ((*it).color == c && (*it).left == b + 1)
        aa = *it;
        b = (*it).right;
    if (it != ranges.begin())
     if ((*it).color == c && (*it).right == a - 1)
        bb = *it;
        a = (*it).left;
    ranges.erase(aa);
   ranges.erase(bb);
   ranges.insert({a, {b, c}});
 void upd(int a, int b, int c) // pinta o intervalo [a, b] com a cor c
   del(a, b);
   add(a, b, c);
};
struct segtree
 vector<int> seq;
 vector<int> lazy;
  segtree(int n)
    seg.resize(4 * n, 0);
   lazy.assign(4 * n, 0);
  int single(int x)
   return x;
  int neutral()
   return 0:
 int merge(int a, int b)
   return a + b;
 void add(int i, int 1, int r, int diff)
    seg[i] += (r - 1 + 1) * diff;
   if(1 != r)
     lazy[i << 1] += diff;</pre>
     lazy[(i << 1) | 1] += diff;</pre>
```

```
lazy[i] = 0;
  void update(int i, int l, int r, int ql, int qr, int diff)
    if (lazy[i])
     add(i, l, r, lazy[i]);
    if (1 > r || 1 > qr || r < q1)
     return;
    if (1 >= q1 && r <= qr)
     add(i, 1, r, diff);
      return;
    int mid = (1 + r) >> 1;
    update(i << 1, 1, mid, q1, qr, diff);
   update((i << 1) | 1, mid + 1, r, ql, qr, diff);
   seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
  int query(int 1, int r, int q1, int qr, int i)
   if (lazy[i])
     add(i, 1, r, lazy[i]);
    if (1 > r || 1 > qr || r < q1)
     return neutral();
    if (1 >= q1 && r <= qr)
     return seg[i];
    int mid = (1 + r) >> 1;
   return merge(query(1, mid, q1, qr, i << 1), query(mid + 1, r, q1, qr, (i <<</pre>
        1) | 1));
};
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n, q;
 cin >> n >> q;
 color_upd c = color_upd(n);
  segtree s = segtree(n);
  for (int i = 0; i < n; i++)
   c.upd(i, i, i + 1);
  while (q--)
   int t;
   cin >> t;
   if (t == 1)
     int 1, r, x;
      cin >> 1 >> r >> x;
     1--, r--;
      c.upd(l, r, x);
      for (auto const &i : c.erased)
        s.update(1, 0, n - 1, i.left, i.right, abs(x - i.color));
    else
     int 1, r;
     cin >> 1 >> r;
     cout << s.query(0, n - 1, 1, r, 1) << endl;
  return 0;
// https://codeforces.com/contest/444/problem/C
```

#### 18.4 fenwick

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
```

```
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 998244353
struct fenw
  int n;
  vector<int> bit;
  fenw() {}
  fenw(int sz)
   n = sz;
    bit.assign(sz + 1, 0);
  int qry(int r) // query de prefixo a[0] + a[1] + ... a[r]
    int ret = 0;
    for (int i = r + 1; i > 0; i -= i & -i)
     ret += bit[i];
    return ret;
  void upd(int r, int x) // a[r] += x
    for (int i = r + 1; i \le n; i += i \& -i)
     bit[i] += x;
  int bs(int x) // retorna o maior indice i (i < n) tal que: qry(i) < x
    int i = 0, k = 0;
    while (1 << (k + 1) <= n)
      k++;
    while (k >= 0)
      int nxt i = i + (1 << k);
      if (nxt_i <= n && bit[nxt_i] < x)</pre>
        i = nxt i;
        x -= bit[i];
      k--;
    return i - 1:
};
```

#### 18.5 fenwick2

```
// fenwick com update pro range [1, r]
// complexidade O(q * log(n)) com a criacao de duas bits ao inves de uma
#include <bits/stdc++.h>
using namespace std;
#define PI acos(-1)
#define int long long int
#define pb push_back
#define mp make_pair
#define pi pair<string, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100001
#define MAXL 20
#define mod 998244353
int n;
vector<int> bit, bit2;
```

```
void add1(int idx, int delta)
  for (; idx < n; idx = idx | (idx + 1))
    bit[idx] += delta;
void add2(int idx, int delta)
  for (; idx < n; idx = idx | (idx + 1))
    bit2[idx] += delta;
void update_range(int val, int l, int r)
  add1(1, val);
  add1(r + 1, -val);
  add2(1, val * (1 - 1));
  add2(r + 1, -val * r);
int sum1(int r)
  int ret = 0;
  for (; r \ge 0; r = (r \& (r + 1)) - 1)
    ret += bit[r];
  return ret:
int sum2(int r)
  int ret = 0;
  for (; r >= 0; r = (r & (r + 1)) - 1)
    ret += bit2[r];
  return ret;
int sum(int x)
  return (sum1(x) * x) - sum2(x);
int range_sum(int 1, int r)
  return sum(r) - sum(1 - 1);
int main()
  bit.assign(MAXN, 0); // inicializar sempre
 bit2.assign(MAXN, 0); // inicializar sempre
update_range(x, 1, r); // pra cada elemento em [1, r] += x
range_sum(1, r); // soma de [1, r]
```

### 18.6 fenwick2D

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1025
#define mod 1000000007
int b[MAXN][MAXN];
int vv[MAXN][MAXN];
int qry(int x, int y)
```

```
int sum = 0:
  for (; x >= 0; x = (x & (x + 1)) - 1)
    for (int yy = y; yy >= 0; yy = (yy & (yy + 1)) - 1)
     sum += b[x][yy];
 return sum;
void add(int x, int y, int v)
 for (; x < MAXN; x = x | (x + 1))
    for (int yy = y; yy < MAXN; yy = yy | (yy + 1))
     b[x][yy] += v;
int qry2(int x1, int y1, int x2, int y2)
 return qry(x2, y2) - qry(x2, y1 - 1) - qry(x1 - 1, y2) + qry(x1 - 1, y1 - 1);
void add2(int x1, int y1, int x2, int y2, int v)
 add(x1, y1, v);
 add(x1, y2 + 1, -v);
  add (x2 + 1, y1, -v);
 add(x^2 + 1, y^2 + 1, v);
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int q;
 cin >> q;
 while (q--)
    int n;
    for (int i = 0; i < n; i++) // reseta</pre>
      for (int j = 0; j < n; j++)
        add(i, j, -vv[i][j]);
        vv[i][j] = 0;
    while (1)
      string s;
      cin >> s;
      if (s == "SET")
        int a, b, c;
        cin >> a >> b >> c;
        add(a, b, -vv[a][b]);
        vv[a][b] = c;
        add(a, b, vv[a][b]);
      else if (s == "SUM")
        int a, b, c, d;
        cin >> a >> b >> c >> d; // c >= a e d >= b
        cout << qry2(a, b, c, d) << endl;</pre>
      else
        break;
 return 0;
// to test: https://www.spoj.com/problems/MATSUM/
```

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
//#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 200005
#define mod 1000000007
int v[MAXN];
namespace bit
  ordered_set<int> bit[MAXN];
  int query(int r, int a, int b)
    int ret = 0, curr = r;
    for (; r >= 0; r = (r & (r + 1)) - 1)
     ret += (bit[r].order_of_key(b + 1) - bit[r].order_of_key(a));
    return ret;
  void add(int idx, int delta)
    for (; idx < MAXN; idx = idx | (idx + 1))
     bit[idx].insert(delta);
  void rem(int idx, int delta)
    for (; idx < MAXN; idx = idx | (idx + 1))
     bit[idx].erase(delta);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
// ideia da merge sort tree na bit (fica mais rapido)
// so fazer uma bit de ordered set ou vector(se nao tiver update)
// add -> adiciona o numero delta na posicao idx
// rem -> remove o numero delta na posicao idx
// query -> retorna o numero de elementos tal que posicao <= r && (a <= num <= b
```

## 18.8 implicit seg

```
#define MAXN 200005
#define mod 998244353
struct implcit_seg
 int 1, r;
 int sum, lazy;
  implcit_seg *left_child = nullptr;
 implcit_seg *right_child = nullptr;
  implcit_seg(int 1, int r) : 1(1), r(r)
    sum = 0;
    lazy = 0;
 void check_childs()
    if (!left_child && 1 != r)
      int mid = (1 + r) >> 1;
      left_child = new implcit_seg(l, mid);
      right_child = new implcit_seg(mid + 1, r);
 void add(int x)
    sum += (r - 1 + 1) * x;
    if (1 != r)
     check_childs();
     left_child->lazy += x;
     right_child->lazy += x;
    lazy = 0;
 void upd(int ql, int qr, int x)
    add(lazy);
   if (l > r || l > qr || r < ql)</pre>
     return;
    if (1 >= q1 && r <= qr)</pre>
      add(x);
      return;
    check_childs();
    left_child->upd(ql, qr, x);
    right_child->upd(ql, qr, x);
    sum = left_child->sum + right_child->sum;
 void upd(int k, int x)
    sum += x;
    check_childs();
    if (left_child)
      if (k <= left_child->r)
        left_child->upd(k, x);
        right_child->upd(k, x);
 int qry(int ql, int qr)
    add(lazy);
   if (1 > r || 1 > qr || r < q1)</pre>
     return 0;
   if (1 >= q1 && r <= qr)</pre>
     return sum;
    check_childs();
   return left_child->qry(ql, qr) + right_child->qry(ql, qr);
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n, q;
```

```
cin >> n >> q;
implcit_seg *s = new implcit_seg(0, n - 1);
while (q--)
 int t;
  cin >> t;
  if (t == 1)
    int 1, r, x;
    cin >> 1 >> r >> x;
    if (1 == r - 1) // point update
      s \rightarrow upd(1, x);
    else // range update
      s \rightarrow upd(1, r - 1, x);
  else
    int 1, r;
    cin >> 1 >> r;
    cout << s->qry(1, r - 1) << endl; // range sum
return 0;
```

### 18.9 lower bound segtree

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 300005
#define mod 998244353
const int inf = 1e18;
struct segtree
  int n;
  vector<int> v;
  vector<int> seq;
  segtree (vector<int> &vv)
   n = v.size();
   seg.assign(4 * n, 0);
   build (0, n - 1, 1);
  int single(int x)
   return x:
  int neutral()
    return 1e18;
  int merge(int a, int b)
   return max(a, b);
  void update(int i, int 1, int r, int q, int x)
```

```
if (1 == r)
     seg[i] = single(x);
     return;
    int mid = (1 + r) >> 1;
    if (q <= mid)
     update(i << 1, 1, mid, q, x);
     update((i << 1) | 1, mid + 1, r, q, x);
    seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
 int query(int 1, int r, int q1, int qr, int i, int x)
    if (1 == r)
     return (seg[i] >= x && l >= ql) ? l : -1;
    int mid = (1 + r) >> 1, at = -1;
   if (seg[i << 1] >= x && mid >= ql)
     at = query(1, mid, q1, qr, i \ll 1, x);
   if (at == -1)
     at = query(mid + 1, r, ql, qr, (i << 1) | 1, x);
    return at;
 void build(int 1, int r, int i)
    if (1 == r)
     seg[i] = single(v[1]);
     return;
   int mid = (1 + r) >> 1;
   build(1, mid, i << 1);
   build (mid + 1, r, (i << 1) | 1);
   seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
 int gry(int x, int 1)
   return query(0, n - 1, 1, n - 1, 1, x);
 void upd(int x, int 1)
    update(1, 0, n - 1, x, 1);
signed main()
 int n, q;
 cin >> n >> q;
 vector<int> v(n);
 for (int i = 0; i < n; i++)
   cin >> v[i];
  segtree st(v);
 while (q--)
   int t;
   cin >> t:
   if (t == 2)
     int x, 1;
     cin >> x >> 1; // find the minimum index j such that j >= 1 and v[j] >= x
     cout << st.qry(x, 1) << endl;
    else
     int a, b;
     cin >> a >> b; // v[a] = b;
     st.upd(a, b);
```

## 18.10 mergesorttree

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
               tree_order_statistics_node_update>;
 #define int long long int
 #define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
 #define fir first
 #define sec second
#define DEBUG 0
 #define MAXN 100001
#define mod 1000000007
vector<int> seg[4 * MAXN];
int v[MAXN];
void update(int i, int 1, int r, int q, int x)
      if (1 == r)
            seg[i].clear();
            seg[i].pb(x);
            return;
       int mid = (1 + r) >> 1;
       if (q <= mid)</pre>
           update(i << 1, 1, mid, q, x);
            update((i << 1) | 1, mid + 1, r, q, x);
       // a merge do c++ une os dois vectors, deixando ele ordenado em O(n)
      merge(seg[i << 1].begin(), seg[i << 1].end(), seg[(i << 1) | 1].begin(), seg[(i << 1) | 2].begin(), seg[(i << 1) | 3].begin(), seg[(i << 1) | 4].begin(), seg[(i << 1) | 5].begin(), 
                      i << 1) | 1].end(), back_inserter(seg[i]));</pre>
int query(int 1, int r, int q1, int qr, int i, int x)
      int mid = (1 + r) >> 1;
      if (1 > r || 1 > qr || r < q1)</pre>
            return 0;
       if (1 >= q1 && r <= qr) // quantidade de elementos maiores do que x no range
            return seg[i].end() - upper_bound(seg[i].begin(), seg[i].end(), x);
      return query(1, mid, q1, qr, i << 1, x) + query(mid + 1, r, q1, qr, (i << 1) |
                        1, x);
void build(int 1, int r, int i)
      if (1 == r)
            seg[i].pb(v[1]);
            return;
       int mid = (1 + r) >> 1;
      build(1, mid, i << 1);
      build (mid + 1, r, (i << 1) | 1);
       // a merge do c++ une os dois vectors, deixando ele ordenado em O(n)
      merge(seg[i << 1].begin(), seg[i << 1].end(), seg[(i << 1) | 1].begin(), 
                      i << 1) | 1].end(), back_inserter(seg[i]));</pre>
signed main()
      ios_base::sync_with_stdio(false);
       cin.tie(NULL);
      return 0;
 // merge sort tree
// a segment tree with ordered vectors in range nodes
 // example:
 // number of elements > x in a range [1, r]
```

```
// memory: O(n * log n)
// query: O(log^2 n)
```

## 18.11 min queue

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 1005
#define mod 998244353
namespace min_queue
  deque<pi> q;
  int 1, r;
  void init()
    1 = r = 1:
    q.clear();
  void push(int v)
    while (!q.empty() \&\& v < q.back().fir)
      q.pop_back();
    q.pb({v, r});
    r++;
  void pop()
    if (!q.empty() && q.front().sec == 1)
      q.pop_front();
    1++;
  int getmin()
    return q.front().fir;
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, m;
  cin >> n >> m;
 vector<int> v(n);
for (int i = 0; i < n; i++)</pre>
   cin >> v[i];
  int 1 = 0, r = m - 1;
cout << 1 << " " << r << endl;</pre>
  for (int i = 1; i <= r; i++)</pre>
   min_queue::push(v[i]);
  cout << min_queue::getmin() << " ";</pre>
  1++, r++;
  while (r < n)
    min_queue::pop();
    min_queue::push(v[r]);
    cout << min_queue::getmin() << " ";</pre>
    1++, r++;
  cout << endl;
```

```
return 0;
}
// minimum of each subarray of length m (m <= n)</pre>
```

#### 18.12 mo

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500005
#define mod 1000000007
int n, q;
int v[MAXN];
namespace mo
  struct query
   int idx, 1, r;
  };
  int block;
  vector<query> queries;
  vector<int> ans;
  bool cmp(query x, query y)
   if (x.1 / block != y.1 / block)
      return x.1 / block < y.1 / block;
    return x.r < y.r;</pre>
  void run()
    block = (int)sqrt(n);
    sort(queries.begin(), queries.end(), cmp);
    ans.resize(queries.size());
    int c1 = 0, cr = -1, sum = 0;
    auto add = [&](int x)
      sum += x;
    auto rem = [&](int x)
      sum -= x;
    for (int i = 0; i < queries.size(); i++)</pre>
      while (cl > queries[i].1)
        add(v[cl]);
      while (cr < queries[i].r)</pre>
        cr++:
        add(v[cr]);
      while (cl < queries[i].l)</pre>
        rem(v[cl]);
        cl++;
```

```
while (cr > queries[i].r)
        rem(v[cr]);
        cr--;
      ans[queries[i].idx] = sum;
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
  cin >> n >> q;
 for (int i = 0; i < n; i++)
   cin >> v[i];
 for (int i = 0; i < q; i++)
   mo::query curr;
   cin >> curr.1 >> curr.r;
   curr.r--;
   curr idx = i;
   mo::queries.pb(curr);
 mo::run();
  for (auto const &i : mo::ans)
   cout << i << endl;</pre>
// to test: https://judge.yosupo.jp/problem/static_range_sum
```

## 18.13 mo update

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
int n, q;
int v[MAXN];
int vv[MAXN];
namespace mo
  struct query
    int idx, 1, r, t;
  };
  struct update
    int i, prevx, x;
  };
  int block;
  vector<query> queries;
  vector<update> updates;
  vector<int> ans;
 bool cmp(query x, query y)
```

```
if (x.1 / block != y.1 / block)
     return x.1 / block < y.1 / block;
    if (x.r / block != y.r / block)
     return x.r / block < y.r / block;
    return x.t < y.t;</pre>
 void run()
   block = 3153; // (2 * n) ^ 0.666
   sort(queries.begin(), queries.end(), cmp);
    ans.resize(queries.size());
   int c1 = 0, cr = -1, sum = 0, t = 0;
   auto add = [&](int x)
     sum += x;
    };
    auto rem = [&](int x)
      sum -= x;
    for (int i = 0; i < queries.size(); i++)</pre>
      while (cl > queries[i].l)
       add(v[cl]);
      while (cr < queries[i].r)</pre>
       cr++;
        add(v[cr]);
      while (cl < queries[i].l)</pre>
        rem(v[cl]);
        cl++;
      while (cr > queries[i].r)
        rem(v[cr]);
        cr--;
      while (t > queries[i].t)
        if (queries[i].1 <= updates[t].i && queries[i].r >= updates[t].i)
          rem(updates[t].x);
          add(updates[t].prevx);
        v[updates[t].i] = updates[t].prevx;
      while (t < queries[i].t)</pre>
        if (queries[i].1 <= updates[t].i && queries[i].r >= updates[t].i)
          rem(updates[t].prevx);
          add(updates[t].x);
        v[updates[t].i] = updates[t].x;
        t++;
      ans[queries[i].idx] = sum;
 }
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 cin >> n >> q;
 for (int i = 0; i < n; i++)
   cin >> v[i];
   vv[i] = v[i];
  for (int i = 0; i < q; i++)
```

int type;

```
cin >> type;
    if (type == 1)
     mo::update curr;
     cin >> curr.i >> curr.x;
      curr.prevx = vv[curr.i];
      vv[curr.i] = curr.x;
     mo::updates.pb(curr);
    else
     mo::query curr;
     cin >> curr.l >> curr.r;
     curr.r--;
      curr.idx = mo::queries.size();
     curr.t = mo::updates.size();
     mo::queries.pb(curr);
 mo::run();
 for (auto const &i : mo::ans)
   cout << i << endl;</pre>
// to test: https://codeforces.com/edu/course/2/lesson/4/1/practice/contest
    /273169/problem/A
// 1 i v - set the element with index i to v
// 2 1 r - calculate the sum of elements with indices from 1 to r - 1
// n, q <= 100000
// runs in 467ms
```

## 18.14 persistent seg

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 100003
#define mod 1000000007
int v[MAXN];
namespace seg
  struct node
    int item, lazy, lazy_status, l, r;
    node(int 1, int r, int lazy, int lazy_status, int item) : 1(1), r(r), lazy(
        lazy), lazy_status(lazy_status), item(item) {}
  };
  vector<node> seg;
  vector<int> roots;
  void init()
    seq.resize(1);
  int neutral()
    return 0;
  int merge(int a, int b)
```

```
return a + b:
    int newleaf(int vv)
        int p = seq.size();
        seg.pb(node(0, 0, 0, 0, vv));
        return p;
    int newparent(int 1, int r)
        int p = seq.size();
        seg.pb(node(1, r, 0, 0, merge(seg[1].item, seg[r].item)));
        return p;
    int newkid(int i, int diff, int l, int r)
         int p = seq.size();
        seg.pb(node(seg[i].l, seg[i].r, seg[i].lazy + diff, 1, seg[i].item + ((r - l).seg[i].item + ((r - l).seg[i].lazy + diff, 1, seg[i].item + ((r - l).seg[i].seg[i].lazy + diff, 1, seg[i].item + ((r - l).seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg[i].seg
                      + 1) * diff());
        return p;
    void add(int i, int l, int r)
        if (!seq[i].lazy_status)
             return;
         if (1 != r)
             int mid = (1 + r) >> 1;
             seg[i].l = newkid(seg[i].l, seg[i].lazy, l, mid);
             seg[i].r = newkid(seg[i].r, seg[i].lazy, mid + 1, r);
        seg[i].lazy = 0;
        seg[i].lazy_status = 0;
    int update(int i, int l, int r, int gl, int gr, int diff)
        if (1 > r || 1 > qr || r < ql)
             return i;
         if (1 >= ql && r <= qr)
             return newkid(i, diff, l, r);
         add(i, 1, r);
         int mid = (1 + r) >> 1;
         return newparent (update (seg[i].1, 1, mid, ql, qr, diff), update (seg[i].r,
                    mid + 1, r, ql, qr, diff));
    int query (int 1, int r, int q1, int qr, int i)
        if (1 > r || 1 > qr || r < ql)
            return neutral();
         if (1 >= q1 && r <= qr)
             return seg[i].item;
         add(i, 1, r);
         int mid = (1 + r) >> 1;
         return merge(query(1, mid, q1, qr, seg[i].1), query(mid + 1, r, q1, qr, seg[
                    i].r));
    int build(int 1, int r)
        if (1 == r)
             return newleaf(v[1]);
         int mid = (1 + r) >> 1;
        return newparent(build(1, mid), build(mid + 1, r));
signed main()
    ios_base::sync_with_stdio(false);
    cin.tie(NULL);
    int n, q;
    cin >> n >> q;
    for (int i = 0; i < n; i++)
       cin >> v[i];
     seg::init();
    int root = seg::build(0, n - 1);
    seg::roots.pb(root);
    while (q--)
```

```
char t;
    cin >> t:
    if (t == 'C')
     int 1, r, d;
     cin >> 1 >> r >> d;
      1--, r--;
      int root = seq::update(seq::roots.back(), 0, n - 1, 1, r, d);
      seq::roots.pb(root);
    else if (t == 'O')
      int 1, r;
      cin >> 1 >> r;
      1--, r--;
     cout << seg::query(0, n - 1, 1, r, seg::roots.back()) << endl;</pre>
    else if (t == 'H')
      int 1, r, d;
     cin >> 1 >> r >> d;
      1--, r--;
     cout << seq::query(0, n - 1, 1, r, seq::roots[d]) << endl;</pre>
    else
      int d;
     cin >> d:
      while (seg::roots.size() > d + 1)
        seg::roots.pop_back();
 return 0;
// https://www.spoj.com/problems/TTM/
// rollback segtree to a time stamp t
```

### 18.15 persistent seg2

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
//#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 1000000007
#define PI acos (-1)
struct node
 int item, 1, r;
 node() {}
 node(int 1, int r, int item) : 1(1), r(r), item(item) {}
int n, q;
vector<node> seq;
vector<int> roots;
void init()
  seg.resize(1);
int newleaf(int vv)
```

```
int p = seg.size();
  seg.pb(node(0, 0, vv));
  return p;
int newpar(int 1, int r)
  int p = seg.size();
  seq.pb(node(l, r, seq[l].item + seq[r].item));
  return p;
int upd(int i, int 1, int r, int pos)
  if (1 == r)
   return newleaf(seg[i].item + 1);
  int mid = (1 + r) >> 1;
  if (pos <= mid)</pre>
    return newpar(upd(seg[i].1, 1, mid, pos), seg[i].r);
  return newpar(seg[i].1, upd(seg[i].r, mid + 1, r, pos));
int build(int 1, int r)
  if (1 == r)
   return newleaf(0);
  int mid = (1 + r) >> 1;
  return newpar(build(1, mid), build(mid + 1, r));
int qry(int v1, int vr, int 1, int r, int k)
  if (1 == r)
   return 1;
  int mid = (1 + r) >> 1;
  int c = seg[seg[vr].1].item - seg[seg[vl].1].item;
  if (c >= k)
    return qry(seg[v1].1, seg[vr].1, 1, mid, k);
  return qry(seg[v1].r, seg[vr].r, mid + 1, r, k - c);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  cin >> n >> q;
  vector<int> v(n);
  set<int> vals:
  for (int i = 0; i < n; i++)
    cin >> v[i];
   vals.insert(v[i]);
  int mx = 1;
  map<int, int> mp, mpr;
  for (auto const &i : vals)
   mpr[mx] = i;
   mx++;
 init();
  roots.pb(build(0, mx));
  for (auto const &i : v)
    roots.pb(upd(roots.back(), 0, mx, mp[i]));
  while (q--)
    char c;
    cin >> c;
    if (c == 'Q')
      int 1, r, k;
     cin >> 1 >> r >> k;
      1--, r--;
      cout << mpr[qry(roots[1], roots[r + 1], 0, mx, k)] << endl;
    else
      int x;
      cin >> x;
      swap(v[x], v[x + 1]);
      int a = upd(roots[x], 0, mx, mp[v[x]]);
      int b = upd(a, 0, mx, mp[v[x + 1]]);
```

```
roots[x + 1] = a, roots[x + 2] = b;
}
return 0;
}
// https://neps.academy/br/exercise/127
// queries de k-esimo menor em um range
// e fazer um swap entre v[i] e v[i + 1]
```

### 18.16 rmq

```
#include <bits/stdc++.h>
using namespace std;
#define lli long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 998244353
struct rmq
 vector<int> v:
 bool is_max;
 int n;
 static const int b = 30;
 vector<int> mask, t;
 int op(int x, int y)
   if (is_max)
     return v[x] >= v[y] ? x : y;
    return v[x] \leftarrow v[y] ? x : y;
 int msb(int x) { return builtin clz(1) - builtin clz(x); }
 int small(int r, int sz = b) { return r - msb(mask[r] & ((1 << sz) - 1)); }</pre>
 rmq() {}
 rmq(vector < int > &v_, bool flag) : v(v_), n(v.size()), mask(n), t(n), is_max()
       flag)
    for (int i = 0, at = 0; i < n; mask[i++] = at |= 1)</pre>
     at = (at << 1) & ((1 << b) - 1);
      while (at and op(i - msb(at \& -at), i) == i)
       at ^= at & -at;
    for (int i = 0; i < n / b; i++)</pre>
     t[i] = small(b * i + b - 1);
    for (int j = 1; (1 << j) <= n / b; j++)
      for (int i = 0; i + (1 << j) <= n / b; i++)
        t[n / b * j + i] = op(t[n / b * (j - 1) + i], t[n / b * (j - 1) + i + (1)
             << (j - 1))]);
  int qry(int 1, int r)
   if (r - 1 + 1 \le b)
     return small(r, r - 1 + 1);
    int x = 1 / b + 1, y = r / b - 1;
    if (x > y)
     return op(small(1 + b - 1), small(r));
    int j = msb(y - x + 1);
    int ans = op(small(1 + b - 1), op(t[n / b * j + x], t[n / b * j + y - (1 <<
         j) + 1]));
    return op(ans, small(r));
 int query(int 1, int r) { return v[qry(1, r)]; }
signed main()
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
```

```
int n;
 cin >> n;
  vector<int> v(n);
  for (int i = 0; i < n; i++)
   cin >> v[i];
  rmq r_min(v, 0);
  rmq r_max(v, 1);
  11i ans = 0;
    vector<pi> q;
    q.pb({0, n-1});
    while (!q.empty())
      int 1 = q.back().fir;
      int r = q.back().sec;
      int max_pos = r_max.qry(1, r);
      q.pop_back();
      11i qt = (max_pos - 1 + 1) * 111 * (r - max_pos + 1);
      ans += (v[max_pos] * 111 * qt);
      if (\max pos > 1)
        q.pb({1, max_pos - 1});
      if (max_pos < r)</pre>
        q.pb((max_pos + 1, r));
    vector<pi> q;
    q.pb({0, n - 1});
    while (!q.empty())
      int l = q.back().fir;
      int r = q.back().sec;
      int min_pos = r_min.qry(1, r);
      q.pop_back();
      11i qt = (min_pos - 1 + 1) * 111 * (r - min_pos + 1);
      ans -= (v[min_pos] * 111 * qt);
      if (min_pos > 1)
        q.pb({1, min_pos - 1});
      if (min_pos < r)</pre>
        q.pb({min_pos + 1, r});
  cout << ans << endl;</pre>
  return 0;
// https://github.com/brunomaletta/Biblioteca/blob/master/Codigo/Estruturas/rmg.
// O(n) pra buildar, query O(1)
// qry(1, r) \rightarrow retorna o indice do menor elemento no range [1, r]
// query(1, r) -> retorna o menor elemento no range [1, r]
// problema exemplo: https://codeforces.com/contest/817/problem/D
```

## 18.17 SegTree

```
#include <bits/stdc++.h>
using namespace std;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define mp make_pair
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 100001
#define MAXL 100
#define mod 1000000007
vector<int> seg;
vector<int> v;
int single(int x)
```

```
return x;
int neutral()
 return 0;
int merge(int a, int b)
 return a + b;
void update(int i, int 1, int r, int q, int x)
 if (1 == r)
   seg[i] = single(x);
   return;
 int mid = (1 + r) >> 1;
 if (q <= mid)
   update(i << 1, 1, mid, q, x);
 else
   update((i << 1) | 1, mid + 1, r, q, x);
 seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
int query(int 1, int r, int q1, int qr, int i)
 int mid = (1 + r) >> 1;
 if (1 > r || 1 > qr || r < q1)</pre>
   return neutral();
  if (1 >= q1 \&\& r <= qr)
   return seq[i];
 return merge(query(l, mid, ql, qr, i << 1), query(mid + 1, r, ql, qr, (i << 1)
        | 1));
void build(int 1, int r, int i)
 if (1 == r)
    seg[i] = single(v[1]);
   return;
 int mid = (1 + r) >> 1;
 build(1, mid, i << 1);
 build(mid + 1, r, (i << 1) | 1);
 seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n, q;
 cin >> n >> q;
 v.resize(n);
  seg.resize(4 * n);
  for (int i = 0; i < n; i++)</pre>
   cin >> v[i];
 build(0, n - 1, 1);
 while (q--)
    int 1, r;
   int t;
    cin >> t >> 1 >> r;
    if (t == 2)
     cout << query(0, n - 1, 1, r - 1, 1) << endl;
     update(1, 0, n - 1, 1, r);
```

# 18.18 Segtree2

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
```

```
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 500001
#define mod 1000000007
struct segtree
  int n;
  vector<int> seq;
  int neutral()
    return 0;
  int merge(int a, int b)
    return a + b;
  void build(vector<int> &v)
   n = 1;
    while (n < v.size())</pre>
     n <<= 1;
    seg.assign(n << 1, neutral());</pre>
    for (int i = 0; i < v.size(); i++)</pre>
     seq[i + n] = v[i];
    for (int i = n - 1; i; i--)
      seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
  void upd(int i, int value)
    seg[i += n] += value;
    for (i >>= 1; i; i >>= 1)
      seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
  int qry(int 1, int r)
    int ansl = neutral(), ansr = neutral();
    for (1 += n, r += n + 1; 1 < r; 1 >>= 1, r >>= 1)
      if (1 & 1)
       ansl = merge(ansl, seg[l++]);
      if (r & 1)
        ansr = merge(seg[--r], ansr);
    return merge(ansl, ansr);
};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  return 0;
// iterative segtree without lazy propagation
```

### 18.19 segtree2d

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
```

```
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1003
#define mod 1000000007
struct segtree2d
  int n, m;
  vector<vector<int>> seg;
  int neutral()
    return 0;
  int merge(int a, int b)
    return a + b;
  segtree2d(int nn, int mm)
    n = nn, m = mm;
    seg = vector<vector<int>>(2 * n, vector<int>(2 * m, neutral()));
  int qry(int x1, int y1, int x2, int y2)
    int ret = neutral();
    int y3 = y1 + m, y4 = y2 + m;
    for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2)
      for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /= 2)
        if (x1 % 2 == 1 and y1 % 2 == 1)
          ret = merge(ret, seg[x1][y1]);
        if (x1 \% 2 == 1 \text{ and } y2 \% 2 == 0)
          ret = merge(ret, seg[x1][y2]);
        if (x2 \% 2 == 0 \text{ and } y1 \% 2 == 1)
          ret = merge(ret, seg[x2][y1]);
        if (x2 \% 2 == 0 \text{ and } y2 \% 2 == 0)
          ret = merge(ret, seg[x2][y2]);
    return ret;
  void upd(int x, int y, int val)
    int y2 = y += m;
    for (x += n; x; x /= 2, y = y2)
      if (x >= n)
        seg[x][y] = val;
      else
        seg[x][y] = merge(seg[2 * x][y], seg[2 * x + 1][y]);
      while (y /= 2)
        seg[x][y] = merge(seg[x][2 * y], seg[x][2 * y + 1]);
  }
signed main()
  ios base::sync with stdio(false);
  cin.tie(NULL);
  int q;
  cin >> q;
  while (q--)
    int n;
    segtree2d st(n, n); // matriz NxN
    while (1)
```

```
{
    string s;
    cin >> s;
    if (s == "SET")
    {
        int a, b, c;
        cin >> a >> b >> c;
        st.upd(a, b, c);
    }
    else if (s == "SUM")
    {
        int a, b, c, d;
        cin >> a >> b >> c >> d; // c >= a e d >= b
        cout << st.qry(a, b, c, d) << endl;
    }
    else
    {
        break;
    }
}
return 0;
}
// to test: https://www.spoj.com/problems/MATSUM/</pre>
```

### 18.20 segtree lazy

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define endl '\n
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 200005
#define mod 998244353
struct segtree
 int n;
  vector<int> v;
  vector<int> seg;
  vector<int> lazy;
  segtree(int sz)
   n = sz;
    seg.assign(4 * n, 0);
   lazy.assign(4 * n, 0);
// v = vv; // for build
    // build(0, n - 1, 1); // for build
  int single(int x)
    return x;
  int neutral()
    return 0;
  int merge(int a, int b)
   return a + b;
  void add(int i, int l, int r, int diff)
```

```
seg[i] += (r - 1 + 1) * diff;
   if (1 != r)
     lazy[i << 1] += diff;
     lazy[(i << 1) | 1] += diff;</pre>
   lazy[i] = 0;
 void update(int i, int l, int r, int ql, int qr, int diff)
   if (lazy[i])
     add(i, 1, r, lazy[i]);
   if (1 > r || 1 > qr || r < ql)</pre>
     return;
   if (1 >= ql && r <= qr)
     add(i, 1, r, diff);
     return;
   int mid = (1 + r) >> 1;
   update(i << 1, 1, mid, q1, qr, diff);
   update((i << 1) | 1, mid + 1, r, ql, qr, diff);
   seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
 int query(int 1, int r, int q1, int qr, int i)
   if (lazy[i])
     add(i, 1, r, lazy[i]);
   if (1 > r || 1 > qr || r < q1)
     return neutral();
   if (l >= ql \&\& r <= qr)
     return seg[i];
   int mid = (1 + r) >> 1;
   return merge(query(1, mid, q1, qr, i << 1), query(mid + 1, r, q1, qr, (i <<</pre>
        1) | 1));
 void build(int 1, int r, int i)
   if (1 == r)
     seg[i] = single(v[1]);
     return;
   int mid = (1 + r) >> 1;
   build(1, mid, i << 1);
   build (mid + 1, r, (i << 1) | 1);
   seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
 int qry(int 1, int r)
   return query(0, n - 1, 1, r, 1);
 void upd(int 1, int r, int x)
   update(1, 0, n - 1, 1, r, x);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 int n, q;
 cin >> n >> q;
 segtree s(n);
 while (q--)
   int t;
   if (t == 1)
     int 1, r, x;
     cin >> 1 >> r >> x;
     s.upd(1, r, x);
   else
```

```
int 1, r;
    cin >> 1 >> r;
    cout << s.qry(1, r) << endl;
}
return 0;</pre>
```

## 18.21 segtree max seg sum

```
#include <bits/stdc++.h>
using namespace std;
#define int long long int
#define endl '\n
#define pb push back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 200005
#define mod 1000000007
struct item
     int sg_max, pref_max, suf_max, sg_min, pref_min, suf_min, sum;
};
struct segtree
     int n:
     vector<item> seg;
     item single(int x)
          return {max(011, x), max(011, x), max(011, x), min(011, x
                        (011, x), x;
     item neutral()
          return {0, 0, 0, 0, 0, 0, 0};
     segtree() {}
     segtree (int sz)
          seg.assign(4 * n, neutral());
     item merge(item a, item b)
          ret.sq_max = max({a.sq_max, b.sq_max, a.suf_max + b.pref_max});
           ret.pref_max = max(a.pref_max, a.sum + b.pref_max);
           ret.suf_max = max(b.suf_max, b.sum + a.suf_max);
           ret.sg_min = min({a.sg_min, b.sg_min, a.suf_min + b.pref_min});
           ret.pref_min = min(a.pref_min, a.sum + b.pref_min);
           ret.suf_min = min(b.suf_min, b.sum + a.suf_min);
          ret.sum = a.sum + b.sum;
           return ret;
     void update(int i, int 1, int r, int q, int x)
           if (1 == r)
               seg[i] = single(x);
                return;
           int mid = (1 + r) >> 1;
           if (q <= mid)
               update(i << 1, 1, mid, q, x);
               update((i << 1) | 1, mid + 1, r, q, x);
          seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
```

item query(int 1, int r, int q1, int qr, int i)

```
{
  int mid = (1 + r) >> 1;
  if (1 > r || 1 > qr || r < ql)
    return neutral();
  if (1 >= ql && r <= qr)
    return seg[i];
  return merge(query(1, mid, ql, qr, i << 1), query(mid + 1, r, ql, qr, (i << 1) | 1));
  }
};
signed main()
{
  return 0;
}
// segtree for maximum segment sum
// me enrolo pra codar toda vez, e bom deixar na lib</pre>
```

### 18.22 SegTree pa

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 100005
#define mod 1000000007
struct lazy_node
  int n, a, d;
  int sum()
    int an = a + (d * (n - 1));
    return ((a + an) * n) >> 1;
  void merge(lazy_node to_add)
    a += to add.a:
    d += to_add.d;
struct segtree
  vector<int> seq;
  vector<lazy_node> lazy;
  vector<bool> lazy_status;
  segtree(int n)
    seg.resize(4 * n);
    lazy.resize(4 * n);
    lazy_status.resize(4 * n);
    build(0, n - 1, 1);
  int single(int x)
    return x;
  int neutral()
    return 0;
  int merge(int a, int b)
```

```
return a + b;
  void add(int i, int l, int r, lazy_node to_add)
    seq[i] += to_add.sum();
    if (1 != r)
      int mid = (1 + r) >> 1;
      lazy[i << 1].merge({mid - 1 + 1, to_add.a, to_add.d});</pre>
      lazy_status[i << 1] = 1;
      int diff = (mid + 1) - 1, a = to_add.a, d = to_add.d;
      lazy[(i << 1) | 1].merge({r - (mid + 1) + 1, a + (d * diff), d});
      lazy_status[(i << 1) | 1] = 1;
    lazy[i] = {r - 1 + 1, 0, 0};
    lazy_status[i] = 0;
  void update(int i, int 1, int r, int q1, int qr, lazy_node to_add)
    if (lazy_status[i])
     add(i, 1, r, lazy[i]);
    if (1 > r || 1 > qr || r < q1)
      return;
    if (l >= ql \&\& r <= qr)
      int diff = 1 - ql, a = to_add.a, d = to_add.d;
      lazy_node curr = \{r - 1 + 1, a + (d * diff), d\};
      add(i, 1, r, curr);
      return;
    int mid = (1 + r) >> 1;
    update(i << 1, 1, mid, q1, qr, to_add);
    update((i << 1) | 1, mid + 1, r, ql, qr, to_add);
    seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
  int query(int 1, int r, int q1, int qr, int i)
    if (lazy_status[i])
     add(i, l, r, lazy[i]);
    if (1 > r || 1 > qr || r < q1)
     return neutral();
    if (1 >= q1 \&\& r <= qr)
     return seg[i];
    int mid = (1 + r) \gg 1;
    return merge(query(1, mid, q1, qr, i << 1), query(mid + 1, r, q1, qr, (i <<</pre>
  void build(int 1, int r, int i)
    seg[i] = 0;
    lazv status[i] = 0;
    lazy[i] = {r - 1 + 1, 0, 0};
    if (1 == r)
     return;
    int mid = (1 + r) >> 1;
   build(1, mid, i << 1);
   build (mid + 1, r, (i << 1) | 1);
};
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
  int n, q;
  cin >> n >> q;
  segtree s(n);
  while (q--)
    int t;
    cin >> t;
    if (t == 1)
     int 1, r, a, d;
cin >> 1 >> r >> a >> d;
      s.update(1, 0, n - 1, 1, r, {r - 1 + 1, a, d});
```

```
}
else
{
    int x;
    cin >> x;
    x--;
    cout << s.query(0, n - 1, x, x, 1) << endl;
}
return 0;
}
// queries of:
// add an arithmetic progression to a segment [1, r]
// print current value of a given element</pre>
```

## 18.23 sparsetable

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define PI acos(-1)
#define pb push_back
#define int long long int
#define pi pair<int, int>
#define pii pair<int, pair<int, pi>>
#define fir first
#define sec second
#define DEBUG 0
#define MAXN 10005
#define mod 1000000007
struct rmq
 bool is_min;
 vector<vector<int>> st;
 vector<int> log;
  int f(int a, int b) { return (is_min) ? min(a, b) : max(a, b); }
  int qry(int 1, int r) { return f(st[1][log[r-1+1]], st[r-(1 << log[r-1])]
        + 1]) + 1][log[r - 1 + 1]]); }
  rmq(vector<int> &v, bool flag)
    is_min = flag;
    int n = v.size();
    log.resize(n + 1);
    \log[1] = 0;
    for (int i = 2; i <= n; i++)
     log[i] = log[i / 2] + 1;
    int m = log[n] + 2;
    st.assign(n + 1, vector<int>(m, 0));
    for (int i = 0; i < n; i++)
     st[i][0] = v[i];
    for (int j = 1; j < m; j++)
      for (int i = 0; i + (1 << j) <= n; i++)
        st[i][j] = f(st[i][j-1], st[i+(1 << (j-1))][j-1]);
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
```

## 18.24 treap

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
// #define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1000005
#define mod 1000000007
struct treap
  int data, priority;
  int sz, lazy2;
  bool lazy;
  treap *1, *r, *parent;
int size(treap *node)
  return (!node) ? 0 : node->sz;
void recalc(treap *node)
  if (!node)
    return;
  node->sz = 1;
  node->parent = 0;
  if (node->1)
   node->sz += node->l->sz, node->l->parent = node;
  if (node->r)
   node->sz += node->r->sz, node->r->parent = node;
void lazy_propagation(treap *node)
  if (node == NULL)
   return;
  if (node->lazy2)
    if (node->1)
     node->1->lazy2 += node->lazy2;
    if (node->r)
     node->r->lazy2 += node->lazy2;
    node->data += node->lazy2;
   node \rightarrow lazy2 = 0;
  if (node->lazy)
    swap(node->1, node->r);
    if (node->1)
     node->1->lazy = !node->1->lazy;
    if (node->r)
     node->r->lazy = !node->r->lazy;
   node -> lazy = 0;
void split(treap *t, treap *&l, treap *&r, int n)
  if (!t)
   return void(l = r = 0);
  lazy_propagation(t);
  if (size(t->1) >= n)
    split(t->1, 1, t->1, n), r = t;
    split(t->r, t->r, r, n - size(t->1) - 1), 1 = t;
  recalc(t);
```

```
void merge(treap *&t, treap *1, treap *r)
  lazy_propagation(1);
  lazy_propagation(r);
  if (!1)
   t = r;
  else if (!r)
  else if (l->priority > r->priority)
    merge (1->r, 1->r, r), t = 1;
  else
   merge (r->1, 1, r->1), t = r;
  recalc(t);
void troca(treap *&t, int 1, int r, int 11, int rr)
  treap *a0, *a1, *b0, *b1, *c0, *c1, *d0, *d1;
  split(t, a0, a1, 1);
  split(a1, b0, b1, r - 1 + 1);
  11 -= (r + 1);
  rr -= (r + 1);
  split(b1, c0, c1, l1);
  split(c1, d0, d1, rr - 11 + 1);
  merge(t, a0, d0);
  merge(t, t, c0);
 merge(t, t, b0);
 merge(t, t, d1);
void add(treap *&t, int 1, int r)
  treap *a0, *a1, *b0, *b1;
  split(t, a0, a1, 1);
  split(a1, b0, b1, r - 1 + 1);
  b0 \rightarrow lazy = 1;
  b0->lazy2 += 1;
 merge(t, a0, b0);
  merge(t, t, b1);
void solve(int x)
  x = x % 26;
  char c = x + 'a';
 cout << c;
void dfs(treap *t)
 if (!t)
   return;
  lazy_propagation(t);
  dfs(t->1);
  solve(t->data);
 dfs(t->r);
treap *create_node(int data, int priority)
  treap *ret = new treap;
 ret->data = data;
  ret->priority = priority;
  ret \rightarrow 1 = 0;
  ret->r = 0;
  ret->sz = 1;
  ret->lazy = 0;
  ret \rightarrow lazy2 = 0;
  ret->parent = 0;
void goup(treap *&ans, treap *t) // vai pra raiz da arvore
  if (!t->parent)
   ans = t;
    return;
  goup(ans, t->parent);
signed main()
  ios_base::sync_with_stdio(false);
  cin.tie(NULL);
```

```
srand(time(NULL));
 int q;
  cin >> q;
 while (q--)
   int n, m;
   string s;
    cin >> s >> m;
   n = s.size();
    treap *t = 0;
    for (auto const &i : s)
      int x = i - 'a';
     merge(t, t, create_node(x, rand()));
    while (m--)
     int a, b, c, d;
     cin >> a >> b >> c >> d;
      a--, b--, c--, d--;
     add(t, a, b);
     add(t, c, d);
     troca(t, a, b, c, d);
   dfs(t);
   cout << endl;
 return 0;
// https://vjudge.net/contest/478186#problem/E
// - lazy propagation
// - reverse range with lazy propagation
// - swap ranges with equal lenght
// extra:
// - save node parent
```

## 18.25 treap2

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define endl '\n'
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 1000005
#define mod 1000000007
vector<int> ans;
struct treap
  int data, priority;
  int sz;
  bool lazy;
  treap *1, *r;
int size(treap *node)
  return (!node) ? 0 : node->sz;
void recalc(treap *node)
  if (!node)
```

```
return;
  node->sz = 1:
  if (node->1)
   node->sz += node->l->sz;
  if (node->r)
   node->sz += node->r->sz;
void lazy_propagation(treap *node)
 if (!node || !(node->lazy))
    return;
  swap(node->1, node->r);
  if (node->1)
    node->l->lazy ^= 1;
  if (node->r)
   node \rightarrow r \rightarrow lazy = 1;
 node -> lazy = 0;
void merge(treap *&t, treap *1, treap *r)
  lazy_propagation(1);
  lazy_propagation(r);
  if (!1)
   t = r:
  else if (!r)
    t = 1:
  else if (1->priority > r->priority)
   merge (1->r, 1->r, r), t = 1;
  else
   merge(r->1, 1, r->1), t = r;
 recalc(t);
void split(treap *t, treap *&l, treap *&r, int n)
   return void(1 = r = 0);
  lazy_propagation(t);
  if (size(t->1) >= n)
    split(t->1, 1, t->1, n), r = t;
   split(t->r, t->r, r, n - size(t->1) - 1), 1 = t;
  recalc(t);
void reverse(treap *&t, int 1, int r)
  treap *a0, *a1, *b0, *b1;
  split(t, a0, a1, 1);
  split(a1, b0, b1, r - 1 + 1);
  b0->1azy = 1;
  merge(t, a0, b0);
  merge(t, t, b1);
void shift(treap *&t, int 1, int r)
  treap *a0, *a1, *b0, *b1, *c0, *c1;
  split(t, a0, a1, 1);
  \tilde{split} (al, b0, b1, r - 1 + 1);
  split(b0, c0, c1, r - 1);
  merge(t, a0, c1);
 merge(t, t, c0);
  merge(t, t, b1);
void dfs(treap *t)
  if (!t)
   return;
  lazy_propagation(t);
  dfs(t->1);
  ans.pb(t->data);
  dfs(t->r);
treap *create_node(int data, int priority)
 treap *ret = new treap;
  ret->data = data;
 ret->priority = priority;
 ret -> \bar{1} = 0;
  ret->r = 0;
```

```
ret->sz = 1;
 ret->lazy = 0;
  return ret;
signed main()
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 srand(time(NULL));
  treap *t = 0;
 int n, m, q;
 cin >> n >> q >> m;
  for (int i = 0; i < n; i++)</pre>
   int k;
   cin >> k:
   merge(t, t, create_node(k, rand()));
  while (q--)
   int ty, 1, r;
   cin >> ty >> 1 >> r;
   1--, r--;
    (ty == 1) ? shift(t, 1, r) : reverse(t, 1, r);
 dfs(t);
  while (m--)
   int i;
   cin >> i;
   i--;
   cout << ans[i] << " ";
 cout << endl;
 return 0;
```

## 19 Theorems and Formulas

### 19.1 binomial theorem

```
## Binomial Theorem
### Theorem
(x + y)^n = \sum_{k=0}^{n} \{n \in k\} x^{n-k}y^{k}
in addition, we have:
(x
  -y)^n = \sum_{k=0}^{n} (-1)^k \{n \in k\} x^{n-k}y^{k}
55
(1 + x)^n = \sum_{k=0}^{n} \{n \in k\} x^{k}
$$
### Cool Problem
[Fibonacci Fever] (https://codeforces.com/gym/104412/problem/F)
Given n and k you're asked to compute (mod $10^9 + 7$):
\sum_{i=1}^{n} f_i^k
$$
where $f n$ is the n-th fibonacci number.
Recall that:
$$
```

## 19.2 chicken mcnugget

```
## Chicken McNugget Theorem
For any two coprime numbers (n > 0, m > 0), the greatest integer that cannot be
     written in the form:
an + bm, (a >= 0, b >= 0)
is (n \times m) - n - m
## Consequence of the theorem
That there are exactly ((n-1) \  \  \  \  \  \  \  \  \  \  \  ) / 2 positive integers which cannot
    be expressed in the form an + bm, (a \ge 0, b \ge 0)
## Generalization
If n and m are not coprime, so all numbers that are not multiples of qcd(n, m)
     cannot be expressed in the form an + bm, (a >= 0, b >= 0)
in addition, you can consider n = (n / gcd(n, m)) and m = (m / gcd(n, m)), to
     find how many multiples of gcd(n, m) cannot be expressed, or to find the
     greatest multiple of gcd(n, m) that cannot be expressed
## Considering a > 0, b > 0
Considering (n > 0, m > 0), n and m are coprime:
let y = ((n \times m) + min(n, m)) - 1
The number of positive integers which cannot be expressed increases by (y / n)
The number of positive integers which cannot be expressed increases by (y / m)
you must not count the multiples of (n \ m) more than once, just decrease
     number of positive integers which cannot be expressed by (y / (n \times m))
## Problems
```

## 19.3 graph notes

problem/I)

- ## Bipartite Graph
- A bipartite graph is a graph that does not contain any odd-length cycles.

- [Forming Compounds] (https://codeforces.com/group/XrhoJtxCjm/contest/422716/

- ## Directed acyclic graph (DAG)
- Is a directed graph with no directed cycles.
- ## Independent Set
- Is a set of vertices in a graph, no two of which are adjacent. That is, it is a set S of vertices such that for every two vertices in S, there is no edge connecting the two.
- ## Clique
- Is a subset of vertices of an undirected graph such that every two distinct vertices in the clique are adjacent.
- ## Vertex Cover
- Is a set of vertices that includes at least one endpoint of every edge of the graph.
- ## Edge Cover
- Is a set of edges such that every vertex of the graph is incident to at least one edge of the set.
- ## Path Cover
- Given a directed graph G = (V, E), a path cover is a set of directed paths such that every vertex v belongs to at least one path.
- ## Koning's Theorem
- In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover.
- ## Properties
- Every tree is a bipartite graph.
- Any NxM grid is a bipartite graph.
- A set of vertices is a vertex cover if and only if its complement is an independent set.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- In bipartite graphs, the size of the minimum edge cover is equal to the size of the maximum independent set
- In bipartite graphs, the size of the minimum edge cover plus the size of the minimum vertex cover is equal to the number of vertices.
- In bipartite graphs, maximum clique size is two.
- ## Min-cut
- The smallest total weight of the edges which if removed would disconnect the source from the sink.
- ## Max-flow min-cut theorem
- In a flow network, the maximum amount of flow passing from the source to the sink is equal to the total weight of the edges in a minimum cut.
- ## Maximum flow with vertex capacities
- In other words, the amount of flow passing through a vertex cannot exceed its capacity. To find the maximum flow, we can transform the problem into the maximum flow problem by expanding the network. Each vertex v is replaced by v-in and v-out, where v-in is connected by edges going into v and v-out is connected to edges coming out from v. Then assign capacity c(v) to the edge connecting v-in and v-out.
- ## Undirected edge-disjoint paths problem
- We are given an undirected graph G = (V, E) and two vertices s and t, and we have to find the maximum number of edge-disjoint s-t paths in G.
- ## Undirected vertex-disjoint paths problem
- We are given an undirected graph G = (V, E) and two vertices s and t, and we have to find the maximum number of vertex-disjoint (except for s and t) paths in G.

#### ## Menger's theorem

The maximum number of edge-disjoint s-t paths in an undirected graph is equal to the minimum number of edges in an s-t cut-set.

#### ## Undirected vertex-disjoint paths solution

We can construct a network N=(V,E) from G with vertex capacities, where the capacities of all vertices and all edges are 1. Then the value of the maximum flow is equal to the maximum number of independent paths from s to t.

#### ## Minimum vertex-disjoint path cover in directed acyclic graph (DAG)

Given a directed acyclic graph G=(V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V. We can construct a bipartite graph G' from G. Each vertex v is replaced by v-in and v-out, where v-in is connected by edges going into v and v-out is connected to edges coming out from v. Then it can be shown that G' has a matching M of size m if and only if G has a vertex-disjoint path cover C of containing m edges and n-m paths.

#### ## Minimum general path cover in directed acyclic graph (DAG)

A general path cover is a path cover where a vertex can belong to more than one path. A minimum general path cover may be smaller than a minimum vertex-disjoint path cover. A minimum general path cover can be found almost like a minimum vertex-disjoint path cover. It suffices to add some new edges to the matching graph so that there is an edge a - b always when there is a path from a to b in the original graph.

#### ## Dilworth's theorem and maximum antichain

An antichain is a set of nodes of a graph such that there is no path from any node to another node using the edges of the graph. Dilworth's theorem states that in a directed acyclic graph, the size of a minimum general path cover equals the size of a maximum antichain.

Or in other words: For a DAG G that if has edges from vertex  $i \to vertex\ j$  and vertex  $j \to k$ , then it also has a edge from vertex  $i \to vertex\ k$ , the size of a minimum path cover is equal to the size of a maximum independent set.

### ## Maximum weighted antichain

(https://atcoder.jp/contests/abc354/tasks/abc354\_g) In this problem, each vertex has a cost a[i]. The cost of an antichain is equal to the sum of the costs of the vertices present in it. We need to find the maximum cost of a antichain. We can construct the same bipartite of the maximum antichain problem from a dag G, these edges have an infinite capacity. We also need to create a source vertex and a sink, and we need to add edges source -> v-in with capacity a[v] and v-out -> sink with capacity a[v]. The answer is equal to the sum of all a[i] minus the maximum flow on this network.

#### ## Hall's Theorem

Hall's theorem can be used to find out whether a bipartite graph has a matching that contains all left or right nodes. Assume that we want to find a matching that contains all left nodes. Let X be any set of left nodes and let f(X) be the set of their neighbors. According to Hall's theorem, a matching that contains all left nodes exists exactly when for each X, the condition |X| <= |f(X)| holds.

#### ## References

- [Competitive Programmer's Handbook] (https://cses.fi/book/book.pdf)
- [(Graph Theory) Wikipedia](https://en.wikipedia.org/wiki/Graph\_theory)
- [(Medium Article) Solving Minimum Path Cover on a DAG](https://towardsdatascience.com/solving-minimum-path-cover-on-a-dag-21b16cal1ac0)

#### ## Extra (Getting Confidence Trick)

[2019-2020 ACM-ICPC Brazil Subregional Programming Contest, problem G] (https:// codeforces.com/gym/102346/problem/G)

If you need to maximize a number x = (a \* b \* c \* ...), then you can write it as  $x = (e^{\hat{1}}\log(a) * e^{\hat{1}}\log(b) * e^{\hat{1}}\log(c) * ...)$ , and then the number is  $x = e^{\hat{1}}\log(a) + \log(b) + \log(c) + ...)$ , and the problem now becomes a problem of maximizing the sum of  $(\log(a) + \log(b) + \log(c) + ...) \cdot p/>$ 

```
Use exp() and log() C++ functions :)
```

## 19.4 manhattan and chebyshev

# 20 ufmg forked

# 21 ufmg forked/DP

## 21.1 dcDp

```
// Divide and Conquer DP
// Particiona o array em k subarrays
// minimizando o somatorio das queries
// O(k n log n), assumindo quer query(l, r) eh O(1)
11 dp[MAX][2];
void solve(int k, int l, int r, int lk, int rk) {
        if (1 > r) return;
        int m = (1+r)/2, p = -1;
        auto& ans = dp[m][k&1] = LINF;
        for (int i = max(m, lk); i <= rk; i++) {</pre>
                ll at = dp[i+1][~k&1] + query(m, i);
                if (at < ans) ans = at, p = i;
        solve(k, l, m-1, lk, p), solve(k, m+1, r, p, rk);
11 DC(int n, int k) {
        dp[n][0] = dp[n][1] = 0;
        for (int i = 0; i < n; i++) dp[i][0] = LINF;</pre>
        for (int i = 1; i \le k; i++) solve(i, 0, n-i, 0, n-i);
        return dp[0][k&1];
```

### 21.2 lcs

```
// Longest Common Subsequence
//
// Computa a LCS entre dois arrays usando
// o algoritmo de Hirschberg para recuperar
//
// O(n*m), O(n+m) de memoria
int lcs_s[MAX], lcs_t[MAX];
```

```
int dp[2][MAX];
// dp[0][j] = max lcs(s[li...ri], t[lj, lj+j])
void dp_top(int li, int ri, int lj, int rj) {
        memset(dp[0], 0, (rj-lj+1)*sizeof(dp[0][0]));
        for (int i = li; i <= ri; i++) {
                 for (int j = rj; j >= 1j; j--)
                         \frac{dp[0][j'-1j] = \max(dp[0][j-1j],}{(lcs_s[i] == lcs_t[j]) + (j > 1j ? dp[0][j-1-1j] : 0))}
                 for (int j = 1j+1; j <= rj; j++)</pre>
                          dp[0][j-1j] = max(dp[0][j-1j], dp[0][j-1-lj]);
// dp[1][j] = max lcs(s[li...ri], t[lj+j, rj])
void dp_bottom(int li, int ri, int lj, int rj)
        memset(dp[1], 0, (rj-lj+1)*sizeof(dp[1][0]));
        for (int i = ri; i >= li; i--) {
                for (int j = 1j; j <= rj; j++)
    dp[1][j - 1j] = max(dp[1][j - 1j],</pre>
                          (lcs_s[i] == lcs_t[j]) + (j < rj ? dp[1][j+1 - lj] : 0)
                 for (int j = rj-1; j >= 1j; j--)
                          dp[1][j-1j] = max(dp[1][j-1j], dp[1][j+1-1j]);
void solve(vector<int>& ans, int li, int ri, int lj, int rj) {
        if (li == ri) {
                 for (int j = lj; j <= rj; j++)</pre>
                          if (lcs_s[li] == lcs_t[j]){
                                  ans.push_back(lcs_t[j]);
                 return;
        if (lj == rj) {
                 for (int i = li; i <= ri; i++) {</pre>
                         if (lcs_s[i] == lcs_t[lj]){
                                  ans.push_back(lcs_s[i]);
                                  break:
                 return:
        int mi = (li+ri)/2;
        dp_top(li, mi, lj, rj), dp_bottom(mi+1, ri, lj, rj);
        int j_{-} = 0, mx = -1;
        for (int j = lj-1; j <= rj; j++) {</pre>
                 int val = 0;
                 if (j >= lj) val += dp[0][j - lj];
                 if (j < rj) val += dp[1][j+1 - lj];</pre>
                if (val >= mx) mx = val, j_ = j;
        if (mx == -1) return;
        solve(ans, li, mi, lj, j_), solve(ans, mi+1, ri, j_+1, rj);
vector<int> lcs(const vector<int>& s, const vector<int>& t) {
        for (int i = 0; i < s.size(); i++) lcs_s[i] = s[i];</pre>
        for (int i = 0; i < t.size(); i++) lcs_t[i] = t[i];</pre>
        vector<int> ans;
        solve(ans, 0, s.size()-1, 0, t.size()-1);
        return ans;
```

### 21.3 mochila

```
// Mochila
//
// Resolve mochila, recuperando a resposta
//
```

```
// O(n * cap), O(n + cap) de memoria
int v[MAX], w[MAX]; // valor e peso
int dp[2][MAX_CAP];
// DP usando os itens [1, r], com capacidade = cap
void get_dp(int x, int 1, int r, int cap)
         memset(dp[x], 0, (cap+1)*sizeof(dp[x][0]));
for (int i = 1; i \le r; i++) for (int j = cap; j \ge 0; j--)
                  if (j - w[i] \ge 0) dp[x][j] = max(dp[x][j], v[i] + dp[x][j - w[i])
void solve(vector<int>& ans, int 1, int r, int cap) {
         if (1 == r) {
                  if (w[1] <= cap) ans.push_back(1);</pre>
                  return;
         int m = (1+r)/2;
         get_dp(0, 1, m, cap), get_dp(1, m+1, r, cap);
int left_cap = -1, opt = -INF;
         for (int j = 0; j <= cap; j++)
    if (int at = dp[0][j] + dp[1][cap - j]; at > opt)
                           opt = at, left_cap = j;
         solve(ans, 1, m, left_cap), solve(ans, m+1, r, cap - left_cap);
vector<int> knapsack(int n, int cap) {
         vector<int> ans;
         solve(ans, 0, n-1, cap);
         return ans;
```

### 21.4 sosDP

### 21.5 subsetSum

```
// Subset sum
//
// Retorna max(x <= t tal que existe subset de w que soma x)
//
// O(n * max(w))
// O(max(w)) de memoria

int subset_sum(vector<int> w, int t) {
    int pref = 0, k = 0;
    while (k < w.size() and pref + w[k] <= t) pref += w[k++];
    if (k == w.size()) return pref;</pre>
```

# 22 ufmg forked/Estruturas

### 22.1 bit

```
// BIT
// BIT de soma 0-based
// upper_bound(x) retorna o menor p tal que pref(p) > x
// Complexidades:
// build - O(n)
// update - 0(log(n))
// query - O(\log(n))
// upper_bound - O(log(n))
struct Bit {
         vector<11> bit;
         Bit (int _n=0) : n(_n), bit (n + 1) {}
         Bit (vector \leq int > \& v) : n(v.size()), bit (n + 1) {
                 for (int i = 1; i <= n; i++) {</pre>
                          bit[i] += v[i - 1];
                          int j = i + (i & -i);
if (j <= n) bit[j] += bit[i];</pre>
         void update(int i, ll x) { // soma x na posicao i
                 for (i++; i <= n; i += i & -i) bit[i] += x;
         11 pref(int i) { // soma [0, i]
                 for (i++; i; i -= i & -i) ret += bit[i];
                 return ret;
         11 query(int 1, int r) { // soma [1, r]
                 return pref(r) - pref(1 - 1);
         int upper_bound(ll x) {
                 int p = 0;
                 for (int i = __lg(n); i+1; i--)
                          if (p + (1 << i) <= n \text{ and } bit[p + (1 << i)] <= x)
                                  x -= bit[p += (1 << i)];
                 return p;
};
```

### 22.2 bit2d

```
// BIT 2D
//
// BIT de soma, update incrementa posicao
// Tem que construir com um vetor com todos os pontos
// que vc quer um dia atualizar (os pontos q vc vai chamar update)
```

```
// Complexidades:
// construir - O(n log(n))
// update e query - O(log^2(n))
template<class T = int> struct bit2d {
        vector<T> X;
        vector<vector<T>> Y, t;
        int ub(vector<T>& v, T x) {
                return upper_bound(v.begin(), v.end(), x) - v.begin();
        bit2d(vector<pair<T, T>> v) {
                for (auto [x, y] : v) X.push_back(x);
                sort(X.begin(), X.end());
                X.erase(unique(X.begin(), X.end()), X.end());
                t.resize(X.size() + 1);
                Y.resize(t.size());
                sort(v.begin(), v.end(), [](auto a, auto b) {
                        return a.second < b.second; });</pre>
                for (auto [x, y] : v) for (int i = ub(X, x); i < t.size(); i +=
                     i&-i)
                        if (!Y[i].size() or Y[i].back() != y) Y[i].push_back(y);
                for (int i = 0; i < t.size(); i++) t[i].resize(Y[i].size() + 1);</pre>
        }
        void update(T x, T y, T v) {
                for (int i = ub(X, x); i < t.size(); i += i&-i)</pre>
                        for (int j = ub(Y[i], y); j < t[i].size(); j += j&-j) t[
                             i][j] += v;
        T query(T x, T y) {
                T ans = 0;
                for (int i = ub(X, x); i; i = i&-i)
                        for (int j = ub(Y[i], y); j; j = j\&-j) ans += t[i][j];
        T query (T x1, T y1, T x2, T y2) {
                return query (x2, y2) -query (x2, y1-1) -query (x1-1, y2) +query (x1-1,
                      y_{1}-1);
};
```

## 22.3 bitRange

```
// BIT com update em range
// Operacoes 0-based
// query(l, r) retorna a soma de v[l..r]
// update(1, r, x) soma x em v[1..r]
// Complexidades:
// build - O(n)
// query - O(\log(n))
// update - 0(log(n))
namespace bit {
        11 bit[2][MAX+2];
        int n;
        void build(int n2, int* v) {
                 for (int i = 1; i <= n; i++)</pre>
                         bit[1] [min(n+1, i+(i&-i))] += bit[1][i] += v[i-1];
        11 get(int x, int i) {
                 11 ret = 0;
                for (; i; i -= i&-i) ret += bit[x][i];
                return ret;
        void add(int x, int i, ll val) {
                 for (; i <= n; i += i&-i) bit[x][i] += val;</pre>
```

```
}
11 get2(int p) {
            return get(0, p) * p + get(1, p);
}
11 query(int 1, int r) {
            return get2(r+1) - get2(1);
}
void update(int 1, int r, 11 x) {
            add(0, 1+1, x), add(0, r+2, -x);
            add(1, 1+1, -x*1), add(1, r+2, x*(r+1));
}
};
```

### 22.4 bitSortTree

```
// BIT-Sort Tree
// Tipo uma MergeSort Tree usando Bit
// Apesar da complexidade ser pior, fica melhor na pratica.
// query(1, r, k) retorna o numero de elementos menores que k
// no intervalo [1, r]
// Usa O(n log(n)) de memoria
// Complexidades:
// construir - O(n log^2(n))
// query - O(log^2(n))
template<typename T> struct ms_bit {
    int n;
    vector<vector<T>> bit;
    ms_bit(vector<T>& v) : n(v.size()), bit(n+1) {
        for (int i = 0; i < n; i++)
            for (int j = i+1; j \le n; j += j\&-j)
               bit[j].push_back(v[i]);
        for (int i = 1; i <= n; i++)</pre>
            sort(bit[i].begin(), bit[i].end());
    int p_query(int i, T k) {
        int ret = 0;
        for (i++; i; i -= i\&-i)
            ret += lower_bound(bit[i].begin(), bit[i].end(), k) - bit[i].begin()
        return ret;
    int query(int 1, int r, T k) {
        return p_query(r, k) - p_query(l-1, k);
};
```

### 22.5 cht

```
// Convex Hull Trick Estatico
// adds tem que serem feitos em ordem de slope
// queries tem que ser feitas em ordem de x
//
// add O(1) amortizado, get O(1) amortizado

struct CHT {
    int it;
    vector<11> a, b;
    CHT():it(0){}
    ll eval(int i, ll x){
        return a[i]*x + b[i];
    }
    bool useless(){
        int sz = a.size();
        int r = sz-1, m = sz-2, l = sz-3;
```

```
#warning cuidado com overflow!
                return (b[1] - b[r]) * (a[m] - a[1]) <
                         (b[1] - b[m]) * (a[r] - a[1]);
        void add(l1 A, l1 B) {
                a.push_back(A); b.push_back(B);
                while (!a.empty()){
                         if ((a.size() < 3) || !useless()) break;</pre>
                         a.erase(a.end() - 2);
                        b.erase(b.end() - 2);
                it = min(it, int(a.size()) - 1);
        11 get(11 x){
                while (it+1 < a.size()) {
                         if (eval(it+1, x) > eval(it, x)) it++;
                         else break;
                return eval(it, x);
};
```

### 22.6 chtDinamico

```
// Convex Hull Trick Dinamico
// para double, use LINF = 1/.0, div(a, b) = a/b
// update(x) atualiza o ponto de intersecao da reta x
// overlap(x) verifica se a reta x sobrepoe a proxima
// add(a, b) adiciona reta da forma ax + b
// query(x) computa maximo de ax + b para entre as retas
// O(log(n)) amortizado por insercao
// O(log(n)) por query
struct Line {
        mutable 11 a, b, p;
        bool operator<(const Line& o) const { return a < o.a; }</pre>
        bool operator<(11 x) const { return p < x; }</pre>
};
struct dynamic hull : multiset<Line, less<>>> {
        ll div(ll a, ll b) {
                return a / b - ((a ^ b) < 0 and a % b);
        void update(iterator x) {
                if (next(x) == end()) x -> p = LINF;
                else if (x->a == next(x)->a) x->p = x->b >= next(x)->b ? LINF :
                     -LINF:
                else x - p = div(next(x) - b - x - b, x - a - next(x) - a);
        bool overlap(iterator x) {
                update(x);
                if (next(x) == end()) return 0;
                if (x->a == next(x)->a) return x->b >= next(x)->b;
                return x->p >= next(x)->p;
        void add(ll a, ll b) {
                auto x = insert({a, b, 0});
                while (overlap(x)) erase(next(x)), update(x);
                if (x != begin() and !overlap(prev(x))) x = prev(x), update(x);
                while (x != begin() and overlap(prev(x)))
                        x = prev(x), erase(next(x)), update(x);
        11 query(11 x) {
                assert(!empty());
                auto 1 = *lower_bound(x);
#warning cuidado com overflow!
                return 1.a * x + 1.b;
};
```

### 22.7 dsu

```
// DSU
// Une dois conjuntos e acha a qual conjunto um elemento pertence por seu id
// find e unite: O(a(n)) ~= O(1) amortizado
struct dsu {
        vector<int> id, sz;
        dsu(int n) : id(n), sz(n, 1) { iota(id.begin(), id.end(), 0); }
        int find(int a) { return a == id[a] ? a : id[a] = find(id[a]); }
        void unite(int a, int b) {
                a = find(a), b = find(b);
                if (a == b) return;
                if (sz[a] < sz[b]) swap(a, b);</pre>
                sz[a] += sz[b], id[b] = a;
};
// DSU de bipartido
// Une dois vertices e acha a qual componente um vertice pertence
// Informa se a componente de um vertice e bipartida
// find e unite: O(log(n))
struct dsu {
        vector<int> id, sz, bip, c;
        dsu(int n) : id(n), sz(n, 1), bip(n, 1), c(n) {
                iota(id.begin(), id.end(), 0);
        int find(int a) { return a == id[a] ? a : find(id[a]); }
        int color(int a) { return a == id[a] ? c[a] : c[a] ^ color(id[a]); }
        void unite(int a, int b) {
                bool change = color(a) == color(b);
                a = find(a), b = find(b);
                if (a == b) {
                        if (change) bip[a] = 0;
                        return;
                if (sz[a] < sz[b]) swap(a, b);
                if (change) c[b] = 1;
                sz[a] += sz[b], id[b] = a, bip[a] &= bip[b];
};
// DSU Persistente
// Persistencia parcial, ou seja, tem que ir
// incrementando o 't' no une
// find e unite: O(log(n))
struct dsu {
        vector<int> id, sz, ti;
        dsu(int n) : id(n), sz(n, 1), ti(n, -INF) {
                iota(id.begin(), id.end(), 0);
        int find(int a, int t) {
                if (id[a] == a or ti[a] > t) return a;
                return find(id[a], t);
        void unite(int a, int b, int t) {
                a = find(a, t), b = find(b, t);
```

```
if (a == b) return;
                if (sz[a] < sz[b]) swap(a, b);</pre>
                sz[a] += sz[b], id[b] = a, ti[b] = t;
};
// DSU com rollback
// checkpoint(): salva o estado atual de todas as variaveis
// rollback(): retorna para o valor das variaveis para
// o ultimo checkpoint
// Sempre que uma variavel muda de valor, adiciona na stack
// find e unite: O(log(n))
// checkpoint: O(1)
// rollback: O(m) em que m e o numero de vezes que alguma
// variavel mudou de valor desde o ultimo checkpoint
struct dsu {
        vector<int> id, sz;
        stack<stack<pair<int&, int>>> st;
        dsu(int n) : id(n), sz(n, 1) {
                iota(id.begin(), id.end(), 0), st.emplace();
        void save(int &x) { st.top().emplace(x, x); }
        void checkpoint() { st.emplace(); }
        void rollback() {
                while(st.top().size()) {
                        auto [end, val] = st.top().top(); st.top().pop();
                        end = val;
                st.pop();
        int find(int a) { return a == id[a] ? a : find(id[a]); }
        void unite(int a, int b)
                a = find(a), b = find(b);
                if (a == b) return;
                if (sz[a] < sz[b]) swap(a, b);</pre>
                save(sz[a]), save(id[b]);
                sz[a] += sz[b], id[b] = a;
};
```

### 22.8 lichao

```
// Li-Chao Tree
// Adiciona retas (ax+b), e computa o minimo entre as retas
// em um dado 'x'
// Cuidado com overflow!
// Se tiver overflow, tenta comprimir o 'x' ou usar
// convex hull trick
// O(log(MA-MI)), O(n) de memoria
template<11 MI = 11(-1e9), 11 MA = 11(1e9) > struct lichao {
        struct line {
                11 a, b;
                array<int, 2> ch;
                line(11 a_{-} = 0, 11 b_{-} = LINF):
                        a(a_), b(b_), ch(\{-1, -1\})  {}
                11 operator ()(11 x) { return a*x + b; }
        vector<line> ln;
        int ch(int p, int d) {
                if (ln[p].ch[d] == -1) {
                        ln[p].ch[d] = ln.size();
                        ln.emplace_back();
```

```
return ln[p].ch[d];
        lichao() { ln.emplace back(); }
        void add(line s, ll l=MI, ll r=MA, int p=0) {
                11 m = (1+r)/2;
                bool L = s(1) < ln[p](1);
                bool M = s(m) < ln[p](m);
                bool R = s(r) < ln[p](r);
                if (M) swap(ln[p], s), swap(ln[p].ch, s.ch);
                if (s.b == LINF) return;
                if (L != M) add(s, 1, m-1, ch(p, 0));
                else if (R != M) add(s, m+1, r, ch(p, 1));
        11 query(int x, 11 1=MI, 11 r=MA, int p=0) {
                11 m = (1+r)/2, ret = ln[p](x);
                if (ret == LINF) return ret;
                if (x < m) return min(ret, query(x, 1, m-1, ch(p, 0)));
                return min(ret, query(x, m+1, r, ch(p, 1)));
};
```

## 22.9 lichaoLazy

```
// Li-Chao Tree - Lazy
// Sendo N = MA-MI:
// insert({a, b}) minimiza tudo com ax+b - O(log N)
// insert(\{a, b\}, 1, r) minimiza com ax+b no range [1, r] - O(\log^2 N)
// shift({a, b}) soma ax+b em tudo - O(1)
// shift({a, b}, 1, r) soma ax+b no range [1, r] - O(log^2 N)
// query(x) retorna o valor da posicao x - O(log N)
// No inicio eh tudo LINF, se inserir {0, 0} fica tudo 0
// O(n log N) de memoria ; O(n) de memoria se nao usar as operacoes de range
template<int MI = int(-1e9), int MA = int(1e9)> struct lichao {
        struct line {
                 11 a, b;
                 11 la, lb; // lazy
                 array<int, 2> ch;
                 line(ll a_ = 0, ll b_ = LINF):
    a(a_), b(b_), la(0), lb(0), ch({-1, -1}) {}
                 11 operator ()(11 x) { return a*x + b; }
        vector<line> ln;
        int ch(int p, int d) {
                 if (\ln[p].ch[d] == -1) {
                          ln[p].ch[d] = ln.size();
                          ln.emplace_back();
                 return ln[p].ch[d];
        lichao() { ln.emplace_back(); }
        void prop(int p, int 1, int r) {
    if (ln[p].la == 0 and ln[p].lb == 0) return;
                 ln[p].a += ln[p].la, ln[p].b += ln[p].lb;
                 if (1 != r) {
                          int pl = ch(p, 0), pr = ch(p, 1);
                          ln[pl].la += ln[p].la, ln[pl].lb += ln[p].lb;
                          ln[pr].la += ln[p].la, ln[pr].lb += ln[p].lb;
                 ln[p].la = ln[p].lb = 0;
        11 query(int x, int p=0, int l=MI, int r=MA) {
                 prop(p, 1, r);
                 ll ret = ln[p](x);
                 if (\ln[p]. \operatorname{ch}[0] == -1 \text{ and } \ln[p]. \operatorname{ch}[1] == -1) return ret;
                 int m = 1 + (r-1)/2;
                 if (x <= m) return min(ret, query(x, ch(p, 0), 1, m));</pre>
```

```
return min(ret, query(x, ch(p, 1), m+1, r));
        void push(line s, int p, int 1, int r) {
                 prop(p, 1, r);
                 int m = 1 + (r-1)/2;
                 bool L = s(1) < ln[p](1);
                bool M = s(m) < ln[p](m);
bool R = s(r) < ln[p](r);
                 if (M) swap(ln[p].a, s.a), swap(ln[p].b, s.b);
                 if (s.b == LINF) return;
                 if (L != M) push(s, ch(p, 0), 1, m);
                 else if (R != M) push(s, ch(p, 1), m+1, r);
        void insert(line s, int a=MI, int b=MA, int p=0, int l=MI, int r=MA) {
                 prop(p, 1, r);
                 if (a <= 1 and r <= b) return push(s, p, 1, r);</pre>
                 if (b < 1 \text{ or } r < a) return;
                 int m = 1 + (r-1)/2;
                 insert(s, a, b, ch(p, 0), 1, m);
                 insert(s, a, b, ch(p, 1), m+1, r);
        void shift(line s, int a=MI, int b=MA, int p=0, int l=MI, int r=MA) {
                 prop(p, 1, r);
int m = 1 + (r-1)/2;
                 if (a \le 1 \text{ and } r \le b) {
                         ln[p].la += s.a, ln[p].lb += s.b;
                          return;
                 if (b < 1 or r < a) return;
                 if (ln[p].b != LINF) {
                          push(ln[p], ch(p, 0), 1, m);
                          push(ln[p], ch(p, 1), m+1, r);
                          ln[p].a = 0, ln[p].b = LINF;
                 shift(s, a, b, ch(p, 0), l, m);
                 shift(s, a, b, ch(p, 1), m+1, r);
};
```

## 22.10 mergeSortTree

```
// MergeSort Tree
// Se for construida sobre um array:
                count(i, j, a, b) retorna quantos
                elementos de v[i..j] pertencem a [a, b]
                report (i, j, a, b) retorna os indices dos
                elementos de v[i..j] que pertencem a [a, b]
                 retorna o vetor ordenado
   Se for construida sobre pontos (x, y):
                count(x1, x2, y1, y2) retorna quantos pontos
                pertencem ao retangulo (x1, y1), (x2, y2)
                report (x1, x2, y1, y2) retorna os indices dos pontos que
                pertencem ao retangulo (x1, y1), (x2, y2)
retorna os pontos ordenados lexicograficamente
                 (assume \ x1 <= x2, \ y1 <= y2)
// kth(y1, y2, k) retorna o indice do ponto com k-esimo menor
// x dentre os pontos que possuem y em [y1, y2] (0 based)
// Se quiser usar para achar k-esimo valor em range, construir
// com ms_tree t(v, true), e chamar kth(l, r, k)
// Usa O(n log(n)) de memoria
// Complexidades:
// construir - O(n log(n))
// count - O(\log(n))
// report - O(log(n) + k) para k indices retornados
// kth - O(log(n))
template <typename T = int> struct ms_tree {
        vector<tuple<T, T, int>> v;
        int n;
```

```
vector<vector<tuple<T, T, int>>> t; // {y, idx, left}
vector<T> vv:
ms_tree(vector<pair<T, T>>& vv) : n(vv.size()), t(4*n), vy(n) {
        for (int i = 0; i < n; i++) v.push_back({vv[i].first, vv[i].</pre>
            second, i});
        sort(v.begin(), v.end());
        build(1, 0, n-1);
        for (int i = 0; i < n; i++) vy[i] = get<0>(t[1][i+1]);
ms_tree(vector<T>& vv, bool inv = false) { // inv: inverte indice e
    valor
        vector<pair<T, T>> v2;
        for (int i = 0; i < vv.size(); i++)</pre>
                inv ? v2.push_back({vv[i], i}) : v2.push_back({i, vv[i]})
                    ] } ) ;
        *this = ms_tree(v2);
void build(int p, int 1, int r) {
        t[p].push_back({get<0>(v[1]), get<0>(v[r]), 0}); // {min_x,
             max_x, 0
        if (1 == r) return t[p].push_back({get<1>(v[1]), get<2>(v[1]),
            0 } ) ;
        int m = (1+r)/2;
        build(2*p, 1, m), build(2*p+1, m+1, r);
        int L = 0, R = 0;
        while (t[p].size() \le r-1+1)
                int left = get<2>(t[p].back());
                if (L > m-1) or (R+m+1 \le r) and t[2*p+1][1+R] \le t[2*p][1+r]
                     L])) {
                        t[p].push_back(t[2*p+1][1 + R++]);
                        get<2>(t[p].back()) = left;
                        continue;
                t[p].push_back(t[2*p][1 + L++]);
                get<2>(t[p].back()) = left+1;
        }
int get_1(T y) { return lower_bound(vy.begin(), vy.end(), y) - vy.begin
int get_r(T y) { return upper_bound(vy.begin(), vy.end(), y) - vy.begin
     (); }
int count(T x1, T x2, T y1, T y2) {
        function<int(int, int, int)> dfs = [&](int p, int 1, int r)
                if (1 == r \text{ or } x2 < get<0>(t[p][0]) \text{ or } get<1>(t[p][0]) <
                     x1) return 0;
                if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2)
                int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
                return dfs(2*p, nl, nr) + dfs(2*p+1, 1-nl, r-nr);
        return dfs(1, get_1(y1), get_r(y2));
vector<int> report(T x1, T x2, T y1, T y2) {
        vector<int> ret;
        function<void(int, int, int)> dfs = [&](int p, int 1, int r) {
                if (1 == r \text{ or } x2 < get<0>(t[p][0]) \text{ or } get<1>(t[p][0]) <
                     x1) return:
                if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2) {
                        for (int i = 1; i < r; i++) ret.push_back(get</pre>
                              <1>(t[p][i+1]));
                int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
                dfs(2*p, nl, nr), dfs(2*p+1, l-nl, r-nr);
        dfs(1, get_1(y1), get_r(y2));
        return ret;
int kth(T y1, T y2, int k)
        function<int(int, int, int)> dfs = [&](int p, int 1, int r) {
                if (k >= r-1) {
                        k -= r-1:
                        return -1;
                if (r-l == 1) return get<1>(t[p][l+1]);
```

```
int nl = get<2>(t[p][l]), nr = get<2>(t[p][r]);
    int left = dfs(2*p, nl, nr);
    if (left != -1) return left;
    return dfs(2*p+1, l-nl, r-nr);
};
return dfs(1, get_l(y1), get_r(y2));
};
};
```

## 22.11 minqueueDeque

## 22.12 minqueueStack

```
// Min queue - stack
// Tudo O(1) amortizado
template < class T > struct minstack {
        stack<pair<T, T>> s;
        void push(T x) {
                if (!s.size()) s.push({x, x});
                else s.emplace(x, std::min(s.top().second, x));
       T top() { return s.top().first; }
        T pop() {
                T ans = s.top().first;
                s.pop();
                return ans;
        int size() { return s.size(); }
        T min() { return s.top().second; }
template<class T> struct minqueue {
       minstack<T> s1, s2;
        void push(T x) { sl.push(x); }
        void move() {
                if (s2.size()) return;
                while (s1.size()) {
                        T x = s1.pop();
                        s2.push(x);
        T front() { return move(), s2.top(); }
        T pop() { return move(), s2.pop(); }
        int size() { return s1.size()+s2.size(); }
        T min() +
                if (!sl.size()) return s2.min();
                else if (!s2.size()) return s1.min();
                return std::min(s1.min(), s2.min());
```

};

### 22.13 orderStatisticSet

```
// Order Statistic Set
// Funciona do C++11 pra cima
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <class T>
        using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
        tree_order_statistics_node_update>;
// para declarar:
// ord_set<int> s;
// coisas do set normal funcionam:
// for (auto i : s) cout << i << endl;
// cout << s.size() << endl;
// k-esimo maior elemento O(log|s|):
// k=0: menor elemento
// cout << *s.find_by_order(k) << endl;</pre>
// quantos sao menores do que k O(log|s|):
// cout << s.order_of_key(k) << endl;</pre>
// Para fazer um multiset, tem que
// usar ord_set<pair<int, int>> com o
// segundo parametro sendo algo para diferenciar
// os ementos iquais.
// s.order_of_key({k, -INF}) vai retornar o
// numero de elementos < k
```

## 22.14 priorityQueueDs

```
// Priority Queue DS
// Mantem updates aplicados em uma estrutura de dados
// que permita rollback e nao seja amortizada.
// Cada update possui uma prioridade,
// sendo possivel remover o update com maior prioridade.
// Os updates devem ser comutativos, ou seja, o estado
// da estrutura deve ser o mesmo independente da ordem
// que eles sejam aplicados.
// Complexidades:
// update - O(log(n) + T(n))
// query - T(n)
// pop - O(log(n) * T(n)) amortizado
// onde T(n) eh a complexidade do update
// assumes all priorities are distinct
template<typename DS, typename UPD> struct priority_queue_ds {
        vector<tuple<UPD, int, int>> upd; // {u, p, idx_in_pos}
        set<pair<int, int>> st;
        vector<int> pos;
        priority_queue_ds(int n) : D(n) {}
        void update(UPD u, int p) {
               D.update(u);
                st.emplace(p, pos.size());
                upd.emplace_back(u, p, pos.size());
                pos.push_back(upd.size() - 1);
        int query(int a) {
                return D.find(a);
```

```
void pop() {
                int k = 1, min_p; // k = number of pops we will do
                vector<tuple<UPD, int, int>> small, big;
                auto it = st.end();
                for (int qt = 0; qt++ < (k+1)/2;) {
                        it--;
                        min_p = it->first;
                        int i = pos[it->second];
                        if (qt > 1) big.push_back(upd[i]);
                        k = max < int > (k, upd.size() - i);
                for (int i = 0; i < k; i++) {
                        D.rollback();
                        auto [u, p, idx] = upd.rbegin()[i];
                        if (p < min_p) small.emplace_back(u, p, idx);</pre>
                st.erase(prev(st.end()));
                upd.erase(upd.end() - k, upd.end());
                small.insert(small.end(), big.rbegin(), big.rend());
                for (auto [u, p, idx] : small) {
                        D.update(u);
                        upd.emplace_back(u, p, idx);
                        pos[idx] = upd.size() - 1;
};
```

## 22.15 rangeColor

```
// Range color
// update(l, r, c) colore o range [l, r] com a cor c,
// e retorna os ranges que foram coloridos {1, r, cor}
// query(i) returna a cor da posicao i
// Complexidades (para q operacoes):
// update - O(log(q)) amortizado
// query - 0(log(q))
template<typename T> struct color {
        set<tuple<int, int, T>> se;
        vector<tuple<int, int, T>> update(int 1, int r, T val) {
                auto it = se.upper_bound({r, INF, val});
                if (it != se.begin() and get<1>(*prev(it)) > r) {
                        auto [L, R, V] = *--it;
                        se.erase(it);
                        se.emplace(L, r, V), se.emplace(r+1, R, V);
                it = se.lower_bound({1, -INF, val});
                if (it != se.begin() and get<1>(*prev(it)) >= 1) {
                        auto [L, R, V] = \star--it;
                        se.erase(it);
                        se.emplace(L, l-1, V), it = se.emplace(l, R, V).first;
                vector<tuple<int, int, T>> ret;
                for (; it != se.end() and qet<0>(*it) <= r; it = se.erase(it))
                        ret.push_back(*it);
                se.emplace(l, r, val);
                return ret;
        T query(int i) {
                auto it = se.upper_bound({i, INF, T()});
                if (it == se.begin() or get<1>(\star--it) < i) return -1; // nao tem
                return get<2>(*it);
};
```

```
// RMQ <O(n), O(1)> - min queue
// O(n) pra buildar, query O(1)
// Se tiver varios minimos, retorna
// o de menor indice
template<typename T> struct rmq {
         vector<T> v;
         int n; static const int b = 30;
         vector<int> mask, t;
         int op(int x, int y) { return v[x] \leftarrow v[y] ? x : y; }
         int msb(int x) { return __builtin_clz(1) -__builtin_clz(x); }
         int small(int r, int sz = b) { return r-msb(mask[r]&((1<sz)-1)); }
         rmq (const vector<T>& v_) : v(v_), n(v.size()), mask(n), t(n) {
                 for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {</pre>
                         at = (at << 1) & ((1 << b) -1);
                          while (at and op(i-msb(at&-at), i) == i) at \hat{} at&-at;
                 for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
                 for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0; i+(1<<<math>j) <=
                      n/b; i++)
                         t[n/b*j+i] = op(t[n/b*(j-1)+i], t[n/b*(j-1)+i+(1<<(j-1)))
                              ]);
        int index_query(int 1, int r) {
    if (r-l+1 <= b) return small(r, r-l+1);</pre>
                 int x = 1/b+1, y = r/b-1;
                 if (x > y) return op(small(l+b-1), small(r));
                 int j = msb(y-x+1);
                 int ans = op(small(1+b-1), op(t[n/b*j+x], t[n/b*j+y-(1<<j)+1]));
                 return op(ans, small(r));
         T query(int 1, int r) { return v[index_query(1, r)]; }
};
```

## 22.17 slopeTrick

```
// SlopeTrick
// Armazena uma estrutura convexa piecewise linear
// Permite adicionar slopes sem peso e realizar query de minimo
// Comentarios acima das funcoes para explicar o que cada uma faz
template<typename T> struct SlopeTrick {
        T inf = numeric_limits<T>::max() / 3;
        priority_queue<T, vector<T>, less<>> L;
        priority_queue<T, vector<T>, greater<>>> R;
        T add_1, add_r;
        T top_R() {
                if (R.empty()) return inf;
                else return R.top() + add_r;
        T pop_R() {
                T val = top_R();
                if (R.size()) R.pop();
                return val;
        T top_L() {
                if (L.empty()) return -inf;
                else return L.top() + add_l;
        T pop_L() {
                T val = top_L();
                if (L.size()) L.pop();
                return val;
```

```
size_t size() {
        return L.size() + R.size();
SlopeTrick() : min_f(0), add_l(0), add_r(0) {};
// return \{min\ f(x),\ lx,\ rx\}
// Em que [lx, rx] eh o intervalo que atinge o minimo
array<T, 3> query() {
        return {min_f, top_L(), top_R()};
// f(x) += a
void add_all(T a) {
        min f += a;
// add \_
// f(x) += max(a - x, 0)
void add_a_minus_x(T a) {
        \min_{f} += \max(T(0), a - top_R());
        R.push(a - add_r);
        L.push(pop_R() - add_l);
// add _/
// f(x) += max(x - a, 0)
void add_x_minus_a(T a) {
        min_f += max(T(0), top_L() - a);
        L.push(a - add_l);
        R.push(pop_L() - add_r);
// add \/
// f(x) += abs(x - a)
void add_abs(T a) {
        add_a_minus_x(a);
        add_x_minus_a(a);
// \/ -> \_
// f_{new}(x) = min f(y) (y <= x)
void clear_right() {
        while (R.size()) R.pop();
// \ \ // \rightarrow \_/
// f_{new} (x) = min f(y) (y \ge x)
void clear_left() {
        while (L.size()) L.pop();
// \/ -> \_/
// f_{new}(x) = min f(y) (x-b \le y \le x-a)
void shift(T a, T b) {
        assert (a <= b);
        add_1 += a;
        add_r += b;
// \/. -> .\/
// f_{new}(x) = f(x - a)
void shift(T a) {
        shift(a, a);
// Retorna f(x)
// O(size)
T get (T x) {
        auto L2 = L;
        auto R2 = R;
        T ret = min_f;
        while (L.size()) {
                 ret += max(T(0), pop_L() - x);
        while (R.size()) {
                 ret += \max(T(0), x - pop_R());
```

```
}
L = L2, R = R2;
return ret;
}

// O(min(size, st.size))
void merge(SlopeTrick &st) {
    if (st.size() > size()) {
        swap(*this, st);
    }
    while (st.R.size()) {
            add_x_minus_a(st.pop_R());
    }
    while (st.L.size()) {
            add_a_minus_x(st.pop_L());
    }
    min_f += st.min_f;
};
```

## 22.18 sparseTable

# 22.19 sparseTableDisjunta

```
// Sparse Table Disjunta
// Resolve qualquer operacao associativa
// MAX2 = log(MAX)
// Complexidades:
// build - O(n log(n))
// query - 0(1)
namespace sparse {
         int m[MAX2][2*MAX], n, v[2*MAX];
int op(int a, int b) { return min(a, b); }
         void build(int n2, int* v2) {
                   for (int i = 0; i < n; i++) v[i] = v2[i]; while (n\&(n-1)) n++;
                   for (int j = 0; (1<<j) < n; j++) {
                             int len = 1<<j;</pre>
                             for (int c = len; c < n; c += 2*len) {
    m[j][c] = v[c], m[j][c-1] = v[c-1];</pre>
                                       for (int i = c+1; i < c+len; i++) m[j][i] = op(</pre>
                                            m[j][i-1], v[i]);
                                       for (int i = c-2; i >= c-len; i--) m[j][i] = op(
                                             v[i], m[j][i+1]);
```

```
}
int query(int 1, int r) {
    if (1 == r) return v[1];
    int j = __builtin_clz(1) - __builtin_clz(1^r);
    return op(m[j][1], m[j][r]);
}
```

## 22.20 splaytree

```
// Splay Tree
// SEMPRE QUE DESCER NA ARVORE, DAR SPLAY NO
// NODE MAIS PROFUNDO VISITADO
// Todas as operacoes sao O(log(n)) amortizado
// Se guiser colocar mais informação no node,
// mudar em 'update'
template<typename T> struct splaytree {
        struct node {
                node *ch[2], *p;
                 int sz;
                T val;
                node(T v) {
                         ch[0] = ch[1] = p = NULL;
                         sz = 1;
                         val = v;
                 void update() {
                         for (int i = 0; i < 2; i++) if (ch[i]) {
                                  sz += ch[i] -> sz;
        };
        node* root;
        splaytree() { root = NULL; }
        splaytree (const splaytree& t) {
                throw logic_error("Nao copiar a splaytree!");
         splaytree() {
                 vector<node*> q = {root};
                 while (q.size()) {
                         node* x = q.back(); q.pop_back();
                         if (!x) continue;
                         q.push_back(x->ch[0]), q.push_back(x->ch[1]);
        void rotate(node* x) { // x vai ficar em cima
                node *p = x->p, *pp = p->p;
if (pp) pp->ch[pp->ch[1] == p] = x;
                bool d = p->ch[0] == x;
                p \rightarrow ch[!d] = x \rightarrow ch[d], x \rightarrow ch[d] = p;
                if (p->ch[!d]) p->ch[!d]->p = p;
                x->p = pp, p->p = x;
                p->update(), x->update();
        node* splay(node* x) {
                 if (!x) return x;
                 root = x:
                 while (x->p) {
                         node *p = x->p, *pp = p->p;
                         if (!pp) return rotate(x), x; // zig
                         if ((pp->ch[0] == p)^(p->ch[0] == x))
                                 rotate(x), rotate(x); // zigzag
                         else rotate(p), rotate(x); // zigzig
                 return x;
```

```
node* insert(T v, bool lb=0) {
        if (!root) return lb ? NULL : root = new node(v);
        node *x = root, *last = NULL;;
        while (1) {
                bool d = x -> val < v;
                if (!d) last = x;
                if (x->val == v) break;
                if (x->ch[d]) x = x->ch[d];
                else {
                        if (lb) break;
                        x->ch[d] = new node(v);
x->ch[d]->p = x;
                        x = x \rightarrow ch[d];
                        break;
        splay(x);
        return 1b ? splay(last) : x;
int size() { return root ? root->sz : 0; }
int count(T v) { return insert(v, 1) and root->val == v; }
node* lower_bound(T v) { return insert(v, 1); }
void erase(T v) {
        if (!count(v)) return;
        node *x = root, *1 = x -> ch[0];
        if (!1) {
                root = x->ch[1];
                if (root) root->p = NULL;
                return delete x;
        root = 1, 1->p = NULL;
        while (1->ch[1]) 1 = 1->ch[1];
        splay(1);
        1->ch[1] = x->ch[1];
        if (1->ch[1]) 1->ch[1]->p = 1;
        delete x:
        1->update();
int order_of_key(T v) {
        if (!lower_bound(v)) return root ? root->sz : 0;
        return root->ch[0] ? root->ch[0]->sz : 0;
node* find_by_order(int k) {
        if (k >= size()) return NULL;
        node* x = root;
        while (1) {
                if (x->ch[0] and x->ch[0]->sz >= k+1) x = x->ch[0];
                        if (x->ch[0]) k -= x->ch[0]->sz;
                        if (!k) return splay(x);
                        k--, x = x->ch[1];
T min() {
        node* x = root;
        while (x->ch[0]) x = x->ch[0]; // max -> ch[1]
        return splay(x)->val;
```

## 22.21 splaytreeImplicita

};

```
T val, sub, lazy;
        bool rev:
        node(T v) {
                 ch[0] = ch[1] = p = NULL;
                 sz = 1;
                 sub = val = v;
                 lazy = 0;
                 rev = false;
         void prop() {
                 if (lazy) {
                           val += lazy, sub += lazy*sz;
                          if (ch[0]) ch[0]->lazy += lazy;
                          if (ch[1]) ch[1]->lazy += lazy;
                           swap(ch[0], ch[1]);
                          if (ch[0]) ch[0]->rev ^= 1;
if (ch[1]) ch[1]->rev ^= 1;
                  lazy = 0, rev = 0;
        void update() {
                 sz = 1, sub = val;
                  for (int i = 0; i < 2; i++) if (ch[i]) {
                          ch[i]->prop();
                          sz += ch[i] -> sz;
                          sub += ch[i]->sub;
        }
};
node* root;
splay() { root = NULL; }
splay(node* x) {
         root = x;
        if (root) root->p = NULL;
splay(vector<T> v) { // O(n)
         root = NULL;
         for (T i : v) {
                 node* x = new node(i);
                 x->ch[0] = root;
                 if (root) root->p = x;
                 root = x;
                 root->update();
splay(const splay& t) {
        throw logic_error("Nao copiar a splay!");
 splay() {
        vector<node*> q = {root};
        while (q.size()) {
                 node* x = q.back(); q.pop_back();
                 if (!x) continue;
                  q.push_back(x->ch[0]), q.push_back(x->ch[1]);
                 delete x:
int size(node* x) { return x ? x->sz : 0; }
void rotate(node* x) { // x vai ficar em cima
         node *p = x->p, *pp = p->p;
if (pp) pp->ch[pp->ch[1] == p] = x;
        bool d = p->ch[0] == x;
p->ch[!d] = x->ch[d], x->ch[d] = p;
         if (p->ch[!d]) p->ch[!d]->p = p;
        x->p = pp, p->p = x;
p->update(), x->update();
node* splaya(node* x) {
        if (!x) return x;
         root = x, x->update();
         while (x->p) {
                 node *p = x->p, *pp = p->p;
if (!pp) return rotate(x), x; // zig
                  if ((pp->ch[0] == p)^(p->ch[0] == x))
```

```
rotate(x), rotate(x); // zigzag
                 else rotate(p), rotate(x); // zigzig
        return x;
node* find(int v) {
        if (!root) return NULL;
        node *x = root;
        int key = 0;
        while (1) {
                 x->prop();
                 bool d = key + size(x->ch[0]) < v;
                 if (key + size(x->ch[0]) != v and x->ch[d]) {
                         if (d) key += size(x->ch[0])+1;
                         x = x \rightarrow ch[d];
                 } else break:
        return splaya(x);
int size() { return root ? root->sz : 0; }
void join(splay<T>& 1) { // assume que 1 < *this</pre>
        if (!size()) swap(root, l.root);
        if (!size() or !l.size()) return;
        node* x = 1.root;
        while (1) {
                 x->prop();
                 if (!x->ch[1]) break;
                 x = x \rightarrow ch[1];
        1.splaya(x), root->prop(), root->update();
        x - ch[1] = root, x - ch[1] - p = x;
        root = 1.root, 1.root = NULL;
        root->update();
node* split(int v) { // retorna os elementos < v
    if (v <= 0) return NULL;</pre>
        if (v >= size()) {
                 node* ret = root;
                 root = NULL;
                 ret->update();
                 return ret;
        find(v);
        node * 1 = root -> ch[0];
        root->ch[0] = NULL;
        if (1) 1->p = NULL;
        root->update();
        return 1;
T& operator [](int i) {
        find(i);
        return root->val;
void push_back(T v) { // O(1)
        node* r = new node(v);
        r->ch[0] = root;
        if (root) root->p = r;
        root = r, root->update();
T query(int 1, int r) {
        splay<T> M(split(r+1));
        splay<T> L(M.split(1));
        T ans = M.root->sub;
        M.join(L), join(M);
        return ans;
void update(int 1, int r, T s) {
        splay<T> M(split(r+1));
        splay<T> L(M.split(l));
        M.root->lazy += s;
M.join(L), join(M);
void reverse(int 1, int r) {
        splay<T> M(split(r+1));
        splay<T> L(M.split(1));
        M.root->rev ^= 1;
        M.join(L), join(M);
void erase(int 1, int r) {
```

```
splay<T> M(split(r+1));
splay<T> L(M.split(l));
join(L);
}
```

## 22.22 splitMergeSet

```
// Split-Merge Set
// Representa um conjunto de inteiros nao negativos
// Todas as operacoes custam O(log(N)),
// em que N = maior elemento do set,
// exceto o merge, que custa O(log(N)) amortizado
// Usa O(min(N, n log(N))) de memoria, sendo 'n' o
// numero de elementos distintos no set
template<typename T, bool MULTI=false, typename SIZE T=int> struct sms {
        struct node {
                node *1, *r;
                SIZE_T cnt;
                node() : 1(NULL), r(NULL), cnt(0) {}
                void update() {
                         if (1) cnt += 1->cnt;
                         if (r) cnt += r->cnt;
        };
        node* root;
        T N;
        sms() : root(NULL), N(0) {}
        sms(T v) : sms() { while (v >= N) N = 2*N+1; }
        sms(const sms& t) : root(NULL), N(t.N) {
    for (SIZE_T i = 0; i < t.size(); i++) {</pre>
                         T at = t[i];
                         SIZE_T qt = t.count(at);
                         insert(at, qt);
                         i += qt-1;
        sms(initializer_list<T> v) : sms() { for (T i : v) insert(i); }
                 vector<node*> q = {root};
                while (q.size()) {
                         node* x = q.back(); q.pop_back();
                         if (!x) continue;
                         q.push_back(x->1), q.push_back(x->r);
                         delete x:
        friend void swap(sms& a, sms& b) {
                swap(a.root, b.root), swap(a.N, b.N);
        sms& operator = (const sms& v) {
                sms tmp = v;
swap(tmp, *this);
                return *this;
        SIZE_T size() const { return root ? root->cnt : 0; }
        SIZE_T count(node* x) const { return x ? x->cnt : 0; }
        void clear() {
                sms tmp;
                swap(*this, tmp);
        void expand(T v) {
                 for (; N < v; N = 2*N+1) if (root) {
                         node* nroot = new node();
                         nroot -> 1 = root;
                         root = nroot;
                         root->update();
```

```
node* insert(node* at, T idx, SIZE_T qt, T 1, T r) {
        if (!at) at = new node();
       if (1 == r) {
                at->cnt += qt;
                if (!MULTI) at->cnt = 1;
                return at;
        T m = 1 + (r-1)/2;
        if (idx <= m) at->1 = insert(at->1, idx, qt, 1, m);
        else at->r = insert(at->r, idx, qt, m+1, r);
       return at->update(), at;
void insert(T v, SIZE_T qt=1) { // insere 'qt' ocorrencias de 'v'
       if (qt <= 0) return erase(v, -qt);</pre>
       assert (v >= 0);
       expand(v);
        root = insert(root, v, qt, 0, N);
node* erase(node* at, T idx, SIZE_T qt, T 1, T r) {
        if (!at) return at;
       if (1 == r) at->cnt = at->cnt < qt ? 0 : at->cnt - qt;
                T m = 1 + (r-1)/2;
                if (idx <= m) at->1 = erase(at->1, idx, qt, 1, m);
                else at->r = erase(at->r, idx, qt, m+1, r);
                at->update();
       if (!at->cnt) delete at, at = NULL;
       return at;
void erase(T v, SIZE_T qt=1) { // remove 'qt' ocorrencias de 'v'
        if (v < 0 \text{ or } v > N \text{ or } !qt) return;
        if (qt < 0) insert(v, -qt);
       root = erase(root, v, qt, 0, N);
void erase_all(T v) { // remove todos os 'v'
        if (v < 0 \text{ or } v > N) return;
        root = erase(root, v, numeric_limits<SIZE_T>::max(), 0, N);
}
SIZE_T count(node* at, T a, T b, T l, T r) const {
        if (!at or b < 1 or r < a) return 0;
        if (a <= 1 and r <= b) return at->cnt;
        T m = 1 + (r-1)/2;
       return count(at->1, a, b, 1, m) + count(at->r, a, b, m+1, r);
SIZE_T count(T v) const { return count(root, v, v, 0, N); }
SIZE_T order_of_key(T v) { return count(root, 0, v-1, 0, N); }
SIZE_T lower_bound(T v) { return order_of_key(v); }
const T operator [](SIZE_T i) const { // i-esimo menor elemento
       assert(i >= 0 and i < size());</pre>
        node* at = root;
        T 1 = 0, r = N;
        while (1 < r)
                T m = 1 + (r-1)/2;
                if (count(at->1) > i) at = at->1, r = m;
                        i -= count(at->1);
                        at = at -> r; 1 = m+1;
       return 1;
node* merge(node* 1, node* r) {
        if (!l or !r) return 1 ? 1 : r;
        if (!1->1 and !1->r) { // folha
                if (MULTI) 1->cnt += r->cnt;
                delete r;
                return 1;
        1->1 = merge(1->1, r->1), 1->r = merge(1->r, r->r);
        1->update(), delete r;
        return 1;
void merge(sms& s) { // mergeia dois sets
```

```
if (N > s.N) swap(*this, s);
        expand(s.N);
        root = merge(root, s.root);
        s.root = NULL;
node* split(node*& x, SIZE_T k) {
         if (k <= 0 or !x) return NULL;</pre>
        node* ret = new node();
        if (!x\rightarrow l \text{ and } !x\rightarrow r) x\rightarrow cnt -= k, ret\rightarrow cnt += k;
        else {
                 if (k \le count(x->1)) ret->1 = split(x->1, k);
                 else {
                          ret->r = split(x->r, k - count(x->l));
                          swap (x->1, ret->1);
                  ret->update(), x->update();
        if (!x->cnt) delete x, x = NULL;
         return ret;
void split(SIZE_T k, sms& s) { // pega os 'k' menores
        s.clear();
        s.root = split(root, min(k, size()));
        s.N = N;
// pega os menores que 'k'
void split_val(T k, sms& s) { split(order_of_key(k), s); }
```

## 22.23 splitMergeSetLazy

};

```
// Split-Merge Set - Lazy [noprint]
// Representa um conjunto de inteiros nao negativos
// Todas as operacoes custam O(log(N)),
// em que N = maior elemento do set,
// exceto o merge e o insert_range, que custa O(log(N)) amortizado
// Usa O(min(N, n log(N))) de memoria, sendo 'n' o
// numero de elementos distintos no set
template<typename T> struct sms {
        struct node {
                node *1, *r;
                 int cnt;
                node(): 1(NULL), r(NULL), cnt(0), flip(0) {}
                void update() {
                         if (1) cnt += 1->cnt;
                         if (r) cnt += r->cnt;
        };
        void prop(node* x, int size) {
                 if (!x or !x->flip) return;
                x \rightarrow flip = 0;
                x\rightarrow cnt = size - x\rightarrow cnt;
                 if (size > 1) {
                         if (!x->1) x->1 = new node();
                         if (!x->r) x->r = new node();
                         x->1->flip ^= 1;
x->r->flip ^= 1;
        }
        node* root:
        T N;
        sms() : root(NULL), N(0) {}
        sms(T v) : sms() { while (v >= N) N = 2*N+1; }
        sms(sms& t) : root(NULL), N(t.N) {
                for (int i = 0; i < t.size(); i++) insert(t[i]);</pre>
        sms(initializer_list<T> v) : sms() { for (T i : v) insert(i); }
```

```
void destroy(node* r) {
        vector<node*> q = {r};
        while (q.size()) {
                node* x = q.back(); q.pop_back();
                if (!x) continue;
                q.push_back(x->1), q.push_back(x->r);
                 delete x;
~sms() { destroy(root); }
friend void swap(sms& a, sms& b) {
        swap(a.root, b.root), swap(a.N, b.N);
sms& operator = (const sms& v) {
        sms tmp = v;
        swap(tmp, *this);
        return *this;
int count(node* x, T size) {
        if (!x) return 0;
        prop(x, size);
        return x->cnt;
int size() { return count(root, N+1); }
void clear() {
        sms tmp;
       swap(*this, tmp);
void expand(T v) {
        for (; N < v; N = 2*N+1) if (root) {
                prop(root, N+1);
                node* nroot = new node();
                nroot->1 = root;
                root = nroot;
                root->update();
node* insert(node* at, T idx, T l, T r) {
        if (!at) at = new node();
        else prop(at, r-1+1);
        if (1 == r) {
                at->cnt = 1;
                return at;
        T m = 1 + (r-1)/2;
        if (idx <= m) at->1 = insert(at->1, idx, 1, m);
        else at->r = insert(at->r, idx, m+1, r);
        return at->update(), at;
void insert(T v) {
        assert (v >= 0);
        expand(v);
        root = insert(root, v, 0, N);
node* erase(node* at, T idx, T l, T r) {
        if (!at) return at;
        prop(at, r-l+1);
        if (1 == r) at->cnt = 0;
        else {
                 T m = 1 + (r-1)/2;
                 if (idx <= m) at->1 = erase(at->1, idx, 1, m);
                 else at->r = erase(at->r, idx, m+1, r);
                at->update();
        return at;
void erase(T v) {
        if (v < 0 \text{ or } v > N) return;
        root = erase(root, v, 0, N);
int count(node* at, T a, T b, T l, T r) {
    if (!at or b < l or r < a) return 0;</pre>
        prop(at, r-l+1);
        if (a <= 1 and r <= b) return at->cnt;
        T m = 1 + (r-1)/2;
```

```
return count (at->1, a, b, 1, m) + count (at->r, a, b, m+1, r);
int count(T v) { return count(root, v, v, 0, N); }
int order_of_key(T v) { return count(root, 0, v-1, 0, N); }
int lower_bound(T v) { return order_of_key(v); }
const T operator [](int i) { // i-esimo menor elemento
        assert(i >= 0 and i < size());
        node* at = root;
        T 1 = 0, r = N;
        while (1 < r) {
                prop(at, r-l+1);
                T m = 1 + (r-1)/2;
                if (count(at->1, m-1+1) > i) at = at->1, r = m;
                else {
                        i -= count(at->1, r-m);
                        at = at -> r; 1 = m+1;
        return 1;
node* merge(node* a, node* b, T tam) {
        if (!a or !b) return a ? a : b;
        prop(a, tam), prop(b, tam);
        if (b->cnt == tam) swap(a, b);
        if (tam == 1 or a->cnt == tam) {
                destroy(b);
                return a;
        a->1 = merge(a->1, b->1, tam>>1), a->r = merge(a->r, b->r, tam
            >>1);
        a->update(), delete b;
        return a;
void merge(sms& s) { // mergeia dois sets
        if (N > s.N) swap(*this, s);
        expand(s.N);
        root = merge(root, s.root, N+1);
        s.root = NULL;
}
node* split(node*& x, int k, T tam) {
        if (k <= 0 or !x) return NULL;</pre>
        prop(x, tam);
        node* ret = new node();
        if (tam == 1) x->cnt = 0, ret->cnt = 1;
                if (k \le count(x->1, tam>>1)) ret->1 = split(x->1, k,
                     tam>>1);
                else {
                        ret->r = split(x->r, k - count(x->l, tam>>1),
                             tam>>1);
                        swap(x->1, ret->1);
                ret->update(), x->update();
        return ret;
void split(int k, sms& s) { // pega os 'k' menores
        s.clear();
        s.root = split(root, min(k, size()), N+1);
        s.N = N;
// pega os menores que 'k'
void split_val(T k, sms& s) { split(order_of_key(k), s); }
void flip(node*& at, T a, T b, T l, T r) {
        if (!at) at = new node();
        else prop(at, r-l+1);
        if (a \le 1 \text{ and } r \le b)
                at->flip ^= 1;
                prop(at, r-1+1);
                return:
        if (r < a or b < 1) return;
        T m = 1 + (r-1)/2;
        flip(at->1, a, b, 1, m), flip(at->r, a, b, m+1, r);
        at->update();
```

```
void flip(T l, T r) { // flipa os valores em [l, r]
        assert (1 >= 0 \text{ and } 1 <= r);
        expand(r);
        flip(root, 1, r, 0, N);
// complemento considerando que o universo eh [0, lim]
void complement(T lim) {
        assert(lim >= 0);
        if (lim > N) expand(lim);
        flip(root, 0, lim, 0, N);
        sms tmp;
        split val(lim+1, tmp);
        swap(*this, tmp);
void insert_range(T 1, T r) { // insere todo os valores em [1, r]
        sms tmp;
        tmp.flip(l, r);
        merge(tmp);
```

### 22.24 sqrtTree

};

```
// SQRT Tree
// RMQ em O(log log n) com O(n log log n) pra buildar
// Funciona com qualquer operacao associativa
// Tao rapido quanto a sparse table, mas usa menos memoria
// (log log (le9) < 5, entao a query eh praticamente O(1))
// build - O(n log log n)
// query - O(log log n)
namespace sgrtTree {
        int n, *v;
        int pref[4][MAX], sulf[4][MAX], getl[4][MAX], entre[4][MAX], sz[4];
        int op(int a, int b) { return min(a, b); }
        inline int getblk(int p, int i) { return (i-getl[p][i])/sz[p]; }
        void build(int p, int 1, int r) {
                if (1+1 >= r) return;
                for (int i = 1; i <= r; i++) get1[p][i] = 1;</pre>
                 for (int L = 1; L <= r; L += sz[p]) {
                         int R = min(L+sz[p]-1, r);
                         pref[p][L] = v[L], sulf[p][R] = v[R];
for (int i = L+1; i <= R; i++) pref[p][i] = op(pref[p][i]</pre>
                              -1], v[i]);
                         for (int i = R-1; i >= L; i--) sulf[p][i] = op(v[i],
                              sulf[p][i+1]);
                         build(p+1, L, R);
                 for (int i = 0; i <= sz[p]; i++) {
                         int at = entre[p][l+i*sz[p]+i] = sulf[p][l+i*sz[p]];
                         for (int j = i+1; j \le sz[p]; j++) entre[p][l+i*sz[p]+j]
                               = at =
                                          op(at, sulf[p][l+j*sz[p]]);
        void build(int n2, int* v2) {
                n = n2, v = v2;
                 for (int p = 0; p < 4; p++) sz[p] = n2 = sqrt(n2);
                build(0, 0, n-1);
        int query(int 1, int r) {
                if (l+1 >= r) return l == r ? v[l] : op(v[l], v[r]);
                int p = 0;
                 while (getblk(p, 1) == getblk(p, r)) p++;
                 int ans = sulf[p][1], a = getblk(p, 1)+1, b = getblk(p, r)-1;
                if (a \le b) ans = op(ans, entre[p][getl[p][l]+a*sz[p]+b]);
                return op(ans, pref[p][r]);
```

## 22.25 treap

```
// Treap
// Todas as operacoes custam
// O(log(n)) com alta probabilidade, exceto meld
// meld custa O(log^2 n) amortizado com alta prob.,
// e permite unir duas treaps sem restricao adicional
// Na pratica, esse meld tem constante muito boa e
// o pior caso eh meio estranho de acontecer
mt19937 rng((int) chrono::steady_clock::now().time_since_epoch().count());
template<typename T> struct treap {
        struct node {
                node *1, *r;
                int p, sz;
                node(T \ v) : 1(NULL), r(NULL), p(rnq()), sz(1), val(v), mi(v) {}
                void update() {
                        sz = 1;
                         mi = val;
                         if (1) sz += 1->sz, mi = min(mi, 1->mi);
                         if (r) sz += r->sz, mi = min(mi, r->mi);
        };
        node* root;
        treap() { root = NULL; }
        treap(const treap& t) {
                throw logic_error("Nao copiar a treap!");
         ~treap() {
                vector<node*> q = {root};
                while (q.size()) {
                         node* x = q.back(); q.pop_back();
                         if (!x) continue;
                         q.push_back(x->1), q.push_back(x->r);
                         delete x;
        int size(node* x) { return x ? x->sz : 0; }
        int size() { return size(root); }
        {f void} join(node* 1, node* r, node*& i) { // assume que 1 < r
                if (!1 or !r) return void(i = 1 ? 1 : r);
                if (1->p > r->p) join(1->r, r, 1->r), i = 1;
                else join(1, r->1, r->1), i = r;
                i->update();
        void split(node* i, node*& l, node*& r, T v) {
                if (!i) return void(r = 1 = NULL);
                if (i->val < v) split(i->r, i->r, r, v), l = i;
else split(i->l, l, i->l, v), r = i;
                i->update();
        void split_leq(node* i, node*& 1, node*& r, T v) {
                if (!i) return void(r = l = NULL);
                if (i->val <= v) split_leq(i->r, i->r, r, v), 1 = i;
                else split_leq(i\rightarrow 1, l, i\rightarrow 1, v), r = i;
                i->update();
        int count(node* i, T v) {
                if (!i) return 0;
                if (i->val == v) return 1;
                if (v < i->val) return count(i->1, v);
                return count(i->r, v);
        void index_split(node* i, node*& 1, node*& r, int v, int key = 0) {
                if (!i) return void(r = l = NULL);
                if (key + size(i->1) < v) index_split(i->r, i->r, r, v, key+size
                     (i->1)+1), 1 = i;
                 else index_split(i->1, 1, i->1, v, key), r = i;
                i->update();
        int count(T v) {
                 return count(root, v);
```

```
void insert(T v) {
        if (count(v)) return;
        node *L, *R;
        split(root, L, R, v);
        node* at = new node(v);
        join(L, at, L);
        join(L, R, root);
void erase(T v) {
        node *L, *M, *R;
        split_leq(root, M, R, v), split(M, L, M, v);
        if (M) delete M;
        M = NULL;
        join(L, R, root);
void meld(treap& t) { // segmented merge
        node *L = root, *R = t.root;
        root = NULL;
        while (L or R) {
                if (!L or (L and R and L->mi > R->mi)) std::swap(L, R);
                if (!R) join(root, L, root), L = NULL;
                else if (L->mi == R->mi) {
                        node* LL;
                        split(L, LL, L, R->mi+1);
delete LL;
                } else {
                        node* LL;
                        split(L, LL, L, R->mi);
                        join(root, LL, root);
        t.root = NULL:
```

## 22.26 treapImplicita

};

```
// Treap Implicita
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
mt19937 rng((int) chrono::steady clock::now().time since epoch().count());
template<typename T> struct treap {
        struct node {
                node *1, *r;
                int p, sz;
                T val, sub, lazy;
                bool rev;
                node(T \ v) : l(NULL), r(NULL), p(rng()), sz(1), val(v), sub(v),
                    lazy(0), rev(0) {}
                void prop() {
                        if (lazy) {
                                val += lazy, sub += lazy*sz;
                                if (1) 1->lazy += lazy;
                                if (r) r->lazy += lazy;
                        if (rev) {
                                swap(1, r);
                                if (1) 1->rev ^= 1;
                                if (r) r->rev ^= 1;
                        lazy = 0, rev = 0;
                void update() {
                        sz = 1, sub = val;
                        if (1) 1->prop(), sz += 1->sz, sub += 1->sub;
                        if (r) r->prop(), sz += r->sz, sub += r->sub;
        };
        node* root;
```

```
treap() { root = NULL; }
treap(const treap& t) {
        throw logic_error("Nao copiar a treap!");
~treap() {
        vector<node*> q = {root};
        while (q.size()) {
                 node* x = q.back(); q.pop_back();
                 if (!x) continue;
                 q.push_back(x->1), q.push_back(x->r);
                 delete x;
int size(node* x) { return x ? x->sz : 0; }
int size() { return size(root); }
void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
        if (!1 or !r) return void(i = 1 ? 1 : r);
        1->prop(), r->prop();
        if (1->p > r->p) join(1->r, r, 1->r), i = 1;
        else join(1, r \rightarrow 1, r \rightarrow 1), i = r;
        i->update();
void split(node* i, node*& l, node*& r, int v, int key = 0) {
        if (!i) return void(r = 1 = NULL);
        i->prop();
         if (\text{key} + \text{size}(i\rightarrow 1) < v) split(i\rightarrow r, i\rightarrow r, r, v, \text{key+size}(i\rightarrow 1)
             +\bar{1}), 1 = i;
        else split(i->1, 1, i->1, v, key), r = i;
        i->update();
void push_back(T v) {
        node* i = new node(v);
         join(root, i, root);
T query(int 1, int r) {
        node *L, *M, *R;
         split(root, M, R, r+1), split(M, L, M, 1);
        T ans = M->sub;
         join(L, M, M), join(M, R, root);
         return ans;
void update(int 1, int r, T s) {
        node *L, *M, *R;
         split(root, M, R, r+1), split(M, L, M, 1);
        M->lazy += s;
        join(L, M, M), join(M, R, root);
void reverse(int 1, int r) {
        node *L, *M, *R;
         split(root, M, R, r+1), split(M, L, M, 1);
        M->rev ^= 1;
        join(L, M, M), join(M, R, root);
```

# 22.27 treapPersistent

};

```
// Treap Persistent Implicita
//
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
mt19937_64 rng((int) chrono::steady_clock::now().time_since_epoch().count());
struct node {
    node *1, *r;
    l1 sz, val, sub;
    node(11 v) : 1(NULL), r(NULL), sz(1), val(v), sub(v) {}
    node(node* x) : 1(x->1), r(x->r), sz(x->sz), val(x->val), sub(x->sub) {}
    void update() {
        sz = 1, sub = val;
        if (1) sz += 1->sz, sub += 1->sub;
        if (r) sz += r->sz, sub += r->sub;
        sub %= MOD;
```

```
};
11 size(node* x) { return x ? x->sz : 0; }
void update(node* x) { if (x) x->update(); }
node* copy(node* x) { return x ? new node(x) : NULL; }
node* join(node* 1, node* r)
        if (!1 or !r) return 1 ? copy(1) : copy(r);
        if (rng() % (size(l) + size(r)) < size(l)) {</pre>
                 ret = copy(1);
                 ret->r = join(ret->r, r);
        } else {
                 ret = copy(r);
                 ret->1 = join(1, ret->1);
        return update(ret), ret;
void split(node* x, node*& 1, node*& r, 11 v, 11 key = 0) {
        if (!x) return void(l = r = NULL);
        if (\text{key} + \text{size}(x\rightarrow 1) < v) {
                 1 = copy(x);
                 split(1->r, 1->r, r, v, key+size(1->1)+1);
        } else {
                 r = copy(x);
                 split(r\rightarrow 1, 1, r\rightarrow 1, v, key);
        update(1), update(r);
vector<node*> treap;
void init(const vector<11>& v) {
        treap = {NULL};
        for (auto i : v) treap[0] = join(treap[0], new node(i));
```

### 22.28 waveletTree

```
// Wavelet Tree
// Usa O(sigma + n log(sigma)) de memoria,
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
// count(i, j, x, y) retorna o numero de elementos de
// v[i, j) que pertencem a [x, y]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v(i, j), se ele fosse ordenado
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
// Complexidades:
// build - O(n log(sigma))
// count - O(log(sigma))
// kth - O(log(sigma))
// sum - O(log(sigma))
// sumk - O(log(sigma))
int n, v[MAX];
vector<int> esq[4*(MAXN-MINN)], pref[4*(MAXN-MINN)];
void build(int b = 0, int e = n, int p = 1, int l = MINN, int r = MAXN) {
        int m = (1+r)/2; esq[p].push_back(0); pref[p].push_back(0);
        for (int i = b; i < e; i++)
                esg[p].push back(esg[p].back()+(v[i]<=m));</pre>
                pref(p).push_back(pref(p).back()+v(i));
        if (1 == r) return;
        int m2 = stable_partition(v+b, v+e, [=](int i){return i <= m;}) - v;</pre>
        build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
```

```
int count (int i, int j, int x, int y, int p = 1, int l = MINN, int r = MAXN) {
        if (y < 1 \text{ or } r < x) return 0;
        if (x \le 1 \text{ and } r \le y) return j-i;
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        return count (ei, ej, x, y, 2*p, 1, m) +count (i-ei, j-ej, x, y, 2*p+1, m
int kth(int i, int j, int k, int p=1, int l = MINN, int r = MAXN) {
        if (1 == r) return 1;
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
        return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
int sum(int i, int j, int x, int y, int p = 1, int 1 = MINN, int r = MAXN) {
        if (y < 1 \text{ or } r < x) return 0;
        if (x <= 1 and r <= y) return pref[p][j]-pref[p][i];</pre>
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei, j-ej, x, y, 2*p+1, m+1,
int sumk(int i, int j, int k, int p = 1, int 1 = MINN, int r = MAXN) {
        if (1 == r) return 1*k;
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
        if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);</pre>
        return pref[2*p][ej]-pref[2*p][ei]+sumk(i-ei, j-ej, k-(ej-ei), 2*p+1, m
```

# 23 ufmg forked/Estruturas/Segtree

## 23.1 segTreap

```
// SegTreap
// Muda uma posicao do plano, e faz query de operacao
// associativa e comutativa em retangulo
// Mudar ZERO e op
// Esparso nas duas coordenadas, inicialmente eh tudo ZERO
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Valores no X tem que ser de 0 ateh NX
// Para q operacoes, usa O(q log(NX)) de memoria, e as
// operacoes custa O(log(q) log(NX))
const int ZERO = INF;
const int op(int 1, int r) { return min(1, r); }
mt19937 rng((int) chrono::steady_clock::now().time_since_epoch().count());
template<typename T> struct treap {
        struct node {
                node *1, *r;
                int p;
                pair<11, 11> idx; // {y, x}
                T val, mi;
                node(ll x, ll y, T val_) : l(NULL), r(NULL), p(rng()),
                        idx(pair(y, x)), val(val_), mi(val) {}
                void update() {
                        if (1) mi = op (mi, 1->mi);
                        if (r) mi = op(mi, r->mi);
        };
        node* root;
        treap() { root = NULL; }
```

### 23.2 segTree

```
// SegTree
// Recursiva com Lazy Propagation
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Pode usar a sequinte funcao para indexar os nohs:
// f(1, r) = (1+r)/(1!=r), usando 2N de memoria
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - 0(log(n))
        11 seg[4*MAX], lazy[4*MAX];
        int n, *v;
        11 build(int p=1, int l=0, int r=n-1) {
                lazy[p] = 0;
                if (1 == r) return seg[p] = v[1];
                int m = (1+r)/2;
                return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1, r);
        void build(int n2, int* v2) {
                n = n2, v = v2;
                build();
        void prop(int p, int 1, int r) {
                seg[p] += lazy[p]*(r-l+1);
                if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
                lazy[p] = 0;
        11 query(int a, int b, int p=1, int l=0, int r=n-1) {
                prop(p, 1, r);
                 if (a <= 1 and r <= b) return seg[p];</pre>
                if (b < 1 or r < a) return 0;
                int m = (1+r)/2;
                return query (a, b, 2*p, 1, m) + query (a, b, 2*p+1, m+1, r);
        11 update(int a, int b, int x, int p=1, int l=0, int r=n-1) {
                prop(p, 1, r);
                if (a \le 1 \text{ and } r \le b) {
                        lazy[p] += x;
                         prop(p, 1, r);
                         return seg[p];
                if (b < 1 or r < a) return seg[p];</pre>
                int m = (1+r)/2;
                return seg[p] = update(a, b, x, 2*p, 1, m) +
                         update(a, b, x, 2*p+1, m+1, r);
};
// Se tiver uma seg de max, da pra descobrir em O(log(n))
// o primeiro e ultimo elemento >= val numa range:
// primeira posicao >= val em [a, b] (ou -1 se nao tem)
int get_left(int a, int b, int val, int p=1, int l=0, int r=n-1) {
        prop(p, 1, r);
if (b < 1 or r < a or seq[p] < val) return -1;</pre>
        if (r == 1) return 1;
        int m = (1+r)/2;
        int x = \text{get\_left}(a, b, \text{val}, 2*p, 1, m);
        if (x != -1) return x;
        return get_left(a, b, val, 2*p+1, m+1, r);
// ultima posicao >= val em [a, b] (ou -1 se nao tem)
int get_right(int a, int b, int val, int p=1, int l=0, int r=n-1) {
        prop(p, 1, r);
        if (b < 1 or r < a or seg[p] < val) return -1;</pre>
        if (r == 1) return 1;
        int m = (1+r)/2;
        int x = get_right(a, b, val, 2*p+1, m+1, r);
        if (x != -1) return x;
```

```
return get_right(a, b, val, 2*p, 1, m);
}

// Se tiver uma seg de soma sobre um array nao negativo v, da pra
// descobrir em O(log(n)) o maior j tal que v[i]+v[i+1]+...+v[j-1] < val
int lower_bound(int i, 11& val, int p, int 1, int r) {
    prop(p, 1, r);
    if (r < i) return n;
    if (i <= 1 and seg[p] < val) {
        val -= seg[p];
        return n;
    }
    if (l == r) return 1;
    int m = (1+r)/2;
    int x = lower_bound(i, val, 2*p, 1, m);
    if (x != n) return x;
    return lower_bound(i, val, 2*p+1, m+1, r);
}</pre>
```

## $23.3 \quad segTree2D$

```
// SegTree 2D Iterativa
// Consultas 0-based
// Um valor inicial em (x, y) deve ser colocado em seg[x+n][y+n]
// Query: soma do retangulo ((x1, y1), (x2, y2))
// Update: muda o valor da posicao (x, y) para val
// Nao pergunte como que essa coisa funciona
// Para query com distancia de manhattan <= d, faca
// nx = x+y, ny = x-y
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
// Complexidades:
// build - O(n^2)
// query - O(log^2(n))
// update - O(log^2(n))
int seq[2*MAX][2*MAX], n;
void build() {
         for (int x = 2*n; x; x--) for (int y = 2*n; y; y--) {
                  if (x < n) seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
if (y < n) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
int query(int x1, int y1, int x2, int y2) {
         int ret = 0, y3 = y1 + n, y4 = y2 + n;
         for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2)
for (y1 = y3, y2 = y4; y1 <= y2; ++y1 /= 2, --y2 /= 2) {
                           if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } += \text{seq}[x1][y1];
                           if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x1][y2];
                           if (x2\%2 == 0 and y1\%2 == 1) ret += seg[x2][y1];
                           if (x2\%2 == 0 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x2][y2];
         return ret;
void update(int x, int y, int val) {
         int y2 = y += n;
         for (x += n; x; x /= 2, y = y2) {
                  if (x \ge n) seg[x][y] = val;
                  else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
                  while (y /= 2) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
```

## 23.4 segTreeBeats

```
// SegTree Beats
// query(a, b) - {{min(v[a..b]), max(v[a..b])}, sum(v[a..b])}
// updatemin(a, b, x) faz com que v[i] <- min(v[i], x),
// para i em [a, b]
// updatemax faz o mesmo com max, e updatesum soma x
// em todo mundo do intervalo [a, b]
// Complexidades:
// build - O(n)
// query - O(\log(n))
// update - O(log^2 (n)) amortizado
// (se nao usar updatesum, fica log(n) amortizado)
#define f first
#define s second
namespace beats {
        struct node
                int tam;
                11 sum, lazy; // lazy pra soma
                11 mi1, mi2, mi; // mi = #mi1
                11 ma1, ma2, ma; // ma = #ma1
                node(11 x = 0) {
                        sum = mi1 = ma1 = x;
                        mi2 = LINF, ma2 = -LINF;
                        mi = ma = tam = 1;
                        lazy = 0;
                node(const node& 1, const node& r) {
                        sum = 1.sum + r.sum, tam = 1.tam + r.tam;
                        lazy = 0;
                        if (1.mi1 > r.mi1) {
                                mi1 = r.mi1, mi = r.mi;
                                mi2 = min(1.mi1, r.mi2);
                        } else if (1.mi1 < r.mi1) {</pre>
                                mi1 = 1 mi1, mi = 1 mi;
                                mi2 = min(r.mi1, 1.mi2);
                                mi1 = 1.mi1. mi = 1.mi+r.mi:
                                mi2 = min(1.mi2, r.mi2);
                        if (l.ma1 < r.ma1) {
                                ma1 = r.ma1, ma = r.ma;
                                ma2 = max(1.ma1, r.ma2);
                        } else if (1.ma1 > r.ma1) {
                                ma1 = 1.ma1, ma = 1.ma;
                                ma2 = max(r.ma1, 1.ma2);
                        } else {
                                ma1 = 1.ma1, ma = 1.ma+r.ma;
                                ma2 = max(1.ma2, r.ma2);
                void setmin(ll x) {
                        if (x >= mal) return;
                        sum += (x - ma1) *ma;
                        if (mi1 == ma1) mi1 = x;
                        if (mi2 == ma1) mi2 = x;
                        ma1 = x:
                void setmax(ll x) {
                        if (x <= mil) return;</pre>
                        sum += (x - mi1) *mi;
                        if (ma1 == mi1) ma1 = x;
                        if (ma2 == mi1) ma2 = x;
                        mi1 = x:
                void setsum(11 x) {
                        mi1 += x, mi2 += x, ma1 += x, ma2 += x;
                        sum += x*tam;
                        lazv += x;
        };
        node seg[4*MAX];
```

```
int n, *v;
node build(int p=1, int l=0, int r=n-1) {
        if (l == r) return seg[p] = {v[1]};
        int m = (1+r)/2;
        return seg[p] = \{build(2*p, 1, m), build(2*p+1, m+1, r)\};
void build(int n2, int* v2) {
       n = n2, v = v2;
       build();
void prop(int p, int 1, int r) {
        if (1 == r) return;
        for (int k = 0; k < 2; k++) {
                if (seg[p].lazy) seg[2*p+k].setsum(seg[p].lazy);
                seg[2*p+k].setmin(seg[p].mal);
                seg[2*p+k].setmax(seg[p].mi1);
        seq[p].lazy = 0;
pair<pair<11, 11>, 11> query(int a, int b, int p=1, int 1=0, int r=n-1)
        if (b < 1 or r < a) return {{LINF, -LINF}, 0};</pre>
        if (a \leq 1 and r \leq b) return {{seg[p].mi1, seg[p].ma1}, seg[p].
            sum};
        prop(p, 1, r);
        int m = (1+r)/2;
        auto L = query(a, b, 2*p, 1, m), R = query(a, b, 2*p+1, m+1, r);
        return {{min(L.f.f, R.f.f), max(L.f.s, R.f.s)}, L.s+R.s};
node updatemin(int a, int b, 11 x, int p=1, int 1=0, int r=n-1) {
        if (b < 1 or r < a or seg[p].ma1 <= x) return seg[p];</pre>
        if (a \leq 1 and r \leq b and seg[p].ma2 \leq x) {
                seg[p].setmin(x);
                return seg[p];
        prop(p, 1, r);
int m = (1+r)/2;
        return seg[p] = \{updatemin(a, b, x, 2*p, 1, m),
                                          updatemin(a, b, x, 2*p+1, m+1, r
node updatemax(int a, int b, ll x, int p=1, int l=0, int r=n-1) {
        if (b < 1 or r < a or seg[p].mi1 >= x) return seg[p];
        if (a \le 1 \text{ and } r \le b \text{ and } seg[p].mi2 > x) {
                seq[p].setmax(x);
                return seg[p];
        prop(p, l, r);
        int m = (1+r)/2;
        return seg[p] = \{updatemax(a, b, x, 2*p, 1, m),
                                          updatemax(a, b, x, 2*p+1, m+1, r
                                              ) } ;
node updatesum(int a, int b, ll x, int p=1, int l=0, int r=n-1) {
        if (b < 1 or r < a) return seg[p];</pre>
        if (a \le 1 \text{ and } r \le b) {
                seg[p].setsum(x);
                return seq[p];
        prop(p, 1, r);
        int m = (1+r)/2;
        return seg[p] = \{updatesum(a, b, x, 2*p, 1, m),
                                          updatesum(a, b, x, 2*p+1, m+1, r
                                               ) };
```

## 23.5 segTreeColor

};

```
// SegTree Colorida
//
// Cada posicao tem um valor e uma cor
// O construtor receve um vector de {valor, cor}
// e o numero de cores (as cores devem estar em [0, c-1])
```

```
// query(c, a, b) retorna a soma dos valores
// de todo mundo em [a, b] que tem cor c
// update(c, a, b, x) soma x em todo mundo em
// [a, b] que tem cor c
// paint(c\bar{1}, c\bar{2}, a, b) faz com que todo mundo
// em [a, b] que tem cor cl passe a ter cor c2
// Complexidades:
// construir - O(n log(n)) espaco e tempo
// query - O(log(n))
// update - 0(log(n))
// paint - O(log(n)) amortizado
struct seg_color {
        struct node {
                node *1, *r;
                int cnt;
                ll val, lazy;
                node() : 1(NULL), r(NULL), cnt(0), val(0), lazy(0) {}
                void update() {
                         cnt = 0, val = 0;
                         for (auto i : {1, r}) if (i) {
                                 i->prop();
                                 cnt += i->cnt, val += i->val;
                void prop() {
                         if (!lazv) return;
                         val += lazy*(11)cnt;
                         for (auto i : {1, r}) if (i) i->lazy += lazy;
                         lazy = 0;
        };
        int n;
        vector<node*> seq;
        seg_color(vector<pair<int, int>>& v, int c) : n(v.size()), seg(c, NULL)
                for (int i = 0; i < n; i++)
                         seg[v[i].second] = insert(seg[v[i].second], i, v[i].
                              first, 0, n-1);
        ~seg_color() {
                 queue<node*> q;
                 for (auto i : seq) q.push(i);
                while (q.size()) {
                         auto i = q.front(); q.pop();
                         if (!i) continue;
                         q.push(i->1), q.push(i->r);
                         delete i;
        node* insert(node* at, int idx, int val, int l, int r) {
                 if (!at) at = new node();
                if (1 == r) return at->cnt = 1, at->val = val, at;
                 int m = (1+r)/2;
                if (idx <= m) at->1 = insert(at->1, idx, val, 1, m);
                else at->r = insert(at->r, idx, val, m+1, r);
                return at->update(), at;
        11 query(node* at, int a, int b, int l, int r) {
                if (!at or b < 1 or r < a) return 0;</pre>
                 at->prop();
                if (a <= 1 and r <= b) return at->val;
                int m = (1+r)/2;
                return query(at->1, a, b, 1, m) + query(at->r, a, b, m+1, r);
        11 query(int c, int a, int b) { return query(seg[c], a, b, 0, n-1); }
void update(node* at, int a, int b, int x, int 1, int r) {
                if (!at or b < 1 or r < a) return;</pre>
                at->prop();
                if (a <= 1 and r <= b) {</pre>
                         at->lazy += x;
                         return void(at->prop());
                 int m = (1+r)/2;
                update(at->1, a, b, x, 1, m), update(at->r, a, b, x, m+1, r);
```

```
at->update();
        void update(int c, int a, int b, int x) { update(seg[c], a, b, x, 0, n
            -1);
        void paint(node*& from, node*& to, int a, int b, int 1, int r) {
                if (to == from or !from or b < 1 or r < a) return;</pre>
                from->prop();
                if (to) to->prop();
                if (a \le 1 \text{ and } r \le b) {
                        if (!to) {
                                 to = from;
                                 from = NULL;
                                 return;
                         int m = (1+r)/2;
                        paint(from->1, to->1, a, b, 1, m), paint(from->r, to->r,
                              a, b, m+1, r);
                        to->update();
                        delete from;
                        from = NULL;
                        return;
                if (!to) to = new node();
                int m = (1+r)/2;
                paint(from->1, to->1, a, b, 1, m), paint(from->r, to->r, a, b, m
                     +1, r);
                from->update(), to->update();
        void paint(int c1, int c2, int a, int b) { paint(seg[c1], seg[c2], a, b,
              0, n-1); }
};
```

## 23.6 segTreeEsparsa

```
// SegTree Esparsa - Lazy
// Ouery: soma do range [a, b]
// Update: flipa os valores de [a, b]
// O MAX tem q ser Q log N para Q updates
// Complexidades:
// build - 0(1)
// query - 0(log(n))
// update - 0(log(n))
namespace seg {
        int seg[MAX], lazy[MAX], R[MAX], L[MAX], ptr;
        int get_l(int i) {
                if (L[i] == 0) L[i] = ptr++;
                return L[i];
        int get_r(int i) {
                if (R[i] == 0) R[i] = ptr++;
                return R[i];
        void build() { ptr = 2; }
        void prop(int p, int 1, int r) {
                if (!lazy[p]) return;
                seg[p] = r-1+1 - seg[p];
                if (l != r) lazy[get_l(p)]^=lazy[p], lazy[get_r(p)]^=lazy[p];
        int query(int a, int b, int p=1, int l=0, int r=N-1) {
                prop(p, 1, r);
                if (b < 1 or r < a) return 0;</pre>
                if (a <= 1 and r <= b) return seg[p];</pre>
                int m = (1+r)/2;
                return query(a, b, get_l(p), l, m)+query(a, b, get_r(p), m+1, r)
                     ;
```

```
int update(int a, int b, int p=1, int l=0, int r=N-1) {
    prop(p, l, r);
    if (b < l or r < a) return seg[p];
    if (a <= l and r <= b) {
        lazy[p] ^= l;
        prop(p, l, r);
        return seg[p];
    }
    int m = (l+r)/2;
    return seg[p] = update(a, b, get_l(p), l, m)+update(a, b, get_r(p), m+1, r);
};</pre>
```

## 23.7 segTreeEsparsa2

```
// SegTree Esparsa - O(q) memoria
// Query: min do range [a, b]
// Update: troca o valor de uma posicao
// Usa O(q) de memoria para q updates
// Complexidades:
// query - O(log(n))
// update - 0(log(n))
template<typename T> struct seg {
        struct node {
                node* ch[2];
                 char d;
                 T v;
                 T mi:
                node(int d_, T v_, T val) : d(d_), v(v_) {
    ch[0] = ch[1] = NULL;
                         mi = val;
                 node (node * x) : d(x->d), v(x->v), mi(x->mi) {
                         ch[0] = x - ch[0], ch[1] = x - ch[1];
                 void update() {
                         mi = numeric_limits<T>::max();
                         for (int i = 0; i < 2; i++) if (ch[i])
                                  mi = min(mi, ch[i]->mi);
        };
        node* root;
        char n;
        seg() : root(NULL), n(0) {}
         ~seg() {
                 std::vector<node*> q = {root};
                 while (q.size()) {
                         node* x = q.back(); q.pop_back();
                         if (!x) continue;
                         q.push\_back(x->ch[0]), q.push\_back(x->ch[1]);
                         delete x;
                 }
        char msb(T v, char l, char r) { // msb in range (l, r]
                 for (char i = r; i > 1; i--) if (v>>i&1) return i;
                 return -1;
        void cut(node* at, T v, char i) {
    char d = msb(v ^ at->v, at->d, i);
                 if (d == -1) return; // no need to split
                 node* nxt = new node(at);
                 at->ch[v>>d&1] = NULL;
                 at->ch[!(v>>d&1)] = nxt;
                 at->d = d;
```

```
node* update(node* at, T idx, T val, char i) {
                 if (!at) return new node(-1, idx, val);
                 cut(at, idx, i);
                 if (at->d == -1) { // leaf
                         at->mi = val;
                         return at;
                 bool dir = idx>>at->d&1;
                 at->ch[dir] = update(at->ch[dir], idx, val, at->d-1);
                 at->update();
                 return at;
        void update(T idx, T val) {
                 while (idx>>n) n++;
                 root = update(root, idx, val, n-1);
        T query(node* at, T a, T b, T l, T r, char i) {
    if (!at or b < l or r < a) return numeric_limits<T>::max();
                 if (a <= 1 and r <= b) return at->mi;
                 T m = 1 + (r-1)/2;
                 if (at->d < i) {
                         if ((at->v>>i&1) == 0) return query (at, a, b, 1, m, i-1)
                          else return query (at, a, b, m+1, r, i-1);
                 return min(query(at->ch[0], a, b, 1, m, i-1), query(at->ch[1], a
                      , b, m+1, r, i-1));
        T query (T 1, T r) { return query (root, 1, r, 0, (T(1) << n) -1, n-1); }
};
```

## 23.8 segTreeIterativa

```
// SegTree Iterativa
// Consultas 0-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b]
// Update: muda o valor da posicao p para x
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
int seg[2 * MAX];
int n;
void build() {
        for (int i = n - 1; i; i--) seg[i] = seg[2*i] + seg[2*i+1];
int query(int a, int b) {
        int ret = 0;
        for (a += n, b += n; a <= b; ++a /= 2, --b /= 2) {
                if (a % 2 == 1) ret += seg[a];
                if (b % 2 == 0) ret += seg[b];
        return ret;
void update(int p, int x) {
        seg[p += n] = x;
        while (p \neq 2) seg[p] = seg[2*p] + seg[2*p+1];
```

## 23.9 segTreeIterativaComLazy

```
// SegTree Iterativa com Lazy Propagation
//
// Query: soma do range [a, b]
```

```
// Update: soma x em cada elemento do range [a, b]
// Para mudar, mudar as funcoes junta, poe e query
// LOG = ceil(log2(MAX))
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
namespace seg {
        11 seg[2*MAX], lazy[2*MAX];
        int n;
        11 junta(11 a, 11 b) {
                return a+b;
        // soma x na posicao p de tamanho tam
        void poe(int p, 11 x, int tam, bool prop=1) {
                seq[p] += x*tam;
                if (prop and p < n) lazy[p] += x;</pre>
        // atualiza todos os pais da folha p
        void sobe(int p) {
                for (int tam = 2; p /= 2; tam *= 2) {
                        seg[p] = junta(seg[2*p], seg[2*p+1]);
                        poe(p, lazy[p], tam, 0);
        // propaga o caminho da raiz ate a folha p
        void prop(int p) {
                int tam = 1 << (LOG-1);</pre>
                for (int s = LOG; s; s--, tam /= 2) {
                        int i = p \gg s;
                        if (lazy[i]) {
                                 poe(2*i, lazy[i], tam);
                                 poe(2*i+1, lazy[i], tam);
                                 lazy[i] = 0;
        void build(int n2, int* v) {
                for (int i = 0; i < n; i++) seg[n+i] = v[i];</pre>
                for (int i = n-1; i; i--) seg[i] = junta(seg[2*i], seg[2*i+1]);
                for (int i = 0; i < 2*n; i++) lazy[i] = 0;
        11 query(int a, int b) {
                 11 \text{ ret} = 0;
                for (prop(a+=n), prop(b+=n); a \le b; ++a/=2, --b/=2) {
                        if (a%2 == 1) ret = junta(ret, seg[a]);
                        if (b%2 == 0) ret = junta(ret, seg[b]);
                return ret;
        void update(int a, int b, int x) {
                int a2 = a += n, b2 = b += n, tam = 1;
                for (; a <= b; ++a/=2, --b/=2, tam *= 2) {
                        if (a \% 2 == 1) poe(a, x, tam);
                        if (b\%2 == 0) poe(b, x, tam);
                sobe (a2), sobe (b2);
};
```

# 23.10 segTreePa

```
// SegTree PA
//
// Segtree de PA
```

```
// update_set(1, r, A, R) seta [1, r] para PA(A, R),
// update_add soma PA(A, R) em [1, r]
// query(\overline{1}, r) retorna a soma de [1, r]
// PA(A, R) eh a PA: [A+R, A+2R, A+3R, ...]
// Complexidades:
// construir - O(n)
// update_set, update_add, query - O(log(n))
struct seg pa {
        struct Data {
               11 sum;
                11 set_a, set_r, add_a, add_r;
                Data() : sum(0), set_a(LINF), set_r(0), add_a(0), add_r(0) {}
        vector<Data> seq;
        int n:
        seq pa(int n ) {
               n = n;
                seg = vector < Data > (4*n);
        void prop(int p, int 1, int r) {
                int tam = r-1+1:
                ll &sum = seg[p].sum, &set_a = seg[p].set_a, &set_r = seg[p].
                     set r,
                        &add_a = seg[p].add_a, &add_r = seg[p].add_r;
                if (set_a != LINF) {
                        set_a += add_a, set_r += add_r;
                        sum = set_a * tam + set_r * tam * (tam + 1) / 2;
                        if (l != r) {
                                 int m = (1+r)/2;
                                 seg[2*p].set_a = set_a;
                                 seg[2*p].set_r = set_r;
                                 seg[2*p].add_a = seg[2*p].add_r = 0;
                                 seg[2*p+1].set_a = set_a + set_r * (m-l+1);
                                 seg[2*p+1].set_r = set_r;
                                 seg[2*p+1].add_a = seg[2*p+1].add_r = 0;
                        set_a = LINF, set_r = 0;
                        add_a = add_r = 0;
                } else if (add_a or add_r) {
                        sum += add_a*tam + add_r*tam*(tam+1)/2;
                        if (1 != r) {
                                 int m = (1+r)/2;
                                 seg[2*p].add_a += add_a;
                                 seq[2*p].add_r += add_r;
                                 seg[2*p+1].add_a += add_a + add_r * (m-1+1);
                                 seg[2*p+1].add_r += add_r;
                        add a = add r = 0;
        int inter(pair<int, int> a, pair<int, int> b) {
                if (a.first > b.first) swap(a, b);
                return max(0, min(a.second, b.second) - b.first + 1);
        il set(int a, int b, ll aa, ll rr, int p, int l, int r) {
                prop(p, 1, r);
                if (b < 1 or r < a) return seg[p].sum;</pre>
                if (a \le 1 \text{ and } r \le b) {
                        seq[p].set a = aa;
                        seg[p].set_r = rr;
                        prop(p, 1, r);
                        return seq[p].sum;
                int m = (1+r)/2;
                int tam_1 = inter({1, m}, {a, b});
                return seg[p].sum = set(a, b, aa, rr, 2*p, 1, m) +
                        set(a, b, aa + rr * tam_1, rr, 2*p+1, m+1, r);
```

```
void update_set(int 1, int r, 11 aa, 11 rr) {
        set(1, r, aa, rr, 1, 0, n-1);
ill add(int a, int b, ll aa, ll rr, int p, int l, int r) {
        prop(p, 1, r);
        if (b < 1 or r < a) return seg[p].sum;</pre>
        if (a \le 1 \text{ and } r \le b) {
                seg[p].add_a += aa;
                seg[p].add_r += rr;
                 prop(p, 1, r);
                 return seg[p].sum;
        int m = (1+r)/2;
        int tam_1 = inter({1, m}, {a, b});
        return seg[p].sum = add(a, b, aa, rr, 2*p, 1, m) +
                add(a, b, aa + rr * tam_l, rr, 2*p+1, m+1, r);
void update_add(int 1, int r, 11 aa, 11 rr) {
        add(l, r, aa, rr, 1, 0, n-1);
ill query(int a, int b, int p, int l, int r) {
        prop(p, 1, r);
if (b < 1 or r < a) return 0;
        if (a <= 1 and r <= b) return seg[p].sum;</pre>
        int m = (1+r)/2;
        return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
11 query(int 1, int r) { return query(1, r, 1, 0, n-1); }
```

## 23.11 segTreePersistent

};

```
// SegTree Persistente
// SegTree de soma, update de somar numa posicao
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, x, t) faz um update v[a]+=x a partir da
// versao de t, criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
// build - O(n)
// query - O(log(n))
// update - 0(log(n))
const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
const int MAXS = 2*MAX+UPD*LOG;
namespace perseg {
        ll seg[MAXS];
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
        int n, *v;
        11 build(int p, int 1, int r) {
                if (1 == r) return seg[p] = v[1];
                L[p] = cnt++, R[p] = cnt++;
                int m = (1+r)/2;
                return seg[p] = build(L[p], 1, m) + build(R[p], m+1, r);
        void build(int n2, int* v2) {
                n = n2, v = v2;
                rt[0] = cnt++;
                build(0, 0, n-1);
        11 query(int a, int b, int p, int l, int r) {
                if (b < 1 or r < a) return 0;
                if (a <= 1 and r <= b) return seg[p];</pre>
                int m = (1+r)/2;
                return query(a, b, L[p], 1, m) + query(a, b, R[p], m+1, r);
        11 query(int a, int b, int tt) {
                return query(a, b, rt[tt], 0, n-1);
        11 update(int a, int x, int lp, int p, int l, int r) {
                if (l == r) return seg[p] = seg[lp]+x;
```

## 23.12 segTreePersistentComLazy

```
// SegTree Persistente com Lazy
// Nao propaga, meio estranho de mexer, mas da
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, b, x, t) faz um update v[a..b]+=x a partir da
// versao de t, criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
// build - O(n)
// query - O(log(n))
// update - 0(log(n))
const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
const int MAXS = 2*MAX + 4*UPD*LOG;
namespace perseg {
        int seg[MAXS];
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
        int n, *v;
        int build(int p, int 1, int r) {
                if (l == r) return seg[p] = v[l];
                L[p] = cnt++, R[p] = cnt++;
                int m = (1+r)/2;
                return seg[p] = max(build(L[p], 1, m), build(R[p], m+1, r));
        void build(int n2, int *v2) {
                n = n2, v = v2;
                rt[0] = cnt++;
                build(0, 0, n-1);
        int query(int a, int b, int p, int 1, int r) {
                if (b < 1 or r < a) return -INF;</pre>
                if (a <= 1 and r <= b) return lazy[p] + seg[p];</pre>
                int m = (1+r)/2;
                int ret = lazy[p] + max(query(a, b, L[p], l, m), query(a, b, R[p
                     ], m+1, r));
                return ret;
        int query(int a, int b, int tt) {
                return query(a, b, rt[tt], 0, n-1);
        int update(int a, int b, int x, int lp, int p, int l, int r) {
                tie(seg[p], lazy[p], L[p], R[p]) = {seg[lp], lazy[lp], L[lp], R[
                if (b < 1 or r < a) return seg[p] + lazy[p];</pre>
                if (a \le 1 \text{ and } r \le b) \text{ return } seq[p] + (lazy[p] += x);
                int m = (1+r)/2;
                seg[p] = max(update(a, b, x, L[lp], L[p] = cnt++, l, m),
                                          update(a, b, x, R[lp], R[p] = cnt++, m
                                               +1, r));
                lazy[p] = lazy[lp];
                return seg[p] + lazy[p];
        int update(int a, int b, int x, int tt=t) {
                assert(tt <= t);</pre>
                update(a, b, x, rt[tt], rt[++t]=cnt++, 0, n-1);
                return t;
```

# 24 ufmg forked/Extra

## **24.1** debug

};

```
void debug_out(string s, int line) { cerr << endl; }
template<typename H, typename... T>
void debug_out(string s, int line, H h, T... t) {
   if (s[0] != ',') cerr << "Line(" << line << ") ";
   do { cerr << s[0]; s = s.substr(1);
   } while (s.size() and s[0] != ',');
   cerr << " = " << h;
   debug_out(s, line, t...);
}
#ifdef DEBUG
#define debug(...) debug_out(#_VA_ARGS__, __LINE__, __VA_ARGS__)
#else
#define debug(...) 42
#endif</pre>
```

### 24.2 fastIO

```
int read_int() {
    bool minus = false;
    int result = 0;
    char ch;
    ch = getchar();
    while (1) {
        if (ch == '-') break;
if (ch >= '0' && ch <= '9') break;</pre>
        ch = getchar();
    if (ch == '-') minus = true;
    else result = ch-'0';
    while (1) {
        ch = getchar();
        if (ch < '0' || ch > '9') break;
        result = result *10 + (ch - '0');
    if (minus) return -result;
    else return result;
```

### 24.3 hash

```
# Para usar (hash das linhas [11, 12]):
# bash hash.sh arquivo.cpp 11 12
sed -n $2','$3' p' $1 | sed '/^#w/d' | cpp -dD -P -fpreprocessed | tr -d '[:
    space:]' | md5sum | cut -c-6
```

## 24.4 pragma

```
// Otimizacoes agressivas, pode deixar mais rapido ou mais devagar
#pragma GCC optimize("Ofast")
// Auto explicativo
#pragma GCC optimize("unroll-loops")
// Vetorizacao
#pragma GCC target("avx2")
// Para operacoes com bits
#pragma GCC target("bmi,bmi2,popent,lzent")
```

### 24.5 rand

```
mt19937 rng((int) chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r){
          uniform_int_distribution<int> uid(1, r);
          return uid(rng);
}
```

### 24.6 stress

```
P=a
make {P} {P}2 gen || exit 1
for ((i = 1; ; i++)) do
        ./gen $i > in
        ./${P} < in > out
        ./${P}2 < in > out2
        if (! cmp -s out out2) then
                echo "--> entrada:"
                cat in
                echo "--> saida1:"
                cat out
                echo "--> saida2:"
                cat out2
                break;
        fi
        echo $i
done
```

## 24.7 template

### 24.8 timer

```
// timer T; T() -> retorna o tempo em ms desde que declarou
using namespace chrono;
struct timer : high_resolution_clock {
        const time_point start;
        timer(): start(now()) {}
        int operator()() {
            return duration_cast<milliseconds>(now() - start).count();
        }
};
```

# 25 ufmg forked/Grafos

### 25.1 articulationPoints

```
// Articulation Points
// Computa os pontos de articulação (vertices criticos) de um grafo
// art[i] armazena o numero de novas componentes criadas ao deletar vertice i
// se art[i] >= 1, entao vertice i eh ponto de articulação
// O(n+m)
vector<vector<int>> g;
stack<int> s;
vector<int> id, art;
int dfs_art(int i, int& t, int p = -1) {
        int lo = id[i] = t++;
        s.push(i);
        for (int j : g[i]) if (j != p) {
                if (id[j] == -1) {
                        int val = dfs_art(j, t, i);
                        lo = min(lo, val);
                        if (val >= id[i]) {
                                art[i]++;
                                while (s.top() != j) s.pop();
                                s.pop();
                         // if (val > id[i]) aresta i-j eh ponte
                else lo = min(lo, id[j]);
        if (p == -1 and art[i]) art[i]--;
        return lo;
void compute_art_points() {
        id = vector<int>(n, -1);
        art = vector<int>(n, 0);
        int t = 0;
        for (int i = 0; i < n; i++) if (id[i] == -1)</pre>
                dfs_art(i, t, -1);
```

## 25.2 bellmanFord

```
// Bellman-Ford
// Calcula a menor distancia
// entre a e todos os vertices e
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
// O(nm)
int n, m;
int d[MAX];
vector<pair<int, int>> ar; // vetor de arestas
vector<int> w;
                             // peso das arestas
bool bellman_ford(int a) {
        for (int i = 0; i < n; i++) d[i] = INF;</pre>
        d[a] = 0;
        for (int i = 0; i \le n; i++)
                for (int j = 0; j < m; j++) {
    if (d[ar[j].second] > d[ar[j].first] + w[j]) {
                                  if (i == n) return 1;
                                  d[ar[j].second] = d[ar[j].first] + w[j];
        return 0;
```

### 25.3 blockCutTree

```
// Block-Cut Tree
// Cria a block-cut tree, uma arvore com os blocos
// e os pontos de articulação
// Blocos sao componentes 2-vertice-conexos maximais
// Uma 2-coloração da arvore eh tal que uma cor sao
// os blocos, e a outra cor sao os pontos de art.
// Funciona para grafo nao conexo
// art[i] responde o numero de novas componentes conexas
// criadas apos a remocao de i do grafo g
// Se art[i] >= 1, i eh ponto de articulação
// Para todo i <= blocks.size()</pre>
// blocks[i] eh uma componente 2-vertce-conexa maximal
// edgblocks[i] sao as arestas do bloco i
// tree[i] eh um vertice da arvore que corresponde ao bloco i
// pos[i] responde a qual vertice da arvore vertice i pertence
// Arvore tem no maximo 2n vertices
// O(n+m)
struct block_cut_tree {
        vector<vector<int>> g, blocks, tree;
        vector<vector<pair<int, int>>> edgblocks;
        stack<int> s;
        stack<pair<int, int>> s2;
        vector<int> id, art, pos;
        block_cut_tree(vector<vector<int>> q_) : q(q_) {
                int n = q.size();
                id.resize(n, -1), art.resize(n), pos.resize(n);
                build();
        int dfs(int i, int& t, int p = -1) {
                int lo = id[i] = t++;
                s.push(i);
                if (p != -1) s2.emplace(i, p);
for (int j : g[i]) if (j != p and id[j] != -1) s2.emplace(i, j);
                for (int j : g[i]) if (j != p) {
                        if (id[j] == -1) {
                                 int val = dfs(j, t, i);
                                 lo = min(lo, val);
                                 if (val >= id[i]) {
                                         art[i]++;
                                         blocks.emplace_back(1, i);
                                         while (blocks.back().back() != j)
                                                 blocks.back().push_back(s.top())
                                                      , s.pop();
                                         edgblocks.emplace_back(1, s2.top()), s2.
                                              pop();
                                         while (edgblocks.back().back() != pair()
                                              , i))
                                                 edgblocks.back().push back(s2.
                                                      top()), s2.pop();
                                 // if (val > id[i]) aresta i-j eh ponte
                        else lo = min(lo, id[j]);
                if (p == -1 and art[i]) art[i]--;
                return lo;
        void build() {
                int t = 0;
                for (int i = 0; i < g.size(); i++) if (id[i] == -1) dfs(i, t,
```

### 25.4 blossom

// Blossom

```
// Matching maximo em grafo geral
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
vector<int> g[MAX];
int match[MAX]; // match[i] = com quem i esta matchzado ou -1
int n, pai[MAX], base[MAX], vis[MAX];
queue<int> q;
void contract(int u, int v, bool first = 1) {
        static vector<bool> bloss;
        static int 1;
        if (first) {
                bloss = vector<bool>(n, 0);
                vector<bool> teve(n, 0);
                int k = u; l = v;
                while (1) {
                        teve[k = base[k]] = 1;
                        if (match[k] == -1) break;
                        k = pai[match[k]];
                while (!teve[l = base[l]]) l = pai[match[l]];
        while (base[u] != 1) {
                bloss[base[u]] = bloss[base[match[u]]] = 1;
                pai[u] = v;
                v = match[u];
                u = pai[match[u]];
        if (!first) return;
        contract(v, u, 0);
        for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
               base[i] = 1;
                if (!vis[i]) q.push(i);
                vis[i] = 1;
int getpath(int s) {
        for (int i = 0; i < n; i++) base[i] = i, pai[i] = -1, vis[i] = 0;
        vis[s] = 1; q = queue<int>(); q.push(s);
        while (q.size()) {
                int u = q.front(); q.pop();
                for (int i : g[u]) {
                        if (base[i] == base[u] or match[u] == i) continue;
                        if (i == s or (match[i] != -1 and pai[match[i]] != -1))
                                contract(u, i);
                        else if (pai[i] == -1) {
                                pai[i] = u;
                                if (match[i] == -1) return i;
                                i = match[i];
                                vis[i] = 1; q.push(i);
```

```
return -1;
int blossom() {
         int ans = 0;
         memset(match, -1, sizeof(match));
         for (int i = 0; i < n; i++) if (match[i] == -1)</pre>
                  for (int j : g[i]) if (match[j] == -1) {
    match[i] = j;
                            match[j] = i;
                            ans++;
                            break;
         for (int i = 0; i < n; i++) if (match[i] == -1) {</pre>
                  int j = getpath(i);
if (j == -1) continue;
                   ans++;
                   while (j != -1) {
                            int p = pai[j], pp = match[p];
match[p] = j;
                            match[j] = p;
                            j = pp;
         return ans;
```

### 25.5 center

```
// Centro de arvore
// Retorna o diametro e o(s) centro(s) da arvore
// Uma arvore tem sempre um ou dois centros e estes estao no meio do diametro
// O(n)
vector<int> g[MAX];
int d[MAX], par[MAX];
pair<int, vector<int>> center() {
         int f, df;
        function<void(int)> dfs = [&] (int v) {
    if (d[v] > df) f = v, df = d[v];
                  for (int u : g[v]) if (u != par[v])
                          d[u] = d[v] + 1, par[u] = v, dfs(u);
         };
         f = df = par[0] = -1, d[0] = 0;
        dfs(0);
        int root = f;
         f = df = par[root] = -1, d[root] = 0;
         dfs(root);
         vector<int> c;
                 if (d[f] == df/2 \text{ or } d[f] == (df+1)/2) \text{ c.push_back}(f);
                  f = par[f];
         return {df, c};
```

### 25.6 centroid

```
// Centroid
//
// Computa os 2 centroids da arvore
//
// O(n)
int n, subsize[MAX];
vector<int> g[MAX];
```

## 25.7 centroidDecomp

```
// Centroid decomposition
// decomp(0, k) computa numero de caminhos com 'k' arestas
// Mudar depois do comentario
// O(n log(n))
vector<int> g[MAX];
int sz[MAX], rem[MAX];
void dfs(vector<int>& path, int i, int l=-1, int d=0) {
        path.push_back(d);
        for (int j : q[i]) if (j != 1 and !rem[j]) dfs(path, j, i, d+1);
int dfs_sz(int i, int l=-1) {
        sz[i] = 1;
        for (int j : g[i]) if (j != 1 and !rem[j]) sz[i] += dfs_sz(j, i);
        return sz[i];
int centroid(int i, int 1, int size) {
        for (int j : q[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
                return centroid(j, i, size);
        return i;
11 decomp(int i, int k) {
        int c = centroid(i, i, dfs_sz(i));
        rem[c] = 1;
        // gasta O(n) aqui - dfs sem ir pros caras removidos
        11 \text{ ans} = 0;
        vector<int> cnt(sz[i]);
        cnt[0] = 1;
        for (int j : g[c]) if (!rem[j]) {
                vector<int> path;
                dfs(path, j);
                for (int d : path) if (0 \le k-d-1) and k-d-1 \le sz[i])
                        ans += cnt[k-d-1];
                for (int d : path) cnt[d+1]++;
        for (int j : q[c]) if (!rem[j]) ans += decomp(j, k);
        rem[c] = 0;
        return ans;
```

### 25.8 centroidTree

```
// Centroid Tree
// Constroi a centroid tree
// p[i] eh o pai de i na centroid-tree
// dist[i][k] = distancia na arvore original entre i
// e o k-esimo ancestral na arvore da centroid
// O(n log(n)) de tempo e memoria
vector<int> q[MAX], dist[MAX];
int sz[MAX], rem[MAX], p[MAX];
int dfs_sz(int i, int l=-1) {
        sz[i] = 1;
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j, i);
        return sz[i];
int centroid(int i, int l, int size) {
        for (int j : g[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
                 return centroid(j, i, size);
        return i;
void dfs_dist(int i, int l, int d=0) {
        dist[i].push_back(d);
        for (int j : g[i]) if (j != l and !rem[j])
                dfs_dist(j, i, d+1);
}
void decomp(int i, int l = -1) {
    int c = centroid(i, i, dfs_sz(i));
        rem[c] = 1, p[c] = 1;
        dfs_dist(c, c);
        for (int j : g[c]) if (!rem[j]) decomp(j, c);
void build(int n) {
        for (int i = 0; i < n; i++) rem[i] = 0, dist[i].clear();</pre>
        for (int i = 0; i < n; i++) reverse(dist[i].begin(), dist[i].end());</pre>
```

### 25.9 cover

```
// Vertex cover
// Encontra o tamanho do vertex cover minimo
// Da pra alterar facil pra achar os vertices
// Parece rodar com < 2 s pra N = 90
// O(n * 1.38^n)
namespace cover {
        const int MAX = 96;
        vector<int> g[MAX];
        bitset<MAX> bs[MAX];
        int n:
        void add(int i, int j) {
                if (i == j) return;
                n = \max(\{n, i+1, j+1\});
                bs[i][j] = bs[j][i] = 1;
        int rec(bitset<MAX> m) {
                int ans = 0;
                for (int x = 0; x < n; x++) if (m[x]) {
                        bitset < MAX > comp;
                        function<void(int)> dfs = [&](int i) {
                                 comp[i] = 1, m[i] = 0;
                                 for (int j : g[i]) if (m[j]) dfs(j);
                        };
```

```
dfs(x);
                 int ma, deg = -1, cyc = 1;
for (int i = 0; i < n; i++) if (comp[i]) {</pre>
                          int d = (bs[i]&comp).count();
                          if (d <= 1) cyc = 0;
                          if (d > deg) deg = d, ma = i;
                 if (deg <= 2) { // caminho ou ciclo</pre>
                          ans += (comp.count() + cyc) / 2;
                          continue;
                 comp[ma] = 0;
                 // ou ta no cover, ou nao ta no cover
                 ans += min(1 + rec(comp), deg + rec(comp & ~bs[ma]));
        return ans:
int solve() {
        bitset < MAX > m;
        for (int i = 0; i < n; i++) {
                m[i] = 1;
                 for (int j = 0; j < n; j++)
                          if (bs[i][j]) g[i].push_back(j);
        return rec(m);
```

## 25.10 dijkstra

```
// Dijkstra
// encontra menor distancia de x
// para todos os vertices
// se ao final do algoritmo d[i] = LINF,
// entao x nao alcanca i
// O(m \log(n))
11 d[MAX];
vector<pair<int, int>> q[MAX]; // {vizinho, peso}
void dijkstra(int v) {
        for (int i = 0; i < n; i++) d[i] = LINF;</pre>
        d[v] = 0;
        priority_queue<pair<11, int>> pq;
        pq.emplace(0, v);
        while (pq.size()) {
                auto [ndist, u] = pq.top(); pq.pop();
                if (-ndist > d[u]) continue;
                for (auto [idx, w] : q[u]) if (d[idx] > d[u] + w) {
                        d[idx] = d[u] + w;
                        pq.emplace(-d[idx], idx);
```

## 25.11 dinitz

```
// Dinitz
//
// O(min(m * max_flow, n^2 m))
// Grafo com capacidades 1: O(min(m sqrt(m), m * n^(2/3)))
// Todo vertice tem grau de entrada ou saida 1: O(m sqrt(n))
struct dinitz {
```

```
const bool scaling = false; // com scaling -> O(nm log(MAXCAP)),
        int lim;
                                     // com constante alta
        struct edge {
                int to, cap, rev, flow;
                bool res;
                edge(int to_, int cap_, int rev_, bool res_)
                        : to(to_), cap(cap_), rev(rev_), flow(0), res(res_) {}
        };
        vector<vector<edge>> g;
        vector<int> lev, beg;
        11 F;
        dinitz(int n) : g(n), F(0) {}
        void add(int a, int b, int c) {
                g[a].emplace_back(b, c, g[b].size(), false);
                g[b].emplace_back(a, 0, g[a].size()-1, true);
        bool bfs(int s, int t) {
                lev = vector<int>(g.size(), -1); lev[s] = 0;
                beg = vector<int>(g.size(), 0);
                queue<int> q; q.push(s);
                while (q.size()) {
                         int u = q.front(); q.pop();
                        for (auto& i : g[u]) {
                                 if (lev[i.to] != -1 or (i.flow == i.cap))
                                      continue;
                                 if (scaling and i.cap - i.flow < lim) continue;</pre>
                                 lev[i.to] = lev[u] + 1;
                                 q.push(i.to);
                return lev[t] != -1;
        int dfs(int v, int s, int f = INF) {
    if (!f or v == s) return f;
                for (int& i = beg[v]; i < g[v].size(); i++) {</pre>
                         auto& e = g[v][i];
                         if (lev[e.to] != lev[v] + 1) continue;
                        int foi = dfs(e.to, s, min(f, e.cap - e.flow));
                         if (!foi) continue;
                        e.flow += foi, g[e.to][e.rev].flow -= foi;
                        return foi;
                return 0;
        11 max_flow(int s, int t) {
                for (lim = scaling ? (1<<30) : 1; lim; lim /= 2)</pre>
                        while (bfs(s, t)) while (int ff = dfs(s, t)) F += ff;
                return F;
// Recupera as arestas do corte s-t
vector<pair<int, int>> get_cut(dinitz& g, int s, int t) {
        q.max_flow(s, t);
        vector<pair<int, int>> cut;
        vector<int> vis(g.g.size(), 0), st = {s};
        vis[s] = 1;
        while (st.size()) {
                int u = st.back(); st.pop_back();
                for (auto e : g.g[u]) if (!vis[e.to] and e.flow < e.cap)</pre>
                        vis[e.to] = 1, st.push_back(e.to);
        for (int i = 0; i < q.q.size(); i++) for (auto e : q.q[i])
                if (vis[i] and !vis[e.to] and !e.res) cut.emplace_back(i, e.to);
        return cut;
```

#### 25.12directedMst

};

}

```
// AGM Direcionada
// Fala o menor custo para selecionar arestas tal que
// o vertice 'r' alcance todos
// Se nao tem como, retorna LINF
```

```
,,
// O(m log(n))
struct node {
        pair<ll, int> val;
        ll lazy;
        node *1, *r;
        node() {}
        node(pair<int, int> v) : val(v), lazy(0), l(NULL), r(NULL) {}
        void prop() {
                val.first += lazy;
                if (1) 1->lazy += lazy;
                if (r) r->lazy += lazy;
                lazy = 0;
void merge(node*& a, node* b) {
        if (!a) swap(a, b);
        if (!b) return;
        a->prop(), b->prop();
        if (a->val > b->val) swap(a, b);
        merge (rand() 2? a \rightarrow 1: a \rightarrow r, b;
pair<11, int> pop(node*& R) {
        R->prop();
        auto ret = R->val;
        node* tmp = R;
        merge(R->1, R->r);
        R = R -> 1;
        if (R) R->lazy -= ret.first;
        delete tmp;
        return ret:
void apaga(node* R) { if (R) apaga(R->1), apaga(R->r), delete R; }
11 dmst(int n, int r, vector<pair<pair<int, int>, int>>& ar) {
        vector<int> p(n); iota(p.begin(), p.end(), 0);
        function<int(int)> find = [&](int k) { return p[k] == k?k:p[k] = find(p[k]);
              };
        vector<node*> h(n);
        for (auto e : ar) merge(h[e.first.second], new node({e.second, e.first.
             first }));
        vector<int> pai(n, -1), path(n);
        pai[r] = r;
        11 \text{ ans} = 0;
        for (int i = 0; i < n; i++) { // vai conectando todo mundo
                int u = i, at = 0;
                while (pai[u] == -1)
                         if (!h[u]) { // nao tem
                                 for (auto i : h) apaga(i);
                                 return LINF;
                         path[at++] = u, pai[u] = i;
                         auto [mi, v] = pop(h[u]);
                         ans += mi;
                         if (pai[u = find(v)] == i) { // ciclo
                                 while (find(v = path[--at]) != u)
                                         merge(h[u], h[v]), h[v] = NULL, p[find(v)]
                                              ) ] = u;
                                 pai[u] = -1;
        for (auto i : h) apaga(i);
        return ans;
```

### 25.13 dominatorTree

```
// Dominator Tree
// Codigo do Kawakami. Se vira pra usar ai
```

```
// dominates - O(1)
int n;
namespace d_tree {
       vector<int> g[MAX];
        // The dominator tree
        vector<int> tree[MAX];
        int dfs_l[MAX], dfs_r[MAX];
        // Auxiliary data
        vector<int> rg[MAX], bucket[MAX];
        int idom[MAX], sdom[MAX], prv[MAX], pre[MAX];
        int ancestor[MAX], label[MAX];
        vector<int> preorder;
        void dfs(int v) {
                static int t = 0;
                pre[v] = ++t;
                sdom[v] = label[v] = v;
                preorder.push_back(v);
                for (int nxt: g[v]) {
                        if (sdom[nxt] == -1) {
                                prv[nxt] = v;
                                dfs(nxt);
                        rg[nxt].push_back(v);
        int eval(int v) {
                if (ancestor[v] == -1) return v;
                if (ancestor[ancestor[v]] == -1) return label[v];
                int u = eval(ancestor[v]);
                if (pre[sdom[u]] < pre[sdom[label[v]]]) label[v] = u;</pre>
                ancestor[v] = ancestor[u];
                return label[v];
        void dfs2(int v) {
                static int t = 0;
                dfs_1[v] = t++;
                for (int nxt: tree[v]) dfs2(nxt);
                dfs r[v] = t++;
        void build(int s) {
                for (int i = 0; i < n; i++) {
                        sdom[i] = pre[i] = ancestor[i] = -1;
                        rg[i].clear();
                        tree[i].clear();
                        bucket[i].clear();
                preorder.clear();
                dfs(s);
                if (preorder.size() == 1) return;
                for (int i = int(preorder.size()) - 1; i >= 1; i--) {
                        int w = preorder[i];
                        for (int v: rg[w]) {
                                int u = eval(v);
                                if (pre[sdom[u]] < pre[sdom[w]]) sdom[w] = sdom[</pre>
                                     ul;
                        bucket[sdom[w]].push_back(w);
                        ancestor[w] = prv[w];
                        for (int v: bucket[prv[w]]) {
                                int u = eval(v);
                                idom[v] = (u == v) ? sdom[v] : u;
                        bucket[prv[w]].clear();
                for (int i = 1; i < preorder.size(); i++) {</pre>
                        int w = preorder[i];
                        if (idom[w] != sdom[w]) idom[w] = idom[idom[w]];
                        tree[idom[w]].push_back(w);
                idom[s] = sdom[s] = -1;
                dfs2(s);
```

// build - O(m log(n))

```
// Whether every path from s to v passes through u
bool dominates(int u, int v) {
    if (pre[v] == -1) return 1; // vacuously true
        return dfs_l[u] <= dfs_l[v] && dfs_r[v] <= dfs_r[u];
}
};</pre>
```

### 25.14 eulerPath

```
// Euler Path / Euler Cycle
// Para declarar: 'euler<true> E(n);' se quiser
// direcionado e com 'n' vertices
// As funcoes retornam um par com um booleano
// indicando se possui o cycle/path que voce pediu,
// e um vector de {vertice, id da aresta para chegar no vertice}
// Se for get_path, na primeira posicao o id vai ser -1
// get_path(src) tenta achar um caminho ou ciclo euleriano
// comecando no vertice 'src'.
// Se achar um ciclo, o primeiro e ultimo vertice serao 'src'.
// Se for um P3, um possiveo retorno seria [0, 1, 2, 0]
// get_cycle() acha um ciclo euleriano se o grafo for euleriano.
// Se for um P3, um possivel retorno seria [0, 1, 2]
// (vertie inicial nao repete)
// O(n+m)
template<bool directed=false> struct euler {
        int n:
        vector<vector<pair<int, int>>> q;
        vector<int> used;
       euler(int n_) : n(n_), g(n) {}
void add(int a, int b) {
                int at = used.size();
                used.push back(0);
                g[a].emplace_back(b, at);
                if (!directed) g[b].emplace_back(a, at);
#warning chamar para o src certo!
       pair<bool, vector<pair<int, int>>> get_path(int src) {
                if (!used.size()) return {true, {}};
                vector<int> beg(n, 0);
                for (int& i: used) i = 0;
                // {{vertice, anterior}, label}
                vector<pair<int, int>, int>> ret, st = {{{src, -1}, -1}};
                while (st.size()) {
                        int at = st.back().first.first;
                        int& it = beg[at];
                        while (it < g[at].size() and used[g[at][it].second]) it</pre>
                        if (it == g[at].size()) {
                                if (ret.size() and ret.back().first.second != at
                                        return {false, {}};
                                ret.push_back(st.back()), st.pop_back();
                        } else {
                                st.push_back({{g[at][it].first, at}, g[at][it].
                                    second });
                                used[g[at][it].second] = 1;
                if (ret.size() != used.size()+1) return {false, {}};
                vector<pair<int, int>> ans;
                for (auto i : ret) ans.emplace_back(i.first.first, i.second);
                reverse(ans.begin(), ans.end());
                return {true, ans};
        pair<bool, vector<pair<int, int>>> get_cycle() {
                if (!used.size()) return {true, {}};
                int src = 0;
                while (!g[src].size()) src++;
                auto ans = get_path(src);
                if (!ans.first or ans.second[0].first != ans.second.back().first
                        return {false, {}};
```

```
ans.second[0].second = ans.second.back().second;
ans.second.pop_back();
return ans;
}
```

### 25.15 eulerTourTree

```
// Euler Tour Tree
// Mantem uma floresta enraizada dinamicamente
// e permite queries/updates em sub-arvore
// Chamar ETT E(n, v), passando n = numero de vertices
// e v = vector com os valores de cada vertice (se for vazio,
// constroi tudo com 0
// link(v, u) cria uma aresta de v pra u, de forma que u se torna
// o pai de v (eh preciso que v seja raiz anteriormente)
// cut(v) corta a resta de v para o pai
// query(v) retorna a soma dos valores da sub-arvore de v
// update(v, val) soma val em todos os vertices da sub-arvore de v
// update_v(v, val) muda o valor do vertice v para val
// is_in_subtree(v, u) responde se o vertice u esta na sub-arvore de v
// Tudo O(log(n)) com alta probabilidade
mt19937 rng((int) chrono::steady_clock::now().time_since_epoch().count());
template<typename T> struct ETT {
        // treap
        struct node {
                node *1, *r, *p;
                int pr, sz;
                T val, sub, lazy;
                int id;
                bool f; // se eh o 'first'
                int qt_f; // numero de firsts na subarvore
                node(int id_, T v, bool f_ = 0) : 1(NULL), r(NULL), p(NULL), pr(
                        sz(1), val(v), sub(v), lazy(), id(id_), f(f_), qt_f(f_)
                             { }
                void prop()
                        if (lazy != T()) {
                                 if (f) val += lazv;
                                 sub += lazy*sz;
                                 if (1) 1->lazy += lazy;
                                 if (r) r->lazy += lazy;
                         lazy = T();
                void update() {
                         sz = 1, sub = val, qt_f = f;
                         if (1) 1->prop(), sz += 1->sz, sub += 1->sub, qt_f += 1
                             ->qt_f;
                         if (r) r->prop(), sz += r->sz, sub += r->sub, qt_f += r
                              ->qt_f;
        };
        node* root:
        int size(node* x) { return x ? x->sz : 0; }
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
                if (!1 or !r) return void(i = 1 ? 1 : r);
                1->prop(), r->prop();
                if (1->pr > r->pr) join(1->r, r, 1->r), 1->r->p = i = 1;
                else join(l, r->l, r->l), r->l->p = i = r;
                i->update();
        void split(node* i, node*& l, node*& r, int v, int key = 0) {
                if (!i) return void(r = 1 = NULL);
                i->prop();
                if (\text{key} + \text{size}(i\rightarrow 1) < v) {
                         split(i->r, i->r, r, v, key+size(i->l)+1), l = i;
                         if (r) r \rightarrow p = NULL;
```

```
if (i->r) i->r->p = i;
        } else {
                 split(i->1, 1, i->1, v, key), r = i;
                if (1) 1->p = NULL;
                if (i->1) i->1->p = i;
        i->update();
int get_idx(node* i) {
        int ret = size(i->1);
        for (; i->p; i = i->p) {
                node* pai = i->p;
                if (i != pai->1) ret += size(pai->1) + 1;
        return ret;
node* get_min(node* i) {
        if (!i) return NULL;
        return i->1 ? get_min(i->1) : i;
node* get_max(node* i)
        if (!i) return NULL;
        return i->r ? get_max(i->r) : i;
// fim da treap
vector<node*> first, last;
ETT(int n, vector<T> v = {}) : root(NULL), first(n), last(n) {
        if (!v.size()) v = vector<T>(n);
        for (int i = 0; i < n; i++) {
                first[i] = last[i] = new node(i, v[i], 1);
join(root, first[i], root);
ETT(const ETT& t) { throw logic_error("Nao copiar a ETT!"); }
~ETT() {
        vector<node*> q = {root};
        while (q.size()) {
                node* x = q.back(); q.pop_back();
                 if (!x) continue;
                q.push_back(x->1), q.push_back(x->r);
                delete x;
pair<int, int> get_range(int i) {
        return {get_idx(first[i]), get_idx(last[i])};
void link(int v, int u) { // 'v' tem que ser raiz
        auto [lv, rv] = get_range(v);
        int ru = get_idx(last[u]);
        node* V;
        node *L, *M, *R;
        split(root, M, R, rv+1), split(M, L, M, lv);
        V = M;
        join(L, R, root);
        split(root, L, R, ru+1);
        join(L, V, L);
         join(L, last[u] = new node(u, T() /* elemento neutro */), L);
        join(L, R, root);
void cut(int v) {
        auto [1, r] = get_range(v);
        node *L, *M, *R;
        split(root, M, R, r+1), split(M, L, M, 1);
        node *LL = get_max(L), *RR = get_min(R);
if (LL and RR and LL->id == RR->id) { // remove duplicata
                 if (last[RR->id] == RR) last[RR->id] = LL;
                 node *A, *B;
                  split(R, A, B, 1);
                  delete A;
                 R = B:
        join(L, R, root);
```

join(root, M, root);

```
T query(int v) {
        auto [1, r] = get_range(v);
        node *L, *M, *R;
        split(root, M, R, r+1), split(M, L, M, 1);
        T ans = M->sub;
        join(L, M, M), join(M, R, root);
        return ans;
void update(int v, T val) { // soma val em todo mundo da subarvore
    auto [1, r] = get_range(v);
        node *L, *M, *R;
        split(root, M, R, r+1), split(M, L, M, 1);
        M->lazy += val;
        join(L, M, M), join(M, R, root);
void update_v(int v, T val) { // muda o valor de v pra val
        int 1 = get_idx(first[v]);
        node *L, *M, *R;
        split(root, M, R, 1+1), split(M, L, M, 1);
        M->val = M->sub = val;
        join(L, M, M), join(M, R, root);
bool is_in_subtree(int v, int u) { // se u ta na subtree de v
        auto [lv, rv] = get_range(v);
        auto [lu, ru] = get_range(u);
        return lv <= lu and ru <= rv;
void print(node* i) {
        if (!i) return;
        print(i->1);
        cout << i->id+1 << " ";
        print(i->r);
void print() { print(root); cout << endl; }</pre>
```

## 25.16 floydWarshall

};

```
// Floyd-Warshall
// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta
// (i, j) de peso w, INF caso contrario
// O(n^3)
int n;
int d[MAX][MAX];
bool floyd_warshall() {
        for (int k = 0; k < n; k++)
        for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
                d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
        for (int i = 0; i < n; i++)
                if (d[i][i] < 0) return 1;</pre>
        return 0;
}
```

## 25.17 functionalGraph

```
// Functional Graph
//
// rt[i] fala o ID da raiz associada ao vertice i
// d[i] fala a profundidade (0 sse ta no ciclo)
```

```
// pos[i] fala a posicao de i no array que eh a concat. dos ciclos
// build(f, val) recebe a funcao f e o custo de ir de
// i para f[i] (por default, val = f)
// f_k(i, k) fala onde i vai parar se seguir k arestas
// path(i, k) fala o custo (soma) seguir k arestas a partir de i
// Se quiser outra operacao, da pra alterar facil o codigo
// Codigo um pouco louco, tenho que admitir
// build - O(n)
// f_k - O(\log(\min(n, k)))
// path - O(log(min(n, k)))
namespace func_graph {
        int f[MAX], vis[MAX], d[MAX];
        int p[MAX], pp[MAX], rt[MAX], pos[MAX];
        int sz[MAX], comp;
        vector<vector<int>> ciclo;
        11 val[MAX], jmp[MAX], seg[2*MAX];
        11 op(11 a, 11 b) { return a+b; }; // mudar a operacao aqui
        void dfs(int i, int t = 2) {
                vis[i] = t;
                if (vis[f[i]] \ge 2) { // comeca ciclo - f[i] eh o rep.
                         d[i] = 0, rt[i] = comp;
sz[comp] = t - vis[f[i]] + 1;
                         p[i] = pp[i] = i, jmp[i] = val[i];
                         ciclo emplace back();
                         ciclo.back().push_back(i);
                 } else {
                         if (!vis[f[i]]) dfs(f[i], t+1);
                         rt[i] = rt[f[i]];
                         if (sz[comp]+1) { // to no ciclo
                                 d[i] = 0;
                                 p[i] = pp[i] = i, jmp[i] = val[i];
                                  ciclo.back().push_back(i);
                         } else { // nao to no ciclo
                                 d[i] = d[f[i]]+1, p[i] = f[i];
pp[i] = 2*d[pp[f[i]]] == d[pp[pp[f[i]]]]+d[f[i]]
                                       ? pp[pp[f[i]]] : f[i];
                                  jmp[i] = pp[i] == f[i] ? val[i] : op(val[i], op(
                                       jmp[f[i]], jmp[pp[f[i]]]));
                if (f[ciclo[rt[i]][0]] == i) comp++; // fim do ciclo
        void build(vector<int> f_, vector<int> val_ = {}) {
                n = f_size(), comp = 0;
                 if (!val_.size()) val_ = f_;
                 for (int i = 0; i < n; i++)
                         f[i] = f_{i}, val[i] = val_{i}, vis[i] = 0, sz[i] = -1;
                 ciclo.clear();
                 for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
                 int t = 0;
                 for (auto& c : ciclo) {
                         reverse(c.begin(), c.end());
                         for (int j : c) {
                                 pos[j] = t;
                                 seg[n+t] = val[j];
                                 t++;
                 for (int i = n-1; i; i--) seq[i] = op(seq[2*i], seq[2*i+1]);
        int f_k(int i, ll k) {
                 while (d[i] and k) {
                         int big = d[i] - d[pp[i]];
if (big <= k) k -= big, i = pp[i];
                         else k--, i = p[i];
                 if (!k) return i;
                return ciclo[rt[i]][(pos[i] - pos[ciclo[rt[i]][0]] + k) % sz[rt[
        11 path(int i, 11 k) {
                 auto query = [&](int 1, int r) {
```

```
11 q = 0;
                         for (1 += n, r += n; 1 <= r; ++1/=2, --r/=2) {
                                 if (1\%2 == 1) q = op(q, seg[1]);
                                 if (r%2 == 0) q = op(q, seg[r]);
                         return q;
                 11 \text{ ret} = 0;
                 while (d[i] and k) {
                         int big = d[i] - d[pp[i]];
                         if (big <= k) k -= big, ret = op(ret, jmp[i]), i = pp[i</pre>
                         else k--, ret = op(ret, val[i]), i = p[i];
                 if (!k) return ret;
                int first = pos[ciclo[rt[i]][0]], last = pos[ciclo[rt[i]].back()
                     1;
                 // k/sz[rt[i]] voltas completas
                 if (k/sz[rt[i]]) ret = op(ret, k/sz[rt[i]] * query(first, last))
                 k %= sz[rt[i]];
                 if (!k) return ret;
                 int l = pos[i], r = first + (pos[i] - first + k - 1) % sz[rt[i]
                     ]];
                 if (1 <= r) return op(ret, query(1, r));</pre>
                 return op(ret, op(query(l, last), query(first, r)));
}
```

# 25.18 hopcroftKarp

```
// Hopcroft Karp
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// O(|E| * sqrt(|V|)) com constante baixa
// Para grafos esparsos gerados aleatoriamente, roda em O(|E| * log(|V|))
// com alta probabilidade
mt19937 rng((int) chrono::steady_clock::now().time_since_epoch().count());
struct hopcroft_karp {
        int n, m;
        vector<vector<int>> q;
        vector<int> dist, nxt, ma, mb;
        hopcroft_karp(int n_, int m_) : n(n_), m(m_), g(n),
                dist(n), nxt(n), ma(n, -1), mb(m, -1) {}
        void add(int a, int b) { g[a].push_back(b); }
        bool dfs(int i) {
                for (int &id = nxt[i]; id < g[i].size(); id++) {</pre>
                         int j = g[i][id];
                         if (mb[j] == -1 \text{ or } (dist[mb[j]] == dist[i] + 1 \text{ and } dfs(
                              mb[j]))) {
                                 ma[i] = j, mb[j] = i;
                                 return true;
                return false;
        bool bfs() {
                for (int i = 0; i < n; i++) dist[i] = n;</pre>
                queue<int> q;
                for (int i = 0; i < n; i++) if (ma[i] == -1) {
                        dist[i] = 0;
                        q.push(i);
```

```
bool rep = 0;
                 while (q.size()) {
                         int i = q.front(); q.pop();
for (int j : g[i]) {
                                  if (mb[j] == -1) rep = 1;
                                  else if (dist[mb[j]] > dist[i] + 1) {
                                           dist[mb[j]] = dist[i] + 1;
                                           q.push(mb[j]);
                 return rep;
        int matching() {
                 int ret = 0;
                 for (auto& i : g) shuffle(i.begin(), i.end(), rng);
                 while (bfs()) {
                          for (int i = 0; i < n; i++) nxt[i] = 0;</pre>
                          for (int i = 0; i < n; i++)
                                  if (ma[i] == -1 and dfs(i)) ret++;
                 return ret;
};
```

### 25.19 johnson

```
// Johnson
// funciona iqual ao Floyd-Warshall
// encontra o menor caminho entre todo
// par de vertices e retorna 1 sse tem
// ciclo negativo no grafo
// O(nm log(m))
vector<pair<int, 11>> q[MAX]; // {vizinho, peso}
11 d[MAX][MAX];
bool johnson(int n) {
        vector<ll> h(n, 0);
        for (int i = 0; i <= n; i++)</pre>
                for (int v = 0; v < n; v++)
                        for (auto [u, w] : g[v]) if (h[u] > h[v] + w) {
                                 if (i == n) return 1;
                                 h[u] = h[v] + w;
        for (int i = 0; i < n; i++) {</pre>
                for (int j = 0; j < n; j++) d[i][j] = LINF;</pre>
                d[i][i] = 0;
                priority_queue<pair<11, int>> pq;
                pq.emplace(0, i);
                while (pq.size()) {
                        auto [ndist, v] = pq.top(); pq.pop();
                        if (-ndist > d[i][v]) continue;
                         for (auto [u, w] : g[v]) {
                                 w += h[v] - h[u];
                                 if (d[i][u] > d[i][v] + w) {
                                         d[i][u] = d[i][v] + w;
                                         pq.emplace(-d[i][u], u);
                for (int j = 0; j < n; j++)
                        d[i][j] += h[j] - h[i];
        return 0;
```

### 25.20 kosaraju

```
// Kosaraju
// O(n + m)
int n;
vector<int> g[MAX];
vector<int> gi[MAX]; // grafo invertido
int vis[MAX];
stack<int> S;
int comp[MAX]; // componente conexo de cada vertice
void dfs(int k) {
        vis[k] = 1;
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
                 if (!vis[g[k][i]]) dfs(g[k][i]);
        S.push(k);
void scc(int k, int c) {
        vis[k] = 1;
comp[k] = c;
        for (int i = 0; i < (int) gi[k].size(); i++)</pre>
                 if (!vis[gi[k][i]]) scc(gi[k][i], c);
void kosaraju() {
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
        while (S.size()) {
                 int u = S.top();
                 S.pop();
                 if (!vis[u]) scc(u, u);
}
```

#### 25.21 kruskal

```
// Kruskal
// Gera e retorna uma AGM e seu custo total a partir do vetor de arestas (edg)
// do grafo
// O(m log(m) + m a(m))
// 864875
vector<tuple<int, int, int>> edg; // {peso, [x, y]}
// DSU em O(a(n))
void dsu build();
int find(int a);
void unite(int a, int b);
pair<11, vector<tuple<int, int, int>>> kruskal(int n) {
        dsu build(n);
        sort(edg.begin(), edg.end());
        11 cost = 0:
        vector<tuple<int, int, int>> mst;
        for (auto [w,x,y] : edg) if (find(x) != find(y)) {
                mst.emplace_back(w, x, y);
                cost += w;
                unite(x,y);
        return {cost, mst};
```

### 25.22 kuhn

```
// Kuhn
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// Para recuperar o matching, basta olhar 'ma' e 'mb'
// 'recover' recupera o min vertex cover como um par de
// {caras da particao A, caras da particao B}
// O(|V| * |E|)
// Na pratica, parece rodar tao rapido quanto o Dinitz
mt19937 rng((int) chrono::steady_clock::now().time_since_epoch().count());
struct kuhn {
        int n, m;
        vector<vector<int>> q;
        vector<int> vis, ma, mb;
        kuhn(int n_, int m_) : n(n_), m(m_), g(n),
                vis (n+m), ma (n, -1), mb (m, -1) {}
        void add(int a, int b) { g[a].push_back(b); }
        bool dfs(int i) {
                 vis[i] = 1;
                 for (int j : g[i]) if (!vis[n+j]) {
                         vis[n+j] = 1;
                         if (mb[j] == -1 or dfs(mb[j])) {
                                  ma[i] = j, mb[j] = i;
                                  return true;
                return false:
        int matching() {
                 int ret = 0, aum = 1;
                 for (auto& i : q) shuffle(i.begin(), i.end(), rng);
                 while (aum)
                         for (int j = 0; j < m; j++) vis[n+j] = 0;
                         aum = 0:
                         for (int i = 0; i < n; i++)</pre>
                                  if (ma[i] == -1 and dfs(i)) ret++, aum = 1;
                 return ret;
pair<vector<int>, vector<int>> recover(kuhn& K) {
        K.matching();
        int n = K.n, m = K.m;
        for (int i = 0; i < n+m; i++) K.vis[i] = 0;
for (int i = 0; i < n; i++) if (K.ma[i] == -1) K.dfs(i);</pre>
        vector<int> ca, cb;
        for (int i = 0; i < n; i++) if (!K.vis[i]) ca.push_back(i);</pre>
        for (int i = 0; i < m; i++) if (K.vis[n+i]) cb.push_back(i);</pre>
        return {ca, cb};
```

#### 25.23 linetree

```
// Line Tree
//
// Reduz min-query em arvore para RMQ
// Se o grafo nao for uma arvore, as queries
// sao sobre a arvore geradora maxima
// Queries de minimo
//
// build - O(n log(n))
// query - O(log(n))
int n;
namespace linetree {
```

```
int id[MAX], seg[2*MAX], pos[MAX];
vector<int> v[MAX], val[MAX];
vector<pair<int, pair<int, int> > ar;
void add(int a, int b, int p) { ar.push_back({p, {a, b}}); }
void build() {
        sort(ar.rbegin(), ar.rend());
        for (int i = 0; i < n; i++) id[i] = i, v[i] = {i}, val[i].clear</pre>
        for (auto i : ar) {
                int a = id[i.second.first], b = id[i.second.second];
                if (a == b) continue;
                if (v[a].size() < v[b].size()) swap(a, b);</pre>
                for (auto j : v[b]) id[j] = a, v[a].push_back(j);
                val[a].push_back(i.first);
                for (auto j : val[b]) val[a].push_back(j);
                v[b].clear(), val[b].clear();
        vector<int> vv:
        for (int i = 0; i < n; i++) for (int j = 0; j < v[i].size(); j
            ++) {
                pos[v[i][j]] = vv.size();
                if (j + 1 < v[i].size()) vv.push_back(val[i][j]);</pre>
                else vv.push_back(0);
        for (int i = n; i < 2*n; i++) seg[i] = vv[i-n];
        for (int i = n-1; i; i--) seg[i] = min(seg[2*i], seg[2*i+1]);
int query(int a, int b) {
        if (id[a] != id[b]) return 0; // nao estao conectados
        a = pos[a], b = pos[b];
        if (a > b) swap(a, b);
       h--:
        int ans = INF;
        for (a += n, b += n; a <= b; ++a/=2, --b/=2) ans = min(\{ans, seq\})
             [a], seg[b]});
        return ans;
```

### 25.24 lowerBoundMaxFlow

};

```
// Max flow com lower bound
// add(a, b, l, r):
       adiciona aresta de a pra b, onde precisa passar f de fluxo, l <= f <= r
// add(a, b, c):
       adiciona aresta de a pra b com capacidade c
// Mesma complexidade do Dinitz
struct lb_max_flow : dinitz {
        vector<int> d;
        lb_{max_flow(int n)} : dinitz(n + 2), d(n, 0) {}
        void add(int a, int b, int 1, int r) {
               d[a] -= 1;
                d[b] += 1;
                dinitz::add(a, b, r - 1);
        void add(int a, int b, int c) {
                dinitz::add(a, b, c);
        bool has_circulation() {
               int n = d.size();
                11 cost = 0;
                for (int i = 0; i < n; i++) {
                        if (d[i] > 0) {
                                cost += d[i];
                                dinitz::add(n, i, d[i]);
                        } else if (d[i] < 0) {</pre>
                                dinitz::add(i, n+1, -d[i]);
                return (dinitz::max_flow(n, n+1) == cost);
```

#### 25.25 minCostMaxFlow

```
// MinCostMaxFlow
// min_cost_flow(s, t, f) computa o par (fluxo, custo)
// com max(fluxo) <= f que tenha min(custo)
// min_cost_flow(s, t) -> Fluxo maximo de custo minimo de s pra t
// Se for um dag, da pra substituir o SPFA por uma DP pra nao
// pagar O(nm) no comeco
// Se nao tiver aresta com custo negativo, nao precisa do SPFA
// O(nm + f * m log n)
template<typename T> struct mcmf {
        struct edge {
                 int to, rev, flow, cap; // para, id da reversa, fluxo,
                      capacidade
                 bool res; // se eh reversa
                 T cost; // custo da unidade de fluxo
                 edge() : to(0), rev(0), flow(0), cap(0), cost(0), res(false) {}
                 edge(int to_, int rev_, int flow_, int cap_, T cost_, bool res_)
                         : to(to_), rev(rev_), flow(flow_), cap(cap_), res(res_),
                                cost(cost ) {}
        };
        vector<vector<edge>> q;
        vector<int> par_idx, par;
        T inf;
        vector<T> dist:
        mcmf(int n) : g(n), par_idx(n), par(n), inf(numeric_limits<T>::max()/3)
        void add(int u, int v, int w, T cost) { // de u pra v com cap w e custo
                 edge a = edge(v, g[v].size(), 0, w, cost, false);
                 edge b = edge(u, g[u].size(), 0, 0, -cost, true);
                 g[u].push_back(a);
                 g[v].push_back(b);
        vector<T> spfa(int s) { // nao precisa se nao tiver custo negativo
                 deque<int> q;
                 vector<bool> is_inside(q.size(), 0);
                 dist = vector<T>(g.size(), inf);
                 dist[s] = 0;
                 q.push_back(s);
                 is_inside[s] = true;
                 while (!q.empty()) {
                         int v = q.front();
                         q.pop_front();
                          is_inside[v] = false;
                          for (int i = 0; i < g[v].size(); i++) {</pre>
                                  auto [to, rev, flow, cap, res, cost] = g[v][i];
if (flow < cap and dist[v] + cost < dist[to]) {</pre>
                                           dist[to] = dist[v] + cost;
                                           if (is_inside[to]) continue;
                                           if (!q.empty() and dist[to] > dist[q.
                                                front()]) q.push_back(to);
```

```
else q.push_front(to);
                                  is_inside[to] = true;
        return dist;
bool dijkstra(int s, int t, vector<T>& pot) {
        priority_queue<pair<T, int>, vector<pair<T, int>>, greater<>> q;
        dist = vector<T>(g.size(), inf);
        dist[s] = 0;
        q.emplace(0, s);
        while (q.size()) {
                 auto [d, v] = q.top();
                 a.pop();
                 if (dist[v] < d) continue;</pre>
                 for (int i = 0; i < g[v].size(); i++) {</pre>
                         auto [to, rev, flow, cap, res, cost] = g[v][i];
cost += pot[v] - pot[to];
                         if (flow < cap and dist[v] + cost < dist[to]) {</pre>
                                  dist[to] = dist[v] + cost;
                                  q.emplace(dist[to], to);
                                  par_idx[to] = i, par[to] = v;
        return dist[t] < inf;</pre>
pair<int, T> min_cost_flow(int s, int t, int flow = INF) {
        vector<T> pot(g.size(), 0);
        pot = spfa(s); // mudar algoritmo de caminho minimo aqui
        int f = 0;
        T ret = 0;
        while (f < flow and dijkstra(s, t, pot)) {
                 for (int i = 0; i < g.size(); i++)</pre>
                         if (dist[i] < inf) pot[i] += dist[i];</pre>
                 int mn_flow = flow - f, u = t;
                 while (u != s) {
                         mn_flow = min(mn_flow,
    g[par[u]][par_idx[u]].cap - g[par[u]][
                                       par_idx[u]].flow);
                         u = par[u];
                 ret += pot[t] * mn_flow;
                 u = t;
                 while (u != s) {
                         g[par[u]][par_idx[u]].flow += mn_flow;
                         g[u][g[par[u]][par_idx[u]].rev].flow -= mn_flow;
                         u = par[u];
                 f += mn_flow;
        return make_pair(f, ret);
// Opcional: retorna as arestas originais por onde passa flow = cap
vector<pair<int,int>> recover() {
        vector<pair<int,int>> used;
        for (int i = 0; i < g.size(); i++) for (edge e : g[i])</pre>
                 if(e.flow == e.cap && !e.res) used.push_back({i, e.to});
        return used;
```

# 25.26 prufer

};

```
// Prufer code
```

```
// Traduz de lista de arestas para prufer code
// e vice-versa
// Os vertices tem label de 0 a n-1
// Todo array com n-2 posicoes e valores de
// 0 a n-1 sao prufer codes validos
// O(n)
vector<int> to_prufer(vector<pair<int, int>> tree) {
        int n = tree.size()+1;
        vector<int> d(n, 0);
        vector<vector<int>> g(n);
        for (auto [a, b] : tree) d[a]++, d[b]++,
                g[a].push_back(b), g[b].push_back(a);
        vector<int> pai(n, -1);
        queue<int> q; q.push(n-1);
        while (q.size()) {
                int u = q.front(); q.pop();
for (int v : g[u]) if (v != pai[u])
                        pai[v] = u, q.push(v);
        int idx, x;
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
        vector<int> ret;
        for (int i = 0; i < n-2; i++) {
                int y = pai[x];
                ret.push_back(y);
                if (-d[y] == 1 and y < idx) x = y;
                else idx = x = find(d.begin()+idx+1, d.end(), 1) - d.begin();
        return ret;
vector<pair<int, int>> from_prufer(vector<int> p) {
        int n = p.size()+2;
        vector<int> d(n, 1);
        for (int i : p) d[i]++;
        p.push_back(n-1);
        int idx, x;
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
        vector<pair<int, int>> ret;
        for (int y : p) {
                ret.push_back({x, y});
                if (--d[y] == 1 \text{ and } y < idx) x = y;
                else idx = x = find(d.begin()+idx+1, d.end(), 1) - d.begin();
        return ret;
```

## 25.27 sack

```
// Sack (DSU em arvores)
// Responde queries de todas as sub-arvores
// offline
// O(n log(n))
int sz[MAX], cor[MAX], cnt[MAX];
vector<int> g[MAX];
void build(int k, int d=0) {
        sz[k] = 1;
        for (auto& i : g[k]) {
                build(i, d+1); sz[k] += sz[i];
                if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
}
void compute(int k, int x, bool dont=1) {
        cnt[cor[k]] += x;
        for (int i = dont; i < g[k].size(); i++)</pre>
                compute(g[k][i], x, 0);
void solve(int k, bool keep=0) {
```

### 25.28 stableMarriage

```
// Stable Marriage
// Emparelha todos os elementos de A com elementos de B
// de forma que nao exista um par x \in A, y \in B
// e x nao pareado com y tal que x prefira parear com y
// e y prefira parear com x.
// a[i] contem os elementos de B ordenados por preferencia de i
// b[j] contem os elementos de A ordenados por preferencia de j
// |A| <= |B|
// Retorna um vetor v de tamanho |A| onde v[i] guarda o match de i.
// O(|A| * |B|)
vector<int> stable_marriage(vector<vector<int>> &a, vector<vector<int>> &b) {
        int n = a.size(), m = b.size();
        assert(a[0].size() == m and b[0].size() == n and n \le m;
        vector<int> match(m, -1), it(n, 0);
        vector inv_b(m, vector<int>(n));
        for (int i = 0; i < m; i++) for (int j = 0; j < n; j++)
                inv_b[i][b[i][j]] = j;
        queue<int> q;
        for (int i = 0; i < n; i++) q.push(i);
        while (q.size()) {
                int i = q.front(); q.pop();
                int j = a[i][it[i]];
                if (match[j] == -1) match[j] = i;
                else if (inv_b[j][i] < inv_b[j][match[j]]) {
                        q.emplace(match[j]);
                        it[match[j]]++;
                        match[j] = i;
                } else q.emplace(i), it[i]++;
        vector<int> ret(n);
        for (int i = 0; i < m; i++) if (match[i] != -1) ret[match[i]] = i;</pre>
        return ret:
```

# **25.29** tarjan

```
// Tarjan para SCC
//
// O(n + m)

vector<int> g[MAX];
stack<int> s;
int vis[MAX], comp[MAX];
int id[MAX];

// se quiser comprimir ciclo ou achar ponte em grafo nao direcionado,
// colocar um if na dfs para nao voltar pro pai da DFS tree
int dfs(int i, int& t) {
    int lo = id[i] = t++;
        s.push(i);
        vis[i] = 2;
```

```
for (int j : g[i]) {
            if (!vis[j]) lo = min(lo, dfs(j, t));
            else if (vis[j] == 2) lo = min(lo, id[j]);
}

// aresta de i pro pai eh uma ponte (no caso nao direcionado)
if (lo == id[i]) while (1) {
            int u = s.top(); s.pop();
                vis[u] = 1, comp[u] = i;
                 if (u == i) break;
}

return lo;
}

void tarjan(int n) {
        int t = 0;
        for (int i = 0; i < n; i++) vis[i] = 0;

for (int i = 0; i < n; i++) if (!vis[i]) dfs(i, t);
}</pre>
```

# 25.30 topoSort

```
// Topological Sort
// Retorna uma ordenacaoo topologica de g
// Se g nao for DAG retorna um vetor vazio
// O(n + m)
vector<int> g[MAX];
vector<int> topo_sort(int n) {
        vector<int> ret(n,-1), vis(n,0);
        int pos = n-1, dag = 1;
        function<void(int) > dfs = [&](int v) {
                vis[v] = 1;
                for (auto u : g[v]) {
                        if (vis[u] == 1) dag = 0;
                        else if (!vis[u]) dfs(u);
                ret[pos--] = v, vis[v] = 2;
        };
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
        if (!dag) ret.clear();
        return ret;
```

### 25.31 treeIsomorf

```
// Isomorfismo de arvores
//
// thash() retorna o hash da arvore (usando centroids como vertices especiais).
// Duas arvores sao isomorfas sse seu hash eh o mesmo
//
// O(|V|.log(|V|))
map<vector<int>, int> mphash;
struct tree {
    int n;
    vector<vector<int>> g;
    vector<int> sz, cs;
    tree(int n_) : n(n_), g(n_), sz(n_) {}
    void dfs_centroid(int v, int p) {
```

```
sz[v] = 1;
                bool cent = true;
                for (int u : g[v]) if (u != p) {
                         dfs_centroid(u, v), sz[v] += sz[u];
                         if(sz[u] > n/2) cent = false;
                if (cent and n - sz[v] \le n/2) cs.push_back(v);
        int fhash(int v, int p) {
                vector<int> h;
                for (int u : g[v]) if (u != p) h.push_back(fhash(u, v));
                sort(h.begin(), h.end());
                if (!mphash.count(h)) mphash[h] = mphash.size();
                return mphash[h];
        il thash() {
                cs.clear();
                dfs_centroid(0, -1);
if (cs.size() == 1) return fhash(cs[0], -1);
                11 h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1], cs[0]);
                return (min(h1, h2) << 30) + max(h1, h2);</pre>
};
// Versao mais rapida com hash, ideal para hash de floresta.
// subtree_hash(v, p) retorna o hash da subarvore enraizada em v com pai p.
// tree hash() retorna o hash da arvore.
// forest_hash() retorna o hash da floresta.
// use o vetor forb[] para marcar vertices que nao podem ser visitados.
// O(|V|.log(|V|))
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
int uniform(ll 1, ll r) {
        uniform_int_distribution<ll> uid(l, r);
        return uid(rng);
const int MOD = 1e9 + 7;
const int H = 13;
const int P = uniform(1, MOD-1);
const int P2 = uniform(1, MOD-1);
struct tree {
        int fn;
        vector<vector<int>> g;
        vector<int> sz, cs;
        vector<bool> forb;
        tree(int n_) : fn(n_), g(n_), sz(n_), forb(n_) {}
        void dfs_size(int v, int p) {
                sz[v] = 1;
                for (int u : g[v]) if (u != p and !forb[u]) {
                        dfs_size(u, v), sz[v] += sz[u];
        void dfs_centroid(int v, int p, int n) {
                bool cent = true;
                for (int u : g[v]) if (u != p and !forb[u]) {
                         dfs_centroid(u, v, n);
                        if(sz[u] > n/2) cent = false;
                if (cent and n - sz[v] <= n/2) cs.push_back(v);</pre>
        int subtree_hash(int v, int p) {
                int h = H;
                for (int u : g[v]) if (u != p and !forb[u]) {
                        h = 11(h) * (P + subtree_hash(u, v)) % MOD;
                return h;
        int tree_hash(int v=0) {
                cs.clear();
                dfs_size(v, -1);
                dfs_centroid(v, -1, sz[v]);
```

```
if (cs.size() == 1) return subtree_hash(cs[0], -1);
    assert (cs.size() == 2);
    int h1 = subtree_hash(cs[0], cs[1]);
    int h2 = subtree_hash(cs[1], cs[0]);
    return l1(P + h1) * (P + h2) * MOD;
}
int forest_hash() {
    fill(sz.begin(), sz.end(), 0);
    int hash = 1;
    for (int v = 0; v < fn; v++) if (!sz[v] and !forb[v]) {
        hash = hash * l1(P2 + tree_hash(v)) * MOD;
    }
    return hash;
};</pre>
```

#### 25.32 virtualTree

```
// Virtual Tree
// Comprime uma arvore dado um conjunto S de vertices, de forma que
// o conjunto de vertices da arvore comprimida contenha S e seja
// minimal e fechado sobre a operacao de LCA
// Se |S| = k, a arvore comprimida tem menos que 2k vertices
// As arestas de virt possuem a distancia do vertice ate o vizinho
// Retorna a raiz da virtual tree
// lca::pos deve ser a ordem de visitacao no dfs
// voce pode usar o LCAcomHLD, por exemplo
// O(k \log(k))
vector<pair<int, int>> virt[MAX];
#warning lembrar de buildar o LCA antes
int build_virt(vector<int> v) +
        auto cmp = [&](int i, int j) { return lca::pos[i] < lca::pos[j]; };</pre>
        sort(v.begin(), v.end(), cmp);
        for (int i = v.size()-1; i; i--) v.push_back(lca::lca(v[i], v[i-1]));
        sort(v.begin(), v.end(), cmp);
        v.erase(unique(v.begin(), v.end()), v.end());
        for (int i = 0; i < v.size(); i++) virt[v[i]].clear();</pre>
        for (int i = 1; i < v.size(); i++) virt[lca::lca(v[i-1], v[i])].clear(); for (int i = 1; i < v.size(); i++) {
                int parent = lca::lca(v[i-1], v[i]);
                int d = lca::dist(parent, v[i]);
#warning soh to colocando aresta descendo
                virt[parent].emplace_back(v[i], d);
        return v[0];
```

# 26 ufmg forked/Grafos/LCA-HLD

#### 26.1 hldAresta

```
// HLD - aresta
//
// SegTree de soma
// query / update de soma das arestas
//
// Complexidades:
// build - O(n)
// query_path - O(log^2 (n))
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
```

```
namespace hld {
        vector<pair<int, int> > g[MAX];
        int pos[MAX], sz[MAX];
        int sobe[MAX], pai[MAX];
        int h[MAX], v[MAX], t;
        void build_hld(int k, int p = -1, int f = 1) {
                 v[pos[k] = t++] = sobe[k]; sz[k] = 1;
                for (auto& i : g[k]) if (i.first != p) {
                         auto [u, w] = i;
                         sobe[u] = w; pai[u] = k;
                         h[u] = (i == g[k][0] ? h[k] : u);
                         build_hld(u, k, f); sz[k] += sz[u];
                         if (sz[u] > sz[q[k][0].first] or q[k][0].first == p)
                                 swap(i, g[k][0]);
                if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
        void build(int root = 0) {
                t = 0;
                build_hld(root);
                seq::build(t, v);
        il query_path(int a, int b) {
    if (a == b) return 0;
                if (pos[a] < pos[b]) swap(a, b);</pre>
                if (h[a] == h[b]) return seg::query(pos[b]+1, pos[a]);
                return seg::query(pos[h[a]], pos[a]) + query_path(pai[h[a]], b);
        void update_path(int a, int b, int x) {
                if (a == b) return;
                if (pos[a] < pos[b]) swap(a, b);</pre>
                if (h[a] == h[b]) return (void) seq::update(pos[b]+1, pos[a], x);
                seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]], b, x);
        11 query_subtree(int a) {
                if (sz[a] == 1) return 0;
                 return seg::query(pos[a]+1, pos[a]+sz[a]-1);
        void update_subtree(int a, int x) {
                if (sz[a] == 1) return;
                seg::update(pos[a]+1, pos[a]+sz[a]-1, x);
        int lca(int a, int b) {
                if (pos[a] < pos[b]) swap(a, b);</pre>
                return h[a] == h[b] ? b : lca(pai[h[a]], b);
}
```

# 26.2 hldSemUpdate

```
// HLD sem Update
// query de min do caminho
// Complexidades:
// build - O(n)
// guery_path - 0(log(n))
namespace hld {
        vector<pair<int, int> > g[MAX];
        int pos[MAX], sz[MAX];
        int sobe[MAX], pai[MAX];
        int h[MAX], v[MAX], t;
int men[MAX], seg[2*MAX];
        void build_hld(int k, int p = -1, int f = 1) {
                v[pos[k] = t++] = sobe[k]; sz[k] = 1;
                 for (auto& i : g[k]) if (i.first != p) {
                         sobe[i.first] = i.second; pai[i.first] = k;
                         h[i.first] = (i == q[k][0]? h[k] : i.first);
                         men[i.first] = (i == g[k][0] ? min(men[k], i.second) : i
                              .second);
```

### 26.3 hldVertice

};

```
// HLD - vertice
// SegTree de soma
// query / update de soma dos vertices
// Complexidades:
// build - O(n)
// query_path - 0(log^2 (n))
// update_path - 0(log^2 (n))
// query_subtree - O(log(n))
// update subtree - O(log(n))
// namespace seg { ... }
namespace hld {
        vector<int> g[MAX];
        int pos[MAX], sz[MAX];
        int peso[MAX], pai[MAX];
        int h[MAX], v[MAX], t;
        void build_hld(int k, int p = -1, int f = 1) {
                v[pos[k] = t++] = peso[k]; sz[k] = 1;
                for (auto& i : g[k]) if (i != p) {
   pai[i] = k;
                        h[i] = (i == g[k][0] ? h[k] : i);
                        build_hld(i, k, f); sz[k] += sz[i];
                        if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i, g[k]
                             ][0]);
                if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
        void build(int root = 0) {
                t = 0;
                build hld(root);
                seg::build(t, v);
        11 query_path(int a, int b) {
                if (pos[a] < pos[b]) swap(a, b);
                if (h[a] == h[b]) return seg::query(pos[b], pos[a]);
                return seg::query(pos[h[a]], pos[a]) + query_path(pai[h[a]], b);
        void update_path(int a, int b, int x) {
                if (pos[a] < pos[b]) swap(a, b);
                if (h[a] == h[b]) return (void) seq::update(pos[b], pos[a], x);
                seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]], b, x);
```

```
11 query_subtree(int a) {
    return seg::query(pos[a], pos[a]+sz[a]-1);
}
void update_subtree(int a, int x) {
    seg::update(pos[a], pos[a]+sz[a]-1, x);
}
int lca(int a, int b) {
    if (pos[a] < pos[b]) swap(a, b);
    return h[a] == h[b] ? b : lca(pai[h[a]], b);
}</pre>
```

#### 26.4 lca

```
// LCA com RMO
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// dist(a, b) retorna a distancia entre a e b
// Complexidades:
// build - O(n)
// lca - 0(1)
// dist - 0(1)
template<typename T> struct rmq {
        vector<T> v;
        int n; static const int b = 30;
        vector<int> mask, t;
        int op(int x, int y) { return v[x] < v[y] ? x : y; }
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
        rmq(const \ vector<T>\& \ v_) : v(v_), n(v.size()), mask(n), t(n) {
                for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
                        at = (at << 1) & ((1 << b) -1);
                         while (at and op(i, i-msb(at&-at)) == i) at ^= at&-at;
                for (int i = 0; i < n/b; i++) t[i] = b*i+b-1-msb(mask[b*i+b-1]);
                for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0; i+(1<<<math>j) <=
                     n/b; i++)
                        t[n/b*j+i] = op(t[n/b*(j-1)+i], t[n/b*(j-1)+i+(1<<(j-1))
                             ]);
        int small(int r, int sz = b) { return r-msb(mask[r]&((1<<sz)-1)); }
        T query(int 1, int r) {
                if (r-1+1 <= b) return small(r, r-1+1);</pre>
                int ans = op(small(1+b-1), small(r));
                int x = 1/b+1, y = r/b-1;
                if (x <= y) {
                        int j = msb(y-x+1);
                        ans = op(ans, op(t[n/b*j+x], t[n/b*j+y-(1<<j)+1]));
                return ans;
};
namespace lca {
        vector<int> q[MAX];
        int v[2*MAX], pos[MAX], dep[2*MAX];
        int t;
        rmq<int> RMQ;
        void dfs(int i, int d = 0, int p = -1) {
                v[t] = i, pos[i] = t, dep[t++] = d;
                for (int j : g[i]) if (j != p) {
                        dfs(j, d+1, i);
                        v[t] = i, dep[t++] = d;
        void build(int n, int root) {
                t = 0;
                dfs(root);
                RMQ = rmq < int > (vector < int > (dep, dep+2*n-1));
```

# 26.5 lcaComBinaryLifting

```
// LCA com binary lifting
// Assume que um vertice eh ancestral dele mesmo, ou seja, // se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
// Complexidades:
// build - O(n log(n))
// 1 ca - O(log(n))
vector<vector<int> > q(MAX);
int n, p;
int pai[MAX2][MAX];
int in[MAX], out[MAX];
void dfs(int k) {
        in[k] = p++;
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
                if (in[g[k][i]] == -1) {
                         pai[0][g[k][i]] = k;
                         dfs(g[k][i]);
        out[k] = p++;
void build(int raiz) {
        for (int i = 0; i < n; i++) pai[0][i] = i;</pre>
        p = 0, memset(in, -1, sizeof in);
        dfs(raiz);
        // pd dos pais
        for (int k = 1; k < MAX2; k++) for (int i = 0; i < n; i++)</pre>
                pai[k][i] = pai[k - 1][pai[k - 1][i]];
bool anc(int a, int b) { // se a eh ancestral de b
        return in[a] <= in[b] and out[a] >= out[b];
int lca(int a, int b) {
        if (anc(a, b)) return a;
        if (anc(b, a)) return b;
        // sobe a
        for (int k = MAX2 - 1; k >= 0; k--)
                if (!anc(pai[k][a], b)) a = pai[k][a];
        return pai[0][a];
// Alternativamente:
// 'binary lifting' gastando O(n) de memoria
// Da pra add folhas e fazer queries online
// 3 vezes o tempo do binary lifting normal
// build - 0(n)
// kth, lca, dist - O(log(n))
int d[MAX], p[MAX], pp[MAX];
void set_root(int i) { p[i] = pp[i] = i, d[i] = 0; }
void add_leaf(int i, int u)
        p[i] = u, d[i] = d[u]+1;
```

```
pp[i] = 2*d[pp[u]] == d[pp[pp[u]]]+d[u] ? pp[pp[u]] : u;
int kth(int i, int k) {
        int dd = max(0, d[i]-k);
        while (d[i] > dd) i = d[pp[i]] >= dd ? pp[i] : p[i];
        return i;
int lca(int a, int b)
        if (d[a] < d[b]) swap(a, b);</pre>
        while (d[a] > d[b]) a = d[pp[a]] >= d[b] ? pp[a] : p[a];
        while (a != b) {
                if (pp[a] != pp[b]) a = pp[a], b = pp[b];
                else a = p[a], b = p[b];
        return a;
int dist(int a, int b) { return d[a]+d[b]-2*d[lca(a,b)]; }
vector<int> g[MAX];
void build(int i, int pai=-1) {
        if (pai == -1) set_root(i);
        for (int j : g[i]) if (j != pai) {
                add_leaf(j, i);
                build(j, i);
```

#### 26.6 lcaComHld

```
// LCA com HLD
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// Para buildar pasta chamar build(root)
// anc(a, b) responde se 'a' eh ancestral de 'b'
// Complexidades:
// build - O(n)
// 1 ca - O(log(n))
// anc - O(1)
vector<int> g[MAX];
int pos[MAX], h[MAX], sz[MAX];
int pai[MAX], t;
void build(int k, int p = -1, int f = 1) {
        pos[k] = t++; sz[k] = 1;
        for (int& i : g[k]) if (i != p) {
                pai[i] = k;
                h[i] = (i == g[k][0] ? h[k] : i);
                build(i, k, f); sz[k] += sz[i];
                if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i, g[k][0]);
        if (p*f == -1) t = 0, h[k] = k, build (k, -1, 0);
int lca(int a, int b) {
        if (pos[a] < pos[b]) swap(a, b);</pre>
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
bool anc(int a, int b) {
        return pos[a] <= pos[b] and pos[b] <= pos[a]+sz[a]-1;</pre>
```

# 27 ufmg forked/Grafos/LCT

#### 27.1 lct

```
// Link-cut Tree
// Link-cut tree padrao
// Todas as operacoes sao O(log(n)) amortizado
namespace lct {
        struct node {
                int p, ch[2];
                node() \{ p = ch[0] = ch[1] = -1; \}
        };
        node t[MAX];
        bool is_root(int x) {
                return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].ch[1]
                       ! = x);
        void rotate(int x)
                int p = t[x].p, pp = t[p].p;
                if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
                bool d = t[p].ch[0] == x;
                t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
                if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
                t[x].p = pp, t[p].p = x;
        void splay(int x) {
                while (!is_root(x)) {
                         int p = t[x].p, pp = t[p].p;
if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
                              == x) ? x : p);
                         rotate(x);
        int access(int v) {
                int last = -1;
                for (int w = v; w+1; last = w, splay(v), w = t[v].p)
                         splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
                 return last;
        int find_root(int v)
                access(v);
                while (t[v].ch[0]+1) v = t[v].ch[0];
                return splay(v), v;
        void link(int v, int w) { // v deve ser raiz
                access(v);
                t[v].p = w;
        void cut(int v) { // remove aresta de v pro pai
                access(v);
                t[v].ch[0] = t[t[v].ch[0]].p = -1;
        int lca(int v, int w) {
                return access(v), access(w);
```

#### 27.2 lctAresta

```
// Link-cut Tree - aresta
//
// Valores nas arestas
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nas arestas do caminho v--w
//
// Todas as operacoes sao O(log(n)) amortizado
```

```
namespace lct {
        struct node {
                  int p, ch[2];
                  11 val, sub;
                 bool rev;
                  int sz, ar;
                 11 lazy;
                 node() {}
                 node(int v, int ar_) :
                 p(-1), val(v), sub(v), rev(0), sz(ar_), ar(ar_), lazy(0) {
                          ch[0] = ch[1] = -1;
         };
         node t[2*MAX]; // MAXN + MAXQ
         map<pair<int, int>, int> aresta;
         int sz;
         void prop(int x) {
                 if (t[x].lazy) {
                           if (\bar{t}[x].ar) t[x].val += t[x].lazy;
                           t[x].sub += t[x].lazy*t[x].sz;
                           if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
                           if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
                  if (t[x].rev) {
                          if (t[x].ch[0], t[x].ch[1]);
if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
                 t[x].lazy = 0, t[x].rev = 0;
         void update(int x) {
                  t[x].sz = t[x].ar, t[x].sub = t[x].val;
                  for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
    prop(t[x].ch[i]);</pre>
                           t[x].sz += t[t[x].ch[i]].sz;
                           t[x].sub += t[t[x].ch[i]].sub;
         bool is_root(int x) {
                  return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].ch[1]
                        ! = x);
         void rotate(int x) {
                 int p = t[x].p, pp = t[p].p;
if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
                  bool d = t[p].ch[0] == x;
                 t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
                 t[x].p = pp, t[p].p = x;
                 update(p), update(x);
         int splay(int x) {
                  while (!is_root(x)) {
                           int p = t[x].p, pp = t[p].p;
                           if (!is_root(p)) prop(pp);
                           prop(p), prop(x);
                           if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
                                == x) ? x : p);
                           rotate(x);
                 return prop(x), x;
         int access(int v) {
                  int last = -1:
                  for (int w = v; w+1; update(last = w), splay(v), w = t[v].p)
                          splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
                  return last;
         void make_tree(int v, int w=0, int ar=0) { t[v] = node(w, ar); }
         int find_root(int v) {
                 access(v), prop(v);
while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
                 return splay(v);
         bool conn(int v, int w) {
                 access(v), access(w);
                  return v == w ? true : t[v].p != -1;
```

```
void rootify(int v) {
        access(v);
        t[v].rev ^= 1;
11 query(int v, int w) {
         rootify(w), access(v);
         return t[v].sub;
void update(int v, int w, int x) {
         rootify(w), access(v);
        t[v].lazy += x;
void link_(int v, int w) {
        rootify(w);
        t[w].p = v;
void link(int v, int w, int x) { // v--w com peso x
    int id = MAX + sz++;
        aresta[make_pair(v, w)] = id;
        make_tree(id, x, 1);
        link_(v, id), link_(id, w);
void cut_(int v, int w) {
        rootify(w), access(v);
t[v].ch[0] = t[t[v].ch[0]].p = -1;
void cut(int v, int w) {
        int id = aresta[make_pair(v, w)];
        cut_(v, id), cut_(id, w);
int lca(int v, int w) {
        access(v);
         return access(w);
```

#### 27.3 lctVertice

```
// Link-cut Tree - vertice
// Valores nos vertices
// make_tree(v, w) cria uma nova arvore com um
// vertice soh com valor 'w'
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nos vertices do caminho v--w
// Todas as operacoes sao O(log(n)) amortizado
namespace lct {
        struct node {
                 int p, ch[2];
                 ll val, sub;
                 bool rev;
                 int sz;
                 11 lazy;
                 node() {}
                 node(int \ v) : p(-1), val(v), sub(v), rev(0), sz(1), lazy(0) 
                         ch[0] = ch[1] = -1;
        };
        node t[MAX];
        void prop(int x) {
                 if (t[x].lazy) {
                          t[x].val += t[x].lazy, t[x].sub += t[x].lazy*t[x].sz;
                         if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
                 if (t[x].rev) {
                          swap(t[x].ch[0], t[x].ch[1]);
                         if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
                         if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
```

```
t[x].lazy = 0, t[x].rev = 0;
void update(int x) {
        t[x].sz = 1, t[x].sub = t[x].val;
for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {</pre>
                prop(t[x].ch[i]);
                t[x].sz += t[t[x].ch[i]].sz;
                t[x].sub += t[t[x].ch[i]].sub;
bool is_root(int x) {
        return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].ch[1]
void rotate(int x) {
        int p = t[x].p, pp = t[p].p;
        if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
        bool d = t[p].ch[0] == x;
t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
        if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
        t[x].p = pp, t[p].p = x;
        update(p), update(x);
int splay(int x) {
        while (!is_root(x)) {
                int p = t[x].p, pp = t[p].p;
                if (!is_root(p)) prop(pp);
                prop(p), prop(x);
                 if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
                    == x) ? x : p);
                rotate(x);
        return prop(x), x;
int access(int v) {
        int last = -1;
        for (int w = v; w+1; update(last = w), splay(v), w = t[v].p)
                 splay(w), t[w] \cdot ch[1] = (last == -1 ? -1 : v);
void make_tree(int v, int w) { t[v] = node(w); }
int find root(int v) {
        access(v), prop(v);
        while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
        return splay(v);
bool connected(int v, int w) {
        access(v), access(w);
        return v == w ? true : t[v].p != -1;
void rootify(int v) {
        access(v);
t[v].rev ^= 1;
11 query(int v, int w) {
        rootify(w), access(v);
        return t[v].sub;
void update(int v, int w, int x) {
        rootify(w), access(v);
        t[v].lazy += x;
void link(int v, int w) {
        rootify(w);
        t[w].p = v;
void cut(int v, int w) {
        rootify(w), access(v);
        t[v].ch[0] = t[t[v].ch[0]].p = -1;
int lca(int v, int w) {
        access(v);
        return access(w);
```

# 28 ufmg forked/Matematica

#### 28.1 2sat

```
// 2-SAT
// solve() retorna um par, o first fala se eh possivel
// atribuir, o second fala se cada variavel eh verdadeira
// O(|V|+|E|) = O(\#variaveis + \#restricoes)
struct sat {
        int n, tot;
        vector<vector<int>> q;
        vector<int> vis, comp, id, ans;
        stack<int> s;
        sat(int n) : n(n), tot(n), q(2*n) {}
        int dfs(int i, int& t) {
                 int lo = id[i] = t++;
                 s.push(i), vis[i] = 2;
                 for (int j : g[i]) {
                         if (!vis[j]) lo = min(lo, dfs(j, t));
                         else if (vis[j] == 2) lo = min(lo, id[j]);
                 if (lo == id[i]) while (1) {
                         int u = s.top(); s.pop();
vis[u] = 1, comp[u] = i;
                         if ((u>>1) < n and ans[u>>1] == -1) ans[u>>1] = ~u&1;
                         if (u == i) break;
                return lo;
        void add_impl(int x, int y) { // x \rightarrow y = !x ou y
                x = x >= 0 ? 2*x : -2*x-1;
                 y = y >= 0 ? 2*y : -2*y-1;
                g[x].push_back(y);
                g[y^1].push_back(x^1);
        void add_cl(int x, int y) { // x ou y
                add_impl(~x, y);
        void add_xor(int x, int y) { // x xor y
                 add_cl(x, y), add_cl(~x, ~y);
        void add_eq(int x, int y) { // x = y
                add_xor(~x, y);
        void add_true(int x) { // x = T
                add_impl(~x, x);
        void at_most_one(vector<int> v) { // no max um verdadeiro
                 q.resize(2*(tot+v.size()));
                 for (int i = 0; i < v.size(); i++) {</pre>
                         add_impl(tot+i, ~v[i]);
                         if (i) {
                                  add_impl(tot+i, tot+i-1);
                                  add_impl(v[i], tot+i-1);
                 tot += v.size();
        pair<bool, vector<int>> solve() {
                ans = vector < int > (n, -1);
                 int t = 0;
                vis = comp = id = vector<int>(2*tot, 0);
for (int i = 0; i < 2*tot; i++) if (!vis[i]) dfs(i, t);</pre>
                 for (int i = 0; i < tot; i++)
                         if (comp[2*i] == comp[2*i+1]) return {false, {}};
                 return {true, ans};
```

# 28.2 berlekampMassey

```
// Berlekamp-Massey
// guess_kth(s, k) chuta o k-esimo (0-based) termo
// de uma recorrencia linear que gera s
// Para uma rec. lin. de ordem x, se passar 2x termos
// vai gerar a certa
// Usar aritmetica modular
// Pro fast_evaluate, precisa de ntt e divmod (powerSeries.cpp)
// Complexidades: (n = |s|)
// evaluate: 0(n^2 log k)
// fast_evaluate: O(n log n log k)
// berlekampMassey: O(n^2 + O(evaluate))
template<typename T> T evaluate(vector<T> c, vector<T> s, 11 k) {
         int n = c.size();
         assert(c.size() <= s.size());</pre>
         auto mul = [&](const vector<T> &a, const vector<T> &b) {
                 vector<T> ret(a.size() + b.size() - 1);
                 for (int i = 0; i < a.size(); i++) for (int j = 0; j < b.size();</pre>
                          ret[i+j] += a[i] * b[j];
                 for (int i = ret.size()-1; i >= n; i--) for (int j = n-1; j >=
                      0; j--)
                          ret[i-j-1] += ret[i] * c[j];
                 ret.resize(min<int>(ret.size(), n));
                 return ret:
         };
         vector < T > a = n == 1 ? vector < T > ({c[0]}) : vector < T > ({0, 1}), x = {1};
                 if(k&1) x = mul(x, a);
                 a = mul(a, a), k >>= 1;
         x.resize(n);
         T ret = 0:
         for (int i = 0; i < n; i++) ret += x[i] * s[i];
         return ret:
mint fast_evaluate(poly c, poly s, ll k) {
        if (k < s.size()) return s[k];</pre>
        int n = c.size();
        assert(c.size() <= s.size());</pre>
         auto f = poly(n + 1, 1);
         for (int i = 0; i < n; i++) f[i] = -c[n-i-1];
         poly a = n == 1 ? poly({c[0]}) : poly({0, 1}), x = {1};
         while (k) {
                 if (k&1) x = divmod(convolution(x, a), f).second;
                 a = divmod(convolution(a, a), f).second, k >>= 1;
         mint ret = 0;
        for (int i = 0; i < n; i++) ret += x[i] * s[i];
         return ret;
template<typename T> vector<T> berlekamp_massey(vector<T> s) {
         int n = s.size(), 1 = 0, m = 1;
        vector<T> b(n), c(n);
T ld = b[0] = c[0] = 1;
         for (int i = 0; i < n; i++, m++) {
                 T d = s[i];
                 for (int j = 1; j \le 1; j++) d += c[j] * s[i-j];
                 if (d == 0) continue;
                 vector<T> temp = c;
                 T coef = d / ld;
                 for (int j = m; j < n; j++) c[j] -= coef * b[j-m];
if (2 * 1 <= i) 1 = i + 1 - 1, b = temp, ld = d, m = 0;</pre>
```

```
c.resize(1 + 1);
c.erase(c.begin());
for (T& x : c) x = -x;
return c;
}

template<typename T> T guess_kth(const vector<T>& s, ll k) {
    auto c = berlekamp_massey(s);
    return evaluate(c, s, k);
}
```

#### 28.3 chinese

```
// Teorema Chines do Resto
// Combina equacoes modulares lineares: x = a \pmod{m}
// O m final eh o lcm dos m's, e a resposta eh unica mod o lcm
// Os m nao precisam ser coprimos
// Se nao tiver solucao, o 'a' vai ser -1
template<typename T> tuple<T, T, T> ext_gcd(T a, T b) {
    if (!a) return {b, 0, 1};
    auto [g, x, y] = ext_gcd(b%a, a);
    return {g, y - b/a*x, x};
template<typename T = 11> struct crt {
        crt() : a(0), m(1) {}
        crt(T a_, T m_) : a(a_), m(m_) {}
        crt operator * (crt C) {
                 auto [g, x, y] = ext_gcd(m, C.m);
if ((a - C.a) % g) a = -1;
                 if (a == -1 or C.a == -1) return crt(-1, 0);
                 T lcm = m/q*C.m;
                 T \text{ ans} = a + (x*(C.a-a)/g % (C.m/g))*m;
                 return crt((ans % lcm + lcm) % lcm, lcm);
};
```

#### 28.4 convolution

```
// FFT
// Chamar convolution com vector<complex<double>> para FFT
// Precisa do mint para NTT
// O(n log(n))
// Para FFT
void get_roots(bool f, int n, vector<complex<double>>& roots) {
       const static double PI = acosl(-1);
        for (int i = 0; i < n/2; i++) {
                double alpha = i*((2*PI)/n);
                if (f) alpha = -alpha;
                roots[i] = {cos(alpha), sin(alpha)};
// Para NTT
template<int p>
void get_roots(bool f, int n, vector<mod_int<p>>& roots) {
       mod_int r;
        int ord;
       if (p == 998244353) {
                r = 102292;
                ord = (1 << 23);
        } else if (p == 754974721) {
                r = 739831874;
                ord = (1 << 24);
```

```
} else if (p == 167772161) {
                r = 243:
                ord = (1 << 25);
        } else assert(false);
        if (f) r = r^(p - 1 - ord/n);
        else r = r^{(ord/n)};
        roots[0] = 1;
        for (int i = 1; i < n/2; i++) roots[i] = roots[i-1] *r;
template<typename T> void fft(vector<T>& a, bool f, int N, vector<int>& rev) {
        for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
        int 1, r, m;
        vector<T> roots(N);
        for (int n = 2; n \le N; n \ne 2) {
                get_roots(f, n, roots);
                for (int pos = 0; pos < N; pos += n) {</pre>
                         1 = pos + 0, r = pos + n/2, m = 0;
                         while (m < n/2) {
                                 auto t = roots[m] * a[r];
                                 a[r] = a[1] - t;
                                 a[1] = a[1] + t;
                                 1++, r++, m++;
        if (f) {
                auto invN = T(1) / T(N);
                for (int i = 0; i < N; i++) a[i] = a[i] * invN;</pre>
template<typename T> vector<T> convolution(vector<T>& a, vector<T>& b) {
        vector<T> l(a.begin(), a.end()), r(b.begin(), b.end());
        int N = 1.size()+r.size()-1;
        int n = 1, log_n = 0;
        while (n \le N) n \ne 2, log_n++;
        vector<int> rev(n);
        for (int i = 0; i < n; i++) {</pre>
                rev[i] = 0;
                for (int j = 0; j < log_n; j++) if (i >> j \& 1)
                         rev[i] = 1 << (loq_n-1-j);
        assert(N <= n);
        1.resize(n);
        r.resize(n);
        fft(l, false, n, rev);
        fft(r, false, n, rev);
        for (int i = 0; i < n; i++) l[i] *= r[i];</pre>
        fft(l, true, n, rev);
        1.resize(N);
        return 1;
template<int p, typename T> vector<mod_int<p>> ntt(vector<T>& a, vector<T>& b) {
        vector<mod_int<p>> A(a.begin(), a.end()), B(b.begin(), b.end());
        return convolution (A, B);
// Convolucao de inteiro
// Precisa do CRT
// Tabela de valores:
// [0,1]
              - <int, 1>
// [-1e5, 1e5] - <11, 2>
// [-1e9, 1e9] - <__int128, 3>
template<typename T, int mods>
vector<T> int convolution(vector<int>& a, vector<int>& b) {
        static const int M1 = 998244353, M2 = 754974721, M3 = 167772161;
        auto c1 = ntt<M1>(a, b);

auto c2 = (mods >= 2 ? ntt<M2>(a, b) : vector<mod_int<M2>>());
        auto c3 = (mods >= 3 ? ntt<M3>(a, b) : vector<mod_int<M3>>());
        vector<T> ans;
```

```
for (int i = 0; i < c1.size(); i++) {
    crt<T> at(c1[i].v, M1);
    if (mods >= 2) at = at * crt<T>(c2[i].v, M2);
    if (mods >= 3) at = at * crt<T>(c3[i].v, M3);
    ans.push_back(at.a);
    if (at.a > at.m/2) ans.back() -= at.m;
}
return ans;
```

# 28.5 coprimeBasis

```
// Coprime Basis
// Dado um conjunto de elementos A constroi uma base B
// de fatores coprimos tal que todo elemento A[i]
// pode ser fatorado como A[i] = \prod B[j]^p_ij
// Sendo n o numero de inserts, a complexidade esperada fica
// O(n*(n*loglog(MAX) + log(MAX)^2))
// No pior caso, podemos trocar n*loglog(MAX) por
// 8n, se MAX <= 1e6
// 10n, se MAX <= 1e9
// 16n, se MAX <= 1e18
// 26n, se MAX <= 1e36
template <typename T> struct coprime_basis {
        vector<T> basis;
        coprime basis() {}
        coprime_basis(vector<T> v) { for (T i : v) insert(i); }
        void insert(T z) {
                int n = basis.size();
                basis.push_back(z);
                for (int i = n; i < basis.size(); i++) {</pre>
                         for (int j = (i != n) ? i+1 : 0; j < basis.size(); j++)</pre>
                                 if (i == j) continue;
                                 T \&x = basis[i];
                                 if (x == 1)
                                         j = INF;
                                         continue;
                                 T \& y = basis[j];
                                 T g = gcd(x, y);
                                 if (g == 1) continue;
                                 y /= g, x /= g;
                                 basis.push_back(q);
                basis.erase(remove(basis.beqin(), basis.end(), 1), basis.end());
        vector<int> factor(T x) {
                vector<int> fat(basis.size());
                for (int i = 0; i < basis.size(); i++) {</pre>
                        while (x % basis[i] == 0) x /= basis[i], fat[i]++;
                return fat;
};
```

#### 28.6 crivo

```
// Crivo de Eratosthenes
// "O" crivo
//
// Encontra maior divisor primo
// Um numero eh primo sse divi[x] == x
```

```
// fact fatora um numero <= lim
// A fatoracao sai ordenada
// crivo - O(n log(log(n)))
// fact - O(\log(n))
int divi[MAX];
void crivo(int lim) {
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
         for (int i = 2; i <= lim; i++) if (divi[i] == 1)</pre>
                 for (int j = i; j <= lim; j += i) divi[j] = i;</pre>
#warning A funcao fact ira adicionar o 1 no vetor se voce tentar fatorar
     especificamente o numero 1
void fact(vector<int>& v, int n) {
   if (n != divi[n]) fact(v, n/divi[n]);
         v.push back(divi[n]);
// Crivo linear
// Mesma coisa que o de cima, mas tambem
// calcula a lista de primos
// O(n)
int divi[MAX];
vector<int> primes;
void crivo(int lim) {
        divi[1] = 1;
        for (int i = 2; i <= lim; i++) {
    if (divi[i] == 0) divi[i] = i, primes.push_back(i);</pre>
                 for (int j : primes) {
                          if (j > divi[i] or i*j > lim) break;
                          divi[i*j] = j;
// Crivo de divisores
// Encontra numero de divisores
// ou soma dos divisores
// O(n log(n))
int divi[MAX];
void crivo(int lim) {
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
         for (int i = 2; i <= lim; i++)</pre>
                 for (int j = i; j <= lim; j += i) {</pre>
                          // para numero de divisores
                          divi[j]++;
                          // para soma dos divisores
                          divi[j] += i;
// Crivo de totiente
// Encontra o valor da funcao
// totiente de Euler
// O(n log(log(n)))
int tot[MAX];
void crivo(int lim) {
         for (int i = 1; i <= lim; i++) {</pre>
                 tot[i] += i;
                 for (int j = 2*i; j <= lim; j += i)</pre>
                          tot[j] -= tot[i];
```

```
// Crivo de funcao de mobius
// O(n log(log(n)))
char meb[MAX];
void crivo(int lim) {
         for (int i = 2; i <= lim; i++) meb[i] = 2;</pre>
         meb[1] = 1;
         for (int i = 2; i <= lim; i++) if (meb[i] == 2)</pre>
                 for (int j = i; j <= lim; j += i) if (meb[j]) {
    if (meb[j] == 2) meb[j] = 1;</pre>
                          meb[j] *= j/i%i ? -1 : 0;
// Crivo linear de funcao multiplicativa
// Computa f(i) para todo 1 <= i <= n, sendo f
// uma funcao multiplicativa (se gcd(a,b) = 1,
// entao f(a*b) = f(a)*f(b)
// f_prime tem que computar f de um primo, e
// add_prime tem que computar f(p^n(k+1)) dado f(p^nk) e p // Se quiser computar f(p^nk) dado p e k, usar os comentarios
// O(n)
vector<int> primes;
int f[MAX], pot[MAX];
//int expo[MAX];
void sieve(int lim) {
         // Funcoes para soma dos divisores:
         auto f_prime = [](int p) { return p+1; };
         auto add_prime = [](int fpak, int p) { return fpak*p+1; };
         //auto f_pak = [](int p, int k) {};
         f[1] = 1;
         for (int i = 2; i <= lim; i++) {</pre>
                 if (!pot[i]) {
                          primes.push_back(i);
                          f[i] = f_prime(i), pot[i] = i;
                          //expo[i] = 1;
                  for (int p : primes) {
                          if (i*p > lim) break;
                          if (i%p == 0) {
                                   f[i*p] = f[i / pot[i]] * add_prime(f[pot[i]], p)
                                    // se for descomentar, tirar a linha de cima
                                        tambem
                                    //f[i*p] = f[i / pot[i]] * f_pak(p, expo[i]+1);
                                    //\exp(i*p) = \exp(i)+1;
                                    pot[i*p] = pot[i] * p;
                                    break;
                           } else {
                                    f[i*p] = f[i] * f[p];
                                   pot[i*p] = p;
//expo[i*p] = 1;
         }
```

### 28.7 cycleDetection

```
// Deteccao de ciclo - Tortoise and Hare
//
// Linear no tanto que tem que andar pra ciclar,
// O(1) de memoria
// Retorna um par com o tanto que tem que andar
// do f0 ate o inicio do ciclo e o tam do ciclo
```

```
pair<11, 11> find_cycle() {
        11 \text{ tort} = f(f0);
        ll hare = f(f(f0));
        11 t = 0;
        while (tort != hare) {
                tort = f(tort);
                hare = f(f(hare));
                 t++;
        11 st = 0;
        tort = f0;
        while (tort != hare) {
                tort = f(tort);
                hare = f(hare);
                 st++;
        11 len = 1;
hare = f(tort);
        while (tort != hare) {
                hare = f(hare);
                len++;
        return {st, len};
```

#### 28.8 diofantina

```
// Equacao Diofantina Linear
// Encontra o numero de solucoes de a*x + b*y = c,
// em que x \in [lx, rx] e y \in [ly, ry]
// Usar o comentario para recuperar as solucoes
// (note que o b ao final eh b/gcd(a, b))
// Cuidado com overflow! Tem que caber o quadrado dos valores
// O(log(min(a, b)))
template<typename T> tuple<11, T, T> ext_gcd(11 a, 11 b) {
    if (!a) return {b, 0, 1};
    auto [g, x, y] = ext_gcd<T>(b%a, a);
return {g, y - b/a*x, x};
// numero de solucoes de a*[lx, rx] + b*[ly, ry] = c
template<typename T = 11> // usar __int128 se for ate 1e18
11 diophantine(11 a, 11 b, 11 c, 11 lx, 11 rx, 11 ly, 11 ry) {
         if (lx > rx or ly > ry) return 0;
         if (a == 0 and b == 0) return c ? 0 : (rx-lx+1)*(ry-ly+1);
         auto [g, x, y] = ext_gcd<T>(abs(a), abs(b));
         if (c % g != 0) return 0;
         if (a == 0) return (rx-lx+1)*(ly <= c/b and c/b <= ry);
         if (b == 0) return (ry-ly+1)*(lx <= c/a and c/a <= rx);
         x *= a/abs(a) * c/g, y *= b/abs(b) * c/g, a /= g, b /= g;
         auto shift = [&](T qt) { x += qt*b, y -= qt*a; };
auto test = [&](T& k, ll mi, ll ma, ll coef, int t) {
                 shift((mi - k)*t / coef);
                  if (k < mi) shift(coef > 0 ? t : -t);
                  if (k > ma) return pair<T, T>(rx+2, rx+1);
                  T x1 = x;
                  shift((ma - k) *t / coef);
                  if (k > ma) shift(coef > 0 ? -t : t);
                  return pair<T, T>(x1, x);
         };
         auto [11, r1] = test(x, lx, rx, b, 1);
         auto [12, r2] = test(y, ly, ry, a, -1);
         if (12 > r2) swap(12, r2);
         T 1 = max(11, 12), r = min(r1, r2);
         if (1 > r) return 0;
         11 k = (r-1) / abs(b) + 1;
         return k; // solucoes: x = 1 + [0, k)*|b|
```

#### 28.9 divisionTrick

```
// Division Trick
//
// Gera o conjunto n/i, pra todo i, em O(sqrt(n))
// copiei do github do tfg50

for(int l = 1, r; l <= n; l = r + 1) {
    r = n / (n / 1);
    // n / i has the same value for l <= i <= r
}</pre>
```

## 28.10 evalInterpol

```
// Avaliacao de Interpolacao
// Dado 'n' pontos (i, y[i]), i \in [0, n),
// avalia o polinomio de grau n-1 que passa
// por esses pontos em 'x
// Tudo modular, precisa do mint
// O(n)
mint evaluate_interpolation(int x, vector<mint> y) {
        int n = y.size();
        vector<mint> sulf(n+1, 1), fat(n, 1), ifat(n);
        for (int i = n-1; i >= 0; i--) sulf[i] = sulf[i+1] * (x - i); for (int i = 1; i < n; i++) fat[i] = fat[i-1] * i;
         ifat[n-1] = 1/fat[n-1];
        for (int i = n-2; i >= 0; i--) ifat[i] = ifat[i+1] * (i + 1);
        mint pref = 1, ans = 0;
        for (int i = 0; i < n; pref *= (x - i++)) {
                 mint num = pref * sulf[i+1];
                 mint den = ifat[i] * ifat[n-1 - i];
                 if ((n-1 - i) \% 2) den *= -1;
                 ans += y[i] * num * den;
         return ans;
```

## 28.11 fastPow

28.12 fwht

```
// Fast Walsh Hadamard Transform
// FWHT<'|'>(f) eh SOS DP
// FWHT<'&'>(f) eh soma de superset DP
// Se chamar com ^, usar tamanho potencia de 2!!
// O(n log(n))
template<char op, class T> vector<T> FWHT(vector<T> f, bool inv = false) {
        int n = f.size();
        for (int k = 0; (n-1) >> k; k++) for (int i = 0; i < n; i++) if (i>> k&1) {
                int j = i^(1 << k);
                if (op == '^{\prime})' f[j] += f[i], f[i] = f[j] - 2*f[i];
                if (op == '|') f[i] += (inv ? -1 : 1) * f[j];
                if (op == '&') f[j] += (inv ? -1 : 1) * f[i];
        if (op == '^' and inv) for (auto& i : f) i /= n;
        return f;
// Generalizacao de FWHT de Xor
// Convolucao de soma mod B, usar tamanho potencia de B!!
// Precisa definir o tipo T e a raiz primitiva q
// satisfazendo g^b == g
// Se possivel, hardcodar a multiplicacao de matriz
// feita em cada iteracao faz ficar bem mais rapido
// O(n b log_b(n))
template < class T>
vector<T> FWHT(vector<T> f, int b, T q, bool inv = false) {
       int n = f.size();
        vector<T> w(b);
        w[1] = g;
        for (int i = 2; i < b; i++) w[i] = w[i-1] * g;
        w[0] = w[b - 1] * q;
        if (inv) reverse(w.begin() + 1, w.end());
        for (int pot = 1; pot < n; pot *= b) +</pre>
                for (int i = 0; i < n; i++) if (!(i / pot % b)) {</pre>
                        vector<T> res(b);
                        for (int j = 0; j < b; j++) {
                                for (int k = 0; k < b; k++)
                                        res[j] = res[j] + w[j * k % b] * f[i + k
                                               * pot];
                                if (inv) res[j] = res[j] / b;
                        for (int j = 0; j < b; j++) f[i + j * pot] = res[j];
        return f;
// Exemplos da FWHT Generalizada:
// mod 7, resposta mod 998244353:
// T = mint, g = 14553391
// mod 3, resposta cabe em um long long:
// T = array<11, 2>, g = {0, 1};
// using T = array<11, 2>;
// T operator +(const T& a, const T& b) {
        return T\{a[0] + b[0], a[1] + b[1]\};
// T operator *(const T& a, const T& b) {
       return T\{a[0] * b[0] - a[1] * b[1],
                     a[0] * b[1] + a[1] * b[0] - a[1] * b[1];
// T operator /(const T& a, const int& b) {
        return T{a[0] / b, a[1] / b};
```

### 28.13 gauss

```
// Gauss
// Resolve sistema linear
// Retornar um par com o numero de solucoes
// e alguma solucao, caso exista
// O(n^2 * m)
template<typename T>
pair<int, vector<T>> gauss(vector<T>> a, vector<T> b) {
        const double eps = 1e-6;
        int n = a.size(), m = a[0].size();
        for (int i = 0; i < n; i++) a[i].push_back(b[i]);</pre>
        vector<int> where(m, -1);
        for (int col = 0, row = 0; col < m and row < n; col++) {</pre>
                 int sel = row;
                 for (int i=row; i<n; ++i)</pre>
                         if (abs(a[i][col]) > abs(a[sel][col])) sel = i;
                 if (abs(a[sel][col]) < eps) continue;</pre>
                 for (int i = col; i <= m; i++)
    swap(a[sel][i], a[row][i]);</pre>
                 where[col] = row;
                 for (int i = 0; i < n; i++) if (i != row) {</pre>
                         T c = a[i][col] / a[row][col];
                          for (int j = col; j <= m; j++)</pre>
                                  a[i][j] = a[row][j] * c;
                 row++;
        vector<T> ans(m, 0);
        for (int i = 0; i < m; i++) if (where[i] != -1)</pre>
                ans[i] = a[where[i]][m] / a[where[i]][i];
        for (int i = 0; i < n; i++) {</pre>
                 T sum = 0;
                 for (int j = 0; j < m; j++)
                          sum += ans[j] * a[i][j];
                 if (abs(sum - a[i][m]) > eps)
                          return pair(0, vector<T>());
        for (int i = 0; i < m; i++) if (where [i] == -1)
                 return pair(INF, ans);
        return pair(1, ans);
```

# 28.14 gaussZ2

```
bool add(bitset<D> v) {
        in++;
        bitset<D> k;
        for (int i = D - 1; i >= 0; i--) if (v[i]) {
    if (basis[i][i]) v ^= basis[i], k ^= keep[i];
                          k[i] = true, id[i] = in, keep[i] = k;
                         basis[i] = v, rk++;
                         return true;
        return false;
pair<bool, bitset<D>> coord(bitset<D> v) {
        bitset<D> c;
        for (int i = D - 1; i >= 0; i--) if (v[i]) {
                 if (basis[i][i]) v ^= basis[i], c[i] = true;
                 else return {false, bitset<D>() };
        return {true, c};
pair<bool, vector<int>> recover(bitset<D> v) {
        auto [span, bc] = coord(v);
        if (not span) return {false, {}};
        bitset <D> aux;
        for (int i = D - 1; i >= 0; i--) if (bc[i]) aux ^= keep[i];
        vector<int> oc;
        for (int i = D - 1; i >= 0; i--) if (aux[i]) oc.push_back(id[i])
        return {true, oc};
```

# 28.15 gcdEstendido

};

```
// Euclides estendido
//
Acha x e y tal que ax + by = mdc(a, b) (nao eh unico)
// Assume a, b >= 0
//
// O(log(min(a, b)))

tuple<11, 11, 11> ext_gcd(11 a, 11 b) {
    if (!a) return {b, 0, 1};
    auto [g, x, y] = ext_gcd(b%a, a);
    return {g, y - b/a*x, x};
}
```

# 28.16 gcdLcmConvolution

### 28.17 integral

```
// Integracao Numerica
//
// Metodo de Simpson 3/8
// Integra f no intervalo [a, b], erro cresce proporcional a (b - a)^5

const int N = 3*100; // multiplo de 3
ld integrate(ld a, ld b, function<ld(ld)> f) {
   ld s = 0, h = (b - a)/N;
   for (int i = 1; i < N; i++) s += f(a + i*h)*(i*3 ? 3 : 2);
   return (f(a) + s + f(b))*3*h/8;
}</pre>
```

#### 28.18 karatsuba

```
// Karatsuba
// Os pragmas podem ajudar
// Para n ~ 2e5, roda em < 1 s
// O(n^1.58)
//#pragma GCC optimize("Ofast")
//#pragma GCC target ("avx,avx2")
template<typename T> void kar(T* a, T* b, int n, T* r, T* tmp) {
        if (n <= 64)
                for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)
                        r[i+j] += a[i] * b[j];
        int mid = n/2;
        T *atmp = tmp, *btmp = tmp+mid, *E = tmp+n;
        memset(E, 0, sizeof(E[0])*n);
        for (int i = 0; i < mid; i++) {</pre>
                atmp[i] = a[i] + a[i+mid];
                btmp[i] = b[i] + b[i+mid];
        kar(atmp, btmp, mid, E, tmp+2*n);
        kar(a, b, mid, r, tmp+2*n);
        kar(a+mid, b+mid, mid, r+n, tmp+2*n);
        for (int i = 0; i < mid; i++) {</pre>
                T \text{ temp} = r[i+mid];
                r[i+mid] += E[i] - r[i] - r[i+2*mid];
                r[i+2*mid] += E[i+mid] - temp - r[i+3*mid];
        }
template<typename T> vector<T> karatsuba(vector<T> a, vector<T> b) {
        int n = max(a.size(), b.size());
        while (n&(n-1)) n++;
        a.resize(n), b.resize(n);
        vector<T> ret(2*n), tmp(4*n);
        kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
        return ret;
```

# 28.19 logDiscreto

```
// Logaritmo Discreto
// Resolve logaritmo discreto com o algoritmo baby step giant step
// Encontra o menor x tal que a^x = b (mod m)
// Se nao tem, retorna -1
// O(sqrt(m) * log(sqrt(m))
int dlog(int b, int a, int m) {
        if (a == 0) return b ? -1 : 1; // caso nao definido
        a %= m, b %= m;
        int k = 1, shift = 0;
        while (1)
                int g = gcd(a, m);
                if (g == 1) break;
                if (b == k) return shift;
                if (b % q) return -1;
                b \neq g, m \neq g, shift++;
                k = (11) k * a / g % m;
        int sq = sqrt(m) + 1, giant = 1;
        for (int i = 0; i < sq; i++) giant = (11) giant * a % m;</pre>
        vector<pair<int, int>> baby;
        for (int i = 0, cur = b; i <= sq; i++) {
                baby.emplace_back(cur, i);
                cur = (11) cur * a % m;
        sort(baby.begin(), baby.end());
        for (int j = 1, cur = k; j \le sq; j++) {
                cur = (11) cur * giant % m;
                auto it = lower_bound(baby.begin(), baby.end(), pair(cur, INF));
                if (it != baby.begin() and (--it)->first == cur)
                        return sq * j - it->second + shift;
        return -1;
}
```

#### 28.20 millerRabin

```
// Miller-Rabin
// Testa se n eh primo, n <= 3 * 10^18
// O(log(n)), considerando multiplicacao
// e exponenciacao constantes
11 mul(l1 a, l1 b, l1 m) {
        11 ret = a*b - 11((long double)1/m*a*b+0.5)*m;
        return ret < 0 ? ret+m : ret;</pre>
11 pow(11 x, 11 y, 11 m) {
        if (!v) return 1;
        11 ans = pow(mul(x, x, m), y/2, m);
        return y%2 ? mul(x, ans, m) : ans;
bool prime(ll n) {
        if (n < 2) return 0;
        if (n <= 3) return 1;
        if (n % 2 == 0) return 0;
        ll r = \underline{builtin\_ctzll(n - 1)}, d = n >> r;
        // com esses primos, o teste funciona garantido para n <= 2^64
        // funciona para n <= 3*10^24 com os primos ate 41
```

### 28.21 modInverse

```
// Inverso Modular
//
// Computa o inverso de a modulo b
// Se b eh primo, basta fazer
// a^(b-2)

11 inv(11 a, 11 b) {
    return a > 1 ? b - inv(b%a, a)*b/a : 1;
}

// computa o inverso modular de 1..MAX-1 modulo um primo
11 inv[MAX]:
inv[1] = 1;
for (int i = 2; i < MAX; i++) inv[i] = MOD - MOD/i*inv[MOD%i]%MOD;</pre>
```

#### 28.22 mulmod

# 28.23 multipointEvaluation

```
// Multipoint Evaluation And Interpolation
// Evaluation:
        Avalia o polinomio f(x) nos pontos p[0], p[1], ..., p[n-1]
// Interpolation:
        Retorna o polinomio f(x) de grau n que
        satisfaz f(x) = y pra o conjunto de pontos x, y
// Precisa do ntt e
// - do divmod pro evaluate
// - da derivada pro interpolate
// O divmod e a derivada estao no arquivo powerSeries.cpp
// O(n log^2(n))
namespace multipoint {
        vector<poly> tree;
        void build(vector<mint>& p) {
                int n = p.size();
                tree.resize(2*n);
                for (int i = 0; i < n; i++) tree[n + i] = {-p[i], 1};</pre>
                for (int i = n - 1; i > 0; i--)
                        tree[i] = convolution(tree[2*i], tree[2*i + 1]);
```

```
vector<mint> evaluate(poly& f, vector<mint>& p) {
                build(p);
                int n = p.size();
                vector<poly> ans(2 * n);
                ans[1] = divmod(f, tree[1]).second;
                for (int i = 2; i < 2 * n; i++)
                        ans[i] = divmod(ans[i/2], tree[i]).second;
                vector<mint> results(n);
                for (int i = 0; i < n; i++) results[i] = ans[n + i][0];</pre>
                return results;
        poly prod(vector<mint>& p, int l, int r) {
                if (1 == r) return {-p[1], 1};
                int m = (1 + r) / 2;
                return convolution(prod(p, 1, m), prod(p, m + 1, r));
        poly interpolate(vector<mint>& x, vector<mint>& y) {
                int n = x.size();
                poly p = D(prod(x, 0, n - 1));
                auto d = evaluate(p, x);
                vector<poly> ans(2 * n);
                for (int i = 0; i < n; i++) ans[n + i] = {y[i] / d[i]};
                for (int i = n - 1; i > 0; i--) {
                        poly p1 = convolution(tree[2*i], ans[2*i + 1]);
                        poly p2 = convolution(tree[2*i + 1], ans[2*i]);
                        ans[i] = p1;
                        for (int j = 0; j < p1.size(); j++) ans[i][j] += p2[j];</pre>
                return ans[1];
}
```

#### 28.24 ntt

```
// NTT
// Precisa do mint (primitivas de aritmetica modular)
// O(n log (n))
const int MOD = 998244353;
typedef mod_int<MOD> mint;
void ntt(vector<mint>& a, bool rev) {
        int n = a.size(); auto b = a;
        assert(!(n&(n-1)));
        mint g = 1;
        while ((g^(MOD / 2)) == 1) g += 1;
        if (rev) g = 1 / g;
        for (int step = n / 2; step; step /= 2) {
                 mint w = g^(MOD / (n / step)), wn = 1;
for (int i = 0; i < n/2; i += step) {
                         for (int j = 0; j < step; j++) {
                                  auto u = a[2 * i + j], v = wn * a[2 * i + j +
                                       step];
                                  b[i+j] = u + v; b[i + n/2 + j] = u - v;
                         wn = wn * w;
                 swap(a, b);
        if (rev) {
                 auto n1 = mint(1) / n;
                 for (auto& x : a) x *= n1;
vector<mint> convolution(const vector<mint>& a, const vector<mint>& b) {
        vector<mint> l(a.begin(), a.end()), r(b.begin(), b.end());
        int N = 1.size()+r.size()-1, n = 1;
        while (n \le N) n \ne 2;
        1.resize(n);
        r.resize(n);
        ntt(1, false);
        ntt(r, false);
```

```
for (int i = 0; i < n; i++) l[i] *= r[i];
ntt(l, true);
l.resize(N);
return l;</pre>
```

### 28.25 pollardrho

```
// Pollard's Rho Alg
// Usa o algoritmo de deteccao de ciclo de Floyd
// com uma otimizacao na qual o gcd eh acumulado
// A fatoracao nao sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
// Complexidades (considerando mul constante):
// rho - esperado O(n^(1/4)) no pior caso
// fact - esperado menos que O(n^{\hat{}}(1/4) \log(n)) no pior caso
11 mul(11 a, 11 b, 11 m) {
        11 ret = a*b - 11((long double)1/m*a*b+0.5)*m;
        return ret < 0 ? ret+m : ret;</pre>
ll pow(ll x, ll y, ll m) {
        if (!y) return 1;
        ll ans = pow(mul(x, x, m), y/2, m);
        return y%2 ? mul(x, ans, m) : ans;
bool prime(ll n) {
        if (n < 2) return 0;
        if (n <= 3) return 1;
        if (n % 2 == 0) return 0;
        ll r = \underline{\quad builtin\_ctzll(n - 1), d = n >> r;}
        for (int a: {2, 325, 9375, 28178, 450775, 9780504, 1795265022}) {
                11 x = pow(a, d, n);
                if (x == 1 \text{ or } x == n - 1 \text{ or } a \% n == 0) continue;
                 for (int j = 0; j < r - 1; j++) {
                         x = mul(x, x, n);
                         if (x == n - 1) break;
                if (x != n - 1) return 0;
        return 1;
11 rho(11 n) {
        if (n == 1 or prime(n)) return n;
        auto f = [n](11 x) \{ return mul(x, x, n) + 1; \};
        11 x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
        while (t % 40 != 0 or gcd(prd, n) == 1) {
                if (x==y) x = ++x0, y = f(x);
                q = mul(prd, abs(x-y), n);
                if (q != 0) prd = q;
                x = f(x), y = f(f(y)), t++;
        return gcd(prd, n);
vector<ll> fact(ll n) {
        if (n == 1) return {};
        if (prime(n)) return {n};
        11 d = rho(n);
        vector<11>1 = fact(d), r = fact(n / d);
        l.insert(l.end(), r.begin(), r.end());
        return 1;
```

### 28.26 powerSeries

```
// Operacoes em Polinomios e Series de Potencias
// Precisa do NTT
// O exp nao foi bem testado
// Fonte: github.com/celiopassos/competitive-programming/blob/master/algorithms/
    mathematics/formal_power_series.hpp
// inv, divmod, log e exp: O(n log(n))
using poly = vector<mint>;
const int MAGIC = 512:
poly D(poly p) {
        if (p.empty()) return p;
        for (int i = 0; i + 1 < p.size(); i++)
                p[i] = (i + 1) * p[i + 1];
        p.pop_back();
poly I(poly p) {
        int n = p.size();
        p.push_back(0);
        for (int i = n - 1; i >= 0; i--)
                p[i + 1] = p[i] / (i + 1);
        p[0] = 0;
        return p;
poly inv(poly p) {
        assert(!p.empty() && p[0] == 1);
        poly q = {mint(1) / p[0]};
        int n = p.size(), k = 1;
        while (k < n) {
               k \star = 2;
                q.resize(2 * k);
                ntt(q, false);
                poly p0(2 * k);
                copy_n(p.begin(), min(k, n), p0.begin());
                ntt(p0, false);
                for (int i = 0; i < 2 * k; i++)
                        q[i] *= 2 - p0[i] * q[i];
                ntt(q, true);
                q.resize(k);
        q.resize(n);
        return q;
pair<poly, poly> divslow(const poly& a, const poly& b) {
        poly q, r = a;
        while (r.size() >= b.size()) {
                q.push_back(r.back() / b.back());
                if (q.back() != 0)
                        for (int i = 0; i < b.size(); i++)</pre>
                                 r.end()[-i-1] = q.back() * b.end()[-i-1];
                r.pop_back();
        reverse(q.begin(), q.end());
        return {q, r};
// retorna (q, r) : a(x) = b(x) * q(x) + r(x)
pair<poly, poly> divmod(const poly& a, const poly& b) {
        if (a.size() < b.size()) return {{}, a};</pre>
        if (max(b.size(), a.size() - b.size()) < MAGIC) return divslow(a, b);</pre>
        poly ra = poly(a.rbegin(), a.rend());
        poly rb = poly(b.rbegin(), b.rend());
        int k = a.size() - b.size() + 1;
        rb.resize(k);
        poly irb = inv(move(rb)), q = convolution(ra, irb);
        q = poly(q.rend() - k, q.rend());
        poly r = convolution (move(q), b);
```

```
for (int i = 0; i < r.size(); i++) r[i] = a[i] - r[i];</pre>
        while (r.size() > 1 && r.back() == 0) r.pop_back();
        return {q, r};
poly log(poly p) {
        assert(!p.empty() && p[0] == 1);
        int n = p.size();
        auto d = D(p), i = inv(p);
        auto r = convolution(d, i);
        r.resize(n - 1);
        return I (move(r));
poly exp(poly p) {
        assert(p.empty() || p[0] == 0);
        poly q = \{1\};
        int n = p.size(), k = 1;
        while (k < n) {
                k *= 2;
                q.resize(k);
                poly b = log(q);
                 for (int i = 0; i < k; i++) b[i] *= -1;
                b[0] += 1;
                for (int i = 0; i < min(n, k); i++) b[i] += p[i];</pre>
                q = convolution(q, b);
                q.resize(k);
        q.resize(n);
        return q;
```

## 28.27 probabilityBinomial

```
// Binomial Distribution
//
binom(n, k, p) retorna a probabilidade de k sucessos
// numa binomial(n, p)

double logfact[MAX];

void calc() {
        logfact[0] = 0;
        for (int i = 1; i < MAX; i++) logfact[i] = logfact[i-1] + log(i);
}

double binom(int n, int k, double p) {
        return exp(logfact[n] - logfact[k] - logfact[n-k] + k * log(p) + (n-k) * log(1 - p));
}</pre>
```

# 28.28 simplex

```
// Simplex
//
// Maximiza c^T x s.t. Ax <= b, x >= 0
//
// O(2^n), porem executa em O(n^3) no caso medio

const double eps = le-7;

namespace Simplex {
    vector<vector<double>> T;
    int n, m;
    vector<int> X, Y;

    void pivot(int x, int y) {
        swap(X[y], Y[x-1]);
        for (int i = 0; i <= m; i++) if (i != y) T[x][i] /= T[x][y];
        for (int i = 0; i <= n; i++) if (i != x and abs(T[i][y]) > eps)
}
```

```
for (int j = 0; j <= m; j++) if (j != y) T[i][j] -= T[i</pre>
                     ][y] * T[x][j];
                 T[i][y] = -T[i][y] * T[x][y];
        }
// Retorna o par (valor maximo, vetor solucao)
pair<double, vector<double>> simplex(
                vector<vector<double>> A, vector<double> b, vector<
                      double> c) {
        n = b.size(), m = c.size();
        T = vector(n + 1, vector < double > (m + 1));
        X = vector<int>(m);
        Y = vector<int>(n);
        for (int i = 0; i < m; i++) X[i] = i; for (int i = 0; i < n; i++) Y[i] = i+m;
        for (int i = 0; i < m; i++) T[0][i] = -c[i];
        for (int i = 0; i < n; i++) {
                 for (int j = 0; j < m; j++) T[i+1][j] = A[i][j];
                 T[i+1][m] = b[i];
        while (true) {
                 int x = -1, y = -1;
                 double mn = -eps;
                 for (int i = 1; i <= n; i++) if (T[i][m] < mn) mn = T[i
                     [m], x = i;
                 if (x < 0) break;
                 for (int i = 0; i < m; i++) if (T[x][i] < -eps) { y = i;}
                       break; }
                 if (y < 0) return {-1e18, {}}; // sem solucao para Ax</pre>
                      \leq b
                pivot(x, y);
        while (true) {
                 int x = -1, y = -1;
                 double mn = -eps;
                 for (int i = 0; i < m; i++) if (T[0][i] < mn) mn = T[0][</pre>
                     i], y = i;
                 if (y < 0) break;
                 mn = 1e200;
                 for (int i = 1; i <= n; i++) if (T[i][y] > eps and T[i][ m] / T[i][y] < mn)
                         mn = T[i][m] / T[i][y], x = i;
                 if (x < 0) return {1e18, {}}; // c^T x eh ilimitado</pre>
                 pivot(x, y);
        vector<double> r (m);
        for (int i = 0; i < n; i++) if (Y[i] < m) r[Y[i]] = T[i+1][m];
        return {T[0][m], r};
```

#### 28.29 totiente

```
// Totiente
//
// O(sqrt(n))
int tot(int n) {
    int ret = n;

    for (int i = 2; i*i <= n; i++) if (n % i == 0) {
        while (n % i == 0) n /= i;
        ret -= ret / i;
    }
    if (n > 1) ret -= ret / n;
    return ret;
}
```

# 29 ufmg forked/Primitivas

### 29.1 bigint

```
// Big Integer
// Complexidades: (para n digitos)
// Soma, subtracao, comparacao - O(n)
// Multiplicacao - O(n log(n))
// Divisao, resto - O(n^2)
struct bint {
        static const int BASE = 1e9;
        vector<int> v;
        bool neg;
        bint() : neg(0) {}
        bint(int val) : bint() { *this = val; }
       bint(long long val) : bint() { *this = val; }
        void trim()
                while (v.size() and v.back() == 0) v.pop_back();
                if (!v.size()) neg = 0;
        // converter de/para string | cin/cout
        bint(const char* s) : bint() { from_string(string(s)); }
        bint(const string& s) : bint() { from_string(s); }
        void from_string(const string& s) {
                v.clear(), neg = 0;
                int ini = 0;
                while (ini < s.size() and (s[ini] == '-' or s[ini] == '+' or s[
                    ini] == '0'))
                        if (s[ini++] == '-') neg = 1;
                for (int i = s.size()-1; i >= ini; i -= 9) {
                        int at = 0;
                        for (int j = max(ini, i - 8); j <= i; j++) at = 10*at +
                            (s[i]-'0');
                        v.push_back(at);
                if (!v.size()) neg = 0;
        string to_string() const {
                if (!v.size()) return "0";
                string ret;
                if (neg) ret += '-';
                for (int i = v.size()-1; i >= 0; i--) {
                        string at = ::to_string(v[i]);
                        int add = 9 - at.size();
                        if (i+1 < v.size()) for (int j = 0; j < add; j++) ret +=</pre>
                             ′0′;
                        ret += at;
                return ret;
        friend istream& operator>>(istream& in, bint& val) {
               string s; in >> s;
               val = s:
                return in;
        friend ostream& operator<<(ostream& out, const bint& val) {</pre>
               string s = val.to_string();
               out << s:
                return out;
        // operators
        friend bint abs(bint val) {
               val.neg = 0;
               return val;
        friend bint operator-(bint val) {
                if (val != 0) val.neg ^= 1;
                return val;
        bint& operator=(const bint& val) { v = val.v, neg = val.neg; return *
```

```
this; }
bint& operator=(long long val) {
        v.clear(), neg = 0;
        if (val < 0) neg = 1, val *= -1;
for (; val; val /= BASE) v.push_back(val % BASE);</pre>
        return *this;
int cmp(const bint& r) const { // menor: -1 | igual: 0 | maior: 1
        if (neg != r.neg) return neg ? -1 : 1;
        if (v.size() != r.v.size()) {
                int ret = v.size() < r.v.size() ? -1 : 1;</pre>
                return neg ? -ret : ret;
        for (int i = int(v.size())-1; i >= 0; i--) {
                if (v[i] != r.v[i]) {
                         int ret = v[i] < r.v[i] ? -1 : 1;</pre>
                         return neg ? -ret : ret;
        return 0;
friend bool operator<(const bint& 1, const bint& r) { return 1.cmp(r) ==</pre>
friend bool operator>(const bint& 1, const bint& r) { return 1.cmp(r) ==
     1; }
friend bool operator <= (const bint& 1, const bint& r) { return 1.cmp(r)
    <= 0; }
friend bool operator>=(const bint& 1, const bint& r) { return 1.cmp(r)
    >= 0;}
friend bool operator==(const bint& 1, const bint& r) { return 1.cmp(r)
friend bool operator!=(const bint& 1, const bint& r) { return l.cmp(r)
    != 0; }
bint& operator +=(const bint& r) {
        if (!r.v.size()) return *this;
        if (neg != r.neg) return *this -= -r;
        for (int i = 0, c = 0; i < r.v.size() or c; i++) {</pre>
                if (i == v.size()) v.push_back(0);
                v[i] += c + (i < r.v.size() ? r.v[i] : 0);
                if ((c = v[i] >= BASE)) v[i] -= BASE;
        return *this;
friend bint operator+(bint a, const bint& b) { return a += b; }
bint& operator -= (const bint& r) {
        if (!r.v.size()) return *this;
        if (neg != r.neg) return *this += -r;
        if ((!neg and *this < r) or (neg and r < *this)) {</pre>
                *this = r - *this;
                neg ^= 1;
                return *this;
        for (int i = 0, c = 0; i < r.v.size() or c; i++) {
                v[i] = c + (i < r.v.size() ? r.v[i] : 0);
                if ((c = v[i] < 0)) v[i] += BASE;
        trim();
        return *this;
friend bint operator-(bint a, const bint& b) { return a -= b; }
// operators de * / %
bint& operator *=(int val) {
        if (val < 0) val *= -1, neg ^= 1;</pre>
        for (int i = 0, c = 0; i < v.size() or c; i++) {
                if (i == v.size()) v.push_back(0);
                long long at = (long long) v[i] * val + c;
                v[i] = at % BASE;
                c = at / BASE;
        trim();
        return *this;
friend bint operator *(bint a, int b) { return a *= b; }
friend bint operator *(int a, bint b) { return b *= a; }
using cplx = complex<double>;
void fft(vector<cplx>& a, bool f, int N, vector<int>& rev) const {
```

```
for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i], a[rev[i])</pre>
        vector<cplx> roots(N);
        for (int n = 2; n \le N; n *= 2) {
                const static double PI = acos(-1);
                for (int i = 0; i < n/2; i++) {
                         double alpha = (2*PI*i)/n;
                         if (f) alpha = -alpha;
                         roots[i] = cplx(cos(alpha), sin(alpha));
                for (int pos = 0; pos < N; pos += n)
                         for (int 1 = pos, r = pos+n/2, m = 0; m < n/2; 1
                              ++, r++, m++) {
                                 auto t = roots[m] *a[r];
                                 a[r] = a[1] - t;
                                 a[1] = a[1] + t;
        if (!f) return;
        auto invN = cplx(1)/cplx(N);
        for (int i = 0; i < N; i++) a[i] *= invN;</pre>
vector<long long> convolution(const vector<int>& a, const vector<int>& b
    ) const {
        vector<cplx> l(a.begin(), a.end()), r(b.begin(), b.end());
        int ln = 1.size(), rn = r.size(), N = ln+rn+1, n = 1, log_n = 0;
        while (n <= N) n <<= 1, log_n++;</pre>
        vector<int> rev(n);
        for (int i = 0; i < n; i++) {</pre>
                rev[i] = 0;
                for (int j = 0; j < log_n; j++) if (i>>j&1)
                        rev[i] |= 1 << (log_n-1-j);
        l.resize(n), r.resize(n);
        fft(1, false, n, rev), fft(r, false, n, rev);
        for (int i = 0; i < n; i++) l[i] *= r[i];</pre>
        fft(1, true, n, rev);
        vector<long long> ret;
        for (auto& i : 1) ret.push_back(round(i.real()));
        return ret;
vector<int> convert_base(const vector<int>& a, int from, int to) const {
        static vector<long long> pot(10, 1);
        if (pot[1] == 1) for (int i = 1; i < 10; i++) pot[i] = 10*pot[i</pre>
             -11;
        vector<int> ret;
        long long at = 0;
        int digits = 0;
        for (int i : a) {
                at += i * pot[digits];
                digits += from;
                while (digits >= to) {
                        ret.push_back(at % pot[to]);
                        at /= pot[to];
digits -= to;
        ret.push_back(at);
        while (ret.size() and ret.back() == 0) ret.pop_back();
        return ret;
bint operator*(const bint& r) const { // O(n log(n))
        bint ret;
        ret.neg = neg ^ r.neg;
        auto conv = convolution(convert_base(v, 9, 4), convert_base(r.v,
              9, 4));
        long long c = 0;
        for (auto i : conv) {
                long long at = i+c;
ret.v.push_back(at % 10000);
                c = at / 10000;
        for (; c; c /= 10000) ret.v.push_back(c%10000);
        ret.v = convert_base(ret.v, 4, 9);
        if (!ret.v.size()) ret.neg = 0;
        return ret;
bint& operator*=(const bint& r) { return *this = *this * r; };
```

```
bint& operator/=(int val) {
        if (val < 0) neg ^= 1, val *= -1;</pre>
        for (int i = int(v.size())-1, c = 0; i >= 0; i--) {
                long long at = v[i] + c * (long long) BASE;
                v[i] = a\bar{t} / val;
                c = at % val;
        trim();
        return *this;
friend bint operator/(bint a, int b) { return a /= b; }
int operator %=(int val) {
        if (val < 0) val *= -1;
        long long at = 0;
        for (int i = int(v.size())-1; i >= 0; i--)
               at = (BASE * at + v[i]) % val;
        if (neg) at \star = -1;
        return at:
friend int operator%(bint a, int b) { return a %= b; }
friend pair<bint, bint> divmod(const bint& a_, const bint& b_) { // O(n
       if (a_ == 0) return {0, 0};
        int norm = BASE / (b_.v.back() + 1);
        bint a = abs(a_) * norm;
        bint b = abs(b_) * norm;
        bint q, r;
        for (int i = a.v.size() - 1; i >= 0; i--) {
                r \star = BASE, r += a.v[i];
                long long upper = b.v.size() < r.v.size() ? r.v[b.v.size</pre>
                     ()]:0;
                int lower = b.v.size() - 1 < r.v.size() ? r.v[b.v.size()</pre>
                     - 1] : 0;
                int d = (upper * BASE + lower) / b.v.back();
                r \rightarrow b*d;
                while (r < 0) r += b, d--; // roda O(1) vezes
                q.v.push_back(d);
        reverse(q.v.begin(), q.v.end());
        q.neg = a_.neg ^ b_.neg;
        r.neg = a_.neg;
        q.trim(), r.trim();
        return {q, r / norm};
bint operator/(const bint& val) { return divmod(*this, val).first; }
bint& operator/=(const bint& val) { return *this = *this / val; }
bint operator% (const bint& val) { return divmod(*this, val) second; }
bint& operator%=(const bint& val) { return *this = *this % val; }
```

#### 29.2 calendario

};

```
// Calendario
//
// Congruencia de Zeller
//
// Os dias da semana correspondem aos restos % 7
// Segunda=0, Terca=1, ..., Domingo=6

int get_id(int d, int m, int y) {
    if (m < 3) y--, m += 12;
        return 365 * y + y / 4 - y / 100 + y / 400 + (153 * (m - 3) + 2) / 5 + d -
        307;
}

tuple<int, int, int> date(int id) {
    int x = id + 1789995, n = 4 * x / 146097, i, j, d, m, y;
    x -= (146097 * n + 3) / 4;
    i = (4000 * (x + 1)) / 1461001;
    x -= 1461 * i / 4 - 31;
    j = 80 * x / 2447, d = x - 2447 * j / 80;
    x = j / 11;
    m = j + 2 - 12 * x, y = 100 * (n - 49) + i + x;
    return {d, m, y};
```

#### 29.3 frac

```
// Fracao
// Funciona com o Big Int
template<typename T = int> struct frac {
        T num, den:
        template<class U, class V>
        frac(U num_ = 0, V den_ = 1) : num(num_), den(den_) {
                assert (den != 0);
                if (den < 0) num *= -1, den *= -1;
                T g = gcd(abs(num), den);
                num /= g, den /= g;
        friend bool operator<(const frac& 1, const frac& r) {</pre>
                return 1.num * r.den < r.num * 1.den;</pre>
        friend frac operator+(const frac& 1, const frac& r) {
                return {1.num*r.den + 1.den*r.num, 1.den*r.den};
        friend frac operator-(const frac& 1, const frac& r) {
                return {1.num*r.den - 1.den*r.num, 1.den*r.den};
        friend frac operator*(const frac& 1, const frac& r) {
                return {1.num*r.num, 1.den*r.den};
        friend frac operator/(const frac& 1, const frac& r) {
                return {1.num*r.den, 1.den*r.num};
        friend ostream& operator<<(ostream& out, frac f) {</pre>
                out << f.num << '/' << f.den;
                return out;
};
```

# 29.4 geometria

```
// Geometria
typedef double 1d;
const 1d DINF = 1e18;
const 1d pi = acos(-1.0);
const 1d eps = 1e-9;
#define sq(x) ((x)*(x))
bool eq(ld a, ld b) {
        return abs(a - b) <= eps;
struct pt { // ponto
        1d x, y;
pt(ld x_ = 0, ld y_ = 0) : x(x_), y(y_) {}
        bool operator < (const pt p) const {</pre>
                 if (!eq(x, p.x)) return x < p.x;
                 if (!eq(y, p.y)) return y < p.y;</pre>
                 return 0;
        bool operator == (const pt p) const {
                 return eq(x, p.x) and eq(y, p.y);
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y);
        pt operator * (const ld c) const { return pt(x*c , y*c ); }
         pt operator / (const ld c) const { return pt(x/c , y/c ); }
        Id operator ^ (const pt p) const { return x*p.x + y*p.y; } ld operator ^ (const pt p) const { return x*p.y - y*p.x; }
         friend istream& operator >> (istream& in, pt& p) {
                 return in >> p.x >> p.y;
};
```

```
struct line { // reta
        pt p, q;
        line() {}
        line(pt p_, pt q_) : p(p_), q(q_) {}
        friend istream& operator >> (istream& in, line& r) {
                return in >> r.p >> r.q;
};
// PONTO & VETOR
ld dist(pt p, pt q) { // distancia
        return hypot(p.y - q.y, p.x - q.x);
ld dist2(pt p, pt q) { // quadrado da distancia
        return sq(p.x - q.x) + sq(p.y - q.y);
ld norm(pt v) { // norma do vetor
        return dist(pt(0, 0), v);
ld angle(pt v) { // angulo do vetor com o eixo x
        ld ang = atan2(v.y, v.x);
        if (ang < 0) ang += 2*pi;</pre>
        return ang;
ld sarea(pt p, pt q, pt r) { // area com sinal
        return ((q-p)^(r-q))/2;
bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
        return eq(sarea(p, q, r), 0);
bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
        return sarea(p, q, r) > eps;
pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
        return pt (p.x * cos(th) - p.v * sin(th),
                         p.x * sin(th) + p.y * cos(th));
pt rotate90(pt p) { // rotaciona 90 graus
        return pt(-p.y, p.x);
// RETA
bool isvert(line r) { // se r eh vertical
        return eq(r.p.x, r.q.x);
bool isinseg(pt p, line r) { // se p pertence ao seg de r
        pt a = r.p - p, b = r.q - p;
return eq((a ^ b), 0) and (a * b) < eps;</pre>
ld get_t(pt v, line r) { // retorna t tal que t*v pertence a reta r
        return (r.p^r.q) / ((r.p-r.q)^v);
pt proj(pt p, line r) { // projecao do ponto p na reta r
        if (r.p == r.q) return r.p;
        r.q = r.q - r.p; p = p - r.p;
        pt proj = r.q * ((p*r.q) / (r.q*r.q));
        return proj + r.p;
pt inter(line r, line s) { // r inter s
if (eq((r.p - r.q) ^ (s.p - s.q), 0)) return pt(DINF, DINF);
        r.q = r.q - r.p, s.p = s.p - r.p, s.q = s.q - r.p;
        return r.q * get_t(r.q, s) + r.p;
}
```

```
bool interseg(line r, line s) { // se o seg de r intersecta o seg de s
        if (isinseg(r.p, s) or isinseg(r.q, s)
                or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
                        ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
ld disttoline(pt p, line r) { // distancia do ponto a reta
        return 2 * abs(sarea(p, r.p, r.q)) / dist(r.p, r.q);
ld disttoseg(pt p, line r) { // distancia do ponto ao seg
        if ((r.q - r.p) * (p - r.p) < 0) return dist(r.p, p);
        if ((r.p - r.q) * (p - r.q) < 0) return dist(r.q, p);
        return disttoline(p, r);
ld distseg(line a, line b) { // distancia entre seg
        if (interseg(a, b)) return 0;
        ld ret = DINF;
        ret = min(ret, disttoseg(a.p, b));
        ret = min(ret, disttoseg(a.q, b));
        ret = min(ret, disttoseg(b.p, a));
        ret = min(ret, disttoseg(b.q, a));
        return ret:
// POLIGONO
// corta poligono com a reta r deixando os pontos p tal que
// ccw(r.p, r.q, p)
vector<pt> cut_polygon(vector<pt> v, line r) { // O(n)
        vector<pt> ret;
        for (int j = 0; j < v.size(); j++) {</pre>
                if (ccw(r.p, r.q, v[j])) ret.push_back(v[j]);
                if (v.size() == 1) continue;
                line s(v[j], v[(j+1) %v.size()]);
                pt p = inter(r, s);
                if (isinseg(p, s)) ret.push_back(p);
        ret.erase(unique(ret.begin(), ret.end()), ret.end());
        if (ret.size() > 1 and ret.back() == ret[0]) ret.pop_back();
        return ret;
// distancia entre os retangulos a e b (lados paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
ld dist_rect(pair<pt, pt> a, pair<pt, pt> b) {
        1d hor = 0, vert = 0;
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
        return dist(pt(0, 0), pt(hor, vert));
ld polarea(vector<pt> v) { // area do poligono
        1d ret = 0;
        for (int i = 0; i < v.size(); i++)</pre>
                ret += sarea(pt(0, 0), v[i], v[(i + 1) % v.size()]);
        return abs(ret);
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
int inpol(vector<pt>& v, pt p) { // O(n)
        int qt = 0;
        for (int i = 0; i < v.size(); i++) {
                if (p == v[i]) return 2;
                int j = (i+1)%v.size();
                if (eq(p.y, v[i].y) and eq(p.y, v[j].y)) {
                        if ((v[i]-p)*(v[j]-p) < eps) return 2;</pre>
                        continue;
                bool baixo = v[i].y+eps < p.y;</pre>
                if (baixo == (v[j].y+eps < p.y)) continue;</pre>
```

```
auto t = (p-v[i])^(v[j]-v[i]);
                 if (eq(t, 0)) return 2;
                 if (baixo == (t > eps)) qt += baixo ? 1 : -1;
         return qt != 0;
bool interpol(vector<pt> v1, vector<pt> v2) { // se dois poligonos se
     intersectam - O(n*m)
         int n = v1.size(), m = v2.size();
        for (int i = 0; i < n; i++) if (inpol(v2, v1[i])) return 1;
for (int i = 0; i < n; i++) if (inpol(v1, v2[i])) return 1;</pre>
         for (int i = 0; i < n; i++) for (int j = 0; j < m; j++)
                 if (interseg(line(v1[i], v1[(i+1)%n]), line(v2[j], v2[(j+1)%m]))
                      ) return 1:
         return 0:
ld distpol(vector<pt> v1, vector<pt> v2) { // distancia entre poligonos
         if (interpol(v1, v2)) return 0;
        ld ret = DINF;
         for (int i = 0; i < v1.size(); i++) for (int j = 0; j < v2.size(); j++)
                 ret = min(ret, distseg(line(v1[i], v1[(i + 1) % v1.size()]),
                                           line(v2[j], v2[(j + 1) % v2.size()])));
         return ret;
vector<pt> convex_hull (vector<pt> v) { // convex hull - O(n log(n))
         sort(v.begin(), v.end());
         v.erase(unique(v.begin(), v.end()), v.end());
         if (v.size() <= 1) return v;</pre>
         vector<pt> 1, u;
        for (int i = 0; i < v.size(); i++) {
   while (l.size() > 1 and !ccw(l.end()[-2], l.end()[-1], v[i]))
                         1.pop_back();
                 1.push_back(v[i]);
         for (int i = v.size() - 1; i >= 0; i--) {
                 while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1], v[i]))
                          u.pop_back();
                 u.push_back(v[i]);
         1.pop_back(); u.pop_back();
         for (pt i : u) l.push_back(i);
         return 1:
struct convex_pol {
        vector<pt> pol;
         // nao pode ter ponto colinear no convex hull
         convex_pol() {}
         convex_pol(vector<pt> v) : pol(convex_hull(v)) {}
         // se o ponto ta dentro do hull - O(log(n))
         bool is_inside(pt p) {
                 if (pol.size() == 0) return false;
                 if (pol.size() == 1) return p == pol[0];
                 int 1 = 1, r = pol.size();
                 while (1 < r) {
                          int m = (1+r)/2;
                          if (ccw(p, pol[0], pol[m])) l = m+1;
                          else r = m;
                 if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
                 if (1 == pol.size()) return false;
                 return !ccw(p, pol[1], pol[1-1]);
         // ponto extremo em relacao a cmp(p, q) = p mais extremo q
         // (copiado de https://github.com/gustavoM32/caderno-zika)
         int extreme(const function<bool(pt, pt)>& cmp) {
                 int n = pol.size();
                 auto extr = [&](int i, bool& cur_dir) {
                          \operatorname{cur\_dir} = \operatorname{cmp}(\operatorname{pol}[(i+1) %n], \operatorname{pol}[i]);
                          return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
                 bool last_dir, cur_dir;
```

```
if (extr(0, last_dir)) return 0;
               int 1 = 0, r = n;
                while (1+1 < r) {
                       int m = (1+r)/2;
                       if (extr(m, cur_dir)) return m;
                       bool rel_dir = cmp(pol[m], pol[l]);
                       if ((!last_dir and cur_dir) or
                                        (last_dir == cur_dir and rel_dir ==
                                            cur_dir)) {
                               last dir = cur dir;
                        } else r = m;
               return 1;
        int max_dot(pt v) {
                return extreme([&](pt p, pt q) { return p*v > q*v; });
        pair<int, int> tangents(pt p) {
                auto L = [&] (pt q, pt r) { return ccw(p, r, q); };
                auto R = [&] (pt q, pt r) { return ccw(p, q, r); };
               return {extreme(L), extreme(R)};
};
// CIRCUNFERENCIA
b = (a + b) / 2;
        c = (a + c) / 2;
        return inter(line(b, b + rotate90(a - b)),
                       line(c, c + rotate90(a - c)));
vector<pt> circ_line_inter(pt a, pt b, pt c, ld r) { // intersecao da circunf (c
    , r) e reta ab
        vector<pt> ret;
        b = b-a, a = a-c;
        1d A = b*b:
        1d B = a*b;
        1d C = a*a - r*r;
        1d D = B*B - A*C;
       if (D < -eps) return ret;</pre>
        ret.push back(c+a+b*(-B+sqrt(D+eps))/A);
        if (D > eps) ret.push_back(c+a+b*(-B-sqrt(D))/A);
        return ret;
vector<pt> circ_inter(pt a, pt b, ld r, ld R) { // intersecao da circunf (a, r)
    e (b, R)
        vector<pt> ret;
        1d d = dist(a, b);
        if (d > r+R or d+min(r, R) < max(r, R)) return ret;</pre>
        1d x = (d*d-R*R+r*r)/(2*d);
        1d y = sqrt(r*r-x*x);
        pt v = (b-a)/d;
        ret.push_back(a+v*x + rotate90(v)*y);
        if (y > 0) ret.push_back(a+v*x - rotate90(v)*y);
        return ret;
bool operator <(const line& a, const line& b) { // comparador pra reta</pre>
        // assume que as retas tem p < q
        pt v1 = a.q - a.p, v2 = b.q - b.p;
        if (!eq(angle(v1), angle(v2))) return angle(v1) < angle(v2);</pre>
        return ccw(a.p, a.q, b.p); // mesmo angulo
bool operator ==(const line& a, const line& b) {
        return !(a < b) and !(b < a);
// comparador pro set pra fazer sweep line com segmentos
struct cmp_sweepline {
       bool operator () (const line& a, const line& b) const {
                // assume que os segmentos tem p < q
                if (a.p == b.p) return ccw(a.p, a.q, b.q);
                if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or a.p.x+eps < b.p.x</pre>
                    ))
                       return ccw(a.p, a.q, b.p);
```

```
return ccw(a.p, b.q, b.p);
};

// comparador pro set pra fazer sweep angle com segmentos
pt dir;
struct cmp_sweepangle {
            bool operator () (const line& a, const line& b) const {
                return get_t(dir, a) + eps < get_t(dir, b);
            }
};</pre>
```

### 29.5 geometria3d

```
// Geometria 3D
typedef double ld;
const 1d DINF = 1e18;
const ld eps = 1e-9;
#define sq(x) ((x)*(x))
bool eq(ld a, ld b) {
        return abs(a - b) <= eps;</pre>
struct pt { // ponto
        ld x, y, z;
        pt(1d x_{=} 0, 1d y_{=} 0, 1d z_{=} 0) : x(x_{)}, y(y_{)}, z(z_{)} {}
        bool operator < (const pt p) const {
                if (!eq(x, p.x)) return x < p.x;</pre>
                if (!eq(y, p.y)) return y < p.y;</pre>
                if (!eq(z, p.z)) return z < p.z;</pre>
                return 0;
        bool operator == (const pt p) const {
                return eq(x, p.x) and eq(y, p.y) and eq(z, p.z);
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y, z+p.z);
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y, z-p.z); }
        pt operator * (const ld c) const {
                                            return pt(x*c , y*c , z*c ); }
        pt operator / (const ld c) const { return pt(x/c , y/c , z/c ); }
        ld operator * (const pt p) const { return x*p.x + y*p.y + z*p.z; }
        pt operator ^ (const pt p) const { return pt(y*p.z - z*p.y, z*p.x - x*p.
             z, x*p.y - y*p.x); }
        friend istream& operator >> (istream& in, pt& p) {
                return in >> p.x >> p.y >> p.z;
};
struct line { // reta
        pt p, q;
        line() {}
        line(pt p_, pt q_) : p(p_), q(q_) {}
        friend istream& operator >> (istream& in, line& r) {
                return in >> r.p >> r.q;
};
struct plane { // plano
        array<pt, 3> p; // pontos que definem o plano
        array<ld, 4> eq; // equacao do plano
        plane() {}
        plane(pt p_, pt q_, pt r_) : p(\{p_, q_, r_\}) \{ build(); \}
        friend istream& operator >> (istream& in, plane& P) {
                return in >> P.p[0] >> P.p[1] >> P.p[2];
                P.build();
        void build() {
                pt dir = (p[1] - p[0]) ^ (p[2] - p[0]);
                eq = {dir.x, dir.y, dir.z, dir*p[0]*(-1)};
};
```

```
// converte de coordenadas polares para cartesianas
// (angulos devem estar em radianos)
// phi eh o angulo com o eixo z (cima) theta eh o angulo de rotacao ao redor de
pt convert(ld rho, ld th, ld phi) {
        return pt(sin(phi) * cos(th), sin(phi) * sin(th), cos(phi)) * rho;
// projecao do ponto p na reta r
pt proj(pt p, line r) {
        if (r.p == r.q) return r.p;
        r.q = r.q - r.p; p = p - r.p;
        pt proj = r.q * ((p*r.q) / (r.q*r.q));
        return proj + r.p;
// projecao do ponto p no plano P
pt proj(pt p, plane P) {
        p = p - P.p[0], P.p[1] = P.p[1] - P.p[0], P.p[2] = P.p[2] - P.p[0];
        pt norm = P.p[1] ^ P.p[2];
        pt proj = p - (norm * (norm * p) / (norm*norm));
        return proj + P.p[0];
// distancia
ld dist(pt a, pt b) {
        return sqrt(sq(a.x-b.x) + sq(a.y-b.y) + sq(a.z-b.z));
// distancia ponto reta
ld distline(pt p, line r) {
        return dist(p, proj(p, r));
// distancia de ponto para segmento
ld distseg(pt p, line r) {
        if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
        if ((r.p - r.q) * (p - r.q) < 0) return dist(r.q, p);
        return distline(p, r);
// distancia de ponto a plano com sinal
ld sdist(pt p, plane P) {
        return P.eq[0]*p.x + P.eq[1]*p.y + P.eq[2]*p.z + P.eq[3];
// distancia de ponto a plano
ld distplane(pt p, plane P) {
        return abs(sdist(p, P));
// se ponto pertence a reta
bool isinseg(pt p, line r) {
        return eq(distseg(p, r), 0);
// se ponto pertence ao triangulo definido por P.p
bool isinpol(pt p, vector<pt> v) {
        assert(v.size() >= 3);
        pt norm = (v[1]-v[0]) ^ (v[2]-v[1]);
        bool inside = true;
        int sign = -1;
        for (int i = 0; i < v.size(); i++) {</pre>
                line r(v[(i+1)%3], v[i]);
                if (isinseq(p, r)) return true;
                pt ar = v[(i+1)%3] - v[i];
                if (sign == -1) sign = ((ar^(p-v[i]))*norm > 0);
                else if (((ar^(p-v[i]))*norm > 0) != sign) inside = false;
        return inside;
// distancia de ponto ate poligono
ld distpol(pt p, vector<pt> v) {
        pt p2 = proj(p, plane(v[0], v[1], v[2]));
        if (isinpol(p2, v)) return dist(p, p2);
        ld ret = DINF;
```

for (int i = 0; i < v.size(); i++) {</pre>

```
int j = (i+1)%v.size();
                ret = min(ret, distseg(p, line(v[i], v[j])));
        return ret:
// intersecao de plano e segmento
// BOTH = o segmento esta no plano
// ONE = um dos pontos do segmento esta no plano
// PARAL = segmento paralelo ao plano
// CONCOR = segmento concorrente ao plano
enum RETCODE {BOTH, ONE, PARAL, CONCOR};
pair<RETCODE, pt> intersect(plane P, line r) {
    1d d1 = sdist(r.p, P);
    1d d2 = sdist(r.q, P);
    if (eq(d1, 0)) and eq(d2, 0)
                return pair(BOTH, r.p);
    if (eq(d1, 0))
                return pair(ONE, r.p);
    if (eq(d2, 0))
                return pair(ONE, r.q);
    if ((d1 > 0 \text{ and } d2 > 0) \text{ or } (d1 < 0 \text{ and } d2 < 0))
        if (eq(d1-d2, 0)) return pair(PARAL, pt());
        return pair(CONCOR, pt());
    1d frac = d1 / (d1 - d2);
    pt res = r.p + ((r.q - r.p) * frac);
    return pair(ONE, res);
// rotaciona p ao redor do eixo u por um angulo a
return u * (u * p) + (u ^ p ^ u) * cos(a) + (u ^ p) * sin(a);
```

### 29.6 geometriaInt

```
// Geometria - inteiro
#define sq(x) ((x)*(ll)(x))
struct pt { // ponto
        int x, y;
        pt (int x_{-} = 0, int y_{-} = 0) : x(x_{-}), y(y_{-}) {}
        bool operator < (const pt p) const {
                 if (x != p.x) return x < p.x;
                 return y < p.y;</pre>
        bool operator == (const pt p) const {
                 return x == p.x and y == p.y;
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
         pt operator * (const int c) const { return pt(x*c, y*c); }
         11 operator * (const pt p) const { return x*(11)p.x + y*(11)p.y; }
         11 operator ^ (const pt p) const { return x*(11)p.y - y*(11)p.x; }
        friend istream& operator >> (istream& in, pt& p) {
                 return in >> p.x >> p.y;
};
struct line { // reta
        pt p, q;
         line() {}
        line(pt p_, pt q_) : p(p_), q(q_) {} friend istream& operator >> (istream& in, line& r) {
                 return in >> r.p >> r.q;
// PONTO & VETOR
11 dist2(pt p, pt q) { // quadrado da distancia
        return sq(p.x - q.x) + sq(p.y - q.y);
```

```
11 sarea2(pt p, pt q, pt r) { // 2 * area com sinal
         return (q-p) ^ (r-q);
bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
         return sarea2(p, q, r) == 0;
bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
         return sarea2(p, q, r) > 0;
int quad(pt p) { // quadrante de um ponto
         return (p.x<0) ^3*(p.y<0);
\begin{tabular}{ll} \textbf{bool} & \texttt{compare\_angle}(\texttt{pt} & \texttt{p}, & \texttt{pt} & \texttt{q}) & \texttt{f} & \texttt{//} & \texttt{retorna} & \texttt{se} & \texttt{ang}(\texttt{p}) & \texttt{f} & \texttt{ang}(\texttt{q}) \\ \end{tabular}
         if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
         return ccw(q, pt(0, 0), p);
pt rotate90(pt p) { // rotaciona 90 graus
         return pt(-p.y, p.x);
// RETA
bool isinseg(pt p, line r) { // se p pertence ao seg de r
         pt a = r.p - p, b = r.q - p;
return (a ^ b) == 0 and (a * b) <= 0;
bool interseg(line r, line s) { // se o seg de r intersecta o seg de s
         if (isinseq(r.p, s) or isinseq(r.q, s)
                  or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
         return ccw(r.p, r.q, s.p) := ccw(r.p, r.q, s.q) and
                           ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
int segpoints(line r) { // numero de pontos inteiros no segmento
         return 1 + __gcd(abs(r.p.x - r.q.x), abs(r.p.y - r.q.y));
double get_t(pt v, line r) { // retorna t tal que t*v pertence a reta r
         return (r.p^r.q) / (double) ((r.p-r.q)^v);
// POLIGONO
// quadrado da distancia entre os retangulos a e b (lados paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
11 dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
         int hor = 0, vert = 0;
         if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
         else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
         if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
         else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
         return sq(hor) + sq(vert);
11 polarea2(vector<pt> v) { // 2 * area do poligono
         11 \text{ ret} = 0;
         for (int i = 0; i < v.size(); i++)</pre>
                  ret += sarea2(pt(0, 0), v[i], v[(i + 1) % v.size()]);
         return abs(ret);
// se o ponto ta dentro do poligono: retorna 0 se ta fora,
// 1 se ta no interior e 2 se ta na borda
int inpol(vector<pt>& v, pt p) { // O(n)
         int qt = 0;
         for (int i = 0; i < v.size(); i++) {</pre>
                  if (p == v[i]) return 2;
                  int j = (i+1)%v.size();
                  if (p.y == v[i].y and p.y == v[j].y) {
                           if ((v[i]-p)*(v[j]-p) <= 0) return 2;</pre>
```

```
continue;
                bool baixo = v[i].y < p.y;</pre>
                if (baixo == (v[j].y < p.y)) continue;</pre>
                auto t = (p-v[i])^(v[j]-v[i]);
                if (!t) return 2;
                if (baixo == (t > 0)) qt += baixo ? 1 : -1;
        return qt != 0;
vector<pt> convex_hull (vector<pt> v) { // convex hull - O(n log(n))
        sort(v.begin(), v.end());
        v.erase(unique(v.begin(), v.end()), v.end());
        if (v.size() <= 1) return v;</pre>
        vector<pt> 1, u;
        for (int i = 0; i < v.size(); i++) {</pre>
                while (1.size() > 1 \text{ and } !ccw(1.end() [-2], 1.end() [-1], v[i]))
                        1.pop_back();
                1.push back(v[i]);
        for (int i = v.size() - 1; i >= 0; i--) {
                while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1], v[i]))
                        u.pop_back();
                u.push_back(v[i]);
        1.pop_back(); u.pop_back();
        for (pt i : u) l.push_back(i);
        return 1;
ll interior_points(vector<pt> v) { // pontos inteiros dentro de um poligono
     simples
        11 b = 0;
        for (int i = 0; i < v.size(); i++)</pre>
                b += segpoints(line(v[i], v[(i+1)%v.size()])) - 1;
        return (polarea2(v) - b) / 2 + 1;
struct convex_pol {
        vector<pt> pol;
        // nao pode ter ponto colinear no convex hull
        convex pol() {}
        convex_pol(vector<pt> v) : pol(convex_hull(v)) {}
        // se o ponto ta dentro do hull - O(log(n))
        bool is_inside(pt p) {
                if (pol.size() == 0) return false;
                if (pol.size() == 1) return p == pol[0];
                int 1 = 1, r = pol.size();
                while (1 < r) {
                        int m = (1+r)/2;
                        if (ccw(p, pol[0], pol[m])) 1 = m+1;
                        else r = m;
                if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
                if (1 == pol.size()) return false;
                return !ccw(p, pol[1], pol[1-1]);
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
        int extreme(const function<bool(pt, pt)>& cmp) {
                int n = pol.size();
                auto extr = [&](int i, bool& cur_dir) {
                        cur_dir = cmp(pol[(i+1)%n], pol[i]);
                        return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
                bool last_dir, cur_dir;
                if (extr(0, last dir)) return 0;
                int 1 = 0, r = n;
                while (1+1 < r) {
                        int m = (1+r)/2;
                        if (extr(m, cur_dir)) return m;
                        bool rel_dir = cmp(pol[m], pol[l]);
                        if ((!last_dir and cur_dir) or
                                         (last_dir == cur_dir and rel_dir ==
                                              cur_dir)) {
                                 1 = m:
```

```
last_dir = cur_dir;
                         } else r = m:
                return 1;
        int max_dot(pt v) {
                return extreme([&](pt p, pt q) { return p*v > q*v; });
        pair<int, int> tangents(pt p) {
                 auto L = [\&] (pt q, pt r) { return ccw(p, r, q); };
                 auto R = [&] (pt q, pt r) { return ccw(p, q, r); };
                return {extreme(L), extreme(R)};
        }
};
bool operator < (const line& a, const line& b) { // comparador pra reta
        // assume que as retas tem p < q
        pt v1 = a.q - a.p, v2 = b.q - b.p;
bool b1 = compare_angle(v1, v2), b2 = compare_angle(v2, v1);
        if (b1 or b2) return b1;
        return ccw(a.p, a.q, b.p); // mesmo angulo
bool operator == (const line& a, const line& b) {
        return !(a < b) and !(b < a);
// comparador pro set pra fazer sweep line com segmentos
struct cmp_sweepline {
        bool operator () (const line& a, const line& b) const {
                 // assume que os segmentos tem p < q
                 if (a.p == b.p) return ccw(a.p, a.q, b.q);
                if (a.p.x != a.q.x and (b.p.x == b.q.x or a.p.x < b.p.x))
                        return ccw(a.p, a.q, b.p);
                 return ccw(a.p, b.q, b.p);
};
// comparador pro set pra fazer sweep angle com segmentos
struct cmp_sweepangle {
    bool operator () (const line& a, const line& b) const {
        return get_t(dir, a) < get_t(dir, b);</pre>
};
```

#### 29.7 matrix

```
// Matriz
#define MODULAR false
template<typename T> struct matrix : vector<vector<T>> {
        int n, m;
        void print() {
                for (int i = 0; i < n; i++) {
                        for (int j = 0; j < m; j++) cout << (*this)[i][j] << " "</pre>
                        cout << endl:
        matrix(int n_, int m_, bool ident = false) :
                        vector<vector<T>> (n_, vector<T> (m_, 0)), n(n_), m(m_) {
                if (ident) {
                        assert(n == m):
                        for (int i = 0; i < n; i++) (*this)[i][i] = 1;</pre>
        matrix(const vector<vector<T>>& c) : vector<vector<T>> (c),
                n(c.size()), m(c[0].size()) {}
        matrix(const initializer_list<initializer_list<T>>& c) {
                vector<vector<T>> val;
                for (auto& i : c) val.push_back(i);
                *this = matrix(val);
```

```
matrix<T> operator*(matrix<T>& r) {
                 assert (m == r.n);
                 matrix<T> M(n, r.m);
                 for (int i = 0; i < n; i++) for (int k = 0; k < m; k++)
                         for (int j = 0; j < r.m; j++) {
                                 T \text{ add} = (*this)[i][k] * r[k][j];
#if MODULAR
#warning Usar matrix<11> e soh colocar valores em [0, MOD) na matriz!
                                 M[i][j] += add%MOD;
                                 if (M[i][j] >= MOD) M[i][j] -= MOD;
#else
                                 M[i][j] += add;
#endif
                return M;
        matrix<T> operator^(11 e) {
                matrix<T> M(n, n, true), at = *this;
                 while (e) {
                         if (e&1) M = M*at;
                         e >>= 1:
                         at = at*at;
                 return M;
        void apply_transform(matrix M, 11 e){
                 auto& v = *this;
                 while (e) {
                         if (e\&1) v = M*v;
                         e >>= 1;
                         M = M \star M:
};
```

#### 29.8 matroid

```
// Matroid
// Matroids de Grafo e Particao
// De modo geral, toda Matroid contem um build() linear
// e uma funcao constante oracle()
// oracle(i) responde se o conjunto continua independente
// apos adicao do elemento i
// oracle(i, j) responde se o conjunto continua indepente
// apos trocar o elemento i pelo elemento j
// Intersecao sem peso O(r^2 n)
// em que n eh o tamanho do conjunto e r eh o tamanho da resposta
// Matroid Grafica
// Matroid das florestas de um grafo
// Um conjunto de arestas eh independente se formam uma floresta
// build() : O(n)
// oracle() : 0(1)
struct graphic_matroid {
        int n, m, t;
        vector<array<int, 2>> edges;
        vector<vector<int>> q;
        vector<int> comp, in, out;
        graphic_matroid(int n_, vector<array<int, 2>> edges_)
                : n(n_{-}), m(edges_{-}size()), edges(edges_{-}), g(n), comp(n), in(n),
                    out (n) {}
        void dfs(int u) {
                in[u] = t++;
                for (auto v : g[u]) if (in[v] == -1)
                       comp[v] = comp[u], dfs(v);
                out[u] = t;
        void build(vector<int> I) {
                t = 0;
                for (int u = 0; u < n; u++) q[u].clear(), in[u] = -1;
                for (int e : I) {
                        auto [u, v] = edges[e];
```

```
bool is_ancestor(int u, int v) {
                return in[u] <= in[v] and in[v] < out[u];</pre>
        bool oracle(int e) {
                return comp[edges[e][0]] != comp[edges[e][1]];
        bool oracle(int e, int f) {
                if (oracle(f)) return true;
                int u = edges[e][in[edges[e][0]] < in[edges[e][1]]];</pre>
                return is_ancestor(u, edges[f][0]) != is_ancestor(u, edges[f]
};
// Matroid de particao ou cores
// Um conjunto eh independente se a quantidade de elementos
// de cada cor nao excede a capacidade da cor
// Quando todas as capacidades sao 1, um conjunto eh independente
// se todas as suas cores sao distintas
// build() : O(n)
// oracle() : O(1)
struct partition_matroid {
        vector<int> cap, color, d;
        partition_matroid(vector<int> cap_, vector<int> color_)
                : cap(cap_), color(color_), d(cap.size()) {}
        void build(vector<int> I) {
                fill(d.begin(), d.end(), 0);
                for (int u : I) d[color[u]]++;
        bool oracle(int u) {
                return d[color[u]] < cap[color[u]];</pre>
        bool oracle(int u, int v) {
                return color[u] == color[v] or oracle(v);
};
// Intersecao de matroid sem pesos
// Dadas duas matroids M1 e M2 definidas sobre o mesmo
// conjunto I, retorna o maior subconjunto de I
// que eh independente tanto para M1 quanto para M2
// O(r^2*n)
// Matroid "pesada" deve ser a M2
template<typename Matroid1, typename Matroid2>
vector<int> matroid_intersection(int n, Matroid1 M1, Matroid2 M2) {
        vector<bool> b(n);
        vector<int> I[2];
        bool converged = false;
        while (!converged) {
                I[0].clear(), I[1].clear();
                for (int u = 0; u < n; u++) I[b[u]].push_back(u);
                M1.build(I[1]), M2.build(I[1]);
                vector<bool> target(n), pushed(n);
                queue<int> q;
                for (int u : I[0]) {
                        target[u] = M2.oracle(u);
                        if (M1.oracle(u)) pushed[u] = true, q.push(u);
                vector<int> p(n, -1);
                converged = true;
                while (q.size()) {
                        int u = q.front(); q.pop();
                        if (target[u]) {
                                 converged = false;
                                 for (int v = u; v != -1; v = p[v]) b[v] = !b[v];
                                break:
                        for (int v : I[!b[u]]) if (!pushed[v]) {
                                 if ((b[u] and M1.oracle(u, v)) or (b[v] and M2.
```

g[u].push\_back(v), g[v].push\_back(u);

for (int u = 0; u < n; u++) if (in[u] == -1)

comp[u] = u, dfs(u);

#### 29.9 modular Arithmetic

```
// Aritmetica Modular
// O mod tem q ser primo
template<int p> struct mod int {
        11 expo(11 b, 11 e) {
                11 ret = 1;
                while (e) {
                        if (e % 2) ret = ret * b % p;
                        e /= 2, b = b * b % p;
                return ret;
        11 inv(11 b) { return expo(b, p-2); }
        using m = mod_int;
        int v;
        mod_int() : v(0) {}
        mod_int(ll v_) {
                if (v_ >= p or v_ <= -p) v_ %= p;
                if (v_ < 0) v_ += p;
        m& operator += (const m& a) {
                v += a.v;
                if (v >= p) v -= p;
                return *this;
        m& operator -= (const m& a) {
                v -= a.v;
                if (v < 0) v += p;
                return *this;
        m& operator *=(const m& a) {
                v = v * ll(a.v) % p;
                return *this;
        m& operator /= (const m& a) {
                v = v * inv(a.v) % p;
                return *this;
        m operator -() { return m(-v); }
        m& operator ^=(11 e) {
                if (e < 0) {
                        v = inv(v);
                        e = -e;
                v = expo(v, e);
                // possivel otimizacao:
                // cuidado com 0^0
                // v = \exp((v, e^*(p-1)));
                return *this;
        bool operator == (const m& a) { return v == a.v; }
        bool operator !=(const m& a) { return v != a.v; }
        friend istream& operator >> (istream& in, m& a) {
                11 val; in >> val;
                a = m(val);
                return in;
        friend ostream& operator <<(ostream& out, m a) {</pre>
                return out << a.v;</pre>
        friend m operator + (m a, m b) { return a += b; }
        friend m operator - (m a, m b) { return a -= b;
        friend m operator * (m a, m b) { return a *= b;
        friend m operator / (m a, m b) { return a /= b; }
        friend m operator ^(m a, ll e) { return a ^= e; }
};
typedef mod_int<(int)1e9+7> mint;
```

# 30 ufmg forked/Problemas

#### 30.1 additionChain

```
// Shortest Addition Chain
// Computa o menor numero de adicoes para construir
// cada valor, comecando com 1 (e podendo salvar variaveis)
// Retorna um par com a dp e o pai na arvore
// A arvore en tao que o taminho da raiz (1) ate x
// contem os valores que devem ser criados para gerar x
// A profundidade de x na arvore eh dp[x]
// DP funciona para ateh 300, mas a arvore soh funciona
// para ateh 148
// recuperacao certa soh ateh 148 (erra para 149, 233, 298)
pair<vector<int>, vector<int>> addition_chain() {
        int MAX = 301;
         vector<int> dp(MAX), p(MAX);
         for (int n = 2; n < MAX; n++) {
                  pair<int, int> val = {INF, -1};
                 for (int i = 1; i < n; i++) for (int j = i; j; j = p[j])
    if (j == n-i) val = min(val, pair(dp[i]+1, i));</pre>
                 tie(dp[n], p[n]) = val;
                 if (n == 9) p[n] = 8;
                 if (n == 149 \text{ or } n == 233) \text{ dp}[n] --;
         return {dp, p};
```

### 30.2 angleRange

```
// Angle Range Intersection
// Computa intersecao de angulos
// Os angulos (arcos) precisam ter comprimeiro < pi
// (caso contrario a intersecao eh estranha)
// Tudo 0(1)
struct angle_range {
        static constexpr ld ALL = 1e9, NIL = -1e9;
        angle_range() : 1(ALL), r(ALL) {}
        angle_range(ld l_, ld r_) : l(l_), r(r_) { fix(l), fix(r); }
        void fix(ld& theta) {
                if (theta == ALL or theta == NIL) return;
                if (theta > 2*pi) theta -= 2*pi;
                if (theta < 0) theta += 2*pi;</pre>
        bool empty() { return 1 == NIL; }
        bool contains(ld q) {
                fix(q);
                if (1 == ALL) return true;
                if (1 == NIL) return false;
                if (1 < r) return 1 < q and q < r;
                return q > 1 or q < r;
        friend angle_range operator &(angle_range p, angle_range q) {
                if (p.1 == ALL or q.1 == NIL) return q;
                if (q.1 == ALL or p.1 == NIL) return p;
                if (p.1 > p.r) and q.1 > q.r) return \{\max(p.1, q.1), \min(p.r, q.1)\}
                r)};
if (q.1 > q.r) swap(p.1, q.1), swap(p.r, q.r);
                if (p.1 > p.r) {
                        if (q.r > p.l) return {max(q.l, p.l) , q.r};
                        else if (q.1 < p.r) return {q.1, min(q.r, p.r)};</pre>
                        return {NIL, NIL};
                if (max(p.l, q.l) > min(p.r, q.r)) return {NIL, NIL};
                return {max(p.l, q.l), min(p.r, q.r)};
```

## 30.3 areaHistograma

};

```
// Area Maxima de Histograma
// Assume que todas as barras tem largura 1,
// e altura dada no vetor v
// O(n)
11 area(vector<int> v) {
        11 ret = 0;
        stack<int> s;
        // valores iniciais pra dar tudo certo
        v.insert(v.begin(), -1);
        v.insert(v.end(), -1);
        s.push(0);
        for(int i = 0; i < (int) v.size(); i++) {</pre>
                while (v[s.top()] > v[i]) {
                         11 h = v[s.top()]; s.pop();
                        ret = max(ret, h * (i - s.top() - 1));
                s.push(i);
        return ret;
```

### 30.4 areaUniaoRetangulo

```
// Area da Uniao de Retangulos
// O(n log(n))
// 5d8d2f
namespace seg {
        pair<int, 11> seg[4*MAX];
        11 lazy[4*MAX], *v;
        int n;
        pair<int, 11> merge(pair<int, 11> 1, pair<int, 11> r) {
                if (1.second == r.second) return {1.first+r.first, 1.second};
                else if (1.second < r.second) return 1;</pre>
                else return r;
        pair<int, 11> build(int p=1, int 1=0, int r=n-1) {
                lazy[p] = 0;
                if (1 == r) return seg[p] = {1, v[1]};
                int m = (1+r)/2;
                return seq[p] = merge(build(2*p, 1, m), build(2*p+1, m+1, r));
        void build(int n2, l1* v2) {
                n = n2, v = v2;
                build();
        void prop(int p, int 1, int r)
                seq[p].second += lazy[p];
                if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
                lazy[p] = 0;
        pair<int, 11> query(int a, int b, int p=1, int 1=0, int r=n-1) {
                prop(p, 1, r);
                if (a <= 1 and r <= b) return seg[p];</pre>
                if (b < 1 or r < a) return {0, LINF};</pre>
                int m = (1+r)/2;
                return merge (query (a, b, 2*p, 1, m), query (a, b, 2*p+1, m+1, r))
        pair<int, 11> update(int a, int b, int x, int p=1, int 1=0, int r=n-1) {
```

```
prop(p, 1, r);
if (a <= 1 and r <= b) {</pre>
                          lazy[p] += x;
                          prop(p, 1, r);
                          return seg[p];
                 if (b < l or r < a) return seg[p];</pre>
                 int m = (1+r)/2;
                 return seg[p] = merge(update(a, b, x, 2*p, 1, m),
                                   update(a, b, x, 2*p+1, m+1, r));
};
11 seg_vec[MAX];
11 area_sq(vector<pair<pair<int, int>, pair<int, int>>> &sq) {
         vector<pair<int, int>, pair<int, int>>> up;
         for (auto it : sq) {
                 int x1, y1, x2, y2;
                 tie(x1, y1) = it.first;
                 tie(x2, y2) = it.second;
                 up.push_back(\{\{x1+1, 1\}, \{y1, y2\}\}\});
                 up.push_back(\{\{x2+1, -1\}, \{y1, y2\}\}\);
         sort(up.begin(), up.end());
         memset(seg_vec, 0, sizeof seg_vec);
         11 \text{ H\_MAX} = \text{MAX};
         seq::build(H_MAX-1, seg_vec);
         auto it = up.begin();
         11 \text{ ans} = 0;
         while (it != up.end()){
                  11 L = (*it).first.first;
                 while (it != up.end() && (*it).first.first == L) {
                          int x, inc, y1, y2;
                          tie(x, inc) = it->first;
tie(y1, y2) = it->second;
                          seg::update(y1+1, y2, inc);
                          it++;
                 if (it == up.end()) break;
                 11 R = (*it) .first.first;
                 11 W = R-L;
                 auto jt = seg::query(0, H_MAX-1);
                  11 H = H_MAX - 1;
                 if (jt.second == 0) H -= jt.first;
                 ans += W*H;
         return ans;
```

#### 30.5 binomial

```
// Binomial modular
// Computa C(n, k) mod m em O(m + log(m) log(n))
// = O(rapido)

11 divi[MAX];

11 expo(ll a, ll b, ll m) {
    if (!b) return 1;
        ll ans = expo(a*a*m, b/2, m);
    if (b*2) ans *= a;
    return ans*m;
}

11 inv(ll a, ll b) {
    return 1<a ? b - inv(b*a,a)*b/a : 1;
}

template<typename T> tuple<T, T, T> ext_gcd(T a, T b) {
    if (!a) return {b, 0, 1};
    auto [g, x, y] = ext_gcd(b*a, a);
    return {g, y - b/a*x, x};
}
```

```
template<typename T = 11> struct crt {
        T a, m;
        crt() : a(0), m(1) {}
        crt(T a_, T m_) : a(a_), m(m_) {}
        crt operator * (crt C) {
                auto [g, x, y] = ext_gcd(m, C.m);
if ((a - C.a) % g) a = -1;
                 if (a == -1 or C.a == -1) return crt(-1, 0);
                 T lcm = m/g*C.m;
                 T \text{ ans} = a + (x*(C.a-a)/g % (C.m/g))*m;
                 return crt((ans % lcm + lcm) % lcm, lcm);
};
pair<11, 11> divide_show(11 n, int p, int k, int pak) {
        if (n == 0) return {0, 1};
        11 blocos = n/pak, falta = n%pak;
        11 periodo = divi[pak], resto = divi[falta];
        11 r = expo(periodo, blocos, pak)*resto%pak;
        auto rec = divide_show(n/p, p, k, pak);
        11 y = n/p + rec.first;
        r = r*rec.second % pak;
        return {y, r};
11 solve_pak(11 n, 11 x, int p, int k, int pak) {
        divi[0] = 1;
        for (int i = 1; i <= pak; i++) {</pre>
                divi[i] = divi[i-1];
                 if (i%p) divi[i] = divi[i] * i % pak;
        auto dn = divide_show(n, p, k, pak), dx = divide_show(x, p, k, pak),
                 dnx = \overline{divide\_show}(n-x, p, k, pak);
        11 y = dn.first-dx.first-dnx.first, r =
                 (dn.second*inv(dx.second, pak)%pak)*inv(dnx.second, pak)%pak;
        return expo(p, y, pak) * r % pak;
11 solve(ll n, ll x, int mod) {
        vector<pair<int, int>> f;
        int mod2 = mod;
        for (int i = 2; i*i <= mod2; i++) if (mod2%i==0) {</pre>
                int c = 0;
                while (mod2\%i==0) mod2 /= i, c++;
                f.push_back({i, c});
        if (mod2 > 1) f.push_back({mod2, 1});
        crt ans(0, 1);
        for (int i = 0; i < f.size(); i++) {</pre>
                int pak = 1;
                 for (int j = 0; j < f[i].second; j++) pak *= f[i].first;
                ans = ans * crt(solve_pak(n, x, f[i].first, f[i].second, pak),
        return ans.a;
```

### 30.6 closestPairOfPoints

```
// Closest pair of points
//
// O(nlogn)

pair<pt, pt> closest_pair_of_points(vector<pt> v) {
    int n = v.size();
    sort(v.begin(), v.end());
    for (int i = 1; i < n; i++) if (v[i] == v[i-1]) return {v[i-1], v[i]};
    auto cmp_y = [&] (const pt &l, const pt &r) {
        if (l.y != r.y) return l.y < r.y;
        return l.x < r.x;
    };</pre>
```

```
set<pt, decltype(cmp_y) > s(cmp_y);
int 1 = 0, r = -1;
11 d2_min = numeric_limits<11>::max();
pt pl, pr;
const int magic = 5;
while (r+1 < n) {
        auto it = s.insert(v[++r]).first;
        int cnt = magic/2;
        while (cnt-- and it != s.begin()) it--;
        cnt = 0;
        while (cnt++ < magic and it != s.end()) {</pre>
                if (!((*it) == v[r])) {
                         11 d2 = dist2(*it, v[r]);
                         if (d2_min > d2) {
                                 d2_min = d2;
                                 p1 = *it;
                                 pr = v[r];
                it++;
        while (1 < r \text{ and } sq(v[1].x-v[r].x) > d2_min) s.erase(v[1++]);
return {pl, pr};
```

### 30.7 conectividadeDinamica

```
// Conectividade Dinamica DC
// Offline com Divide and Conquer e
// DSU com rollback
// O(n log^2(n))
typedef pair<int, int> T;
namespace data {
        int n, ans;
        int p[MAX], sz[MAX];
        stack<int> S;
        void build(int n2) {
                n = n2:
                for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
                ans = n;
        int find(int k) {
                while (p[k] != k) k = p[k];
                return k;
        void add(T x) {
                int a = x.first, b = x.second;
                a = find(a), b = find(b);
                if (a == b) return S.push(-1);
                ans--:
                if (sz[a] > sz[b]) swap(a, b);
                S.push(a);
                sz[b] += sz[a];
                p[a] = b;
        int query() {
                return ans;
        void rollback() {
                int u = S.top(); S.pop();
                if (u == -1) return;
                sz[p[u]] = sz[u];
                p[u] = u;
                ans++;
int ponta[MAX]; // outra ponta do intervalo ou -1 se for query
int ans[MAX], n, q;
T qu[MAX];
```

```
void solve(int 1 = 0, int r = q-1) {
    if (1 >= r) {
        ans[1] = data::query(); // agora a estrutura ta certa
        return;
    }
    int m = (l+r)/2, qnt = 1;
    for (int i = m+1; i <= r; i++) if (ponta[i]+1 and ponta[i] < l)
        data::add(qu[i]), qnt++;
    solve(1, m);
    while (--qnt) data::rollback();
    for (int i = l; i <= m; i++) if (ponta[i]+1 and ponta[i] > r)
        data::add(qu[i]), qnt++;
    solve(m+1, r);
    while (qnt--) data::rollback();
}
```

#### 30.8 conectividadeDinamica2

```
// Conectividade Dinamica LCT
// Offline com link-cut trees
// O(n log(n))
namespace lct {
        struct node {
                 int p, ch[2];
                int val, sub;
                bool rev;
                node() {}
                node(int \ v) : p(-1), \ val(v), \ sub(v), \ rev(0) \{ ch[0] = ch[1] = ch[1] = ch[1] = ch[1] = ch[1]
                     -1; }
        };
        node t[2*MAX]; // MAXN + MAXQ
        map<pair<int, int>, int> aresta;
        int sz;
        void prop(int x) {
                if (t[x].rev) {
                         swap(t[x].ch[0], t[x].ch[1]);
                         if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
                         if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
                t[x].rev = 0;
        void update(int x) {
                t[x].sub = t[x].val;
                for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
                         prop(t[x].ch[i]);
                         t[x].sub = min(t[x].sub, t[t[x].ch[i]].sub);
        bool is_root(int x) {
                return t[x].p == -1 or (t[t[x].p].ch[0] != x and t[t[x].p].ch[1]
                      ! = x);
        void rotate(int x)
                int p = t[x].p, pp = t[p].p;
                if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
                bool d = t[p].ch[0] == x;
                if (t[p].ch[!d]+1) t[t[p].ch[!d]).p = p;
                t[x].p = pp, t[p].p = x;
                update(p), update(x);
        int splay(int x) {
                while (!is_root(x)) {
                         int p = t[x].p, pp = t[p].p;
                         if (!is_root(p)) prop(pp);
                         prop(p), prop(x);
                         if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
                             == x) ? x : p);
                         rotate(x);
                return prop(x), x;
```

```
int access(int v) {
                int last = -1;
                 for (int w = v; w+1; update(last = w), splay(v), w = t[v].p)
                         splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
                 return last;
        void make_tree(int v, int w=INF) { t[v] = node(w); }
        bool conn(int v, int w) {
                access(v), access(w);
                 return v == w ? true : t[v].p != -1;
        void rootify(int v) {
                access(v);
t[v].rev ^= 1;
        int query(int v, int w) {
                 rootify(w), access(v);
                 return t[v].sub;
        void link (int v, int w) {
                 rootify(w);
                 t[w].p = v;
        void link(int v, int w, int x) { // v--w com peso x
                 int id = MAX + sz++;
                 aresta[make_pair(v, w)] = id;
                 make_tree(id, x);
                 link_(v, id), link_(id, w);
        void cut_(int v, int w) {
                 rootify(w), access(v);
                 t[v].ch[0] = t[t[v].ch[0]].p = -1;
        void cut(int v, int w) {
                 int id = aresta[make_pair(v, w)];
                cut_(v, id), cut_(id, w);
void dyn_conn() {
        int n, q; cin >> n >> q;
        vector<int> p(2*q, -1); // outra ponta do intervalo
for (int i = 0; i < n; i++) lct::make_tree(i);</pre>
        vector<pair<int, int>> qu(q);
        map<pair<int, int>, int> m;
        for (int i = 0; i < q; i++) {
                 char c; cin >> c;
                 if (c == '?') continue;
                 int a, b; cin >> a >> b; a--, b--;
                 if (a > b) swap(a, b);
                 qu[i] = \{a, b\};
                 if (c == '+') {
                         p[i] = i+q, p[i+q] = i;
                         m[make_pair(a, b)] = i;
                 } else {
                         int j = m[make_pair(a, b)];
                         p[i] = j, p[j] = i;
        int ans = n;
        for (int i = 0; i < q; i++) {
                if (p[i] == -1) {
                         cout << ans << endl; // numero de comp conexos</pre>
                         continue:
                 int a = qu[i].first, b = qu[i].second;
                 if (p[i] > i) { // +
                         if (lct::conn(a, b)) {
                                  int mi = lct::query(a, b);
                                  if (p[i] < mi) {
    p[p[i]] = p[i];</pre>
                                          continue;
                                  lct::cut(qu[p[mi]].first, qu[p[mi]].second), ans
                                       ++;
                                  p[mi] = mi;
                         lct::link(a, b, p[i]), ans--;
                 } else if (p[i] != i) lct::cut(a, b), ans++; // -
```

# 30.9 deBrujin

```
// Sequencia de de Brujin
// Se passar sem o terceiro parametro, gera um vetor com valores
// em [0, k) de tamanho k^n de forma que todos os subarrays ciclicos
// de tamanho n ocorrem exatamente uma vez
// Se passar com um limite lim, gera o menor vetor com valores
// em [0, k) que possui lim subarrays de tamanho n distintos
// (assume que lim <= k^n)
// Linear no tamanho da resposta
vector<int> de_brujin(int n, int k, int lim = INF) {
        if (k == 1) return vector<int>(lim == INF ? 1 : n, 0);
        vector<int> 1 = {0}, ret; // 1 eh lyndon word
        while (true) {
                if (1.size() == 0) {
    if (lim == INF) break;
                        1.push_back(0);
                if (n % 1.size() == 0) for (int i : 1) {
                        ret.push_back(i);
                        if (ret.size() == n+lim-1) return ret;
                int p = 1.size();
                while (l.size() < n) l.push_back(l[l.size()%p]);</pre>
                while (1.size() and 1.back() == k-1) 1.pop_back();
                if (1.size()) 1.back()++;
        return ret;
```

### 30.10 delaunay

```
// Triangulação de Delaunav
// Computa a triangulacao de Delaunay, o dual
// do diagrama de Voronoi (a menos de casos degenerados)
// Retorna um grafo indexado pelos indices dos pontos, e as arestas
// sao as arestas da triangulação
// As arestas partindo de um vertice ja vem ordenadas por angulo,
// ou seja, se o vertice v nao esta no convex hull, (v, v_i, v_{i+1})
// eh um triangulo da triangulacao, em que v_i eh o i-esimo vizinho
// Usa o alg d&c, precisa representar MAX_COOR^4, por isso __int128
// pra aguentar valores ateh 1e9
// Propriedades:
// 1 - 0 grafo tem no max 3n-6 arestas
// 2 - Para todo triangulo, a circunf. que passa pelos 3 pontos
     nao contem estritamente nenhum ponto
// 3 - A MST euclidiana eh subgrafo desse grafo
// 4 - Cada ponto eh vizinho do ponto mais proximo dele
// O(n log n)
typedef struct QuadEdge* Q;
struct QuadEdge {
       int id;
        pt o:
       Q rot, nxt;
       bool used;
       QuadEdge(int id_ = -1, pt o_ = pt(INF, INF)) :
                id(id_), o(o_), rot(nullptr), nxt(nullptr), used(false) {}
       Q rev() const { return rot->rot; }
       0 next() const { return nxt; }
       Q prev() const { return rot->next()->rot; }
```

```
pt dest() const { return rev()->o; }
};
Q edge(pt from, pt to, int id_from, int id_to) {
        Q e1 = new QuadEdge(id_from, from);
        Q e2 = new QuadEdge(id_to, to);
        Q e3 = new QuadEdge;
        Q e4 = new QuadEdge;
        tie(e1->rot, e2->rot, e3->rot, e4->rot) = \{e3, e4, e2, e1\};
        tie(e1->nxt, e2->nxt, e3->nxt, e4->nxt) = \{e1, e2, e4, e3\};
void splice(Q a, Q b) {
        swap(a->nxt->rot->nxt, b->nxt->rot->nxt);
        swap(a->nxt, b->nxt);
void del_edge(Q& e, Q ne) { // delete e and assign e <- ne</pre>
        splice(e, e->prev());
        splice(e->rev(), e->rev()->prev());
        delete e->rev()->rot, delete e->rev();
        delete e->rot; delete e;
        e = ne;
Q conn(Q a, Q b) {
        Q e = edge(a->dest(), b->o, a->rev()->id, b->id);
        splice(e, a->rev()->prev());
        splice(e->rev(), b);
        return e;
bool in_c(pt a, pt b, pt c, pt p) { // p ta na circunf. (a, b, c) ?
        __int128 p2 = p*p, A = a*a - p2, B = b*b - p2, C = c*c - p2;
return sarea2(p, a, b) * C + sarea2(p, b, c) * A + sarea2(p, c, a) * B >
pair<Q, Q> build_tr(vector<pt>& p, int 1, int r) {
        if (r-1+1 \le 3) {
                 Q = edge(p[1], p[1+1], 1, 1+1), b = edge(p[1+1], p[r], 1+1, r)
                 if (r-1+1 == 2) return \{a, a->rev()\};
                 splice(a->rev(), b);
                 11 \text{ ar} = \text{sarea2}(p[1], p[1+1], p[r]);
                 Q c = ar ? conn(b, a) : 0;
                 if (ar >= 0) return {a, b->rev()};
                 return {c->rev(), c};
        int m = (1+r)/2;
        auto [la, ra] = build_tr(p, l, m);
        auto [lb, rb] = build_tr(p, m+1, r);
        while (true) {
                 if (ccw(lb->o, ra->o, ra->dest())) ra = ra->rev()->prev();
                 else if (ccw(lb->o, ra->o, lb->dest())) lb = lb->rev()->next();
                 else break:
        0 b = conn(lb->rev(), ra);
        auto valid = [&](Q e) { return ccw(e->dest(), b->o); };
        if (ra->o == la->o) la = b->rev();
        if (lb->o == rb->o) rb = b;
        while (true) {
                 Q L = b->rev()->next();
                 if (valid(L)) while (in_c(b->dest(), b->o, L->dest(), L->next()
                      ->dest()))
                         del_edge(L, L->next());
                 OR = b \rightarrow prev();
                 if (valid(R)) while (in_c(b->dest(), b->o, R->dest(), R->prev()
                      ->dest()))
                         del_edge(R, R->prev());
                 if (!valid(L) and !valid(R)) break;
                 if (!valid(L) or (valid(R) and in_c(L->dest(), L->o, R->o, R->
                      dest())))
                         b = conn(R, b\rightarrow rev());
                 else b = conn(b->rev(), L->rev());
        return {la, rb};
```

```
vector<vector<int>> delaunay(vector<pt> v) {
        int n = v.size();
        auto tmp = v;
        vector<int> idx(n);
        iota(idx.begin(), idx.end(), 0);
        sort(idx.begin(), idx.end(), [&](int l, int r) { return v[l] < v[r]; });
        for (int i = 0; i < n; i++) v[i] = tmp[idx[i]];</pre>
        assert(unique(v.begin(), v.end()) == v.end());
        vector<vector<int>> g(n);
        bool col = true;
        for (int i = 2; i < n; i++) if (sarea2(v[i], v[i-1], v[i-2])) col =
             false;
        if (col) {
                 for (int i = 1; i < n; i++)
                         g[idx[i-1]].push_back(idx[i]), g[idx[i]].push_back(idx[i
                              -11);
                return q;
        0 = build tr(v, 0, n-1).first;
        vector<Q> edg = {e};
        for (int i = 0; i < edg.size(); e = edg[i++]) {</pre>
                for (Q at = e; !at->used; at = at->next()) {
                         at->used = true;
                         g[idx[at->id]].push_back(idx[at->rev()->id]);
edg.push_back(at->rev());
        return q;
```

#### 30.11 distinct

```
// Distinct Range Query
// build - O(n (log n + log(sigma)))
// query - O(log(sigma))
namespace perseq { };
int qt[MAX];
void build(vector<int>& v) {
        int n = v.size();
        perseq::build(n);
        map<int, int> last;
        int at = 0;
        for (int i = 0; i < n; i++) {</pre>
                if (last.count(v[i])) {
                        perseg::update(last[v[i]], -1);
                         at++;
                perseg::update(i, 1);
                qt[i] = ++at;
                last[v[i]] = i;
int query(int 1, int r) {
        return perseg::query(l, r, qt[r]);
```

# 30.12 distinctUpdate

```
// Distinct Range Query com Update
//
// build - O(n log(n))
// query - O(log^2(n))
// update - O(log^2(n))

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
```

```
using namespace __gnu_pbds;
template <class T>
         using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
         tree_order_statistics_node_update>;
int v[MAX], n, nxt[MAX], prv[MAX];
map<int, set<int> > ocor;
namespace bit {
        ord_set<pair<int, int>> bit[MAX];
         void build() {
                 for (int i = 1; i <= n; i++) bit[i].insert({nxt[i-1], i-1});</pre>
                 for (int i = 1; i <= n; i++) {</pre>
                          int j = i + (i\&-i);
                          if (j <= n) for (auto x : bit[i]) bit[j].insert(x);</pre>
         int pref(int p, int x) {
                 int ret = 0;
                 for (; p; p \rightarrow p\&-p) ret += bit[p].order_of_key({x, -INF});
                 return ret;
         int query(int 1, int r, int x) {
                 return pref(r+1, x) - pref(l, x);
         void update(int p, int x) {
                 int p2 = p;
                 for (p++; p \le n; p += p\&-p) {
                          bit[p] erase({nxt[p2], p2});
                          bit[p].insert({x, p2});
void build() {
        for (int i = 0; i < n; i++) nxt[i] = INF;
for (int i = 0; i < n; i++) prv[i] = -INF;</pre>
         vector<pair<int, int>> t;
         for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
         sort(t.begin(), t.end());
        for (int i = 0; i < n; i++) {
    if (i and t[i].first == t[i-1].first)</pre>
                          prv[t[i].second] = t[i-1].second;
                 if (i+1 < n and t[i].first == t[i+1].first)</pre>
                          nxt[t[i].second] = t[i+1].second;
         for (int i = 0; i < n; i++) ocor[v[i]].insert(i);</pre>
        bit::build();
void muda(int p, int x) {
        bit::update(p, x);
        nxt[p] = x;
int query(int a, int b) {
         return b-a+1 - bit::query(a, b, b+1);
void update(int p, int x) { // mudar valor na pos. p para x
         if (prv[p] > -INF) muda(prv[p], nxt[p]);
         if (nxt[p] < INF) prv[nxt[p]] = prv[p];</pre>
         ocor[v[p]].erase(p);
         if (!ocor[x].size()) {
                 muda(p, INF);
                 prv[p] = -INF;
         } else if (*ocor[x].rbegin() < p) {</pre>
                 int i = *ocor[x].rbegin();
                 prv[p] = i;
                 muda(p, INF);
                 muda(i, p);
         } else {
                 int i = *ocor[x].lower_bound(p);
                 if (prv[i] > -INF) {
                          muda(prv[i], p);
```

```
prv[p] = prv[i];
} else prv[p] = -INF;
prv[i] = p;
muda(p, i);
}
v[p] = x; ocor[x].insert(p);
}
```

#### 30.13 dominacao3d

```
// DP de Dominacao 3D
// Computa para todo ponto i,
// dp[i] = 1 + max_{i} dominado por i dp[i]
// em que ser dominado eh ter as 3 coordenadas menores
// Da pra adaptar facil para outras dps
// O(n log^2 n), O(n) de memoria
void lis2d(vector<vector<tuple<int, int, int>>>& v, vector<int>& dp, int 1, int
    r) {
        if (1 == r) {
                for (int i = 0; i < v[1].size(); i++) {</pre>
                        int ii = get<2>(v[1][i]);
                        dp[ii] = max(dp[ii], 1);
                return;
        int m = (1+r)/2;
        lis2d(v, dp, 1, m);
        vector<tuple<int, int, int>> vv[2];
        vector<int> Z;
        for (int i = 1; i <= r; i++) for (auto it : v[i]) {
                vv[i > m].push_back(it);
                Z.push_back(get<1>(it));
        sort(vv[0].begin(), vv[0].end());
        sort(vv[1].begin(), vv[1].end());
        sort(Z.begin(), Z.end());
        auto get_z = [&](int z) { return lower_bound(Z.begin(), Z.end(), z) - Z.
             begin(); };
        vector<int> bit(Z.size());
        int i = 0;
        for (auto [y, z, id] : vv[1]) {
                while (i < vv[0].size() and qet<0>(vv[0][i]) < y) {
                        auto [y2, z2, id2] = vv[0][i++];
                        for (int p = get_z(z2)+1; p <= Z.size(); p += p&-p)</pre>
                                bit[p-1] = max(bit[p-1], dp[id2]);
                int q = 0;
                for (int p = get_z(z); p; p -= p\&-p) q = max(q, bit[p-1]);
                dp[id] = max(dp[id], q + 1);
        lis2d(v, dp, m+1, r);
vector<int> solve(vector<tuple<int, int, int>> v) {
        int n = v.size();
        vector<tuple<int, int, int, int>> vv;
        for (int i = 0; i < n; i++)
                auto [x, y, z] = v[i];
                vv.emplace_back(x, y, z, i);
        sort(vv.begin(), vv.end());
        vector<vector<tuple<int, int, int>>> V;
        for (int i = 0; i < n; i++) {
                int j = i;
                V emplace_back();
                while (j < n \text{ and } get<0>(vv[j]) == get<0>(vv[i])) {
                        auto [x, y, z, id] = vv[j++];
                        V.back().emplace_back(y, z, id);
                i = j-1;
```

```
}
vector<int> dp(n);
lis2d(V, dp, 0, V.size()-1);
return dp;
```

### 30.14 dominatorPoints

```
// Dominator Points
// Se um ponto A tem ambas as coordenadas >= B, dizemos
// que A domina B
// is_dominated(p) fala se existe algum ponto no conjunto
// que domina p
// insert(p) insere p no conjunto
// (se p for dominado por alguem, nao vai inserir)
// o multiset 'quina' guarda informacao sobre os pontos
// nao dominados por um elemento do conjunto que nao dominam
// outro ponto não dominado por um elemento do conjunto
// No caso, armazena os valores de x+y esses pontos
// Complexidades:
// is dominated - O(log(n))
// insert - O(log(n)) amortizado
// query - O(1)
struct dominator_points {
        set<pair<int, int>> se;
        multiset<int> quina;
        bool is_dominated(pair<int, int> p) {
                auto it = se.lower_bound(p);
                if (it == se.end()) return 0;
                return it->second >= p.second;
        void mid(pair<int, int> a, pair<int, int> b, bool rem) {
                pair<int, int> m = {a.first+1, b.second+1};
                int val = m.first + m.second;
                if (!rem) quina.insert(val);
                else quina.erase(quina.find(val));
        bool insert(pair<int, int> p) {
                if (is_dominated(p)) return 0;
                auto it = se.lower_bound(p);
                if (it != se.begin() and it != se.end())
                        mid(*prev(it), *it, 1);
                while (it != se.begin()) {
                        if (it->second > p.second) break;
                        if (it != se begin()) mid(*prev(it), *it, 1);
                        it = se.erase(it);
                it = se.insert(p).first;
                if (it != se.begin()) mid(*prev(it), *it, 0);
                if (next(it) != se.end()) mid(*it, *next(it), 0);
                return 1;
        int query()
                if (!quina.size()) return INF;
                return *quina.begin();
};
```

# 30.15 dynamicHull

```
// Convex Hull Dinamico
//
insert - O(log n) amortizado
// is_inside - O(log n)

struct upper {
    set<pt> se;
```

```
set<pt>::iterator it;
        int is_under(pt p) { // 1 -> inside ; 2 -> border
                it = se.lower_bound(p);
                if (it == se.end()) return 0;
                if (it == se.begin()) return p == *it ? 2 : 0;
                if (ccw(p, *it, *prev(it))) return 1;
                return ccw(p, *prev(it), *it) ? 0 : 2;
        void insert(pt p) {
                if (is_under(p)) return;
                if (it != se.end()) while (next(it) != se.end() and !ccw(*next()
                      it), *it, p))
                         it = se.erase(it);
                 if (it != se.begin()) while (--it != se.begin() and !ccw(p, *it,
                       *prev(it)))
                         it = se.erase(it);
                se.insert(p);
};
struct dyn_hull {
        upper U, L;
        int is_inside(pt p) {
                int u = U.is_under(p), 1 = L.is_under({-p.x, -p.y});
if (!u or !1) return 0;
                return max(u, 1);
        void insert(pt p) {
                U.insert(p);
                L.insert(\{-p.x, -p.y\});
        int size() {
                int ans = U.se.size() + L.se.size();
                return ans <= 2 ? ans/2 : ans-2;
};
```

# 30.16 graphTriangles

```
// Triangulos em Grafos
// get_triangles(i) encontra todos os triangulos ijk no grafo
// Custo nas arestas
// retorna {custo do triangulo, {j, k}}
// O(m sqrt(m) log(n)) se chamar para todos os vertices
vector<pair<int, int>> g[MAX]; // {para, peso}
#warning o 'g' deve estar ordenado
vector<pair<int, pair<int, int>>> get_triangles(int i) {
       vector<pair<int, pair<int, int>>> tri;
        for (pair<int, int> j : g[i]) {
                int a = i, b = j.first;
                if (g[a].size() > g[b].size()) swap(a, b);
                for (pair<int, int> c : g[a]) if (c.first != b and c.first > j.
                     first) {
                        auto it = lower_bound(g[b].begin(), g[b].end(),
                            make_pair(c.first, -INF));
                        if (it == g[b].end() or it->first != c.first) continue;
                        tri.push_back({j.second+c.second+it->second, {a == i ? b
                             : a, c.first}});
        return tri;
```

# 30.17 grayCode

```
// Gray Code
//
// Gera uma permutacao de 0 a 2^n-1, de forma que
// duas posicoes adjacentes diferem em exatamente 1 bit
//
// O(2^n)
vector<int> gray_code(int n) {
    vector<int> ret(1<<n);
    for (int i = 0; i < (1<<n); i++) ret[i] = i^(i>>1);
    return ret;
}
```

## 30.18 halfPlaneIntersection

```
// Half-plane intersection
// Computa a regiao convexa formada pela intersecao de n half-planes
// Cada half-plane eh identificado por uma reta e a regiao ccw a ela
// O(n log n)
vector<pt> hp_intersection(vector<line> &v) {
       deque < pt > dq = {\{INF, INF\}, {-INF, INF}, {-INF, -INF}\}; }
#warning considerar trocar por compare_angle
        sort(v.begin(), v.end(), [&](line r, line s) { return angle(r.q-r.p) <</pre>
            angle(s.q-s.p); });
        for(int i = 0; i < v.size() and dq.size() > 1; i++) {
                pt p1 = dq.front(), p2 = dq.back();
                while (dq.size() and !ccw(v[i].p, v[i].q, dq.back()))
                        p1 = dq.back(), dq.pop_back();
                while (dq.size() and !ccw(v[i].p, v[i].q, dq.front()))
                        p2 = dq.front(), dq.pop_front();
                if (!dq.size()) break;
                if (p1 == dq.front() and p2 == dq.back()) continue;
                dq.push_back(inter(v[i], line(dq.back(), p1)));
                dq.push_front(inter(v[i], line(dq.front(), p2)));
                if (dq.size() > 1 and dq.back() == dq.front()) dq.pop_back();
        return vector<pt>(dq.begin(), dq.end());
```

# 30.19 heapSort

# 30.20 hungarian

```
// Hungaro
```

```
// Resolve o problema de assignment (matriz n x n)
// Colocar os valores da matriz em 'a' (pode < 0)
// assignment() retorna um par com o valor do
// assignment minimo, e a coluna escolhida por cada linha
// O(n^3)
template<typename T> struct hungarian {
        int n;
         vector<vector<T>> a;
        vector<T> u, v;
        vector<int> p, way;
        T inf:
         hungarian (int n_{-}): n(n_{-}), u(n+1), v(n+1), p(n+1), way (n+1) {
                 a = vector<vector<T>>(n, vector<T>(n));
                 inf = numeric limits<T>::max();
        pair<T, vector<int>> assignment() {
                 for (int i = 1; i <= n; i++) {
                          p[0] = i;
                          int j0 = 0;
                          vector<T> minv(n+1, inf);
                          vector<int> used(n+1, 0);
                                   used[j0] = true;
                                   int i0 = p[j0], j1 = -1;
                                   T delta = inf;
                                   for (int j = 1; j <= n; j++) if (!used[j]) {
   T cur = a[i0-1][j-1] - u[i0] - v[j];</pre>
                                            if (cur < minv[j]) minv[j] = cur, way[j]</pre>
                                                  = j0;
                                            if (minv[j] < delta) delta = minv[j], j1</pre>
                                                  = j;
                                   for (int j = 0; j <= n; j++)
    if (used[j]) u[p[j]] += delta, v[j] -=</pre>
                                                 delta;
                                            else minv[j] -= delta;
                                   i0 = i1;
                           } while (p[j0] != 0);
                          do {
                                   int j1 = way[j0];
                                   p[j0] = p[j1];
                                    j0 = j1;
                           } while (j0);
                  vector<int> ans(n);
                  for (int j = 1; j \le n; j++) ans[p[j]-1] = j-1;
                 return make_pair(-v[0], ans);
};
```

# 30.21 intervalGraphColoring

```
// Coloracao de Grafo de Intervalo
// Colore os intervalos com o numero minimo
// de cores de tal forma que dois intervalos
// que se interceptam tem cores diferentes
// As cores vao de 1 ate n
// O(n log(n))
vector<int> coloring(vector<pair<int, int>>& v) {
        int n = v.size();
        vector<pair<int, pair<int, int>>> ev;
        for (int i = 0; i < n; i++) {
                ev.push_back({v[i].first, {1, i}});
                ev.push_back({v[i].second, {0, i}});
        sort(ev.begin(), ev.end());
        vector<int> ans(n), avl(n);
        for (int i = 0; i < n; i++) avl.push_back(n-i);</pre>
        for (auto i : ev) {
```

## 30.22 intervalGraphIndSet

```
// Conj. Indep. Maximo com Peso em Grafo de Intervalo
// Retorna os indices ordenados dos intervalos selecionados
// Se tiver empate, retorna o que minimiza o comprimento total
// O(n log(n))
vector<int> ind_set(vector<tuple<int, int, int>>& v) {
        vector<tuple<int, int, int>> w;
        for (int i = 0; i < v.size(); i++) {</pre>
                w.push_back(tuple(get<0>(v[i]), 0, i));
                w.push_back(tuple(get<1>(v[i]), 1, i));
        sort(w.begin(), w.end());
        vector<int> nxt(v.size());
        vector<pair<11, int>> dp(v.size());
        int last = -1;
        for (auto [fim, t, i] : w) {
                if (t == 0) {
                        nxt[i] = last;
                        continue;
                if (last != -1) dp[i] = max(dp[i], dp[last]);
                pair<11, int> pega = {get<2>(v[i]), -(get<1>(v[i]) - get<0>(v[i
                     ]) + 1)};
                if (nxt[i] != -1) pega.first += dp[nxt[i]].first, pega.second +=
                     dp[nxt[i]] second;
                if (pega > dp[i]) dp[i] = pega;
                else nxt[i] = last;
                last = i;
        pair<11, int> ans = {0, 0};
        int idx = -1;
        for (int i = 0; i < v.size(); i++) if (dp[i] > ans) ans = dp[i], idx = i
        vector<int> ret;
        while (idx != -1) {
                if (\text{get}<2>(\text{v[idx]}) > 0 and
                         (nxt[idx] == -1 \text{ or } get<1>(v[nxt[idx]]) < get<0>(v[idx]))
                             ) ret.push_back(idx);
                idx = nxt[idx];
        sort(ret.begin(), ret.end());
        return ret;
```

## 30.23 inversionCount

```
// Inversion Count
//
// Computa o numero de inversoes para transformar
// 1 em r (se nao tem como, retorna -1)
//
// O(n log(n))

template<typename T> 11 inv_count(vector<T> 1, vector<T> r = {}) {
    if (!r.size()) {
        r = 1;
        sort(r.begin(), r.end());
    }
}
```

```
int n = l.size();
vector<int> v(n), bit(n);
vector<pair<T, int>> w;
for (int i = 0; i < n; i++) w.push_back({r[i], i+1});
sort(w.begin(), w.end());
for (int i = 0; i < n; i++) {
        auto it = lower_bound(w.begin(), w.end(), make_pair(l[i], 0));
        if (it == w.end() or it->first != l[i]) return -l; // nao da
        v[i] = it->second;
        it->second = -1;
}

ll ans = 0;
for (int i = n-1; i >= 0; i--) {
        for (int j = v[i]-1; j; j -= j&-j) ans += bit[j];
        for (int j = v[i]; j < n; j += j&-j) bit[j]++;
}
return ans;
}</pre>
```

## 30.24 lis

```
// LIS - recupera
// Calcula e retorna uma LIS
// O(n.log(n))
template<typename T> vector<T> lis(vector<T>& v) {
        int n = v.size(), m = -1;
        vector<T> d(n+1, INF);
        vector<int> 1(n);
        d[0] = -INF;
        for (int i = 0; i < n; i++) {</pre>
                // Para non-decreasing use upper_bound()
                int t = lower_bound(d.begin(), d.end(), v[i]) - d.begin();
                d[t] = v[i], l[i] = t, m = max(m, t);
        int p = n;
        vector<T> ret;
        while (p--) if (l[p] == m) {
                ret.push_back(v[p]);
        reverse (ret.begin(), ret.end());
        return ret;
```

## 30.25 lis2

#### 30.26 maxDist

```
// Distancia maxima entre dois pontos
// \max_{dist2(v)} - O(n \log(n))
// max_dist_manhattan - O(n)
// Quadrado da Distancia Euclidiana (precisa copiar convex_hull, ccw e pt)
11 max_dist2(vector<pt> v) {
        v = convex_hull(v);
        if (v.size() <= 2) return dist2(v[0], v[1%v.size()]);</pre>
        11 \text{ ans} = 0;
        int n = v.size(), j = 0;
for (int i = 0; i < n; i++) {</pre>
                 while (!ccw(v[(i+1)%n]-v[i], pt(0, 0), v[(j+1)%n]-v[j])) j = (j)
                 ans = \max(\{ans, dist2(v[i], v[j]), dist2(v[(i+1)%n], v[j])\});
        return ans:
// Distancia de Manhattan
template<typename T> T max_dist_manhattan(vector<pair<T, T>> v) {
        T min_sum, max_sum, min_dif, max_dif;
        min_sum = max_sum = v[0].first + v[0].second;
        min_dif = max_dif = v[0].first - v[0].second;
        for (auto [x, y] : v) {
                min_sum = min(min_sum, x+y);
                 max_sum = max(max_sum, x+y);
                 min_dif = min(min_dif, x-y);
                 max_dif = max(max_dif, x-y);
        return max(max sum - min sum, max dif - min dif);
}
```

### 30.27 minCirc

```
// Minimum Enclosing Circle
// O(n) com alta probabilidade
const double EPS = 1e-12;
mt19937 rng((int) chrono::steady_clock::now() time_since_epoch().count());
struct pt {
        double x, y;
        pt (double x_{-} = 0, double y_{-} = 0) : x(x_{-}), y(y_{-}) {}
        pt operator + (const pt& p) const { return pt(x+p.x, y+p.y); }
        pt operator - (const pt& p) const { return pt(x-p.x, y-p.y); }
        pt operator * (double c) const { return pt(x*c, y*c); }
        pt operator / (double c) const { return pt(x/c, y/c); }
};
double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
pt center(pt p, pt q, pt r) {
        pt a = p-r, b = q-r;
        pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
        return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x, b.x), c)) / cross(a, b);
struct circle {
        pt cen;
        double r:
        circle(pt cen_, double r_) : cen(cen_), r(r_) {}
        circle(pt a, pt b, pt c) {
                cen = center(a, b, c);
                r = dist(cen, a);
        bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
};
circle minCirc(vector<pt> v) {
```

### 30.28 minkowski

```
// Minkowski Sum
// Computa A+B = \{a+b : a \setminus in A, b \setminus in B\}, em que
// A e B sao poligonos convexos
// A+B eh um poligono convexo com no max |A|+|B| pontos
// O(|A|+|B|)
vector<pt> minkowski(vector<pt> p, vector<pt> q) {
        auto fix = [](vector<pt>& P) {
                 rotate(P.begin(), min_element(P.begin(), P.end()), P.end());
                 P.push_back(P[0]), P.push_back(P[1]);
        fix(p), fix(q);
        vector<pt> ret;
        int i = 0, j = 0;
        while (i < p.size()-2 \text{ or } j < q.size()-2) {
                ret.push_back(p[i] + q[j]);

auto c = ((p[i+1] - p[i]) ^ (q[j+1] - q[j]));
                 if (c >= 0) i = min<int>(i+1, p.size()-2);
                 if (c <= 0) j = min<int>(j+1, q.size()-2);
        return ret;
ld dist_convex(vector<pt> p, vector<pt> q) {
        for (pt& i : p) i = i * -1;
        auto s = minkowski(p, q);
        if (inpol(s, pt(0, 0))) return 0;
        ld ans = DINF;
        for (int i = 0; i < s.size(); i++) ans = min(ans,
                         disttoseg(pt(0, 0), line(s[(i+1)%s.size()], s[i])));
        return ans;
```

#### 30.29 mo

```
// MO
//
// Para ter o bound abaixo, escolher
// SQ = n / sqrt(q)
//
// O(n * sqrt(q))

const int MAX = 1e5+10;
const int SQ = sqrt(MAX);
int v[MAX];

int ans, freq[MAX];

inline void insert(int p) {
   int o = v[p];
   freq[o]++;
   ans += (freq[o] == 1);
}

inline void erase(int p) {
```

```
int o = v[p];
        ans -= (freq[o] == 1);
        freq[o]--;
inline 11 hilbert(int x, int y) {
        static int N = 1 << (__builtin_clz(0) - __builtin_clz(MAX));</pre>
        int rx, ry, s;
        11 d = 0:
        for (s = N/2; s > 0; s /= 2) {
                 rx = (x \& s) > 0, ry = (y \& s) > 0;
                 d += s * 11(s) * ((3 * rx) ^ ry);
                 if (ry == 0) {
                          if (rx == 1) x = N-1 - x, y = N-1 - y;
                          swap(x, y);
        return d;
#define HILBERT true
vector<int> MO(vector<pair<int, int>> &q) {
        ans = 0;
        int m = q.size();
        vector<int> ord(m);
        iota(ord.begin(), ord.end(), 0);
#if HILBERT
         vector<ll> h (m);
        for (int i = 0; i < m; i++) h[i] = hilbert(q[i].first, q[i].second);
        sort(ord.begin(), ord.end(), [&](int 1, int r) { return h[1] < h[r]; });</pre>
#6196
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
    if (q[1].first / SQ != q[r].first / SQ) return q[1].first < q[r]</pre>
                      1.first:
                 if ((q[1].first / SQ) % 2) return q[1].second > q[r].second;
                 return q[1].second < q[r].second;</pre>
        });
#endif
         vector<int> ret(m);
        int 1 = 0, r = -1;
        for (int i : ord) {
                 int ql, qr;
                 tie(ql, qr) = q[i];
                 while (r < qr) insert(++r);</pre>
                 while (1 > q1) insert (--1);
                 while (1 < q1) erase(1++);
                 while (r > qr) erase (r--);
                 ret[i] = ans;
        return ret;
```

## $30.30 \quad moDsu$

```
// MO - DSU
//
// Dado uma lista de arestas de um grafo, responde
// para cada query(1, r), quantos componentes conexos
// o grafo tem se soh considerar as arestas 1, l+1, ..., r
// Da pra adaptar pra usar MO com qualquer estrutura rollbackavel
//
// O(m sqrt(q) log(n))

struct dsu {
    int n, ans;
    vector<int> p, sz;
    stack<int> S;

    dsu(int n_) : n(n_), ans(n), p(n), sz(n) {
        for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;
    }
    int find(int k) {
        while (p[k] != k) k = p[k];
        return k;
    }
}</pre>
```

```
void add(pair<int, int> x) {
                int a = x.first, b = x.second;
                 a = find(a), b = find(b);
                if (a == b) return S.push(-1);
                if (sz[a] > sz[b]) swap(a, b);
                S.push(a);
                sz[b] += sz[a];
                p[a] = b;
        int query() { return ans; }
        void rollback() {
                int u = S.top(); S.pop();
                if (u == -1) return;
                sz[p[u]] = sz[u];
                p[u] = u;
                ans++;
};
int n;
vector<pair<int, int>> ar;
// 9d242b
vector<int> MO(vector<pair<int, int>> &q) {
        int SQ = sqrt(q.size()) + 1;
        int m = q.size();
        vector<int> ord(m);
        iota(ord.begin(), ord.end(), 0);
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
                         if (q[l].first / SQ != q[r].first / SQ) return q[l].
                              first < q[r].first;
                         return q[1].second < q[r].second;</pre>
        vector<int> ret(m);
        dsu small(n);
        for (int i = 0; i < m; i++) {
                auto [1, r] = q[ord[i]];
                if (1 / SQ == r / SQ) {
                         for (int k = 1; k <= r; k++) small.add(ar[k]);</pre>
                         ret[ord[i]] = small.query();
                         for (int k = 1; k <= r; k++) small.rollback();</pre>
        for (int i = 0; i < m; i++) {</pre>
                 dsu D(n);
                 int fim = q[ord[i]].first/SQ*SQ + SQ - 1;
                int last_r = fim;
                 int j = i-1;
                while (j+1 < m and q[ord[j+1]].first / SQ == q[ord[i]].first /</pre>
                     SQ) {
                         auto [1, r] = q[ord[++j]];
                         if (1 / SQ == r / SQ) continue;
                         while (last_r < r) D.add(ar[++last_r]);</pre>
                         for (int k = 1; k <= fim; k++) D.add(ar[k]);</pre>
                         ret[ord[j]] = D.query();
                         for (int k = 1; k <= fim; k++) D.rollback();</pre>
                i = j;
        return ret;
```

### 30.31 moOnTrees

```
// MO em Arvores
//
// Problema que resolve: https://www.spoj.com/problems/COT2/
//
// Complexidade sendo c = O(update) e SQ = sqrt(n):
```

```
// O((n + q) * sqrt(n) * c)
const int MAX = 40010, SQ = 400;
vector<int> q[MAX];
namespace LCA { ... }
int in[MAX], out[MAX], vtx[2 * MAX];
bool on[MAX];
int dif, freq[MAX];
vector<int> w;
void dfs(int v, int p, int &t) {
        vtx[t] = v, in[v] = t++;
        for (int u : g[v]) if (u != p) {
                dfs(u, v, t);
        vtx[t] = v, out[v] = t++;
void update(int p) { // faca alteracoes aqui
        int v = vtx[p];
        if (not on[v]) { // insere vtx v
                 dif += (freq[w[v]] == 0);
                 freq[w[v]]++;
        else { // retira o vertice v
                 dif -= (freq[w[v]] == 1);
                 freq[w[v]]--;
        on[v] = not on[v];
vector<tuple<int, int, int>> build_queries(const vector<pair<int, int>>& q) {
        LCA::build(0);
        vector<tuple<int, int, int>> ret;
        for (auto [1, r] : q) {
                 if (in[r] < in[l]) swap(l, r);</pre>
                 int p = LCA::lca(l, r);
                 int init = (p == 1) ? in[1] : out[1];
                 ret.emplace_back(init, in[r], in[p]);
        return ret;
vector<int> mo_tree(const vector<pair<int, int>>& vq) {
        int t = 0;
        dfs(0, -1, t);
        auto q = build_queries(vq);
        vector<int> ord(q.size());
        iota(ord.begin(), ord.end(), 0);
        sort(ord.begin(), ord.end(), [&] (int 1, int r) {
   int bl = get<0>(q[1]) / SQ, br = get<0>(q[r]) / SQ;
                 if (bl != br) return bl < br;</pre>
                 else if (b1 % 2 == 1) return get<1>(q[1]) < get<1>(q[r]);
                 else return get<1>(q[1]) > get<1>(q[r]);
        memset(freq, 0, sizeof freq);
        dif = 0;
        vector<int> ret(q.size());
        int 1 = 0, r = -1;
        for (int i : ord) {
                 auto [ql, qr, qp] = q[i];
                 while (r < qr) update(++r);
while (l > ql) update(--l);
                 while (1 < q1) update(1++);</pre>
                 while (r > qr) update(r--);
                 if (qp < 1 or qp > r) { // se LCA estah entre as pontas
                          update(qp);
                          ret[i] = dif;
                          update(qp);
```

```
else ret[i] = dif;
}
return ret;
```

## 30.32 palindromicFactorization

```
// Palindromic Factorization
// Precisa da eertree
// Computa o numero de formas de particionar cada
// prefixo da string em strings palindromicas
// O(n log n), considerando alfabeto O(1)
struct eertree { ... };
11 factorization(string s) {
       int n = s.size(), sz = 2;
        eertree PT(n);
        vector<int> diff(n+2), slink(n+2), sans(n+2), dp(n+1);
        for (int i = 1; i <= n; i++) {
                PT.add(s[i-1]);
                if (PT.size()+2 > sz)
                        diff[sz] = PT.len[sz] - PT.len[PT.link[sz]];
                        if (diff[sz] == diff[PT.link[sz]])
                                slink[sz] = slink[PT.link[sz]];
                        else slink[sz] = PT.link[sz];
                        sz++;
                for (int v = PT.last; PT.len[v] > 0; v = slink[v])
                        sans[v] = dp[i - (PT.len[slink[v]] + diff[v])];
                        if (diff[v] == diff[PT.link[v]])
                                sans[v] = (sans[v] + sans[PT.link[v]]) % MOD;
                        dp[i] = (dp[i] + sans[v]) % MOD;
        return dp[n];
```

# 30.33 parsing

```
// Parsing de Expressao
// Operacoes associativas a esquerda por default
// Para mudar isso, colocar em r_assoc
// Operacoes com maior prioridade sao feitas primeiro
bool blank(char c) {
        return c == ' ';
bool is_unary(char c) {
        return c == '+' or c == '-';
bool is_op(char c) {
        if (is_unary(c)) return true;
        return c == '*' or c == '/' or c == '+' or c == '-';
bool r_assoc(char op) {
        // operator unario - deve ser assoc. a direita
        return op < 0;
int priority(char op) {
        // operator unario - deve ter precedencia maior
        if (op < 0) return INF;</pre>
        if (op == '*' or op == '/') return 2;
```

```
if (op == '+' or op == '-') return 1;
        return -1;
void process_op(stack<int>& st, stack<int>& op) {
        char o = op.top(); op.pop();
        if (o < 0) {

    ★= -1;

                int 1 = st.top(); st.pop();
                if (o == '+') st.push(1);
                if (o == '-') st.push(-1);
        } else {
                int r = st.top(); st.pop();
                int 1 = st.top(); st.pop();
                if (o == '*') st.push(l * r);
                if (o == '/') st.push(1 / r);
                if (o == '+') st.push(1 + r);
                if (o == '-') st.push(1 - r);
int eval(string& s) {
        stack<int> st, op;
        bool un = true;
        for (int i = 0; i < s.size(); i++) {</pre>
                if (blank(s[i])) continue;
                if (s[i] == '(') {
                        op.push('(');
                        un = true;
                } else if (s[i] == ')')
                        while (op.top() != '(') process_op(st, op);
                        op.pop();
                        un = false;
                } else if (is_op(s[i])) {
                        char o = s[i];
                        if (un and is_unary(o)) o *= -1;
                        while (op.size() and (
                                                 (!r_assoc(o) and priority(op.top
                                                      ()) >= priority(o)) or
                                                  (r_assoc(o) and priority(op.top
                                                      ()) > priority(o))))
                                 process_op(st, op);
                        op.push(o);
                        un = true;
                } else {
                        int val = 0;
                        while (i < s.size() and isalnum(s[i]))</pre>
                                val = val * 10 + s[i++] - '0';
                        st.push(val);
                        un = false;
        while (op.size()) process_op(st, op);
        return st.top();
```

# 30.34 rmqOffline

```
// RMQ com Divide and Conquer
//
// Responde todas as queries em
// O(n log(n))

typedef pair<pair<int, int>, int> iii;
#define f first
#define s second

int n, q, v[MAX];
iii qu[MAX];
int ans[MAX], pref[MAX], sulf[MAX];

void solve(int l=0, int r=n-1, int ql=0, int qr=q-1) {
```

## 30.35 segmentIntersection

```
// Segment Intersection
// Verifica, dado n segmentos, se existe algum par de segmentos
// que se intersecta
// O(n log n)
bool operator < (const line& a, const line& b) { // comparador pro sweepline</pre>
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
        if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps < b.p.x)
                return ccw(a.p, a.q, b.p);
        return ccw(a.p, b.q, b.p);
bool has_intersection(vector<line> v) {
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
                return interseg(a.first, b.first);
        };
        vector<pair<pt, pair<int, int>>> w;
        for (int i = 0; i < v.size(); i++) {</pre>
                if (v[i].q < v[i].p) swap(v[i].p, v[i].q);</pre>
                w.push_back({v[i].p, {0, i}});
                w.push_back({v[i].q, {1, i}});
        sort(w.begin(), w.end());
        set<pair<line, int>> se;
        for (auto i : w) {
                line at = v[i.second.second];
                if (i.second.first == 0) {
                        auto nxt = se.lower_bound({at, i.second.second});
                        if (nxt != se.end() and intersects(*nxt, {at, i.second.
                             second})) return 1;
                        if (nxt != se.begin() and intersects(*(--nxt), {at, i.
                             second.second()) return 1;
                        se.insert({at, i.second.second});
                } else
                        auto nxt = se.upper_bound({at, i.second.second}), cur =
                             nxt, prev = --cur;
                        if (nxt != se.end() and prev != se.begin()
                                and intersects(*nxt, *(--prev))) return 1;
                        se.erase(cur);
        return 0:
```

# 30.36 simplePolygon

```
// Simple Polygon
//
// Verifica se um poligono com n pontos eh simples
//
// O(n log n)
bool operator < (const line& a, const line& b) { // comparador pro sweepline</pre>
```

```
if (a.p == b.p) return ccw(a.p, a.q, b.q);
        if (!eq(a.p.x, a.q.x) and (eq(b.p.x, b.q.x) or a.p.x+eps < b.p.x))</pre>
                 return ccw(a.p, a.q, b.p);
        return ccw(a.p, b.q, b.p);
bool simple(vector<pt> v) {
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
                 if ((a.second+1)%v.size() == b.second or
                          (b.second+1) %v.size() == a.second) return false;
                 return interseg(a.first, b.first);
        vector<line> seq;
        vector<pair<pt, pair<int, int>>> w;
for (int i = 0; i < v.size(); i++) {</pre>
                 pt at = v[i], nxt = v[(i+1)%v.size()];
                 if (nxt < at) swap(at, nxt);</pre>
                 seg.push_back(line(at, nxt));
w.push_back({at, {0, i}});
                 w.push_back({nxt, {1, i}});
                 // casos degenerados estranhos
                 if (isinseg(v[(i+2)%v.size()], line(at, nxt))) return 0;
                 if (isinseg(v[(i+v.size()-1)%v.size()], line(at, nxt))) return
        sort(w.begin(), w.end());
         set<pair<line, int>> se;
        for (auto i : w) {
                 line at = seg[i.second.second];
                 if (i.second.first == 0) {
                          auto nxt = se.lower_bound({at, i.second.second});
                          if (nxt != se.end() and intersects(*nxt, {at, i.second.
                               second})) return 0;
                          if (nxt != se.begin() and intersects(*(--nxt), {at, i.
                               second.second})) return 0;
                          se.insert({at, i.second.second});
                 } else {
                          auto nxt = se.upper_bound({at, i.second.second}), cur =
                              nxt, prev = --cur;
                          if (nxt != se.end() and prev != se.begin()
                                  and intersects(*nxt, *(--prev))) return 0;
                          se.erase(cur);
        return 1:
```

## 30.37 steinerTree

```
// Steiner Tree
// steiner: retorna o peso da menor arvore que cobre os vertices S
// get_steiner: retorna o valor minimo e as arestas de uma solucao
// se nao tiver solucao retorna LINF
// grafo nao pode ter pesos negativos
// se so tiver peso nas arestas/vertices pode deletar os vw/w no codigo
//k = |S|
// O(3^k * n + 2^k * m \log m)
// otimizacao: joga um vertice x do S fora e peque a resposta em dp[\dots][x] e
    reconstrua a arvore a partir dele
// ta comentado no codigo as mudancas necessarias
int n; // numero de vertices
vector<pair<int, int>> g[MAX]; // {vizinho, peso}
11 d[1 << K][MAX]; // dp[mask][v] = arvore minima com o subconjunto mask de S e
    o vertice v
11 vw[MAX]; // peso do vertice
11 steiner(const vector<int> &S) {
        int k = S.size(); // k--;
        for (int mask = 0; mask < (1 << k); mask++) for (int v = 0; v < n; v++) d
             [mask][v] = LINF;
```

```
for (int v = 0; v < n; v++) d[0][v] = vw[v];
        for (int i = 0; i < k; ++i) d[1 << i][S[i]] = vw[S[i]];
        for (int mask = 1; mask < (1 << k); mask++) {</pre>
                for (int a = (mask - 1) & mask; a; a = (a - 1) & mask) {
   int b = mask ^ a;
                        if (b > a) break;
                        for (int v = 0; v < n; v++)
                                 d[mask][v] = min(d[mask][v], d[a][v] + d[b][v] -
                priority_queue<pair<11, int>> pq;
                for (int v = 0; v < n; v++) {
                        if (d[mask][v] == LINF) continue;
                        pq.emplace(-d[mask][v], v);
                while (pq.size()) {
                         auto [ndist, u] = pq.top(); pq.pop();
                        if (-ndist > d[mask][u]) continue;
                         for (auto [idx, w] : g[u]) if (d[mask][idx] > d[mask][u]
                               + w + vw[idx]) 
                                 d[mask][idx] = d[mask][u] + w + vw[idx];
                                 pq.emplace(-d[mask][idx], idx);
        return d[(1 << k) - 1][S[0]]; // S[k]
#warning se k=1 a solucao eh a folha isolada e a funcao retorna edg = {}
#warning se k=0 crasha
pair<11, vector<pair<int, int>>> get_steiner(const vector<int> &S) {
        int k = S.size(); // k--;
        11 ans = steiner(S);
        vector<pair<int,int>> edg;
        stack<pair<int,int>> stk;
        stk.emplace((1 << k) - 1, S[0]); // S[k]
        while (!stk.empty()) {
                bool cont = 0;
                auto [mask,u] = stk.top();stk.pop();
                if ((__builtin_popcount(mask) == 1 and u == S[__bit_width(mask)
                     - 1])) continue;
                for (auto [idx, w] : g[u]) {
    if (d[mask][u] == d[mask][idx] + w + vw[u]) {
                                 edg.emplace back(u, idx);
                                 stk.emplace(mask, idx);
                                 cont = true;
                                 break;
                if (cont) continue;
                for (int a = (mask - 1) & mask; a; a = (a - 1) & mask) {
                        int b = mask ^ a;
                        if (d[mask][u] == d[a][u] + d[b][u] - vw[u]) {
                                 stk.emplace(a, u);
                                 stk.emplace(b, u);
                                 cont = true;
                                 break;
                assert(!mask || cont);
        return {ans, edg};
```

## 30.38 sweepDirection

```
// Sweep Direction
//
// Passa por todas as ordenacoes dos pontos definitas por "direcoes"
// Assume que nao existem pontos coincidentes
//
// O(n^2 log n)

void sweep_direction(vector<pt> v) {
    int n = v.size();
```

```
sort(v.begin(), v.end(), [](pt a, pt b) {
        if (a.x != b.x) return a.x < b.x;</pre>
        return a.y > b.y;
});
vector<int> at(n);
iota(at.begin(), at.end(), 0);
vector<pair<int, int>> swapp;
for (int i = 0; i < n; i++) for (int j = i+1; j < n; j++)
        swapp.push_back({i, j}), swapp.push_back({j, i});
sort(swapp.begin(), swapp.end(), [&](auto a, auto b) {
        pt A = rotate90(v[a.first] - v[a.second]);
        pt B = rotate90(v[b.first] - v[b.second]);
        if (quad(A) == quad(B) and !sarea2(pt(0, 0), A, B)) return a < b</pre>
        return compare_angle(A, B);
for (auto par : swapp) {
          assert(abs(at[par.first] - at[par.second]) == 1);
        int 1 = min(at[par.first], at[par.second]),
                 r = n-1 - max(at[par.first], at[par.second]);
        // l e r sao quantos caras tem de cada lado do par de pontos
        // (cada par eh visitado duas vezes)
        swap(v[at[par.first]], v[at[par.second]]);
        swap(at[par.first], at[par.second]);
```

# 31 ufmg forked/Strings

## 31.1 ahocorasick

```
// Aho-corasick
// query retorna o somatorio do numero de matches de
// todas as stringuinhas na stringona
// insert - O(|s| log(SIGMA))
// build - O(N), onde N = somatorio dos tamanhos das strings
// query - 0(|s|)
namespace aho {
        map<char, int> to[MAX];
        int link[MAX], idx, term[MAX], exit[MAX], sobe[MAX];
        void insert(string& s) {
                int at = 0;
                for (char c : s) {
                        auto it = to[at].find(c);
                        if (it == to[at].end()) at = to[at][c] = ++idx;
                        else at = it->second;
                term[at]++, sobe[at]++;
#warning nao esquece de chamar build() depois de inserir
        void build() {
                queue<int> q;
                q.push(0);
                link[0] = exit[0] = -1;
                while (q.size()) {
                        int i = q.front(); q.pop();
                        for (auto [c, j] : to[i]) {
    int l = link[i];
                                 while (1 != -1 and !to[1].count(c)) 1 = link[1];
                                 link[j] = 1 == -1 ? 0 : to[1][c];
                                 exit[j] = term[link[j]] ? link[j] : exit[link[j
                                 if (exit[j]+1) sobe[j] += sobe[exit[j]];
                                 q.push(j);
        int query(string& s) {
```

```
int at = 0, ans = 0;
for (char c : s) {
    while (at != -1 and !to[at].count(c)) at = link[at];
        at = at == -1 ? 0 : to[at][c];
        ans += sobe[at];
}
return ans;
```

# 31.2 dynamicSuffixArray

```
// Suffix Array Dinamico
// Mantem o suffix array, lcp e rank de uma string,
// premitindo push_front e pop_front
// O operador [i] return um par com sa[i] e lcp[i]
// lcp[i] tem o lcp entre sa[i] e sa[i-1] (lcp[0] = 0)
// Complexidades:
// Construir sobre uma string de tamanho n: O(n log n)
// push_front e pop_front: O(log n) amortizado
struct dyn_sa {
        struct node {
                int sa, lcp;
                node *1, *r, *p;
                int sz, mi;
                node(int sa_, int lcp_, node* p_) : sa(sa_), lcp(lcp_),
                        1(NULL), r(NULL), p(p_), sz(1), mi(lcp) {}
                void update() {
                        sz = 1, mi = 1cp;
                        if (1) sz += 1->sz, mi = min(mi, 1->mi);
                        if (r) sz += r->sz, mi = min(mi, r->mi);
        };
        node* root;
        vector<11> tag; // tag of a suffix (reversed id)
        string s; // reversed
        dyn_sa() : root(NULL) {}
        dyn_sa(string s_) : dyn_sa() {
                reverse(s_.begin(), s_.end());
                for (char c : s_) push_front(c);
        dyn_sa() {
                vector<node*> q = {root};
                while (q.size()) {
                        node* x = q.back(); q.pop_back();
                        if (!x) continue;
                        q.push_back(x->1), q.push_back(x->r);
                        delete x;
        int size(node* x) { return x ? x->sz : 0; }
        int mirror(int i) { return s.size()-1 - i; }
        bool cmp(int i, int j) {
                if (s[i] != s[j]) return s[i] < s[j];</pre>
                if (i == 0 \text{ or } j == 0) return i < j;
                return tag[i-1] < tag[j-1];</pre>
        void fix_path(node* x) { while (x) x->update(), x = x->p; }
        void flatten(vector<node*>& v, node* x) {
                if (!x) return;
                flatten(v, x \rightarrow 1);
                v.push_back(x);
                flatten(v, x->r);
        void build(vector<node*>& v, node*& x, node* p, int L, int R, 11 1, 11 r
                if (L > R) return void(x = NULL);
                int M = (L+R)/2;
                11 m = (1+r)/2;
```

```
x = v[M];
        x->p = p;
         tag[x->sa] = m;
        build(v, x->1, x, L, M-1, 1, m-1), build(v, x->r, x, M+1, R, m
             +1, r);
         x->update();
void fix(node*& x, node* p, 11 1, 11 r) {
    if (3*max(size(x->1), size(x->r)) <= 2*size(x)) return x->update
              ();
         vector<node*> v;
         flatten(v, x);
        build(v, x, p, 0, v.size()-1, 1, r);
node* next(node* x) {
        if (x->r) {
                 x = x->r;
                  while (x->1) x = x->1;
                 return x;
         while (x->p and x->p->r == x) x = x->p;
        return x->p;
node* prev(node* x) {
        if (x->1) {
                  while (x->r) x = x->r;
                 return x;
         while (x->p \text{ and } x->p->l == x) x = x->p;
         return x->p;
int get_lcp(node* x, node* y) {
        if (!x or !y) return 0; // change defaut value here
if (s[x->sa] != s[y->sa]) return 0;
         if (x->sa == 0 \text{ or } y->sa == 0) return 1;
         return 1 + query (mirror (x->sa-1), mirror (y->sa-1));
void add_suf(node*& x, node* p, int id, ll l, ll r) {
        if (!x) {
                 x = new node(id, 0, p);
                  node *prv = prev(x), *nxt = next(x);
                  int lcp_cur = get_lcp(prv, x), lcp_nxt = get_lcp(x, nxt)
                  if (nxt) nxt->lcp = lcp_nxt, fix_path(nxt);
                 x->lcp = lcp_cur;
                  tag[id] = (1+r)/2;
                  x->update();
                 return;
         if (cmp(id, x->sa)) add_suf(x->1, x, id, 1, tag[x->sa]-1);
         else add_suf(x\rightarrow r, x, id, tag[x\rightarrow sa]+1, r);
         fix(x, p, 1, r);
void push_front(char c) {
         s += c;
         tag.push_back(-1);
         add_suf(root, NULL, s.size() - 1, 0, 1e18);
void rem_suf(node*& x, int id) {
         if (x->sa != id) {
                 if (tag[id] < tag[x->sa]) return rem_suf(x->1, id);
                 return rem_suf(x->r, id);
         node* nxt = next(x);
         if (nxt) nxt \rightarrow lcp = min(nxt \rightarrow lcp, x \rightarrow lcp), fix path(nxt);
         node *p = x->p, *tmp = x;
        if (!x->1 or !x->r) {
                 x = x->1 ? x->1 : x->r;
                  if (x) x->p = p;
         } else {
                  for (tmp = x->1, p = x; tmp->r; tmp = tmp->r) p = tmp;
                  x->sa = tmp->sa, x->lcp = tmp->lcp;
                  if (tmp->1) tmp->1->p = p;
                  if (p->1 == tmp) p->1 = tmp->1;
                  else p \rightarrow r = tmp \rightarrow 1;
```

```
fix_path(p);
        delete tmp;
void pop_front() {
        if (!s.size()) return;
        s.pop_back();
        rem_suf(root, s.size());
        tag.pop_back();
int query(node* x, 11 1, 11 r, 11 a, 11 b) {
        if (!x \text{ or } tag[x->sa] == -1 \text{ or } r < a \text{ or } b < 1) \text{ return } s.size();
        if (a <= 1 and r <= b) return x->mi;
        int ans = s.size();
        if (a \le tag[x->sa] and tag[x->sa] \le b) and tag[x->sa]
        ans = min(ans, query(x->1, 1, tag[x->sa]-1, a, b));
        ans = min(ans, query(x->r, tag[x->sa]+1, r, a, b));
        return ans:
int query(int i, int j) { // lcp(s[i..], s[j..])
        if (i == j) return s.size() - i;
        11 a = tag[mirror(i)], b = tag[mirror(j)];
        int ret = query(root, 0, 1e18, min(a, b)+1, max(a, b));
// optional: get rank[i], sa[i] and lcp[i]
int rank(int i) {
        i = mirror(i);
        node* x = root;
        int ret = 0;
        while (x) {
                if (tag[x->sa] < tag[i]) {
                         ret += size(x->1)+1;
                         x = x->r;
                } else x = x -> 1;
        return ret;
pair<int, int> operator[](int i) {
        node* x = root;
        while (1) {
                if (i < size(x->1)) x = x->1;
                else {
                         i \rightarrow size(x\rightarrow 1);
                         if (!i) return {mirror(x->sa), x->lcp};
                         i--, x = x->r;
        }
}
```

### 31.3 eertree

};

```
// eertree
// Constroi a eertree, caractere a caractere
// Inicializar com a quantidade de caracteres maxima
// size() retorna a quantidade de substrings pal. distintas
// depois de chamar propagate(), cada substring palindromica
// ocorre qt[i] vezes. O propagate() retorna o numero de
// substrings pal. com repeticao
// O(n) amortizado, considerando alfabeto O(1)
struct eertree {
       vector<vector<int>> t;
       int n, last, sz;
       vector<int> s, len, link, gt;
        eertree(int N) {
               t = vector(N+2, vector(26, int()));
                s = len = link = qt = vector<int>(N+2);
               s[0] = -1;
               link[0] = 1, len[0] = 0, link[1] = 1, len[1] = -1;
                sz = 2, last = 0, n = 1;
```

```
void add(char c) {
        s[n++] = c -= 'a';
        while (s[n-len[last]-2] != c) last = link[last];
        if (!t[last][c]) {
                 int prev = link[last];
                 while (s[n-len[prev]-2] != c) prev = link[prev];
                 link[sz] = t[prev][c];
                 len[sz] = len[last] + 2;
                 t[last][c] = sz++;
        qt[last = t[last][c]]++;
int size() { return sz-2; }
11 propagate() {
        11 \text{ ret} = 0;
        for (int i = n; i > 1; i--) {
    qt[link[i]] += qt[i];
                 ret += qt[i];
        return ret:
```

## 31.4 hashing

};

```
// String Hashing
// Complexidades:
// construtor - O(|s|)
// operator() - 0(1)
mt19937 rng((int) chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r) {
        uniform_int_distribution<int> uid(1, r);
        return uid(rng);
template<int MOD> struct str_hash { // 116fcb
        static int P;
        vector<ll> h, p;
        str_hash(string s) : h(s.size()), p(s.size()) {
                 p[0] = 1, h[0] = s[0];
for (int i = 1; i < s.size(); i++)
                         p[i] = p[i - 1] *P%MOD, h[i] = (h[i - 1] *P + s[i])%MOD;
        11 operator()(int 1, int r) { // retorna hash s[1...r]
                 11 \text{ hash} = h[r] - (1 ? h[1 - 1] *p[r - 1 + 1] *MOD : 0);
                 return hash < 0 ? hash + MOD : hash;</pre>
};
template<int MOD> int str_hash<MOD>::P = uniform(256, MOD - 1); // 1 > |sigma|
```

# 31.5 hashingLargeMod

```
ans = (ans\&MOD) + (ans>>61), ans = (ans\&MOD) + (ans>>61);
        return ans - 1:
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
11 uniform(11 1, 11 r) {
        uniform_int_distribution<ll> uid(l, r);
        return uid(rng);
struct str_hash {
        static 11 P;
        vector<ll> h, p;
        str_hash(string s) : h(s.size()), p(s.size()) {
                p[0] = 1, h[0] = s[0];
                 for (int i = 1; i < s.size(); i++)</pre>
                         p[i] = mulmod(p[i-1], P), h[i] = (mulmod(h[i-1], P))
                              + s[i])%MOD;
        11 operator()(int 1, int r) { // retorna hash s[1...r]
                11 \text{ hash} = h[r] - (1 ? \text{ mulmod}(h[1 - 1], p[r - 1 + 1]) : 0);
                return hash < 0 ? hash + MOD : hash;</pre>
};
11 str hash::P = uniform(256, MOD - 1); // 1 > /sigma/
```

## 31.6 kmp

```
// KMP
// matching(s, t) retorna os indices das ocorrencias
// de s em t
// autKMP constroi o automato do KMP
// Complexidades:
// pi - O(n)
// match - O(n + m)
// construir o automato - O(|sigma|*n)
// n = |padrao| e m = |texto|
template<typename T> vector<int> pi(T s) {
        vector<int> p(s.size());
        for (int i = 1, j = 0; i < s.size(); i++) {</pre>
                 while (j \text{ and } s[j] != s[i]) j = p[j-1];
                 if (s[j] == s[i]) j++;
                p[i] = j;
        return p:
template<typename T> vector<int> matching(T& s, T& t) {
        vector<int> p = pi(s), match;
        for (int i = 0, j = 0; i < t.size(); i++) {</pre>
                 while (j \text{ and } s[j] != t[i]) j = p[j-1];
                 if (s[j] == t[i]) j++;
                if (j == s.size()) match.push_back((i-j+1), j = p[j-1];
        return match;
struct KMPaut : vector<vector<int>> {
        KMPaut (string& s) : vector<vector<int>>(26, vector<int>(s.size()+1)) {
                 vector<int> p = pi(s);
                 auto& aut = *this;
                 aut[s[0]-'a'][0] = 1;
                 for (char c = 0; c < 26; c++)
                         for (int i = 1; i <= s.size(); i++)</pre>
                                 aut[c][i] = s[i] - 'a' == c ? i+1 : aut[c][p[i]
                                       -1]];
};
```

#### 31.7 manacher

```
// Manacher
// manacher recebe um vetor de T e retorna o vetor com tamanho dos palindromos
// ret[2*i] = tamanho do maior palindromo centrado em i
// ret[2*i+1] = tamanho maior palindromo centrado em i e i+1
// Complexidades:
// manacher - O(n)
// palindrome - <0(n), 0(1)>
// pal_end - 0(n)
template<typename T> vector<int> manacher(const T& s) {
        int 1 = 0, r = -1, n = s.size();
        vector < int > d1(n), d2(n);
        for (int i = 0; i < n; i++) {
                int k = i > r ? 1 : min(d1[l+r-i], r-i);
                while (i+k < n \&\& i-k >= 0 \&\& s[i+k] == s[i-k]) k++;
                d1[i] = k--;
                if (i+k > r) l = i-k, r = i+k;
        1 = 0, r = -1;
        for (int i = 0; i < n; i++) {
                int k = i > r ? 0 : min(d2[1+r-i+1], r-i+1); k++;
                while (i+k \le n \&\& i-k \ge 0 \&\& s[i+k-1] == s[i-k]) k++;
                d2[i] = --k;
if (i+k-1 > r) l = i-k, r = i+k-1;
        vector<int> ret(2*n-1);
        for (int i = 0; i < n; i++) ret[2*i] = 2*d1[i]-1;
        for (int i = 0; i < n-1; i++) ret[2*i+1] = 2*d2[i+1];
        return ret;
// verifica se a string s[i..j] eh palindromo
template<typename T> struct palindrome {
        vector<int> man;
        palindrome(const T& s) : man(manacher(s)) {}
        bool query(int i, int j) {
                return man[i+j] >= j-i+1;
};
// tamanho do maior palindromo que termina em cada posicao
template<typename T> vector<int> pal_end(const T& s) {
        vector<int> ret(s.size());
        palindrome<T> p(s);
        ret[0] = 1;
        for (int i = 1; i < s.size(); i++) {</pre>
                ret[i] = min(ret[i-1]+2, i+1);
                while (!p.query(i-ret[i]+1, i)) ret[i]--;
        return ret;
```

# 31.8 minMaxSuffixCyclic

```
// Min/max suffix/cyclic shift
//
// Computa o indice do menor/maior sufixo/cyclic shift
// da string, lexicograficamente
//
// O(n)

template<typename T> int max_suffix(T s, bool mi = false) {
    s.push_back(*min_element(s.begin(), s.end())-1);
    int ans = 0;
    for (int i = 1; i < s.size(); i++) {
        int j = 0;
        while (ans+j < i and s[i+j] == s[ans+j]) j++;
        if (s[i+j] > s[ans+j]) {
```

# 31.9 suffixArray

```
// Suffix Array - O(n)
// Rapidao
// Computa o suffix array em 'sa', o rank em 'rnk'
// e o lcp em 'lcp'
// query(i, j) retorna o LCP entre s[i..n-1] e s[j..n-1]
// Complexidades
// O(n) para construir
// query - 0(1)
template<typename T> struct rmq {
        int n; static const int b = 30;
        vector<int> mask, t;
        int op(int x, int y) { return v[x] \leftarrow v[y] ? x : y; }
        int msb(int x) { return __builtin_clz(1) -__builtin_clz(x); }
        int small(int r, int sz = b) { return r-msb(mask[r]&((1<<sz)-1)); }</pre>
        rmq(const \ vector<T>\& v_) : v(v_), n(v.size()), mask(n), t(n) {
                for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
                         at = (at << 1) & ((1 << b) -1);
                         while (at and op(i-msb(at&-at), i) == i) at ^= at&-at;
                for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
                for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0; i+(1<<<math>j) <=
                     n/b; i++)
                         t[n/b*j+i] = op(t[n/b*(j-1)+i], t[n/b*(j-1)+i+(1<<(j-1))
                             ]);
        int index_query(int 1, int r) {
                if (r-l+1 <= b) return small(r, r-l+1);</pre>
                int x = 1/b+1, y = r/b-1;
                if (x > y) return op (small(1+b-1), small(r));
                int j = msb(y-x+1);
                int ans = op(small(1+b-1), op(t[n/b*j+x], t[n/b*j+y-(1<<j)+1]));
                return op(ans, small(r));
        T query(int 1, int r) { return v[index_query(1, r)]; }
};
struct suffix_array {
        string s;
        int n;
        vector<int> sa, cnt, rnk, lcp;
        rmq<int> RMQ;
        bool cmp(int a1, int b1, int a2, int b2, int a3=0, int b3=0) {
                return a1 != b1 ? a1 < b1 : (a2 != b2 ? a2 < b2 : a3 < b3);
```

```
template<typename T> void radix(int* fr, int* to, T* r, int N, int k) {
        cnt = vector<int>(k+1, 0);
        for (int i = 0; i < N; i++) cnt[r[fr[i]]]++;</pre>
        for (int i = 1; i <= k; i++) cnt[i] += cnt[i-1];</pre>
        for (int i = N-1; i+1; i--) to[--cnt[r[fr[i]]]] = fr[i];
void rec(vector<int>& v, int k) {
        auto &tmp = rnk, &m0 = lcp;
        int N = v.size()-3, sz = (N+2)/3, sz2 = sz+N/3;
        vector<int> R(sz2+3);
        for (int i = 1, j = 0; j < sz2; i += i*3) R[j++] = i;
        radix(&R[0], &tmp[0], &v[0]+2, sz2, k);
        radix(\&tmp[0], \&R[0], \&v[0]+1, sz2, k);
        radix(&R[0], &tmp[0], &v[0]+0, sz2, k);
        int dif = 0;
        int 10 = -1, 11 = -1, 12 = -1;
        for (int i = 0; i < sz2; i++) {</pre>
                if (v[tmp[i]] != 10 or v[tmp[i]+1] != 11 or v[tmp[i]+2]
                        10 = v[tmp[i]], 11 = v[tmp[i]+1], 12 = v[tmp[i]
                             ]+2], dif++;
                if (tmp[i]%3 == 1) R[tmp[i]/3] = dif;
                else R[tmp[i]/3+sz] = dif;
        if (dif < sz2) {
                rec(R, dif);
                for (int i = 0; i < sz2; i++) R[sa[i]] = i+1;
        } else for (int i = 0; i < sz2; i++) sa[R[i]-1] = i;
        for (int i = 0, j = 0; j < sz2; i++) if (sa[i] < sz) tmp[j++] =
             3*sa[i];
        radix(&tmp[0], &m0[0], &v[0], sz, k);
        for (int i = 0; i < sz2; i++)
                sa[i] = sa[i] < sz ? 3*sa[i]+1 : 3*(sa[i]-sz)+2;
        int at = sz2+sz-1, p = sz-1, p2 = sz2-1;
        while (p >= 0 and p2 >= 0) {
                if ((sa[p2]%3==1 and cmp(v[m0[p]], v[sa[p2]], R[m0[p
                        R[sa[p2]/3+sz])) or (sa[p2]%3==2 and cmp(v[m0[p
                             ]], v[sa[p2]],
                         v[m0[p]+1], v[sa[p2]+1], R[m0[p]/3+sz], R[sa[p2]+1]
                             ]/3+1])))
                        sa[at--] = sa[p2--];
                else sa[at--] = m0[p--];
        while (p >= 0) sa[at--] = m0[p--];
        if (N%3==1) for (int i = 0; i < N; i++) sa[i] = sa[i+1];</pre>
suffix_array(const string& s_) : s(s_), n(s.size()), sa(n+3),
                cnt(n+1), rnk(n), lcp(n-1) {
        vector<int> v(n+3);
        for (int i = 0; i < n; i++) v[i] = i;
        radix(&v[0], &rnk[0], &s[0], n, 256);
        int dif = 1;
        for (int i = 0; i < n; i++)</pre>
                v[rnk[i]] = dif += (i and s[rnk[i]] != s[rnk[i-1]]);
        if (n \ge 2) rec(v, dif);
        sa.resize(n);
        for (int i = 0; i < n; i++) rnk[sa[i]] = i;</pre>
        for (int i = 0, k = 0; i < n; i++, k -= !!k) {
                if (rnk[i] == n-1) {
                        k = 0:
                        continue;
                int j = sa[rnk[i]+1];
                while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
                lcp[rnk[i]] = k;
        RMQ = rmq<int>(lcp);
int query(int i, int j) {
```

```
if (i == j) return n-i;
        i = rnk[i], j = rnk[j];
        return RMQ.query(min(i, j), max(i, j)-1);
pair<int, int> next(int L, int R, int i, char c) {
        int 1 = L, r = R+1;
        while (1 < r) {
                int m = (1+r)/2;
                if (i+sa[m] >= n or s[i+sa[m]] < c) l = m+1;</pre>
        if (1 == R+1 or s[i+sa[1]] > c) return {-1, -1};
        L = 1;
        1 = L, r = R+1;
        while (1 < r) {
                int m = (1+r)/2;
                if (i+sa[m] >= n or s[i+sa[m]] <= c) l = m+1;</pre>
                else r = m:
        \hat{R} = 1 - 1;
        return {L, R};
// quantas vezes 't' ocorre em 's' - O(|t| log n)
int count_substr(string& t) {
        int L = 0, R = n-1;
        for (int i = 0; i < t.size(); i++) {</pre>
                tie(L, R) = next(L, R, i, t[i]);
                if (L == -1) return 0;
        return R-L+1;
// exemplo de f que resolve o problema
// https://codeforces.com/edu/course/2/lesson/2/5/practice/contest
     /269656/problem/D
11 f(11 k) { return k*(k+1)/2; }
11 dfs(int L, int R, int p) { // dfs na suffix tree chamado em pre ordem
        int ext = L != R ? RMQ.query(L, R-1) : n - sa[L];
        // Tem 'ext - p' substrings diferentes que ocorrem 'R-L+1' vezes
        // O LCP de todas elas eh 'ext'
        ll ans = (ext-p)*f(R-L+1);
         // L eh terminal, e folha sse L == R
        if (sa[L]+ext == n) L++;
         // se for um SA de varias strings separadas como s#t$u&, usar no
              lugar do if de cima
                 (separadores < 'a', diferentes e inclusive no final)
        // while (L \le R \&\& (sa[L] + ext == n || s[sa[L] + ext] < 'a')) {
        // L++;
        11 }
        while (L \le R)
                int idx = L != R ? RMQ.index_query(L, R-1) : -1;
                if (idx == -1 or lcp[idx] != ext) idx = R;
                ans += dfs(L, idx, ext);
                T_{x} = i dx + 1:
        return ans:
// sum over substrings: computa, para toda substring t distinta de s,
// \sum f(# ocorrencias de t em s) - O (n)
ll sos() { return dfs(0, n-1, 0); }
```

# 31.10 suffixArray2

};

```
// Suffix Array - O(n log n)
//
// kasai recebe o suffix array e calcula lcp[i],
```

```
// o lcp entre s[sa[i],...,n-1] e s[sa[i+1],...,n-1]
// Complexidades:
// suffix_array - O(n log(n))
// kasai - O(n)
vector<int> suffix_array(string s) {
        int n = s.size(), N = max(n, 260);
        vector<int> sa(n), ra(n);
        for (int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];
        for (int k = 0; k < n; k ? k *= 2 : k++) {
                vector<int> nsa(sa), nra(n), cnt(N);
                 for (int i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n) %n, cnt[ra[i]
                     ]]++;
                 for(int i = 1; i < N; i++) cnt[i] += cnt[i-1];</pre>
                 for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]]] = nsa[i];
                 for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r += ra[sa[i]] !=</pre>
                         ra[sa[i-1]] or ra[(sa[i]+k)%n] != ra[(sa[i-1]+k)%n];
                 ra = nra;
                if (ra[sa[n-1]] == n-1) break;
        return vector<int>(sa.begin()+1, sa.end());
vector<int> kasai(string s, vector<int> sa) {
        int n = s.size(), k = 0;
        vector<int> ra(n), lcp(n);
        for (int i = 0; i < n; i++) ra[sa[i]] = i;</pre>
        for (int i = 0; i < n; i++, k -= !!k) {
                 if (ra[i] == n-1) { k = 0; continue; }
                 int j = sa[ra[i]+1];
                 while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
                 lcp[ra[i]] = k;
        return lcp;
```

## 31.11 suffixAutomaton

```
// Suffix Automaton
// Automato que aceita os sufixos de uma string
// Todas as funcoes sao lineares
namespace sam {
        int cur, sz, len[2*MAX], link[2*MAX], acc[2*MAX];
        int nxt[2*MAX][26];
        void add(int c) {
                int at = cur;
                len[sz] = len[cur]+1, cur = sz++;
                while (at != -1 and !nxt[at][c]) nxt[at][c] = cur, at = link[at
                    ];
                if (at == -1) { link[cur] = 0; return; }
                int q = nxt[at][c];
                if (len[q] == len[at]+1) { link[cur] = q; return; }
                int qq = sz++;
                len[qq] = len[at]+1, link[qq] = link[q];
                for (int i = 0; i < 26; i++) nxt[qq][i] = nxt[q][i];</pre>
                while (at !=-1 and nxt[at][c] == q) nxt[at][c] = qq, at = link[
                    at];
                link[cur] = link[q] = qq;
        void build(string& s) -
                cur = 0, sz = 0, len[0] = 0, link[0] = -1, sz++;
                for (auto i : s) add(i-'a');
                int at = cur;
                while (at) acc[at] = 1, at = link[at];
        // coisas que da pra fazer:
```

```
11 distinct_substrings() {
        11 ans = 0:
        for (int i = 1; i < sz; i++) ans += len[i] - len[link[i]];</pre>
string longest_common_substring(string& S, string& T) {
        build(S);
        int at = 0, 1 = 0, ans = 0, pos = -1;
for (int i = 0; i < T.size(); i++) {</pre>
                 while (at and !nxt[at][T[i]-'a']) at = link[at], l = len
                     [atl:
                 if (nxt[at][T[i]-'a']) at = nxt[at][T[i]-'a'], l++;
                 else at = 0, 1 = 0;
                 if (1 > ans) ans = 1, pos = i;
        return T.substr(pos-ans+1, ans);
11 dp[2*MAX];
11 paths(int i) {
        auto& x = dp[i];
        if (x) return x;
        for (int j = 0; j < 26; j++) if (nxt[i][j]) x += paths(nxt[i][j</pre>
             ]);
        return x;
void kth_substring(int k, int at=0) { // k=1 : menor substring lexicog.
        for (int i = 0; i < 26; i++) if (k and nxt[at][i]) {
                 if (paths(nxt[at][i]) >= k) {
                         cout << char('a'+i);
                          kth_substring(k-1, nxt[at][i]);
                         return;
                 k -= paths(nxt[at][i]);
        }
```

## 31.12 trie

};

```
// Trie
// trie T() constroi uma trie para o alfabeto das letras minusculas
// trie T(tamanho do alfabeto, menor caracter) tambem pode ser usado
// T.insert(s) - O(|s|*sigma)
// T.erase(s) - O(|s|)
// T.find(s) retorna a posicao, -1 se nao achar - O(|s|)
// T.count_pref(s) numero de strings que possuem s como prefixo - O(|s|)
struct trie {
        vector<vector<int>> to;
        vector<int> end, pref;
        int sigma; char norm;
        trie(int sigma_=26, char norm_='a') : sigma(sigma_), norm(norm_) {
               to = {vector<int>(sigma)};
               end = {0}, pref = {0};
        void insert(string s) {
                int x = 0;
                for (auto c : s) {
                        int &nxt = to[x][c-norm];
                        if (!nxt) {
                                nxt = to.size();
                                to.push_back(vector<int>(sigma));
                                end.push_back(0), pref.push_back(0);
                        x = nxt, pref[x]++;
                end[x]++, pref[0]++;
        void erase(string s) {
                int x = 0;
                for (char c : s) {
                        int &nxt = to[x][c-norm];
```

```
x = nxt, pref[x]--;
if (!pref[x]) nxt = 0;
} end[x]--, pref[0]--;
} int find(string s) {
    int x = 0;
    for (auto c : s) {
        x = to[x][c-norm];
        if (!x) return -1;
    }
    return x;
} int count_pref(string s) {
    int id = find(s);
    return id >= 0 ? pref[id] : 0;
};
```

#### 31.13 z

```
// Z
//
// z[i] = lcp(s, s[i..n))
//
// Complexidades:
// z - O(|s|)
// match - O(|s| + |p|)

vector<int> get_z(string s) {
    int n = s.size();
    vector<int> z(n, 0);

    int 1 = 0, r = 0;
    for (int i = 1; i < n; i++) {
        if (i <= r) z[i] = min(r - i + 1, z[i - 1]);
        while (i + z[i] < n and s[z[i]] == s[i + z[i]]) z[i]++;
        if (i + z[i] - 1 > r) 1 = i, r = i + z[i] - 1;
}

return z;
}
```

## 32 Utils

## 32.1 execution time

```
// https://codeforces.com/blog/entry/57647
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 300005
#define mod 1000000007
signed main()
  ios_base::sync_with_stdio(false);
```

## 32.2 rand

```
// https://codeforces.com/blog/entry/61587
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
     tree_order_statistics_node_update>;
#define int long long int
#define pb push_back
#define pi pair<int, int>
#define pii pair<int, pi>
#define fir first
#define sec second
#define MAXN 300005
#define mod 1000000007
signed main()
    ios_base::sync_with_stdio(false);
    cin.tie(NULL);
    mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
    int n = 10;
    vector<int> v(n);
    for (int i = 0; i < n; i++)
        v[i] = i;
    shuffle(v.begin(), v.end(), rng); // random shuffle
    int x = rng() % 10; // better than rand()
    cout << x << endl;</pre>
    // random integer values on the closed interval [a, b]
    int y = uniform_int_distribution<int>(3, 77)(rng);
    cout << y << endl;
    return 0;
```

## 32.3 runner

```
# This script does the following:
# 1 - Generate a random testcase
# 2 - Run some "naive" code with this input
# 3 - Run your code with this input
# 4 - Compare the outputs
import os
import subprocess

naive = "brute.cpp" # path to naive code
code = "d.cpp" # path to your code
```

```
generator = "g.cpp" # path to test generator
def compile_codes():
    os.system('g++ ' + generator + ' -o generator -02')
os.system('g++ ' + naive + ' -o naive -02')
os.system('g++ ' + code + ' -o code -02')
def generate_case():
    os.system('./generator > in');
def get_naive_output():
    output = os.popen('./naive <in').read()</pre>
    return output
def get_code_output():
     # se tiver um runtime error, vai parar
    xalala = subprocess.run('./code <in', shell=True, text=True, capture_output=</pre>
         True, check=True)
    return xalala.stdout
def main():
    compile_codes()
    while True:
         generate_case()
         naive_output = get_naive_output()
         code_output = get_code_output()
         if naive_output == code_output:
             print ('ACCEPTED')
         else:
             print('FAILED\n')
             print('ANSWER:')
             print (naive_output)
             print('\nCODE OUTPUT:')
             print (code_output)
             break
if __name__ == '__main__':
    main()
```

#### 32.4 runner2

```
# This script does the following:
# 1 - Run a code with all inputs files from a folder
# 2 - Compare the output for each test case with the answer
import os
code = "a.cpp" # Path to your code
input_folder = "input" # Path to folder which the input files are
output_folder = "output" # Path to folder which the output files are
input_prefix = "L_" # prefix of all input files names output_prefix = "L_" # prefix of all input files names
tests = 56 # Number of test cases
def compile_code():
    os.system('q++ ' + code + ' -o code -02')
def get_ans(output):
    out = open(output, "r")
    ret = out.read()
    out.close()
    return ret
```

```
def get_code_output(input):
    output = os.popen('./code <' + input).read()</pre>
    return output
def main():
    compile_code()
    # tests indexed from 1
    for i in range (1, tests + 1):
        ans = get_ans(output_folder + '/' + output_prefix + str(i))
        code_output = get_code_output(input_folder + '/' + input_prefix + str(i)
        print('Case' + str(i) + ': ')
        if ans == code_output:
            print('ACCEPTED')
        else :
            print('FAILED\n')
            print('ANSWER:')
            print(ans)
            print('\nCODE OUTPUT:')
            print (code_output)
        print()
if __name__ == '__main__':
    main()
```

#### $32.5 \quad int 128$

```
// https://codeforces.com/blog/entry/75044
// functions to print and read a __int128 in c++
__int128 read()
   _{int128 x = 0, f = 1;}
  char ch = getchar();
  while (ch < '0' || ch > '9')
    if (ch == '-')
      f = -1;
    ch = getchar();
  while (ch >= '0' && ch <= '9')
    x = x * 10 + ch - '0';
    ch = getchar();
  return x * f;
void print(__int128 x)
  if (x < 0)
    cout << "-";
    x = -x;
  if (x > 9)
   print(x / 10);
  char at = (x % 10) + '0';
  cout << at;
```