

Lists Stacks & Queues

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Comparisons

- Is a linear data structure.
- A List is a collection where the elements are ordered and therefore each element has an index which is the position in the list. Allows duplicates.
- A Set is an unordered collection in which no two elements are identical.
- A Bag is an unordered collection in which can have duplicates.

	Linked List	Array based
Access	$\Theta(n)$	$\Theta(1)$
Insertion	$O(n)$	$O(1)$
Deletion	$O(n)$	$O(n)$

Amortized Analysis

- In **Amortized analysis**, the **time taken to execute a sequence** is **averaged over all the operations executed**.
- **Even though one of the operations in the sequence might take a long time**, the **average time taken over all operations is small**.
- This is **not the same as the average case**.
- This **guarantees the average performance of each operation in the worst-case**.

Stacks

- It's a **list structure** where **all operations occur on one end of the list**, known as the **top of the stack**.
- To **add** an element **is called a push** operation.
- To **remove** an element **is called a pop** operation.
- To **get** the element **at the top of the stack** is called a **top** operation.

Array implementation

- Requires a means of handling array overflow, i.e. double size of array when full.
- `push()` has an amortized complexity of $O(1)$ in the worst case.
- `top()` does not alter the stack at all and simply gives the top element, this is $O(1)$ in the worst case.
- `pop()` only alters the last element, nothing is shifted and therefore has complexity $O(1)$ in the worst case.

Linked-list implementation

- `push()` has a complexity of $O(1)$ in the worst case, this is **NOT** **amortized** due to the lack of array overflow requirement.
- `top()` has $O(1)$ complexity in the worst case.
- `pop()` has $O(1)$ complexity in the worst case.

Parenthesis checking

- Stacks can be used to determine if parenthesis match correctly or not. e.g. $[a(b + c)da/c + e]/b$
- Use a stack to push the left side of the parenthesis and when the right side has been found, pop the parenthesis.
- When an item is popped, it is compared to the found parenthesis and if they are of the same type, then it's matching.

- A **FIFO (First In First Out)** queue is where the item at the front of the queue is used first.
- A new arrival joins the end.
- All insertions are made at one end of the list, known as the **rear of the queue**.
- The **insertion** method is known as an **enqueue** operation.
- All **removals** happen at the other end of the queue known as the **the front of the queue**.
- The **removal** method is known as a **dequeue** operation.

Circular Array implementation

- Representing a queue in a traditional array such that the queue elements march through the array in one direction is not very convenient.
- This keeps track of the front rear points of the queue pointing to their respective positions.
- enqueue adds the element to $\text{rearPos} + 1 \% \text{arrSize}$.
- If the queue length is equal to arrSize, then the array is doubled in size, when copying elements over, they are copied in queue order and the front and rear keep pointing to their respective elements.
- dequeue simply removes the element pointed at by the frontPos and increments the front pointer to $\text{frontPos} = (\text{frontPos} + 1) \% \text{arrSize}$.

Linked List implementation

- Keep track of front and rear nodes.
- **enqueue** Will create a new node and the current rear node points to the new node. The rear is then changed to point at the new node.
- **dequeue** simply removes the node pointed at by front and front points at the next node.

The End