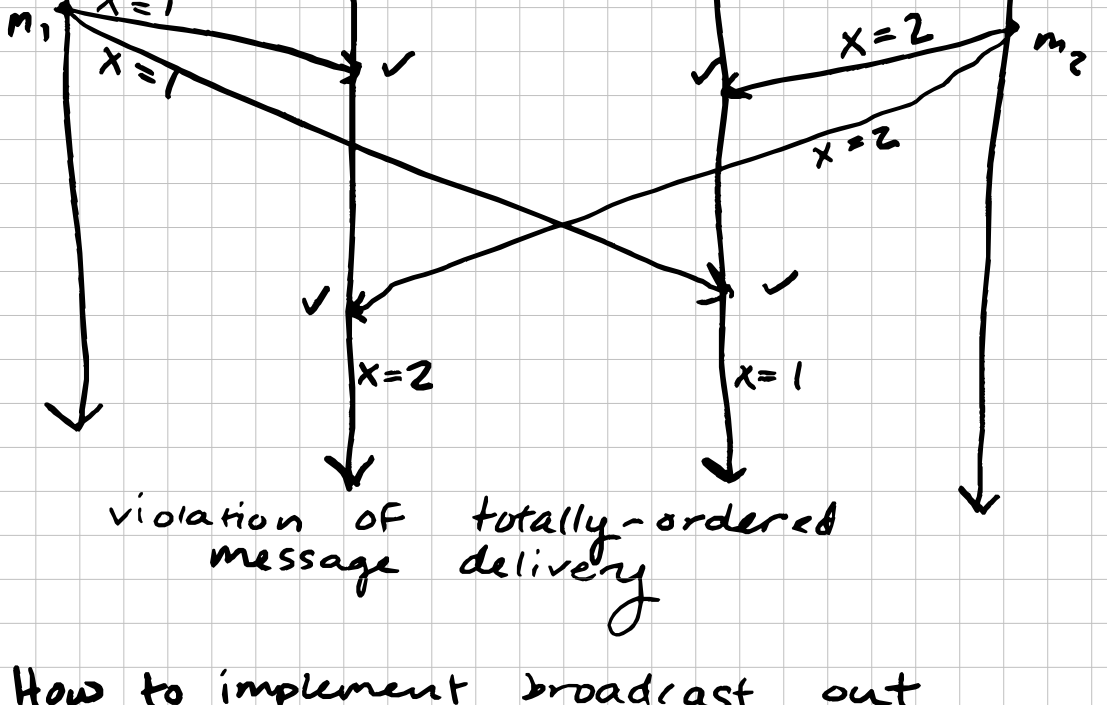
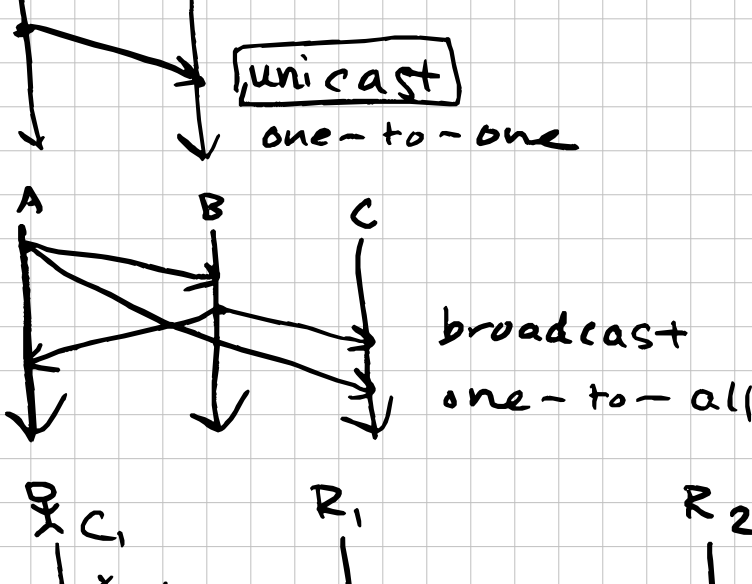
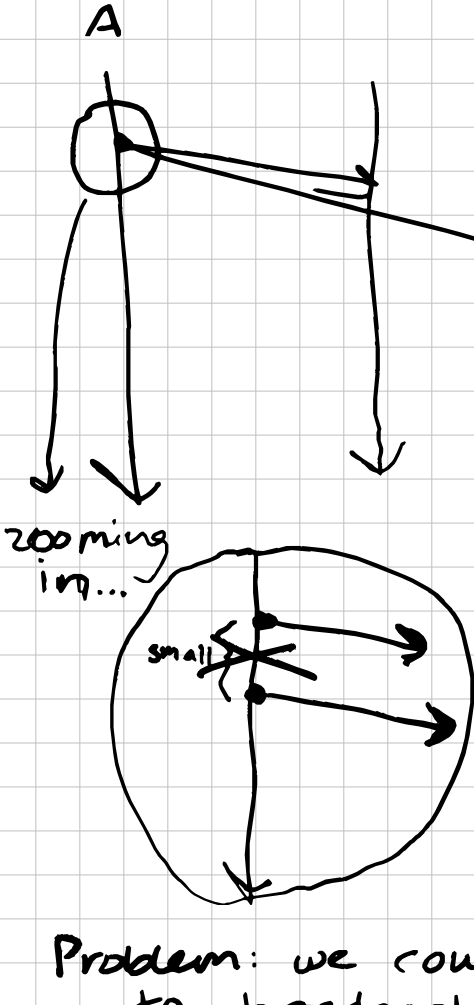


CSE138 Lecture 9

- ✓ - reliable broadcast ← define & implement
- ✓ - intro to replication
- review for the midterm

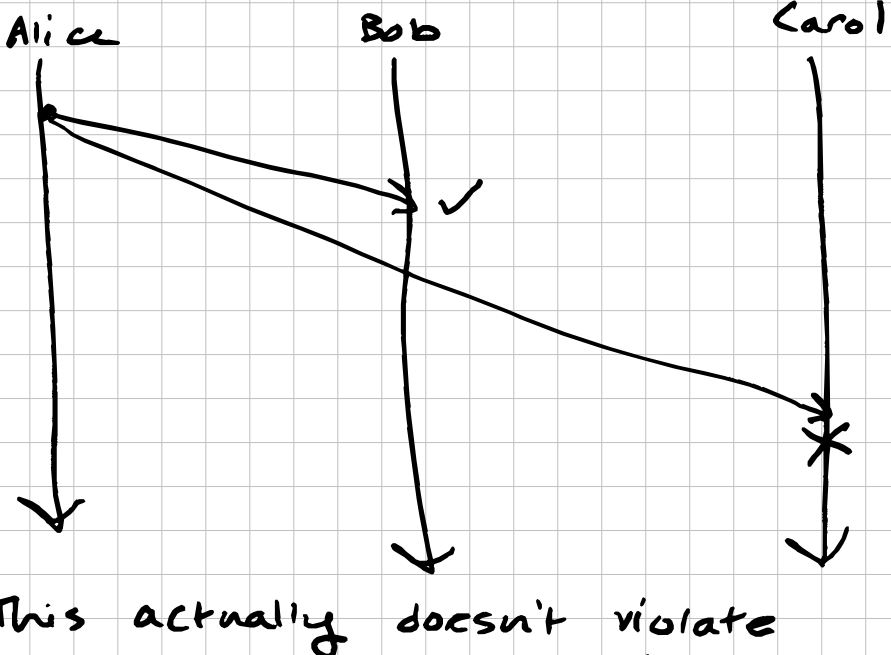


How to implement broadcast out of unicast?



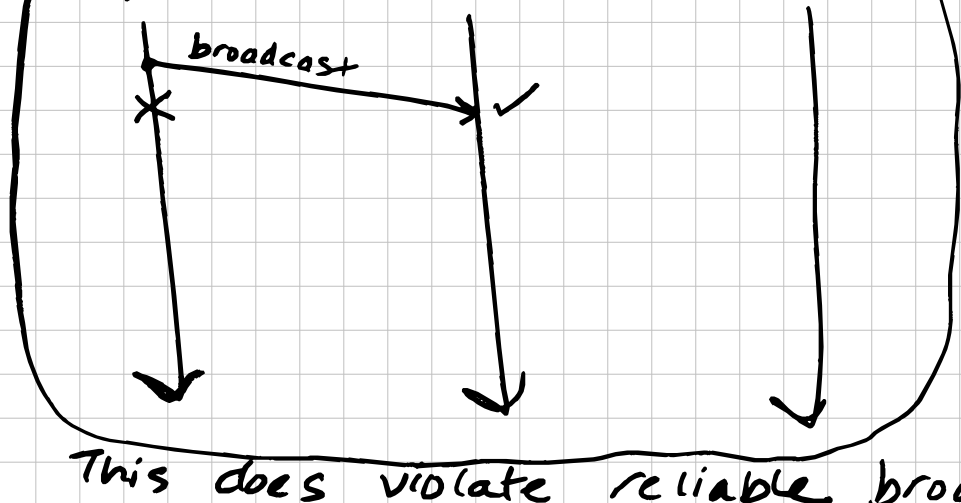
Problem: we could crash while trying to broadcast.

Reliable broadcast : If a correct process delivers a broadcasted message m , then all correct processes deliver m .

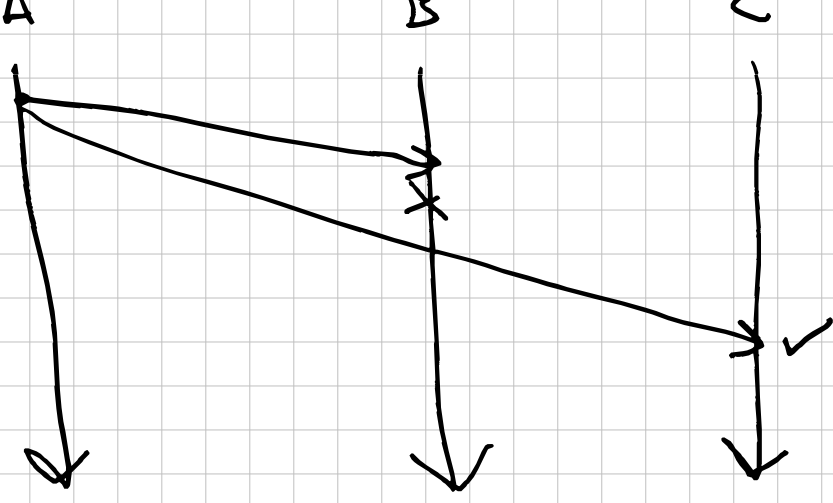


This actually doesn't violate reliable broadcast!

(Carol crashed, so she's not a correct process.)

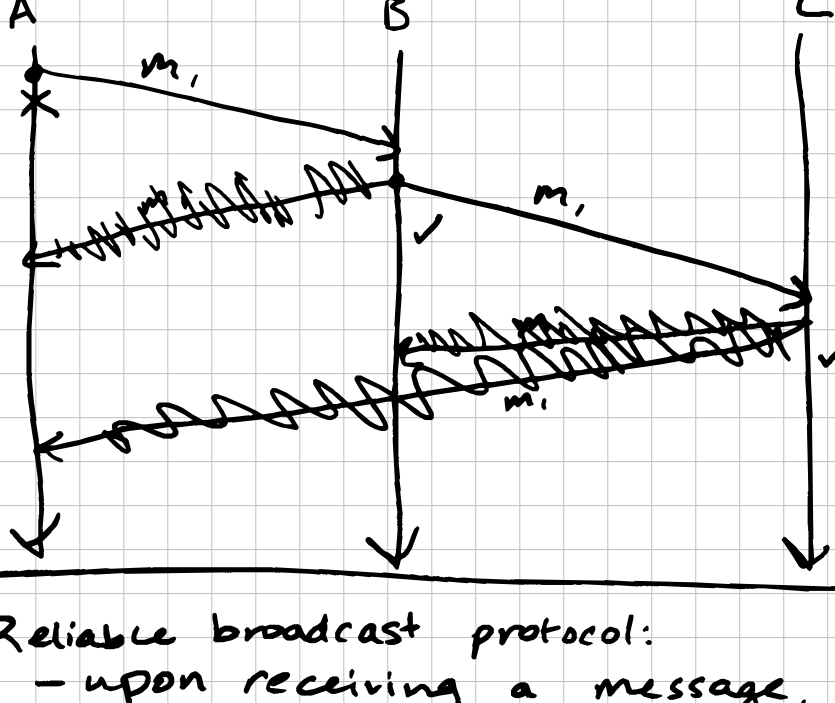


This does violate reliable broadcast.



This doesn't violate reliable broadcast.

How about this plan?

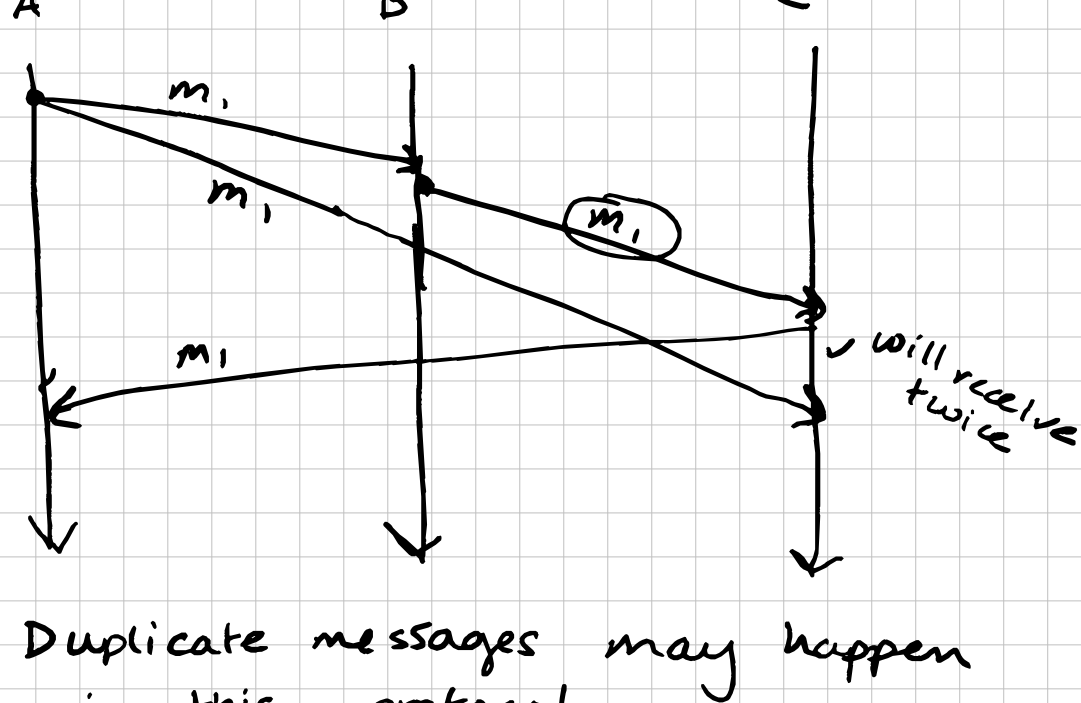


Reliable broadcast protocol:
- upon receiving a message, don't deliver it immediately. instead, broadcast it to everyone else! then deliver it.

* except whoever sent it to you



in this scenario, no correct process delivers m_1 . So, we're not violating reliable broadcast.



Duplicate messages may happen in this protocol.

Reliable broadcast is sometimes called atomic broadcast

all-or-nothing! either everyone delivers the broadcasted message, or nobody does.

Fault tolerance often involves making copies.

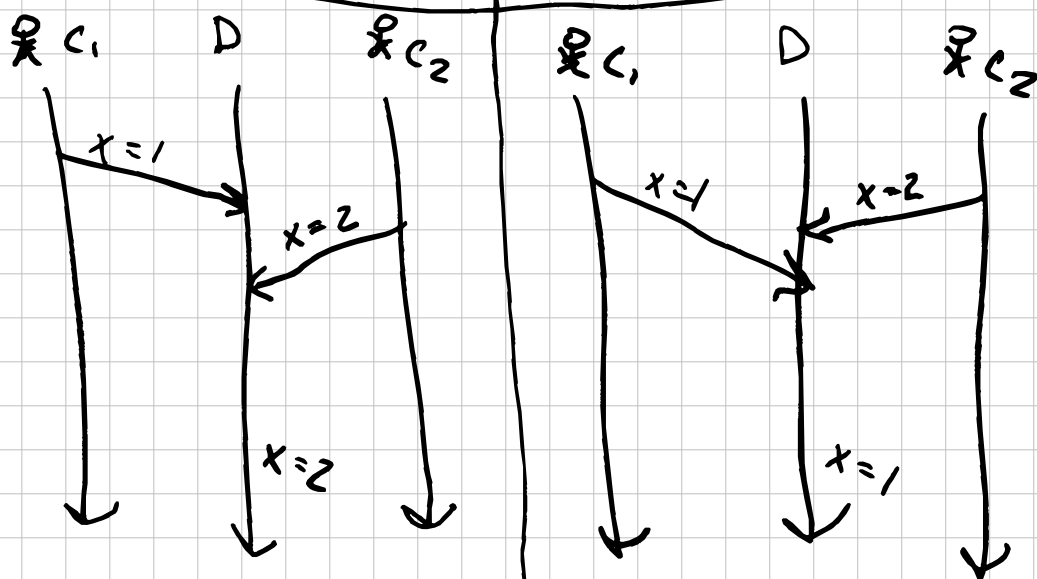
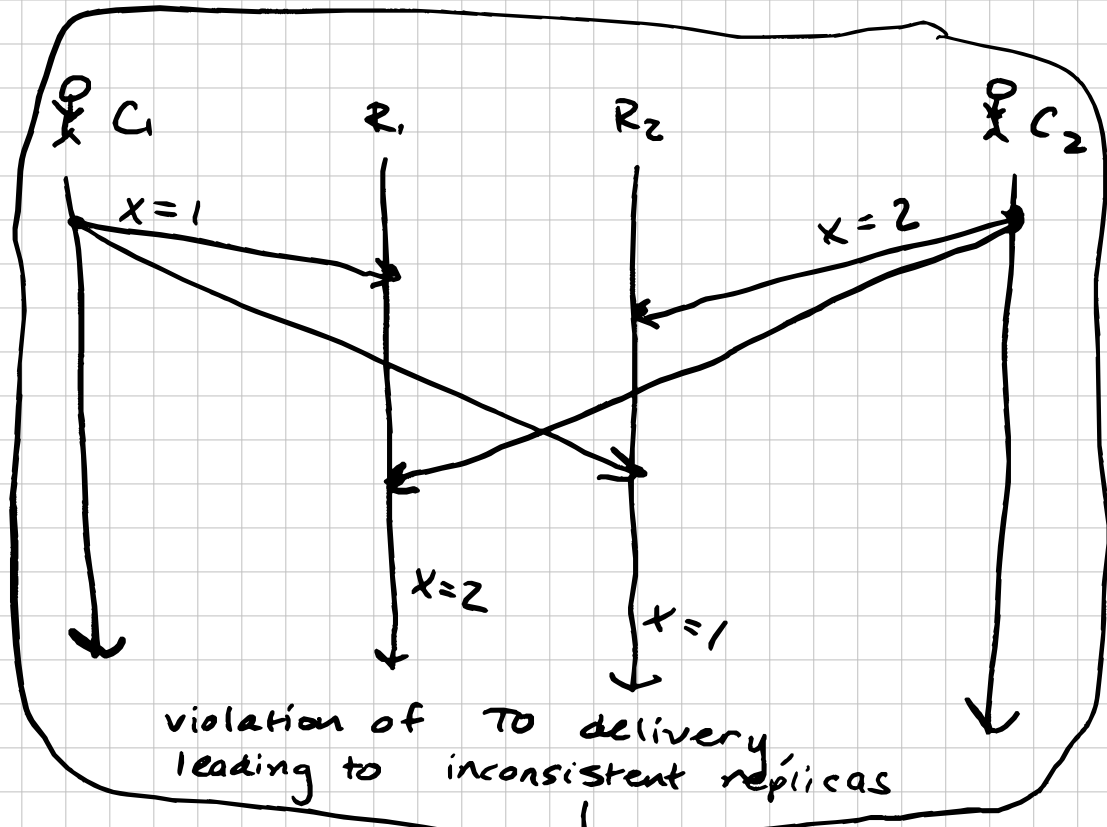
reliable delivery/ reliable broadcast involved copies of messages.

REPLICATION (making copies of data)

Replication : reasons to do it

- guard against data loss
- data locality (keep data close to clients that need it)
- dividing up the work (scalability) : handling more requests at once.

availability



Not a problem of consistency, rather, a problem of determinism.

determinism is a property that relates different runs of a system.

↑ i.e. a hyperproperty

OK, back to consistency.

informally: a replicated storage system is strongly consistent if clients can't tell that the data is replicated.

↑ Somehow we have to make everyone deliver messages in the same order.

if you want a formal definition

linearizability ← Herlihy & Wing