MODULE *TCommitJonRo*

CONSTANT RM The set of participating resource managers

VARIABLE $rmState$ $rmState[rm]$ is the state of resource manager r.

$TCTypeOK \triangleq$

The type-correctness invariant

$rmState \in [RM \to \{\text{“working”}, \text{“prepared”}, \text{“committed”}, \text{“aborted”}\}]$

$TCInit \triangleq \quad rmState = [r \in RM \mapsto \text{“working”}]$

The initial predicate.

$canCommit \triangleq \forall\, r \in RM : rmState[r] \in \{\text{“prepared”}, \text{“committed”}\}$

Can commit iff all RMs are in the “prepared” or “committed” state.

$canAbort \triangleq \forall\, r \in RM : rmState[r] \neq \text{“committed”}$

Can abort iff no RM has decided to commit.

We now define the actions that may be performed by the RMs, and then define the complete next-state action of the specification to be the disjunction of the possible RM actions.

$Prepare(r) \triangleq \land\, rmState[r] = \text{“working”}$

Any RM in the working state can go to the prepared state

$\land\, rmState' = [rmState \text{ EXCEPT } ![r] = \text{“prepared”}]$

$Commit(r) \triangleq \land\, rmState[r] = \text{“prepared”}$

$\land\, canCommit$ Only allowed to commit from prepared if we can commit

$\land\, rmState' = [rmState \text{ EXCEPT } ![r] = \text{“committed”}]$

$Abort(r) \triangleq \land\, rmState[r] \in \{\text{“working”}, \text{“prepared”}\}$

$\land\, canAbort$ Only allowed to abort from working or prepared if we can abort

$\land\, rmState' = [rmState \text{ EXCEPT } ![r] = \text{“aborted”}]$

$TCNext \triangleq \exists\, r \in RM : Prepare(r) \lor Commit(r) \lor Abort(r)$

The next-state action.

$TCSpec \triangleq TCInit \land \Box[TCNext]_{rmState}$

The complete specification of the protocol.

We now assert invariance properties of the specification.

$TCConsistent \triangleq$

A state predicate asserting that two RMs have not arrived at conflicting decisions.

$\forall\, r1, r2 \in RM : \neg \land\, rmState[r1] = \text{“aborted”}$
$\land\, rmState[r2] = \text{“committed”}$

THEOREM $TCSpec \Rightarrow \Box(TCTypeOK \land TCConsistent)$

Asserts that $TCTypeOK$ and $TCInvariant$ are invariants of the protocol.