

Calvinballs



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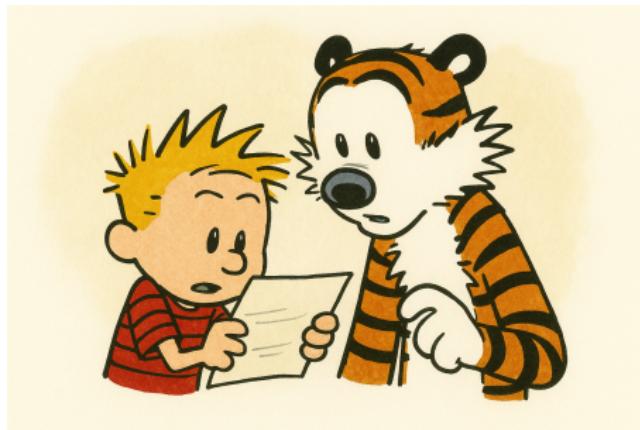


Do Calvin and Hobbes ever have a chance of doing well in this game?

Calvinballs

Do Calvin and Hobbes ever have a chance of doing well in this game?

Yes! ... sometimes ... but they have to read our paper.



TRACKING SOLUTIONS OF TIME-VARYING VARIATIONAL INEQUALITIES

Based on joint work with



Hédi Hadji



Cristóbal Guzmań

November 6, 2025

Overview

- Our objective: Time-varying Variational Inequalities.
- Special cases:
 - Dynamic Regret Guarantees for Online Learning;
 - Tracking Equilibria in Time-Varying Games;
- Tracking Guarantees for:
 - Tame Time-Varying Variational Inequalities;
 - Periodic Time-Varying Variational Inequalities.
- Negative Results.
- Open Questions.

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Variational Inequalities

Definition (Variational Inequality Problem)

Denote

- $\mathcal{C} \subseteq \mathbb{R}^d$, closed and convex set,
- $F : \mathcal{C} \rightarrow \mathcal{C}$, L -Lipschitz continuous operator.

The finite-dimensional Variational Inequality Problem VIP(F, \mathcal{C}) is defined as:

$$\text{find } Z^* \in \mathcal{C} \text{ such that } \langle F(Z^*), Z - Z^* \rangle \geq 0 \text{ for all } Z \in \mathcal{C}. \quad (\text{VIP})$$

Assumption: The VIP has a unique solution Z^* .

Special Cases:

- Computing minimizers of convex optimization problems;
- Equilibrium computation in concave games.

Time-Varying Variational Inequality Problems

Protocol: For each iteration $t \in [T]$:

- Learner chooses $Z_t \in \mathcal{C}$ using algorithm \mathcal{A} ;
- Nature chooses $F_t : \mathcal{C} \rightarrow \mathcal{C}$;
- Learner observes $F_t(Z_t)$.

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Definition (Tracking Error)

Tracking Error: Let $Z_t^* \in \mathcal{C}$ denote the solution for $\text{VIP}(F_t, \mathcal{C})$,

$$\tau_T(\mathcal{A}) = \sum_{t=1}^T \|Z_t - Z_t^*\|^2.$$

Lemma: For any algorithm \mathcal{A} , there exists a sequence (F_t) , such that $\tau_T(\mathcal{A}) \gtrsim T$.

Special Case I: Dynamic Regret

Remark (Relation Tracking Error and Dynamic Regret.)

Dynamic regret:

$$\text{dynReg}_T \left((Z_t^*)_{t \in [T]} \right) = \sum_{t=1}^T f_t(Z_t) - f_t(Z_t^*) .$$

Assume: For all $t \in [T]$, f_t 's are

- μ -strongly convex,
- differentiable,
- ∇f_t are L -Lipschitz, and
- $Z_t^* \in \text{int}(\mathcal{C})$.

Then

$$\frac{\mu}{2} \tau_T(\mathcal{A}) \leq \text{dynReg}_T \left((Z_t^*)_{t \in [T]} \right) \leq \frac{L}{2} \tau_T(\mathcal{A}) .$$

Tame Time-Varying VIP

Definition (Tame Time-Varying VIP)

We call a time-varying VIP α -*tame* if the quadratic path length is sublinear: For $\alpha \in [0, 1)$

$$P_T^* = \sum_{t=2}^T \|Z_{t-1}^* - Z_t^*\|^2 \lesssim T^\alpha.$$

Existing results:

Definition inspired by Duvocelle et al. [DMSV23].

Special Case II: Equilibrium Tracking for Time-varying Games

Two-Player Zero-Sum Time-Varying Games (Duvocelle et al. [DMSV23]):

Protocol: For each iteration $t \in [T]$:

- Agents chooses $x_t \in \mathcal{X}$ and $y_t \in \mathcal{Y}$ using algorithm \mathcal{A} ;
- Nature chooses game $\nu_t : \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}$;
- Agents observes first-order feedback based on $\nu_t(x_t, y_t)$.

Lemma

- $\nu : \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}$ the differentiable and convex-concave.

- $F([x, y]^\top) := [\nabla_x \nu(x, y), -\nabla_y \nu(x, y)]$

$$(x^*, y^*) \text{ is a Nash-equilibrium} \quad \Leftrightarrow \quad [x^*, y^*]^\top \text{ is a solution for VIP } (F, \mathcal{X} \times \mathcal{Y}) .$$

Tracking Guarantees for Tame VIP

Definition (Contractive Algorithms)

Let $C \in (0, 1)$. An algorithm \mathcal{A} is said to be C -contractive over a set of operators \mathcal{F} if for all $F \in \mathcal{F}$

$$\forall Z \in \mathcal{C} : \quad \|\mathcal{A}(Z, F) - Z^*\| \leq C \|Z - Z^*\|,$$

where Z^* is the solution for $\text{VIP}(F, \mathcal{C})$.

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Example:

- \mathcal{A} is projected gradient descent with step-size c ;
- $\mathcal{F} := \{\nabla f \mid f : \mathcal{C} \rightarrow \mathbb{R} \text{ and } f \text{ is } L\text{-smooth and } \mu\text{-strongly convex with } \frac{\mu}{L^2} \leq c\}$;
- \mathcal{A} is $(1 - (\mu/L)^2)$ contractive over \mathcal{F} .

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Theorem (Theorem 2.1 [HSG24])

Suppose \mathcal{A} is C -contractive over \mathcal{F} . Then, for any sequence of operators in \mathcal{F} with solutions $(Z_t^*)_{t \in [T]}$, the tracking error is bounded by

$$\tau_T(\mathcal{A}) \lesssim \frac{1}{(1-C)^2} P_T^* + 1 .$$

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- This bound is tight up to constants.
- **Corollary:** For α -tame problems, this gives $\tau_T(\mathcal{A}) \lesssim T^\alpha$.

Comparison to Existing Results

	Rate	Setting
Duvocelle et al. [DMSV23]	$T^{\frac{2+\alpha}{3}}$	Stochastic feedback Strongly monotone time-varying games.
Besbes et al. [BGZ15]	$T^{\frac{1+\alpha}{2}}$	Stochastic feedback Strongly convex time-varying functions.
Mokhtari et al. [MSJR16]	T^α	Deterministic feedback Time-varying strongly convex functions.
Our Result	T^α	Deterministic feedback Time-varying VIPs - any contractive algorithm.

Open Question

- Can our results be generalized for stochastic feedback?
- Can the work by Duvocelle et al. [DMSV23] or Besbes et al. [BGZ15] be generalized for VIP?

Periodic Time-Varying Variational Inequality Problems

Definition (Periodic Time-Varying Variational Inequality Problems)

A time-varying variational inequality problem $\text{VIP}((F_t), \mathcal{C})$ is periodic if there exists $k \in \mathbb{N}$, such that

$$F_{t+k} = F_t$$

for all $t \in \mathbb{N}$.

- **Example:** Seasonal shifting market.

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- **Example:** Seasonal shifting market.
- **Remark:** Periodic Time-Varying VIPs are not tame.

Assumption (Strong Monotonicity)

Operator F is strongly monotone if

$$\exists \mu > 0 : \forall Z, Z' \in \mathcal{C} : \quad \langle F(Z) - F(Z'), Z - Z' \rangle \geq \mu \|Z - Z'\|^2.$$

Cyclic Forward Method

Algorithm 1: $\mathcal{A}_{\text{cFW}}^{(k)}$, Cyclic Forward Method

Input: Period length k , step-size schedule (η_t) .

Initialization: Initial iterates $(Z_{1,1}, Z_{2,1}, \dots, Z_{k,1}) = (Z_1, Z_1, \dots, Z_1)$ with $Z_1 \in \mathcal{C}$.

1 **for** $t = 1, \dots, T$ **do**

2 Pick current index $n = (t \bmod k) + 1$;

3 Play $Z_t = Z_{n,t}$, receive $F_t(Z_t)$;

4 Update for all $i \in \{1, \dots, k\}$

$$Z_{i,t+1} = \begin{cases} \text{Proj}_{\mathcal{C}}(Z_t - \eta_t F_t(Z_t)) & \text{for } i = n \\ Z_{i,t} & \text{otherwise.} \end{cases}$$

Tracking Guarantees

Theorem (Corollary 3.1 [HSG24])

Let (F_t) be a k -periodic sequence. Assume

- F_t 's are μ -strongly monotone operators,
- there exists a constant G such that $\forall Z \in \mathcal{C} : \|F_t(Z)\| \leq G$.

Set $\eta_t = k/(\mu t)$. Then

$$\tau_T(\mathcal{A}_{\text{cFW}}^{(k)}) \lesssim k \left(\frac{G}{\mu} \right)^2 \left(\log \left(\frac{T}{k} \right) + 1 \right).$$

Problem: We do not necessarily know k .

Meta-Algorithm: Intuition

Suppose we know an upper bound $K \geq k$.

$$\mathcal{A}_{\text{cFW}}^{(1)}, \quad \mathcal{A}_{\text{cFW}}^{(2)}, \quad \dots \quad \mathcal{A}_{\text{cFW}}^{(k-1)}, \quad \mathcal{A}_{\text{cFW}}^{(k)}, \quad \mathcal{A}_{\text{cFW}}^{(k+1)}, \quad \dots \quad \mathcal{A}_{\text{cFW}}^{(K-1)}, \quad \mathcal{A}_{\text{cFW}}^{(K)}$$

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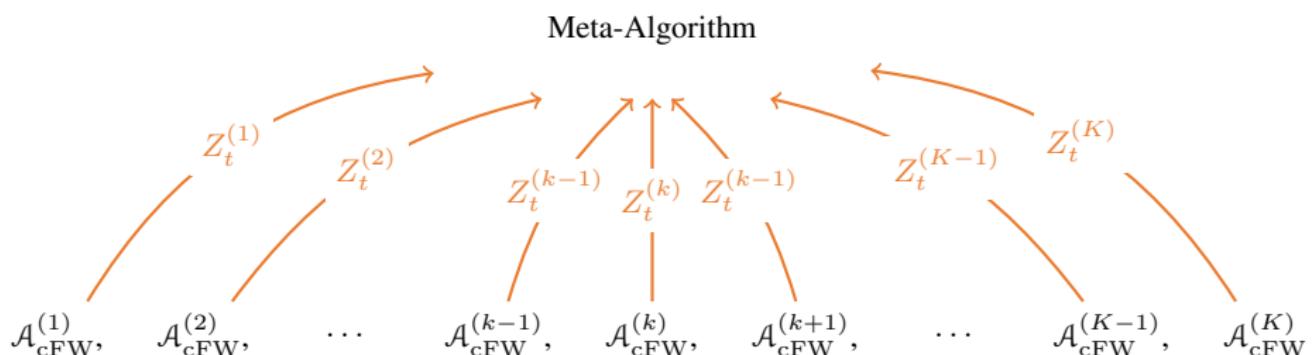
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Nature

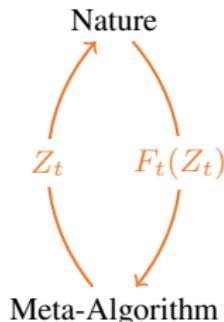
Meta-Algorithm

Compute Z_t , based on $Z_t^{(1)}, \dots, Z_t^{(K)}$

$\mathcal{A}_{\text{cFW}}^{(1)}, \quad \mathcal{A}_{\text{cFW}}^{(2)}, \quad \dots \quad \mathcal{A}_{\text{cFW}}^{(k-1)}, \quad \mathcal{A}_{\text{cFW}}^{(k)}, \quad \mathcal{A}_{\text{cFW}}^{(k+1)}, \quad \dots \quad \mathcal{A}_{\text{cFW}}^{(K-1)}, \quad \mathcal{A}_{\text{cFW}}^{(K)}$

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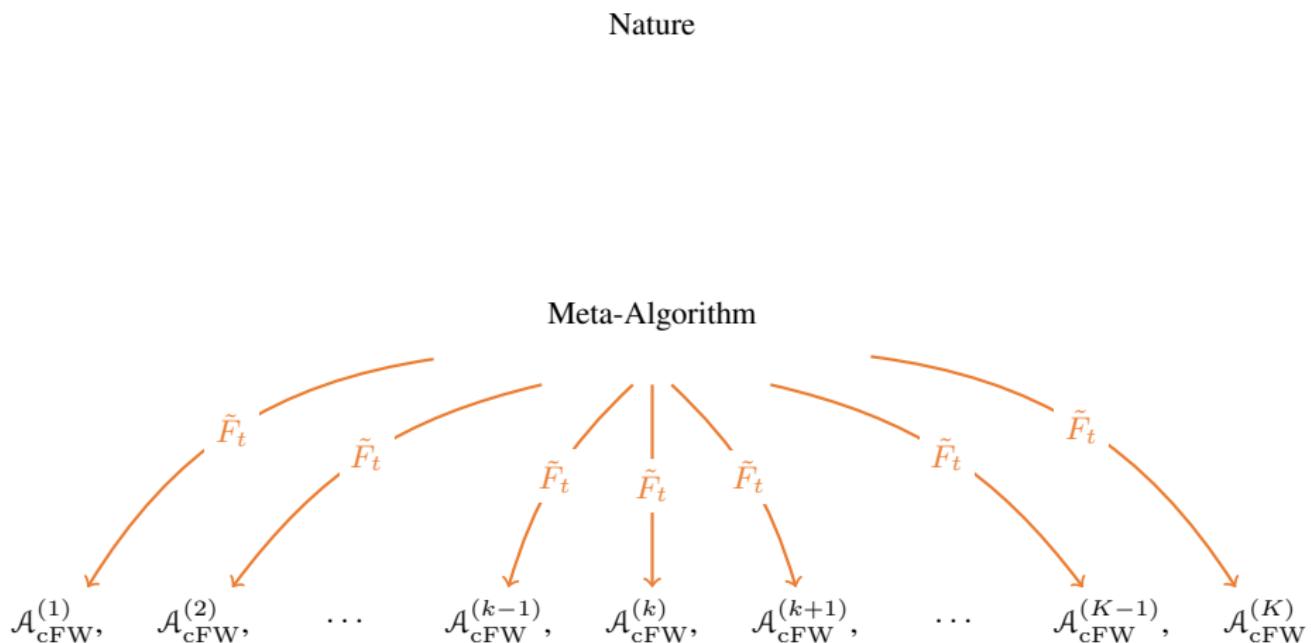
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Meta-Algorithm: Intuition

Suppose we know an upper bound $K \geq k$.



Meta-Algorithm

Algorithm 2: Meta-Algorithm $\mathcal{A}_{\text{meta}}$

Data: Maximum period length $K \in \mathbb{N}$,
prior distribution $p_1 \in \Delta_K$ over expert algorithms.

- 1 **for** $t = 1, \dots, T$ **do**
 - 2 Receive $(Z_t^{(1)}, \dots, Z_t^{(K)})$ from expert algorithms $\mathcal{A}_{\text{cFW}}^{(1)}, \dots, \mathcal{A}_{\text{cFW}}^{(K)}$;
 - 3 Play $Z_t = p_{t,1}Z_t^{(1)} + \dots + p_{t,K}Z_t^{(K)}$;
 - 4 Receive $F_t(Z_t)$;
 - 5 Update p_t using exponential weights;
 - 6 Send surrogate operator \tilde{F}_t to experts $\mathcal{A}_{\text{cFW}}^{(1)}, \dots, \mathcal{A}_{\text{cFW}}^{(K)}$;
-

Remark: The algorithm only requires one evaluation of F_t .

Tracking Guarantees

Theorem (Corollary 3.2 [HSG24])

Let (F_t) be a k -periodic sequence. Assume

- F_t 's are μ -strongly monotone operators,
- diameter(\mathcal{C}) is bounded, and
- $\forall Z \in \mathcal{C} : \|F_t(Z)\| \leq G$.

If upper bound $K \geq k$ is known, then

$$\tau_T(\mathcal{A}_{\text{meta}}) \lesssim \kappa^2 D^2 \left(k \left(\log \left(\frac{T}{k} \right) + 1 \right) + \log K \right),$$

where $\kappa = L/\mu$.

Remark: Tracking bound does not depend on path length P_T^* .

Open Research Questions

- Can we have similar results for periodic time-varying VIPs with stochastic feedback?
- Can we obtain tracking guarantees for periodic time-varying VIPs with a truly distributed meta-algorithm?
- Can we show tracking guarantees for 'almost' periodic problems?

Let's Ignore the Theory

Can we just ignore time variance?



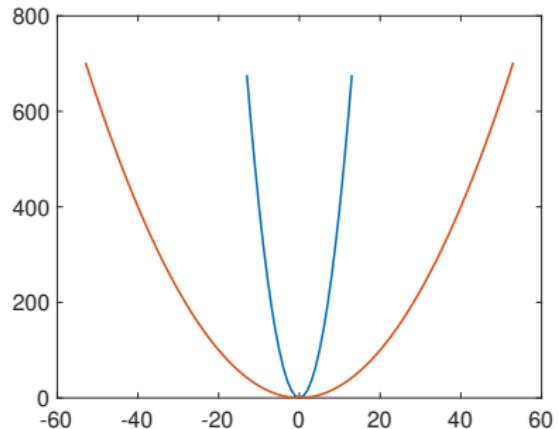
Negative Results: Setting

$$\forall t \in 2\mathbb{N} : f_t(x) = \log(1 + e^{4x^2/2}),$$

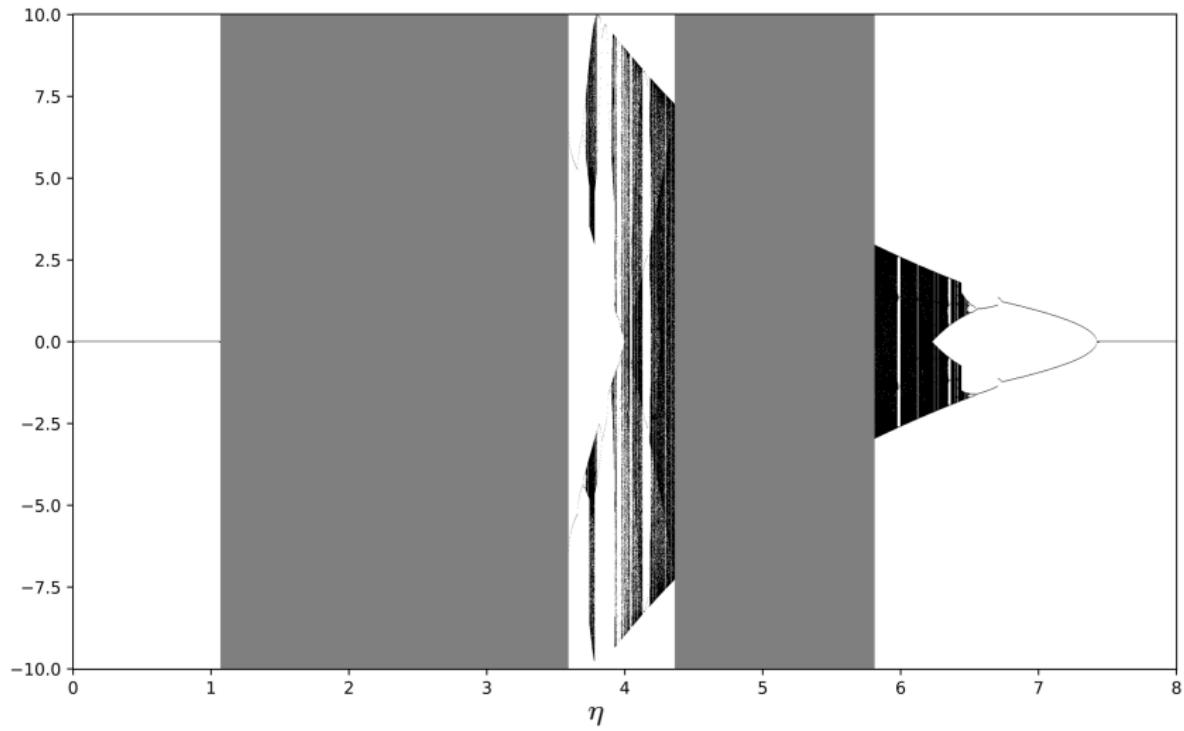
$$\forall t \notin 2\mathbb{N} : f_t(x) = \log(1 + e^{\frac{1}{4}x^2/2}).$$

We use gradient descent:

$$x \mapsto x - \eta f'_t(x).$$



Negative Results



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Appendix: Proof of Theorem 2.1 [HSG24]

Use $\|u + v\|^2 \leq (1 + \alpha)\|u\|^2 + (1 + \alpha^{-1})\|v\|^2$ for any $u, v \in \mathbb{R}^d$ and $\alpha > 0$:

$$\underbrace{\|Z_{t+1} - Z_{t+1}^*\|^2}_{=\|u+v\|^2} \leq \underbrace{\left(1 + \left(\frac{1}{C} - 1\right)\right)}_{=\alpha} \underbrace{\|Z_{t+1} - Z_t^*\|^2}_{=\|u\|^2} + \underbrace{\left(1 + \frac{1}{1/C - 1}\right)}_{=\|v\|^2} \underbrace{\|Z_{t+1}^* - Z_t^*\|^2}_{=\|v\|^2}$$

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Summing over $T - 1$ steps:

$$\sum_{t=1}^{T-1} \|Z_{t+1} - Z_{t+1}^*\|^2 \leq C \sum_{t=1}^{T-1} \|Z_t - Z_t^*\|^2 + \frac{1}{1-C} \sum_{t=1}^{T-1} \|Z_{t+1}^* - Z_t^*\|^2.$$

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Reorganizing and adding $\|Z_1 - Z_1^*\|^2$:

$$\sum_{t=1}^T \|Z_t - Z_t^*\|^2 \leq \frac{1}{(1-C)^2} \sum_{t=1}^{T-1} \|Z_{t+1}^* - Z_t^*\|^2 + \frac{1}{1-C} \|Z_1 - Z_1^*\|^2.$$

Appendix: Details on Surrogate Loss

Aggregator Method:

- Use exponential weights updates:

$$p_{t,i} \propto \exp\left(-\lambda \sum_{s=1}^{t-1} \ell_{s,i}\right).$$

- Obtain loss

$$\ell_{s,i} = \langle F_s(Z_s), Z_s^{(i)} \rangle + \frac{\mu}{2} \left\| Z_s^{(i)} - Z_s \right\|^2.$$

Surrogate Operator \tilde{F}_t :

$$\tilde{F}_t : z \mapsto F_t(Z_t) + \mu(z - Z_t).$$

Appendix: Distributed Setting

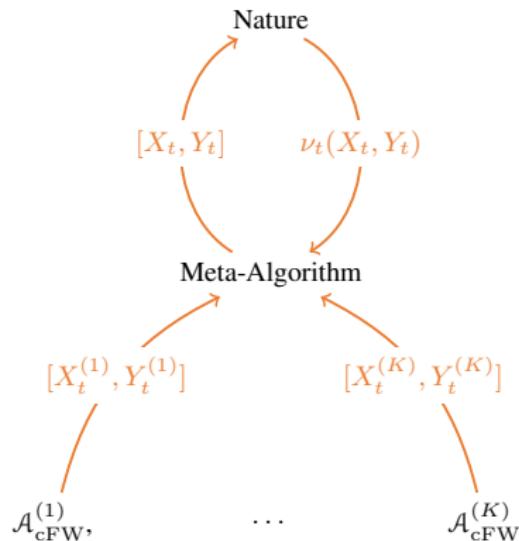
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Can we obtain tracking guarantees for periodic VIPs with a distributed meta-algorithm?

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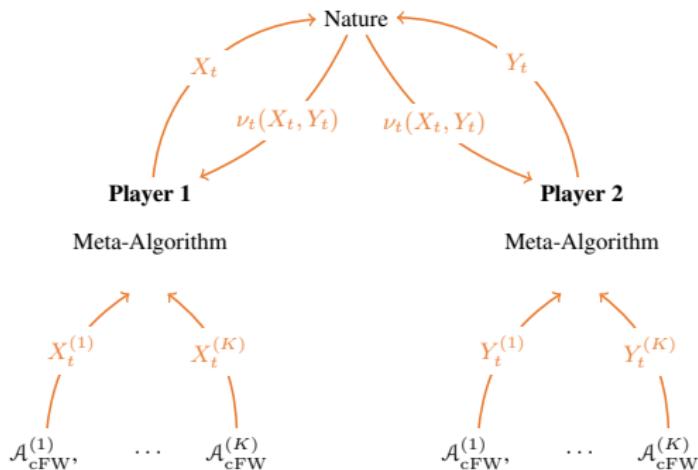
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