

On Kummer Lines With Full Rational 2-torsion and Their Usage in Cryptography

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A paper by Karati and Sarkar at Asiacrypt'17 has pointed out the potential for Kummer lines in genus one, by observing that its SIMD-friendly arithmetic is competitive with the status quo. A more recent preprint explores the connection with (twisted) Edwards curves. In this paper we extend this work and significantly simplify their treatment. We show that their Kummer line is the x -line of a Montgomery curve translated by a point of order two, and exhibit a natural isomorphism to a twisted Edwards curve. Moreover, we show that the Kummer line presented by Gaudry and Lubicz can be obtained via the action of a point of order two on the y -line of an Edwards curve. The maps connecting these curves and lines are all very simple. As a result, a cryptographic implementation can use the arithmetic that is optimal for its instruction set at negligible cost.

CCS Concepts: • **Security and privacy** → **Cryptography**; **Public key encryption**.

Additional Key Words and Phrases: Montgomery curves, Edwards curves, Kummer lines, Montgomery ladder, Digital signatures

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1 INTRODUCTION

A decade after the introduction of public-key cryptography by Diffie and Hellman [1976] it was observed (independently) by Miller [1986] and Koblitz [1987] that one can instantiate protocols based on the hardness of the discrete logarithm problem with the group of rational points of an elliptic curve E defined over a finite field. Moreover, it was immediately noted by Miller that one can do a full key exchange by solely relying on the line of x -coordinates of points. That is, one can identify points with their inverses and as a result only work with points up to sign. In other words, one can work on the corresponding Kummer line $K = E/\{\pm 1\}$, possibly simplifying the arithmetic. Recently it was shown that one can also directly use K for digital signatures very efficiently with the qDSA scheme [Renes and Smith 2017, §2]. In short, Kummer lines are a very interesting topic of study from a cryptographic perspective.

Because a reduction in the number of field operations needed for a scalar multiplication directly affects the efficiency of the cryptographic scheme, there have been multiple proposals for Kummer lines. Probably the most available example is Curve25519 [Bernstein 2006a], which is the Kummer line of a Montgomery curve. One can show that every Montgomery curve is birationally equivalent to a twisted Edwards curve [Bernstein et al. 2008, Theorem 3.2], which currently needs the least

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number of field operations to perform group operations [Hisil et al. 2008] and underlies the very efficient FourQ curve [Costello and Longa 2015]. As a result, the Kummer lines of Montgomery and twisted Edwards curves are strongly related, and one can move easily from one to the other [Bernstein et al. 2008; Castryck et al. 2008]. Through the usage of theta functions Gaudry and Lubicz [2009, §6] derived yet another Kummer line. We shall refer to this as the *canonical* Kummer line, following the terminology of the genus 2 analogue presented by Renes and Smith [2017, §4]. By squaring its coefficients we arrive on a different variety, which we refer to as the *squared* Kummer line (again c.f. the genus 2 analogue [Bernstein 2006b; Chudnovsky and Chudnovsky 1986]). Although Gaudry and Lubicz only presented arithmetic on the canonical line, the differential addition formulae on the squared Kummer line are well-known [Bernstein and Lange 2015]. The squared Kummer line has the advantage that it is easier to find suitable small parameters, and it was shown by Karati and Sarkar [2017b] that its arithmetic leads to very efficient implementations when single-instruction multiple-data (SIMD) instructions are available.

In a follow-up paper [Karati and Sarkar 2017a] the same authors present connections to twisted Edwards curves. This requires the associated Legendre curve to be put in Montgomery form or have a rational point of order 4, or otherwise relies on the usage of a 2-isogeny. Consequently, there are case distinctions and one must deal with the doubling induced by moving through a 2-isogeny and its dual. In [Karati and Sarkar 2017a, Table 7] they present the possibility of birational maps and isogenies between the Legendre form for certain choices of small constants.

In this paper we significantly simplify the connections between the various Kummer lines. Since the field of definition of the canonical and squared Kummer lines corresponds to their rational 2-torsion, we shall assume all points of order 2 to be rational. In that case, we show that the squared (resp. canonical) Kummer arises as the x -line (resp. y -line) of a Montgomery (resp. Edwards) curve translated by a suitable point of order 2. Moreover, a third Kummer line (referred to as the *intermediate* Kummer) appears as the y -line of a *twisted* Edwards curve via a translation by a point of order 2. These observations induce very simple isomorphisms between them. Furthermore, the respective translations by a point of order 2 lead to fast isomorphisms (in fact, involutions) with the well-known x -lines (or y -lines) of Montgomery, Edwards and twisted Edwards curves. As a result, we unify the most popular Kummer lines in the literature and conclude that their usage is completely interchangeable on an implementation level. For example, we can directly use the squared Kummer line in the qDSA scheme through its connection with a Montgomery curve [Renes and Smith 2017, §3]. Moreover, although there exist efficient implementations of Montgomery curves based on 4-way SIMD parallelization by optimizing the field arithmetic [Faz-Hernández and López 2015], it is unclear how to optimally parallelize instructions 4-way on the level of the x -line [Chou 2015]. This is straightforward on the squared Kummer line, and therefore by extension becomes trivial on Montgomery curves with full rational 2-torsion by moving through the isomorphism. Of course, if desired, one can also do arithmetic on the full group of points of the twisted Edwards curve (as also noted by Karati and Sarkar [2017a]). In particular, we provide isomorphic Montgomery and twisted Edwards models for all the Kummer lines present in [Karati and Sarkar 2017a, Table 7] (see Table 1 in §3.2).

2 PRELIMINARIES

Let k be a field such that $\text{char}(k) \neq 2$. This assumption is implicit in the whole document unless mentioned otherwise. An elliptic curve is a smooth projective curve of genus 1 with a specified base point O , and it is said to be defined over k if E is defined over k and $O \in E(k)$ [Silverman 2009, §III.3]. Its points form an abelian group with neutral element O . One can show [Silverman 2009,

Proposition III.3.1] that any elliptic curve defined over k can be put in *Weierstrass form*¹

$$y^2 + a_1xy + a_3y = x^3 + a_2x^2 + a_4x + a_6 \subset \mathbb{P}^2$$

with $O = (0 : 1 : 0)$, but the curves we consider in this paper are not necessarily in this standard model. If we want to emphasize the base point of the curve we are working with, we shall write (E, O) .

2.0.1 Montgomery curves. Let $A, B \in k$ such that $B(A^2 - 4) \neq 0$ and let

$$M/k : By^2 = x^3 + Ax^2 + x$$

be (the affine part of) a smooth projective curve of genus one. We may write $M_{A,B}$ for M to emphasize the coefficients of the curve we are referring to. Denoting $O_M = (0 : 1 : 0)$, the elliptic curve (M, O_M) is commonly referred to as a *Montgomery curve* [Montgomery 1987] and is ubiquitous in elliptic-curve-based cryptography protocols (see e.g. [Bernstein 2006a]). By projecting the points to \mathbb{P}^1 via the (surjective) map

$$\mathbf{x} : M \rightarrow \mathbb{P}^1$$

$$(X : Y : Z) \mapsto \begin{cases} (X : Z) & \text{if } Z \neq 0 \\ (1 : 0) & \text{if } Z = 0 \end{cases},$$

the projective line inherits a pseudo-group structure. That is, by viewing \mathbb{P}^1 as the image of (M, O_M) under \mathbf{x} we obtain the well-known scalar multiplication on \mathbb{P}^1 as first used by Montgomery [1987, §10]. Since inversion on (affine points of) (M, O_M) is simply negation of the y -coordinate, it is immediate that \mathbf{x} is well-defined on $(M, O_M)/\{\pm 1\}$ and induces a bijection between $(M, O_M)/\{\pm 1\}$ and \mathbb{P}^1 . We denote the projective line with the induced pseudo-group structure of (M, O_M) by $K_M^{O_M}$, and refer to it as its Kummer line.

Now suppose that $T \in (M, O_M)$ is a point such that $[2]T = O_M$. Then the translation-by- T map

$$\begin{aligned} \tau_T : (M, O_M) &\rightarrow (M, T) \\ P &\mapsto P + T \end{aligned}$$

is an isomorphism of elliptic curves. Moreover, the map \mathbf{x} is again well-defined on $(M, T)/\{\pm 1\}$ and we denote its Kummer line by K_M^T . In summary, we have a commutative diagram

$$\begin{array}{ccc} (M, O_M) & \xleftarrow{\tau_T} & (M, T) \\ \downarrow \mathbf{x} & & \downarrow \mathbf{x} \\ K_M^{O_M} & \xleftarrow{\bar{\tau}_T} & K_M^T \end{array}$$

where $\bar{\tau}_T$ is the induced isomorphism (again, involution) between the corresponding Kummer lines. For example, we obtain the map $\bar{\tau}_{(0,0)} : (X : Z) \mapsto (Z : X)$. Since $\#(M, O_M)[2] = 4$, there are at most two other points of order 2. This gives rise to only a single non-trivial action on the Kummer line $K_M^{O_M}$, since the other is simply the composition with $\bar{\tau}_{(0,0)}$.

¹ We shall in many cases talk about *affine* curves and maps for simplicity, but always mean their projective counterparts. This depends on the particular embedding of the affine curve into projective space, but it should be clear from context what is meant. In particular, we always embed Montgomery curves into \mathbb{P}^2 while (twisted) Edwards curves are embedded into \mathbb{P}^3 (as opposed to $\mathbb{P}^1 \times \mathbb{P}^1$, which is also commonplace).

2.0.2 *Twisted Edwards curves.* Let $\alpha, \delta \in k$ such that $\alpha\delta(\alpha - \delta) \neq 0$ and let

$$\alpha x^2 + y^2 = 1 + \delta x^2 y^2$$

be (the affine part of) a smooth projective curve of genus one. This is commonly referred to as the *twisted Edwards* model [Bernstein et al. 2008], where the base point is chosen as $O_E = (0, 1)$. It is closely related to a Montgomery curve via a birational map [Bernstein et al. 2008, Theorem 3.2(i)].

Embedding the curve into \mathbb{P}^2 via $(x, y) \mapsto (x : y : 1)$ gives two singularities at $(1 : 0 : 0)$ and $(0 : 1 : 0)$. We can resolve these by blowing up (see e.g. [Hisil 2010, §2.3.4] or [Galbraith 2012, Lemma 9.12.18]) to obtain the curve

$$E/k = V(\alpha X^2 + Y^2 - Z^2 - \delta T^2, XY - TZ) \subset \mathbb{P}^3 \quad (1)$$

and embedding $(x, y) \mapsto (xy : x : y : 1)$. When referring to twisted Edwards curves, we will mean its embedding into \mathbb{P}^3 conform (1) and may write $E_{\alpha, \delta}$ to emphasize its coefficients. For affine points we will sometimes use the affine notation and expect that this should not cause confusion. Note that this is a purely theoretical tool, since once all is set and done the cryptographically relevant arithmetic is performed in a prime order subgroup in which all points are affine. On E there exist 4 points at infinity

$$\begin{aligned} \Omega_1 &= (1 : \sqrt{\delta/\alpha} : 0 : 0), & \Omega_2 &= (1 : -\sqrt{\delta/\alpha} : 0 : 0), \\ \omega_1 &= (1 : 0 : \sqrt{\delta} : 0), & \omega_2 &= (1 : 0 : -\sqrt{\delta} : 0), \end{aligned}$$

where Ω_1, Ω_2 have order 2 and ω_1, ω_2 have order 4 on (E, O_E) .

Similar to the x -map for Montgomery curves arising as a projection away from O , we have a y -map

$$\begin{aligned} y : E &\rightarrow \mathbb{P}^1 \\ (T : X : Y : Z) &\mapsto (Y : Z), \\ (\Omega_1, \Omega_2) &\mapsto ((1 : \sqrt{\delta/\alpha}), (1 : -\sqrt{\delta/\alpha})) \end{aligned}$$

corresponding to projection away from Ω_1 . Since inversion in this case is negation of the x -coordinate, this map is well-defined on $(E, O_E)/\{\pm 1\}$ and we denote the Kummer line by $K_E^{O_E}$. The point $S = (0 : 0 : -1 : 1)$ of order 2 induces the commutative diagram

$$\begin{array}{ccc} (E, O_E) & \xleftarrow{\tau_S} & (E, S) \\ \downarrow y & & \downarrow y \\ K_E^{O_E} & \xleftarrow{\bar{\tau}_S} & K_E^S \end{array} \quad (2)$$

where $\bar{\tau}_S : (Y : Z) \mapsto (-Y : Z)$. The two other 2-torsion points induce one other non-trivial translation (analogous to the Montgomery model).

2.0.3 *Edwards curves.* Let $c \in \bar{k}$ such that $c^5 \neq c$ and

$$\mathcal{E} : x^2 + y^2 = c^2(1 + x^2 y^2)$$

again a smooth curve of genus 1. This is technically only a subset of the set of curves of the form $x^2 + y^2 = c^2(1 + dx^2 y^2)$ originally defined as Edwards curves by Bernstein and Lange [2007]. But in this paper we are only concerned with the case $d = 1$, which corresponds to the form introduced

by Edwards [2007], who observed that its arithmetic with respect to the base point $O_E = (0, c)$ is extremely symmetric. As above, we use the smooth model inside \mathbb{P}^3 containing the elements

$$\begin{aligned}\Theta_1 &= (1 : c : 0 : 0), & \Theta_2 &= (1 : -c : 0 : 0) \\ \theta_1 &= (1 : 0 : c : 0), & \theta_2 &= (1 : 0 : -c : 0),\end{aligned}$$

where Θ_1, Θ_2 resp. θ_1, θ_2 have order 2 resp. 4 on (\mathcal{E}, O_E) . Again, we have a projection to \mathbb{P}^1

$$\begin{aligned}\mathbf{y} : \mathcal{E} &\rightarrow \mathbb{P}^1 \\ (T : X : Y : Z) &\mapsto (Y : Z) \\ (\Theta_1, \Theta_2) &\mapsto ((1 : c), (1 : -c)) .\end{aligned}$$

We denote the Kummer line of (\mathcal{E}, O_E) obtained by projection through \mathbf{y} by $K_E^{O_E}$. For any point 2-torsion point R of (\mathcal{E}, O_E) we obtain a commutative diagram as in (2) by translation by R and denote the Kummer line by K_E^R .

2.0.4 Rationality and quadratic twists. Suppose that q is a prime power and $k = \mathbb{F}_q$ is a finite field. Then any elliptic curve E defined over \mathbb{F}_q will have a quadratic twist, i.e. an elliptic curve E^t which is \mathbb{F}_{q^2} -isomorphic but not \mathbb{F}_q -isomorphic to E . This is unique up to \mathbb{F}_q -isomorphism (hence why we talk about *the* quadratic twist).

In all the curve models (c.f. the above) that we consider there is an immediate connection between \mathbb{F}_q -rational points on the Kummer line \mathcal{K}_E of E , and \mathbb{F}_q -rational points of E and E^t . As such, when thinking about Kummer lines it is natural not to distinguish these (i.e., to consider everything up to \mathbb{F}_{q^2} -isomorphism). As a result, although some maps may only be defined over \mathbb{F}_{q^2} , this will at most induce a twist. Since we are only concerned with the \mathbb{F}_q -rational points of the Kummer line, this is not an issue. In all that follows we *could* easily make everything defined over \mathbb{F}_q , but as we shall see in §4 this may limit us when finding instantiations.

3 MAPS BETWEEN KUMMER LINES

In this section we present the theoretical basis. We observe first that many Kummer lines have appeared in the literature; the work of Gaudry and Lubicz [2009] present the so-called *canonical* Kummer line, while Karati and Sarkar use² the *squared* Kummer line [Karati and Sarkar 2017b]. Moreover, there is the *x*-line of Montgomery curve (e.g. Curve25519 [Bernstein 2006a] by Bernstein) and the *y*-line of a (twisted) Edwards curve [Castrick et al. 2008; Farashahi and Hosseini 2017]. It is not immediately clear how these are all connected, in particular the relation between the (canonical and squared) Kummer lines and Montgomery and (twisted) Edwards curves is not clear. Though a recent paper by Karati and Sarkar [2017a] provides some connections, this is not completely satisfying. For instance, it relies on having rational points or using 2-isogeny, and does not give a unique connection.

In this section we settle this and, in essence, show that they are all the same up to isomorphism. These isomorphisms are natural and simple (including computationally) and lead to natural connections between all the above Kummer lines. The core is summarized in Theorem 3.4, and a more complete overview is shown in Appendix A.

3.1 Models with rational 2-torsion

It is immediate (through their description via theta functions) that the canonical and squared Kummer lines are projections of curves that have full rational 2-torsion. As such, we shall always

² The formulas for this model had already appeared in the Explicit-Formulas Database [Bernstein and Lange 2015] referring to a discussion between Bernstein, Kohel and Lange and contributing the main idea to Gaudry [2006].

assume to have this. We begin by showing that this allows a nice parametrization of Montgomery curves.

PROPOSITION 3.1. *Let k be a field such that $\text{char}(k) \neq 2$ and let $(M_{A,B}, O_M)$ be a Montgomery curve such that $M_{A,B}[2] \subset M_{A,B}(k)$. Then there exist $a, b \in \bar{k}^*$ such that $ab(a^4 - b^4) \neq 0$ and $a^2/b^2 \in k$ such that*

$$A = -\frac{a^4 + b^4}{a^2b^2}, \quad \Delta_M = 16B^6 \cdot \frac{(a^4 - b^4)^2}{a^4b^4}.$$

Moreover, its points of order 2 are $(0 : 0 : 1)$, $(a^2 : 0 : b^2)$ and $(b^2 : 0 : a^2)$.

PROOF. As $M_{A,B}[2] \subset M_{A,B}(k)$, the polynomial $x^2 + Ax + 1$ splits over k and thus $\sqrt{A^2 - 4} \in k$. Now fix any $b \in \bar{k}^*$ and take $a \in \bar{k}^*$ such that $a^2/b^2 = (\sqrt{A^2 - 4} - A)/2$. Note that $\sqrt{A^2 - 4} - A \neq 0, \pm 2$ because $\text{char}(k) \neq 2$. Moreover $a^4 - b^4 = 0 \iff a^4/b^4 - 1 = 0 \iff a^2/b^2 = \pm 1$. Again, this is not possible since $\text{char}(k) \neq 2$. The statements for A, Δ_M and the 2-torsion points are simple calculations, recalling that M has discriminant $\Delta_M = 16B^6(A^2 - 4)$. \square

For simplicity we would like to have $B = 1$. Note that the curve $M_{A,B}$ is isomorphic to the curve $M_{A,1} : y^2 = x^3 + Ax^2 + x$ over \bar{k} , but not necessarily over k . Therefore, by making the assumption that $B = 1$ we are working *up to twist*. In what follows this shall not give rise to any issues, and as remarked earlier it does not impact the k -rational points of the Kummer line (even though it does change the k -rational point of the curve itself). So from this point on we consider

$$M/k : y^2 = x^3 - \frac{a^4 + b^4}{a^2b^2}x^2 + x,$$

where $a, b \in \bar{k}^*$ such that $ab(a^4 - b^4) \neq 0$ and $a^2/b^2 \in k$.

Given this model we can define a *dual* curve. For this purpose, we define $\hat{a}, \hat{b} \in \bar{k}^*$ such that

$$2\hat{a}^2 = a^2 + b^2, \quad 2\hat{b}^2 = a^2 - b^2.$$

It is easily checked that $\hat{a}^2/\hat{b}^2 \in k^*$ and that $\hat{a}\hat{b}(\hat{a}^4 - \hat{b}^4) \neq 0$. Therefore

$$\widehat{M} : y^2 = x^3 - \frac{\hat{a}^4 + \hat{b}^4}{\hat{a}^2\hat{b}^2}x^2 + x, \quad \Delta_{\widehat{M}} = 16 \cdot \frac{(\hat{a}^4 - \hat{b}^4)^2}{\hat{a}^4\hat{b}^4}$$

is a Montgomery curve whose elements of order 2 are $(0 : 0 : 1)$, $(\hat{a}^2 : 0 : \hat{b}^2)$ and $(\hat{b}^2 : 0 : \hat{a}^2)$. We call \widehat{M} the *dual* of M . More generally, for any curve model we call the action of swapping a resp. b by \hat{a} resp. \hat{b} (and vice versa) *dualizing* (c.f. Renes and Smith [2017, §4.1]). The curves M and \widehat{M} are 2-isogenous via a 2-isogeny $\phi : M \rightarrow \widehat{M}$, and the kernel of both ϕ and $\widehat{\phi}$ is generated by the point $(0 : 0 : 1)$ on the respective curves [Renes 2018, Remark 6]. This leads to a decomposition of the doubling map [2], which we use to construct the following sequence of maps.

PROPOSITION 3.2. *Let $a, b \in \bar{k}^*$ with $ab(a^4 - b^4) \neq 0$ and $a^2/b^2 \in k^*$ and*

$$M/k : y^2 = x^3 - \frac{a^4 + b^4}{a^2b^2}x^2 + x.$$

Then there exists a commutative diagram³ of isogenies (over \bar{k})

$$\begin{array}{ccccc}
 & & (\widehat{E}, \mathcal{O}_{\widehat{E}}) & \xrightarrow{\phi_5} & (M, \mathcal{O}_M) \\
 & \nearrow \phi_4 & & \nearrow \widehat{\phi} & \searrow \phi_0 \\
 (\widehat{E}, \mathcal{O}_{\widehat{E}}) & & & & (E, \mathcal{O}_E) \\
 & \nwarrow \phi_3 & & \nwarrow \phi & \nearrow \phi_1 \\
 & & (\widehat{M}, \mathcal{O}_{\widehat{M}}) & \xleftarrow{\phi_2} & (E, \mathcal{O}_E)
 \end{array} \tag{3}$$

where

$$E/k : -x^2 + y^2 = 1 - \frac{(a^2 - b^2)^2}{(a^2 + b^2)^2} x^2 y^2, \quad E/k : x^2 + y^2 = \frac{a^2 - b^2}{a^2 + b^2} (1 + x^2 y^2)$$

and \widehat{E} and \widehat{E} are their respective duals. The maps ϕ_2 and ϕ_5 are 2-isogenies with

$$\ker(\phi_2) = \langle (0 : 0 : -\hat{b} : \hat{a}) \rangle, \quad \ker(\phi_5) = \langle (0 : 0 : -b : a) \rangle,$$

while the maps ϕ_0, ϕ_1, ϕ_3 and ϕ_4 are isomorphisms.

PROOF. We define

$$\phi_0 : (x, y) \mapsto \left(\frac{2\hat{a}^2 x}{aby}, \frac{x+1}{x-1} \right), \quad \phi_0^{-1} : (x, y) \mapsto \left(\frac{y+1}{y-1}, \frac{2\hat{a}^2(y+1)}{abx(y-1)} \right).$$

Note that this is a priori only a birational map, but naturally becomes an isomorphism when (canonically) extended to the smooth \mathbb{P}^3 model, see e.g. [Silverman 2009, Proposition II.2.1]. In particular, $\phi_0 : \mathcal{O}_M \mapsto \mathcal{O}_E, (0 : 0 : 1) \mapsto (0 : 0 : -1 : 1)$. It is similar to the maps used by Bernstein et al. [2008, Theorem 3.2(i)] and by Castryck et al. [2008], but composed with the map by Hisil et al. [2008, §3.1] to ensure a twisted Edwards curve $E_{\alpha, \delta}$ with $\alpha = -1$ that is well-defined everywhere. Moreover, we tweak it such that it acts as an involution (i.e. a Hadamard transformation) on the Kummer line. We define the isomorphism ϕ_1 as

$$\phi_1 : (x, y) \mapsto \left(-\frac{i\hat{b}}{\hat{a}}x, \frac{\hat{b}}{\hat{a}}y \right), \quad \phi_1^{-1} : (x, y) \mapsto \left(\frac{i\hat{a}}{\hat{b}}x, \frac{\hat{a}}{\hat{b}}y \right),$$

where $i \in \bar{k}$ is such that $i^2 = -1$. Then we set $\phi_2 = \phi \circ \phi_0^{-1} \circ \phi_1^{-1}$. It follows that

$$\ker(\phi_2) = \langle \phi_1 \phi_0(0, 0) \rangle = \langle (0 : 0 : -\hat{b} : \hat{a}) \rangle.$$

A completely analogous construction can be made for ϕ_3, ϕ_4 and ϕ_5 . □

REMARK 1. Note that one can argue that the above construction can be done for any sequence of isomorphisms starting at M . Indeed this is the case, but the above choice is a natural one and gives rise to nice arithmetic on the Kummer lines. Moreover, it is a choice that allows to explain the connection between Montgomery curves and the genus-1 Kummer lines arising from theta functions (i.e. [Gaudry and Lubicz 2009, §6.2] and [Karati and Sarkar 2017b, §2.4]).

The maps behave very nicely on the Kummer lines.

³ The diagram is drawn in the shape of a hexagon because its induced diagram on the Kummer lines after translations by points of order 2 is the genus-1 analogue of the hexagon in genus 2 by Renes and Smith [2017, Figure 1].

COROLLARY 3.3. *There is an induced commutative diagram of Kummer lines*

$$\begin{array}{ccccc}
 & & K_{\widehat{E}}^{O_{\widehat{E}}} & \xrightarrow{\bar{\phi}_5} & K_M^{O_M} \\
 & \nearrow \bar{\phi}_4 & & & \searrow \bar{\phi}_0 \\
 K_{\widehat{E}}^{O_{\widehat{E}}} & & & & K_E^{O_E} \\
 & \nwarrow \bar{\phi}_3 & & & \nearrow \bar{\phi}_1 \\
 & & K_M^{O_{\widehat{M}}} & \xleftarrow{\bar{\phi}_2} & K_E^{O_E}
 \end{array} \tag{4}$$

such that

$$\begin{aligned}
 \bar{\phi}_0 : (X : Z) &\mapsto (X + Z : X - Z), \\
 \bar{\phi}_1 : (X : Z) &\mapsto (\hat{b}X : \hat{a}Z), \\
 \bar{\phi}_2 : (X : Z) &\mapsto (\hat{b}^2X^2 - \hat{a}^2Z^2 : \hat{a}^2X^2 - \hat{b}^2Z^2),
 \end{aligned}$$

while $\bar{\phi}_3 = \bar{\phi}_0$ and $\bar{\phi}_4$ resp. $\bar{\phi}_5$ are obtained from $\bar{\phi}_1$ resp. $\bar{\phi}_2$ by dualizing.

PROOF. Apply the respective x and y projection maps to the curves in (3). \square

This provides clear connections between the x - and y - lines of Montgomery and (twisted) Edwards curves with full rational 2-torsion. We now show that we can use these 2-torsion points to obtain simple isomorphisms to the canonical and squared Kummer lines.

3.2 Actions of points of order 2

First recall from §2 that we have points of order 2

$$\begin{aligned}
 T &= (a^2 : 0 : b^2) \in (M, O_M), & \widehat{T} &= (\hat{a}^2 : 0 : \hat{b}^2) \in (\widehat{M}, O_{\widehat{M}}), \\
 \Omega_1 &= (\hat{a}^2 : \hat{b}^2 : 0 : 0) \in (E, O_E), & \widehat{\Omega}_1 &= (a^2 : b^2 : 0 : 0) \in (\widehat{E}, O_{\widehat{E}}), \\
 \Theta_1 &= (\hat{a} : \hat{b} : 0 : 0) \in (E, O_E), & \widehat{\Theta}_1 &= (a : b : 0 : 0) \in (\widehat{E}, O_{\widehat{E}}).
 \end{aligned}$$

One can check that these are all respective images of one another under the ϕ_i . They correspond to translations⁴ τ by the respective points which commute with the projection maps to \mathbb{P}^1 . As a result, we obtain induced involutions $\bar{\tau}$ on the Kummer lines. More concretely, we can show that

$$\begin{aligned}
 \bar{\tau}_T : (X : Z) &\mapsto (a^2X - b^2Z : b^2X - a^2Z), & \bar{\tau}_{\widehat{\Omega}_1} : (X : Z) &\mapsto (a^2Z : b^2X), \\
 \bar{\tau}_{\widehat{T}} : (X : Z) &\mapsto (\hat{a}^2X - \hat{b}^2Z : \hat{b}^2X - \hat{a}^2Z), & \bar{\tau}_{\Theta_1} : (X : Z) &\mapsto (Z : X), \\
 \bar{\tau}_{\Omega_1} : (X : Z) &\mapsto (\hat{a}^2Z : \hat{b}^2X), & \bar{\tau}_{\widehat{\Theta}_1} : (X : Z) &\mapsto (Z : X).
 \end{aligned}$$

Note that we could apply the maps τ to the diagram (3), but that requires keeping track of multiple coordinates and is somewhat tedious. Instead, for simplicity, we will focus on the Kummer lines. Applying the maps $\bar{\tau}$ to (4), we obtain the following result.

⁴ Translations are morphisms [Silverman 2009, Theorem 3.6] and are therefore isogenies if and only if they send the base point of the domain curve to the base point of the co-domain curve. For example, $\tau_T : (M, O_M) \rightarrow (M, T)$ is an isogeny. As such, it is a group homomorphism.

THEOREM 3.4. *The diagram⁵*

$$\begin{array}{ccccccc}
 & & K_E^{\Omega_1} & \xleftarrow{\bar{\psi}_4} & K_E^{\Theta_1} & \xrightarrow{\bar{\psi}_5} & K_M^T \\
 & \nearrow \bar{\tau}_{\Omega_1} & & & \nearrow \bar{\tau}_{\Theta_1} & & \nearrow \bar{\tau}_T \\
 K_E^{O_{\hat{E}}} & \xleftarrow{\bar{\psi}_3} & K_E^{O_{\hat{E}}} & \xrightarrow{\bar{\phi}_4} & K_E^{O_{\hat{E}}} & \xrightarrow{\bar{\phi}_5} & K_M^{O_M} \\
 & \downarrow \bar{\phi}_3 & & & \downarrow \bar{\phi}_0 & & \downarrow \bar{\psi}_0 \\
 & & K_M^T & \xleftarrow{\bar{\psi}_2} & K_E^{\Theta_1} & \xrightarrow{\bar{\psi}_1} & K_E^{\Omega_1} \\
 & \nearrow \bar{\tau}_{\hat{T}} & & & \nearrow \bar{\tau}_{\Theta_1} & & \nearrow \bar{\tau}_{\Omega_1} \\
 K_M^{O_{\hat{M}}} & \xleftarrow{\bar{\phi}_2} & K_E^{O_{\hat{E}}} & \xrightarrow{\bar{\phi}_1} & K_E^{O_E} & &
 \end{array} \tag{5}$$

is commutative, where

$$\begin{aligned}
 \bar{\psi}_0 : (X : Z) &\mapsto (X + Z : X - Z), & \bar{\psi}_3 : (X : Z) &\mapsto (X + Z : X - Z), \\
 \bar{\psi}_1 : (X : Z) &\mapsto (\hat{b}X : \hat{a}Z), & \bar{\psi}_4 : (X : Z) &\mapsto (bX : aZ), \\
 \bar{\psi}_2 : (X : Z) &\mapsto (X^2 : Z^2), & \bar{\psi}_5 : (X : Z) &\mapsto (X^2 : Z^2),
 \end{aligned}$$

and every \leftrightarrow is an isomorphism.

PROOF. This is the diagram from (3) translated by corresponding points of order 2 through the different τ , projected to their respective Kummer lines. We construct

$$\bar{\psi}_0 = \bar{\tau}_{\Omega_1} \circ \bar{\phi}_0 \circ \bar{\tau}_T$$

and proceed similarly for the other $\bar{\psi}_i$. \square

Recall that (the duals of) $K_M^{O_M}, K_E^{O_E}$ resp. $K_E^{O_{\hat{E}}}$ are the Kummer lines of (the duals of) a Montgomery, twisted Edwards resp. Edwards curve. Hence it remains to identify $K_M^T, K_E^{\Omega_1}$ and $K_E^{\Theta_1}$ (and their duals). Since they are all simply \mathbb{P}^1 as an algebraic variety, we analyze their (pseudo-)addition formulae.

First note that Proposition 3.1 tells us that moving through the sequence $\bar{\phi}_0, \dots, \bar{\phi}_5$ corresponds to the [2] map (starting at any of the $\bar{\phi}_i$). Since the $\bar{\tau}$ are isomorphisms, the same is true for $\bar{\psi}_0, \dots, \bar{\psi}_5$. In other words, for example

$$\begin{aligned}
 [2] &= \bar{\psi}_5 \circ \dots \circ \bar{\psi}_0 \text{ on } K_M^T, \\
 [2] &= \bar{\psi}_4 \circ \dots \circ \bar{\psi}_0 \circ \bar{\psi}_5 \text{ on } K_E^{\Theta_1}.
 \end{aligned}$$

Comparing these with the algorithm from Gaudry and Lubicz [2009, §6.2] (and the formulas also appear in [Bernstein and Lange 2015]) reveals that these are the doubling formulae for the squared and canonical Kummer lines. One readily⁶ verifies that the same is true for the differential addition formulae. The third Kummer line $K_E^{\Omega_1}$ has not appeared to our knowledge, and has similar arithmetic to the squared Kummer line. We refer to it as the *intermediate* Kummer, c.f. [Renes and Smith 2017, §4.3]. Interestingly, it appears as the y -line of a twisted Edwards curve where the coefficient of

⁵ This is no longer a hexagon due to the authors' inability to draw it in a readable way.

⁶ This can be done by using the known addition formulae on the elliptic curves whose identities are at infinity, and composing with the translation and projection maps. This is somewhat tedious, but is relatively straightforward by using a computer algebra package [Bosma et al. 1997; The Sage Developers 2018].

Table 1. Kummer lines over a finite field \mathbb{F}_q and their associated (i) squared Kummer ($a^2 : b^2$) (ii) Montgomery A (iii) twisted Edwards δ and (iv) Edwards c^2 constants.

q	$(a^2 : b^2)$	$(A : 1)$	$(\delta : 1)$	$(c^2 : 1)$
$2^{251} - 9$	(81 : 20)	(-6961 : 1620)	(-3721 : 10201)	(61 : 101)
$2^{251} - 9$	(186 : 175)	(-65221 : 130200)	(-121 : 130221)	(11 : 361)
$2^{255} - 19$	(82 : 77)	(-12653 : 6314)	(-25 : 25281)	(5 : 159)
$2^{266} - 3$	(260 : 139)	(-86921 : 36140)	(-14641 : 159201)	(121 : 399)

x^2 is -1 , in which case the optimal formulas by Hisil et al. [2008] are available. For completeness, we summarize the associated curve constants for the instances provided by Karati and Sarkar in Table 1, connecting the squared Kummer line to the Kummer lines of Montgomery and twisted Edwards models via isomorphisms (as opposed to birational maps or isogenies).

REMARK 2. *We reiterate that only the intermediate Kummer line is new, while all the others have already appeared in the literature and are well-known. However, there had been little work in providing explicit maps between them, and this is exactly what we provide.*

3.3 Hybrid Kummer lines

Since the arithmetic on these Kummer lines is generally well-studied, the (cryptographic) value of this study does not come from improved operation counts. Beside its theoretical contribution, we ease the problem of selecting which curves to use for best performance (e.g. for standardization). That is, the simplicity of the isomorphisms gives quasi-cost-free transformations that allow interchangeable usage of any of the models. This is similar to the usage of a birational map to move between the Montgomery and twisted Edwards model, but we extend it with the squared Kummer line. We summarize this in Figure 1. In particular, Karati and Sarkar [2017b] show the benefits of the squared Kummer line on platforms where SIMD instructions are available.

$$\begin{array}{ccccc}
 & & & & E \\
 & & & & \downarrow y \\
 K_M^T & \xleftarrow{(a^2X - b^2Z : b^2X - a^2Z)} & K_M^{O_M} & \xleftarrow{(X+Z : X-Z)} & K_E^{O_E}
 \end{array}$$

Fig. 1. The squared Kummer line, the x -line of a Montgomery curve and the y -line of a twisted Edwards curve E , connected by involutions.

REMARK 3. *Recall that all the above works under the assumption of having full rational 2-torsion. Although Montgomery and (twisted) Edwards curves always have a group order divisible by 4, it does not necessarily mean that they have full 2-torsion (i.e. they could have a point of order 4). Note that standardized curves such as Curve25519 and Curve448 do not have full 2-torsion, so this theory does not directly apply.*

Moreover, results from the well-studied Montgomery model immediately carry over to the squared Kummer line. For example, we can straightforwardly fit a (squared) Kummer line into the qDSA signature scheme. For signature verification, given $\mathbf{x}(P), \mathbf{x}(Q), \mathbf{x}(R) \in K_M^T$ we must be able to

Table 2. Comparison of an implementation of the qDSA signature scheme based on Curve25519 and the Montgomery model of the squared Kummer line defined by $(a^2, b^2) = (159, 5)$, where the memory is measured in bytes (B).

Ref.	Object	Constant	Clock cycles	Stack	Code
RS [2017]	Curve25519	$(A + 2 : 4) = (121666 : 1)$	3 889 116 (sign) 6 793 695 (verify)	660 B 788 B	18 443 B
This	$K_M^{O_M}$	$(A + 2 : 4) = (-5929 : 795)$	3 916 879 (sign) 6 857 007 (verify)	660 B 788 B	18 391 B
This	K_M^T	$(a^2 : b^2) = (159 : 5)$	3 824 857 (sign) 6 673 039 (verify)	660 B 788 B	18 557 B

check whether $\mathbf{x}(R) = \mathbf{x}(P \pm Q)$. Although this can certainly be directly defined on K_M^T , we note that it is equivalent to checking whether

$$\bar{\tau}_T(\mathbf{x}(R)) = \bar{\tau}_T(\mathbf{x}(P \pm Q)).$$

This is simply the function $\text{Check}(\bar{\tau}_T(\mathbf{x}(P)), \bar{\tau}_T(\mathbf{x}(Q)), \bar{\tau}_T(\mathbf{x}(R)))$, where Check is defined in [Renes and Smith 2017, Algorithm 2].

To demonstrate feasibility of this approach, we extend the publicly available Curve25519-based instantiation of qDSA from Renes and Smith [2017] on the ARM Cortex M0 architecture⁷. For this purpose we choose a squared Kummer line over $\mathbb{F}_{2^{255}-19}$, allowing field arithmetic to remain essentially unchanged. A notable exception to this is an efficient assembly implementation of 16×256 -bit field multiplication, which is used for the multiplications by the line constants. This replaces the highly optimized multiplication by 121666 from D  ll et al. [2015]. We select $(a^2, b^2) = (159, 5)$, so that the squared Kummer line K_M^T corresponds to the *dual*⁸ of KL25519(82, 77) presented and implemented by Karati and Sarkar [2017b]. This implies the Montgomery constant of the curve above the line $K_M^{O_M}$ to be $(A + 2 : 4) = (-5929 : 795)$. We summarize the implementation results in Table 2. We emphasize that the point of this work is not to provide the most efficient implementation for this given platform, but rather to show the close connection between the different Kummer lines. Although on this platform results differ only by a minimal margin, the difference can be much larger on other devices (in particular, when SIMD instructions are available). Our isomorphisms allow an implementer to select the model that is most appropriate for a given architecture.

REMARK 4. *The implementations that we present are constant-time, and all standard countermeasures (e.g. projective blinding, scalar blinding [Coron 1999, §5]) against more advanced side-channel and fault attacks can be applied if required. In particular, as mentioned by the authors, the recent fault attack by Takahashi et al. [2018] can (cheaply) be thwarted by requiring nonces to be multiples of the cofactor (i.e. by “clamping”). However, such countermeasures are only necessary when an implementation is used in a context where fault attacks are considered part of the attacker model. We emphasize that our implementation is intended as a reference and not for production use.*

⁷ All code (including reference implementations in C) is available in the public domain at <http://www.cs.ru.nl/~jrenes/>.

⁸ The constants $(a^2, b^2) = (88, 77)$ lead to $(A + 2 : 4) = (-25 : 25256)$ which has slightly larger constants on $K_M^{O_M}$ than its dual. However, results should be very similar.

4 ISOMORPHISM CLASSES OVER FINITE FIELDS

For cryptographic purposes, we are mostly concerned with the case that $k = \mathbb{F}_q$, for some prime (power) q . As using extension fields is generally expensive, we would like to set things up such that all computation is performed in \mathbb{F}_q . Whether or not we can do this in a way such that constants remain small, depends on the number of Kummer lines that exist. Following earlier studies on the number of isomorphism classes for certain curve models [Bernstein et al. 2008; Farashahi et al. 2012; Farashahi and Shparlinski 2010], we provide counts for the canonical, squared and intermediate Kummer lines.

4.1 Identifying Kummer lines

For this purpose it is interesting to ask when two Kummer lines should be considered to be the same. Given two Kummer lines $K_1 = E_1/\{\pm 1\}$ and $K_2 = E_2/\{\pm 1\}$ of elliptic curves E_1, E_2 defined over \mathbb{F}_q , it could be natural to identify K_1 with K_2 whenever E_1 is \mathbb{F}_q -isomorphic to E_2 . However, as noted in §2, the arithmetic on the \mathbb{F}_q -rational points of the Kummer lines will be identical whenever E_1 is \mathbb{F}_{q^2} -isomorphic to E_2 (i.e. E_2 is the quadratic twist of E_1). Since the curves are defined over \mathbb{F}_q , this will happen if and only if $j(E_1) = j(E_2)$. As such, we equate the number of Kummer lines with the number of elliptic curves defined over \mathbb{F}_q up to $\bar{\mathbb{F}}_q$ -isomorphism.

Recall that we parametrize Kummer lines by $a, b \in \bar{\mathbb{F}}_q$ such that $ab(a^4 - b^4) \neq 0$ and $a^2/b^2 \in \mathbb{F}_q$. Since $b \neq 0$, a Kummer line is defined by the fraction a/b or, equivalently, by the point $(a : b) \in \mathbb{P}^1$. Again, since $b \neq 0$ we can therefore simply assume $b = 1$. As such, we can consider $a \in \bar{\mathbb{F}}_q$ such that $a^2 \in \mathbb{F}_q$ and $a^5 - a \neq 0$.

4.2 Canonical Kummer lines

We begin by considering the canonical Kummer line from Gaudry and Lubicz [2009] defined by some a as above. Recall that it corresponds to the y -line of the curve

$$\widehat{E}/\mathbb{F}_q : x^2 + y^2 = \frac{1}{a^2} (1 + x^2 y^2) .$$

with identity $\widehat{\Omega}_1 = (a : 1 : 0 : 0)$, whose image in \mathbb{P}^1 is $(a : 1)$. Therefore, we certainly require that $a \in \mathbb{F}_q$. It is easily seen that $\hat{a}^2, \hat{b}^2 \in \mathbb{F}_q$ and that this is enough to perform all arithmetic with \mathbb{F}_q operations.

Now note that $(\widehat{E}, \widehat{\Omega}_1)$ is \mathbb{F}_q -isomorphic to $(\widehat{E}, O_{\widehat{E}})$ via $\tau_{\widehat{\Omega}_1}$, which is an Edwards curve if and only if $a \in \mathbb{F}_q$ and $1/a^5 \neq 1/a$. The first is true by assumption, while the latter follows from $a^5 \neq a$. Therefore we simply count the number of Edwards curves defined over \mathbb{F}_q up to $\bar{\mathbb{F}}_q$ -isomorphism. A result by Farashahi and Shparlinski [2010, Theorem 5] shows that there are exactly

$$\begin{cases} \left\lfloor \frac{q+23}{24} \right\rfloor & \text{if } q \equiv 1, 9, 13, 17 \pmod{24}, \\ \left\lfloor \frac{q-5}{24} \right\rfloor & \text{if } q \equiv 5 \pmod{24}, \\ \left\lfloor \frac{q+1}{8} \right\rfloor & \text{if } q \equiv 3 \pmod{4}. \end{cases}$$

Thus, in general there will be no problem to find Kummer lines with the desired security properties. However, it may not be easy to find them such that its constants are small. For that reason, we look towards the squared and intermediate Kummer lines.

4.3 Squared and intermediate Kummer lines

If we use canonical Kummer lines, we restrict ourselves to $a \in \mathbb{F}_q$ for all of the arithmetic to be in \mathbb{F}_q . This (seemingly) limits the number of Kummer lines that we can use. This is no longer the case on squared and intermediate Kummer lines; it suffices to only have $a^2 \in \mathbb{F}_q$. Note that this implies that $a \in \mathbb{F}_{q^2}$.

Since the j -invariants of M , E and \mathcal{E} and their duals are all equal, we can count the number of curves up to isomorphism of the form

$$\widehat{\mathcal{E}} : x^2 + y^2 = \frac{1}{a^2} (1 + x^2 y^2)$$

such that $a^5 \neq a$ (but note that $\widehat{\mathcal{E}}$ is not necessarily an Edwards curve over \mathbb{F}_q). There are exactly $q - 3$ such curves, so it remains to determine how many are in the same \mathbb{F}_q -isomorphism class. This question has already been considered by Edwards [2007, Proposition 6.1], whose statement implies that two Edwards curves determined by $a^2, \bar{a}^2 \in \mathbb{F}_q$ have the same j -invariant whenever \bar{a}^2 is one of the following:

$$\pm a^2, \pm \frac{1}{a^2}, \pm \left(\frac{a-1}{a+1}\right)^2, \pm \left(\frac{a+1}{a-1}\right)^2, \pm \left(\frac{a-i}{a+i}\right)^2, \pm \left(\frac{a+i}{a-i}\right)^2. \quad (6)$$

If $q \equiv 1 \pmod{4}$, then $i^q = i$ and a straightforward computation show that

$$\pm \left(\frac{a-1}{a+1}\right)^2, \pm \left(\frac{a+1}{a-1}\right)^2, \pm \left(\frac{a-i}{a+i}\right)^2, \pm \left(\frac{a+i}{a-i}\right)^2 \in \mathbb{F}_q \iff a \in \mathbb{F}_q.$$

If $q \equiv 3 \pmod{4}$, then $i^q = -i$ and a similar computation shows that

$$\begin{aligned} \pm \left(\frac{a-1}{a+1}\right)^2, \pm \left(\frac{a+1}{a-1}\right)^2 &\in \mathbb{F}_q \iff a \in \mathbb{F}_q, \\ \pm \left(\frac{a-i}{a+i}\right)^2, \pm \left(\frac{a+i}{a-i}\right)^2 &\in \mathbb{F}_q \iff i \cdot a \in \mathbb{F}_q. \end{aligned}$$

Given that either $a \in \mathbb{F}_q$ or $i \cdot a \in \mathbb{F}_q$, while half the elements of \mathbb{F}_q are squares, we closely approximate⁹ that the number of isomorphism classes is

$$\approx \begin{cases} \left\lfloor \left(\left(\frac{1}{4} + \frac{1}{12} \right) \frac{q}{2} \right) \right\rfloor = \left\lfloor \frac{q}{6} \right\rfloor & \text{if } q \equiv 1 \pmod{4}, \\ \left\lfloor \left(\left(\frac{1}{8} + \frac{1}{8} \right) \frac{q}{2} \right) \right\rfloor = \left\lfloor \frac{q}{8} \right\rfloor & \text{if } q \equiv 3 \pmod{4}. \end{cases}$$

A more careful analysis c.f. [Farashahi and Shparlinski 2010] could be done, but such a close estimate suffices for our purposes. Interestingly, for $q \equiv 3 \pmod{4}$ the number of canonical and squared Kummer lines is about the same. Thus although $a^2 \in \mathbb{F}_q$ is a weaker restriction than $a \in \mathbb{F}_q$, it does not actually lead to more Kummer lines (up to isomorphism). This is explained by the fact that -1 is a non-square since $q \equiv 3 \pmod{4}$, hence exactly one of a^2 or $-a^2$ must be a square in \mathbb{F}_q , while their corresponding Edwards curves are isomorphic. For $q \equiv 1 \pmod{4}$ there is a clear difference in the number of Kummer lines, so in that case there is a significant advantage in finding small parameters for a squared or intermediate Kummer line over a canonical Kummer line.

⁹This statement is exact up to the observation that some of the elements in (6) can be the same, which happens only exceptionally.

5 CONCLUSION

This paper fills in the gaps between several existing results in the literature. We show that Kummer lines are very closely related to the traditional Montgomery and (twisted) Edwards models via translations by (rational) points of order 2. In particular, we demonstrate that the use of the cryptographically unconventional and sub-optimal Legendre model (as by e.g. [Karati and Sarkar 2017a] and [Gaudry and Lubicz 2009]) is unnecessary and can be replaced by the arithmetically more efficient Montgomery and (twisted) Edwards models. Concretely, we provide the constants for these models for all Kummer lines appearing in the work of Karati and Sarkar [2017a] in Table 1 and work out simple and efficient maps to move between them in Theorem 3.4. They provide guidance for implementers and significantly simplify their work. As an example, we provide details of an implementation of the qDSA signature scheme based on a squared Kummer line in §3.3.

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