

# Security on the Line: Modern Curve-based Cryptography

Joost Renes

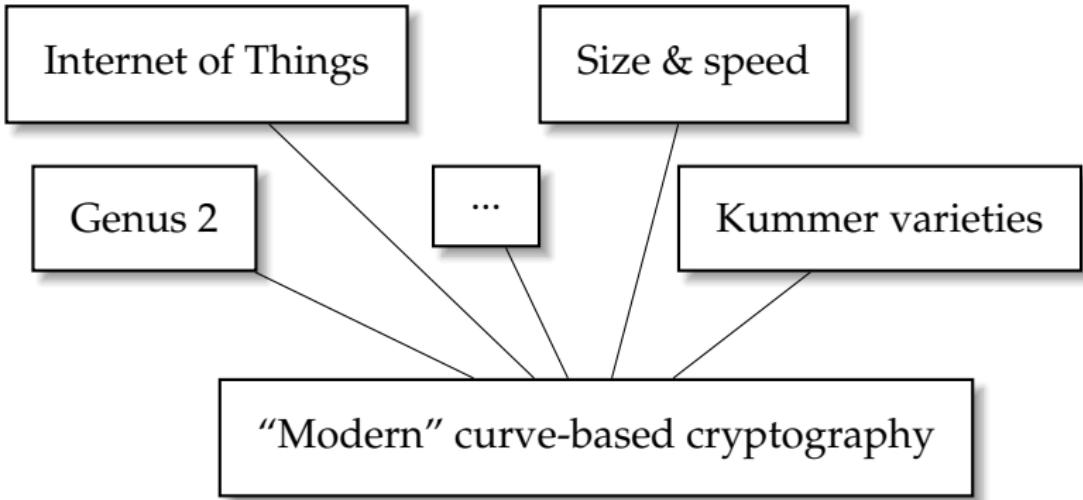
SCA Workshop

18 June 2019

# Modern curve-based cryptography

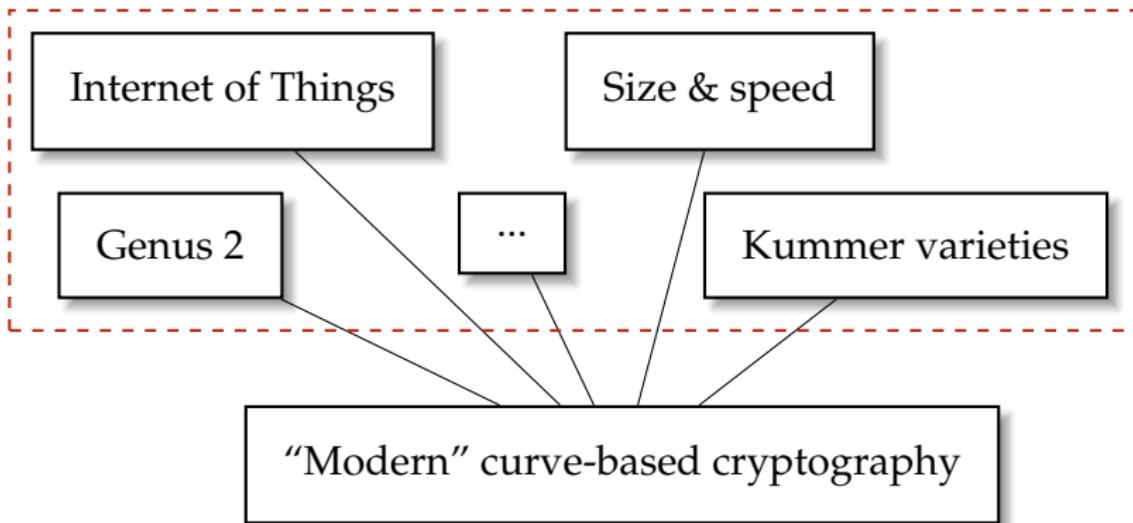
“Modern” curve-based cryptography

# Modern curve-based cryptography



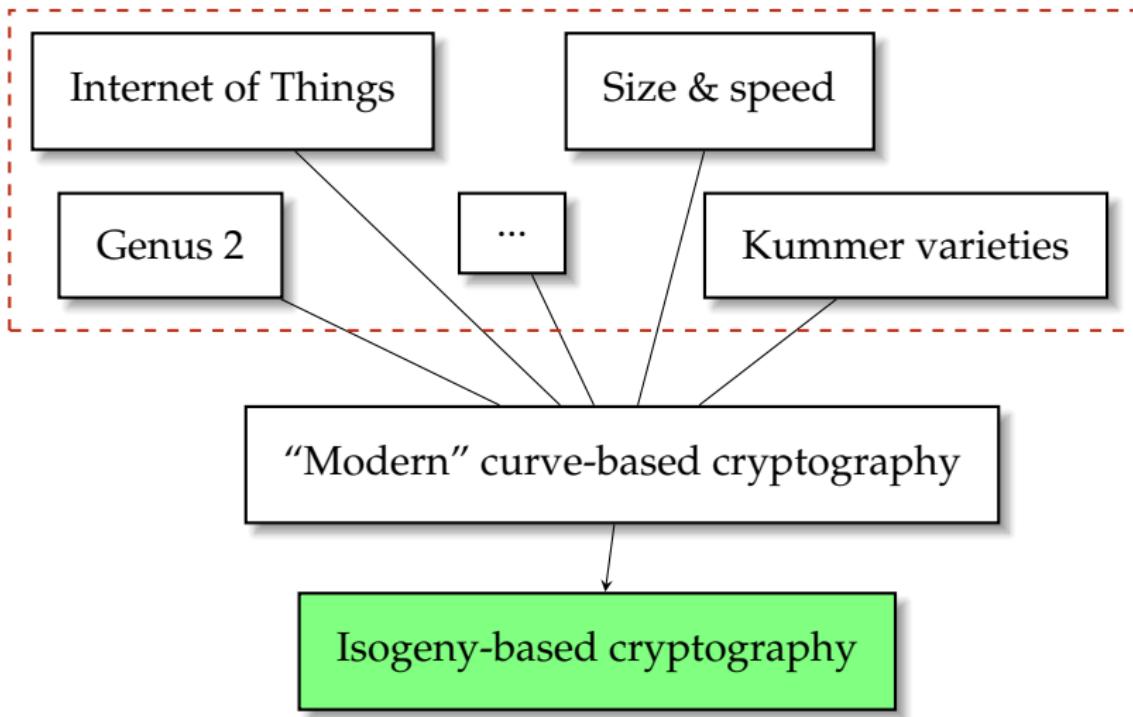
# Modern curve-based cryptography

*Classical setting (Ben Smith's talk)*



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*Classical setting (Ben Smith's talk)*



# Elliptic curves in cryptography

Discrete-log-based elliptic-curve  
cryptography [Mil86; Kob87]



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Ordinary isogeny-based group  
actions [Cou06; RS06; DKS18]



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Supersingular isogeny-based  
cryptography /  $\mathbb{F}_{p^2}$  [CLG09; JF11]



# Elliptic curves in cryptography

Discrete-log-based elliptic-curve cryptography [Mil86; Kob87]

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Ordinary isogeny-based group actions [Cou06; RS06; DKS18]

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Supersingular isogeny-based cryptography /  $\mathbb{F}_{p^2}$  [CLG09; JF11]

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Supersingular isogeny-based group actions /  $\mathbb{F}_p$  [Cas+18]



# Elliptic curves & isogenies (1)

**Fixed:** prime  $p$


$$E_{a,b} : y^2 = x^3 + ax + b$$

# Elliptic curves & isogenies (1)

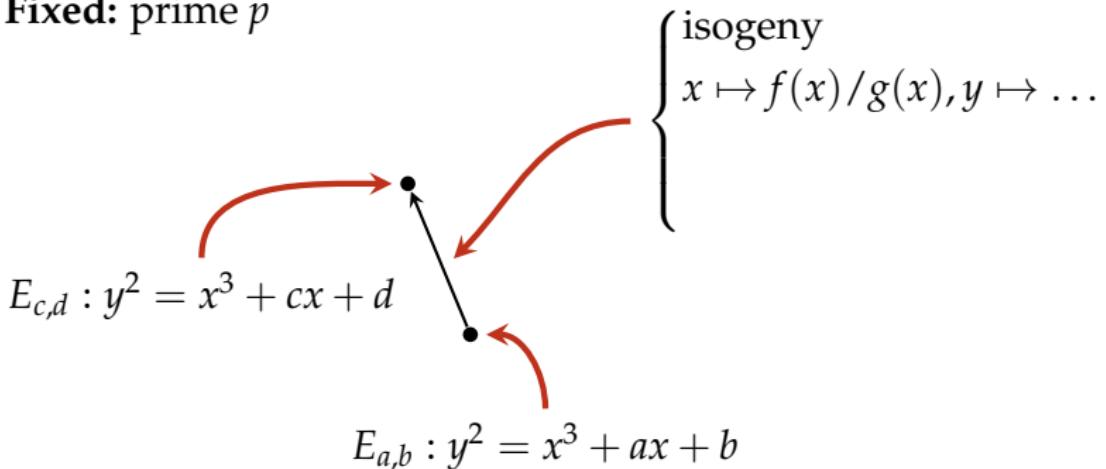
**Fixed:** prime  $p$

The diagram illustrates the relationship between two elliptic curves,  $E_{c,d}$  and  $E_{a,b}$ . The curve  $E_{c,d}$  is defined by the equation  $y^2 = x^3 + cx + d$ . The curve  $E_{a,b}$  is defined by the equation  $y^2 = x^3 + ax + b$ . Two points on the curves are connected by red arrows: one point on  $E_{c,d}$  has an arrow pointing to a point on  $E_{a,b}$ , and another point on  $E_{a,b}$  has an arrow pointing back to a point on  $E_{c,d}$ .

$$E_{c,d} : y^2 = x^3 + cx + d$$
$$E_{a,b} : y^2 = x^3 + ax + b$$

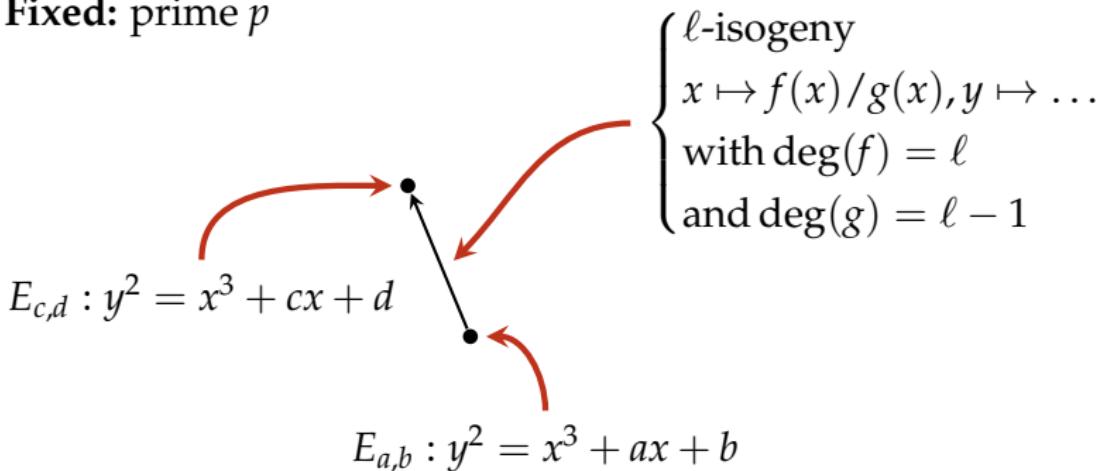
# Elliptic curves & isogenies (1)

Fixed: prime  $p$



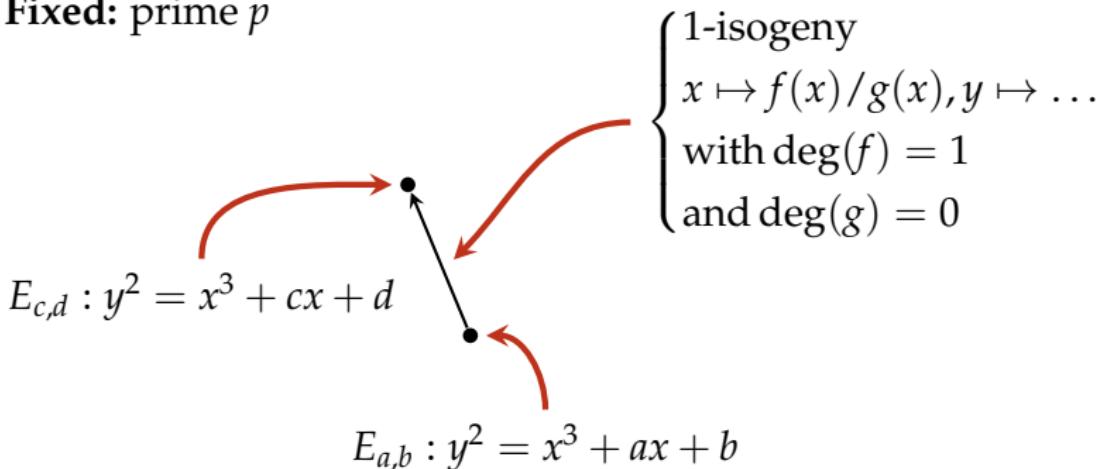
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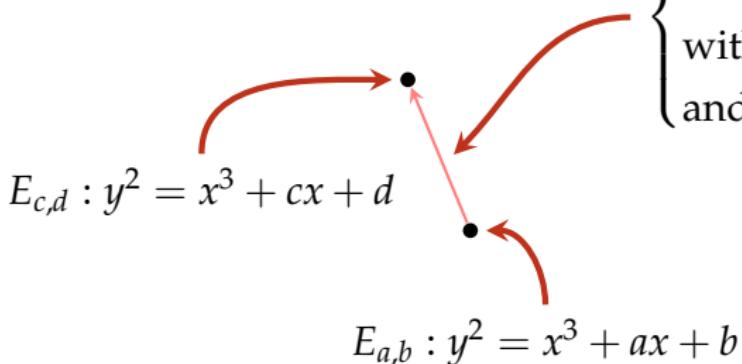
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# Elliptic curves & isogenies (1)

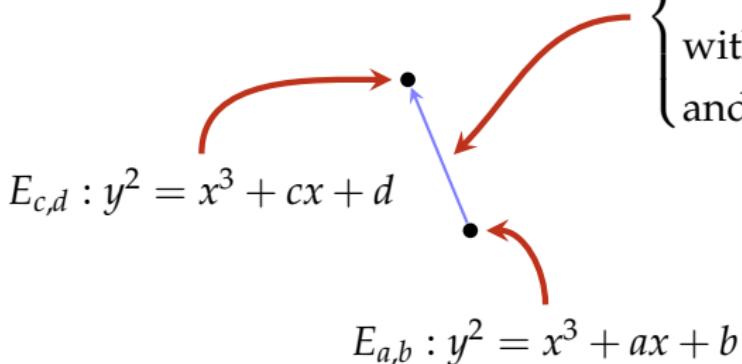
Fixed: prime  $p$



2-isogeny  
 $x \mapsto f(x)/g(x), y \mapsto \dots$   
with  $\deg(f) = 2$   
and  $\deg(g) = 1$

# Elliptic curves & isogenies (1)

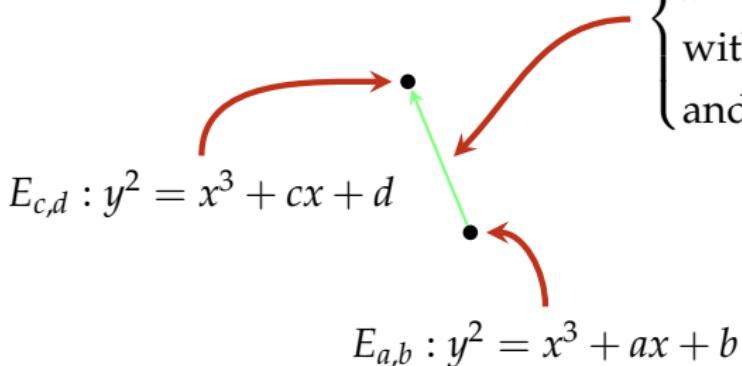
Fixed: prime  $p$



**3-isogeny**  
 $x \mapsto f(x)/g(x), y \mapsto \dots$   
with  $\deg(f) = 3$   
and  $\deg(g) = 2$

# Elliptic curves & isogenies (1)

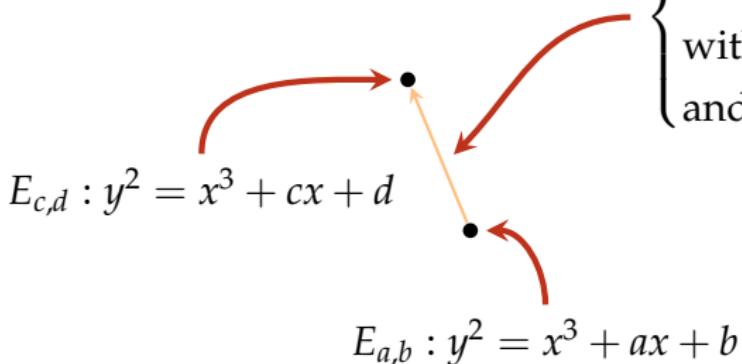
Fixed: prime  $p$



5-isogeny  
 $x \mapsto f(x)/g(x), y \mapsto \dots$   
with  $\deg(f) = 5$   
and  $\deg(g) = 4$

# Elliptic curves & isogenies (1)

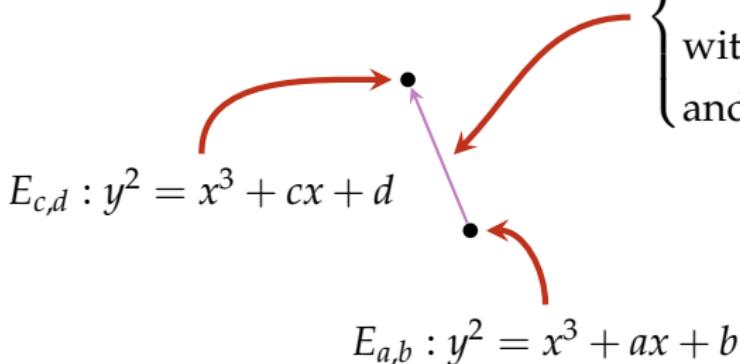
Fixed: prime  $p$



$\begin{cases} \text{7-isogeny} \\ x \mapsto f(x)/g(x), y \mapsto \dots \\ \text{with } \deg(f) = 7 \\ \text{and } \deg(g) = 6 \end{cases}$

# Elliptic curves & isogenies (1)

Fixed: prime  $p$



11-isogeny  
 $x \mapsto f(x)/g(x), y \mapsto \dots$   
with  $\deg(f) = 11$   
and  $\deg(g) = 10$

# Elliptic curves & isogenies (1)

Fixed: prime  $p$

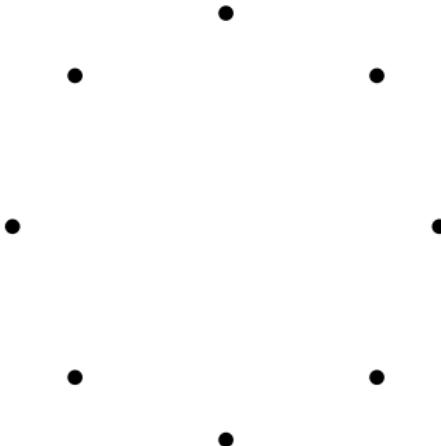
The diagram illustrates a 13-isogeny between two elliptic curves. It features two black dots representing points on the curves. A vertical dashed line connects these points. Red curved arrows point from the top curve to the bottom curve and from the bottom curve to the top curve, indicating the direction of the isogeny. Below each curve is its respective equation:

$$E_{c,d} : y^2 = x^3 + cx + d$$
$$E_{a,b} : y^2 = x^3 + ax + b$$

$\left\{ \begin{array}{l} \text{13-isogeny} \\ x \mapsto f(x)/g(x), y \mapsto \dots \\ \text{with } \deg(f) = 13 \\ \text{and } \deg(g) = 12 \end{array} \right.$

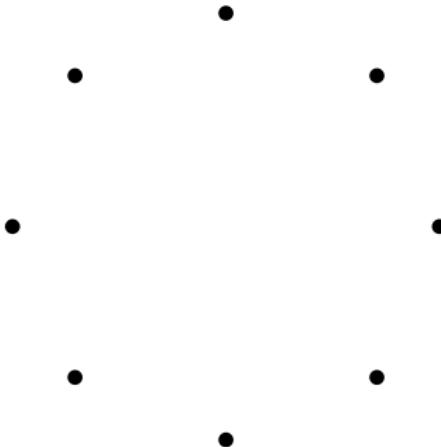
# Elliptic curves & isogenies (1)

**Fixed:** prime  $p$ ,  $\text{End}_{\mathbb{F}_p}(E_{a,b}) = \mathcal{O}_{\mathbb{Q}(\pi)}$



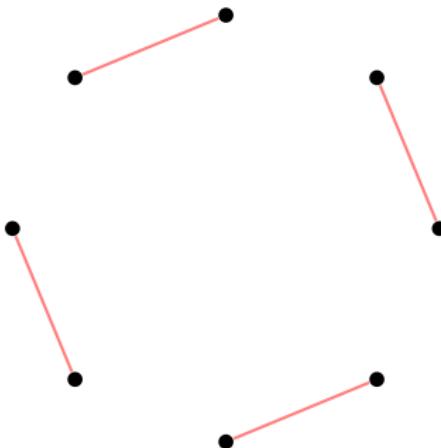
# Elliptic curves & isogenies (1)

**Fixed:** prime  $p$ ,  $\text{End}_{\mathbb{F}_p}(E_{a,b}) = \mathcal{O}_{\mathbb{Q}(\pi)}$ ,  $\ell = 2$



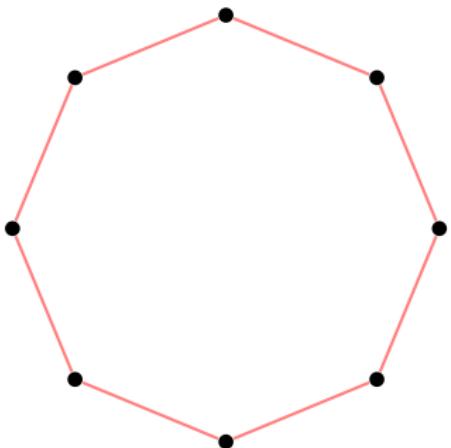
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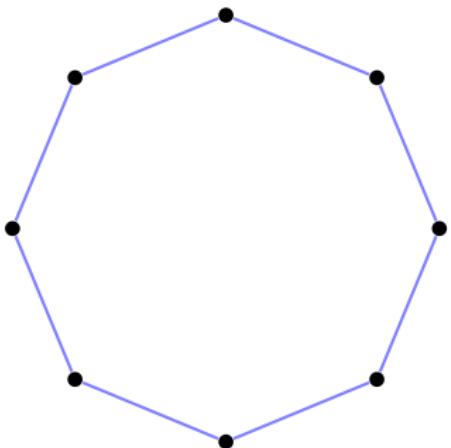
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**Fixed:** prime  $p$ ,  $\text{End}_{\mathbb{F}_p}(E_{a,b}) = \mathcal{O}_{\mathbb{Q}(\pi)}$ ,  $\ell = 2$



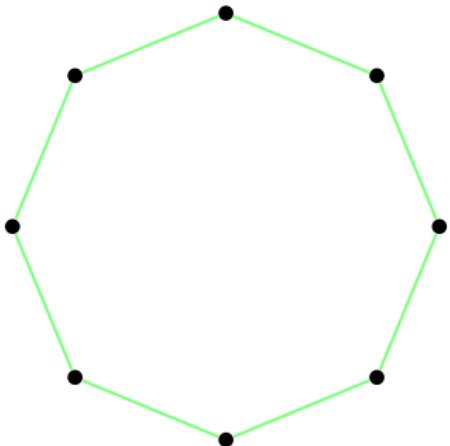
# Elliptic curves & isogenies (1)

**Fixed:** prime  $p$ ,  $\text{End}_{\mathbb{F}_p}(E_{a,b}) = \mathcal{O}_{\mathbb{Q}(\pi)}$ ,  $\ell = 3$



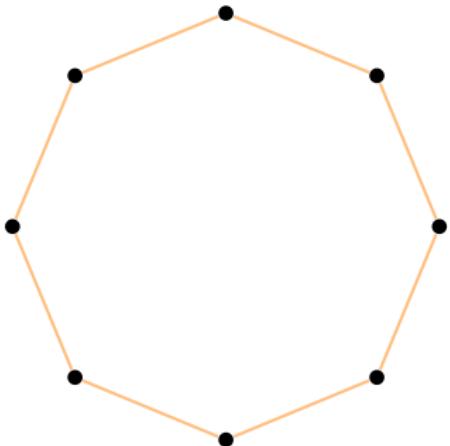
# Elliptic curves & isogenies (1)

**Fixed:** prime  $p$ ,  $\text{End}_{\mathbb{F}_p}(E_{a,b}) = \mathcal{O}_{\mathbb{Q}(\pi)}$ ,  $\ell = 5$



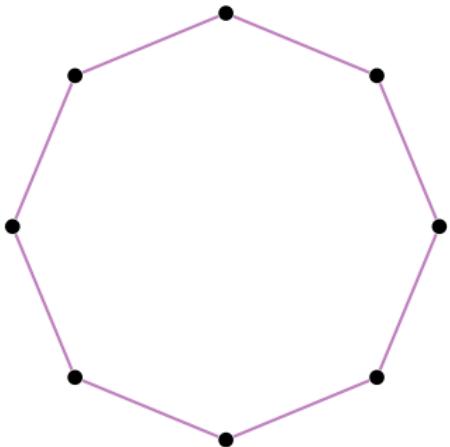
# Elliptic curves & isogenies (1)

**Fixed:** prime  $p$ ,  $\text{End}_{\mathbb{F}_p}(E_{a,b}) = \mathcal{O}_{\mathbb{Q}(\pi)}$ ,  $\ell = 7$



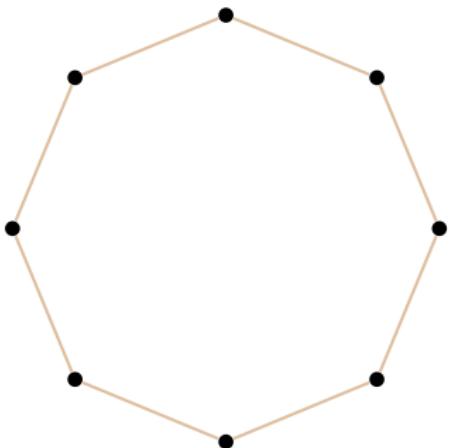
# Elliptic curves & isogenies (1)

**Fixed:** prime  $p$ ,  $\text{End}_{\mathbb{F}_p}(E_{a,b}) = \mathcal{O}_{\mathbb{Q}(\pi)}$ ,  $\ell = 11$

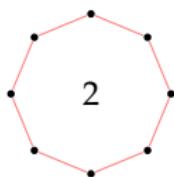


# Elliptic curves & isogenies (1)

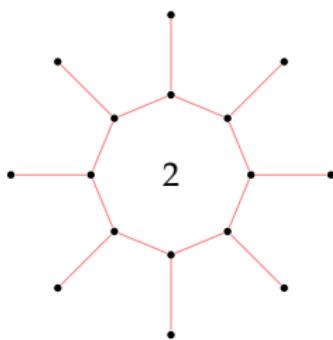
**Fixed:** prime  $p$ ,  $\text{End}_{\mathbb{F}_p}(E_{a,b}) = \mathcal{O}_{\mathbb{Q}(\pi)}$ ,  $\ell = 13$



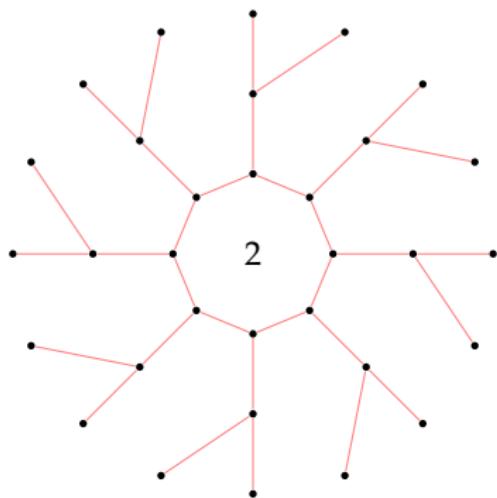
# Elliptic curves & isogenies (2)



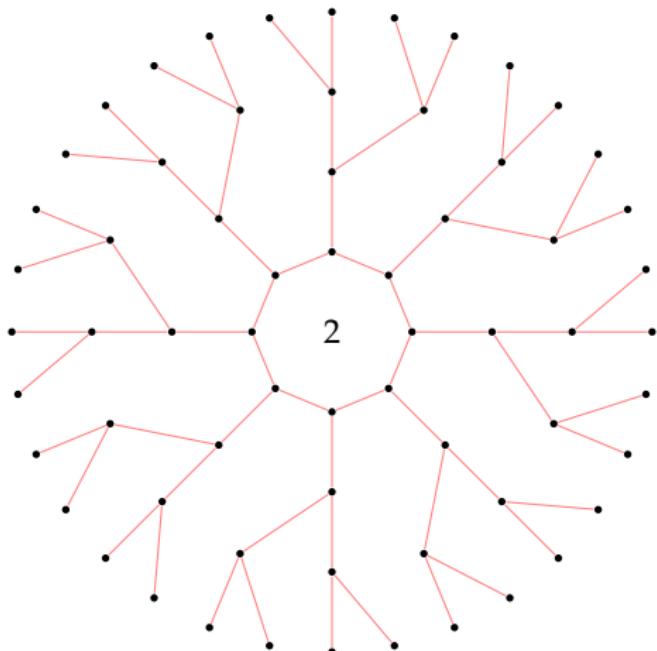
## Elliptic curves & isogenies (2)



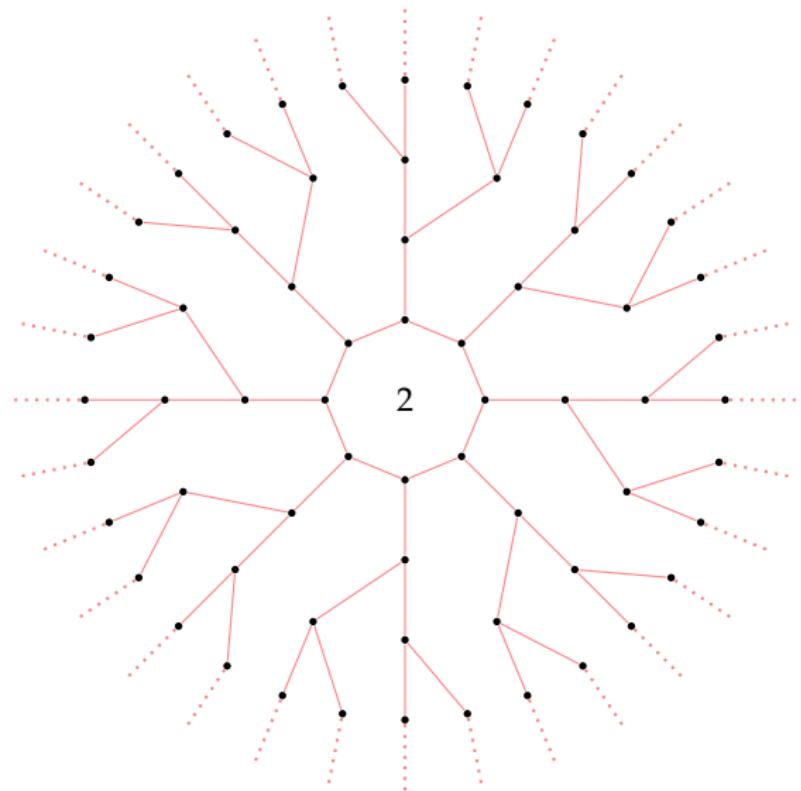
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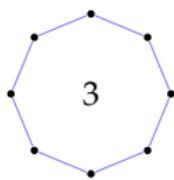
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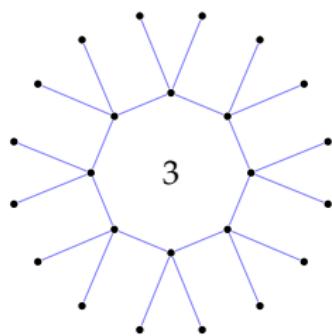
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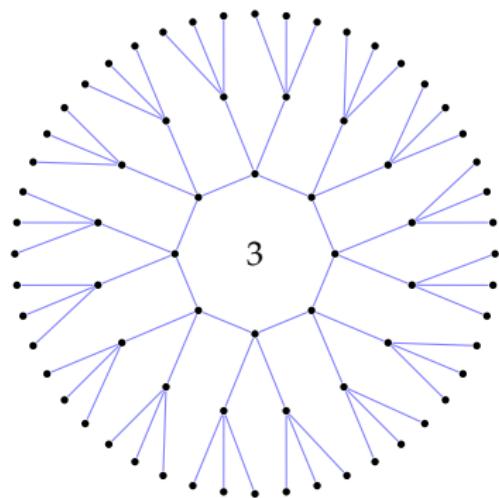
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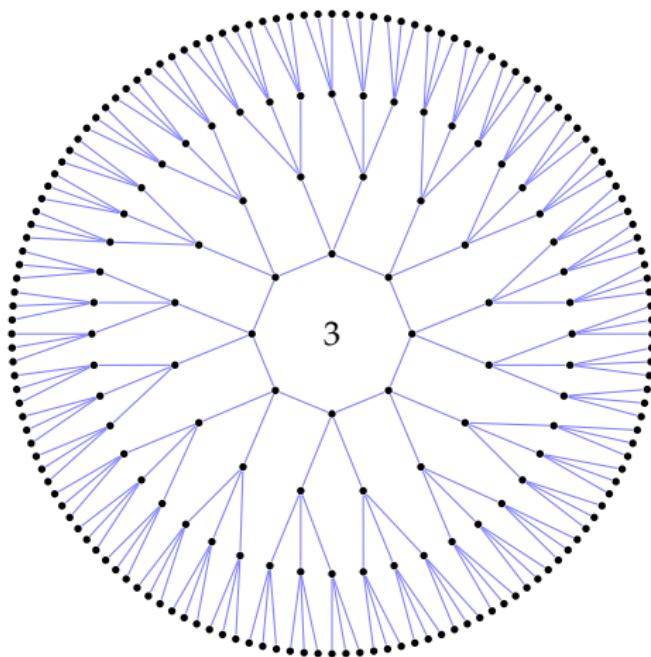
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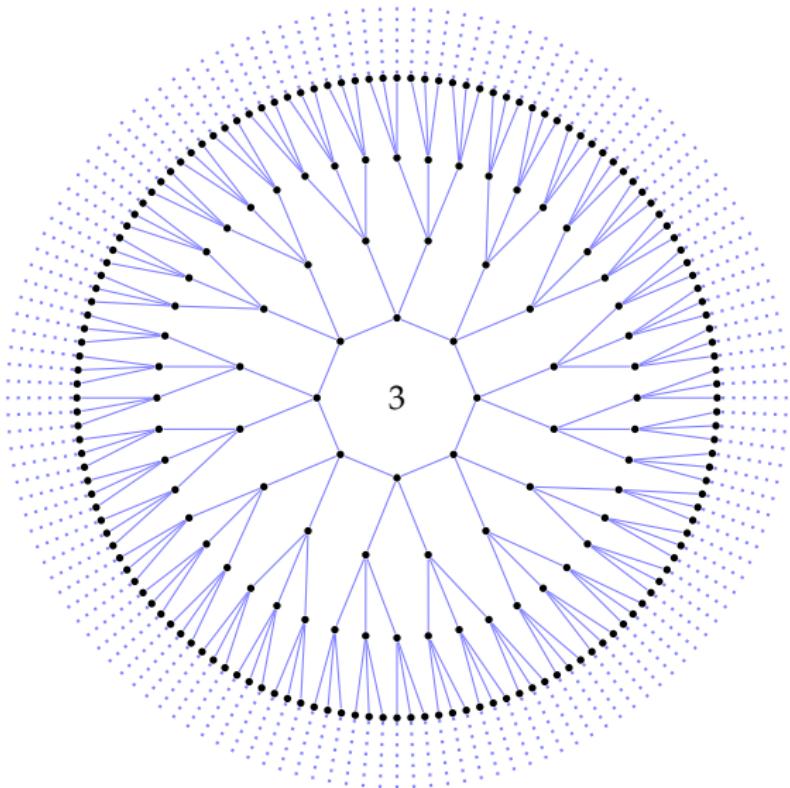
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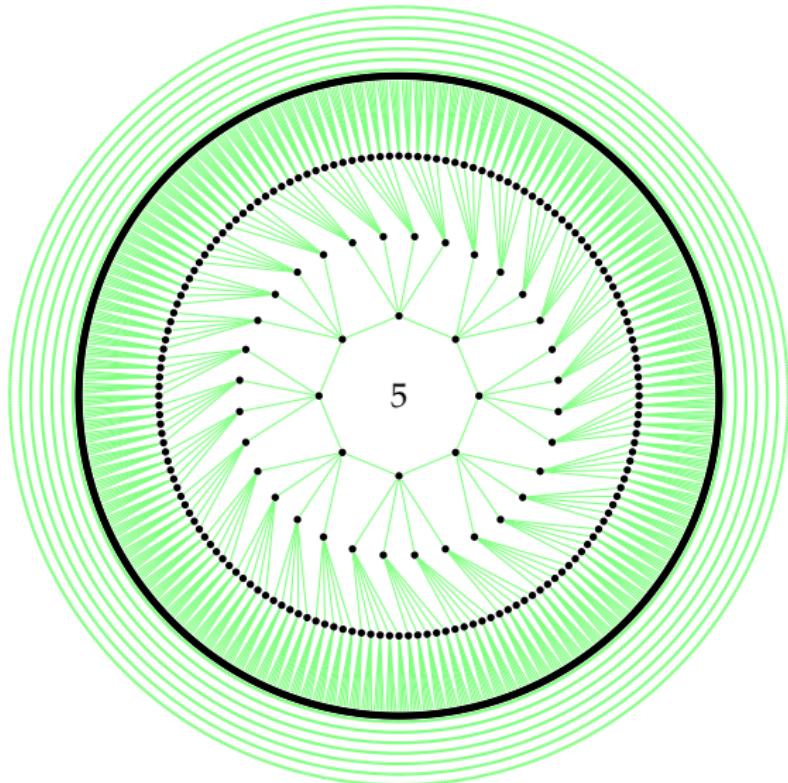
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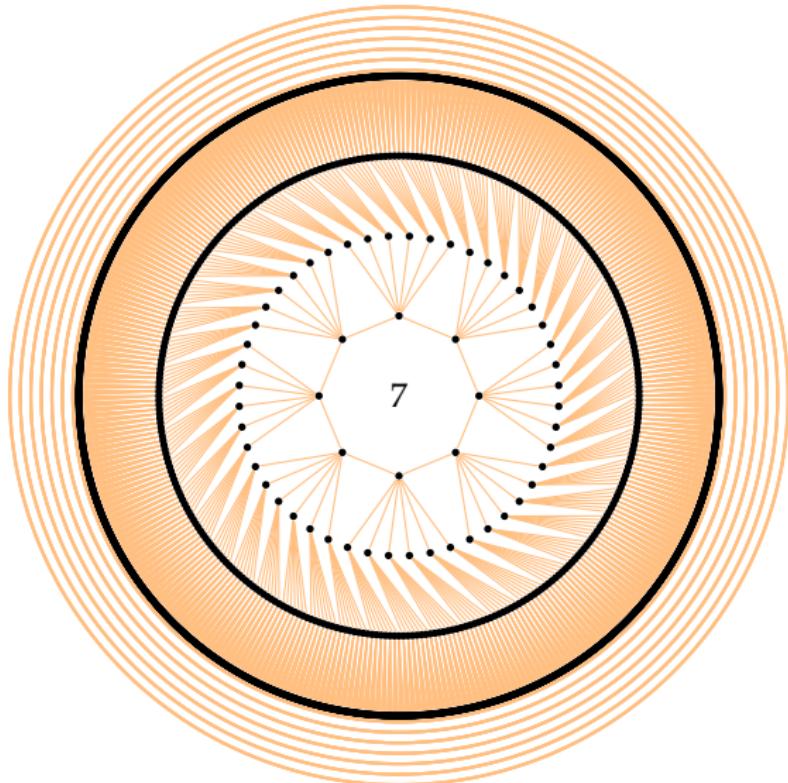
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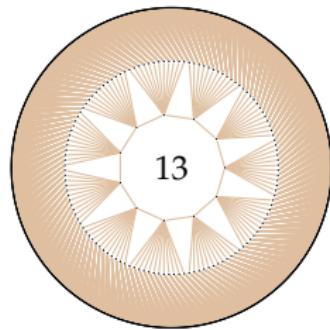
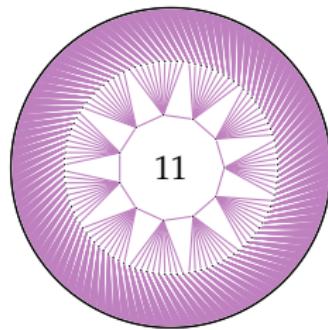
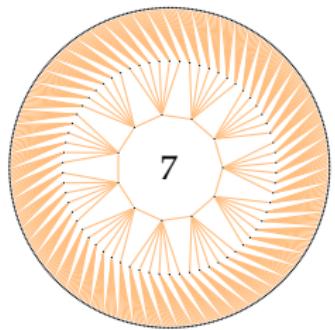
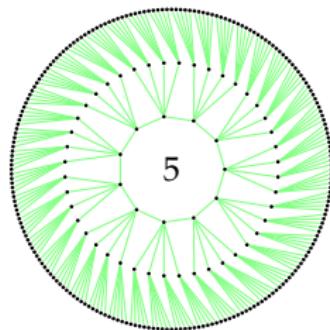
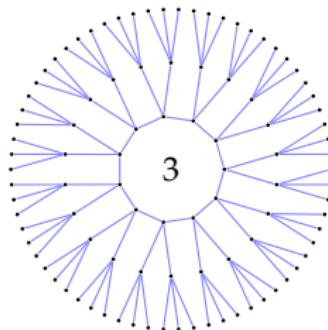
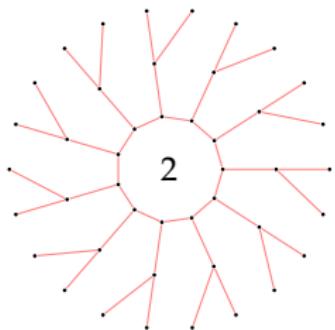
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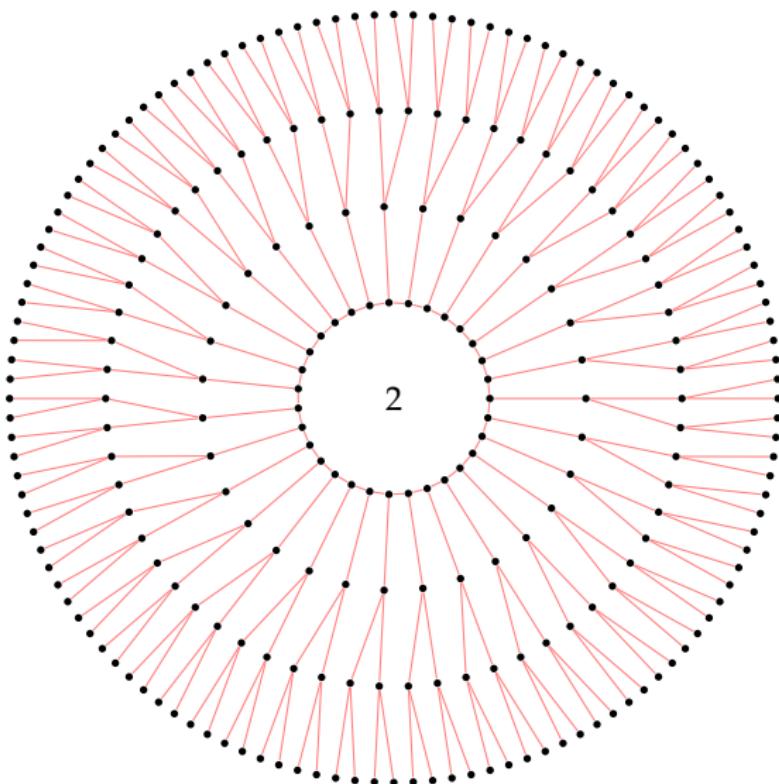
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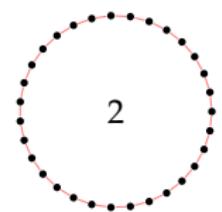
# Isogeny volcanoes



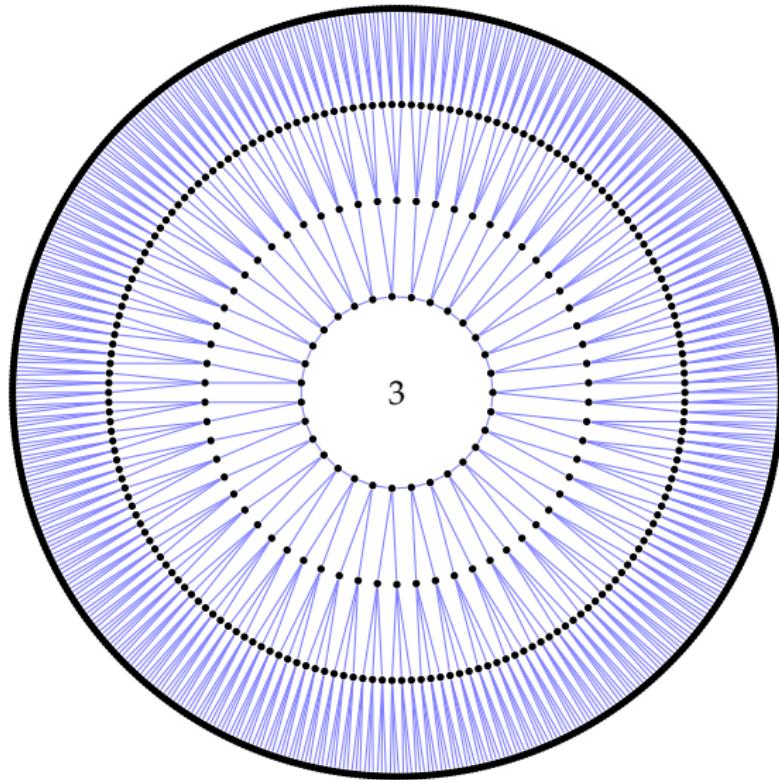
# Isogeny-based cryptography (1)



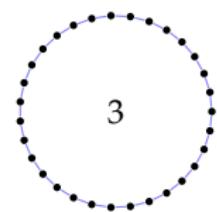
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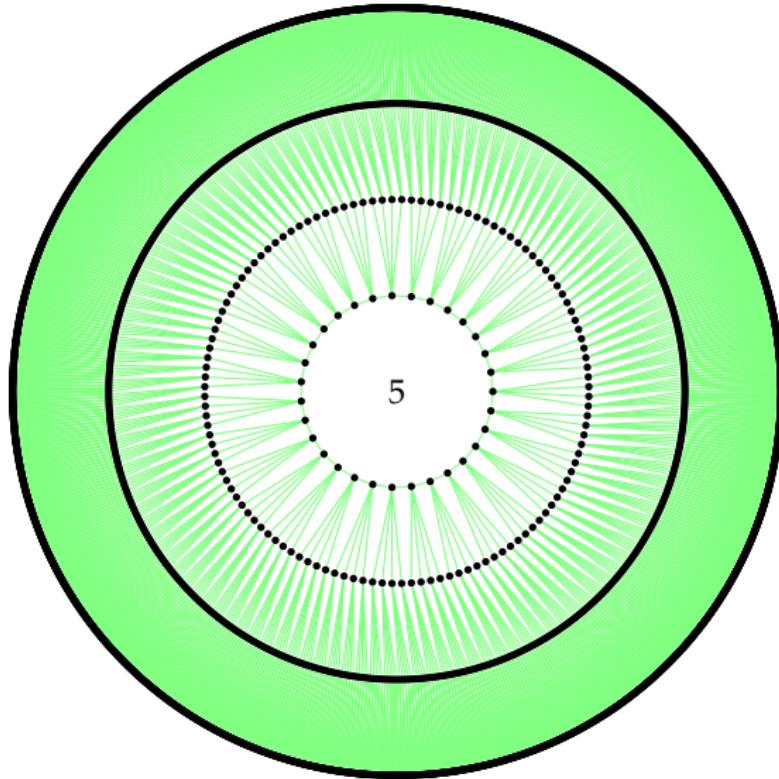
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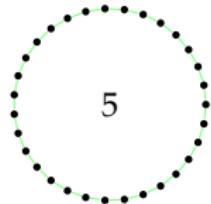
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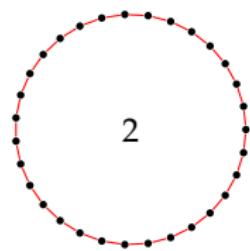
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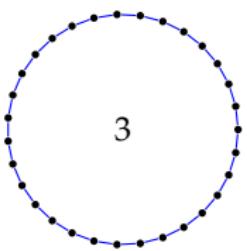
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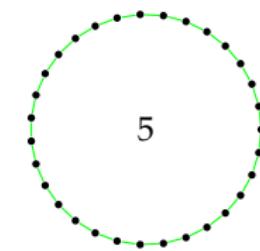
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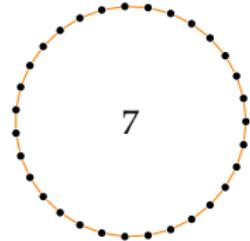
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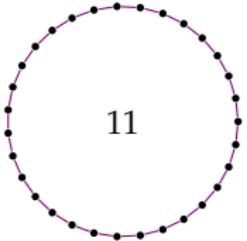
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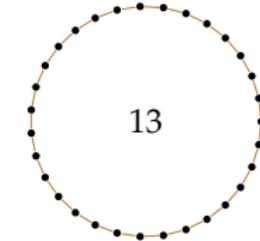
5



7

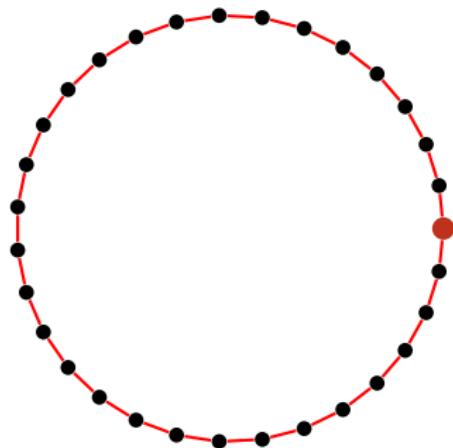


11



13

# Isogeny-based cryptography (1)



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# primes:

1

Work (per prime):

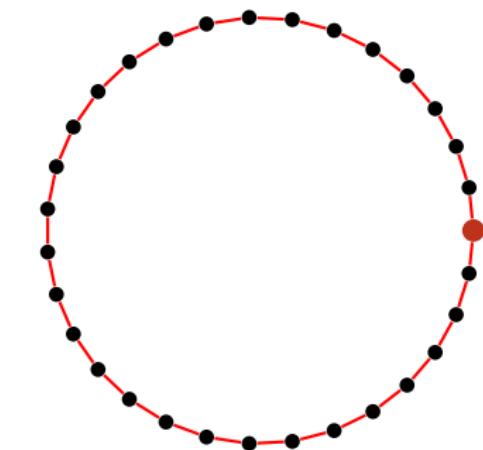
$\leq t$

Work (total):

$\leq t$

Entropy:

$t$



# Isogeny-based cryptography (1)

# primes:

1

Work (per prime):

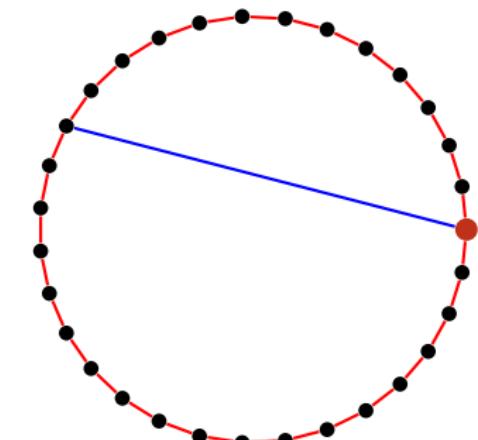
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# Isogeny-based cryptography (1)

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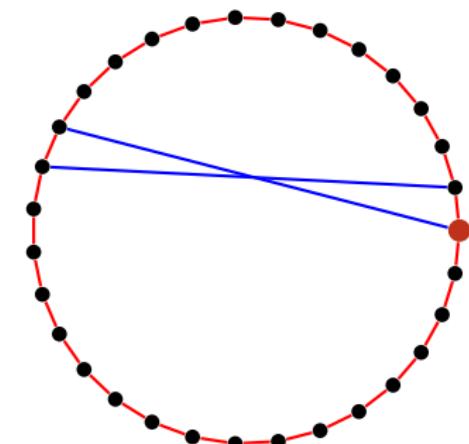
$\leq t$

Work (total):

$\leq t$

Entropy:

$t$



# Isogeny-based cryptography (1)

# primes:

$$2$$

Work (per prime):

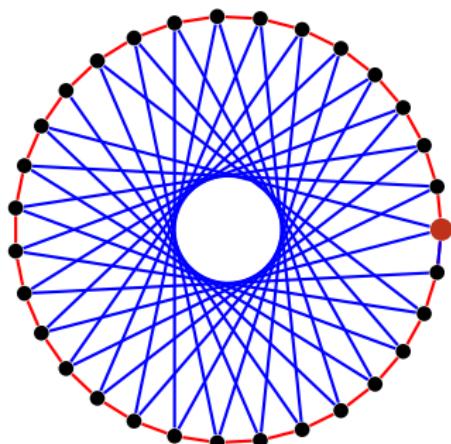
$$\leq t$$

Work (total):

$$\leq 2 \cdot t$$

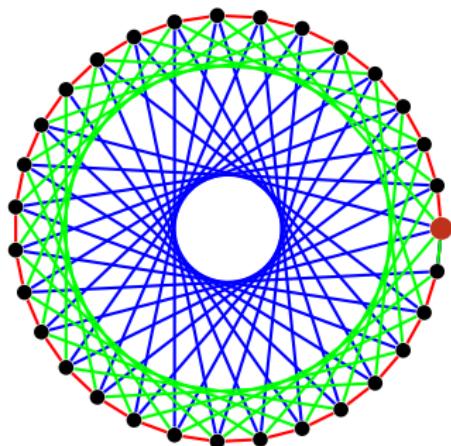
Entropy:

$$t^2$$



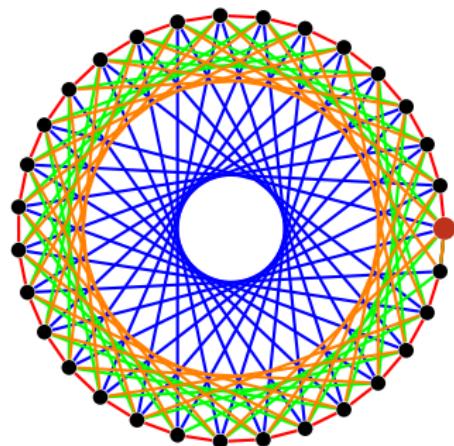
# Isogeny-based cryptography (1)

# primes:	3
Work (per prime):	$\leq t$
Work (total):	$\leq 3 \cdot t$
Entropy:	$t^3$



# Isogeny-based cryptography (1)

# primes:	4
Work (per prime):	$\leq t$
Work (total):	$\leq 4 \cdot t$
Entropy:	$t^4$



# Isogeny-based cryptography (1)

# primes:

5

Work (per prime):

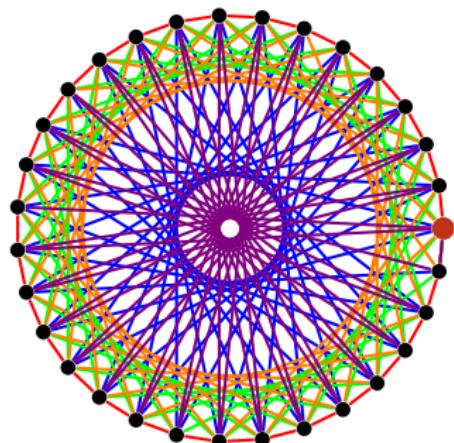
$\leq t$

Work (total):

$\leq 5 \cdot t$

Entropy:

$t^5$



# Isogeny-based cryptography (1)

# primes:

6

Work (per prime):

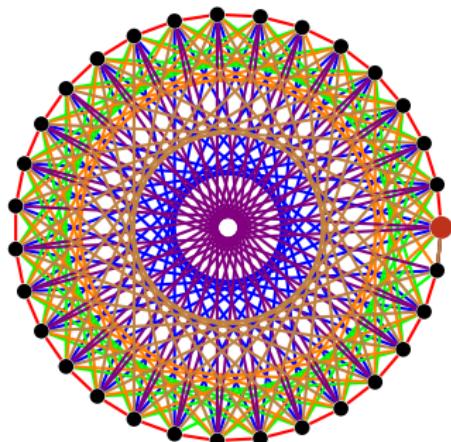
$\leq t$

Work (total):

$\leq 6 \cdot t$

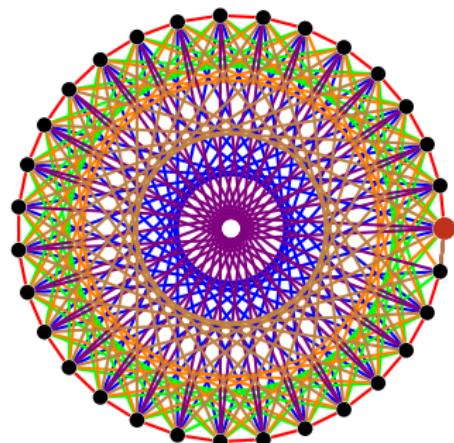
Entropy:

$t^6$



# Isogeny-based cryptography (1)

$$\begin{aligned}\text{\# primes:} & L \\ \text{Work (per prime):} & \leq t \\ \text{Work (total):} & \leq L \cdot t \\ \text{Entropy:} & t^L\end{aligned}$$



# OIDH & CSIDH

Two different ways to instantiate;

1. *Ordinary* isogeny Diffie–Hellman (OIDH)
2. *Supersingular* isogeny Diffie–Hellman (CSIDH)

The idea for OIDH first by Couveignes in '96 [Cou06]

- ⇒ Post-quantum security with very small keys [DKS18]
- ⇒ CSIDH almost identical but easier to instantiate [Cas+18]

# State of CSIDH

(~ NIST level I security)

## 1. CSIDH key exchange

- ▶ *Non-interactive* with 64-byte public keys
- ▶ ~ 80 ms for full exchange (not constant-time)

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( $\sim$  NIST level I security)

1. **CSIDH key exchange**
  - ▶ *Non-interactive* with 64-byte public keys
  - ▶  $\sim 80$  ms for full exchange (not constant-time)
2. Constant-time implementations [MCR18] (at  $\sim 246$  ms)

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1. **CSIDH** key exchange
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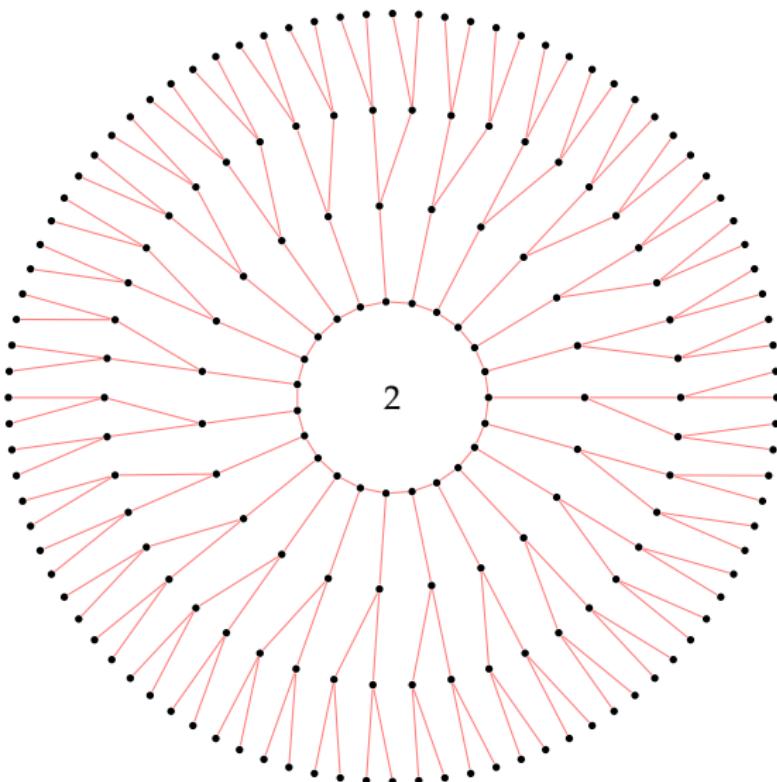
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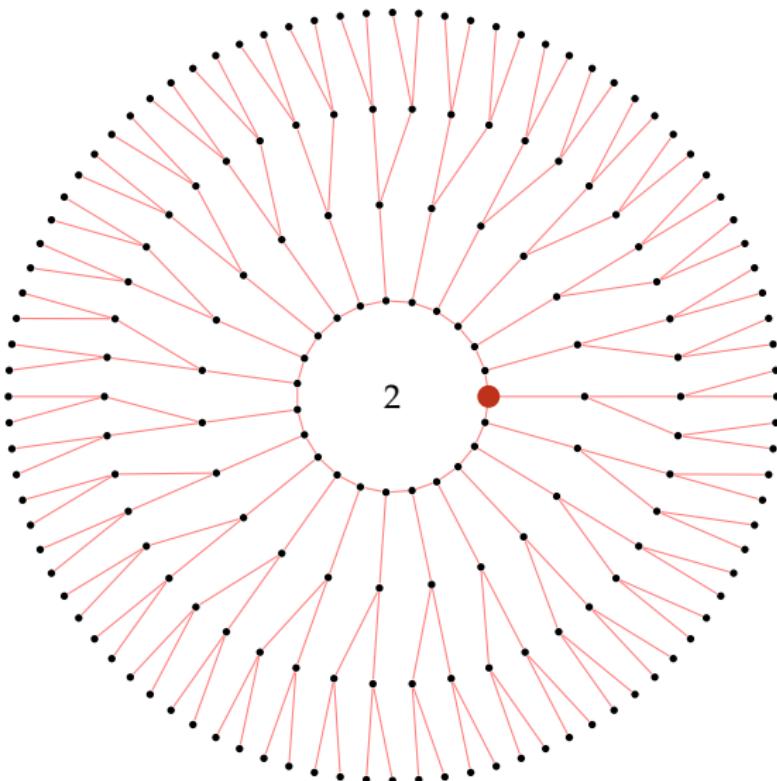
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Lots of stuff coming out!

## Isogeny-based cryptography (2)



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# primes: 1 ( $\ell = 2$ )

Work (per prime): 1

Work (total): 1

Entropy: 3



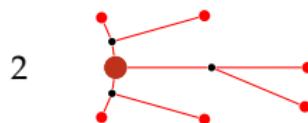
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# primes: 1 ( $\ell = 2$ )

Work (per prime): 2

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Entropy:  $3 \cdot 2$



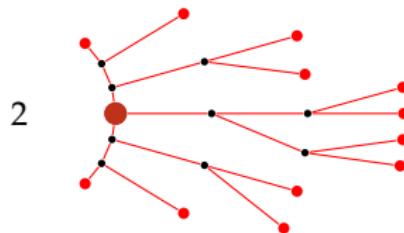
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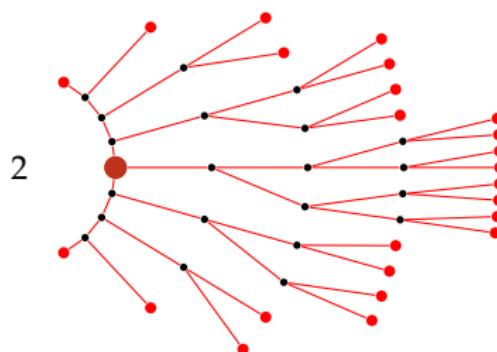
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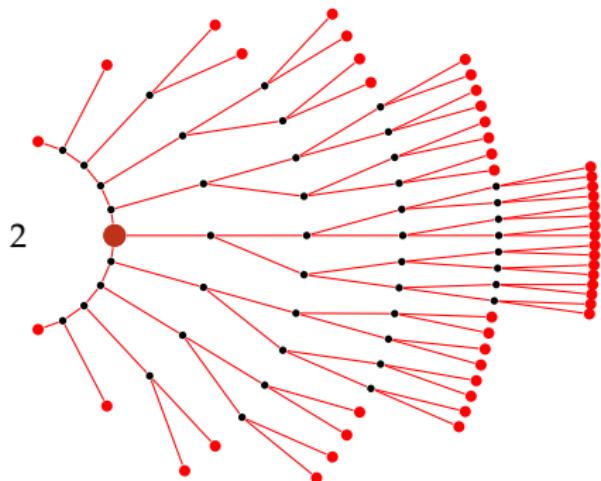
5

Work (total):

5

Entropy:

$3 \cdot 2^4$



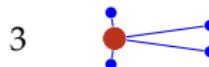
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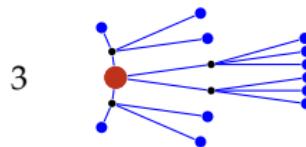
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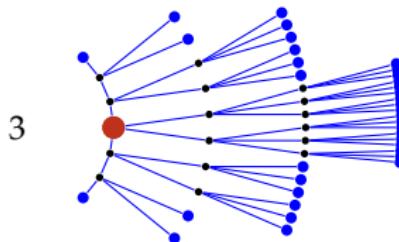
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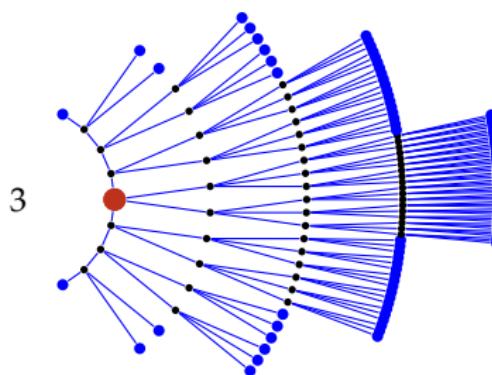
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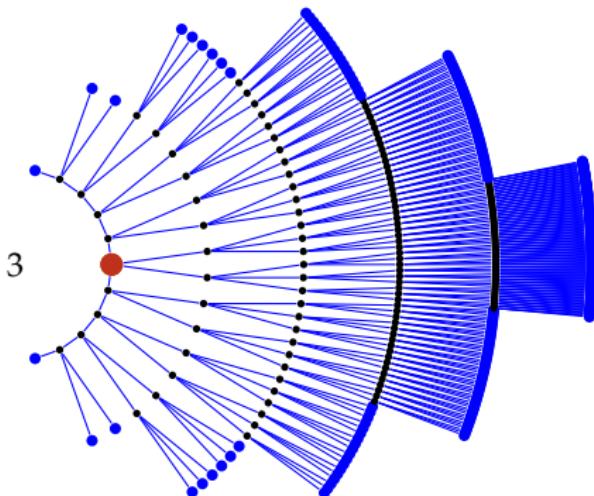
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# primes: 1 ( $\ell = 3$ )

Work (per prime): 5

Work (total): 5

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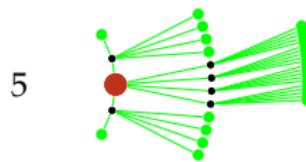
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# primes: 1 ( $\ell = 5$ )

Work (per prime): 2

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Entropy:  $6 \cdot 5$



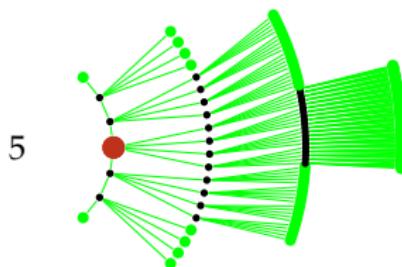
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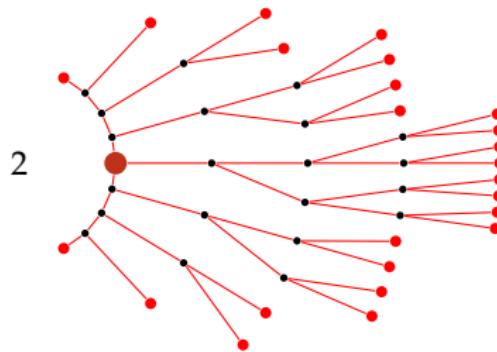
# Isogeny-based cryptography (2)

# primes: 1  
Work (per prime):  $t$   
Work (total):  $t$   
Entropy:  $\sim \ell^t$

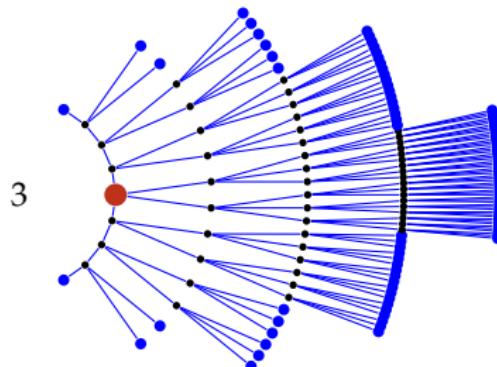


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Alice:

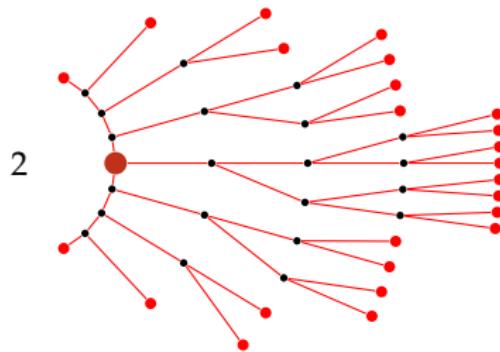


Bob:

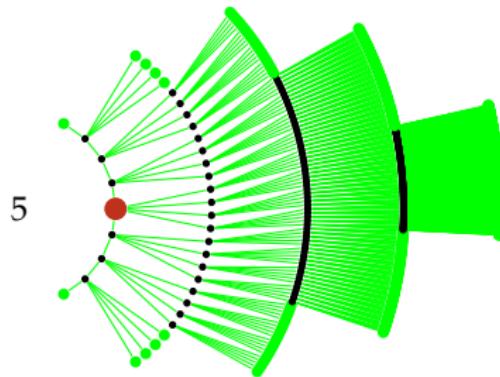


# Isogeny-based cryptography (2)

Alice:



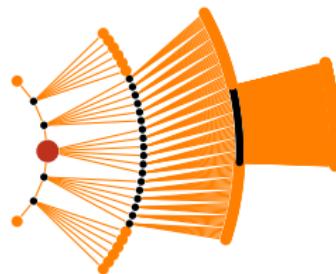
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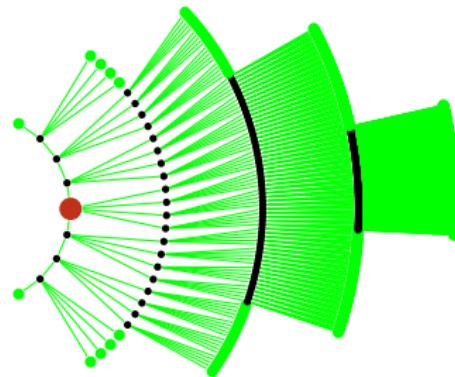
Alice:

7



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5



# State of SIDH / SIKE

(NIST level I security)

## 1. SIKE key encapsulation [Jao+]

- ▶ ~ 330-byte public keys / ciphertexts
- ▶ ~ 6.3 ms for enc + dec (constant-time)

(Recall CSIDH has 64-byte pk and ~ 246 ms exchange)

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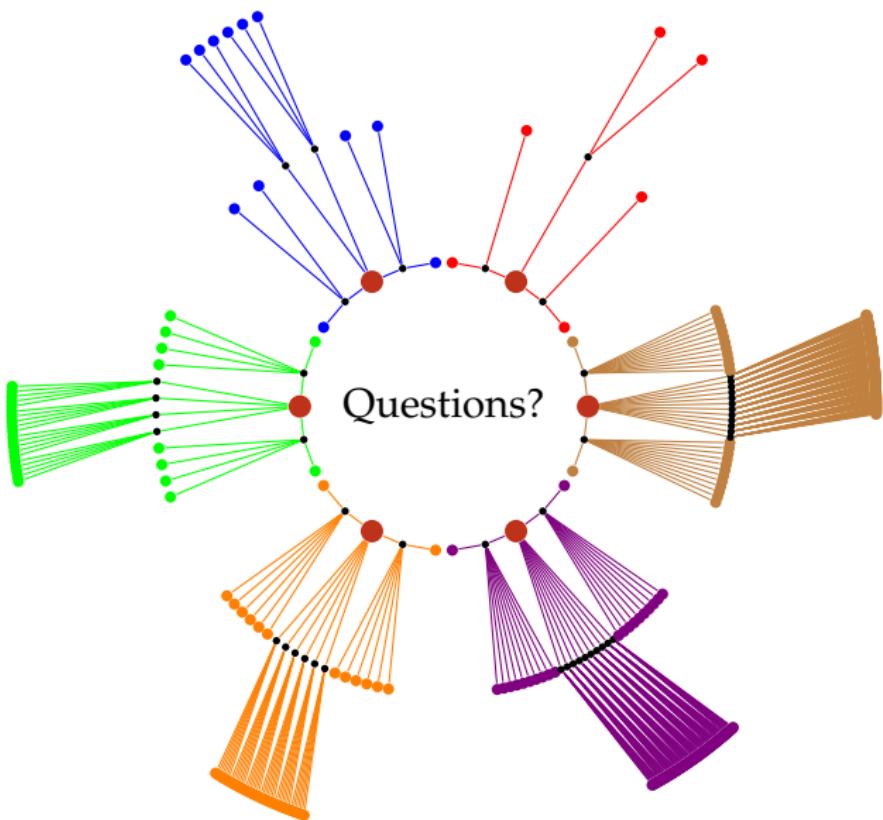
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Thanks!



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