## A TYPE THEORY WITH SCOPED EQUALITY REFLECTION

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Feel free to change the stuff below as our thinking evolves. This is just a starting point. We have the following judgments:

$\vdash \Sigma \text{ sig}$	$\Sigma$ is a valid signature
$\mathcal{H} \vdash_{\Sigma} \Gamma \operatorname{ctx}$	$\Gamma$ is a valid context
$\Gamma \vdash_{\Sigma} \mathcal{H} \text{ hints}$	$\mathcal{H}$ is a valid hints set
$\Gamma; \mathcal{H} \vdash_{\Sigma} M \Rightarrow A$	M infers type $A$
$\Gamma; \mathcal{H} \vdash_{\Sigma} M \Leftarrow A$	M checks type $A$
$\Gamma; \mathcal{H} \vdash_{\Sigma} M \equiv N : A$	M equals $N$ at type $A$

The rules for the theory are given inductive-recursively below.

$$(Signatures) \qquad \overline{\vdash} \diamond sig \qquad \frac{\vdash \Sigma \text{ sig}}{\vdash \Sigma, c : A \text{ sig}} \qquad \Sigma \# c \qquad \qquad \overline{\vdash} \Sigma \text{ sig} \qquad \qquad \Sigma \# c \qquad \qquad \overline{\vdash} \Sigma, c : A \text{ sig} \qquad \qquad \overline{\vdash} \Sigma, c : A \text{$$

$$\begin{array}{c} \Gamma; \mathcal{H} \vdash_{\Sigma} p \Rightarrow \mathsf{Id}_{A}(M;N) \\ \Gamma; \mathcal{H}, M \equiv N : A \vdash_{\Sigma} C \Leftarrow \mathbb{U} \\ \Gamma; \mathcal{H}, M \equiv N : A \vdash_{\Sigma} e \Leftarrow C \end{array}$$
 (Reflect) 
$$\begin{array}{c} \Gamma; \mathcal{H} \vdash_{\Sigma} \mathsf{reflect} \ p \ \mathsf{in} \ e \Leftarrow C \end{array}$$

(Binders) 
$$\begin{array}{c} \Gamma; \mathcal{H} \vdash_{\Sigma} A \Leftarrow \mathbb{U} \\ \Gamma, x : A; \mathcal{H} \vdash_{\Sigma} [x]B \Leftarrow \mathbb{U} \end{array} \quad Q \in \{\Sigma, \Pi\} \\ \hline \Gamma; \mathcal{H} \vdash_{\Sigma} Q_{A}[x]B \Rightarrow \mathbb{U} \end{array}$$

(Sigma)

$$\begin{array}{ll} \Gamma; \mathcal{H} \vdash_{\Sigma} a \Leftarrow A & \Gamma; \mathcal{H} \vdash_{\Sigma} p \Rightarrow \Sigma_{A}[x]B \\ \Gamma; \mathcal{H} \vdash_{\Sigma} b \Leftarrow [a/x]B & \Gamma, u : \Sigma_{A}[x]B; \mathcal{H} \vdash_{\Sigma} [u]C \Leftarrow \mathbb{U} \\ \Gamma; \mathcal{H} \vdash_{\Sigma} \langle a, b \rangle \Leftarrow \Sigma_{A}[x]B & \Gamma; \mathcal{H} \vdash_{\Sigma} [v/x]B; \mathcal{H} \vdash_{\Sigma} [v, w]M \Leftarrow [\langle v, w \rangle / u]C \\ \hline \Gamma; \mathcal{H} \vdash_{\Sigma} \langle a, b \rangle \Leftarrow \Sigma_{A}[x]B & \Gamma; \mathcal{H} \vdash_{\Sigma} \mathrm{split}([u]C; [v, w]M; p) \Rightarrow [p/u]C \end{array}$$

$$(\text{Pi}) \qquad \frac{\Gamma, x : A; \mathcal{H} \vdash_{\Sigma} [x]M \Leftarrow [x]B}{\Gamma; \mathcal{H} \vdash_{\Sigma} \Lambda[x]M \Leftarrow \Pi_{A}[x]B} \qquad \frac{\Gamma; \mathcal{H} \vdash_{\Sigma} f \Rightarrow \Pi_{A}[x]B}{\Gamma; \mathcal{H} \vdash_{\Sigma} c \Leftarrow A}$$

$$\begin{array}{ccc} \Gamma; \mathcal{H} \vdash_{\Sigma} A \Leftarrow \mathbb{U} \\ \Gamma; \mathcal{H} \vdash_{\Sigma} M \Leftarrow A \\ \Gamma; \mathcal{H} \vdash_{\Sigma} N \Leftarrow A \\ \hline \Gamma; \mathcal{H} \vdash_{\Sigma} N \Leftarrow A \\ \hline \Gamma; \mathcal{H} \vdash_{\Sigma} \operatorname{Id}_{A}(M; N) \Rightarrow \mathbb{U} \end{array} \qquad \begin{array}{c} \Gamma; \mathcal{H} \vdash_{\Sigma} M \Leftarrow A \\ \hline \Gamma; \mathcal{H} \vdash_{\Sigma} \operatorname{Id}_{A}(M; M) \Rightarrow \mathbb{U} \end{array}$$

$$\begin{split} &\Gamma; \mathcal{H} \vdash_{\Sigma} \mathbf{Q}_{A}[x]B \Leftarrow \mathbb{U} \\ &\Gamma; \mathcal{H} \vdash_{\Sigma} \mathbf{Q}_{A}'[x]B' \Leftarrow \mathbb{U} \\ &\Gamma; \mathcal{H} \vdash_{\Sigma} M \Leftarrow \mathsf{Id}_{\mathbb{U}}(A;A') \\ &\Gamma, u: A; \mathcal{H}, A \equiv A': \mathbb{U} \vdash_{\Sigma} N \Leftarrow \mathsf{Id}_{\mathbb{U}}([u/x]B; [u/x]B') \end{split} \qquad \mathbf{Q} \in \{\Pi, \Sigma\}$$

 $(BinderEq) \qquad \qquad \Gamma; \mathcal{H} \vdash_{\Sigma} \mathsf{binderEq}(M; [y]N) \Leftarrow \mathsf{Id}_{\mathbb{U}}(Q_A[x]B; Q_{A'}[x]B')$ 

$$\begin{array}{c} \Gamma; \mathcal{H} \vdash_{\Sigma} f \Leftarrow \Pi_{A}[x]B \\ \Gamma; \mathcal{H} \vdash_{\Sigma} g \Leftarrow \Pi_{A}[x]B \\ \Gamma, u : A; \mathcal{H} \vdash_{\Sigma} M \Leftarrow \mathsf{Id}_{[u/x]B}(f @ u; g @ u) \\ \hline \Gamma; \mathcal{H} \vdash_{\Sigma} \mathsf{funext}([u]M) \Leftarrow \mathsf{Id}_{\Pi_{A}[x]B}(f; g) \end{array}$$
 (Funext)