# Intro to Interactive Theorem Proving

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http://github.com/joom/presentations

# Why?

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- Proofs are too long to check carefully!
- Seeing what we proved so far helps during the proof process.

#### **Examples**

Four color theorem (required checking 1,936 cases)
 Appel & Haken 1976 with computer assistance
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   Appel & Haken 1976 with computer assistance
   Werner & Gonthier 2005 with Coq
- Kepler conjecture (hexagonal close packing of spheres)
   Hales 1998 with computer assistance (C++ etc.)
   Hales 2014 with HOL Light and Isabelle (Flyspeck)

How?

• Curry-Howard isomorphism

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- Curry-Howard isomorphism
- Martin-Löf type theory

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$$\frac{\Gamma, x : A \vdash t : B}{\Gamma \vdash \lambda x.t : A \rightarrow B}$$

1

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$$\frac{\Gamma \vdash A \qquad \Gamma \vdash B}{\Gamma \vdash A \land B} \qquad \frac{\Gamma \vdash a : A \qquad \Gamma \vdash b : B}{\Gamma \vdash (a, b) : A \times B}$$

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## Martin-Löf type theory

- Extension of Curry-Howard to a full intuitionistic logic
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#### Note

MLTT is constructive, but you can add the law of excluded middle as an axiom and prove classical theorems.

#### Some Proof Assistants

- Isabelle (1986) (tactics!)
- Coq (1989) (tactics!)
- Agda (2007)
- Idris (2007) (actually mainly a general purpose language)