

Lists Stacks & Queues

Jonathan Windle

University of East Anglia

J.Windle@uea.ac.uk

May 21, 2017

Overview I

1 Lists

- Comparisons
- Amortized Analysis

2 Stacks

- Intro
- Array implementation
- Linked-list implementation
- Parenthesis checking

3 Queues

- Intro
- Circular Array implementation

Comparisons

- Is a **linear data structure**.
- A List is a collection where the elements are **ordered** and therefore each element has an index which is the position in the list. Allows duplicates.
- A Set is an **unordered** collection in which no two elements are identical.
- A Bag is an **unordered** collection in which can have duplicates.

	Linked List	Array based
Access	$\Theta(n)$	$\Theta(1)$
Insertion	$O(n)$	$O(1)$
Deletion	$O(n)$	$O(n)$

Amortized Analysis

- In **Amortized analysis**, the time taken to execute a sequence is averaged over all the operations executed.
- Even though one of the operations in the sequence might take a long time, the average time taken over all operations is small.
- This is not the same as the average case.
- This guarantees the average performance of each operation in the worst-case.

- It's a list structure where all operations occur on one end of the list, known as the **top of the stack**.
- To add an element is called a **push** operation.
- To remove an element is called a **pop** operation.
- To get the element at the **top of the stack** is called a **top** operation.

Array implementation

- Requires a means of handling **array overflow**, i.e. double size of array when full.
- `push()` has an **amortized** complexity of $O(1)$ in the worst case.
- `top()` does not alter the stack at all and simply gives the top element, this is $O(1)$ in the worst case.
- `pop()` only alters the last element, nothing is shifted and therefore has complexity $O(1)$ in the worst case.

Linked-list implementation

- `push()` has a complexity of $O(1)$ in the worst case, this is **NOT** **amortized** due to the lack of array overflow requirement.
- `top()` has $O(1)$ complexity in the worst case.
- `pop()` has $O(1)$ complexity in the worst case.

Parenthesis checking

- Stacks can be used to determine if parenthesis match correctly or not. e.g. $[a(b + c)da/c + e]/b$
- Use a stack to push the left side of the parenthesis and when the right side has been found, pop the parenthesis.
- When an item is popped, it is compared to the found parenthesis and if they are of the same type, then it's matching.

- A **FIFO (First In First Out)** queue is where the item at the front of the queue is used first.
- A new arrival joins the end.
- All insertions are made at one end of the list, known as the **rear of the queue**.
- The insertion method is known as an **enqueue** operation.
- All removals happen at the other end of the queue known as the **the front of the queue**.
- The removal method is known as a **dequeue** operation.

Circular Array implementation

- Representing a queue in a traditional array such that the queue elements march through the array in one direction is not very convenient.
- This keeps track of the front rear points of the queue pointing to their respective positions.
- **enqueue** adds the element to $\text{rearPos} + 1 \% \text{arrSize}$.
- If the queue length is equal to arrSize , then the array is doubled in size, when copying elements over, they are copied in queue order and the front and rear keep pointing to their respective elements.
- **dequeue** simply removes the element pointed at by the frontPos and increments the front pointer to $\text{frontPos} = (\text{frontPos} + 1) \% \text{arrSize}$.

Linked List implementation

- Keep track of front and rear nodes.
- **enqueue** Will create a new node and the current rear node points to the new node. The rear is then changed to point at the new node.
- **dequeue** simply removes the node pointed at by front and front points at the next node.

The End