

Programming Assignment 1

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Statement of Integrity: I, Joni Vrapı, attempted to answer each question honestly and to the best of my abilities. I cited any, and all, help that I received in completing this assignment.

Problem 2a. This algorithm will check every possible pair of points exactly once, calculating their distance apart from each other, then populate an array of pairs. Finally, it will sort that array of pairs by distance, and return the first m pairs in the array.

M-CLOSEST-POINTS(P, m)

```
1  pairs = [ ]
2  pointA = null
3  pointB = null
4  for i = 1; i < P.length; i++
5      pointA = P[i]
6      for j = i + 1; j < P.length; j++
7          pointB = P[j]
8          distance =  $\sqrt{(pointA.x - pointB.x)^2 + (pointA.y - pointB.y)^2}$ 
9          pairs[pairs.length + 1] = (pointA, pointB, distance)
10     return pairs.timSort(key = distance)[:m]
```

Problem 2b. The worst case running time for this algorithm is $O(n^2)$. There is a nested loop, which does not execute n times for each iteration of the outer loop. It executes only the number of times required to check only not-previously-checked points, relative to the outer loop. This roughly translates to $\frac{n^2}{2}$ iterations, however asymptotically this is still n^2 .

Problem 3. A problem similar to this (find the closest pair in a set P) is able to run in $O(n \log(n))$ worst case complexity because it utilizes a divide and conquer algorithm which is able to reduce the problem size by half upon each iteration. In this problem, however, the requirement to return the m closest pairs made it difficult for me to come up with an $O(n \log(n))$ solution to this problem. I have a sneaking suspicion that it is still possible, however. As it pertains to improving the worst case running time of this algorithm, I am not

sure how to significantly improve it. As you can see in line 10, after populating the pairs array with all the pairs (which include their distances), I use TimSort to sort the array by distance, then return the first part of the array (up to m). I believe a performance improvement could be made here by inserting new elements into the array in such a way as to preserve their order, negating the need to sort them at the end. This would remove a factor of $O(n\log(n))$, which would increase the performance of this function, however the worst case scenario would still be $O(n^2)$.

Analysis. After running the *charts.py* file as described in the README.md file, you will see that the produced output exactly mirrors an exponential curve. This was expected as per the asymptotic analysis performed above.