

Universidade Federal de Alagoas

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ICPC World Finals 2025

September 4, 2025

1	Contest	1
2	Mathematics	1
3	Data structures	6
4	Graph	9
5	Geometry	16
6	Strings	20
7	Various	22

Contest (1)

```
template.cpp
```

25 lines

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 1000006
#define mod 998244353
signed main() {
  ios_base::sync_with_stdio(false);
 cin.tie(NULL);
  mt19937 64 rng(chrono::steady clock::now().time since epoch()
 return 0;
```

runner.py

44 lines

```
import os
import subprocess
naive = "brute.cpp" # path to naive code
code = "d.cpp" # path to your code
generator = "g.cpp" # path to test generator
def compile_codes():
   os.system('q++ ' + generator + ' -o generator -O2')
    os.system('g++ ' + naive + ' -o naive -02')
   os.system('g++ ' + code + ' -o code -O2')
def generate_case():
    os.system('./generator > in');
def get_naive_output():
```

```
output = os.popen('./naive <in').read()</pre>
    return output
def get_code_output():
    # se tiver um runtime error, vai parar
    xalala = subprocess.run('./code <in', shell=True, text=True</pre>
         , capture_output=True, check=True)
    return xalala.stdout
def main():
    compile_codes()
    while True:
        generate_case()
        naive_output = get_naive_output()
        code_output = get_code_output()
        if naive_output == code_output:
            print ('ACCEPTED')
        else :
            print('FAILED\n')
            print ('ANSWER:')
            print (naive output)
            print('\nCODE OUTPUT:')
            print(code_output)
            break
if __name__ == '__main__':
    main()
pragma.cpp
                                                             5 lines
```

```
#include <iostream>
using namespace std;
#pragma GCC target ("avx2")
#pragma GCC optimize("03")
#pragma GCC optimize("unroll-loops")
```

Mathematics (2)

modint.cpp

78 lines

```
struct modint {
 int val:
 modint(int v = 0) \{ val = v \% mod; \}
 int pow(int y) {
   modint x = val;
   modint z = 1:
    while (y) {
     if (y \& 1) z *= x;
     x *= x;
     y >>= 1;
   return z.val;
 int inv() { return pow(mod - 2); }
 void operator=(int o) { val = o % mod; }
  void operator=(modint o) { val = o.val % mod; }
 void operator+=(modint o) { *this = *this + o; }
 void operator = (modint o) { *this = *this - o; }
 void operator*=(modint o) { *this = *this * o; }
  void operator/=(modint o) { *this = *this / o; }
 bool operator==(modint o) { return val == o.val; }
 bool operator!=(modint o) { return val != o.val; }
 int operator*(modint o) { return ((val * o.val) % mod); }
 int operator/(modint o) { return (val * o.inv()) % mod; }
 int operator+(modint o) { return (val + o.val) % mod; }
 int operator-(modint o) { return (val - o.val + mod) % mod; }
```

```
modint f[MAXN]:
modint inv[MAXN];
modint invfat[MAXN];
void calc() {
 f[0] = 1;
  for (int i = 1; i < MAXN; i++) {
   f[i] = f[i - 1] * i;
 inv[1] = 1;
  for (int i = 2; i < MAXN; ++i) {
   int val = mod / i;
   val = (inv[mod % i] * val) % mod;
   val = mod - val;
   inv[i] = val;
 invfat[0] = 1;
 invfat[MAXN - 1] = modint(f[MAXN - 1]).inv();
 for (int i = MAXN - 2; i >= 1; i--) {
   invfat[i] = invfat[i + 1] * (i + 1);
modint ncr(int n, int k) // combinacao
 modint ans = f[n] * invfat[k];
 ans *= invfat[n - k];
 return ans;
modint arr(int n, int k) // arranjo
 modint ans = f[n] * invfat[n - k];
 return ans;
modint ncr(int n, int k) {
 // calcular combinacao para n grande
 // nesse problema n <= 10^12
 // em O(k)
 modint num = 1;
  modint den = 1;
  for (int i = 0; i < k; i++) {
   num = num * modint(n - i);
    den = den * modint(i + 1);
 modint ans = num / den;
  return ans:
modint stars and bars(int n, int k) {
 // para pares de inteiros n e k
 // enconte a quantidade de k-tuplas com soma == n
 // x1 + x2 + ... + xk = n
 return ncr (n + k - 1, k - 1);
```

divisiontrick.cpp

for (int l = 1, r; $l \le n$; l = r + 1) { // todos os numeros i no intervalo [1, r] possuem (n / i) ==

r = n / (n / 1);int x = (n / 1);

crivo.cpp

prime.set();

bitset < MAXN > prime; vector<int> nxt(MAXN); vector<int> factors; void crivo() {

```
prime[0] = false, prime[1] = false;
  for (int i = 2; i < MAXN; i++) {
   if (prime[i]) {
     nxt[i] = i;
     for (int j = 2; j * i < MAXN; j++) {
       prime[j * i] = false;
       nxt[j * i] = i;
void fact(int n) {
 factors.clear();
 while (n > 1) {
   factors.pb(nxt[n]);
   n = n / nxt[n];
```

gaussianelimination.cpp

```
#define EPS 1e-9
vector<double> ans;
int gauss (vector<vector<double>> a) {
  int n = a.size(), m = a[0].size() - 1, ret = 1;
  ans.assign(m, 0);
  vector<int> where (m, -1);
  for (int col = 0, row = 0; col < m && row < n; col++, row++)
    int sel = row;
   for (int i = row; i < n; i++)
     if (abs(a[i][col]) > abs(a[sel][col])) sel = i;
    if (abs(a[sel][col]) < EPS) continue;
    for (int i = col; i \le m; i++) swap(a[sel][i], a[row][i]);
    where[col] = row;
    for (int i = 0; i < n; i++) {
     if (i != row) {
       double c = a[i][col] / a[row][col];
       for (int j = col; j \le m; j++) a[i][j] -= a[row][j] * c
  for (int i = 0; i < m; i++) {
   if (where[i] != -1)
     ans[i] = (a[where[i]][m] / a[where[i]][i]);
   else
     ret = 2:
  for (int i = 0; i < n; i++) {
   double sum = 0;
   for (int j = 0; j < m; j++) sum += (ans[j] * a[i][j]);
   if (abs(sum - a[i][m]) > EPS) ret = 0;
  return ret; // 0 = nao existe solucao, 1 = existe uma
      solucao, 2 = existem multiplas solucoes
```

gaussianbitset.cpp

31 lines

```
bitset<MAXN> ans;
int gauss (vector < bitset < MAXN >> & a) {
  ans reset ():
  int n = a.size(), m = a[0].size() - 1, ret = 1;
  vector<int> where (m, -1);
  for (int col = 0, row = 0; col < m && row < n; col++) {
   for (int i = row; i < n; i++) {
     if (a[i][col]) {
       swap(a[i], a[row]);
```

```
break;
  if (!a[row][col]) continue;
  where[col] = row;
  for (int i = 0; i < n; i++)
   if (i != row && a[i][col]) a[i] ^= a[row];
for (int i = 0; i < m; i++) {
  if (where[i] != -1)
   ans[i] = (a[where[i]][m] / a[where[i]][i]);
   ret = 2;
for (int i = 0; i < n; i++) {
  double sum = 0;
  for (int j = 0; j < m; j++) sum += (ans[j] * a[i][j]);
  if (abs(sum - a[i][m]) > EPS) ret = 0;
return ret;
```

lagrange.cpp

```
struct lagrange {
 int n;
 vector<modint> den;
 vector<modint> v;
 vector<modint> fat:
 vector<modint> inv fat;
 lagrange(vector<modint> &v) // f(i) = v[i], gera um
      polinomio de grau n - 1
   n = v.size();
   calc(n):
   calc den(n);
   y = v;
 void calc_den(int n) {
   den.resize(n);
   for (int i = 0; i < n; i++) {
     den[i] = inv_fat[n - i - 1] * inv_fat[i];
     if ((n - i - 1) % 2 == 1) {
       int x = (mod - den[i].val) % mod;
       den[i] = x;
 void calc(int n) {
   fat.resize(n + 1);
   inv_fat.resize(n + 1);
   fat[0] = 1;
   inv_fat[0] = 1;
   for (int i = 1; i \le n; i++) {
     fat[i] = fat[i - 1] * i;
     inv_fat[i] = fat[i].inv();
 modint get_val(int x) // complexidade: O(n)
   x %= mod;
   vector<modint> l(n);
   vector<modint> r(n);
   1[0] = 1, r[n - 1] = 1;
   for (int i = 1; i < n; i++) {
     modint cof = (x - (i - 1) + mod);
     l[i] = l[i - 1] * cof;
```

```
for (int i = n - 2; i >= 0; i--) {
     modint cof = (x - (i + 1) + mod);
     r[i] = r[i + 1] * cof;
   modint ans = 0;
    for (int i = 0; i < n; i++) {
     modint cof = l[i] * r[i];
     ans += modint(cof * y[i]) * den[i];
   return ans:
 vector<modint> find coefs() // encontra os coeficientes do
      polinomio
   int nn = n;
   int d = nn - 1;
   vector<modint> c(nn, 0);
    for (int i = 0; i < y.size(); i++) {
     c[d] += (y[i] * den[i]);
    for (int p = nn - 2; p >= 0; p--) {
     nn--;
     calc den(nn);
      for (int i = 0; i \le p; i++) {
       y[i] = (c[p + 1] * modint(i).pow(d));
       c[p] += (y[i] * den[i]);
      d--;
    return c;
};
```

```
berlekampmassev.cpp
// berlekamp massey
// mas precisa que o mod seja primo (para poder achar inverso)
// dado os n primeiros termos de uma recorrencia linear
// a[0], a[1], a[2], ..., a[n-1]
// ele acha a reccorrencia linear mais curta que da matching
    com os n primeiros valores
vector<modint> berlekamp_massey(vector<modint> x) {
 vector<modint> ls. cur:
 int lf. ld:
 for (int i = 0; i < x.size(); i++) {
   modint t = 0;
   for (int j = 0; j < cur.size(); j++) {
     t += (x[i - j - 1] * cur[j]);
   if (modint(t - x[i]).val == 0) continue;
   if (cur.empty()) {
     cur.resize(i + 1);
     lf = i;
     ld = (t - x[i]) % mod;
     continue;
   modint k = -(x[i] - t);
   k *= modint(ld).inv();
    vector<modint> c(i - lf - 1);
    for (auto const &j : ls) {
     modint curr = modint(j.val * -1) * k;
     c.pb(curr);
    if (c.size() < cur.size()) c.resize(cur.size());</pre>
    for (int j = 0; j < cur.size(); j++) {
     c[j] = c[j] + cur[j];
    if (i - lf + ls.size() >= cur.size()) {
```

crt crttrick diophantine extendedeuclidean ntt fft

```
tie(ls, lf, ld) = make_tuple(cur, i, t - x[i]);
   cur = c:
 return cur;
modint get_nth(vector<modint> rec, vector<modint> dp, int n) {
 int m = rec.size();
  vector<modint> s(m), t(m);
  s[0] = 1;
  if (m != 1)
  t[1] = 1;
  else
   t[0] = rec[0];
  auto mul = [&rec](vector<modint> v, vector<modint> w) {
   vector<modint> ans(2 * v.size());
    for (int j = 0; j < v.size(); j++) {
     for (int k = 0; k < v.size(); k++) ans[j + k] += v[j] * w
    for (int j = 2 * v.size() - 1; j >= v.size(); j--) {
     for (int k = 1; k \le v.size(); k++) ans[j - k] += ans[<math>j]
          * rec[k - 1];
   ans.resize(v.size());
   return ans;
  while (n) {
   if (n \& 1) s = mul(s, t);
   t = mul(t, t);
   n >>= 1;
  modint ret = 0;
  for (int i = 0; i < m; i++) ret += s[i] * dp[i];
  return ret:
modint guess_nth_term(vector<modint> x, int n) {
 if (n < x.size()) return x[n];
  vector<modint> coef = berlekamp massey(x); // coeficientes
      da recorrencia
  if (coef.empty()) return 0;
 return get_nth(coef, x, n);
```

crt.cpp

```
namespace crt {
vector<pi> eq;
int gcd(int a, int b, int &x, int &y) {
 if (b == 0) {
   x = 1, y = 0;
   return a;
  int x1, y1, d = gcd(b, a % b, x1, y1);
  x = y1, y = x1 - y1 * (a / b);
 return d;
pi crt() {
  int a1 = eq[0].fir, m1 = eq[0].sec;
  for (int i = 1; i < eq.size(); i++) {
    int a2 = eq[i].fir, m2 = eq[i].sec;
    int q = \underline{\hspace{0.5cm}} \gcd(m1, m2);
    if (a1 % g != a2 % g) return \{-1, -1\};
    int p, q;
    gcd(m1 / g, m2 / g, p, q);
    int mod = m1 / g * m2;
    int x = (a1 * (m2 / q) % mod * q % mod + a2 * (m1 / q) %
         mod * p % mod) % mod;
```

```
a1 = x;
   if (a1 < 0) a1 += mod;
   m1 = mod;
 return {a1, m1};
} // namespace crt
// o menor inteiro a que satisfaz:
// a mod p1 = x1
// a mod p2 = x2
// a funcao crt retorna um pair {a, mod}
// dai a solucao pode ser descrita como
// x = a % mod
// entao os valores possiveis sao:
// a, (a + mod), a + (2 * mod), a + (3 * mod), ...
// cuidado com overflow!
crttrick.cpp
                                                           19 lines
vector<pi> eq;
map<int, int> by_mod;
```

```
// quero checar se existe solucao para o sistema das equacoes
    que ja adicionei
// junto da equacao curr
// geralmente da pra fazer algo como if (check(curr)) { add(
    curr):
bool check(pi curr) {
 if (by_mod.find(curr.sec) != by_mod.end()) {
   return by_mod[curr.sec] == curr.fir;
 for (auto [x, mod] : eq) {
   if ((curr.fir - x) % __gcd(curr.sec, mod)) return 0;
 return 1;
void add(pi curr) { // [valor, mod]
 eq.pb(curr);
 by_mod[curr.sec] = curr.fir;
// quando tem algo de sgrt mods distintos ou algo do tipo
```

diophantine.cpp

```
namespace dio {
vector<pi> sols;
int gcd(int a, int b, int &x, int &y) {
 if (b == 0) {
   x = 1, y = 0;
   return a:
 int x1, y1, d = gcd(b, a % b, x1, y1);
 x = y1, y = x1 - y1 * (a / b);
 return d;
void one_sol(int a, int b, int c) {
 int x0, y0, g;
 q = gcd(abs(a), abs(b), x0, y0);
 if (c % g) return;
 x0 *= (c / g);
 y0 \star = (c / q);
 if (a < 0) \times 0 *= -1;
 if (b < 0) y0 \star = -1;
 sols.pb({x0, y0});
void more_sols(int a, int b, int c) {
 int g = \underline{gcd(a, b)};
 int x0 = sols[0].fir, y0 = sols[0].sec;
 for (int k = -200000; k \le 200000; k++) {
   int x = x0 + k * (b / q);
```

```
int y = y0 - k * (a / q);
   sols.pb({x, y});
} // namespace dio
// equacoes do tipo: ax + by = c
```

extendedeuclidean.cpp

const int MOD = 998244353;

```
14 lines
int gcd(int a, int b, int &x, int &v) {
 if (b == 0) {
   x = 1;
   y = 0;
    return a;
 int x1, y1;
 int d = gcd(b, a % b, x1, y1);
 x = v1;
 y = x1 - y1 * (a / b);
 return d:
// achar os numeros x e y tal que:
// a * x + b * y = gcd(a, b)
```

ntt.cpp

```
typedef mod int<MOD> mint;
void ntt(vector<mint>& a, bool rev) {
 int n = a.size():
 auto b = a;
 assert(!(n & (n - 1)));
 mint \alpha = 1:
 while ((g ^ (MOD / 2)) == 1) g += 1;
 if (rev) q = 1 / q;
 for (int step = n / 2; step; step /= 2) {
   mint w = g ^ (MOD / (n / step)), wn = 1;
    for (int i = 0; i < n / 2; i += step) {
     for (int j = 0; j < step; j++) {
       auto u = a[2 * i + j], v = wn * a[2 * i + j + step];
       b[i + j] = u + v;
       b[i + n / 2 + j] = u - v;
     wn = wn * w;
   swap(a, b);
 if (rev) {
   auto n1 = mint(1) / n;
   for (auto& x : a) x *= n1;
vector<mint> convolution(const vector<mint>& a, const vector<
 vector<mint> l(a.begin(), a.end()), r(b.begin(), b.end());
 int N = 1.size() + r.size() - 1, n = 1 << __lg(2 * N - 1);
 l.resize(n);
 r.resize(n);
 ntt(1, false);
 ntt(r, false);
 for (int i = 0; i < n; i++) l[i] *= r[i];
 ntt(1, true);
 l.resize(N);
 return 1;
```

fft.cpp

32 lines

#define PI acos(-1)

othermodint fraction totient pollardrho

```
#define cd complex<double>
namespace fft {
int n;
void fft(vector<cd> &a, bool invert) {
  int n = a.size();
  for (int i = 1, j = 0; i < n; i++) {
   int bit = n \gg 1;
    for (; j & bit; bit >>= 1) j ^= bit;
    j ^= bit;
   if (i < j) swap(a[i], a[j]);
  for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 : 1);
    cd wlen(cos(ang), sin(ang));
    for (int i = 0; i < n; i += len) {
     cd w(1);
      for (int j = 0; j < len / 2; j++) {
       cd u = a[i + j], v = a[i + j + len / 2] * w;
       a[i + j] = u + v;
       a[i + j + len / 2] = u - v;
       w \star = wlen;
    }
  if (invert)
    for (cd &x : a) x \neq n;
vector<int> mul(vector<int> a, vector<int> b) {
  vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());
  while (n < a.size() + b.size()) n <<= 1;
  fa.resize(n);
  fb.resize(n);
  fft(fa, false);
  fft(fb, false);
  for (int i = 0; i < n; i++) fa[i] *= fb[i];
  fft(fa, true);
  vector<int> ans(n);
  for (int i = 0; i < n; i++) ans[i] = round(fa[i].real());
  return ans:
} // namespace fft
```

othermodint.cpp

64 lines

```
template <int p>
struct mod int {
 ll expo(ll b, ll e) {
   ll ret = 1;
   while (e) {
     if (e % 2) ret = ret * b % p;
     e /= 2, b = b * b % p;
   return ret:
  11 inv(11 b) { return expo(b, p - 2); }
  using m = mod_int;
  int v:
  mod int() : v(0) {}
  mod_int(ll v_) {
   if (v_ >= p or v_ <= -p) v_ %= p;
   if (v_{-} < 0) v_{-} += p;
   v = v_;
  m& operator+=(const m& a) {
   v += a.v;
   if (v \ge p) v = p;
    return *this;
```

```
m& operator-=(const m& a) {
  v -= a.v;
  if (v < 0) v += p;
  return *this;
m& operator *= (const m& a) {
  v = v * 11(a.v) % p;
  return *this;
m& operator/=(const m& a) {
  v = v * inv(a.v) % p;
  return *this;
m operator-() { return m(-v); }
m& operator^=(ll e) {
  if (e < 0) {
    v = inv(v);
    e = -e;
  v = expo(v, e);
  // possivel otimizacao:
  // cuidado com 0^0
  // v = \exp((v, e^{(p-1)});
  return *this;
bool operator==(const m& a) { return v == a.v; }
bool operator!=(const m& a) { return v != a.v; }
friend istream& operator>>(istream& in, m& a) {
  ll val:
  in >> val:
  a = m(val);
  return in;
friend ostream& operator << (ostream& out, m a) { return out <<
friend m operator+(m a, m b) { return a += b; }
friend m operator-(m a, m b) { return a -= b; }
friend m operator*(m a, m b) { return a *= b; }
friend m operator/(m a, m b) { return a /= b; }
friend m operator (m a, ll e) { return a '= e; }
```

fraction.cpp

struct fraction { int x, y; // x / y fraction() {} fraction(int x, int y) : x(x), y(y) {} bool operator == (fraction o) { return (x * o.y == o.x * y); } bool operator!=(fraction o) { return (x * o.y != o.x * y); } bool operator>(fraction o) { return (x * o.y > o.x * y); } bool operator>=(fraction o) { return $(x * o.y >= o.x * y); }$ bool operator<(fraction o) { return (x * o.y < o.x * y); }</pre> bool operator<=(fraction o) { return (x * o.y <= o.x * y); }</pre> fraction operator+(fraction o) { fraction ans: ans.y = (y == o.y) ? y : y * o.y; ans. $\bar{x} = (\bar{x}) * (ans.y / y) + (o.x) * (ans.y / o.y);$ // ans.simplify(); return ans: fraction operator*(fraction o) { fraction ans: ans.x = x * o.x;ans.y = y * o.y;// ans.simplify(); return ans; fraction inv() {

```
fraction ans = fraction(x, y);
    swap(ans.x, ans.y);
    return ans;
 fraction neg() {
    fraction ans = fraction(x, y);
    ans.x \star = -1;
   return ans:
 void simplify() {
    if (abs(x) > 1e9 \mid | abs(y) > 1e9) // slow simplification
      int g = \underline{gcd}(y, x);
      x /= g;
      y /= g;
 // substraction and division can be easily done
};
```

totient.cpp

20 lines

```
int phi[MAXN];
void calc() {
 for (int i = 0; i < MAXN; i++) phi[i] = i;
 for (int i = 2; i < MAXN; i++) {
   if (phi[i] == i) {
     for (int j = i; j < MAXN; j += i) phi[j] -= phi[j] / i;
int calc_phi(int n) {
 int ans = n;
 for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) {
     while (n \% i == 0) n /= i;
      ans -= ans / i;
 if (n > 1) ans -= ans / n;
 return ans:
```

pollardrho.cpp

```
#define int __int128
namespace pollard rho {
int multiplicate(int x, int y, int m) { return (x * y) % m; }
int modpow(int x, int y, int m) {
  int z = 1;
  while (v) {
    if (y \& 1) z = (z * x) % m;
    x = (x * x) % m;
    y >>= 1;
  return z;
bool is_composite(int n, int a, int d, int s) {
  int x = modpow(a, d, n);
  if (x == 1 \text{ or } x == n - 1) return false;
  for (int r = 1; r < s; r++) {
    x = multiplicate(x, x, n);
    if (x == n - 1LL) return false;
 return true;
int miller rabin(int n) {
 if (n < 2) return false;
 int r = 0, d = n - 1LL;
```

```
while ((d \& 1LL) == 0) {
   d >>= 1;
  for (int a: {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}) {
   if (n == a) return true;
   if (is_composite(n, a, d, r)) return false;
  return true:
int f(int x, int m) { return multiplicate(x, x, m) + 1; }
int rho(int n) {
  int x0 = 1, t = 0, prd = 2;
  int x = 0, y = 0, q;
  while (t % 40 \mid | \_gcd(prd, n) == 1) {
   if (x == y) {
     x0++:
     x = x0;
     y = f(x, n);
   q = multiplicate(prd, max(x, y) - min(x, y), n);
   if (q != 0) prd = q;
   x = f(x, n);
   y = f(y, n);
   y = f(y, n);
   t++;
  return __gcd(prd, n);
vector<int> fact(int n) {
 if (n == 1) return {};
  if (miller_rabin(n)) return {n};
  int x = rho(n);
  auto l = fact(x), r = fact(n / x);
  l.insert(l.end(), r.begin(), r.end());
} // namespace pollard_rho
```

primefactors.cpp

```
vector<int> facts;
void primefactors(int n) {
  while (n % 2 == 0) {
    facts.pb(2);
    n = n / 2;
  }
  for (int i = 3; i <= sqrt(n); i += 2) {
    while (n % i == 0) {
      facts.pb(i);
      n = n / i;
    }
  }
  if (n > 2) facts.pb(n);
}
```

xorbasis.cpp

int modpow(int a, int b) {
 int res = 1;
 while (b > 0) {
 if (b & 1) res = (res * a) % mod;
 a = (a * a) % mod;
 b >>= 1;
 }
 return res;
}
int all, qt;
int dp[33];

```
void add(int x) {
 all++;
 for (int i = 32; i >= 0; i--) {
   if (x & (111 << i)) {
      if (dp[i] == 0) {
       dp[i] = x;
        at++;
       return:
     x ^= dp[i];
int get(int x) // gual o x-esimo menor valor de xor de uma
    subsequencia
 int tot = (111 \ll gt), ans = 0;
 for (int i = 32; i >= 0; i--) {
   if (dp[i] > 0) {
     int d = tot / 2;
      if (d < x && !(ans & (111 << i)))
       ans ^= dp[i];
      else if (d >= x && (ans & (111 << i)))
       ans ^= dp[i];
     if (d < x) x = d;
     tot /= 2;
 return ans;
bool check(int x) // se existe pelo menos uma subsequencia com
 for (int i = 32; i >= 0; i--) {
   if (x & (111 << i)) {
     if (!dp[i]) return 0;
      x ^= dp[i];
 return 1;
int count(int x) // quantas subsequencias tem xor x
 if (!check(x)) return 0;
 return modpow(2, all - qt);
int x = get(111 \ll qt); // maior xor possivel de uma
    subsequencia
int y = get(1);
                         // maior xor possivel != 0 (o 0 sempre
     eh possivel - subsequencia vazia)
```

mobius.cpp

57 lines

```
int lpf[MAXN];
int mobius[MAXN];
int mp[MAXN];
vectorint> d[MAXN];
void calc_lpf() {
  for (int i = 2; i < MAXN; i++) {
    if (!!pf[i]) {
      for (int j = i; j < MAXN; j += i) {
        if (!lpf[j]) lpf[j] = i;
      }
    }
}
void calc() {
  for (int i = 2; i < MAXN; i++) // divisores</pre>
```

```
for (int j = i; j < MAXN; j += i) d[j].pb(i);
 calc lpf();
 mobius[1] = 1;
 for (int i = 2; i < MAXN; i++) {
   if (lpf[i / lpf[i]] == lpf[i])
     mobius[i] = 0;
      mobius[i] = -1 * mobius[i / lpf[i]];
void add(int x, int dd) // adiciona dd em todos os val[i] que
    gcd(x, i) > 1
 for (auto const &i : d[x]) mp[i] += dd;
int sum(int x) // valor de val[x]
 int ans = 0:
 for (auto const &i : d[x]) ans += (mobius[i] * -1 * mp[i]);
 return ans;
// mobius/inclusao-exclusao com os fatores primos
// a funcao de mobius eh definida como:
// mi(n) = 1, se n eh um square-free com um numero par de
    fatores primos
// mi(n) = -1, se n eh um square-free com um numero impar de
    fatores primos
// mi(n) = 0, caso nenhum dos dois
// square-free = nenhum fator primo aparece duas vezes ou mais
```

segmentedsieve.cpp

40.11

```
vector<int> prime;
void segmentedsieve(int 1, int r) {
 int lim = sqrt(r);
 vector<bool> mark(lim + 1, false);
 vector<int> primes;
 for (int i = 2; i \le \lim_{t \to 0} ++i) {
   if (!mark[i]) {
     primes.pb(i);
      for (int j = i * i; j \le lim; j += i) mark[j] = true;
 vector<bool> isprime(r - 1 + 1, true);
 for (int i : primes)
   for (int j = \max(i * i, (l + i - 1) / i * i); j <= r; j +=
        i) isprime[j - 1] = false;
 if (1 == 1) isprime[0] = false;
 for (int i = 0; i < isprime.size(); i++)
   if (isprime[i]) prime.pb(i + 1);
```

xortrie.cpp

```
struct node {
  int me, cnt, id;
  int down[2];
  node(int c = 0) : me(c) {
    cnt = 0;
    id = -1;
    fill(begin(down), end(down), -1);
  }
};
struct trie_xor {
  vector<node> t;
  trie_xor() { t.resize(1); }
  void add(int n, int id) {
    int v = 0;
```

```
t[v].cnt++;
    for (int i = 30; i >= 0; i--) {
     int bit = (n & (1 << i)) ? 1 : 0;
     if (t[v].down[bit] == -1) {
       t[v].down[bit] = t.size();
       t.emplace_back(bit);
     v = t[v].down[bit];
     t[v].cnt++;
   t[v].id = id;
  void rem(int n, int id) {
   int v = 0;
   t[v].cnt--;
    for (int i = 30; i >= 0; i--) {
     int bit = (n & (1 << i)) ? 1 : 0;
     v = t[v].down[bit];
     t[v].cnt--;
  int qry(int n) // maximum xor with n
    if (t[0].cnt == 0) // no element
     return -1;
    int v = 0;
    for (int i = 30; i >= 0; i--) {
     int bit = (n & (1 << i)) ? 0 : 1;
     int nxt = t[v].down[bit];
     if (nxt != -1 \&\& t[nxt].cnt > 0)
       v = nxt;
     else
       v = t[v].down[bit ^ 1];
    return t[v].id;
};
```

matrix.cpp

```
namespace matrix {
vector<vector<int>> ans;
int multi(int x, int y) { return (x * y) % mod; }
int sum(int a, int b) { return (a + b \ge mod) ? a + b - mod : a
vector<vector<int>> multiply(vector<vector<int>> a, vector<</pre>
    vector<int>> b) {
  vector<vector<int>> res(a[0].size(), vector<int>(b[0].size())
  for (int i = 0; i < a.size(); i++) {
    for (int j = 0; j < b[0].size(); j++) {
      res[i][j] = 0;
      for (int k = 0; k < a[0].size(); k++) res[i][j] = sum(res
          [i][j], multi(a[i][k], b[k][j]));
  return res;
vector<vector<int>> expo(vector<vector<int>> mat, int m) {
  ans = vector<vector<int>>(mat.size(), vector<int>(mat[0].size
       ()));
  for (int i = 0; i < mat.size(); i++)
    for (int j = 0; j < mat[0].size(); j++) ans[i][j] = (i == j
        );
  while (m > 0) {
   if (m & 1) ans = multiply(ans, mat);
   m = m / 2;
   mat = multiply(mat, mat);
```

```
matrix fwht teoremas segtree segtreelazy
 return ans;
} // namespace matrix
fwht.cpp
vector<modint> fwht(char op, vector<modint> f, bool inv = 0) {
 int n = f.size();
 for (int k = 0; (n - 1) >> k; k++) {
   for (int i = 0; i < n; i++) {
     if (i >> k & 1) {
       int j = i ^ (1 << k);
       if (op == '^') f[j] += f[i], f[i] = f[j] - modint(2) *
       if (op == '|') f[i] += modint(inv ? -1 : 1) * f[j];
       if (op == '&') f[j] += modint(inv ? -1 : 1) * f[i];
 if (op == '^' and inv) {
   for (auto &i : f) i /= n;
 return f;
vector<modint> conv(char op, vector<modint> a, vector<modint> b
 a = fwht(op, a, 0);
 b = fwht(op, b, 0);
 for (int i = 0; i < a.size(); i++) {
   a[i] *= b[i];
 return fwht(op, a, 1);
teoremas.md
                                                         39 lines
```

```
# Teorema de Lucas
sejam m e n numeros inteiros nao negativos e p um numero primo
desenvolver n e m na base p
ou seja:
m = m[k] *p^{(k)} + m[k - 1] *p^{(k - 1)} + ... + m[0] *p^{(0)}
n = n[k] *p^{(k)} + n[k - 1] *p^{(k - 1)} + ... + n[0] *p^{(0)}
ncr(m, n) \mod p = produtorio de (ncr(m[i], n[i]) \mod p)
dai pra generalizar pro mod 2 eh deboas, pq se tiver um bit
     setado em n[i] que nao ta setado em m[i], entao miou, vai
     dar zero
# Manhattan and Chebyshev distances equivalences
It is well known that given points (x, y) and you need to
     calculate the Manhattan distances between them, instead of
|x1-x2|+|v1-v2|
you can first convert all points (x, y) into (x+y, x-y) (rotate
      45 degrees) and the distances will become max(|x1-x2|, |
     y1-y2|) (also known as Chebyshev distance).
# Chicken McNugget Theorem
For any two coprime numbers (n > 0, m > 0), the greatest
     integer that cannot be written in the form:
an + bm, (a >= 0, b >= 0)
is (n \times m) - n - m
## Consequence of the theorem
That there are exactly ((n-1) \times (m-1)) / 2 positive
     integers which cannot be expressed in the form an + bm, (a
     >= 0, b >= 0)
## Generalization
If n and m are not coprime, so all numbers that are not
     multiples of gcd(n, m) cannot be expressed in the form an
     + bm, (a >= 0, b >= 0)
```

```
in addition, you can consider n = (n / gcd(n, m)) and m = (m / gcd(n, m))
     gcd(n, m)), to find how many multiples of gcd(n, m) cannot
     be expressed, or to find the greatest multiple of gcd(n,
    m) that cannot be expressed
## Considering a > 0, b > 0
Considering (n > 0, m > 0), n and m are coprime:
let v = ((n \times m) + min(n, m)) - 1
The number of positive integers which cannot be expressed
    increases by (y / n)
The number of positive integers which cannot be expressed
    increases by (y / m)
you must not count the multiples of (n \times m) more than once,
     just decrease number of positive integers which cannot be
     expressed by (v / (n \times m))
# Binomial Theorem
### Theorem
SS
(x + y)^n = \sum_{k=0}^{n} \{n \in k\} x^{n-k}y^{k}
$$
in addition, we have:
SS
(x - y)^n = \sum_{k = 0}^{n} (-1)^k \{n \in k\} x^{n - k}y^{k}
$$
(1 + x)^n = \sum_{k=0}^{n} \{n \in k\} x^{k}
```

Data structures (3)

segtree.cpp

25 lines

```
struct segtree {
 int n;
 vector<int> seq;
 int neutral() { return 0; }
 int merge(int a, int b) { return a + b; }
  void build(vector<int> &v) {
    n = 1:
    while (n < v.size()) n <<= 1;
    seq.assign(n << 1, neutral());</pre>
    for (int i = 0; i < v.size(); i++) seg[i + n] = v[i];
    for (int i = n - 1; i; i--) seq[i] = merge(seq[i << 1], seq
         [(i << 1) | 11);</pre>
 void upd(int i, int value) {
    seg[i += n] += value;
    for (i >>= 1; i; i >>= 1) seg[i] = merge(seg[i << 1], seg[(
        i << 1) | 11);
  int gry(int 1, int r) {
   int ansl = neutral(), ansr = neutral();
    for (1 += n, r += n + 1; 1 < r; 1 >>= 1, r >>= 1) {
     if (l \& 1) ansl = merge(ansl, seg[l++]);
     if (r \& 1) ansr = merge(seg[--r], ansr);
    return merge(ansl, ansr);
};
```

segtreelazy.cpp

```
struct segtree {
 int n;
 vector<int> v;
 vector<int> seg;
 vector<int> lazv;
  segtree(int sz) {
    n = sz;
```

```
seq.assign(4 * n, 0);
  lazy.assign(4 * n, 0);
int single(int x) { return x; }
int neutral() { return 0; }
int merge(int a, int b) { return a + b; }
void add(int i, int l, int r, int diff) {
  seg[i] += (r - 1 + 1) * diff;
 if (1 != r) {
   lazy[i << 1] += diff;
   lazy[(i << 1) | 1] += diff;
 lazy[i] = 0;
void update(int i, int l, int r, int ql, int qr, int diff) {
 if (lazy[i]) add(i, l, r, lazy[i]);
 if (1 > r || 1 > qr || r < ql) return;
 if (1 >= q1 \&\& r <= qr) {
   add(i, l, r, diff);
   return;
  int mid = (1 + r) >> 1;
 update(i << 1, 1, mid, ql, qr, diff);
 update((i << 1) | 1, mid + 1, r, ql, qr, diff);
  seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
int query(int 1, int r, int q1, int qr, int i) {
 if (lazv[i]) add(i, l, r, lazv[i]);
  if (l > r \mid\mid l > qr \mid\mid r < ql) return neutral();
 if (1 >= ql \&\& r <= qr) return seg[i];
 int mid = (1 + r) >> 1;
  return merge(query(1, mid, ql, qr, i << 1), query(mid + 1,
      r, ql, qr, (i << 1) | 1));
void build(int 1, int r, int i) {
 if (1 == r) {
   seq[i] = single(v[l]);
   return:
  int mid = (1 + r) >> 1;
 build(1, mid, i << 1);
 build(mid + 1, r, (i << 1) | 1);
 seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
int qry(int 1, int r) \{ return query(0, n - 1, 1, r, 1); \}
void upd(int 1, int r, int x) { update(1, 0, n - 1, 1, r, x);
```

fenwick.cpp

33 lines

```
struct fenw {
  int n;
  vector<int> bit;
  fenw() {}
  fenw(int sz) {
    n = sz;
    bit.assign(sz + 1, 0);
  }
  int qry(int r) // query de prefixo a[0] + a[1] + ... a[r]
  {
    int ret = 0;
    for (int i = r + 1; i > 0; i -= i & -i) ret += bit[i];
    return ret;
  }
  void upd(int r, int x) // a[r] += x
  {
    for (int i = r + 1; i <= n; i += i & -i) bit[i] += x;
  }
}</pre>
```

```
int bs(int x) // retorna o maior indice i (i < n) tal que:
      qry(i) < x
    int i = 0, k = 0;
    while (1 << (k + 1) <= n) k++;
    while (k \ge 0) {
     int nxt_i = i + (1 << k);
      if (nxt_i <= n && bit[nxt_i] < x) {
       i = nxt_i;
       x -= bit[i]:
      k--;
    return i - 1;
treap.cpp
                                                           93 lines
struct treap {
 int data, priority;
 int sz, lazy2;
 bool lazy;
 treap *1, *r, *parent;
int size(treap *node) { return (!node) ? 0 : node->sz; }
void recalc(treap *node) {
 if (!node) return;
 node -> sz = 1;
  node -> parent = 0:
  if (node->1) node->sz += node->1->sz, node->1->parent = node;
  if (node->r) node->sz += node->r->sz, node->r->parent = node;
void lazy propagation(treap *node) {
 if (node == NULL) return;
 if (node->lazv2) {
   if (node->1) node->1->lazy2 += node->lazy2;
    if (node->r) node->r->lazy2 += node->lazy2;
    node->data += node->lazy2;
   node \rightarrow lazv2 = 0;
 if (node->lazy) {
    swap(node->1, node->r);
    if (node->1) node->1->lazv = !node->1->lazv;
    if (node->r) node->r->lazy = !node->r->lazy;
    node \rightarrow lazy = 0;
void split(treap *t, treap *&l, treap *&r, int n) {
 if (!t) return void(1 = r = 0);
 lazy_propagation(t);
 if (size(t->1) >= n)
    split(t->1, 1, t->1, n), r = t;
    split(t->r, t->r, r, n - size(t->1) - 1), 1 = t;
  recalc(t);
void merge(treap *&t, treap *1, treap *r) {
  lazy propagation(1);
  lazy_propagation(r);
 if (!1)
   t = r;
 else if (!r)
   t = 1;
  else if (l->priority > r->priority)
    merge(1->r, 1->r, r), t=1;
    merge(r->1, 1, r->1), t = r;
  recalc(t);
```

```
void troca(treap *&t, int 1, int r, int 11, int rr) // sap de
    ranges
 treap *a0, *a1, *b0, *b1, *c0, *c1, *d0, *d1;
  split(t, a0, a1, 1);
  split(a1, b0, b1, r - 1 + 1);
 11 -= (r + 1);
 rr -= (r + 1);
  split(b1, c0, c1, l1);
  split(c1, d0, d1, rr - l1 + 1);
 merge(t, a0, d0);
 merge(t, t, c0);
 merge(t, t, b0);
 merge(t, t, d1);
void add(treap *&t, int 1, int r) {
 treap *a0, *a1, *b0, *b1;
 split(t, a0, a1, 1);
  split(a1, b0, b1, r - 1 + 1);
 b0 -> lazy ^= 1;
 b0 -> lazy2 += 1;
 merge(t, a0, b0);
 merge(t, t, b1);
void dfs(treap *t) {
 if (!t) return;
 lazy propagation(t);
 dfs(t->1):
  solve(t->data);
  dfs(t->r);
treap *create_node(int data, int priority) {
 treap *ret = new treap;
  ret->data = data;
 ret->priority = priority;
 ret->1 = 0;
  ret->r = 0:
  ret->sz = 1;
  ret->lazy = 0;
  ret->lazy2 = 0;
  ret->parent = 0;
  return ret;
```

colorupdate.cpp

```
const int inf = 1e15;
struct color upd {
#define left fir
#define right sec.fir
#define color sec.sec
 set<pii> ranges;
 vector<pii> erased;
  color_upd(int n) // inicialmente, todo mundo pintado com a
       cor inf
    // nao usar cores negativas!!!!!!!!
    ranges.insert(\{0, \{n-1, inf\}\}\);
 int get(int i) {
    auto it = ranges.upper_bound({i, {1e18, 1e18}});
    if (it == ranges.begin()) return -1;
    it--;
    return (*(it)).color;
 void del(int 1, int r) // apaga o intervalo [1, r]
    erased.clear();
```

```
auto it = ranges.upper_bound(\{1, \{0, 0\}\});
    if (it != ranges.begin()) {
     it--;
    while (it != ranges.end()) {
     if ((*(it)).left > r)
       break;
      else if ((*(it)).right >= 1)
       erased.push_back(*it);
      it++;
    if (erased.size() > 0) {
      int sz = erased.size();
      auto it = ranges.lower_bound({erased[0].left, {0, 0}});
      auto it2 = ranges.lower_bound({erased[sz - 1].left, {0, 0}
      pii ini = *it, fim = *it2;
      it2++;
      ranges.erase(it, it2);
      pii upd1 = \{ini.left, \{l-1, ini.color\}\};
      pii upd2 = \{r + 1, \{fim.right, fim.color\}\};
      erased[0].left = max(erased[0].left, 1);
      erased[sz - 1].right = min(erased[sz - 1].right, r);
      if (upd1.left <= upd1.right) ranges.insert(upd1);</pre>
      if (upd2.left <= upd2.right) ranges.insert(upd2);
  void add(int a, int b, int c) {
    auto it = ranges.lower_bound({a, {b, 0}});
    pii aa = \{-1, \{-1, -1\}\};
   pii bb = \{-1, \{-1, -1\}\};
    if (it != ranges.end()) {
     if ((*it).color == c && (*it).left == b + 1) {
       aa = *it;
        b = (*it).right;
    if (it != ranges.begin()) {
     if ((*it).color == c && (*it).right == a - 1) {
       bb = *it;
        a = (*it).left;
    ranges.erase(aa);
    ranges.erase(bb);
    ranges.insert({a, {b, c}});
  void upd(int a, int b, int c) // pinta o intervalo [a, b]
      com a cor c
    del(a, b);
    add(a, b, c);
};
```

bit2d.cpp

51 lines

```
struct bit2d {
  vector<int> ord;
  vector<vector<int>> t;
  vector<vector<int>> t;
  vector<vector<int>> t;
  vector<vector<int>> pass  // recebe todos os pontos que vao ser
        inseridos pra construir, mas nao insere eles
  {
    sort(pts.begin(), pts.end());
    for (auto const &a : pts) {
        if (ord.empty() || a.fir != ord.back()) ord.pb(a.fir);
    }
}
```

```
t.resize(ord.size() + 1);
   coord.resize(t.size());
   for (auto &a : pts) {
     swap(a.fir, a.sec);
   sort(pts.begin(), pts.end());
   for (auto &a : pts) {
     swap(a.fir, a.sec);
     for (int on = upper_bound(ord.begin(), ord.end(), a.fir)
           - ord.begin(); on < t.size(); on += on & -on) {
       if (coord[on].empty() || coord[on].back() != a.sec)
            coord[on].push_back(a.sec);
    for (int i = 0; i < t.size(); i++) t[i].assign(coord[i].
        size() + 1, 0);
 void add(int x, int y, int v) // v[a][b] += v
   for (int xx = upper_bound(ord.begin(), ord.end(), x) - ord.
        begin(); xx < t.size(); xx += xx & -xx) {
     for (int yy = upper_bound(coord[xx].begin(), coord[xx].
          end(), y) - coord[xx].begin(); yy < t[xx].size();
          yy += yy \& -yy)
       t[xx][yy] += v;
 int qry(int x, int y) // soma de todos os v[a][b] com (a <=
      x && b <= y)
   int ans = 0;
    for (int xx = upper_bound(ord.begin(), ord.end(), x) - ord.
        begin(); xx > 0; xx -= xx & -xx) {
     for (int yy = upper_bound(coord[xx].begin(), coord[xx].
          end(), y) - coord[xx].begin(); yy > 0; yy -= yy & -
       ans += t[xx][yy];
   return ans;
 int qry2(int x1, int y1, int x2, int y2) {
   return qry(x2, y2) - qry(x2, y1 - 1) - qry(x1 - 1, y2) +
        qry(x1 - 1, y1 - 1);
 void add2(int x1, int y1, int x2, int y2, int v) {
   add(x1, y1, v);
   add(x1, y2 + 1, -v);
   add(x2 + 1, y1, -v);
   add(x2 + 1, y2 + 1, y);
};
```

mo.cpp

// }

namespace mo {

return x.r > y.r;

bool cmp(query &x, query &y) {

struct query {
 int idx, 1, r;
};
int block;
vector<query> queries;
vector<int> ans;
// bool cmp(query &x, query &y) { essa funcao de ordenacao pode
 funcionar em caso de TLE
// int ablock = x.l / MAGIC, bblock = y.l / MAGIC;
// if (ablock != bblock) return ablock < bblock;
// if (ablock & 1) return x.r < y.r;</pre>

```
if (x.1 / block != y.1 / block) return x.1 / block < y.1 /
      block:
 return x.r < v.r;
void run() {
 block = (int)sgrt(n);
 sort(queries.begin(), queries.end(), cmp);
  ans.resize(queries.size());
 int c1 = 0, cr = -1, sum = 0;
  auto add = [\&] (int x) { sum += x; };
  auto rem = [\&] (int x) { sum -= x; };
  for (int i = 0; i < queries.size(); i++) {
    while (cl > queries[i].1) {
      cl--;
      add(v[cl]);
    while (cr < queries[i].r) {
      cr++:
      add(v[cr]);
    while (cl < queries[i].1) {
      rem(v[cl]);
      cl++;
    while (cr > queries[i].r) {
      rem(v[cr]);
    ans[queries[i].idx] = sum;
} // namespace mo
```

segtree2d.cpp

if (x >= n)

seg[x][y] = val;

45 lines

```
struct segtree2d {
 int n, m;
  vector<vector<int>> sea:
  int neutral() { return 0; }
 int merge(int a, int b) { return a + b; }
  segtree2d(int nn, int mm) {
    n = nn, m = mm;
    seg = vector<vector<int>>(2 * n, vector<int>(2 * m, neutral
  int gry(int x1, int y1, int x2, int y2) {
    int ret = neutral();
    int y3 = y1 + m, y4 = y2 + m;
    for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2) {
      for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /= 2) {
         if (x1 \% 2 == 1 \text{ and } y1 \% 2 == 1) \text{ ret } = \text{merge}(\text{ret, seq}[
              x1][y1]);
         if (x1 \% 2 == 1 \text{ and } y2 \% 2 == 0) \text{ ret } = \text{merge}(\text{ret, seg}[
              x1][y2]);
         if (x2 \% 2 == 0 \text{ and } y1 \% 2 == 1) \text{ ret} = merge(ret, seg[}
         if (x2 \% 2 == 0 \text{ and } y2 \% 2 == 0) \text{ ret } = \text{merge}(\text{ret, seg}[
              x2][y2]);
    return ret;
 void upd(int x, int y, int val) {
    int v2 = v += m;
    for (x += n; x; x /= 2, y = y2) {
```

seg[x][y] = merge(seg[2 * x][y], seg[2 * x + 1][y]);

```
while (y \neq 2) seg[x][y] = merge(seg[x][2 * y], seg[x][2
          * y + 1]);
};
persistentseg.cpp
                                                          37 lines
struct node {
  int item, 1, r;
 node (int 1, int r, int item) : l(1), r(r), item(item) {}
int n, q;
vector<node> seg;
vector<int> roots;
void init() { seg.resize(1); }
int newleaf(int vv) {
  int p = seq.size();
  seg.pb(node(0, 0, vv));
  return p;
int newpar(int 1, int r) {
  int p = seq.size();
  seg.pb(node(l, r, seg[l].item + seg[r].item));
  return p;
int upd(int i, int l, int r, int pos) {
  if (1 == r) return newleaf(seg[i].item + 1);
  int mid = (1 + r) >> 1;
  if (pos <= mid) return newpar(upd(seg[i].1, 1, mid, pos), seg
  return newpar(seg[i].1, upd(seg[i].r, mid + 1, r, pos));
int build(int 1, int r) {
 if (1 == r) return newleaf(0);
  int mid = (1 + r) >> 1;
  return newpar(build(1, mid), build(mid + 1, r));
int qry(int vl, int vr, int l, int r, int k) {
 if (l == r) return l;
  int mid = (1 + r) >> 1;
  int c = seg[seg[vr].1].item - seg[seg[vl].1].item;
 if (c \ge k) return qry(seg[v1].1, seg[vr].1, 1, mid, k);
 return qry(seg[v1].r, seg[vr].r, mid + 1, r, k - c);
rmq.cpp
  int n;
```

```
struct rmq {
  vector<vector<pi>> m;
  vector<int> log;
  rmq() {}
  rmq(vector<pi> &v) {
   n = v.size();
   log.resize(n + 1);
    log[1] = 0;
    for (int i = 2; i \le n; i++) log[i] = log[i / 2] + 1;
    int sz = log[n] + 2;
    m = vector<vector<pi>>>(sz, vector<pi>(n + 1));
    for (int i = 0; i < n; i++) {
     m[0][i] = v[i];
    for (int j = 1; j < sz; j++) {
      for (int i = 0; i + (1 << j) <= n; i++) m[j][i] = min(m[j])
            -1[i], m[j -1][i + (1 << (j -1))]);
```

```
int qry(int a, int b) { return min(m[log[b - a + 1]][a], m[
      log[b - a + 1]][b - (1 << log[b - a + 1]) + 1]).second;
};
binarylifting.cpp
```

```
item st[MAXN][21];
for (int i = 0; i < n; i++) {
 st[i][0].nxt = min(i + 1, n - 1);
 st[i][0].sum = v[st[i][0].nxt];
for (int i = 1; i < 21; i++) {
 for (int v = 0; v < n; v++) {
   st[v][i].nxt = st[st[v][i - 1].nxt][i - 1].nxt;
    st[v][i].sum = st[v][i - 1].sum + st[st[v][i - 1].nxt][i -
while (q--) {
 int l, r;
 cin >> 1 >> r;
 int ans = v[1], len = r - 1;
 for (int i = 20; i >= 0; i--) {
   if (len & (1 << i)) {
     ans += st[1][i].sum;
     l = st[l][i].nxt;
 cout << ans << endl;
```

minqueue.cpp

namespace min_queue {

```
deque<pi> q;
int l, r;
void init() {
 1 = r = 1;
 q.clear();
void push (int v) {
 while (!q.empty() && v < q.back().fir) q.pop_back();</pre>
 q.pb({v, r});
void pop() {
 if (!q.empty() && q.front().sec == 1) q.pop_front();
int getmin() { return q.front().fir; }
} // namespace min_queue
```

Graph (4)

floydwarshall.cpp

```
int dist[MAXN][MAXN];
void floyd_warshall() {
 for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
      for (int j = 0; j < n; j++) {
        dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j]);
```

```
void initialize() {
 for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++) {
     if (i == j) {
       dist[i][j] = 0;
     } else {
       dist[i][j] = INF;
```

centroiddecomp.cpp

```
int n, k, resp;
vector<int> adj[MAXN];
vector<int> cnt;
namespace cd {
int sz;
vector<int> subtree_size;
vector<bool> visited;
void dfs(int s, int f) {
 subtree\_size[s] = 1;
 for (auto const &v : adj[s]) {
   if (v != f && !visited[v]) {
      dfs(v, s);
      subtree_size[s] += subtree_size[v];
int get_centroid(int s, int f) {
 bool is_centroid = true;
 int heaviest_child = -1;
  for (auto const &v : adj[s]) {
    if (v != f && !visited[v]) {
      if (subtree_size[v] > sz / 2) is_centroid = false;
      if (heaviest_child == -1 || subtree_size[v] >
          subtree_size[heaviest_child]) heaviest_child = v;
 return (is_centroid && sz - subtree_size[s] <= sz / 2) ? s :
      get_centroid(heaviest_child, s);
void dfs2(int s, int f, int d) {
 while (d >= cnt.size()) cnt.pb(0);
 cnt[d]++;
  for (auto const &v : adj[s]) {
    if (v != f \&\& !visited[v]) dfs2(v, s, d + 1);
void solve(int s) {
 vector<int> tot;
  for (auto const &v : adj[s]) {
    if (visited[v]) continue;
    cnt.clear();
    dfs2(v, s, 1);
    for (int i = 1; i < cnt.size(); i++) {
      if (k - i < tot.size() \&\& k - i >= 1) resp += (cnt[i] *
          tot[k - i]);
    for (int i = 1; i < cnt.size(); i++) {
      while (i \ge tot.size()) tot.pb(0);
      tot[i] += cnt[i];
 if (k < tot.size()) resp += tot[k];
int decompose_tree(int s) {
```

sz = 0;

dfs(s, s);

dsu dsubipartido cycledetection blockcuttree bridges

```
int cend_tree = get_centroid(s, s);
  visited[cend_tree] = true;
  solve(cend_tree);
  for (auto const &v : adj[cend_tree]) {
   if (!visited[v]) decompose_tree(v);
  return cend_tree;
void init() {
  subtree_size.resize(n);
  visited.resize(n);
  decompose tree(0);
} // namespace cd
dsu.cpp
                                                           24 lines
struct dsu {
  int tot;
  vector<int> parent;
  vector<int> sz;
  dsu(int n) {
   parent.resize(n);
    sz.resize(n);
    tot = n;
    for (int i = 0; i < n; i++) {
     parent[i] = i;
     sz[i] = 1;
  int find_set(int i) { return parent[i] = (parent[i] == i) ? i
       : find set(parent[i]); }
  void make_set(int x, int y) {
   x = find set(x), y = find set(y);
   if (x != y) {
     if (sz[x] > sz[y]) swap(x, y);
     parent[x] = y;
     sz[y] += sz[x];
     tot--:
};
dsubipartido.cpp
                                                           41 lines
struct dsu {
  vector<pi> parent;
  vector<int> rank;
```

```
vector<int> bipartite;
dsu(int n) {
 parent.resize(n);
 rank.resize(n);
 bipartite.resize(n);
  for (int v = 0; v < n; v++) {
   parent[v] = \{v, 0\};
   rank[v] = 0;
   bipartite[v] = 1;
dsu() {}
pi find_set(int v) {
 if (v != parent[v].fir) {
   int parity = parent[v].sec;
   parent[v] = find_set(parent[v].fir);
   parent[v].sec ^= parity;
  return parent[v];
```

```
void add_edge(int a, int b) {
   pi pa = find set(a);
   a = pa.fir;
   int x = pa.sec;
   pi pb = find_set(b);
   b = pb.fir;
   int y = pb.sec;
   if (a == b) {
     if (x == y) bipartite[a] = 0;
     if (rank[a] < rank[b]) swap(a, b);</pre>
     parent[b] = \{a, x^y^1\};
     bipartite[a] &= bipartite[b];
     if (rank[a] == rank[b]) rank[a]++;
 bool is_bipartite(int v) { return bipartite[find_set(v).fir];
cycledetection.cpp
int n, m, idx;
vector<int> cycles[MAXN];
```

```
vector<int> adj[MAXN];
int color[MAXN];
int parent[MAXN];
int ans[MAXN];
void dfs(int u, int p) { // chama dfs a partir de qm tem cor 0
 if (color[u] == 2) return;
 if (color[u] == 1) {
    idx++;
    int curr = p;
    ans[curr] = idx;
    cvcles[idx].pb(curr);
    while (curr != u) {
     curr = parent[curr];
     cycles[idx].pb(curr);
     ans[curr] = idx;
   return;
 parent[u] = p;
 color[u] = 1;
 for (auto const &v : adj[u])
   if (v != parent[u]) dfs(v, u);
```

blockcuttree.cpp

struct block cut tree {

int n = q.size();

build();

color[u] = 2;

```
// Se art[i] >= 1, i eh ponto de articulação
// tree - eh a propia block-cut tree
// pos[i] responde a qual vertice da arvore vertice i
     pertence
vector<vector<int>> g, blocks, tree;
vector<vector<pi>>> edgblocks; // sao as arestas do bloco i
stack<int> s:
stack<pi> s2;
vector<int> id, art, pos;
```

block_cut_tree(vector<vector<int>> g_) : g(g_) {

int dfs(int i, int &t, int p = -1) {

id.resize(n, -1), art.resize(n), pos.resize(n);

```
if (j != p) {
       if (id[j] == -1) {
         int val = dfs(j, t, i);
         lo = min(lo, val);
         if (val >= id[i]) {
            art[i]++;
            blocks.emplace_back(1, i);
            while (blocks.back().back() != j) {
             blocks.back().pb(s.top());
            edgblocks.emplace_back(1, s2.top());
            s2.pop();
            pi aux = {j, i};
            while (edgblocks.back().back() != aux) {
              edgblocks.back().pb(s2.top());
              s2.pop();
          // if (val > id[i]) aresta i-j eh ponte
          lo = min(lo, id[j]);
    if (p == -1 \text{ and } art[i]) {
     art[i]--;
   return lo;
 void build() {
    for (int i = 0; i < q.size(); i++) {
     if (id[i] == -1) dfs(i, t, -1);
   tree.resize(blocks.size());
    for (int i = 0; i < q.size(); i++) {
     if (art[i]) pos[i] = tree.size(), tree.emplace_back();
    for (int i = 0; i < blocks.size(); i++) {
      for (int j : blocks[i]) {
       if (!art[j])
         pos[j] = i;
         tree[i].pb(pos[j]), tree[pos[j]].pb(i);
};
```

if (j != p and id[j] != -1) s2.emplace(i, j);

int lo = id[i] = t++;

s2.emplace(i, p);

for (int j : g[i]) {

for (int j : g[i]) {

s.push(i);

if (p != -1) {

bridges.cpp

```
nt n, m, timer;
vector<pi> edges;
vector<bool> is bridge;
vector<pi> adj[MAXN];
int tin[MAXN];
int low[MAXN];
                          // memset -1
                          // memset -1
bool vis[MAXN];
void dfs(int v, int p) { // chama de quem nao foi vis ainda
```

```
vis[v] = true;
tin[v] = timer, low[v] = timer++;
for (auto const &u : adj[v]) {
 if (u.fir == p) continue;
 if (vis[u.fir]) {
   low[v] = min(low[v], tin[u.fir]);
   continue;
 dfs(u.fir, v);
 low[v] = min(low[v], low[u.fir]);
 if (low[u.fir] > tin[v]) is_bridge[u.sec] = 1;
```

```
dinic.cpp
#define INF 1e9
struct edge {
 int to, from, flow, capacity, id;
struct dinic {
 int n, src, sink;
  vector<vector<edge>> adj;
  vector<int> level;
  vector<int> ptr;
  dinic(int sz) {
   n = sz;
    adj.resize(n);
   level.resize(n);
   ptr.resize(n);
  void add_edge(int a, int b, int c, int id) {
    adj[a].pb({b, (int)adj[b].size(), c, c, id});
   adj[b].pb({a, (int)adj[a].size() - 1, 0, 0, id});
 bool bfs() {
    level.assign(n, -1);
   level[src] = 0;
   queue<int> q;
   q.push(src);
    while (!q.empty()) {
     int u = q.front();
     for (auto at : adj[u]) {
       if (at.flow && level[at.to] == -1) {
          q.push(at.to);
          level[at.to] = level[u] + 1;
    return level[sink] != -1;
  int dfs(int u, int flow) {
    if (u == sink || flow == 0) return flow;
    for (int p = ptr[u]; p < adj[u].size(); p++) {
      edge &at = adj[u][p];
     if (at.flow && level[u] == level[at.to] - 1) {
       int kappa = dfs(at.to, min(flow, at.flow));
       at.flow -= kappa;
       adj[at.to][at.from].flow += kappa;
       if (kappa != 0) return kappa;
    return 0;
  int run() {
    int max_flow = 0;
    while (bfs()) {
```

```
ptr.assign(n, 0);
     while (1) {
       int flow = dfs(src, INF);
       if (flow == 0) break;
       max_flow += flow;
   return max flow:
 vector<pii> cut_edges() // arestas do corte minimo
   bfs();
   vector<pii> ans;
   for (int i = 0; i < n; i++) {
     for (auto const &j : adj[i]) {
        if (level[i] != -1 && level[j.to] == -1 && j.capacity >
             0) ans.pb({j.capacity, {i, j.to}});
   return ans;
 vector<int> flow_edges(int n, int m) // fluxo em cada aresta
      , na ordem da entrada
   vector<int> ans(m);
   for (int i = 0; i < n; i++) {
     for (auto const &j : adj[i])
       if (!j.capacity) ans[j.id] = j.flow;
   return ans;
};
```

hopcroftkarp.cpp

```
137 lines
#define INF 1e9
struct hopcroft karp {
 vector<int> match:
 vector<int> dist;
 vector<vector<int>> adj;
 int n, m, t;
 hopcroft_karp(int a, int b) {
   n = a, m = b;
   t = n + m + 1;
   match.assign(t, n + m);
   dist.assign(t, 0);
   adj.assign(t, vector<int>{});
 void add_edge(int u, int v) {
   adj[u].pb(v);
   adj[v].pb(u);
 bool bfs() {
   queue<int> q;
    for (int u = 0; u < n; u++) {
     if (match[u] == n + m)
       dist[u] = 0, q.push(u);
     else
        dist[u] = INF;
   dist[n + m] = INF;
   while (!q.empty()) {
     int u = q.front();
     q.pop();
     if (dist[u] < dist[n + m]) {
       for (auto const &v : adj[u]) {
         if (dist[match[v]] == INF) {
            dist[match[v]] = dist[u] + 1;
           q.push(match[v]);
```

```
return dist[n + m] < INF;
bool dfs(int u) {
  if (u < n + m) {
    for (auto const &v : adj[u]) {
      if (dist[match[v]] == dist[u] + 1 && dfs(match[v])) {
        match[v] = u;
        match[u] = v;
        return true;
    dist[u] = INF;
    return false;
  return true;
vector<pi> run() {
  int cnt = 0;
  while (bfs())
    for (int u = 0; u < n; u++)
     if (match[u] == n + m && dfs(u)) cnt++;
  vector<pi> ans;
  for (int v = n; v < n + m; v++)
    if (match[v] < n + m) ans.pb(\{match[v], v\});
vector<int> mvc() // minimum vertex cover
  vector<pi> ans = run();
  vector<bool> vis(n + m, 0);
  for (int i = 0; i < n; i++) {
    if (match[i] == n + m) {
      queue<int> q;
      q.push(i);
      while (!q.emptv()) {
        int x = q.front();
        q.pop();
        vis[x] = 1;
        for (auto const &y : adj[x]) {
          if (!vis[v]) {
            vis[y] = 1;
            q.push(match[y]);
  vector<int> vc;
  for (int i = 0; i < n; i++) {
    if (!vis[i]) vc.pb(i);
  for (int i = n; i < n + m; i++) {
    if (vis[i]) vc.pb(i);
  return vc:
vector<pi> mec() // minimum edge cover
  vector<pi> ans = run();
  for (int i = 0; i < n + m; i++) {
    if (match[i] == n + m \&\& adj[i].size() > 0) {
      if (i < n)
        ans.pb({i, adj[i][0]});
        ans.pb({adj[i][0], i});
```

11

```
return ans:
};
// minimum path cover on dag
\ensuremath{//} minimum set of paths such that each of the vertices belongs
    to exactly one path
vector<vector<int>> mpc(int n, vector<pi> &e) {
  hopcroft_karp h(n, n);
  for (auto const &i : e) h.add_edge(i.fir, n + i.sec);
  vector<pi> mat = h.run();
  vector<int> prv(n, -1);
  vector<int> nxt(n, -1);
  for (int i = 0; i < mat.size(); i++) {
   nxt[mat[i].fir] = mat[i].sec - n;
   prv[mat[i].sec - n] = mat[i].fir;
  vector<vector<int>> ans;
  for (int i = 0; i < n; i++) {
    if (prv[i] == -1 && nxt[i] == -1) {
     ans.pb({i});
   } else if (prv[i] == -1) {
     vector<int> curr;
     int x = i;
     while (1) {
       curr.pb(x);
       if (nxt[x] == -1) break;
       x = nxt[x];
     }
      ans.pb(curr);
  return ans;
```

pushrelabel.cpp

```
93 lines
#define INF 1e9
struct edge {
  int dest, back, f, c, id;
struct push_relabel {
  int n:
  vector<vector<edge>> g;
  vector<int> ec;
  vector<edge *> cur;
  vector<vector<int>> hs;
  vector<int> H:
  push_relabel(int sz) : g(sz), ec(sz), cur(sz), hs(2 * sz), H(
      sz) { n = sz; }
  void add_edge(int s, int t, int cap, int rcap, int id) {
   if (s == t) return;
    g[s].pb({t, (int)g[t].size(), 0, cap, id});
   g[t].pb({s, (int)g[s].size() - 1, 0, rcap, -1});
  void add_flow(edge &e, int f) {
    edge &back = g[e.dest][e.back];
   if (!ec[e.dest] && f) hs[H[e.dest]].push_back(e.dest);
   e.f += f;
   e.c -= f;
    ec[e.dest] += f;
   back.f -= f;
   back.c += f;
    ec[back.dest] -= f;
  int calc(int s, int t) {
    int v = q.size();
   H[s] = v;
```

```
ec[t] = 1;
    vector<int> co(2 * v);
    co[0] = v - 1;
    for (int i = 0; i < v; i++) cur[i] = g[i].data();
    for (edge &e : q[s]) add_flow(e, e.c);
    for (int hi = 0;;) {
      while (hs[hi].empty())
        if (!hi--) return -ec[s];
      int u = hs[hi].back();
      hs[hi].pop_back();
      while (ec[u] > 0) {
        if (cur[u] == g[u].data() + g[u].size()) {
          H[u] = INF;
          for (edge &e : g[u])
            if (e.c && H[u] > H[e.dest] + 1) H[u] = H[e.dest] +
                  1, cur[u] = &e;
          if (++co[H[u]], !--co[hi] && hi < v)
            for (int i = 0; i < v; i++)
              if (hi < H[i] && H[i] < v) --co[H[i]], H[i] = v +
          hi = H[u];
        } else if (\operatorname{cur}[u] \rightarrow \operatorname{c \&\& H}[u] == \operatorname{H}[\operatorname{cur}[u] \rightarrow \operatorname{dest}] + 1)
          add_flow(*cur[u], min(ec[u], cur[u]->c));
        else
          ++cur[u];
 vector<int> flow_edges(int m) // fluxo em cada aresta
    vector<int> ans(m);
    for (int i = 0; i < n; i++) {
     for (auto const &j : g[i]) {
        if (j.id != -1) ans[j.id] = j.f;
    }
    return ans:
struct flow_with_demands {
 push_relabel pr;
 vector<int> in, out;
 int n;
  flow_with_demands(int sz) : n(sz), pr(sz + 2), in(sz), out(sz
 void add_edge(int u, int v, int cap, int dem, int id) {
    pr.add_edge(u, v, cap - dem, 0, id);
    out[u] += dem, in[v] += dem;
 int run(int s, int t) {
    pr.add_edge(t, s, INF, 0, -1);
    for (int i = 0; i < n; i++) {
     pr.add_edge(n, i, in[i], 0, -1);
      pr.add_edge(i, n + 1, out[i], 0, -1);
    return pr.calc(n, n + 1);
 bool check() // todas as constraints foram satisfeitas?
    for (auto const &i : pr.g[n]) {
     if (i.c > 0) return 0;
    return 1;
};
```

```
hld.cpp
```

65 lines

12

```
struct hld {
 int n, cur_pos;
 segtree seg;
 vector<vector<int>> adj;
 vector<int> parent, depth, heavy, head, pos, sz;
 int dfs(int s) {
   int size = 1, max_c_size = 0;
   for (auto const &c : adj[s]) {
     if (c != parent[s]) {
       parent[c] = s;
       depth[c] = depth[s] + 1;
       int c_size = dfs(c);
       size += c size;
       if (c_size > max_c_size) max_c_size = c_size, heavy[s]
            = c;
   return sz[s] = size;
 void decompose(int s, int h) {
   head[s] = h;
   pos[s] = cur_pos++;
   if (heavy[s] != -1) decompose(heavy[s], h);
   for (int c : adj[s]) {
     if (c != parent[s] && c != heavy[s]) decompose(c, c);
 hld(vector<vector<int>> &g) {
   n = g.size();
   adj = q;
   seg = segtree(n);
   parent.assign(n, -1);
   depth.assign(n, -1);
   heavy.assign(n, -1);
   head.assign(n, -1);
   pos.assign(n, -1);
   sz.assign(n, 1);
   cur_pos = 0;
   dfs(0);
   decompose(0, 0);
 int query_path(int a, int b) {
   int res = 0:
   for (; head[a] != head[b]; b = parent[head[b]]) {
     if (depth[head[a]] > depth[head[b]]) swap(a, b);
     res += seq.query(0, n - 1, pos[head[b]], pos[b], 1);
   if (depth[a] > depth[b]) swap(a, b);
   res += seq.query(0, n - 1, pos[a], pos[b], 1);
   return res;
 void update_path(int a, int b, int x) {
   for (; head[a] != head[b]; b = parent[head[b]]) {
     if (depth[head[a]] > depth[head[b]]) swap(a, b);
     seg.update(1, 0, n - 1, pos[head[b]], pos[b], x);
   if (depth[a] > depth[b]) swap(a, b);
   seg.update(1, 0, n - 1, pos[a], pos[b], x);
 void update_subtree(int a, int x) { seq.update(1, 0, n - 1,
      pos[a], pos[a] + sz[a] - 1, x); }
 int query_subtree(int a) { return seq.query(0, n - 1, pos[a],
       pos[a] + sz[a] - 1, 1); }
 int lca(int a, int b) {
   if (pos[a] < pos[b]) swap(a, b);
   return (head[a] == head[b]) ? b : lca(parent[head[a]], b);
```

hldedge.cpp

```
};
```

```
67 lines
namespace hld {
int cur_pos;
vector<int> parent, depth, heavy, head, pos, sz, up;
int dfs(int s) {
 int size = 1, max c size = 0;
  for (auto const &c : adj[s]) {
    if (c.fir != parent[s]) {
     parent[c.fir] = s;
     depth[c.fir] = depth[s] + 1;
     int c size = dfs(c.fir);
     size += c size;
     if (c_size > max_c_size) max_c_size = c_size, heavy[s] =
  return sz[s] = size;
void decompose (int s, int h) {
  head[s] = h;
  pos[s] = cur_pos++;
  seq::v[pos[s]] = up[s];
  for (auto const &c : adi[s]) {
    if (c.fir != parent[s] && c.fir == heavy[s]) {
     up[c.fir] = c.sec;
      decompose(heavy[s], h);
  for (auto const &c : adj[s]) {
    if (c.fir != parent[s] && c.fir != heavy[s]) {
     up[c.fir] = c.sec;
     decompose(c.fir, c.fir);
void init() {
  parent.assign(MAXN, -1);
  depth.assign(MAXN, -1);
  heavy.assign(MAXN, -1);
  head.assign(MAXN, -1);
  pos.assign(MAXN, -1);
  sz.assign(MAXN, 1);
  up.assign(MAXN, 0);
  cur_pos = 0;
  dfs(0);
  decompose(0, 0);
  seg::build(0, n - 1, 1);
int query_path(int a, int b) {
  int res = -1;
  for (; head[a] != head[b]; b = parent[head[b]]) {
    if (depth[head[a]] > depth[head[b]]) swap(a, b);
    res = max(res, seg::query(0, n - 1, pos[head[b]], pos[b],
        1));
  if (depth[a] > depth[b]) swap(a, b);
  res = max(res, seg::query(0, n - 1, pos[a] + 1, pos[b], 1));
  return res;
void update_path(int a, int b, int x) {
  for (; head[a] != head[b]; b = parent[head[b]]) {
    if (depth[head[a]] > depth[head[b]]) swap(a, b);
    seg::update(1, 0, n - 1, pos[head[b]], pos[b], x);
  if (depth[a] > depth[b]) swap(a, b);
  seg::update(1, 0, n - 1, pos[a] + 1, pos[b], x);
```

```
void update_subtree(int a, int x) { seq::update(1, 0, n - 1,
    pos[a] + 1, pos[a] + sz[a] - 1, x); }
int query_subtree(int a, int x) { return seg::query(0, n - 1,
    pos[a] + 1, pos[a] + sz[a] - 1, 1); }
} // namespace hld
strongorientation.cpp
// encontrar uma orientacao para as arestas tal que o numero
// minimo de scc eh o menor possivel
// numero minimo de scc = numero de componentes conexas +
    numero de pontes
int n, m, timer, comps, bridges;
vector<pi> edges;
vector<pi> adj[MAXN];
int tin[MAXN]; // memset -1
int low[MAXN]; // memset -1
bool vis[MAXN];
char orient[MAXN];
void find_bridges(int v) { // chama se tem tin == -1
 low[v] = timer, tin[v] = timer++;
 for (auto const &p : adj[v]) {
   if (vis[p.sec]) continue;
   vis[p.sec] = true;
    orient[p.sec] = (v == edges[p.sec].first) ? '>' : '<';
    if (tin[p.fir] == -1) {
      find_bridges(p.fir);
     low[v] = min(low[v], low[p.fir]);
     if (low[p.fir] > tin[v]) bridges++;
     low[v] = min(low[v], low[p.fir]);
twosat.cpp
                                                          74 lines
struct two sat {
 int n;
 vector<vector<int>> q, qr; // gr is the reversed graph
 vector<int> comp, ord, ans; // comp[v]: ID of the SCC
      containing node v
 vector<bool> vis:
 two sat() {}
 two_sat(int sz) {
   n = sz;
   g.assign(2 * n, vector<int>());
   gr.assign(2 * n, vector<int>());
   comp.resize(2 * n);
   vis.resize(2 * n);
   ans.resize(2 * n);
 void add_edge(int u, int v) {
   q[u].push back(v);
   gr[v].push_back(u);
 // int x, bool val: if 'val' is true, we take the variable to
       be x. Otherwise we take it to be x's complement (not x)
 void implies(int i, bool f, int j, bool g) // a -> b
    add_edge(i + (f ? 0 : n), j + (g ? 0 : n));
   add_edge(j + (g ? n : 0), i + (f ? n : 0));
 void add_clause_or(int i, bool f, int j, bool g) // At least
       one of them is true
    add_edge(i + (f ? n : 0), j + (g ? 0 : n));
```

```
void add_clause_xor(int i, bool f, int j, bool g) // only
       one of them is true
    add_clause_or(i, f, j, g);
    add_clause_or(i, !f, j, !g);
  void add_clause_and(int i, bool f, int j, bool g) // both of
        them have the same value
    add_clause_xor(i, !f, j, g);
  void set(int i, bool f) // Set a variable
    add_clause_or(i, f, i, f);
 void top_sort(int u) {
    vis[u] = 1;
    for (auto const &v : g[u]) {
      if (!vis[v]) top_sort(v);
    ord.push back(u);
 void scc(int u, int id) {
    vis[u] = 1;
    comp[u] = id;
    for (auto const &v : gr[u]) {
      if (!vis[v]) scc(v, id);
 bool solve() {
    fill(vis.begin(), vis.end(), 0);
    for (int i = 0; i < 2 * n; i++) {
      if (!vis[i]) top_sort(i);
    fill(vis.begin(), vis.end(), 0);
    reverse(ord.begin(), ord.end());
    int id = 0;
    for (const auto &v : ord) {
      if (!vis[v]) scc(v, id++);
    for (int i = 0; i < n; i++) {
      if (comp[i] == comp[i + n]) return 0;
      ans[i] = (comp[i] > comp[i + n]) ? 1 : 0;
    return 1;
};
sack.cpp
vector<int> adj[MAXN];
vector<int> v[MAXN];
int c[MAXN];
int cnt[MAXN];
int sz[MAXN];
void dfs_sz(int x, int p) {
 sz[x] = 1;
  for (auto const &i : adj[x]) {
```

if (i != p) { dfs_sz(i, x);

sz[x] += sz[i];

void modify(int c, int val) { cnt[c] += val; }

void dfs(int x, int p, bool keep) {
 int best = -1, big_child = -1;

 $add_edge(j + (g ? n : 0), i + (f ? 0 : n));$

```
for (auto const &i : adj[x]) {
 if (i != p && sz[i] > best) {
   best = sz[i];
   big_child = i;
for (auto const &i : adj[x]) {
 if (i != p && i != big_child) dfs(i, x, 0);
if (big child !=-1) {
 dfs(big_child, x, 1);
 swap(v[x], v[big_child]); // 0(1)
v[x].pb(x);
modify(c[x], 1); // adiciona
for (auto const &i : adj[x]) {
 if (i != p && i != big_child) {
   for (auto const &j : v[i]) {
     v[x].pb(j);
     modify(c[j], 1); // adiciona
 }
// a cor c aparece cnt[c] vezes na subtree de x
// dai vc pode fazer algo tendo essa informacao
// seja responser queries ou algo do tipo aqui
  for (auto const &i : v[x]) modify(c[i], -1); // remove
```

reroot.cpp

```
23 lines
int n;
vector<int> adj[MAXN];
int sz[MAXN];
int dp[MAXN];
int dfs(int u, int v) {
  sz[u] = 1;
  for (auto const &i : adj[u])
   if (i != v) sz[u] += dfs(i, u);
  return sz[u];
void reroot(int u, int v) {
  for (auto const &i : adj[u]) {
   if (i != v) {
     int a = sz[u], b = sz[i];
     dp[i] = dp[u];
     dp[i] = sz[u], dp[i] = sz[i];
     sz[u] = sz[i], sz[i] = n;
     dp[i] += sz[u], dp[i] += sz[i];
      reroot(i, u);
      sz[u] = a, sz[i] = b;
```

hungarian.cpp

```
struct hungarian { // declarar algo como hungarian h(n), n de
    cada lado
  int n, inf;
  vector<vector<int>> a;
  vector<int> u, v, p, way;
  hungarian(int n_{-}): n(n_{-}), u(n + 1), v(n + 1), p(n + 1), way(
   a = vector<vector<int>>(n, vector<int>(n));
   inf = numeric_limits<int>::max();
```

```
void add_edge(int x, int y, int c) { a[x][y] = c; }
 pair<int, vector<int>> run() {
    for (int i = 1; i \le n; i++) {
     p[0] = i;
      int j0 = 0;
      vector<int> minv(n + 1, inf);
      vector < int > used(n + 1, 0);
      do {
       used[j0] = true;
       int i0 = p[j0], j1 = -1;
       int delta = inf;
       for (int j = 1; j \le n; j++) {
         if (!used[j]) {
           int cur = a[i0 - 1][j - 1] - u[i0] - v[j];
           if (cur < minv[j]) minv[j] = cur, way[j] = j0;
            if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
        for (int j = 0; j \le n; j++) {
         if (used[j])
           u[p[j]] += delta, v[j] -= delta;
            minv[j] -= delta;
        j0 = j1;
      } while (p[j0] != 0);
      do {
       int j1 = wav[j0];
       p[j0] = p[j1];
        j0 = j1;
     } while (j0);
    vector<int> ans(n);
    for (int j = 1; j \le n; j++) ans[p[j] - 1] = j - 1;
    return make_pair(-v[0], ans);
};
```

mincostmaxflow.cpp

```
namespace mcf {
struct edge {
 int to, capacity, cost, res;
int source, destiny;
vector<edge> adj[MAXN];
vector<int> dist;
vector<int> parent;
vector<int> edge_index;
vector<bool> in_queue;
void add_edge(int a, int b, int c, int d) {
 adj[a].pb({b, c, d, (int)adj[b].size()});
 adj[b].pb({a, 0, -d, (int)adj[a].size() - 1});
bool dijkstra(int s) {
 dist.assign(MAXN, INF);
  parent.assign(MAXN, -1);
  edge_index.assign(MAXN, -1);
 in_queue.assign(MAXN, false);
  dist[s] = 0;
  queue<int> q;
  q.push(s);
  while (!q.empty()) {
    int u = q.front(), idx = 0;
    q.pop();
    in_queue[u] = false;
    for (auto const &v : adj[u]) {
      if (v.capacity && dist[v.to] > dist[u] + v.cost) {
        dist[v.to] = dist[u] + v.cost;
```

```
parent[v.to] = u;
       edge_index[v.to] = idx;
       if (!in_queue[v.to]) {
         in_queue[v.to] = true;
         q.push(v.to);
     idx++;
 return dist[destiny] != INF;
int get_cost() {
 int flow = 0, cost = 0;
 while (dijkstra(source)) {
   int curr_flow = INF, curr = destiny;
   while (curr != source) {
     int p = parent[curr];
     curr_flow = min(curr_flow, adj[p][edge_index[curr]].
          capacity);
     curr = p;
    flow += curr flow;
   cost += curr_flow * dist[destiny];
   curr = destiny;
   while (curr != source) {
     int p = parent[curr];
     int res_idx = adj[p][edge_index[curr]].res;
     adj[p][edge_index[curr]].capacity -= curr_flow;
     adj[curr][res_idx].capacity += curr_flow;
 return cost;
} // namespace mcf
```

scc.cpp

```
int n, m;
bool vis[MAXN];
int root[MAXN];
vector<int> order;
vector<int> roots;
vector<int> comp;
vector<vector<int>> comps;
vector<int> adj[MAXN];
vector<int> adj rev[MAXN];
vector<int> adj_scc[MAXN];
void dfs(int v) {
 vis[v] = true;
 for (auto const &u : adj[v])
   if (!vis[u]) dfs(u);
 order.pb(v);
void dfs2(int v) {
 comp.pb(v);
 vis[v] = true;
 for (auto const &u : adj rev[v])
    if (!vis[u]) dfs2(u);
//...
cin >> n >> m;
for (int i = 0; i < m; i++) {
 int a, b;
 cin >> a >> b;
  adj[a].pb(b);
 adj_rev[b].pb(a);
```

articulationpoints lca topsort eulerian virtualtree notes

```
for (int i = 0; i < n; i++) {
 if (!vis[i]) dfs(i);
reverse(order.begin(), order.end());
memset(vis, false, sizeof(vis));
for (auto const &v : order) {
 if (!vis[v]) {
    comp.clear();
   dfs2(v);
    comps.pb(comp);
   // making condensation graph
    int r = comp.back();
    for (auto const &u : comp) root[u] = r;
    roots.push back(r);
// making condensation graph
for (int v = 0; v < n; v++) {
  for (auto const &u : adj[v]) {
   int root_v = roots[v];
   int root_u = roots[u];
   if (root_u != root_v) adj_scc[root_v].pb(root_u);
}
```

articulationpoints.cpp

```
int n, m, timer;
vector<int> adj[MAXN];
bool is_cutpoint[MAXN];
int tin[MAXN]; // memset -1
int low[MAXN]; // memset -1
bool vis[MAXN];
void dfs(int v, int p) { // chama pros nao vis
  vis[v] = true;
  tin[v] = timer, low[v] = timer++;
  int childs = 0;
  for (auto const &u : adj[v]) {
   if (u == p) continue;
   if (vis[u]) {
     low[v] = min(low[v], tin[u]);
    } else {
     dfs(u, v);
     low[v] = min(low[v], low[u]);
     if (low[u] >= tin[v] && p != -1) is_cutpoint[v] = true;
 if (p == -1 && childs > 1) is_cutpoint[v] = true;
```

lca.cpp

int n: vector<int> adi[MAXN]; namespace lca { int 1, timer; vector<int> tin, tout, depth; vector<vector<int>> up; void dfs(int v, int p) { tin[v] = ++timer;up[v][0] = p;for (int i = 1; $i \le 1$; i + 1) up[v][i] = up[up[v][i - 1]][i - 1]for (auto const &u : adj[v]) { if (p == u) continue; depth[u] = depth[v] + 1;dfs(u, v);

```
tout[v] = ++timer;
bool is ancestor(int u, int v) { return tin[u] <= tin[v] &&
    tout[u] >= tout[v]; }
int binary_lifting(int u, int v) {
 if (is_ancestor(u, v)) return u;
 if (is_ancestor(v, u)) return v;
 for (int i = 1; i >= 0; --i)
   if (!is\_ancestor(up[u][i], v)) u = up[u][i];
 return up[u][0];
void init() {
 tin.resize(n);
 tout.resize(n);
 depth.resize(n);
 timer = 0:
 l = ceil(log2(n));
 up.assign(n, vector<int>(1 + 1));
 dfs(0, 0);
int dist(int s, int v) {
 int at = binary_lifting(s, v);
 return (depth[s] + depth[v] - 2 * depth[at]);
} // namespace lca
```

topsort.cpp

vector<bool> visited;

```
vector<int> ans:
void dfs(int v) {
 visited[v] = true;
 for (int u : adj[v]) {
   if (!visited[u]) {
     dfs(u);
 ans.push_back(v);
void topological_sort() {
 visited.assign(n, false);
 ans.clear();
 for (int i = 0; i < n; ++i) {
   if (!visited[i]) {
     dfs(i);
 reverse(ans.begin(), ans.end());
```

eulerian.cpp

39 lines

```
void dfs2(int s) { // caminho euleriano
 while (a[s].size() > 0) {
   auto v = a[s].back();
   a[s].pop_back();
   if (!vis[v.sec]) {
     vis[v.sec] = 1;
     dfs2(v.fir);
 path.pb(s);
void dfs(int i) { // ajeitar para que todo mundo tenha grau
 vis[i] = 1;
 for (auto const &j : adj[i]) {
   if (!vis[j.fir]) {
     dfs(j.fir);
```

```
if (deg[j.fir]) {
       ans.pb(edges[j.sec]);
       deg[j.fir] ^= 1;
       deg[i] ^= 1;
// se eu cham dfs(0) no final checar se o deg[0] ta safe
```

virtualtree.cpp

29 lines

```
// virtual tree
// dado um conjunto de vertices v
// montar uma arvore comprimida
// tal que escolhendo dois vertices do conjunto v[i] e v[j]
// lca(v[i], v[j]) tambem ta na arvore
// se o conjunto v tem k vertices
// entao a arvore comprimida tem menos do que 2k vertices
// O(k log(k)), sem considerar a complexidade de achar lca
int build virt(vector<int> v) {
 auto cmp = [&](int i, int j) { return lca::tin[i] < lca::tin[</pre>
       j]; };
 sort(v.begin(), v.end(), cmp);
 for (int i = v.size() - 1; i > 0; i--) {
   v.pb(lca::find_lca(v[i], v[i - 1]));
 sort(v.begin(), v.end(), cmp);
 v.erase(unique(v.begin(), v.end()), v.end());
 for (int i = 0; i < v.size(); i++) {
   virt[v[i]].clear();
 for (int i = 1; i < v.size(); i++) {
   virt[lca::find lca(v[i - 1], v[i])].clear();
 for (int i = 1; i < v.size(); i++) {
   int parent = lca::find lca(v[i - 1], v[i]);
   int d = lca::dist(parent, v[i]);
   virt[parent].pb({v[i], d});
 return v[0];
```

notes.md

Edge Cover

Path Cover

least one path.

Koning's Theorem

21 lines

Bipartite Graph A bipartite graph is a graph that does not contain any odd-

```
length cycles.
## Directed acyclic graph (DAG)
Is a directed graph with no directed cycles.
## Independent Set
Is a set of vertices in a graph, no two of which are adjacent.
    That is, it is a set S of vertices such that for every two
     vertices in S, there is no edge connecting the two.
Is a subset of vertices of an undirected graph such that every
    two distinct vertices in the clique are adjacent.
## Vertex Cover
Is a set of vertices that includes at least one endpoint of
    every edge of the graph.
```

Is a set of edges such that every vertex of the graph is

Given a directed graph G = (V, E), a path cover is a set of

directed paths such that every vertex \boldsymbol{v} belongs to at

incident to at least one edge of the set.

convexhull convexhullpointlocation

```
In any bipartite graph, the number of edges in a maximum
    matching equals the number of vertices in a minimum vertex
```

Properties

- Every tree is a bipartite graph.

- Any NxM grid is a bipartite graph.
- A set of vertices is a vertex cover if and only if its complement is an independent set.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent
- In bipartite graphs, the size of the minimum edge cover is equal to the size of the maximum independent set
- In bipartite graphs, the size of the minimum edge cover plus the size of the minimum vertex cover is equal to the number of vertices.
- In bipartite graphs, maximum clique size is two.

Min-cut

The smallest total weight of the edges which if removed would disconnect the source from the sink.

Max-flow min-cut theorem

In a flow network, the maximum amount of flow passing from the source to the sink is equal to the total weight of the edges in a minimum cut.

Maximum flow with vertex capacities

In other words, the amount of flow passing through a vertex cannot exceed its capacity. To find the maximum flow, we can transform the problem into the maximum flow problem by expanding the network. Each vertex v is replaced by v-in and v-out, where v-in is connected by edges going into v and v-out is connected to edges coming out from v. Then assign capacity c(v) to the edge connecting v-in and v-out

Undirected edge-disjoint paths problem

We are given an undirected graph G = (V, E) and two vertices s and t, and we have to find the maximum number of edgedisjoint s-t paths in G.

Undirected vertex-disjoint paths problem

We are given an undirected graph G = (V, E) and two vertices s and t, and we have to find the maximum number of vertexdisjoint (except for s and t) paths in G.

Menger's theorem

The maximum number of edge-disjoint s-t paths in an undirected graph is equal to the minimum number of edges in an s-t cut-set.

Undirected vertex-disjoint paths solution

We can construct a network N=(V,E) from G with vertex capacities, where the capacities of all vertices and all edges are 1. Then the value of the maximum flow is equal to the maximum number of independent paths from s to t.

Minimum vertex-disjoint path cover in directed acyclic graph

Given a directed acyclic graph G=(V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V. We can construct a bipartite graph G' from G. Each vertex v is replaced by v-in and v-out, where v-in is connected by edges going into v and v-out is connected to edges coming out from v. Then it can be shown that G' has a matching M of size m if and only if G has a vertex-

disjoint path cover C of containing m edges and n-m paths. ## Minimum general path cover in directed acyclic graph (DAG) A general path cover is a path cover where a vertex can belong to more than one path. A minimum general path cover may be smaller than a minimum vertex-disjoint path cover. A minimum general path cover can be found almost like a minimum vertex-disjoint path cover. It suffices to add some new edges to the matching graph so that there is an edge a to b always when there is a path from a to b in the original graph.

```
## Dilworths theorem and maximum antichain
An antichain is a set of nodes of a graph such that there is no
     path from any node to another node using the edges of the
     graph. Dilworths theorem states that in a directed
     acyclic graph, the size of a minimum general path cover
     equals the size of a maximum antichain.
Or in other words: For a DAG G that if has edges from vertex i
     to vertex i and vertex i to k, then it also has a edge
     from vertex i to vertex k, the size of a minimum path
     cover is equal to the size of a maximum independent set.
## Maximum weighted antichain
In this problem, each vertex has a cost a[i]. The cost of an
     antichain is equal to the sum of the costs of the vertices
     present in it. We need to find the maximum cost of a
     antichain. We can construct the same bipartite of the
     maximum antichain problem from a dag G, these edges have
     an infinite capacity. We also need to create a source
     vertex and a sink, and we need to add edges source to v-in
     with capacity a[v] and v-out to sink with capacity a[v].
     The answer is equal to the sum of all a[i] minus the
     maximum flow on this network.
## Halls Theorem
Halls theorem can be used to find out whether a bipartite graph
     has a matching that contains all left or right nodes.
     Assume that we want to find a matching that contains all
     left nodes. Let X be any set of left nodes and let f(X) be
     the set of their neighbors. According to Halls theorem, a
     matching that contains all left nodes exists exactly when
each X, the condition |X| \le |f(X)| holds.
## Extra (Getting Confidence Trick)
If you need to maximize a number x = (a * b * c * ...), then
     you can write it as x = (e^{\log(a)} * e^{\log(b)} * e^{\log(c)} *
     ...), and then the number is x = e^{(\log(a) + \log(b) + \log a)}
     (c) + \dots), and the problem now becomes a problem of
     maximizing the sum of (\log(a) + \log(b) + \log(c) + \ldots).<p
Use exp() and log() C++ functions :)
```

Geometry (5)

convexhull.cpp

```
47 lines
struct point {
 int x, y, id;
 point(int x, int y, int id) : x(x), y(y), id(id) {}
 point() {}
 point operator-(point const &o) const { return {x - o.x, y -
      o.y, -1; }
 bool operator < (point const &o) const {
   if (x == o.x) return y < o.y;
   return x < o.x;
 int operator (point const &o) const { return x * o.y - y * o.
int ccw(point const &a, point const &b, point const &x) {
 auto p = (b - a) ^ (x - a);
 return (p > 0) - (p < 0);
vector<point> convex_hull(vector<point> P) // sem colineares
 sort(P.begin(), P.end());
 vector<point> L, U;
 for (auto p : P) {
   while (L.size() \geq 2 && ccw(L.end()[-2], L.end()[-1], p) ==
         -1) L.pop_back();
   L.push_back(p);
```

```
reverse(P.begin(), P.end());
  for (auto p : P) {
    while (U.size() \ge 2 \&\& ccw(U.end()[-2], U.end()[-1], p) = 
         -1) U.pop_back();
    U.push_back(p);
 L.insert(L.end(), U.begin(), U.end() - 1);
 return L;
vector<point> convex_hull_no_collinears(vector<point> P) //
    com colineares
 sort(P.begin(), P.end());
 vector<point> L, U;
  for (auto p : P) {
    while (L.size() \geq 2 && ccw(L.end()[-2], L.end()[-1], p) <=
         0) L.pop_back();
    L.push_back(p);
  reverse(P.begin(), P.end());
  for (auto p : P) {
    while (U.size() \ge 2 \&\& ccw(U.end()[-2], U.end()[-1], p) \le
         0) U.pop_back();
    U.push_back(p);
 L.insert(L.end(), U.begin(), U.end() - 1);
 return L:
```

convexhullpointlocation.cpp

44 lines

```
struct pt {
 int x, v;
 pt operator+(pt p) { return \{x + p.x, y + p.y\}; }
  pt operator-(pt p) { return {x - p.x, y - p.y}; }
  bool operator==(pt p) { return (x == p.x && y == p.y); }
  int cross(pt p) { return x * p.y - y * p.x; }
  int cross(pt a, pt b) { return (a - *this).cross(b - *this);
 int dot(pt p) { return x * p.x + y * p.y; }
bool cmp_x(pt a, pt b) {
 if (a.x != b.x) return a.x < b.x;
 return a.y < b.y;
// acha o convex hull
vector<pt> convex hull(vector<pt> pts) {
 if (pts.size() <= 1) return pts;
 sort(pts.begin(), pts.end(), cmp_x);
 vector<pt> h(pts.size() + 1);
 int s = 0, t = 0;
  for (int it = 2; it--; s = --t, reverse(pts.begin(), pts.end
       ())) {
    for (auto const &p : pts) {
      while (t >= s + 2 \&\& h[t - 2].cross(h[t - 1], p) <= 0) t
      h[t++] = p;
 return \{h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])\};
int sqn(int x) { return (x > 0) - (x < 0); }
int side_of(pt s, pt e, pt p) { return sgn(s.cross(e, p)); }
bool on_segment(pt s, pt e, pt p) { return p.cross(s, e) == 0
    && (s - p).dot(e - p) \le 0;}
// retorna se o ponto p esta dentro ou nao do convex hull l
// caso strict = true, entao considera true se tiver na borda
```

// caso strict = false, entao considera false se tiver na borda

linetrick polygonarea geometria

```
bool is_hull(vector<pt> &1, pt p, bool strict = true) {
  int a = 1, b = 1.size() - 1, r = !strict;
  if (1.size() < 3) return r && on segment(1[0], 1.back(), p);
  if (side_of(1[0], 1[a], 1[b]) > 0) swap(a, b);
  if (side_of(1[0], 1[a], p) >= r || side_of(1[0], 1[b], p) <=
      -r) return false;
  while (abs(a - b) > 1) {
   int c = (a + b) / 2;
    (side_of(1[0], 1[c], p) > 0 ? b : a) = c;
  return sqn(l[a].cross(l[b], p)) < r;
linetrick.cpp
                                                           13 lines
pi get_line(pi x, pi y) { // um jeito normalizado de
    representar a reta entre 2 pontos
  int xx = x.fir - y.fir;
  int yy = x.sec - y.sec;
  int g = \underline{gcd(abs(xx), abs(yy))};
  if (g != 0) {
   xx /= g, yy /= g;
  if (xx < 0) {
   xx \star = -1;
   yy \star = -1;
  return {xx, yy};
polygonarea.cpp
                                                           19 lines
double area(vector<pi> fig) {
  double res = 0;
  for (unsigned i = 0; i < fig.size(); i++) {
   pi p = i ? fiq[i - 1] : fiq.back();
   pi q = fig[i];
   res += (p.fir - q.fir) * (p.sec + q.sec);
  return fabs(res) / 2;
int cross(pi a, pi b) { return a.fir * b.sec - a.sec * b.fir; }
double area2(vector<pi> fig) {
  double res = 0;
  for (unsigned i = 0; i < fig.size(); i++) {
   pi p = i ? fiq[i - 1] : fiq.back();
   pi q = fiq[i];
```

geometria.cpp

res += cross(p, q);

return fabs(res) / 2;

```
typedef double ld;
const ld DINF = le18;
const ld pi = acos(-1.0);
const ld eps = le-9;
#define sq(x) ((x) * (x))
bool eq(ld a, ld b) { return abs(a - b) <= eps; }
struct pt { // ponto
    ld x, y;
    pt(ld x_ = 0, ld y_ = 0) : x(x_), y(y_) {}
    bool operator<(const pt p) const {
        if (!eq(x, p.x)) return x < p.x;
        if (!eq(y, p.y)) return y < p.y;
        return 0;
}</pre>
```

```
bool operator == (const pt p) const { return eq(x, p.x) and eq(
      y, p.y); }
  pt operator+(const pt p) const { return pt(x + p.x, y + p.y);
  pt operator-(const pt p) const { return pt(x - p.x, y - p.y);
  pt operator*(const ld c) const { return pt(x * c, y * c); }
  pt operator/(const ld c) const { return pt(x / c, y / c); }
  ld operator*(const pt p) const { return x * p.x + y * p.y; }
  ld operator^(const pt p) const { return x * p.y - y * p.x; }
  friend istream& operator>>(istream& in, pt& p) { return in >>
       p.x >> p.y; }
struct line { // reta
 pt p, q;
 line() {}
 line(pt p_, pt q_) : p(p_), q(q_) {}
 friend istream& operator>>(istream& in, line& r) { return in
      >> r.p >> r.q; }
// PONTO & VETOR
ld dist(pt p, pt q) { // distancia
 return hypot (p.y - q.y, p.x - q.x);
ld dist2(pt p, pt q) { // quadrado da distancia
 return sq(p.x - q.x) + sq(p.y - q.y);
ld norm(pt v) { // norma do vetor
 return dist(pt(0, 0), v);
ld angle(pt v) { // angulo do vetor com o eixo x
 1d ang = atan2(v.y, v.x);
 if (ang < 0) ang += 2 * pi;
 return ang;
ld sarea(pt p, pt q, pt r) { // area com sinal
 return ((q - p) ^ (r - q)) / 2;
bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
 return eq(sarea(p, q, r), 0);
bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
 return sarea(p, q, r) > eps;
pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
 return pt(p.x * cos(th) - p.y * sin(th), p.x * sin(th) + p.y
      * cos(th));
pt rotate90(pt p) { // rotaciona 90 graus
 return pt(-p.y, p.x);
// RETA
bool isvert(line r) { // se r eh vertical
 return eq(r.p.x, r.q.x);
bool isinseq(pt p, line r) { // se p pertence ao seq de r
 pt a = r.p - p, b = r.q - p;
 return eq((a ^ b), 0) and (a * b) < eps;
ld get t(pt v, line r) { // retorna t tal gue t*v pertence a
    reta r
 return (r.p ^ r.q) / ((r.p - r.q) ^ v);
pt proj(pt p, line r) { // projecao do ponto p na reta r
 if (r.p == r.q) return r.p;
 r.q = r.q - r.p;
 p = p - r.p;
 pt proj = r.q * ((p * r.q) / (r.q * r.q));
  return proj + r.p;
```

```
pt inter(line r, line s) { // r inter s
 if (eq((r.p - r.q) ^ (s.p - s.q), 0)) return pt(DINF, DINF);
 r.q = r.q - r.p, s.p = s.p - r.p, s.q = s.q - r.p;
 return r.q * qet_t(r.q, s) + r.p;
bool interseg(line r, line s) { // se o seg de r intersecta o
 if (isinseg(r.p, s) or isinseg(r.q, s) or isinseg(s.p, r) or
       isinseg(s.q, r)) return 1;
 return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and ccw(s.p, r.q, s.q)
       s.q, r.p) != ccw(s.p, s.q, r.q);
ld disttoline(pt p, line r) { // distancia do ponto a reta
 return 2 * abs(sarea(p, r.p, r.q)) / dist(r.p, r.q);
ld disttoseg(pt p, line r) { // distancia do ponto ao seg
 if ((r.q - r.p) * (p - r.p) < 0) return dist(r.p, p);
 if ((r.p - r.q) * (p - r.q) < 0) return dist(r.q, p);
 return disttoline(p, r);
ld distseg(line a, line b) { // distancia entre seg
 if (interseg(a, b)) return 0;
 ld ret = DINF;
 ret = min(ret, disttoseg(a.p, b));
 ret = min(ret, disttoseg(a.g, b));
 ret = min(ret, disttoseg(b.p, a));
 ret = min(ret, disttoseg(b.g, a));
 return ret;
// POLIGONO
// corta poligono com a reta r deixando os pontos p tal que
// ccw(r.p, r.q, p)
vector<pt> cut_polygon(vector<pt> v, line r) { // O(n)
 vector<pt> ret;
  for (int j = 0; j < v.size(); j++) {
   if (ccw(r.p, r.q, v[i])) ret.push back(v[i]);
    if (v.size() == 1) continue;
    line s(v[j], v[(j + 1) % v.size()]);
    pt p = inter(r, s);
    if (isinseg(p, s)) ret.push_back(p);
  ret.erase(unique(ret.begin(), ret.end()), ret.end());
 if (ret.size() > 1 and ret.back() == ret[0]) ret.pop_back();
// distancia entre os retangulos a e b (lados paralelos aos
// assume que ta representado (inferior esquerdo, superior
ld dist rect(pair<pt, pt> a, pair<pt, pt> b) {
 ld hor = 0, vert = 0;
 if (a.second.x < b.first.x)
    hor = b.first.x - a.second.x;
  else if (b.second.x < a.first.x)
   hor = a.first.x - b.second.x;
 if (a.second.y < b.first.y)</pre>
   vert = b.first.v - a.second.v;
  else if (b.second.v < a.first.v)</pre>
   vert = a.first.y - b.second.y;
 return dist(pt(0, 0), pt(hor, vert));
ld polarea(vector<pt> v) { // area do poligono
 1d ret = 0:
  for (int i = 0; i < v.size(); i++) ret += sarea(pt(0, 0), v[i]
      ], v[(i + 1) % v.size()]);
  return abs(ret);
```

1d x = (d * d - R * R + r * r) / (2 * d);

```
// se o ponto ta dentro do poligono: retorna 0 se ta fora,
// 1 se ta no interior e 2 se ta na borda
int inpol(vector<pt>& v, pt p) { // O(n)
  int at = 0:
  for (int i = 0; i < v.size(); i++) {
   if (p == v[i]) return 2;
   int j = (i + 1) % v.size();
   if (eq(p.y, v[i].y) and eq(p.y, v[j].y)) {
     if ((v[i] - p) * (v[j] - p) < eps) return 2;
     continue;
   bool baixo = v[i].y + eps < p.y;
   if (baixo == (v[j].y + eps < p.y)) continue;</pre>
    auto t = (p - v[i]) ^ (v[j] - v[i]);
   if (eq(t, 0)) return 2;
   if (baixo == (t > eps)) qt += baixo ? 1 : -1;
  return qt != 0;
bool interpol(vector<pt> v1, vector<pt> v2) { // se dois
    poligonos se intersectam - O(n*m)
  int n = v1.size(), m = v2.size();
  for (int i = 0; i < n; i++)
   if (inpol(v2, v1[i])) return 1;
  for (int i = 0; i < n; i++)
   if (inpol(v1, v2[i])) return 1;
  for (int i = 0; i < n; i++)
    for (int j = 0; j < m; j++)
     if (interseg(line(v1[i], v1[(i + 1) % n]), line(v2[j], v2
          [(j + 1) % m]))) return 1;
  return 0;
ld distpol(vector<pt> v1, vector<pt> v2) { // distancia entre
  if (interpol(v1, v2)) return 0;
 ld ret = DINF;
  for (int i = 0; i < v1.size(); i++)
    for (int j = 0; j < v2.size(); j++)
      ret = min(ret, distseg(line(v1[i], v1[(i + 1) % v1.size()
          ]), line(v2[j], v2[(j + 1) % v2.size()])));
 return ret;
// CIRCUNFERENCIA
pt getcenter(pt a, pt b, pt c) { // centro da circunf dado 3
    pontos
 b = (a + b) / 2;
 c = (a + c) / 2;
  return inter(line(b, b + rotate90(a - b)), line(c, c +
      rotate90(a - c)));
vector<pt> circ_line_inter(pt a, pt b, pt c, ld r) { //
    intersecao da circunf (c, r) e reta ab
  vector<pt> ret:
 b = b - a, a = a - c;
 1d A = b * b;
 1d B = a * b;
 1d C = a * a - r * r;
 1d D = B * B - A * C;
 if (D < -eps) return ret;
  ret.push_back(c + a + b \star (-B + sqrt(D + eps)) / A);
 if (D > eps) ret.push_back(c + a + b * (-B - sqrt(D)) / A);
 return ret;
vector<pt> circ_inter(pt a, pt b, ld r, ld R) { // intersecao
    da circunf (a, r) e (b, R)
  vector<pt> ret;
 ld d = dist(a, b);
 if (d > r + R \text{ or } d + min(r, R) < max(r, R)) return ret;
```

```
1d v = sart(r * r - x * x);
  pt v = (b - a) / d;
  ret.push_back(a + v * x + rotate90(v) * y);
  if (y > 0) ret.push_back(a + v * x - rotate90(v) * y);
  return ret;
geometriaint.cpp
                                                          105 lines
#define sq(x) ((x) * (ll)(x))
struct pt { // ponto
 int x, y;
  pt(int x_{-} = 0, int y_{-} = 0) : x(x_{-}), y(y_{-}) {}
  bool operator<(const pt p) const {
    if (x != p.x) return x < p.x;
    return y < p.y;
  bool operator == (const pt p) const { return x == p.x and y ==
      p.v; }
  pt operator+(const pt p) const { return pt(x + p.x, y + p.y);
  pt operator-(const pt p) const { return pt(x - p.x, y - p.y);
  pt operator*(const int c) const { return pt(x * c, y * c); }
  11 operator*(const pt p) const { return x * (ll)p.x + y * (ll
      (;v.a(
  11 operator (const pt p) const { return x * (11)p.v - v * (11
  friend istream& operator>>(istream& in, pt& p) { return in >>
       p.x >> p.y; }
struct line { // reta
  pt p, q;
 line() {}
 line(pt p_, pt q_) : p(p_), q(q_) {}
  friend istream& operator>>(istream& in, line& r) { return in
      >> r.p >> r.q; }
// PONTO & VETOR
11 dist2(pt p, pt q) { // quadrado da distancia
 return sq(p.x - q.x) + sq(p.y - q.y);
ll sarea2(pt p, pt q, pt r) { // 2 * area com sinal
 return (q - p) ^ (r - q);
bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
 return sarea2(p, q, r) == 0;
bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
 return sarea2(p, q, r) > 0;
int quad(pt p) { // quadrante de um ponto
 return (p.x < 0) ^ 3 * (p.y < 0);
bool compare_angle(pt p, pt q) { // retorna se ang(p) < ang(q)</pre>
 if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
 return ccw(q, pt(0, 0), p);
pt rotate90(pt p) { // rotaciona 90 graus
  return pt(-p.y, p.x);
// RETA
bool isinseg(pt p, line r) { // se p pertence ao seg de r
  pt a = r.p - p, b = r.q - p;
 return (a ^{\circ} b) == 0 and (a * b) <= 0;
bool interseg(line r, line s) { // se o seg de r intersecta o
    seg de s
```

```
if (isinseg(r.p, s) or isinseg(r.q, s) or isinseg(s.p, r) or
      isinseq(s.q, r)) return 1;
 return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and ccw(s.p, r.q, s.q)
      s.q, r.p) != ccw(s.p, s.q, r.q);
int segpoints(line r) { // numero de pontos inteiros no
 return 1 + __gcd(abs(r.p.x - r.q.x), abs(r.p.y - r.q.y));
double get_t(pt v, line r) { // retorna t tal que t*v pertence
     a reta r
 return (r.p ^ r.q) / (double) ((r.p - r.q) ^ v);
// POLIGONO
// quadrado da distancia entre os retangulos a e b (lados
    paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior
    direito)
11 dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
 int hor = 0, vert = 0;
 if (a.second.x < b.first.x)
   hor = b.first.x - a.second.x;
  else if (b.second.x < a.first.x)
   hor = a.first.x - b.second.x;
  if (a.second.y < b.first.y)</pre>
   vert = b.first.y - a.second.y;
  else if (b.second.v < a.first.v)</pre>
   vert = a.first.y - b.second.y;
  return sq(hor) + sq(vert);
11 polarea2(vector<pt> v) { // 2 * area do poligono
 11 \text{ ret} = 0;
  for (int i = 0; i < v.size(); i++) ret += sarea2(pt(0, 0), v[
      i], v(i + 1) % v.size());
 return abs(ret);
// se o ponto ta dentro do poligono: retorna 0 se ta fora,
// 1 se ta no interior e 2 se ta na borda
int inpol(vector<pt>& v, pt p) { // O(n)
 int qt = 0;
  for (int i = 0; i < v.size(); i++) {
    if (p == v[i]) return 2;
    int j = (i + 1) % v.size();
    if (p.y == v[i].y \text{ and } p.y == v[j].y) {
      if ((v[i] - p) * (v[j] - p) \le 0) return 2;
      continue;
    bool baixo = v[i].v < p.v;
    if (baixo == (v[j].y < p.y)) continue;
    auto t = (p - v[i]) ^ (v[j] - v[i]);
    if (!t) return 2;
    if (baixo == (t > 0)) gt += baixo ? 1 : -1;
 return qt != 0;
11 interior_points(vector<pt> v) { // pontos inteiros dentro
    de um poligono simples
 11 b = 0;
 for (int i = 0; i < v.size(); i++) b += segpoints(line(v[i],
      v(i + 1) % v.size())) - 1;
 return (polarea2(v) - b) / 2 + 1;
```

geometria3d.cpp

132 line

```
typedef double ld;
const ld DINF = 1e18;
const ld eps = 1e-9;
```

```
\#define sq(x) ((x) * (x))
bool eq(ld a, ld b) { return abs(a - b) <= eps; }
struct pt { // ponto
  ld x, y, z;
  pt(ld x_{=} = 0, ld y_{=} = 0, ld z_{=} = 0) : x(x_{-}), y(y_{-}), z(z_{-}) {}
  bool operator<(const pt p) const {
    if (!eq(x, p.x)) return x < p.x;
    if (!eq(y, p.y)) return y < p.y;
   if (!eq(z, p.z)) return z < p.z;
   return 0:
  bool operator==(const pt p) const { return eq(x, p.x) and eq(
      y, p.y) and eq(z, p.z); }
  pt operator+(const pt p) const { return pt(x + p.x, y + p.y,
      z + p.z); }
  pt operator-(const pt p) const { return pt(x - p.x, y - p.y,
      z - p.z); }
  pt operator*(const ld c) const { return pt(x * c, y * c, z *
  pt operator/(const ld c) const { return pt(x / c, y / c, z /
  ld operator*(const pt p) const { return x * p.x + y * p.y + z
  pt operator^(const pt p) const { return pt(y * p.z - z * p.y,
       z * p.x - x * p.z, x * p.y - y * p.x); }
  friend istream& operator>>(istream& in, pt& p) { return in >>
       p.x >> p.y >> p.z; }
struct line { // reta
 pt p, q;
  line() {}
  line(pt p_, pt q_) : p(p_), q(q_) {}
  friend istream& operator>>(istream& in, line& r) { return in
      >> r.p >> r.q; }
struct plane {
                   // plano
  array<pt, 3> p; // pontos que definem o plano
  array<ld, 4> eq; // equacao do plano
  plane() {}
  plane(pt p_, pt q_, pt r_) : p({p_, q_, r_}) { build(); }
  friend istream& operator>>(istream& in, plane& P) {
    return in >> P.p[0] >> P.p[1] >> P.p[2];
   P.build();
  void build() {
   pt dir = (p[1] - p[0]) ^ (p[2] - p[0]);
    eq = \{dir.x, dir.y, dir.z, dir * p[0] * (-1)\};
}:
// converte de coordenadas polares para cartesianas
// (angulos devem estar em radianos)
// phi eh o angulo com o eixo z (cima) theta eh o angulo de
    rotacao ao redor de z
pt convert(ld rho, ld th, ld phi) { return pt(sin(phi) * cos(th
    ), sin(phi) * sin(th), cos(phi)) * rho; }
// projecao do ponto p na reta r
pt proj(pt p, line r) {
  if (r.p == r.q) return r.p;
  r.q = r.q - r.p;
  p = p - r.p;
  pt proj = r.q * ((p * r.q) / (r.q * r.q));
  return proj + r.p;
// projecao do ponto p no plano P
pt proj(pt p, plane P) {
  p = p - P.p[0], P.p[1] = P.p[1] - P.p[0], P.p[2] = P.p[2] - P
      .p[0];
  pt norm = P.p[1] ^ P.p[2];
```

```
pt proj = p - (norm * (norm * p) / (norm * norm));
 return proj + P.p[0];
// distancia
ld dist(pt a, pt b) { return sqrt(sq(a.x - b.x) + sq(a.y - b.y)
      + sq(a.z - b.z)); }
// distancia ponto reta
ld distline(pt p, line r) { return dist(p, proj(p, r)); }
// distancia de ponto para segmento
ld distseg(pt p, line r) {
  if ((r.q - r.p) * (p - r.p) < 0) return dist(r.p, p);
 if ((r.p - r.q) * (p - r.q) < 0) return dist(r.q, p);
 return distline(p, r);
// distancia de ponto a plano com sinal
ld sdist(pt p, plane P) { return P.eq[0] * p.x + P.eq[1] * p.y
    + P.eq[2] * p.z + P.eq[3]; }
// distancia de ponto a plano
ld distplane(pt p, plane P) { return abs(sdist(p, P)); }
// se ponto pertence a reta
bool isinseq(pt p, line r) { return eq(distseq(p, r), 0); }
// se ponto pertence ao triangulo definido por P.p
bool isinpol(pt p, vector<pt> v) {
 assert(v.size() >= 3);
  pt norm = (v[1] - v[0]) ^ (v[2] - v[1]);
  bool inside = true;
  int sign = -1;
  for (int i = 0; i < v.size(); i++) {
    line r(v[(i + 1) % 3], v[i]);
    if (isinseq(p, r)) return true;
    pt ar = v[(i + 1) % 3] - v[i];
    if (sign == -1)
      sign = ((ar ^ (p - v[i])) * norm > 0);
    else if (((ar ^ (p - v[i])) * norm > 0) != sign)
      inside = false;
  return inside;
// distancia de ponto ate poligono
ld distpol(pt p, vector<pt> v) {
  pt p2 = proj(p, plane(v[0], v[1], v[2]));
  if (isinpol(p2, v)) return dist(p, p2);
  ld ret = DINF;
  for (int i = 0; i < v.size(); i++) {
    int j = (i + 1) % v.size();
    ret = min(ret, distseq(p, line(v[i], v[j])));
 return ret:
// intersecao de plano e segmento
// BOTH = o segmento esta no plano
// ONE = um dos pontos do segmento esta no plano
// PARAL = segmento paralelo ao plano
// CONCOR = segmento concorrente ao plano
enum RETCODE { BOTH, ONE, PARAL, CONCOR };
pair<RETCODE, pt> intersect(plane P, line r) {
  1d d1 = sdist(r.p, P);
 1d d2 = sdist(r.q, P);
  if (eq(d1, 0) \text{ and } eq(d2, 0)) return pair(BOTH, r.p);
  if (eq(d1, 0)) return pair(ONE, r.p);
  if (eq(d2, 0)) return pair(ONE, r.q);
  if ((d1 > 0 \text{ and } d2 > 0) \text{ or } (d1 < 0 \text{ and } d2 < 0)) {}
    if (eq(d1 - d2, 0)) return pair(PARAL, pt());
    return pair (CONCOR, pt());
  1d frac = d1 / (d1 - d2);
  pt res = r.p + ((r.q - r.p) * frac);
  return pair (ONE, res);
```

```
// rotaciona p ao redor do eixo u por um angulo a
pt rotate(pt p, pt u, ld a) {
    u = u / dist(u, pt());
    return u * (u * p) + (u ^ p ^ u) * cos(a) + (u ^ p) * sin(a);
}
```

19

```
kdtree.cpp
struct pt {
 int x, y, id;
 pt() {}
  pt(int xx, int yy) { x = xx, y = yy; }
  pt operator-(pt p) const { return pt(x - p.x, y - p.y); }
 bool operator<(pt p) const { return x < p.x; }</pre>
 int dist() const { return x * x + y * y; }
bool on_x(const pt &a, const pt &b) { return a.x < b.x; }
bool on_y(const pt &a, const pt &b) { return a.y < b.y; }
struct node {
 pt pp;
  int id:
  int x0 = inf, x1 = -inf, y0 = inf, y1 = -inf;
  node *first = 0, *second = 0;
 int distance(const pt &p) {
   int x = (p.x < x0 ? x0 : p.x > x1 ? x1 : p.x);
    int y = (p.y < y0 ? y0 : p.y > y1 ? y1 : p.y);
    return (pt(x, y) - p).dist();
 node(vector<pt> &&vp) : pp(vp[0]) {
    for (pt p : vp) {
      x0 = min(x0, p.x);
      x1 = max(x1, p.x);
      v0 = min(v0, p.v);
     y1 = max(y1, p.y);
    if (vp.size() > 1) {
      sort(vp.begin(), vp.end(), x1 - x0 >= y1 - y0 ? on_x :
          on_y);
      int half = vp.size() / 2;
      first = new node({vp.begin(), vp.begin() + half});
      second = new node({vp.begin() + half, vp.end()});
};
struct kd_tree {
 node *root:
  kd_tree(const vector<pt> &vp) : root(new node({vp.begin(), vp
       .end()})) {}
 pi search (node *n, const pt &p) {
   if (!n->first) {
      if (n->pp.x == p.x && n->pp.y == p.y) return make_pair(
          inf, n->pp.id); // distancia infinita pra pontos
          iquais
      return make_pair((p - n->pp).dist(), n->pp.id);
    node *f = n- first, *s = n- second;
    int bfirst = f->distance(p), bsec = s->distance(p);
    if (bfirst > bsec) swap(bsec, bfirst), swap(f, s);
    auto best = search(f, p);
    if (bsec < best.first || (!f->first)) best = min(best,
         search(s, p));
    return best;
 pi nearest(const pt &p) { return search(root, p); }
```

20

halfplaneintersection.cpp 75 lines const long double eps = 1e-9; const long double inf = 1e9; struct pt { long double x, y; pt(long double x = 0, long double y = 0) : x(x), y(y) {} friend pt operator+(pt p, pt q) { return pt(p.x + q.x, p.y + friend pt operator-(pt p, pt q) { return pt(p.x - q.x, p.y q.y); } friend pt operator*(pt p, long double k) { return pt(p.x * k, $p.y * k); }$ friend long double dot(pt p, pt q) { return p.x * q.x + p.y * friend long double cross(pt p, pt q) { return p.x * q.y - p.y * q.x; } struct halfplane { pt p, pq; long double angle; halfplane() {} halfplane(pt a, pt b) : p(a), pq(b - a) { angle = atan21(pq.y bool out (const pt &r) { return cross(pq, r - p) < -eps; } bool operator<(halfplane e) const { return angle < e.angle; }</pre> friend pt inter(halfplane s, halfplane t) { long double alpha = cross((t.p - s.p), t.pq) / cross(s.pq,return s.p + (s.pq * alpha); }; vector<pt> hp_intersect(vector<halfplane> &h) { $pt box[4] = {pt(inf, inf), pt(-inf, inf), pt(-inf, -inf), pt(}$ inf, -inf)}; // Bounding box in CCW order for (int i = 0; i < 4; i++) { halfplane aux(box[i], box[(i + 1) % 4]);h.pb(aux); sort(h.begin(), h.end()); deque<halfplane> dq; int len = 0; for (int i = 0; i < h.size(); i++) { while (len > 1 && h[i].out(inter(dq[len - 1], dq[len - 2])) dq.pop_back(); while (len > 1 && h[i].out(inter(dq[0], dq[1]))) { dq.pop_front(); --len; if (len > 0 && fabsl(cross(h[i].pq, dq[len - 1].pq)) < eps)if (dot(h[i].pq, dq[len - 1].pq) < 0.0) { return vector<pt>(); if (h[i].out(dq[len - 1].p)) { dq.pop_back(); --len; } else {

while (len > 2 && dq[0].out(inter(dq[len - 1], dq[len - 2])))

continue;

dq.push_back(h[i]);

++len;

```
dq.pop_back();
   --len:
 while (len > 2 && dq[len - 1].out(inter(dq[0], dq[1]))) {
   dq.pop_front();
   --len;
 if (len < 3) {
   return vector<pt>();
 vector<pt> ret(len);
 for (int i = 0; i + 1 < len; i++) {
   ret[i] = inter(dq[i], dq[i + 1]);
 ret.back() = inter(dq[len - 1], dq[0]);
 return ret;
// half-plane - regiao planar que consiste de todos os pontos
    que estao de um lado de uma reta
// e dai pros half-planes, considerando que eh a regiao da
    esquerda em relacao ao vetor de direcao
```

Strings (6)

ahocorasick.cpp namespace aho {

v = ac[v].down[c];

ac[v].is_leaf++;

```
int go(int v, char ch);
const int K = 26: // tamanho do alfabeto
struct trie {
 char me:
                       // char correspondente ao no atual
                       // proximo vertice que eu devo ir
 int go[K]:
      estando em um estado (v, c)
 int down[K];
                       // proximo vertice da trie
 int is leaf = 0;
                       // se o vertice atual da trie eh uma
      folha (fim de uma ou mais strings)
 int parent = -1:
                      // no ancestral do no atual
 int link = -1;
                       // link de sufixo do no atual (outro no
      com o maior matching de sufixo)
 int exit_link = -1; // folha mais proxima que pode ser
      alcancada a partir de v usando links de sufixo
 trie(int p = -1, char ch = '$') : parent(p), me(ch) {
    fill (begin (go), end (go), -1);
    fill (begin (down), end (down), -1);
};
vector<trie> ac;
void init() // criar a raiz da trie
 ac.resize(1);
void add_string(string s) // adicionar string na trie
 int v = 0;
 for (auto const &ch : s) {
   int c = ch - 'a';
   if (ac[v].down[c] == -1) {
     ac[v].down[c] = ac.size();
     ac.emplace_back(v, ch);
```

int get_link(int v) // pegar o suffix link saindo de v

: go(get_link(ac[v].parent), ac[v].me);

if (ac[v].link == -1) ac[v].link = (!v || !ac[v].parent) ? 0

```
return ac[v].link;
int go(int v, char ch) // proximo estado saindo do estado(v,
 int c = ch - 'a';
 if (ac[v].go[c] == -1) {
   if (ac[v].down[c] != -1)
     ac[v].go[c] = ac[v].down[c];
     ac[v].go[c] = (!v) ? 0 : go(get_link(v), ch);
 return ac[v].go[c];
int get_exit_link(int v) // suffix link mais proximo de v que
    seja uma folha
 if (ac[v].exit_link == -1) {
   int curr = get_link(v);
   if (!v || !curr)
     ac[v].exit_link = 0;
   else if (ac[curr].is_leaf)
     ac[v].exit link = curr;
     ac[v].exit_link = get_exit_link(curr);
 return ac[v].exit_link;
int query(string s) // query O(n + ans)
 int ans = 0, curr = 0, at;
 for (auto const &i : s) {
   curr = qo(curr, i);
   ans += ac[curr].is_leaf;
   at = get_exit_link(curr);
   while (at) {
     ans += ac[at].is_leaf;
     at = get_exit_link(at);
 return ans;
  // namespace aho
```

kmp.cpp

```
string s;
int n, m;
string a, b;
int c[MAXN][26];
vector<int> kmp(string &s) {
  int n = s.size();
  vector<int> p(n);
  for (int i = 1; i < n; i++) {
    int j = p[i - 1];
    while (j > 0 \&\& s[i] != s[j]) j = p[j - 1];
    if (s[i] == s[j]) j++;
    p[i] = j;
  return p;
void compute(string s) {
  s.pb('*');
  vector < int > p = kmp(s);
  for (int i = 0; i < s.size(); i++) {
    for (int cc = 0; cc < 26; cc++) {
      int j = i;
      while (j > 0 \&\& 'a' + cc != s[j]) j = p[j - 1];
      if ('a' + cc == s[j]) j++;
```

```
c[i][cc] = j;
manacher.cpp
                                                            22 lines
```

```
vector<int> d1:
vector<int> d2:
void manacher(string s) {
 dl.resize(s.size());
  d2.resize(s.size());
  int 1 = 0, r = -1;
  for (int i = 0; i < s.size(); i++) {
   int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
   while (0 \le i - k \&\& i + k \le s.size() \&\& s[i - k] == s[i + k]
        k]) k++;
   d1[i] = k;
   k = k - 1;
   if (i + k > r) l = i - k, r = i + k;
 1 = 0, r = -1;
  for (int i = 0; i < s.size(); i++) {
   int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1);
   while (0 \le i - k - 1 \&\& i + k < s.size() \&\& s[i - k - 1]
        == s[i + k]) k++;
   d2[i] = k;
   k = k - 1;
   if (i + k > r) l = i - k - 1, r = i + k;
```

minsuffix.cpp

```
int max_suffix(string s, bool mi = false) {
  s.push_back(*min_element(s.begin(), s.end()) - 1);
  int ans = 0;
  for (int i = 1; i < s.size(); i++) {
   int j = 0;
   while (ans + j < i and s[i + j] == s[ans + j]) j++;
   if (s[i + j] > s[ans + j]) {
     if (!mi or i != s.size() - 2) ans = i;
   } else if (j)
     i += j - 1;
  return ans;
int min_suffix(string s) {
  for (auto &i : s) i \star = -1;
  s.push_back(*max_element(s.begin(), s.end()) + 1);
  return max_suffix(s, true);
int max_cyclic_shift(string s) {
  int n = s.size();
  for (int i = 0; i < n; i++) s.pb(s[i]);
  return max_suffix(s);
int min_cyclic_shift(string s) {
  for (auto &i : s) i \star = -1;
  return max_cyclic_shift(s);
// retorna a posicao de inicio menor/maior sufixo/shift de uma
```

hashing.cpp

int mulmod(int a, int b) {

const int MOD = (111 << 61) - 1;

```
const static int LOWER = (111 << 30) - 1, GET31 = (111 << 31)
  int 11 = a \& LOWER, h1 = a >> 30, 12 = b \& LOWER, h2 = b >>
      30;
  int m = 11 * h2 + 12 * h1, h = h1 * h2;
  int ans = 11 * 12 + (h >> 1) + ((h & 1) << 60) + (m >> 31) +
       ((m & GET31) << 30) + 1;
 ans = (ans & MOD) + (ans >> 61), ans = (ans & MOD) + (ans >>
      61);
 return ans - 1;
mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().
int uniform(int 1, int r) {
 uniform_int_distribution<int> uid(l, r);
 return uid(rng);
struct string_hashing -
 vector<int> h, p;
  string_hashing() {}
  string_hashing(string s) : h(s.size()), p(s.size()) {
   p[0] = 1, h[0] = s[0];
    for (int i = 1; i < s.size(); i++) p[i] = mulmod(p[i - 1],
        P), h[i] = (mulmod(h[i - 1], P) + s[i]) % MOD;
  int get(int 1, int r) {
    int hash = h[r] - (1 ? mulmod(h[1 - 1], p[r - 1 + 1]) : 0);
    return hash < 0 ? hash + MOD : hash;
  int append(int h, int hb, int blen) { return (hb + mulmod(h,
      p[blen])) % MOD; }
// lembrar do P = uniform(256, MOD - 1);
```

suffixarray.cpp

ra[sa[i]] = i;

```
vector<int> suffix_array(string s) {
 s += "$"; // menor do que todos os chars da string st
 int n = s.size(), N = max(n + 1, 26011);
 vector<int> sa(n), ra(n);
 for (int i = 0; i < n; i++) {
   sa[i] = i, ra[i] = s[i];
 for (int k = 0; k < n; k ? k *= 2 : k++) {
   vector<int> nsa(sa), nra(n), cnt(N);
   for (int i = 0; i < n; i++) {
     nsa[i] = (nsa[i] - k + n) % n;
     cnt[ra[i]]++;
   for (int i = 1; i < N; i++) {
     cnt[i] += cnt[i - 1];
   for (int i = n - 1; i + 1; i--) {
     sa[--cnt[ra[nsa[i]]]] = nsa[i];
   for (int i = 1, r = 0; i < n; i++) {
     nra[sa[i]] = r += (ra[sa[i]] != ra[sa[i - 1]] || ra[(sa[i
          ] + k) % n] != ra[(sa[i - 1] + k) % n]);
   ra = nra;
   if (ra[sa[n - 1]] == n - 1) break;
 return vector<int>(sa.begin() + 1, sa.end());
vector<int> kasai(string s, vector<int> sa) {
 int n = s.size(), k = 0;
 vector<int> ra(n), lcp(n);
 for (int i = 0; i < n; i++) {
```

```
for (int i = 0; i < n; i++, k -= !!k) {
 if (ra[i] == n - 1) {
    k = 0;
    continue;
  int j = sa[ra[i] + 1];
  while (i + k < n \text{ and } j + k < n \text{ and } s[i + k] == s[j + k]) k
  lcp[ra[i]] = k;
return lcp;
```

zfunction.cpp

```
vector<int> z_function(string &s) {
 int n = s.size();
 vector<int> z(n);
 z[0] = n;
 for (int i = 1, l = 0, r = 0; i < n; i++) {
   if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
   while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]]) z[i]++;
   if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
 return z;
// z[i] = o tamanho de lcp(s, s.substr(i, n - i))
```

suffixautomaton.cpp

```
namespace sa {
struct state {
 int len, suf link;
 map<char, int> nxt;
vector<int> term;
state st[2 * MAXN];
int dp[2 * MAXN];
int sz, last;
void init() {
 memset(dp, -1, sizeof(dp));
 st[0].len = 0;
 st[0].suf_link = -1;
 sz++;
 last = 0;
void get_link(int curr, int p, char c) {
 while (p != -1 \&\& !st[p].nxt.count(c)) {
    st[p].nxt[c] = curr;
   p = st[p].suf_link;
 if (p == -1) {
    st[curr].suf_link = 0;
    return;
  int q = st[p].nxt[c];
 if (st[p].len + 1 == st[q].len) {
    st[curr].suf_link = q;
    return;
  int clone = sz;
  st[clone].len = st[p].len + 1;
  st[clone].nxt = st[q].nxt;
  st[clone].suf_link = st[q].suf_link;
  while (p != -1 \&\& st[p].nxt[c] == q) {
    st[p].nxt[c] = clone;
    p = st[p].suf_link;
```

```
st[a].suf link = clone;
  st[curr].suf link = clone;
void build(string &s) {
  for (auto const &c : s) {
    int curr = sz;
   st[curr].len = st[last].len + 1;
    get link(curr, last, c);
   last = curr;
  // achar os estados terminais
  // um estado terminal eh aquele que representa um sufixo da
      string s
  int p = last;
  while (p != -1) {
   term.pb(p);
    p = st[p].suf_link;
void dfs2(int v) {
 if (dp[v] != -1) return;
  dp[v] = 1;
  for (auto const &u : st[v].nxt) {
   if (!u.sec) continue;
   dfs2(u.sec);
   dp[v] += dp[u.sec];
void dfs(int v, int k, int &at, string &curr) {
  if (at == k) return;
  for (auto const &u : st[v].nxt) {
   if (!u.sec) continue;
    if (at + dp[u.sec] < k) {
     at += dp[u.sec];
     continue;
    curr.pb(u.fir);
    at++;
    dfs(u.sec, k, at, curr);
    if (at == k) return;
    curr.pop_back();
void find_kth(int k) {
  int at = 0;
  string curr = "";
  dfs(0, k, at, curr);
  cout << curr << endl;
} // namespace sa
// chamar sa::init(); e sa::build(s);
```

eertree.cpp

44 lines

```
struct eertree {
  vector<vector<int>> t;
  int n, last, sz;
  vector<int> s, len, link, qt;
  eertree(int N) {
   t = vector(N + 2, vector<int>(26, 0));
   s = len = link = qt = vector < int > (N + 2);
   link[0] = 1, len[0] = 0, link[1] = 1, len[1] = -1;
   sz = 2, last = 0, n = 1;
  int add(char c) {
   int ret = -1;
```

```
s[n++] = c -= 'a';
   while (s[n - len[last] - 2] != c) last = link[last];
   if (!t[last][c]) {
     int prev = link[last];
     while (s[n - len[prev] - 2] != c) prev = link[prev];
     link[sz] = t[prev][c];
     len[sz] = len[last] + 2;
     ret = len[sz];
     t[last][c] = sz++;
   qt[last = t[last][c]]++;
   return ret;
 int size() { return sz - 2; }
 int propagate() {
   int ret = 0;
   for (int i = n; i > 1; i--) {
     qt[link[i]] += qt[i];
     ret += qt[i];
   return ret;
 vector<pi> get palindromes(string &s) {
   vector<pi> ans;
   for (int i = 0; i < s.size(); i++) {
     int len = add(s[i]);
     if (len != -1) ans.pb((i - len + 1, i));
   return ans;
};
```

Various (7)

```
binarysearch.cpp
```

```
26 lines
// busca ternaria para int, usando busca binaria:
int 1 = 0, r = 1e9;
while (1 < r) {
 int mid = (1 + r) >> 1;
 (calc(mid) < calc(mid + 1)) ? r = mid : 1 = mid + 1;
return calc(1);
// first element >= x
int lower(int l, int r, int x) // first element >= x
 while (1 < r) {
   // int mid = 1 + (r - 1) / 2; se tiver numero negativo
   int mid = (1 + r) >> 1;
    (x \le k[mid]) ? r = mid : 1 = mid + 1;
 return k[1];
// last element <= x
vector<int> k(MAXN);
int upper(int 1, int r, int x) {
 while (l < r) {
   int mid = (1 + r + 1) >> 1;
    (k[mid] \le x) ? 1 = mid : r = mid - 1;
 return k[1];
```

ternarysearch.cpp

```
double eps = 1e-9;
while (r - 1 > eps) {
 double m1 = 1 + (r - 1) / 3;
```

```
double m2 = r - (r - 1) / 3;
  double f1 = f(m1);
 double f2 = f(m2);
 if (f1 < f2)
   1 = m1;
 else
   r = m2:
return f(1); // o maximo de f(1)
```

parallelbinarysearch.cpp

```
vector<pii> qry(q);
for (int i = 0; i < q; i++) {
 cin >> qry[i].sec.fir >> qry[i].sec.sec >> qry[i].fir;
 qry[i].sec.fir--, qry[i].sec.sec--;
vector<int> l(n);
vector<int> r(n);
vector<vector<int>> on(q);
for (int i = 0; i < n; i++) {
 1[i] = 0;
 r[i] = q;
while (1) {
 bool ok = 1;
 for (int i = 0; i < n; i++) {
   if (l[i] < r[i]) {
      ok = 0;
      int mid = (l[i] + r[i]) >> 1;
      on[mid].pb(i);
 if (ok) break;
  reset();
  for (int mid = 0; mid < q; mid++) {
    upd(qry[mid].sec.fir, qry[mid].sec.sec, qry[mid].fir);
    for (auto const &j : on[mid]) {
      int val = 0;
      for (auto const &k : adj[j]) {
       val += sum(k);
       if (val >= need[i]) break;
      (val >= need[j]) ? r[j] = mid : l[j] = mid + 1;
    on[mid].clear();
// a resposta eh o l[i] para cada busca
```

alienstrick.cpp

```
pi solve (vector<int> &v, int lambda) {
 // associar um custo lambda para ser subtraido quando
       realizamos uma operacao
 // dp[i] - melhor profit que tivemos considerando as i
      primeiras posicoes
  // cnt[i] - quantas operacoes utilizamos para chegar no valor
        de dp[i]
 vector < int > dp(n + 1);
  vector<int> cnt(n + 1);
 dp[0] = 0;
 cnt[0] = 0;
  for (int i = 1; i \le n; i++) {
    dp[i] = dp[i - 1];
    cnt[i] = cnt[i - 1];
    int id = i - 1;
    dp[i] += v[id];
    int lo = \max(011, id - 1 + 1);
```

if (s > dp[i]) {

dp[i] = s;

```
cnt[i] = cnt[lo] + 1;
  return {dp[n], cnt[n]};
int aliens_trick(vector<int> &v) {
  int 1 = 0, r = n;
  while (l < r) {
   int mid = (1 + r) >> 1;
   pi ans = solve(v, mid);
    (ans.sec > k) ? 1 = mid + 1 : r = mid;
  pi ans = solve(v, 1);
  return ans.fir + (1 * k);
cht.cpp
struct line {
  mutable int m, b, p;
  bool operator<(const line &o) const {
   if (m != o.m) return m < o.m;
    return b < o.b;
  bool operator<(const int x) const { return p < x; }
  int eval(int x) const { return m * x + b; }
  int inter(const line &o) const {
   int x = b - o.b, y = o.m - m;
    return (x / y) - ((x ^ y) < 0 && x % y);
};
struct cht {
  int INF = 1e18;
  multiset<line, less<>> 1;
  void add(int m, int b) {
    auto y = l.insert({m, b, INF});
    auto z = next(y);
    if (z != 1.end() \&\& y->m == z->m) {
     l.erase(y);
     return:
    if (y != l.begin()) {
     auto x = prev(y);
     if (x->m == y->m) x = 1.erase(x);
    while (1) {
     if (z == 1.end()) {
       y->p = INF;
       break:
      y->p = y->inter(*z);
     if (y->p < z->p)
       break:
      else
        z = 1.erase(z);
    if (y == 1.begin()) return;
    z = y;
    auto x = --y;
    while (1) {
      int ninter = x \rightarrow inter(*z);
     if (ninter <= x->p)
       x->p = ninter;
      else {
       l.erase(z);
       break;
```

int s = dp[lo] + (id - lo + 1) - lambda;

```
if (x == 1.begin()) break;
     y = x;
     x--;
      if (x->p < y->p)
       break;
      else
       l.erase(y);
 int get(int x) {
   if (l.empty()) return 0;
   return l.lower_bound(x)->eval(x);
};
expected valued p.cpp
                                                          26 lines
double dp[MAXN][MAXN][MAXN];
// sushi problem
double solve(int i, int j, int k) {
 if (!i \&\& !j \&\& !k) return dp[i][j][k] = 0;
 if (dp[i][j][k] != -1) return dp[i][j][k];
 It is well-known from statistics that for the geometric
  (counting number of trials before a success, where each
      independent trial is probability p)
 the expected value is i / p
 double p = ((double)(i + j + k) / n);
 double ret = 1 / p; // expected number of trials before a
      success
 if (i) {
   double prob = (double)i / (i + j + k); // probabilidade de
         ser um prato com um sushi
   ret += (solve(i - 1, j, k) * prob);
 if (j) {
   double prob = (double) j / (i + j + k); // probabilidade de
         ser um prato com dois sushis
   ret += (solve(i + 1, j - 1, k) * prob);
 if (k) {
   double prob = (double)k / (i + j + k); // probabilidade de
         ser um prato com tres sushis
   ret += (solve(i, j + 1, k - 1) * prob);
 return dp[i][j][k] = ret;
kadane.cpp
int kadane (vector<int> v) {
 int n = v.size(), ans = 0, max_here = 0;
 for (int i = 0; i < n; i++) {
   max_here += v[i];
   if (ans < max_here) ans = max_here;
   if (\max here < 0) \max here = 0;
 return ans;
int kadane_circular(vector<int> v) {
 int n = v.size(), max kadane = kadane(v);
 int max_wrap = 0, i;
 for (i = 0; i < n; i++) {
   max_wrap += v[i];
   v[i] = -v[i];
```

```
max wrap += kadane(v);
 return max(max_wrap, max_kadane);
divideandconquerdp.cpp
                                                          16 lines
void compute(int 1, int r, int opt1, int optr, int i) {
 if (1 > r) return;
 int mid = (1 + r) >> 1;
 pair<int, int> ans = \{1e18, -1\}; // dp, k
 for (int q = optl; q <= min(mid, optr); q++) {
     ans = min(ans, \{dp[i-1][q-1] + cost(q, mid), q\});
     ans = min(ans, {cost(q, mid), q});
 dp[i][mid] = ans.fir;
 compute(1, mid - 1, optl, ans.sec, i);
 compute(mid + 1, r, ans.sec, optr, i);
// dp[i][j] = min(dp[i - 1][k] + c(k, j)), para algum k <= j
// O(k * n * log(n))
sosdp.cpp
// nesse caso, f[x] eh a funcaoo que soma:
// todos os a[i], tal que, (x & i) == i)
for (int mask = 0; mask < (1 << m); mask++) {
 f[mask] = a[mask];
for (int i = 0; i < m; ++i) {
 for (int mask = 0; mask < (1 << m); mask++) {
    if (mask & (1 << i)) f[mask] += f[mask ^ (1 << i)];
// nesse caso, f[x] eh a funcao que soma:
// todos os a[i], tal que, (x & i) == x)
for (int mask = 0; mask < (1 << m); mask++) {
 f[mask] = a[mask];
for (int i = 0; i < m; ++i) {
 for (int mask = 0; mask < (1 << m); mask++) {
    if (!(mask & (1 << i))) f[mask] += f[mask ^ (1 << i)];
lis.cpp
int lis() {
 for (int i = 0; i < v.size(); i++) {
    vector<int>::iterator it = lower bound(q.begin(), q.end(),
        v[i]);
   if (it == q.end())
     q.pb(v[i]);
   else
      *it = v[i];
 for (int i = 0; i < q.size(); i++) cout << q[i] << " ";
 cout << endl;
 return q.size();
largestsquareofones.cpp
for (int i = 0; i < n; i++) {
 for (int j = 0; j < n; j++) {
   dp[i][j] = v[i][j];
```

```
if (i \&\& j \&\& dp[i][j]) dp[i][j] = min(\{dp[i][j-1], dp[i])
        -1][j], dp[i -1][j -1]}) + 1;
    ans = max(ans, dp[i][j]);
cout << ans * ans << endl:
```

steinertree.cpp

16 lines

```
for (int i = 0; i \le 100; i++)
 for (int j = 0; j < (1 << 9); j++) dp[i][j] = 1e9;
for (int i = 0; i < k; i++)
 for (int j = 0; j < v.size(); j++) dp[j][1 << i] = d(cap[i],
for (int mask = 2; mask < (1 << k); mask++) {
 for (int i = 0; i < v.size(); i++) {
   for (int mask2 = 1; mask2 < mask; mask2++) {
     if (!((mask & mask2) == mask2)) continue;
     int mask3 = mask ^ mask2;
     dp[i][mask] = min(dp[i][mask], dp[i][mask2] + dp[i][mask3]
   for (int j = 0; j < v.size(); j++) dp[j][mask] = min(dp[j][
        mask], dp[i][mask] + d(v[i], v[j]));
double ans = 1e9:
for (int i = 0; i \le 100; i++) ans = min(ans, dp[i][(1 << k) -
    1]);
```

subsetsum.cpp

89 lines

```
// subset sum com bitset de tamanho variado
// with n <= 10^6
template <int len = 1>
int subset_sum(int n, int h) {
 if (n \ge len) {
   return subset_sum<std::min(len * 2, (int)MAXN)>(n, h);
 bitset<len> dp;
  dp[0] = 1;
  for (auto const &x : w) {
   dp = dp \mid (dp \ll x);
  return dp._Find_next(max(011, h - 1)); // retorna o proximo
      bit setado apos a posicao passada como parametro
int solve(vector<int> &w, int tot, int h) {
 // tot -> soma de todos os elementos de w
 // h -> valor desejado
  // quero retornar o menor valor x \ge h, tal que existe um
      subset com soma x em w
  if (!w.size()) return 0;
  sort(w.rbegin(), w.rend());
  if (w[0] * 2 >= tot) return w[0];
  int n = w.size();
  w.pb(0);
  vector<int> aux;
 int p = 0;
  for (int i = 1; i \le n; i++) {
   if (w[i] != w[i - 1]) {
     int cnt = i - p;
     int x = w[i - 1];
     int j = 1;
     while (j < cnt) {
       aux.pb(x * j);
       cnt -= j;
       j *= 2;
```

```
aux.pb(x * cnt);
     p = i;
 swap(aux, w);
 return subset sum(tot, h);
int f[MAXN];
                  // f[i] -> quantos "itens" com valor i tem
bitset<MAXN> dp; // dp[i] = 1, se existe um subset com soma i
// garantir que a soma de todo mundo seja < MAXN
void subset_sum(vector<int> &v) {
 for (auto const &i : v) {
   f[i]++;
 dp[0] = 1;
  for (int i = 1; i < MAXN; i++) {
   while (f[i] > 2) {
     f[i * 2]++;
     f[i] -= 2;
    while (f[i]--) dp \mid= (dp << i);
 Given N non-negative integer weights w and a non-negative
 computes the maximum S \le t such that S is the sum of some
     subset of the weights.
 O(N * max(w[i]))
*/
int knapsack(vector<int> w, int t) {
 int a = 0, b = 0;
 while (b < w.size() \&\& a + w[b] <= t) {
   a += w[b++];
 if (b == w.size()) {
    return a;
 int m = *max element(w.begin(), w.end());
  vector<int> u, v(2 * m, -1);
 v[a + m - t] = b;
  for (int i = b; i < w.size(); i++) {
    for (int x = 0; x < m; x++) {
      v[x + w[i]] = max(v[x + w[i]], u[x]);
    for (int x = 2 * m; --x > m;) {
      for (int j = max(011, u[x]); j < v[x]; j++) v[x - w[j]] =
            max(v[x - w[i]], i);
 a = t;
  while (v[a + m - t] < 0) {
   a--:
 return a:
bitmasks.cpp
                                                          19 lines
```

```
// quantidade de bits setados na mask
cout << __builtin_popcount(mask) << endl;</pre>
// quando eh necessario percorrer todas as submasks ate (1 << n
// e fazer algo com todas as submasks dessa mask O(3^n)
for (int m = 0; m < (1 << n); m++) {
 for (int s = m; s; s = (s - 1) \& m) {
    // alguma coisa aqui sabendo que mask s eh uma submask de m
```

```
// comprimindo as masks de um vector baseada em uma mask
    qualquer
for (int i = 0; i < masks.size(); i++) {
 int compressed = 0, curr_bit = 0;
 for (int j = 0; j < n; j++) {
   if (!(mask & (1LL << j))) continue;
    if (masks[i] & (1LL << j)) compressed |= (1LL << curr_bit);</pre>
    curr_bit++;
 // alguma coisa sabendo que a mask compressed eh a mask
      comprimida da mask atual
```

hanoi.cpp

18 lines

```
vector<pair<char, char>> ans;
void solve(int n, char a, char b, char c) {
 if (n == 0) return;
 solve(n - 1, a, c, b);
 ans.pb({a, b});
 solve(n - 1, c, b, a);
// chamar pra sol
solve(n, 'A', 'C', 'B');
// 3 pilhas, sendo a pilha A com n discos e as outras duas
    pilhas vazias
// quero mover todo mundo da pilha A para a pilha C
// em cada movimento, vc tira o disco do topo de uma pilha e
    poe no topo de outra pilha
// desde que o raio do disco seja menor do que o raio do disco
    que ta no topo da outra pilha
// os n discos tem raios distintos aos pares
// e printar os movimentos em um par [to, from]
// numero minimo pra resolver pros primeiros n
// 1, 3, 7, 15, 31, 63, 127, 255
// f(n) = 2^n - 1
```

stacktrick.cpp

9 lines

```
stack<pi> s:
for (int i = n - 1; i >= 0; i--) {
 while (!s.empty() && s.top().fir <= v[i]) s.pop();
 (!s.empty()) ? ans[i] = s.top().sec : ans[i] = -1;
 s.push(\{v[i], i\});
// for each index (0 \le i \le n), find another index (0 \le j \le n)
// which v[j] > v[i] and j > i and j is as close as possible to
// if this index does not exist, print -1
```

spraguegrundy.cpp

```
vector<int> v = \{2, 3, 4, 5, 6\};
unordered map<int, bool> vis;
unordered_map<int, int> dp;
int q(int x) // achar o grundy number na marra
 if (x == 0) return 0;
 vector<bool> ok(4, 0);
 int mex = 0;
 for (auto const &i : v) {
   int curr = q(x / i);
   if (curr < 4) ok[curr] = 1;
    while (ok[mex]) mex++;
 vis[x] = 1;
 return dp[x] = mex;
```

int n;

cin >> n:

int x = 0;

int k; cin >> k;

for (int i = 0; i < n; i++) {

rectangleunion prefixsum2d bellmanfordinequalities

```
x ^= solve(k);
(x > 0) ? cout << "Henry\n" : cout << "Derek\n";
// nim classico -> o jogador que comeca ganha se o xor dos
    tamanhos das pilhas for != 0
// teorema sprague-grundy (transformar um jogo qualquer em nim)
rectangleunion.cpp
vector<int> x_vals;
struct segtree {
  vector<int> seq, taq;
  segtree() {
    seg.assign(8 * x_vals.size(), 0);
    tag.assign(8 * x vals.size(), 0);
  void add(int gl, int gr, int x, int v, int l, int r) {
    if (qr \le 1 | | r \le ql) {
      return;
    if (ql <= 1 && r <= qr) {
     tag[v] += x;
     if (tag[v] == 0) {
       if (1 != r)
          seg[v] = seg[v << 1] + seg[(v << 1) | 1];
          seq[v] = 0;
      } else {
        seq[v] = x_vals[r] - x_vals[l];
    } else {
     int mid = (1 + r) >> 1;
      add(ql, qr, x, (v << 1), 1, mid);
      add(q1, qr, x, ((v << 1) | 1), mid, r);
     if (tag[v] == 0 \&\& 1 != r) seg[v] = seg[v << 1] + seg[(v
          << 1) | 1];
  int qry() { return seg[1]; }
  void upd(int 1, int r, int x) { add(1, r, x, 1, 0, x_vals.
      size()); }
struct rect {
 int x1, y1, x2, y2;
struct event {
  int time, 1, r, type;
  bool operator<(const event &b) {
    if (time != b.time) return time < b.time;
    return type > b.type;
const int inf = 1e9;
signed main() {
  int n;
  cin >> n;
  vector<rect> v(n);
  for (int i = 0; i < n; i++) {
    cin >> v[i].x1 >> v[i].y1 >> v[i].x2 >> v[i].y2;
    x_{vals.pb}(v[i].x1);
    x_vals.pb(v[i].x2);
  // comprime o x
```

```
sort(x_vals.begin(), x_vals.end());
 x_vals.erase(unique(x_vals.beqin(), x_vals.end()), x_vals.end
 vector<event> ev;
 for (int i = 0; i < n; i++) {
   v[i].x1 = lower\_bound(x\_vals.begin(), x\_vals.end(), v[i].x1
        ) - x vals.begin();
   v[i].x2 = lower_bound(x_vals.begin(), x_vals.end(), v[i].x2
        ) - x_vals.begin();
   ev.pb({v[i].y1, v[i].x1, v[i].x2, 0}); // adicao
    ev.pb({v[i].y2, v[i].x1, v[i].x2, 1}); // remocao
 seatree s:
 sort(ev.begin(), ev.end());
 int area = 0, 1 = -inf;
 for (auto const &i : ev) {
   if (l == -inf) {
     l = i.time;
     s.upd(i.l, i.r, 1);
   } else if (i.type == 1) {
     int curr = s.gry();
     s.upd(i.l, i.r, -1);
     if (s.gry() != curr) {
       int new_t = (s.qry() == 0) ? -inf : i.time;
       int lo = 1, hi = i.time - 1;
       area += ((hi - lo + 1) * curr);
       1 = new_t;
   } else {
     int curr = s.gry();
     s.upd(i.1, i.r, 1);
     if (s.qry() != curr) {
       int lo = 1, hi = i.time - 1;
       area += ((hi - lo + 1) * curr);
       l = i.time:
 cout << area << endl;
 return 0:
// n <= 5 * 10^5
// 0 <= x, y <= 10^9
// comprime coordenada no x pra montar a segtree dos valores de
// faz o line sweep pelo y
prefixsum2d.cpp
                                                           9 lines
int v[1001][1001];
int p[1001][1001];
int qry(int x1, int y1, int x2, int y2) { return p[x2 + 1][y2 +
     1] - p[x2 + 1][y1] - p[x1][y2 + 1] + p[x1][y1]; }
for (int i = 0; i < n; i++) {
 for (int j = 0; j < n; j++) {
   p[i + 1][j + 1] = p[i][j + 1] + p[i + 1][j] - p[i][j];
   p[i + 1][j + 1] += v[i][j];
bellmanfordinequalities.cpp
struct solver {
 const int inf = 1e18;
 int n, src;
 vector<int> d;
 vector<pii> edges;
```

solver(int _n) // recebe o numero de variaveis, indexadas de

0 ate n - 1

```
src = _n; // aqui denotaremos _n como o source
    n = n + 1:
    for (int i = 0; i < _n; i++) // arestas de source para
        cada um dos vertices com custo 0
      edges.pb({{src, i}, 0});
  bool solve() // bellman ford
    d.assign(n, inf);
    d[src] = 0;
    int x;
    for (int i = 0; i < n; i++) {
     x = -1;
      for (auto const &e : edges) {
       auto [a, b] = e.fir;
       int cost = e.sec;
       if (d[a] < inf) {
          if (d[b] > d[a] + cost) {
            d[b] = max(-inf, d[a] + cost);
            x = b;
    return (x == -1); // false se tem ciclo negativo
 void add_constraint_leg(int i, int j, int c) // value_i -
      value_j <= c
    edges.pb({{j, i}, c});
  void add_constraint_geq(int i, int j, int c) // value_i -
      value_j >= c
    edges.pb({{i, j}, -c});
  void add_constraint_eq(int i, int j, int c) // value_i -
      value_j = c
    add_constraint_leg(i, j, c);
    add_constraint_geq(i, j, c);
 };
};
```