

In the movie Die Hard 3, the heros must obtain exactly 4 gallons of water using a 5 gallon jug, a 3 gallon jug, and a water faucet. Our goal: to get *TLC* to solve the problem for us.

First, we write a spec that describes all allowable behaviors of our heros.

EXTENDS *Integers*

This statement imports the definitions of the ordinary operators on integers, such as +.

This statement defines the sizes of the jugs as well as the target amount of water to defuse the bomb.

CONSTANT	<i>BigJug</i> ,	The capacity of the bigger jug
	<i>SmallJug</i> ,	The capacity of the smaller jug
	<i>Goal</i>	The target amount of water

We next declare the specification's variables.

VARIABLES	<i>big</i> ,	The number of gallons of water in the 5 gallon jug.
	<i>small</i>	The number of gallons of water in the 3 gallon jug.

Now we define of the initial predicate, that specifies the initial values of the variables. I like to name this predicate *Init*, but the name doesn't matter.

Note: TLA+ uses the convention that a list of formulas bulleted by  $\wedge$  or  $\vee$  denotes the conjunction or disjunction of those formulas. Indentation of subitems is significant, allowing one to eliminate lots of parentheses. This makes a large formula much easier to read. However, it does mean that you have to be careful with your indentation.

$$\begin{aligned} Init &\triangleq \wedge big = 0 \\ &\quad \wedge small = 0 \end{aligned}$$

Now we define the actions that our hero can perform. There are three things they can do:

- Pour water from the faucet into a jug.
- Pour water from a jug onto the ground.
- Pour water from one jug into another

We now consider the first two. Since the jugs are not calibrated, partially filling or partially emptying a jug accomplishes nothing. So, the first two possibilities yield the following four possible actions.

$$\begin{aligned} FillSmall &\triangleq \wedge small' = SmallJug \\ &\quad \wedge big' = big \end{aligned}$$

$$\begin{aligned} FillBig &\triangleq \wedge big' = BigJug \\ &\quad \wedge small' = small \end{aligned}$$

$$\begin{aligned} EmptySmall &\triangleq \wedge small' = 0 \\ &\quad \wedge big' = big \end{aligned}$$

$$\begin{aligned} EmptyBig &\triangleq \wedge big' = 0 \\ &\quad \wedge small' = small \end{aligned}$$

We now consider pouring water from one jug into another. Again, since the jugs are not callibrated, when pouring from jug  $A$  to jug  $B$ , it makes sense only to either fill  $B$  or empty  $A$ . And there's no point in emptying  $A$  if this will cause  $B$  to overflow, since that could be accomplished by the two actions of first filling  $B$  and then emptying  $A$ . So, pouring water from  $A$  to  $B$  leaves  $B$  with the lesser of (i) the water contained in both jugs and (ii) the volume of  $B$ . To express this mathematically, we first define  $Min(m, n)$  to equal the minimum of the numbers  $m$  and  $n$ .

$$Min(m, n) \triangleq \text{IF } m < n \text{ THEN } m \text{ ELSE } n$$

Now we define the last two pouring actions. From the observation above, these definitions should be clear.

$$\begin{aligned} SmallToBig &\triangleq \text{LET } poured \triangleq Min(big + small, BigJug) - big \\ &\quad \text{IN } \begin{array}{l} \wedge big' = big + poured \\ \wedge small' = small - poured \end{array} \end{aligned}$$

$$\begin{aligned} BigToSmall &\triangleq \\ &\text{LET } poured \triangleq Min(big + small, SmallJug) - small \\ &\text{IN } \begin{array}{l} \wedge big' = big - poured \\ \wedge small' = small + poured \end{array} \end{aligned}$$

We define the next-state relation, which I like to call *Next*. A *Next* step is a step of one of the six actions defined above. Hence, *Next* is the disjunction of those actions.

$$\begin{aligned} Next &\triangleq \vee FillSmall \\ &\quad \vee FillBig \\ &\quad \vee EmptySmall \\ &\quad \vee EmptyBig \\ &\quad \vee SmallToBig \\ &\quad \vee BigToSmall \end{aligned}$$

We now define *TypeOK* to be the type invariant, asserting that the value of each variable is an element of the appropriate set. A type invariant like this is not part of the specification, but it's generally a good idea to include it because it helps the reader understand the spec. Moreover, having *TLC* check that it is an invariant of the spec catches errors that, in a typed language, are caught by type checking.

$$\begin{aligned} TypeOK &\triangleq \wedge small \in 0 \dots SmallJug \\ &\quad \wedge big \in 0 \dots BigJug \end{aligned}$$

Remember that our heros must measure out 4 gallons of water. Obviously, those 4 gallons must be in the 5 gallon jug. So, they have solved their problem when they reach a state with  $big = 4$ . We find a solution by having *TLC* check if  $big \neq 4$  is an invariant, which will cause it to print out an "error trace" consisting of a behavior ending in a states where  $big \neq 4$  is false—that is, ending in a state in which  $big = 4$ . Such a behavior is the desired solution. (Because *TLC* uses a breadth-first search, it will find the shortest solution.)