1 The Ring of Formal Power Series

1.1 Basic Ring Properties

For now, let F be any arbitrary field (not an arbitrary ring). One might notice that if we take the so called **Fibonacci Power Series**

$$1 + x + 2x^2 + 3x^3 + 5x^4 + 8x^5 + 13x^6 + \cdots$$

we can claim a certain "formal" equality

$$1 + x + 2x^2 + 3x^3 + 5x^4 + \dots = \frac{1}{1 - x - x^2}.$$

Is this equation true in any sense? We might notice if we multiply by $1-x-x^2$ on both sides and join common monomial terms we get the trivial equality 1 = 1. But what exactly does this equality mean? Note that this equation is true in the traditional power series sense. It has a non-zero radius of convergence. But what about other power series and formulae?

The point of introducing formal power series is to make equalities like the above actual equalities in a certain very well-defined sense, without worrying at all about issues like convergence, especially if the field we are working under is not the field of real or complex numbers.

Definition 1. Let F be a field. The formal power series F[[x]] consists of sequences $(a_0, a_1, a_2, ...)$ where the $a_i \in F$. We pretty much always represent these sequences by $a_0 + a_1x + a_2x^2 + \cdots$. Given two formal power series

$$a = \sum_{n=0}^{\infty} a_n x^n$$

and

$$b = \sum_{n=0}^{\infty} b_n x^n,$$

define their sum a + b to be

$$a+b = \sum_{n=0}^{\infty} (a_n + b_n)x^n$$

and define their product ab to be

$$ab = \sum_{n=0}^{\infty} \left(\sum_{k=0}^{n} a_k b_{n-k} \right) x^n$$

We can also define things like the formal derivative, which is unambiguous no matter the field F we are working in, since n can be interpreted as $1 + 1 + \cdots + 1$ repeated n times.

Proposition 1. The set of formal power series F[[x]] forms a commutative ring (with unity) under the operations defined above.

Proof. Hopefully the following claims are fairly clear:

- The formal "zero" $0 + 0x + 0x^2 + \cdots$ is a zero element.
- The formal "one" $1 + 0x + 0x^2 + \cdots$ is a multiplicative identity.
- Every element of a formal power series has an additive inverse.
- Addition and multiplication in F[[x]] are associative and commutative. This follows from commutativity and associativity of the base field F.

The first, fourth and third claims make clear that this ring is an abelian group under addition. The second and fourth items make clear the sense that the multiplication operation in F[[x]] has the appropriate ring structure. Distributivity is straightforward, though a little tedious, to check.

In this ring, what is the so called group of units (that is, elements with a multiplicative inverse)?

Proposition 2. The formal power series $a_0 + a_1x + a_2x^2 + \cdots$ has no multiplicative inverse if $a_0 = 0$ and has a unique multiplicative inverse otherwise.

Proof. Given any such formal power series $a_0 + a_1x + a_2x^2 + \cdots$ we will suppose that such an inverse exists. Then the relations

$$a_0b_0 = 1$$

$$a_0b_1 + a_1b_0 = 0$$

$$\vdots$$

$$\sum_{k=0}^{n} a_k b_{n-k} = 0$$

$$\vdots$$

must hold. It is clear from the first equation that no such inverse exists if $a_0 = 0$. Now suppose $a_0 \neq 0$. Then from the general equation for n we can solve for b_n as

$$b_n = \begin{cases} \frac{1}{a_0} & n = 0\\ -\frac{1}{a_0} \sum_{k=1}^n a_k b_{n-k} & \text{otherwise.} \end{cases}$$

It follows inductively that if $a_0 \neq 0$ then b_n is always well defined in terms of the a_i and thus a well-defined inverse for $a_0 + a_1x + a_2x^2 + \cdots$ exists. Hence the claim follows, as desired.

The following proposition shows that F[[x]] is an integral domain, which is a particularly nice property.

Theorem 1. If f, g, and h are formal power series in R[[x]], where R is an integral domain, then

$$fh = gh \implies f = g.$$

So R[[x]] is an integral domain.

Proof. It suffices to prove that if f and g are non-zero power series then $fg \neq 0$ (for then, in the language of our original statement, we would have (f-g)h = 0 with $h \neq 0$, implying f = g). Let $f = a_0 + a_1x + a_2x^2 + \cdots$ and let $g = b_0 + b_1x + b_2x^2 + \cdots$. Let i be the least index such that $a_i \neq 0$ and let j be the least index such that $b_j \neq 0$. Then it follows that the coefficient of x^{i+j} of gh is

$$\sum_{n=0}^{i+j} a_k b_{i+j-k} = \sum_{k=0}^{i-1} a_k b_{i+j-k} + a_i b_j + \sum_{k=i+1}^{i+j} a_k b_{i+j-k}$$
$$= \sum_{k=0}^{i-1} a_k b_{i+j-k} + a_i b_j + \sum_{k=0}^{j-1} a_{i+j-k} b_k$$
$$= 0 + a_i b_i + 0 \neq 0$$

since $a_i \neq 0$, $b_j \neq 0$, and $a_k = 0$ for all k < i and $b_l = 0$ for all l < j (by definition). So fg is not the zero element, as desired.

1.2 Convergence of formal power series and an underlying topology on F[[x]]

One deficit of the ring of formal power series is that unlike the usual notion of power series, which have radius of convergence, there might not be a useful way to evaluate a power series outside of x = 0. Such evaluation might not make sense if F is not \mathbb{R} or \mathbb{C} . However, we still do want to say something about "formal" convergence of a power series. As an example, we would like to say that the sequence of power series $a_0, a_0 + a_1x, a_0 + a_1x + a_2x^2, \ldots$ converges to the associated power series $a_0 + a_1x + a_2x^2 + \cdots$. To make this notion precise we will precisely define the notion of convergence (which will hopefully be enough for our purposes, since strictly speaking convergence of sequences does not determine a topology if your underlying space is not metrizable or first countable I might flesh this out a little bit later, considering that the convergence I've defined is just R^{ω} where R gets the discrete topology and we endow the product with the product topology). We will also use this notion of convergence to give a criterion about whether or not an infinite sum of formal power series converges.

Definition 2. We will say a sequence r_1, r_2, r_3, \ldots of elements in F converges sharply to some limit s if r_n is eventually constant and equal to s.

Definition 3. A sequence of formal power series f_1, f_2, \ldots converges to a formal power series g if the sequence (r_n^i) obtained by taking the x^i coefficient of each f_n converges sharply to the x^i coefficient in g.

For example, the sequence of formal power series'

$$1 + x + x^{2} + x^{3} + \cdots$$

$$1 + 2x + 2x^{2} + 2x^{3} + \cdots$$

$$1 + 2x + 3x^{2} + 3x^{3} + \cdots$$

$$1 + 2x + 3x^{2} + 4x^{3} + \cdots$$

converges to the formal power series $1 + 2x + 3x^2 + 4x^3 + 5x^4 + \cdots$

Lemma 1. In F[[x]] the sequence $a_0, a_0 + a_1x, a_0 + a_1x + a_2x^2, \ldots$ converges to the formal power series $f(x) = a_0 + a_1 + a_2x^2 + \cdots$.

The proof is very simple.

Proof. For any $k \geq 0$, the degree k term of $a_0 + \cdots + a_n x^n$ is equal to the degree term of f(x) for sufficiently large n (specifically if $n \geq k$).

Theorem 2. Let f_1, f_2, f_3, \ldots be formal power series. Then the infinite series $f_1 + f_2 + f_3 + \cdots$ exists if and only if $\operatorname{ord}(f_k) \to \infty$ where

$$\operatorname{ord}(f) = \min(k \mid a_k \neq 0)$$

(provided that $f = a_0 + a_1x + a_2x^2 + \cdots$).

Proof. Let $p_n(f_k)$ denote the coefficient of x^n in f_k .

First suppose $\operatorname{ord}(f_k) \to \infty$. This means that for any $n \geq 0$ there exists N > 0 such that $p_n(f_k) = 0$ for all $k \geq N$. It follows that $p_n(f_1 + \cdots + f_k)$ is constant for all $k \geq N$, implying for all n that each coefficient converges sharply. Hence by definition $f_1 + f_2 + f_3 + \cdots$ exists.

Conversely, suppose that the series $f_1+f_2+f_3+\cdots$ exists. Then for all $n\geq 0$ there exists $N_n\geq 0$ such that $p_n(f_1+\cdots+f_k)$ is constant for all $k\geq N_n$. But this implies that $p_n(f_k)=0$ for all $k>N_n$ (because p_n is actually a ring homomorphism). Taking $N=\max_{1\leq i\leq n}(N_i)$, we see that $\operatorname{ord}(f_k)>n$ for $k\geq N$. It follows that $\operatorname{ord}(f_k)\to\infty$ as desired.

Some examples which illustrate this theorem is appropriate. The sum of series

$$1 + (x + x^2 + x^3 + \cdots) + (x + x^2 + x^3 + \cdots)^3 + \cdots$$

converges, since we can see that $\operatorname{ord}(f_k) = \operatorname{ord}(f_{k-1})$ for $k \geq 1$. However, the series

$$1 + \frac{1}{2} + \frac{1}{4} + \cdots$$

does not, which is a potential limitation of the way we defined convergence. However, for most things we need these power series for (combinatorics) these limitations do not come up in practice. **Theorem 3.** Suppose $a_i \to a$ and $b_i \to b$, where a_i, b_i, a, b live in F[[x]]. Then we have that

$$a_i + b_i \rightarrow a + b$$

and

$$a_i b_i \to ab$$
.

It follows that F[[x]] is a topological ring: that is, the addition and multiplication operators are continuous operations that respect convergence.

Proof. Omitted for now.

Theorem 4. Any series which is convergent in F[[x]] is absolutely convergent. In particular, given any convergent infinite sum of power series we can rearrange any number of terms without changing the equality.

Proof. Omitted for now.

With the above facts, we are ready to do some interesting things with formal power series which are completely rigorous.

Theorem 5. Suppose $f \in F[[x]]$ is some power series where the constant term is 0 (so ord(f) = 1). Then

$$\frac{1}{1-f} = 1 + f + f^2 + f^3 + \cdots$$

Proof. First observe by hypothesis $(\operatorname{ord}(f) = 1 \text{ and } f(0) = 0)$ that the terms on the left hand and right hand sides are well defined formal power series. By definition, g is the limit of the power series

$$g_n(x) = \sum_{k=0}^{n} (f(x))^k.$$

We observe that

$$(1-f)(g_n) = 1 - f^{n+1}.$$

Since ord $f \ge 1$ we can deduce that $f^{n+1} \to 0$ and hence $(1-f)g_n \to 1$. But since $g_n \to g$ it follows that (1-f)g = 1, as desired.

Alternatively, using absolute convergence we have that

$$(1-f)g = (1-f) + (f-f^2) + (f^2 - f^3) + \dots = 1,$$

as desired.

From this theorem we can deduce some interesting patterns. For example, we obtain the usual geometric series formula

$$\frac{1}{1-x} = 1 + x + x^2 + x^3 + \cdots$$

We can use this to generate power series for the inverses of $(1-x)^n$.

$$\frac{1}{1-x} = 1+x+ \qquad x^2+x^3+ \qquad \cdots$$

$$\frac{1}{(1-x)^2} = 1+2x+ \quad 3x^2+4x^3+ \quad \cdots$$

$$\frac{1}{(1-x)^3} = 1+3x+ \quad 6x^2+10x^3+\cdots$$

$$\frac{1}{(1-x)^3} = 1+4x+10x^2+20x^3+\cdots$$

Something very interesting is going on here! It appears as if the terms of these expansions line up with binomial coefficients. But why are the expansions like this?

2 Proof of the Formal Binomial Theorem

In this section we will prove what is known as the formal binomial theorem for power series. (To fill in later we will consider why we only consider power series of the form (1 + xf(x)) and why we can only hope to define a general exponentiation there).

For any formal power series f(x) define

$$(1+xf(x))*r = 1 + rxf(x) + \left(\frac{r(r-1)}{2}\right)x^2f(x)^2 + \left(\frac{r(r-1)(r-2)}{3!}\right)x^3f(x)^3 + \cdots$$

Theorem 6. The following statements are all true:

- 1. The right hand side converges in the ring of formal power series.
- 2. When $r \ge 0$ is an integer, then $(1 + xf(x)) * r = (1 + xf(x))^r$.
- 3. We have that for all r and s,

$$((1+xf(x))*r)((1+xf(x))*s) = (1+xf(x))*(r+s).$$

- 4. When r < 0 is an integer, then $(1 + xf(x)) * r = (1 + f(x))^r$.
- 5. For all r and s,

$$((1 + xf(x)) * r) * s = (1 + xf(x)) * (rs)$$

Proof. The first statement follows readily from the fact that each succeeding term has order greater than the previous. The second statement is a reformulation of the usual binomial theorem.

For the third statement, for any integer $k \geq 0$, the coefficient of $f(x)^k x^k$ in the absolute product is a polynomial $P_k(r,s)$ of degree k in r and s. For example, when k=2, we have that this coefficient is

$$1 \cdot \left(\frac{s(s-1)}{2}\right) + rs + \left(\frac{r(r-1)}{2}\right) \cdot 1.$$

Likewise the coefficient of $x^k f(x)^k$ in (1+xf(x))*(r+s) is a polynomial $Q_k(r,s)$ of degree k in r and s. We know that $P_k(r,s) = Q_k(r,s)$ when $r,s \in \mathbb{N}$. So $P_k = Q_k$ for infinitely many points. So we use a k+1 point equality argument 3 times in order to conclude that $P_k = Q_k$ in general. This implies equality of the power series.

The fourth statement follows directly from the second and third statements.

The fifth statement has an argument which is the exact same reasoning as the third statement.

By this theorem, we can introduce the convention of just writing g(x) * r as $g(x)^r$ for any g with g(0) = 1. In particular we have

$$(1-x)^{-m} = 1 + (-m)(-x) + \left(\frac{(-m)(-m-1)}{2}\right)(-x)^2 + \left(\frac{-m(-m-1)(-m-2)}{3!}\right)(-x)^3 + \cdots$$

$$= 1 + mx + \frac{m(m+1)}{2}x^2 + \frac{m(m+1)(m+2)}{6}x^3 + \cdots$$

$$= \sum_{n=0}^{\infty} {m+n-1 \choose n} x^n.$$

This explains the connection with expansions of $(1-x)^{-m}$ and binomial coefficients that we saw earlier.

3 More Topology; distances between power series

Now let's consider the case when $F = \mathbb{R}$. We will define a distance function between f and g to be 2^n if the first disagreement of coefficients is at the x^n monomial. Then define the norm of a formal power series to be its distance from 0. This distance function satisfies the inequality

$$d(f,h) \le \max(d(f,g),d(f,h)),$$

this is called a supermetric/ultrametric inequality. We will apply this distance function to study the equation

$$D(x) = 1 + x[D(x)]^2$$

(which is the generating function relation for the Catalan Numbers).

Theorem 7. The map $f \mapsto 1 + xf^2$ has a unique fixed point in the ring of formal power series. Then, for any power series g_0 , if we define g_0 , $g_w = xg_{w-1}^2$, then $g_w \to D(x)$ in the ring of formal power series.

Proof. Only a sketch. First, we observe that the subset of f such that f(0) = 1 is a contains the image of the map in question. Then we can use the contraction mapping principle to show that a fixed point exists. To show it is unique, then we observe that if f and g both have constant term 1, then $d(1 + xf^2, 1 + xg^2)$ decreases.

Now, we might consider the spurious solution to the formal power series of D(x). Can we find an initial position which convergest to it? If we take the initial solution 1/x and apply the map above, we don't actually get convergence. However, the map $f \mapsto \frac{1}{f} + \frac{1}{xf}$ seems to work fine. Now what about

$$E(x) = 1 + x(E(x))^3$$

(I think this is the generating function for the Fuss-Catalan numbers of parameter 2 or 3). And what about its spurious solutions? It turns out that the spurious solutions are Laurent series in the variable \sqrt{x} .

4 Homogeneous Linear Recurrence Equations and the magic of Linear Algebra

One of the most basic combinatorial questions that one can ask is "What is the number/how many objects are there?" In the last section we set up the machinery of *generating functions*, where we put a parametrized family into a sequence and we operated on it using a so called "generating function" approach.

Here is an example of the power of this approach. Consider the problem of finding a formula for the sequence

$$1, 2, 5, 14, 41, \dots$$

satisfying f(n+2) = 4f(n+1) - 3f(n) for $n \ge 0$.

Setting up the generating function

$$F(x) = \sum_{k=0}^{\infty} f(n)x^{n}$$

we rearrange and we find

$$F(x) = \sum_{k=0}^{\infty} f(n)x^{n}$$

$$= f(0) + f(1)x + \sum_{k=2}^{\infty} f(n)x^{n}$$

$$= f(0) + f(1)x + \sum_{k=0}^{\infty} (4f(n+1) - 3f(n))x^{n+2}$$

$$= f(0) + f(1)x - 4x + 4x \sum_{k=0}^{\infty} f(n)x^{n} - 3x^{2} \sum_{k=0}^{\infty} f(n)x^{n}$$

$$= f(0) + (f(1) - 4)x + (4x - 3x^{2})F(x)$$

which implies that

$$F(x)(1 - 4x + 3x^{2}) = f(0) + (f(1) - 4)x.$$

Conversely, any formal power series of the form $\frac{ax+b}{1-4x+3x^2}$ is a solution to the recurrence relation. Since $1-4x+3x^2=(1-x)(1-3x)$, by partial fraction decomposition we have

$$\frac{ax+b}{1-4x+3x^2} = \frac{c}{1-x} + \frac{d}{1-3x}.$$

By the geometric series formula of last chapter, this is equal to $c(1 + x + x^2 + x^3 + \cdots) + d(1 + 3x + 9x^2 + 27x^3 + \cdots)$. We conclude that

$$f(n) = c + d3^n$$

for arbitrary c and d are all the solutions to this recurrence relation.

In this section, instead of using the ring structure of the formal power series F[[x]] we will put a vector space structure on the set of infinite sequences.

Observe that the equation f(n+2) = 4f(n+1) - 3f(n) is **linear**. That is, if f_1 , and f_2 satisfy this recurrence relation, so does $g = Af_1 + bf_2$, where A and B are arbitrary constants. This implies that the set of solutions to this linear recurrence relation is a vector space (namely, a subspace of the \mathbb{F}^{ω} sequence space). Here are some elements of this vector space:

$$(1, 1, 1, 1, 1, \dots)$$

$$(1, 3, 9, 27, \dots)$$

$$(1, 0, -3, -12, -39, \dots)$$

$$(0, 1, 4, 13, 40, \dots)$$

This vector space is 2-dimensional: every sequence is completely determined by the first two elements. One might notice that either the first two elements above or the last two elements above form a basis. The first two elements are nice in that they have simple formulas, namely 1^n and 3^n . The last two are nice because they can be seen to be like a standard basis for the vector space, and you can easily express any other sequence in the vector space in terms of a linear combination of the two.

Definition 4. Any recurrence relation of the form

$$\sum_{j=0}^{k} A_j f(n+j) = 0$$

where the A_j are constants is called a kth order homogeneous linear recurrence relation with constant coefficients.

There are certainly important recurrence relations where the coefficients are non-constant, consider the simple example

$$(n+1)! = (n+1)(n!).$$

For the rest of this section, let S be the vector space of all sequences

$$f = (f(0), f(1), f(2), \dots)$$

called the sequence space. For any sequence f, define T(f) (or Tf) as the sequence whose value at n is f(n+1). In other words, T is the "left shift" operator. It is easy to see that T is linear. If we view it as an infinite diagonal matrix, then T is zero everywhere except 1 on the diagonal just above the main diagonal.

Now observe that the recurrence relation

$$f(n+2) - 4f(n+1) + 3f(n) = 0$$

can be written as

$$T^{2}(f) - 4T(f) + 3f = 0.$$

Note that the 0 in the above equation is the zero sequence $(0,0,0,\ldots)$, and not the zero of the field. T^2 denotes $T \circ T$.

This act of representing recurrence relations as polynomials of this shift operator is very powerful, as we shall see. Let's see what T - 3I does to a vector of the form $f(n) = r^n$. We have that

$$(T-3I)f(n) = Tf(n) - 3If(n)$$
$$= r^{n+1} - 3r^n$$
$$= (3-r)r^n.$$

That is, f is an eigenvector of T - 3I with eigenvalue r - 3. In particular, if r = 3, then f is annihilated by T - 3I. Observe that (T - I)(T - 3I) annihilates $f(n) = 3^n$ and (T - 3I)(T - I) annihilates $f(n) = 1^n$.

But by linearity, both these operators are equal! Hence any linear combination of $f_2(n) = 3^n$ and $f_1(n) = 1$ is in the kernel of $T^2 - 4T + 3I$, and therefore satisfies the recurrence relation we were interested in above.

Let's introduce some new terminology. The operator T-I (also sometimes written T-1) is called the *difference operator* (compare with the derivative operator). The equation p(T)f = 0 for any polynomial p is called a *difference equation*.

Here is the first result of a more general theorem:

Theorem 8. Let $p(t) = a_d t^d + a_{d-1} t^{d-1} + \cdots + a_0$, and consider the equation p(T)f = 0, a homogeneous linear recurrence equation with constant coefficients.

Suppose p(t) has d distinct roots. Then the general solution to the linear recurrence relation p(T) f = 0 is of the form

$$f(n) = A_1 r_1^n + A_2 r_2^n + \dots + A_d r_d^n,$$

where the r_i are the roots of p(t).

Proof. It is easy to check that $f(n) = r_i^n$ is a solution to the linear recurrence equation we are considering. These d functions are linearly independent, since they are all eigenvectors of the left shift operator T with distinct eigenvalues. Since we know that the subspace of the sequence space consisting of $\ker(p(T))$ is d-dimensional, we are done.

Note that even if the terms of the sequence f(n) are integers, the numbers r_i and A_i need not be.

But what if our polynomial doesn't have repeated roots? For example, take $p(t) = t^2 - 2t + 1$. Using a standard method (such as the one in the beginning of this section), we get a 2 parameter family of generating functions

$$\frac{ax+b}{x^2-2x+1}.$$

When a = -1 and b = 1 we can simplify to get the formal power series

$$1 + x + x^2 + x^3 + \cdots$$

Put a = 1, b = 0 to get

$$\frac{x}{1 - 2x + x^2} = x(1 - x)^{-2} = x(1 + 2x + 3x^2 + 4x^3 + \dots) = x + 2x^2 = 3x^3 + \dots$$

then we have two "fundamental solutions" f(n) = 1, g(n) = n.

In general, it turns out that if $p(T) = (T - rI)^d$, there are d fundamental solutions $f(n) = n^k r^n$ for k = 0 to d - 1. When r = 1, these fundamental solutions are polynomials and the general solution is a polynomial of degree less than or equal to d.

Theorem 9. If p(t) is degree d, with leading coefficient and constant term non-zero, the solutions of p(T)f = 0 form a d-dimensional space with basis elements of the form $n^k r^n$, where r is any root of p(t) = 0 and k is any non-negative integer less than the multiplicity of r in p.

Proof. omitted for now

Theorem 10. If f satisfies a homogeneous linear recurrence equation of order d, then $f, Tf, T^2f, \ldots, T^df$ are linearly dependent.

Conversely, if f, Tf, T^2f, \ldots span a d-dimensional sequence space, then f satisfies a linear recurrence equation of order d.

Proof. Exercise to reader; follows from basic properties of dimension in finite dimensional vector spaces.

5 More comments on HLREs

The sequence $\{F_n\}$ of fibonacci numbers is annihilated by the operator

$$T^2 - T - I$$

where T is the left shift operator. The generating function has denominator $1 - x - x^2$. Note that the coefficients appear flipped. This is true in genreal, as the following theorem indicates. **Theorem 11.** Suppose a_0, \ldots, a_d are non-zero coefficients, and that f is some sequence. The following statements are equivalent:

• f satisfies the equation

$$(a_d T^d + a_{d-1} T^{d-1} + \dots + a_1 T + a_0) f = 0.$$

• The generating function for f can be expressed as a rational function whose numerator is of degree less than d, with denominator

$$a_d + a_{d-1}x + \cdots + a_0x^d$$

• f(n) is expressible as a linear combination of functions of the form $n^j r^n$ where r is a root fo $a_d t^d + \cdots + a_1 t + a_0 = 0$ and where $j \ge 0$ is less than the multiplicity of r.

Observe the following: if the roots of such a polynomial are ± 1 , then the solutions to such a recurrence relation are simply polynomial equations in the variable n.

Let's consider some of these results in action. Suppose we'd like to find either a recurrence or closed form formula for some sequence, say,

$$1, 2, 5, 12, 29, 70, \cdots$$

satisfying a linear recurrence relation.

If we suspect that such a relation is second order, we can take the system of equations

$$5 = A(2) + B(1)$$
$$12 = A(5) + B(2)$$

From which we get A = 2 and B = 1. This gives us the candidate linear recurrence relation f(n) = 2f(n-1) + f(n-2), which can be seen to work. But what if we suspected that the relation was third order instead? Then we get the linear recurrence

$$f(n) = 3f(n-1) - f(n-2) - f(n-3).$$

Is there any relation between these two recurrence relations? It seems like the former is "simpler" because it involves less previous terms. It turns out that this intuition is correct. Here are some facts about the kinds of sequences you get.

Proposition 3. 1. Suppose p and q are polynomials such that p(T)f = 0 and $p \mid q$. Then q(T)f = 0.

- 2. For any sequence f satisfying an LRE with constant coefficients, there exists a unique monic polynomial p(t) such that p(T)f = 0 and q(T)f = 0 implies that $p \mid q$.
- 3. Any sequence satisfying an LRE satisfies a linear recurrence relation satisfies a recurrence relation with a minimal number of terms and which is unique up to scalar multiples.

Proof. The first item is more or less obvious. The second item follows from the fact that the set of polynomials p such that p(T)f = 0 form an ideal in the polynomial ring and there is a unique monic polynomial which is the generator for said ideal. The third item more or less follows from the second item.

For example, in the above example, the polynomial operators for each recurrence relation are

$$T^{2} - 2T - I$$

$$T^{3} - 3T^{2} + T + I = (T^{2} - 2T - I)(T - I)$$

Let's consider a less trivial example. Let's find a recurrence relation for the sequence $g(n) = F_n + 2^n$, where F_n is the *n*th Fibonacci number. If we suspect that this sequence satisfies a third order linear recurrence relation, we can solve a 3 dimensional system to get that

$$g(n) = 3g(n-1) - g(n-2) - 2g(n-3).$$

One can confirm that g(n) indeed satisfies the recurrence relation above.

Here is a slicker method: We might use linearity, and observe that $(T^2 - T - I)(T - 2I)$ annihilates the g(n). This gives us the desired recurrence relation. Or similarly, we add

$$\frac{1}{1-x-x^2} + \frac{1}{2x} = \frac{2-3x-x^2}{1-3x+x^2+2x^3}.$$

To test the minimality of this recurrence relation, we can show that $gcd(2 - 3x - x^2, 1 - 3x + x^2 + 2x^3) = 1$, or we can show that the determinant of cycles of the first equation is non-zero (that there are no non-trivial dependence relations in this recurrence, showing it is minimal).

Finally, we conclude this section with some combinatorics: we will tile a $2 \times n$ rectangle by dominos (known as a domino tiling).

This is a classic problem. We can solve it with the following observation. Either the $2 \times n$ rectangle begins with a vertical rectangle or two horizontal rectangles. Thus for $n \geq 2$ the number of rectangles D_n satisfies the recurrence relation

$$D_n = D_{n-1} + D_{n-2}.$$

It follows that the number of tilings is counted by the Fibonacci numbers.

What about a $3 \times n$ rectangle? (to be illustrated later).

5.1 A short comment on various bases

Let's consider the subspace of solutions to say the equation $(T - rI)^m f = 0$ where T is the left shift operator. One such basis is given by

$$f(n) = r^n, f(n) = nr^n, \dots, f(n) = n^{m-1}r^n.$$

Instead, taking the generating function point of view, A natural basis for the space of generating functions these sequences is

$$\frac{1}{(1-rx)^m}, \frac{x}{(1-rx)^m}, \frac{x^2}{(1-rx)^m}, \dots, \frac{x^{m-1}}{(1-rx)^m}.$$

Another natural basis is

$$\frac{1}{1-rx}, \frac{1}{(1-rx)^2}, \dots, \frac{1}{(1-rx)^m}.$$

These basis elements correspond to the solutions $f_{r,k}(n) = \binom{n+k-1}{n} r^n$.

6 The Calculus of Finite Differences

In this section we investigate some special properties of the basis

$$1, x, \frac{x(x-1)}{2}, \frac{x(x-1)(x-2)}{6}, \dots$$

of polynomial space.

Suppose for the first example that we're given a sequence

$$0, 1, 5, 15, 34, 65, 111, \dots$$

and we suspect its formula is given by a cubic polynomial. How do we find it?

One primitive way is to use undetermined coefficients, or Lagrange interpolation. Here is another way. First, we observe that if f(n) is a polynomial of degree d, then (T-I)f is degree d-1. So we can construct a difference table.

(difference table here).

Conversely, if the nth difference is 0, then f is an nth degree polynomial and lower. We observe the terms on the very right of the difference table. If we look at the difference tables for the polynomials above, we find that we can represent our cubic polynomial using this basis in a natural way. It will follow that

$$p(x) = 0 \cdot 1 + 1 \cdot x + 3\left(\frac{x(x-1)}{2}\right) + 3\left(\frac{x(x-1)(x-2)}{6}\right) = \frac{x^3 + x}{2}.$$

This basis above is called the basis of falling factorials. We'll denote the kth element (k starts index at 0) as $(x)_k$.

The falling factorial polynomials form an integral basis in P[x] with the property that p(n) is an integer for all integers n. In other words, any polynomial p(n) can be represented as a linear combination of finitely many basis polynomials where the coefficients are integers.

Proof. The converse is hopefully obvious, we'll prove the forward direction. Suppose p(n) is an integer for all n. Then

$$p(x) = a_0(x)_0 + a_1(x)_1 + \dots + a_n(x)_n.$$

The idea is that plugging in x = 0 we will find that $a_0 \in \mathbb{Z}$. Then we continue inductively to deduce that $a_j \in \mathbb{Z}$ for all coefficients a_j .

6.1 Application of the falling factorial basis: Dobinski's formula for the Bell Numbers

Let B_n be the number of **partitions** of an n-element set (called Bell numbers). This is the same thing as dividing the set into chunks such that every element is in exactly one chunk. For example, when n = 3, $S = \{a, b, c\}$ we can represent the partitions as follows:

$$\{\{a,b,c\}\},\{\{a,b\},\{c\}\},\{\{a,c\},\{b\}\},\{\{c,b\},\{a\}\},\{\{a\},\{b\},\{c\}\}.$$

It is easy to prove that I've listed all of them.

So it's simple to show that $B_3 = 5$, $B_2 = 2$, and $B_1 = 1$. We would like to try and find a formula for the B_n . This will involve a couple of magic tricks.

The first thing to do is to consider a seemingly unrelated question: given an n element set N and an m element set M, how many functions $f: N \to M$ are there? We can use a left-right approach here. First we notice that the answer is obviously m^n . Alternatively, we might notice that given any function $f: N \to M$ we can create a partition π of N by taking the non-empty preimages $f^{-1}(\{i\}), i \in M$.

Let $S_{n,k}$ represent the number of partitions of N of size k. Note that any such partition does not uniquely determine a function f. We also need to determine the values in M which each partition is sent to, there are m!/(m-k)! such ways we can do this for a partition of size k. Hence we obtain the identity

$$m^n = \sum_{k=1}^n S_{n,k} \frac{m!}{(m-k)!}.$$

Since this is true for all positive integer values of m and it is a polynomial identity we conclude it actually holds for general x and hence we have the formula

$$x^n = \sum_{k=1}^n S_{n,k}[x]_k,$$

where $[x]_0 = 1, [x]_k = x(x-1)(\cdots)(x-k+1)$.

We observe that the set of polynomials $[x]_k$ actually form a basis for the set of real valued polynomials (note that as defined $[x]_k$ has degree k). So we might define a linear functional $L: \mathbb{R}[x] \to \mathbb{R}$ as being the unique functional with the property that

$$L([x]_k) = 1$$

for all k. Applying L to the formula above we get

$$L(x^n) = \sum_{k=1}^n S_{n,k} = B_n,$$

for the $S_{n,k}$ counts the number of partitions of N of size k, and we are summing over all the possible values of k.

Here is where the real magic happens: We observe that

$$L([x]_n) = 1 = \frac{1}{e}e = \frac{1}{e}\sum_{k=0}^{\infty} \frac{1}{k!} = \frac{1}{e}\sum_{k=0}^{\infty} \frac{[k]_n}{k!}.$$

To justify the last equality just observe that the infinite sum is just the previous one shifted over n places to the right (using the cancellation of $[k]_n$ and k!). It follows by the properties of the basis and the fact that any polynomial p(x) is a finite sum of the basis elements $[x_k]$ that we get

$$L(p(x)) = \frac{1}{e} \sum_{k=0}^{\infty} \frac{p(k)}{k!}.$$

In particular, when $p(x) = x^n$, we get

$$B_n = L(x^n) = \frac{1}{e} \sum_{k=0}^{\infty} \frac{k^n}{k!}$$

which is called *Dobinski's formula* for the Bell numbers.

7 Weighted Enumeration

Previously, we looked at the number of ways to tile a $2 \times n$ rectangle with dominoes, and we found that the number of such tilings was a Fibonacci number. Here is another question: If we tile a $2 \times n$ rectangle with dominoes, then what fraction of the tiles will be vertical? This question is not exactly formed, as it depends on the tiling. Here is a better question: What is the expected number of tilings? Here's the magic technique that we will use: we will count with polynomials.

Definition 5. Given any $2 \times n$ tiling, give it weight w^k if it has exactly k vertical dominoes. The **number** of tilings by dominoes is defined as the sum of the weights. We denote this polynomial by $p_n(w)$.

For example, for a 2×3 tiling the number of tilings is $w^3 + 2w$. If we evaluate using w = 1 we get the actual *number* of tilings. Here are the first few examples of tilings.

$$p_0(w) = 1(?)$$

 $p_1(w) = w$
 $p_2(w) = w^2 + 1$
 $p_3(2) = w^3 + 2w$.

Can we find a recurrence for $p_n(w)$ similarly to how we found a recurrence for ordinary tilings of the $2 \times n$ rectangle? It turns out that we can. Here is the reasoning: We can decompose any $2 \times n$ tiling as a tiling where there is one vertical tile on the left or 2 horizontal tiles on the right. But for each tiling of a $2 \times (n-1)$ rectangle the corresponding

 $2 \times n$ tiling will have one more vertical tile. For the $2 \times (n-2)$ rectangles no new vertical tiles are added. It follows that the recurrence for these polynomials is

$$p_n(w) = wp_{n-1}(w) + p_{n-2}(w).$$

This is the essence of weighted enumeration.

We have observed above that $p_n(1)$ is the total number of tilings. Observe that $p'_n(1)$, the formal derivative of $p_n(w)$ evaluated at 1, is the total number of **all** vertical tiles in all the $2 \times n$ tilings. To see this we need only observe that w^k represents a tiling with k vertical dominoes. When you take the derivative you get kw^{k-1} , which evaluated at w = 1 gives you the number of vertical dominoes in that tiling.

It follows that the average or expected number of vertical tiles in a $2 \times n$ tiling is $\frac{p'_n(1)}{p_n(1)}$. We already know a formula for $p_n(1)$ (as a Fibonacci number). What about the numerator? Here is the magic trick.

Let

$$P(w,x) = \sum_{n=0}^{\infty} p_n(w)x^n$$

= 1 + wx + (w² + 1)x² + (w³ + 2w)x³ + ...

as a formal power series in $\mathbb{Q}[w][[x]]$. Using the recurrence for the weighted polynomials we can calculate that

$$P(w,x)(q - wx - x^2) = 1 \implies P(w,x) = \frac{1}{1 - wx - x^2}$$

as the generating function for P(w, x). Observe now (by differentiation with respect to w in the formal power series) that

$$\sum_{n=0}^{\infty} p'_n(w)x^n = \frac{d}{dw} \sum_{n=0}^{\infty} p_n(w)x^n$$
$$= \frac{d}{dw} P(w, x)$$
$$= \frac{x}{(1 - wx - x^2)^2}.$$

So we get a formula for $p'_n(1)$. Letting α and β be the reciprocal roots for $1-x-x^2$ we get a formula of the form $A\alpha^n + B\beta^n + A'n\alpha^n + B'n\beta^n$ for $p'_n(1)$. Letting α be the larger root (in magnitude) this means $p'_n(1) \sim Cn\alpha^n$ for some constant C. It follows that approximately,

$$\frac{p_n'(1)}{p_n(1)} = \frac{C_1 n \alpha^n}{C_2 \alpha^n} = Cn$$

for some constant C.

Now we consider the general problem of tiling the $2 \times n$ grid where each horizontal domino is given weight xt and each vertical domino is given weight yt, and the weight of any given

tiling is the product of the tile weights. Then what is the sum of all the weights of all the tilings?

For example, when n = 1, the sum of the weights is just yt, when n = 2 it is $(x^2 + y^2)t^2$ and when n = 3 it is $(x^2y + x^2y + y^3)t^3$. Now instead of grouping the sum of tiles by the exponent of t we will instead use the number of blocks, where a block consists of either a vertical domino or two horizontal dominoes. By the multiplication principle for generating functions the generating function looks like

$$\sum_{n=0}^{\infty} (x^2t^2 + yt)^n = \frac{1}{1 - x^2t^2 - yt}.$$

For example, if x = 1 and y = 1 then the generating function is as expected (the fibonacci numbers).

8 Path Enumeration, Directed Acyclic Graphs

Here is another problem to consider: How many length n sequences of 0's and 1's are there such that no two ones appear consecutively? We might calculate them out for short sequences and we might discover that they are equal to a shift in the Fibonacci numbers. We will solve this by considering paths through a network. A sequence is specified by paths through the network. It should be clear that each sequence is specified by a path in the graph.

(add the graph here)

We will solve this essentially using dynamic programming. Suppose we would like to calculate the paths from the source vertices on the left to the sink vertex s. We can first calculate the number of paths from the vertices directly adjacent to s, for each of those vertices there is only one path. Then for the vertices directly left of those, for the top vertex we can go to either vertex, so there are 1+1 paths to s. For the bottom vertex we can only travel to the top right vertex so there is only 1 path to s. And so on.

We might notice that we can represent the following operation by the matrix multiplication

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

and we get the number of paths of the second right most vertices to s and to t respectively. Then we can iterate by using the matrix multiplication

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^k \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

and the total number of paths to either s or t is given by

$$\begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix}.$$

In this section we will generalize this result sufficiently to work with graphs that are similar to the one above.

Definition 6. A finite directed acyclic graph (DAG) is a directed graph with only a finite number of paths between vertices and no infinite paths (no cycles). For vertices x, y in G we denote of paths from x to y by N(x, y). Note that in general $N(x, y) \neq N(y, x)$ (and in fact, if $N(x, y) \neq 0$ it follows that N(y, x) = 0).

Theorem 12. Suppose G is a finite DAG and $x, y \in G$. Suppose that B is a non-empty set of vertices with the property that every path from x to y passes through exactly one vertex in B. Then

$$N(x,y) = \sum_{b \in B} N(x,b)N(b,y)$$

More generally, suppose A is the set of sources and C is the set of sinks of some DAG, and that some set B has the property such that for every path from A to C passes exactly once through B. Define the gransfer matrix N(A, B) as the matrix whose (a, b)th entry is N(a, b) and define N(B, C) and N(A, C) respectively. Then

$$N(A,C) = N(A,B)N(B,C).$$

Proof. The first result follows from the divide and conquer principle and the principle of multiplication. The second claim directly follows from the first essentially by the definition of matrix multiplication.

Theorem 13. Suppose the vertex set of the DAG G is of the form $V(G) = V_0 \cup V_1 \cup \cdots \cup V_m$ where V_i and V_j are disjoint for $i \neq j$ and every edge goes from V_i to V_{i+1} for some i. Suppose also that each path from V_0 to V_m passes through exactly one vertex in each of the sets V_j , and that the pattern of connections from V_i to V_{i+1} is independent of i. That is, there exists a graph isomorphism $f: V \to (V_1 \cup V_2 \cup \cdots)$, such that

$$V_0 \mapsto V_1 \mapsto V_2 \mapsto \cdots$$
.

Then the transfer matrix from V_i to V_{i+1} is independent of i, and the transfer matrix from V_0 to V_m is equal to the mth power of this matrix.

Theorem 14. Let V be a semi-infinite (ie, extending infinitely rightwards but not leftwards) with the symmetry property of the last theorem. Fix a vertex x in V_0 and y_0, y_1, y_2 where $y_j \in V_j$, respectively, which are all symmetrical vertices. The $N(x, y_0), N(x, y_1), \ldots$ is a sequence that satisfies a dth order linear recurrence equation where $d = |V_i|$.

To prove this we will use the Cayley Hamilton

Theorem 15. Let M be a matrix with characteristic polynomial $p(x) = \det(xI - M)$. Then p(M) = 0.

We will not prove this theorem, it is a fundamental theorem of linear algebra. A nice computation to do is to directly calculate it in the case when M is a 2×2 matrix. Here is a corollary which can be thought of as a combinatorial interpretation of the Cayley Hamilton Theorem.

Corollary 1. Write $p(t) = t^m + a_{m-1}t^{m-1} + \cdots + a_0$ where p(t) is as above. Then the sequence of matrices I, M, M^2, M^3, \ldots satisfy a linear recurrence relation

$$M^{n+m} + a_{m-1}M^{n+m-1} + \dots + a_0M^n = 0$$

for all $n \geq 0$.

It can clearly be seen that each entry of the matrix as a component also satisfies this recurrence relation. Now we are ready to prove the theorems above:

Proof. Let M be the transfer matrix of connections between V_0 and V_1 (M is also the transfer matrix from V_i to V_{i+1}). Then $N(x, y_k)$ is the (x, y)th entry of M^k (by definition of the transfer matrix). Applying the corollary above gives us the desired results.

9 Change of Basis and the Stirling Numbers

In these notes we have introduced (one implicitly, one explicitly) two different bases for the vector space of polynomials P[x]. First we have the standard monomial basis

$$1, x, x^2, x^3, \dots$$

and then we introduced the falling factorial basis

$$1, (x), (x)_2, (x)_3, \dots$$

where

$$(x)_k = x(x-1)(x-2)\cdots(x-k+1).$$

It is clear that both sets of polynomials form bases. Moreover, due to the nature of the monomial basis, the change of basis "matrix" (in the sense of an infinite matrix) is upper triangular (so working with it we have no analytic difficulties). It makes sense to consider the coefficients of a change of basis formula. We first define s(n,k) to be the coefficients such that

$$(x)_n = \sum_{k=0}^{\infty} s(n,k)x^k$$

is true. Note that s(n, k) = 0 if k > n. Similarly, we can write

$$x^n = \sum_{k=0}^{\infty} S(n,k)(x)_k.$$

Note as well that S(n,k) = 0 if k > n. If we substitute the first equation into the second we get

$$x^{n} = \sum_{k=0}^{\infty} S(n,k)(x)_{k}$$
$$= \sum_{k=0}^{\infty} S(n,k) \sum_{j=0}^{\infty} s(k,j) x^{j}$$
$$= \sum_{j=0}^{\infty} \left[\sum_{k=0}^{\infty} S(n,k) s(k,j) \right] x^{j}.$$

By linear independence it follows that

$$\sum_{k=0}^{\infty} S(n,k)s(k,j) = \begin{cases} 1 & n=j\\ 0 & \text{otherwise} \end{cases}.$$

This is perhaps not so surprising on its own: This actually always happens when you consider any change of basis from one basis back to itself. What will be surprising is that s and S have combinatorial meaning. For example, when we derived a formula for the Bell Numbers, we realized S(n,k) as the number of ways of partitioning $\{1,\ldots,n\}$ into exactly k disjoint subsets. But what about s(n,k)?

First, we must observe that s(n,k) is not always positive. It has precisely the same sign as $(-1)^{n-k}$, and we call it a *signed Stirling number of the first kind*. On this note we will denote c(n,k) to be the coefficient of x^k in the polynomial

$$r_n(x) = x(x+1)(x+2)(\cdots)(x+n-1).$$

Proposition 4. For all positive integers n, k,

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k).$$

Proof. From the fact that

$$r_n(x) = r_{n-1}(x)(x+n-1) = xr_{n-1}(x) + (n-1)r_{n-1}(x)$$

this fact easily follows by degree counting.

Theorem 16. Let S_n be the symmetric group on $\{1, 2, ..., n\}$. Then c(n, k) equals the number of elements of S_n with exactly k cycles.

Proof. Let d(n,k) be the number of elements of S_n with exactly k cycles. We will show that d(n,k) follows the same recurrence as c(n,k) with the same initial conditions.

When n = 0, k = 0 we have that c(n, k) = 1. When $k = 0, n \neq 0$ we have that c(n, k) = 0. Similarly, when n = 0, k = 0 there is only one permutation with 0 cycles, which is the empty permutation, so d(0,0) = 1. If $n \neq 0$ every permutation has at least one cycle so that d(n,0) = 1, as desired.

Suppose that σ is any permutation of [n-1] with k cycles. Then there are n-1 ways to add the element n into an existing cycle (We add n to a cycle such that $\tilde{\sigma}(n) = i$ for each

 $1 \le i \le n-1$. Also, there's a unique way to create a unique cycle consisting of just the element n in any permutation of [n-1] with k-1 cycles. This describes every permutation of [n] with k cycles. It follows that

$$d(n,k) = (n-1)d(n-1,k) + d(n-1,k-1).$$

Now we will discuss the Stirling number of the second kind S(n, k).

Proposition 5. We have for all n and k that S(n,k) = S(n-1,k-1) + kS(n-1,k).

Proof. Given a partition of [n] into k blocks, removing n gives a partition of [n-1] either with k blocks or with k-1 blocks. Conversely, every partition of [n-1] into k-1 blocks gives rise to a partition of [n] into k blocks (by adding $\{n\}$) and every partition of [n-1] into k blocks gives rise to k partitions of [n] into k blocks. The recurrence is seen to hold.

Here is a claim about the generating function of the Stirling numbers of the second kind:

Proposition 6. For $k \geq 0$ we have that

$$\sum_{n \ge k} S(n-k)x^n = \frac{x^k}{(1-x)(1-2x)\cdots(1-kx)}.$$

Proof. Sketch: The proposition is true for k = 1. If we denote $S_k(x) = \sum_{n \geq k} S(n-k)x^n$, we can use the recurrence relation to show that

$$S_k(x) = xS_{k-1}(x) + kxS_k(x).$$

Rewriting this as

$$S_k(x) = \frac{x}{1 - kx} S_{k-1}(x)$$

the inductive hypothesis will prove the claim.

As a product of simple looking generating functions the above proposition might have a simpler combinatorial proof. We will see how this will work out in the next section.

9.1 The Combinatorial Meaning of Stirling Orthogonality

Above we deduced the formula

$$\sum_{k=0}^{\infty} S(n,k)s(k,j) = \begin{cases} 1 & n=j\\ 0 & \text{otherwise} \end{cases}$$

For general n, m, and k. In this subsection we will show that this formula has combinatorial meaning.

To show the sum is 0, we will use the following principle: suppose that S is a set with a weight function on S and there is some involution $\Phi: S \to S$ which satisfies $W(\Phi(s)) = -W(s)$ for all $s \in S$. Then

$$\sum_{s \in S} W(s) = 0.$$

To show Stirling orthogonality, then, we will find some sign reversing involution between a combinatorial object.

Definition 7. A gadget of type (n, m, k) (where $n \ge m \ge k$) is a tuple consisting of a partition of [n] into m blocks B_1, \ldots, B_m and a permutation of $\{B_1, \ldots, B_m\}$ into k cycles. We will define the weight of such a gadget to be $(-1)^{m-k}$.

So in this way, gadgets realize the Stirling Numbers of both kinds in one setup. With the weight we will find a specific sign reversing involution; then we will be done.

First, we observe that the sum of the weights of all the (n, m, k) gadgets is S(n, m)s(m, k). This follows from the fact that the weight of any (n, m, k) gadget is specified by the sign of the stirling number s(m, k). For the sign-reversing involution, we say a cycle in the gadget γ is *ample* if it involves more than one element of [n], that is, either the cycle contains greater than 1 block or it contains a block with more than 1 element.

Since n > k there is at least one ample cycle in γ . Let x be the smallest element of [n] that lies in an ample cycle. If x is in a block with other elements of [n] rip x off to form a new block of size n, and put it before the old block of the cycle. For example, if γ is of the form

$$\gamma = \dots (\dots \{x, a, b, \dots \} \dots) \dots$$

then we rip off x to create the cycle

$$\gamma' = \dots (\dots \{x\}\{a, b, \dots\} \dots) \dots$$

Otherwise, if x is in a block by itself, then merge $\{x\}$ with whatever follows it in its cycle.

Then this defines a sign reversing involution, for it changes the number of blocks by 1, so the weight of each will be negative. So we're done in this case.

10 The addition and multiplication principle for Generating Functions

In this section we'll discuss the addition and multiplication principle for generating functions. They are quite simple.

Suppose some sequence a_n of a type A is either of type B or type C (but not both). Then $a_n = b_n + c_n$ for all n. It follows that if A(x), B(x), C(x) are their generating functions then A(x) = B(x) + C(x). This is the Addition Principle for Generating Functions.

Similarly, suppose some sequence a_n of type A can be written in a unique way as a b_k sequence followed by a c_{n-k} sequence. Then A(x) = B(x)C(x). This is the Multiplication Principle for Generating Functions.

These considerations suggest a combinatorial proof of the last proposition of the previous section:

Proof. Given a partition of [n] into k parts, let B_1 be the part that contains 1, B_2 be the part that contains the smallest number not in B_1 , and so on. Every partition of [n] with exactly k parts can be specified by a sequence of length n where the ith symbol is j if i belongs to B_j .

For example, $\{\{1,3\},\{2,5\},\{4\}\}$ can be associated with the sequence (1,2,1,3,2). Conversely, every sequence of length n corresponds to a unique partition provided that $1,\ldots,k$ are used at least once and j+1 does not appear in the sequence until the first j appears.

Then we can break up each sequence into k parts, where the kth part is everything from the first k up to and right before the first k+1 (if k+1 appears at all). For example, in the sequence (1,2,2,1,3,2,4,2,5,1,2,3,6,5,6) This becomes (1|2,2,1|3,2|4,2|5,1,3|6,5,6). We observe that the jth part consists of j followed by elements between 1 and j inclusive. The number of possible strings of this type of length n is j^{m-1} . Hence, the generating function for this is x/(1-jx). By the Multiplication principle for generating functions, the generating function for these sequences is

$$\frac{x}{(1-x)}\cdots\frac{x}{(1-kx)},$$

proving the claim.

11 The Catalan Numbers and their generating function

We will introduce the Catalan numbers in terms of one of the simplest formulations of object that they count. Let's say that a sequence of n 1's and -1's is **suitable** if every partial sum is greater than or equal to 0. Call such a sequence

- Trivial, if n = 0,
- Primitive if the partial sums are greater than 0 except for the empty and total sums,
- Composite if some partial sum is 0.

Let s(n), p(n), and c(n) denote these sequences respectively, and let S(x), P(x), and C(x) (bad notation, I know) denote the generating functions for these sequences. We have that

$$s(n) = p(n) + c(n)$$

for n > 1. Then it follows that

- 1. S(x) = 1 + P(x) + C(x) (where the 1 is added to account for the trivial sequence.)
- 2. We have

$$C(x) = P(x)^{2} + P(x)^{3} + P(x)^{4} + \dots = P(x)[P(x) + C(x)]$$

as the consideration that any composite sequence can be thought about as a concatenation of primitive sequences. The latter equality also follows from the fact that each composite sequence starts with either a prime sequence and then follows with either a prime or composite sequence. Using this latter equation we deduce the equality

$$C(x) = \frac{P(x)^2}{1 - P(x)}.$$

3. Observe that P(x) = xS(x), since prime sequences are essentially determined by the middle portions. It follows that p(n) = s(n-1) for $n \ge 1$.

From these 3 considerations we conclude that

$$S(x) = 1 + P(x) + \frac{P(x)^2}{1 - P(x)}$$
$$= 1 + \frac{P(x)}{1 - P(x)}$$
$$= \frac{1}{1 - xS(x)}.$$

Rearranging we get the quadratic equation

$$xS(x)^{2} - S(x) + 1 = 0$$

for which if we solve it we get

$$S(x) = \frac{1 - \sqrt{1 - 4x}}{2x}$$

(for if we take the + root S(x) is no longer a power series but a Laurent series). So

$$S(x) = \frac{1 - (1 - 2x - 2x^2 - 4x^3 - \dots)}{2x}.$$

We can do better and get an explicit formula: (derivation goes here)

and hence we see that $C_n = \frac{1}{n+1} {2n \choose n}$. This combinatorial coefficient here suggests we can also prove this formula directly using some kind of equivalence class.

12 Enumeration using the trace function

The principle here follows from this linear algebraic result:

Theorem 17. Let A be a square matrix. Define $A_{ij}(n)$ denote the (i, j)th entry of A^n . Define the generating function

$$F_{ij}(t) = \sum_{n>0} A_{ij}(n)t^n.$$

Then

$$F_{ij}(t) = \frac{(-1)^{i+j} \det(I - TA; j, i)}{\det(I - tA)}$$

where (B; j, i) is the matrix obtained by removing the jth row and the ith column in B. In particular, $F_{ij}(t)$ is a rational function of t.

Proof. Use Cramer's rule on the matrix product $(I - tA)^{-1}$. I haven't figured out the details yet.

Now we will introduce the trace theorem, which has numerous counting applications.

Theorem 18. Define

$$T(t) = \sum_{n>1} \operatorname{Tr}(A^n) t^n.$$

Then

$$T(t) = \frac{-tQ'(t)}{Q(t)}$$

where $Q(t) = \det(I - tA)$.

Proof. Let $\omega_1, \ldots, \omega_r$ be the non-zero eigenvalues of A. Then

$$\operatorname{Tr}(A^n) = \omega_1^n + \dots + \omega_r^n.$$

So

$$T(t) = \frac{\omega_1}{1 - \omega_1 t} + \dots + \frac{\omega_r}{1 - \omega_r t}.$$

When you divide this by $Q(t) = (1 - \omega_1 t) \cdots (1 - \omega_r t)$ the numerator becomes -tQ'(t).

We might use this theorem, for example, to count the domino tilings of a $2 \times n$ cylinder using the codeword method. We might introduce a cylinder with an identifying vertical line and count the trace (the number of ways that things can go back to each other and whatnot).

13 Eventually Recurrent Sequences and a Characterization of Rational Generating Functions

In the previous section, the trace formula had a rational generating function where the degree of the denominator was equal to that of the degree of the numerator. In this section, we will emphasize that such a function cannot satisfy a linear recurrence relation. However, such a sequence is actually "eventually recurrent" which we will show below in general with rational generating functions. It boils down to these theorems:

Theorem 19. A sequence $f(0), f(1), f(2), \ldots$ which satisfies a linear recurrence equation of the form

$$f(n+d) + a_1 f(n+d-1) + \dots + a_d f(n) = 0$$

for all $n \geq 0$ has a generating function of the form

$$\sum_{n\geq 0} f(n)x^n = \frac{P(x)}{Q(x)}$$

Where $Q(x) = 1 + a_1x + \cdots + a_dx^d$ and $\deg(P(x)) < d$.

This was already proven. Here is the other rather trivial theorem:

Theorem 20. A sequence $f(0), f(1), f(2), \ldots$ that satisfies f(n) = 0 for all but finitely many values for n has a generating function of the form $\sum_{n\geq 0} f(n)x^n = P(x)$ for some polynomial P with degree being the index of the maximum non-zero element of the sequence.

Combining, we obtain the following result.

Theorem 21. Let $f: \mathbb{N} \to \mathbb{C}$ be a sequence and suppose that

$$\sum_{n>0} f(n)x^n = \frac{P(x)}{Q(x)}$$

for polynomials P and Q.

Then there is a function $g: \mathbb{N} \to \mathbb{C}$ that agrees with f outside a finite set of values (which we denote by the exceptional set) such that

$$\sum_{n>0} g(n)x^n = \frac{R(x)}{Q(x)}$$

with deg $R < \deg Q$. If the exceptional set is non-empty, then its largest element is deg P – deg Q, and the generating function for f(n) - g(n) is a polynomial of degree deg(P) – deg(Q).

We say that such a sequence f is "eventually recurrent".

Applying this theorem to $f(n) = \text{Tr}(A^n)$ as we have there is a unique way to define f(0) so that it satisfies the linear recurrence equation for n = 0 as well as n > 0. Namely, this number is the number of non-zero eigenvalues of A (including multiplicity). Does this number have any combinatorial significance?

14 LREs indexed over the integers and Popoviciu's Lemma

Given a linear recurrence relation, we can run it backwards over the integers to make it over all the integers, instead of just the natural numbers. For example, for the Fibonacci sequence it looks like

$$\dots, 5, -3, 2, -1, 1, 0, 1, 1, 2, 3, 5, \dots$$

Notice that the left half looks suspiciously similar to the right half. And in fact, for any linear recurrence relation it will always turn out like this! This is known as **Popoviciu's Lemma**:

Theorem 22. Let d be a positive integer and a_1, \ldots, a_d be complex coefficients with $a_d \neq 0$. Suppose $f : \mathbb{Z} \to \mathbb{C}$ satisfies

$$f(n+d) + a_1 f(n+d-1) + \dots + a_d f(n) = 0$$

for all $n \in \mathbb{Z}$. Then

$$F(x) = \sum_{n>0} f(n)x^n$$

and

$$G(x) = \sum_{n \ge 1} f(-n)xn$$

are both rational functions satisfying

$$G(x) = -F\left(\frac{1}{x}\right).$$

For example, consider $f(n) = c^n$ for $c \neq 0$. Then

$$F(x) = 1 + cx + c^{2}x^{2} + \dots = \frac{1}{1 - cx},$$

$$G(x) = c^{-1}x + c^{-2}x^{2} + \dots = \frac{c^{-1}x}{1 - c^{-1}x} = \frac{1}{\frac{c}{x} - 1} = -\frac{1}{1 - c\left(\frac{1}{x}\right)} = -F\left(\frac{1}{x}\right).$$

As another example, we might consider $f(n) = \binom{n}{k}$ for $k \geq 0$.

Recall that

$$\binom{n}{k} - \binom{n-1}{k} = \binom{n-1}{k-1}.$$

If we multiply by x^n and sum over $n \geq 0$ we get

$$F_k(x) - xF_k(x) = xF_{k-1}(x)$$

Or in other words

$$F_k(x) = \frac{x}{1-x} F_{k-1}(x) \implies F_k(x) = \frac{x^k}{(1-x)^{k+1}}.$$

We have as well that

$$g(n) = \binom{-n}{k} = (-1)^k \binom{n+k-1}{k}.$$

If we multiply by x^n and sum over all $n \geq 1$ we get

$$G(x) = (-1)^k \frac{x}{(1-x)^{k+1}}$$

from which we can verify Popoviciu's lemma easily.

In general, the general case sort of follows from these two examples. For as in the theorem, f, F, and G are mutually determining in a linear way, and the determining relations are also linearly related. So it suffices to prove the lemma where $F(x) = \frac{1}{(1-cx)^k}$ for $c \neq 0, k \geq 1$ since these form a basis for the space of rational functions of x of the form $\frac{P(x)}{Q(x)}$ and $\deg(P) < \deg(Q)$. Or one could see page 206 of Stanley's Enumerative Combinatorics.

15 The Reflection Principle

In this section we describe the *reflection principle*, which is a bijective principle that comes up when studying graph routings and is very visual and illuminative.

First, let's consider again one of the ways to count the Catalan Numbers. Let C_n be the number of Dyck paths from (0,0) to (2n,0) with steps (1,1) and (1,-1) which do not touch or cross the line y=-1.

If we do not have this restriction of intersecting the line y = -1 then this number is $\binom{2n}{n}$. Then the number of non-Dyck paths would be

$$\binom{2n}{n} - \frac{1}{n+1} \binom{2n}{n} = \frac{n}{n+1} \frac{(2n)!}{n!n!} = \binom{2n}{n-1}.$$

Motivated by the principle "when there's a nice answer, there's a nice reason", we will find a purely combinatorial proof of the Catalan Number sequence.

Consider any non-Dyck path, it is characterized by the fact that it intersects the line y = -1 eventually. Let P be the first point on the path that lies on y = -1. Then create a new path by reflecting everything before P over the line y = -1. Then we get a new path, but this time from (0, -2) to (2n, 0). This reflection operation is clearly injective, for the operation of reflecting everything before P over y = -1 is its own inverse. Conversely, any path from (0, -2) to (2n, 0) must cross the line y = -1. So it follows that there is a non-Dyck path for which its reflection would be the arbitrary path being considered. It follows that the number of non-Dyck paths is counted by $\binom{2n}{n-1}$ (this is the number of total paths from (0, -2) to (2n, 0), as such a path must use n - 1 - 1's.

16 Tilings and Routings, the Exchange Principle

It's well known that given an $m \times n$ grid, the number of paths where one can travel along the basis vectors (1,0) and (0,1) is $\binom{m+n}{n}$. This approach is also easily generalized for lattice paths in three dimensions: the number of lattice paths of length (a,b,c) from (0,0,0) to (a,b,c) is the multinomial coefficient

$$\binom{a+b+c}{a,b,c} = \frac{(a+b+c)!}{a!b!c!}.$$

But when we move to three space, we have an extra dimension to work with. So now we might want to consider the number of lattice surfaces. Consider the following problem: what is the number of lattice surfaces of area ab + bc + ac whose boundary is the non-planar hexagon with vertices

$$(a, 0, 0), (a, b, 0), (0, b, 0), (0, b, c), (0, 0, c), (a, 0, c)$$

As a special case for example, what is the number of Art of Problem Solving Default Avatars? We will prove later that the number of such surfaces is

$$\frac{H(a+b+c)H(a)H(b)H(c)}{H(a+b)H(a+c)H(b+c)}$$

where H(n) is the hyperfactorial $1!2!3!\cdots(n-1)!$. We'll prove this by bijection to a two dimensional picture (much like the AoPS default avatars) and we will show that this is the number of ways of tiling an (a, b, c, a, b, c) hexagon with unit rhombuses. For c = 0 the number of tilings is 1 (since the placement of rhombuses is forced).

When c=1, the number of tilings is $\frac{(a+b)!}{a!b!}$. This looks very familiar! This can be illustrated by putting some dots through the vertical sections of the tilings. If a dot hits a vertical line, then we draw a path through it. In this way, we actually always get a lattice path.

With $2 \times n$ domino tilings, we can place alternating dots in a certain way and we will so that the tilings count lattice paths of the directed acyclic graph (put the graph here).

For the $3 \times n$ domino tilings, we might place our dots like so (here) in which case the tilings correspond to paths from s to t.

But alternatively for this tiling, the tilings instead correspond to non-intersecting lattice paths (called 2-routings). The number of such paths will be $N(s_1, t_1)N(s_2, t_2)-N(s_1, t_2)N(s_2, t_1)$. To prove this we can use the reflection principle to show that bad rountings where paths intersect can be directly bijected with paths which cross over each other.

Here is the exchange principle, formally stated:

Theorem 23. Suppose s_1 and s_2 are sources in a directed acyclic graph and t_1 and t_2 are sinks with the property that a path that joins s_1 to t_2 and a path that joins s_2 to t_1 must have a vertex in common.

Then the number of 2-routings that join s_1 to t_1 and s_2 to t_2 with non-intersecting paths is equal to $N_{11}N_{22} - N_{12}N_{21}$ where N_{ij} is the number of paths from s_i to t_j .

17 Coins in a Fountain

Let's return to the subject of Dyck paths. First, let's weigh each Dyck path according to the area between it and the horizontal axis. More precisely, we say that a Dyck path has weight $q^m x^n$ if its length is 2n and the area under the path is m.

Let F(x,q) be the sum of all the weights of all the Dyck paths. Then

$$F(x,q) = 1 + qx + (q^2 + q^4)x^2 + (q^3 + 2q^5 + q^7 + q^9)x^3 + \cdots$$

We claim that

$$F(x,q) = q + xq(F(xq^2, q)F(x, q).$$

Proof. If a Dyck path from (0,0) to (2k,0) has weight $q^m x^k$, then when it is shifted to give a path from (1,1) to (2k+1,1) it contributes $q^{2k}q^m x^k$ to the weight of a longer Dyck path.

From the above equation we obtain the identity

$$F(x,q) = \frac{1}{1 - xqF(xq^2,q)}, F(xq^2,q) = \frac{1}{1 - xq^3F(xq^4,q)}, \dots$$

giving us a continued fraction identity.

18 Lattice Animals, Polyominoes, Dimers

In this section we will highlight some results made in the combinatorial study of *lattice* animals.

Definition 8. A polyomino is a non-empty union of lattice squares with connected interior.

Theorem 24. Let a_n be the number of polyominos of area n. Then $a_n^{1/n}$ convergest to some limit μ (called Klarner's constant).

By numerical experiements it seems like $\mu \approx 4.06$. The best known bounds for μ are $4.002 < \mu < 4.65$.

From a polyomino we can associate it with a lattice animal: (image here).

Definition 9. A directed animal is an animal in which there exists a special node called the source such that every other point can be reached from the source using a NE path (a path which goes only north or east).

We know better the asymptotics for directed animals, by this theorem of Dhar.

Theorem 25. Let b_n be the number of directed animals with n nodes. Then $b_n^{1/n} \to 3$.

This was proven by Mélou and Rechnitzer. The idea is to associate these directed animals with "dimer stacks" and let them fall by gravity. These dimers form a heap.

Definition 10. A pyramid is a non-empty heap of dimers with a unique minimal element. A half-pyramid is a pyramid in which no dimer lies to the right of the dimer on the bottom.

Theorem 26. This (where this will be explained eventually) mapping gives a bijection between directed animals and pyramids of dimers.

Theorem 27. Let

$$P(x) = x + 2x^2 + 5x^3 + 13x^4 + 35x^5 + \cdots$$

be the generating function for nonempty pyramids of n dimers, and

$$Q(x) = x + x^2 + 2x^3 + 4x^4 + 9x^5$$

be the generating function for half pyramids. Then

$$P(x) = Q(x) + P(x)Q(x)$$

and

$$Q(x) = x + xQ(x) + xQ(x)^{2}.$$

Then we can solve and we find that

$$Q(x) = \frac{1 - x - \sqrt{(1+x)(1-3x)}}{2x},$$
$$P(x) = \frac{\sqrt{\frac{1+x}{1-3x}} - 1}{2}.$$

Proof. For the first equation, we observe that any pyramid is either a half pyramid, or a pyramid which is sitting on top of a half pyramid. For the second equation we decompose the half pyramids. Either we have one dimer, we have one dimer and sitting on top of it is a half pyramid, or sitting on top of it is a pyramid. The third row right entry must be a half pyramid (or else the pyramid would not be a half pyramid). These give the generating functional relations as desired.

Finally, we use the following theorem from complex analysis:

Theorem 28. If $a_0 + a_1x + a_2x^2 + \cdots$ is the Taylor expansion about z = 0 of f(z) that is analytic with a pole at |z| = r, then the series has radius of convergence r and $a_n^{1/n} \to 1/r$ as $n \to \infty$.

We note that P(x) has a branch point at x = -1 and a pole at x = 1/3. Dhar's Theorem immediately follows, as desired.

19 Short recurrences and contractions

Returning to Catalan Numbers, we will consider the Triangulation formulation of the Catalan Numbers. Let T(n) be the number of triangulations of a convex (n+2)-gon. Directly considering the formula $\frac{1}{n+1}\binom{2n}{n}$ we get the **short recurrence**

$$(n+1)T(n) = (4n-2)T(n-1).$$

We can interpret this result using combinatorial arguments. First, given any triangulation, we define an *edge* as any side or diagonal. When n=3 there are 7 edges. In general there are 2n+1 edges. We will distinguish one side of the (n+2)-gon as the base. The left hand side counts triangulations of the (n+2)-gon with one side marked other than the base. The right hand side counts triangulations of an (n+1)-gon with one of 2n-1 edges marked and oriented.

The bijection is as follows: Given the marked edge which is not the base we "contract" it so that the marked line itself is identified as one point. We also identify the lines in the triangle that the red triangle shares with two unmarked lines. We identify the other two lines together and direct it in the direction of the vertex which used to be the red line.

Conversely, given a directed (n + 1)-gon, we create an edge out of the target vertex and split the directed line into two. One can check that these maps are well defined and are inverses of each other, and hence give a bijection between the two. (images needed).

To further show the power of the short recurrence method, we will use it to produce a recurrence for lattice paths with q-weight.

Given a lattice path from (0,0) to (m,n) then we define its weight as q to the power of the total area under the path (this is well defined and integer valued). We denote the sum of weights as $P_{m,n}(q)$. For example, $P_{2,2}(q) = q^0 + q^1 + 2q^2 + q^3 + q^4$.

To produce a short recurrence for the q-weighted lattice paths we will consider an ordinary binomial identity:

$$\binom{a}{b}b = a\binom{a-1}{b-1}.$$

This can be proven by committee leader enumeration. If a = m + n and b = m we can rewrite this identity as

$$\binom{m+n}{m}m = (m+n)\binom{m+n-1}{m-1}.$$

So we have that

$$mP_{m,n}(1) = (m+n)P_{m-1,n}(1).$$

This is a special case of the more general short recurrence.

Theorem 29. We have

$$[m]P_{m,n}(q) = [m+n]P_{m-1,n}(q)$$

where $[n] = 1 + q + q^2 + \cdots + q^{n-1}$. This is called the q-analogue of n.

Proof. We will count the left and right sides and exhibit a bijection.

The left hand side counts lattice paths from (0,0) to (m,n) with a marked horizontal edge where the area associated with that marked path is defined as the ordinary area plus the x-coordinate of the left endpoint of the chosen edge.

The right hand side counts lattice paths from (0,0) to (m-1,n) with a marked vertex where the area associated with a marked path is defined as the ordinary area plus the sum of the coordinates of the marked vertex.

The bijection is as follows: we contract the edge to a vertex, and expand the vertex right-wards to create an edge. Clearly it is a bijection, and moreover, it is **weight-preserving**. This prove the theorem.

From this theorem we get the short recurrence

$$P_{m,n} = \frac{[n+m]}{[m]} P_{m-1,n}.$$

Expanding this out we get

$$P_{m,n} = \frac{[n+m]}{[m]} \cdot \frac{[n+m-1]}{[m-1]} \cdot \dots \cdot \frac{[n+1]}{[1]}.$$

There are further questions about symmetrical lattice path analogs and q-enumerations. We can also consider the relationship between bracket numbers and cyclotomic polynomials. These will be added here if I think about them some more.

We can deduce some properties about $P_{m,n}$ through elementary or geometrical considerations. For example, $P_{m,n} = P_{n,m}$. This is clear because rotation by 45 degrees diagonally down will give you an appropriate lattice path. We also have that $P_{m,n}$ is palindromic, this is clear because rotation by 180 will give you a lattice path with area mn minus the original area.