COMP/LING 445/645 Problem Set 5

There are several types of questions below. For programming questions, please put your answers into a file called lastname-firstname.clj. Be careful to follow the instructions exactly and be sure that all of your function definitions use the exact names, number of inputs, input types, number of outputs, and output types as requested in each question. Please make sure all other variables are also named as indicated.

To do the computational problems, we recommend that you install Clojure on your local machine and write and debug the answers to each problem in a local copy of lastname-firstname.clj. You can find information about installing and using Clojure here https://clojure.org/.

For the code portion of the assignment, it is crucial to submit a standalone file that runs. Before you submit lastname-firstname.clj to us via email, make sure that your code executes correctly without any errors when run at the command line by typing clojure lastname-firtname.clj at a terminal prompt. We cannot grade any code that does not run correctly as a standalone file, and if the preceding command like produces an error, the code portion of the assignment will receive a 0.

For questions involving answers in English or mathematics or a combination of the two, put your answers to the question in an **Answer** section like in the example below. You can find more information about LaTeX here https://www.latex-project.org/.

Once you have answered the question, please compile your copy of this LATEX document into a PDF and submit (i) the compiled PDF (ii) the raw LATEX file and (iii) your lastname-firstname.clj file via email to both timothy.odonnell@mcgill.ca and savanna.willerton@mail.mcgill.ca. The problem is set is due before 16:05 on Wednesday, December 18, 2019.

Problem 0: This is an example question using some fake math like this $L = \sum_{0}^{\infty} \mathcal{G} \delta_{x}$.

Answer 0: Put your answer right under the question like this $L = \sum_{0}^{\infty} \mathcal{G} \delta_{x}$.

Background: Throughout this problem set we will be working with the Hidden Markov Model. In an HMM, there is a sequence of hidden states $c^{(1)},...,c^{(n)}$, and a sequence of observed words, $w^{(1)},...,w^{(n)}$. The set of all possible hidden states is S, and the set of all possible words is O. Thus $c^{(i)} \in S$ and $w^{(i)} \in O$ for all i. We defined the distribution on hidden states and words as follows in class:

$$\Pr(W^{(1)}, \dots, W^{(k)}, C^{(1)}, \dots, C^{(k)}) = \prod_{i=1}^{k} \Pr(W^{(i)} \mid C^{(i)} = c^{(i)}, \vec{\theta}_{O, c^{(i)}}) \Pr(C^{(i)} \mid C^{(i-1)} = c^{(i-1)}, \vec{\theta}_{T, c^{(i-1)}})$$

$$\tag{1}$$

Note that we are assuming (for the time being) that the parameters $\vec{\theta}_{T,c^{(i-1)}}$ and $\vec{\theta}_{O,c^{(i)}}$ are fixed, i.e. there is no uncertainty about their value.

We will assume throughout that there are three hidden states (categories) and three words. These are defined by:

```
(def hidden-states '(Start N V))
(def vocabulary '(Call me Ishmael))
```

Note that, to make the exercises less painful, we are not including a stop state. There are three state transition distributions that we need to define, one for each of the hidden states: $\Pr(C^{(i)} \mid C^{(i-1)} = Start)$, $\Pr(C^{(i)} \mid C^{(i-1)} = N)$, and $\Pr(C^{(i)} \mid C^{(i-1)} = V)$. As stated in the definition of the HMM probabilities above, these distributions are specified by the parameters $\vec{\theta}_{T,Start}$, $\vec{\theta}_{T,N}$, and $\vec{\theta}_{T,V}$:

```
(def theta-transition-Start '(0 0.9 0.1))
(def theta-transition-N '(0 0.3 0.7))
(def theta-transition-V '(0 0.8 0.2))
(def theta-transition-dists-1
  (list theta-transition-Start theta-transition-N theta-transition-V))
```

This means that the hidden state N has probability 0.3 of transitioning to hidden state N, and probability 0.7 of transitioning to V. Note that none of the states ever transition to the **start** hidden state. This means that the **start** state can only occur as the first hidden state in a sequence, and never afterwards.

We also need to specify the observation distribution for the hidden states N and V (the hidden state start cannot produce observations, so it is ignored here). They are defined by:

This means, for example, that the hidden state N has probability 0.4 of emitting the word Ishmael. The following helper functions will be useful for writing your code.

Problem 1: Suppose that at time t the hidden state is $c^{(t)}$. Then we can compute the probability that the next hidden state is $c^{(t+1)}$ and the next observed word is $w^{(t+1)}$. This probability is equal to:

$$\begin{split} &\Pr(W^{(t+1)} = w^{(t+1)}, C^{(t+1)} = c^{(t+1)} \mid C^{(t)} = c^{(t)}) \\ &= \Pr(W^{(t+1)} = w^{(t+1)} \mid C^{(t+1)} = c^{(t+1)}) \Pr(C^{(t+1)} = c^{(t+1)} \mid C^{(t)} = c^{(t)}) \end{split}$$

Write a procedure score-next-state-word which computes the log of this probability. It should take five arguments: the current hidden state, the next hidden state, and next observed word, the hidden state transition distributions, and the observation distributions. A call to (score-next-state-word current-hidden next-hidden next-observed theta-transition-dists theta-observation-dists) should return the log probability of the next hidden state and next observed word, given the current hidden state. Hint: use the functions dist-lookup and logscore-categorical defined above.

Answer 1: Please put the answer in lastname-firstname.clj.

Problem 2: Suppose that at time t the hidden state is $c^{(t)}$. Then we can compute the *marginal* probability of the word $w^{(t+1)}$ at time t+1 given the hidden state at time t. This probability is equal to:

$$\begin{aligned} & \Pr(W^{(t+1)} = w^{(t+1)} \mid C^{(t)} = c^{(t)}) \\ &= \sum_{c^{(t+1)}} \Pr(W^{(t+1)} = w^{(t+1)}, C^{(t+1)} = c^{(t+1)} \mid C^{(t)} = c^{(t)}) \\ &= \sum_{c^{(t+1)}} \Pr(W^{(t+1)} = w^{(t+1)} \mid C^{(t+1)} = c^{(t+1)}) \Pr(C^{(t+1)} = c^{(t+1)} \mid C^{(t)} = c^{(t)}) \end{aligned}$$

That is, we sum over all of the possible hidden states that could appear at time t+1.

Using the procedure score-next-state-word that you wrote in Problem 1, write a function compute-next-observation-marginal, which takes four arguments: current-state, next-observation, theta-transition-dists, and theta-observation-dists. When (compute-next-observation-marginal current-state next-observation theta-transition-dists theta-observation-dists) is called, it should return the marginal probability of the next observed word given the current hidden state.

Hint: use logsumexp.

Answer 2: Please put the answer in lastname-firstname.clj.

Problem 3: Suppose that at time t the hidden state is $c^{(t)}$. Then we can compute the probability of the next k hidden states and next k observed words given the hidden state at time t. This probability is equal to:

$$\Pr(W^{(t+1)}, \dots, W^{(t+k)}, C^{(t+1)}, \dots, C^{(t+k)} \mid C^{(t)} = c^{(t)})$$

$$= \prod_{i=t}^{t+k-1} \Pr(W^{(i+1)} = w^{(i+1)} \mid C^{(i+1)} = c^{(i+1)}) \Pr(C^{(i+1)} = c^{(i+1)} \mid C^{(i)} = c^{(i)})$$

$$= \Pr(C^{(t+1)} = c^{(t+1)} \mid C^{(t)} = c^{(t)}) \Pr(W^{(t+1)} = w^{(t+1)} \mid C^{(t+1)} = c^{(t+1)})$$

$$\cdot \Pr(W^{(t+2)}, \dots, W^{(t+k)}, C^{(t+2)}, \dots, C^{(t+k)} \mid C^{(t+1)} = c^{(t+1)})$$

Note that the equation shows that there is a recursive formula for this probability.

Write a procedure score-next-states-words which computes this probability. The function should take five arguments: the current hidden state, a list of k hidden states, a list of k observed words, the transition distributions, and the observation distributions. When (score-next-states-words current-hidden next-hidden-states next-words theta-transition-dists theta-observation-dists) is called, it should return the log probability that the k hidden states and k observed words will appear immediately following the current hidden state.

You should use recursion to write the function.

Answer 3: Please put the answer in lastname-firstname.clj.

Problem 4: Suppose the at time t the hidden state is $c^{(t)}$. Then we can compute the marginal probability of the next k observed words given this hidden state, i.e. summing out all of the possible settings of the hidden states. This probability is equal to:

$$\begin{aligned} & \Pr(W^{(t+1)}, \dots, W^{(t+k)} \mid C^{(t)} = c^{(t)}) \\ &= \sum_{c^{(t+1)}, \dots c^{(t+k)}} \Pr(W^{(t+1)}, \dots, W^{(t+k)}, C^{(t+1)}, \dots, C^{(t+k)} \mid C^{(t)} = c^{(t)}) \\ &= \sum_{c^{(t+1)}, \dots c^{(t+k)}} \prod_{i=t}^{t+k} \Pr(W^{(i+1)} = w^{(i+1)} \mid C^{(i+1)} = c^{(i+1)}) \Pr(C^{(i+1)} = c^{(i+1)} \mid C^{(i)} = c^{(i)}) \\ &= \sum_{c^{(t+1)}} \Pr(C^{(t+1)} = c^{(t+1)} \mid C^{(t)} = c^{(t)}) \Pr(W^{(t+1)} = w^{(t+1)} \mid C^{(t+1)} = c^{(t+1)}) \\ &\cdot \Pr(W^{(t+2)}, \dots, W^{(t+k)}) \mid C^{(t+1)} = c^{(t+1)}) \end{aligned}$$

The final equation is very important. It shows that the marginal probability of the words following hidden state $c^{(t)}$ can be computed in a recursive manner. Write a function compute-next-words-marginal, which computes the log marginal probability of a sequence of observed words given the current hidden state. The function should take four arguments: the current hidden state, a list of words, the transition distributions, and the observation distributions. When (compute-next-words-marginal current-hidden next-words theta-transition-dists theta-observation-dists) is called, it should return the log of the marginal probability of the list of words given the hidden state, as defined in the equation above.

You should use recursion to write your function.

Answer 4: Please put the answer in lastname-firstname.clj.

Problem 5: Find a sequence of two words $w^{(1)}, w^{(2)}$ such that the marginal probability of this sequence is different than the marginal probability of its reverse. That is, find $w^{(1)}, w^{(2)}$ such that:

$$\Pr(W^{(1)} = w^{(1)}, W^{(2)} = w^{(2)} \mid C^{(0)} = Start) \neq \Pr(W^{(1)} = w^{(2)}, W^{(2)} = w^{(1)} \mid C^{(0)} = Start)$$

Use your solution to Problem 4 to confirm that your sequences have different probabilities, and explain why this is true.

Answer 5: Calling compute-next-words-marginal on the following sequences of words yields the respective values:

 $(\mathrm{def\ seq 1\ '(Call\ me)})\ 2.8479969065549495\ (\mathrm{def\ seq 2\ '(me\ Call)})\ 2.9068905956085187$

The values differ due to the different theta values of the observed words in the given observation distributions. The word 'Call' will be multiplied by the observed values 0.1 and 0.8, while 'me' will be multiplied by the observed values 0.5 and 0.1 for the hidden states N and V respectively. Though the theta values of the transition states are marginalized over (and should return the same values when using dist-lookup) the theta values of the observations will differ.

Problem 6: In the next several problems, we will be using your solution to Problem 4 to perform Bayesian inference in the HMM. Our goal will be to compute the posterior distribution over hidden states given a sequence of words. Recall that Bayes' Rule tells us how to calculate this posterior distribution:

$$\frac{\Pr(C^{(1)} = c^{(1)}, \cdots, C^{(t)} = c^{(t)} \mid W^{(1)} = w^{(1)}, \dots, W^{(t)} = w^{(t)}) = }{\Pr(C^{(1)} = c^{(1)}, \cdots, C^{(t)} = c^{(t)}) \Pr(W^{(1)} = w^{(1)}, \dots, W^{(t)} = w^{(t)} \mid C^{(1)} = c^{(1)}, \cdots, C^{(t)} = c^{(t)})}{\Pr(W^{(1)} = w^{(1)}, \dots, W^{(t)} = w^{(t)})}$$

Write a procedure compute-hidden-prior, which takes two arguments: a list of k hidden states and the hidden state transition distributions. When (compute-hidden-prior hidden-states theta-transition-dists) is called, it should return the log prior probability of the hidden state sequence, i.e. the log value of

$$\Pr(C^{(1)} = c^{(1)}, \dots, C^{(k)} = c^{(k)})$$

.

Answer 6: Please put the answer in lastname-firstname.clj.

Problem 7: In this problem, we will continue with the Bayesian calculation from Problem 6. Write a procedure compute-likelihood-of-words, which takes three arguments: a list of k hidden states, a list of k words, and the observation distributions. When (compute-likelihood-of-words hidden-states words theta-observation-dists) is called, it should return the log probability of the k words being generated from the hidden states, i.e. the log value of $\Pr(W^{(1)} = w^{(1)}, \dots, W^{(k)} = w^{(k)} \mid C^{(1)} = c^{(1)}, \dots, C^{(k)} = c^{(k)})$.

Answer 7: Please put the answer in lastname-firstname.clj.

Problem 8: We are finally ready to solve our Bayesian inference problem. Write a procedure compute-hidden-posterior, which takes four arguments: a list of k hidden states, a list of k words, the hidden state transition distributions, and the observation distributions. When (compute-hidden-posterior hidden-states words theta-transition-dists theta-observation-dists) is called, it should return the log posterior probability of the k hidden states given the k words, i.e. it should return the log value of $\Pr(C^{(1)} = c^{(1)}, \dots, C^{(k)} = c^{(k)} \mid W^{(1)} = w^{(1)}, \dots, W^{(k)} = w^{(k)})$.

Hint: Use your solutions to Problems 4, 6, and 7.

Answer 8: Please put the answer in lastname-firstname.clj.

Problem 9: Suppose that we use the function compute-next-words-marginal that you wrote in Problem 4, in order to compute the marginal probability of a sequence of k observed words. How many times is compute-next-words-marginal called during the execution of the program? What does this mean about the run-time of the program?

Hint: your formula should involve the number of hidden states and the length k of the sequence of words.

Answer 9: Calling compute-next-words-marginal will trigger one call for each possible next hidden state at each word transition recursively. Every word transition will have n^m calls, where m is the m-th word in the sequence and n is the number of possible next hidden states. So, for n possible hidden states and a sequence of k words, the recursive function will be called $n^0 + n^1 + n^2 + ... + n^k$ times. So, the max run time will be bound by $O(n^k)$.

Problem 10: In class, we introduced *memoization* as a method for sharing computation and/or reducing the run-time of programs. Memoization refers to a technique where we intercept calls to a function with some arguments, check if the function has been called with those arguments before and, if not, call the function on those arguments, get the result, store it in a table and return it. On subsequent calls to that function with those arguments, we simply look up the result in the table and avoid recomputing the value. In Clojure, functions can be memoized by a call to **memoize**.

Using memoize, define a memoized version of your function compute-next-words-marginal. Name this function compute-next-words-marginal-mem.

Answer 10: Please put the answer in lastname-firstname.clj.

Problem 11: Suppose that we use the memoized function compute-next-words-marginal-mem that you wrote in the previous problem, in order to compute the marginal probability of a sequence of k observed words. How many times is compute-next-words-marginal-mem called during the execution of the program? How does this differ from the non-memoized version of the procedure?

Answer 11: The memoized version of compute-next-words-marginal will save the value for each hidden state at each word transition. So, the function will memoize a version of the function called on subsequent subsets of the (length k) sequence of words for each possible hidden state n. This means that each branch of the recursion tree that repeats is cut off. So each function for each word transition after the first call will only be called twice, once for each hidden state and the total number of calls to compute-next-words-marginal-mem is [(k-1)*n]+1. So, the run time will be bound by O(nk).