A TYPE THEORY WITH SCOPED EQUALITY REFLECTION

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Feel free to change the stuff below as our thinking evolves. This is just a starting point. We have the following judgments:

 $\begin{array}{lll} \vdash \Sigma \ \mathrm{sig} & \Sigma \ \mathrm{is \ a \ valid \ signature} \\ \vdash^{\mathcal{H}}_{\Sigma} \Gamma \ \mathrm{ctx} & \Gamma \ \mathrm{is \ a \ valid \ context} \\ \Gamma \vdash_{\Sigma} \mathcal{H} \ \mathrm{hints} & \mathcal{H} \ \mathrm{is \ a \ valid \ hints \ set} \\ \Gamma \vdash^{\mathcal{H}}_{\Sigma} M : A & M \ \mathrm{has \ type} \ A \\ \Gamma \vdash^{\mathcal{H}}_{\Sigma} M \equiv N : A & M \ \mathrm{equals} \ N \ \mathrm{at \ type} \ A \end{array}$

The rules for the theory are given inductive-recursively below.

$$(\text{Signatures}) \qquad \qquad \frac{\vdash \Sigma \text{ sig} \quad \diamond \vdash_{\Sigma}^{\diamond} A : \mathbb{U} \quad \Sigma \# c}{\vdash \Sigma, c : A \text{ sig}}$$

$$(Contexts) \qquad \qquad \frac{\vdash^{\mathcal{H}}_{\Sigma} \Gamma \operatorname{ctx} \quad \Gamma \vdash^{\mathcal{H}}_{\Sigma} A : \mathbb{U} \quad \Gamma \# x}{\vdash^{\mathcal{H}}_{\Sigma} \Gamma, x : A \operatorname{ctx}}$$

$$\frac{\Gamma \vdash_{\Sigma} \mathcal{H} \text{ hints} \quad \Gamma \vdash_{\Sigma}^{\mathcal{H}} A : \mathbb{U} \quad \Gamma \vdash_{\Sigma}^{\mathcal{H}} a, b : A}{\Gamma \vdash_{\Sigma} \mathcal{H}, a \equiv b : A \text{ hints}}$$

$$(Projection) \qquad \frac{x: A \in \Gamma}{\Gamma \vdash_{\Sigma}^{\mathcal{H}} x: A} \qquad \frac{c: A \in \Sigma}{\Gamma \vdash_{\Sigma}^{\mathcal{H}} c: A} \qquad \frac{a \equiv b: A \in \mathcal{H}}{\Gamma \vdash_{\Sigma}^{\mathcal{H}} a \equiv b: A}$$

$$\frac{\Gamma \vdash^{\mathcal{H}}_{\Sigma} p : \operatorname{Id}_{A} a \ b \quad \Gamma \vdash^{\mathcal{H}, a \equiv b : A}_{\Sigma} e : C}{\Gamma \vdash^{\mathcal{H}}_{\Sigma} \operatorname{reflect} p \text{ in } e : C}$$