

# Universidade Federal de Alagoas

# Floyd (WA)rshall

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```
3 Data structures
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4 Graph
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5 Geometry
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6 Strings
7 Various
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Contest (1)
template.cpp
                                                         25 lines
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 1000006
#define mod 998244353
signed main() {
  ios_base::sync_with_stdio(false);
  mt19937 64 rng(chrono::steady clock::now().time since epoch()
 return 0;
runner.py
                                                         44 lines
import os
import subprocess
naive = "brute.cpp" # path to naive code
code = "d.cpp" # path to your code
generator = "g.cpp" # path to test generator
def compile_codes():
   os.system('q++ ' + generator + ' -o generator -O2')
    os.system('g++ ' + naive + ' -o naive -02')
   os.system('g++ ' + code + ' -o code -02')
def generate_case():
    os.system('./generator > in');
def get_naive_output():
```

1

1

1 Contest

2 Mathematics

```
modint.cpp
    output = os.popen('./naive <in').read()</pre>
    return output
def get_code_output():
    # se tiver um runtime error, vai parar
    xalala = subprocess.run('./code <in', shell=True, text=True</pre>
         , capture_output=True, check=True)
    return xalala.stdout
def main():
    compile_codes()
    while True:
        generate_case()
        naive_output = get_naive_output()
        code_output = get_code_output()
        if naive_output == code_output:
            print('ACCEPTED')
        else:
            print('FAILED\n')
            print ('ANSWER:')
            print (naive output)
            print('\nCODE OUTPUT:')
            print(code_output)
            break
if __name__ == '__main__':
    main()
int128.cpp
                                                           22 lines
__int128 read() {
  __int128 x = 0, f = 1;
  char ch = getchar();
  while (ch < '0' || ch > '9') {
    if (ch == '-') f = -1;
    ch = getchar();
  while (ch >= '0' && ch <= '9') {
    x = x * 10 + ch - '0';
    ch = getchar();
  return x * f;
void print(__int128 x) {
 if (x < 0) {
    cout << "-";
    x = -x;
 if (x > 9) print (x / 10);
  char at = (x % 10) + '0';
  cout << at;
pragma.cpp
                                                            5 lines
#include <iostream>
using namespace std;
#pragma GCC target("avx2")
#pragma GCC optimize("03")
#pragma GCC optimize("unroll-loops")
Mathematics (2)
```

```
78 lines
struct modint {
 int val;
  modint(int v = 0) { val = v % mod; }
  int pow(int y) {
    modint x = val;
    modint z = 1;
    while (y) {
     if (y \& 1) z *= x;
     x *= x;
      y >>= 1;
    return z.val;
 int inv() { return pow(mod - 2); }
  void operator=(int o) { val = 0 % mod; }
  void operator=(modint o) { val = o.val % mod; }
  void operator+=(modint o) { *this = *this + o; }
  void operator = (modint o) { *this = *this - o; }
  void operator*=(modint o) { *this = *this * o; }
  void operator/=(modint o) { *this = *this / o; }
  bool operator==(modint o) { return val == o.val; }
 bool operator!=(modint o) { return val != o.val; }
 int operator*(modint o) { return ((val * o.val) % mod); }
 int operator/(modint o) { return (val * o.inv()) % mod; }
 int operator+(modint o) { return (val + o.val) % mod; }
 int operator-(modint o) { return (val - o.val + mod) % mod; }
modint f[MAXN];
modint inv[MAXN];
modint invfat[MAXN];
void calc() {
 f[0] = 1;
 for (int i = 1; i < MAXN; i++) {
   f[i] = f[i - 1] * i;
 inv[1] = 1;
 for (int i = 2; i < MAXN; ++i) {
   int val = mod / i;
   val = (inv[mod % i] * val) % mod;
   val = mod - val;
   inv[i] = val;
 invfat[0] = 1;
 invfat[MAXN - 1] = modint(f[MAXN - 1]).inv();
 for (int i = MAXN - 2; i >= 1; i--) {
    invfat[i] = invfat[i + 1] * (i + 1);
modint ncr(int n, int k) // combinacao
 modint ans = f[n] * invfat[k];
 ans *= invfat[n - k];
 return ans;
modint arr(int n, int k) // arranjo
 modint ans = f[n] * invfat[n - k];
 return ans;
modint ncr(int n, int k) {
 // calcular combinacao para n grande
 // nesse problema n <= 10^12
 // em O(k)
 modint num = 1;
 modint den = 1;
 for (int i = 0; i < k; i++) {
    num = num * modint(n - i);
```

```
den = den * modint(i + 1);
 modint ans = num / den;
 return ans;
modint stars and bars(int n, int k) {
 // para pares de inteiros n e k
  // enconte a quantidade de k-tuplas com soma == n
 // x1 + x2 + ... + xk = n
 return ncr(n + k - 1, k - 1);
```

## divisiontrick.cpp

```
for (int l = 1, r; l <= n; l = r + 1) {
 // todos os numeros i no intervalo [1, r] possuem (n / i) ==
 r = n / (n / 1);
 int x = (n / 1);
```

## crivo.cpp

```
23 lines
bitset < MAXN > prime:
vector<int> nxt(MAXN);
vector<int> factors;
void crivo() {
 prime.set();
 prime[0] = false, prime[1] = false;
  for (int i = 2; i < MAXN; i++) {
   if (prime[i]) {
     nxt[i] = i;
     for (int j = 2; j * i < MAXN; j++) {
       prime[j * i] = false;
       nxt[j * i] = i;
void fact(int n) {
 factors.clear();
  while (n > 1) {
   factors.pb(nxt[n]);
   n = n / nxt[n];
```

#### gaussianelimination.cpp

```
33 lines
#define EPS 1e-9
vector<double> ans:
int gauss(vector<vector<double>> a) {
 int n = a.size(), m = a[0].size() - 1, ret = 1;
  ans.assign(m, 0);
  vector<int> where (m, -1);
  for (int col = 0, row = 0; col < m && row < n; col++, row++)
     {
   int sel = row;
   for (int i = row; i < n; i++)
     if (abs(a[i][col]) > abs(a[sel][col])) sel = i;
    if (abs(a[sel][col]) < EPS) continue;
    for (int i = col; i \le m; i++) swap(a[sel][i], a[row][i]);
   where[col] = row;
    for (int i = 0; i < n; i++) {
     if (i != row) {
       double c = a[i][col] / a[row][col];
        for (int j = col; j \le m; j++) a[i][j] -= a[row][j] * c
```

```
for (int i = 0; i < m; i++) {
 if (where[i] != -1)
   ans[i] = (a[where[i]][m] / a[where[i]][i]);
   ret = 2;
for (int i = 0; i < n; i++) {
 double sum = 0;
  for (int j = 0; j < m; j++) sum += (ans[j] * a[i][j]);
  if (abs(sum - a[i][m]) > EPS) ret = 0;
return ret; // 0 = nao existe solucao, 1 = existe uma
     solucao, 2 = existem multiplas solucoes
```

## gaussianbitset.cpp

5 lines

```
bitset < MAXN > ans;
int gauss(vector<bitset<MAXN>>& a) {
 ans.reset();
 int n = a.size(), m = a[0].size() - 1, ret = 1;
 vector<int> where (m, -1);
 for (int col = 0, row = 0; col < m && row < n; col++) {
    for (int i = row; i < n; i++) {
     if (a[i][col]) {
       swap(a[i], a[row]);
       break;
    if (!a[row][col]) continue;
    where[col] = row;
   for (int i = 0; i < n; i++)
     if (i != row && a[i][col]) a[i] ^= a[row];
 for (int i = 0; i < m; i++) {
   if (where[i] != -1)
     ans[i] = (a[where[i]][m] / a[where[i]][i]);
    else
     ret = 2;
 for (int i = 0; i < n; i++) {
   double sum = 0;
    for (int j = 0; j < m; j++) sum += (ans[j] * a[i][j]);
   if (abs(sum - a[i][m]) > EPS) ret = 0;
 return ret;
```

# Data structures (3)

#### segtree.cpp

```
25 lines
struct segtree {
 int n;
 vector<int> seq;
 int neutral() { return 0; }
 int merge(int a, int b) { return a + b; }
 void build(vector<int> &v) {
   n = 1;
   while (n < v.size()) n <<= 1;
   seq.assign(n << 1, neutral());</pre>
    for (int i = 0; i < v.size(); i++) seg[i + n] = v[i];
    for (int i = n - 1; i; i--) seg[i] = merge(seg[i << 1], seg
         [(i << 1) | 1]);
```

```
void upd(int i, int value) {
    seg[i += n] += value;
    for (i >>= 1; i; i >>= 1) seg[i] = merge(seg[i << 1], seg[(
        i << 1) | 11);
 int qry(int 1, int r) {
   int ansl = neutral(), ansr = neutral();
    for (1 += n, r += n + 1; 1 < r; 1 >>= 1, r >>= 1) {
     if (1 & 1) ansl = merge(ansl, seg[1++]);
     if (r \& 1) ansr = merge(seg[--r], ansr);
    return merge(ansl, ansr);
};
```

```
segtreelazy.cpp
struct segtree {
  vector<int> v;
  vector<int> seq;
  vector<int> lazy;
  segtree(int sz) {
    seq.assign(4 * n, 0);
    lazy.assign(4 * n, 0);
 int single(int x) { return x; }
  int neutral() { return 0; }
  int merge(int a, int b) { return a + b; }
  void add(int i, int l, int r, int diff) {
    seg[i] += (r - 1 + 1) * diff;
    if (1 != r) {
     lazv[i << 1] += diff;
     lazv[(i << 1) | 1] += diff;
    lazy[i] = 0;
 void update(int i, int l, int r, int gl, int gr, int diff) {
    if (lazv[i]) add(i, l, r, lazv[i]);
    if (1 > r \mid | 1 > qr \mid | r < ql) return;
    if (1 >= q1 \&\& r <= qr) {
      add(i, l, r, diff);
      return:
    int mid = (1 + r) >> 1;
    update(i << 1, 1, mid, q1, qr, diff);
    update((i << 1) | 1, mid + 1, r, ql, qr, diff);
    seq[i] = merge(seq[i << 1], seq[(i << 1) | 1]);
 int query(int 1, int r, int q1, int qr, int i) {
    if (lazy[i]) add(i, l, r, lazy[i]);
    if (1 > r \mid \mid 1 > qr \mid \mid r < ql) return neutral();
    if (1 >= ql \&\& r <= qr) return seq[i];
    int mid = (1 + r) >> 1;
    return merge(query(1, mid, q1, qr, i << 1), query(mid + 1,
         r, ql, qr, (i << 1) | 1));
 void build(int 1, int r, int i) {
   if (1 == r) {
      seg[i] = single(v[l]);
      return:
    int mid = (1 + r) >> 1;
    build(1, mid, i << 1);
    build(mid + 1, r, (i << 1) | 1);
    seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
```

int  $qry(int 1, int r) \{ return query(0, n - 1, 1, r, 1); \}$ 

```
void upd(int 1, int r, int x) { update(1, 0, n - 1, 1, r, x); };
```

## fenwick.cpp

33 lines

```
struct fenw {
 int n;
  vector<int> bit;
  fenw() {}
  fenw(int sz) +
   n = sz;
   bit.assign(sz + 1, 0);
  int qry(int r) // query de prefixo a[0] + a[1] + ... a[r]
   for (int i = r + 1; i > 0; i -= i \& -i) ret += bit[i];
    return ret;
  void upd(int r, int x) // a[r] += x
   for (int i = r + 1; i \le n; i += i \& -i) bit[i] += x;
  int bs(int x) // retorna o maior indice i (i < n) tal que:
      qrv(i) < x
   int i = 0, k = 0;
   while (1 << (k + 1) <= n) k++;
    while (k \ge 0) {
     int nxt i = i + (1 << k);
     if (nxt_i <= n && bit[nxt_i] < x) {
       i = nxt_i;
       x -= bit[i];
     k--;
    return i - 1;
};
```

#### treap.cpp

93 lines

```
struct treap {
  int data, priority;
  int sz, lazy2;
 bool lazy;
 treap *1, *r, *parent;
int size(treap *node) { return (!node) ? 0 : node->sz; }
void recalc(treap *node) {
  if (!node) return;
  node -> sz = 1;
  node->parent = 0;
  if (node->1) node->sz += node->1->sz, node->1->parent = node;
  if (node->r) node->sz += node->r->sz, node->r->parent = node;
void lazy_propagation(treap *node) {
  if (node == NULL) return;
  if (node->lazy2) {
    if (node->1) node->1->lazy2 += node->lazy2;
    if (node->r) node->r->lazy2 += node->lazy2;
   node->data += node->lazy2;
   node \rightarrow lazv2 = 0;
  if (node->lazy) {
    swap(node->1, node->r);
    if (node->1) node->1->lazy = !node->1->lazy;
    if (node->r) node->r->lazy = !node->r->lazy;
```

```
node -> lazy = 0;
void split(treap *t, treap *&l, treap *&r, int n) {
 if (!t) return void(1 = r = 0);
 lazy_propagation(t);
 if (size(t->1) >= n)
    split(t->1, 1, t->1, n), r = t;
    split(t->r, t->r, r, n - size(t->1) - 1), 1 = t;
 recalc(t);
void merge(treap *&t, treap *1, treap *r) {
 lazy propagation(1);
 lazy_propagation(r);
 if (!1)
   t = r;
 else if (!r)
   t = 1:
 else if (l->priority > r->priority)
   merge(1->r, 1->r, r), t=1;
   merge(r->1, 1, r->1), t = r;
 recalc(t);
void troca(treap *&t, int 1, int r, int 11, int rr) // sap de
 treap *a0, *a1, *b0, *b1, *c0, *c1, *d0, *d1;
 split(t, a0, a1, 1);
 split(a1, b0, b1, r - 1 + 1);
 11 -= (r + 1);
 rr -= (r + 1);
 split(b1, c0, c1, l1);
 split(c1, d0, d1, rr - l1 + 1);
 merge(t, a0, d0);
 merge(t, t, c0);
 merge(t, t, b0);
 merge(t, t, d1);
void add(treap *&t, int 1, int r) {
 treap *a0, *a1, *b0, *b1;
 split(t, a0, a1, 1);
 split(a1, b0, b1, r - 1 + 1);
 b0 -> lazy ^= 1;
 b0 -> lazy2 += 1;
 merge(t, a0, b0);
 merge(t, t, b1);
void dfs(treap *t) {
 if (!t) return;
 lazy propagation(t);
 dfs(t->1);
 solve(t->data);
 dfs(t->r);
treap *create_node(int data, int priority) {
 treap *ret = new treap;
 ret->data = data;
 ret->priority = priority;
 ret->1 = 0;
 ret->r = 0;
 ret->sz = 1;
 ret->lazy = 0;
 ret -> lazy2 = 0;
 ret->parent = 0;
 return ret;
```

```
| colorupdate.cpp
```

74 lines const int inf = 1e15; struct color\_upd { #define left fir #define right sec.fir #define color sec.sec set<pii> ranges; vector<pii> erased; color\_upd(int n) // inicialmente, todo mundo pintado com a // nao usar cores negativas!!!!!!!! ranges.insert( $\{0, \{n-1, inf\}\}\);$ int get(int i) { auto it = ranges.upper bound({i, {1e18, 1e18}}); if (it == ranges.begin()) return -1; return (\*(it)).color; void del(int 1, int r) // apaga o intervalo [1, r] erased.clear(); auto it = ranges.upper\_bound({1, {0, 0}}); if (it != ranges.begin()) { while (it != ranges.end()) { if ((\*(it)).left > r) break: else if ((\*(it)).right >= 1)erased.push back(\*it); it++; if (erased.size() > 0) { int sz = erased.size(); auto it = ranges.lower\_bound({erased[0].left, {0, 0}}); auto it2 = ranges.lower\_bound({erased[sz - 1].left, {0, 0} }}); pii ini = \*it, fim = \*it2; it2++; ranges.erase(it, it2); pii upd1 =  $\{ini.left, \{l-1, ini.color\}\};$ pii upd2 = {r + 1, {fim.right, fim.color}}; erased[0].left = max(erased[0].left, 1); erased[sz - 1].right = min(erased[sz - 1].right, r); if (upd1.left <= upd1.right) ranges.insert(upd1);</pre> if (upd2.left <= upd2.right) ranges.insert(upd2);</pre> void add(int a, int b, int c) { auto it = ranges.lower\_bound({a, {b, 0}}); pii aa =  $\{-1, \{-1, -1\}\};$ pii bb =  $\{-1, \{-1, -1\}\};$ if (it != ranges.end()) { if  $((*it).color == c && (*it).left == b + 1) {$ aa = \*it;b = (\*it).right;if (it != ranges.begin()) { if  $((*it).color == c && (*it).right == a - 1) {$ bb = \*it;a = (\*it).left;ranges.erase(aa);

```
ranges.erase(bb);
ranges.insert({a, {b, c}});
}
void upd(int a, int b, int c) // pinta o intervalo [a, b]
        com a cor c
{
    del(a, b);
    add(a, b, c);
}
};
```

## bit2d.cpp

```
51 lines
struct bit2d {
 vector<int> ord;
 vector<vector<int>> t:
 vector<vector<int>> coord;
  bit2d(vector<pi> &pts) // recebe todos os pontos que vao ser
       inseridos pra construir, mas nao insere eles
   sort(pts.begin(), pts.end());
   for (auto const &a : pts) {
     if (ord.empty() || a.fir != ord.back()) ord.pb(a.fir);
   t.resize(ord.size() + 1);
   coord.resize(t.size());
   for (auto &a : pts) {
     swap(a.fir, a.sec);
   sort(pts.begin(), pts.end());
   for (auto &a : pts) {
     swap(a.fir, a.sec);
     for (int on = upper_bound(ord.begin(), ord.end(), a.fir)
          - ord.begin(); on < t.size(); on += on & -on) {
       if (coord[on].empty() || coord[on].back() != a.sec)
            coord[on].push_back(a.sec);
   for (int i = 0; i < t.size(); i++) t[i].assign(coord[i].
        size() + 1, 0);
  void add(int x, int y, int v) // v[a][b] += v
   for (int xx = upper_bound(ord.begin(), ord.end(), x) - ord.
        begin(); xx < t.size(); xx += xx & -xx) {
     for (int yy = upper_bound(coord[xx].begin(), coord[xx].
          end(), y) - coord[xx].begin(); yy < t[xx].size();
          yy += yy \& -yy)
       t[xx][yy] += v;
  int qry(int x, int y) // soma de todos os v[a][b] com (a <=
      x && b <= y)
   int ans = 0;
   for (int xx = upper_bound(ord.begin(), ord.end(), x) - ord.
        begin(); xx > 0; xx -= xx & -xx) {
      for (int yy = upper_bound(coord[xx].begin(), coord[xx].
          end(), y) - coord[xx].begin(); yy > 0; yy -= yy & -
          уу)
       ans += t[xx][yy];
   return ans:
  int qry2(int x1, int y1, int x2, int y2) {
   return qry(x2, y2) - qry(x2, y1 - 1) - qry(x1 - 1, y2) +
        qry(x1 - 1, y1 - 1);
  void add2(int x1, int y1, int x2, int y2, int v) {
```

```
add(x1, y1, v);
    add(x1, y2 + 1, -v);
    add(x2 + 1, y1, -v);
    add(x2 + 1, y2 + 1, v);
};
mo.cpp
namespace mo
struct query {
 int idx, l, r;
int block;
vector<query> queries;
vector<int> ans;
// bool cmp(query &x, query &y) { essa funcao de ordenacao pode
      funcionar em caso de TLE
     int ablock = x.1 / MAGIC, bblock = y.1 / MAGIC;
    if (ablock != bblock) return ablock < bblock;
    if (ablock & 1) return x.r < y.r;
    return x.r > y.r;
bool cmp(query &x, query &y) {
 if (x.1 / block != y.1 / block) return x.1 / block < y.1 /
  return x.r < y.r;
void run() {
 block = (int)sqrt(n);
  sort(queries.begin(), queries.end(), cmp);
  ans.resize(queries.size());
  int c1 = 0, cr = -1, sum = 0;
  auto add = [\&] (int x) { sum += x; };
  auto rem = [\&] (int x) { sum -= x; };
  for (int i = 0; i < queries.size(); i++) {
    while (cl > queries[i].l) {
      cl--;
      add(v[cl]);
    while (cr < queries[i].r) {
      cr++:
      add(v[crl);
    while (cl < queries[i].1) {
     rem(v[cl]);
      c1++:
    while (cr > queries[i].r) {
     rem(v[cr]);
    ans[queries[i].idx] = sum;
} // namespace mo
segtree2d.cpp
                                                          33 lines
struct segtree2d {
 int n, m;
  vector<vector<int>> seq;
  int neutral() { return 0; }
  int merge(int a, int b) { return a + b; }
  segtree2d(int nn, int mm) {
    seg = vector<vector<int>>(2 * n, vector<int>(2 * m, neutral
         ()));
```

```
int qry(int x1, int y1, int x2, int y2) {
    int ret = neutral();
    int y3 = y1 + m, y4 = y2 + m;
    for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2) {
       for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /= 2) {
        if (x1 \% 2 == 1 \text{ and } y1 \% 2 == 1) \text{ ret = merge(ret, seg[}
        if (x1 \% 2 == 1 \text{ and } y2 \% 2 == 0) \text{ ret } = \text{merge}(\text{ret, seg}[
             x1][y2]);
        if (x2 \% 2 == 0 \text{ and } y1 \% 2 == 1) \text{ ret } = \text{merge}(\text{ret, seg}[
              x2][y1]);
        if (x2 \% 2 == 0 \text{ and } y2 \% 2 == 0) \text{ ret} = merge(ret, seg[}
             x2][y2]);
    return ret;
  void upd(int x, int y, int val) {
    int y2 = y += m;
    for (x += n; x; x /= 2, y = y2) {
      if (x >= n)
        seg[x][y] = val;
         seg[x][y] = merge(seg[2 * x][y], seg[2 * x + 1][y]);
      while (y \neq 2) seg[x][y] = merge(seg[x][2 * y], seg[x][2
           * y + 1]);
};
persistentseg.cpp
                                                               37 lines
struct node {
 int item, 1, r;
  node() {}
 node (int 1, int r, int item) : l(1), r(r), item(item) {}
int n, q;
vector<node> seg;
vector<int> roots;
void init() { seq.resize(1); }
int newleaf(int vv) {
 int p = seq.size();
  seg.pb(node(0, 0, vv));
  return p;
int newpar(int 1, int r) {
  int p = seg.size();
  seg.pb(node(l, r, seg[l].item + seg[r].item));
  return p;
int upd(int i, int l, int r, int pos) {
  if (1 == r) return newleaf(seg[i].item + 1);
  int mid = (1 + r) >> 1;
  if (pos <= mid) return newpar(upd(seg[i].1, 1, mid, pos), seg</pre>
  return newpar(seg[i].1, upd(seg[i].r, mid + 1, r, pos));
int build(int 1, int r) {
  if (1 == r) return newleaf(0);
  int mid = (1 + r) >> 1;
  return newpar(build(1, mid), build(mid + 1, r));
```

int qry(int vl, int vr, int l, int r, int k) {

int c = seg[seg[vr].1].item - seg[seg[vl].1].item;

return qry(seg[v1].r, seg[vr].r, mid + 1, r, k - c);

if  $(c \ge k)$  return gry(seg[v1].1, seg[vr].1, 1, mid, k);

if (1 == r) return 1;

int mid = (1 + r) >> 1;

```
rmq.cpp
struct rmq {
 vector<vector<pi>> m;
 vector<int> log;
 rmq() {}
  rmg(vector<pi> &v) {
   n = v.size();
   log.resize(n + 1);
   log[1] = 0;
   for (int i = 2; i \le n; i++) log[i] = log[i / 2] + 1;
   int sz = log[n] + 2;
   m = vector<vector<pi>>>(sz, vector<pi>(n + 1));
   for (int i = 0; i < n; i++) {
     m[0][i] = v[i];
   for (int j = 1; j < sz; j++) {
     for (int i = 0; i + (1 << j) <= n; i++) m[j][i] = min(m[j])
           -1[i], m[j -1][i + (1 << (j -1))]);
 int qry(int a, int b) { return min(m[log[b - a + 1]][a], m[
      log[b - a + 1]][b - (1 << log[b - a + 1]) + 1]).second;
```

## binarylifting.cpp

};

```
item st[MAXN][21];
for (int i = 0; i < n; i++) {
 st[i][0].nxt = min(i + 1, n - 1);
  st[i][0].sum = v[st[i][0].nxt];
for (int i = 1; i < 21; i++) {
  for (int v = 0; v < n; v++) {
   st[v][i].nxt = st[st[v][i - 1].nxt][i - 1].nxt;
   st[v][i].sum = st[v][i - 1].sum + st[st[v][i - 1].nxt][i -
        1].sum;
while (q--) {
  int l, r;
 cin >> 1 >> r;
  int ans = v[1], len = r - 1;
  for (int i = 20; i >= 0; i--) {
   if (len & (1 << i)) {
     ans += st[1][i].sum;
     l = st[l][i].nxt;
 cout << ans << endl;
```

# Graph (4)

## floydwarshall.cpp

int dist[MAXN] [MAXN];
void floyd\_warshall() {
 for (int k = 0; k < n; k++) {
 for (int i = 0; i < n; i++) {
 for (int j = 0; j < n; j++) {
 dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j]);
 }
}</pre>

```
}
}
void initialize() {
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
        if (i == j) {
            dist[i][j] = 0;
        } else {
            dist[i][j] = INF;
        }
    }
}
</pre>
```

## centroiddecomp.cpp

int n, k, resp;

```
vector<int> adj[MAXN];
vector<int> cnt;
namespace cd {
int sz;
vector<int> subtree_size;
vector<bool> visited;
void dfs(int s, int f) {
 subtree\_size[s] = 1;
  for (auto const &v : adj[s]) {
    if (v != f && !visited[v]) {
     dfs(v.s):
      subtree_size[s] += subtree_size[v];
int get_centroid(int s, int f) {
 bool is_centroid = true;
 int heaviest_child = -1;
 for (auto const &v : adj[s]) {
    if (v != f && !visited[v]) {
      if (subtree_size[v] > sz / 2) is_centroid = false;
      if (heaviest_child == -1 || subtree_size[v] >
           subtree_size[heaviest_child]) heaviest_child = v;
 return (is_centroid && sz - subtree_size[s] <= sz / 2) ? s:
      get_centroid(heaviest_child, s);
void dfs2(int s, int f, int d) {
 while (d >= cnt.size()) cnt.pb(0);
 cnt[d]++;
 for (auto const &v : adj[s]) {
    if (v != f && !visited[v]) dfs2(v, s, d + 1);
void solve(int s) {
 vector<int> tot;
 for (auto const &v : adj[s]) {
    if (visited[v]) continue;
    cnt.clear();
    dfs2(v, s, 1);
    for (int i = 1; i < cnt.size(); i++) {
      if (k - i < tot.size() \&\& k - i >= 1) resp += (cnt[i] *
          tot[k - i]);
    for (int i = 1; i < cnt.size(); i++) {
     while (i >= tot.size()) tot.pb(0);
      tot[i] += cnt[i];
```

if (k < tot.size()) resp += tot[k];

```
int decompose_tree(int s) {
    sz = 0;
    dfs(s, s);
    int cend_tree = get_centroid(s, s);
    visited[cend_tree] = true;
    solve(cend_tree);
    for (auto const &v : adj[cend_tree]) {
        if (!visited[v]) decompose_tree(v);
    }
    return cend_tree;
}
void init() {
    subtree_size.resize(n);
    visited.resize(n);
    decompose_tree(0);
}
} // namespace cd
```

## dsu.cpp

24 line

```
struct dsu {
 int tot;
  vector<int> parent;
  vector<int> sz;
 dsu(int n) {
    parent.resize(n);
    sz.resize(n);
    tot = n;
    for (int i = 0; i < n; i++) {
     parent[i] = i;
      sz[i] = 1;
  int find_set(int i) { return parent[i] = (parent[i] == i) ? i
       : find_set(parent[i]); }
  void make_set(int x, int y) {
   x = find_set(x), y = find_set(y);
    if (x != y) {
      if (sz[x] > sz[y]) swap(x, y);
      parent[x] = y;
      sz[y] += sz[x];
      tot--;
};
```

## dsubipartido.cpp

41 lines

```
struct dsu {
 vector<pi> parent;
 vector<int> rank;
 vector<int> bipartite;
 dsu(int n) {
   parent.resize(n);
   rank.resize(n);
   bipartite.resize(n);
    for (int v = 0; v < n; v++) {
     parent[v] = \{v, 0\};
     rank[v] = 0;
     bipartite[v] = 1;
 dsu() {}
 pi find_set(int v) {
   if (v != parent[v].fir) {
     int parity = parent[v].sec;
     parent[v] = find_set(parent[v].fir);
     parent[v].sec ^= parity;
```

```
return parent[v];
void add_edge(int a, int b) {
 pi pa = find_set(a);
 a = pa.fir;
 int x = pa.sec;
 pi pb = find_set(b);
 b = pb.fir;
  int y = pb.sec;
  if (a == b) {
   if (x == y) bipartite[a] = 0;
   if (rank[a] < rank[b]) swap(a, b);</pre>
   parent[b] = \{a, x^y ^1\};
   bipartite[a] &= bipartite[b];
   if (rank[a] == rank[b]) rank[a]++;
bool is_bipartite(int v) { return bipartite[find_set(v).fir];
```

## cycledetection.cpp

```
int n, m, idx;
vector<int> cycles[MAXN];
vector<int> adj[MAXN];
int color[MAXN];
int parent[MAXN];
int ans[MAXN];
void dfs(int u, int p) { // chama dfs a partir de qm tem cor 0
  if (color[u] == 2) return;
  if (color[u] == 1) {
    idx++;
    int curr = p;
    ans[curr] = idx;
    cycles[idx].pb(curr);
    while (curr != u) {
     curr = parent[curr];
     cvcles[idx].pb(curr);
     ans[curr] = idx;
   return;
  parent[u] = p;
  color[u] = 1;
  for (auto const &v : adi[u])
   if (v != parent[u]) dfs(v, u);
 color[u] = 2;
```

## blockcuttree.cpp

74 lines

```
struct block cut tree {
 // Se art[i] >= 1, i eh ponto de articulação
 // tree - eh a propia block-cut tree
 // pos[i] responde a qual vertice da arvore vertice i
      pertence
  vector<vector<int>> q, blocks, tree;
  vector<vector<pi>> edgblocks; // sao as arestas do bloco i
  stack<int> s;
  stack<pi> s2;
 vector<int> id, art, pos;
  block_cut_tree(vector<vector<int>> g_) : g(g_) {
   int n = q.size();
   id.resize(n, -1), art.resize(n), pos.resize(n);
   build();
```

```
int dfs(int i, int &t, int p = -1) {
   int lo = id[i] = t++;
   s.push(i);
   if (p != -1) {
     s2.emplace(i, p);
    for (int j : g[i]) {
     if (j != p \text{ and } id[j] != -1) s2.emplace(i, j);
   for (int j : g[i]) {
    if (j != p) {
       if (id[j] == -1) {
         int val = dfs(j, t, i);
         lo = min(lo, val);
         if (val >= id[i]) {
           art[i]++;
           blocks.emplace_back(1, i);
           while (blocks.back().back() != j) {
             blocks.back().pb(s.top());
             s.pop();
            edgblocks.emplace back(1, s2.top());
            s2.pop();
           pi aux = {j, i};
            while (edgblocks.back().back() != aux) {
             edgblocks.back().pb(s2.top());
             s2.pop();
         // if (val > id[i]) aresta i-j eh ponte
         lo = min(lo, id[j]);
    if (p == -1 \text{ and } art[i]) {
     art[i]--;
   return lo;
 void build() {
    for (int i = 0; i < q.size(); i++) {
     if (id[i] == -1) dfs(i, t, -1);
   tree.resize(blocks.size());
    for (int i = 0; i < q.size(); i++) {
     if (art[i]) pos[i] = tree.size(), tree.emplace back();
   for (int i = 0; i < blocks.size(); i++) {
     for (int j : blocks[i]) {
       if (!art[i])
         pos[j] = i;
         tree[i].pb(pos[j]), tree[pos[j]].pb(i);
};
```

## bridges.cpp

21 lines nt n, m, timer; vector<pi> edges; vector<bool> is\_bridge; vector<pi> adj[MAXN]; int tin[MAXN]; int low[MAXN]; // memset -1

```
bool vis[MAXN];
                         // memset -1
void dfs(int v, int p) { // chama de quem nao foi vis ainda
 vis[v] = true;
 tin[v] = timer, low[v] = timer++;
 for (auto const &u : adj[v]) {
   if (u.fir == p) continue;
   if (vis[u.fir]) {
     low[v] = min(low[v], tin[u.fir]);
     continue;
   dfs(u.fir, v);
   low[v] = min(low[v], low[u.fir]);
   if (low[u.fir] > tin[v]) is_bridge[u.sec] = 1;
```

## dinic.cpp

int run() {

6

```
#define INF 1e9
struct edge {
 int to, from, flow, capacity, id;
struct dinic {
 int n, src, sink;
 vector<vector<edge>> adj;
 vector<int> level;
 vector<int> ptr;
 dinic(int sz) {
   n = sz:
    adj.resize(n);
   level.resize(n);
   ptr.resize(n);
 void add_edge(int a, int b, int c, int id) {
    adj[a].pb({b, (int)adj[b].size(), c, c, id});
   adj[b].pb({a, (int)adj[a].size() - 1, 0, 0, id});
 bool bfs() {
   level.assign(n, -1);
   level[src] = 0;
    queue<int> q;
   q.push(src);
    while (!q.empty()) {
     int u = q.front();
     q.pop();
      for (auto at : adj[u]) {
       if (at.flow && level[at.to] == -1) {
         q.push(at.to);
         level[at.to] = level[u] + 1;
   return level[sink] != -1;
 int dfs(int u, int flow) {
   if (u == sink || flow == 0) return flow;
    for (int p = ptr[u]; p < adj[u].size(); p++) {
      edge &at = adi[u][p];
     if (at.flow && level[u] == level[at.to] - 1) {
       int kappa = dfs(at.to, min(flow, at.flow));
       at.flow -= kappa;
       adj[at.to][at.from].flow += kappa;
       if (kappa != 0) return kappa;
   return 0;
```

```
int max flow = 0;
    while (bfs()) {
     ptr.assign(n, 0);
     while (1) {
       int flow = dfs(src, INF);
       if (flow == 0) break;
       max flow += flow;
   return max flow:
  vector<pii> cut_edges() // arestas do corte minimo
   bfs();
   vector<pii> ans;
    for (int i = 0; i < n; i++) {
     for (auto const &j : adj[i]) {
       if (level[i] != -1 && level[j.to] == -1 && j.capacity >
             0) ans.pb({j.capacity, {i, j.to}});
    }
    return ans;
  vector<int> flow_edges(int n, int m) // fluxo em cada aresta
      , na ordem da entrada
   vector<int> ans(m);
    for (int i = 0; i < n; i++) {
     for (auto const &j : adj[i])
       if (!j.capacity) ans[j.id] = j.flow;
    return ans;
};
```

## hopcroftkarp.cpp

```
137 lines
#define INF 1e9
struct hopcroft karp {
  vector<int> match;
  vector<int> dist;
 vector<vector<int>> adj;
  int n, m, t;
  hopcroft_karp(int a, int b) {
   n = a, m = b;
   t = n + m + 1;
   match.assign(t, n + m);
   dist.assign(t, 0);
   adj.assign(t, vector<int>{});
  void add_edge(int u, int v) {
    adj[u].pb(v);
   adj[v].pb(u);
 bool bfs() {
    queue<int> q;
    for (int u = 0; u < n; u++) {
     if (match[u] == n + m)
       dist[u] = 0, q.push(u);
     else
        dist[u] = INF;
    dist[n + m] = INF;
    while (!q.emptv()) {
     int u = q.front();
     if (dist[u] < dist[n + m]) {
        for (auto const &v : adj[u]) {
         if (dist[match[v]] == INF) {
```

```
dist[match[v]] = dist[u] + 1;
          q.push(match[v]);
  return dist[n + m] < INF;
bool dfs(int u) {
  if (u < n + m) {
    for (auto const &v : adj[u]) {
      if (dist[match[v]] == dist[u] + 1 && dfs(match[v])) {
        match[v] = u;
        match[u] = v;
        return true;
    dist[u] = INF;
    return false;
  return true;
vector<pi> run() {
  int cnt = 0;
  while (bfs())
    for (int u = 0; u < n; u++)
      if (match[u] == n + m && dfs(u)) cnt++;
  vector<pi> ans;
  for (int v = n; v < n + m; v++)
    if (match[v] < n + m) ans.pb(\{match[v], v\});
vector<int> mvc() // minimum vertex cover
  vector<pi> ans = run();
  vector<bool> vis(n + m, 0);
  for (int i = 0; i < n; i++) {
    if (match[i] == n + m) {
      queue<int> q;
      q.push(i);
      while (!q.empty()) {
        int x = q.front();
        q.pop();
        vis[x] = 1;
        for (auto const &y : adj[x]) {
          if (!vis[y]) {
            vis[v] = 1;
            q.push(match[y]);
  vector<int> vc:
  for (int i = 0; i < n; i++) {
   if (!vis[i]) vc.pb(i);
  for (int i = n; i < n + m; i++) {
    if (vis[i]) vc.pb(i);
  return vc;
vector<pi> mec() // minimum edge cover
  vector<pi> ans = run();
  for (int i = 0; i < n + m; i++) {
    if (match[i] == n + m \&\& adj[i].size() > 0) {
      if (i < n)
        ans.pb({i, adj[i][0]});
```

```
else
         ans.pb({adj[i][0], i});
   return ans;
// minimum path cover on dag
// minimum set of paths such that each of the vertices belongs
    to exactly one path
vector<vector<int>> mpc(int n, vector<pi> &e) {
 hopcroft_karp h(n, n);
  for (auto const &i : e) h.add_edge(i.fir, n + i.sec);
 vector<pi> mat = h.run();
 vector<int> prv(n, -1);
 vector<int> nxt(n, -1);
 for (int i = 0; i < mat.size(); i++) {
   nxt[mat[i].fir] = mat[i].sec - n;
   prv[mat[i].sec - n] = mat[i].fir;
 vector<vector<int>> ans;
 for (int i = 0; i < n; i++) {
   if (prv[i] == -1 && nxt[i] == -1) {
     ans.pb(\{i\});
   } else if (prv[i] == -1) {
     vector<int> curr;
     int x = i;
     while (1) {
       curr.pb(x);
       if (nxt[x] == -1) break;
       x = nxt[x];
      ans.pb(curr);
 return ans;
```

## pushrelabel.cpp

```
#define INF 1e9
struct edge {
 int dest, back, f, c, id;
struct push_relabel {
 int n;
 vector<vector<edge>> g;
 vector<int> ec;
 vector<edge *> cur;
 vector<vector<int>> hs;
 vector<int> H:
 push_relabel(int sz) : q(sz), ec(sz), cur(sz), hs(2 * sz), H(
      sz) { n = sz; }
 void add_edge(int s, int t, int cap, int rcap, int id) {
   if (s == t) return;
   g[s].pb({t, (int)g[t].size(), 0, cap, id});
   g[t].pb({s, (int)g[s].size() - 1, 0, rcap, -1});
 void add_flow(edge &e, int f) {
    edge &back = q[e.dest][e.back];
   if (!ec[e.dest] && f) hs[H[e.dest]].push_back(e.dest);
   e.f += f;
   e.c -= f;
   ec[e.dest] += f;
   back.f -= f;
   back.c += f;
   ec[back.dest] -= f;
 int calc(int s, int t) {
```

```
int v = q.size();
    H[s] = v;
    ec[t] = 1;
    vector<int> co(2 * v);
    co[0] = v - 1;
    for (int i = 0; i < v; i++) cur[i] = g[i].data();
    for (edge &e : g[s]) add_flow(e, e.c);
    for (int hi = 0;;) {
     while (hs[hi].empty())
       if (!hi--) return -ec[s];
     int u = hs[hi].back();
     hs[hi].pop_back();
     while (ec[u] > 0) {
       if (cur[u] == g[u].data() + g[u].size()) {
         H[u] = INF;
          for (edge &e : g[u])
           if (e.c && H[u] > H[e.dest] + 1) H[u] = H[e.dest] +
                 1, cur[u] = &e;
          if (++co[H[u]], !--co[hi] && hi < v)
            for (int i = 0; i < v; i++)
              if (hi < H[i] && H[i] < v) --co[H[i]], H[i] = v +
         hi = H[u];
       } else if (cur[u]->c && H[u] == H[cur[u]->dest] + 1)
          add_flow(*cur[u], min(ec[u], cur[u]->c));
        else
          ++cur[u];
  vector<int> flow_edges(int m) // fluxo em cada aresta
   vector<int> ans(m);
    for (int i = 0; i < n; i++) {
     for (auto const &j : q[i]) {
       if (j.id != -1) ans[j.id] = j.f;
   }
    return ans;
};
struct flow with demands {
  push_relabel pr;
  vector<int> in, out;
  int n:
  flow with demands(int sz) : n(sz), pr(sz + 2), in(sz), out(sz
  void add edge (int u, int v, int cap, int dem, int id) {
   pr.add edge(u, v, cap - dem, 0, id);
   out[u] += dem, in[v] += dem;
  int run(int s, int t) {
   pr.add edge(t, s, INF, 0, -1);
   for (int i = 0; i < n; i++) {
     pr.add_edge(n, i, in[i], 0, -1);
     pr.add_edge(i, n + 1, out[i], 0, -1);
   return pr.calc(n, n + 1);
  bool check() // todas as constraints foram satisfeitas?
    for (auto const &i : pr.q[n]) {
     if (i.c > 0) return 0;
   return 1;
};
```

```
hld.cpp
                                                          65 lines
struct hld {
 int n, cur_pos;
 segtree seg;
 vector<vector<int>> adj;
 vector<int> parent, depth, heavy, head, pos, sz;
 int dfs(int s) {
   int size = 1, max_c_size = 0;
   for (auto const &c : adi[s]) {
     if (c != parent[s]) {
       parent[c] = s;
       depth[c] = depth[s] + 1;
       int c_size = dfs(c);
       size += c_size;
       if (c_size > max_c_size) max_c_size = c_size, heavy[s]
   return sz[s] = size;
 void decompose(int s, int h) {
   head[s] = h;
   pos[s] = cur pos++;
   if (\text{heavy}[s] != -1) decompose(\text{heavy}[s], h);
   for (int c : adj[s]) {
     if (c != parent[s] && c != heavy[s]) decompose(c, c);
 hld(vector<vector<int>> &g) {
   n = q.size();
   adj = q;
   seg = segtree(n);
   parent.assign(n, -1);
   depth.assign(n, -1);
   heavy.assign(n, -1);
   head.assign(n, -1);
   pos.assign(n, -1);
   sz.assign(n, 1);
   cur pos = 0;
   dfs(0);
   decompose(0, 0);
 int guery path (int a, int b) {
   int res = 0;
   for (; head[a] != head[b]; b = parent[head[b]]) {
     if (depth[head[a]] > depth[head[b]]) swap(a, b);
     res += seq.query(0, n - 1, pos[head[b]], pos[b], 1);
   if (depth[a] > depth[b]) swap(a, b);
   res += seq.query(0, n - 1, pos[a], pos[b], 1);
   return res;
 void update_path(int a, int b, int x) {
   for (; head[a] != head[b]; b = parent[head[b]]) {
     if (depth[head[a]] > depth[head[b]]) swap(a, b);
     seg.update(1, 0, n - 1, pos[head[b]], pos[b], x);
   if (depth[a] > depth[b]) swap(a, b);
   seq.update(1, 0, n - 1, pos[a], pos[b], x);
 void update_subtree(int a, int x) { seq.update(1, 0, n - 1,
      pos[a], pos[a] + sz[a] - 1, x); }
 int query_subtree(int a) { return seg.query(0, n - 1, pos[a],
       pos[a] + sz[a] - 1, 1); }
 int lca(int a, int b) {
   if (pos[a] < pos[b]) swap(a, b);
   return (head[a] == head[b]) ? b : lca(parent[head[a]], b);
```

```
hldedge.cpp
                                                           67 lines
namespace hld {
int cur_pos;
vector<int> parent, depth, heavy, head, pos, sz, up;
int dfs(int s) {
 int size = 1, max c size = 0;
 for (auto const &c : adj[s]) {
   if (c.fir != parent[s]) {
      parent[c.fir] = s;
      depth[c.fir] = depth[s] + 1;
      int c size = dfs(c.fir);
      size += c size;
      if (c_size > max_c_size) max_c_size = c_size, heavy[s] =
          c.fir:
 return sz[s] = size;
void decompose (int s, int h) {
 head[s] = h;
  pos[s] = cur_pos++;
  seq::v[pos[s]] = up[s];
  for (auto const &c : adi[s]) {
    if (c.fir != parent[s] && c.fir == heavy[s]) {
      up[c.fir] = c.sec;
      decompose (heavy[s], h);
  for (auto const &c : adj[s]) {
    if (c.fir != parent[s] && c.fir != heavy[s]) {
      up[c.fir] = c.sec;
      decompose(c.fir, c.fir);
void init() {
 parent.assign(MAXN, -1);
  depth.assign(MAXN, -1);
 heavy.assign(MAXN, -1);
  head.assign(MAXN, -1);
 pos.assign(MAXN, -1);
  sz.assign(MAXN, 1);
 up.assign(MAXN, 0);
 cur_pos = 0;
  dfs(0);
  decompose(0, 0);
 seg::build(0, n - 1, 1);
int query_path(int a, int b) {
 int res = -1;
 for (; head[a] != head[b]; b = parent[head[b]]) {
    if (depth[head[a]] > depth[head[b]]) swap(a, b);
    res = max(res, seg::query(0, n - 1, pos[head[b]], pos[b],
        1));
 if (depth[a] > depth[b]) swap(a, b);
  res = \max(\text{res}, \text{seq}::\text{query}(0, n - 1, pos[a] + 1, pos[b], 1));
  return res;
void update_path(int a, int b, int x) {
 for (; head[a] != head[b]; b = parent[head[b]]) {
    if (depth[head[a]] > depth[head[b]]) swap(a, b);
    seg::update(1, 0, n - 1, pos[head[b]], pos[b], x);
```

if (depth[a] > depth[b]) swap(a, b);

seg::update(1, 0, n - 1, pos[a] + 1, pos[b], x);

} // namespace hld

## strongorientation twosat sack reroot hungarian

```
strongorientation.cpp
                                                          25 lines
// encontrar uma orientacao para as arestas tal que o numero
// minimo de scc eh o menor possivel
// numero minimo de scc = numero de componentes conexas +
    numero de pontes
int n, m, timer, comps, bridges;
vector<pi> edges;
vector<pi> adj[MAXN];
int tin[MAXN]; // memset -1
int low[MAXN]; // memset -1
bool vis[MAXN];
char orient[MAXN];
void find_bridges(int v) { // chama se tem tin == -1
  low[v] = timer, tin[v] = timer++;
  for (auto const &p : adj[v]) {
    if (vis[p.sec]) continue;
   vis[p.sec] = true;
   orient[p.sec] = (v == edges[p.sec].first) ? '>' : '<';
    if (tin[p.fir] == -1) {
     find_bridges(p.fir);
     low[v] = min(low[v], low[p.fir]);
     if (low[p.fir] > tin[v]) bridges++;
     low[v] = min(low[v], low[p.fir]);
```

void update\_subtree(int a, int x) { seq::update(1, 0, n - 1,

int query\_subtree(int a, int x) { return seg::query(0, n - 1,

 $pos[a] + 1, pos[a] + sz[a] - 1, x); }$ 

pos[a] + 1, pos[a] + sz[a] - 1, 1);

#### twosat.cpp

74 lines

if (i != p) { dfs\_sz(i, x);

sz[x] += sz[i];

void modify(int c, int val) { cnt[c] += val; }

void dfs(int x, int p, bool keep) {

int best = -1, big\_child = -1;

```
struct two sat {
 vector<vector<int>> g, gr; // gr is the reversed graph
 vector<int> comp, ord, ans; // comp[v]: ID of the SCC
      containing node v
 vector<bool> vis:
 two sat() {}
  two_sat(int sz) {
   n = sz;
   g.assign(2 * n, vector<int>());
   gr.assign(2 * n, vector<int>());
   comp.resize(2 * n);
   vis.resize(2 * n);
   ans.resize(2 * n);
  void add_edge(int u, int v) {
   q[u].push back(v);
   gr[v].push_back(u);
  // int x, bool val: if 'val' is true, we take the variable to
       be x. Otherwise we take it to be x's complement (not x)
  void implies(int i, bool f, int j, bool g) // a -> b
   add_edge(i + (f ? 0 : n), j + (g ? 0 : n));
   add_edge(j + (g ? n : 0), i + (f ? n : 0));
  void add_clause_or(int i, bool f, int j, bool g) // At least
       one of them is true
   add_edge(i + (f ? n : 0), j + (g ? 0 : n));
```

```
add_edge(j + (g ? n : 0), i + (f ? 0 : n));
 void add_clause_xor(int i, bool f, int j, bool g) // only
      one of them is true
    add_clause_or(i, f, j, g);
    add_clause_or(i, !f, j, !g);
 void add_clause_and(int i, bool f, int j, bool g) // both of
       them have the same value
    add_clause_xor(i, !f, j, g);
 void set(int i, bool f) // Set a variable
    add_clause_or(i, f, i, f);
 void top_sort(int u) {
   vis[u] = 1;
    for (auto const &v : g[u]) {
     if (!vis[v]) top_sort(v);
    ord.push back(u);
 void scc(int u, int id) {
   vis[u] = 1;
    comp[u] = id;
    for (auto const &v : gr[u]) {
     if (!vis[v]) scc(v, id);
 bool solve() {
   fill(vis.begin(), vis.end(), 0);
    for (int i = 0; i < 2 * n; i++) {
     if (!vis[i]) top_sort(i);
    fill(vis.begin(), vis.end(), 0);
    reverse(ord.begin(), ord.end());
    int id = 0;
    for (const auto &v : ord) {
     if (!vis[v]) scc(v, id++);
    for (int i = 0; i < n; i++) {
     if (comp[i] == comp[i + n]) return 0;
      ans[i] = (comp[i] > comp[i + n]) ? 1 : 0;
   return 1;
};
sack.cpp
                                                          47 lines
vector<int> adj[MAXN];
vector<int> v[MAXN];
int c[MAXN];
int cnt[MAXN];
int sz[MAXN];
void dfs_sz(int x, int p) {
 sz[x] = 1;
 for (auto const &i : adj[x]) {
```

```
for (auto const &i : adj[x]) {
  if (i != p && sz[i] > best) {
    best = sz[i];
    big_child = i;
for (auto const &i : adj[x]) {
 if (i != p && i != big_child) dfs(i, x, 0);
if (big child !=-1) {
  dfs(big_child, x, 1);
  swap(v[x], v[big_child]); // O(1)
; (x) dq. [x] v
modify(c[x], 1); // adiciona
for (auto const &i : adj[x]) {
 if (i != p && i != big_child) {
    for (auto const &j : v[i]) {
     v[x].pb(j);
      modify(c[j], 1); // adiciona
// a cor c aparece cnt[c] vezes na subtree de x
// dai vc pode fazer algo tendo essa informacao
// seja responser queries ou algo do tipo aqui
  for (auto const &i : v[x]) modify(c[i], -1); // remove
```

#### reroot.cpp

```
int n;
vector<int> adj[MAXN];
int sz[MAXN];
int dp[MAXN];
int dfs(int u, int v) {
 sz[u] = 1;
 for (auto const &i : adj[u])
   if (i != v) sz[u] += dfs(i, u);
 return sz[u];
void reroot(int u, int v) {
 for (auto const &i : adj[u]) {
   if (i != v) {
     int a = sz[u], b = sz[i];
     dp[i] = dp[u];
     dp[i] = sz[u], dp[i] = sz[i];
     sz[u] = sz[i], sz[i] = n;
     dp[i] += sz[u], dp[i] += sz[i];
     reroot(i, u);
     sz[u] = a, sz[i] = b;
```

#### hungarian.cpp

```
struct hungarian { // declarar algo como hungarian h(n), n de
    cada lado
 int n, inf;
 vector<vector<int>> a;
 vector<int> u, v, p, way;
 hungarian(int n_{-}): n(n_{-}), u(n + 1), v(n + 1), p(n + 1), way(
    a = vector<vector<int>>(n, vector<int>(n));
    inf = numeric_limits<int>::max();
```

## mincostmaxflow scc articulationpoints notes

```
void add_edge(int x, int y, int c) { a[x][y] = c; }
  pair<int, vector<int>> run() {
    for (int i = 1; i \le n; i++) {
     p[0] = i;
      int j0 = 0;
      vector<int> minv(n + 1, inf);
      vector < int > used(n + 1, 0);
      do {
       used[j0] = true;
        int i0 = p[j0], j1 = -1;
        int delta = inf;
        for (int j = 1; j \le n; j++) {
          if (!used[j]) {
            int cur = a[i0 - 1][j - 1] - u[i0] - v[j];
            if (cur < minv[j]) minv[j] = cur, way[j] = j0;
            if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
        for (int j = 0; j \le n; j++) {
         if (used[j])
            u[p[j]] += delta, v[j] -= delta;
          else
            minv[i] -= delta;
        j0 = j1;
      } while (p[j0] != 0);
      do {
       int j1 = wav[j0];
       p[j0] = p[j1];
        j0 = j1;
     } while (j0);
    vector<int> ans(n);
    for (int j = 1; j \le n; j++) ans[p[j] - 1] = j - 1;
    return make_pair(-v[0], ans);
};
```

## mincostmaxflow.cpp

64 lines

scc.cpp

int n, m;

bool vis[MAXN];

cin >> n >> m;

cin >> a >> b;

adj\_rev[b].pb(a);

adj[a].pb(b);

int a, b;

for (int i = 0; i < m; i++) {

```
namespace mcf {
struct edge {
  int to, capacity, cost, res;
int source, destiny;
vector<edge> adj[MAXN];
vector<int> dist;
vector<int> parent;
vector<int> edge_index;
vector<bool> in_queue;
void add_edge(int a, int b, int c, int d) {
  adj[a].pb({b, c, d, (int)adj[b].size()});
  adj[b].pb({a, 0, -d, (int)adj[a].size() - 1});
bool dijkstra(int s) {
  dist.assign(MAXN, INF);
  parent.assign(MAXN, -1);
  edge_index.assign(MAXN, -1);
  in queue.assign(MAXN, false);
  dist[s] = 0;
  queue<int> q;
  q.push(s);
  while (!q.empty()) {
    int u = q.front(), idx = 0;
    q.pop();
    in_queue[u] = false;
    for (auto const &v : adj[u]) {
     if (v.capacity && dist[v.to] > dist[u] + v.cost) {
        dist[v.to] = dist[u] + v.cost;
```

```
parent[v.to] = u;
       edge_index[v.to] = idx;
       if (!in_queue[v.to]) {
         in_queue[v.to] = true;
         q.push(v.to);
     idx++;
 return dist[destiny] != INF;
int get_cost() {
 int flow = 0, cost = 0;
 while (dijkstra(source)) {
    int curr_flow = INF, curr = destiny;
   while (curr != source) {
     int p = parent[curr];
     curr_flow = min(curr_flow, adj[p][edge_index[curr]].
     curr = p;
    flow += curr flow;
   cost += curr_flow * dist[destiny];
   curr = destiny;
   while (curr != source) {
     int p = parent[curr];
     int res_idx = adj[p][edge_index[curr]].res;
     adj[p][edge_index[curr]].capacity -= curr_flow;
     adj[curr][res_idx].capacity += curr_flow;
     curr = p;
 return cost;
} // namespace mcf
```

```
int root[MAXN];
vector<int> order;
vector<int> roots;
vector<int> comp;
vector<vector<int>> comps;
vector<int> adj[MAXN];
vector<int> adj rev[MAXN];
vector<int> adj_scc[MAXN];
void dfs(int v) {
 vis[v] = true;
 for (auto const &u : adj[v])
   if (!vis[u]) dfs(u);
 order.pb(v);
void dfs2(int v) {
 comp.pb(v);
 vis[v] = true;
 for (auto const &u : adj rev[v])
    if (!vis[u]) dfs2(u);
//...
```

```
for (int i = 0; i < n; i++) {
 if (!vis[i]) dfs(i);
reverse(order.begin(), order.end());
memset (vis, false, sizeof (vis));
for (auto const &v : order) {
 if (!vis[v]) {
   comp.clear();
   dfs2(v);
    comps.pb(comp);
   // making condensation graph
   int r = comp.back();
    for (auto const &u : comp) root[u] = r;
   roots.push back(r);
// making condensation graph
for (int v = 0; v < n; v++) {
 for (auto const &u : adj[v]) {
   int root_v = roots[v];
    int root_u = roots[u];
   if (root_u != root_v) adj_scc[root_v].pb(root_u);
```

#### articulationpoints.cpp

23 lines

```
int n, m, timer;
vector<int> adj[MAXN];
bool is_cutpoint[MAXN];
int tin[MAXN]; // memset -1
int low[MAXN]; // memset -1
bool vis[MAXN];
void dfs(int v, int p) { // chama pros nao vis
 vis[v] = true;
 tin[v] = timer, low[v] = timer++;
 int childs = 0;
 for (auto const &u : adj[v]) {
   if (u == p) continue;
   if (vis[u]) {
     low[v] = min(low[v], tin[u]);
    } else {
     dfs(u, v);
     low[v] = min(low[v], low[u]);
     if (low[u] >= tin[v] && p != -1) is_cutpoint[v] = true;
     childs++;
 if (p == -1 && childs > 1) is_cutpoint[v] = true;
```

## Clique

54 lines

```
notes.md
                                                          98 lines
## Bipartite Graph
A bipartite graph is a graph that does not contain any odd-
    length cycles.
## Directed acyclic graph (DAG)
Is a directed graph with no directed cycles.
## Independent Set
Is a set of vertices in a graph, no two of which are adjacent.
    That is, it is a set S of vertices such that for every two
     vertices in S, there is no edge connecting the two.
```

- Is a subset of vertices of an undirected graph such that every two distinct vertices in the clique are adjacent.
- ## Vertex Cover
- Is a set of vertices that includes at least one endpoint of every edge of the graph.
- ## Edge Cover
- Is a set of edges such that every vertex of the graph is incident to at least one edge of the set.
- ## Path Cover
- Given a directed graph G = (V, E), a path cover is a set of directed paths such that every vertex v belongs to at least one path.
- ## Koning's Theorem
- In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover.
- ## Properties
- Every tree is a bipartite graph.
- Any NxM grid is a bipartite graph.
- A set of vertices is a vertex cover if and only if its complement is an independent set.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set
- In bipartite graphs, the size of the minimum edge cover is equal to the size of the maximum independent set
- In bipartite graphs, the size of the minimum edge cover plus the size of the minimum vertex cover is equal to the number of vertices.
- In bipartite graphs, maximum clique size is two.
- ## Min-cut
- The smallest total weight of the edges which if removed would disconnect the source from the sink.
- ## Max-flow min-cut theorem
- ## Maximum flow with vertex capacities
- In other words, the amount of flow passing through a vertex cannot exceed its capacity. To find the maximum flow, we can transform the problem into the maximum flow problem by expanding the network. Each vertex v is replaced by v-in and v-out, where v-in is connected by edges going into v and v-out is connected to edges coming out from v. Then assign capacity c(v) to the edge connecting v-in and v-out
- ## Undirected edge-disjoint paths problem
- We are given an undirected graph  $G=(V,\;E)$  and two vertices s and t, and we have to find the maximum number of edgedisjoint s-t paths in G.

- ## Undirected vertex-disjoint paths problem
- ## Menger's theorem
- The maximum number of edge-disjoint s-t paths in an undirected graph is equal to the minimum number of edges in an s-t cut-set.
- ## Undirected vertex-disjoint paths solution
- We can construct a network N=(V,E) from G with vertex capacities, where the capacities of all vertices and all edges are 1. Then the value of the maximum flow is equal to the maximum number of independent paths from s to t.
- Given a directed acyclic graph G=(V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V. We can construct a bipartite graph G' from G. Each vertex v is replaced by v-in and v-out, where v-in is connected by edges going into v and v-out is connected to edges coming out from v. Then it can be shown that G' has a matching M of size m if and only if G has a vertex-disjoint path cover C of containing m edges and n-m paths.
- ## Minimum general path cover in directed acyclic graph (DAG)
- A general path cover is a path cover where a vertex can belong to more than one path. A minimum general path cover may be smaller than a minimum vertex-disjoint path cover. A minimum general path cover can be found almost like a minimum vertex-disjoint path cover. It suffices to add some new edges to the matching graph so that there is an edge a to b always when there is a path from a to b in the original graph.
- ## Dilworths theorem and maximum antichain
- An antichain is a set of nodes of a graph such that there is no path from any node to another node using the edges of the graph. Dilworths theorem states that in a directed acyclic graph, the size of a minimum general path cover equals the size of a maximum antichain.
- Or in other words: For a DAG G that if has edges from vertex i to vertex j and vertex j to k, then it also has a edge from vertex i to vertex k, the size of a minimum path cover is equal to the size of a maximum independent set.
- ## Maximum weighted antichain
- In this problem, each vertex has a cost a[i]. The cost of an antichain is equal to the sum of the costs of the vertices present in it. We need to find the maximum cost of a antichain. We can construct the same bipartite of the maximum antichain problem from a dag G, these edges have an infinite capacity. We also need to create a source vertex and a sink, and we need to add edges source to v-in with capacity a[v] and v-out to sink with capacity a[v]. The answer is equal to the sum of all a[i] minus the maximum flow on this network.

```
## Halls Theorem
Halls theorem can be used to find out whether a bipartite graph
    has a matching that contains all left or right nodes.
    Assume that we want to find a matching that contains all
    left nodes. Let X be any set of left nodes and let f(X) be
    the set of their neighbors. According to Halls theorem, a
    matching that contains all left nodes exists exactly when
    for
each X, the condition |X| <= |f(X)| holds.

## Extra (Getting Confidence Trick)

<p>If you need to maximize a number x = (a * b * c * ...), then
    you can write it as x = (e^log(a) * e^log(b) * e^log(c) *
        ...), and then the number is x = e^log(a) * log(b) * log
    (c) + ...), and the problem now becomes a problem of
    maximizing the sum of (log(a) + log(b) + log(c) + ...).
Use exp() and log() C++ functions:)
```

47 lines

# $\underline{\text{Geometry}}$ (5)

## convex hull. cpp

```
struct point {
 int x, v, id;
 point(int x, int y, int id) : x(x), y(y), id(id) {}
 point operator-(point const &o) const { return {x - o.x, y -
      o.v, -1; }
 bool operator<(point const &o) const {
   if (x == o.x) return y < o.y;
   return x < o.x;
 int operator^(point const &o) const { return x * o.y - y * o.
int ccw(point const &a, point const &b, point const &x) {
 auto p = (b - a) ^ (x - a);
 return (p > 0) - (p < 0);
vector<point> convex_hull(vector<point> P) // sem colineares
 sort(P.begin(), P.end());
 vector<point> L, U;
 for (auto p : P) {
    while (L.size() \geq 2 && ccw(L.end()[-2], L.end()[-1], p) ==
         -1) L.pop_back();
   L.push_back(p);
 reverse(P.begin(), P.end());
  for (auto p : P) {
   while (U.size() >= 2 \&\& ccw(U.end()[-2], U.end()[-1], p) ==
         -1) U.pop_back();
   U.push_back(p);
 L.insert(L.end(), U.begin(), U.end() - 1);
 return L;
vector<point> convex_hull_no_collinears(vector<point> P) //
    com colineares
 sort(P.begin(), P.end());
 vector<point> L, U;
 for (auto p : P) {
```

12

```
while (L.size() >= 2 && ccw(L.end()[-2], L.end()[-1], p) <=
          0) L.pop_back();
    L.push_back(p);
  reverse(P.begin(), P.end());
  for (auto p : P) {
    while (U.size() \ge 2 \&\& ccw(U.end()[-2], U.end()[-1], p) \le
          0) U.pop_back();
    U.push_back(p);
  L.insert(L.end(), U.begin(), U.end() - 1);
  return L;
linetrick.cpp
                                                           13 lines
pi get_line(pi x, pi y) { // um jeito normalizado de
    representar a reta entre 2 pontos
  int xx = x.fir - y.fir;
  int yy = x.sec - y.sec;
  int g = \underline{gcd(abs(xx), abs(yy))};
  if (g != 0) {
   xx /= g, yy /= g;
  if (xx < 0) {
   xx \star = -1;
   yy \star = -1;
  return {xx, yy};
polygonarea.cpp
                                                           19 lines
double area(vector<pi> fig) {
  double res = 0;
  for (unsigned i = 0; i < fig.size(); i++) {
   pi p = i ? fig[i - 1] : fig.back();
    pi q = fig[i];
    res += (p.fir - q.fir) * (p.sec + q.sec);
  return fabs(res) / 2;
int cross(pi a, pi b) { return a.fir * b.sec - a.sec * b.fir; }
double area2(vector<pi> fig) {
  double res = 0;
  for (unsigned i = 0; i < fig.size(); i++) {
   pi p = i ? fig[i - 1] : fig.back();
   pi q = fig[i];
    res += cross(p, q);
  return fabs(res) / 2;
```

# Strings (6)

## Various (7)