

Universidade Federal de Alagoas

Floyd (WA)rshall

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```
3 Data structures
                                                             6
4 Graph
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5 Geometry
                                                            16
6 Strings
                                                            20
7 Various
                                                            22
Contest (1)
template.cpp
                                                         25 lines
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
#define int long long int
#define pb push back
#define pi pair<int, int>
#define pii pair<pi, int>
#define fir first
#define sec second
#define MAXN 1000006
#define mod 998244353
signed main() {
  ios_base::sync_with_stdio(false);
  mt19937 64 rng(chrono::steady clock::now().time since epoch()
 return 0;
runner.py
                                                         44 lines
import os
import subprocess
naive = "brute.cpp" # path to naive code
code = "d.cpp" # path to your code
generator = "g.cpp" # path to test generator
def compile_codes():
   os.system('q++ ' + generator + ' -o generator -O2')
    os.system('g++ ' + naive + ' -o naive -02')
   os.system('g++ ' + code + ' -o code -02')
def generate_case():
    os.system('./generator > in');
def get_naive_output():
```

1

1

1 Contest

2 Mathematics

```
modint.cpp
    output = os.popen('./naive <in').read()</pre>
    return output
def get_code_output():
    # se tiver um runtime error, vai parar
    xalala = subprocess.run('./code <in', shell=True, text=True</pre>
         , capture_output=True, check=True)
    return xalala.stdout
def main():
    compile_codes()
    while True:
        generate_case()
        naive_output = get_naive_output()
        code_output = get_code_output()
        if naive_output == code_output:
            print('ACCEPTED')
        else:
            print('FAILED\n')
            print ('ANSWER:')
            print (naive output)
            print('\nCODE OUTPUT:')
            print(code_output)
            break
if __name__ == '__main__':
    main()
int128.cpp
                                                           22 lines
__int128 read() {
  __int128 x = 0, f = 1;
  char ch = getchar();
  while (ch < '0' || ch > '9') {
    if (ch == '-') f = -1;
    ch = getchar();
  while (ch >= '0' && ch <= '9') {
    x = x * 10 + ch - '0';
    ch = getchar();
  return x * f;
void print(__int128 x) {
 if (x < 0) {
    cout << "-";
    x = -x;
 if (x > 9) print (x / 10);
  char at = (x % 10) + '0';
  cout << at;
pragma.cpp
                                                            5 lines
#include <iostream>
using namespace std;
#pragma GCC target("avx2")
#pragma GCC optimize("03")
#pragma GCC optimize("unroll-loops")
Mathematics (2)
```

```
78 lines
struct modint {
 int val;
  modint(int v = 0) { val = v % mod; }
  int pow(int y) {
    modint x = val;
    modint z = 1;
    while (y) {
     if (y \& 1) z *= x;
     x *= x;
      y >>= 1;
    return z.val;
 int inv() { return pow(mod - 2); }
  void operator=(int o) { val = 0 % mod; }
  void operator=(modint o) { val = o.val % mod; }
  void operator+=(modint o) { *this = *this + o; }
  void operator = (modint o) { *this = *this - o; }
  void operator*=(modint o) { *this = *this * o; }
  void operator/=(modint o) { *this = *this / o; }
  bool operator==(modint o) { return val == o.val; }
 bool operator!=(modint o) { return val != o.val; }
 int operator*(modint o) { return ((val * o.val) % mod); }
 int operator/(modint o) { return (val * o.inv()) % mod; }
 int operator+(modint o) { return (val + o.val) % mod; }
 int operator-(modint o) { return (val - o.val + mod) % mod; }
modint f[MAXN];
modint inv[MAXN];
modint invfat[MAXN];
void calc() {
 f[0] = 1;
 for (int i = 1; i < MAXN; i++) {
   f[i] = f[i - 1] * i;
 inv[1] = 1;
 for (int i = 2; i < MAXN; ++i) {
   int val = mod / i;
   val = (inv[mod % i] * val) % mod;
   val = mod - val;
   inv[i] = val;
 invfat[0] = 1;
 invfat[MAXN - 1] = modint(f[MAXN - 1]).inv();
 for (int i = MAXN - 2; i >= 1; i--) {
    invfat[i] = invfat[i + 1] * (i + 1);
modint ncr(int n, int k) // combinacao
 modint ans = f[n] * invfat[k];
 ans *= invfat[n - k];
 return ans;
modint arr(int n, int k) // arranjo
 modint ans = f[n] * invfat[n - k];
 return ans;
modint ncr(int n, int k) {
 // calcular combinacao para n grande
 // nesse problema n <= 10^12
 // em O(k)
 modint num = 1;
 modint den = 1;
 for (int i = 0; i < k; i++) {
    num = num * modint(n - i);
```

```
den = den * modint(i + 1);
} modint ans = num / den;
return ans;
}
modint stars_and_bars(int n, int k) {
   // para pares de inteiros n e k
   // enconte a quantidade de k-tuplas com soma == n
   // x1 + x2 + ... + xk = n
   return ncr(n + k - 1, k - 1);
}
```

divisiontrick.cpp

crivo.cpp

itset<MANN prime:

```
bitset<MAXN> prime:
vector<int> nxt(MAXN);
vector<int> factors:
void crivo() {
 prime.set();
 prime[0] = false, prime[1] = false;
  for (int i = 2; i < MAXN; i++) {
   if (prime[i]) {
     nxt[i] = i;
     for (int j = 2; j * i < MAXN; j++) {
       prime[j * i] = false;
       nxt[j * i] = i;
void fact(int n) {
 factors.clear();
  while (n > 1) {
   factors.pb(nxt[n]);
   n = n / nxt[n];
```

gaussianelimination.cpp

fine EPS 1e-9

```
#define EPS 1e-9
vector<double> ans:
int gauss(vector<vector<double>> a) {
 int n = a.size(), m = a[0].size() - 1, ret = 1;
 ans.assign(m, 0);
 vector<int> where (m, -1);
  for (int col = 0, row = 0; col < m && row < n; col++, row++)
     {
   int sel = row;
   for (int i = row; i < n; i++)
     if (abs(a[i][col]) > abs(a[sel][col])) sel = i;
   if (abs(a[sel][col]) < EPS) continue;
   for (int i = col; i \le m; i++) swap(a[sel][i], a[row][i]);
   where[col] = row;
   for (int i = 0; i < n; i++) {
     if (i != row) {
       double c = a[i][col] / a[row][col];
       for (int j = col; j \le m; j++) a[i][j] -= a[row][j] * c
```

```
}
}
for (int i = 0; i < m; i++) {
    if (where[i] != -1)
        ans[i] = (a[where[i]][m] / a[where[i]][i]);
    else
        ret = 2;
}
for (int i = 0; i < n; i++) {
    double sum = 0;
    for (int j = 0; j < m; j++) sum += (ans[j] * a[i][j]);
    if (abs(sum - a[i][m]) > EPS) ret = 0;
}
return ret; // 0 = nao existe solucao, 1 = existe uma
        solucao, 2 = existem multiplas solucoes
```

gaussianbitset.cpp

5 lines

bitset<MAXN> ans;

```
int gauss (vector < bitset < MAXN >> & a) {
 ans.reset();
 int n = a.size(), m = a[0].size() - 1, ret = 1;
 vector<int> where (m, -1);
 for (int col = 0, row = 0; col < m && row < n; col++) {
   for (int i = row; i < n; i++) {
     if (a[i][col]) {
       swap(a[i], a[row]);
       break;
    if (!a[row][col]) continue;
    where[col] = row;
    for (int i = 0; i < n; i++)
     if (i != row && a[i][col]) a[i] ^= a[row];
 for (int i = 0; i < m; i++) {
   if (where[i] != -1)
     ans[i] = (a[where[i]][m] / a[where[i]][i]);
   else
     ret = 2;
 for (int i = 0; i < n; i++) {
   double sum = 0;
   for (int j = 0; j < m; j++) sum += (ans[j] * a[i][j]);
```

if (abs(sum - a[i][m]) > EPS) ret = 0;

lagrange.cpp

den.resize(n);

return ret:

33 lines

```
struct lagrange {
  int n;
  vector<modint> den;
  vector<modint> fat;
  vector<modint> inv_fat;
  lagrange(vector<modint> &v) // f(i) = v[i], gera um
      polinomio de grau n - 1
  {
      n = v.size();
      calc(n);
      calc_den(n);
      y = v;
  }
  void calc_den(int n) {
```

```
for (int i = 0; i < n; i++) {
     den[i] = inv_fat[n - i - 1] * inv_fat[i];
     if ((n - i - 1) % 2 == 1) {
       int x = (mod - den[i].val) % mod;
       den[i] = x;
 void calc(int n) {
   fat.resize(n + 1):
    inv_fat.resize(n + 1);
    fat[0] = 1;
    inv_fat[0] = 1;
    for (int i = 1; i \le n; i++) {
     fat[i] = fat[i - 1] * i;
     inv_fat[i] = fat[i].inv();
 modint get_val(int x) // complexidade: O(n)
   x %= mod:
   vector<modint> 1(n);
    vector<modint> r(n);
   1[0] = 1, r[n - 1] = 1;
    for (int i = 1; i < n; i++) {
     modint cof = (x - (i - 1) + mod);
     l[i] = l[i - 1] * cof;
    for (int i = n - 2; i >= 0; i--) {
     modint cof = (x - (i + 1) + mod);
     r[i] = r[i + 1] * cof;
   modint ans = 0;
    for (int i = 0; i < n; i++) {
     modint cof = l[i] * r[i];
     ans += modint(cof * y[i]) * den[i];
   return ans;
 vector<modint> find_coefs() // encontra os coeficientes do
      polinomio
   int nn = n;
   int d = nn - 1;
    vector<modint> c(nn, 0);
    for (int i = 0; i < y.size(); i++) {
     c[d] += (v[i] * den[i]);
    for (int p = nn - 2; p >= 0; p--) {
     nn--:
     calc_den(nn);
     for (int i = 0; i \le p; i++) {
       y[i] = (c[p + 1] * modint(i).pow(d));
       c[p] += (y[i] * den[i]);
     d--;
   return c;
};
```

berlekampmassey.cpp

31 lines

74 lines

73 line

```
// berlekamp massey
// mas precisa que o mod seja primo (para poder achar inverso)
// dado os n primeiros termos de uma recorrencia linear
// a[0], a[1], a[2], ..., a[n - 1]
// ele acha a reccorrencia linear mais curta que da matching
com os n primeiros valores
```

```
vector<modint> berlekamp_massey(vector<modint> x) {
  vector<modint> ls, cur;
  int lf, ld;
  for (int i = 0; i < x.size(); i++) {
   modint t = 0;
   for (int j = 0; j < cur.size(); j++) {
     t += (x[i - j - 1] * cur[j]);
    if (modint(t - x[i]).val == 0) continue;
    if (cur.empty()) {
     cur.resize(i + 1);
     lf = i;
     ld = (t - x[i]) % mod;
     continue:
    modint k = -(x[i] - t);
   k *= modint(ld).inv();
    vector<modint> c(i - lf - 1);
    for (auto const &j : ls) {
     modint curr = modint(j.val * -1) * k;
     c.pb(curr);
    if (c.size() < cur.size()) c.resize(cur.size());
    for (int j = 0; j < cur.size(); j++) {
     c[j] = c[j] + cur[j];
    if (i - lf + ls.size() >= cur.size()) {
     tie(ls, lf, ld) = make_tuple(cur, i, t - x[i]);
   cur = c;
  return cur;
modint get_nth(vector<modint> rec, vector<modint> dp, int n) {
 int m = rec.size();
  vector<modint> s(m), t(m);
  s[0] = 1;
  if (m != 1)
   t[1] = 1;
   t[0] = rec[0];
  auto mul = [&rec] (vector<modint> v, vector<modint> w) {
   vector<modint> ans(2 * v.size());
    for (int j = 0; j < v.size(); j++) {
     for (int k = 0; k < v.size(); k++) ans[j + k] += v[j] * w
           [k];
    for (int j = 2 * v.size() - 1; j >= v.size(); j--) {
     for (int k = 1; k \le v.size(); k++) ans[j - k] += ans[j]
          * rec[k - 1];
   ans.resize(v.size());
   return ans:
  while (n) {
   if (n \& 1) s = mul(s, t);
   t = mul(t, t);
   n >>= 1;
  modint ret = 0:
  for (int i = 0; i < m; i++) ret += s[i] * dp[i];
  return ret;
modint guess_nth_term(vector<modint> x, int n) {
 if (n < x.size()) return x[n];
  vector<modint> coef = berlekamp_massey(x); // coeficientes
      da recorrencia
 if (coef.empty()) return 0;
```

```
return get_nth(coef, x, n);
crt.cpp
namespace crt {
vector<pi> eq;
int gcd(int a, int b, int &x, int &y) {
 if (b == 0) {
   x = 1, y = 0;
   return a;
 int x1, y1, d = gcd(b, a % b, x1, y1);
  x = y1, y = x1 - y1 * (a / b);
  return d;
pi crt() {
 int a1 = eq[0].fir, m1 = eq[0].sec;
  a1 %= m1;
  for (int i = 1; i < eq.size(); i++) {
    int a2 = eq[i].fir, m2 = eq[i].sec;
    int q = \underline{\hspace{0.5cm}} \gcd(m1, m2);
    if (a1 % g != a2 % g) return {-1, -1};
    gcd(m1 / g, m2 / g, p, q);
    int mod = m1 / q * m2;
    int x = (a1 * (m2 / g) % mod * q % mod + a2 * (m1 / g) %
         mod * p % mod) % mod;
    if (a1 < 0) a1 += mod;
    m1 = mod;
  return {a1, m1};
} // namespace crt
// o menor inteiro a que satisfaz:
// a mod p1 = x1
// a mod p2 = x2
// a funcao crt retorna um pair {a, mod}
// dai a solucao pode ser descrita como
// x = a \% mod
// entao os valores possiveis sao:
// a, (a + mod), a + (2 * mod), a + (3 * mod), ...
// cuidado com overflow!
crttrick.cpp
                                                           19 lines
vector<pi> ea:
map<int, int> by_mod;
// quero checar se existe solucao para o sistema das equacoes
    que ja adicionei
// junto da equacao curr
// geralmente da pra fazer algo como if (check(curr)) { add(
    curr); }
bool check(pi curr) {
 if (by_mod.find(curr.sec) != by_mod.end()) {
    return by_mod[curr.sec] == curr.fir;
  for (auto [x, mod] : eq) {
    if ((curr.fir - x) % __gcd(curr.sec, mod)) return 0;
  return 1:
void add(pi curr) { // [valor, mod]
  eq.pb(curr);
  by_mod[curr.sec] = curr.fir;
// quando tem algo de sqrt mods distintos ou algo do tipo
```

```
diophantine.cpp
                                                           32 lines
namespace dio {
vector<pi> sols;
int gcd(int a, int b, int &x, int &y) {
if (b == 0) {
    x = 1, v = 0;
    return a:
 int x1, y1, d = gcd(b, a % b, x1, y1);
 x = y1, y = x1 - y1 * (a / b);
 return d;
void one_sol(int a, int b, int c) {
 int x0, y0, q;
 g = gcd(abs(a), abs(b), x0, y0);
 if (c % g) return;
 x0 *= (c / q);
 v0 *= (c / a);
 if (a < 0) \times 0 *= -1;
 if (b < 0) y0 \star = -1;
 sols.pb({x0, y0});
void more sols(int a, int b, int c) {
 int g = \underline{gcd(a, b)};
 int x0 = sols[0].fir, y0 = sols[0].sec;
 for (int k = -200000; k \le 200000; k++) {
   int x = x0 + k * (b / q);
    int y = y0 - k * (a / q);
    sols.pb({x, y});
} // namespace dio
// equacoes do tipo: ax + by = c
extendedeuclidean.cpp
int gcd(int a, int b, int &x, int &y) {
 if (b == 0) {
   x = 1:
    y = 0;
    return a:
 int x1, y1;
 int d = gcd(b, a % b, x1, y1);
 x = y1;
 y = x1 - y1 * (a / b);
 return d;
// achar os numeros x e y tal que:
// a * x + b * y = gcd(a, b)
ntt.cpp
const int MOD = 998244353;
typedef mod int<MOD> mint;
void ntt(vector<mint>& a, bool rev) {
 int n = a.size();
 auto b = a;
  assert(!(n & (n - 1)));
 mint q = 1;
  while ((g ^ (MOD / 2)) == 1) g += 1;
 if (rev) q = 1 / q;
  for (int step = n / 2; step; step /= 2) {
    mint w = q ^ (MOD / (n / step)), wn = 1;
    for (int i = 0; i < n / 2; i += step) {
      for (int j = 0; j < step; j++) {
```

auto u = a[2 * i + j], v = wn * a[2 * i + j + step];

b[i + j] = u + v;

b[i + n / 2 + j] = u - v;

```
wn = wn * w:
    swap(a, b);
  if (rev) {
   auto n1 = mint(1) / n;
   for (auto& x : a) x *= n1;
vector<mint> convolution(const vector<mint>& a, const vector<
    mint>& b) {
  vector<mint> l(a.begin(), a.end()), r(b.begin(), b.end());
 int N = 1.size() + r.size() - 1, n = 1 << lq(2 * N - 1);
 l.resize(n);
  r.resize(n):
  ntt(1, false);
  ntt(r, false);
  for (int i = 0; i < n; i++) l[i] *= r[i];
  ntt(l, true);
 l.resize(N);
 return 1;
```

fft.cpp

```
43 lines
#define PI acos(-1)
#define cd complex<double>
namespace fft {
int n;
void fft (vector<cd> &a, bool invert) {
  int n = a.size();
  for (int i = 1, j = 0; i < n; i++) {
    int bit = n \gg 1;
    for (; j & bit; bit >>= 1) j ^= bit;
    j ^= bit;
    if (i < j) swap(a[i], a[i]);</pre>
  for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 : 1);
    cd wlen(cos(ang), sin(ang));
    for (int i = 0; i < n; i += len)
     cd w(1);
      for (int i = 0; i < len / 2; i++) {
       cd u = a[i + j], v = a[i + j + len / 2] * w;
       a[i + j] = u + v;
       a[i + j + len / 2] = u - v;
       w \star = wlen;
  if (invert)
    for (cd &x : a) x /= n;
vector<int> mul(vector<int> a, vector<int> b) {
  vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());
  n = 1:
  while (n < a.size() + b.size()) n <<= 1;
  fa.resize(n);
  fb.resize(n):
  fft(fa, false);
  fft(fb, false);
  for (int i = 0; i < n; i++) fa[i] *= fb[i];
  fft(fa, true);
  vector<int> ans(n);
  for (int i = 0; i < n; i++) ans[i] = round(fa[i].real());
  return ans;
  // namespace fft
```

```
othermodint.cpp
```

```
template <int p>
struct mod int {
 ll expo(ll b, ll e) {
   ll ret = 1;
   while (e) {
     if (e % 2) ret = ret * b % p;
     e /= 2, b = b * b % p;
   return ret;
 11 inv(11 b) { return expo(b, p - 2); }
 using m = mod_int;
 mod_int() : v(0) {}
 mod int(ll v ) {
   if (v_ >= p or v_ <= -p) v_ %= p;
   if (v < 0) v += p;
   v = v;
 m& operator+=(const m& a) {
   v += a.v;
   if (v >= p) v -= p;
   return *this;
 m& operator -= (const m& a) {
   v -= a.v;
   if (v < 0) v += p;
   return *this;
 m& operator *= (const m& a) {
   v = v * ll(a.v) % p;
   return *this:
 m& operator/=(const m& a) {
   v = v * inv(a.v) % p;
   return *this;
 m operator-() { return m(-v); }
 m& operator^=(11 e) {
   if (e < 0) {
     v = inv(v);
     e = -e;
   v = expo(v, e);
   // possivel otimizacao:
   // cuidado com 0^0
   // v = \exp(v, e^{(p-1)});
   return *this;
 bool operator==(const m& a) { return v == a.v; }
 bool operator!=(const m& a) { return v != a.v; }
 friend istream& operator>>(istream& in, m& a) {
   ll val;
   in >> val:
   a = m(val);
   return in;
 friend ostream& operator << (ostream& out, m a) { return out <<
 friend m operator+(m a, m b) { return a += b; }
 friend m operator-(m a, m b) { return a -= b; }
 friend m operator* (m a, m b) { return a *= b; }
 friend m operator/(m a, m b) { return a /= b; }
 friend m operator (m a, ll e) { return a '= e; }
```

```
fraction.cpp
```

44 lines

```
struct fraction {
 int x, y; // x / y
 fraction() {}
 fraction(int x, int y) : x(x), y(y) {}
 bool operator==(fraction o) { return (x * o.y == o.x * y); }
 bool operator!=(fraction o) { return (x * o.y != o.x * y); }
 bool operator>(fraction o) { return (x * o.y > o.x * y); }
 bool operator>=(fraction o) { return (x * o.y >= o.x * y); }
 bool operator<(fraction o) { return (x * o.y < o.x * y); }</pre>
 bool operator<=(fraction o) { return (x * o.v \le o.x * v); }
 fraction operator+(fraction o) {
   fraction ans;
   ans.y = (y == o.y) ? y : y * o.y;
   ans.x = (x) * (ans.y / y) + (o.x) * (ans.y / o.y);
   // ans.simplify();
   return ans:
 fraction operator*(fraction o) {
   fraction ans;
   ans.x = x * o.x;
   ans.v = v * o.v;
   // ans.simplifv();
   return ans;
 fraction inv() {
   fraction ans = fraction(x, y);
   swap(ans.x, ans.y);
   return ans;
 fraction neg() {
   fraction ans = fraction(x, y);
   ans.x \star = -1;
   return ans:
 void simplify() {
   if (abs(x) > 1e9 \mid \mid abs(y) > 1e9) // slow simplification
     int g = \underline{gcd}(y, x);
     x /= g;
     y /= g;
 // substraction and division can be easily done
};
```

totient.cpp

```
int phi[MAXN];
void calc() {
 for (int i = 0; i < MAXN; i++) phi[i] = i;
 for (int i = 2; i < MAXN; i++) {
   if (phi[i] == i) {
      for (int j = i; j < MAXN; j += i) phi[j] -= phi[j] / i;
int calc_phi(int n) {
 int ans = n;
 for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) {
     while (n % i == 0) n /= i;
     ans -= ans / i;
 if (n > 1) ans -= ans / n;
 return ans;
```

pollardrho.cpp #define int int128 namespace pollard rho { int multiplicate(int x, int y, int m) { return $(x * y) % m; }$ int modpow(int x, int y, int m) { while (y) { if (y & 1) z = (z * x) % m;x = (x * x) % m;y >>= 1; return z; bool is composite(int n, int a, int d, int s) { int x = modpow(a, d, n);if (x == 1 or x == n - 1) return false; for (int r = 1; r < s; r++) { x = multiplicate(x, x, n);if (x == n - 1LL) return false; return true; int miller rabin(int n) { if (n < 2) return false; int r = 0, d = n - 1LL; while $((d \& 1LL) == 0) {$ d >>= 1;r++; for (int a: {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}) { if (n == a) return true; if (is_composite(n, a, d, r)) return false; return true; int f(int x, int m) { return multiplicate(x, x, m) + 1; } int rho(int n) { int x0 = 1, t = 0, prd = 2; int x = 0, y = 0, q; while (t $\% \ 40 \ || \ \underline{\ \ \ } gcd(prd, n) == 1) {$ if (x == y) { x0++; x = x0;y = f(x, n);q = multiplicate(prd, max(x, y) - min(x, y), n);if (q != 0) prd = q;x = f(x, n);y = f(y, n);y = f(y, n);t++; return __gcd(prd, n); vector<int> fact(int n) { if (n == 1) return {}; if (miller_rabin(n)) return {n}; int x = rho(n); auto l = fact(x), r = fact(n / x); l.insert(l.end(), r.begin(), r.end()); return 1; } // namespace pollard_rho primefactors.cpp 14 lines

vector<int> facts; void primefactors(int n) { while $(n % 2 == 0) {$

```
facts.pb(2);
  n = n / 2;
for (int i = 3; i \le sqrt(n); i += 2) {
  while (n % i == 0) {
   facts.pb(i);
   n = n / i;
if (n > 2) facts.pb(n);
```

```
xorbasis.cpp
int modpow(int a, int b) {
 int res = 1;
 while (b > 0) {
   if (b & 1) res = (res * a) % mod;
   a = (a * a) % mod;
 return res;
int all, qt;
int dp[33];
void add(int x) {
 for (int i = 32; i >= 0; i--) {
    if (x & (111 << i)) {
      if (dp[i] == 0) {
        dp[i] = x;
        qt++;
        return;
      x ^= dp[i];
int get(int x) // gual o x-esimo menor valor de xor de uma
    subsequencia
 int tot = (111 \ll gt), ans = 0;
 for (int i = 32; i >= 0; i--) {
   if (dp[i] > 0) {
     int d = tot / 2;
     if (d < x && !(ans & (111 << i)))
       ans ^= dp[i];
      else if (d >= x && (ans & (111 << i)))
       ans ^= dp[i];
     if (d < x) x = d;
     tot /= 2;
 return ans:
bool check(int x) // se existe pelo menos uma subsequencia com
 for (int i = 32; i >= 0; i--) {
   if (x & (111 << i)) {
     if (!dp[i]) return 0;
     x ^= dp[i];
 return 1;
int count(int x) // quantas subsequencias tem xor x
```

if (!check(x)) return 0;

```
return modpow(2, all - gt);
int x = get(111 \ll gt); // maior xor possivel de uma
    subsequencia
int y = get(1);
                         // maior xor possivel != 0 (o 0 sempre
     eh possivel - subsequencia vazia)
```

```
mobius.cpp
int lpf[MAXN];
int mobius[MAXN];
int mp[MAXN];
vector<int> d[MAXN];
void calc_lpf() {
 for (int i = 2; i < MAXN; i++) {
    if (!lpf[i]) {
     for (int j = i; j < MAXN; j += i) {
       if (!lpf[j]) lpf[j] = i;
 for (int i = 2; i < MAXN; i++) // divisores
    for (int j = i; j < MAXN; j += i) d[j].pb(i);
 calc_lpf();
  mobius[1] = 1;
  for (int i = 2; i < MAXN; i++) {
   if (lpf[i / lpf[i]] == lpf[i])
     mobius[i] = 0;
      mobius[i] = -1 * mobius[i / lpf[i]];
void add(int x, int dd) // adiciona dd em todos os val[i] que
    gcd(x, i) > 1
 for (auto const &i : d[x]) mp[i] += dd;
int sum(int x) // valor de val[x]
 int ans = 0;
 for (auto const &i : d[x]) ans += (mobius[i] * -1 * mp[i]);
// mobius/inclusao-exclusao com os fatores primos
// a funcao de mobius eh definida como:
// mi(n) = 1, se n eh um square-free com um numero par de
    fatores primos
// mi(n) = -1, se n eh um square-free com um numero impar de
    fatores primos
// mi(n) = 0, caso nenhum dos dois
// square-free = nenhum fator primo aparece duas vezes ou mais
```

segmentedsieve.cpp

```
vector<int> prime;
void segmentedsieve(int 1, int r) {
 int lim = sqrt(r);
 vector<bool> mark(lim + 1, false);
 vector<int> primes;
 for (int i = 2; i \le \lim_{t \to 0} t + i) {
    if (!mark[i]) {
      primes.pb(i);
      for (int j = i * i; j \le lim; j += i) mark[j] = true;
```

```
UFAL
  vector<bool> isprime(r - 1 + 1, true);
  for (int i : primes)
   for (int j = \max(i * i, (1 + i - 1) / i * i); j <= r; j +=
        i) isprime[j - 1] = false;
  if (1 == 1) isprime[0] = false;
  for (int i = 0; i < isprime.size(); i++)
    if (isprime[i]) prime.pb(i + 1);
xortrie.cpp
                                                           51 lines
struct node {
  int me, cnt, id;
  int down[2];
  node(int c = 0) : me(c) {
   cnt = 0;
   id = -1;
    fill(begin(down), end(down), -1);
};
struct trie_xor {
  vector<node> t:
  trie xor() { t.resize(1); }
  void add(int n, int id) {
   int v = 0;
    t[v].cnt++;
    for (int i = 30; i >= 0; i--) {
     int bit = (n & (1 << i)) ? 1 : 0;
     if (t[v].down[bit] == -1) {
       t[v].down[bit] = t.size();
       t.emplace back(bit);
     v = t[v].down[bit];
     t[v].cnt++;
   t[v].id = id;
```

```
teoremas.md
```

};

void rem(int n, int id) {

v = t[v].down[bit];

for (int i = 30; i >= 0; i--) {

int bit = (n & (1 << i)) ? 1 : 0;

int gry(int n) // maximum xor with n

if (t[0].cnt == 0) // no element

for (int i = 30; i >= 0; i--) {

int nxt = t[v].down[bit];

 $v = t[v].down[bit ^ 1];$

int bit = (n & (1 << i)) ? 0 : 1;

if (nxt != -1 && t[nxt].cnt > 0)

int v = 0;

t[v].cnt--;

t[v].cnt--;

return -1;

v = nxt:

return t[v].id;

int v = 0:

Teorema de Lucas

66 lines

Theorem

sejam m e n numeros inteiros nao negativos e p um numero primo desenvolver n e m na base p

```
m = m[k] *p^{(k)} + m[k - 1] *p^{(k - 1)} + ... + m[0] *p^{(0)}
n = n[k] *p^{(k)} + n[k - 1] *p^{(k - 1)} + ... + n[0] *p^{(0)}
ncr(m, n) \mod p = produtorio de (ncr(m[i], n[i]) \mod p)
dai pra generalizar pro mod 2 eh deboas, pq se tiver um bit
     setado em n[i] que nao ta setado em m[i], entao miou, vai
     dar zero
# Manhattan and Chebyshev distances equivalences
It is well known that given points (x, y) and you need to
    calculate the Manhattan distances between them, instead of
     using:
|x1-x2|+|y1-y2|
you can first convert all points (x, y) into (x+y, x-y) (rotate
     45 degrees) and the distances will become max(|x1-x2|, |
    y1-y2|) (also known as Chebyshev distance).
# Chicken McNugget Theorem
For any two coprime numbers (n > 0, m > 0), the greatest
    integer that cannot be written in the form:
an + bm, (a >= 0, b >= 0)
is (n \times m) - n - m
## Consequence of the theorem
That there are exactly ((n-1) \times (m-1)) / 2 positive
    integers which cannot be expressed in the form an + bm, (a
     >= 0, b >= 0)
## Generalization
If n and m are not coprime, so all numbers that are not
    multiples of gcd(n, m) cannot be expressed in the form an
    + bm, (a >= 0, b >= 0)
in addition, you can consider n = (n / gcd(n, m)) and m = (m / gcd(n, m))
    gcd(n, m)), to find how many multiples of gcd(n, m) cannot
     be expressed, or to find the greatest multiple of gcd(n,
    m) that cannot be expressed
## Considering a > 0, b > 0
Considering (n > 0, m > 0), n and m are coprime:
let v = ((n \times m) + min(n, m)) - 1
The number of positive integers which cannot be expressed
    increases by (v / n)
The number of positive integers which cannot be expressed
    increases by (y / m)
you must not count the multiples of (n \times m) more than once,
    just decrease number of positive integers which cannot be
     expressed by (v / (n \times m))
# Binomial Theorem
```

Data structures (3)

```
segtree.cpp
```

25 lines

```
struct seatree {
 int n;
  vector<int> seq;
  int neutral() { return 0; }
  int merge(int a, int b) { return a + b; }
  void build(vector<int> &v) {
    n = 1:
    while (n < v.size()) n <<= 1;
    seq.assign(n << 1, neutral());</pre>
    for (int i = 0; i < v.size(); i++) seg[i + n] = v[i];
    for (int i = n - 1; i; i--) seg[i] = merge(seg[i <math>\lt \lt 1], seg
         [(i << 1) | 11);
 void upd(int i, int value) {
    seg[i += n] += value;
    for (i >>= 1; i; i >>= 1) seg[i] = merge(seg[i << 1], seg[(
        i << 1) | 11);
 int qry(int 1, int r) {
    int ansl = neutral(), ansr = neutral();
    for (1 += n, r += n + 1; 1 < r; 1 >>= 1, r >>= 1) {
      if (1 & 1) ansl = merge(ansl, seg[l++]);
     if (r \& 1) ansr = merge(seg[--r], ansr);
    return merge (ansl, ansr);
};
```

segtreelazy.cpp

53 li

```
struct segtree {
 int n:
 vector<int> v:
 vector<int> seq;
 vector<int> lazy;
 segtree(int sz) {
   n = sz;
    seg.assign(4 * n, 0);
   lazy.assign(4 * n, 0);
 int single(int x) { return x; }
 int neutral() { return 0; }
 int merge(int a, int b) { return a + b; }
 void add(int i, int l, int r, int diff) {
   seg[i] += (r - 1 + 1) * diff;
   if (1 != r) {
     lazy[i << 1] += diff;
     lazy[(i << 1) | 1] += diff;
   lazy[i] = 0;
```

```
void update(int i, int l, int r, int ql, int qr, int diff) {
  if (lazy[i]) add(i, l, r, lazy[i]);
 if (1 > r || 1 > qr || r < ql) return;
 if (1 >= ql \&\& r <= qr) {
   add(i, l, r, diff);
   return:
  int mid = (1 + r) >> 1;
  update(i << 1, 1, mid, q1, qr, diff);
 update((i << 1) | 1, mid + 1, r, ql, qr, diff);
  seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
int query(int 1, int r, int q1, int qr, int i) {
 if (lazy[i]) add(i, 1, r, lazy[i]);
 if (1 > r \mid \mid 1 > qr \mid \mid r < ql) return neutral();
 if (1 >= ql \&\& r <= qr) return seg[i];
 int mid = (1 + r) >> 1;
  return merge(query(1, mid, ql, qr, i << 1), query(mid + 1,
      r, ql, qr, (i << 1) | 1));
void build(int 1, int r, int i) {
 if (1 == r) {
   seq[i] = single(v[l]);
   return:
  int mid = (1 + r) >> 1;
 build(1, mid, i << 1);
 build(mid + 1, r, (i << 1) | 1);
  seg[i] = merge(seg[i << 1], seg[(i << 1) | 1]);
int qry(int 1, int r) \{ return query(0, n - 1, 1, r, 1); \}
void upd(int 1, int r, int x) { update(1, 0, n - 1, 1, r, x);
```

fenwick.cpp

```
33 lines
struct fenw {
  int n;
  vector<int> bit:
  fenw() {}
  fenw(int sz) {
   n = sz:
   bit.assign(sz + 1, 0);
  int qry(int r) // query de prefixo a[0] + a[1] + ... a[r]
   for (int i = r + 1; i > 0; i -= i \& -i) ret += bit[i];
   return ret;
  void upd(int r, int x) // a[r] += x
    for (int i = r + 1; i \le n; i += i \& -i) bit[i] += x;
  int bs(int x) // retorna o maior indice i (i < n) tal que:
      qry(i) < x
    int i = 0, k = 0;
   while (1 << (k + 1) <= n) k++;
    while (k \ge 0) {
     int nxt_i = i + (1 << k);
     if (nxt_i <= n && bit[nxt_i] < x) {
       i = nxt_i;
       x \rightarrow bit[i];
     k--;
    return i - 1;
```

```
};
treap.cpp
                                                           93 lines
struct treap {
 int data, priority;
 int sz, lazy2;
 bool lazv;
 treap *1, *r, *parent;
int size(treap *node) { return (!node) ? 0 : node->sz; }
void recalc(treap *node) {
 if (!node) return;
 node -> sz = 1;
  node->parent = 0;
  if (node->1) node->sz += node->1->sz, node->1->parent = node;
 if (node->r) node->sz += node->r->sz, node->r->parent = node;
void lazy_propagation(treap *node) {
 if (node == NULL) return;
 if (node->lazy2) {
    if (node->1) node->1->lazy2 += node->lazy2;
    if (node->r) node->r->lazy2 += node->lazy2;
    node->data += node->lazy2;
    node \rightarrow lazv2 = 0;
 if (node->lazy) {
    swap(node->1, node->r);
    if (node->1) node->1->lazv = !node->1->lazv;
    if (node->r) node->r->lazy = !node->r->lazy;
    node \rightarrow lazy = 0;
void split(treap *t, treap *&l, treap *&r, int n) {
 if (!t) return void(l = r = 0);
 lazy propagation(t);
 if (size(t->1) >= n)
    split(t->1, 1, t->1, n), r = t;
    split(t->r, t->r, r, n - size(t->1) - 1), 1 = t;
 recalc(t);
void merge(treap *&t, treap *l, treap *r) {
  lazy propagation(1);
 lazy_propagation(r);
 if (!1)
   t = r;
 else if (!r)
   t = 1;
  else if (l->priority > r->priority)
   merge(1->r, 1->r, r), t=1;
    merge(r->1, 1, r->1), t = r;
 recalc(t);
void troca(treap *&t, int 1, int r, int 11, int rr) // sap de
 treap *a0, *a1, *b0, *b1, *c0, *c1, *d0, *d1;
  split(t, a0, a1, 1);
 split(a1, b0, b1, r - 1 + 1);
 11 -= (r + 1);
 rr -= (r + 1);
  split(b1, c0, c1, l1);
  split(c1, d0, d1, rr - 11 + 1);
 merge(t, a0, d0);
  merge(t, t, c0);
 merge(t, t, b0);
```

```
merge(t, t, d1);
void add(treap *&t, int l, int r) {
 treap *a0, *a1, *b0, *b1;
 split(t, a0, a1, 1);
  split(a1, b0, b1, r - 1 + 1);
 b0->lazy ^= 1;
 b0 -> lazy2 += 1;
 merge(t, a0, b0);
 merge(t, t, b1);
void dfs(treap *t) {
 if (!t) return;
 lazy propagation(t);
 dfs(t->1);
  solve(t->data);
 dfs(t->r);
treap *create_node(int data, int priority) {
 treap *ret = new treap;
  ret->data = data;
 ret->priority = priority;
 ret->1 = 0;
  ret->r = 0:
 ret->sz = 1;
  ret->lazy = 0;
  ret->lazy2 = 0;
 ret->parent = 0;
 return ret;
```

colorupdate.cpp

```
74 lines
const int inf = 1e15;
struct color upd {
#define left fir
#define right sec.fir
#define color sec.sec
 set<pii> ranges;
 vector<pii> erased;
  color_upd(int n) // inicialmente, todo mundo pintado com a
      cor inf
    // nao usar cores negativas!!!!!!!!
    ranges.insert(\{0, \{n-1, inf\}\}\);
 int get(int i) {
    auto it = ranges.upper_bound({i, {1e18, 1e18}});
    if (it == ranges.begin()) return -1;
    return (*(it)).color;
 void del(int 1, int r) // apaga o intervalo [1, r]
    erased.clear();
    auto it = ranges.upper_bound({1, {0, 0}});
    if (it != ranges.begin()) {
      it--;
    while (it != ranges.end()) {
     if ((*(it)).left > r)
       break:
      else if ((*(it)).right >= 1)
        erased.push back(*it);
      it++;
    if (erased.size() > 0) {
      int sz = erased.size();
      auto it = ranges.lower_bound({erased[0].left, {0, 0}});
```

```
auto it2 = ranges.lower_bound({erased[sz - 1].left, {0, 0}
          }}):
      pii ini = *it, fim = *it2;
      it2++;
      ranges.erase(it, it2);
      pii upd1 = {ini.left, {1 - 1, ini.color}};
     pii upd2 = \{r + 1, \{fim.right, fim.color\}\};
      erased[0].left = max(erased[0].left, 1);
      erased[sz - 1].right = min(erased[sz - 1].right, r);
      if (upd1.left <= upd1.right) ranges.insert(upd1);</pre>
      if (upd2.left <= upd2.right) ranges.insert(upd2);
  void add(int a, int b, int c) {
    auto it = ranges.lower_bound({a, {b, 0}});
    pii aa = \{-1, \{-1, -1\}\};
    pii bb = \{-1, \{-1, -1\}\};
    if (it != ranges.end()) {
     if ((*it).color == c && (*it).left == b + 1) {
       aa = *it;
       b = (*it).right;
    if (it != ranges.begin()) {
     if ((*it).color == c && (*it).right == a - 1) {
       bb = *it;
        a = (*it).left;
    ranges.erase(aa);
    ranges.erase(bb);
    ranges.insert({a, {b, c}});
  void upd(int a, int b, int c) // pinta o intervalo [a, b]
    del(a, b);
    add(a, b, c);
};
```

bit2d.cpp

```
51 lines
struct bit2d {
  vector<int> ord;
  vector<vector<int>> t;
  vector<vector<int>> coord;
  bit2d(vector<pi> &pts) // recebe todos os pontos que vao ser
       inseridos pra construir, mas nao insere eles
    sort(pts.begin(), pts.end());
    for (auto const &a : pts) {
     if (ord.empty() || a.fir != ord.back()) ord.pb(a.fir);
   t.resize(ord.size() + 1);
    coord.resize(t.size());
    for (auto &a : pts) {
     swap(a.fir, a.sec);
    sort(pts.begin(), pts.end());
    for (auto &a : pts) {
     swap(a.fir, a.sec);
      for (int on = upper_bound(ord.begin(), ord.end(), a.fir)
          - ord.begin(); on < t.size(); on += on & -on) {
       if (coord[on].empty() || coord[on].back() != a.sec)
            coord[on].push_back(a.sec);
```

```
for (int i = 0; i < t.size(); i++) t[i].assign(coord[i].
        size() + 1, 0);
 void add(int x, int y, int v) // v[a][b] += v
   for (int xx = upper_bound(ord.begin(), ord.end(), x) - ord.
        begin(); xx < t.size(); xx += xx & -xx) {
     for (int yy = upper_bound(coord[xx].begin(), coord[xx].
          end(), y) - coord[xx].begin(); yy < t[xx].size();</pre>
           yy += yy & -yy
       t[xx][yy] += v;
 int qry(int x, int y) // soma de todos os v[a][b] com (a <=
      x && b <= y)
   int ans = 0;
    for (int xx = upper_bound(ord.begin(), ord.end(), x) - ord.
        begin(); xx > 0; xx -= xx & -xx) {
     for (int yy = upper_bound(coord[xx].begin(), coord[xx].
          end(), y) - coord[xx].begin(); yy > 0; yy -= yy & -
       ans += t[xx][yy];
    return ans;
 int gry2(int x1, int y1, int x2, int y2) {
   return qry (x2, y2) - qry (x2, y1 - 1) - qry (x1 - 1, y2) +
        qry(x1 - 1, y1 - 1);
 void add2(int x1, int y1, int x2, int y2, int v) {
   add(x1, y1, v);
   add(x1, y2 + 1, -v);
   add(x2 + 1, y1, -v);
    add(x2 + 1, y2 + 1, v);
};
```

mo.cpp

cl--;

45 lines namespace mo { struct query { int idx, l, r; int block; vector<query> queries; vector<int> ans; // bool cmp(query &x, query &y){ essa funcao de ordenacao pode funcionar em caso de TLE int ablock = x.1 / MAGIC, bblock = y.1 / MAGIC; if (ablock != bblock) return ablock < bblock; if (ablock & 1) return x.r < y.r; // return x.r > y.r; // } bool cmp(query &x, query &y) { if (x.1 / block != y.1 / block) return x.1 / block < y.1 / block: return x.r < y.r; void run() { block = (int)sqrt(n); sort(queries.begin(), queries.end(), cmp); ans.resize(queries.size()); int c1 = 0, cr = -1, sum = 0; auto add = [&] (int x) { sum += x; }; auto rem = [&] (int x) { sum -= x; }; for (int i = 0; i < queries.size(); i++) { while (cl > queries[i].1) {

```
add(v[cl]);
    while (cr < queries[i].r) {
     cr++;
     add(v[cr]);
    while (cl < queries[i].l) {
      rem(v[cl]);
     cl++;
    while (cr > queries[i].r) {
     rem(v[cr]);
    ans[queries[i].idx] = sum;
} // namespace mo
```

moupdate.cpp

```
63 lines
int n, q;
int v[MAXN];
int vv[MAXN];
namespace mo {
struct query {
 int idx, l, r, t;
struct update {
 int i, prevx, x;
int block:
vector<query> queries;
vector<update> updates;
vector<int> ans;
bool cmp(query x, query y) {
 if (x.1 / block != y.1 / block) return x.1 / block < y.1 /
  if (x.r / block != y.r / block) return x.r / block < y.r /
      block:
  return x.t < y.t;
void run() {
  block = 3153; // (2 * n) ^ 0.666
  sort(queries.begin(), queries.end(), cmp);
  ans.resize(queries.size());
  int c1 = 0, cr = -1, sum = 0, t = 0;
  auto add = [\&] (int x) { sum += x; };
  auto rem = [\&] (int x) { sum -= x; };
  for (int i = 0; i < queries.size(); i++) {
    while (cl > queries[i].1) {
      cl--;
      add(v[cl]);
    while (cr < queries[i].r) {
      cr++;
      add(v[cr]);
    while (cl < queries[i].1) {
      rem(v[cl]);
      cl++;
    while (cr > queries[i].r) {
      rem(v[cr]);
      cr--;
    while (t > queries[i].t) {
```

```
if (queries[i].1 <= updates[t].i && queries[i].r >=
          updates[t].i) {
       rem(updates[t].x);
       add(updates[t].prevx);
     v[updates[t].i] = updates[t].prevx;
   while (t < queries[i].t) {
     if (queries[i].1 <= updates[t].i && queries[i].r >=
          updates[t].i) {
       rem(updates[t].prevx);
       add(updates[t].x);
     v[updates[t].i] = updates[t].x;
    ans[queries[i].idx] = sum;
} // namespace mo
segtree2d.cpp
struct segtree2d {
 int n, m;
 vector<vector<int>> seg;
 int neutral() { return 0; }
 int merge(int a, int b) { return a + b; }
  segtree2d(int nn, int mm) {
   n = nn, m = mm;
  int qry(int x1, int y1, int x2, int y2) {
   int ret = neutral();
   int y3 = y1 + m, y4 = y2 + m;
   for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2) {
```

```
seg = vector<vector<int>>(2 * n, vector<int>(2 * m, neutral
      for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /= 2) {
        if (x1 \% 2 == 1 \text{ and } y1 \% 2 == 1) \text{ ret } = \text{merge}(\text{ret, seg}[
              x1][y1]);
         if (x1 \% 2 == 1 \text{ and } y2 \% 2 == 0) ret = merge(ret, seg[
              x1][y2]);
         if (x2 % 2 == 0 and y1 % 2 == 1) ret = merge(ret, seg[
              x2][y1]);
        if (x2 \% 2 == 0 \text{ and } y2 \% 2 == 0) \text{ ret } = \text{merge}(\text{ret, seg}[
              x2][y2]);
    return ret;
  void upd(int x, int y, int val) {
    int y2 = y += m;
    for (x += n; x; x /= 2, y = y2) {
      if (x >= n)
        seg[x][y] = val;
        seg[x][y] = merge(seg[2 * x][y], seg[2 * x + 1][y]);
      while (y \neq 2) seg[x][y] = merge(seg[x][2 * y], seg[x][2
            * y + 1]);
};
```

persistentseg.cpp

```
37 lines
struct node {
  int item, l, r;
  node (int 1, int r, int item) : l(1), r(r), item(item) {}
```

```
int n, q;
vector<node> seg;
vector<int> roots;
void init() { seg.resize(1); }
int newleaf(int vv) {
 int p = seg.size();
 seg.pb(node(0, 0, vv));
 return p;
int newpar(int 1, int r) {
 int p = seg.size();
 seg.pb(node(l, r, seg[l].item + seg[r].item));
int upd(int i, int l, int r, int pos) {
 if (l == r) return newleaf(seg[i].item + 1);
 int mid = (1 + r) >> 1;
 if (pos <= mid) return newpar(upd(seg[i].1, 1, mid, pos), seg
 return newpar(seg[i].1, upd(seg[i].r, mid + 1, r, pos));
int build(int 1, int r) {
 if (1 == r) return newleaf(0);
 int mid = (1 + r) >> 1;
 return newpar(build(1, mid), build(mid + 1, r));
int qry(int vl, int vr, int l, int r, int k) {
 if (1 == r) return 1;
 int mid = (1 + r) >> 1;
 int c = seg[seg[vr].1].item - seg[seg[vl].1].item;
 if (c \ge k) return qry(seg[vl].l, seg[vr].l, l, mid, k);
 return qry(seg[v1].r, seg[vr].r, mid + 1, r, k - c);
rmq.cpp
struct rmq {
 int n;
 vector<vector<pi>> m;
 vector<int> log;
 rmq() {}
 rmq(vector<pi> &v) {
   n = v.size();
   log.resize(n + 1);
   log[1] = 0;
    for (int i = 2; i \le n; i++) log[i] = log[i / 2] + 1;
   int sz = log[n] + 2;
   m = vector<vector<pi>>>(sz, vector<pi>(n + 1));
   for (int i = 0; i < n; i++) {
     m[0][i] = v[i];
   for (int j = 1; j < sz; j++) {
     for (int i = 0; i + (1 << j) <= n; i++) m[j][i] = min(m[j])
           -1[i], m[j -1][i + (1 << (j -1))]);
 int qry(int a, int b) { return min(m[log[b - a + 1]][a], m[
      log[b - a + 1]][b - (1 << log[b - a + 1]) + 1]).second;
};
```

```
binarylifting.cpp
```

```
item st[MAXN][21];
for (int i = 0; i < n; i++) {
  st[i][0].nxt = min(i + 1, n - 1);
  st[i][0].sum = v[st[i][0].nxt];
```

```
for (int i = 1; i < 21; i++) {
 for (int v = 0; v < n; v++) {
    st[v][i].nxt = st[st[v][i - 1].nxt][i - 1].nxt;
    st[v][i].sum = st[v][i - 1].sum + st[st[v][i - 1].nxt][i -
         1].sum;
while (q--) {
 int 1, r;
 cin >> 1 >> r;
 int ans = v[1], len = r - 1;
  for (int i = 20; i >= 0; i--) {
    if (len & (1 << i)) {
      ans += st[1][i].sum;
      l = st[l][i].nxt;
  cout << ans << endl;
```

minqueue.cpp

```
namespace min_queue {
deque<pi> q;
int l, r;
void init() {
 1 = r = 1;
 q.clear();
void push(int v) {
 while (!q.empty() && v < q.back().fir) q.pop_back();</pre>
 q.pb(\{v, r\});
void pop() {
 if (!q.empty() && q.front().sec == 1) q.pop_front();
int getmin() { return q.front().fir; }
} // namespace min_queue
```

Graph (4)

floydwarshall.cpp

```
int dist[MAXN][MAXN];
void floyd_warshall() {
 for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
      for (int j = 0; j < n; j++) {
        dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j]);
void initialize() {
 for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
     if (i == j) {
        dist[i][j] = 0;
      } else {
        dist[i][j] = INF;
```

visited.resize(n);

decompose tree(0):

```
centroiddecomp.cpp
int n, k, resp;
vector<int> adj[MAXN];
vector<int> cnt;
namespace cd {
vector<int> subtree size;
vector<bool> visited;
void dfs(int s, int f) {
  subtree size[s] = 1;
  for (auto const &v : adj[s]) {
   if (v != f && !visited[v]) {
     dfs(v, s);
     subtree_size[s] += subtree_size[v];
int get centroid(int s, int f) {
  bool is_centroid = true;
  int heaviest child = -1;
  for (auto const &v : adj[s]) {
    if (v != f && !visited[v]) {
      if (subtree_size[v] > sz / 2) is_centroid = false;
      if (heaviest child == -1 || subtree size[v] >
           subtree_size[heaviest_child]) heaviest_child = v;
  return (is_centroid && sz - subtree_size[s] <= sz / 2) ? s :
       get centroid(heaviest child, s);
void dfs2(int s, int f, int d) {
  while (d >= cnt.size()) cnt.pb(0);
  for (auto const &v : adj[s]) {
    if (v != f \&\& !visited[v]) dfs2(v, s, d + 1);
void solve(int s) {
  vector<int> tot;
  for (auto const &v : adj[s]) {
    if (visited[v]) continue;
    cnt.clear();
    dfs2(v, s, 1);
    for (int i = 1; i < cnt.size(); i++) {
     if (k - i < tot.size() && k - i >= 1) resp += (cnt[i] *
           tot[k - i]);
    for (int i = 1; i < cnt.size(); i++) {
      while (i >= tot.size()) tot.pb(0);
      tot[i] += cnt[i];
  if (k < tot.size()) resp += tot[k];
int decompose_tree(int s) {
  sz = 0;
  dfs(s, s);
  int cend_tree = get_centroid(s, s);
  visited[cend_tree] = true;
  solve(cend_tree);
  for (auto const &v : adj[cend_tree]) {
   if (!visited[v]) decompose_tree(v);
  return cend_tree;
void init() {
  subtree size.resize(n);
```

```
} // namespace cd
dsu.cpp
                                                          24 lines
struct dsu {
 int tot;
 vector<int> parent;
 vector<int> sz;
 dsu(int n) {
   parent.resize(n);
   sz.resize(n);
    for (int i = 0; i < n; i++) {
     parent[i] = i;
     sz[i] = 1;
 int find_set(int i) { return parent[i] = (parent[i] == i) ? i
       : find set(parent[i]); }
 void make_set(int x, int y) {
   x = find_set(x), y = find_set(y);
   if (x != y) {
      if (sz[x] > sz[v]) swap(x, v);
     parent[x] = y;
     sz[y] += sz[x];
     tot--;
};
dsubipartido.cpp
                                                          41 lines
struct dsu {
 vector<pi> parent;
 vector<int> rank:
 vector<int> bipartite;
 dsu(int n) {
   parent.resize(n);
   rank.resize(n);
   bipartite.resize(n);
    for (int v = 0; v < n; v++) {
     parent[v] = \{v, 0\};
     rank[v] = 0;
     bipartite[v] = 1;
 dsu() {}
 pi find_set(int v) {
   if (v != parent[v].fir) {
     int parity = parent[v].sec;
     parent[v] = find_set(parent[v].fir);
     parent[v].sec ^= parity;
   return parent[v];
 void add edge(int a, int b) {
   pi pa = find_set(a);
   a = pa.fir;
   int x = pa.sec;
   pi pb = find_set(b);
   b = pb.fir;
   int y = pb.sec;
   if (a == b) {
     if (x == y) bipartite[a] = 0;
     if (rank[a] < rank[b]) swap(a, b);
```

```
parent[b] = \{a, x^y^1\};
      bipartite[a] &= bipartite[b];
      if (rank[a] == rank[b]) rank[a]++;
  bool is_bipartite(int v) { return bipartite[find_set(v).fir];
cycledetection.cpp
int n, m, idx;
vector<int> cycles[MAXN];
vector<int> adj[MAXN];
int color[MAXN];
int parent[MAXN];
int ans[MAXN];
void dfs(int u, int p) { // chama dfs a partir de qm tem cor 0
  if (color[u] == 2) return;
  if (color[u] == 1) {
    idx++;
    int curr = p;
    ans[curr] = idx;
    cycles[idx].pb(curr);
    while (curr != u) {
      curr = parent[curr];
      cycles[idx].pb(curr);
      ans[curr] = idx;
    return;
  parent[u] = p;
  color[u] = 1;
  for (auto const &v : adj[u])
   if (v != parent[u]) dfs(v, u);
  color[u] = 2;
blockcuttree.cpp
struct block cut tree {
  // Se art[i] >= 1, i eh ponto de articulação
  // tree - eh a propia block-cut tree
  // pos[i] responde a qual vertice da arvore vertice i
      pertence
  vector<vector<int>> g, blocks, tree;
  vector<vector<pi>>> edgblocks; // sao as arestas do bloco i
  stack<int> s;
  stack<pi> s2;
  vector<int> id, art, pos;
  block_cut_tree(vector<vector<int>> g_) : q(g_) {
    int n = q.size();
    id.resize(n, -1), art.resize(n), pos.resize(n);
    build();
  int dfs(int i, int &t, int p = -1) {
    int lo = id[i] = t++;
    s.push(i);
    if (p != -1) {
      s2.emplace(i, p);
    for (int j : q[i]) {
      if (j != p \text{ and } id[j] != -1) s2.emplace(i, j);
    for (int j : g[i]) {
     if (j != p) {
```

if (id[j] == -1) {
 int val = dfs(j, t, i);

```
lo = min(lo, val);
          if (val >= id[i]) {
            art[i]++;
            blocks.emplace_back(1, i);
            while (blocks.back().back() != j) {
             blocks.back().pb(s.top());
             s.pop();
            edgblocks.emplace_back(1, s2.top());
            s2.pop();
            pi aux = {j, i};
            while (edgblocks.back().back() != aux) {
              edgblocks.back().pb(s2.top());
              s2.pop();
          // if (val > id[i]) aresta i-j eh ponte
          lo = min(lo, id[j]);
    if (p == -1 \text{ and } art[i]) {
     art[i]--;
    return lo;
  void build() {
    int t = 0;
    for (int i = 0; i < q.size(); i++) {
     if (id[i] == -1) dfs(i, t, -1);
    tree.resize(blocks.size());
    for (int i = 0; i < g.size(); i++) {
     if (art[i]) pos[i] = tree.size(), tree.emplace_back();
    for (int i = 0; i < blocks.size(); i++) {
      for (int j : blocks[i]) {
       if (!art[i])
         pos[j] = i;
          tree[i].pb(pos[j]), tree[pos[j]].pb(i);
};
```

bridges.cpp

```
21 lines
nt n, m, timer;
vector<pi> edges;
vector<bool> is_bridge;
vector<pi> adj[MAXN];
int tin[MAXN];
int low[MAXN];
                          // memset -1
                          // memset -1
bool vis[MAXN];
void dfs(int v, int p) { // chama de quem nao foi vis ainda
  vis[v] = true;
  tin[v] = timer, low[v] = timer++;
  for (auto const &u : adj[v]) {
   if (u.fir == p) continue;
    if (vis[u.fir]) {
     low[v] = min(low[v], tin[u.fir]);
     continue;
    dfs(u.fir, v);
    low[v] = min(low[v], low[u.fir]);
    if (low[u.fir] > tin[v]) is_bridge[u.sec] = 1;
```

```
dinic.cpp
                                                          83 lines
#define INF 1e9
struct edge {
int to, from, flow, capacity, id;
struct dinic {
 int n, src, sink;
 vector<vector<edge>> adj;
 vector<int> level;
 vector<int> ptr;
 dinic(int sz) {
   n = sz;
    adj.resize(n);
   level.resize(n);
   ptr.resize(n);
 void add edge(int a, int b, int c, int id) {
    adj[a].pb({b, (int)adj[b].size(), c, c, id});
    adj[b].pb({a, (int)adj[a].size() - 1, 0, 0, id});
 bool bfs() {
   level.assign(n, -1);
   level[src] = 0;
    queue<int> q;
   q.push(src);
    while (!a.emptv()) {
     int u = q.front();
     q.pop();
      for (auto at : adj[u]) {
       if (at.flow && level[at.to] == -1) {
         q.push(at.to);
          level[at.to] = level[u] + 1;
    return level[sink] != -1;
 int dfs(int u, int flow) {
   if (u == sink || flow == 0) return flow;
    for (int &p = ptr[u]; p < adj[u].size(); p++) {
      edge &at = adj[u][p];
      if (at.flow && level[u] == level[at.to] - 1) {
       int kappa = dfs(at.to, min(flow, at.flow));
       at.flow -= kappa;
       adj[at.to][at.from].flow += kappa;
       if (kappa != 0) return kappa;
   return 0;
 int run() {
   int max_flow = 0;
    while (bfs()) {
     ptr.assign(n, 0);
      while (1) {
        int flow = dfs(src, INF);
       if (flow == 0) break;
       max_flow += flow;
    return max_flow;
 vector<pii> cut_edges() // arestas do corte minimo
   bfs();
```

```
vector<pii> ans;
    for (int i = 0; i < n; i++) {
      for (auto const &j : adj[i]) {
       if (level[i] != -1 && level[j.to] == -1 && j.capacity >
              0) ans.pb({j.capacity, {i, j.to}});
   return ans;
 vector<int> flow_edges(int n, int m) // fluxo em cada aresta
      , na ordem da entrada
   vector<int> ans(m);
   for (int i = 0; i < n; i++) {
     for (auto const &j : adj[i])
       if (!j.capacity) ans[j.id] = j.flow;
   return ans:
};
```

```
hopcroftkarp.cpp
                                                          137 lines
#define INF 1e9
struct hopcroft_karp {
 vector<int> match;
 vector<int> dist;
 vector<vector<int>> adj;
 int n, m, t;
 hopcroft_karp(int a, int b) {
   n = a, m = b;
    t = n + m + 1;
    match.assign(t, n + m);
    dist.assign(t, 0);
    adj.assign(t, vector<int>{});
 void add_edge(int u, int v) {
    adj[u].pb(v);
    adj[v].pb(u);
 bool bfs() {
    queue<int> q;
    for (int u = 0; u < n; u++) {
      if (match[u] == n + m)
        dist[u] = 0, q.push(u);
      else
        dist[u] = INF;
    dist[n + m] = INF;
    while (!q.empty()) {
     int u = q.front();
      q.pop();
      if (dist[u] < dist[n + m]) {
        for (auto const &v : adj[u]) {
          if (dist[match[v]] == INF) {
            dist[match[v]] = dist[u] + 1;
            q.push(match[v]);
    return dist[n + m] < INF;
 bool dfs(int u) {
    if (u < n + m) {
      for (auto const &v : adj[u]) {
        if (dist[match[v]] == dist[u] + 1 && dfs(match[v])) {
          match[v] = u;
          match[u] = v;
```

```
return true;
     dist[u] = INF;
     return false;
   return true;
  vector<pi> run() {
   int cnt = 0;
   while (bfs())
     for (int u = 0; u < n; u++)
       if (match[u] == n + m \&\& dfs(u)) cnt++;
   vector<pi> ans;
    for (int v = n; v < n + m; v++)
     if (match[v] < n + m) ans.pb(\{match[v], v\});
  vector<int> mvc() // minimum vertex cover
   vector<pi> ans = run();
   vector<bool> vis(n + m, 0);
    for (int i = 0; i < n; i++) {
     if (match[i] == n + m) {
       queue<int> q;
       q.push(i);
        while (!q.empty()) {
         int x = q.front();
         q.pop();
         vis[x] = 1;
          for (auto const &y : adj[x]) {
           if (!vis[y]) {
             vis[y] = 1;
              q.push(match[y]);
    vector<int> vc;
    for (int i = 0; i < n; i++) {
     if (!vis[i]) vc.pb(i);
    for (int i = n; i < n + m; i++) {
     if (vis[i]) vc.pb(i);
   return vc;
  vector<pi> mec() // minimum edge cover
   vector<pi> ans = run();
    for (int i = 0; i < n + m; i++) {
     if (match[i] == n + m \&\& adj[i].size() > 0) {
       if (i < n)
         ans.pb({i, adj[i][0]});
       else
          ans.pb({adj[i][0], i});
    return ans;
// minimum path cover on dag
// minimum set of paths such that each of the vertices belongs
    to exactly one path
vector<vector<int>> mpc(int n, vector<pi> &e) {
 hopcroft_karp h(n, n);
  for (auto const &i : e) h.add_edge(i.fir, n + i.sec);
  vector<pi> mat = h.run();
```

```
vector<int> prv(n, -1);
 vector<int> nxt(n, -1);
 for (int i = 0; i < mat.size(); i++) {
   nxt[mat[i].fir] = mat[i].sec - n;
   prv[mat[i].sec - n] = mat[i].fir;
 vector<vector<int>> ans;
 for (int i = 0; i < n; i++) {
   if (prv[i] == -1 && nxt[i] == -1) {
     ans.pb(\{i\});
   } else if (prv[i] == -1) {
     vector<int> curr;
     int x = i:
     while (1) {
       curr.pb(x);
       if (nxt[x] == -1) break;
       x = nxt[x];
      ans.pb(curr);
 return ans;
pushrelabel.cpp
                                                          93 lines
#define INF 1e9
struct edge {
 int dest, back, f, c, id;
struct push_relabel {
 int n:
 vector<vector<edge>> g;
 vector<int> ec;
 vector<edge *> cur;
 vector<vector<int>> hs;
 vector<int> H;
 push\_relabel(int sz) : q(sz), ec(sz), cur(sz), hs(2 * sz), H(
      sz) { n = sz; }
 void add_edge(int s, int t, int cap, int rcap, int id) {
   if (s == t) return;
   g[s].pb({t, (int)g[t].size(), 0, cap, id});
   g[t].pb({s, (int)g[s].size() - 1, 0, rcap, -1});
 void add_flow(edge &e, int f) {
   edge &back = g[e.dest][e.back];
   if (!ec[e.dest] && f) hs[H[e.dest]].push_back(e.dest);
   e.f += f:
   e.c -= f;
   ec[e.dest] += f;
   back.f -= f;
   back.c += f;
   ec[back.dest] -= f;
 int calc(int s, int t) {
   int v = g.size();
   H[s] = v;
    ec[t] = 1;
   vector<int> co(2 * v);
    co[0] = v - 1;
    for (int i = 0; i < v; i++) cur[i] = g[i].data();
    for (edge &e : q[s]) add_flow(e, e.c);
    for (int hi = 0;;) {
      while (hs[hi].empty())
       if (!hi--) return -ec[s];
      int u = hs[hi].back();
     hs[hi].pop_back();
      while (ec[u] > 0) {
```

if (cur[u] == g[u].data() + g[u].size()) {

```
H[u] = INF;
          for (edge &e : q[u])
            if (e.c && H[u] > H[e.dest] + 1) H[u] = H[e.dest] +
                  1, cur[u] = &e;
          if (++co[H[u]], !--co[hi] && hi < v)
            for (int i = 0; i < v; i++)
              if (hi < H[i] \&\& H[i] < v) --co[H[i]], H[i] = v +
                    1:
          hi = H[u];
        } else if (cur[u] \rightarrow c \&\& H[u] == H[cur[u] \rightarrow dest] + 1)
          add_flow(*cur[u], min(ec[u], cur[u]->c));
          ++cur[u];
 vector<int> flow_edges(int m) // fluxo em cada aresta
    vector<int> ans(m);
    for (int i = 0; i < n; i++) {
      for (auto const &j : q[i]) {
       if (j.id != -1) ans[j.id] = j.f;
    return ans;
struct flow_with_demands {
 push_relabel pr;
  vector<int> in, out;
  flow_with_demands(int sz) : n(sz), pr(sz + 2), in(sz), out(sz
  void add_edge(int u, int v, int cap, int dem, int id) {
    pr.add_edge(u, v, cap - dem, 0, id);
    out[u] += dem, in[v] += dem;
 int run(int s, int t) {
    pr.add_edge(t, s, INF, 0, -1);
    for (int i = 0; i < n; i++) {
      pr.add_edge(n, i, in[i], 0, -1);
      pr.add_edge(i, n + 1, out[i], 0, -1);
    return pr.calc(n, n + 1);
 bool check() // todas as constraints foram satisfeitas?
    for (auto const &i : pr.g[n]) {
     if (i.c > 0) return 0;
    return 1;
};
```

hld.cpp

65 lin

12

```
struct hld {
  int n, cur_pos;
  segtree seg;
  vector<vector<int>> adj;
  vector<vector<int>> depth, heavy, head, pos, sz;
  int dfs(int s) {
    int size = 1, max_c_size = 0;
    for (auto const &c : adj[s]) {
        if (c != parent[s]) {
            parent[c] = s;
            depth[c] = depth[s] + 1;
            int c_size = dfs(c);
    }
}
```

hldedge strongorientation twosat

```
size += c size;
       if (c_size > max_c_size) max_c_size = c_size, heavy[s]
            = c:
   return sz[s] = size;
  void decompose (int s, int h) {
   head[s] = h;
   pos[s] = cur_pos++;
   if (heavy[s] != -1) decompose(heavy[s], h);
    for (int c : adj[s]) {
     if (c != parent[s] && c != heavy[s]) decompose(c, c);
  hld(vector<vector<int>> &g) {
   n = g.size();
   adj = q;
    seg = segtree(n);
   parent.assign(n, -1);
   depth.assign(n, -1);
   heavy.assign(n, -1);
   head.assign(n, -1);
   pos.assign(n, -1);
   sz.assign(n, 1);
    cur_pos = 0;
   dfs(0);
    decompose(0, 0);
  int query_path(int a, int b) {
    int res = 0;
    for (; head[a] != head[b]; b = parent[head[b]]) {
     if (depth[head[a]] > depth[head[b]]) swap(a, b);
     res += seg.query(0, n - 1, pos[head[b]], pos[b], 1);
    if (depth[a] > depth[b]) swap(a, b);
    res += seq.query(0, n - 1, pos[a], pos[b], 1);
  void update_path(int a, int b, int x) {
    for (; head[a] != head[b]; b = parent[head[b]]) {
     if (depth[head[a]] > depth[head[b]]) swap(a, b);
     seg.update(1, 0, n - 1, pos[head[b]], pos[b], x);
    if (depth[a] > depth[b]) swap(a, b);
    seq.update(1, 0, n - 1, pos[a], pos[b], x);
  void update_subtree(int a, int x) { seq.update(1, 0, n - 1,
      pos[a], pos[a] + sz[a] - 1, x); }
  int query_subtree(int a) { return seg.query(0, n - 1, pos[a],
       pos[a] + sz[a] - 1, 1); }
  int lca(int a, int b) {
   if (pos[a] < pos[b]) swap(a, b);
    return (head[a] == head[b]) ? b : lca(parent[head[a]], b);
};
```

hldedge.cpp

namespace hld {

int cur_pos; vector<int> parent, depth, heavy, head, pos, sz, up; int dfs(int s) { int size = 1, max c size = 0; for (auto const &c : adj[s]) { if (c.fir != parent[s]) { parent[c.fir] = s; depth[c.fir] = depth[s] + 1;int c_size = dfs(c.fir);

```
size += c size;
      if (c_size > max_c_size) max_c_size = c_size, heavy[s] =
          c.fir:
 return sz[s] = size;
void decompose(int s, int h) {
 head[s] = h;
 pos[s] = cur_pos++;
  seq::v[pos[s]] = up[s];
 for (auto const &c : adj[s]) {
    if (c.fir != parent[s] && c.fir == heavy[s]) {
      up[c.fir] = c.sec;
      decompose(heavy[s], h);
 for (auto const &c : adj[s]) {
   if (c.fir != parent[s] && c.fir != heavy[s]) {
     up[c.fir] = c.sec;
      decompose(c.fir, c.fir);
 }
void init() {
 parent.assign(MAXN, -1);
 depth.assign(MAXN, -1);
 heavy.assign(MAXN, -1);
 head.assign(MAXN, -1);
 pos.assign(MAXN, -1);
  sz.assign(MAXN, 1);
  up.assign(MAXN, 0);
  cur_pos = 0;
  dfs(0);
  decompose(0, 0);
  seg::build(0, n - 1, 1);
int query_path(int a, int b) {
 int res = -1;
  for (; head[a] != head[b]; b = parent[head[b]]) {
    if (depth[head[a]] > depth[head[b]]) swap(a, b);
    res = max(res, seg::query(0, n - 1, pos[head[b]], pos[b],
  if (depth[a] > depth[b]) swap(a, b);
  res = \max(\text{res, seq}: \text{query}(0, n - 1, pos[a] + 1, pos[b], 1));
  return res;
void update path(int a, int b, int x) {
  for (; head[a] != head[b]; b = parent[head[b]]) {
    if (depth[head[a]] > depth[head[b]]) swap(a, b);
    seg::update(1, 0, n - 1, pos[head[b]], pos[b], x);
 if (depth[a] > depth[b]) swap(a, b);
  seg::update(1, 0, n - 1, pos[a] + 1, pos[b], x);
void update_subtree(int a, int x) { seq::update(1, 0, n - 1,
    pos[a] + 1, pos[a] + sz[a] - 1, x); }
int query_subtree(int a, int x) { return seg::query(0, n - 1,
    pos[a] + 1, pos[a] + sz[a] - 1, 1); }
} // namespace hld
```

strongorientation.cpp

67 lines

// encontrar uma orientacao para as arestas tal que o numero // minimo de scc eh o menor possivel // numero minimo de scc = numero de componentes conexas + numero de pontes int n, m, timer, comps, bridges;

```
vector<pi> edges;
vector<pi> adj[MAXN];
int tin[MAXN]; // memset -1
int low[MAXN]; // memset -1
bool vis[MAXN];
char orient[MAXN];
void find_bridges(int v) { // chama se tem tin == -1
 low[v] = timer, tin[v] = timer++;
 for (auto const &p : adj[v]) {
   if (vis[p.sec]) continue;
   vis[p.sec] = true;
    orient[p.sec] = (v == edges[p.sec].first) ? '>' : '<';
    if (tin[p.fir] == -1) {
     find_bridges(p.fir);
     low[v] = min(low[v], low[p.fir]);
     if (low[p.fir] > tin[v]) bridges++;
      low[v] = min(low[v], low[p.fir]);
```

twosat.cpp

```
74 lines
struct two_sat {
 int n;
 vector<vector<int>> q, gr; // gr is the reversed graph
 vector<int> comp, ord, ans; // comp[v]: ID of the SCC
      containing node v
 vector<bool> vis:
 two sat() {}
 two_sat(int sz) {
   n = sz;
   g.assign(2 * n, vector<int>());
   qr.assign(2 * n, vector<int>());
   comp.resize(2 * n);
   vis.resize(2 * n);
   ans.resize(2 * n);
 void add_edge(int u, int v) {
    q[u].push_back(v);
   gr[v].push_back(u);
 // int x, bool val: if 'val' is true, we take the variable to
       be x. Otherwise we take it to be x's complement (not x)
 void implies(int i, bool f, int j, bool g) // a -> b
   add_edge(i + (f ? 0 : n), j + (g ? 0 : n));
    add_edge(j + (g ? n : 0), i + (f ? n : 0));
 void add_clause_or(int i, bool f, int j, bool g) // At least
       one of them is true
    add_edge(i + (f ? n : 0), j + (g ? 0 : n));
    add_edge(j + (g ? n : 0), i + (f ? 0 : n));
 void add_clause_xor(int i, bool f, int j, bool g) // only
      one of them is true
    add_clause_or(i, f, j, g);
    add_clause_or(i, !f, j, !g);
 void add_clause_and(int i, bool f, int j, bool g) // both of
       them have the same value
    add_clause_xor(i, !f, j, g);
```

void set(int i, bool f) // Set a variable

```
add clause or(i, f, i, f);
 void top_sort(int u) {
   vis[u] = 1;
   for (auto const &v : g[u]) {
    if (!vis[v]) top_sort(v);
   ord.push_back(u);
 void scc(int u, int id) {
   vis[u] = 1;
   comp[u] = id;
   for (auto const &v : gr[u]) {
     if (!vis[v]) scc(v, id);
 bool solve() {
   fill(vis.begin(), vis.end(), 0);
   for (int i = 0; i < 2 * n; i++) {
     if (!vis[i]) top_sort(i);
   fill(vis.begin(), vis.end(), 0);
   reverse(ord.begin(), ord.end());
   int id = 0;
   for (const auto &v : ord) {
     if (!vis[v]) scc(v, id++);
   for (int i = 0; i < n; i++) {
     if (comp[i] == comp[i + n]) return 0;
     ans[i] = (comp[i] > comp[i + n]) ? 1 : 0;
   return 1;
};
```

sack.cpp

```
vector<int> adi[MAXN];
vector<int> v[MAXN];
int c[MAXN];
int cnt[MAXN];
int sz[MAXN];
void dfs_sz(int x, int p) {
  sz[x] = 1;
  for (auto const &i : adj[x]) {
   if (i != p) {
     dfs sz(i, x);
     sz[x] += sz[i];
void modify(int c, int val) { cnt[c] += val; }
void dfs(int x, int p, bool keep) {
 int best = -1, big_child = -1;
  for (auto const &i : adj[x]) {
   if (i != p && sz[i] > best) {
     best = sz[i];
     big child = i;
  for (auto const &i : adj[x]) {
   if (i != p && i != big_child) dfs(i, x, 0);
  if (big_child !=-1) {
   dfs(big_child, x, 1);
   swap(v[x], v[big_child]); // 0(1)
  v[x].pb(x);
```

```
sack reroot hungarian mincostmaxflow
 modify(c[x], 1); // adiciona
 for (auto const &i : adj[x]) {
   if (i != p && i != big_child) {
      for (auto const &j : v[i]) {
       v[x].pb(j);
       modify(c[j], 1); // adiciona
 // a cor c aparece cnt[c] vezes na subtree de x
 // dai vc pode fazer algo tendo essa informacao
 // seja responser queries ou algo do tipo aqui
    for (auto const &i : v[x]) modify(c[i], -1); // remove
reroot.cpp
int n;
vector<int> adj[MAXN];
int sz[MAXN];
int dp[MAXN];
int dfs(int u, int v) {
 sz[u] = 1;
 for (auto const &i : adj[u])
   if (i != v) sz[u] += dfs(i, u);
 return sz[u];
void reroot(int u, int v) {
 for (auto const &i : adj[u]) {
   if (i != v) {
     int a = sz[u], b = sz[i];
      dp[i] = dp[u];
      dp[i] = sz[u], dp[i] = sz[i];
     sz[u] = sz[i], sz[i] = n;
     dp[i] += sz[u], dp[i] += sz[i];
     reroot(i, u);
     sz[u] = a, sz[i] = b;
 }
```

hungarian.cpp

47 lines

```
struct hungarian { // declarar algo como hungarian h(n), n de
    cada lado
 int n, inf;
 vector<vector<int>> a;
 vector<int> u, v, p, way;
 hungarian(int n_) : n(n_{-}), u(n + 1), v(n + 1), p(n + 1), way(
   a = vector<vector<int>>(n, vector<int>(n));
   inf = numeric limits<int>::max();
 void add_edge(int x, int y, int c) { a[x][y] = c; }
 pair<int, vector<int>> run() {
   for (int i = 1; i <= n; i++) {
     p[0] = i;
     int j0 = 0;
     vector<int> minv(n + 1, inf);
     vector<int> used(n + 1, 0);
     qo {
       used[j0] = true;
       int i0 = p[j0], j1 = -1;
       int delta = inf;
       for (int j = 1; j \le n; j++) {
         if (!used[j]) {
           int cur = a[i0 - 1][j - 1] - u[i0] - v[j];
```

```
if (cur < minv[j]) minv[j] = cur, way[j] = j0;
            if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
        for (int j = 0; j \le n; j++) {
         if (used[j])
            u[p[j]] += delta, v[j] -= delta;
            minv[j] -= delta;
        j0 = j1;
      } while (p[j0] != 0);
      do {
       int j1 = way[j0];
        p[j0] = p[j1];
        j0 = j1;
      } while (j0);
    vector<int> ans(n);
    for (int j = 1; j \le n; j++) ans[p[j] - 1] = j - 1;
    return make_pair(-v[0], ans);
};
```

mincostmaxflow.cpp

0.4.11

14

```
namespace mcf {
struct edge {
 int to, capacity, cost, res;
int source, destiny;
vector<edge> adj[MAXN];
vector<int> dist:
vector<int> parent;
vector<int> edge_index;
vector<bool> in_queue;
void add edge(int a, int b, int c, int d) {
 adj[a].pb({b, c, d, (int)adj[b].size()});
 adj[b].pb({a, 0, -d, (int)adj[a].size() - 1});
bool dijkstra(int s) {
 dist.assign(MAXN, INF);
 parent.assign(MAXN, -1);
  edge_index.assign(MAXN, -1);
  in_queue.assign(MAXN, false);
  dist[s] = 0;
  queue<int> q;
  q.push(s);
  while (!q.empty()) {
    int u = q.front(), idx = 0;
    q.pop();
    in_queue[u] = false;
    for (auto const &v : adj[u]) {
      if (v.capacity && dist[v.to] > dist[u] + v.cost) {
        dist[v.to] = dist[u] + v.cost;
        parent[v.to] = u;
        edge_index[v.to] = idx;
        if (!in_queue[v.to]) {
         in queue[v.to] = true;
          q.push(v.to);
      idx++;
 return dist[destiny] != INF;
int get_cost() {
 int flow = 0, cost = 0;
```

while (dijkstra(source)) {

while (curr != source) {

int p = parent[curr];

int curr_flow = INF, curr = destiny;

curr flow = min(curr_flow, adj[p][edge_index[curr]].

scc articulationpoints lca topsort eulerian notes

```
capacity);
     curr = p;
    flow += curr_flow;
    cost += curr flow * dist[destinv];
    curr = destiny;
    while (curr != source) {
     int p = parent[curr];
     int res_idx = adj[p][edge_index[curr]].res;
     adj[p][edge_index[curr]].capacity -= curr_flow;
     adj[curr][res_idx].capacity += curr_flow;
      curr = p;
  return cost;
} // namespace mcf
scc.cpp
int n, m;
bool vis[MAXN];
int root[MAXN];
vector<int> order;
vector<int> roots:
vector<int> comp;
vector<vector<int>> comps;
vector<int> adi[MAXN];
vector<int> adj rev[MAXN];
vector<int> adi scc[MAXN];
void dfs(int v) {
  vis[v] = true;
  for (auto const &u : adj[v])
   if (!vis[u]) dfs(u);
  order.pb(v);
void dfs2(int v) {
  comp.pb(v);
  vis[v] = true;
  for (auto const &u : adj_rev[v])
    if (!vis[u]) dfs2(u);
//...
cin >> n >> m;
for (int i = 0; i < m; i++) {
 int a, b:
  cin >> a >> b;
  adj[a].pb(b);
  adj_rev[b].pb(a);
for (int i = 0; i < n; i++) {
 if (!vis[i]) dfs(i);
reverse(order.begin(), order.end());
memset (vis, false, sizeof (vis));
for (auto const &v : order) {
  if (!vis[v]) {
    comp.clear();
    dfs2(v);
    comps.pb(comp);
    // making condensation graph
    int r = comp.back();
    for (auto const &u : comp) root[u] = r;
    roots.push_back(r);
```

```
// making condensation graph
for (int v = 0; v < n; v++) {
 for (auto const &u : adj[v]) {
    int root_v = roots[v];
    int root_u = roots[u];
    if (root_u != root_v) adj_scc[root_v].pb(root_u);
articulationpoints.cpp
int n, m, timer;
vector<int> adj[MAXN];
bool is_cutpoint[MAXN];
int tin[MAXN]; // memset -1
int low[MAXN]; // memset -1
bool vis[MAXN];
void dfs(int v, int p) { // chama pros nao vis
 vis[v] = true;
 tin[v] = timer, low[v] = timer++;
  int childs = 0;
  for (auto const &u : adj[v]) {
   if (u == p) continue;
    if (vis[u]) {
     low[v] = min(low[v], tin[u]);
      dfs(u, v);
      low[v] = min(low[v], low[u]);
      if (low[u] >= tin[v] \&\& p != -1) is cutpoint[v] = true;
      childs++:
 if (p == -1 && childs > 1) is_cutpoint[v] = true;
lca.cpp
                                                          39 lines
int n;
vector<int> adj[MAXN];
namespace lca {
int 1, timer;
vector<int> tin, tout, depth;
vector<vector<int>> up;
void dfs(int v, int p) {
 tin[v] = ++timer;
 up[v][0] = p;
 for (int i = 1; i \le 1; i++) up[v][i] = up[up[v][i-1]][i-
 for (auto const &u : adj[v]) {
   if (p == u) continue;
   depth[u] = depth[v] + 1;
   dfs(u, v);
 tout[v] = ++timer;
bool is_ancestor(int u, int v) { return tin[u] <= tin[v] &&</pre>
    tout[u] >= tout[v]; }
int binary_lifting(int u, int v) {
 if (is_ancestor(u, v)) return u;
```

if (is_ancestor(v, u)) return v;

if (!is_ancestor(up[u][i], v)) u = up[u][i];

for (int i = 1; i >= 0; --i)

return up[u][0];

tin.resize(n);

tout.resize(n);

void init() {

```
depth.resize(n);
 timer = 0:
 l = ceil(log2(n));
 up.assign(n, vector<int>(1 + 1));
 dfs(0, 0);
int dist(int s, int v) {
 int at = binary_lifting(s, v);
 return (depth[s] + depth[v] - 2 * depth[at]);
} // namespace lca
topsort.cpp
                                                          21 lines
vector<bool> visited:
vector<int> ans:
void dfs(int v) {
 visited[v] = true;
 for (int u : adj[v]) {
   if (!visited[u]) {
     dfs(u);
 ans.push_back(v);
void topological_sort() {
 visited.assign(n, false);
 ans.clear();
 for (int i = 0; i < n; ++i) {
   if (!visited[i]) {
      dfs(i);
 reverse(ans.begin(), ans.end());
eulerian.cpp
void dfs2(int s) { // caminho euleriano
 while (a[s].size() > 0) {
   auto v = a[s].back();
   a[s].pop_back();
   if (!vis[v.sec])
     vis[v.sec] = 1;
     dfs2(v.fir);
 path.pb(s);
void dfs(int i) { // ajeitar para que todo mundo tenha grau
    par
 vis[i] = 1;
 for (auto const &j : adj[i]) {
   if (!vis[j.fir]) {
     dfs(j.fir);
     if (deg[j.fir]) {
       ans.pb(edges[j.sec]);
       deg[j.fir] ^= 1;
       deq[i] ^= 1;
 }
// se eu cham dfs(0) no final checar se o deg[0] ta safe
notes.md
## Bipartite Graph
```

- A bipartite graph is a graph that does not contain any odd-length cycles.
- ## Directed acyclic graph (DAG)
- Is a directed graph with no directed cycles.
- ## Independent Set
- Is a set of vertices in a graph, no two of which are adjacent.

 That is, it is a set S of vertices such that for every two vertices in S, there is no edge connecting the two.
- ## Clique
- Is a subset of vertices of an undirected graph such that every two distinct vertices in the clique are adjacent.
- ## Vertex Cover
- Is a set of vertices that includes at least one endpoint of every edge of the graph.
- ## Edge Cover
- Is a set of edges such that every vertex of the graph is incident to at least one edge of the set.
- ## Path Cover
- Given a directed graph G = (V, E), a path cover is a set of directed paths such that every vertex v belongs to at least one path.
- ## Koning's Theorem
- In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover.
- ## Properties
- Every tree is a bipartite graph.
- Any NxM grid is a bipartite graph.
- A set of vertices is a vertex cover if and only if its complement is an independent set.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- In bipartite graphs, the size of the minimum edge cover is equal to the size of the maximum independent set
- In bipartite graphs, the size of the minimum edge cover plus the size of the minimum vertex cover is equal to the number of vertices.
- In bipartite graphs, maximum clique size is two.
- ## Min-cut
- The smallest total weight of the edges which if removed would disconnect the source from the sink.
- ## Max-flow min-cut theorem
- In a flow network, the maximum amount of flow passing from the source to the sink is equal to the total weight of the edges in a minimum cut.
- ## Maximum flow with vertex capacities

- In other words, the amount of flow passing through a vertex cannot exceed its capacity. To find the maximum flow, we can transform the problem into the maximum flow problem by expanding the network. Each vertex v is replaced by v-in and v-out, where v-in is connected by edges going into v and v-out is connected to edges coming out from v. Then assign capacity c(v) to the edge connecting v-in and v-out
- ## Undirected edge-disjoint paths problem
- We are given an undirected graph $G=(V,\;E)$ and two vertices s and t, and we have to find the maximum number of edgedisjoint s-t paths in G.
- ## Undirected vertex-disjoint paths problem
- We are given an undirected graph $G=(V,\ E)$ and two vertices s and t, and we have to find the maximum number of vertex-disjoint (except for s and t) paths in G.
- ## Menger's theorem
- The maximum number of edge-disjoint s-t paths in an undirected graph is equal to the minimum number of edges in an s-t cut-set.
- ## Undirected vertex-disjoint paths solution
- We can construct a network N=(V,E) from G with vertex capacities, where the capacities of all vertices and all edges are 1. Then the value of the maximum flow is equal to the maximum number of independent paths from s to t.
- Given a directed acyclic graph G=(V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V. We can construct a bipartite graph G' from G. Each vertex v is replaced by v-in and v-out, where v-in is connected by edges going into v and v-out is connected to edges coming out from v. Then it can be shown that G' has a matching M of size m if and only if G has a vertex-disjoint path cover C of containing m edges and n-m paths.
- ## Minimum general path cover in directed acyclic graph (DAG)
- A general path cover is a path cover where a vertex can belong to more than one path. A minimum general path cover may be smaller than a minimum vertex-disjoint path cover. A minimum general path cover can be found almost like a minimum vertex-disjoint path cover. It suffices to add some new edges to the matching graph so that there is an edge a to b always when there is a path from a to b in the original graph.
- ## Dilworths theorem and maximum antichain
- An antichain is a set of nodes of a graph such that there is no path from any node to another node using the edges of the graph. Dilworths theorem states that in a directed acyclic graph, the size of a minimum general path cover equals the size of a maximum antichain.
- Or in other words: For a DAG G that if has edges from vertex i to vertex j and vertex j to k, then it also has a edge from vertex i to vertex k, the size of a minimum path cover is equal to the size of a maximum independent set.

- ## Maximum weighted antichain
- In this problem, each vertex has a cost a[i]. The cost of an antichain is equal to the sum of the costs of the vertices present in it. We need to find the maximum cost of a antichain. We can construct the same bipartite of the maximum antichain problem from a dag G, these edges have an infinite capacity. We also need to create a source vertex and a sink, and we need to add edges source to v-in with capacity a[v] and v-out to sink with capacity a[v]. The answer is equal to the sum of all a[i] minus the maximum flow on this network.
- ## Halls Theorem
- Halls theorem can be used to find out whether a bipartite graph has a matching that contains all left or right nodes. Assume that we want to find a matching that contains all left nodes. Let X be any set of left nodes and let f(X) be the set of their neighbors. According to Halls theorem, a matching that contains all left nodes exists exactly when for
- each X, the condition $|X| \le |f(X)|$ holds.
- ## Extra (Getting Confidence Trick)
- If you need to maximize a number x = (a * b * c * ...), then
 you can write it as $x = (e^{\log(a)} * e^{\log(b)} * e^{\log(c)} *$...), and then the number is $x = e^{\log(a)} + \log(b) + \log(a)$ (c) + ...), and the problem now becomes a problem of
 maximizing the sum of $(\log(a) + \log(b) + \log(c) + ...).<p$ />

47 lines

Use exp() and log() C++ functions :)

$\underline{\text{Geometry}} \ (5)$

convexhull.cpp

```
struct point {
 int x, y, id;
 point(int x, int y, int id) : x(x), y(y), id(id) {}
 point() {}
 point operator-(point const &o) const { return {x - o.x, y -
      o.y, -1; }
 bool operator<(point const &o) const {</pre>
   if (x == o.x) return y < o.y;
   return x < o.x;
 int operator^(point const &o) const { return x * o.y - y * o.
int ccw(point const &a, point const &b, point const &x) {
 auto p = (b - a) ^ (x - a);
 return (p > 0) - (p < 0);
vector<point> convex_hull(vector<point> P) // sem colineares
 sort(P.begin(), P.end());
 vector<point> L, U;
 for (auto p : P) {
   while (L.size() \geq= 2 && ccw(L.end()[-2], L.end()[-1], p) ==
         -1) L.pop_back();
   L.push_back(p);
 reverse(P.begin(), P.end());
 for (auto p : P) {
```

convexhullpointlocation linetrick polygonarea geometria

```
while (U.size() >= 2 \&\& ccw(U.end()[-2], U.end()[-1], p) ==
         -1) U.pop back();
   U.push back(p);
  L.insert(L.end(), U.begin(), U.end() - 1);
 return L:
vector<point> convex_hull_no_collinears(vector<point> P) //
    com colineares
  sort(P.begin(), P.end());
  vector<point> L, U;
  for (auto p : P) {
   while (L.size() >= 2 && ccw(L.end()[-2], L.end()[-1], p) <=
         0) L.pop_back();
   L.push_back(p);
  reverse(P.begin(), P.end());
  for (auto p : P) {
   while (U.size() \ge 2 \&\& ccw(U.end()[-2], U.end()[-1], p) \le
         0) U.pop_back();
   U.push_back(p);
  L.insert(L.end(), U.begin(), U.end() - 1);
  return L;
```

convexhullpointlocation.cpp

44 lines

```
struct pt {
  int x, v;
  pt operator+(pt p) { return {x + p.x, y + p.y}; }
  pt operator-(pt p) { return {x - p.x, y - p.y}; }
  bool operator==(pt p) { return (x == p.x \&\& y == p.y); }
  int cross(pt p) { return x * p.y - y * p.x; }
  int cross(pt a, pt b) { return (a - *this).cross(b - *this);
  int dot(pt p) { return x * p.x + y * p.y; }
bool cmp_x(pt a, pt b) {
  if (a.x != b.x) return a.x < b.x;
  return a.y < b.y;
// acha o convex hull
vector<pt> convex_hull(vector<pt> pts) {
  if (pts.size() <= 1) return pts;
  sort(pts.begin(), pts.end(), cmp_x);
  vector<pt> h(pts.size() + 1);
  int s = 0, t = 0;
  for (int it = 2; it--; s = --t, reverse(pts.begin(), pts.end
       ())) {
    for (auto const &p : pts) {
      while (t \ge s + 2 \&\& h[t - 2].cross(h[t - 1], p) \le 0) t
     h[t++] = p;
    }
  return \{h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])\};
int sqn(int x) { return (x > 0) - (x < 0); }
int side_of(pt s, pt e, pt p) { return sgn(s.cross(e, p)); }
bool on_segment(pt s, pt e, pt p) { return p.cross(s, e) == 0
     && (s - p).dot(e - p) <= 0; }
// retorna se o ponto p esta dentro ou nao do convex hull l
// caso strict = true, entao considera true se tiver na borda
// caso strict = false, entao considera false se tiver na borda
bool is hull(vector<pt> &1, pt p, bool strict = true) {
  int a = 1, b = 1.size() - 1, r = !strict;
  if (1.size() < 3) return r && on segment(1[0], 1.back(), p);
```

```
if (side_of(1[0], 1[a], 1[b]) > 0) swap(a, b);
 if (side_of(1[0], 1[a], p) >= r || side_of(1[0], 1[b], p) <=
      -r) return false;
 while (abs(a - b) > 1) {
   int c = (a + b) / 2;
    (side_of(1[0], 1[c], p) > 0 ? b : a) = c;
 return sqn(l[a].cross(l[b], p)) < r;
linetrick.cpp
                                                            13 lines
pi get_line(pi x, pi y) { // um jeito normalizado de
    representar a reta entre 2 pontos
 int xx = x.fir - y.fir;
 int yy = x.sec - y.sec;
 int g = \underline{gcd(abs(xx), abs(yy))};
 if (g != 0) {
   xx /= g, yy /= g;
 if (xx < 0) {
   xx \star = -1;
   yy \star = -1;
 return {xx, yy};
polygonarea.cpp
                                                            19 lines
```

```
double area(vector<pi> fig) {
   double res = 0;
   for (unsigned i = 0; i < fig.size(); i++) {
      pi p = i ? fig[i - 1] : fig.back();
      pi q = fig[i];
      res += (p.fir - q.fir) * (p.sec + q.sec);
   }
   return fabs(res) / 2;
}
int cross(pi a, pi b) { return a.fir * b.sec - a.sec * b.fir; }
double area2(vector<pi> fig) {
   double res = 0;
   for (unsigned i = 0; i < fig.size(); i++) {
      pi p = i ? fig[i - 1] : fig.back();
      pi q = fig[i];
      res += cross(p, q);
   }
   return fabs(res) / 2;
}</pre>
```

geometria.cpp

207 lines

```
typedef double ld;
const ld DINF = 1e18;
const ld pi = acos(-1.0);
const 1d eps = 1e-9;
\#define sq(x) ((x) * (x))
bool eq(ld a, ld b) { return abs(a - b) <= eps; }
struct pt { // ponto
 ld x, y;
  pt (ld x_{-} = 0, ld y_{-} = 0) : x(x_{-}), y(y_{-}) {}
  bool operator<(const pt p) const {
    if (!eq(x, p.x)) return x < p.x;
    if (!eq(y, p.y)) return y < p.y;
    return 0;
  bool operator==(const pt p) const { return eq(x, p.x) and eq(
       y, p.y); }
  pt operator+(const pt p) const { return pt(x + p.x, y + p.y);
```

```
pt operator-(const pt p) const { return pt(x - p.x, y - p.y);
  pt operator*(const ld c) const { return pt(x * c, y * c); }
 pt operator/(const ld c) const { return pt(x / c, y / c); }
  ld operator*(const pt p) const { return x * p.x + y * p.y; }
 ld operator^(const pt p) const { return x * p.y - y * p.x; }
 friend istream& operator>>(istream& in, pt& p) { return in >>
       p.x >> p.y; }
};
struct line { // reta
 pt p, q;
 line() {}
 line(pt p_, pt q_) : p(p_), q(q_) {}
 friend istream& operator>>(istream& in, line& r) { return in
      >> r.p >> r.q; }
// PONTO & VETOR
ld dist(pt p, pt q) { // distancia
 return hypot(p.y - q.y, p.x - q.x);
ld dist2(pt p, pt q) { // quadrado da distancia
 return sq(p.x - q.x) + sq(p.y - q.y);
ld norm(pt v) { // norma do vetor
 return dist(pt(0, 0), v);
ld angle(pt v) { // angulo do vetor com o eixo x
 1d ang = atan2(v.v, v.x);
 if (ang < 0) ang += 2 * pi;
 return ang;
ld sarea(pt p, pt q, pt r) { // area com sinal
 return ((q - p) ^ (r - q)) / 2;
bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
 return eq(sarea(p, q, r), 0);
bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
 return sarea(p, q, r) > eps;
pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
 return pt(p.x * cos(th) - p.y * sin(th), p.x * sin(th) + p.y
      * cos(th));
pt rotate90(pt p) { // rotaciona 90 graus
 return pt(-p.y, p.x);
// RETA
bool isvert(line r) { // se r eh vertical
 return eq(r.p.x, r.q.x);
bool isinseq(pt p, line r) { // se p pertence ao seq de r
 pt a = r.p - p, b = r.q - p;
 return eq((a ^ b), 0) and (a * b) < eps;
ld get_t(pt v, line r) { // retorna t tal que t*v pertence a
 return (r.p ^ r.q) / ((r.p - r.q) ^ v);
pt proj(pt p, line r) { // projecao do ponto p na reta r
 if (r.p == r.q) return r.p;
 r.q = r.q - r.p;
 p = p - r.p;
 pt proj = r.q * ((p * r.q) / (r.q * r.q));
 return proj + r.p;
pt inter(line r, line s) { // r inter s
 if (eq((r.p - r.q) ^ (s.p - s.q), 0)) return pt(DINF, DINF);
 r.q = r.q - r.p, s.p = s.p - r.p, s.q = s.q - r.p;
```

```
int qt = 0;
  for (int i = 0; i < v.size(); i++) {
   if (p == v[i]) return 2;
    int j = (i + 1) % v.size();
    if (eq(p.y, v[i].y) and eq(p.y, v[j].y)) {
     if ((v[i] - p) * (v[j] - p) < eps) return 2;
     continue;
    bool baixo = v[i].y + eps < p.y;</pre>
    if (baixo == (v[j].y + eps < p.y)) continue;
    auto t = (p - v[i]) ^ (v[j] - v[i]);
    if (eq(t, 0)) return 2;
    if (baixo == (t > eps)) qt += baixo ? 1 : -1;
 return qt != 0;
bool interpol(vector<pt> v1, vector<pt> v2) { // se dois
    poligonos se intersectam - O(n*m)
  int n = v1.size(), m = v2.size();
 for (int i = 0; i < n; i++)
    if (inpol(v2, v1[i])) return 1;
 for (int i = 0; i < n; i++)
   if (inpol(v1, v2[i])) return 1;
 for (int i = 0; i < n; i++)
   for (int j = 0; j < m; j++)
      if (interseg(line(v1[i], v1[(i + 1) % n]), line(v2[i], v2
          [(j + 1) % m]))) return 1;
 return 0;
ld distpol(vector<pt> v1, vector<pt> v2) { // distancia entre
  if (interpol(v1, v2)) return 0;
 ld ret = DINF;
  for (int i = 0; i < v1.size(); i++)
    for (int j = 0; j < v2.size(); j++)
      ret = min(ret, distseg(line(v1[i], v1[(i + 1) % v1.size()
          ]), line(v2[j], v2[(j + 1) % v2.size()])));
 return ret;
// CIRCUNFERENCIA
pt getcenter(pt a, pt b, pt c) { // centro da circunf dado 3
 b = (a + b) / 2;
 c = (a + c) / 2;
 return inter(line(b, b + rotate90(a - b)), line(c, c +
       rotate90(a - c)));
vector<pt> circ_line_inter(pt a, pt b, pt c, ld r) { //
    intersecao da circunf (c, r) e reta ab
 vector<pt> ret;
 b = b - a, a = a - c;
 1d A = b * b;
 1d B = a * b;
 1d C = a * a - r * r;
 1d D = B * B - A * C;
 if (D < -eps) return ret;
 ret.push_back(c + a + b * (-B + sqrt(D + eps)) / A);
 if (D > eps) ret.push_back(c + a + b * (-B - sqrt(D)) / A);
 return ret;
vector<pt> circ_inter(pt a, pt b, ld r, ld R) { // intersecao
    da circunf (a, r) e (b, R)
 vector<pt> ret;
 1d d = dist(a, b);
 if (d > r + R \text{ or } d + min(r, R) < max(r, R)) return ret;
 1d x = (d * d - R * R + r * r) / (2 * d);
 1d y = sqrt(r * r - x * x);
  pt v = (b - a) / d;
```

ret.push_back(a + v * x + rotate90(v) * y);

```
if (y > 0) ret.push back(a + y * x - rotate 90(y) * y);
return ret:
```

18

geometriaint.cpp

```
105 lines
\#define sq(x) ((x) * (11)(x))
struct pt { // ponto
 int x, y;
 pt(int x_{-} = 0, int y_{-} = 0) : x(x_{-}), y(y_{-}) {}
 bool operator<(const pt p) const {
   if (x != p.x) return x < p.x;
    return y < p.y;
 bool operator==(const pt p) const { return x == p.x and y ==
  pt operator+(const pt p) const { return pt(x + p.x, y + p.y);
  pt operator-(const pt p) const { return pt(x - p.x, y - p.y);
 pt operator*(const int c) const { return pt(x * c, y * c); }
  11 operator*(const pt p) const { return x * (11)p.x + y * (11)
  11 operator^(const pt p) const { return x * (11)p.y - y * (11
  friend istream& operator>>(istream& in, pt& p) { return in >>
       p.x >> p.y; }
struct line { // reta
 pt p, q;
 line() {}
 line(pt p_, pt q_) : p(p_), q(q_) {}
 friend istream& operator>>(istream& in, line& r) { return in
       >> r.p >> r.a; }
// PONTO & VETOR
11 dist2(pt p, pt q) { // quadrado da distancia
 return sq(p.x - q.x) + sq(p.y - q.y);
ll sarea2(pt p, pt q, pt r) { // 2 * area com sinal
 return (q - p) ^ (r - q);
bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
 return sarea2(p, q, r) == 0;
bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
 return sarea2(p, q, r) > 0;
int quad(pt p) { // quadrante de um ponto
 return (p.x < 0) ^3 * (p.y < 0);
bool compare_angle(pt p, pt q) { // retorna se ang(p) < ang(q)</pre>
 if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
 return ccw(q, pt(0, 0), p);
pt rotate90(pt p) { // rotaciona 90 graus
 return pt(-p.y, p.x);
bool isinseg(pt p, line r) { // se p pertence ao seg de r
 pt a = r.p - p, b = r.q - p;
 return (a ^{\circ} b) == 0 and (a * b) <= 0;
bool interseg(line r, line s) { // se o seg de r intersecta o
 if (isinseg(r.p, s) or isinseg(r.q, s) or isinseg(s.p, r) or
       isinseq(s.q, r)) return 1;
```

```
return ccw(r.p, r.q, s.p) \mathrel{!=} ccw(r.p, r.q, s.q) and ccw(s.p, r.q, s.q)
      s.q, r.p) != ccw(s.p, s.q, r.q);
int segpoints(line r) { // numero de pontos inteiros no
 return 1 + \underline{gcd(abs(r.p.x - r.q.x), abs(r.p.y - r.q.y))};
double get_t(pt v, line r) { // retorna t tal que t*v pertence
  return (r.p ^ r.q) / (double) ((r.p - r.q) ^ v);
// POLIGONO
// quadrado da distancia entre os retangulos a e b (lados
    paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior
    direito)
11 dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
  int hor = 0, vert = 0;
  if (a.second.x < b.first.x)
   hor = b.first.x - a.second.x;
  else if (b.second.x < a.first.x)
  hor = a.first.x - b.second.x;
  if (a.second.v < b.first.v)</pre>
   vert = b.first.y - a.second.y;
  else if (b.second.y < a.first.y)</pre>
   vert = a.first.y - b.second.y;
  return sq(hor) + sq(vert);
11 polarea2(vector<pt> v) { // 2 * area do poligono
  for (int i = 0; i < v.size(); i++) ret += sarea2(pt(0, 0), v[
      i], v[(i + 1) % v.size()]);
  return abs(ret);
// se o ponto ta dentro do poligono: retorna 0 se ta fora.
// 1 se ta no interior e 2 se ta na borda
int inpol(vector<pt>& v, pt p) { // O(n)
  int qt = 0;
  for (int i = 0; i < v.size(); i++) {
    if (p == v[i]) return 2;
   int j = (i + 1) % v.size();
   if (p.v == v[i].v \text{ and } p.v == v[j].v) {
     if ((v[i] - p) * (v[j] - p) \le 0) return 2;
     continue;
   bool baixo = v[i].y < p.y;
   if (baixo == (v[j].y < p.y)) continue;
    auto t = (p - v[i]) ^ (v[j] - v[i]);
   if (!t) return 2;
   if (baixo == (t > 0)) qt += baixo ? 1 : -1;
 return qt != 0;
11 interior_points(vector<pt> v) { // pontos inteiros dentro
    de um poligono simples
  11 b = 0;
  for (int i = 0; i < v.size(); i++) b += segpoints(line(v[i],
      v(i + 1) % v.size())) - 1;
  return (polarea2(v) - b) / 2 + 1;
```

geometria3d.cpp

132 lines

typedef double ld;
const ld DINF = le18;
const ld eps = le-9;
#define sq(x) ((x) * (x))
bool eq(ld a, ld b) { return abs(a - b) <= eps; }
struct pt { // ponto</pre>

```
ld x, y, z;
 pt(ld x_{-} = 0, ld y_{-} = 0, ld z_{-} = 0) : x(x_{-}), y(y_{-}), z(z_{-}) {}
 bool operator<(const pt p) const {
   if (!eq(x, p.x)) return x < p.x;
    if (!eq(y, p.y)) return y < p.y;
   if (!eq(z, p.z)) return z < p.z;
    return 0;
 bool operator==(const pt p) const { return eq(x, p.x) and eq(
      y, p.y) and eq(z, p.z); }
 pt operator+(const pt p) const { return pt(x + p.x, y + p.y,
      z + p.z); }
 pt operator-(const pt p) const { return pt(x - p.x, y - p.y,
       z - p.z); }
 pt operator*(const ld c) const { return pt(x * c, y * c, z *
 pt operator/(const ld c) const { return pt(x / c, y / c, z /
 ld operator*(const pt p) const { return x * p.x + y * p.y + z
 pt operator^(const pt p) const { return pt(y * p.z - z * p.y,
       z * p.x - x * p.z, x * p.y - y * p.x); }
 friend istream& operator>>(istream& in, pt& p) { return in >>
       p.x >> p.y >> p.z; }
struct line { // reta
 pt p, q;
 line() {}
 line(pt p_, pt q_) : p(p_), q(q_) {}
 friend istream& operator>>(istream& in, line& r) { return in
      >> r.p >> r.q; }
struct plane {
                   // plano
 array<pt, 3> p; // pontos que definem o plano
 array<ld, 4> eq; // equacao do plano
 plane() {}
 plane(pt p_, pt q_, pt r_) : p({p_, q_, r_}) { build(); }
 friend istream& operator>>(istream& in, plane& P) {
    return in >> P.p[0] >> P.p[1] >> P.p[2];
   P.build();
 void build() {
   pt dir = (p[1] - p[0]) ^ (p[2] - p[0]);
    eq = \{dir.x, dir.y, dir.z, dir * p[0] * (-1)\};
}:
// converte de coordenadas polares para cartesianas
// (angulos devem estar em radianos)
// phi eh o angulo com o eixo z (cima) theta eh o angulo de
    rotacao ao redor de z
pt convert(ld rho, ld th, ld phi) { return pt(sin(phi) * cos(th
    ), sin(phi) * sin(th), cos(phi)) * rho; }
// projecao do ponto p na reta r
pt proj(pt p, line r) {
 if (r.p == r.q) return r.p;
 r.q = r.q - r.p;
 p = p - r.p;
 pt proj = r.q * ((p * r.q) / (r.q * r.q));
 return proj + r.p;
// projecao do ponto p no plano P
pt proj(pt p, plane P) {
 p = p - P.p[0], P.p[1] = P.p[1] - P.p[0], P.p[2] = P.p[2] - P
       .p[0];
 pt norm = P.p[1] ^ P.p[2];
 pt proj = p - (norm * (norm * p) / (norm * norm));
 return proj + P.p[0];
```

```
// distancia
ld dist(pt a, pt b) { return sqrt(sq(a.x - b.x) + sq(a.y - b.y)
      + sq(a.z - b.z)); }
// distancia ponto reta
ld distline(pt p, line r) { return dist(p, proj(p, r)); }
// distancia de ponto para segmento
ld distseg(pt p, line r) {
 if ((r.q - r.p) * (p - r.p) < 0) return dist(r.p, p);
 if ((r.p - r.q) * (p - r.q) < 0) return dist(r.q, p);
  return distline(p, r);
// distancia de ponto a plano com sinal
ld sdist(pt p, plane P) { return P.eq[0] * p.x + P.eq[1] * p.y
     + P.eq[2] * p.z + P.eq[3]; }
// distancia de ponto a plano
ld distplane(pt p, plane P) { return abs(sdist(p, P)); }
// se ponto pertence a reta
bool isinseq(pt p, line r) { return eq(distseq(p, r), 0); }
// se ponto pertence ao triangulo definido por P.p
bool isinpol(pt p, vector<pt> v) {
  assert(v.size() >= 3);
  pt norm = (v[1] - v[0]) ^ (v[2] - v[1]);
  bool inside = true;
  int sign = -1;
  for (int i = 0; i < v.size(); i++) {
    line r(v[(i + 1) % 3], v[i]);
    if (isinseq(p, r)) return true;
    pt ar = v(i + 1) % 3 - v[i];
    if (sign == -1)
      sign = ((ar ^ (p - v[i])) * norm > 0);
    else if (((ar ^{\circ} (p - v[i])) * norm > 0) != sign)
      inside = false;
  return inside;
// distancia de ponto ate poligono
ld distpol(pt p, vector<pt> v) {
  pt p2 = proj(p, plane(v[0], v[1], v[2]);
  if (isinpol(p2, v)) return dist(p, p2);
  ld ret = DINF;
  for (int i = 0; i < v.size(); i++) {
    int j = (i + 1) % v.size();
    ret = min(ret, distseq(p, line(v[i], v[j])));
  return ret;
// intersecao de plano e segmento
// BOTH = o segmento esta no plano
// ONE = um dos pontos do segmento esta no plano
// PARAL = segmento paralelo ao plano
// CONCOR = segmento concorrente ao plano
enum RETCODE { BOTH, ONE, PARAL, CONCOR };
pair<RETCODE, pt> intersect(plane P, line r) {
  1d d1 = sdist(r.p, P);
  1d d2 = sdist(r.q, P);
  if (eq(d1, 0) and eq(d2, 0)) return pair(BOTH, r.p);
  if (eq(d1, 0)) return pair(ONE, r.p);
  if (eq(d2, 0)) return pair(ONE, r.q);
  if ((d1 > 0 \text{ and } d2 > 0) \text{ or } (d1 < 0 \text{ and } d2 < 0)) {}
    if (eq(d1 - d2, 0)) return pair(PARAL, pt());
    return pair(CONCOR, pt());
  1d frac = d1 / (d1 - d2);
  pt res = r.p + ((r.q - r.p) * frac);
  return pair (ONE, res);
// rotaciona p ao redor do eixo u por um angulo a
pt rotate(pt p, pt u, ld a) {
  u = u / dist(u, pt());
```

19

kdtree halfplaneintersection ahocorasick

```
kdtree.cpp
struct pt {
  int x, y, id;
  pt() {}
  pt(int xx, int yy) { x = xx, y = yy; }
  pt operator-(pt p) const { return pt(x - p.x, y - p.y); }
 bool operator<(pt p) const { return x < p.x; }</pre>
 int dist() const { return x * x + y * y; }
bool on_x(const pt &a, const pt &b) { return a.x < b.x; }
bool on_y(const pt &a, const pt &b) { return a.y < b.y; }
struct node {
  pt pp;
  int id;
  int x0 = inf, x1 = -inf, y0 = inf, y1 = -inf;
  node *first = 0, *second = 0;
  int distance (const pt &p) {
   int x = (p.x < x0 ? x0 : p.x > x1 ? x1 : p.x);
   int y = (p.y < y0 ? y0 : p.y > y1 ? y1 : p.y);
   return (pt(x, y) - p).dist();
  node(vector<pt> &&vp) : pp(vp[0]) {
    for (pt p : vp) {
     x0 = min(x0, p.x);
     x1 = max(x1, p.x);
     y0 = min(y0, p.y);
     v1 = max(v1, p.v);
    if (vp.size() > 1) {
     sort(vp.begin(), vp.end(), x1 - x0 >= y1 - y0 ? on_x :
          on_y);
     int half = vp.size() / 2;
     first = new node({vp.begin(), vp.begin() + half});
     second = new node({vp.begin() + half, vp.end()});
};
struct kd tree {
 node *root:
  kd_tree(const vector<pt> &vp) : root(new node({vp.begin(), vp
      .end()})) {}
  pi search(node *n, const pt &p) {
   if (!n->first) {
     if (n->pp.x == p.x && n->pp.y == p.y) return make_pair(
          inf, n->pp.id); // distancia infinita pra pontos
          iguais
     return make_pair((p - n->pp).dist(), n->pp.id);
   node *f = n->first, *s = n->second;
    int bfirst = f->distance(p), bsec = s->distance(p);
   if (bfirst > bsec) swap(bsec, bfirst), swap(f, s);
    auto best = search(f, p);
   if (bsec < best.first || (!f->first)) best = min(best,
        search(s, p));
    return hest:
 pi nearest(const pt &p) { return search(root, p); }
```

return $u * (u * p) + (u ^ p ^ u) * cos(a) + (u ^ p) * sin(a);$

```
halfplaneintersection.cpp
```

const long double eps = 1e-9;

```
const long double inf = 1e9;
struct pt {
  long double x, y;
```

75 lines

```
pt(long double x = 0, long double y = 0) : x(x), y(y) {}
 friend pt operator+(pt p, pt q) { return pt(p.x + q.x, p.y +
      q.y); }
 friend pt operator-(pt p, pt q) { return pt(p.x - q.x, p.y -
      q.y); }
 friend pt operator*(pt p, long double k) { return pt(p.x * k,
       p.y * k); }
 friend long double dot(pt p, pt q) { return p.x * q.x + p.y *
 friend long double cross(pt p, pt q) { return p.x * q.y - p.y
struct halfplane {
 pt p, pq;
 long double angle;
 halfplane() {}
 halfplane(pt a, pt b) : p(a), pq(b - a) { angle = atan21(pq.y
      , pq.x); }
 bool out(const pt &r) { return cross(pq, r - p) < -eps; }
 bool operator<(halfplane e) const { return angle < e.angle; }</pre>
 friend pt inter(halfplane s, halfplane t) {
   long double alpha = cross((t.p - s.p), t.pq) / cross(s.pq,
    return s.p + (s.pq * alpha);
};
vector<pt> hp_intersect(vector<halfplane> &h) {
 pt box[4] = {pt(inf, inf), pt(-inf, inf), pt(-inf, -inf), pt(}
      inf, -inf)}; // Bounding box in CCW order
 for (int i = 0; i < 4; i++) {
   halfplane aux(box[i], box[(i + 1) % 4]);
   h.pb(aux);
 sort(h.begin(), h.end());
 deque<halfplane> dq;
 int len = 0;
 for (int i = 0; i < h.size(); i++) {
    while (len > 1 && h[i].out(inter(dq[len - 1], dq[len - 2]))
        ) {
      dq.pop_back();
    while (len > 1 && h[i].out(inter(dq[0], dq[1]))) {
     dq.pop front();
    if (len > 0 \&\& fabsl(cross(h[i].pq, dq[len - 1].pq)) < eps)
      if (dot(h[i].pq, dq[len - 1].pq) < 0.0) {
       return vector<pt>();
      if (h[i].out(dq[len - 1].p)) {
       dq.pop_back();
       --len;
     } else {
       continue:
    dq.push_back(h[i]);
    ++len:
 while (len > 2 && dq[0].out(inter(dq[len - 1], dq[len - 2])))
    dq.pop_back();
    --len:
 while (len > 2 && dq[len - 1].out(inter(dq[0], dq[1]))) {
   dq.pop_front();
    --len;
```

```
if (len < 3) {
   return vector<pt>();
 vector<pt> ret(len);
 for (int i = 0; i + 1 < len; i++) {
   ret[i] = inter(dq[i], dq[i + 1]);
 ret.back() = inter(dq[len - 1], dq[0]);
 return ret:
// half-plane - regiao planar que consiste de todos os pontos
    que estao de um lado de uma reta
// e dai pros half-planes, considerando que eh a regiao da
    esquerda em relacao ao vetor de direcao
```

Strings (6)

ahocorasick.cpp

78 lines

```
namespace aho {
int go (int v, char ch);
const int K = 26; // tamanho do alfabeto
struct trie {
                       // char correspondente ao no atual
 char me;
  int qo[K];
                       // proximo vertice que eu devo ir
      estando em um estado (v, c)
                       // proximo vertice da trie
 int down[K];
  int is leaf = 0;
                      // se o vertice atual da trie eh uma
       folha (fim de uma ou mais strings)
 int parent = -1:
                      // no ancestral do no atual
                       // link de sufixo do no atual (outro no
 int link = -1;
      com o maior matching de sufixo)
  int exit link = -1; // folha mais proxima que pode ser
      alcancada a partir de v usando links de sufixo
 trie(int p = -1, char ch = '$') : parent(p), me(ch) {
    fill (begin (go), end (go), -1);
    fill (begin (down), end (down), -1);
};
vector<trie> ac;
void init() // criar a raiz da trie
 ac.resize(1);
void add_string(string s) // adicionar string na trie
 int v = 0:
 for (auto const &ch : s) {
   int c = ch - 'a';
   if (ac[v].down[c] == -1) {
     ac[v].down[c] = ac.size();
     ac.emplace_back(v, ch);
   v = ac[v].down[c];
 ac[v].is_leaf++;
int get_link(int v) // pegar o suffix link saindo de v
 if (ac[v].link == -1) ac[v].link = (!v || !ac[v].parent) ? 0
      : go(get_link(ac[v].parent), ac[v].me);
 return ac[v].link;
int go(int v, char ch) // proximo estado saindo do estado(v,
 int c = ch - 'a';
```

kmp manacher minsuffix hashing suffixarray

```
if (ac[v].go[c] == -1) {
   if (ac[v].down[c] != -1)
     ac[v].go[c] = ac[v].down[c];
     ac[v].go[c] = (!v) ? 0 : go(get_link(v), ch);
 return ac[v].go[c];
int get_exit_link(int v) // suffix link mais proximo de v que
    seia uma folha
 if (ac[v].exit link == -1) {
   int curr = get_link(v);
   if (!v || !curr)
     ac[v].exit_link = 0;
   else if (ac[curr].is_leaf)
     ac[v].exit link = curr;
     ac[v].exit_link = get_exit_link(curr);
 return ac[v].exit_link;
int query(string s) // query O(n + ans)
 int ans = 0, curr = 0, at;
 for (auto const &i : s) {
   curr = go(curr, i);
   ans += ac[curr].is_leaf;
   at = get_exit_link(curr);
   while (at) {
     ans += ac[at].is_leaf;
     at = get_exit_link(at);
 return ans:
} // namespace aho
```

kmp.cpp

27 lines

```
string s;
int n. m:
string a, b;
int c[MAXN][26];
vector<int> kmp(string &s) {
 int n = s.size();
  vector<int> p(n);
  for (int i = 1; i < n; i++) {
   int j = p[i - 1];
   while (j > 0 \&\& s[i] != s[j]) j = p[j-1];
   if (s[i] == s[j]) j++;
   p[i] = j;
  return p;
void compute(string s) {
  s.pb('*');
  vector < int > p = kmp(s);
  for (int i = 0; i < s.size(); i++) {
    for (int cc = 0; cc < 26; cc++) {
     int j = i;
     while (j > 0 \&\& 'a' + cc != s[j]) j = p[j - 1];
     if ('a' + cc == s[j]) j++;
     c[i][cc] = j;
```

```
manacher.cpp
                                                           22 lines
vector<int> d1;
vector<int> d2;
void manacher(string s) {
 dl.resize(s.size());
 d2.resize(s.size());
 int 1 = 0, r = -1;
 for (int i = 0; i < s.size(); i++) {
    int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
    while (0 \le i - k \&\& i + k \le s.size() \&\& s[i - k] == s[i + k]
        k]) k++;
   d1[i] = k:
   k = k - 1;
   if (i + k > r) l = i - k, r = i + k;
 1 = 0, r = -1;
 for (int i = 0; i < s.size(); i++) {
   int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1);
   while (0 \le i - k - 1 \&\& i + k \le s.size() \&\& s[i - k - 1]
        == s[i + k]) k++;
   d2[i] = k;
   k = k - 1;
   if (i + k > r) 1 = i - k - 1, r = i + k;
```

minsuffix.cpp

28 lines

```
int max suffix(string s, bool mi = false) {
 s.push_back(*min_element(s.begin(), s.end()) - 1);
 int ans = 0;
 for (int i = 1; i < s.size(); i++) {
   int j = 0;
   while (ans + j < i and s[i + j] == s[ans + j]) j++;
   if (s[i + j] > s[ans + j]) {
     if (!mi or i != s.size() - 2) ans = i;
   } else if (i)
     i += j - 1;
 return ans;
int min suffix(string s) {
 for (auto &i : s) i \star = -1;
 s.push_back(*max_element(s.begin(), s.end()) + 1);
 return max_suffix(s, true);
int max_cyclic_shift(string s) {
 int n = s.size();
 for (int i = 0; i < n; i++) s.pb(s[i]);
 return max_suffix(s);
int min_cyclic_shift(string s) {
 for (auto &i : s) i \star = -1;
 return max_cyclic_shift(s);
// retorna a posicao de inicio menor/maior sufixo/shift de uma
    string
```

hashing.cpp

29 lines

continue;

```
int ans = 11 * 12 + (h >> 1) + ((h & 1) << 60) + (m >> 31) +
       ((m & GET31) << 30) + 1;
  ans = (ans \& MOD) + (ans >> 61), ans = (ans \& MOD) + (ans >>
      61):
 return ans - 1:
mt19937 64 rng(chrono::steady clock::now().time since epoch().
    count());
int uniform(int 1, int r) {
 uniform int distribution<int> uid(l, r);
 return uid(rng);
struct string_hashing {
 vector<int> h, p;
 string_hashing() {}
 string_hashing(string s) : h(s.size()), p(s.size()) {
    p[0] = 1, h[0] = s[0];
    for (int i = 1; i < s.size(); i++) p[i] = mulmod(p[i - 1],
        P), h[i] = (mulmod(h[i - 1], P) + s[i]) % MOD;
  int get(int l, int r) {
   int hash = h[r] - (1 ? mulmod(h[1 - 1], p[r - 1 + 1]) : 0);
    return hash < 0 ? hash + MOD : hash;
 int append(int h, int hb, int blen) { return (hb + mulmod(h,
      p[blen])) % MOD; }
// lembrar do P = uniform(256, MOD - 1);
suffixarray.cpp
                                                          44 lines
```

```
vector<int> suffix_array(string s) {
    s += "$"; // menor do que todos os chars da string st
```

```
int n = s.size(), N = max(n + 1, 26011);
      vector<int> sa(n), ra(n);
      for (int i = 0; i < n; i++) {
             sa[i] = i, ra[i] = s[i];
      for (int k = 0; k < n; k ? k *= 2 : k++) {
             vector<int> nsa(sa), nra(n), cnt(N);
             for (int i = 0; i < n; i++) {
                   nsa[i] = (nsa[i] - k + n) % n;
                    cnt[ra[i]]++;
              for (int i = 1; i < N; i++) {
                    cnt[i] += cnt[i - 1];
              for (int i = n - 1; i + 1; i--) {
                    sa[--cnt[ra[nsa[i]]]] = nsa[i];
              for (int i = 1, r = 0; i < n; i++) {
                    nra[sa[i]] = r += (ra[sa[i]] != ra[sa[i - 1]] || ra[(sa[i - 1])] || 
                                   ] + k) % n] != ra[(sa[i - 1] + k) % n]);
             ra = nra;
             if (ra[sa[n - 1]] == n - 1) break;
     return vector<int>(sa.begin() + 1, sa.end());
vector<int> kasai(string s, vector<int> sa) {
     int n = s.size(), k = 0;
      vector<int> ra(n), lcp(n);
      for (int i = 0; i < n; i++) {
            ra[sa[i]] = i;
      for (int i = 0; i < n; i++, k -= !!k) {
            if (ra[i] == n - 1) {
                  k = 0;
```

```
UFAL
    int j = sa[ra[i] + 1];
    while (i + k < n \text{ and } j + k < n \text{ and } s[i + k] == s[j + k]) k
    lcp[ra[i]] = k;
  return lcp;
zfunction.cpp
                                                                 12 lines
  int n = s.size();
  vector<int> z(n);
```

```
vector<int> z_function(string &s) {
  z[0] = n;
  for (int i = 1, l = 0, r = 0; i < n; i++) {
   if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
   while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]]) z[i]++;
   if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
  return z;
// z[i] = o tamanho de lcp(s, s.substr(i, n - i))
```

suffixautomaton.cpp

90 lines

```
namespace sa {
struct state {
  int len, suf link;
 map<char, int> nxt;
vector<int> term;
state st[2 * MAXN];
int dp[2 * MAXN];
int sz. last:
void init() {
  memset(dp, -1, sizeof(dp));
  st[0].len = 0;
  st[0].suf link = -1;
  sz++;
  last = 0;
void get_link(int curr, int p, char c) {
  while (p != -1 \&\& !st[p].nxt.count(c)) {
    st[p].nxt[c] = curr;
    p = st[p].suf_link;
  if (p == -1) {
   st[curr].suf_link = 0;
    return;
  int q = st[p].nxt[c];
  if (st[p].len + 1 == st[q].len) {
    st[curr].suf_link = q;
    return;
  int clone = sz;
  sz++;
  st[clone].len = st[p].len + 1;
  st[clone].nxt = st[q].nxt;
  st[clone].suf_link = st[q].suf_link;
  while (p != -1 \&\& st[p].nxt[c] == q) {
    st[p].nxt[c] = clone;
    p = st[p].suf_link;
  st[q].suf_link = clone;
  st[curr].suf_link = clone;
void build(string &s) {
```

```
for (auto const &c : s) {
   int curr = sz;
   st[curr].len = st[last].len + 1;
    get_link(curr, last, c);
   last = curr;
 // achar os estados terminais
 // um estado terminal eh aquele que representa um sufixo da
 int p = last;
 while (p != -1) {
   term.pb(p);
    p = st[p].suf_link;
void dfs2(int v) {
 if (dp[v] != -1) return;
 dp[v] = 1;
 for (auto const &u : st[v].nxt) {
   if (!u.sec) continue;
   dfs2(u.sec);
    dp[v] += dp[u.sec];
void dfs(int v, int k, int &at, string &curr) {
 if (at == k) return;
 for (auto const &u : st[v].nxt) {
   if (!u.sec) continue;
   if (at + dp[u.sec] < k) {
     at += dp[u.sec];
     continue;
   curr.pb(u.fir);
   dfs(u.sec, k, at, curr);
   if (at == k) return;
    curr.pop_back();
void find_kth(int k) {
 string curr = "";
 dfs(0, k, at, curr);
 cout << curr << endl;
} // namespace sa
// chamar sa::init(); e sa::build(s);
```

Various (7)

binarysearch.cpp

```
// busca ternaria para int, usando busca binaria:
int 1 = 0, r = 1e9;
while (1 < r) {
 int mid = (1 + r) >> 1;
 (calc(mid) < calc(mid + 1)) ? r = mid : 1 = mid + 1;
return calc(1);
// first element >= x
int lower(int l, int r, int x) // first element >= x
 while (l < r) {
   // int mid = 1 + (r - 1) / 2; se tiver numero negativo
    int mid = (1 + r) >> 1;
    (x \le k[mid]) ? r = mid : 1 = mid + 1;
```

```
return k[1];
// last element <= x
vector<int> k(MAXN);
int upper(int 1, int r, int x) {
 while (l < r) {
   int mid = (1 + r + 1) >> 1;
    (k[mid] \le x) ? 1 = mid : r = mid - 1;
 return k[1];
```

ternarysearch.cpp

```
double eps = 1e-9;
while (r - 1 > eps) {
 double m1 = 1 + (r - 1) / 3;
 double m2 = r - (r - 1) / 3;
 double f1 = f(m1);
 double f2 = f(m2);
 if (f1 < f2)
   1 = m1;
 else
   r = m2;
return f(l); // o maximo de f(l)
```

parallelbinarysearch.cpp

37 lines

```
vector<pii> qry(q);
for (int i = 0; i < q; i++) {
 cin >> qry[i].sec.fir >> qry[i].sec.sec >> qry[i].fir;
 qry[i].sec.fir--, qry[i].sec.sec--;
vector<int> 1(n);
vector<int> r(n):
vector<vector<int>> on(q);
for (int i = 0; i < n; i++) {
 1[i] = 0;
 r[i] = q;
while (1) {
 bool ok = 1;
 for (int i = 0; i < n; i++) {
   if (l[i] < r[i]) {
      ok = 0:
      int mid = (1[i] + r[i]) >> 1;
      on[mid].pb(i);
 if (ok) break;
  for (int mid = 0; mid < q; mid++) {
    upd(qry[mid].sec.fir, qry[mid].sec.sec, qry[mid].fir);
    for (auto const &j : on[mid]) {
     int val = 0;
      for (auto const &k : adj[j]) {
       val += sum(k);
        if (val >= need[j]) break;
      (val \ge need[j]) ? r[j] = mid : l[j] = mid + 1;
    on[mid].clear();
// a resposta eh o l[i] para cada busca
```

19 lines

```
alienstrick.cpp
```

```
32 lines
pi solve (vector<int> &v, int lambda) {
  // associar um custo lambda para ser subtraido quando
      realizamos uma operacao
  // dp[i] - melhor profit que tivemos considerando as i
      primeiras posicoes
  // cnt[i] - quantas operacoes utilizamos para chegar no valor
  vector < int > dp(n + 1);
  vector<int> cnt(n + 1);
  dp[0] = 0;
  cnt[0] = 0;
  for (int i = 1; i \le n; i++) {
   dp[i] = dp[i - 1];
   cnt[i] = cnt[i - 1];
   int id = i - 1;
   dp[i] += v[id];
   int lo = \max(011, id - 1 + 1);
   int s = dp[lo] + (id - lo + 1) - lambda;
   if (s > dp[i]) {
     dp[i] = s;
     cnt[i] = cnt[lo] + 1;
 return {dp[n], cnt[n]};
int aliens trick(vector<int> &v) {
 int 1 = 0, r = n;
  while (1 < r) {
   int mid = (1 + r) >> 1;
   pi ans = solve(v, mid);
    (ans.sec > k) ? 1 = mid + 1 : r = mid;
 pi ans = solve(v, 1);
 return ans.fir + (1 * k);
```

cht.cpp

```
63 lines
struct line {
  mutable int m, b, p;
  bool operator<(const line &o) const {
   if (m != o.m) return m < o.m;
    return b < o.b;
  bool operator<(const int x) const { return p < x; }
  int eval(int x) const { return m * x + b; }
  int inter(const line &o) const {
   int x = b - o.b, y = o.m - m;
    return (x / y) - ((x ^ y) < 0 && x % y);
};
struct cht {
  int INF = 1e18;
  multiset<line, less<>> 1;
  void add(int m, int b) {
    auto y = 1.insert({m, b, INF});
    auto z = next(y);
    if (z != 1.end() \&\& y->m == z->m) {
     l.erase(y);
     return;
    if (y != l.begin()) {
     auto x = prev(y);
     if (x->m == y->m) x = 1.erase(x);
    while (1) {
     if (z == 1.end()) {
       y->p = INF;
```

```
break;
      y->p = y->inter(*z);
      if (y->p < z->p)
       hreak:
      else
        z = 1.erase(z);
    if (y == 1.begin()) return;
   z = v;
    auto x = --y;
    while (1) {
     int ninter = x->inter(*z);
     if (ninter \leq x - > p)
       x->p = ninter;
      else {
       l.erase(z);
       break:
      if (x == 1.begin()) break;
     y = x;
     x--;
      if (x->p < y->p)
       break;
      else
       l.erase(y);
 int get(int x) {
   if (1.empty()) return 0;
    return l.lower_bound(x)->eval(x);
};
```

expected valued p.cpp

double dp[MAXN][MAXN][MAXN]; // sushi problem

```
double solve(int i, int j, int k) {
 if (!i && !j && !k) return dp[i][j][k] = 0;
 if (dp[i][j][k] != -1) return dp[i][j][k];
 It is well-known from statistics that for the geometric
 (counting number of trials before a success, where each
      independent trial is probability p)
 the expected value is i / p
 double p = ((double)(i + j + k) / n);
 double ret = 1 / p; // expected number of trials before a
 if (i) {
   double prob = (double)i / (i + j + k); // probabilidade de
         ser um prato com um sushi
   ret += (solve(i - 1, j, k) * prob);
 if (j) {
   double prob = (double) j / (i + j + k); // probabilidade de
         ser um prato com dois sushis
   ret += (solve(i + 1, j - 1, k) * prob);
 if (k) {
   double prob = (double)k / (i + j + k); // probabilidade de
         ser um prato com tres sushis
```

ret += (solve(i, j + 1, k - 1) * prob);

return dp[i][j][k] = ret;

```
int kadane (vector<int> v) {
 int n = v.size(), ans = 0, max\_here = 0;
 for (int i = 0; i < n; i++) {
   max here += v[i];
   if (ans < max here) ans = max here;
   if (max_here < 0) max_here = 0;
 return ans:
int kadane circular(vector<int> v) {
 int n = v.size(), max_kadane = kadane(v);
 int max_wrap = 0, i;
 for (i = 0; i < n; i++) {
   max_wrap += v[i];
   v[i] = -v[i];
 max wrap += kadane(v);
 return max(max_wrap, max_kadane);
```

divideandconquerdp.cpp

16 lines

```
void compute(int 1, int r, int opt1, int optr, int i) {
 if (1 > r) return;
 int mid = (1 + r) >> 1;
 pair<int, int> ans = \{1e18, -1\}; // dp, k
 for (int q = optl; q <= min(mid, optr); q++) {
   if (q > 0)
     ans = min(ans, \{dp[i-1][q-1] + cost(q, mid), q\});
     ans = min(ans, {cost(q, mid), q});
 dp[i][mid] = ans.fir;
 compute(1, mid - 1, optl, ans.sec, i);
 compute (mid + 1, r, ans.sec, optr, i);
// dp[i][j] = min(dp[i - 1][k] + c(k, j)), para algum k <= j
// O(k * n * log(n))
```

sosdp.cpp

26 lines

```
// nesse caso, f[x] eh a funcaoo que soma:
// todos os a[i], tal que, (x & i) == i)
for (int mask = 0; mask < (1 << m); mask++) {
 f[mask] = a[mask];
for (int i = 0; i < m; ++i) {
 for (int mask = 0; mask < (1 << m); mask++) {
   if (mask & (1 << i)) f[mask] += f[mask ^ (1 << i)];
// nesse caso, f[x] eh a funcao que soma:
// todos os a[i], tal que, (x & i) == x)
for (int mask = 0; mask < (1 << m); mask++) {
 f[mask] = a[mask];
for (int i = 0; i < m; ++i) {
 for (int mask = 0; mask < (1 << m); mask++) {
   if (!(mask & (1 << i))) f[mask] += f[mask ^ (1 << i)];
```

lis.cpp

```
int lis() {
 vector<int> q;
 for (int i = 0; i < v.size(); i++) {
```

largestsquareofones.cpp

9 lines

```
int ans = 0;
for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {
    dp[i][j] = v[i][j];
    if (i && j && dp[i][j]) dp[i][j] = min({dp[i][j - 1], dp[i - 1][j], dp[i - 1][j - 1]}) + 1;
    ans = max(ans, dp[i][j]);
  }
}
cout << ans * ans << endl;</pre>
```

steinertree.cpp

```
16 lines
for (int i = 0; i \le 100; i++)
 for (int j = 0; j < (1 << 9); j++) dp[i][j] = 1e9;
for (int i = 0; i < k; i++)
  for (int j = 0; j < v.size(); j++) dp[j][1 << i] = d(cap[i],
      v[j]);
for (int mask = 2; mask < (1 << k); mask++) {
  for (int i = 0; i < v.size(); i++) {
    for (int mask2 = 1; mask2 < mask; mask2++) {
     if (!((mask & mask2) == mask2)) continue;
     int mask3 = mask ^ mask2;
     dp[i][mask] = min(dp[i][mask], dp[i][mask2] + dp[i][mask3]
    for (int j = 0; j < v.size(); j++) dp[j][mask] = min(dp[j][
        mask], dp[i][mask] + d(v[i], v[j]));
double ans = 1e9;
for (int i = 0; i \le 100; i++) ans = min(ans, dp[i][(1 << k) -
    11);
```

subsetsum.cpp

_ ..

```
89 lines
// subset sum com bitset de tamanho variado
// with n <= 10^6
template <int len = 1>
int subset_sum(int n, int h) {
 if (n \ge len) {
    return subset_sum<std::min(len * 2, (int)MAXN)>(n, h);
 bitset<len> dp;
  dp[0] = 1;
  for (auto const &x : w) {
   dp = dp \mid (dp << x);
  return dp._Find_next(max(011, h - 1)); // retorna o proximo
      bit setado apos a posicao passada como parametro
int solve(vector<int> &w, int tot, int h) {
 // tot -> soma de todos os elementos de w
  // h -> valor desejado
```

```
// quero retornar o menor valor x >= h, tal que existe um
      subset com soma x em w
  if (!w.size()) return 0;
  sort(w.rbegin(), w.rend());
  if (w[0] * 2 >= tot) return w[0];
  int n = w.size();
  w.pb(0);
 vector<int> aux;
  int p = 0;
 for (int i = 1; i \le n; i++) {
    if (w[i] != w[i-1]) {
      int cnt = i - p;
      int x = w[i - 1];
      int j = 1;
      while (j < cnt) {
       aux.pb(x * j);
       cnt -= j;
       j *= 2;
      aux.pb(x * cnt);
     p = i;
  swap(aux, w);
 return subset_sum(tot, h);
int f[MAXN];
                 // f[i] -> quantos "itens" com valor i tem
bitset<MAXN> dp; // dp[i] = 1, se existe um subset com soma i
// garantir que a soma de todo mundo seja < MAXN
void subset_sum(vector<int> &v) {
 for (auto const &i : v) {
    f[i]++;
  dp[0] = 1;
  for (int i = 1; i < MAXN; i++) {
    while (f[i] > 2) {
     f[i * 2]++;
      f[i] -= 2;
    while (f[i]--) dp |= (dp << i);
 Given N non-negative integer weights w and a non-negative
 computes the maximum S \le t such that S is the sum of some
     subset of the weights.
 O(N * max(w[i]))
int knapsack(vector<int> w, int t) {
 int a = 0, b = 0;
  while (b < w.size() && a + w[b] \le t) {
   a += w[b++];
 if (b == w.size()) {
   return a:
 int m = *max element(w.begin(), w.end());
  vector<int> u, v(2 * m, -1);
  v[a + m - t] = b;
 for (int i = b; i < w.size(); i++) {
    for (int x = 0; x < m; x++) {
     v[x + w[i]] = max(v[x + w[i]], u[x]);
    for (int x = 2 * m; --x > m;) {
     for (int j = max(011, u[x]); j < v[x]; j++) v[x - w[j]] =
            max(v[x - w[j]], j);
```

```
} a = t;
while (v[a + m - t] < 0) {
   a--;
} return a;
}</pre>
```

bitmasks.cpp

19 lines

```
// quantidade de bits setados na mask
cout << __builtin_popcount(mask) << endl;</pre>
// quando eh necessario percorrer todas as submasks ate (1 << n
// e fazer algo com todas as submasks dessa mask O(3^n)
for (int m = 0; m < (1 << n); m++) {
 for (int s = m; s; s = (s - 1) \& m) {
    // alguma coisa agui sabendo que mask s eh uma submask de m
// comprimindo as masks de um vector baseada em uma mask
    qualquer
for (int i = 0; i < masks.size(); i++) {
 int compressed = 0, curr_bit = 0;
 for (int j = 0; j < n; j++) {
   if (!(mask & (1LL << j))) continue;
    if (masks[i] & (1LL << j)) compressed |= (1LL << curr_bit);
   curr bit++;
 // alguma coisa sabendo que a mask compressed eh a mask
      comprimida da mask atual
```

hanoi.cpp

18 lines

```
vector<pair<char, char>> ans;
void solve(int n, char a, char b, char c) {
 if (n == 0) return;
 solve(n - 1, a, c, b);
 ans.pb({a, b});
 solve(n - 1, c, b, a);
// chamar pra sol
solve(n, 'A', 'C', 'B');
// 3 pilhas, sendo a pilha A com n discos e as outras duas
    pilhas vazias
// quero mover todo mundo da pilha A para a pilha C
// em cada movimento, vc tira o disco do topo de uma pilha e
    poe no topo de outra pilha
// desde que o raio do disco seja menor do que o raio do disco
    que ta no topo da outra pilha
// os n discos tem raios distintos aos pares
// e printar os movimentos em um par [to, from]
// numero minimo pra resolver pros primeiros n
// 1, 3, 7, 15, 31, 63, 127, 255
// f(n) = 2^n - 1
```

stacktrick.cpp

9 lin

```
spraguegrundy.cpp
vector<int> v = \{2, 3, 4, 5, 6\};
unordered_map<int, bool> vis;
unordered_map<int, int> dp;
```

```
int g(int x) // achar o grundy number na marra
  if (x == 0) return 0;
  vector<bool> ok(4, 0);
  int mex = 0:
  for (auto const &i : v) {
    int curr = q(x / i);
    if (curr < 4) ok [curr] = 1;
   while (ok[mex]) mex++;
  vis[x] = 1;
  return dp[x] = mex;
int n:
cin >> n;
int x = 0;
for (int i = 0; i < n; i++) {
 int k;
 cin >> k;
  x ^= solve(k);
(x > 0) ? cout << "Henry\n" : cout << "Derek\n";
// nim classico -> o jogador que comeca ganha se o xor dos
    tamanhos das pilhas for != 0
// teorema sprague-grundy (transformar um jogo qualquer em nim)
```

rectangleunion.cpp

```
94 lines
vector<int> x vals:
struct seatree {
  vector<int> seq, tag;
  segtree() {
    seg.assign(8 * x vals.size(), 0);
   tag.assign(8 * x_vals.size(), 0);
  void add(int ql, int qr, int x, int v, int l, int r) {
    if (qr \le 1 \mid | r \le ql) {
     return;
    if (ql <= 1 && r <= qr) {
     tag[v] += x;
     if (tag[v] == 0) {
       if (1 != r)
          seg[v] = seg[v << 1] + seg[(v << 1) | 1];
       else
          seg[v] = 0;
     } else {
        seq[v] = x_vals[r] - x_vals[1];
    } else {
     int mid = (1 + r) >> 1;
     add(q1, qr, x, (v << 1), 1, mid);
     add(ql, qr, x, ((v << 1) | 1), mid, r);
     if (tag[v] == 0 \&\& 1 != r) seg[v] = seg[v << 1] + seg[(v = r)]
           << 1) | 1];
  int gry() { return seg[1]; }
  void upd(int 1, int r, int x) { add(1, r, x, 1, 0, x_vals.
      size()); }
struct rect {
 int x1, y1, x2, y2;
struct event {
```

```
int time, 1, r, type;
  bool operator<(const event &b) {
   if (time != b.time) return time < b.time;
    return type > b.type;
};
const int inf = 1e9;
signed main() {
 int n;
 cin >> n;
 vector<rect> v(n);
 for (int i = 0; i < n; i++) {
    cin >> v[i].x1 >> v[i].y1 >> v[i].x2 >> v[i].y2;
    x_vals.pb(v[i].x1);
    x_vals.pb(v[i].x2);
 // comprime o x
  sort(x_vals.begin(), x_vals.end());
  x_vals.erase(unique(x_vals.beqin(), x_vals.end()), x_vals.end
       ());
  vector<event> ev;
  for (int i = 0; i < n; i++) {
    v[i].x1 = lower\_bound(x\_vals.begin(), x\_vals.end(), v[i].x1
         ) - x_vals.begin();
    v[i].x2 = lower_bound(x_vals.begin(), x_vals.end(), v[i].x2
        ) - x_vals.begin();
    ev.pb({v[i].yl, v[i].xl, v[i].x2, 0}); // adicao
    ev.pb({v[i].y2, v[i].x1, v[i].x2, 1}); // remocao
  segtree s;
  sort(ev.begin(), ev.end());
  int area = 0, 1 = -inf;
  for (auto const &i : ev) {
    if (1 == -inf) {
      1 = i.time:
      s.upd(i.1, i.r, 1);
    } else if (i.type == 1) {
      int curr = s.qry();
      s.upd(i.l, i.r, -1);
      if (s.gry() != curr) {
        int new_t = (s.qry() == 0) ? -inf : i.time;
        int lo = 1, hi = i.time - 1;
        area += ((hi - lo + 1) * curr);
        1 = \text{new t};
   } else {
      int curr = s.grv();
      s.upd(i.1, i.r, 1);
      if (s.grv() != curr) {
       int lo = 1, hi = i.time - 1;
        area += ((hi - lo + 1) * curr);
        l = i.time;
 cout << area << endl:
 return 0;
// n \le 5 * 10^5
// 0 <= x, y <= 10^9
\ensuremath{//} comprime coordenada no x pra montar a segtree dos valores de
// faz o line sweep pelo y
```

prefixsum2d.cpp

int v[1001][1001];

int p[1001][1001];

```
int qry(int x1, int y1, int x2, int y2) { return p[x2 + 1][y2 +
     1] - p[x2 + 1][y1] - p[x1][y2 + 1] + p[x1][y1]; }
for (int i = 0; i < n; i++) {
 for (int j = 0; j < n; j++) {
   p[i + 1][j + 1] = p[i][j + 1] + p[i + 1][j] - p[i][j];
   p[i + 1][j + 1] += v[i][j];
```

9 lines

48 lines

```
bellmanfordinequalities.cpp
struct solver {
 const int inf = 1e18;
 int n. src:
 vector<int> d;
 vector<pii> edges;
 solver(int _n) // recebe o numero de variaveis, indexadas de
       0 ate n - 1
   src = _n; // aqui denotaremos _n como o source
   n = _n + 1;
    for (int i = 0; i < _n; i++) // arestas de source para
        cada um dos vertices com custo 0
      edges.pb({{src, i}, 0});
 bool solve() // bellman ford
    d.assign(n, inf);
   d[src] = 0;
   int x:
    for (int i = 0; i < n; i++) {
     x = -1;
     for (auto const &e : edges) {
       auto [a, b] = e.fir;
       int cost = e.sec;
       if (d[a] < inf) {
         if (d[b] > d[a] + cost) {
            d[b] = max(-inf, d[a] + cost);
            x = b;
   return (x == -1); // false se tem ciclo negativo
 void add_constraint_leq(int i, int j, int c) // value_i -
      value_j <= c
    edges.pb({{j, i}, c});
 void add_constraint_geq(int i, int j, int c) // value_i -
      value_j >= c
   edges.pb(\{\{i, j\}, -c\});
 void add_constraint_eq(int i, int j, int c) // value_i -
      value_j = c
   add constraint leg(i, j, c);
    add_constraint_geg(i, j, c);
 };
};
```